

A Restricted Definition of the Magic Wand to Soundly Combine Fractions of a Wand

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Abstract

Many separation logics support *fractional permissions* [1] to distinguish between read and write access to a heap location, for instance, to allow concurrent reads while enforcing exclusive writes. The concept has been generalized to fractional assertions [2, 5, 6, 3]. A^p (where A is a separation logic assertion and p a fraction between 0 and 1) represents a fraction p of A . A^p holds in a state σ iff there exists a state σ_A in which A holds and σ is obtained from σ_A by multiplying all permission amounts held by p .

While A^{p+q} can always be split into $A^p * A^q$, recombining $A^p * A^q$ into A^{p+q} is not always sound. We say that A is *combinable* iff the entailment $A^p * A^q \models A^{p+q}$ holds for any two positive fractions p and q such that $p + q \leq 1$. Combinable assertions are particularly useful to reason about concurrent programs, for instance, to combine the post-conditions of parallel branches when they terminate. Unfortunately, the magic wand assertion $A \multimap B$, commonly used to specify properties of partial data structures, is typically *not* combinable.

In this entry, we formalize a novel, restricted definition of the magic wand, described in a paper at CAV 22 [4], which we call the *combinable wand*. We prove some key properties of the combinable wand; in particular, a combinable wand is combinable if its right-hand side is combinable.

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1 State Model with Fractional Permissions

In this section, we define a concrete state model based on fractional permissions [1]. A state is a pair of a permission mask and a partial heap. A permission mask is a total map from heap locations to a rational between 0 and 1 (included), while a partial heap is a partial map from heap locations to values. We also define a partial addition on these states, and show that this state model corresponds to a separation algebra.

1.1 Non-negative rationals (permission amounts)

```
theory PosRat
  imports Main HOL.Rat
begin
```

1.1.1 Definitions

```
typedef prat = { r :: rat | r. r ≥ 0 } by fastforce
```

```
setup-lifting type-definition-prat
```

```
lift-definition pwrite :: prat is 1 by simp
```

```
lift-definition pnone :: prat is 0 by simp
```

```
lift-definition half :: prat is 1 / 2 by fastforce
```

```
lift-definition pgte :: prat ⇒ prat ⇒ bool is (≥) done
```

```
lift-definition pgt :: prat ⇒ prat ⇒ bool is (>) done
```

```
lift-definition lt :: prat ⇒ prat ⇒ bool is (<) done
```

```
lift-definition ppos :: prat ⇒ bool is λp. p > 0 done
```

```
lift-definition pmult :: prat ⇒ prat ⇒ prat is (*) by simp
```

```
lift-definition padd :: prat ⇒ prat ⇒ prat is (+) by simp
```

lift-definition $pdiv :: prat \Rightarrow prat \Rightarrow prat$ is (/) by simp

lift-definition $pmin :: prat \Rightarrow prat \Rightarrow prat$ is (min) by simp

lift-definition $pmax :: prat \Rightarrow prat \Rightarrow prat$ is (max) by simp

1.1.2 Lemmas

lemma $pmin-comm$:

$pmin\ a\ b = pmin\ b\ a$

by (metis Rep-prat-inverse min.commute pmin.rep-eq)

lemma $pmin-greater$:

$pgte\ a\ (pmin\ a\ b)$

by (simp add: pgte.rep-eq pmin.rep-eq)

lemma $pmin-is$:

assumes $pgte\ a\ b$

shows $pmin\ a\ b = b$

by (metis Rep-prat-inject assms min-absorb2 pgte.rep-eq pmin.rep-eq)

lemma $pmax-comm$:

$pmax\ a\ b = pmax\ b\ a$

by (metis Rep-prat-inverse max.commute pmax.rep-eq)

lemma $pmax-smaller$:

$pgte\ (pmax\ a\ b)\ a$

by (simp add: pgte.rep-eq pmax.rep-eq)

lemma $pmax-is$:

assumes $pgte\ a\ b$

shows $pmax\ a\ b = a$

by (metis Rep-prat-inject assms max-absorb-iff1 pgte.rep-eq pmax.rep-eq)

lemma $pmax-is-smaller$:

assumes $pgte\ x\ a$

and $pgte\ x\ b$

shows $pgte\ x\ (pmax\ a\ b)$

proof (cases $pgte\ a\ b$)

case True

then show ?thesis

by (simp add: assms(1) pmax-is)

next

case False

then show ?thesis

using assms(2) pgte.rep-eq pmax.rep-eq by auto

qed

lemma *half-between-0-1*:
ppos half \wedge *pgt pwrite half*
by (*simp add: half.rep-eq pgt.rep-eq ppos.rep-eq pwrite.rep-eq*)

lemma *pgt-implies-pgte*:
assumes *pgt a b*
shows *pgte a b*
by (*meson assms less-imp-le pgt.rep-eq pgte.rep-eq*)

lemma *half-plus-half*:
padd half half = pwrite
by (*metis Rep-prat-inject divide-less-eq-numeral1(1) dual-order.irrefl half.rep-eq less-divide-eq-numeral1(1) linorder-neqE-linordered-idom mult.right-neutral one-add-one padd.rep-eq pwrite.rep-eq ring-class.ring-distrib(1)*)

lemma *padd-comm*:
padd a b = padd b a
by (*metis Rep-prat-inject add.commute padd.rep-eq*)

lemma *padd-asso*:
padd (padd a b) c = padd a (padd b c)
by (*metis Rep-prat-inverse group-cancel.add1 padd.rep-eq*)

lemma *p-greater-exists*:
pgte a b \longleftrightarrow ($\exists r. a = padd\ b\ r$)
proof (*rule iffI*)
show *pgte a b* \implies $\exists r. a = padd\ b\ r$
proof –
assume *asm: pgte a b*
obtain *aa bb* **where** *aa = Rep-prat a bb = Rep-prat b*
by *simp*
then have *aa* \geq *bb* **using** *asm*
using *pgte.rep-eq* **by** *fastforce*
then obtain *rr* **where** *rr = aa - bb*
by *simp*
then have *a = padd b (Abs-prat rr)*
by (*metis Abs-prat-inverse Rep-prat-inject* $\langle aa = Rep-prat\ a \rangle$ $\langle bb = Rep-prat\ b \rangle$ $\langle bb \leq aa \rangle$ *add.commute diff-add-cancel diff-ge-0-iff-ge mem-Collect-eq padd.rep-eq*)
then show $\exists r. a = padd\ b\ r$ **by** *blast*
qed
show $\exists r. a = padd\ b\ r \implies pgte\ a\ b$
using *Rep-prat padd.rep-eq pgte.rep-eq* **by** *force*
qed

lemma *pgte-antisym*:
assumes *pgte a b*

and $pgte\ b\ a$
shows $a = b$
by (*metis Rep-prat-inverse assms(1) assms(2) leD le-less pgte.rep-eq*)

lemma *greater-sum-both*:
assumes $pgte\ a\ (padd\ b\ c)$
shows $\exists a1\ a2. a = padd\ a1\ a2 \wedge pgte\ a1\ b \wedge pgte\ a2\ c$
proof –
obtain $aa\ bb\ cc$ **where** $aa = Rep-prat\ a\ bb = Rep-prat\ b\ cc = Rep-prat\ c$
by *simp*
then have $aa \geq bb + cc$
using *assms padd.rep-eq pgte.rep-eq* **by** *auto*
then obtain x **where** $aa = bb + x\ x \geq cc$
by (*metis add.commute add-le-cancel-left diff-add-cancel*)
then show *?thesis*
by (*metis assms order-refl p-greater-exists padd-asso pgte.rep-eq*)
qed

lemma *padd-cancellative*:
assumes $a = padd\ x\ b$
and $a = padd\ y\ b$
shows $x = y$
by (*metis Rep-prat-inject add-le-cancel-right assms(1) assms(2) leD less-eq-rat-def padd.rep-eq*)

lemma *not-pgte-charact*:
 $\neg pgte\ a\ b \longleftrightarrow pgt\ b\ a$
by (*meson not-less pgt.rep-eq pgte.rep-eq*)

lemma *pgte-pgt*:
assumes $pgt\ a\ b$
and $pgte\ c\ d$
shows $pgt\ (padd\ a\ c)\ (padd\ b\ d)$
using *assms(1) assms(2) padd.rep-eq pgt.rep-eq pgte.rep-eq* **by** *auto*

lemma *pmult-distr*:
 $pmult\ a\ (padd\ b\ c) = padd\ (pmult\ a\ b)\ (pmult\ a\ c)$
by (*metis Rep-prat-inject distrib-left padd.rep-eq pmult.rep-eq*)

lemma *pmult-comm*:
 $pmult\ a\ b = pmult\ b\ a$
by (*metis Rep-prat-inject mult.commute pmult.rep-eq*)

lemma *pmult-special*:
 $pmult\ pwrite\ x = x$
 $pmult\ pnone\ x = pnone$
apply (*metis Rep-prat-inverse comm-monoid-mult-class.mult-1 pmult.rep-eq pwrite.rep-eq*)

by (metis Rep-prat-inverse mult-zero-left pmult.rep-eq pnone.rep-eq)

definition *pinv* where

pinv *p* = *pdiv* *pwrite* *p*

lemma *pinv-double-half*:

assumes *ppos* *p*

shows *pmult* *half* (*pinv* *p*) = *pinv* (*padd* *p* *p*)

proof –

have (*Fract* 1 2) * ((*Fract* 1 1) / (*Rep-prat* *p*)) = (*Fract* 1 1) / (*Rep-prat* *p* + *Rep-prat* *p*)

by (metis (no-types, lifting) One-rat-def comm-monoid-mult-class.mult-1 divide-rat mult-2 mult-rat rat-number-expand(3) times-divide-times-eq)

then show ?thesis

by (metis Rep-prat-inject half.rep-eq mult-2 mult-numeral-1-right numeral-One padd.rep-eq pdiv.rep-eq pinv-def pmult.rep-eq pwrite.rep-eq times-divide-times-eq)

qed

lemma *ppos-mult*:

assumes *ppos* *a*

and *ppos* *b*

shows *ppos* (*pmult* *a* *b*)

using *assms*(1) *assms*(2) *pmult.rep-eq* *ppos.rep-eq* by auto

lemma *padd-zero*:

pnone = *padd* *a* *b* \longleftrightarrow *a* = *pnone* \wedge *b* = *pnone*

by (metis Rep-prat Rep-prat-inject add.right-neutral leD le-add-same-cancel2 less-eq-rat-def mem-Collect-eq padd.rep-eq pnone.rep-eq)

lemma *ppos-add*:

assumes *ppos* *a*

shows *ppos* (*padd* *a* *b*)

by (metis Rep-prat Rep-prat-inject *assms* dual-order.strict-iff-order mem-Collect-eq padd-zero pnone.rep-eq *ppos.rep-eq*)

lemma *pinv-inverts*:

assumes *pgte* *a* *b*

and *ppos* *b*

shows *pgte* (*pinv* *b*) (*pinv* *a*)

proof –

have *Rep-prat* *a* \geq *Rep-prat* *b*

using *assms*(1) *pgte.rep-eq* by auto

then have (*Fract* 1 1) / *Rep-prat* *b* \geq (*Fract* 1 1) / *Rep-prat* *a*

by (metis One-rat-def *assms*(2) frac-le le-numeral-extra(4) *ppos.rep-eq* zero-le-one)

then show ?thesis

by (*simp add: One-rat-def pdiv.rep-eq pgte.rep-eq pinv-def pwrite.rep-eq*)
qed

lemma *pinv-pmult-ok:*

assumes *ppos p*

shows $pmult\ p\ (pinv\ p) = pwrite$

proof –

obtain *r* where $r = Rep-prat\ p$ by *simp*

then have $r * ((Fract\ 1\ 1) / r) = Fract\ 1\ 1$

using *assms ppos.rep-eq* by *force*

then show *?thesis*

by (*metis One-rat-def Rep-prat-inject ‹r = Rep-prat p› pdiv.rep-eq pinv-def pmult.rep-eq pwrite.rep-eq*)

qed

lemma *pinv-pwrite:*

$pinv\ pwrite = pwrite$

by (*metis Rep-prat-inverse div-by-1 pdiv.rep-eq pinv-def pwrite.rep-eq*)

lemma *pmult-ppos:*

assumes *ppos a*

and *ppos b*

shows $ppos\ (pmult\ a\ b)$

using *assms(1) assms(2) pmult.rep-eq ppos.rep-eq* by *auto*

lemma *ppos-inv:*

assumes *ppos p*

shows $ppos\ (pinv\ p)$

by (*metis Rep-prat Rep-prat-inverse assms less-eq-rat-def mem-Collect-eq not-one-le-zero pinv-pmult-ok pmult-comm pmult-special(2) pnone.rep-eq ppos.rep-eq pwrite.rep-eq*)

lemma *pmin-pmax:*

assumes $pgte\ x\ (pmin\ a\ b)$

shows $x = pmin\ (pmax\ x\ a)\ (pmax\ x\ b)$

proof (*cases pgte x a*)

case *True*

then show *?thesis*

by (*metis pmax-is pmax-smaller pmin-comm pmin-is*)

next

case *False*

then show *?thesis*

by (*metis assms not-pgte-charact pgt-implies-pgte pmax-is pmax-smaller pmin-comm pmin-is*)

qed

definition *comp-one* **where**

comp-one $p = (\text{SOME } r. \text{padd } p \ r = \text{pwrite})$

lemma *padd-comp-one*:

assumes *pgte* *pwrite* *x*

shows $\text{padd } x \ (\text{comp-one } x) = \text{pwrite}$

by (*metis* (*mono-tags*, *lifting*) *assms comp-one-def p-greater-exists someI-ex*)

lemma *ppos-eq-pnone*:

$\text{ppos } p \longleftrightarrow p \neq \text{pnone}$

by (*metis* *Rep-prat Rep-prat-inject dual-order.strict-iff-order mem-Collect-eq pnone.rep-eq ppos.rep-eq*)

lemma *pmult-order*:

assumes *pgte* *a* *b*

shows $\text{pgte } (\text{pmult } p \ a) \ (\text{pmult } b \ p)$

using *assms p-greater-exists pmult-comm pmult-distr* **by** *force*

lemma *multiply-smaller-pwrite*:

assumes *pgte* *pwrite* *a*

and *pgte* *pwrite* *b*

shows $\text{pgte } \text{pwrite } (\text{pmult } a \ b)$

by (*metis* *assms(1) assms(2) p-greater-exists padd-asso pmult-order pmult-special(1)*)

lemma *pmult-pdiv-cancel*:

assumes *ppos* *a*

shows $\text{pmult } a \ (\text{pdiv } x \ a) = x$

by (*metis* *Rep-prat-inject assms divide-cancel-right dual-order.strict-iff-order nonzero-mult-div-cancel-left pdiv.rep-eq pmult.rep-eq ppos.rep-eq*)

lemma *pmult-padd*:

$\text{pmult } a \ (\text{padd } (\text{pmult } b \ x) \ (\text{pmult } c \ y)) = \text{padd } (\text{pmult } (\text{pmult } a \ b) \ x) \ (\text{pmult } (\text{pmult } a \ c) \ y)$

by (*metis* *Rep-prat-inject mult.assoc pmult.rep-eq pmult-distr*)

lemma *pdiv-smaller*:

assumes *pgte* *a* *b*

and *ppos* *a*

shows $\text{pgte } \text{pwrite } (\text{pdiv } b \ a)$

proof –

let *?a* = *Rep-prat* *a*

let *?b* = *Rep-prat* *b*

have *?b* / *?a* ≤ 1

by (*meson* *assms(1) assms(2) divide-le-eq-1-pos pgte.rep-eq ppos.rep-eq*)


```

then show ?thesis
  by (simp add: pdiv.rep-eq pgte.rep-eq pwrite.rep-eq)
qed

```

```

lemma sum-coeff:
  assumes ppos a
    and ppos b
  shows padd (pdiv a (padd a b)) (pdiv b (padd a b)) = pwrite
proof -
  let ?a = Rep-prat a
  let ?b = Rep-prat b
  have  $(?a / (?a + ?b)) + (?b / (?a + ?b)) = 1$ 
  by (metis add-divide-distrib add-pos-pos assms(1) assms(2) less-irrefl ppos.rep-eq
right-inverse-eq)
  then show ?thesis
  by (metis Rep-prat-inject padd.rep-eq pdiv.rep-eq pwrite.rep-eq)
qed

```

```

lemma padd-one-ineq-sum:
  assumes padd a b = pwrite
    and pgte x aa
    and pgte x bb
  shows pgte x (padd (pmult a aa) (pmult b bb))
  by (metis (mono-tags, lifting) Rep-prat assms(1) assms(2) assms(3) convex-bound-le
mem-Collect-eq padd.rep-eq pgte.rep-eq pmult.rep-eq pwrite.rep-eq)

```

end

1.2 Permission masks: Maps from heap locations to permission amounts

```

theory Mask
  imports PosRat
begin

```

1.2.1 Definitions

```

type-synonym field = string
type-synonym address = nat
type-synonym heap-loc = address × field

```

```

type-synonym mask = heap-loc ⇒ prat
type-synonym bmask = heap-loc ⇒ bool

```

```

definition null where null = 0

```

```

definition full-mask :: mask where
  full-mask hl = (if fst hl = null then pnone else pwrite)

```

definition *multiply-mask* :: *prat* \Rightarrow *mask* \Rightarrow *mask* **where**
multiply-mask *p* π *hl* = *pmult* *p* (π *hl*)

fun *empty-mask* **where**
empty-mask *hl* = *pnone*

fun *empty-bmask* **where**
empty-bmask *hl* = *False*

fun *add-acc* **where** *add-acc* π *hl* *p* = π (*hl* := *padd* (π *hl*) *p*)

inductive *rm-acc* **where**
 π *hl* = *padd* *p* *r* \Longrightarrow *rm-acc* π *hl* *p* (π (*hl* := *r*))

fun *add-masks* **where**
add-masks π' π *hl* = *padd* (π' *hl*) (π *hl*)

definition *greater-mask* **where**
greater-mask π' π \longleftrightarrow (\exists *r*. π' = *add-masks* π *r*)

fun *uni-mask* **where**
uni-mask *hl* *p* = *empty-mask*(*hl* := *p*)

fun *valid-mask* :: *mask* \Rightarrow *bool* **where**
valid-mask π \longleftrightarrow (\forall *hl*. *pgte* *pwrite* (π *hl*)) \wedge (\forall *f*. π (*null*, *f*) = *pnone*)

definition *valid-null* :: *mask* \Rightarrow *bool* **where**
valid-null π \longleftrightarrow (\forall *f*. π (*null*, *f*) = *pnone*)

definition *equal-on-mask* **where**
equal-on-mask π *h* *h'* \longleftrightarrow (\forall *hl*. *ppos* (π *hl*) \longrightarrow *h* *hl* = *h'* *hl*)

definition *equal-on-bmask* **where**
equal-on-bmask π *h* *h'* \longleftrightarrow (\forall *hl*. π *hl* \longrightarrow *h* *hl* = *h'* *hl*)

definition *big-add-masks* **where**
big-add-masks Π Π' *h* = *add-masks* (Π *h*) (Π' *h*)

definition *big-greater-mask* **where**
big-greater-mask Π Π' \longleftrightarrow (\forall *h*. *greater-mask* (Π *h*) (Π' *h*))

definition *greater-bmask* **where**
greater-bmask *H* *H'* \longleftrightarrow (\forall *h*. *H'* *h* \longrightarrow *H* *h*)

definition *update-dm* **where**
update-dm *dm* π π' *hl* \longleftrightarrow (*dm* *hl* \vee *pgt* (π *hl*) (π' *hl*))

```

fun pre-get-m where pre-get-m  $\varphi = \text{fst } \varphi$ 
fun pre-get-h where pre-get-h  $\varphi = \text{snd } \varphi$ 
fun srm-acc where srm-acc  $\varphi \text{ hl } p = (\text{rm-acc } (\text{pre-get-m } \varphi) \text{ hl } p, \text{pre-get-h } \varphi)$ 

```

```

datatype val = Bool (the-bool: bool) | Address (the-address: address) | Rat (the-rat:
prat)

```

```

definition upper-bounded :: mask  $\Rightarrow$  prat  $\Rightarrow$  bool where
upper-bounded  $\pi p \longleftrightarrow (\forall \text{hl. } \text{pgte } p (\pi \text{hl}))$ 

```

1.2.2 Lemmas

```

lemma ssubsetI:
assumes  $\bigwedge \pi h. (\pi, h) \in A \implies (\pi, h) \in B$ 
shows  $A \subseteq B$ 
using assms by auto

```

```

lemma double-inclusion:
assumes  $A \subseteq B$ 
and  $B \subseteq A$ 
shows  $A = B$ 
using assms by blast

```

```

lemma add-masks-comm:
add-masks a b = add-masks b a
proof (rule ext)
fix x show add-masks a b x = add-masks b a x
by (metis Rep-prat-inverse add.commute add-masks.simps padd.rep-eq)
qed

```

```

lemma add-masks-asso:
add-masks (add-masks a b) c = add-masks a (add-masks b c)
proof (rule ext)
fix x show add-masks (add-masks a b) c x = add-masks a (add-masks b c) x
by (metis Rep-prat-inverse add.assoc add-masks.simps padd.rep-eq)
qed

```

```

lemma minus-empty:
 $\pi = \text{add-masks } \pi \text{ empty-mask}$ 
proof (rule ext)
fix x show  $\pi x = \text{add-masks } \pi \text{ empty-mask } x$ 
by (metis Rep-prat-inverse add.right-neutral add-masks.simps empty-mask.simps
padd.rep-eq pnone.rep-eq)
qed

```

```

lemma add-acc-uni-mask:
add-acc  $\pi \text{ hl } p = \text{add-masks } \pi (\text{uni-mask } \text{hl } p)$ 
proof (rule ext)

```

```

fix x show add-acc  $\pi$  hl p x = add-masks  $\pi$  (uni-mask hl p) x
by (metis (no-types, opaque-lifting) add-acc.simps add-masks.simps fun-upd-apply
minus-empty uni-mask.simps)
qed

```

```

lemma add-masks-equiv-valid-null:
valid-null (add-masks a b)  $\longleftrightarrow$  valid-null a  $\wedge$  valid-null b
by (metis (mono-tags, lifting) add-masks.simps padd-zero valid-null-def)

```

```

lemma valid-maskI:
assumes  $\bigwedge$ hl. pgte pwrite ( $\pi$  hl)
and  $\bigwedge$ f.  $\pi$  (null, f) = pnone
shows valid-mask  $\pi$ 
by (simp add: assms(1) assms(2))

```

```

lemma greater-mask-equiv-def:
greater-mask  $\pi'$   $\pi$   $\longleftrightarrow$  ( $\forall$  hl. pgte ( $\pi'$  hl) ( $\pi$  hl))
(is ?A  $\longleftrightarrow$  ?B)
proof (rule iffI)
show ?A  $\implies$  ?B
proof (clarify)
fix hl assume greater-mask  $\pi'$   $\pi$ 
then obtain r where  $\pi' =$  add-masks  $\pi$  r
using greater-mask-def by blast
then show pgte ( $\pi'$  hl) ( $\pi$  hl)
using Rep-prat padd.rep-eq pgte.rep-eq by auto
qed
show ?B  $\implies$  ?A
proof -
assume ?B
let ?r =  $\lambda$ hl. (SOME p.  $\pi'$  hl = padd ( $\pi$  hl) p)
have  $\pi' =$  add-masks  $\pi$  ?r
proof (rule ext)
fix hl
have  $\pi'$  hl = padd ( $\pi$  hl) (?r hl)
by (meson  $\langle \forall$  hl. pgte ( $\pi'$  hl) ( $\pi$  hl)  $\rangle$  p-greater-exists someI-ex)
then show  $\pi'$  hl = add-masks  $\pi$  ?r hl
by auto
qed
then show ?A
using greater-mask-def by blast
qed
qed

```

```

lemma greater-maskI:
assumes  $\bigwedge$ hl. pgte ( $\pi'$  hl) ( $\pi$  hl)
shows greater-mask  $\pi'$   $\pi$ 
by (simp add: assms greater-mask-equiv-def)

```

lemma *greater-mask-properties*:

```
greater-mask  $\pi$   $\pi$ 
greater-mask  $a$   $b \wedge$  greater-mask  $b$   $c \implies$  greater-mask  $a$   $c$ 
greater-mask  $\pi' \pi \wedge$  greater-mask  $\pi \pi' \implies \pi = \pi'$ 
apply (simp add: greater-maskI pgte.rep-eq)
apply (metis add-masks-asso greater-mask-def)
proof (rule ext)
  fix  $x$  assume greater-mask  $\pi' \pi \wedge$  greater-mask  $\pi \pi'$ 
  show  $\pi x = \pi' x$ 
    by (meson <greater-mask  $\pi' \pi \wedge$  greater-mask  $\pi \pi'$ > greater-mask-equiv-def
    pgte-antisym)
qed
```

lemma *greater-mask-decomp*:

```
assumes greater-mask  $a$  (add-masks  $b$   $c$ )
shows  $\exists a1 a2. a =$  add-masks  $a1 a2 \wedge$  greater-mask  $a1$   $b \wedge$  greater-mask  $a2$   $c$ 
by (metis add-masks-asso assms greater-mask-def greater-mask-properties(1))
```

lemma *valid-empty*:

```
valid-mask empty-mask
by (metis empty-mask.simps le-add-same-cancel1 p-greater-exists padd.rep-eq pgte.rep-eq
pnone.rep-eq valid-mask.simps)
```

lemma *upper-valid-aux*:

```
assumes valid-mask  $a$ 
  and  $a =$  add-masks  $b$   $c$ 
shows valid-mask  $b$ 
proof (rule valid-maskI)
  show  $\bigwedge hl. pgte$  pwrite  $(b$   $hl)$ 
    using assms(1) assms(2) p-greater-exists padd-asso by fastforce
  fix  $f$  show  $b$  (null,  $f$ ) = pnone
    by (metis add-masks-comm assms(1) assms(2) empty-mask.simps greater-mask-def
    greater-mask-equiv-def minus-empty pgte-antisym valid-mask.simps)
qed
```

lemma *upper-valid*:

```
assumes valid-mask  $a$ 
  and  $a =$  add-masks  $b$   $c$ 
shows valid-mask  $b \wedge$  valid-mask  $c$ 
using add-masks-comm assms(1) assms(2) upper-valid-aux by blast
```

lemma *equal-on-bmaskI*:

```
assumes  $\bigwedge hl. \pi$   $hl \implies h$   $hl = h'$   $hl$ 
shows equal-on-bmask  $\pi$   $h$   $h'$ 
using assms equal-on-bmask-def by blast
```

lemma *big-add-greater*:

```
big-greater-mask (big-add-masks  $A$   $B$ )  $B$ 
by (metis add-masks-comm big-add-masks-def big-greater-mask-def greater-mask-def)
```

lemma *big-greater-iff*:
 $big-greater-mask\ A\ B \implies (\exists C. A = big-add-masks\ B\ C)$
proof –
assume *big-greater-mask* $A\ B$
let $?C = \lambda h. SOME\ r. A\ h = add-masks\ (B\ h)\ r$
have $A = big-add-masks\ B\ ?C$
proof (*rule ext*)
fix x
have $A\ x = add-masks\ (B\ x)\ (?C\ x)$
proof (*rule ext*)
fix xa
have $A\ x = add-masks\ (B\ x)\ (SOME\ r. A\ x = add-masks\ (B\ x)\ r)$
by (*metis (mono-tags, lifting) <big-greater-mask A B> big-greater-mask-def greater-mask-def someI-ex*)
then show $A\ x\ xa = add-masks\ (B\ x)\ (SOME\ r. A\ x = add-masks\ (B\ x)\ r)$
 xa
by *auto*
qed
then show $A\ x = big-add-masks\ B\ (\lambda h. SOME\ r. A\ h = add-masks\ (B\ h)\ r)\ x$
by (*metis (no-types, lifting) big-add-masks-def*)
qed
then show $\exists C. A = big-add-masks\ B\ C$
by *fast*
qed

lemma *big-add-masks-asso*:
 $big-add-masks\ A\ (big-add-masks\ B\ C) = big-add-masks\ (big-add-masks\ A\ B)\ C$
proof (*rule ext*)
fix x **show** $big-add-masks\ A\ (big-add-masks\ B\ C)\ x = big-add-masks\ (big-add-masks\ A\ B)\ C\ x$
by (*simp add: add-masks-asso big-add-masks-def*)
qed

lemma *big-add-mask-neutral*:
 $big-add-masks\ \Pi\ (\lambda-. empty-mask) = \Pi$
proof (*rule ext*)
fix x **show** $big-add-masks\ \Pi\ (\lambda-. empty-mask)\ x = \Pi\ x$
by (*metis big-add-masks-def minus-empty*)
qed

lemma *sym-equal-on-mask*:
 $equal-on-mask\ \pi\ a\ b \iff equal-on-mask\ \pi\ b\ a$
proof –
have $\bigwedge a\ b. equal-on-mask\ \pi\ a\ b \implies equal-on-mask\ \pi\ b\ a$
by (*simp add: equal-on-mask-def*)
then show *?thesis* **by** *blast*
qed

lemma *greater-mask-uni-equiv*:
 $greater_mask\ \pi\ (uni_mask\ hl\ r) \longleftrightarrow pgt\ (\pi\ hl)\ r$
by (*metis add-masks-comm fun-upd-apply greater-mask-def greater-mask-equiv-def minus-empty uni-mask.simps*)

lemma *greater-mask-uniI*:
assumes $pgte\ (\pi\ hl)\ r$
shows $greater_mask\ \pi\ (uni_mask\ hl\ r)$
using *greater-mask-uni-equiv assms by metis*

lemma *greater-bmask-refl*:
 $greater_bmask\ H\ H$
by (*simp add: greater-bmask-def*)

lemma *greater-bmask-trans*:
assumes $greater_bmask\ A\ B$
and $greater_bmask\ B\ C$
shows $greater_bmask\ A\ C$
by (*metis assms(1) assms(2) greater-bmask-def*)

lemma *update-dm-same*:
 $update_dm\ dm\ \pi\ \pi = dm$
proof (*rule ext*)
fix x **show** $update_dm\ dm\ \pi\ \pi\ x = dm\ x$
by (*simp add: pgt.rep-eq update-dm-def*)
qed

lemma *update-trans*:
assumes $greater_mask\ \pi\ \pi'$
and $greater_mask\ \pi'\ \pi''$
shows $update_dm\ (update_dm\ dm\ \pi\ \pi')\ \pi'\ \pi'' = update_dm\ dm\ \pi\ \pi''$
proof (*rule ext*)
fix hl **show** $update_dm\ (update_dm\ dm\ \pi\ \pi')\ \pi'\ \pi''\ hl = update_dm\ dm\ \pi\ \pi''\ hl$
proof –
have $update_dm\ (update_dm\ dm\ \pi\ \pi')\ \pi'\ \pi''\ hl \longleftrightarrow (update_dm\ dm\ \pi\ \pi')\ hl \vee pgt\ (\pi'\ hl)\ (\pi'' hl)$
using *update-dm-def by metis*
also have $\dots \longleftrightarrow dm\ hl \vee pgt\ (\pi\ hl)\ (\pi'\ hl) \vee pgt\ (\pi'\ hl)\ (\pi'' hl)$
using *update-dm-def by metis*
moreover have $update_dm\ dm\ \pi\ \pi''\ hl \longleftrightarrow dm\ hl \vee pgt\ (\pi\ hl)\ (\pi'' hl)$
using *update-dm-def by metis*
moreover have $pgt\ (\pi\ hl)\ (\pi'\ hl) \vee pgt\ (\pi'\ hl)\ (\pi'' hl) \longleftrightarrow pgt\ (\pi\ hl)\ (\pi'' hl)$
proof
show $pgt\ (\pi\ hl)\ (\pi'\ hl) \vee pgt\ (\pi'\ hl)\ (\pi'' hl) \implies pgt\ (\pi\ hl)\ (\pi'' hl)$
by (*metis assms(1) assms(2) greater-mask-equiv-def greater-mask-properties(2) not-pgte-charact pgte-antisym*)
show $pgt\ (\pi\ hl)\ (\pi'' hl) \implies pgt\ (\pi\ hl)\ (\pi'\ hl) \vee pgt\ (\pi'\ hl)\ (\pi'' hl)$
by (*metis assms(1) greater-mask-equiv-def not-pgte-charact pgte-antisym*)
qed

ultimately show *?thesis* by *blast*
 qed
 qed

lemma *equal-on-bmask-greater*:
 assumes *equal-on-bmask* $\pi' h h'$
 and *greater-bmask* $\pi' \pi$
 shows *equal-on-bmask* $\pi h h'$
 by (*metis* (*mono-tags*, *lifting*) *assms*(1) *assms*(2) *equal-on-bmask-def* *greater-bmask-def*)

lemma *update-dm-equal-bmask*:
 assumes $\pi = \text{add-masks } \pi' m$
 shows *equal-on-bmask* (*update-dm* $dm \pi \pi'$) $h' h \iff \text{equal-on-mask } m h h' \wedge \text{equal-on-bmask } dm h h'$
proof –
 have *equal-on-bmask* (*update-dm* $dm \pi \pi'$) $h' h \iff (\forall hl. \text{update-dm } dm \pi \pi' hl \longrightarrow h' hl = h hl)$
 by (*simp add: equal-on-bmask-def*)
 moreover have $\bigwedge hl. \text{update-dm } dm \pi \pi' hl \iff dm hl \vee \text{pgt } (\pi hl) (\pi' hl)$
 by (*simp add: update-dm-def*)
 moreover have $(\forall hl. \text{update-dm } dm \pi \pi' hl \longrightarrow h' hl = h hl) \iff \text{equal-on-mask } m h h' \wedge \text{equal-on-bmask } dm h h'$
proof
 show $\forall hl. \text{update-dm } dm \pi \pi' hl \longrightarrow h' hl = h hl \implies \text{equal-on-mask } m h h' \wedge \text{equal-on-bmask } dm h h'$
 by (*simp add: assms equal-on-bmask-def equal-on-mask-def padd.rep-eq pgt.rep-eq ppos.rep-eq update-dm-def*)
 assume *equal-on-mask* $m h h' \wedge \text{equal-on-bmask } dm h h'$
 then have $\bigwedge hl. \text{update-dm } dm \pi \pi' hl \implies h' hl = h hl$
 by (*metis* (*full-types*) *add.right-neutral* *add-masks.simps* *assms dual-order.strict-iff-order equal-on-bmask-def equal-on-mask-def padd.rep-eq pgt.rep-eq pnone.rep-eq ppos-eq-pnone update-dm-def*)
 then show $\forall hl. \text{update-dm } dm \pi \pi' hl \longrightarrow h' hl = h hl$
 by *simp*
 qed
 then show *?thesis*
 by (*simp add: calculation*)
 qed

lemma *const-sum-mask-greater*:
 assumes *add-masks* $a b = \text{add-masks } c d$
 and *greater-mask* $a c$
 shows *greater-mask* $d b$
proof (*rule ccontr*)
 assume $\neg \text{greater-mask } d b$
 then obtain *hl* where $\neg \text{pgte } (d hl) (b hl)$
 using *greater-mask-equiv-def* by *blast*
 then have *pgt* $(b hl) (d hl)$
 using *not-pgte-charact* by *auto*

then have $pgt (padd (a hl) (b hl)) (padd (c hl) (d hl))$
by (*metis* *assms*(2) *greater-mask-equiv-def padd-comm pgte-pgt*)
then show *False*
by (*metis* *add-masks.simps assms*(1) *not-pgte-charact order-refl pgte.rep-eq*)
qed

lemma *add-masks-cancellative*:
assumes *add-masks* $b c = add-masks b d$
shows $c = d$
proof (*rule ext*)
fix x **show** $c x = d x$
by (*metis* *assms*(1) *const-sum-mask-greater greater-mask-properties*(1) *greater-mask-properties*(3))
qed

lemma *equal-on-maskI*:
assumes $\bigwedge hl. ppos (\pi hl) \implies h hl = h' hl$
shows *equal-on-mask* $\pi h h'$
by (*simp* *add: assms equal-on-mask-def*)

lemma *greater-equal-on-mask*:
assumes *equal-on-mask* $\pi' h h'$
and *greater-mask* $\pi' \pi$
shows *equal-on-mask* $\pi h h'$
proof (*rule equal-on-maskI*)
fix hl **assume** *asm*: $ppos (\pi hl)$
then show $h hl = h' hl$
by (*metis* *assms*(1) *assms*(2) *equal-on-mask-def greater-mask-equiv-def less-le-trans pgte.rep-eq ppos.rep-eq*)
qed

lemma *equal-on-mask-sum*:
 $equal-on-mask \pi 1 h h' \wedge equal-on-mask \pi 2 h h' \longleftrightarrow equal-on-mask (add-masks \pi 1 \pi 2) h h'$
proof
show $equal-on-mask (add-masks \pi 1 \pi 2) h h' \implies equal-on-mask \pi 1 h h' \wedge equal-on-mask \pi 2 h h'$
using *add-masks-comm greater-equal-on-mask greater-mask-def* **by** *blast*
assume $equal-on-mask \pi 1 h h' \wedge equal-on-mask \pi 2 h h'$
show $equal-on-mask (add-masks \pi 1 \pi 2) h h'$
proof (*rule equal-on-maskI*)
fix hl **assume** $ppos (add-masks \pi 1 \pi 2 hl)$
then show $h hl = h' hl$
proof (*cases* $ppos (\pi 1 hl)$)
case *True*
then show *?thesis*
by (*meson* $\langle equal-on-mask \pi 1 h h' \wedge equal-on-mask \pi 2 h h' \rangle equal-on-mask-def$)
next
case *False*
then show *?thesis*

by (*metis* $\langle \text{equal-on-mask } \pi 1 h h' \wedge \text{equal-on-mask } \pi 2 h h' \rangle \langle \text{ppos } (\text{add-masks } \pi 1 \pi 2 hl) \rangle \text{add-masks.simps equal-on-mask-def padd-zero ppos-eq-pnone}$)
qed
qed
qed

lemma *valid-larger-mask*:

valid-mask $\pi \longleftrightarrow \text{greater-mask full-mask } \pi$
by (*metis* *fst-eqD full-mask-def greater-maskI greater-mask-def not-one-le-zero not-pgte-charact pgt-implies-pgte pgte.rep-eq pnone.rep-eq pwrite.rep-eq surjective-pairing upper-valid-aux valid-mask.elims(1)*)

lemma *valid-mask-full-mask*:

valid-mask full-mask
using *greater-mask-properties(1) valid-larger-mask* **by** *blast*

lemma *mult-greater*:

assumes *greater-mask a b*
shows *greater-mask (multiply-mask p a) (multiply-mask p b)*
by (*metis (full-types) assms greater-mask-equiv-def multiply-mask-def p-greater-exists pmult-distr*)

lemma *mult-write-mask*:

multiply-mask pwrite $\pi = \pi$
proof (*rule ext*)
fix *x* **show** *multiply-mask pwrite* $\pi x = \pi x$
by (*simp add: multiply-mask-def pmult-special(1)*)
qed

lemma *mult-distr-masks*:

multiply-mask a (add-masks b c) = add-masks (multiply-mask a b) (multiply-mask a c)
proof (*rule ext*)
fix *x* **show** *multiply-mask a (add-masks b c) x = add-masks (multiply-mask a b) (multiply-mask a c) x*
by (*simp add: multiply-mask-def pmult-distr*)
qed

lemma *mult-add-states*:

multiply-mask (padd a b) $\pi = \text{add-masks } (\text{multiply-mask } a \pi) (\text{multiply-mask } b \pi)$
proof (*rule ext*)
fix *x* **show** *multiply-mask (padd a b) $\pi x = \text{add-masks } (\text{multiply-mask } a \pi) (\text{multiply-mask } b \pi) x$*
by (*simp add: multiply-mask-def pmult-comm pmult-distr*)
qed

lemma *upper-boundedI*:

assumes $\bigwedge hl. \text{pgte } p (\pi hl)$

shows *upper-bounded* π p
by (*simp add: assms upper-bounded-def*)

lemma *upper-bounded-smaller*:
assumes *upper-bounded* π a
shows *upper-bounded* (*multiply-mask* p π) (*pmult* p a)
by (*metis assms multiply-mask-def p-greater-exists pmult-distr upper-bounded-def*)

lemma *upper-bounded-bigger*:
assumes *upper-bounded* π a
and *pgte* b a
shows *upper-bounded* π b
by (*meson assms(1) assms(2) order-trans pgte.rep-eq upper-bounded-def*)

lemma *valid-mult*:
assumes *valid-mask* π
and *pgte* *pwrite* p
shows *valid-mask* (*multiply-mask* p π)
proof (*rule valid-maskI*)
have *upper-bounded* π *pwrite*
using *assms(1) upper-bounded-def* **by** *auto*
then have *upper-bounded* (*multiply-mask* p π) (*pmult* p *pwrite*)
by (*simp add: upper-bounded-smaller*)
then show $\bigwedge hl.$ *pgte* *pwrite* (*multiply-mask* p π hl)
by (*metis assms(2) pmult-comm pmult-special(1) upper-bounded-bigger upper-bounded-def*)
show $\bigwedge f.$ *multiply-mask* p π (*null*, f) = *pnone*
by (*metis Rep-prat-inverse add-0-left assms(1) multiply-mask-def padd.rep-eq padd-cancellative pmult-distr pnone.rep-eq valid-mask.elims(1)*)
qed

lemma *valid-sum*:
assumes *valid-mask* a
and *valid-mask* b
and *upper-bounded* a ma
and *upper-bounded* b mb
and *pgte* *pwrite* (*padd* ma mb)
shows *valid-mask* (*add-masks* a b)
and *upper-bounded* (*add-masks* a b) (*padd* ma mb)
proof (*rule valid-maskI*)
show $\bigwedge hl.$ *pgte* *pwrite* (*add-masks* a b hl)
proof –
fix hl
have *pgte* (*padd* ma mb) (*add-masks* a b hl)
by (*metis (mono-tags, lifting) add-masks.simps add-mono-thms-linordered-semiring(1) assms(3) assms(4) padd.rep-eq pgte.rep-eq upper-bounded-def*)
then show *pgte* *pwrite* (*add-masks* a b hl)
by (*meson assms(5) dual-order.trans pgte.rep-eq*)

```

qed
show  $\bigwedge f. \text{add-masks } a \ b \ (\text{null}, f) = \text{pnone}$ 
  by (metis Rep-prat-inverse add-0-left add-masks.simps assms(1) assms(2)
padd.rep-eq pnone.rep-eq valid-mask.simps)
show upper-bounded (add-masks a b) (padd ma mb)
  using add-mono-thms-linordered-semiring(1) assms(3) assms(4) padd.rep-eq
pgte.rep-eq upper-bounded-def by fastforce
qed

```

```

lemma valid-multiply:
  assumes valid-mask a
    and upper-bounded a ma
    and pgte pwrite (pmult ma p)
  shows valid-mask (multiply-mask p a)
  by (metis (no-types, opaque-lifting) assms(1) assms(2) assms(3) multiply-mask-def
pmult-comm pmult-special(2) upper-bounded-bigger upper-bounded-def upper-bounded-smaller
valid-mask.elims(1))

```

```

lemma greater-mult:
  assumes greater-mask a b
  shows greater-mask (multiply-mask p a) (multiply-mask p b)
  by (metis Rep-prat assms greater-mask-equiv-def mem-Collect-eq mult-left-mono
multiply-mask-def pgte.rep-eq pmult.rep-eq)

```

end

1.3 Partial heaps: Partial maps from heap location to values

```

theory PartialHeapSA
  imports Mask Package-logic.PackageLogic
begin

```

1.3.1 Definitions

```

type-synonym heap = heap-loc  $\rightarrow$  val
type-synonym pre-state = mask  $\times$  heap

```

```

definition valid-heap :: mask  $\Rightarrow$  heap  $\Rightarrow$  bool where
  valid-heap  $\pi$  h  $\longleftrightarrow$  ( $\forall \text{hl. } \text{ppos } (\pi \ \text{hl}) \longrightarrow \text{h hl} \neq \text{None}$ )

```

```

fun valid-state :: pre-state  $\Rightarrow$  bool where
  valid-state  $(\pi, h)$   $\longleftrightarrow$  valid-mask  $\pi \wedge$  valid-heap  $\pi$  h

```

```

lemma valid-stateI:
  assumes valid-mask  $\pi$ 
    and  $\bigwedge \text{hl. } \text{ppos } (\pi \ \text{hl}) \Longrightarrow \text{h hl} \neq \text{None}$ 
  shows valid-state  $(\pi, h)$ 
  using assms(1) assms(2) valid-heap-def valid-state.simps by blast

```

```

definition empty-heap where empty-heap hl = None

```

```

lemma valid-pre-unit:
  valid-state (empty-mask, empty-heap)
  using ppone.rep-eq ppos.rep-eq valid-empty valid-stateI by fastforce

typedef state = {  $\varphi$  |  $\varphi$ . valid-state  $\varphi$  }
  using valid-pre-unit by blast

fun get-m :: state  $\Rightarrow$  mask where get-m a = fst (Rep-state a)
fun get-h :: state  $\Rightarrow$  heap where get-h a = snd (Rep-state a)

fun compatible-options where
  compatible-options (Some a) (Some b)  $\longleftrightarrow$  a = b
| compatible-options - -  $\longleftrightarrow$  True

definition compatible-heaps :: heap  $\Rightarrow$  heap  $\Rightarrow$  bool where
  compatible-heaps h h'  $\longleftrightarrow$  ( $\forall$  hl. compatible-options (h hl) (h' hl))

definition compatible :: pre-state  $\Rightarrow$  pre-state  $\Rightarrow$  bool where
  compatible  $\varphi$   $\varphi'$   $\longleftrightarrow$  compatible-heaps (snd  $\varphi$ ) (snd  $\varphi'$ )  $\wedge$  valid-mask (add-masks
(fst  $\varphi$ ) (fst  $\varphi'$ ))

fun add-states :: pre-state  $\Rightarrow$  pre-state  $\Rightarrow$  pre-state where
  add-states ( $\pi$ , h) ( $\pi'$ , h') = (add-masks  $\pi$   $\pi'$ , h ++ h')

definition larger-heap where
  larger-heap h' h  $\longleftrightarrow$  ( $\forall$  hl x. h hl = Some x  $\longrightarrow$  h' hl = Some x)

definition unit :: state where
  unit = Abs-state (empty-mask, empty-heap)

definition plus :: state  $\Rightarrow$  state  $\rightarrow$  state (infixl  $\oplus$  63) where
  a  $\oplus$  b = (if compatible (Rep-state a) (Rep-state b) then Some (Abs-state (add-states
(Rep-state a) (Rep-state b))) else None)

definition core :: state  $\Rightarrow$  state (|-|) where
  core  $\varphi$  = Abs-state (empty-mask, get-h  $\varphi$ )

definition stable :: state  $\Rightarrow$  bool where
  stable  $\varphi$   $\longleftrightarrow$  ( $\forall$  hl. ppos (get-m  $\varphi$  hl)  $\longleftrightarrow$  get-h  $\varphi$  hl  $\neq$  None)

```

1.3.2 Lemmas

```

lemma valid-heapI:
  assumes  $\bigwedge$  hl. ppos ( $\pi$  hl)  $\Longrightarrow$  h hl  $\neq$  None
  shows valid-heap  $\pi$  h
  using assms valid-heap-def by presburger

```

```

lemma valid-state-decompose:

```

```

assumes valid-state (add-masks a b, h)
shows valid-state (a, h)
proof (rule valid-stateI)
  show valid-mask a
    using assms upper-valid-aux valid-state.simps by blast
  fix hl assume ppos (a hl) then show h hl  $\neq$  None
    by (metis add-masks.simps assms ppos-add valid-heap-def valid-state.simps)
qed

```

```

lemma compatible-heapsI:
assumes  $\bigwedge hl$  a b. h hl = Some a  $\implies$  h' hl = Some b  $\implies$  a = b
shows compatible-heaps h h'
by (metis assms compatible-heaps-def compatible-options.elims(3))

```

```

lemma compatibleI-old:
assumes  $\bigwedge hl$  x y. snd  $\varphi$  hl = Some x  $\wedge$  snd  $\varphi'$  hl = Some y  $\implies$  x = y
  and valid-mask (add-masks (fst  $\varphi$ ) (fst  $\varphi'$ ))
shows compatible  $\varphi$   $\varphi'$ 
using assms(1) assms(2) compatible-def compatible-heapsI by presburger

```

```

lemma larger-heap-anti:
assumes larger-heap a b
  and larger-heap b a
shows a = b
proof (rule ext)
  fix x show a x = b x
  proof (cases a x)
    case None
      then show ?thesis
      by (metis assms(1) larger-heap-def not-None-eq)
    next
      case (Some a)
      then show ?thesis
      by (metis assms(2) larger-heap-def)
  qed
qed

```

```

lemma larger-heapI:
assumes  $\bigwedge hl$  x. h hl = Some x  $\implies$  h' hl = Some x
shows larger-heap h' h
by (simp add: assms larger-heap-def)

```

```

lemma larger-heap-refl:
larger-heap h h
using larger-heap-def by blast

```

```

lemma compatible-heaps-comm:
assumes compatible-heaps a b
shows a ++ b = b ++ a

```

```

proof (rule ext)
  fix x show (a ++ b) x = (b ++ a) x
  proof (cases a x)
    case None
      then show ?thesis
      by (simp add: domIff map-add-dom-app-simps(2) map-add-dom-app-simps(3))
    next
      case (Some a)
        then show ?thesis
        by (metis (no-types, lifting) assms compatible-heaps-def compatible-options.elims(2)
map-add-None map-add-dom-app-simps(1) map-add-dom-app-simps(3))
  qed
qed

```

```

lemma larger-heaps-sum-ineq:
  assumes larger-heap a' a
    and larger-heap b' b
    and compatible-heaps a' b'
  shows larger-heap (a' ++ b') (a ++ b)
proof (rule larger-heapI)
  fix hl x assume (a ++ b) hl = Some x
  show (a' ++ b') hl = Some x
  proof (cases a hl)
    case None
      then show ?thesis
      by (metis ⟨(a ++ b) hl = Some x⟩ assms(2) larger-heap-def map-add-SomeD
map-add-find-right)
    next
      case (Some aa)
        then show ?thesis
        by (metis (mono-tags, lifting) ⟨(a ++ b) hl = Some x⟩ assms(1) assms(2)
assms(3) compatible-heaps-comm larger-heap-def map-add-Some-iff)
  qed
qed

```

```

lemma larger-heap-trans:
  assumes larger-heap a b
    and larger-heap b c
  shows larger-heap a c
  by (metis (no-types, opaque-lifting) assms(1) assms(2) larger-heap-def)

```

```

lemma larger-heap-comp:
  assumes larger-heap a b
    and compatible-heaps a c
  shows compatible-heaps b c
proof (rule compatible-heapsI)
  fix hl a ba
  assume b hl = Some a c hl = Some ba
  then show a = ba

```

by (*metis* *assms*(1) *assms*(2) *compatible-heaps-def compatible-options.simps*(1) *larger-heap-def*)

qed

lemma *larger-heap-plus*:

assumes *larger-heap* *a b*

and *larger-heap* *a c*

shows *larger-heap* *a (b ++ c)*

proof (*rule larger-heapI*)

fix *hl x* **assume** $(b ++ c) \text{ hl} = \text{Some } x$

then show $a \text{ hl} = \text{Some } x$

proof (*cases* *b hl*)

case *None*

then show *?thesis*

by (*metis* $\langle (b ++ c) \text{ hl} = \text{Some } x \rangle$ *assms*(2) *larger-heap-def map-add-SomeD*)

next

case (*Some bb*)

then show *?thesis*

by (*metis* $\langle (b ++ c) \text{ hl} = \text{Some } x \rangle$ *assms*(1) *assms*(2) *larger-heap-def map-add-SomeD*)

qed

qed

lemma *compatible-heaps-sum*:

assumes *compatible-heaps* *a b*

and *compatible-heaps* *a c*

shows *compatible-heaps* *a (b ++ c)*

by (*metis* (*no-types, opaque-lifting*) *assms*(1) *assms*(2) *compatible-heaps-def map-add-dom-app-simps*(1) *map-add-dom-app-simps*(3))

lemma *larger-compatible-sum-heaps*:

assumes *larger-heap* *a x*

and *larger-heap* *b y*

and *compatible-heaps* *a b*

shows *compatible-heaps* *x y*

proof (*rule compatible-heapsI*)

fix *hl a b* **assume** $x \text{ hl} = \text{Some } a$ $y \text{ hl} = \text{Some } b$

then show $a = b$

by (*metis* *assms*(1) *assms*(2) *assms*(3) *compatible-heaps-def compatible-options.simps*(1) *larger-heap-def*)

qed

lemma *get-h-m*:

Rep-state *x* = (*get-m* *x*, *get-h* *x*) **by** *simp*

lemma *get-pre*:

get-h *x* = *snd* (*Rep-state* *x*)

get-m *x* = *fst* (*Rep-state* *x*)

by *simp-all*

lemma *plus-ab-defined*:

$\varphi \oplus \varphi' \neq \text{None} \longleftrightarrow \text{compatible-heaps } (\text{get-h } \varphi) (\text{get-h } \varphi') \wedge \text{valid-mask } (\text{add-masks } (\text{get-m } \varphi) (\text{get-m } \varphi'))$
(is $?A \longleftrightarrow ?B$ **)**

proof

show $?A \implies ?B$
by (*metis compatible-def get-pre(1) get-pre(2) plus-def*)

show $?B \implies ?A$
using *compatible-def plus-def* **by** *auto*

qed

lemma *plus-charact*:

assumes $a \oplus b = \text{Some } x$
shows $\text{get-m } x = \text{add-masks } (\text{get-m } a) (\text{get-m } b)$
and $\text{get-h } x = (\text{get-h } a) ++ (\text{get-h } b)$

proof –

have $x = (\text{Abs-state } (\text{add-states } (\text{Rep-state } a) (\text{Rep-state } b)))$
by (*metis assms option.discI option.inject plus-def*)

moreover have *compatible* $(\text{Rep-state } a) (\text{Rep-state } b)$
using *assms(1) plus-def* **by** (*metis option.discI*)

moreover have *valid-state* $(\text{add-states } (\text{Rep-state } a) (\text{Rep-state } b))$

proof –

have *valid-state* $(\text{add-masks } (\text{get-m } a) (\text{get-m } b), (\text{get-h } a) ++ (\text{get-h } b))$

proof (*rule valid-stateI*)

show *valid-mask* $(\text{add-masks } (\text{get-m } a) (\text{get-m } b))$
using *calculation(2) compatible-def* **by** *fastforce*

fix *hl* **assume** *ppos* $(\text{add-masks } (\text{get-m } a) (\text{get-m } b) \text{ hl})$

then show $(\text{get-h } a ++ \text{get-h } b) \text{ hl} \neq \text{None}$

proof (*cases ppos (get-m a hl)*)

case *True*

then show *?thesis*
by (*metis Rep-state get-h-m map-add-None mem-Collect-eq valid-heap-def valid-state.simps*)

next

case *False*

then have *ppos* $(\text{get-m } b \text{ hl})$
using $\langle \text{ppos } (\text{add-masks } (\text{get-m } a) (\text{get-m } b) \text{ hl}) \rangle \text{ padd.rep-eq ppos.rep-eq}$

by *auto*

then show *?thesis*
by (*metis Rep-state get-h-m map-add-None mem-Collect-eq valid-heap-def valid-state.simps*)

qed

qed

then show *?thesis*
using *add-states.simps get-h-m* **by** *presburger*

qed

ultimately show $\text{get-m } x = \text{add-masks } (\text{get-m } a) (\text{get-m } b)$

by (metis Abs-state-inverse add-states.simps fst-conv get-h-m mem-Collect-eq)

show $get-h\ x = (get-h\ a) ++ (get-h\ b)$

by (metis Abs-state-inject CollectI Rep-state Rep-state-inverse ‹valid-state (add-states (Rep-state a) (Rep-state b))› ‹x = Abs-state (add-states (Rep-state a) (Rep-state b))› add-states.simps eq-snd-iff get-h.simps)

qed

lemma commutative:

$$a \oplus b = b \oplus a$$

proof (cases compatible-heaps (get-h a) (get-h b) \wedge valid-mask (add-masks (get-m a) (get-m b)))

case True

then have compatible-heaps (get-h b) (get-h a) \wedge add-masks (get-m a) (get-m b) = add-masks (get-m b) (get-m a)

by (metis add-masks-comm compatible-heapsI compatible-heaps-def compatible-options.simps(1))

then have $(get-h\ a) ++ (get-h\ b) = (get-h\ b) ++ (get-h\ a)$

by (simp add: compatible-heaps-comm)

then show ?thesis

by (metis True ‹compatible-heaps (get-h b) (get-h a) \wedge add-masks (get-m a) (get-m b) = add-masks (get-m b) (get-m a)› add-states.simps get-h-m plus-ab-defined plus-def)

next

case False

then show ?thesis

by (metis add-masks-comm compatible-heapsI compatible-heaps-def compatible-options.simps(1) plus-ab-defined)

qed

lemma asso1:

assumes $a \oplus b = \text{Some } ab \wedge b \oplus c = \text{Some } bc$

shows $ab \oplus c = a \oplus bc$

proof (cases $ab \oplus c$)

case None

then show ?thesis

proof (cases compatible-heaps (get-h ab) (get-h c))

case True

then have \neg valid-mask (add-masks (add-masks (get-m a) (get-m b)) (get-m c))

by (metis None assms plus-ab-defined plus-charact(1))

then show ?thesis

by (metis add-masks-asso assms plus-ab-defined plus-charact(1))

next

case False

then have \neg compatible-heaps (get-h a ++ get-h b) (get-h c)

using assms plus-charact(2) by force

then obtain $l\ x\ y$ where $(get-h\ a ++ get-h\ b)\ l = \text{Some } x\ get-h\ c\ l = \text{Some } y\ x \neq y$

```

    using compatible-heapsI by blast
  then have  $\neg$  compatible-heaps (get-h a) (get-h b ++ get-h c)
  proof (cases get-h a l)
    case None
    then show ?thesis
      by (metis  $\langle$ get-h a ++ get-h b $\rangle$  l = Some x  $\langle$ get-h c l = Some y $\rangle$   $\langle$ x  $\neq$  y $\rangle$ 
  assms compatible-heaps-comm map-add-dom-app-simps(1) map-add-dom-app-simps(3)
  map-add-find-right option.inject option.simps(3) plus-ab-defined)
    next
    case (Some aa)
    then show ?thesis
      by (metis  $\langle$ get-h a ++ get-h b $\rangle$  l = Some x  $\langle$ get-h c l = Some y $\rangle$   $\langle$ x  $\neq$  y $\rangle$ 
  assms commutative compatible-heaps-def compatible-options.elims(2) map-add-find-right
  option.inject option.simps(3) plus-charact(2))
  qed
  then show ?thesis
    by (metis None assms plus-ab-defined plus-charact(2))
  qed
next
case (Some x)
then have compatible-heaps (get-h a ++ get-h b) (get-h c)
  by (metis assms option.simps(3) plus-ab-defined plus-charact(2))
then have compatible-heaps (get-h a) (get-h b ++ get-h c)
  by (metis (full-types) assms compatible-heaps-comm compatible-heaps-def com-
  patible-heaps-sum compatible-options.simps(2) domIff map-add-dom-app-simps(1)
  option.distinct(1) plus-ab-defined)
  moreover have valid-mask (add-masks (get-m a) (add-masks (get-m b) (get-m
  c)))
  by (metis Some add-masks-asso assms option.distinct(1) plus-ab-defined plus-charact(1))
  ultimately obtain y where Some y = a  $\oplus$  bc
  by (metis assms plus-ab-defined plus-charact(1) plus-charact(2) plus-def)
  then show ?thesis
    by (metis (mono-tags, lifting) Some add-masks-asso add-states.simps assms
  get-h-m map-add-assoc option.distinct(1) plus-charact(1) plus-charact(2) plus-def)
  qed

```

lemma asso2:

```

  assumes a  $\oplus$  b = Some ab  $\wedge$  b  $\oplus$  c = None
  shows ab  $\oplus$  c = None
  proof (cases valid-mask (add-masks (get-m b) (get-m c)))
    case True
    then have  $\neg$  compatible-heaps (get-h b) (get-h c)
      using assms plus-ab-defined by blast
    then obtain l x y where get-h b l = Some x get-h c l = Some y x  $\neq$  y
      using compatible-heapsI by blast
    then have get-h ab l = Some x
      by (metis assms map-add-find-right plus-charact(2))
    then show ?thesis
      by (metis  $\langle$ get-h c l = Some y $\rangle$   $\langle$ x  $\neq$  y $\rangle$  compatible-heaps-def compatible-options.simps(1)

```

```

plus-ab-defined)
next
  case False
  then obtain l where  $\neg$  (pgte pwrite (add-masks (get-m b) (get-m c) l))
  by (metis Abs-state-cases Rep-state-cases Rep-state-inverse add-masks-equiv-valid-null
get-h-m mem-Collect-eq valid-mask.simps valid-null-def valid-state.simps)
  then have  $\neg$  (pgte pwrite (add-masks (get-m ab) (get-m c) l))
  proof -
    have pgte (add-masks (get-m ab) (get-m c) l) (add-masks (get-m b) (get-m c)
l)
      using assms p-greater-exists padd-asso padd-comm plus-character(1) by auto
    then show ?thesis
      by (meson  $\langle \neg$  pgte pwrite (add-masks (get-m b) (get-m c) l)  $\rangle$  order-trans
pgte.rep-eq)
  qed
  then show ?thesis
    using plus-ab-defined valid-mask.simps by blast
qed

```

```

lemma core-defined:
  get-h | $\varphi$ | = get-h  $\varphi$ 
  get-m | $\varphi$ | = empty-mask
  using Abs-state-inverse core-def pnone.rep-eq ppos.rep-eq valid-empty valid-stateI
apply force
  by (metis Abs-state-inverse CollectI core-def empty-mask.simps fst-conv get-pre(2)
less-irrefl pnone.rep-eq ppos.rep-eq valid-empty valid-stateI)

```

```

lemma state-ext:
  assumes get-h a = get-h b
  and get-m a = get-m b
  shows a = b
  by (metis Rep-state-inverse assms(1) assms(2) get-h-m)

```

```

lemma core-is-smaller:
  Some x = x  $\oplus$  |x|
proof -
  obtain y where Some y = x  $\oplus$  |x|
  by (metis Rep-state compatible-heapsI core-defined(1) core-defined(2) get-h-m
mem-Collect-eq minus-empty option.collapse option.sel plus-ab-defined valid-state.simps)
  moreover have y = x
  proof (rule state-ext)
    have get-h x = get-h x ++ get-h x
    by (simp add: map-add-subsumed1)
    then show get-h y = get-h x
    using calculation core-defined(1) plus-character(2) by presburger
  show get-m y = get-m x
  by (metis calculation core-defined(2) minus-empty plus-character(1))
  qed
ultimately show ?thesis by blast

```

qed

lemma *core-is-pure*:

Some $|x| = |x| \oplus |x|$

proof –

obtain y **where** *Some* $y = |x| \oplus |x|$

by (*metis core-def core-defined(1) core-is-smaller*)

moreover have $y = |x|$

by (*metis calculation core-def core-defined(1) core-is-smaller option.sel*)

ultimately show *?thesis* **by** *blast*

qed

lemma *core-sum*:

assumes *Some* $c = a \oplus b$

shows *Some* $|c| = |a| \oplus |b|$

proof –

obtain x **where** *Some* $x = |a| \oplus |b|$

by (*metis assms core-defined(1) core-defined(2) minus-empty option.exhaust-sel plus-ab-defined valid-empty*)

moreover have $x = |c|$

by (*metis assms calculation core-defined(1) core-defined(2) minus-empty plus-charact(1) plus-charact(2) state-ext*)

ultimately show *?thesis* **by** *blast*

qed

lemma *core-max*:

assumes *Some* $x = x \oplus c$

shows $\exists r. \textit{Some } |x| = c \oplus r$

proof –

obtain y **where** *Some* $y = c \oplus |x|$

by (*metis assms asso2 core-is-smaller plus-def*)

moreover have $|x| = y$

by (*metis (mono-tags, opaque-lifting) Rep-state-inverse add-masks-cancellative assms calculation commutative core-defined(1) core-sum get-h-m minus-empty option.inject plus-charact(1)*)

ultimately show *?thesis* **by** *blast*

qed

lemma *positivity*:

assumes $a \oplus b = \textit{Some } c$

and *Some* $c = c \oplus c$

shows *Some* $a = a \oplus a$

proof –

obtain x **where** *Some* $x = a \oplus a$

by (*metis assms(1) assms(2) asso2 commutative option.exhaust-sel*)

moreover have $x = a$

by (*metis Rep-state-inverse add-masks-cancellative add-masks-comm assms(1) assms(2) calculation core-defined(1) core-defined(2) core-is-smaller get-h-m greater-mask-def greater-mask-properties(3) option.sel plus-charact(1)*)

ultimately show *?thesis* **by** *blast*
qed

lemma *cancellative*:

assumes *Some a = b ⊕ x*
and *Some a = b ⊕ y*
and $|x| = |y|$
shows $x = y$
by (*metis add-masks-cancellative assms(1) assms(2) assms(3) core-defined(1) plus-charact(1) state-ext*)

lemma *unit-charact*:

get-h unit = empty-heap
get-m unit = empty-mask
proof –
have *valid-state (empty-mask, empty-heap)*
using *valid-pre-unit* **by** *auto*
then show *get-h unit = empty-heap* **using** *unit-def*
by (*simp add: ⟨unit = Abs-state (empty-mask, empty-heap)⟩ Abs-state-inverse*)
show *get-m unit = empty-mask*
using *⟨valid-state (empty-mask, empty-heap)⟩ unit-def Abs-state-inverse*
by *fastforce*
qed

lemma *empty-heap-neutral*:

$a ++ \text{empty-heap} = a$
proof (*rule ext*)
fix x **show** $(a ++ \text{empty-heap}) x = a x$
by (*simp add: domIff empty-heap-def map-add-dom-app-simps(3)*)
qed

lemma *unit-neutral*:

Some a = a ⊕ unit
proof –
obtain x **where** *Some x = a ⊕ unit*
by (*metis Abs-state-cases Rep-state-cases Rep-state-inverse compatible-heapsI empty-heap-def fst-conv get-h-m mem-Collect-eq minus-empty option.distinct(1) option.exhaust-sel plus-ab-defined snd-conv unit-def valid-pre-unit valid-state.simps*)
moreover have $x = a$
proof (*rule state-ext*)
show *get-h x = get-h a*
using *calculation empty-heap-neutral plus-charact(2) unit-charact(1)* **by** *auto*
show *get-m x = get-m a*
by (*metis calculation minus-empty plus-charact(1) unit-charact(2)*)
qed
ultimately show *?thesis* **by** *blast*
qed

lemma *stableI*:

```

assumes  $\bigwedge hl. ppos (get-m \varphi hl) \longleftrightarrow get-h \varphi hl \neq None$ 
shows stable  $\varphi$ 
using assms stable-def by blast

```

lemma *stable-unit*:

```

stable unit
by (metis empty-heap-def stable-def unit-charact(1) unit-charact(2) valid-heap-def
valid-pre-unit valid-state.simps)

```

lemma *stable-sum*:

```

assumes stable a
and stable b
and Some x = a  $\oplus$  b
shows stable x
proof (rule stableI)
fix hl
show  $ppos (get-m x hl) = (get-h x hl \neq None) \text{ (is } ?A \longleftrightarrow ?B)$ 
proof
show  $?A \implies ?B$ 
by (metis add-le-same-cancel2 add-masks.simps assms(1) assms(2) assms(3)
leI less-le-trans map-add-None padd.rep-eq plus-charact(1) plus-charact(2) ppos.rep-eq
stable-def)
show  $?B \implies ?A$ 
by (metis add-masks.simps assms(1) assms(2) assms(3) map-add-None
padd-comm plus-charact(1) plus-charact(2) ppos-add stable-def)
qed
qed

```

lemma *multiply-valid*:

```

assumes pgte pwrite p
shows valid-state (multiply-mask p (get-m  $\varphi$ ), get-h  $\varphi$ )
proof (rule valid-stateI)
show valid-mask (multiply-mask p (get-m  $\varphi$ ))
by (metis Rep-state assms(1) get-h-m mem-Collect-eq valid-mult valid-state.simps)
fix hl show  $ppos (multiply-mask p (get-m \varphi) hl) \implies get-h \varphi hl \neq None$ 
by (metis Abs-state-cases Rep-state-cases Rep-state-inverse get-h-m mem-Collect-eq
multiply-mask-def pmult-comm pmult-special(2) ppos-eq-pnone valid-heap-def valid-state.simps)
qed

```

1.4 This state model corresponds to a separation algebra

global-interpretation *PartialSA*: *package-logic plus core unit stable*

```

defines greater (infixl  $\succeq$  50) = PartialSA.greater
and add-set (infixl  $\otimes$  60) = PartialSA.add-set
and defined (infixl  $\#\#$  60) = PartialSA.defined
and greater-set (infixl  $\gg\gg$  50) = PartialSA.greater-set
and minus (infixl  $|\ominus|$  60) = PartialSA.minus
apply standard
apply (simp add: commutative)

```

```

using asso1 apply blast
using asso2 apply blast
using core-is-smaller apply blast
using core-is-pure apply blast
using core-max apply blast
using core-sum apply blast
using positivity apply blast
using cancellative apply blast
using unit-neutral apply blast
using stable-sum apply blast
using stable-unit by blast

```

lemma *greaterI*:

```

  assumes larger-heap (get-h a) (get-h b)
    and greater-mask (get-m a) (get-m b)
  shows  $a \succeq b$ 
proof -
  let ?m =  $\lambda l. \text{SOME } p. \text{get-m } a \ l = \text{padd } (\text{get-m } b \ l) \ p$ 
  have get-m a = add-masks (get-m b) ?m
  proof (rule ext)
    fix l
    have pgte (get-m a l) (get-m b l)
      by (meson assms(2) greater-mask-equiv-def)
    then have get-m a l = padd (get-m b l) (SOME p. get-m a l = padd (get-m b
l) p)
      by (simp add: p-greater-exists verit-sko-ex')
    then show get-m a l = add-masks (get-m b) ( $\lambda l. \text{SOME } p. \text{get-m } a \ l = \text{padd } (\text{get-m } b \ l) \ p$ ) l
      by simp
  qed
  moreover have valid-state (?m, get-h a)
  proof (rule valid-stateI)
    show valid-mask ( $\lambda l. \text{SOME } p. \text{get-m } a \ l = \text{padd } (\text{get-m } b \ l) \ p$ )
      by (metis (no-types, lifting) Rep-state calculation get-h-m mem-Collect-eq
upper-valid valid-state.simps)
    fix hl
    assume asm0: ppos (SOME p. get-m a hl = padd (get-m b hl) p)
    then have ppos (get-m a hl)
      by (metis (no-types, lifting) add-masks.elims add-masks-comm calculation
greater-mask-def ppos-add)
    then show get-h a hl  $\neq$  None
      by (metis Rep-state get-h.simps get-pre(2) mem-Collect-eq prod.collapse
valid-heap-def valid-state.simps)
  qed
  moreover have compatible-heaps (get-h b) (get-h a)
    by (metis (mono-tags, lifting) assms(1) compatible-heapsI larger-heap-def op-
tion.inject)
  ultimately have (get-m a, get-h a) = add-states (get-m b, get-h b) (?m, get-h

```


a)

```

proof –
  have get-h b ++ get-h a = get-h a
  proof (rule ext)
    fix x show (get-h b ++ get-h a) x = get-h a x
      by (metis assms(1) domIff larger-heap-def map-add-dom-app-simps(1))
map-add-dom-app-simps(3) not-Some-eq
  qed
  then show ?thesis
    by (metis ‹get-m a = add-masks (get-m b) (λl. SOME p. get-m a l = padd
(get-m b l) p› add-states.simps)
  qed
  moreover have compatible-heaps (get-h b) (get-h a) ∧ valid-mask (add-masks
(get-m b) ?m)
  by (metis Rep-state ‹compatible-heaps (get-h b) (get-h a)› ‹get-m a = add-masks
(get-m b) (λl. SOME p. get-m a l = padd (get-m b l) p› get-h-m mem-Collect-eq
valid-state.simps)
  ultimately have Some a = b ⊕ Abs-state (?m, get-h a)
  proof –
    have Rep-state (Abs-state (?m, get-h a)) = (?m, get-h a)
      using Abs-state-inverse ‹valid-state (λl. SOME p. get-m a l = padd (get-m b
l) p, get-h a› by blast
    moreover have compatible (Rep-state b) (?m, get-h a)
      using ‹compatible-heaps (get-h b) (get-h a) ∧ valid-mask (add-masks (get-m
b) (λl. SOME p. get-m a l = padd (get-m b l) p)› compatible-def by auto
    moreover have valid-state (add-states (Rep-state b) (?m, get-h a))
      by (metis Rep-state ‹(get-m a, get-h a) = add-states (get-m b, get-h b) (λl.
SOME p. get-m a l = padd (get-m b l) p, get-h a› get-h-m mem-Collect-eq)
    ultimately show ?thesis
      by (metis (no-types, lifting) Rep-state-inverse ‹(get-m a, get-h a) = add-states
(get-m b, get-h b) (λl. SOME p. get-m a l = padd (get-m b l) p, get-h a› get-h-m
plus-def)
  qed
  then show ?thesis
    by (meson PartialSA.greater-def)
qed

```

```

lemma larger-implies-greater-mask-hl:
  assumes a ≽ b
  shows pgte (get-m a hl) (get-m b hl)
  using PartialSA.greater-def assms p-greater-exists plus-charact(1) by auto

```

```

lemma larger-implies-larger-heap:
  assumes a ≽ b
  shows larger-heap (get-h a) (get-h b)
  by (metis (full-types) PartialSA.greater-equiv assms larger-heapI map-add-find-right
plus-charact(2))

```

```

lemma compatibleI:

```

```

assumes compatible-heaps (get-h a) (get-h b)
  and valid-mask (add-masks (get-m a) (get-m b))
  shows a |#| b
using PartialSA.defined-def assms(1) assms(2) plus-ab-defined by presburger

end

```

2 Combinable Magic Wands

Note that, in this theory, assertions are represented as semantic assertions, i.e., as the set of states in which they hold.

```

theory CombinableWands
  imports PartialHeapSA
begin

```

2.1 Definitions

```

type-synonym sem-assertion = state set

```

```

fun multiply :: prat ⇒ state ⇒ state where
  multiply p φ = Abs-state (multiply-mask p (get-m φ), get-h φ)

```

Because we work in an intuitionistic setting, a fraction of an assertion is defined using the upper-closure operator.

```

fun multiply-sem-assertion :: prat ⇒ sem-assertion ⇒ sem-assertion where
  multiply-sem-assertion p P = PartialSA.upper-closure (multiply p ‘ P)

```

```

definition combinable :: sem-assertion ⇒ bool where
  combinable P ⟷ (∀ α β. ppos α ∧ ppos β ∧ pgte pwrite (padd α β) ⟶
    (multiply-sem-assertion α P) ⊗ (multiply-sem-assertion β P) ⊆ multiply-sem-assertion
    (padd α β) P)

```

```

definition scaled where
  scaled φ = { multiply p φ | p. ppos p ∧ pgte pwrite p }

```

```

definition comp-min-mask :: mask ⇒ (mask ⇒ mask) where
  comp-min-mask b a hl = pmin (a hl) (comp-one (b hl))

```

```

definition scalable where
  scalable w a ⟷ (∀ φ ∈ scaled w. ¬ a |#| φ)

```

```

definition R where
  R a w = (if scalable w a then w else Abs-state (comp-min-mask (get-m a) (get-m
  w), get-h w))

```

```

definition cwand where
  cwand A B = { w | w. ∀ a x. a ∈ A ∧ Some x = R a w ⊕ a ⟶ x ∈ B }

```

definition *wand* :: *sem-assertion* \Rightarrow *sem-assertion* \Rightarrow *sem-assertion* **where**
wand $A B = \{ w \mid w. \forall a x. a \in A \wedge \text{Some } x = w \oplus a \longrightarrow x \in B \}$

definition *intuitionistic* **where**
intuitionistic $A \longleftrightarrow (\forall a b. a \succeq b \wedge b \in A \longrightarrow a \in A)$

definition *binary-mask* :: *mask* \Rightarrow *mask* **where**
binary-mask $\pi l = (\text{if } \pi l = \text{pwrite then pwrite else pnone})$

definition *binary* :: *sem-assertion* \Rightarrow *bool* **where**
binary $A \longleftrightarrow (\forall \varphi \in A. \text{Abs-state } (\text{binary-mask } (\text{get-m } \varphi), \text{get-h } \varphi) \in A)$

2.2 Lemmas

lemma *wand-equiv-def*:

wand $A B = \{ \varphi \mid \varphi. A \otimes \{\varphi\} \subseteq B \}$

proof

show *wand* $A B \subseteq \{ \varphi \mid \varphi. A \otimes \{\varphi\} \subseteq B \}$

proof

fix w **assume** $w \in \text{wand } A B$

have $A \otimes \{w\} \subseteq B$

proof

fix x **assume** $x \in A \otimes \{w\}$

then show $x \in B$

using *PartialSA.add-set-elem* $\langle w \in \text{wand } A B \rangle$ *commutative wand-def* **by**

auto

qed

then show $w \in \{ \varphi \mid \varphi. A \otimes \{\varphi\} \subseteq B \}$

by *simp*

qed

show $\{ \varphi \mid \varphi. A \otimes \{\varphi\} \subseteq B \} \subseteq \text{wand } A B$

proof

fix w **assume** $w \in \{ \varphi \mid \varphi. A \otimes \{\varphi\} \subseteq B \}$

have $\bigwedge a x. a \in A \wedge \text{Some } x = w \oplus a \implies x \in B$

proof –

fix $a x$ **assume** $a \in A \wedge \text{Some } x = w \oplus a$

then have $x \in A \otimes \{w\}$

using *PartialSA.add-set-elem* *PartialSA.commutative* **by** *auto*

then show $x \in B$

using $\langle w \in \{ \varphi \mid \varphi. A \otimes \{\varphi\} \subseteq B \} \rangle$ **by** *blast*

qed

then show $w \in \text{wand } A B$

using *wand-def* **by** *force*

qed

qed

lemma *w-in-scaled*:

$w \in \text{scaled } w$

proof –

have *multiply pwrite* $w = w$
by (*simp add: Rep-state-inverse mult-write-mask*)
then show *?thesis*
by (*metis (mono-tags, lifting) half-between-0-1 half-plus-half mem-Collect-eq not-pgte-charact pgt-implies-pgte ppos-add scaled-def*)
qed

lemma *non-scalable-instantiate*:
assumes \neg *scalable* w a
shows $\exists p. ppos\ p \wedge pgte\ pwrite\ p \wedge a\ |\#|$ *multiply* $p\ w$
using *assms scalable-def scaled-def* **by** *auto*

lemma *compatible-same-mask*:
assumes *valid-mask (add-masks a w)*
shows $w = comp\ min\ mask\ a\ w$
proof (*rule ext*)
fix x
have *pgte pwrite (padd (a x) (w x))*
by (*metis add-masks.simps assms valid-mask.elims(1)*)
moreover have *padd (a x) (comp-one (a x)) = pwrite*
by (*meson assms padd-comp-one upper-valid-aux valid-mask.elims(1)*)
then have *pgte (comp-one (a x)) (w x)*
by (*metis add-le-cancel-left calculation padd.rep-eq pgte.rep-eq*)
then show $w\ x = comp\ min\ mask\ a\ w\ x$
by (*metis comp-min-mask-def pmin-comm pmin-is*)
qed

lemma *R-smaller*:
 $w \succeq R\ a\ w$
proof (*cases scalable w a*)
case *True*
then show *?thesis*
by (*simp add: PartialSA.succ-refl R-def*)
next
case *False*
then have $R\ a\ w = Abs\ state\ (comp\ min\ mask\ (get\ m\ a)\ (get\ m\ w),\ get\ h\ w)$
by (*meson R-def*)
moreover have *greater-mask (get-m w) (comp-min-mask (get-m a) (get-m w))*
proof (*rule greater-maskI*)
fix hl **show** *pgte (get-m w hl) (comp-min-mask (get-m a) (get-m w) hl)*
by (*simp add: comp-min-mask-def pmin-greater*)
qed
ultimately show *?thesis*
by (*metis Abs-state-cases larger-heap-refl Rep-state-cases Rep-state-inverse fst-conv get-h-m greaterI greater-mask-def mem-Collect-eq snd-conv valid-state-decompose*)
qed

lemma *R-compatible-same*:
assumes $a\ |\#|$ w

shows $R\ a\ w = w$
proof –
have $\neg\ \text{scalable}\ w\ a$
using *assms scalable-def w-in-scaled* **by** *blast*
then have $R\ a\ w = \text{Abs-state}\ (\text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w),\ \text{get-h}\ w)$
using *R-def* **by** *auto*
then show *?thesis*
by (*metis PartialSA.defined-def Rep-state-inverse assms compatible-same-mask*
get-h.simps get-m.simps plus-ab-defined prod.collapse)
qed

lemma *in-cwand*:
assumes $\bigwedge a\ x.\ a \in A \wedge \text{Some}\ x = R\ a\ w \oplus a \implies x \in B$
shows $w \in \text{cwand}\ A\ B$
using *assms cwand-def* **by** *force*

lemma *wandI*:
assumes $\bigwedge a\ x.\ a \in A \wedge \text{Some}\ x = a \oplus w \implies x \in B$
shows $w \in \text{wand}\ A\ B$
proof –
have $A \otimes \{w\} \subseteq B$
proof (*rule subsetI*)
fix x **assume** $x \in A \otimes \{w\}$
then obtain a **where** $\text{Some}\ x = a \oplus w\ a \in A$
using *PartialSA.add-set-elim* **by** *auto*
then show $x \in B$
using *assms* **by** *blast*
qed
then show *?thesis*
using *wand-equiv-def* **by** *force*
qed

lemma *non-scalable-R-charact*:
assumes $\neg\ \text{scalable}\ w\ a$
shows $\text{get-m}\ (R\ a\ w) = \text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w) \wedge \text{get-h}\ (R\ a\ w) = \text{get-h}\ w$
proof –
have $R\ a\ w = \text{Abs-state}\ (\text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w),\ \text{get-h}\ w)$
using *R-def assms* **by** *auto*
moreover have $\text{valid-state}\ (\text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w),\ \text{get-h}\ w)$
proof (*rule valid-stateI*)
show $\text{valid-mask}\ (\text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w))$
proof (*rule valid-maskI*)
show $\bigwedge f.\ \text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w)\ (\text{null},\ f) = \text{pnone}$
by (*metis (no-types, opaque-lifting) PartialSA.unit-neutral add-masks.simps*
comp-min-mask-def option.distinct(1) p-greater-exists padd-zero plus-ab-defined pmin-greater
valid-mask.simps)
fix hl **show** $\text{pgte}\ \text{pwrite}\ (\text{comp-min-mask}\ (\text{get-m}\ a)\ (\text{get-m}\ w)\ hl)$
by (*metis PartialSA.unit-neutral comp-min-mask-def greater-mask-def greater-mask-equiv-def*)

option.distinct(1) plus-ab-defined pmin-greater upper-valid-aux valid-mask.simps
qed
fix *hl* **assume** *ppos (comp-min-mask (get-m a) (get-m w) hl)*
show *get-h w hl ≠ None*
by (*metis Rep-state ⟨ppos (comp-min-mask (get-m a) (get-m w) hl)⟩ comp-min-mask-def*
get-h.simps get-pre(2) mem-Collect-eq p-greater-exists pmin-greater ppos-add prod.collapse
valid-heap-def valid-state.simps)
qed
ultimately show *?thesis*
by (*metis Rep-state-cases Rep-state-inverse fst-conv get-h.simps get-m.simps*
mem-Collect-eq snd-conv)
qed

lemma *valid-bin*:
valid-state (binary-mask (get-m a), get-h a)
proof (*rule valid-stateI*)
show *valid-mask (binary-mask (get-m a))*
by (*metis PartialSA.unit-neutral binary-mask-def minus-empty option.discI*
plus-ab-defined unit-charact(2) valid-mask.elims(2) valid-mask.elims(3))
show $\bigwedge hl. ppos (binary-mask (get-m a) hl) \implies get-h a hl \neq None$
by (*metis Rep-prat Rep-state binary-mask-def get-h.simps get-pre(2) leD mem-Collect-eq*
pnone.rep-eq ppos.rep-eq prod.collapse valid-heap-def valid-state.simps)
qed

lemma *in-multiply-sem*:
assumes $x \in multiply\text{-}sem\text{-}assertion\ p\ A$
shows $\exists a \in A. x \succeq multiply\ p\ a$
using *PartialSA.sep-algebra-axioms assms greater-def sep-algebra.upper-closure-def*
by *fastforce*

lemma *get-h-multiply*:
assumes *pgte pwrite p*
shows $get-h (multiply\ p\ x) = get-h\ x$
using *Abs-state-inverse assms multiply-valid by auto*

lemma *in-multiply-refl*:
assumes $x \in A$
shows $multiply\ p\ x \in multiply\text{-}sem\text{-}assertion\ p\ A$
using *PartialSA.succ-refl PartialSA.upper-closure-def assms by fastforce*

lemma *get-m-smaller*:
assumes *pgte pwrite p*
shows $get-m (multiply\ p\ a)\ hl = pmult\ p (get-m\ a\ hl)$
using *Abs-state-inverse assms multiply-mask-def multiply-valid by auto*

lemma *get-m-smaller-mask*:
assumes *pgte pwrite p*
shows $get-m (multiply\ p\ a) = multiply\text{-}mask\ p (get-m\ a)$
using *Abs-state-inverse assms multiply-mask-def multiply-valid by auto*

lemma *multiply-order*:
assumes *pgte pwrite p*
and $a \succeq b$
shows $\text{multiply } p \ a \succeq \text{multiply } p \ b$
proof (*rule greaterI*)
show *larger-heap* (*get-h* (*multiply p a*)) (*get-h* (*multiply p b*))
using *assms(1) assms(2) get-h-multiply larger-implies-larger-heap* **by** *presburger*
show *greater-mask* (*get-m* (*multiply p a*)) (*get-m* (*multiply p b*))
by (*metis assms(1) assms(2) get-m-smaller-mask greater-maskI larger-implies-greater-mask-hl mult-greater*)
qed

lemma *multiply-twice*:
assumes *pgte pwrite a \wedge pgte pwrite b*
shows $\text{multiply } a \ (\text{multiply } b \ x) = \text{multiply } (\text{pmult } a \ b) \ x$
proof –
have *get-h* (*multiply* (*pmult a b*) *x*) = *get-h x*
by (*metis assms get-h-multiply p-greater-exists padd-asso pmult-order pmult-special(1)*)
moreover **have** *get-h* (*multiply a* (*multiply b x*)) = *get-h x*
using *assms get-h-multiply* **by** *presburger*
moreover **have** *get-m* (*multiply a* (*multiply b x*)) = *get-m* (*multiply* (*pmult a b*)
x)
proof (*rule ext*)
fix *l*
have *pgte pwrite* (*pmult a b*) **using** *multiply-smaller-pwrite assms* **by** *simp*
then **have** *get-m* (*multiply* (*pmult a b*) *x*) *l* = *pmult* (*pmult a b*) (*get-m x l*)
using *get-m-smaller* **by** *blast*
then **show** *get-m* (*multiply a* (*multiply b x*)) *l* = *get-m* (*multiply* (*pmult a b*)
x) *l*
by (*metis Rep-prat-inverse assms get-m-smaller mult.assoc pmult.rep-eq*)
qed
ultimately **show** *?thesis*
using *state-ext* **by** *presburger*
qed

lemma *valid-mask-add-comp-min*:
assumes *valid-mask a*
and *valid-mask b*
shows *valid-mask* (*add-masks* (*comp-min-mask b a*) *b*)
proof (*rule valid-maskI*)
show $\bigwedge f. \text{add-masks} \ (\text{comp-min-mask } b \ a) \ b \ (\text{null}, f) = \text{pnone}$
proof –
fix *f*
have *comp-min-mask b a* (*null, f*) = *pnone*
by (*metis assms(1) comp-min-mask-def p-greater-exists padd-zero pmin-greater valid-mask.simps*)
then **show** *add-masks* (*comp-min-mask b a*) *b* (*null, f*) = *pnone*
by (*metis add-masks.simps assms(2) padd-zero valid-mask.simps*)

```

qed
fix hl show pgte pwrite (add-masks (comp-min-mask b a) b hl)
proof (cases pgte (a hl) (comp-one (b hl)))
  case True
  then have add-masks (comp-min-mask b a) b hl = padd (comp-one (b hl)) (b
hl)
    by (simp add: comp-min-mask-def pmin-is)
  then have add-masks (comp-min-mask b a) b hl = pwrite
    by (metis assms(2) padd-comm padd-comp-one valid-mask.simps)
  then show ?thesis
    by (simp add: pgte.rep-eq)
next
  case False
  then have comp-min-mask b a hl = a hl
    by (metis comp-min-mask-def not-pgte-charact pgt-implies-pgte pmin-comm
pmin-is)
  then have add-masks (comp-min-mask b a) b hl = padd (a hl) (b hl)
    by auto
  moreover have pgte (padd (comp-one (b hl)) (b hl)) (padd (a hl) (b hl))
    using False padd.rep-eq pgte.rep-eq by force
  moreover have padd (comp-one (b hl)) (b hl) = pwrite
    by (metis assms(2) padd-comm padd-comp-one valid-mask.simps)
  ultimately show ?thesis by simp
qed
qed

```

2.3 The combinable wand is stronger than the original wand

lemma *cwand-stronger*:

$cwand\ A\ B \subseteq wand\ A\ B$

proof

fix w **assume** $asm0: w \in cwand\ A\ B$

then have $r: \bigwedge a\ x. a \in A \wedge Some\ x = R\ a\ w \oplus a \implies x \in B$

using *cwand-def* **by** *blast*

show $w \in wand\ A\ B$

proof (*rule wandI*)

fix $a\ x$ **assume** $asm1: a \in A \wedge Some\ x = a \oplus w$

then have $R\ a\ w = w$

by (*metis PartialSA.defined-def R-compatible-same option.distinct(1)*)

then show $x \in B$

by (*metis PartialSA.commutative asm1 r*)

qed

qed

2.4 The combinable wand is the same as the original wand when the left-hand side is binary

lemma *binary-same*:

assumes *binary A*


```

    and intuitionistic B
  shows wand A B  $\subseteq$  cwand A B
proof (rule subsetI)
  fix w assume w  $\in$  wand A B
  then have asm0: A  $\otimes$  {w}  $\subseteq$  B
    by (simp add: wand-equiv-def)
  show w  $\in$  cwand A B
proof (rule in-cwand)
  fix a x assume a  $\in$  A  $\wedge$  Some x = R a w  $\oplus$  a
  show x  $\in$  B
proof (cases scalable w a)
  case True
  then show ?thesis
    by (metis PartialSA.commutative PartialSA.defined-def R-def  $\langle$ a  $\in$  A  $\wedge$ 
Some x = R a w  $\oplus$  a $\rangle$  option.distinct(1) scalable-def w-in-scaled)
  next
  case False
  then have get-m (R a w) = comp-min-mask (get-m a) (get-m w)  $\wedge$  get-h (R
a w) = get-h w
    using non-scalable-R-charact by blast
  moreover have Abs-state (binary-mask (get-m a), get-h a)  $\in$  A
    using  $\langle$ a  $\in$  A  $\wedge$  Some x = R a w  $\oplus$  a $\rangle$  assms(1) binary-def by blast
  moreover have greater-mask (add-masks (comp-min-mask (get-m a) (get-m
w)) (get-m a))
    (add-masks (binary-mask (get-m a)) (get-m w))
    proof (rule greater-maskI)
      fix hl show pgte (add-masks (comp-min-mask (get-m a) (get-m w)) (get-m
a) hl) (add-masks (binary-mask (get-m a)) (get-m w) hl)
      proof (cases get-m a hl = pwrite)
        case True
        obtain  $\varphi$  where  $\varphi \in$  scaled w a  $|\#|$   $\varphi$  using False scalable-def[of w a]
          by blast
        then obtain p where ppos p pgte pwrite p multiply p w  $|\#|$  a
          using PartialSA.commutative PartialSA.defined-def mem-Collect-eq
scaled-def by auto
        have get-m w hl = pnone
          proof (rule ccontr)
            assume get-m w hl  $\neq$  pnone
            then have ppos (get-m w hl)
              by (metis less-add-same-cancel1 not-pgte-charact p-greater-exists
padd.rep-eq padd-zero pgt.rep-eq ppos.rep-eq)
            moreover have get-m (multiply p  $\varphi$ ) = multiply-mask p (get-m  $\varphi$ )
              using multiply-valid[of p  $\varphi$ ] multiply.simps[of p  $\varphi$ ]
            by (metis Rep-state-cases Rep-state-inverse  $\langle$ pgte pwrite p $\rangle$  fst-conv
get-pre(2) mem-Collect-eq)
            then have ppos (get-m (multiply p w) hl) using pmult-ppos
              by (metis Rep-state-cases Rep-state-inverse  $\langle$ pgte pwrite p $\rangle$   $\langle$ ppos
p $\rangle$  calculation fst-conv get-pre(2) mem-Collect-eq multiply.simps multiply-mask-def
multiply-valid)

```

then have pgt ($padd$ ($get-m$ ($multiply$ p w) hl) ($get-m$ a hl)) $pwrite$
by ($metis$ $True$ $add-le-same-cancel2$ leD $not-pgte-charact$ $padd.rep-eq$ $pgte.rep-eq$ $ppos.rep-eq$)
then have \neg $valid-mask$ ($add-masks$ ($get-m$ ($multiply$ p w)) ($get-m$ a))
by ($metis$ $add-masks.elims$ $not-pgte-charact$ $valid-mask.elims(1)$)
then show $False$
using $PartialSA.defined-def$ $\langle multiply$ p w $| \# |$ $a \rangle$ $plus-ab-defined$ **by**
 $blast$
qed
then show $?thesis$
by ($metis$ $Rep-prat-inverse$ $add.right-neutral$ $add-masks.simps$ $binary-mask-def$ $p-greater-exists$ $padd.rep-eq$ $padd-comm$ $pnone.rep-eq$)
next
case $False$
then have $add-masks$ ($binary-mask$ ($get-m$ a)) ($get-m$ w) $hl = get-m$ w hl
by ($metis$ $Rep-prat-inject$ $add.right-neutral$ $add-masks.simps$ $binary-mask-def$ $padd.rep-eq$ $padd-comm$ $pnone.rep-eq$)
then show $?thesis$
proof ($cases$ $pgte$ ($get-m$ w hl) ($comp-one$ ($get-m$ a hl)))
case $True$
then have $comp-min-mask$ ($get-m$ a) ($get-m$ w) $hl = comp-one$ ($get-m$ a hl)
using $comp-min-mask-def$ $pmin-is$ **by** $presburger$
then have $add-masks$ ($comp-min-mask$ ($get-m$ a) ($get-m$ w)) ($get-m$ a) $hl = pwrite$
by ($metis$ $PartialSA.unit-neutral$ $add-masks.simps$ $add-masks-comm$ $minus-empty$ $option.distinct(1)$ $padd-comp-one$ $plus-ab-defined$ $unit-charact(2)$ $valid-mask.simps$)
then show $?thesis$
by ($metis$ $PartialSA.unit-neutral$ $\langle add-masks$ ($binary-mask$ ($get-m$ a)) ($get-m$ w) $hl = get-m$ w hl \rangle $minus-empty$ $option.distinct(1)$ $plus-ab-defined$ $unit-charact(2)$ $valid-mask.simps$)
next
case $False$
then have $comp-min-mask$ ($get-m$ a) ($get-m$ w) $hl = get-m$ w hl
by ($metis$ $comp-min-mask-def$ $not-pgte-charact$ $pgt-implies-pgte$ $pmin-comm$ $pmin-is$)
then show $?thesis$
using $\langle add-masks$ ($binary-mask$ ($get-m$ a)) ($get-m$ w) $hl = get-m$ w hl \rangle $p-greater-exists$ **by** $auto$
qed
qed
qed
then have $valid-mask$ ($add-masks$ ($binary-mask$ ($get-m$ a)) ($get-m$ w))
by ($metis$ $\langle a \in A \wedge Some$ $x = R$ a $w \oplus a \rangle$ $calculation(1)$ $greater-mask-def$ $option.distinct(1)$ $plus-ab-defined$ $upper-valid-aux$)
moreover have $compatible-heaps$ ($get-h$ a) ($get-h$ w)
by ($metis$ $PartialSA.commutative$ $\langle a \in A \wedge Some$ $x = R$ a $w \oplus a \rangle$ $\langle get-m$ (R a w) $= comp-min-mask$ ($get-m$ a) ($get-m$ w) $\wedge get-h$ (R a w) $= get-h$ $w \rangle$ $option.simps(3)$ $plus-ab-defined$)

then obtain xx **where** $\text{Some } xx = \text{Abs-state } (\text{binary-mask } (\text{get-m } a), \text{get-h } a) \oplus w$
using $\text{Abs-state-inverse calculation compatible-def fst-conv plus-def valid-bin}$
by auto
then have $xx \in B$ **using** asm0
by $(\text{meson } \text{PartialSA.add-set-elem } \langle \text{Abs-state } (\text{binary-mask } (\text{get-m } a), \text{get-h } a) \in A \rangle \text{ singletonI subset-iff})$
moreover have $x \succeq xx$
proof $(\text{rule } \text{greaterI})$
show $\text{greater-mask } (\text{get-m } x) (\text{get-m } xx)$
using $\text{Abs-state-inverse } \langle \text{Some } xx = \text{Abs-state } (\text{binary-mask } (\text{get-m } a), \text{get-h } a) \oplus w \rangle \langle a \in A \wedge \text{Some } x = R a w \oplus a \rangle \langle \text{greater-mask } (\text{add-masks } (\text{comp-min-mask } (\text{get-m } a) (\text{get-m } w)) (\text{get-m } a)) (\text{add-masks } (\text{binary-mask } (\text{get-m } a)) (\text{get-m } w)) \rangle$
 $\text{calculation(1) plus-charact(1) valid-bin}$ **by** auto
show $\text{larger-heap } (\text{get-h } x) (\text{get-h } xx)$
proof $(\text{rule } \text{larger-heapI})$
fix $hl \ xa$ **assume** $\text{get-h } xx \ hl = \text{Some } xa$
then show $\text{get-h } x \ hl = \text{Some } xa$
by $(\text{metis } \text{PartialSA.commutative Rep-state-cases Rep-state-inverse } \langle \text{Some } xx = \text{Abs-state } (\text{binary-mask } (\text{get-m } a), \text{get-h } a) \oplus w \rangle \langle a \in A \wedge \text{Some } x = R a w \oplus a \rangle \text{ calculation(1) get-h.simps mem-Collect-eq plus-charact(2) snd-conv valid-bin})$
qed
qed
ultimately show $?thesis$
using $\text{assms(2) intuitionistic-def}$ **by** blast
qed
qed
qed

2.5 The combinable wand is combinable

lemma combinableI :

assumes $\bigwedge a \ b. \text{ppos } a \wedge \text{ppos } b \wedge \text{padd } a \ b = \text{pwrite} \implies (\text{multiply-sem-assertion } a \ (\text{cwand } A \ B)) \otimes (\text{multiply-sem-assertion } b \ (\text{cwand } A \ B)) \subseteq \text{cwand } A \ B$

shows $\text{combinable } (\text{cwand } A \ B)$

proof –

have $\bigwedge a \ b. \text{ppos } a \wedge \text{ppos } b \wedge \text{pgte } \text{pwrite } (\text{padd } a \ b) \implies (\text{multiply-sem-assertion } a \ (\text{cwand } A \ B)) \otimes (\text{multiply-sem-assertion } b \ (\text{cwand } A \ B)) \subseteq \text{multiply-sem-assertion } (\text{padd } a \ b) \ (\text{cwand } A \ B)$

proof –

fix $a \ b$ **assume** $\text{asm0}: \text{ppos } a \wedge \text{ppos } b \wedge \text{pgte } \text{pwrite } (\text{padd } a \ b)$

then have $\text{pgte } \text{pwrite } a \wedge \text{pgte } \text{pwrite } b$

using $\text{padd.rep-eq pgte.rep-eq ppos.rep-eq}$ **by** auto

show $(\text{multiply-sem-assertion } a \ (\text{cwand } A \ B)) \otimes (\text{multiply-sem-assertion } b \ (\text{cwand } A \ B)) \subseteq \text{multiply-sem-assertion } (\text{padd } a \ b) \ (\text{cwand } A \ B)$

proof

fix x **assume** $x \in \text{multiply-sem-assertion } a \ (\text{cwand } A \ B) \otimes \text{multiply-sem-assertion } b \ (\text{cwand } A \ B)$

then obtain $xa \ xb$ **where** $\text{Some } x = xa \oplus xb \ xa \in \text{multiply-sem-assertion } a$

```

(cwand A B) xb ∈ multiply-sem-assertion b (cwand A B)
  by (meson PartialSA.add-set-elem)
  then obtain wa wb where wa ∈ cwand A B wb ∈ cwand A B xa ≥ multiply
a wa xb ≥ multiply b wb
  by (meson in-multiply-sem)
  let ?a = pdiv a (padd a b)
  let ?b = pdiv b (padd a b)
  have pgte pwrite ?a ∧ pgte pwrite ?b
  using asm0 p-greater-exists padd-comm pdiv-smaller ppos-add by blast
  have multiply ?a wa |#| multiply ?b wb
  proof (rule compatibleI)
    show compatible-heaps (get-h (multiply (pdiv a (padd a b)) wa)) (get-h
(multiply (pdiv b (padd a b)) wb))
    proof -
      have compatible-heaps (get-h (multiply a wa)) (get-h (multiply b wb))
      by (metis PartialSA.asso2 PartialSA.asso3 PartialSA.greater-equiv
PartialSA.minus-some ⟨Some x = xa ⊕ xb⟩ ⟨xa ≥ multiply a wa⟩ ⟨xb ≥ multiply b
wb⟩ option.simps(3) plus-ab-defined)
      moreover have get-h (multiply (pdiv a (padd a b)) wa) = get-h (multiply
a wa) ∧ get-h (multiply (pdiv b (padd a b)) wb) = get-h (multiply b wb)
      proof -
        have pgte pwrite a ∧ pgte pwrite b
        by (metis asm0 p-greater-exists padd-asso padd-comm)
        moreover have pgte pwrite ?a ∧ pgte pwrite ?b
        using asm0 p-greater-exists padd-comm pdiv-smaller ppos-add by blast
        ultimately show ?thesis
        using get-h-multiply by presburger
      qed
      then show ?thesis
      using calculation by presburger
    qed
    show valid-mask (add-masks (get-m (multiply (pdiv a (padd a b)) wa))
(get-m (multiply (pdiv b (padd a b)) wb)))
    proof (rule valid-maskI)
      show ∧f. add-masks (get-m (multiply (pdiv a (padd a b)) wa)) (get-m
(multiply (pdiv b (padd a b)) wb)) (null, f) = pnone
      by (metis PartialSA.unit-neutral add-masks-equiv-valid-null option.distinct(1)
plus-ab-defined valid-mask.simps valid-null-def)
      fix hl have add-masks (get-m (multiply (pdiv a (padd a b)) wa)) (get-m
(multiply (pdiv b (padd a b)) wb)) hl
      = padd (pmult ?a (get-m wa hl)) (pmult ?b (get-m wb hl))
      proof -
        have get-m (multiply ?a wa) hl = pmult ?a (get-m wa hl)
        using Abs-state-inverse ⟨pgte pwrite (pdiv a (padd a b)) ∧ pgte pwrite
(pdiv b (padd a b))⟩ multiply-mask-def multiply-valid by auto
        moreover have get-m (multiply ?b wb) hl = pmult ?b (get-m wb hl)
        using Abs-state-inverse ⟨pgte pwrite (pdiv a (padd a b)) ∧ pgte pwrite
(pdiv b (padd a b))⟩ multiply-mask-def multiply-valid by auto
        ultimately show ?thesis by simp

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qed
  moreover have  $pgte\ pwrite\ (padd\ (pmult\ ?a\ (get-m\ wa\ hl))\ (pmult\ ?b\ (get-m\ wb\ hl)))$ 
  proof (rule padd-one-ineq-sum)
    show  $pgte\ pwrite\ (get-m\ wa\ hl)$ 
      by (metis PartialSA.unit-neutral option.discI plus-ab-defined upper-valid-aux valid-mask.simps)
    show  $pgte\ pwrite\ (get-m\ wb\ hl)$ 
      by (metis PartialSA.unit-neutral option.discI plus-ab-defined upper-valid-aux valid-mask.simps)
    show  $padd\ (pdiv\ a\ (padd\ a\ b))\ (pdiv\ b\ (padd\ a\ b)) = pwrite$ 
      using asm0 sum-coeff by blast
  qed
  ultimately show  $pgte\ pwrite\ (add-masks\ (get-m\ (multiply\ (pdiv\ a\ (padd\ a\ b))\ wa))\ (get-m\ (multiply\ (pdiv\ b\ (padd\ a\ b))\ wb))\ hl)$ 
    by presburger
  qed
qed
then obtain  $xx$  where  $Some\ xx = multiply\ ?a\ wa \oplus multiply\ ?b\ wb$ 
  using PartialSA.defined-def by auto
moreover have  $(multiply-sem-assertion\ ?a\ (cwand\ A\ B)) \otimes (multiply-sem-assertion\ ?b\ (cwand\ A\ B)) \subseteq cwand\ A\ B$ 
  proof (rule assms)
    show  $ppos\ (pdiv\ a\ (padd\ a\ b)) \wedge ppos\ (pdiv\ b\ (padd\ a\ b)) \wedge padd\ (pdiv\ a\ (padd\ a\ b))\ (pdiv\ b\ (padd\ a\ b)) = pwrite$ 
      using asm0 padd.rep-eq pdiv.rep-eq ppos.rep-eq sum-coeff by auto
  qed
ultimately have  $xx \in cwand\ A\ B$ 
proof –
  have  $multiply\ ?a\ wa \in multiply-sem-assertion\ ?a\ (cwand\ A\ B)$ 
    using  $\langle wa \in cwand\ A\ B \rangle$  in-multiply-refl by presburger
  moreover have  $multiply\ ?b\ wb \in multiply-sem-assertion\ ?b\ (cwand\ A\ B)$ 
    by (meson  $\langle wb \in cwand\ A\ B \rangle$  in-multiply-refl)
  ultimately show  $?thesis$ 
    using PartialSA.add-set-def  $\langle Some\ xx = multiply\ (pdiv\ a\ (padd\ a\ b))\ wa \oplus multiply\ (pdiv\ b\ (padd\ a\ b))\ wb \rangle$   $\langle multiply-sem-assertion\ (pdiv\ a\ (padd\ a\ b))\ (cwand\ A\ B) \otimes multiply-sem-assertion\ (pdiv\ b\ (padd\ a\ b))\ (cwand\ A\ B) \subseteq cwand\ A\ B \rangle$  by fastforce
  qed
moreover have  $x \succeq multiply\ (padd\ a\ b)\ xx$ 
proof (rule greaterI)
  have valid-state  $(multiply-mask\ (padd\ a\ b)\ (get-m\ xx),\ get-h\ xx)$ 
    using asm0 multiply-valid by blast
  show larger-heap  $(get-h\ x)\ (get-h\ (multiply\ (padd\ a\ b)\ xx))$ 
proof –
  have  $get-h\ (multiply\ (padd\ a\ b)\ xx) = get-h\ xx$ 
    using asm0 get-h-multiply by blast
  moreover have  $get-h\ xx = get-h\ wa ++ get-h\ wb$ 
    by (metis  $\langle Some\ xx = multiply\ (pdiv\ a\ (padd\ a\ b))\ wa \oplus multiply\ (pdiv\ b\ (padd\ a\ b))\ wb \rangle$  in-multiply-refl)

```

b ($\text{padd } a \ b$) wb asm0 get-h-multiply $p\text{-greater-exists}$ padd-comm $\text{plus-charact}(2)$ sum-coeff)
moreover have $\text{get-h } x = \text{get-h } xa ++ \text{get-h } xb$
using $\langle \text{Some } x = xa \oplus xb \rangle$ $\text{plus-charact}(2)$ **by** presburger
moreover have $\text{get-h } wa = \text{get-h } (\text{multiply } a \ wa) \wedge \text{get-h } wb = \text{get-h } (\text{multiply } b \ wb)$
by (metis asm0 get-h-multiply order-trans $p\text{-greater-exists}$ padd-comm pgte.rep-eq)
moreover have $\text{larger-heap } (\text{get-h } xa) (\text{get-h } wa) \wedge \text{larger-heap } (\text{get-h } xb) (\text{get-h } wb)$
using $\langle xa \succeq \text{multiply } a \ wa \rangle \langle xb \succeq \text{multiply } b \ wb \rangle$ $\text{calculation}(4)$ $\text{larger-implies-larger-heap}$ **by** presburger
ultimately show $?thesis$
by (metis $\langle \text{Some } x = xa \oplus xb \rangle$ $\text{larger-heaps-sum-ineq}$ $\text{option.simps}(3)$ plus-ab-defined)
qed
show $\text{greater-mask } (\text{get-m } x) (\text{get-m } (\text{multiply } (\text{padd } a \ b) \ xx))$
proof (rule greater-maskI)
fix hl
have $\text{pgte } (\text{get-m } x \ hl) (\text{padd } (\text{get-m } xa \ hl) (\text{get-m } xb \ hl))$
using $\langle \text{Some } x = xa \oplus xb \rangle$ pgte.rep-eq $\text{plus-charact}(1)$ **by** auto
moreover have $\text{pgte } (\text{get-m } xa \ hl) (\text{get-m } (\text{multiply } a \ wa) \ hl) \wedge \text{pgte } (\text{get-m } xb \ hl) (\text{get-m } (\text{multiply } b \ wb) \ hl)$
using $\langle xa \succeq \text{multiply } a \ wa \rangle \langle xb \succeq \text{multiply } b \ wb \rangle$ $\text{larger-implies-greater-mask-hl}$
by blast
moreover have $\text{get-m } (\text{multiply } (\text{padd } a \ b) \ xx) \ hl = \text{pmult } (\text{padd } a \ b) (\text{get-m } xx \ hl)$
by (metis Rep-state-cases Rep-state-inverse $\langle \text{valid-state } (\text{multiply-mask } (\text{padd } a \ b) (\text{get-m } xx), \text{get-h } xx) \rangle$ fst-conv $\text{get-pre}(2)$ mem-Collect-eq multiply.simps multiply-mask-def)
moreover have $\dots = \text{padd } (\text{pmult } (\text{pmult } (\text{padd } a \ b) \ ?a) (\text{get-m } wa \ hl)) (\text{pmult } (\text{pmult } (\text{padd } a \ b) \ ?b) (\text{get-m } wb \ hl))$
proof –
have $\text{get-m } (\text{multiply } ?a \ wa) \ hl = \text{pmult } ?a (\text{get-m } wa \ hl)$
by (metis Abs-state-inverse asm0 fst-conv $\text{get-pre}(2)$ mem-Collect-eq multiply.simps multiply-mask-def multiply-valid $p\text{-greater-exists}$ sum-coeff)
moreover have $\text{get-m } (\text{multiply } ?b \ wb) \ hl = \text{pmult } ?b (\text{get-m } wb \ hl)$
by (metis Abs-state-inverse asm0 fst-conv $\text{get-pre}(2)$ mem-Collect-eq multiply.simps multiply-mask-def multiply-valid $p\text{-greater-exists}$ padd-comm pdiv-smaller ppos-add)
ultimately have $\text{get-m } xx \ hl = \text{padd } (\text{pmult } ?a (\text{get-m } wa \ hl)) (\text{pmult } ?b (\text{get-m } wb \ hl))$
using $\langle \text{Some } xx = \text{multiply } (\text{pdiv } a \ (\text{padd } a \ b)) \ wa \oplus \text{multiply } (\text{pdiv } b \ (\text{padd } a \ b)) \ wb \rangle$ $\text{plus-charact}(1)$ **by** fastforce
then show $?thesis$
by (simp add: pmult-padd)
qed
moreover have $\dots = \text{padd } (\text{pmult } a (\text{get-m } wa \ hl)) (\text{pmult } b (\text{get-m } wb \ hl))$

```

      using asm0 pmult-pdiv-cancel ppos-add by presburger
    moreover have get-m (multiply a wa) hl = pmult a (get-m wa hl)  $\wedge$  get-m
(multiply b wb) hl = pmult b (get-m wb hl)
    proof -
      have valid-mask (multiply-mask a (get-m wa))
    using asm0 mult-add-states multiply-valid upper-valid-aux valid-state.simps
  by blast
    moreover have valid-mask (multiply-mask b (get-m wb))
    using asm0 mult-add-states multiply-valid upper-valid valid-state.simps
  by blast
    ultimately show ?thesis
      by (metis (no-types, lifting) Abs-state-inverse asm0 fst-conv get-pre(2)
mem-Collect-eq multiply.simps multiply-mask-def multiply-valid order-trans p-greater-exists
padd-comm pgte.rep-eq)
    qed
    ultimately show pgte (get-m x hl) (get-m (multiply (padd a b) xx) hl)
      by (simp add: padd.rep-eq pgte.rep-eq)
    qed
    qed
    ultimately show  $x \in \text{multiply-sem-assertion (padd a b) (cwand A B)}$ 
      by (metis PartialSA.up-closed-def PartialSA.upper-closure-up-closed in-multiply-refl
multiply-sem-assertion.simps)
    qed
    qed
    then show ?thesis
      using combinable-def by presburger
  qed

```

```

lemma combinable-cwand:
  assumes combinable B
    and intuitionistic B
  shows combinable (cwand A B)
proof (rule combinableI)
  fix  $\alpha \beta$  assume asm0: ppos  $\alpha \wedge$  ppos  $\beta \wedge$  padd  $\alpha \beta =$  pwrite
  then have pgte pwrite  $\alpha \wedge$  pgte pwrite  $\beta$ 
    by (metis p-greater-exists padd-comm)
  show multiply-sem-assertion  $\alpha$  (cwand A B)  $\otimes$  multiply-sem-assertion  $\beta$  (cwand
A B)  $\subseteq$  cwand A B
  proof
    fix w assume w  $\in$  multiply-sem-assertion  $\alpha$  (cwand A B)  $\otimes$  multiply-sem-assertion
 $\beta$  (cwand A B)
    then obtain xa xb where Some  $w = xa \oplus xb$  xa  $\in$  multiply-sem-assertion  $\alpha$ 
(cwand A B) xb  $\in$  multiply-sem-assertion  $\beta$  (cwand A B)
    by (meson PartialSA.add-set-elem)
    then obtain wa wb where wa  $\in$  cwand A B wb  $\in$  cwand A B xa  $\succeq$  multiply
 $\alpha$  wa xb  $\succeq$  multiply  $\beta$  wb
    by (meson in-multiply-sem)
    then obtain r:  $\bigwedge a x. a \in A \wedge$  Some  $x = R a wa \oplus a \implies x \in B \bigwedge a x. a \in$ 
A  $\wedge$  Some  $x = R a wb \oplus a \implies x \in B$ 

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```

    using cwand-def by blast
  show  $w \in \text{cwand } A \ B$ 
  proof (rule in-cwand)
    fix  $a \ x$  assume asm1:  $a \in A \wedge \text{Some } x = R \ a \ w \oplus \ a$ 
    have  $\neg \text{scalable } w \ a$ 
    proof (rule ccontr)
      assume  $\neg \neg \text{scalable } w \ a$ 
      then have  $R \ a \ w = w \wedge \neg a \ |\#| \ R \ a \ w$ 
        by (simp add: R-def scalable-def w-in-scaled)
      then show False
        using PartialSA.commutative PartialSA.defined-def asm1 by auto
    qed
    then have  $\text{get-h } (R \ a \ w) = \text{get-h } w \wedge \text{get-m } (R \ a \ w) = \text{comp-min-mask}$ 
      (get-m a) (get-m w)
      using non-scalable-R-charact by blast
    moreover obtain  $p$  where  $a \ |\#| \ \text{multiply } p \ w \ \text{ppos } p \wedge \text{pgte } p \ \text{pwrite } p$ 
      using  $\langle \neg \text{scalable } w \ a \rangle$  non-scalable-instantiate by blast
    moreover have  $\neg \text{scalable } w \ a$ 
    proof -
      have  $a \ |\#| \ \text{multiply } (p \ \text{mult } \alpha \ p) \ w$ 
      proof -
        have  $w \succeq x$  using  $\langle \text{Some } w = x \oplus \ x \rangle$  using PartialSA.greater-def by
        blast
        then have  $\text{multiply } p \ w \succeq \text{multiply } p \ x$ 
          using calculation(3) multiply-order by blast
        then have  $\text{multiply } p \ w \succeq \text{multiply } (p \ \text{mult } \alpha \ p) \ w$ 
          proof -
            have  $\text{multiply } p \ w \succeq \text{multiply } p \ (\text{multiply } \alpha \ w)$ 
              using PartialSA.succ-trans  $\langle w \succeq x \rangle \langle x \succeq \text{multiply } \alpha \ w \rangle$  calculation(3)
            multiply-order by blast
            then show ?thesis
              using  $\langle \text{pgte } p \ \text{write } \alpha \wedge \text{pgte } p \ \text{write } \beta \rangle$  calculation(3) multiply-twice
              pmult-comm by auto
          qed
        then show ?thesis
          using PartialSA.asso3 PartialSA.defined-def PartialSA.minus-some
          calculation(2) by fastforce
      qed
      moreover have  $\text{ppos } (p \ \text{mult } \alpha \ p) \wedge \text{pgte } p \ \text{write } (p \ \text{mult } \alpha \ p)$ 
        by (metis Rep-prat-inverse  $\langle \text{ppos } p \wedge \text{pgte } p \ \text{write } p \rangle$  add.right-neutral asm0
dual-order.strict-iff-order padd.rep-eq pgte.rep-eq pmult-comm pmult-ppos pmult-special(2)
pnone.rep-eq ppos.rep-eq ppos-eq-pnone padd-one-ineq-sum)
      ultimately show ?thesis
        using scalable-def scaled-def by auto
    qed
    then have  $\text{get-h } (R \ a \ w) = \text{get-h } w \wedge \text{get-m } (R \ a \ w) = \text{comp-min-mask}$ 
      (get-m a) (get-m w)
      using non-scalable-R-charact by blast
    moreover have  $R \ a \ w \ |\#| \ a$ 

```



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proof (rule compatibleI)
  have larger-heap (get-h w) (get-h xa)  $\wedge$  larger-heap (get-h xa) (get-h wa)
    by (metis PartialSA.commutative PartialSA.greater-equiv  $\langle$ Some w =
  xa  $\oplus$  xb $\rangle$   $\langle$ pgte pwrite  $\alpha$   $\wedge$  pgte pwrite  $\beta$  $\rangle$   $\langle$ xa  $\succeq$  multiply  $\alpha$  wa $\rangle$  get-h-multiply
  larger-implies-larger-heap)
  then show compatible-heaps (get-h (R a wa)) (get-h a)
    by (metis asm1 calculation(1) calculation(4) larger-heap-comp option.distinct(1) plus-ab-defined)
  show valid-mask (add-masks (get-m (R a wa)) (get-m a))
    by (metis PartialSA.unit-neutral calculation(4) minus-empty option.distinct(1) plus-ab-defined unit-charact(2) valid-mask-add-comp-min)
qed
then obtain ba where Some ba = R a wa  $\oplus$  a
  using PartialSA.defined-def by auto

moreover have  $\neg$  scalable wb a
proof –
  have a  $|\#|$  multiply (pmult  $\beta$  p) wb
  proof –
    have w  $\succeq$  xb using  $\langle$ Some w = xa  $\oplus$  xb $\rangle$ 
      using PartialSA.greater-equiv by blast
    then have multiply p w  $\succeq$  multiply p xb
      using calculation(3) multiply-order by blast
    then have multiply p w  $\succeq$  multiply (pmult  $\beta$  p) wb
  proof –
    have multiply p w  $\succeq$  multiply p (multiply  $\beta$  wb)
      using PartialSA.succ-trans  $\langle$ w  $\succeq$  xb $\rangle$   $\langle$ xb  $\succeq$  multiply  $\beta$  wb $\rangle$  calculation(3)
  multiply-order by blast
    then show ?thesis
      using  $\langle$ pgte pwrite  $\alpha$   $\wedge$  pgte pwrite  $\beta$  $\rangle$  calculation(3) multiply-twice
  pmult-comm by auto
    qed
    then show ?thesis
      using PartialSA.asso3 PartialSA.defined-def PartialSA.minus-some
  calculation(2) by fastforce
    qed
    moreover have ppos (pmult  $\beta$  p)  $\wedge$  pgte pwrite (pmult  $\beta$  p)
      by (simp add:  $\langle$ pgte pwrite  $\alpha$   $\wedge$  pgte pwrite  $\beta$  $\rangle$   $\langle$ ppos p  $\wedge$  pgte pwrite p $\rangle$ 
  asm0 multiply-smaller-pwrite pmult-ppos)
    ultimately show ?thesis
      using scalable-def scaled-def by auto
    qed
    then have get-h (R a wb) = get-h wb  $\wedge$  get-m (R a wb) = comp-min-mask
  (get-m a) (get-m wb)
      using non-scalable-R-charact by blast
    moreover have R a wb  $|\#|$  a
  proof (rule compatibleI)
    have larger-heap (get-h w) (get-h xb)  $\wedge$  larger-heap (get-h xb) (get-h wb)
      using  $\langle$ Some w = xa  $\oplus$  xb $\rangle$   $\langle$ pgte pwrite  $\alpha$   $\wedge$  pgte pwrite  $\beta$  $\rangle$   $\langle$ xb  $\succeq$  multiply

```

β wb \langle get-h-multiply larger-heap-def larger-implies-larger-heap plus-charact(2) \rangle **by**
fastforce
then show *compatible-heaps* (get-h (R a wb)) (get-h a)
by (metis *asm1 calculation(1) calculation(6) larger-heap-comp option.simps(3) plus-ab-defined*)
show *valid-mask* (add-masks (get-m (R a wb)) (get-m a))
by (metis *PartialSA.unit-neutral calculation(6) minus-empty option.distinct(1) plus-ab-defined unit-charact(2) valid-mask-add-comp-min*)
qed
then obtain bb **where** *Some* $bb = R a wb \oplus a$
using *PartialSA.defined-def* **by** *auto*

moreover obtain ya **where** *Some* $ya = R a wa \oplus a$
using *calculation(5)* **by** *auto*
then have $ya \in B$
using *asm1 r(1)* **by** *blast*
then have *multiply* $\alpha ya \in$ *multiply-sem-assertion* αB
using *in-multiply-refl* **by** *blast*
moreover obtain yb **where** *Some* $yb = R a wb \oplus a$
using *calculation(7)* **by** *auto*
then have $yb \in B$
using *asm1 r(2)* **by** *blast*
then have *multiply* $\beta yb \in$ *multiply-sem-assertion* βB
using *in-multiply-refl* **by** *blast*
moreover have (*multiply* αya) $|\#|$ (*multiply* βyb)
proof (*rule compatibleI*)
have *get-h* $ya =$ *get-h* $wa ++$ *get-h* a
using \langle *Some* $ya = R a wa \oplus a \rangle$ \langle *get-h* (R a wa) = *get-h* $wa \wedge$ *get-m* (R a wa) = *comp-min-mask* (*get-m* a) (*get-m* wa) \rangle *plus-charact(2)* **by** *presburger*
then have *get-h* (*multiply* αya) = *get-h* $wa ++$ *get-h* a
using \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* $\beta \rangle$ *get-h-multiply* **by** *presburger*
moreover have *get-h* $yb =$ *get-h* $wb ++$ *get-h* a
using \langle *Some* $yb = R a wb \oplus a \rangle$ \langle *get-h* (R a wb) = *get-h* $wb \wedge$ *get-m* (R a wb) = *comp-min-mask* (*get-m* a) (*get-m* wb) \rangle *plus-charact(2)* **by** *presburger*
then have *get-h* (*multiply* βyb) = *get-h* $wb ++$ *get-h* a
using \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* $\beta \rangle$ *get-h-multiply* **by** *presburger*
moreover have *compatible-heaps* (*get-h* wa) (*get-h* wb)
proof (*rule compatible-heapsI*)
fix $hl a b$ **assume** *get-h* $wa hl =$ *Some* a *get-h* $wb hl =$ *Some* b
then have *get-h* $xa hl =$ *Some* a *get-h* $xb hl =$ *Some* b
apply (metis (*full-types*) \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* $\beta \rangle$ \langle $xa \succeq$ *multiply* $\alpha wa \rangle$ *get-h-multiply larger-heap-def larger-implies-larger-heap*)
by (metis \langle *get-h* $wb hl =$ *Some* $b \rangle$ \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* $\beta \rangle$ \langle $xb \succeq$ *multiply* $\beta wb \rangle$ *get-h-multiply larger-heap-def larger-implies-larger-heap*)
moreover have *compatible-heaps* (*get-h* xa) (*get-h* xb)
by (metis \langle *Some* $w = xa \oplus xb \rangle$ *option.simps(3) plus-ab-defined*)
ultimately show $a = b$
by (metis *compatible-heaps-def compatible-options.simps(1)*)
qed

ultimately show *compatible-heaps* (*get-h* (*multiply* α *ya*)) (*get-h* (*multiply* β *yb*))
by (*metis* *PartialSA.commutative* *PartialSA.core-is-smaller* \langle *Some* $ya = R$ a $wa \oplus a$ \rangle \langle *Some* $yb = R$ a $wb \oplus a$ \rangle \langle *get-h* (R a wa) = *get-h* $wa \wedge$ *get-m* (R a wa) = *comp-min-mask* (*get-m* a) (*get-m* wa) \rangle \langle *get-h* (R a wb) = *get-h* $wb \wedge$ *get-m* (R a wb) = *comp-min-mask* (*get-m* a) (*get-m* wb) \rangle *compatible-heaps-sum* *core-defined*(1) *core-defined*(2) *option.distinct*(1) *plus-ab-defined*)
show *valid-mask* (*add-masks* (*get-m* (*multiply* α *ya*)) (*get-m* (*multiply* β *yb*)))
proof (*rule* *valid-maskI*)
show $\bigwedge f. \text{add-masks } (get-m \text{ (multiply } \alpha \text{ ya)}) \text{ (get-m (multiply } \beta \text{ yb)) (null, f) = pnone}$
by (*metis* (*no-types, opaque-lifting*) *PartialSA.core-is-smaller* *add-masks.simps* *core-defined*(2) *minus-empty not-None-eq* *plus-ab-defined* *valid-mask.simps*)
fix *hl*
have *add-masks* (*get-m* (*multiply* α *ya*)) (*get-m* (*multiply* β *yb*)) *hl* = *padd* (*pmult* α (*get-m* *ya* *hl*)) (*pmult* β (*get-m* *yb* *hl*))
using \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* β \rangle *get-m-smaller* **by** *auto*
moreover have *get-m* *ya* *hl* = *padd* (*get-m* (R a wa) *hl*) (*get-m* a *hl*) \wedge *get-m* *yb* *hl* = *padd* (*get-m* (R a wb) *hl*) (*get-m* a *hl*)
using \langle *Some* $ya = R$ a $wa \oplus a$ \rangle \langle *Some* $yb = R$ a $wb \oplus a$ \rangle *plus-charact*(1)
by *auto*
ultimately show *pgte* *pwrite* (*add-masks* (*get-m* (*multiply* α *ya*)) (*get-m* (*multiply* β *yb*)) *hl*)
by (*metis* *PartialSA.unit-neutral* *asm0* *option.distinct*(1) *padd-one-ineq-sum* *plus-ab-defined* *plus-charact*(1) *valid-mask.simps*)
qed
qed
then obtain *y* **where** *Some* $y = \text{multiply } \alpha \text{ ya} \oplus \text{multiply } \beta \text{ yb}$
using *PartialSA.defined-def* **by** *auto*
moreover have $x \succeq y$
proof (*rule* *greaterI*)
have *get-h* *y* = *get-h* *ya* ++ *get-h* *yb*
using \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* β \rangle *calculation*(10) *get-h-multiply* *plus-charact*(2) **by** *presburger*
moreover have *get-h* *ya* = *get-h* *wa* ++ *get-h* *a*
using \langle *Some* $ya = R$ a $wa \oplus a$ \rangle \langle *get-h* (R a wa) = *get-h* $wa \wedge$ *get-m* (R a wa) = *comp-min-mask* (*get-m* a) (*get-m* wa) \rangle *plus-charact*(2) **by** *presburger*
moreover have *get-h* *yb* = *get-h* *wb* ++ *get-h* *a*
using \langle *Some* $yb = R$ a $wb \oplus a$ \rangle \langle *get-h* (R a wb) = *get-h* $wb \wedge$ *get-m* (R a wb) = *comp-min-mask* (*get-m* a) (*get-m* wb) \rangle *plus-charact*(2) **by** *presburger*
moreover have *larger-heap* (*get-h* *x*) (*get-h* *wa*)
proof –
have *larger-heap* (*get-h* *x*) (*get-h* *xa*)
by (*metis* *PartialSA.greater-def* \langle *Some* $w = xa \oplus xb$ \rangle \langle *get-h* (R a w) = *get-h* $w \wedge$ *get-m* (R a w) = *comp-min-mask* (*get-m* a) (*get-m* w) \rangle *asm1* *larger-heap-trans* *larger-implies-larger-heap*)
moreover have *larger-heap* (*get-h* *xa*) (*get-h* *wa*)
by (*metis* \langle *pgte* *pwrite* $\alpha \wedge$ *pgte* *pwrite* β \rangle \langle *xa* \succeq *multiply* α *wa* \rangle)

get-h-multiply larger-implies-larger-heap
ultimately show *?thesis*
using *larger-heap-trans* **by** *blast*
qed
moreover have *larger-heap (get-h x) (get-h wb)*
proof –
have *larger-heap (get-h x) (get-h xb)*
by (*metis PartialSA.greater-def PartialSA.greater-equiv* $\langle \text{Some } w = xa \oplus xb \rangle$ *get-h (R a w) = get-h w \wedge get-m (R a w) = comp-min-mask (get-m a) (get-m w)* *asm1 larger-heap-trans larger-implies-larger-heap*)
moreover have *larger-heap (get-h xb) (get-h wb)*
by (*metis* $\langle \text{pgte pwrite } \alpha \wedge \text{pgte pwrite } \beta \rangle$ $\langle \text{xb} \succeq \text{multiply } \beta \text{ wb} \rangle$)
get-h-multiply larger-implies-larger-heap
ultimately show *?thesis*
using *larger-heap-trans* **by** *blast*
qed
moreover have *larger-heap (get-h x) (get-h a)*
using *PartialSA.greater-equiv asm1 larger-implies-larger-heap* **by** *blast*
ultimately show *larger-heap (get-h x) (get-h y)*
by (*simp add: larger-heap-plus*)
show *greater-mask (get-m x) (get-m y)*
proof (*rule greater-maskI*)
fix *hl*
have *get-m x hl = padd (get-m (R a w) hl) (get-m a hl)*
using *asm1 plus-charact(1)* **by** *auto*
moreover have *get-m y hl = padd (pmult α (padd (get-m (R a wa) hl) (get-m a hl))) (pmult β (padd (get-m (R a wb) hl) (get-m a hl)))*
by (*metis* $\langle \text{Some } y = \text{multiply } \alpha \text{ ya} \oplus \text{multiply } \beta \text{ yb} \rangle$ $\langle \text{Some } \text{ya} = \text{R a wa} \oplus \text{a} \rangle$ $\langle \text{Some } \text{yb} = \text{R a wb} \oplus \text{a} \rangle$ *pgte pwrite $\alpha \wedge$ pgte pwrite β* *add-masks.simps get-m-smaller plus-charact(1)*)

moreover have *padd (pmult α (padd (get-m (R a wa) hl) (get-m a hl))) (pmult β (padd (get-m (R a wb) hl) (get-m a hl))) = padd (padd (pmult α (get-m a hl)) (pmult β (get-m a hl))) (padd (pmult α (get-m (R a wa) hl)) (pmult β (get-m (R a wb) hl)))*
using *padd-asso padd-comm pmult-distr* **by** *force*

have *pgte (get-m (R a w) hl) (padd (pmult α (get-m (R a wa) hl)) (pmult β (get-m (R a wb) hl)))*
proof (*cases pgte (get-m w hl) (comp-one (get-m a hl))*)
case *True*
then have *get-m (R a w) hl = (comp-one (get-m a hl))*
using $\langle \text{get-h (R a w)} = \text{get-h w} \wedge \text{get-m (R a w)} = \text{comp-min-mask (get-m a) (get-m w)} \rangle$ *comp-min-mask-def pmin-is* **by** *presburger*
moreover have *pgte (comp-one (get-m a hl)) (get-m (R a wa) hl)*
by (*metis* $\langle \text{get-h (R a wa)} = \text{get-h wa} \wedge \text{get-m (R a wa)} = \text{comp-min-mask (get-m a) (get-m wa)} \rangle$ *comp-min-mask-def pmin-comm pmin-greater*)
then have *pgte (pmult α (comp-one (get-m a hl))) (pmult α (get-m (R a wa) hl))*

by (metis pmult-comm pmult-order)
 moreover have pgte (comp-one (get-m a hl)) (get-m (R a wb) hl)
 by (metis ⟨get-h (R a wb) = get-h wb ∧ get-m (R a wb) = comp-min-mask
 (get-m a) (get-m wb)⟩ comp-min-mask-def pmin-comm pmin-greater)
 then have pgte (pmult β (comp-one (get-m a hl))) (pmult β (get-m (R
 a wb) hl))
 by (metis pmult-comm pmult-order)
 ultimately show ?thesis
 using ⟨pgte (comp-one (get-m a hl)) (get-m (R a wa) hl)⟩ ⟨pgte (comp-one
 (get-m a hl)) (get-m (R a wb) hl)⟩ asm0 padd-one-ineq-sum by presburger
 next
 case False
 then have get-m (R a w) hl = get-m w hl
 by (metis ⟨get-h (R a w) = get-h w ∧ get-m (R a w) = comp-min-mask
 (get-m a) (get-m w)⟩ comp-min-mask-def not-pgte-charact pgt-implies-pgte pmin-comm
 pmin-is)
 moreover have pgte (get-m w hl) (padd (pmult α (get-m wa hl)) (pmult
 β (get-m wb hl)))
 proof –
 have pgte (get-m w hl) (padd (get-m xa hl) (get-m xb hl))
 using ⟨Some w = xa ⊕ xb⟩ not-pgte-charact pgt-implies-pgte
 plus-charact(1) by auto
 moreover have pgte (get-m xa hl) (pmult α (get-m wa hl))
 by (metis ⟨pgte pwrite α ∧ pgte pwrite β⟩ ⟨xa ≥ multiply α wa⟩
 get-m-smaller larger-implies-greater-mask-hl)
 moreover have pgte (get-m xb hl) (pmult β (get-m wb hl))
 by (metis ⟨pgte pwrite α ∧ pgte pwrite β⟩ ⟨xb ≥ multiply β wb⟩
 get-m-smaller larger-implies-greater-mask-hl)
 ultimately show ?thesis
 by (simp add: padd.rep-eq pgte.rep-eq)
 qed
 moreover have pgte (pmult α (get-m wa hl)) (pmult α (get-m (R a wa)
 hl))
 by (metis R-smaller larger-implies-greater-mask-hl pmult-comm
 pmult-order)
 moreover have pgte (pmult β (get-m wb hl)) (pmult β (get-m (R a wb)
 hl))
 by (metis R-smaller larger-implies-greater-mask-hl pmult-comm
 pmult-order)
 ultimately show ?thesis
 using padd.rep-eq pgte.rep-eq by force
 qed
 moreover have get-m x hl = padd (get-m (R a w) hl) (get-m a hl)
 using calculation(1) by auto
 moreover have get-m y hl = padd (pmult α (get-m ya hl)) (pmult β
 (get-m yb hl))
 using ⟨Some y = multiply α ya ⊕ multiply β yb⟩ ⟨pgte pwrite α ∧ pgte
 pwrite β⟩ get-m-smaller plus-charact(1) by auto
 moreover have padd (pmult α (get-m ya hl)) (pmult β (get-m yb hl)) =

```

padd (pmult  $\alpha$  (padd (get-m (R a wa) hl) (get-m a hl))) (pmult  $\beta$  (padd (get-m (R
a wb) hl) (get-m a hl)))
  using calculation(2) calculation(5) by presburger
  moreover have ... = padd (pmult (padd  $\alpha$   $\beta$ ) (get-m a hl)) (padd (pmult
 $\alpha$  (get-m (R a wa) hl) (pmult  $\beta$  (get-m (R a wb) hl)))
    by (metis <padd (pmult  $\alpha$  (padd (get-m (R a wa) hl) (get-m a hl))) (pmult
 $\beta$  (padd (get-m (R a wb) hl) (get-m a hl))) = padd (padd (pmult  $\alpha$  (get-m a hl))
(pmult  $\beta$  (get-m a hl))) (padd (pmult  $\alpha$  (get-m (R a wa) hl) (pmult  $\beta$  (get-m (R
a wb) hl)))> pmult-comm pmult-distr)
    ultimately show pgte (get-m x hl) (get-m y hl)
      using asm0 p-greater-exists padd-asso padd-comm pmult-special(1) by
force
  qed
  qed
  ultimately have  $y \in \text{multiply-sem-assertion } \alpha B \otimes \text{multiply-sem-assertion}$ 
 $\beta B$ 
    using PartialSA.add-set-elem by blast
    then have  $y \in \text{multiply-sem-assertion } \text{pwrite } B$ 
      by (metis asm0 assms(1) combinable-def not-pgte-charact pgt-implies-pgte
subsetD)
    then obtain  $b$  where  $y \succeq \text{multiply } \text{pwrite } b$   $b \in B$ 
      using in-multiply-sem by blast
    then have  $\text{multiply } \text{pwrite } b = b$ 
      by (metis Rep-state-inverse get-h-m mult-write-mask multiply.simps)
    then have  $y \in B$ 
      by (metis < $b \in B$ > < $y \succeq \text{multiply } \text{pwrite } b$ > assms(2) intuitionistic-def)
    show  $x \in B$ 
      using < $x \succeq y$ > < $y \in B$ > assms(2) intuitionistic-def by blast
  qed
  qed
  qed

```

2.6 Theorems

The following theorem is crucial to use the package logic [4] to automatically compute footprints of combinable wands.

theorem *R-mono-transformer*:

PartialSA.mono-transformer (R a)

proof –

have $R a \text{ unit} = \text{unit}$

by (simp add: PartialSA.succ-antisym PartialSA.unit-smaller R-smaller)

moreover have $\bigwedge \varphi \varphi'. \varphi' \succeq \varphi \implies R a \varphi' \succeq R a \varphi$

proof –

fix $\varphi \varphi'$

assume $\varphi' \succeq \varphi$

show $R a \varphi' \succeq R a \varphi$

proof (cases scalable $\varphi' a$)

case True

then show ?thesis

```

    by (metis PartialSA.succ-trans R-def R-smaller <math>\varphi' \succeq \varphi</math>)
next
case False
then obtain p where ppos p pgte pwrite p multiply p <math>\varphi' \#| a</math>
by (metis PartialSA.commutative PartialSA.defined-def non-scalable-instantiate)
then have multiply p <math>\varphi \#| a</math>
    using PartialSA.smaller-compatible <math>\varphi' \succeq \varphi</math> multiply-order by blast
then have  $\neg$  scalable  $\varphi$  a
    using PartialSA.commutative PartialSA.defined-def <math>\langle \text{pgte } pwrite \ p \rangle \langle \text{ppos } p \rangle</math>
scalable-def scaled-def by auto

    moreover have greater-mask (comp-min-mask (get-m a) (get-m  $\varphi'$ )) (comp-min-mask
(get-m a) (get-m  $\varphi$ ))
    proof (rule greater-maskI)
    fix hl show pgte (comp-min-mask (get-m a) (get-m  $\varphi'$ ) hl) (comp-min-mask
(get-m a) (get-m  $\varphi$ ) hl)
    proof (cases pgte (get-m  $\varphi'$ ) hl) (comp-one (get-m a hl))
    case True
    then show ?thesis
    by (metis comp-min-mask-def pmin-comm pmin-greater pmin-is)
    next
    case False
    then show ?thesis
    by (metis PartialSA.succ-trans R-smaller <math>\varphi' \succeq \varphi</math> calculation comp-min-mask-def
larger-implies-greater-mask-hl non-scalable-R-character not-pgte-character pgt-implies-pgte
pmin-comm pmin-is)
    qed
    qed
    ultimately show ?thesis
    using False <math>\varphi' \succeq \varphi</math> greaterI larger-implies-larger-heap non-scalable-R-character
by presburger
    qed
    qed
    ultimately show ?thesis
    by (simp add: PartialSA.mono-transformer-def)
qed

```

```

theorem properties-of-combinable-wands:
  assumes intuitionistic B
  shows combinable B  $\implies$  combinable (cwand A B)
  and cwand A B  $\subseteq$  wand A B
  and binary A  $\implies$  cwand A B = wand A B
  by (simp-all add: assms combinable-cwand cwand-stronger binary-same dual-order.eq-iff)

```

end

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