# Intersecting Chords Theorem

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#### Abstract

This entry provides a geometric proof of the intersecting chords theorem. The theorem states that when two chords intersect each other inside a circle, the products of their segments are equal.

After a short review of existing proofs in the literature [1, 3-5], I decided to use a proof approach that employs reasoning about lengths of line segments, the orthogonality of two lines and Pythagoras Law. Hence, one can understand the formalized proof easily with the knowledge of a few general geometric facts that are commonly taught in high-school.

This theorem is the 55th theorem of the Top 100 Theorems list.

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## 1 Introduction

The intersecting chords theorem states:

When two chords intersect each other inside a circle, the products of their segments are equal.

To prove this theorem in Isabelle, I reviewed existing formalizations in theorem provers and proofs in the literature [1, 3-5]. At the time of this AFP submission, the formalization of geometry in Isabelle is limited to only a few concepts and theorems. Hence, I selected to formalize the proof approach that fitted best to the already existing geometry formalizations.

The proof in HOL Light [3] simply unfolds the involved geometric predicates and then proves the theorem using only algebraic computations on

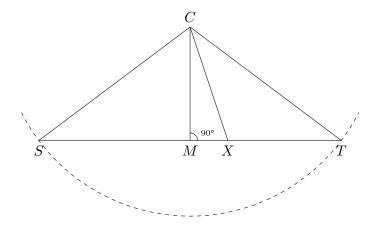


Figure 1: Key Lemma states  $\overline{SX} \cdot \overline{XT} = \overline{SC}^2 - \overline{XC}^2$ 

real numbers. By a quick and shallow inspection of the proof script without executing the proof script step by step in HOL Light, I could not understand the proof script well enough to re-write the proof in Isabelle. As running the script in HOL Light seemed too involved to me, I ignored HOL Light's proof approach and considered the other approaches in the literature.

The first proof approach [5] that I found in the literature employs similarity of triangles, the inscribed angle theorem, and basic reasoning with angles. The intersecting chords theorem only consists of two reasoning steps after stating the geometric observations about angles. However, the proof requires to formalize the concept of similarity of triangles, extend the existing formalization of angles, and prove the inscribed angle theorem. So, I abandoned this proof approach and considered the second proof approach.

The second proof approach [5] needs only basic geometric reasoning about lengths of line segments, the orthogonality of two lines and Pythagoras Law. More specifically, one must prove that the line that goes through the apex and the midpoint of the base in an isosceles triangle is orthogonal to the base. This is easily derived from the property of an isosceles triangle using the congruence properties of triangles, which is already formalized in AFP's Triangle entry [2]. Furthermore, Pythagoras Law is a special case of the Law of Cosines, which is already formalized in AFP's Triangle entry.

Ultimately, I decided to use this second proof approach, which I sketch in more detail in the next subsection.

### 1.1 Informal Proof Sketch

The proof of the intersecting chords theorem relies on the following observation which is depicted in Figure 1.

Instead of considering two arbitrary chords intersecting, consider one arbitrary chord with endpoints S and T on a circle with center C and one arbitrary point X on the chord ST. This point X on the chord creates two line segments on this chord, the left part SX, and the right part XT. Without loss of generality, we can assume that SX is longer that XT, as shown in Figure 1.

The key lemma for the intersecting chords theorem provides a closed expression for the length of these two line segment using the distances of the chord endpoints and the point to the center C, i.e., the lemma states:

$$\overline{SX} \cdot \overline{XT} = \overline{SC}^2 - \overline{XC}^2.$$

To prove this fact, we consider the midpoint M of the chord ST. First, as M is the midpoint,  $\overline{SM}$  and  $\overline{TM}$  are equal. Second, we observe that the lengths of the line segments SX and XT are:

$$\overline{SX} = \overline{SM} + \overline{MX}$$
 and  $\overline{XT} = \overline{TM} - \overline{MX} = \overline{SM} - \overline{MX}$ .

Third, the Pythagoras law for the triangles SMC and XMC states:

$$\overline{SM}^2 + \overline{MC}^2 = \overline{SC}^2$$
 and  $\overline{XM}^2 + \overline{MC}^2 = \overline{XC}^2$ .

Finally, the product can be expressed as:

$$\overline{SX} \cdot \overline{XT} = (\overline{SM} + \overline{MX}) \cdot (\overline{TM} - \overline{MX}) = \overline{SM}^2 - \overline{MX}^2 = (\overline{SC}^2 - \overline{MC}^2) - (\overline{XC}^2 - \overline{MC}^2) = \overline{SC}^2 - \overline{XC}^2.$$

The intersecting chord theorem now follows directly from this lemma: as the distances SC and XC for two arbitrary chords intersecting at X are equal, also the products of the chord segments are equal.

### 2 Intersecting Chord Theorem

theory Chord-Segments imports Triangle. Triangle begin

#### 2.1 Preliminaries

**lemma** between E-if-dist-leq: **fixes**  $A \ B \ X :: 'a::euclidean-space$  **assumes** between  $(A, B) \ X$  **assumes** dist  $A \ X \le dist \ B \ X$  **obtains** u where  $1 \ / \ 2 \le u \ u \le 1$  and  $X = u \ *_R \ A + (1 - u) \ *_R \ B$  **proof** (cases A = B) **assume**  $A \ne B$ **from**  $\langle between \ (A, B) \ X \rangle$  **obtain** u where  $u: u \ge 0 \ u \le 1$  and  $X: X = u \ *_R \ A + (1 - u) \ *_R \ B$ 

**by** (*metis add.commute betweenE between-commute*) from X have  $X = B + u *_R (A - B)$  and  $X = A + (u - 1) *_R (A - B)$ **by** (*simp add: scaleR-diff-left real-vector.scale-right-diff-distrib*)+ from  $\langle X = B + u *_B (A - B) \rangle$  have dist-B: dist  $B X = norm (u *_B (A - B))$ **by** (*auto simp add: dist-norm*) from  $\langle X = A + (u - 1) *_R (A - B) \rangle$  have dist-A: dist A X = norm ((u - 1)) $*_R (A - B))$ by (*auto simp add*: *dist-norm*) from  $\langle A \neq B \rangle$  have norm (A - B) > 0 by auto from this (dist  $A X \leq dist B X$ ) have  $u \geq 1 / 2$ using dist-A dist-B by simp from this  $\langle u \leq 1 \rangle X$  that show thesis by blast next assume A = Bdefine u :: real where u = 1from (between (A, B) X) (A = B) have  $1 / 2 \le u u \le 1 X = u *_R A + (1 - u)$  $u) *_B B$ unfolding *u*-def by auto with that show thesis by blast qed **lemma** *dist-geq-iff-midpoint-in-between*: fixes A B X :: 'a::euclidean-space **assumes** between (A, B) X**shows** dist  $A \ X \leq dist \ B \ X \leftrightarrow between \ (X, \ B) \ (midpoint \ A \ B)$ proof **assume** dist  $A X \leq dist B X$ **from**  $\langle between (A, B) X \rangle$  this obtain u where  $u: 1 / 2 \le u u \le 1$  and  $X: X = u *_R A + (1 - u) *_R B$  $\mathbf{using} \ between E{\text{-}if{\text{-}dist{-}leq}} \ \mathbf{by} \ blast$ have M: midpoint A  $B = (1 / 2) *_R A + (1 / 2) *_R B$ **unfolding** *midpoint-def* **by** (*simp add: scaleR-add-right*) from  $(1 / 2 \le u)$  have 1: midpoint  $A B = (1 / (2 * u)) *_R X + (1 - (1 / (2 * u)))$  $(* u))) *_R B$ proof have (2 - u \* 2) / (2 \* u) = 1 / u - u / uusing u(1) by (simp add: diff-divide-distrib) also have  $\ldots = 1 / u - 1$  using u(1) by *auto* finally have (2 - u \* 2) / (2 \* u) = 1 / u - 1. from  $\langle 1 | 2 \leq u \rangle$  this show ?thesis using X M by (simp add: scaleR-add-right scaleR-add-left[symmetric]) qed moreover from u have  $2: (1 / (2 * u)) \ge 0 (1 / (2 * u)) \le 1$  by auto ultimately show between (X, B) (midpoint A B) using between *I* [of concl: *B* X] by (metis add.commute between-commute)  $\mathbf{next}$ **assume** between (X, B) (midpoint A B) then have between  $(A, midpoint \ A \ B) \ X$ using  $\langle between (A, B) X \rangle$  between-midpoint(1) between-swap by blast

then have dist  $A \ X \leq dist \ A \ (midpoint \ A \ B)$ using between zero-le-dist by force also have dist  $A \ (midpoint \ A \ B) \leq dist \ B \ (midpoint \ A \ B)$ by (simp add: dist-midpoint) also from (between  $(X, B) \ (midpoint \ A \ B)$ ) have dist  $B \ (midpoint \ A \ B) \leq dist \ B \ X$ using between zero-le-dist by (metis add.commute dist-commute le-add-same-cancel1) finally show dist  $A \ X \leq dist \ B \ X$ .

#### qed

### 2.2 Properties of Chord Segments

**lemma** chord-property: fixes S C :: 'a :: euclidean-spaceassumes dist C S = dist C Tassumes between (S, T) Xshows dist  $S X * dist X T = (dist C S)^2 - (dist C X)^2$ proof define M where M = midpoint S Thave between (S, T) M **unfolding** *M*-def by  $(simp \ add: between-midpoint(1))$ have dist T M = dist S M**unfolding** *M*-def **by** (simp add: dist-midpoint) have distances: max (dist S X) (dist X T) = (dist S M) + (dist X M)  $\land$ min (dist S X) (dist X T) = (dist S M) - (dist X M)proof cases assume dist  $S X \leq dist X T$ then have between (X, T) M using  $\langle between (S, T) X \rangle$  M-def **by** (*simp add: dist-geq-iff-midpoint-in-between dist-commute*) then have between (S, M) X**using** (between (S, T) X) (between (S, T) M) between-swap by blast from  $\langle between (X, T) M \rangle$  have dist X T = dist X M + dist M Tusing between by auto **moreover from** (between (S, M) X) have dist S X = dist S M - dist M Xusing between dist-commute by force ultimately show *?thesis* using  $\langle dist \ S \ X \leq dist \ X \ T \rangle \langle dist \ T \ M = dist \ S \ M \rangle$ by (simp add: add.commute dist-commute max-def min-def)  $\mathbf{next}$ assume  $\neg$  (dist  $S X \leq dist X T$ ) then have dist  $T X \leq dist S X$  by (simp add: dist-commute) then have between (S, X) M using  $\langle between (S, T) X \rangle$  M-def **by** (*simp add: dist-geq-iff-midpoint-in-between midpoint-sym between-commute*) then have between (T, M) Xusing  $\langle between (S, T) X \rangle \langle between (S, T) M \rangle$  between-swap between-commute by *metis* 

from (between (S, X) M) have dist S X = dist S M + dist M Xusing between by auto **moreover from** (between (T, M) X) have dist T X = dist T M - dist M Xusing between dist-commute by force ultimately show ?thesis using  $\langle \neg dist \ S \ X \leq dist \ X \ T \rangle \langle dist \ T \ M = dist \ S \ M \rangle$ by (metis dist-commute max-def min-def) qed have orthogonal (C - M) (S - M)using  $\langle dist \ C \ S = dist \ C \ T \rangle$  M-def **by** (*auto simp add: isosceles-triangle-orthogonal-on-midpoint*) have orthogonal (C - M) (X - M)proof have between (S, T) M using *M*-def between-midpoint(1) by blast obtain c where  $(X - M) = c *_R (S - M)$ **proof** (cases S = M) assume  $S \neq M$ then obtain c where  $(X - M) = c *_R (S - M)$ using between-implies-scaled-diff [OF  $\langle between (S, T) X \rangle \langle between (S, T) \rangle$ M ] by metis from this that show thesis by blast  $\mathbf{next}$ assume S = Mfrom this (between (S, T) X) have X = M**by** (simp add: midpoint-between M-def) from  $\langle X = M \rangle \langle S = M \rangle$  have  $(X - M) = \theta *_R (S - M)$  by simp from this that show thesis by blast  $\mathbf{qed}$ from this (orthogonal (C - M) (S - M)) show ?thesis **by** (*auto intro: orthogonal-clauses*(2))  $\mathbf{qed}$ **from** (orthogonal (C - M) (S - M)) (orthogonal (C - M) (X - M)) have  $(dist S M) \ \widehat{2} + (dist M C) \ \widehat{2} = (dist C S) \ \widehat{2}$  $(dist X M) \hat{2} + (dist M C) \hat{2} = (dist C X) \hat{2}$ by (auto simp only: Pythagoras) then have geometric-observation:  $\begin{array}{c} (\operatorname{dist} S \ M) & \widehat{\phantom{a}} 2 = (\operatorname{dist} C \ S) & \widehat{\phantom{a}} 2 - (\operatorname{dist} M \ C) & \widehat{\phantom{a}} 2 \\ (\operatorname{dist} X \ M) & \widehat{\phantom{a}} 2 = (\operatorname{dist} C \ X) & \widehat{\phantom{a}} 2 - (\operatorname{dist} M \ C) & \widehat{\phantom{a}} 2 \end{array}$ by auto have dist S X \* dist X T = max (dist S X) (dist X T) \* min (dist S X) (dist X T)T)**by** (*auto split: split-max*) also have  $\ldots = ((dist \ S \ M) + (dist \ X \ M)) * ((dist \ S \ M) - (dist \ X \ M))$ using distances by simp also have  $\ldots = (dist \ S \ M) \ \widehat{} \ 2 - (dist \ X \ M) \ \widehat{} \ 2$ **by** (*simp add: field-simps power2-eq-square*)

also have  $\ldots = ((dist \ C \ S) \ ^2 - (dist \ M \ C) \ ^2) - ((dist \ C \ X) \ ^2 - (dist \ M \ C) \ ^2)$ using geometric-observation by simp also have  $\ldots = (dist \ C \ S) \ ^2 - (dist \ C \ X) \ ^2$  by simp finally show ?thesis . ged

theorem product-of-chord-segments: fixes  $S_1 \ T_1 \ S_2 \ T_2 \ X \ C$  :: 'a :: euclidean-space assumes between  $(S_1, \ T_1) \ X$  between  $(S_2, \ T_2) \ X$ assumes dist  $C \ S_1 = r \ dist \ C \ T_1 = r$ assumes dist  $C \ S_2 = r \ dist \ C \ T_2 = r$ shows dist  $S_1 \ X \ dist \ X \ T_1 = \ dist \ S_2 \ X \ dist \ X \ T_2$ proof – from  $\langle dist \ C \ S_1 = r \rangle \langle dist \ C \ T_1 = r \ \rangle \langle between \ (S_1, \ T_1) \ X \rangle$ have  $dist \ S_1 \ X \ dist \ X \ T_1 = r \ 2 - (\ dist \ C \ X) \ 2$ by (subst chord-property) auto also from  $\langle dist \ C \ S_2 = r \rangle \langle dist \ C \ T_2 = r \rangle \langle between \ (S_2, \ T_2) \ X \rangle$ have ... = dist  $S_2 \ X \ dist \ X \ T_2$ by (subst chord-property) auto finally show ?thesis . qed

end

### References

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