Chebyshev Polynomials

Manuel Eberl

November 21, 2023

Abstract

The multiple-angle formulas for \cos and \sin state that for any natural number n, the values of $\cos nx$ and $\sin nx$ can be expressed in terms of $\cos x$ and $\sin x$. To be more precise, there are polynomials T_n and U_n such that $\cos nx = T_n(\cos x)$ and $\sin nx = U_n(\cos x)\sin x$. These are called the *Chebyshev polynomials of the first and second kind*, respectively.

This entry contains a definition of these two familes of polynomials in Isabelle/HOL along with some of their most important properties. In particular, it is shown that T_n and U_n are *orthogonal* families of polynomials.

Moreover, we show the well-known result that for any monic polynomial p of degree n > 0, it holds that $\sup_{x \in [-1,1]} |p(x)| \ge 2^{n-1}$, and that this inequality is sharp since equality holds with $p = 2^{1-n}T_n$. This has important consequences in the theory of function interpolation, since it implies that the roots of T_n (also colled the *Chebyshev nodes*) are exceptionally well-suited as interpolation nodes.

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1 Parametricity of polynomial operations

```
theory Polynomial_Transfer
  imports "HOL-Computational_Algebra.Polynomial"
begin
definition rel_poly :: "('a \Rightarrow 'b \Rightarrow bool) \Rightarrow 'a :: zero poly \Rightarrow 'b ::
zero poly \Rightarrow bool" where
  "rel_poly R p q \longleftrightarrow rel_fun (=) R (coeff p) (coeff q)"
lemma\ left\_unique\_rel\_poly\ [transfer\_rule]:\ "left\_unique\ R \implies left\_unique
(rel_poly R)"
  \langle proof \rangle
lemma right_unique_rel_poly [transfer_rule]: "right_unique R \Longrightarrow right_unique
(rel_poly R)"
  \langle proof \rangle
lemma bi_unique_rel_poly [transfer_rule]: "bi_unique R \implies bi_unique
(rel_poly R)"
  \langle proof \rangle
lemma rel_poly_swap: "rel_poly R x y \longleftrightarrow rel_poly (\lambday x. R x y) y x"
  \langle proof \rangle
lemma coeff_transfer [transfer_rule]:
  "rel_fun (rel_poly R) (rel_fun (=) R) coeff coeff"
  \langle proof \rangle
lemma map_poly_transfer:
  assumes "rel_fun R S f g" "f 0 = 0" "g 0 = 0"
           "rel_fun (rel_poly R) (rel_poly S) (map_poly f) (map_poly g)"
  \langle proof \rangle
lemma map_poly_transfer':
  assumes "rel fun R S f g" "rel poly R p q" "f 0 = 0" "g 0 = 0"
           "rel_poly S (map_poly f p) (map_poly g q)"
  shows
  \langle proof \rangle
lemma rel_poly_id: "p = q \implies rel_poly (=) p q"
  \langle proof \rangle
lemma left_total_rel_poly [transfer_rule]:
  assumes "left_total R" "right_unique R" "R 0 0"
  shows
            "left_total (rel_poly R)"
  \langle proof \rangle
```

```
lemma right_total_rel_poly [transfer_rule]:
  assumes "right_total R" "left_unique R" "R 0 0"
  shows
            "right_total (rel_poly R)"
  \langle proof \rangle
lemma bi_total_rel_poly [transfer_rule]:
  assumes "bi_total R" "bi_unique R" "R 0 0"
            "bi_total (rel_poly R)"
  shows
  \langle proof \rangle
lemma \ zero\_poly\_transfer \ [transfer\_rule] \colon "R \ 0 \ 0 \Longrightarrow rel\_poly \ R \ 0 \ 0"
  \langle proof \rangle
lemma one_poly_transfer [transfer_rule]: "R 0 0 ⇒ R 1 1 ⇒ rel_poly
R 1 1"
  \langle proof \rangle
lemma pCons_transfer [transfer_rule]:
  "rel_fun R (rel_fun (rel_poly R) (rel_poly R)) pCons pCons"
  \langle proof \rangle
lemma plus_poly_transfer [transfer_rule]:
  "rel_fun R (rel_fun R R) (+) (+) \Longrightarrow
   rel_fun (rel_poly R) (rel_fun (rel_poly R) (rel_poly R)) (+) (+)"
  \langle proof \rangle
lemma minus_poly_transfer [transfer_rule]:
  "rel_fun R (rel_fun R R) (-) (-) \Longrightarrow
   rel_fun (rel_poly R) (rel_fun (rel_poly R) (rel_poly R)) (-) (-)"
  \langle proof \rangle
lemma uminus_poly_transfer [transfer_rule]:
  "rel_fun R R uminus uminus \Longrightarrow rel_fun (rel_poly R) (rel_poly R) uminus
uminus"
  \langle proof \rangle
lemma smult_transfer [transfer_rule]:
  "rel_fun R (rel_fun R R) (*) (*) \Longrightarrow
   rel_fun R (rel_fun (rel_poly R) (rel_poly R)) smult smult"
  \langle proof \rangle
lemma monom_transfer [transfer_rule]:
  "R 0 0 \Longrightarrow rel_fun R (rel_fun (=) (rel_poly R)) monom monom"
  \langle proof \rangle
lemma pderiv_transfer [transfer_rule]:
  assumes "R 0 0" "rel_fun R (rel_fun R R) (+) (+)"
  shows "rel_fun (rel_poly R) (rel_poly R) pderiv pderiv"
\langle proof \rangle
```

```
lemma If_transfer':
  assumes "P = P'" "P \Longrightarrow R \times x'" "\neg P \Longrightarrow R \times y""
          "R (if P then x else y) (if P' then x' else y')"
  \langle proof \rangle
lemma nth_transfer:
  assumes "list_all2 R xs ys" "i = j" "i < length xs"
           "R (xs ! i) (ys ! j)"
  shows
  \langle proof \rangle
lemma Poly_transfer [transfer_rule]:
  assumes [transfer_rule]: "R 0 0" "bi_unique R"
           "rel_fun (list_all2 R) (rel_poly R) Poly Poly"
  shows
  \langle proof \rangle
lemma poly_of_list_transfer [transfer_rule]:
  assumes [transfer_rule]: "R 0 0" "bi_unique R"
          "rel_fun (list_all2 R) (rel_poly R) poly_of_list poly_of_list"
  \mathbf{shows}
  \langle proof \rangle
lemma degree_transfer [transfer_rule]:
  assumes [transfer_rule]: "R 0 0" "bi_unique R"
  shows
           "rel_fun (rel_poly R) (=) degree degree"
\langle proof \rangle
lemma coeffs_transfer [transfer_rule]:
  assumes [transfer_rule]: "R 0 0" "bi_unique R"
  shows "rel_fun (rel_poly R) (list_all2 R) coeffs coeffs"
\langle proof \rangle
lemma times_poly_transfer [transfer_rule]:
  assumes [transfer_rule]: "rel_fun R (rel_fun R R) (+) (+)"
                              "rel_fun R (rel_fun R R) (*) (*) "R 0 0" "bi_unique
R"
  shows "rel_fun (rel_poly R) (rel_fun (rel_poly R) (rel_poly R)) (*)
(*)"
  \langle proof \rangle
lemma dvd_poly_transfer [transfer_rule]:
  assumes [transfer_rule]: "rel_fun R (rel_fun R R) (+) (+)"
                              "rel_fun R (rel_fun R R) (*) (*) "R 0 0" "bi_unique
R" "bi_total R"
  shows "rel_fun (rel_poly R) (rel_fun (rel_poly R) (=)) (dvd) (dvd)"
  \langle proof \rangle
lemma poly_transfer [transfer_rule]:
  assumes [transfer_rule]: "rel_fun R (rel_fun R R) (+) (+)"
```

```
"rel_fun R (rel_fun R R) (*) (*)" "R 0 0" "bi_unique
R"
  shows "rel_fun (rel_poly R) (rel_fun R R) poly poly"
  \langle proof \rangle
lemma pcompose_transfer [transfer_rule]:
  assumes [transfer_rule]: "rel_fun R (rel_fun R R) (+) (+)"
                             "rel_fun R (rel_fun R R) (*) (*) "R 0 0" "bi_unique
R."
  shows "rel_fun (rel_poly R) (rel_fun (rel_poly R) (rel_poly R)) pcompose
pcompose"
  \langle proof \rangle
lemma order_0_right: "order x 0 = Least (\lambda_. False)"
lemma order_poly_transfer [transfer_rule]:
  assumes [transfer_rule]:
    "rel_fun R (rel_fun R R) (+) (+)" "rel_fun R (rel_fun R R) (*) (*)"
    "rel_fun R R uminus uminus"
    "R 0 0" "R 1 1" "bi_unique R" "bi_total R" "R x y" "rel_poly R p q"
  shows "order x p = order y q"
  \langle proof \rangle
```

2 Missing Library Material

theory Chebyshev_Polynomials_Library
imports "HOL-Computational_Algebra.Polynomial" "HOL-Library.FuncSet"
begin

2.1 Miscellaneous

end

```
lemma bij_betw_Collect:
   assumes "bij_betw f A B" "\xspace x. x \in A \implies Q (f x) \xspace \longrightarrow P x"
   shows "bij_betw f {x\inA. P x} {y\inB. Q y}"
   \xspace \text{proof}\xspace \text{}

lemma induct_nat_012[case_names 0 1 ge2]:
   "P 0 \implies P (Suc 0) \implies (\xspace n) \implies P (Suc x) \implies P n"
   \xspace \text{proof}\xspace \rangle
```

2.2 Lists

lemma distinct_adj_conv_length_remdups_adj:

```
"distinct_adj xs \longleftrightarrow length (remdups_adj xs) = length xs" \langle proof \rangle

lemma successively_conv_nth:
    "successively P xs \longleftrightarrow (\forall i. Suc i < length xs \longrightarrow P (xs ! i) (xs ! Suc i))" \langle proof \rangle

lemma successively_nth: "successively P xs \Longrightarrow Suc i < length xs \Longrightarrow P (xs ! i) (xs ! Suc i)" \langle proof \rangle

lemma distinct_adj_conv_nth:
    "distinct_adj xs \longleftrightarrow (\forall i. Suc i < length xs \longrightarrow xs ! i \neq xs ! Suc i)" \langle proof \rangle

lemma distinct_adj_nth: "distinct_adj xs \Longrightarrow Suc i < length xs \Longrightarrow xs ! i \neq xs ! Suc i" \langle proof \rangle
```

The following two lemmas give a full characterisation of the *filter* function: The list *filter* P xs is the only list ys for which there exists a strictly increasing function $f: \{0, \ldots, |ys|-1\} \rightarrow \{0, \ldots, |xs|-1\}$ such that:

- $ys_i = xs_{f(i)}$
- $P(xs_i) \longleftrightarrow \exists j < n. \ f(j) = i$, i.e. the range of f are precisely the indices of the elements of xs that satisfy P.

lemma filterE:

```
fixes P:: "'a \Rightarrow bool" and xs :: "'a list" assumes "length (filter P xs) = n" obtains f:: "nat \Rightarrow nat" where "strict_mono_on {..<n} f" "\lambda i. i < n \Longrightarrow f i < length xs" "\lambda i. i < n \Longrightarrow filter P xs ! i = xs ! f i" "\lambda i. i < length xs \Longrightarrow P (xs ! i) \longleftrightarrow (\exists j. j < n \lambda f j = i)" \lambda proof\rangle
```

The following lemma shows the uniqueness of the above property. It is very useful for finding a "closed form" for $filter\ P\ xs$ in some concrete situation. For example, if we know that exactly every other element of xs satisfies P,

For example, if we know that exactly every other element of xs satisfies P, we can use it to prove that filter P xs = map ((*) 2) [0..<length xs div 2]

```
lemma filter_eqI:
    fixes f :: "nat \Rightarrow nat" and xs ys :: "'a list"
    defines "n \Rightarrow length ys"
    assumes "strict_mono_on {..<n} f"</pre>
```

```
assumes "\bigwedgei. i < n \Longrightarrow f i < length xs"
  assumes "\bigwedgei. i < n \Longrightarrow ys ! i = xs ! f i"
  assumes "\landi. i < length xs \Longrightarrow P (xs ! i) \longleftrightarrow (\exists j. j < n \land f j = i)"
            "filter P xs = ys"
  \langle proof \rangle
lemma filter_eq_iff:
   "filter P xs = ys \longleftrightarrow
       (\exists f. strict\_mono\_on \{...<length ys} f \land 
             (\forall\: i {<} length \: ys. \: f \: i \: {<} \: length \: xs \: \wedge \: ys \: ! \: i \: {=} \: xs \: ! \: f \: i) \: \wedge
             (\forall i < length \ xs. \ P \ (xs ! i) \longleftrightarrow (\exists j. j < length \ ys \land f \ j = i)))"
   (is "?lhs = ?rhs")
\langle proof \rangle
2.3 Polynomials
lemma poly_of_nat [simp]: "poly (of_nat n) x = of_nat n"
  \langle proof \rangle
lemma poly_of_int [simp]: "poly (of_int n) x = of_int n"
lemma poly_numeral [simp]: "poly (numeral n) x = numeral n"
  \langle proof \rangle
lemma order_gt_0_iff: "p \neq 0 \Longrightarrow order x p > 0 \longleftrightarrow poly p x = 0"
  \langle proof \rangle
lemma order_eq_0_iff: "p \neq 0 \Longrightarrow order x p = 0 \longleftrightarrow poly p x \neq 0"
  \langle proof \rangle
lemma coeff_pcompose_monom_linear [simp]:
  fixes p :: "'a :: comm_ring_1 poly"
  shows "coeff (pcompose p (monom c (Suc 0))) k = c \hat{k} * coeff p k"
  \langle proof \rangle
lemma of_nat_mult_conv_smult: "of_nat n * P = smult (of_nat n) P"
  \langle proof \rangle
lemma numeral_mult_conv_smult: "numeral n * P = smult (numeral n) P"
  \langle proof \rangle
lemma has_field_derivative_poly [derivative_intros]:
  assumes "(f has_field_derivative f') (at x within A)"
              "((\lambda x. poly p (f x)) has_field_derivative
                 (f' * poly (pderiv p) (f x))) (at x within A)"
  \langle proof \rangle
lemma sum order le degree:
```

```
assumes "p \neq 0" shows "(\sum x \mid poly \mid p \mid x = 0. \text{ order } x \mid p) \leq degree \mid p" \mid \langle proof \rangle
```

2.4 Trigonometric functions

lemma sin_multiple_reduce:

```
"sin (x * numeral n :: 'a :: {real_normed_field, banach}) =
      sin x * cos (x * of_nat (pred_numeral n)) + cos x * sin (x * of_nat
(pred_numeral n))"
\langle proof \rangle
lemma cos_multiple_reduce:
  "cos (x * numeral n :: 'a :: {real_normed_field, banach}) =
      cos (x * of_nat (pred_numeral n)) * cos x - sin (x * of_nat (pred_numeral
n)) * sin x"
\langle proof \rangle
lemma \ arccos\_eq\_pi\_iff: \ "x \in \{-1..1\} \implies arccos \ x = pi \longleftrightarrow x = -1"
  \langle proof \rangle
\mathbf{lemma} \ \mathbf{arccos\_eq\_0\_iff:} \ "\mathtt{x} \ \in \ \{\texttt{-1..1}\} \ \Longrightarrow \ \mathbf{arccos} \ \mathtt{x} \ \texttt{=} \ \mathtt{0} \ \longleftrightarrow \ \mathtt{x} \ \texttt{=} \ \mathtt{1"}
  \langle proof \rangle
2.5 Hyperbolic functions
lemma cosh_double_cosh: "cosh (2 * x :: 'a :: {banach, real_normed_field})
= 2 * (cosh x)^2 - 1"
  \langle proof \rangle
lemma sinh_multiple_reduce:
  "sinh (x * numeral n :: 'a :: {real_normed_field, banach}) =
      sinh x * cosh (x * of_nat (pred_numeral n)) + cosh x * sinh (x *
of_nat (pred_numeral n))"
\langle proof \rangle
lemma cosh_multiple_reduce:
  "cosh (x * numeral n :: 'a :: {real_normed_field, banach}) =
      cosh (x * of_nat (pred_numeral n)) * cosh x + sinh (x * of_nat (pred_numeral
n)) * sinh x"
\langle proof \rangle
lemma cosh_arcosh_real [simp]:
  assumes "x \ge (1 :: real)"
            "cosh (arcosh x) = x"
  shows
\langle proof \rangle
lemma arcosh_eq_0_iff_real [simp]: "x \ge 1 \implies arcosh x = 0 \longleftrightarrow x = (1
:: real)"
  \langle proof \rangle
```

```
lemma arcosh_nonneg_real [simp]:
  assumes "x \ge 1"
  shows
            "arcosh (x :: real) \geq 0"
\langle proof \rangle
lemma\ arcosh\_real\_strict\_mono:
  fixes x y :: real
  assumes "1 \le x" "x < y"
  shows "arcosh x < arcosh y"
\langle proof \rangle
lemma arcosh_less_iff_real [simp]:
  fixes x y :: real
  assumes "1 \le x" "1 \le y"
            "arcosh x < arcosh y \longleftrightarrow x < y"
  \mathbf{shows}
  \langle proof \rangle
lemma arcosh_real_gt_1_iff [simp]: "x \ge 1 \Longrightarrow arcosh \ x > 0 \longleftrightarrow x \ne 1
(1 :: real)"
  \langle proof \rangle
lemma sinh_arcosh_real: "x \geq 1 \Longrightarrow sinh (arcosh x) = sqrt (x<sup>2</sup> - 1)"
  \langle proof \rangle
lemma sinh_arsinh_real [simp]: "sinh (arsinh x :: real) = x"
\langle proof \rangle
lemma arsinh_real_strict_mono:
  fixes x y :: real
  assumes "x < y"
  shows
            "arsinh x < arsinh y"
\langle proof \rangle
lemma arsinh_less_iff_real [simp]:
  fixes x y :: real
  shows "arsinh x < arsinh y \longleftrightarrow x < y"
  \langle proof \rangle
lemma arsinh_real_eq_0_iff [simp]: "arsinh x = 0 \longleftrightarrow x = (0 :: real)"
  \langle proof \rangle
lemma \ arsinh\_real\_pos\_iff \ [simp]: \ "arsinh \ x > 0 \longleftrightarrow x > (0 :: real)"
lemma \ arsinh\_real\_neg\_iff \ [simp]: \ "arsinh \ x < 0 \longleftrightarrow x < (0 :: real)"
  \langle proof \rangle
```

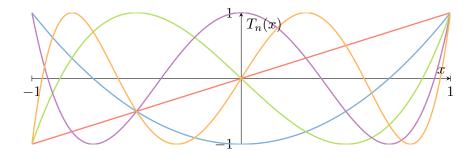


Figure 1: Some of the Chebyshev polynomials of the first kind, T_1 to T_5 .

```
lemma cosh_arsinh_real: "cosh (arsinh x) = sqrt (x^2 + 1)" \langle proof \rangle
```

end

3 Chebyshev Polynomials

```
theory Chebyshev_Polynomials
imports

"HOL-Analysis.Analysis"

"HOL-Real_Asymp.Real_Asymp"

"HOL-Computational_Algebra.Formal_Laurent_Series"

"Polynomial_Interpolation.Ring_Hom_Poly"

"Descartes_Sign_Rule.Descartes_Sign_Rule"

Polynomial_Transfer
   Chebyshev_Polynomials_Library
begin
```

3.1 Definition

We choose the recursive definition of T_n and U_n and do some setup to define both of them at once.

```
locale gen_cheb_poly =
  fixes c :: "'a :: comm_ring_1"
begin

fun f :: "nat ⇒ 'a ⇒ 'a" where
    "f 0 x = 1"
| "f (Suc 0) x = c * x"
| "f (Suc (Suc n )) x = 2 * x * f (Suc n) x - f n x"

fun P :: "nat ⇒ ('a :: comm_ring_1) poly" where
```

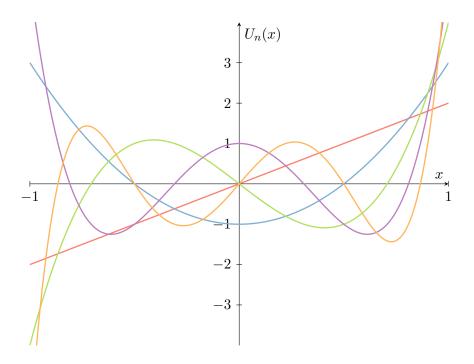


Figure 2: Some of the Chebyshev polynomials of the second kind, U_1 to U_5 .

```
"P \ O = 1"
| "P (Suc 0) = [:0, c:]"
| "P (Suc (Suc n)) = [:0, 2:] * P (Suc n) - P n"
lemma eval [simp]: "poly (P n) x = f n x"
  \langle proof \rangle
lemma eval_0:
  "f n 0 = (if odd n then 0 else (-1) \hat{} (n div 2))"
  \langle proof \rangle
lemma eval_1 [simp]:
  "f n 1 = of_nat n * (c - 1) + 1"
\langle proof \rangle
lemma uminus [simp]: "f n (-x) = (-1) ^ n * f n x"
  \langle proof \rangle
lemma pcompose_minus: "pcompose (P n) (monom (-1) 1) = (-1) ^ n * P n"
lemma degree_le: "degree (P n) \leq n"
\langle proof \rangle
```

```
lemma lead coeff:
  "coeff (P n) n = (if n = 0 then 1 else c * 2 \hat{ } (n - 1))"
\langle proof \rangle
lemma degree_eq:
  "c * 2 ^ (n - 1) \neq 0 \Longrightarrow degree (P n :: 'a poly) = n"
  \langle proof \rangle
lemmas [simp del] = f.simps(3) P.simps(3)
end
The two related constants Cheb_poly and cheb_poly denote the n-th Cheby-
shev polynomial of the first kind T_n and its interpretation as a function.
We make the definition polymorphic so that it works on every commutative
ring; however, many results will only hold for rings (or even only fields) of
characteristic 0.
definition cheb_poly :: "nat ⇒ 'a :: comm_ring_1 ⇒ 'a" where
 "cheb_poly = gen_cheb_poly.f 1"
definition Cheb_poly :: "nat \Rightarrow 'a :: comm_ring_1 poly" where
 "Cheb_poly = gen_cheb_poly.P 1"
interpretation cheb_poly: gen_cheb_poly 1
  rewrites "gen_cheb_poly.f 1 = cheb_poly" and "gen_cheb_poly.P 1 = Cheb_poly"
       and "/x :: 'a. 1 * x = x"
       and "\n". of_nat n * (1 - 1 :: 'a) + 1 = 1"
  \langle proof \rangle
lemmas cheb_poly_simps [code] = cheb_poly.f.simps
lemmas Cheb_poly_simps [code] = cheb_poly.P.simps
lemma Cheb_poly_of_int: "of_int_poly (Cheb_poly n) = Cheb_poly n"
  \langle proof \rangle
lemma degree_Cheb_poly [simp]:
  "degree (Cheb_poly n :: 'a :: {idom, ring_char_0} poly) = n"
  \langle proof \rangle
lemma lead_coeff_Cheb_poly [simp]:
  "lead_coeff (Cheb_poly n :: 'a :: {idom, ring_char_0} poly) = 2 ^ (n-1)"
  \langle proof \rangle
lemma Cheb_poly_nonzero [simp]: "Cheb_poly n \neq 0"
  \langle proof \rangle
lemma continuous_cheb_poly [continuous_intros]:
  fixes f :: "'b :: topological_space ⇒ 'a :: {real_normed_algebra_1,
comm_ring_1}"
```

```
shows "continuous_on A f \Longrightarrow continuous_on A (\lambda x. cheb_poly n (f x))"
  \langle proof \rangle
Similarly, we introduce two constants for U_n.
definition cheb_poly' :: "nat \Rightarrow 'a :: comm_ring_1 \Rightarrow 'a" where
 "cheb_poly' = gen_cheb_poly.f 2"
definition Cheb_poly' :: "nat ⇒ 'a :: comm_ring_1 poly" where
 "Cheb_poly' = gen_cheb_poly.P 2"
interpretation cheb_poly': gen_cheb_poly 2
  rewrites "gen_cheb_poly.f 2 \equiv cheb_poly'" and "gen_cheb_poly.P 2 = cheb_poly
Cheb_poly'"
       and "\nn. of_nat n * (2 - 1 :: 'a) + 1 = of_nat (Suc n)"
  \langle proof \rangle
lemmas cheb_poly'_simps [code] = cheb_poly'.f.simps
lemmas Cheb_poly'_simps [code] = cheb_poly'.P.simps
lemma Cheb_poly'_of_int: "of_int_poly (Cheb_poly' n) = Cheb_poly' n"
  \langle proof \rangle
lemma degree_Cheb_poly' [simp]:
  "degree (Cheb_poly' n :: 'a :: {idom, ring_char_0} poly) = n"
  \langle proof \rangle
lemma lead_coeff_Cheb_poly' [simp]:
  "lead_coeff (Cheb_poly' n :: 'a :: {idom, ring_char_0} poly) = 2 ^ n"
  \langle proof \rangle
lemma Cheb_poly_nonzero' [simp]: "Cheb_poly' n ≠ (0 :: 'a :: {comm_ring_1,
ring_char_0} poly)"
\langle proof \rangle
lemma continuous_cheb_poly' [continuous_intros]:
  fixes f :: "'b :: topological_space >> 'a :: {real_normed_algebra_1,
comm_ring_1}"
  shows "continuous_on A f \Longrightarrow continuous_on A (\lambda x. cheb_poly' n (f x))"
  \langle proof \rangle
```

3.2 Relation to trigonometric functions

Consider the multiple angle formulas for the cosine function:

$$\cos 1x = \cos x$$

$$\cos 2x = 1 + 2\cos^2 x$$

$$\cos 3x = -3\cos x + 4\cos^3 x$$

$$\cos 4x = 1 - 8\cos^2 x + 8\cos^4 x$$

It seems that for any $n \in \mathbb{N}$, we can write $\cos(nx)$ as a sum of powers $\cos^i x$ for $0 \le i \le n$, i.e. as a polynomial in $\cos x$ of degree n. It turns out that this polynomial is exactly T_n . This can also serve as an alternative, trigonometric definition of T_n .

Proving it is a simple induction:

```
lemma cheb_poly_cos [simp]:
  fixes x :: "'a :: {banach, real_normed_field}"
    shows "cheb_poly n (cos x) = cos (of_nat n * x)"
    ⟨proof⟩
```

If we look at the multiple angular formulae for the sine function, we see a similar pattern:

```
\sin 1x = \sin x
\sin 2x = 2\sin x \cos x
\sin 3x = \sin x(-1 + 4\cos^2 x)
\sin 4x = \sin x(-4\cos x + 8\cos^3 x)
```

It seems that $\sin nx/\sin x$ can be expressed as a polynomial in $\cos x$ of degree n-1. This polynomial turns out to be exactly U_{n-1} .

```
lemma cheb_poly'_cos:
  fixes x :: "'a :: {banach, real_normed_field}"
  shows "cheb_poly' n (cos x) * \sin x = \sin (of_nat (n+1) * x)"
\langle proof \rangle
lemma cheb_poly_conv_cos:
  assumes ||x::real| \le 1||
  shows
            "cheb_poly n x = cos (n * arccos x)"
  \langle proof \rangle
lemma cheb_poly'_cos':
  fixes x :: "'a :: {real_normed_field, banach}"
  shows "\sin x \neq 0 \implies cheb\_poly' n (cos x) = sin (of\_nat (n+1) * x)
/ sin x"
  \langle proof \rangle
lemma cheb_poly'_conv_cos:
  assumes "|x::real| < 1"
            "cheb_poly' n x = sin (real (n+1) * arccos x) / sqrt (1 - x^2)"
  shows
\langle proof \rangle
lemma cos_multiple:
  fixes x :: "'a :: {banach, real_normed_field}"
  shows "cos (numeral n * x) = poly (Cheb_poly (numeral n)) (cos x)"
  \langle proof \rangle
```

```
lemma sin_multiple:
  fixes x :: "'a :: {banach, real_normed_field}"
  shows "sin (numeral n * x) = sin x * poly (Cheb_poly' (pred_numeral
n)) (cos x)"
  \langle proof \rangle
Example application: quadruple-angle formulas for sin and cos:
lemma cos_quadruple:
  fixes x :: "'a :: {banach, real_normed_field}"
  shows "\cos (4 * x) = 8 * \cos x ^4 - 8 * \cos x ^2 + 1"
  \langle proof \rangle
lemma sin_quadruple:
  fixes x :: "'a :: {banach, real_normed_field}"
  shows "\sin (4 * x) = \sin x * (8 * \cos x ^3 - 4 * \cos x)"
  \langle proof \rangle
3.3 Relation to hyperbolic functions
lemma cheb_poly_cosh [simp]:
  fixes x :: "'a :: {banach, real_normed_field}"
  shows "cheb_poly n (cosh x) = cosh (of_nat n * x)"
\langle proof \rangle
lemma cheb_poly'_cosh:
  fixes x :: "'a :: {real_normed_field, banach}"
  shows "cheb_poly' n (cosh x) * sinh x = sinh (of_nat (n+1) * x)"
\langle proof \rangle
lemma cheb_poly_conv_cosh:
  assumes "(x :: real) \ge 1"
  shows
           "cheb_poly n x = \cosh (n * arcosh x)"
  \langle proof \rangle
lemma cheb_poly'_cosh':
  fixes x :: "'a :: {real_normed_field, banach}"
  shows "sinh x \neq 0 \implies cheb\_poly' n (cosh x) = sinh (of\_nat (n+1) *
x) / sinh x"
  \langle proof \rangle
lemma cheb_poly'_conv_cosh:
  assumes "x > (1 :: real)"
          "cheb_poly' n x = sinh (real (n+1) * arcosh x) / sqrt (x^2 -
1)"
\langle proof \rangle
```

3.4 Roots

 T_n has n distinct real roots, namely:

$$x_k = \cos\left(\frac{2k+1}{2n}\pi\right)$$

These are called the *Chebyshev nodes* of degree n.

```
definition cheb_node :: "nat \Rightarrow nat \Rightarrow real" where
  "cheb_node n k = cos (real (2*k+1) / real (2*n) * pi)"
lemma cheb_poly_cheb_node [simp]:
  assumes "k < n"
  shows
            "cheb_poly n (cheb_node n k) = 0"
\langle proof \rangle
lemma strict_antimono_cheb_node: "monotone_on {..<n} (<) (>) (cheb_node
n)"
  \langle proof \rangle
lemma cheb node pos iff:
  assumes k: "k < n"
            "cheb node n k > 0 \longleftrightarrow k < n div 2"
  shows
\langle proof \rangle
lemma cheb_poly_roots_bij_betw:
  "bij_betw (cheb_node n) \{..< n\} \{x. cheb_poly n x = 0\}"
lemma card_cheb_poly_roots: "card {x::real. cheb_poly n x = 0} = n"
  \langle proof \rangle
It is easy to see that all the Chebyshev nodes have order 1 as roots of T_n.
lemma order_Cheb_poly_cheb_node [simp]:
  assumes "k < n"
  shows
            "order (cheb_node n k) (Cheb_poly n) = 1"
\langle proof \rangle
lemma order Cheb poly [simp]:
  assumes "poly (Cheb_poly n) (x :: real) = 0"
            "order x (Cheb_poly n) = 1"
  shows
\langle proof \rangle
```

This also means that T_n is square-free. We only show this for the case where we view T_n as a real polynomial, but this also holds in every other reasonable ring since \mathbb{R} is a splitting field of T_n (as we have just shown). However, we chose not to do this here.

lemma rsquarefree_Cheb_poly_real: "rsquarefree (Cheb_poly n :: real poly)"

```
\langle proof \rangle
```

Similarly, the n distinct real roots of U_n are:

$$y_i = \cos\left(\frac{k+1}{n+1}\pi\right)$$

```
definition cheb_node' :: "nat \Rightarrow nat \Rightarrow real" where
  "cheb_node' n k = cos (real (k+1) / real (n+1) * pi)"
lemma cheb_poly'_cheb_node' [simp]:
  assumes "k < n"
  shows
          "cheb_poly' n (cheb_node' n k) = 0"
\langle proof \rangle
lemma strict_antimono_cheb_node': "monotone_on {..<n} (<) (>) (cheb_node'
n)"
  \langle proof \rangle
lemma cheb_node'_pos_iff:
  assumes k: "k < n"
  shows
            "cheb_node' n k > 0 \longleftrightarrow k < n div 2"
\langle proof \rangle
lemma cheb_poly'_roots_bij_betw:
  "bij_betw (cheb_node' n) \{..< n\} \{x. cheb_poly' n x = 0\}"
\langle proof \rangle
lemma card_cheb_poly'_roots: "card {x::real. cheb_poly' n x = 0} = n"
lemma order_Cheb_poly'_cheb_node' [simp]:
  assumes "k < n"
  shows
            "order (cheb_node' n k) (Cheb_poly' n) = 1"
\langle proof \rangle
lemma order_Cheb_poly' [simp]:
  assumes "poly (Cheb_poly' n) (x :: real) = 0"
            "order x (Cheb_poly' n) = 1"
\langle proof \rangle
lemma rsquarefree_Cheb_poly'_real: "rsquarefree (Cheb_poly' n :: real
poly)"
  \langle proof \rangle
```

3.5 Generating functions

 T_n and U_n have the following rational generating functions:

$$\sum_{n=0}^{\infty} T_n(x)t^n = \frac{1-tx}{1-2tx+t^2} \qquad \sum_{n=0}^{\infty} U_n(x)t^n = \frac{1}{1-2tx+t^2}$$

This is a simple consequence of the linear recurrence equations they satisfy (which we used as their definitions).

Due to some limitations coming from the type class structure, we cannot currently write this down nicely as an equation, but the following form is almost as good.

```
theorem Abs\_fps\_Cheb\_poly:
    fixes F \ X \ T :: "real fps \ fps"
    defines "X \equiv fps\_const \ fps\_X" and "T \equiv fps\_X"
    defines "F \equiv Abs\_fps \ (fps\_of\_poly \circ Cheb\_poly)"
    shows "F * (1 - 2 * T * X + T^2) = 1 - T * X"

\langle proof \rangle

theorem Abs\_fps\_Cheb\_poly':
    fixes F \ X \ T :: "real fps \ fps"
    defines "X \equiv fps\_const \ fps\_X" and "T \equiv fps\_X"
    defines "F \equiv Abs\_fps \ (fps\_of\_poly \circ Cheb\_poly')"
    shows "F * (1 - 2 * T * X + T^2) = 1"

\langle proof \rangle
```

3.6 Optimality with respect to the ∞ -norm

We now turn towards a property of T_n that explains why they are interesting for interpolating smooth functions. If $f:[0,1] \to \mathbb{R}$ is a smooth function on the unit interval, the approximation error attained when interpolating f with a polynomial P of degree n at the interpolation points x_1, \ldots, x_n is

$$\frac{f^{(n)}(\xi)}{n!} \prod_{i=1}^{n} (x - x_i) .$$

Therefore, it makes sense to choose the interpolation points such that $\prod_{i=1}^{n} (x-x_i)$ is minimal.

We will show below results that imply that this product cannot be smaller than 2^{1-n} , and it is easy to see that if we choose x_i to be the Chebyshev nodes then the product becomes exactly 2^{1-n} and thus optimal.

Out first result is now the following: The ∞ -norm of a monic polynomial of degree n on the unit interval [-1,1] is at least 2^{1-n} . This gives us a kind of lower bound on the "oscillation" of polynomials: a monic polynomial of degree n cannot stay closer than 2^{1-n} to 0 at every point of the unit interval.

```
lemma Sup\_abs\_poly\_bound\_aux:
    fixes p :: "real poly"
    assumes "lead\_coeff p = 1"
    shows "\exists x \in \{-1..1\}. |poly\ p\ x| \ge 1\ /\ 2 ^ (degree p - 1)"
    \langle proof \rangle

lemma Sup\_abs\_poly\_bound\_unit\_ivl:
    fixes p :: "real poly"
    shows "(SUP\ x \in \{-1..1\}. |poly\ p\ x|) \ge |lead\_coeff\ p|\ /\ 2 ^ (degree p - 1)"
    \langle proof \rangle
```

Using an appropriate change of variables, we obtain the following bound in the most general form for a non-constant polynomial P(x) on some non-empty interval [a, b]:

$$\sup_{x \in [a,b]} |P(x)| \ge 2 \cdot \mathrm{lc}(p) \cdot \left(\frac{b-a}{4}\right)^{\mathrm{deg}(p)}$$

where lc(p) denotes the leading coefficient of p.

```
theorem Sup\_abs\_poly\_bound: fixes p :: "real poly" assumes "a < b" and "degree p > 0" shows "(SUP \ x \in \{a..b\}. |poly \ p \ x|) \geq 2 * |lead\_coeff \ p| * ((b - a) / 4) ^ degree p" \langle proof \rangle
```

If we scale T_n with a factor of 2^{1-n} , it exactly attains the lower bound we just derived. The Chebyshev polynomials of the first kind are, in that sense, the polynomials that stay closest to 0 within the unit interval.

With some more work (that we will not do), one can see that T_n is in fact the *only* polynomial that attains this minimal deviation (see e.g. Corollary 3.4B in Mason & Handscomb [1]). This fact, however, requires proving the Equioscillation Theorem, which is not so easy and beyond the scope of this entry.

```
lemma abs_cheb_poly_le_1: assumes "(x :: real) \in {-1..1}" shows "|cheb_poly n x| \leq 1" \langle proof \rangle

theorem Sup_abs_poly_bound_sharp: fixes n :: nat and p :: "real poly" defines "p \equiv smult (1 / 2 ^ (n - 1)) (Cheb_poly n)" shows "degree p = n" and "lead_coeff p = 1" and "(SUP x\in{-1..1}. |poly p x|) = 1 / 2 ^ (n - 1)" \langle proof \rangle
```

A related fact: among all the real polynomials of degree n whose absolute value is bounded by 1 within the unit interval, T_n is the one that grows fastest *outside* the unit interval.

```
theorem cheb_poly_fastest_growth: fixes p:: "real poly" defines "n \equiv degree p" assumes p_bounded: "\landx. |x| \leq 1 \Longrightarrow |poly \ p \ x| \leq 1" assumes x: "x \notin {-1<..<1}" shows "|cheb_poly n x| \geq |poly p x|" \langle proof \rangle
```

3.7 Some basic equations

We first set up a mechanism to allow us to prove facts about Chebyshev polynomials on any ring with characteristic 0 by proving them for Chebyshev polynomials over \mathbb{R} .

```
definition rel_ring_int :: "'a :: ring_1 \Rightarrow 'b :: ring_1 \Rightarrow bool" where
  "rel_ring_int x y \longleftrightarrow (\exists n::int. x = of_int n \land y = of_int n)"
lemma rel_ring_int_0: "rel_ring_int 0 0"
  \langle proof \rangle
lemma rel_ring_int_1: "rel_ring_int 1 1"
  \langle proof \rangle
lemma rel_ring_int_add:
  "rel_fun rel_ring_int (rel_fun rel_ring_int rel_ring_int) (+) (+)"
  \langle proof \rangle
lemma rel ring int mult:
  "rel_fun rel_ring_int (rel_fun rel_ring_int rel_ring_int) (*) (*)"
  \langle proof \rangle
lemma rel_ring_int_minus:
  "rel_fun rel_ring_int (rel_fun rel_ring_int rel_ring_int) (-) (-)"
  \langle proof \rangle
lemma rel_ring_int_uminus:
  "rel_fun rel_ring_int rel_ring_int uminus uminus"
  \langle proof \rangle
lemma sgn_of_int: "sgn (of_int n :: 'a :: linordered_idom) = of_int (sgn
n)"
  \langle proof \rangle
lemma rel_ring_int_sgn:
  "rel_fun rel_ring_int (rel_ring_int :: 'a :: linordered_idom \Rightarrow 'b ::
linordered_idom ⇒ bool) sgn sgn"
```

```
\langle proof \rangle
lemma bi_unique_rel_ring_int:
  "bi_unique (rel_ring_int :: 'a :: ring_char_0 ⇒ 'b :: ring_char_0 ⇒
bool)"
  \langle proof \rangle
lemmas rel_ring_int_transfer =
  rel_ring_int_0 rel_ring_int_1 rel_ring_int_add rel_ring_int_mult rel_ring_int_minus
  rel_ring_int_uminus bi_unique_rel_ring_int
lemma rel_poly_rel_ring_int:
  "rel_poly rel_ring_int p q \longleftrightarrow (\exists r. p = of_int_poly r \land q = of_int_poly
r)"
\langle proof \rangle
lemma Cheb_poly_transfer:
  "rel_fun (=) (rel_poly rel_ring_int) Cheb_poly Cheb_poly"
\langle proof \rangle
lemma Cheb_poly'_transfer:
  "rel_fun (=) (rel_poly rel_ring_int) Cheb_poly' Cheb_poly'"
\langle proof \rangle
context
  fixes T :: "'a :: {idom, ring_char_0} itself"
  notes [transfer_rule] = rel_ring_int_transfer [where ?'a = real and
?'b = 'a
                             Cheb_poly_transfer[where ?'a = real and ?'b
= 'a]
                             Cheb_poly'_transfer[where ?'a = real and ?'b
= 'a]
                             transfer_rule_of_nat transfer_rule_numeral
begin
The following rule allows us to prove an equality of real polynomials P = Q
by proving that P(\cos x) = Q(\cos x) for all x \in (0, \alpha) for some \alpha > 0.
This holds because there are infinitely many such \cos x, but P-Q, being a
polynomial, can only have finitely many roots if P \neq 0.
lemma Cheb_poly_equalities_aux:
  fixes p q :: "real poly"
  assumes "a > 0"
  assumes "\xspace x. x \in \{0 < ... < a\} \implies poly p (cos x) = poly q (cos x)"
  shows
          p = q''
\langle proof \rangle
First, we show that T_n(x) = nU_{n-1}(x):
lemma pderiv_Cheb_poly: "pderiv (Cheb_poly n) = of_nat n * (Cheb_poly'
(n - 1) :: 'a poly)"
```

```
\langle proof \rangle
```

 $\langle proof \rangle$

Next, we show that:

$$U'_n(x) = \frac{1}{x^2 - 1}((n+1)T_{n+1}(x) - xU_n(x))$$

```
lemma pderiv_Cheb_poly':
  "pderiv (Cheb_poly' n) * [:-1, 0, 1 :: 'a:] =
     of_nat (n+1) * Cheb_poly (n+1) - [:0,1:] * Cheb_poly' n"
Next, we have T_n(x) = \frac{1}{2}(U_n(x) - U_{n-2}(x)).
lemma Cheb_poly_rec:
  assumes n: "n > 2"
  shows "2 * Cheb_poly n = Cheb_poly' n - (Cheb_poly' (n - 2) :: 'a poly)"
\langle proof \rangle
lemma cheb_poly_rec:
  assumes n: "n \ge 2"
  shows "2 * cheb_poly n x = cheb_poly' n x - cheb_poly' (n - 2) (x::'a)"
Next, we have U_n(x) = xU_{n-1}(x) + T_n(x).
lemma Cheb poly' rec:
  assumes n: "n > 0"
           "Cheb_poly' n = [:0,1::'a:] * Cheb_poly' (n - 1) + Cheb_poly
  shows
n"
\langle proof \rangle
lemma cheb_poly'_rec:
  assumes n: "n > 0"
  shows "cheb_poly' n x = x * cheb_poly' (n-1) x + cheb_poly n (x::'a)"
  \langle proof \rangle
Next, T_n(x) = xT_{n-1}(x) + (x^2 - 1)U_{n-2}(x).
lemma Cheb_poly_rec':
  assumes n: "n \ge 2"
  shows "Cheb_poly n = [:0,1::'a:] * Cheb_poly (n-1) + [:-1,0,1:] * Cheb_poly'
(n-2)"
\langle proof \rangle
lemma cheb_poly_rec':
  assumes n: "n \ge 2"
  shows "cheb_poly n x = x * cheb_poly (n-1) x + (x^2 - 1) * cheb_poly'
(n-2) (x::'a)"
```

 T_n and U_{-1} are a solution to a Pell-like equation on polynomials:

$$T_n(x)^2 + (1 - x^2)U_{n-1}(x)^2 = 1$$

```
lemma Cheb_poly_Pell:
   assumes n: "n > 0"
   shows "Cheb_poly n ^2 + [:1, 0, -1::'a:] * Cheb_poly' (n - 1) ^2 = 1"
\langle proof \rangle
lemma cheb_poly_Pell:
   assumes n: "n > 0"
   shows "cheb_poly n \times 2 + (1 - x^2) * cheb_poly' (n-1) \times 2 = (1::'a)"
\langle proof \rangle
```

The following Turán-style equation also holds:

$$T_{n+1}(x)^2 - T_{n+2}(x)T_n(x) = 1 - x^2$$

lemma Cheb_poly_Turan:

"Cheb_poly (n+1) ^ 2 - Cheb_poly (n+2) * Cheb_poly n = [:1,0,-1::'a:]" $\langle proof \rangle$

lemma cheb_poly_Turan:

"cheb_poly (n+1) x ^ 2 - cheb_poly (n+2) x * cheb_poly n x = (1 - x ^ 2 :: 'a)"
$$\langle proof \rangle$$

And, the analogous one for U_n :

$$U_{n+1}(x)^2 - U_{n+2}(x)U_n(x) = 1$$

lemma Cheb_poly'_Turan:

```
"Cheb_poly' (n+1) ^ 2 - Cheb_poly' (n+2) * Cheb_poly' n = (1 :: 'a poly)" \langle proof \rangle
```

lemma cheb_poly'_Turan:

```
"cheb_poly' (n+1) x ^ 2 - cheb_poly' (n+2) x * cheb_poly' n x = (1 :: 'a)"  \langle proof \rangle
```

There is also a nice formula for the product of two Chebyshev polynomials of the first kind:

$$T_m(x)T_n(x) = \frac{1}{2}(T_{m+n}(x) + T_{m-n}(x))$$

lemma Cheb_poly_prod:

```
assumes "n \le m" shows "2 * Cheb_poly m * Cheb_poly n = Cheb_poly (m + n) + (Cheb_poly (m - n) :: 'a poly)"
```

 $\langle proof \rangle$

```
lemma cheb_poly_prod':
   assumes "n \le m"
   shows "2 * cheb_poly m x * cheb_poly n x = cheb_poly (m + n) x + cheb_poly (m - n) (x :: 'a)"
   \langle proof \rangle
```

In particular, this leads to a divide-and-conquer-style recurrence relation for T_n for even and odd n:

$$T_{2n}(x) = 2T_n(x)^2 - 1$$
$$T_{2n+1} = 2T_n(x)T_{n+1}(x) - x$$

lemma Cheb_poly_even:

```
"Cheb_poly (2 * n) = 2 * Cheb_poly n ^ 2 - (1 :: 'a poly)" \langle proof \rangle
```

lemma cheb_poly_even:

```
"cheb_poly (2 * n) x = 2 * cheb_poly n x ^ 2 - (1 :: 'a)" \langle proof \rangle
```

lemma Cheb_poly_odd:

```
"Cheb_poly (2 * n + 1) = 2 * Cheb_poly n * Cheb_poly (Suc n) - [:0,1::'a:]" \langle proof \rangle
```

lemma cheb_poly_odd:

```
"cheb_poly (2 * n + 1) x = 2 * cheb_poly n x * cheb_poly (Suc n) x - (x :: 'a)" \langle proof \rangle
```

Remarkably, we also have the following formula for the composition of two Chebyshev polynomials of the first kind:

$$T_{mn}(x) = T_m(T_n(x))$$

theorem Cheb_poly_mult:

```
"(Cheb_poly (m * n) :: 'a poly) = pcompose (Cheb_poly m) (Cheb_poly n)" \langle proof \rangle
```

```
corollary cheb_poly_mult: "cheb_poly m (cheb_poly n x) = cheb_poly (m * n) (x :: 'a)" \langle proof \rangle
```

For the Chebyshev polynomials of the second kind, the following more complicated relationship holds:

$$U_{mn-1}(x) = U_{m-1}(T_n(x)) \cdot U_{n-1}(x)$$

The following two lemmas tell tell us that

$$U'_n(1) = 2\binom{n+2}{3} = \frac{1}{3}n(n+1)(n+2)$$
$$U'_n(-1) = (-1)^{n+1}2\binom{n+2}{3} = \frac{(-1)^{n+1}}{3}n(n+1)(n+2)$$

This is good to know because our formula for U'_n has a "division by zero" at ± 1 , so we cannot use it to establish these values.

```
lemma poly_pderiv_Cheb_poly'_1:    "3 * poly (pderiv (Cheb_poly' n) :: 'a poly) 1 = of_nat ((n + 2) * (n + 1) * n)"    \langle proof \rangle
```

```
lemma poly_pderiv_Cheb_poly'_neg_1:

"3 * poly (pderiv (Cheb_poly' n) :: 'a poly) (-1) = (-1)^Suc n * of_nat ((n + 2) * (n + 1) * n)"

\langle proof \rangle
```

Another alternative definition of T_n and U_n is as the solutions of the ordinary differential equations

$$(1 - x^2)T_n'' - xT_n' + n^2T_n = 0$$
$$(1 - x^2)U_n'' - 3xU_n' + n(n+2)U_n = 0$$

```
lemma Cheb_poly_ODE:
    fixes n :: nat
    defines "p \equiv (Cheb\_poly \ n :: 'a poly)"
    shows "[:1,0,-1:] * (pderiv ^^ 2) p - [:0,1:] * pderiv p + of_nat
    n ^ 2 * p = 0"
    \langle proof \rangle

lemma Cheb_poly'_ODE:
    fixes n :: nat
    defines "p \equiv (Cheb\_poly' \ n :: 'a poly)"
```

```
"[:1,0,-1:] * (pderiv ^^ 2) p - [:0,3:] * pderiv p + of_nat
(n*(n+2)) * p = 0"
\langle proof \rangle
end
lemma cheb_poly_prod:
  fixes x :: "'a :: field_char_0"
  assumes "n \leq m"
          "cheb_poly m x * cheb_poly n x = (cheb_poly (m + n) x + cheb_poly
  \mathbf{shows}
(m - n) x) / 2"
  \langle proof \rangle
lemma has_field_derivative_cheb_poly [derivative_intros]:
  assumes "(f has_field_derivative f') (at x within A)"
          "((\lambda x. cheb poly n (f x)) has field derivative
               (of_nat n * cheb_poly' (n-1) (f x) * f')) (at x within
A)"
  \langle proof \rangle
lemma has_field_derivative_cheb_poly' [derivative_intros]:
  "(cheb_poly' n has_field_derivative
     (if x = 1 then of_nat ((n + 2) * (n + 1) * n) / 3
      else if x = -1 then (-1)^Suc n * of_nat ((n + 2) * (n + 1) * n)
      else (of_nat (n+1) * cheb_poly (Suc n) x - x * cheb_poly' n x) /
(x^2 - 1)))
   (at x within A)" (is "(_ has_field_derivative ?f') (at _ within _)")
\langle proof \rangle
lemmas has_field_derivative_cheb_poly'' [derivative_intros] =
  DERIV_chain'[OF _ has_field_derivative_cheb_poly']
```

3.8 Signs of the coefficients

Since $T_n(-x) = (-1)^n T_n(x)$ and analogously for U_n , the Chebyshev polynomials are even functions when n is even and odd functions when n is odd. Consequently, when n is even, the coefficients of X^k for any odd k are 0 and analogously when n is odd.

```
lemma coeff_Cheb_poly_eq_0:
   assumes "odd (n + k)"
   shows "coeff (Cheb_poly n :: 'a :: {idom,ring_char_0} poly) k = 0"
   ⟨proof⟩

lemma coeff_Cheb_poly'_eq_0:
   assumes "odd (n + k)"
   shows "coeff (Cheb_poly' n :: 'a :: {idom,ring_char_0} poly) k = 0"
   ⟨proof⟩
```

Next, we analyse the behaviour of the signs of the coefficients of T_n and U_n more generally and show that:

- The leading coefficient is positive.
- After that, every second coefficient is 0.
- The remaining coefficients are non-zero and their signs alternate.

In conclusion, we have

$$\operatorname{sgn}([X^k] T_n(X)) = \operatorname{sgn}([X^k] U_n(X)) = \begin{cases} 0 & \text{if } k > n \text{ or } (n+k) \text{ is odd} \\ (-1)^{\frac{n-k}{2}} & \text{otherwise} \end{cases}$$

The proof works using Descartes' rule of signs: We know that T_n and U_n have n distinct real roots and $\lfloor \frac{n}{2} \rfloor$ of them are positive. By Descartes' rule of signs, this implies that the coefficient sequences of T_n and U_n must have at least $\lfloor \frac{n}{2} \rfloor$ sign alternations. However, we also already know that every other coefficient of T_n and U_n starting with $\lfloor X^{n-1} \rfloor$ is 0, so the number of sign alternations must be $exactly \lfloor \frac{n}{2} \rfloor$.

```
lemma sgn_coeff_Cheb_poly_aux:
  fixes n :: nat and P :: "real poly"
  assumes "degree P = n"
  assumes "\landi. odd (n + i) \Longrightarrow coeff P i = 0"
  assumes "card \{x. x > 0 \land poly P x = 0\} = n \ div 2"
  assumes "rsquarefree P"
  assumes "coeff P n > 0"
  shows "sgn (coeff P i) = (if i > n \lor odd (n + i) then 0 else (-1) ^
((n - i) div 2))"
\langle proof \rangle
theorem sgn_coeff_Cheb_poly:
  "sgn (coeff (Cheb_poly n) i :: 'a :: linordered_idom) =
     (if i > n \lor odd (n + i) then 0 else (-1) ^ ((n - i) div 2))"
\langle proof \rangle
theorem sgn_coeff_Cheb_poly':
  "sgn (coeff (Cheb_poly' n) i :: 'a :: linordered_idom) =
     (if i > n \lor odd (n + i) then 0 else (-1) ^ ((n - i) div 2))"
\langle proof \rangle
```

3.9 Orthogonality and integrals

context

```
fixes n :: nat and x :: "nat \Rightarrow real" defines "x \equiv (\lambda k. cos (real (Suc (2 * k)) / real (2 * n) * pi))" begin
```

lemma cheb_poly_orthogonality_discrete_aux: assumes "1 \in {0<..<2*n}" shows "(\sum k<n. cos (real 1 * real (Suc (2 * k)) / real (2 * n) * pi)) = 0" $\langle proof \rangle$

For k = 0, ..., n-1 let $x_k = \cos(\frac{2k+1}{2n}\pi)$ be the Chebyshev nodes of order n, i.e. the roots of T_n . Then the following discrete orthogonality relation holds for the Chebyshev polynomials of the first kind (for any i, j < n):

$$\sum_{k=0}^{n-1} T_i(x_k) T_j(x_k) = \begin{cases} n & \text{if } i = j = 0\\ \frac{n}{2} & \text{if } i = j \neq 0\\ 0 & \text{if } i \neq j \end{cases}$$

theorem cheb_poly_orthogonality_discrete:

```
fixes i j :: nat assumes "i < n" "j < n" shows "(\sum k < n. cheb_poly i (x k) * cheb_poly j (x k)) = (if i = j then if i = 0 then n else n / 2 else 0)" \langle proof \rangle
```

A similar relation holds for the Chebyshev polynomials of the second kind:

$$\sum_{k=0}^{n-1} U_i(x_k) U_j(x_k) (1 - x_k^2) = \begin{cases} n & \text{if } i = j = n - 1\\ \frac{n}{2} & \text{if } i = j \neq 0\\ 0 & \text{if } i \neq j \end{cases}$$

theorem cheb_poly'_orthogonality_discrete:

```
fixes i j :: nat
assumes "i < n" "j < n"
shows "(\sum k < n. cheb\_poly' i (x k) * cheb\_poly' j (x k) * (1 - x k ^ 2)) =

(if i = j then if i = n - 1 then n else n / 2 else 0)"
\langle proof \rangle
```

end

We now show the continuous orthogonality relations. For the polynomials of the first kind, the relation is:

$$\int_{-1}^{1} \frac{T_m(x)T_n(x)}{\sqrt{1-x^2}} dx = \begin{cases} \pi & \text{if } m = n = 0\\ \frac{\pi}{2} & \text{if } m = n \neq 0\\ 0 & \text{if } m \neq n \end{cases}$$

The proof works by a change of variables $x = \cos \theta$, which converts the integral to the easier form $\int_0^{\pi} \cos(mt) \cos(nt) dx$, which can then be solved by a computing an indefinite integral (with appropriate case distinctions on m and n).

theorem cheb_poly_orthogonality:

```
fixes m n :: nat defines "I \equiv if m = n then if m = 0 then pi else pi / 2 else 0" shows "((\lambda x. cheb_poly m x * cheb_poly n x / sqrt (1 - x^2)) has_integral I) {-1..1}" \langle proof \rangle
```

For the polynomials of the second kind, the relation is:

$$\int_{-1}^{1} U_m(x)U_n(x)\sqrt{1-x^2} \, \mathrm{d}x = \begin{cases} \frac{\pi}{2} & \text{if } m=n\\ 0 & \text{if } m \neq n \end{cases}$$

The proof works the same as before.

theorem cheb_poly'_orthogonality: fixes m n :: nat defines "I \equiv if m = n then pi / 2 else 0" shows "((λ x. cheb_poly' m x * cheb_poly' n x * sqrt (1 - x²)) has_integral I) {-1..1}" $\langle proof \rangle$

We additionally show the following property about the integral from -1 to 1:

$$\int_{-1}^{1} T_n(x) \, \mathrm{d}x = \frac{1 + (-1)^n}{1 - n^2}$$

theorem cheb_poly_integral_neg1_1: "(cheb_poly n has_integral ((1 + (-1)^n) / (1 - n^2))) {-1..1::real}"

And, for the polynomials of the second kind:

$$\int_{-1}^{1} U_n(x) \, \mathrm{d}x = \frac{1 + (-1)^n}{n+1}$$

theorem cheb_poly'_integral_neg1_1: "(cheb_poly' n has_integral (1 + (-1) ^ n) / (n+1)) {-1..1::real}" $\langle proof \rangle$

3.10 Clenshaw's algorithm

Clenshaw's algorithm allows us to efficiently evaluate a weighted sum of Chebyshev polynomials of the first kind, i.e.

$$\sum_{i=0}^{n} w_i \cdot T_i(x) .$$

This is useful when evaluating interpolations.

```
locale clenshaw =
  fixes g :: "nat ⇒ 'a :: comm_ring_1"
  fixes a b :: "nat \Rightarrow 'a"
  assumes g_rec: "\n. g (Suc (Suc n)) = a n * g (Suc n) + b n * g n"
begin
context
  fixes N :: nat and c :: "nat <math>\Rightarrow 'a"
begin
function clenshaw_aux where
  "n \geq N \Longrightarrow clenshaw_aux n = 0"
| "n < N \Longrightarrow clenshaw_aux n =
     c (Suc n) + a n * clenshaw_aux (n+1) + b (Suc n) * clenshaw_aux (n+2)"
  \langle proof \rangle
termination \langle proof \rangle
lemma clenshaw_aux_correct_aux:
  assumes "n < N"
  shows "g n * c n + g (Suc n) * clenshaw_aux n + b n * g n * clenshaw_aux
(Suc n) = (\sum k=n..N. c k * g k)"
  \langle proof \rangle
fun clenshaw_aux' where
  "clenshaw_aux' 0 acc1 acc2 = g 0 * c 0 + g 1 * acc1 + b 0 * g 0 * acc2"
| "clenshaw_aux' (Suc n) acc1 acc2 = clenshaw_aux' n (c (Suc n) + a n
* acc1 + b (Suc n) * acc2) acc1"
lemma clenshaw_aux'_correct: "clenshaw_aux' N 0 0 = (\sum k \le N. \ c \ k * g
k)"
\langle proof \rangle
lemmas [simp del] = clenshaw_aux'.simps
end
lemma clenshaw_aux'_cong:
  "(\land k. \ k \le n \Longrightarrow c \ k = c' \ k) \Longrightarrow clenshaw_aux' \ c \ n \ acc1 \ acc2 = clenshaw_aux'
c' n acc1 acc2"
  \langle proof \rangle
```

```
definition clenshaw where "clenshaw N c = clenshaw_aux' c N 0 0"
theorem clenshaw_correct: "clenshaw N c = (\sum k \le N. \ c \ k * g \ k)"
  \langle proof \rangle
\quad \mathbf{end} \quad
definition cheb_eval :: "'a :: comm_ring_1 list \Rightarrow 'a \Rightarrow 'a" where
  "cheb_eval cs x = (\sum k < length cs. cs ! k * cheb_poly k x)"
interpretation cheb_poly: clenshaw "\lambdan. cheb_poly n x" "\lambda_. 2 * x" "\lambda_.
-1"
  \langle proof \rangle
fun cheb_eval_aux where
  "cheb_eval_aux 0 cs x acc1 acc2 = hd cs + x * acc1 - acc2"
| "cheb_eval_aux (Suc n) cs x acc1 acc2 =
     cheb_eval_aux n (tl cs) x (hd cs + 2 * x * acc1 - acc2) acc1"
lemma cheb_eval_aux_altdef:
  "length cs = Suc n \Longrightarrow
      cheb_eval_aux n cs x acc1 acc2 =
     cheb_poly.clenshaw_aux' x (\lambda k. rev cs ! k) n acc1 acc2"
\langle proof \rangle
lemmas [simp del] = cheb_eval_aux.simps
lemma cheb_eval_code [code]:
  "cheb_eval [] x = 0"
  "cheb_eval [c] x = c"
  "cheb_eval (c1 \# c2 \# cs) x =
     cheb_eval_aux (Suc (length cs)) (rev (c1 # c2 # cs)) x 0 0"
\langle proof \rangle
end
```

References

[1] J. Mason and D. Handscomb. *Chebyshev Polynomials*. CRC Press, 2002.