Cardinality of Multisets

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Abstract

This entry provides three lemmas to count the number of multisets of a given size and finite carrier set. The first lemma provides a cardinality formula assuming that the multiset's elements are chosen from the given carrier set. The latter two lemmas provide formulas assuming that the multiset's elements also cover the given carrier set, i.e., each element of the carrier set occurs in the multiset at least once.

The proof of the first lemma uses the argument of the recurrence relation for counting multisets [1]. The proof of the second lemma is straightforward, and the proof of the third lemma is easily obtained using the first cardinality lemma. A challenge for the formalization is the derivation of the required induction rule, which is a special combination of the induction rules for finite sets and natural numbers. The induction rule is derived by defining a suitable inductive predicate and transforming the predicate's induction rule.

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1 Cardinality of Multisets

theory Card-Multisets imports HOL-Library.Multiset begin

1.1 Additions to Multiset Theory

lemma *mset-set-set-mset-subseteq*:

```
mset\text{-}set (set\text{-}mset M) \subseteq \# M
proof (induct M)
 case empty
 show ?case by simp
next
  case (add \ x \ M)
 from this show ?case
 proof (cases x \in \# M)
   assume x \in \# M
   from this have mset-set (set-mset (M + \{\#x\#\})) = mset-set (set-mset M)
     by (simp add: insert-absorb)
   from this add.hyps show ?thesis
     using subset-mset.trans by fastforce
 \mathbf{next}
   assume \neg x \in \# M
   from this add.hyps have \{\#x\#\} + mset\text{-set }(set\text{-}mset M) \subseteq \#M + \{\#x\#\}
     by (simp add: insert-subset-eq-iff)
   from this \langle \neg x \in \# M \rangle show ?thesis by simp
 qed
qed
lemma size-mset-set-eq-card:
 assumes finite A
```

```
shows size (mset-set A) = card A
using assms by (induct A) auto
```

lemma card-set-mset-leq: card (set-mset M) \leq size M**by** (induct M) (auto simp add: card-insert-le-m1)

1.2 Lemma to Enumerate Sets of Multisets

lemma *set-of-multisets-eq*: assumes $x \notin A$ shows $\{M. \text{ set-mset } M \subseteq \text{ insert } x \land A \land \text{ size } M = Suc \ k\} =$ $\{M. \text{ set-mset } M \subseteq A \land \text{ size } M = Suc \ k\} \cup$ $(\lambda M. M + \{\#x\#\})$ ' $\{M. set\text{-mset } M \subseteq insert \ x \ A \land size \ M = k\}$ proof – **from** $\langle x \notin A \rangle$ have $\{M. set\text{-mset } M \subseteq insert \ x \ A \land size \ M = Suc \ k\} =$ $\{M. \text{ set-mset } M \subseteq A \land \text{ size } M = Suc \ k\} \cup$ $\{M. \text{ set-mset } M \subseteq \text{ insert } x \land A \land \text{ size } M = Suc \land X \in \# M\}$ by *auto* **moreover have** $\{M. \text{ set-mset } M \subseteq \text{ insert } x \land A \land \text{ size } M = Suc \land A \land \# M\} =$ $(\lambda M. M + \{\#x\#\})$ ' $\{M. set\text{-mset } M \subseteq insert \ x \ A \land size \ M = k\}$ (is ?S =?T) proof show $?S \subseteq ?T$ proof fix M

assume $M \in ?S$ from this have $M = M - \{\#x\#\} + \{\#x\#\}$ by *auto* **moreover have** $M - \{\#x\#\} \in \{M. \text{ set-mset } M \subseteq \text{ insert } x \land A \land \text{ size } M = k\}$ proof – have set-mset $(M - \{\#x\#\} + \{\#x\#\}) \subseteq insert \ x \ A$ using $\langle M \in ?S \rangle$ by force moreover have size $(M - \{\#x\#\} + \{\#x\#\}) = Suc \ k \land x \in \# M \{\#x\#\} + \{\#x\#\}$ using $\langle M \in ?S \rangle$ by force ultimately show ?thesis by force qed ultimately show $M \in ?T$ by *auto* qed \mathbf{next} show $?T \subseteq ?S$ by force qed ultimately show ?thesis by auto qed

1.3 Derivation of Suitable Induction Rule

```
context
begin
private inductive R :: 'a \ set \Rightarrow nat \Rightarrow bool
where
 finite A \Longrightarrow R A \theta
\mid R \mid k
| finite A \Longrightarrow x \notin A \Longrightarrow R A (Suc k) \Longrightarrow R (insert x A) k \Longrightarrow R (insert x A)
(Suc \ k)
private lemma R-eq-finite:
  R \ A \ k \longleftrightarrow finite \ A
proof
 assume R A k
 from this show finite A by cases auto
\mathbf{next}
  assume finite A
  from this show R A k
  proof (induct A)
   case empty
   from this show ?case by (rule R.intros(2))
  \mathbf{next}
   case insert
   from this show ?case
   proof (induct k)
     case \theta
     from this show ?case
       by (intro R.intros(1) finite.insertI)
```

```
next
case Suc
from this show ?case
by (metis R.simps Zero-neq-Suc diff-Suc-1)
qed
qed
qed
```

```
lemma finite-set-and-nat-induct[consumes 1, case-names zero empty step]:

assumes finite A

assumes \bigwedge A. finite A \Longrightarrow P \land 0

assumes \bigwedge k. P \ \} k

assumes \bigwedge A k x. finite A \Longrightarrow x \notin A \Longrightarrow P \land (Suc \ k) \Longrightarrow P (insert x \land k \land k \Longrightarrow)

P (insert x \land A) (Suc k)

shows P \land k

proof -

from \langle finite \ A \rangle have R \land k by (subst R-eq-finite)

from this assms(2-4) show ?thesis by (induct \land k) auto

qed
```

end

1.4 Finiteness of Sets of Multisets

```
lemma finite-multisets:

assumes finite A

shows finite \{M. set\text{-mset } M \subseteq A \land size M = k\}

using assms

proof (induct A \ k rule: finite-set-and-nat-induct)

case zero

from this show ?case by auto

next

case empty

from this show ?case by auto

next

case (step A \ k x)

from this show ?case

using set-of-multisets-eq[OF \ \langle x \notin A \rangle] by simp

qed
```

1.5 Cardinality of Multisets

lemma card-multisets: **assumes** finite A **shows** card $\{M. set-mset M \subseteq A \land size M = k\} = (card A + k - 1)$ choose k **using** assms **proof** (induct A k rule: finite-set-and-nat-induct) **case** (zero A) **assume** finite (A :: 'a set) **have** $\{M. set-mset M \subseteq A \land size M = 0\} = \{\{\#\}\}$ by auto

from this show card $\{M. \text{ set-mset } M \subseteq A \land \text{ size } M = 0\} = card A + 0 - 1$ choose 0by simp \mathbf{next} **case** (empty k)show card {M. set-mset $M \subseteq \{\} \land size M = k\} = card \{\} + k - 1 choose k$ by (cases k) (auto simp add: binomial-eq-0) \mathbf{next} **case** (step A k x) let $?S_1 = \{M. \text{ set-mset } M \subseteq A \land \text{ size } M = Suc \ k\}$ and $?S_2 = \{M. \text{ set-mset } M \subseteq \text{ insert } x \land A \land \text{ size } M = k\}$ **assume** hyps1: card $?S_1 = card A + Suc k - 1$ choose Suc k **assume** hyps2: card $?S_2 = card$ (insert x A) + k - 1 choose khave finite-sets: finite $?S_1$ finite $((\lambda M. M + \{\#x\#\}) `?S_2)$ **using** $\langle finite | A \rangle$ by (auto simp add: finite-multisets) have inj: inj-on $(\lambda M. M + \{\#x\#\})$?S₂ by (rule inj-onI) auto have card $\{M. \text{ set-mset } M \subseteq \text{ insert } x \land A \land \text{ size } M = Suc \ k\} =$ card $(?S_1 \cup (\lambda M. M + \{\#x\#\}) `?S_2)$ using set-of-multisets-eq $\langle x \notin A \rangle$ by fastforce **also have** ... = card $?S_1$ + card (($\lambda M. M + \{\#x\#\}$) ' $?S_2$) using finite-sets $\langle x \notin A \rangle$ by (subst card-Un-disjoint) auto also have $\ldots = card ?S_1 + card ?S_2$ using *inj* by (*auto intro: card-image*) also have $\ldots = (card A + Suc k - 1 choose Suc k) + (card (insert x A) + k - 1)$ 1 choose k) using hyps1 hyps2 by simp also have $\ldots = card$ (insert x A) + Suc k - 1 choose Suc kusing $\langle x \notin A \rangle$ (finite $A \rangle$ by simp finally show ?case . qed **lemma** card-too-small-multisets-covering-set: assumes finite A assumes k < card Ashows card $\{M. \text{ set-mset } M = A \land \text{ size } M = k\} = 0$ proof from $\langle k < card A \rangle$ have eq: $\{M. set-mset M = A \land size M = k\} = \{\}$ using card-set-mset-leq Collect-empty-eq leD by auto **from** this **show** ?thesis **by** (metis card.empty) qed **lemma** card-multisets-covering-set: assumes finite A assumes card $A \leq k$ shows card $\{M. \text{ set-mset } M = A \land \text{ size } M = k\} = (k - 1) \text{ choose } (k - \text{ card } A)$ proof – have $\{M. \text{ set-mset } M = A \land \text{ size } M = k\} = (\lambda M. M + \text{ mset-set } A)$ $\{M. \text{ set-mset } M \subseteq A \land \text{ size } M = k - \text{ card } A\}$ (is ?S = ?f `?T) proof

```
show ?S \subseteq ?f `?T
   proof
     fix M
     assume M \in ?S
     from this have M = M - mset\text{-set } A + mset\text{-set } A
      by (auto simp add: mset-set-mset-subseteq subset-mset.diff-add)
     moreover from \langle M \in ?S \rangle have M - mset\text{-set } A \in ?T
    by (auto simp add: mset-set-mset-subseteq size-Diff-submset size-mset-set-eq-card
in-diffD)
     ultimately show M \in ?f '? T by auto
   qed
 \mathbf{next}
   from (finite A) (card A \leq k) show ?f (?T \subseteq ?S
     by (auto simp add: size-mset-set-eq-card)+
 qed
 moreover have inj-on ?f ?T by (rule inj-onI) auto
 ultimately have card ?S = card ?T by (simp add: card-image)
 also have \ldots = card A + (k - card A) - 1 choose (k - card A)
   using (finite A) by (simp only: card-multisets)
 also have \ldots = (k - 1) choose (k - card A)
   using \langle card \ A \leq k \rangle by auto
 finally show ?thesis .
qed
```

 \mathbf{end}

References

[1] Wikipedia. Multiset — wikipedia, the free encyclopedia, 2016. [Online; accessed 23-June-2016].