The Factorization Algorithm of Berlekamp and Zassenhaus *

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Abstract

We formalize the Berlekamp-Zassenhaus algorithm for factoring square-free integer polynomials in Isabelle/HOL. We further adapt an existing formalization of Yun's square-free factorization algorithm to integer polynomials, and thus provide an efficient and certified factorization algorithm for arbitrary univariate polynomials.

The algorithm first performs a factorization in the prime field GF(p) and then performs computations in the integer ring modulo p^k, where both p and k are determined at runtime. Since a natural modeling of these structures via dependent types is not possible in Isabelle/HOL, we formalize the whole algorithm using Isabelle's recent addition of local type definitions.

Through experiments we verify that our algorithm factors polynomials of degree 100 within seconds.

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1 Introduction

Modern algorithms to factor integer polynomials – following Berlekamp and Zassenhaus – work via polynomial factorization over prime fields GF(p) and quotient rings $\mathbb{Z}/p^k\mathbb{Z}$ [2, 3]. Algorithm 1 illustrates the basic structure of such an algorithm.[1]

Algorithm 1: A modern factorization algorithm

Input: Square-free integer polynomial f.
Output: Irreducible factors $f_1, \ldots, f_n$ such that $f = f_1 \cdot \ldots \cdot f_n$.

4 Choose a suitable prime p depending on f.
5 Factor f in GF(p): $f \equiv g_1 \cdot \ldots \cdot g_m \pmod{p}$.
6 Determine a suitable bound d on the degree, depending on $g_1, \ldots, g_m$. Choose an exponent k such that every coefficient of a factor of a given multiple of f in $\mathbb{Z}$ with degree at most d can be uniquely represent by a number below p^k.
7 From step 5 compute the unique factorization $f \equiv h_1 \cdot \ldots \cdot h_m \pmod{p^k}$ via the Hensel lifting.
8 Construct a factorization $f = f_1 \cdot \ldots \cdot f_n$ over the integers where each f_i corresponds to the product of one or more h_j.

In previous work on algebraic numbers [12], we implemented Algorithm 1 in Isabelle/HOL [11] as a function of type *int poly* ⇒ *int poly list*, where we chose Berlekamp's algorithm in step 5. However, the algorithm was available only as an oracle, and thus a validity check on the result factorization had to be performed.

In this work we fully formalize the correctness of our implementation.

[1] Our algorithm starts with step 4, so that section numbers and step-numbers coincide.

Theorem 1 (Berlekamp-Zassenhaus' Algorithm)

assumes *square_free* (f :: *int poly*)
  and *degree* $f \neq 0$
  and *berlekamp_zassenhaus_factorization* $f = fs$
shows f = *prod_list fs*
  and $\forall f_i \in$ *set fs*. *irreducible* f_i

To obtain Theorem 1 we perform the following tasks.

- We introduce two formulations of GF(p) and $\mathbb{Z}/p^k\mathbb{Z}$. We first define a type to represent these domains, employing ideas from HOL multivariate analysis. This is essential for reusing many type-based algorithms from the Isabelle distribution and the AFP (archive of formal proofs). At some points in our developement, the type-based setting is still too restrictive. Hence we also introduce a second formulation which is *locale-based*.

- The prime p in step 4 must be chosen so that f remains square-free in GF(p). For the termination of the algorithm, we prove that such a prime always exists.

- We explain Berlekamp's algorithm that factors polynomials over prime fields, and formalize its correctness using the type-based representation. Since Isabelle's code generation does not work for the type-based representation of prime fields, we define an implementation of Berlekamp's algorithm which avoids type-based polynomial algorithms and type-based prime fields. The soundness of this implementation is proved via the transfer package [5]: we transform the type-based soundness statement of Berlekamp's algorithm into a statement which speaks solely about integer polynomials. Here, we crucially rely upon local type definitions [9] to eliminate the presence of the type for the prime field GF(p).

- For step 6 we need to find a bound on the coefficients of the factors of a polynomial. For this purpose, we formalize Mignotte's factor bound. During this formalization task we detected a bug in our previous oracle implementation, which computed improper bounds on the degrees of factors.

- We formalize the Hensel lifting. As for Berlekamp's algorithm, we first formalize basic operations in the type-based setting. Unfortunately, however, this result cannot be extended to the full Hensel lifting. Therefore, we model the Hensel lifting in a locale-based way so that modulo operation is explicitly applied on polynomials.

- For the reconstruction in step 8 we closely follow the description of Knuth [7, page 452]. Here, we use the same representation of polynomials over $\mathbb{Z}/p^k\mathbb{Z}$ as for the Hensel lifting.
- We adapt an existing square-free factorization algorithm from $\mathbb{Q}$ to $\mathbb{Z}$. In combination with the previous results this leads to a factorization algorithm for arbitrary integer and rational polynomials.

To our knowledge, this is the first formalization of the Berlekamp-Zassenhaus algorithm. For instance, Barthe et al. report that there is no formalization of an efficient factorization algorithm over GF(p) available in Coq [1, Section 6, note 3 on formalization].

Some key theorems leading to the algorithm have already been formalized in Isabelle or other proof assistants. In ACL2, for instance, polynomials over a field are shown to be a unique factorization domain (UFD) [4]. A more general result, namely that polynomials over UFD are also UFD, was already developed in Isabelle/HOL for implementing algebraic numbers [12] and an independent development by Eberl is now available in the Isabelle distribution.

An Isabelle formalization of Hensel's lemma is provided by Kobayashi et al. [8], who defined the valuations of polynomials via Cauchy sequences, and used this setup to prove the lemma. Consequently, their result requires a 'valuation ring' as precondition in their formalization. While this extra precondition is theoretically met in our setting, we did not attempt to reuse their results, because the type of polynomials in their formalization (from HOL-Algebra) differs from the polynomials in our development (from HOL/Library). Instead, we formalize a direct proof for Hensel's lemma. Our formalizations are incomparable: On the one hand, Kobayashi et al. did not consider only integer polynomials as we do. On the other hand, we additionally formalize the quadratic Hensel lifting [13], extend the lifting from binary to n-ary factorizations, and prove a uniqueness result, which is required for proving the soundness of Theorem 1.

A Coq formalization of Hensel's lemma is also available, which is used for certifying integral roots and 'hardest-to-round computation' [10]. If one is interested in certifying a factorization, rather than a certified algorithm that performs it, it suffices to test that all the found factors are irreducible. Kirkels [6] formalized a sufficient criterion for this test in Coq: when a polynomial is irreducible modulo some prime, it is also irreducible in $\mathbb{Z}$. Both formalizations are in Coq, and we did not attempt to reuse them.

2 Finite Rings and Fields

We start by establishing some preliminary results about finite rings and finite fields

2.1 Finite Rings

theory *Finite-Field*
imports
HOL−Computational-Algebra.Primes
HOL−Number-Theory.Residues
HOL−Library.Cardinality
Subresultants.Binary-Exponentiation
Polynomial-Interpolation.Ring-Hom-Poly
begin

typedef (*'a*::*finite*) *mod-ring* = {*0*..<*int CARD*(*'a*)} **by** *auto*

setup-lifting *type-definition-mod-ring*

lemma *CARD-mod-ring*[*simp*]: *CARD*(*'a mod-ring*) = *CARD*(*'a*::*finite*)
proof −
have *card* {*y*. ∃ *x*∈{*0*..<*int CARD*(*'a*)}. (*y*::*'a mod-ring*) = *Abs-mod-ring x*} = *card* {*0*..<*int CARD*(*'a*)}
proof (*rule bij-betw-same-card*)
have *inj-on Rep-mod-ring* {*y*. ∃ *x*∈{*0*..<*int CARD*(*'a*)}. *y* = *Abs-mod-ring x*}
by (*meson Rep-mod-ring-inject inj-onI*)
moreover have *Rep-mod-ring* ' {*y*. ∃ *x*∈{*0*..<*int CARD*(*'a*)}. (*y*::*'a mod-ring*) = *Abs-mod-ring x*} = {*0*..<*int CARD*(*'a*)}
proof (*auto simp add*: *image-def Rep-mod-ring-inject*)
fix *xb* **show** *0* ≤ *Rep-mod-ring* (*Abs-mod-ring xb*)
using *Rep-mod-ring atLeastLessThan-iff* **by** *blast*
assume *xb1*: *0* ≤ *xb* **and** *xb2*: *xb* < *int CARD*(*'a*)
thus *Rep-mod-ring* (*Abs-mod-ring xb*) < *int CARD*(*'a*)
by (*metis Abs-mod-ring-inverse Rep-mod-ring atLeastLessThan-iff le-less-trans linear*)
have *xb*: *xb* ∈ {*0*..<*int CARD*(*'a*)} **using** *xb1 xb2* **by** *simp*
show ∃ *xa*::*'a mod-ring*. (∃ *x*∈{*0*..<*int CARD*(*'a*)}. *xa* = *Abs-mod-ring x*) ∧ *xb* = *Rep-mod-ring xa*
by (*rule exI*[*of* - *Abs-mod-ring xb*], *auto simp add*: *xb1 xb2*, *rule Abs-mod-ring-inverse*[*OF xb*, *symmetric*])
qed
ultimately show *bij-betw Rep-mod-ring*
{*y*. ∃ *x*∈{*0*..<*int CARD*(*'a*)}. (*y*::*'a mod-ring*) = *Abs-mod-ring x*}
{*0*..<*int CARD*(*'a*)}
by (*simp add*: *bij-betw-def*)
qed
thus *?thesis*
unfolding *type-definition.univ*[*OF type-definition-mod-ring*]
unfolding *image-def* **by** *auto*
qed

instance *mod-ring* :: (*finite*) *finite*
proof (*intro-classes*)
show *finite* (*UNIV*::*'a mod-ring set*)

unfolding *type-definition.univ*[*OF type-definition-mod-ring*]
using *finite* **by** *simp*
qed

instantiation *mod-ring* :: (*finite*) *equal*
begin
lift-definition *equal-mod-ring* :: *'a mod-ring* ⇒ *'a mod-ring* ⇒ *bool* **is** (=) **.**
instance by (*intro-classes*, *transfer*, *auto*)
end

instantiation *mod-ring* :: (*finite*) *comm-ring*
begin

lift-definition *plus-mod-ring* :: *'a mod-ring* ⇒ *'a mod-ring* ⇒ *'a mod-ring* **is**
λ *x y*. (*x* + *y*) *mod int* (*CARD*(*'a*)) **by** *simp*

lift-definition *uminus-mod-ring* :: *'a mod-ring* ⇒ *'a mod-ring* **is**
λ *x*. *if x* = *0 then 0 else int* (*CARD*(*'a*)) − *x* **by** *simp*

lift-definition *minus-mod-ring* :: *'a mod-ring* ⇒ *'a mod-ring* ⇒ *'a mod-ring* **is**
λ *x y*. (*x* − *y*) *mod int* (*CARD*(*'a*)) **by** *simp*

lift-definition *times-mod-ring* :: *'a mod-ring* ⇒ *'a mod-ring* ⇒ *'a mod-ring* **is**
λ *x y*. (*x* ∗ *y*) *mod int* (*CARD*(*'a*)) **by** *simp*

lift-definition *zero-mod-ring* :: *'a mod-ring* **is** *0* **by** *simp*

instance
by *standard*
(*transfer*; *auto simp add*: *mod-simps algebra-simps intro*: *mod-diff-cong*)+

end

lift-definition *to-int-mod-ring* :: *'a*::*finite mod-ring* ⇒ *int* **is** λ *x*. *x* **.**

lift-definition *of-int-mod-ring* :: *int* ⇒ *'a*::*finite mod-ring* **is**
λ *x*. *x mod int* (*CARD*(*'a*)) **by** *simp*

interpretation *to-int-mod-ring-hom*: *inj-zero-hom to-int-mod-ring*
by (*unfold-locales*; *transfer*, *auto*)

lemma *int-nat-card*[*simp*]: *int* (*nat CARD*(*'a*::*finite*)) = *CARD*(*'a*) **by** *auto*

interpretation *of-int-mod-ring-hom*: *zero-hom of-int-mod-ring*
by (*unfold-locales*, *transfer*, *auto*)

lemma *of-int-mod-ring-to-int-mod-ring*[*simp*]:
of-int-mod-ring (*to-int-mod-ring x*) = *x* **by** (*transfer*, *auto*)

lemma *to-int-mod-ring-of-int-mod-ring*[*simp*]: $0 \leq x \Longrightarrow x < int\ CARD('a :: finite) \Longrightarrow$
  to-int-mod-ring (*of-int-mod-ring x* :: *'a mod-ring*) = *x*
  by (*transfer*, *auto*)

lemma *range-to-int-mod-ring*:
  range (*to-int-mod-ring* :: (*'a* :: *finite mod-ring* ⇒ *int*)) = {*0* ..< *CARD*(*'a*)}
  apply (*intro equalityI subsetI*)
  apply (*elim rangeE*, *transfer*, *force*)
  by (*auto intro*!: *range-eqI to-int-mod-ring-of-int-mod-ring*[*symmetric*])

2.2 Nontrivial Finite Rings

class *nontriv* = **assumes** *nontriv*: *CARD*(*'a*) > *1*

subclass(**in** *nontriv*) *finite* **by**(*intro-classes*,*insert nontriv*,*auto intro*:*card-ge-0-finite*)

instantiation *mod-ring* :: (*nontriv*) *comm-ring-1*
begin

lift-definition *one-mod-ring* :: *'a mod-ring* **is** *1* **using** *nontriv*[**where** *?'a*=*'a*] **by** *auto*

instance by (*intro-classes*; *transfer*, *simp*)

end

interpretation *to-int-mod-ring-hom*: *inj-one-hom to-int-mod-ring*
  by (*unfold-locales*, *transfer*, *simp*)

lemma *of-nat-of-int-mod-ring* [*code-unfold*]:
  of-nat = *of-int-mod-ring o int*
proof (*rule ext*, *unfold o-def*)
  show *of-nat n* = *of-int-mod-ring* (*int n*) **for** *n*
  proof (*induct n*)
    case (*Suc n*)
    show *?case*
      by (*simp only*: *of-nat-Suc Suc*, *transfer*) (*simp add*: *mod-simps*)
  qed *simp*
qed

lemma *of-nat-card-eq-0*[*simp*]: (*of-nat* (*CARD*(*'a*::*nontriv*)) :: *'a mod-ring*) = *0*
  by (*unfold of-nat-of-int-mod-ring*, *transfer*, *auto*)

lemma *of-int-of-int-mod-ring*[*code-unfold*]: *of-int* = *of-int-mod-ring*
proof (*rule ext*)
  fix *x* :: *int*
  obtain *n1 n2* **where** *x*: *x* = *int n1* − *int n2* **by** (*rule int-diff-cases*)

show *of-int x = of-int-mod-ring x*
unfolding *x of-int-diff of-int-of-nat-eq of-nat-of-int-mod-ring o-def*
by (*transfer*, *simp add*: *mod-diff-right-eq mod-diff-left-eq*)
qed

unbundle *lifting-syntax*

lemma *pcr-mod-ring-to-int-mod-ring*: *pcr-mod-ring* = (λ*x y*. *x* = *to-int-mod-ring y*)
unfolding *mod-ring.pcr-cr-eq* **unfolding** *cr-mod-ring-def to-int-mod-ring.rep-eq*
..

lemma [*transfer-rule*]:
((=) ===> *pcr-mod-ring*) (λ *x*. *int x mod int* (*CARD*('*a* :: *nontriv*))) (*of-nat* :: *nat* ⇒ '*a mod-ring*)
by (*intro rel-funI*, *unfold pcr-mod-ring-to-int-mod-ring of-nat-of-int-mod-ring*, *transfer*, *auto*)

lemma [*transfer-rule*]:
((=) ===> *pcr-mod-ring*) (λ *x*. *x mod int* (*CARD*('*a* :: *nontriv*))) (*of-int* :: *int* ⇒ '*a mod-ring*)
by (*intro rel-funI*, *unfold pcr-mod-ring-to-int-mod-ring of-int-of-int-mod-ring*, *transfer*, *auto*)

lemma *one-mod-card* [*simp*]: *1 mod CARD*('*a*::*nontriv*) = *1*
using *mod-less nontriv* **by** *blast*

lemma *Suc-0-mod-card* [*simp*]: *Suc 0 mod CARD*('*a*::*nontriv*) = *1*
using *one-mod-card* **by** *simp*

lemma *one-mod-card-int* [*simp*]: *1 mod int CARD*('*a*::*nontriv*) = *1*
proof −
from *nontriv* [**where** *?'a* = '*a*] **have** *int* (*1 mod CARD*('*a*::*nontriv*)) = *1*
by *simp*
then show *?thesis*
using *of-nat-mod* [*of 1 CARD*('*a*), **where** *?'a* = *int*] **by** *simp*
qed

lemma *pow-mod-ring-transfer*[*transfer-rule*]:
(*pcr-mod-ring* ===> (=) ===> *pcr-mod-ring*)
(λ*a*::*int*. λ*n*. *a*^*n mod CARD*('*a*::*nontriv*)) ((^)::'*a mod-ring* ⇒ *nat* ⇒ '*a mod-ring*)
unfolding *pcr-mod-ring-to-int-mod-ring*
proof (*intro rel-funI*,*simp*)
fix *x*::'*a mod-ring* **and** *n*
show *to-int-mod-ring x* ^ *n mod int CARD*('*a*) = *to-int-mod-ring* (*x* ^ *n*)
proof (*induct n*)
case *0*
thus *?case* **by** *auto*
next

```
    case (Suc n)
    have to-int-mod-ring (x ^ Suc n) = to-int-mod-ring (x * x ^ n) by auto
    also have ... = to-int-mod-ring x * to-int-mod-ring (x ^ n) mod CARD('a)
      unfolding to-int-mod-ring-def using times-mod-ring.rep-eq by auto
   also have ... = to-int-mod-ring x * (to-int-mod-ring x ^ n mod CARD('a)) mod
CARD('a)
      using Suc.hyps by auto
    also have ... = to-int-mod-ring x ^ Suc n mod int CARD('a)
      by (simp add: mod-simps)
    finally show ?case ..
  qed
qed

lemma dvd-mod-ring-transfer[transfer-rule]:
((pcr-mod-ring :: int ⇒ 'a :: nontriv mod-ring ⇒ bool) ===>
  (pcr-mod-ring :: int ⇒ 'a mod-ring ⇒ bool) ===> (=))
  (λ i j. ∃ k ∈ {0..<int CARD('a)}. j = i * k mod int CARD('a)) (dvd)
proof (unfold pcr-mod-ring-to-int-mod-ring, intro rel-funI iffI)
  fix x y :: 'a mod-ring and i j
  assume i: i = to-int-mod-ring x and j: j = to-int-mod-ring y
  { assume x dvd y
    then obtain z where y = x * z by (elim dvdE, auto)
    then have j = i * to-int-mod-ring z mod CARD('a) by (unfold i j, transfer)
    with range-to-int-mod-ring
    show ∃ k ∈ {0..<int CARD('a)}. j = i * k mod CARD('a) by auto
  }
  assume ∃ k ∈ {0..<int CARD('a)}. j = i * k mod CARD('a)
   then obtain k where k: k ∈ {0..<int CARD('a)} and dvd: j = i * k mod
CARD('a) by auto
  from k have to-int-mod-ring (of-int k :: 'a mod-ring) = k by (transfer, auto)
  also from dvd have j = i * ... mod CARD('a) by auto
 finally have y = x * (of-int k :: 'a mod-ring) unfolding i j using k by (transfer,
auto)
  then show x dvd y by auto
qed

lemma Rep-mod-ring-mod[simp]: Rep-mod-ring (a :: 'a :: nontriv mod-ring) mod
CARD('a) = Rep-mod-ring a
  using Rep-mod-ring[where 'a = 'a] by auto
```

2.3 Finite Fields

When the domain is prime, the ring becomes a field

```
class prime-card = assumes prime-card: prime (CARD('a))
begin
lemma prime-card-int: prime (int (CARD('a))) using prime-card by auto

subclass nontriv using prime-card prime-gt-1-nat by (intro-classes,auto)
end
```

```
instance bool :: prime-card
  by standard auto

instantiation mod-ring :: (prime-card) field
begin

definition inverse-mod-ring :: 'a mod-ring ⇒ 'a mod-ring where
  inverse-mod-ring x = (if x = 0 then 0 else x ^ (nat (CARD('a) - 2)))

definition divide-mod-ring :: 'a mod-ring ⇒ 'a mod-ring ⇒ 'a mod-ring  where
  divide-mod-ring x y = x * ((λc. if c = 0 then 0 else c ^ (nat (CARD('a) - 2)))
y)

instance
proof
  fix a b c::'a mod-ring
  show inverse 0 = (0::'a mod-ring) by (simp add: inverse-mod-ring-def)
  show a div b = a * inverse b
    unfolding inverse-mod-ring-def by (transfer', simp add: divide-mod-ring-def)
  show a ≠ 0 ⟹ inverse a * a = 1
  proof (unfold inverse-mod-ring-def, transfer)
    let ?p=CARD('a)
    fix x
    assume x: x ∈ {0..<int CARD('a)} and x0: x ≠ 0
    have p0': 0≤?p by auto
    have ¬ ?p dvd x
      using x x0 zdvd-imp-le by fastforce
    then have ¬ CARD('a) dvd nat |x|
      by simp
    with x have ¬ CARD('a) dvd nat x
      by simp
    have rw: x ^ nat (int (?p - 2)) * x = x ^ nat (?p - 1)
    proof -
      have p2: 0 ≤ int (?p-2) using x by simp
      have card-rw: (CARD('a) - Suc 0) = nat (1 + int (CARD('a) - 2))
        using nat-eq-iff x x0 by auto
      have x ^ nat (?p - 2)*x = x ^ (Suc (nat (?p - 2))) by simp
      also have ... = x ^ (nat (?p - 1))
        using Suc-nat-eq-nat-zadd1[OF p2] card-rw by auto
      finally show ?thesis .
    qed
    have [int (nat x ^ (CARD('a) - 1)) = int 1] (mod CARD('a))
      using fermat-theorem [OF prime-card ‹¬ CARD('a) dvd nat x›]
      by (simp only: cong-def cong-def of-nat-mod [symmetric])
    then have *: [x ^ (CARD('a) - 1) = 1] (mod CARD('a))
      using x by auto
    have x ^ (CARD('a) - 2) mod CARD('a) * x mod CARD('a)
      = (x ^ nat (CARD('a) - 2) * x) mod CARD('a) by (simp add: mod-simps)
```

```
  also have ... = (x ^ nat (?p - 1) mod ?p) unfolding rw by simp
  also have ... = (x ^(nat ?p - 1) mod ?p) using p0' by (simp add: nat-diff-distrib')
  also have ... = 1
    using * by (simp add: cong-def)
  finally show (if x = 0 then 0 else x ^nat (int (CARD('a) - 2)) mod CARD('a))
* x mod CARD('a) = 1
    using x0 by auto
 qed
qed
end

instantiation mod-ring :: (prime-card) {normalization-euclidean-semiring, euclidean-ring}
begin

definition modulo-mod-ring :: 'a mod-ring ⇒ 'a mod-ring ⇒ 'a mod-ring where
modulo-mod-ring x y = (if y = 0 then x else 0)
definition normalize-mod-ring :: 'a mod-ring ⇒ 'a mod-ring where normalize-mod-ring
x = (if x = 0 then 0 else 1)
definition unit-factor-mod-ring :: 'a mod-ring ⇒ 'a mod-ring where unit-factor-mod-ring
x = x
definition euclidean-size-mod-ring :: 'a mod-ring ⇒ nat where euclidean-size-mod-ring
x = (if x = 0 then 0 else 1)

instance
proof (intro-classes)
 fix a :: 'a mod-ring show a ≠ 0 ⟹ unit-factor a dvd 1
   unfolding dvd-def unit-factor-mod-ring-def by (intro exI[of - inverse a], auto)
qed (auto simp: normalize-mod-ring-def unit-factor-mod-ring-def modulo-mod-ring-def
    euclidean-size-mod-ring-def field-simps)
end

instantiation mod-ring :: (prime-card) euclidean-ring-gcd
begin

definition gcd-mod-ring :: 'a mod-ring ⇒ 'a mod-ring ⇒ 'a mod-ring where
gcd-mod-ring = Euclidean-Algorithm.gcd
definition lcm-mod-ring :: 'a mod-ring ⇒ 'a mod-ring ⇒ 'a mod-ring where
lcm-mod-ring = Euclidean-Algorithm.lcm
definition Gcd-mod-ring :: 'a mod-ring set ⇒ 'a mod-ring where Gcd-mod-ring
= Euclidean-Algorithm.Gcd
definition Lcm-mod-ring :: 'a mod-ring set ⇒ 'a mod-ring where Lcm-mod-ring
= Euclidean-Algorithm.Lcm

instance by (intro-classes, auto simp: gcd-mod-ring-def lcm-mod-ring-def Gcd-mod-ring-def
Lcm-mod-ring-def)
end

instantiation mod-ring :: (prime-card) unique-euclidean-ring
begin
```

definition [*simp*]: *division-segment-mod-ring* (*x* :: *'a mod-ring*) = (*1* :: *'a mod-ring*)

instance by *intro-classes* (*auto simp*: *euclidean-size-mod-ring-def split*: *if-splits*)

end

instance *mod-ring* :: (*prime-card*) *field-gcd*
by *intro-classes auto*

lemma *surj-of-nat-mod-ring*: ∃ *i*. *i* < *CARD*(*'a* :: *prime-card*) ∧ (*x* :: *'a mod-ring*)
= *of-nat i*
by (*rule exI*[*of - nat* (*to-int-mod-ring x*)], *unfold of-nat-of-int-mod-ring o-def*,
subst nat-0-le, *transfer*, *simp*, *simp*, *transfer*, *auto*)

lemma *of-nat-0-mod-ring-dvd*: **assumes** *x*: *of-nat x* = (*0* :: *'a* ::*prime-card mod-ring*)
shows *CARD*(*'a*) *dvd x*
proof −
let *?x* = *of-nat x* :: *int*
from *x* **have** *of-int-mod-ring ?x* = (*0* :: *'a mod-ring*) **by** (*fold of-int-of-int-mod-ring*,
simp)
hence *?x mod CARD*(*'a*) = *0* **by** (*transfer*, *auto*)
hence *x mod CARD*(*'a*) = *0* **by** *presburger*
thus *?thesis* **unfolding** *mod-eq-0-iff-dvd* **.**
qed

lemma *semiring-char-mod-ring* [*simp*]:
CHAR(*'n* :: *nontriv mod-ring*) = *CARD*(*'n*)
proof (*rule CHAR-eq-posI*)
fix *x* **assume** *x* > *0 x* < *CARD*(*'n*)
thus *of-nat x* ≠ (*0* :: *'n mod-ring*)
by *transfer auto*
qed *auto*

The following Material was contributed by Manuel Eberl

instance *mod-ring* :: (*prime-card*) *finite-field*
by *standard simp-all*

instantiation *mod-ring* :: (*prime-card*) *enum-finite-field*
begin

definition *enum-finite-field-mod-ring* :: *nat* ⇒ *'a mod-ring* **where**
enum-finite-field-mod-ring n = *of-int-mod-ring* (*int n*)

instance proof
interpret *type-definition Rep-mod-ring* :: *'a mod-ring* ⇒ *int Abs-mod-ring* {*0*..<*CARD*(*'a*)}
by (*rule type-definition-mod-ring*)
have *enum-finite-field* ‘ {..<*CARD*(*'a mod-ring*)} = *of-int-mod-ring* ‘ *int* ‘ {..<*CARD*(*'a*

```
mod-ring)}
    unfolding enum-finite-field-mod-ring-def by (simp add: image-image o-def)
  also have int ` {..<CARD('a mod-ring)} = {0..<int CARD('a mod-ring)}
    by (simp add: image-atLeastZeroLessThan-int)
  also have of-int-mod-ring ` ... = (Abs-mod-ring ` ... :: 'a mod-ring set)
    by (intro image-cong refl) (auto simp: of-int-mod-ring-def)
  also have ... = (UNIV :: 'a mod-ring set)
    using Abs-image by simp
  finally show enum-finite-field ` {..<CARD('a mod-ring)} = (UNIV :: 'a mod-ring
set) .
qed

end

typedef (overloaded) 'a :: semiring-1 ring-char = if CHAR('a) = 0 then UNIV
else {0..<CHAR('a)}
  by auto

lemma CARD-ring-char [simp]: CARD ('a :: semiring-1 ring-char) = CHAR('a)
proof −
  let ?A = if CHAR('a) = 0 then UNIV else {0..<CHAR('a)}
  interpret type-definition Rep-ring-char :: 'a ring-char ⇒ nat Abs-ring-char ?A
    by (rule type-definition-ring-char)
  from card show ?thesis
    by auto
qed

instance ring-char :: (semiring-prime-char) nontriv
proof
  show CARD('a ring-char) > 1
    using prime-nat-iff by auto
qed

instance ring-char :: (semiring-prime-char) prime-card
proof
  from CARD-ring-char show prime CARD('a ring-char)
    by auto
qed

lemma to-int-mod-ring-add:
  to-int-mod-ring (x + y :: 'a :: finite mod-ring) = (to-int-mod-ring x + to-int-mod-ring
y) mod CARD('a)
  by transfer auto

lemma to-int-mod-ring-mult:
  to-int-mod-ring (x * y :: 'a :: finite mod-ring) = (to-int-mod-ring x * to-int-mod-ring
y) mod CARD('a)
  by transfer auto
```

lemma *of-nat-mod-CHAR* [*simp*]: *of-nat* (*x mod CHAR*('*a* :: *semiring-1*)) = (*of-nat* *x* :: '*a*)
by (*metis* (*no-types*, *opaque-lifting*) *comm-monoid-add-class.add-0 div-mod-decomp* *mult-zero-right of-nat-CHAR of-nat-add of-nat-mult*)

lemma *of-int-mod-CHAR* [*simp*]: *of-int* (*x mod int CHAR*('*a* :: *ring-1*)) = (*of-int* *x* :: '*a*)
by (*simp add*: *of-int-eq-iff-cong-CHAR*)

end

3 Arithmetics via Records

We create a locale for rings and fields based on a record that includes all the necessary operations.

theory *Arithmetic-Record-Based*
imports
HOL−Library.More-List
HOL−Computational-Algebra.Euclidean-Algorithm
begin
datatype '*a arith-ops-record* = *Arith-Ops-Record*
(*zero* : '*a*)
(*one* : '*a*)
(*plus* : '*a* ⇒ '*a* ⇒ '*a*)
(*times* : '*a* ⇒ '*a* ⇒ '*a*)
(*minus* : '*a* ⇒ '*a* ⇒ '*a*)
(*uminus* : '*a* ⇒ '*a*)
(*divide* : '*a* ⇒ '*a* ⇒ '*a*)
(*inverse* : '*a* ⇒ '*a*)
(*modulo* : '*a* ⇒ '*a* ⇒ '*a*)
(*normalize* : '*a* ⇒ '*a*)
(*unit-factor* : '*a* ⇒ '*a*)
(*of-int* : *int* ⇒ '*a*)
(*to-int* : '*a* ⇒ *int*)
(*DP* : '*a* ⇒ *bool*)

hide-const (**open**)
zero
one
plus
times
minus
uminus
divide
inverse
modulo
normalize
unit-factor

of-int
to-int
DP

fun *listprod-i* :: *'i arith-ops-record* ⇒ *'i list* ⇒ *'i* **where**
listprod-i ops (*x* # *xs*) = *arith-ops-record.times ops x* (*listprod-i ops xs*)
| *listprod-i ops* [] = *arith-ops-record.one ops*

locale *arith-ops* = **fixes** *ops* :: *'i arith-ops-record* (**structure**)
begin

abbreviation (*input*) *zero* **where** *zero* ≡ *arith-ops-record.zero ops*
abbreviation (*input*) *one* **where** *one* ≡ *arith-ops-record.one ops*
abbreviation (*input*) *plus* **where** *plus* ≡ *arith-ops-record.plus ops*
abbreviation (*input*) *times* **where** *times* ≡ *arith-ops-record.times ops*
abbreviation (*input*) *minus* **where** *minus* ≡ *arith-ops-record.minus ops*
abbreviation (*input*) *uminus* **where** *uminus* ≡ *arith-ops-record.uminus ops*
abbreviation (*input*) *divide* **where** *divide* ≡ *arith-ops-record.divide ops*
abbreviation (*input*) *inverse* **where** *inverse* ≡ *arith-ops-record.inverse ops*
abbreviation (*input*) *modulo* **where** *modulo* ≡ *arith-ops-record.modulo ops*
abbreviation (*input*) *normalize* **where** *normalize* ≡ *arith-ops-record.normalize ops*
abbreviation (*input*) *unit-factor* **where** *unit-factor* ≡ *arith-ops-record.unit-factor ops*
abbreviation (*input*) *DP* **where** *DP* ≡ *arith-ops-record.DP ops*

partial-function (*tailrec*) *gcd-eucl-i* :: *'i* ⇒ *'i* ⇒ *'i* **where**
gcd-eucl-i a b = (*if b* = *zero*
then normalize a else gcd-eucl-i b (*modulo a b*))

partial-function (*tailrec*) *euclid-ext-aux-i* :: *'i* ⇒ *'i* ⇒ *'i* ⇒ *'i* ⇒ *'i* ⇒ *'i* ⇒ (*'i* × *'i*) × *'i* **where**
euclid-ext-aux-i s' s t' t r' r = (
if r = *zero then let c* = *divide one* (*unit-factor r'*) *in* ((*times s' c*, *times t' c*), *normalize r'*)
else let q = *divide r' r*
in euclid-ext-aux-i s (*minus s'* (*times q s*)) *t* (*minus t'* (*times q t*)) *r* (*modulo r' r*))

abbreviation (*input*) *euclid-ext-i* :: *'i* ⇒ *'i* ⇒ (*'i* × *'i*) × *'i* **where**
euclid-ext-i ≡ *euclid-ext-aux-i one zero zero one*

end

declare *arith-ops.gcd-eucl-i.simps*[*code*]
declare *arith-ops.euclid-ext-aux-i.simps*[*code*]

unbundle *lifting-syntax*

locale *ring-ops = arith-ops ops* **for** *ops :: 'i arith-ops-record +*
  fixes *R :: 'i ⇒ 'a :: comm-ring-1 ⇒ bool*
  assumes *bi-unique*[*transfer-rule*]: *bi-unique R*
  and *right-total*[*transfer-rule*]: *right-total R*
  and *zero*[*transfer-rule*]: *R zero 0*
  and *one*[*transfer-rule*]: *R one 1*
  and *plus*[*transfer-rule*]: *(R ===> R ===> R) plus (+)*
  and *minus*[*transfer-rule*]: *(R ===> R ===> R) minus (−)*
  and *uminus*[*transfer-rule*]: *(R ===> R) uminus Groups.uminus*
  and *times*[*transfer-rule*]: *(R ===> R ===> R) times ((∗))*
  and *eq*[*transfer-rule*]: *(R ===> R ===> (=)) (=) (=)*
  and *DPR*[*transfer-domain-rule*]: *Domainp R = DP*
begin
lemma *left-right-unique*[*transfer-rule*]: *left-unique R right-unique R*
  using *bi-unique* **unfolding** *bi-unique-def left-unique-def right-unique-def* **by** *auto*

lemma *listprod-i*[*transfer-rule*]: *(list-all2 R ===> R) (listprod-i ops) prod-list*
proof *(intro rel-funI, goal-cases)*
  case *(1 xs ys)*
  thus *?case*
  proof *(induct xs ys rule: list-all2-induct)*
    case *(Cons x xs y ys)*
    note *[transfer-rule] = this*
    show *?case* **by** *simp transfer-prover*
  qed *(simp add: one)*
qed
end

locale *idom-ops = ring-ops ops R* **for** *ops :: 'i arith-ops-record* **and**
  R :: 'i ⇒ 'a :: idom ⇒ bool

locale *idom-divide-ops = idom-ops ops R* **for** *ops :: 'i arith-ops-record* **and**
  R :: 'i ⇒ 'a :: idom-divide ⇒ bool +
  assumes *divide*[*transfer-rule*]: *(R ===> R ===> R) divide Rings.divide*

locale *euclidean-semiring-ops = idom-ops ops R* **for** *ops :: 'i arith-ops-record* **and**
  R :: 'i ⇒ 'a :: {idom,normalization-euclidean-semiring} ⇒ bool +
  assumes *modulo*[*transfer-rule*]: *(R ===> R ===> R) modulo (mod)*
    and *normalize*[*transfer-rule*]: *(R ===> R) normalize Rings.normalize*
    and *unit-factor*[*transfer-rule*]: *(R ===> R) unit-factor Rings.unit-factor*
begin
lemma *gcd-eucl-i* [*transfer-rule*]: *(R ===> R ===> R) gcd-eucl-i Euclidean-Algorithm.gcd*

proof *(intro rel-funI, goal-cases)*
  case *(1 x X y Y)*
  thus *?case*
  proof *(induct X Y arbitrary: x y rule: Euclidean-Algorithm.gcd.induct)*
    case *(1 X Y x y)*

```
    note [transfer-rule] = 1(2-)
    note simps = gcd-eucl-i.simps[of x y] Euclidean-Algorithm.gcd.simps[of X Y]
    have eq: (y = zero) = (Y = 0) by transfer-prover
    show ?case
    proof (cases Y = 0)
      case True
      hence *: y = zero using eq by simp
      have R (normalize x) (Rings.normalize X) by transfer-prover
      thus ?thesis unfolding simps unfolding True * by simp
    next
      case False
      with eq have yz: y ≠ zero by simp
      have R (gcd-eucl-i y (modulo x y)) (Euclidean-Algorithm.gcd Y (X mod Y))
        by (rule 1(1)[OF False], transfer-prover+)
      thus ?thesis unfolding simps using False yz by simp
    qed
  qed
qed
end

locale euclidean-ring-ops = euclidean-semiring-ops ops R for ops :: 'i arith-ops-record
and
  R :: 'i ⇒ 'a :: {idom,euclidean-ring-gcd} ⇒ bool +
  assumes divide[transfer-rule]: (R ===> R ===> R) divide (div)
begin
lemma euclid-ext-aux-i[transfer-rule]:
  (R ===> R ===> R ===> R ===> R ===> R ===> rel-prod (rel-prod
R R) R) euclid-ext-aux-i euclid-ext-aux
proof (intro rel-funI, goal-cases)
  case (1 z Z a A b B c C x X y Y)
  thus ?case
  proof (induct Z A B C X Y arbitrary: z a b c x y rule: euclid-ext-aux.induct)
    case (1 Z A B C X Y z a b c x y)
    note [transfer-rule] = 1(2-)
    note simps = euclid-ext-aux-i.simps[of z a b c x y] euclid-ext-aux.simps[of Z A
B C X Y]
    have eq: (y = zero) = (Y = 0) by transfer-prover
    show ?case
    proof (cases Y = 0)
      case True
      hence *: (y = zero) = True (Y = 0) = True using eq by auto
      show ?thesis unfolding simps unfolding * if-True
        by transfer-prover
    next
      case False
      hence *: (y = zero) = False (Y = 0) = False using eq by auto
      have XY: R (modulo x y) (X mod Y) by transfer-prover
        have YA: R (minus z (times (divide x y) a)) (Z - X div Y * A) by
transfer-prover
```

```
      have YC: R (minus b (times (divide x y) c)) (B - X div Y * C) by
transfer-prover
    note [transfer-rule] = 1(1)[OF False refl 1(3) YA 1(5) YC 1(7) XY]

    show ?thesis unfolding simps * if-False Let-def by transfer-prover
   qed
  qed
qed

lemma euclid-ext-i [transfer-rule]:
  (R ===> R ===> rel-prod (rel-prod R R) R) euclid-ext-i euclid-ext
  by transfer-prover

end

locale field-ops = idom-divide-ops ops R + euclidean-semiring-ops ops R for ops
:: 'i arith-ops-record and
  R :: 'i ⇒ 'a :: {field-gcd} ⇒ bool +
  assumes inverse[transfer-rule]: (R ===> R) inverse Fields.inverse


lemma nth-default-rel[transfer-rule]: (S ===> list-all2 S ===> (=) ===> S)
nth-default nth-default
proof (intro rel-funI, clarify, goal-cases)
  case (1 x y xs ys - n)
  from 1(2) show ?case
  proof (induct arbitrary: n)
    case Nil
    thus ?case using 1(1) by simp
  next
    case (Cons x y xs ys n)
    thus ?case by (cases n, auto)
  qed
qed

lemma strip-while-rel[transfer-rule]:
  ((A ===> (=)) ===> list-all2 A ===> list-all2 A) strip-while strip-while
  unfolding strip-while-def[abs-def] by transfer-prover

lemma list-all2-last[simp]: list-all2 A (xs @ [x]) (ys @ [y]) ⟷ list-all2 A xs ys ∧
A x y
proof (cases length xs = length ys)
  case True
  show ?thesis by (simp add: list-all2-append[OF True])
next
  case False
  note len = list-all2-lengthD[of A]
  from len[of xs ys] len[of xs @ [x] ys @ [y]] False
  show ?thesis by auto
```

qed

end

3.1 Finite Fields

We provide four implementations for $GF(p)$ – the field with p elements for some prime p – one by int, one by integers, one by 32-bit numbers and one 64-bit implementation. Correctness of the implementations is proven by transfer rules to the type-based version of $GF(p)$.

theory *Finite-Field-Record-Based*
imports
Finite-Field
Arithmetic-Record-Based
Native-Word.Uint32
Native-Word.Uint64
HOL−Library.Code-Target-Numeral
Native-Word.Code-Target-Int-Bit
begin

definition *mod-nonneg-pos* :: *integer* ⇒ *integer* ⇒ *integer* **where**
x ≥ *0* ⟹ *y* > *0* ⟹ *mod-nonneg-pos x y* = (*x mod y*)

code-printing — FIXME illusion of partiality
constant *mod-nonneg-pos* ⇀
(*SML*) *IntInf.mod*/ (-,/ -)
and (*Eval*) *IntInf.mod*/ (-,/ -)
and (*OCaml*) *Z.rem*
and (*Haskell*) *Prelude.mod*/ (-)/ (-)
and (*Scala*) !((*k*: *BigInt*) => (*l*: *BigInt*) =>/ (*k* `%` *l*))

definition *mod-nonneg-pos-int* :: *int* ⇒ *int* ⇒ *int* **where**
mod-nonneg-pos-int x y = *int-of-integer* (*mod-nonneg-pos* (*integer-of-int x*) (*integer-of-int y*))

lemma *mod-nonneg-pos-int*[*simp*]: *x* ≥ *0* ⟹ *y* > *0* ⟹ *mod-nonneg-pos-int x y* = (*x mod y*)
unfolding *mod-nonneg-pos-int-def* **using** *mod-nonneg-pos-def* **by** *simp*

context
fixes *p* :: *int*
begin
definition *plus-p* :: *int* ⇒ *int* ⇒ *int* **where**
plus-p x y ≡ *let z* = *x* + *y in if z* ≥ *p then z* − *p else z*

definition *minus-p* :: *int* ⇒ *int* ⇒ *int* **where**
minus-p x y ≡ *if y* ≤ *x then x* − *y else x* + *p* − *y*

definition *uminus-p* :: *int* ⇒ *int* **where**
uminus-p x = (*if x* = *0 then 0 else p* − *x*)

definition *mult-p* :: *int* ⇒ *int* ⇒ *int* **where**
mult-p x y = (*mod-nonneg-pos-int* (*x* ∗ *y*) *p*)

fun *power-p* :: *int* ⇒ *nat* ⇒ *int* **where**
power-p x n = (*if n* = *0 then 1 else*
let (*d*,*r*) = *Euclidean-Rings.divmod-nat n 2*;
rec = *power-p* (*mult-p x x*) *d in*
if r = *0 then rec else mult-p rec x*)

In experiments with Berlekamp-factorization (where the prime p is usually small), it turned out that taking the below implementation of inverse via exponentiation is faster than the one based on the extended Euclidean algorithm.

definition *inverse-p* :: *int* ⇒ *int* **where**
inverse-p x = (*if x* = *0 then 0 else power-p x* (*nat* (*p* − *2*)))

definition *divide-p* :: *int* ⇒ *int* ⇒ *int* **where**
divide-p x y = *mult-p x* (*inverse-p y*)

definition *finite-field-ops-int* :: *int arith-ops-record* **where**
finite-field-ops-int ≡ *Arith-Ops-Record*
0
1
plus-p
mult-p
minus-p
uminus-p
divide-p
inverse-p
(λ *x y* . *if y* = *0 then x else 0*)
(λ *x* . *if x* = *0 then 0 else 1*)
(λ *x* . *x*)
(λ *x* . *x*)
(λ *x* . *x*)
(λ *x*. *0* ≤ *x* ∧ *x* < *p*)

end

context
fixes *p* :: *uint32*
begin
definition *plus-p32* :: *uint32* ⇒ *uint32* ⇒ *uint32* **where**
plus-p32 x y ≡ *let z* = *x* + *y in if z* ≥ *p then z* − *p else z*

definition *minus-p32* :: *uint32* ⇒ *uint32* ⇒ *uint32* **where**

minus-p32 x y ≡ *if y* ≤ *x then x* − *y else* (*x* + *p*) − *y*

definition *uminus-p32* :: *uint32* ⇒ *uint32* **where**
uminus-p32 x = (*if x* = *0 then 0 else p* − *x*)

definition *mult-p32* :: *uint32* ⇒ *uint32* ⇒ *uint32* **where**
mult-p32 x y = (*x* ∗ *y mod p*)

lemma *int-of-uint32-shift*: *int-of-uint32* (*drop-bit k n*) = (*int-of-uint32 n*) *div* (*2* ^ *k*)
apply *transfer*
apply *transfer*
apply (*simp add*: *take-bit-drop-bit min-def*)
apply (*simp add*: *drop-bit-eq-div*)
done

lemma *int-of-uint32-0-iff*: *int-of-uint32 n* = *0* ⟷ *n* = *0*
by (*transfer*, *rule uint-0-iff*)

lemma *int-of-uint32-0*: *int-of-uint32 0* = *0* **unfolding** *int-of-uint32-0-iff* **by** *simp*

lemma *int-of-uint32-ge-0*: *int-of-uint32 n* ≥ *0*
by (*transfer*, *auto*)

lemma *two-32*: *2* ^ *LENGTH*(*32*) = (*4294967296* :: *int*) **by** *simp*

lemma *int-of-uint32-plus*: *int-of-uint32* (*x* + *y*) = (*int-of-uint32 x* + *int-of-uint32 y*) *mod 4294967296*
by (*transfer*, *unfold uint-word-ariths two-32*, *rule refl*)

lemma *int-of-uint32-minus*: *int-of-uint32* (*x* − *y*) = (*int-of-uint32 x* − *int-of-uint32 y*) *mod 4294967296*
by (*transfer*, *unfold uint-word-ariths two-32*, *rule refl*)

lemma *int-of-uint32-mult*: *int-of-uint32* (*x* ∗ *y*) = (*int-of-uint32 x* ∗ *int-of-uint32 y*) *mod 4294967296*
by (*transfer*, *unfold uint-word-ariths two-32*, *rule refl*)

lemma *int-of-uint32-mod*: *int-of-uint32* (*x mod y*) = (*int-of-uint32 x mod int-of-uint32 y*)
by (*transfer*, *unfold uint-mod two-32*, *rule refl*)

lemma *int-of-uint32-inv*: *0* ≤ *x* ⟹ *x* < *4294967296* ⟹ *int-of-uint32* (*uint32-of-int x*) = *x*
by *transfer* (*simp add*: *take-bit-int-eq-self unsigned-of-int*)

context
includes *bit-operations-syntax*
begin

```
function power-p32 :: uint32 ⇒ uint32 ⇒ uint32 where
  power-p32 x n = (if n = 0 then 1 else
    let rec = power-p32 (mult-p32 x x) (drop-bit 1 n) in
    if n AND 1 = 0 then rec else mult-p32 rec x)
  by pat-completeness auto

termination
proof −
  {
    fix n :: uint32
    assume n ≠ 0
    with int-of-uint32-ge-0[of n] int-of-uint32-0-iff[of n] have int-of-uint32 n > 0
by auto
    hence 0 < int-of-uint32 n int-of-uint32 n div 2 < int-of-uint32 n by auto
  } note ∗ = this
  show ?thesis
   by (relation measure (λ (x,n). nat (int-of-uint32 n)), auto simp: int-of-uint32-shift
∗)
qed

end
```

In experiments with Berlekamp-factorization (where the prime p is usually small), it turned out that taking the below implementation of inverse via exponentiation is faster than the one based on the extended Euclidean algorithm.

```
definition inverse-p32 :: uint32 ⇒ uint32 where
  inverse-p32 x = (if x = 0 then 0 else power-p32 x (p − 2))

definition divide-p32 :: uint32 ⇒ uint32 ⇒ uint32  where
  divide-p32 x y = mult-p32 x (inverse-p32 y)

definition finite-field-ops32 :: uint32 arith-ops-record where
  finite-field-ops32 ≡ Arith-Ops-Record
      0
      1
      plus-p32
      mult-p32
      minus-p32
      uminus-p32
      divide-p32
      inverse-p32
      (λ x y . if y = 0 then x else 0)
      (λ x . if x = 0 then 0 else 1)
      (λ x . x)
      uint32-of-int
      int-of-uint32
      (λ x. 0 ≤ x ∧ x < p)
```

end

lemma *shiftr-uint32-code* [*code-unfold*]: *drop-bit 1 x* = (*uint32-shiftr x 1*)
by (*simp add*: *uint32.shiftr-def*)

3.1.1 Transfer Relation

locale *mod-ring-locale* =
fixes *p* :: *int* **and** *ty* :: *'a* :: *nontriv itself*
assumes *p*: *p* = *int CARD*('*a*)
begin
lemma *nat-p*: *nat p* = *CARD*('*a*) **unfolding** *p* **by** *simp*
lemma *p2*: *p* ≥ *2* **unfolding** *p* **using** *nontriv*[**where** *'a* = *'a*] **by** *auto*
lemma *p2-ident*: *int* (*CARD*('*a*) − *2*) = *p* − *2* **using** *p2* **unfolding** *p* **by** *simp*

definition *mod-ring-rel* :: *int* ⇒ *'a mod-ring* ⇒ *bool* **where**
mod-ring-rel x x' = (*x* = *to-int-mod-ring x'*)

lemma *Domainp-mod-ring-rel* [*transfer-domain-rule*]:
Domainp (*mod-ring-rel*) = (λ *v*. *v* ∈ {*0* ..< *p*})
proof −
{
fix *v* :: *int*
assume ∗: *0* ≤ *v v* < *p*
have *Domainp mod-ring-rel v*
proof
show *mod-ring-rel v* (*of-int-mod-ring v*) **unfolding** *mod-ring-rel-def* **using** ∗ *p* **by** *auto*
qed
} note ∗ = *this*
show *?thesis*
by (*intro ext iffI*, *insert range-to-int-mod-ring*[**where** *'a* = *'a*] ∗, *auto simp*: *mod-ring-rel-def p*)
qed

lemma *bi-unique-mod-ring-rel* [*transfer-rule*]:
bi-unique mod-ring-rel left-unique mod-ring-rel right-unique mod-ring-rel
unfolding *mod-ring-rel-def bi-unique-def left-unique-def right-unique-def*
by *auto*

lemma *right-total-mod-ring-rel* [*transfer-rule*]: *right-total mod-ring-rel*
unfolding *mod-ring-rel-def right-total-def* **by** *simp*

3.1.2 Transfer Rules

lemma *mod-ring-0*[*transfer-rule*]: *mod-ring-rel 0 0* **unfolding** *mod-ring-rel-def* **by** *simp*

lemma *mod-ring-1*[*transfer-rule*]: *mod-ring-rel 1 1* **unfolding** *mod-ring-rel-def* **by** *simp*

lemma *plus-p-mod-def*: **assumes** *x*: $x \in \{0 ..< p\}$ **and** *y*: $y \in \{0 ..< p\}$
shows *plus-p p x y* = $((x + y) \bmod p)$
proof (*cases* $p \leq x + y$)
case *False*
thus *?thesis* **using** *x y* **unfolding** *plus-p-def Let-def* **by** *auto*
next
case *True*
from *True x y* **have** ∗: $p > 0$ $0 \leq x + y - p$ $x + y - p < p$ **by** *auto*
from *True* **have** *id*: *plus-p p x y* = $x + y - p$ **unfolding** *plus-p-def* **by** *auto*
show *?thesis* **unfolding** *id* **using** ∗ **using** *mod-pos-pos-trivial* **by** *fastforce*
qed

lemma *mod-ring-plus*[*transfer-rule*]: (*mod-ring-rel* ===> *mod-ring-rel* ===> *mod-ring-rel*) (*plus-p p*) (+)
proof −
{
fix *x y* :: *'a mod-ring*
have *plus-p p* (*to-int-mod-ring x*) (*to-int-mod-ring y*) = *to-int-mod-ring* ($x + y$)
by (*transfer*, *subst plus-p-mod-def*, *auto*, *auto simp*: *p*)
} **note** ∗ = *this*
show *?thesis*
by (*intro rel-funI*, *auto simp*: *mod-ring-rel-def* ∗)
qed

lemma *minus-p-mod-def*: **assumes** *x*: $x \in \{0 ..< p\}$ **and** *y*: $y \in \{0 ..< p\}$
shows *minus-p p x y* = $((x - y) \bmod p)$
proof (*cases* $x - y < 0$)
case *False*
thus *?thesis* **using** *x y* **unfolding** *minus-p-def Let-def* **by** *auto*
next
case *True*
from *True x y* **have** ∗: $p > 0$ $0 \leq x - y + p$ $x - y + p < p$ **by** *auto*
from *True* **have** *id*: *minus-p p x y* = $x - y + p$ **unfolding** *minus-p-def* **by** *auto*
show *?thesis* **unfolding** *id* **using** ∗ **using** *mod-pos-pos-trivial* **by** *fastforce*
qed

lemma *mod-ring-minus*[*transfer-rule*]: (*mod-ring-rel* ===> *mod-ring-rel* ===> *mod-ring-rel*) (*minus-p p*) (−)
proof −
{
fix *x y* :: *'a mod-ring*
have *minus-p p* (*to-int-mod-ring x*) (*to-int-mod-ring y*) = *to-int-mod-ring* ($x - y$)

```
    by (transfer, subst minus-p-mod-def, auto simp: p)
  } note * = this
  show ?thesis
    by (intro rel-funI, auto simp: mod-ring-rel-def *)
qed


lemma mod-ring-uminus[transfer-rule]: (mod-ring-rel ===> mod-ring-rel) (uminus-p
p) uminus
proof −
  {
    fix x :: 'a mod-ring
    have uminus-p p (to-int-mod-ring x) = to-int-mod-ring (uminus x)
      by (transfer, auto simp: uminus-p-def p)
  } note * = this
  show ?thesis
    by (intro rel-funI, auto simp: mod-ring-rel-def *)
qed


lemma mod-ring-mult[transfer-rule]: (mod-ring-rel ===> mod-ring-rel ===>
mod-ring-rel) (mult-p p) ((*))
proof −
  {
    fix x y :: 'a mod-ring
    have mult-p p (to-int-mod-ring x) (to-int-mod-ring y) = to-int-mod-ring (x *
y)
      by (transfer, auto simp: mult-p-def p)
  } note * = this
  show ?thesis
    by (intro rel-funI, auto simp: mod-ring-rel-def *)
qed


lemma mod-ring-eq[transfer-rule]: (mod-ring-rel ===> mod-ring-rel ===> (=))
(=) (=)
  by (intro rel-funI, auto simp: mod-ring-rel-def)


lemma mod-ring-power[transfer-rule]: (mod-ring-rel ===> (=) ===> mod-ring-rel)
(power-p p) (^)
proof (intro rel-funI, clarify, unfold binary-power[symmetric], goal-cases)
  fix x y n
  assume xy: mod-ring-rel x y
  from xy show mod-ring-rel (power-p p x n) (binary-power y n)
  proof (induct y n arbitrary: x rule: binary-power.induct)
    case (1 x n y)
    note 1(2)[transfer-rule]
    show ?case
```

```
    proof (cases n = 0)
      case True
      thus ?thesis by (simp add: mod-ring-1)
    next
      case False
      obtain d r where id: Euclidean-Rings.divmod-nat n 2 = (d,r) by force
      let ?int = power-p p (mult-p p y y) d
      let ?gfp = binary-power (x * x) d
      from False have id': ?thesis = (mod-ring-rel
         (if r = 0 then ?int else mult-p p ?int y)
         (if r = 0 then ?gfp else ?gfp * x))
       unfolding power-p.simps[of - - n] binary-power.simps[of - n] Let-def id split
by simp
      have [transfer-rule]: mod-ring-rel ?int ?gfp
        by (rule 1(1)[OF False refl id[symmetric]], transfer-prover)
      show ?thesis unfolding id' by transfer-prover
    qed
  qed
qed

declare power-p.simps[simp del]

lemma ring-finite-field-ops-int: ring-ops (finite-field-ops-int p) mod-ring-rel
  by (unfold-locales, auto simp:
  finite-field-ops-int-def
  bi-unique-mod-ring-rel
  right-total-mod-ring-rel
  mod-ring-plus
  mod-ring-minus
  mod-ring-uminus
  mod-ring-mult
  mod-ring-eq
  mod-ring-0
  mod-ring-1
  Domainp-mod-ring-rel)
end

locale prime-field = mod-ring-locale p ty for p and ty :: 'a :: prime-card itself
begin

lemma prime: prime p unfolding p using prime-card[where 'a = 'a] by simp


lemma mod-ring-mod[transfer-rule]:
 (mod-ring-rel ===> mod-ring-rel ===> mod-ring-rel) ((λ x y. if y = 0 then x
else 0)) (mod)
proof −
  {
    fix x y :: 'a mod-ring
```

```
    have (if to-int-mod-ring y = 0 then to-int-mod-ring x else 0) = to-int-mod-ring
(x mod y)
      unfolding modulo-mod-ring-def by auto
  } note * = this
  show ?thesis
    by (intro rel-funI, auto simp: mod-ring-rel-def *[symmetric])
qed


lemma mod-ring-normalize[transfer-rule]: (mod-ring-rel ===> mod-ring-rel) ((λ
x. if x = 0 then 0 else 1)) normalize
proof −
  {
    fix x :: 'a mod-ring
    have (if to-int-mod-ring x = 0 then 0 else 1) = to-int-mod-ring (normalize x)
      unfolding normalize-mod-ring-def by auto
  } note * = this
  show ?thesis
    by (intro rel-funI, auto simp: mod-ring-rel-def *[symmetric])
qed


lemma mod-ring-unit-factor[transfer-rule]: (mod-ring-rel ===> mod-ring-rel) (λ
x. x) unit-factor
proof −
  {
    fix x :: 'a mod-ring
    have to-int-mod-ring x = to-int-mod-ring (unit-factor x)
      unfolding unit-factor-mod-ring-def by auto
  } note * = this
  show ?thesis
    by (intro rel-funI, auto simp: mod-ring-rel-def *[symmetric])
qed


lemma mod-ring-inverse[transfer-rule]: (mod-ring-rel ===> mod-ring-rel) (inverse-p
p) inverse
proof (intro rel-funI)
  fix x y
  assume [transfer-rule]: mod-ring-rel x y
  show mod-ring-rel (inverse-p p x) (inverse y)
    unfolding inverse-p-def inverse-mod-ring-def
    apply (transfer-prover-start)
    apply (transfer-step)+
    apply (unfold p2-ident)
    apply (rule refl)
    done
qed
```

lemma *mod-ring-divide*[*transfer-rule*]: (*mod-ring-rel* ===> *mod-ring-rel* ===> *mod-ring-rel*)
(*divide-p p*) (/)
unfolding *divide-p-def*[*abs-def*] *divide-mod-ring-def*[*abs-def*] *inverse-mod-ring-def*[*symmetric*]
by *transfer-prover*

lemma *mod-ring-rel-unsafe*: **assumes** $x < CARD('a)$
shows *mod-ring-rel* (*int x*) (*of-nat x*) $0 < x \Longrightarrow$ *of-nat x* $\neq$ (*0* :: *'a mod-ring*)
proof −
have *id*: *of-nat x* = (*of-int* (*int x*) :: *'a mod-ring*) **by** *simp*
show *mod-ring-rel* (*int x*) (*of-nat x*) $0 < x \Longrightarrow$ *of-nat x* $\neq$ (*0* :: *'a mod-ring*)
unfolding *id*
unfolding *mod-ring-rel-def*
proof (*auto simp add*: *assms of-int-of-int-mod-ring*)
assume $0 < x$ **with** *assms*
have *of-int-mod-ring* (*int x*) $\neq$ (*0* :: *'a mod-ring*)
by (*metis* (*no-types*) *less-imp-of-nat-less less-irrefl of-nat-0-le-iff of-nat-0-less-iff to-int-mod-ring-hom.hom-zero to-int-mod-ring-of-int-mod-ring*)
thus *of-int-mod-ring* (*int x*) = (*0* :: *'a mod-ring*) $\Longrightarrow$ *False* **by** *blast*
qed
qed

lemma *finite-field-ops-int*: *field-ops* (*finite-field-ops-int p*) *mod-ring-rel*
by (*unfold-locales*, *auto simp*:
finite-field-ops-int-def
bi-unique-mod-ring-rel
right-total-mod-ring-rel
mod-ring-divide
mod-ring-plus
mod-ring-minus
mod-ring-uminus
mod-ring-inverse
mod-ring-mod
mod-ring-unit-factor
mod-ring-normalize
mod-ring-mult
mod-ring-eq
mod-ring-0
mod-ring-1
Domainp-mod-ring-rel)

end

Once we have proven the soundness of the implementation, we do not care any longer that *'a mod-ring* has been defined internally via lifting. Disabling the transfer-rules will hide the internal definition in further applications of transfer.

lifting-forget *mod-ring.lifting*

For soundness of the 32-bit implementation, we mainly prove that this implementation implements the int-based implementation of the mod-ring.

context *mod-ring-locale*
begin

context fixes *pp* :: *uint32*
  assumes *ppp*: *p* = *int-of-uint32 pp*
  and *small*: *p* ≤ *65535*
begin

lemmas *uint32-simps* =
  int-of-uint32-0
  int-of-uint32-plus
  int-of-uint32-minus
  int-of-uint32-mult

definition *urel32* :: *uint32* ⇒ *int* ⇒ *bool* **where** *urel32 x y* = (*y* = *int-of-uint32 x* ∧ *y* < *p*)

definition *mod-ring-rel32* :: *uint32* ⇒ *'a mod-ring* ⇒ *bool* **where**
  mod-ring-rel32 x y = (∃ *z*. *urel32 x z* ∧ *mod-ring-rel z y*)

lemma *urel32-0*: *urel32 0 0* **unfolding** *urel32-def* **using** *p2* **by** (*simp*, *transfer*, *simp*)

lemma *urel32-1*: *urel32 1 1* **unfolding** *urel32-def* **using** *p2* **by** (*simp*, *transfer*, *simp*)

lemma *le-int-of-uint32*: (*x* ≤ *y*) = (*int-of-uint32 x* ≤ *int-of-uint32 y*)
  by (*transfer*, *simp add*: *word-le-def*)

lemma *urel32-plus*: **assumes** *urel32 x y urel32 x' y'*
  shows *urel32* (*plus-p32 pp x x'*) (*plus-p p y y'*)
proof −
  let *?x* = *int-of-uint32 x*
  let *?x'* = *int-of-uint32 x'*
  let *?p* = *int-of-uint32 pp*
  from *assms int-of-uint32-ge-0* **have** *id*: *y* = *?x y'* = *?x'*
    and *rel*: *0* ≤ *?x ?x* < *p*
      0 ≤ *?x' ?x'* ≤ *p* **unfolding** *urel32-def* **by** *auto*
  have *le*: (*pp* ≤ *x* + *x'*) = (*?p* ≤ *?x* + *?x'*) **unfolding** *le-int-of-uint32*
    using *rel small* **by** (*auto simp*: *uint32-simps*)
  show *?thesis*
  proof (*cases ?p* ≤ *?x* + *?x'*)
    case *True*
    hence *True*: (*?p* ≤ *?x* + *?x'*) = *True* **by** *simp*
    show *?thesis* **unfolding** *id*
      using *small rel* **unfolding** *plus-p32-def plus-p-def Let-def urel32-def*

```
    unfolding ppp le True if-True
    using True by (auto simp: uint32-simps)
  next
    case False
    hence False: (?p ≤ ?x + ?x') = False by simp
    show ?thesis unfolding id
      using small rel unfolding plus-p32-def plus-p-def Let-def urel32-def
      unfolding ppp le False if-False
      using False by (auto simp: uint32-simps)
  qed
qed

lemma urel32-minus: assumes urel32 x y urel32 x' y'
  shows urel32 (minus-p32 pp x x') (minus-p p y y')
proof −
  let ?x = int-of-uint32 x
  let ?x' = int-of-uint32 x'
  from assms int-of-uint32-ge-0 have id: y = ?x y' = ?x'
    and rel: 0 ≤ ?x ?x < p
      0 ≤ ?x' ?x' ≤ p unfolding urel32-def by auto
  have le: (x' ≤ x) = (?x' ≤ ?x) unfolding le-int-of-uint32
    using rel small by (auto simp: uint32-simps)
  show ?thesis
  proof (cases ?x' ≤ ?x)
    case True
    hence True: (?x' ≤ ?x) = True by simp
    show ?thesis unfolding id
      using small rel unfolding minus-p32-def minus-p-def Let-def urel32-def
      unfolding ppp le True if-True
      using True by (auto simp: uint32-simps)
  next
    case False
    hence False: (?x' ≤ ?x) = False by simp
    show ?thesis unfolding id
      using small rel unfolding minus-p32-def minus-p-def Let-def urel32-def
      unfolding ppp le False if-False
      using False by (auto simp: uint32-simps)
  qed
qed

lemma urel32-uminus: assumes urel32 x y
  shows urel32 (uminus-p32 pp x) (uminus-p p y)
proof −
  let ?x = int-of-uint32 x
  from assms int-of-uint32-ge-0 have id: y = ?x
    and rel: 0 ≤ ?x ?x < p
      unfolding urel32-def by auto
  have le: (x = 0) = (?x = 0) unfolding int-of-uint32-0-iff
    using rel small by (auto simp: uint32-simps)
```

```
  show ?thesis
  proof (cases ?x = 0)
    case True
    hence True: (?x = 0) = True by simp
    show ?thesis unfolding id
      using small rel unfolding uminus-p32-def uminus-p-def Let-def urel32-def
      unfolding ppp le True if-True
      using True by (auto simp: uint32-simps)
  next
    case False
    hence False: (?x = 0) = False by simp
    show ?thesis unfolding id
      using small rel unfolding uminus-p32-def uminus-p-def Let-def urel32-def
      unfolding ppp le False if-False
      using False by (auto simp: uint32-simps)
  qed
qed

lemma urel32-mult: assumes urel32 x y urel32 x' y'
  shows urel32 (mult-p32 pp x x') (mult-p p y y')
proof −
  let ?x = int-of-uint32 x
  let ?x' = int-of-uint32 x'
  from assms int-of-uint32-ge-0 have id: y = ?x y' = ?x'
    and rel: 0 ≤ ?x ?x < p
      0 ≤ ?x' ?x' < p unfolding urel32-def by auto
  from rel have ?x * ?x' < p * p by (metis mult-strict-mono')
  also have ... ≤ 65536 * 65536
    by (rule mult-mono, insert p2 small, auto)
  finally have le: ?x * ?x' < 4294967296 by simp
  show ?thesis unfolding id
      using small rel unfolding mult-p32-def mult-p-def Let-def urel32-def
      unfolding ppp
    by (auto simp: uint32-simps, unfold int-of-uint32-mod int-of-uint32-mult,
        subst mod-pos-pos-trivial[of - 4294967296], insert le, auto)
qed

lemma urel32-eq: assumes urel32 x y urel32 x' y'
  shows (x = x') = (y = y')
proof −
  let ?x = int-of-uint32 x
  let ?x' = int-of-uint32 x'
  from assms int-of-uint32-ge-0 have id: y = ?x y' = ?x'
    unfolding urel32-def by auto
  show ?thesis unfolding id by (transfer, transfer) rule
qed

lemma urel32-normalize:
assumes x: urel32 x y
```

shows *urel32* (*if x = 0 then 0 else 1*) (*if y = 0 then 0 else 1*)
unfolding *urel32-eq*[*OF x urel32-0*] **using** *urel32-0 urel32-1* **by** *auto*

lemma *urel32-mod*:
assumes *x*: *urel32 x x'* **and** *y*: *urel32 y y'*
shows *urel32* (*if y = 0 then x else 0*) (*if y' = 0 then x' else 0*)
unfolding *urel32-eq*[*OF y urel32-0*] **using** *urel32-0 x* **by** *auto*

lemma *urel32-power*: *urel32 x x'* ⟹ *urel32 y* (*int y'*) ⟹ *urel32* (*power-p32 pp x y*) (*power-p p x' y'*)
including *bit-operations-syntax* **proof** (*induct x' y' arbitrary*: *x y rule*: *power-p.induct*[*of - p*])
case (*1 x' y' x y*)
note *x = 1*(*2*) **note** *y = 1*(*3*)
show *?case*
proof (*cases y' = 0*)
case *True*
hence *y*: *y = 0* **using** *urel32-eq*[*OF y urel32-0*] **by** *auto*
show *?thesis* **unfolding** *y True* **by** (*simp add*: *power-p.simps urel32-1*)
next
case *False*
hence *id*: (*y = 0*) = *False* (*y' = 0*) = *False* **using** *urel32-eq*[*OF y urel32-0*] **by** *auto*
from *y* **have** ‹*int y' = int-of-uint32 y*› ‹*int y' < p*›
by (*simp-all add*: *urel32-def*)
obtain *d' r'* **where** *dr'*: *Euclidean-Rings.divmod-nat y' 2* = (*d'*,*r'*) **by** *force*
from *Euclidean-Rings.divmod-nat-def*[*of y' 2*, *unfolded dr'*]
have *r'*: *r' = y' mod 2* **and** *d'*: *d' = y' div 2* **by** *auto*
have *urel32* (*y AND 1*) *r'*
using ‹*int y' < p*› *small*
apply (*simp add*: *urel32-def and-one-eq r'*)
apply (*auto simp add*: *ppp and-one-eq*)
apply (*simp add*: *of-nat-mod int-of-uint32.rep-eq modulo-uint32.rep-eq uint-mod* ‹*int y' = int-of-uint32 y*›)
done
from *urel32-eq*[*OF this urel32-0*]
have *rem*: (*y AND 1 = 0*) = (*r' = 0*) **by** *simp*
have *div*: *urel32* (*drop-bit 1 y*) (*int d'*) **unfolding** *d'* **using** *y* **unfolding** *urel32-def* **using** *small*
unfolding *ppp*
apply *transfer*
apply *transfer*
apply (*auto simp add*: *drop-bit-Suc take-bit-int-eq-self*)
done
note *IH* = *1*(*1*)[*OF False refl dr'*[*symmetric*] *urel32-mult*[*OF x x*] *div*]
show *?thesis* **unfolding** *power-p.simps*[*of - - y'*] *power-p32.simps*[*of - - y*] *dr' id if-False rem*
using *IH urel32-mult*[*OF IH x*] **by** (*auto simp*: *Let-def*)
qed

qed

lemma *urel32-inverse*: **assumes** *x*: *urel32 x x'*
shows *urel32* (*inverse-p32 pp x*) (*inverse-p p x'*)
proof −
have *p*: *urel32* (*pp* − *2*) (*int* (*nat* (*p* − *2*))) **using** *p2 small* **unfolding** *urel32-def*
unfolding *ppp*
by (*simp add*: *int-of-uint32.rep-eq minus-uint32.rep-eq uint-sub-if'*)
show *?thesis*
unfolding *inverse-p32-def inverse-p-def urel32-eq*[*OF x urel32-0*] **using** *urel32-0*
urel32-power[*OF x p*]
by *auto*
qed

lemma *mod-ring-0-32*: *mod-ring-rel32 0 0*
using *urel32-0 mod-ring-0* **unfolding** *mod-ring-rel32-def* **by** *blast*

lemma *mod-ring-1-32*: *mod-ring-rel32 1 1*
using *urel32-1 mod-ring-1* **unfolding** *mod-ring-rel32-def* **by** *blast*

lemma *mod-ring-uminus32*: (*mod-ring-rel32* ===> *mod-ring-rel32*) (*uminus-p32*
pp) *uminus*
using *urel32-uminus mod-ring-uminus* **unfolding** *mod-ring-rel32-def rel-fun-def*
by *blast*

lemma *mod-ring-plus32*: (*mod-ring-rel32* ===> *mod-ring-rel32* ===> *mod-ring-rel32*)
(*plus-p32 pp*) (+)
using *urel32-plus mod-ring-plus* **unfolding** *mod-ring-rel32-def rel-fun-def* **by**
blast

lemma *mod-ring-minus32*: (*mod-ring-rel32* ===> *mod-ring-rel32* ===> *mod-ring-rel32*)
(*minus-p32 pp*) (−)
using *urel32-minus mod-ring-minus* **unfolding** *mod-ring-rel32-def rel-fun-def* **by**
blast

lemma *mod-ring-mult32*: (*mod-ring-rel32* ===> *mod-ring-rel32* ===> *mod-ring-rel32*)
(*mult-p32 pp*) ((∗))
using *urel32-mult mod-ring-mult* **unfolding** *mod-ring-rel32-def rel-fun-def* **by**
blast

lemma *mod-ring-eq32*: (*mod-ring-rel32* ===> *mod-ring-rel32* ===> (=)) (=)
(=)
using *urel32-eq mod-ring-eq* **unfolding** *mod-ring-rel32-def rel-fun-def* **by** *blast*

lemma *urel32-inj*: *urel32 x y* ⟹ *urel32 x z* ⟹ *y* = *z*
using *urel32-eq*[*of x y x z*] **by** *auto*

lemma *urel32-inj'*: *urel32 x z* ⟹ *urel32 y z* ⟹ *x* = *y*

using *urel32-eq*[*of x z y z*] **by** *auto*

lemma *bi-unique-mod-ring-rel32*:
bi-unique mod-ring-rel32 left-unique mod-ring-rel32 right-unique mod-ring-rel32
using *bi-unique-mod-ring-rel urel32-inj'*
unfolding *mod-ring-rel32-def bi-unique-def left-unique-def right-unique-def*
by (*auto simp*: *urel32-def*)

lemma *right-total-mod-ring-rel32*: *right-total mod-ring-rel32*
unfolding *mod-ring-rel32-def right-total-def*
proof
fix *y* :: *'a mod-ring*
from *right-total-mod-ring-rel*[*unfolded right-total-def*, *rule-format*, *of y*]
obtain *z* **where** *zy*: *mod-ring-rel z y* **by** *auto*
hence *zp*: *0* ≤ *z z* < *p* **unfolding** *mod-ring-rel-def p* **using** *range-to-int-mod-ring*[**where** *'a* = *'a*] **by** *auto*
hence *urel32* (*uint32-of-int z*) *z* **unfolding** *urel32-def* **using** *small* **unfolding** *ppp*
by (*auto simp*: *int-of-uint32-inv*)
with *zy* **show** ∃ *x z*. *urel32 x z* ∧ *mod-ring-rel z y* **by** *blast*
qed

lemma *Domainp-mod-ring-rel32*: *Domainp mod-ring-rel32* = (λ*x*. *0* ≤ *x* ∧ *x* < *pp*)
proof
fix *x*
show *Domainp mod-ring-rel32 x* = (*0* ≤ *x* ∧ *x* < *pp*)
unfolding *Domainp.simps*
unfolding *mod-ring-rel32-def*
proof
let *?i* = *int-of-uint32*
assume ∗: *0* ≤ *x* ∧ *x* < *pp*
hence *0* ≤ *?i x* ∧ *?i x* < *p* **using** *small* **unfolding** *ppp*
by (*transfer*, *auto simp*: *word-less-def*)
hence *?i x* ∈ {*0* ..< *p*} **by** *auto*
with *Domainp-mod-ring-rel*
have *Domainp mod-ring-rel* (*?i x*) **by** *auto*
from *this*[*unfolded Domainp.simps*]
obtain *b* **where** *b*: *mod-ring-rel* (*?i x*) *b* **by** *auto*
show ∃ *a b*. *x* = *a* ∧ (∃ *z*. *urel32 a z* ∧ *mod-ring-rel z b*)
proof (*intro exI*, *rule conjI*[*OF refl*], *rule exI*, *rule conjI*[*OF - b*])
show *urel32 x* (*?i x*) **unfolding** *urel32-def* **using** *small* ∗ **unfolding** *ppp*
by (*transfer*, *auto simp*: *word-less-def*)
qed
next
assume ∃ *a b*. *x* = *a* ∧ (∃ *z*. *urel32 a z* ∧ *mod-ring-rel z b*)
then obtain *b z* **where** *xz*: *urel32 x z* **and** *zb*: *mod-ring-rel z b* **by** *auto*
hence *Domainp mod-ring-rel z* **by** *auto*
with *Domainp-mod-ring-rel* **have** *0* ≤ *z z* < *p* **by** *auto*

with *xz* **show** $0 \leq x \wedge x < pp$ **unfolding** *urel32-def* **using** *small* **unfolding** *ppp*
by (*transfer*, *auto simp*: *word-less-def*)
qed
qed

lemma *ring-finite-field-ops32*: *ring-ops* (*finite-field-ops32 pp*) *mod-ring-rel32*
by (*unfold-locales*, *auto simp*:
finite-field-ops32-def
bi-unique-mod-ring-rel32
right-total-mod-ring-rel32
mod-ring-plus32
mod-ring-minus32
mod-ring-uminus32
mod-ring-mult32
mod-ring-eq32
mod-ring-0-32
mod-ring-1-32
Domainp-mod-ring-rel32)
end
end

context *prime-field*
begin
context fixes *pp* :: *uint32*
assumes ∗: p = *int-of-uint32 pp* $p \leq 65535$
begin

lemma *mod-ring-normalize32*: (*mod-ring-rel32* ===> *mod-ring-rel32*) (λ*x*. *if x* = *0 then 0 else 1*) *normalize*
using *urel32-normalize*[*OF* ∗] *mod-ring-normalize* **unfolding** *mod-ring-rel32-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-mod32*: (*mod-ring-rel32* ===> *mod-ring-rel32* ===> *mod-ring-rel32*) (λ*x y*. *if y* = *0 then x else 0*) (*mod*)
using *urel32-mod*[*OF* ∗] *mod-ring-mod* **unfolding** *mod-ring-rel32-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-unit-factor32*: (*mod-ring-rel32* ===> *mod-ring-rel32*) (λ*x*. *x*) *unit-factor*
using *mod-ring-unit-factor* **unfolding** *mod-ring-rel32-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-inverse32*: (*mod-ring-rel32* ===> *mod-ring-rel32*) (*inverse-p32 pp*) *inverse*
using *urel32-inverse*[*OF* ∗] *mod-ring-inverse* **unfolding** *mod-ring-rel32-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-divide32*: (*mod-ring-rel32* ===> *mod-ring-rel32* ===> *mod-ring-rel32*)

(*divide-p32 pp*) (/)
using *mod-ring-inverse32 mod-ring-mult32*[*OF* ∗]
unfolding *divide-p32-def divide-mod-ring-def inverse-mod-ring-def*[*symmetric*]
rel-fun-def **by** *blast*

lemma *finite-field-ops32*: *field-ops* (*finite-field-ops32 pp*) *mod-ring-rel32*
by (*unfold-locales*, *insert ring-finite-field-ops32*[*OF* ∗], *auto simp*:
ring-ops-def
finite-field-ops32-def
mod-ring-divide32
mod-ring-inverse32
mod-ring-mod32
mod-ring-normalize32)

end
end

context
fixes *p* :: *uint64*
begin
definition *plus-p64* :: *uint64* ⇒ *uint64* ⇒ *uint64* **where**
plus-p64 x y ≡ *let z* = *x* + *y in if z* ≥ *p then z* − *p else z*

definition *minus-p64* :: *uint64* ⇒ *uint64* ⇒ *uint64* **where**
minus-p64 x y ≡ *if y* ≤ *x then x* − *y else* (*x* + *p*) − *y*

definition *uminus-p64* :: *uint64* ⇒ *uint64* **where**
uminus-p64 x = (*if x* = *0 then 0 else p* − *x*)

definition *mult-p64* :: *uint64* ⇒ *uint64* ⇒ *uint64* **where**
mult-p64 x y = (*x* ∗ *y mod p*)

lemma *int-of-uint64-shift*: *int-of-uint64* (*drop-bit k n*) = (*int-of-uint64 n*) *div* (2 ^ *k*)
apply *transfer*
apply *transfer*
apply (*simp add*: *take-bit-drop-bit min-def*)
apply (*simp add*: *drop-bit-eq-div*)
done

lemma *int-of-uint64-0-iff*: *int-of-uint64 n* = *0* ⟷ *n* = *0*
by (*transfer*, *rule uint-0-iff*)

lemma *int-of-uint64-0*: *int-of-uint64 0* = *0* **unfolding** *int-of-uint64-0-iff* **by** *simp*

lemma *int-of-uint64-ge-0*: *int-of-uint64 n* ≥ *0*
by (*transfer*, *auto*)

```
lemma two-64: 2 ^ LENGTH(64) = (18446744073709551616 :: int) by simp

lemma int-of-uint64-plus: int-of-uint64 (x + y) = (int-of-uint64 x + int-of-uint64
y) mod 18446744073709551616
  by (transfer, unfold uint-word-ariths two-64, rule refl)

lemma int-of-uint64-minus: int-of-uint64 (x - y) = (int-of-uint64 x - int-of-uint64
y) mod 18446744073709551616
  by (transfer, unfold uint-word-ariths two-64, rule refl)

lemma int-of-uint64-mult: int-of-uint64 (x * y) = (int-of-uint64 x * int-of-uint64
y) mod 18446744073709551616
  by (transfer, unfold uint-word-ariths two-64, rule refl)

lemma int-of-uint64-mod: int-of-uint64 (x mod y) = (int-of-uint64 x mod int-of-uint64
y)
  by (transfer, unfold uint-mod two-64, rule refl)

lemma int-of-uint64-inv: 0 ≤ x ⟹ x < 18446744073709551616 ⟹ int-of-uint64
(uint64-of-int x) = x
  by transfer (simp add: take-bit-int-eq-self unsigned-of-int)

context
  includes bit-operations-syntax
begin

function power-p64 :: uint64 ⇒ uint64 ⇒ uint64 where
  power-p64 x n = (if n = 0 then 1 else
    let rec = power-p64 (mult-p64 x x) (drop-bit 1 n) in
    if n AND 1 = 0 then rec else mult-p64 rec x)
  by pat-completeness auto

termination
proof -
  {
    fix n :: uint64
    assume n ≠ 0
    with int-of-uint64-ge-0[of n] int-of-uint64-0-iff[of n] have int-of-uint64 n > 0
by auto
    hence 0 < int-of-uint64 n int-of-uint64 n div 2 < int-of-uint64 n by auto
  } note * = this
  show ?thesis
   by (relation measure (λ (x,n). nat (int-of-uint64 n)), auto simp: int-of-uint64-shift
*)
qed

end
```

In experiments with Berlekamp-factorization (where the prime p is usually small), it turned out that taking the below implementation of inverse

via exponentiation is faster than the one based on the extended Euclidean algorithm.

definition *inverse-p64* :: *uint64* ⇒ *uint64* **where**
inverse-p64 x = (*if x* = *0 then 0 else power-p64 x* (*p* − *2*))

definition *divide-p64* :: *uint64* ⇒ *uint64* ⇒ *uint64* **where**
divide-p64 x y = *mult-p64 x* (*inverse-p64 y*)

definition *finite-field-ops64* :: *uint64 arith-ops-record* **where**
finite-field-ops64 ≡ *Arith-Ops-Record*
0
1
plus-p64
mult-p64
minus-p64
uminus-p64
divide-p64
inverse-p64
(λ *x y* . *if y* = *0 then x else 0*)
(λ *x* . *if x* = *0 then 0 else 1*)
(λ *x* . *x*)
uint64-of-int
int-of-uint64
(λ *x*. *0* ≤ *x* ∧ *x* < *p*)
end

lemma *shiftr-uint64-code* [*code-unfold*]: *drop-bit 1 x* = (*uint64-shiftr x 1*)
by (*simp add*: *uint64.shiftr-def*)

For soundness of the 64-bit implementation, we mainly prove that this implementation implements the int-based implementation of GF(p).

context *mod-ring-locale*
begin

context fixes *pp* :: *uint64*
assumes *ppp*: *p* = *int-of-uint64 pp*
and *small*: *p* ≤ *4294967295*
begin

lemmas *uint64-simps* =
int-of-uint64-0
int-of-uint64-plus
int-of-uint64-minus
int-of-uint64-mult

definition *urel64* :: *uint64* ⇒ *int* ⇒ *bool* **where** *urel64 x y* = (*y* = *int-of-uint64 x* ∧ *y* < *p*)

definition *mod-ring-rel64* :: *uint64* ⇒ *'a mod-ring* ⇒ *bool* **where**
mod-ring-rel64 x y = (∃ *z*. *urel64 x z* ∧ *mod-ring-rel z y*)

lemma *urel64-0*: *urel64 0 0* **unfolding** *urel64-def* **using** *p2* **by** (*simp*, *transfer*, *simp*)

lemma *urel64-1*: *urel64 1 1* **unfolding** *urel64-def* **using** *p2* **by** (*simp*, *transfer*, *simp*)

lemma *le-int-of-uint64*: (*x* ≤ *y*) = (*int-of-uint64 x* ≤ *int-of-uint64 y*)
by (*transfer*, *simp add*: *word-le-def*)

lemma *urel64-plus*: **assumes** *urel64 x y urel64 x' y'*
shows *urel64* (*plus-p64 pp x x'*) (*plus-p p y y'*)
proof −
let *?x* = *int-of-uint64 x*
let *?x'* = *int-of-uint64 x'*
let *?p* = *int-of-uint64 pp*
from *assms int-of-uint64-ge-0* **have** *id*: *y* = *?x y'* = *?x'*
and *rel*: *0* ≤ *?x ?x* < *p*
0 ≤ *?x' ?x'* ≤ *p* **unfolding** *urel64-def* **by** *auto*
have *le*: (*pp* ≤ *x* + *x'*) = (*?p* ≤ *?x* + *?x'*) **unfolding** *le-int-of-uint64*
using *rel small* **by** (*auto simp*: *uint64-simps*)
show *?thesis*
proof (*cases ?p* ≤ *?x* + *?x'*)
case *True*
hence *True*: (*?p* ≤ *?x* + *?x'*) = *True* **by** *simp*
show *?thesis* **unfolding** *id*
using *small rel* **unfolding** *plus-p64-def plus-p-def Let-def urel64-def*
unfolding *ppp le True if-True*
using *True* **by** (*auto simp*: *uint64-simps*)
next
case *False*
hence *False*: (*?p* ≤ *?x* + *?x'*) = *False* **by** *simp*
show *?thesis* **unfolding** *id*
using *small rel* **unfolding** *plus-p64-def plus-p-def Let-def urel64-def*
unfolding *ppp le False if-False*
using *False* **by** (*auto simp*: *uint64-simps*)
qed
qed

lemma *urel64-minus*: **assumes** *urel64 x y urel64 x' y'*
shows *urel64* (*minus-p64 pp x x'*) (*minus-p p y y'*)
proof −
let *?x* = *int-of-uint64 x*
let *?x'* = *int-of-uint64 x'*
from *assms int-of-uint64-ge-0* **have** *id*: *y* = *?x y'* = *?x'*
and *rel*: *0* ≤ *?x ?x* < *p*
0 ≤ *?x' ?x'* ≤ *p* **unfolding** *urel64-def* **by** *auto*

```
  have le: (x' ≤ x) = (?x' ≤ ?x) unfolding le-int-of-uint64
    using rel small by (auto simp: uint64-simps)
  show ?thesis
  proof (cases ?x' ≤ ?x)
    case True
    hence True: (?x' ≤ ?x) = True by simp
    show ?thesis unfolding id
      using small rel unfolding minus-p64-def minus-p-def Let-def urel64-def
      unfolding ppp le True if-True
      using True by (auto simp: uint64-simps)
  next
    case False
    hence False: (?x' ≤ ?x) = False by simp
    show ?thesis unfolding id
      using small rel unfolding minus-p64-def minus-p-def Let-def urel64-def
      unfolding ppp le False if-False
      using False by (auto simp: uint64-simps)
  qed
qed

lemma urel64-uminus: assumes urel64 x y
  shows urel64 (uminus-p64 pp x) (uminus-p p y)
proof −
  let ?x = int-of-uint64 x
  from assms int-of-uint64-ge-0 have id: y = ?x
    and rel: 0 ≤ ?x ?x < p
      unfolding urel64-def by auto
  have le: (x = 0) = (?x = 0) unfolding int-of-uint64-0-iff
    using rel small by (auto simp: uint64-simps)
  show ?thesis
  proof (cases ?x = 0)
    case True
    hence True: (?x = 0) = True by simp
    show ?thesis unfolding id
      using small rel unfolding uminus-p64-def uminus-p-def Let-def urel64-def
      unfolding ppp le True if-True
      using True by (auto simp: uint64-simps)
  next
    case False
    hence False: (?x = 0) = False by simp
    show ?thesis unfolding id
      using small rel unfolding uminus-p64-def uminus-p-def Let-def urel64-def
      unfolding ppp le False if-False
      using False by (auto simp: uint64-simps)
  qed
qed

lemma urel64-mult: assumes urel64 x y urel64 x' y'
  shows urel64 (mult-p64 pp x x') (mult-p p y y')
```

```
proof −
  let ?x = int-of-uint64 x
  let ?x' = int-of-uint64 x'
  from assms int-of-uint64-ge-0 have id: y = ?x y' = ?x'
    and rel: 0 ≤ ?x ?x < p
      0 ≤ ?x' ?x' < p unfolding urel64-def by auto
  from rel have ?x * ?x' < p * p by (metis mult-strict-mono')
  also have ... ≤ 4294967296 * 4294967296
    by (rule mult-mono, insert p2 small, auto)
  finally have le: ?x * ?x' < 18446744073709551616 by simp
  show ?thesis unfolding id
      using small rel unfolding mult-p64-def mult-p-def Let-def urel64-def
      unfolding ppp
    by (auto simp: uint64-simps, unfold int-of-uint64-mod int-of-uint64-mult,
        subst mod-pos-pos-trivial[of - 18446744073709551616], insert le, auto)
qed

lemma urel64-eq: assumes urel64 x y urel64 x' y'
  shows (x = x') = (y = y')
proof −
  let ?x = int-of-uint64 x
  let ?x' = int-of-uint64 x'
  from assms int-of-uint64-ge-0 have id: y = ?x y' = ?x'
    unfolding urel64-def by auto
  show ?thesis unfolding id by (transfer, transfer) rule
qed

lemma urel64-normalize:
assumes x: urel64 x y
shows urel64 (if x = 0 then 0 else 1) (if y = 0 then 0 else 1)
 unfolding urel64-eq[OF x urel64-0] using urel64-0 urel64-1 by auto

lemma urel64-mod:
assumes x: urel64 x x' and y: urel64 y y'
shows urel64 (if y = 0 then x else 0) (if y' = 0 then x' else 0)
  unfolding urel64-eq[OF y urel64-0] using urel64-0 x by auto

lemma urel64-power: urel64 x x' ⟹ urel64 y (int y') ⟹ urel64 (power-p64 pp
x y) (power-p p x' y')
including bit-operations-syntax proof (induct x' y' arbitrary: x y rule: power-p.induct[of
- p])
  case (1 x' y' x y)
  note x = 1(2) note y = 1(3)
  show ?case
  proof (cases y' = 0)
    case True
    hence y: y = 0 using urel64-eq[OF y urel64-0] by auto
    show ?thesis unfolding y True by (simp add: power-p.simps urel64-1)
  next
```

```
    case False
    hence id: (y = 0) = False (y' = 0) = False using urel64-eq[OF y urel64-0]
by auto
    from y have ‹int y' = int-of-uint64 y› ‹int y' < p›
      by (simp-all add: urel64-def)
    obtain d' r' where dr': Euclidean-Rings.divmod-nat y' 2 = (d',r') by force
    from Euclidean-Rings.divmod-nat-def[of y' 2, unfolded dr']
    have r': r' = y' mod 2 and d': d' = y' div 2 by auto
    have urel64 (y AND 1) r'
      using ‹int y' < p› small
      apply (simp add: urel64-def and-one-eq r')
      apply (auto simp add: ppp and-one-eq)
    apply (simp add: of-nat-mod int-of-uint64.rep-eq modulo-uint64.rep-eq uint-mod
‹int y' = int-of-uint64 y›)
      done
    from urel64-eq[OF this urel64-0]
    have rem: (y AND 1 = 0) = (r' = 0) by simp
     have div: urel64 (drop-bit 1 y) (int d') unfolding d' using y unfolding
urel64-def using small
      unfolding ppp
      apply transfer
      apply transfer
      apply (auto simp add: drop-bit-Suc take-bit-int-eq-self)
      done
    note IH = 1(1)[OF False refl dr'[symmetric] urel64-mult[OF x x] div]
    show ?thesis unfolding power-p.simps[of - - y'] power-p64.simps[of - - y] dr'
id if-False rem
      using IH urel64-mult[OF IH x] by (auto simp: Let-def)
  qed
qed


lemma urel64-inverse: assumes x: urel64 x x'
  shows urel64 (inverse-p64 pp x) (inverse-p p x')
proof −
  have p: urel64 (pp − 2) (int (nat (p − 2))) using p2 small unfolding urel64-def
unfolding ppp
    by (simp add: int-of-uint64.rep-eq minus-uint64.rep-eq uint-sub-if')
  show ?thesis
   unfolding inverse-p64-def inverse-p-def urel64-eq[OF x urel64-0] using urel64-0
urel64-power[OF x p]
    by auto
qed

lemma mod-ring-0-64: mod-ring-rel64 0 0
  using urel64-0 mod-ring-0 unfolding mod-ring-rel64-def by blast

lemma mod-ring-1-64: mod-ring-rel64 1 1
  using urel64-1 mod-ring-1 unfolding mod-ring-rel64-def by blast
```

lemma *mod-ring-uminus64*: (*mod-ring-rel64* ===> *mod-ring-rel64*) (*uminus-p64 pp*) *uminus*
using *urel64-uminus mod-ring-uminus* **unfolding** *mod-ring-rel64-def rel-fun-def*
by *blast*

lemma *mod-ring-plus64*: (*mod-ring-rel64* ===> *mod-ring-rel64* ===> *mod-ring-rel64*) (*plus-p64 pp*) (+)
using *urel64-plus mod-ring-plus* **unfolding** *mod-ring-rel64-def rel-fun-def* **by** *blast*

lemma *mod-ring-minus64*: (*mod-ring-rel64* ===> *mod-ring-rel64* ===> *mod-ring-rel64*) (*minus-p64 pp*) (−)
using *urel64-minus mod-ring-minus* **unfolding** *mod-ring-rel64-def rel-fun-def* **by** *blast*

lemma *mod-ring-mult64*: (*mod-ring-rel64* ===> *mod-ring-rel64* ===> *mod-ring-rel64*) (*mult-p64 pp*) ((∗))
using *urel64-mult mod-ring-mult* **unfolding** *mod-ring-rel64-def rel-fun-def* **by** *blast*

lemma *mod-ring-eq64*: (*mod-ring-rel64* ===> *mod-ring-rel64* ===> (=)) (=) (=)
using *urel64-eq mod-ring-eq* **unfolding** *mod-ring-rel64-def rel-fun-def* **by** *blast*

lemma *urel64-inj*: *urel64 x y* ⟹ *urel64 x z* ⟹ *y* = *z*
using *urel64-eq*[*of x y x z*] **by** *auto*

lemma *urel64-inj'*: *urel64 x z* ⟹ *urel64 y z* ⟹ *x* = *y*
using *urel64-eq*[*of x z y z*] **by** *auto*

lemma *bi-unique-mod-ring-rel64*:
bi-unique mod-ring-rel64 left-unique mod-ring-rel64 right-unique mod-ring-rel64
using *bi-unique-mod-ring-rel urel64-inj'*
unfolding *mod-ring-rel64-def bi-unique-def left-unique-def right-unique-def*
by (*auto simp*: *urel64-def*)

lemma *right-total-mod-ring-rel64*: *right-total mod-ring-rel64*
unfolding *mod-ring-rel64-def right-total-def*
proof
fix *y* :: *'a mod-ring*
from *right-total-mod-ring-rel*[*unfolded right-total-def*, *rule-format*, *of y*]
obtain *z* **where** *zy*: *mod-ring-rel z y* **by** *auto*
hence *zp*: *0* ≤ *z z* < *p* **unfolding** *mod-ring-rel-def p* **using** *range-to-int-mod-ring*[**where** *'a* = *'a*] **by** *auto*
hence *urel64* (*uint64-of-int z*) *z* **unfolding** *urel64-def* **using** *small* **unfolding** *ppp*
by (*auto simp*: *int-of-uint64-inv*)
with *zy* **show** ∃ *x z*. *urel64 x z* ∧ *mod-ring-rel z y* **by** *blast*

qed

lemma *Domainp-mod-ring-rel64*: *Domainp mod-ring-rel64* = (λ*x*. *0* ≤ *x* ∧ *x* < *pp*)
proof
  fix *x*
  show *Domainp mod-ring-rel64 x* = (*0* ≤ *x* ∧ *x* < *pp*)
    unfolding *Domainp.simps*
    unfolding *mod-ring-rel64-def*
  proof
    let *?i* = *int-of-uint64*
    assume ∗: *0* ≤ *x* ∧ *x* < *pp*
    hence *0* ≤ *?i x* ∧ *?i x* < *p* **using** *small* **unfolding** *ppp*
      by (*transfer*, *auto simp*: *word-less-def*)
    hence *?i x* ∈ {*0* ..< *p*} **by** *auto*
    with *Domainp-mod-ring-rel*
    have *Domainp mod-ring-rel* (*?i x*) **by** *auto*
    from *this*[*unfolded Domainp.simps*]
    obtain *b* **where** *b*: *mod-ring-rel* (*?i x*) *b* **by** *auto*
    show ∃ *a b*. *x* = *a* ∧ (∃ *z*. *urel64 a z* ∧ *mod-ring-rel z b*)
    proof (*intro exI*, *rule conjI*[*OF refl*], *rule exI*, *rule conjI*[*OF - b*])
      show *urel64 x* (*?i x*) **unfolding** *urel64-def* **using** *small* ∗ **unfolding** *ppp*
        by (*transfer*, *auto simp*: *word-less-def*)
    qed
  next
    assume ∃ *a b*. *x* = *a* ∧ (∃ *z*. *urel64 a z* ∧ *mod-ring-rel z b*)
    then obtain *b z* **where** *xz*: *urel64 x z* **and** *zb*: *mod-ring-rel z b* **by** *auto*
    hence *Domainp mod-ring-rel z* **by** *auto*
    with *Domainp-mod-ring-rel* **have** *0* ≤ *z z* < *p* **by** *auto*
    with *xz* **show** *0* ≤ *x* ∧ *x* < *pp* **unfolding** *urel64-def* **using** *small* **unfolding** *ppp*
      by (*transfer*, *auto simp*: *word-less-def*)
  qed
qed

lemma *ring-finite-field-ops64*: *ring-ops* (*finite-field-ops64 pp*) *mod-ring-rel64*
  by (*unfold-locales*, *auto simp*:
  finite-field-ops64-def
  bi-unique-mod-ring-rel64
  right-total-mod-ring-rel64
  mod-ring-plus64
  mod-ring-minus64
  mod-ring-uminus64
  mod-ring-mult64
  mod-ring-eq64
  mod-ring-0-64
  mod-ring-1-64
  Domainp-mod-ring-rel64)
end

end

context *prime-field*
begin
context fixes *pp* :: *uint64*
assumes ∗: *p* = *int-of-uint64 pp p* ≤ *4294967295*
begin

lemma *mod-ring-normalize64*: (*mod-ring-rel64* ===> *mod-ring-rel64*) (λ*x*. *if x* = *0 then 0 else 1*) *normalize*
using *urel64-normalize*[*OF* ∗] *mod-ring-normalize* **unfolding** *mod-ring-rel64-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-mod64*: (*mod-ring-rel64* ===> *mod-ring-rel64* ===> *mod-ring-rel64*) (λ*x y*. *if y* = *0 then x else 0*) (*mod*)
using *urel64-mod*[*OF* ∗] *mod-ring-mod* **unfolding** *mod-ring-rel64-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-unit-factor64*: (*mod-ring-rel64* ===> *mod-ring-rel64*) (λ*x*. *x*) *unit-factor*
using *mod-ring-unit-factor* **unfolding** *mod-ring-rel64-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-inverse64*: (*mod-ring-rel64* ===> *mod-ring-rel64*) (*inverse-p64 pp*) *inverse*
using *urel64-inverse*[*OF* ∗] *mod-ring-inverse* **unfolding** *mod-ring-rel64-def*[*OF* ∗] *rel-fun-def* **by** *blast*

lemma *mod-ring-divide64*: (*mod-ring-rel64* ===> *mod-ring-rel64* ===> *mod-ring-rel64*) (*divide-p64 pp*) (/)
using *mod-ring-inverse64 mod-ring-mult64*[*OF* ∗]
unfolding *divide-p64-def divide-mod-ring-def inverse-mod-ring-def*[*symmetric*] *rel-fun-def* **by** *blast*

lemma *finite-field-ops64*: *field-ops* (*finite-field-ops64 pp*) *mod-ring-rel64*
by (*unfold-locales*, *insert ring-finite-field-ops64*[*OF* ∗], *auto simp*:
ring-ops-def
finite-field-ops64-def
mod-ring-divide64
mod-ring-inverse64
mod-ring-mod64
mod-ring-normalize64)
end
end

context
fixes *p* :: *integer*
begin

```
definition plus-p-integer :: integer ⇒ integer ⇒ integer where
  plus-p-integer x y ≡ let z = x + y in if z ≥ p then z − p else z

definition minus-p-integer :: integer ⇒ integer ⇒ integer where
  minus-p-integer x y ≡ if y ≤ x then x − y else (x + p) − y

definition uminus-p-integer :: integer ⇒ integer where
  uminus-p-integer x = (if x = 0 then 0 else p − x)

definition mult-p-integer :: integer ⇒ integer ⇒ integer where
  mult-p-integer x y = (x * y mod p)


context
  includes bit-operations-syntax
begin

function power-p-integer :: integer ⇒ integer ⇒ integer where
  power-p-integer x n = (if n ≤ 0 then 1 else
    let rec = power-p-integer (mult-p-integer x x) (drop-bit 1 n) in
    if n AND 1 = 0 then rec else mult-p-integer rec x)
  by pat-completeness auto

termination
proof −
  include integer.lifting
  have *: ‹nat-of-integer (n div 2) < nat-of-integer n› if ‹0 < n› for n
    using that by transfer simp
  show ?thesis
    by (relation ‹measure (nat-of-integer ∘ snd)›)
      (simp-all add: not-le drop-bit-Suc *)
qed

end
```

In experiments with Berlekamp-factorization (where the prime p is usually small), it turned out that taking the below implementation of inverse via exponentiation is faster than the one based on the extended Euclidean algorithm.

```
definition inverse-p-integer :: integer ⇒ integer where
  inverse-p-integer x = (if x = 0 then 0 else power-p-integer x (p − 2))

definition divide-p-integer :: integer ⇒ integer ⇒ integer  where
  divide-p-integer x y = mult-p-integer x (inverse-p-integer y)

definition finite-field-ops-integer :: integer arith-ops-record where
  finite-field-ops-integer ≡ Arith-Ops-Record
     0
     1
```

plus-p-integer
mult-p-integer
minus-p-integer
uminus-p-integer
divide-p-integer
inverse-p-integer
$(\lambda\ x\ y\ .\ if\ y = 0\ then\ x\ else\ 0)$
$(\lambda\ x\ .\ if\ x = 0\ then\ 0\ else\ 1)$
$(\lambda\ x\ .\ x)$
integer-of-int
int-of-integer
$(\lambda\ x.\ 0 \leq x \wedge x < p)$
end

For soundness of the integer implementation, we mainly prove that this implementation implements the int-based implementation of GF(p).

context *mod-ring-locale*
begin

context fixes *pp* :: *integer*
assumes *ppp*: $p =$ *int-of-integer pp*
begin

lemma *integer-simps*:
‹*int-of-integer 0* $= 0$›
‹*int-of-integer* $(x + y) =$ *int-of-integer* $x +$ *int-of-integer* y›
‹*int-of-integer* $(x - y) =$ *int-of-integer* $x -$ *int-of-integer* y›
‹*int-of-integer* $(x * y) =$ *int-of-integer* $x *$ *int-of-integer* y›
by *simp-all*

definition *urel-integer* :: *integer* $\Rightarrow$ *int* $\Rightarrow$ *bool* **where** *urel-integer* $x\ y = (y =$ *int-of-integer* $x \wedge y \geq 0 \wedge y < p)$

definition *mod-ring-rel-integer* :: *integer* $\Rightarrow$ $'a$ *mod-ring* $\Rightarrow$ *bool* **where**
mod-ring-rel-integer $x\ y = (\exists\ z.$ *urel-integer* $x\ z \wedge$ *mod-ring-rel* $z\ y)$

lemma *urel-integer-0*: *urel-integer 0 0* **unfolding** *urel-integer-def* **using** *p2* **by** *simp*

lemma *urel-integer-1*: *urel-integer 1 1* **unfolding** *urel-integer-def* **using** *p2* **by** *simp*

lemma *le-int-of-integer*: $(x \leq y) =$ (*int-of-integer* $x \leq$ *int-of-integer* y)
by (*rule less-eq-integer.rep-eq*)

lemma *urel-integer-plus*: **assumes** *urel-integer* $x\ y$ *urel-integer* $x'\ y'$
shows *urel-integer* (*plus-p-integer pp* $x\ x'$) (*plus-p* $p\ y\ y'$)
proof −
let *?x* $=$ *int-of-integer* x

```
  let ?x' = int-of-integer x'
  let ?p = int-of-integer pp
  from assms have id: y = ?x y' = ?x'
    and rel: 0 ≤ ?x ?x < p
      0 ≤ ?x' ?x' ≤ p unfolding urel-integer-def by auto
  have le: (pp ≤ x + x') = (?p ≤ ?x + ?x') unfolding le-int-of-integer
    using rel by auto
  show ?thesis
  proof (cases ?p ≤ ?x + ?x')
    case True
    hence True: (?p ≤ ?x + ?x') = True by simp
    show ?thesis unfolding id
      using rel unfolding plus-p-integer-def plus-p-def Let-def urel-integer-def
      unfolding ppp le True if-True
      using True by auto
  next
    case False
    hence False: (?p ≤ ?x + ?x') = False by simp
    show ?thesis unfolding id
      using rel unfolding plus-p-integer-def plus-p-def Let-def urel-integer-def
      unfolding ppp le False if-False
      using False by auto
  qed
qed

lemma urel-integer-minus: assumes urel-integer x y urel-integer x' y'
  shows urel-integer (minus-p-integer pp x x') (minus-p p y y')
proof −
  let ?x = int-of-integer x
  let ?x' = int-of-integer x'
  from assms have id: y = ?x y' = ?x'
    and rel: 0 ≤ ?x ?x < p
      0 ≤ ?x' ?x' ≤ p unfolding urel-integer-def by auto
  have le: (x' ≤ x) = (?x' ≤ ?x) unfolding le-int-of-integer
    using rel by auto
  show ?thesis
  proof (cases ?x' ≤ ?x)
    case True
    hence True: (?x' ≤ ?x) = True by simp
    show ?thesis unfolding id
      using rel unfolding minus-p-integer-def minus-p-def Let-def urel-integer-def
      unfolding ppp le True if-True
      using True by auto
  next
    case False
    hence False: (?x' ≤ ?x) = False by simp
    show ?thesis unfolding id
      using rel unfolding minus-p-integer-def minus-p-def Let-def urel-integer-def
      unfolding ppp le False if-False
```

```
    using False by auto
  qed
qed

lemma urel-integer-uminus: assumes urel-integer x y
  shows urel-integer (uminus-p-integer pp x) (uminus-p p y)
proof −
  include integer.lifting
  let ?x = int-of-integer x
  from assms have id: y = ?x
    and rel: 0 ≤ ?x ?x < p
      unfolding urel-integer-def by auto
  have le: (x = 0) = (?x = 0)
    by transfer rule
  show ?thesis
  proof (cases ?x = 0)
    case True
    hence True: (?x = 0) = True by simp
    show ?thesis unfolding id
     using rel unfolding uminus-p-integer-def uminus-p-def Let-def urel-integer-def

      unfolding ppp le True if-True
      using True by auto
  next
    case False
    hence False: (?x = 0) = False by simp
    show ?thesis unfolding id
     using rel unfolding uminus-p-integer-def uminus-p-def Let-def urel-integer-def

      unfolding ppp le False if-False
      using False by auto
  qed
qed

lemma pp-pos: int-of-integer pp > 0
  using ppp nontriv[where 'a = 'a]  unfolding p
  by (simp add: less-integer.rep-eq)

lemma urel-integer-mult: assumes urel-integer x y urel-integer x' y'
  shows urel-integer (mult-p-integer pp x x') (mult-p p y y')
proof −
  let ?x = int-of-integer x
  let ?x' = int-of-integer x'
  from assms  have id: y = ?x y' = ?x'
    and rel: 0 ≤ ?x ?x < p
      0 ≤ ?x' ?x' < p unfolding urel-integer-def by auto
  from rel(1,3) have xx: 0 ≤ ?x * ?x' by simp
  show ?thesis unfolding id
    using rel unfolding mult-p-integer-def mult-p-def Let-def urel-integer-def
```

unfolding *ppp mod-nonneg-pos-int*[*OF xx pp-pos*] **using** *xx pp-pos* **by** *simp*

qed

lemma *urel-integer-eq*: **assumes** *urel-integer x y urel-integer x′ y′*
shows $(x = x') = (y = y')$
proof −
let *?x* = *int-of-integer x*
let *?x′* = *int-of-integer x′*
from *assms* **have** *id*: *y* = *?x y′* = *?x′*
unfolding *urel-integer-def* **by** *auto*
show *?thesis* **unfolding** *id integer-eq-iff* **..**
qed

lemma *urel-integer-normalize*:
assumes *x*: *urel-integer x y*
shows *urel-integer* (*if x* = *0 then 0 else 1*) (*if y* = *0 then 0 else 1*)
unfolding *urel-integer-eq*[*OF x urel-integer-0*] **using** *urel-integer-0 urel-integer-1*
by *auto*

lemma *urel-integer-mod*:
assumes *x*: *urel-integer x x′* **and** *y*: *urel-integer y y′*
shows *urel-integer* (*if y* = *0 then x else 0*) (*if y′* = *0 then x′ else 0*)
unfolding *urel-integer-eq*[*OF y urel-integer-0*] **using** *urel-integer-0 x* **by** *auto*

lemma *urel-integer-power*: *urel-integer x x′* ⟹ *urel-integer y* (*int y′*) ⟹ *urel-integer* (*power-p-integer pp x y*) (*power-p p x′ y′*)
including *bit-operations-syntax* **proof** (*induct x′ y′ arbitrary*: *x y rule*: *power-p.induct*[*of - p*])
case (*1 x′ y′ x y*)
note *x* = *1*(*2*) **note** *y* = *1*(*3*)
show *?case*
proof (*cases y′* ≤ *0*)
case *True*
hence *y*: *y* = *0 y′* = *0* **using** *urel-integer-eq*[*OF y urel-integer-0*] **by** *auto*
show *?thesis* **unfolding** *y True* **by** (*simp add*: *power-p.simps urel-integer-1*)
next
case *False*
hence *id*: (*y* ≤ *0*) = *False* (*y′* = *0*) = *False* **using** *False y*
by (*auto simp add*: *urel-integer-def not-le*) (*metis of-int-integer-of of-int-of-nat-eq of-nat-0-less-iff*)
obtain *d′ r′* **where** *dr′*: *Euclidean-Rings.divmod-nat y′ 2* = (*d′*,*r′*) **by** *force*
from *Euclidean-Rings.divmod-nat-def*[*of y′ 2*, *unfolded dr′*]
have *r′*: *r′* = *y′ mod 2* **and** *d′*: *d′* = *y′ div 2* **by** *auto*
have *aux*: ⋀ *y′*. *int* (*y′ mod 2*) = *int y′ mod 2* **by** *presburger*
have *urel-integer* (*y AND 1*) *r′* **unfolding** *r′* **using** *y* **unfolding** *urel-integer-def*

unfolding *ppp*
apply (*auto simp add*: *and-one-eq*)

```
    apply (simp add: of-nat-mod)
    done
  from urel-integer-eq[OF this urel-integer-0]
  have rem: (y AND 1 = 0) = (r' = 0) by simp
  have div: urel-integer (drop-bit 1 y) (int d')
    unfolding d' using y unfolding urel-integer-def
    unfolding ppp by (auto simp add: of-nat-div drop-bit-Suc)
  from id have y' ≠ 0 by auto
  note IH = 1(1)[OF this refl dr'[symmetric] urel-integer-mult[OF x x] div]
 show ?thesis unfolding power-p.simps[of - - y'] power-p-integer.simps[of - - y]
dr' id if-False rem
    using IH urel-integer-mult[OF IH x] by (auto simp: Let-def)
 qed
qed


lemma urel-integer-inverse: assumes x: urel-integer x x'
 shows urel-integer (inverse-p-integer pp x) (inverse-p p x')
proof −
 have p: urel-integer (pp − 2) (int (nat (p − 2))) using p2 unfolding urel-integer-def
unfolding ppp
   by auto
 show ?thesis
  unfolding inverse-p-integer-def inverse-p-def urel-integer-eq[OF x urel-integer-0]
using urel-integer-0 urel-integer-power[OF x p]
   by auto
qed

lemma mod-ring-0--integer: mod-ring-rel-integer 0 0
 using urel-integer-0 mod-ring-0 unfolding mod-ring-rel-integer-def by blast

lemma mod-ring-1--integer: mod-ring-rel-integer 1 1
 using urel-integer-1 mod-ring-1 unfolding mod-ring-rel-integer-def by blast

lemma mod-ring-uminus-integer: (mod-ring-rel-integer ===> mod-ring-rel-integer)
(uminus-p-integer pp) uminus
 using urel-integer-uminus mod-ring-uminus unfolding mod-ring-rel-integer-def
rel-fun-def by blast

lemma mod-ring-plus-integer: (mod-ring-rel-integer ===> mod-ring-rel-integer
===> mod-ring-rel-integer) (plus-p-integer pp) (+)
 using urel-integer-plus mod-ring-plus unfolding mod-ring-rel-integer-def rel-fun-def
by blast

lemma mod-ring-minus-integer: (mod-ring-rel-integer ===> mod-ring-rel-integer
===> mod-ring-rel-integer) (minus-p-integer pp) (−)
 using urel-integer-minus mod-ring-minus unfolding mod-ring-rel-integer-def rel-fun-def
by blast
```

lemma *mod-ring-mult-integer*: (*mod-ring-rel-integer* ===> *mod-ring-rel-integer* ===> *mod-ring-rel-integer*) (*mult-p-integer pp*) ((∗))
using *urel-integer-mult mod-ring-mult* **unfolding** *mod-ring-rel-integer-def rel-fun-def*
by *blast*

lemma *mod-ring-eq-integer*: (*mod-ring-rel-integer* ===> *mod-ring-rel-integer* ===> (=)) (=) (=)
using *urel-integer-eq mod-ring-eq* **unfolding** *mod-ring-rel-integer-def rel-fun-def*
by *blast*

lemma *urel-integer-inj*: *urel-integer x y* ⟹ *urel-integer x z* ⟹ *y* = *z*
using *urel-integer-eq*[*of x y x z*] **by** *auto*

lemma *urel-integer-inj'*: *urel-integer x z* ⟹ *urel-integer y z* ⟹ *x* = *y*
using *urel-integer-eq*[*of x z y z*] **by** *auto*

lemma *bi-unique-mod-ring-rel-integer*:
bi-unique mod-ring-rel-integer left-unique mod-ring-rel-integer right-unique mod-ring-rel-integer
using *bi-unique-mod-ring-rel urel-integer-inj'*
unfolding *mod-ring-rel-integer-def bi-unique-def left-unique-def right-unique-def*
by (*auto simp*: *urel-integer-def*)

lemma *right-total-mod-ring-rel-integer*: *right-total mod-ring-rel-integer*
unfolding *mod-ring-rel-integer-def right-total-def*
proof
fix *y* :: *'a mod-ring*
from *right-total-mod-ring-rel*[*unfolded right-total-def*, *rule-format*, *of y*]
obtain *z* **where** *zy*: *mod-ring-rel z y* **by** *auto*
hence *zp*: *0* ≤ *z z* < *p* **unfolding** *mod-ring-rel-def p* **using** *range-to-int-mod-ring*[**where** *'a* = *'a*] **by** *auto*
hence *urel-integer* (*integer-of-int z*) *z* **unfolding** *urel-integer-def* **unfolding** *ppp*

by *auto*
with *zy* **show** ∃ *x z*. *urel-integer x z* ∧ *mod-ring-rel z y* **by** *blast*
qed

lemma *Domainp-mod-ring-rel-integer*: *Domainp mod-ring-rel-integer* = (λ*x*. *0* ≤ *x* ∧ *x* < *pp*)
proof
fix *x*
show *Domainp mod-ring-rel-integer x* = (*0* ≤ *x* ∧ *x* < *pp*)
unfolding *Domainp.simps*
unfolding *mod-ring-rel-integer-def*
proof
let *?i* = *int-of-integer*
assume ∗: *0* ≤ *x* ∧ *x* < *pp*
hence *0* ≤ *?i x* ∧ *?i x* < *p* **unfolding** *ppp*
by (*simp add*: *le-int-of-integer less-integer.rep-eq*)
hence *?i x* ∈ {*0* ..< *p*} **by** *auto*

```
    with Domainp-mod-ring-rel
    have Domainp mod-ring-rel (?i x) by auto
    from this[unfolded Domainp.simps]
    obtain b where b: mod-ring-rel (?i x) b by auto
    show ∃ a b. x = a ∧ (∃ z. urel-integer a z ∧ mod-ring-rel z b)
    proof (intro exI, rule conjI[OF refl], rule exI, rule conjI[OF - b])
     show urel-integer x (?i x) unfolding urel-integer-def using ∗ unfolding ppp
       by (simp add: le-int-of-integer less-integer.rep-eq)
    qed
  next
    assume ∃ a b. x = a ∧ (∃ z. urel-integer a z ∧ mod-ring-rel z b)
    then obtain b z where xz: urel-integer x z and zb: mod-ring-rel z b by auto
    hence Domainp mod-ring-rel z  by auto
    with Domainp-mod-ring-rel have 0 ≤ z z < p by auto
    with xz show 0 ≤ x ∧ x < pp unfolding urel-integer-def unfolding ppp
      by (simp add: le-int-of-integer less-integer.rep-eq)
  qed
qed

lemma ring-finite-field-ops-integer: ring-ops (finite-field-ops-integer pp) mod-ring-rel-integer
  by (unfold-locales, auto simp:
  finite-field-ops-integer-def
  bi-unique-mod-ring-rel-integer
  right-total-mod-ring-rel-integer
  mod-ring-plus-integer
  mod-ring-minus-integer
  mod-ring-uminus-integer
  mod-ring-mult-integer
  mod-ring-eq-integer
  mod-ring-0--integer
  mod-ring-1--integer
  Domainp-mod-ring-rel-integer)
end
end

context prime-field
begin
context fixes pp :: integer
  assumes ∗: p = int-of-integer pp
begin

lemma mod-ring-normalize-integer: (mod-ring-rel-integer ===> mod-ring-rel-integer)
(λx. if x = 0 then 0 else 1) normalize
  using urel-integer-normalize[OF ∗] mod-ring-normalize unfolding mod-ring-rel-integer-def[OF
∗] rel-fun-def by blast

lemma mod-ring-mod-integer: (mod-ring-rel-integer ===> mod-ring-rel-integer
===> mod-ring-rel-integer) (λx y. if y = 0 then x else 0) (mod)
  using urel-integer-mod[OF ∗] mod-ring-mod unfolding mod-ring-rel-integer-def[OF
```

$*$] *rel-fun-def* **by** *blast*

lemma *mod-ring-unit-factor-integer*: (*mod-ring-rel-integer* ===> *mod-ring-rel-integer*) (λ*x*. *x*) *unit-factor*
using *mod-ring-unit-factor* **unfolding** *mod-ring-rel-integer-def*[*OF* $*$] *rel-fun-def*
by *blast*

lemma *mod-ring-inverse-integer*: (*mod-ring-rel-integer* ===> *mod-ring-rel-integer*) (*inverse-p-integer pp*) *inverse*
using *urel-integer-inverse*[*OF* $*$] *mod-ring-inverse* **unfolding** *mod-ring-rel-integer-def*[*OF* $*$] *rel-fun-def* **by** *blast*

lemma *mod-ring-divide-integer*: (*mod-ring-rel-integer* ===> *mod-ring-rel-integer* ===> *mod-ring-rel-integer*) (*divide-p-integer pp*) (/)
using *mod-ring-inverse-integer mod-ring-mult-integer*[*OF* $*$]
unfolding *divide-p-integer-def divide-mod-ring-def inverse-mod-ring-def*[*symmetric*] *rel-fun-def* **by** *blast*

lemma *finite-field-ops-integer*: *field-ops* (*finite-field-ops-integer pp*) *mod-ring-rel-integer*
by (*unfold-locales*, *insert ring-finite-field-ops-integer*[*OF* $*$], *auto simp*:
ring-ops-def
finite-field-ops-integer-def
mod-ring-divide-integer
mod-ring-inverse-integer
mod-ring-mod-integer
mod-ring-normalize-integer)
end
end

context *prime-field*
begin

thm
finite-field-ops64
finite-field-ops32
finite-field-ops-integer
finite-field-ops-int
end

context *mod-ring-locale*
begin

thm
ring-finite-field-ops64
ring-finite-field-ops32
ring-finite-field-ops-integer
ring-finite-field-ops-int
end

end

3.2 Matrix Operations in Fields

We use our record based description of a field to perform matrix operations.

theory *Matrix-Record-Based*
imports
Jordan-Normal-Form.Gauss-Jordan-Elimination
Jordan-Normal-Form.Gauss-Jordan-IArray-Impl
Arithmetic-Record-Based
begin

definition *mat-rel* :: ('*a* ⇒ '*b* ⇒ *bool*) ⇒ '*a mat* ⇒ '*b mat* ⇒ *bool* **where**
mat-rel R A B ≡ *dim-row A* = *dim-row B* ∧ *dim-col A* = *dim-col B* ∧
(∀ *i j*. *i* < *dim-row B* ⟶ *j* < *dim-col B* ⟶ *R* (*A* $$ (*i*,*j*)) (*B* $$ (*i*,*j*)))

lemma *right-total-mat-rel*: *right-total R* ⟹ *right-total* (*mat-rel R*)
unfolding *right-total-def*
proof
fix *B*
assume ∀ *y*. ∃ *x*. *R x y*
from *choice*[*OF this*] **obtain** *f* **where** *f*: ⋀ *x*. *R* (*f x*) *x* **by** *auto*
show ∃ *A*. *mat-rel R A B*
by (*rule exI*[*of* - *map-mat f B*], *unfold mat-rel-def*, *auto simp*: *f*)
qed

lemma *left-unique-mat-rel*: *left-unique R* ⟹ *left-unique* (*mat-rel R*)
unfolding *left-unique-def mat-rel-def mat-eq-iff* **by** (*auto*, *blast*)

lemma *right-unique-mat-rel*: *right-unique R* ⟹ *right-unique* (*mat-rel R*)
unfolding *right-unique-def mat-rel-def mat-eq-iff* **by** (*auto*, *blast*)

lemma *bi-unique-mat-rel*: *bi-unique R* ⟹ *bi-unique* (*mat-rel R*)
using *left-unique-mat-rel*[*of R*] *right-unique-mat-rel*[*of R*]
unfolding *bi-unique-def left-unique-def right-unique-def* **by** *blast*

lemma *mat-rel-eq*: ((*R* ===> *R* ===> (=))) (=) (=) ⟹
((*mat-rel R* ===> *mat-rel R* ===> (=))) (=) (=)
unfolding *mat-rel-def rel-fun-def mat-eq-iff* **by** (*auto*, *blast*+)

definition *vec-rel* :: ('*a* ⇒ '*b* ⇒ *bool*) ⇒ '*a vec* ⇒ '*b vec* ⇒ *bool* **where**
vec-rel R A B ≡ *dim-vec A* = *dim-vec B* ∧ (∀ *i*. *i* < *dim-vec B* ⟶ *R* (*A* $ *i*)
(*B* $ *i*))

lemma *right-total-vec-rel*: *right-total R* ⟹ *right-total* (*vec-rel R*)
unfolding *right-total-def*
proof
fix *B*

assume $\forall\ y.\ \exists\ x.\ R\ x\ y$
from *choice*[*OF this*] **obtain** f **where** f: $\bigwedge x.\ R\ (f\ x)\ x$ **by** *auto*
show $\exists\ A.$ *vec-rel R A B*
by (*rule exI*[*of - map-vec f B*], *unfold vec-rel-def*, *auto simp*: f)
qed

lemma *left-unique-vec-rel*: *left-unique R* $\Longrightarrow$ *left-unique* (*vec-rel R*)
unfolding *left-unique-def vec-rel-def vec-eq-iff* **by** *auto*

lemma *right-unique-vec-rel*: *right-unique R* $\Longrightarrow$ *right-unique* (*vec-rel R*)
unfolding *right-unique-def vec-rel-def vec-eq-iff* **by** *auto*

lemma *bi-unique-vec-rel*: *bi-unique R* $\Longrightarrow$ *bi-unique* (*vec-rel R*)
using *left-unique-vec-rel*[*of R*] *right-unique-vec-rel*[*of R*]
unfolding *bi-unique-def left-unique-def right-unique-def* **by** *blast*

lemma *vec-rel-eq*: $((R$ ===> R ===> $(=)))\ (=)\ (=) \Longrightarrow$
$((vec\text{-}rel\ R$ ===> $vec\text{-}rel\ R$ ===> $(=)))\ (=)\ (=)$
unfolding *vec-rel-def rel-fun-def vec-eq-iff* **by** (*auto*, *blast+*)

lemma *multrow-transfer*[*transfer-rule*]: $((R$ ===> R ===> $R)$ ===> $(=)$ ===>
R
===> *mat-rel R* ===> *mat-rel R*) *mat-multrow-gen mat-multrow-gen*
unfolding *mat-rel-def*[*abs-def*] *mat-multrow-gen-def*[*abs-def*]
by (*intro rel-funI conjI allI impI eq-matI*, *auto simp*: *rel-fun-def*)

lemma *swap-rows-transfer*: *mat-rel R A B* $\Longrightarrow$ $i <$ *dim-row B* $\Longrightarrow$ $j <$ *dim-row*
B $\Longrightarrow$
mat-rel R (*mat-swaprows i j A*) (*mat-swaprows i j B*)
unfolding *mat-rel-def mat-swaprows-def*
by (*intro rel-funI conjI allI impI eq-matI*, *auto*)

lemma *pivot-positions-gen-transfer*: **assumes** [*transfer-rule*]: $(R$ ===> R ===>
$(=))\ (=)\ (=)$
shows
$(R$ ===> *mat-rel R* ===> $(=))$ *pivot-positions-gen pivot-positions-gen*
proof (*intro rel-funI*, *goal-cases*)
case (*1 ze ze' A A'*)
note *trans*[*transfer-rule*] = *1*
from *1* **have** *dim*: *dim-row A* = *dim-row A' dim-col A* = *dim-col A'* **unfolding**
mat-rel-def **by** *auto*
obtain *i j* **where** *id*: $i = 0\ j = 0$ **and** *ij*: $i \leq$ *dim-row A'* $j \leq$ *dim-col A'* **by**
auto
have *pivot-positions-main-gen ze A* (*dim-row A*) (*dim-col A*) *i j* =
pivot-positions-main-gen ze' A' (*dim-row A'*) (*dim-col A'*) *i j*
using *ij*
proof (*induct i j rule*: *pivot-positions-main-gen.induct*[*of dim-row A' dim-col A'*
A' ze'])

```
    case (1 i j)
    note simps[simp] = pivot-positions-main-gen.simps[of - - - - i j]
    show ?case
    proof (cases i < dim-row A' ∧ j < dim-col A')
      case False
      with dim show ?thesis by auto
    next
      case True
      hence ij: i < dim-row A' j < dim-col A' and j: Suc j ≤ dim-col A' by auto
      note IH = 1(1−2)[OF ij - - j]
      from ij True trans have [transfer-rule]:R (A $$ (i,j)) (A' $$ (i,j))
        unfolding mat-rel-def by auto
      have eq: (A $$ (i,j) = ze) = (A' $$ (i,j) = ze') by transfer-prover
      show ?thesis
      proof (cases A' $$ (i,j) = ze')
        case True
        from ij have i ≤ dim-row A' by auto
        note IH = IH(1)[OF True this]
        thus ?thesis using True ij dim eq by simp
      next
        case False
        from ij have Suc i ≤ dim-row A' by auto
        note IH = IH(2)[OF False this]
        thus ?thesis using False ij dim eq by simp
      qed
    qed
  qed
  thus pivot-positions-gen ze A = pivot-positions-gen ze' A'
    unfolding pivot-positions-gen-def id .
qed

lemma set-pivot-positions-main-gen:
  set (pivot-positions-main-gen ze A nr nc i j) ⊆ {0 ..< nr} × {0 ..< nc}
proof (induct i j rule: pivot-positions-main-gen.induct[of nr nc A ze])
  case (1 i j)
  note [simp] = pivot-positions-main-gen.simps[of - - - - i j]
  from 1 show ?case
    by (cases i < nr ∧ j < nc, auto)
qed

lemma find-base-vectors-transfer: assumes [transfer-rule]: (R ===> R ===>
(=)) (=) (=)
  shows ((R ===> R) ===> R ===> R ===> mat-rel R
  ===> list-all2 (vec-rel R)) find-base-vectors-gen find-base-vectors-gen
proof (intro rel-funI, goal-cases)
  case (1 um um' ze ze' on on' A A')
  note trans[transfer-rule] = 1 pivot-positions-gen-transfer[OF assms]
  from 1(4) have dim: dim-row A = dim-row A' dim-col A = dim-col A' unfolding
mat-rel-def by auto
```

```
  have id: pivot-positions-gen ze A = pivot-positions-gen ze' A' by transfer-prover
  obtain xs where xs: map snd (pivot-positions-gen ze' A') = xs by auto
  obtain ys where ys: [j←[0..<dim-col A'] . j ∉ set xs] = ys by auto
  show list-all2 (vec-rel R) (find-base-vectors-gen um ze on A)
    (find-base-vectors-gen um' ze' on' A')
   unfolding find-base-vectors-gen-def Let-def id xs list-all2-conv-all-nth length-map
ys dim
  proof (intro conjI[OF refl] allI impI)
    fix i
    assume i: i < length ys
    define y where y = ys ! i
   from i have y: y < dim-col A' unfolding y-def ys[symmetric] using nth-mem
by fastforce
    let ?map = map-of (map prod.swap (pivot-positions-gen ze' A'))
    {
      fix i
      assume i: i < dim-col A'
      and neq: i ≠ y
      have R (case ?map i of None ⇒ ze | Some j ⇒ um (A $$ (j, y)))
          (case ?map i of None ⇒ ze' | Some j ⇒ um' (A' $$ (j, y)))
      proof (cases ?map i)
        case None
        with trans(2) show ?thesis by auto
      next
        case (Some j)
        from map-of-SomeD[OF this] have (j,i) ∈ set (pivot-positions-gen ze' A')
by auto
      from subsetD[OF set-pivot-positions-main-gen this[unfolded pivot-positions-gen-def]]
        have j: j < dim-row A' by auto
       with trans(4) y have [transfer-rule]: R (A $$ (j,y)) (A' $$ (j,y)) unfolding
mat-rel-def by auto
        show ?thesis unfolding Some by (simp, transfer-prover)
      qed
    } note main = this
    show vec-rel R (map (non-pivot-base-gen um ze on A (pivot-positions-gen ze'
A')) ys ! i)
          (map (non-pivot-base-gen um' ze' on' A' (pivot-positions-gen ze' A')) ys !
i)
      unfolding y-def[symmetric] nth-map[OF i]
      unfolding non-pivot-base-gen-def Let-def dim vec-rel-def
      by (intro conjI allI impI, force, insert main, auto simp: trans(3))
  qed
qed


lemma eliminate-entries-gen-transfer: assumes *[transfer-rule]: (R ===> R ===>
R) ad ad'
  (R ===> R ===> R) mul mul'
  and vs: ⋀ j. j < dim-row B' ⟹ R (vs j) (vs' j)
```

and *i*: *i* < *dim-row B′*
and *B*: *mat-rel R B B′*
shows *mat-rel R*
(*eliminate-entries-gen ad mul vs B i j*)
(*eliminate-entries-gen ad′ mul′ vs′ B′ i j*)
proof −
note *BB* = *B*[*unfolded mat-rel-def*]
show *?thesis* **unfolding** *mat-rel-def dim-eliminate-entries-gen*
proof (*intro conjI impI allI*)
fix *i′ j′*
assume *ij′*: *i′* < *dim-row B′ j′* < *dim-col B′*
with *BB* **have** *ij*: *i′*< *dim-row B j′* < *dim-col B* **by** *auto*
have [*transfer-rule*]: *R* (*B* $$ (*i′*, *j′*)) (*B′* $$ (*i′*, *j′*)) **using** *BB ij′* **by** *auto*
have [*transfer-rule*]: *R* (*B* $$ (*i*, *j′*)) (*B′* $$ (*i*, *j′*)) **using** *BB ij′ i* **by** *auto*
have [*transfer-rule*]: *R* (*vs i′*) (*vs′ i′*) **using** *ij′ vs*[*of i′*] **by** *auto*
show *R* (*eliminate-entries-gen ad mul vs B i j* $$ (*i′*, *j′*))
(*eliminate-entries-gen ad′ mul′ vs′ B′ i j* $$ (*i′*, *j′*))
unfolding *eliminate-entries-gen-def index-mat*(*1*)[*OF ij*] *index-mat*(*1*)[*OF ij′*]
split
by *transfer-prover*
qed (*insert BB*, *auto*)
qed

context
fixes *ops* :: *′i arith-ops-record* (**structure**)
begin
private abbreviation (*input*) *zero* **where** *zero* ≡ *arith-ops-record.zero ops*
private abbreviation (*input*) *one* **where** *one* ≡ *arith-ops-record.one ops*
private abbreviation (*input*) *plus* **where** *plus* ≡ *arith-ops-record.plus ops*
private abbreviation (*input*) *times* **where** *times* ≡ *arith-ops-record.times ops*
private abbreviation (*input*) *minus* **where** *minus* ≡ *arith-ops-record.minus ops*
private abbreviation (*input*) *uminus* **where** *uminus* ≡ *arith-ops-record.uminus ops*
private abbreviation (*input*) *divide* **where** *divide* ≡ *arith-ops-record.divide ops*
private abbreviation (*input*) *inverse* **where** *inverse* ≡ *arith-ops-record.inverse ops*
private abbreviation (*input*) *modulo* **where** *modulo* ≡ *arith-ops-record.modulo ops*
private abbreviation (*input*) *normalize* **where** *normalize* ≡ *arith-ops-record.normalize ops*

definition *eliminate-entries-gen-zero* :: (*′a* ⇒ *′a* ⇒ *′a*) ⇒ (*′a* ⇒ *′a* ⇒ *′a*) ⇒ *′a* ⇒ (*integer* ⇒ *′a*) ⇒ *′a mat* ⇒ *nat* ⇒ *nat* ⇒ *′a mat* **where**
eliminate-entries-gen-zero minu time z v A I J = *mat* (*dim-row A*) (*dim-col A*) (λ (*i*, *j*).
if v (*integer-of-nat i*) ≠ *z* ∧ *i* ≠ *I then minu* (*A* $$ (*i*,*j*)) (*time* (*v* (*integer-of-nat i*)) (*A* $$ (*I*,*j*))) *else A* $$ (*i*,*j*))

definition *eliminate-entries-i* **where** *eliminate-entries-i* ≡ *eliminate-entries-gen-zero*

minus times zero
definition *multrow-i* **where** *multrow-i* ≡ *mat-multrow-gen times*

lemma *dim-eliminate-entries-gen-zero*[*simp*]:
dim-row (*eliminate-entries-gen-zero mm tt z v B i as*) = *dim-row B*
dim-col (*eliminate-entries-gen-zero mm tt z v B i as*) = *dim-col B*
unfolding *eliminate-entries-gen-zero-def* **by** *auto*

partial-function (*tailrec*) *gauss-jordan-main-i* :: *nat* ⇒ *nat* ⇒ *'i mat* ⇒ *nat* ⇒ *nat* ⇒ *'i mat* **where**
[*code*]: *gauss-jordan-main-i nr nc A i j* = (
if i < *nr* ∧ *j* < *nc then let aij* = *A* $$ (*i*,*j*) *in if aij* = *zero then*
(*case* [*i'* . *i'* <− [*Suc i* ..< *nr*], *A* $$ (*i'*,*j*) ≠ *zero*]
of [] ⇒ *gauss-jordan-main-i nr nc A i* (*Suc j*)
| (*i'* # -) ⇒ *gauss-jordan-main-i nr nc* (*swaprows i i' A*) *i j*)
else if aij = *one then let*
v = (λ *i*. *A* $$ (*nat-of-integer i*,*j*)) *in*
gauss-jordan-main-i nr nc
(*eliminate-entries-i v A i j*) (*Suc i*) (*Suc j*)
else let iaij = *inverse aij*; *A'* = *multrow-i i iaij A*;
v = (λ *i*. *A'* $$ (*nat-of-integer i*,*j*))
in gauss-jordan-main-i nr nc (*eliminate-entries-i v A' i j*) (*Suc i*) (*Suc j*)
else A)

definition *gauss-jordan-single-i* :: *'i mat* ⇒ *'i mat* **where**
gauss-jordan-single-i A ≡ *gauss-jordan-main-i* (*dim-row A*) (*dim-col A*) *A 0 0*

definition *find-base-vectors-i* :: *'i mat* ⇒ *'i vec list* **where**
find-base-vectors-i A ≡ *find-base-vectors-gen uminus zero one A*
end

context *field-ops*
begin

lemma *right-total-poly-rel*[*transfer-rule*]: *right-total* (*mat-rel R*)
using *right-total-mat-rel*[*of R*] *right-total* **.**

lemma *bi-unique-poly-rel*[*transfer-rule*]: *bi-unique* (*mat-rel R*)
using *bi-unique-mat-rel*[*of R*] *bi-unique* **.**

lemma *eq-mat-rel*[*transfer-rule*]: (*mat-rel R* ===> *mat-rel R* ===> (=)) (=) (=)
by (*rule mat-rel-eq*[*OF eq*])

lemma *multrow-i*[*transfer-rule*]: ((=) ===> *R* ===> *mat-rel R* ===> *mat-rel R*)

(*multrow-i ops*) *multrow*
using *multrow-transfer*[*of R*] *times* **unfolding** *multrow-i-def rel-fun-def* **by** *blast*

lemma *eliminate-entries-gen-zero*[*simp*]:
assumes *mat-rel R A A' I* < *dim-row A'* **shows**
eliminate-entries-gen-zero minus times zero v A I J = *eliminate-entries-gen minus times* (*v o integer-of-nat*) *A I J*
unfolding *eliminate-entries-gen-def eliminate-entries-gen-zero-def*
proof(*standard*,*goal-cases*)
case (*1 i j*)
have *d1*:*DP* (*A* $$ (*I*, *j*)) **and** *d2*:*DP* (*A* $$ (*i*, *j*)) **using** *assms DPR 1*
unfolding *mat-rel-def dim-col-mat dim-row-mat*
by (*metis Domainp.DomainI*)+
have *e1*:⋀ *x*. (*0*::'*a*) * *x* = *0* **and** *e2*:⋀ *x*. *x* − (*0*::'*a*) = *x* **by** *auto*
from *e1*[*untransferred*,*OF d1*] *e2*[*untransferred*,*OF d2*] *1* **show** *?case* **by** *auto*
qed *auto*

lemma *eliminate-entries-i*: **assumes**
vs: ⋀ *j*. *j* < *dim-row B'* ⟹ *R* (*vs* (*integer-of-nat j*)) (*vs' j*)
and *i*: *i* < *dim-row B'*
and *B*: *mat-rel R B B'*
shows *mat-rel R* (*eliminate-entries-i ops vs B i j*)
(*eliminate-entries vs' B' i j*)
unfolding *eliminate-entries-i-def eliminate-entries-gen-zero*[*OF B i*]
by (*rule eliminate-entries-gen-transfer*, *insert assms*, *auto simp*: *plus times minus*)

lemma *gauss-jordan-main-i*:
nr = *dim-row A'* ⟹ *nc* = *dim-col A'* ⟹ *mat-rel R A A'* ⟹ *i* ≤ *nr* ⟹ *j* ≤ *nc* ⟹
mat-rel R (*gauss-jordan-main-i ops nr nc A i j*) (*fst* (*gauss-jordan-main A' B' i j*))
proof −
obtain *P* **where** *P*: *P* = (*A'*,*i*,*j*) **by** *auto*
let *?Rel* = *measures* [λ (*A'* :: '*a mat*,*i*,*j*). *nc* − *j*, λ (*A'*,*i*,*j*). *if A'* $$ (*i*,*j*) = *0 then 1 else 0*]
have *wf*: *wf ?Rel* **by** *simp*
show *nr* = *dim-row A'* ⟹ *nc* = *dim-col A'* ⟹ *mat-rel R A A'* ⟹ *i* ≤ *nr* ⟹ *j* ≤ *nc* ⟹
mat-rel R (*gauss-jordan-main-i ops nr nc A i j*) (*fst* (*gauss-jordan-main A' B' i j*))
using *P*
proof (*induct P arbitrary*: *A' B' A i j rule*: *wf-induct*[*OF wf*])
case (*1 P A' B' A i j*)
note *prems* = *1*(*2*−*6*)
note *P* = *1*(*7*)
note *A*[*transfer-rule*] = *prems*(*3*)
note *IH* = *1*(*1*)[*rule-format*, *OF - - - - - - refl*]

```
    note simps = gauss-jordan-main-code[of A' B' i j, unfolded Let-def, folded
prems(1−2)]
     gauss-jordan-main-i.simps[of ops nr nc A i j] Let-def if-True if-False
   show ?case
   proof (cases i < nr ∧ j < nc)
     case False
     hence id: (i < nr ∧ j < nc) = False by simp
     show ?thesis unfolding simps id by simp transfer-prover
   next
     case True note ij' = this
    hence id: (i < nr ∧ j < nc) = True ⋀ x y z. (if x = x then y else z) = y by
auto
     from True prems have ij [transfer-rule]:R (A $$ (i,j)) (A' $$ (i,j))
       unfolding mat-rel-def by auto
     from True prems have i: i < dim-row A' j < dim-col A' and i': i < nr j <
nc by auto
     {
       fix i
       assume i < dim-row A'
       with i True prems have R[transfer-rule]:R (A $$ (i,j)) (A' $$ (i,j))
         unfolding mat-rel-def by auto
       have (A $$ (i,j) = zero) = (A' $$ (i,j) = 0) by transfer-prover
       note this R
     } note eq-gen = this
     have eq: (A $$ (i,j) = zero) = (A' $$ (i,j) = 0)
       (A $$ (i,j) = one) = (A' $$ (i,j) = 1)
       by transfer-prover+
     show ?thesis
     proof (cases A' $$ (i, j) = 0)
       case True
       hence eq: A $$ (i,j) = zero using eq by auto
       let ?is = [ i' . i' <− [Suc i ..< nr],  A $$ (i',j) ≠ zero]
       let ?is' = [ i' . i' <− [Suc i ..< nr],  A' $$ (i',j) ≠ 0]
       define xs where xs = [Suc i..<nr]
        have xs: set xs ⊆ {0 ..< dim-row A'} unfolding xs-def using prems by
auto
       hence id': ?is = ?is' unfolding xs-def[symmetric]
         by (induct xs, insert eq-gen, auto)
       show ?thesis
       proof (cases ?is')
         case Nil
         have ?thesis = (mat-rel R (gauss-jordan-main-i ops nr nc A i (Suc j))
           (fst (gauss-jordan-main A' B' i (Suc j))))
           unfolding True simps id eq unfolding Nil id'[unfolded Nil] by simp
         also have ...
           by (rule IH, insert i prems P, auto)
         finally show ?thesis .
       next
         case (Cons i' idx')
```

```
        from arg-cong[OF this, of set] i
        have i': i' < nr A' $$ (i', j) ≠ 0 by auto
       with ij' prems(1−2) have ∗: i' < dim-row A' i < dim-row A' j < dim-col
A' by auto
        have rel: ((swaprows i i' A', i, j), P) ∈ ?Rel
          by (simp add: P True ∗ i')
        have ?thesis = (mat-rel R (gauss-jordan-main-i ops nr nc (swaprows i i'
A) i j)
          (fst (gauss-jordan-main (swaprows i i' A') (swaprows i i' B') i j)))
           unfolding True simps id eq Cons id'[unfolded Cons] by simp
        also have ...
        by (rule IH[OF rel - - swap-rows-transfer], insert i i' prems P True, auto)
        finally show ?thesis .
       qed
     next
       case False
       from False eq have neq: (A $$ (i, j) = zero) = False (A' $$ (i, j) = 0) =
False by auto
       {
        fix B B' i
        assume B[transfer-rule]: mat-rel R B B' and dim: dim-col B' = nc and
i: i < dim-row B'
        from dim i True have j < dim-col B' by simp
        with B i have R (B $$ (i,j)) (B' $$ (i,j))
          by (simp add: mat-rel-def)
       } note vec-rel = this
      from prems have dim: dim-row A = dim-row A' unfolding mat-rel-def by
auto
       show ?thesis
       proof (cases A' $$ (i, j) = 1)
         case True
          from True eq have eq: (A $$ (i,j) = one) = True (A' $$ (i,j) = 1) =
True by auto
         note rel = vec-rel[OF A]
         show ?thesis unfolding simps id neq eq
            by (rule IH[OF - - - eliminate-entries-i], insert rel prems ij i P dim,
auto)
       next
         case False
          from False eq have eq: (A $$ (i,j) = one) = False (A' $$ (i,j) = 1) =
False by auto
         show ?thesis unfolding simps id neq eq
         proof (rule IH, goal-cases)
           case 4
           have A': mat-rel R (multrow-i ops i (inverse (A $$ (i, j))) A)
             (multrow i (inverse-class.inverse (A' $$ (i, j))) A') by transfer-prover
           note rel = vec-rel[OF A']
           show ?case
             by (rule eliminate-entries-i[OF - - A'], insert rel prems i dim, auto)
```

```
        qed (insert prems i P, auto)
      qed
    qed
  qed
 qed
qed

lemma gauss-jordan-i[transfer-rule]:
 (mat-rel R ===> mat-rel R) (gauss-jordan-single-i ops) gauss-jordan-single
proof (intro rel-funI)
 fix A A'
 assume A: mat-rel R A A'
 show mat-rel R (gauss-jordan-single-i ops A) (gauss-jordan-single A')
  unfolding gauss-jordan-single-def gauss-jordan-single-i-def gauss-jordan-def
  by (rule gauss-jordan-main-i[OF - - A], insert A, auto simp: mat-rel-def)
qed

lemma find-base-vectors-i[transfer-rule]:
 (mat-rel R ===> list-all2 (vec-rel R)) (find-base-vectors-i ops) find-base-vectors
 unfolding find-base-vectors-i-def[abs-def]
 using find-base-vectors-transfer[OF eq] uminus zero one
 unfolding rel-fun-def by blast

end

lemma list-of-vec-transfer[transfer-rule]: (vec-rel A ===> list-all2 A) list-of-vec
list-of-vec
 unfolding rel-fun-def vec-rel-def vec-eq-iff list-all2-conv-all-nth
 by auto

lemma IArray-sub'[simp]: i < IArray.length a ⟹ IArray.sub' (a, integer-of-nat
i) = IArray.sub a i
 by auto

lift-definition eliminate-entries-i2 ::
 'a ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ (integer ⇒ 'a) ⇒ 'a mat-impl ⇒
integer ⇒ 'a mat-impl is
 λ z mminus ttimes v (nr, nc, a) i'.
  (nr,nc,let ai' = IArray.sub' (a, i') in (IArray.tabulate (integer-of-nat nr, λ i. let
ai = IArray.sub' (a, i) in
    if i = i' then ai else
    let vi'j = v i
    in if vi'j = z then ai
       else
            IArray.tabulate (integer-of-nat nc, λ j. mminus (IArray.sub' (ai, j))
(ttimes vi'j
           (IArray.sub' (ai', j))))
     )))
proof(goal-cases)
```

case (*1 z mm tt vec prod nat2*)
thus *?case* **by**(*cases prod*;*cases snd* (*snd prod*);*auto simp*:*Let-def*)
qed

lemma *eliminate-entries-gen-zero* [*simp*]:
assumes *i*<(*dim-row A*) *j*<(*dim-col A*) **shows**
eliminate-entries-gen-zero mminus ttimes z v A I J $$ (*i*, *j*) =
(*if v* (*integer-of-nat i*) = *z* ∨ *i* = *I then A* $$ (*i*,*j*) *else mminus* (*A* $$ (*i*,*j*)) (*ttimes* (*v* (*integer-of-nat i*)) (*A* $$ (*I*,*j*))))
using *assms* **unfolding** *eliminate-entries-gen-zero-def* **by** *auto*

lemma *eliminate-entries-gen* [*simp*]:
assumes *i*<(*dim-row A*) *j*<(*dim-col A*) **shows**
eliminate-entries-gen mminus ttimes v A I J $$ (*i*, *j*) =
(*if i* = *I then A* $$ (*i*,*j*) *else mminus* (*A* $$ (*i*,*j*)) (*ttimes* (*v i*) (*A* $$ (*I*,*j*))))
using *assms* **unfolding** *eliminate-entries-gen-def* **by** *auto*

lemma *dim-mat-impl* [*simp*]:
dim-row (*mat-impl x*) = *dim-row-impl x*
dim-col (*mat-impl x*) = *dim-col-impl x*
by (*cases Rep-mat-impl x*;*auto simp*:*mat-impl.rep-eq dim-row-def dim-col-def dim-row-impl.rep-eq dim-col-impl.rep-eq*)+

lemma *dim-eliminate-entries-i2* [*simp*]:
dim-row-impl (*eliminate-entries-i2 z mm tt v m i*) = *dim-row-impl m*
dim-col-impl (*eliminate-entries-i2 z mm tt v m i*) = *dim-col-impl m*
by (*transfer*, *auto*)+

lemma *tabulate-nth*: *i* < *n* ⟹ *IArray.tabulate* (*integer-of-nat n*, *f*) !! *i* = *f* (*integer-of-nat i*)
using *of-fun-nth*[*of i n*] **by** *auto*

lemma *eliminate-entries-i2*[*code*]:*eliminate-entries-gen-zero mm tt z v* (*mat-impl m*) *i j*
= (*if i* < *dim-row-impl m*
then mat-impl (*eliminate-entries-i2 z mm tt v m* (*integer-of-nat i*))
else (*Code.abort* (*STR ''index out of range in eliminate-entries''*)
(λ -. *eliminate-entries-gen-zero mm tt z v* (*mat-impl m*) *i j*)))
proof (*cases i* < *dim-row-impl m*)
case *True*
hence *id*: (*i* < *dim-row-impl m*) = *True* **by** *simp*
show *?thesis* **unfolding** *id if-True*
proof (*standard*;*goal-cases*)
case (*1 i j*)
have *dims*: *i* < *dim-row* (*mat-impl m*) *j* < *dim-col* (*mat-impl m*) **using** *1* **by** (*auto simp*:*eliminate-entries-i2.rep-eq*)
then show *?case* **unfolding** *eliminate-entries-gen-zero*[*OF dims*] **using** *True*
proof(*transfer*, *goal-cases*)

```
      case (1 i m j ia v z mm tt)
      obtain nr nc M where m: m = (nr,nc,M) by (cases m)
      note 1 = 1[unfolded m, simplified]
      have mk: ⋀ f. mk-mat nr nc f (i,j) = f (i,j)
         ⋀ f. mk-mat nr nc f (ia,j) = f (ia,j)
        using 1 unfolding mk-mat-def mk-vec-def by auto
     note of-fun = of-fun-nth[OF 1(2)] of-fun-nth[OF 1(3)] tabulate-nth[OF 1(2)]
tabulate-nth[OF 1(3)]
      let ?c1 = v (integer-of-nat i) = z
      show ?case
      proof (cases ?c1 ∨ i = ia)
        case True
        hence id: (if ?c1 ∨ i = ia then x else y) = x
          (if integer-of-nat i = integer-of-nat ia then x else if ?c1 then x else y) = x
for x y
          by auto
       show ?thesis unfolding id m o-def Let-def split snd-conv mk of-fun by (auto
simp: 1)
      next
        case False
        hence id: ?c1 = False  (integer-of-nat i = integer-of-nat ia) = False (False
∨ i = ia) = False
          by (auto simp add: integer-of-nat-eq-of-nat)
       show ?thesis unfolding m o-def Let-def split snd-conv mk of-fun id if-False
          by (auto simp: 1)
      qed
    qed
  qed (auto simp:eliminate-entries-i2.rep-eq)
qed auto

end
theory More-Missing-Multiset
  imports
    HOL−Combinatorics.Permutations
    Polynomial-Factorization.Missing-Multiset
begin

lemma rel-mset-free:
  assumes rel: rel-mset rel X Y and xs: mset xs = X
  shows ∃ys. mset ys = Y ∧ list-all2 rel xs ys
proof−
  from rel[unfolded rel-mset-def] obtain xs′ ys′
    where xs′: mset xs′ = X and ys′: mset ys′ = Y and xsys′: list-all2 rel xs′ ys′
by auto
  from xs′ xs have mset xs = mset xs′ by auto
  from mset-eq-permutation[OF this]
  obtain f where perm: f permutes {..<length xs′} and xs′: permute-list f xs′ =
xs.
  then have [simp]: length xs′ = length xs by auto
```

from *permute-list-nth*[*OF perm*, *unfolded xs'*] **have** ∗: $\bigwedge i.\ i < length\ xs \Longrightarrow xs\ !\ i = xs'\ !\ f\ i$ **by** *auto*
note [*simp*] = *list-all2-lengthD*[*OF xsys'*,*symmetric*]
note [*simp*] = *atLeast0LessThan*[*symmetric*]
note *bij* = *permutes-bij*[*OF perm*]
define *ys* **where** *ys* ≡ *map* (*nth ys'* ∘ *f*) [*0*..<*length ys'*]
then have [*simp*]: *length ys* = *length ys'* **by** *auto*
have *mset ys* = *mset* (*map* (*nth ys'*) (*map f* [*0*..<*length ys'*]))
unfolding *ys-def* **by** *auto*
also have ... = *image-mset* (*nth ys'*) (*image-mset f* (*mset* [*0*..<*length ys'*]))
by (*simp add*: *multiset.map-comp*)
also have (*mset* [*0*..<*length ys'*]) = *mset-set* {*0*..<*length ys'*}
by (*metis mset-sorted-list-of-multiset sorted-list-of-mset-set sorted-list-of-set-range*)

also have *image-mset f* (...) = *mset-set* (*f* ‘ {..<*length ys'*})
using *subset-inj-on*[*OF bij-is-inj*[*OF bij*]] **by** (*subst image-mset-mset-set*, *auto*)
also have ... = *mset* [*0*..<*length ys'*] **using** *perm* **by** (*simp add*: *permutes-image*)
also have *image-mset* (*nth ys'*) ... = *mset ys'* **by**(*fold mset-map*, *unfold map-nth*, *auto*)
finally have *mset ys* = *Y* **using** *ys'* **by** *auto*
moreover have *list-all2 rel xs ys*
proof(*rule list-all2-all-nthI*)
fix *i* **assume** *i*: *i* < *length xs*
with ∗ **have** *xs* ! *i* = *xs'* ! *f i* **by** *auto*
also from *i permutes-in-image*[*OF perm*]
have *rel* (*xs'* ! *f i*) (*ys'* ! *f i*) **by** (*intro list-all2-nthD*[*OF xsys'*], *auto*)
finally show *rel* (*xs* ! *i*) (*ys* ! *i*) **unfolding** *ys-def* **using** *i* **by** *simp*
qed *simp*
ultimately show *?thesis* **by** *auto*
qed

lemma *rel-mset-split*:
assumes *rel*: *rel-mset rel* (*X1*+*X2*) *Y*
shows ∃ *Y1 Y2*. *Y* = *Y1* + *Y2* ∧ *rel-mset rel X1 Y1* ∧ *rel-mset rel X2 Y2*
proof−
obtain *xs1* **where** *xs1*: *mset xs1* = *X1* **using** *ex-mset* **by** *auto*
obtain *xs2* **where** *xs2*: *mset xs2* = *X2* **using** *ex-mset* **by** *auto*
from *xs1 xs2* **have** *mset* (*xs1* @ *xs2*) = *X1* + *X2* **by** *auto*
from *rel-mset-free*[*OF rel this*] **obtain** *ys*
where *ys*: *mset ys* = *Y list-all2 rel* (*xs1* @ *xs2*) *ys* **by** *auto*
then obtain *ys1 ys2*
where *ys12*: *ys* = *ys1* @ *ys2*
and *xs1ys1*: *list-all2 rel xs1 ys1*
and *xs2ys2*: *list-all2 rel xs2 ys2*
using *list-all2-append1* **by** *blast*
from *ys12 ys* **have** *Y* = *mset ys1* + *mset ys2* **by** *auto*
moreover from *xs1 xs1ys1* **have** *rel-mset rel X1* (*mset ys1*) **unfolding** *rel-mset-def* **by** *auto*
moreover from *xs2 xs2ys2* **have** *rel-mset rel X2* (*mset ys2*) **unfolding** *rel-mset-def*

by *auto*
ultimately show *?thesis* **by** (*subst exI*[*of - mset ys1*], *subst exI*[*of - mset ys2*],*auto*)
qed

lemma *rel-mset-OO*:
assumes *AB*: *rel-mset R A B* **and** *BC*: *rel-mset S B C*
shows *rel-mset* (*R OO S*) *A C*
proof−
from *AB* **obtain** *as bs* **where** *A-as*: *A* = *mset as* **and** *B-bs*: *B* = *mset bs* **and** *as-bs*: *list-all2 R as bs*
by (*auto simp*: *rel-mset-def*)
from *rel-mset-free*[*OF BC*] *B-bs* **obtain** *cs* **where** *C-cs*: *C* = *mset cs* **and** *bs-cs*: *list-all2 S bs cs*
by *auto*
from *list-all2-trans*[*OF - as-bs bs-cs*, *of R OO S*] *A-as C-cs*
show *?thesis* **by** (*auto simp*: *rel-mset-def*)
qed

lemma *ex-mset-zip-right*:
assumes *length xs* = *length ys mset ys′* = *mset ys*
shows ∃ *xs′*. *length ys′* = *length xs′* ∧ *mset* (*zip xs′ ys′*) = *mset* (*zip xs ys*)
using *assms*
proof (*induct xs ys arbitrary*: *ys′ rule*: *list-induct2*)
case *Nil*
thus *?case*
by *auto*
next
case (*Cons x xs y ys ys′*)
obtain *j* **where** *j-len*: *j* < *length ys′* **and** *nth-j*: *ys′* ! *j* = *y*
by (*metis Cons.prems in-set-conv-nth list.set-intros*(*1*) *mset-eq-setD*)

define *ysa* **where** *ysa* = *take j ys′* @ *drop* (*Suc j*) *ys′*
have *mset ys′* = {#*y*#} + *mset ysa*
unfolding *ysa-def* **using** *j-len nth-j*
by (*metis Cons-nth-drop-Suc union-mset-add-mset-right add-mset-remove-trivial add-diff-cancel-left′*
append-take-drop-id mset.simps(*2*) *mset-append*)
hence *ms-y*: *mset ysa* = *mset ys*
by (*simp add*: *Cons.prems*)
then obtain *xsa* **where**
len-a: *length ysa* = *length xsa* **and** *ms-a*: *mset* (*zip xsa ysa*) = *mset* (*zip xs ys*)
using *Cons.hyps*(*2*) **by** *blast*

define *xs′* **where** *xs′* = *take j xsa* @ *x* # *drop j xsa*
have *ys′*: *ys′* = *take j ysa* @ *y* # *drop j ysa*
using *ms-y j-len nth-j Cons.prems ysa-def*
by (*metis append-eq-append-conv append-take-drop-id diff-Suc-Suc Cons-nth-drop-Suc*

length-Cons
length-drop size-mset)
have *j-len'*: $j \leq$ *length ysa*
using *j-len ys' ysa-def*
by (*metis add-Suc-right append-take-drop-id length-Cons length-append less-eq-Suc-le not-less*)
have *length ys' = length xs'*
unfolding *xs'-def* **using** *Cons.prems len-a ms-y*
by (*metis add-Suc-right append-take-drop-id length-Cons length-append mset-eq-length*)
moreover have *mset* (*zip xs' ys'*) = *mset* (*zip* (*x # xs*) (*y # ys*))
unfolding *ys' xs'-def*
apply (*rule HOL.trans*[*OF mset-zip-take-Cons-drop-twice*])
using *j-len'* **by** (*auto simp*: *len-a ms-a*)
ultimately show *?case*
by *blast*
qed

lemma *list-all2-reorder-right-invariance*:
assumes *rel*: *list-all2 R xs ys* **and** *ms-y*: *mset ys' = mset ys*
shows $\exists$ *xs'*. *list-all2 R xs' ys'* $\wedge$ *mset xs' = mset xs*
proof −
have *len*: *length xs = length ys*
using *rel list-all2-conv-all-nth* **by** *auto*
obtain *xs'* **where**
len': *length xs' = length ys'* **and** *ms-xy*: *mset* (*zip xs' ys'*) = *mset* (*zip xs ys*)
using *len ms-y* **by** (*metis ex-mset-zip-right*)
have *list-all2 R xs' ys'*
using *assms*(*1*) *len' ms-xy* **unfolding** *list-all2-iff* **by** (*blast dest*: *mset-eq-setD*)
moreover have *mset xs' = mset xs*
using *len len' ms-xy map-fst-zip mset-map* **by** *metis*
ultimately show *?thesis*
by *blast*
qed

lemma *rel-mset-via-perm*: *rel-mset rel* (*mset xs*) (*mset ys*) $\longleftrightarrow$ ($\exists$ *zs*. *mset xs = mset zs* $\wedge$ *list-all2 rel zs ys*)
proof (*unfold rel-mset-def*, *intro iffI*, *goal-cases*)
case *1*
then obtain *zs ws* **where** *zs*: *mset zs = mset xs* **and** *ws*: *mset ws = mset ys* **and** *zsws*: *list-all2 rel zs ws* **by** *auto*
note *list-all2-reorder-right-invariance*[*OF zsws ws*[*symmetric*], *unfolded zs*]
then show *?case* **by** (*auto dest*: *sym*)
next
case *2*
from *this* **show** *?case* **by** *force*
qed

end
theory *Unique-Factorization*

imports
Polynomial-Interpolation.Ring-Hom-Poly
Polynomial-Factorization.Polynomial-Irreducibility
HOL−Combinatorics.Permutations
HOL−Computational-Algebra.Euclidean-Algorithm
Containers.Containers-Auxiliary
More-Missing-Multiset
HOL−Algebra.Divisibility
begin

hide-const(**open**)
Divisibility.prime
Divisibility.irreducible

hide-fact(**open**)
Divisibility.irreducible-def
Divisibility.irreducibleI
Divisibility.irreducibleD
Divisibility.irreducibleE

hide-const (**open**) *Rings.coprime*

lemma *irreducible-uminus* [*simp*]:
fixes *a*::*'a*::*idom*
shows *irreducible* (−*a*) ⟷ *irreducible a*
using *irreducible-mult-unit-left*[*of* −*1*::*'a*] **by** *auto*

context *comm-monoid-mult* **begin**

definition *coprime* :: *'a* ⇒ *'a* ⇒ *bool*
where *coprime-def'*: *coprime p q* ≡ ∀ *r*. *r dvd p* ⟶ *r dvd q* ⟶ *r dvd 1*

lemma *coprimeI*:
assumes ⋀*r*. *r dvd p* ⟹ *r dvd q* ⟹ *r dvd 1*
shows *coprime p q* **using** *assms* **by** (*auto simp*: *coprime-def'*)

lemma *coprimeE*:
assumes *coprime p q*
and (⋀*r*. *r dvd p* ⟹ *r dvd q* ⟹ *r dvd 1*) ⟹ *thesis*
shows *thesis* **using** *assms* **by** (*auto simp*: *coprime-def'*)

lemma *coprime-commute* [*ac-simps*]:
coprime p q ⟷ *coprime q p*
by (*auto simp add*: *coprime-def'*)

lemma *not-coprime-iff-common-factor*:
¬ *coprime p q* ⟷ (∃ *r*. *r dvd p* ∧ *r dvd q* ∧ ¬ *r dvd 1*)
by (*auto simp add*: *coprime-def'*)

end

lemma (**in** *algebraic-semidom*) *coprime-iff-coprime* [*simp*, *code*]:
coprime = *Rings.coprime*
by (*simp add*: *fun-eq-iff coprime-def coprime-def'*)

lemma (**in** *comm-semiring-1*) *coprime-0* [*simp*]:
coprime p 0 ⟷ *p dvd 1 coprime 0 p* ⟷ *p dvd 1*
by (*auto intro*: *coprimeI elim*: *coprimeE dest*: *dvd-trans*)

lemma *dvd-rewrites*: *dvd.dvd* ((∗)) = (*dvd*) **by** (*unfold dvd.dvd-def dvd-def*, *rule*)

3.3 Interfacing UFD properties

hide-const (**open**) *Divisibility.irreducible*

context *comm-monoid-mult-isom* **begin**
lemma *coprime-hom*[*simp*]: *coprime* (*hom x*) *y'* ⟷ *coprime x* (*Hilbert-Choice.inv hom y'*)
proof−
show *?thesis* **by** (*unfold coprime-def'*, *fold ball-UNIV*, *subst surj*[*symmetric*], *simp*)
qed
lemma *coprime-inv-hom*[*simp*]: *coprime* (*Hilbert-Choice.inv hom x'*) *y* ⟷ *coprime x'* (*hom y*)
proof−
interpret *inv*: *comm-monoid-mult-isom Hilbert-Choice.inv hom*..
show *?thesis* **by** *simp*
qed
end

3.3.1 Original part

lemma *dvd-dvd-imp-smult*:
fixes *p q* :: *'a* :: *idom poly*
assumes *pq*: *p dvd q* **and** *qp*: *q dvd p* **shows** ∃ *c*. *p* = *smult c q*
proof (*cases p* = *0*)
case *True* **then show** *?thesis* **by** *auto*
next
case *False*
from *qp* **obtain** *r* **where** *r*: *p* = *q* ∗ *r* **by** (*elim dvdE*, *auto*)
with *False qp* **have** *r0*: *r* ≠ *0* **and** *q0*: *q* ≠ *0* **by** *auto*
with *divides-degree*[*OF pq*] *divides-degree*[*OF qp*] *False*
have *degree p* = *degree q* **by** *auto*
with *r degree-mult-eq*[*OF q0 r0*] **have** *degree r* = *0* **by** *auto*
from *degree-0-id*[*OF this*] **obtain** *c* **where** *r* = [:*c*:] **by** *metis*

from *r*[*unfolded this*] **show** *?thesis* **by** *auto*
qed

lemma *dvd-const*:
assumes *pq*: (*p*::'*a*::*semidom poly*) *dvd q* **and** *q0*: *q* ≠ *0* **and** *degq*: *degree q* = *0*
shows *degree p* = *0*
proof−
from *dvdE*[*OF pq*] **obtain** *r* **where** ∗: *q* = *p* ∗ *r***.**
with *q0* **have** *p* ≠ *0 r* ≠ *0* **by** *auto*
from *degree-mult-eq*[*OF this*] *degq* ∗ **show** *degree p* = *0* **by** *auto*
qed

context *Rings.dvd* **begin**
abbreviation *ddvd* (**infix** ‹*ddvd*› *40*) **where** *x ddvd y* ≡ *x dvd y* ∧ *y dvd x*
lemma *ddvd-sym*[*sym*]: *x ddvd y* ⟹ *y ddvd x* **by** *auto*
end

context *comm-monoid-mult* **begin**
lemma *ddvd-trans*[*trans*]: *x ddvd y* ⟹ *y ddvd z* ⟹ *x ddvd z* **using** *dvd-trans*
by *auto*
lemma *ddvd-transp*: *transp* (*ddvd*) **by** (*intro transpI*, *fact ddvd-trans*)
end

context *comm-semiring-1* **begin**

definition *mset-factors* **where** *mset-factors F p* ≡
F ≠ {#} ∧ (∀*f*. *f* ∈# *F* ⟶ *irreducible f*) ∧ *p* = *prod-mset F*

lemma *mset-factorsI*[*intro*!]:
assumes ⋀*f*. *f* ∈# *F* ⟹ *irreducible f* **and** *F* ≠ {#} **and** *prod-mset F* = *p*
shows *mset-factors F p*
unfolding *mset-factors-def* **using** *assms* **by** *auto*

lemma *mset-factorsD*:
assumes *mset-factors F p*
shows *f* ∈# *F* ⟹ *irreducible f* **and** *F* ≠ {#} **and** *prod-mset F* = *p*
using *assms*[*unfolded mset-factors-def*] **by** *auto*

lemma *mset-factorsE*[*elim*]:
assumes *mset-factors F p*
and (⋀*f*. *f* ∈# *F* ⟹ *irreducible f*) ⟹ *F* ≠ {#} ⟹ *prod-mset F* = *p* ⟹
thesis
shows *thesis*
using *assms*[*unfolded mset-factors-def*] **by** *auto*

lemma *mset-factors-imp-not-is-unit*:
assumes *mset-factors F p*
shows ¬ *p dvd 1*
proof(*cases F*)

case *empty* **with** *assms* **show** *?thesis* **by** *auto*
next
case (*add f F*)
with *assms* **have** ¬ *f dvd 1 p* = *f* ∗ *prod-mset F* **by** (*auto intro*!: *irreducible-not-unit*)
then show *?thesis* **by** *auto*
qed

definition *primitive-poly* **where** *primitive-poly f* ≡ ∀ *d*. (∀ *i*. *d dvd coeff f i*) ⟶ *d dvd 1*

end

lemma(**in** *semidom*) *mset-factors-imp-nonzero*:
assumes *mset-factors F p*
shows *p* ≠ *0*
proof
assume *p* = *0*
moreover from *assms* **have** *prod-mset F* = *p* **by** *auto*
ultimately obtain *f* **where** *f* ∈# *F f* = *0* **by** *auto*
with *assms* **show** *False* **by** *auto*
qed

class *ufd* = *idom* +
assumes *mset-factors-exist*: ⋀*x*. *x* ≠ *0* ⟹ ¬ *x dvd 1* ⟹ ∃ *F*. *mset-factors F x*
and *mset-factors-unique*: ⋀*x F G*. *mset-factors F x* ⟹ *mset-factors G x* ⟹ *rel-mset* (*ddvd*) *F G*

3.3.2 Connecting to HOL/Divisibility

context *comm-semiring-1* **begin**

abbreviation *mk-monoid* ≡ (|*carrier* = *UNIV* − {*0*}, *mult* = (∗), *one* = *1*|)

lemma *carrier-0*[*simp*]: *x* ∈ *carrier mk-monoid* ⟷ *x* ≠ *0* **by** *auto*

lemmas *mk-monoid-simps* = *carrier-0 monoid.simps*

abbreviation *irred* **where** *irred* ≡ *Divisibility.irreducible mk-monoid*
abbreviation *factor* **where** *factor* ≡ *Divisibility.factor mk-monoid*
abbreviation *factors* **where** *factors* ≡ *Divisibility.factors mk-monoid*
abbreviation *properfactor* **where** *properfactor* ≡ *Divisibility.properfactor mk-monoid*

lemma *factors*: *factors fs y* ⟷ *prod-list fs* = *y* ∧ *Ball* (*set fs*) *irred*
proof −
have *prod-list fs* = *foldr* (∗) *fs 1* **by** (*induct fs*, *auto*)
thus *?thesis* **unfolding** *factors-def* **by** *auto*
qed

lemma *factor*: *factor x y* ⟷ (∃ *z*. *z* ≠ *0* ∧ *x* ∗ *z* = *y*) **unfolding** *factor-def*

by *auto*

lemma *properfactor-nz*:
shows $(y :: 'a) \neq 0 \Longrightarrow$ *properfactor x y* $\longleftrightarrow$ *x dvd y* $\wedge \neg$ *y dvd x*
by (*auto simp*: *properfactor-def factor-def dvd-def*)

lemma *mem-Units*[*simp*]: *y* $\in$ *Units mk-monoid* $\longleftrightarrow$ *y dvd 1*
unfolding *dvd-def Units-def* **by** (*auto simp*: *ac-simps*)

end

context *idom* **begin**
lemma *irred-0*[*simp*]: *irred* (*0*::'*a*) **by** (*unfold Divisibility.irreducible-def*, *auto simp*: *factor properfactor-def*)
lemma *factor-idom*[*simp*]: *factor* (*x*::'*a*) *y* $\longleftrightarrow$ (*if y = 0 then x = 0 else x dvd y*)
by (*cases y = 0*; *auto intro*: *exI*[*of - 1*] *elim*: *dvdE simp*: *factor*)

lemma *associated-connect*[*simp*]: $(\sim_{mk\text{-}monoid}) = (ddvd)$ **by** (*intro ext*, *unfold associated-def*, *auto*)

lemma *essentially-equal-connect*[*simp*]:
essentially-equal mk-monoid fs gs $\longleftrightarrow$ *rel-mset* (*ddvd*) (*mset fs*) (*mset gs*)
by (*auto simp*: *essentially-equal-def rel-mset-via-perm*)

lemma *irred-idom-nz*:
assumes *x0*: (*x*::'*a*) $\neq$ *0*
shows *irred x* $\longleftrightarrow$ *irreducible x*
using *x0* **by** (*auto simp*: *irreducible-altdef Divisibility.irreducible-def properfactor-nz*)

lemma *dvd-dvd-imp-unit-mult*:
assumes *xy*: *x dvd y* **and** *yx*: *y dvd x*
shows $\exists z.$ *z dvd 1* $\wedge$ *y* = *x* $*$ *z*
proof(*cases x = 0*)
case *True* **with** *xy* **show** *?thesis* **by** (*auto intro*: *exI*[*of - 1*])
next
case *x0*: *False*
from *xy* **obtain** *z* **where** *z*: *y* = *x* $*$ *z* **by** (*elim dvdE*, *auto*)
from *yx* **obtain** *w* **where** *w*: *x* = *y* $*$ *w* **by** (*elim dvdE*, *auto*)
from *z w* **have** *x* $*$ (*z* $*$ *w*) = *x* **by** (*auto simp*: *ac-simps*)
then have *z* $*$ *w* = *1* **using** *x0* **by** *auto*
with *z* **show** *?thesis* **by** (*auto intro*: *exI*[*of - z*])
qed

lemma *irred-inner-nz*:
assumes *x0*: *x* $\neq$ *0*

shows $(\forall b.\ b\ dvd\ x \longrightarrow \neg\ x\ dvd\ b \longrightarrow b\ dvd\ 1) \longleftrightarrow (\forall a\ b.\ x = a * b \longrightarrow a\ dvd\ 1 \vee b\ dvd\ 1)$ (**is** *?l* $\longleftrightarrow$ *?r*)
proof (*intro iffI allI impI*)
assume *l*: *?l*
fix *a b*
assume *xab*: $x = a * b$
then have *ax*: *a dvd x* **and** *bx*: *b dvd x* **by** *auto*
{ **assume** *a1*: $\neg$ *a dvd 1*
with *l ax* **have** *xa*: *x dvd a* **by** *auto*
from *dvd-dvd-imp-unit-mult*[*OF ax xa*] **obtain** *z* **where** *z1*: *z dvd 1* **and** *xaz*: $x = a * z$ **by** *auto*
from *xab x0* **have** $a \neq 0$ **by** *auto*
with *xab xaz* **have** $b = z$ **by** *auto*
with *z1* **have** *b dvd 1* **by** *auto*
}
then show *a dvd 1* $\vee$ *b dvd 1* **by** *auto*
next
assume *r*: *?r*
fix *b* **assume** *bx*: *b dvd x* **and** *xb*: $\neg$ *x dvd b*
then obtain *a* **where** *xab*: $x = a * b$ **by** (*elim dvdE*, *auto simp*: *ac-simps*)
with *r* **consider** *a dvd 1* | *b dvd 1* **by** *auto*
then show *b dvd 1*
proof(*cases*)
case *2* **then show** *?thesis* **by** *auto*
next
case *1*
then obtain *c* **where** *ac1*: $a * c = 1$ **by** (*elim dvdE*, *auto*)
from *xab* **have** $x * c = b * (a * c)$ **by** (*auto simp*: *ac-simps*)
with *ac1* **have** $x * c = b$ **by** *auto*
then have *x dvd b* **by** *auto*
with *xb* **show** *?thesis* **by** *auto*
qed
qed

lemma *irred-idom*[*simp*]: *irred x* $\longleftrightarrow$ $x = 0 \vee$ *irreducible x*
by (*cases* $x = 0$; *simp add*: *irred-idom-nz irred-inner-nz irreducible-def*)

lemma assumes $x \neq 0$ **and** *factors fs x* **and** $f \in$ *set fs* **shows** $f \neq 0$
using *assms* **by** (*auto simp*: *factors*)

lemma *factors-as-mset-factors*:
assumes *x0*: $x \neq 0$ **and** *x1*: $x \neq 1$
shows *factors fs x* $\longleftrightarrow$ *mset-factors* (*mset fs*) *x* **using** *assms*
by (*auto simp*: *factors prod-mset-prod-list*)

end

context *ufd* **begin**

```
 interpretation comm-monoid-cancel: comm-monoid-cancel mk-monoid::'a monoid
   apply (unfold-locales)
   apply simp-all
   using mult-left-cancel
   apply (auto simp: ac-simps)
   done
 lemma factors-exist:
   assumes a ≠ 0
   and ¬ a dvd 1
   shows ∃fs. set fs ⊆ UNIV − {0} ∧ factors fs a
 proof−
   from mset-factors-exist[OF assms]
   obtain F where mset-factors F a by auto
   also from ex-mset obtain fs where F = mset fs by metis
   finally have fs: mset-factors (mset fs) a.
   then have factors fs a using assms by (subst factors-as-mset-factors, auto)
  moreover have set fs ⊆ UNIV − {0} using fs by (auto elim!: mset-factorsE)
   ultimately show ?thesis by auto
 qed

 lemma factors-unique:
   assumes fs: factors fs a
     and gs: factors gs a
     and a0: a ≠ 0
     and a1: ¬ a dvd 1
   shows rel-mset (ddvd) (mset fs) (mset gs)
 proof−
   from a1 have a ≠ 1 by auto
  with a0 fs gs have mset-factors (mset fs) a mset-factors (mset gs) a by (unfold
factors-as-mset-factors)
   from mset-factors-unique[OF this] show ?thesis.
 qed

 lemma factorial-monoid: factorial-monoid (mk-monoid :: 'a monoid)
   by (unfold-locales; auto simp add: factors-exist factors-unique)

end

lemma (in idom) factorial-monoid-imp-ufd:
 assumes factorial-monoid (mk-monoid :: 'a monoid)
 shows class.ufd ((∗) :: 'a ⇒ -) 1 (+) 0 (−) uminus
proof (unfold-locales)
 interpret factorial-monoid mk-monoid :: 'a monoid by (fact assms)
 {
   fix x assume x: x ≠ 0 ¬ x dvd 1
   note ∗ = factors-exist[simplified, OF this]
    with x show ∃F. mset-factors F x by (subst(asm) factors-as-mset-factors,
auto)
 }
```

fix *x F G* **assume** *FG*: *mset-factors F x mset-factors G x*
with *mset-factors-imp-not-is-unit* **have** *x1*: ¬ *x dvd 1* **by** *auto*
from *FG*(*1*) **have** *x0*: *x* ≠ *0* **by** (*rule mset-factors-imp-nonzero*)
obtain *fs gs* **where** *fsgs*: *F* = *mset fs G* = *mset gs* **using** *ex-mset* **by** *metis*
note *FG* = *FG*[*unfolded this*]
then have *0*: *0* ∉ *set fs 0* ∉ *set gs* **by** (*auto elim*!: *mset-factorsE*)
from *x1* **have** *x* ≠ *1* **by** *auto*
note *FG*[*folded factors-as-mset-factors*[*OF x0 this*]]
from *factors-unique*[*OF this*, *simplified*, *OF x0 x1*, *folded fsgs*] *0*
show *rel-mset* (*ddvd*) *F G* **by** *auto*
qed

3.4 Preservation of Irreducibility

locale *comm-semiring-1-hom* = *comm-monoid-mult-hom hom* + *zero-hom hom*
for *hom* :: *'a* :: *comm-semiring-1* ⇒ *'b* :: *comm-semiring-1*

locale *irreducibility-hom* = *comm-semiring-1-hom* +
assumes *irreducible-imp-irreducible-hom*: *irreducible a* ⟹ *irreducible* (*hom a*)
begin
lemma *hom-mset-factors*:
assumes *F*: *mset-factors F p*
shows *mset-factors* (*image-mset hom F*) (*hom p*)
proof (*unfold mset-factors-def*, *intro conjI allI impI*)
from *F* **show** *hom p* = *prod-mset* (*image-mset hom F*) *image-mset hom F* ≠ {#} **by** (*auto simp*: *hom-distribs*)
fix *f'* **assume** *f'* ∈# *image-mset hom F*
then obtain *f* **where** *f*: *f* ∈# *F* **and** *f'f*: *f'* = *hom f* **by** *auto*
with *F irreducible-imp-irreducible-hom* **show** *irreducible f'* **unfolding** *f'f* **by** *auto*
qed
end

locale *unit-preserving-hom* = *comm-semiring-1-hom* +
assumes *is-unit-hom-if*: ⋀*x*. *hom x dvd 1* ⟹ *x dvd 1*
begin
lemma *is-unit-hom-iff*[*simp*]: *hom x dvd 1* ⟷ *x dvd 1* **using** *is-unit-hom-if hom-dvd* **by** *force*

lemma *irreducible-hom-imp-irreducible*:
assumes *irr*: *irreducible* (*hom a*) **shows** *irreducible a*
proof (*intro irreducibleI*)
from *irr* **show** *a* ≠ *0* **by** *auto*
from *irr* **show** ¬ *a dvd 1* **by** (*auto dest*: *irreducible-not-unit*)
fix *b c* **assume** *a* = *b* ∗ *c*
then have *hom a* = *hom b* ∗ *hom c* **by** (*simp add*: *hom-distribs*)
with *irr* **have** *hom b dvd 1* ∨ *hom c dvd 1* **by** (*auto dest*: *irreducibleD*)
then show *b dvd 1* ∨ *c dvd 1* **by** *simp*
qed

end

locale *factor-preserving-hom* = *unit-preserving-hom* + *irreducibility-hom*
begin
lemma *irreducible-hom*[*simp*]: *irreducible* (*hom a*) ⟷ *irreducible a*
using *irreducible-hom-imp-irreducible irreducible-imp-irreducible-hom* **by** *metis*
end

lemma *factor-preserving-hom-comp*:
assumes *f*: *factor-preserving-hom f* **and** *g*: *factor-preserving-hom g*
shows *factor-preserving-hom* (*f o g*)
proof−
interpret *f*: *factor-preserving-hom f* **by** (*rule f*)
interpret *g*: *factor-preserving-hom g* **by** (*rule g*)
show *?thesis* **by** (*unfold-locales*, *auto simp*: *hom-distribs*)
qed

context *comm-semiring-isom* **begin**
sublocale *unit-preserving-hom* **by** (*unfold-locales*, *auto*)
sublocale *factor-preserving-hom*
proof (*standard*)
fix *a* :: *'a*
assume *irreducible a*
note *a* = *this*[*unfolded irreducible-def*]
show *irreducible* (*hom a*)
proof (*rule ccontr*)
assume ¬ *irreducible* (*hom a*)
from *this*[*unfolded Factorial-Ring.irreducible-def*,*simplified*] *a*
obtain *hb hc* **where** *eq*: *hom a* = *hb* ∗ *hc* **and** *nu*: ¬ *hb dvd 1* ¬ *hc dvd 1*
by *auto*
from *bij* **obtain** *b* **where** *hb*: *hb* = *hom b* **by** (*elim bij-pointE*)
from *bij* **obtain** *c* **where** *hc*: *hc* = *hom c* **by** (*elim bij-pointE*)
from *eq*[*unfolded hb hc*, *folded hom-mult*] **have** *a* = *b* ∗ *c* **by** *auto*
with *nu hb hc* **have** *a* = *b* ∗ *c* ¬ *b dvd 1* ¬ *c dvd 1* **by** *auto*
with *a* **show** *False* **by** *auto*
qed
qed
end

3.4.1 Back to divisibility

lemma(**in** *comm-semiring-1*) *mset-factors-mult*:
assumes *F*: *mset-factors F a*
and *G*: *mset-factors G b*
shows *mset-factors* (*F*+*G*) (*a*∗*b*)
proof(*intro mset-factorsI*)
fix *f* **assume** *f* ∈# *F* + *G*
then consider *f* ∈# *F* | *f* ∈# *G* **by** *auto*
then show *irreducible f* **by**(*cases*, *insert F G*, *auto*)

qed (*insert F G*, *auto*)

lemma(**in** *ufd*) *dvd-imp-subset-factors*:
assumes *ab*: *a dvd b*
and *F*: *mset-factors F a*
and *G*: *mset-factors G b*
shows $\exists\, G'.\ G' \subseteq\# G \wedge$ *rel-mset* (*ddvd*) *F G'*
proof−
from *F G* **have** *a0*: $a \neq 0$ **and** *b0*: $b \neq 0$ **by** (*simp-all add*: *mset-factors-imp-nonzero*)
from *ab* **obtain** *c* **where** *c*: $b = a * c$ **by** (*elim dvdE*, *auto*)
with *b0* **have** *c0*: $c \neq 0$ **by** *auto*
show *?thesis*
proof(*cases c dvd 1*)
case *True*
show *?thesis*
proof(*cases F*)
case *empty* **with** *F* **show** *?thesis* **by** *auto*
next
case (*add f F'*)
with *F*
have *a*: $f *$ *prod-mset F'* $= a$
and *F'*: $\bigwedge f.\ f \in\# F' \Longrightarrow$ *irreducible f*
and *irrf*: *irreducible f* **by** *auto*
from *irrf* **have** *f0*: $f \neq 0$ **and** *f1*: $\neg f$ *dvd 1* **by** (*auto dest*: *irreducible-not-unit*)
from *a c* **have** $(f * c) *$ *prod-mset F'* $= b$ **by** (*auto simp*: *ac-simps*)
moreover **{**
have *irreducible* $(f * c)$ **using** *True irrf* **by** (*subst irreducible-mult-unit-right*)
with *F' irrf* **have** $\bigwedge f'.\ f' \in\# F' + \{\#f * c\#\} \Longrightarrow$ *irreducible f'* **by** *auto*
}
ultimately have *mset-factors* $(F' + \{\#f * c\#\})$ *b* **by** (*intro mset-factorsI*, *auto*)
from *mset-factors-unique*[*OF this G*]
have *F'G*: *rel-mset* (*ddvd*) $(F' + \{\#f * c\#\})$ *G*.
from *True add* **have** *FF'*: *rel-mset* (*ddvd*) *F* $(F' + \{\#f * c\#\})$
by (*auto simp add*: *multiset.rel-refl intro*!: *rel-mset-Plus*)
have *rel-mset* (*ddvd*) *F G*
apply(*rule transpD*[*OF multiset.rel-transp*[*OF transpI*] *FF' F'G*])
using *ddvd-trans*.
then show *?thesis* **by** *auto*
qed
next
case *False*
from *mset-factors-exist*[*OF c0 this*] **obtain** *H* **where** *H*: *mset-factors H c* **by** *auto*
from *c mset-factors-mult*[*OF F H*] **have** *mset-factors* $(F + H)$ *b* **by** *auto*
note *mset-factors-unique*[*OF this G*]
from *rel-mset-split*[*OF this*] **obtain** *G1 G2*

```
      where G = G1 + G2 rel-mset (ddvd) F G1 rel-mset (ddvd) H G2 by auto
    then show ?thesis by (intro exI[of - G1], auto)
  qed
qed

lemma(in idom) irreducible-factor-singleton:
  assumes a: irreducible a
  shows mset-factors F a ⟷ F = {#a#}
proof(cases F)
  case empty with mset-factorsD show ?thesis by auto
next
  case (add f F')
  show ?thesis
  proof
    assume F: mset-factors F a
    from add mset-factorsD[OF F] have *: a = f * prod-mset F' by auto
    then have fa: f dvd a by auto
    from * a have f0: f ≠ 0 by auto
    from add have f ∈# F by auto
    with F have f: irreducible f by auto
    from add have F' ⊆# F by auto
    then have unitemp: prod-mset F' dvd 1 ⟹ F' = {#}
    proof(induct F')
      case empty then show ?case by auto
    next
      case (add f F')
        from add have f ∈# F by (simp add: mset-subset-eq-insertD)
        with F irreducible-not-unit have ¬ f dvd 1 by auto
        then have ¬ (prod-mset F' * f) dvd 1 by simp
        with add show ?case by auto
    qed
    show F = {#a#}
    proof(cases a dvd f)
      case True
        then obtain r where f = a * r by (elim dvdE, auto)
        with * have f = (r * prod-mset F') * f by (auto simp: ac-simps)
        with f0 have r * prod-mset F' = 1 by auto
        then have prod-mset F' dvd 1 by (metis dvd-triv-right)
        with unitemp * add show ?thesis by auto
    next
      case False with fa a f show ?thesis by (auto simp: irreducible-altdef)
    qed
  qed (insert a, auto)
qed


lemma(in ufd) irreducible-dvd-imp-factor:
  assumes ab: a dvd b
      and a: irreducible a
```

```
    and G: mset-factors G b
  shows ∃ g ∈# G. a ddvd g
proof−
  from a have mset-factors {#a#} a by auto
  from dvd-imp-subset-factors[OF ab this G]
  obtain G' where G'G: G' ⊆# G and rel: rel-mset (ddvd) {#a#} G' by auto
  with rel-mset-size size-1-singleton-mset size-single
  obtain g where gG': G' = {#g#} by fastforce
  from rel[unfolded this rel-mset-def]
  have a ddvd g by auto
  with gG' G'G show ?thesis by auto
qed

lemma(in idom) prod-mset-remove-units:
  prod-mset F ddvd prod-mset {# f ∈# F. ¬f dvd 1 #}
proof(induct F)
  case (add f F) then show ?case by (cases f = 0, auto)
qed auto

lemma(in comm-semiring-1) mset-factors-imp-dvd:
  assumes mset-factors F x and f ∈# F shows f dvd x
  using assms by (simp add: dvd-prod-mset mset-factors-def)

lemma(in ufd) prime-elem-iff-irreducible[iff]:
  prime-elem x ⟷ irreducible x
proof (intro iffI, fact prime-elem-imp-irreducible, rule prime-elemI)
  assume r: irreducible x
  then show x0: x ≠ 0 and x1: ¬ x dvd 1 by (auto dest: irreducible-not-unit)
  from irreducible-factor-singleton[OF r]
  have ∗: mset-factors {#x#} x by auto
  fix a b
  assume x dvd a ∗ b
  then obtain c where abxc: a ∗ b = x ∗ c by (elim dvdE, auto)
  show x dvd a ∨ x dvd b
  proof(cases c = 0 ∨ a = 0 ∨ b = 0)
    case True with abxc show ?thesis by auto
  next
    case False
    then have a0: a ≠ 0 and b0: b ≠ 0 and c0: c ≠ 0 by auto
    from x0 c0 have xc0: x ∗ c ≠ 0 by auto
    from x1 have xc1: ¬ x ∗ c dvd 1 by auto
    show ?thesis
    proof (cases a dvd 1 ∨ b dvd 1)
      case False
      then have a1: ¬ a dvd 1 and b1: ¬ b dvd 1 by auto
      from mset-factors-exist[OF a0 a1]
      obtain F where Fa: mset-factors F a by auto
      then have F0: F ≠ {#} by auto
      from mset-factors-exist[OF b0 b1]
```

```
    obtain G where Gb: mset-factors G b by auto
    then have G0: G ≠ {#} by auto
    from mset-factors-mult[OF Fa Gb]
    have FGxc: mset-factors (F + G) (x * c) by (simp add: abxc)
    show ?thesis
    proof (cases c dvd 1)
      case True
      from r irreducible-mult-unit-right[OF this] have irreducible (x*c) by simp
      note irreducible-factor-singleton[OF this] FGxc
      with F0 G0 have False by (cases F; cases G; auto)
      then show ?thesis by auto
    next
      case False
      from mset-factors-exist[OF c0 this] obtain H where mset-factors H c by
auto
      with * have xHxc: mset-factors (add-mset x H) (x * c) by force
      note rel = mset-factors-unique[OF this FGxc]
      obtain hs where mset hs = H using ex-mset by auto
      then have mset (x#hs) = add-mset x H by auto
      from rel-mset-free[OF rel this]
      obtain jjs where jjsGH: mset jjs = F + G and rel: list-all2 (ddvd) (x #
hs) jjs by auto
      then obtain j js where jjs: jjs = j # js by (cases jjs, auto)
      with rel have xj: x ddvd j by auto
    from jjs jjsGH have j: j ∈ set-mset (F + G) by (intro union-single-eq-member,
auto)
      from j consider j ∈# F | j ∈# G by auto
      then show ?thesis
      proof(cases)
        case 1
        with Fa have j dvd a by (auto intro: mset-factors-imp-dvd)
        with xj dvd-trans have x dvd a by auto
        then show ?thesis by auto
      next
        case 2
        with Gb have j dvd b by (auto intro: mset-factors-imp-dvd)
        with xj dvd-trans have x dvd b by auto
        then show ?thesis by auto
      qed
    qed
  next
    case True
    then consider a dvd 1 | b dvd 1 by auto
    then show ?thesis
    proof(cases)
      case 1
      then obtain d where ad: a * d = 1 by (elim dvdE, auto)
      from abxc have x * (c * d) = a * b * d by (auto simp: ac-simps)
      also have ... = a * d * b by (auto simp: ac-simps)
```

finally have *x dvd b* **by** (*intro dvdI*, *auto simp*: *ad*)
then show *?thesis* **by** *auto*
next
case *2*
then obtain *d* **where** *bd*: *b * d = 1* **by** (*elim dvdE*, *auto*)
from *abxc* **have** *x * (c * d) = a * b * d* **by** (*auto simp*: *ac-simps*)
also have *... = (b * d) * a* **by** (*auto simp*: *ac-simps*)
finally have *x dvd a* **by** (*intro dvdI*, *auto simp*:*bd*)
then show *?thesis* **by** *auto*
qed
qed
qed
qed

3.5 Results for GCDs etc.

lemma *prod-list-remove1*: (*x* :: *'b* :: *comm-monoid-mult*) ∈ *set xs* ⟹ *prod-list* (*remove1 x xs*) * *x* = *prod-list xs*
by (*induct xs*, *auto simp*: *ac-simps*)

class *comm-monoid-gcd* = *gcd* + *comm-semiring-1* +
assumes *gcd-dvd1*[*iff*]: *gcd a b dvd a*
and *gcd-dvd2*[*iff*]: *gcd a b dvd b*
and *gcd-greatest*: *c dvd a* ⟹ *c dvd b* ⟹ *c dvd gcd a b*
begin

lemma *gcd-0-0*[*simp*]: *gcd 0 0 = 0*
using *gcd-greatest*[*OF dvd-0-right dvd-0-right*, *of 0*] **by** *auto*

lemma *gcd-zero-iff*[*simp*]: *gcd a b = 0* ⟷ *a = 0* ∧ *b = 0*
proof
assume *gcd a b = 0*
from *gcd-dvd1*[*of a b*, *unfolded this*] *gcd-dvd2*[*of a b*, *unfolded this*]
show *a = 0* ∧ *b = 0* **by** *auto*
qed *auto*

lemma *gcd-zero-iff'*[*simp*]: *0 = gcd a b* ⟷ *a = 0* ∧ *b = 0*
using *gcd-zero-iff* **by** *metis*

lemma *dvd-gcd-0-iff*[*simp*]:
shows *x dvd gcd 0 a* ⟷ *x dvd a* (**is** *?g1*)
and *x dvd gcd a 0* ⟷ *x dvd a* (**is** *?g2*)
proof−
have *a dvd gcd a 0 a dvd gcd 0 a* **by** (*auto intro*: *gcd-greatest*)
with *dvd-refl* **show** *?g1 ?g2* **by** (*auto dest*: *dvd-trans*)
qed

lemma *gcd-dvd-1*[*simp*]: *gcd a b dvd 1* ⟷ *coprime a b*

using *dvd-trans*[*OF gcd-greatest*[*of - a b*], *of - 1*]
by (*cases a = 0 ∧ b = 0*) (*auto intro*!: *coprimeI elim*: *coprimeE*)

lemma *dvd-imp-gcd-dvd-gcd*: *b dvd c ⟹ gcd a b dvd gcd a c*
by (*meson gcd-dvd1 gcd-dvd2 gcd-greatest dvd-trans*)

definition *listgcd* :: *'a list ⇒ 'a* **where**
listgcd xs = foldr gcd xs 0

lemma *listgcd-simps*[*simp*]: *listgcd [] = 0 listgcd* (*x # xs*) = *gcd x* (*listgcd xs*)
by (*auto simp*: *listgcd-def*)

lemma *listgcd*: *x ∈ set xs ⟹ listgcd xs dvd x*
proof (*induct xs*)
case (*Cons y ys*)
show *?case*
proof (*cases x = y*)
case *False*
with *Cons* **have** *dvd*: *listgcd ys dvd x* **by** *auto*
thus *?thesis* **unfolding** *listgcd-simps* **using** *dvd-trans* **by** *blast*
next
case *True*
thus *?thesis* **unfolding** *listgcd-simps* **using** *dvd-trans* **by** *blast*
qed
qed *simp*

lemma *listgcd-greatest*: (⋀ *x*. *x ∈ set xs ⟹ y dvd x*) ⟹ *y dvd listgcd xs*
by (*induct xs arbitrary*:*y*, *auto intro*: *gcd-greatest*)

end

context *Rings.dvd* **begin**

definition *is-gcd x a b* ≡ *x dvd a ∧ x dvd b ∧* (∀ *y*. *y dvd a ⟶ y dvd b ⟶ y dvd x*)

definition *some-gcd a b* ≡ *SOME x*. *is-gcd x a b*

lemma *is-gcdI*[*intro*!]:
assumes *x dvd a x dvd b* ⋀*y*. *y dvd a ⟹ y dvd b ⟹ y dvd x*
shows *is-gcd x a b* **by** (*insert assms*, *auto simp*: *is-gcd-def*)

lemma *is-gcdE*[*elim*!]:
assumes *is-gcd x a b*
and *x dvd a ⟹ x dvd b ⟹* (⋀*y*. *y dvd a ⟹ y dvd b ⟹ y dvd x*) ⟹ *thesis*
shows *thesis* **by** (*insert assms*, *auto simp*: *is-gcd-def*)

lemma *is-gcd-some-gcdI*:
assumes $\exists x.$ *is-gcd x a b* **shows** *is-gcd* (*some-gcd a b*) *a b*
by (*unfold some-gcd-def*, *rule someI-ex*[*OF assms*])

end

context *comm-semiring-1* **begin**

lemma *some-gcd-0*[*intro*!]: *is-gcd* (*some-gcd a 0*) *a 0 is-gcd* (*some-gcd 0 b*) *0 b*
by (*auto intro*!: *is-gcd-some-gcdI intro*: *exI*[*of - a*] *exI*[*of - b*])

lemma *some-gcd-0-dvd*[*intro*!]:
some-gcd a 0 dvd a some-gcd 0 b dvd b **using** *some-gcd-0* **by** *auto*

lemma *dvd-some-gcd-0*[*intro*!]:
a dvd some-gcd a 0 b dvd some-gcd 0 b **using** *some-gcd-0*[*of a*] *some-gcd-0*[*of b*] **by** *auto*

end

context *idom* **begin**

lemma *is-gcd-connect*:
assumes $a \neq 0$ $b \neq 0$ **shows** *isgcd mk-monoid x a b* $\longleftrightarrow$ *is-gcd x a b*
using *assms* **by** (*force simp*: *isgcd-def*)

lemma *some-gcd-connect*:
assumes $a \neq 0$ **and** $b \neq 0$ **shows** *somegcd mk-monoid a b = some-gcd a b*
using *assms* **by** (*auto intro*!: *arg-cong*[*of - - Eps*] *simp*: *is-gcd-connect some-gcd-def somegcd-def*)
end

context *comm-monoid-gcd*
begin
lemma *is-gcd-gcd*: *is-gcd* (*gcd a b*) *a b* **using** *gcd-greatest* **by** *auto*
lemma *is-gcd-some-gcd*: *is-gcd* (*some-gcd a b*) *a b* **by** (*insert is-gcd-gcd*, *auto intro*!: *is-gcd-some-gcdI*)
lemma *gcd-dvd-some-gcd*: *gcd a b dvd some-gcd a b* **using** *is-gcd-some-gcd* **by** *auto*
lemma *some-gcd-dvd-gcd*: *some-gcd a b dvd gcd a b* **using** *is-gcd-some-gcd* **by** (*auto intro*: *gcd-greatest*)
lemma *some-gcd-ddvd-gcd*: *some-gcd a b ddvd gcd a b* **by** (*auto intro*: *gcd-dvd-some-gcd some-gcd-dvd-gcd*)
lemma *some-gcd-dvd*: *some-gcd a b dvd d* $\longleftrightarrow$ *gcd a b dvd d d dvd some-gcd a b* $\longleftrightarrow$ *d dvd gcd a b*
using *some-gcd-ddvd-gcd*[*of a b*] **by** (*auto dest*:*dvd-trans*)

end

class *idom-gcd* = *comm-monoid-gcd* + *idom*
begin

interpretation *raw*: *comm-monoid-cancel mk-monoid* :: *'a monoid*
by (*unfold-locales*, *auto intro*: *mult-commute mult-assoc*)

interpretation *raw*: *gcd-condition-monoid mk-monoid* :: *'a monoid*
by (*unfold-locales*, *auto simp*: *is-gcd-connect intro*!: *exI*[*of - gcd - -*] *dest*: *gcd-greatest*)

lemma *gcd-mult-ddvd*:
*d * gcd a b ddvd gcd (d * a) (d * b)*
proof (*cases d = 0*)
case *True* **then show** *?thesis* **by** *auto*
next
case *d0*: *False*
show *?thesis*
proof (*cases a = 0 ∨ b = 0*)
case *False*
note *some-gcd-ddvd-gcd*[*of a b*]
with *d0* **have** *d * gcd a b ddvd d * some-gcd a b* **by** *auto*
also have *d * some-gcd a b ddvd some-gcd (d * a) (d * b)*
using *False d0 raw.gcd-mult* **by** (*simp add*: *some-gcd-connect*)
also note *some-gcd-ddvd-gcd*
finally show *?thesis*.
next
case *True*
with *d0* **show** *?thesis*
apply (*elim disjE*)
apply (*rule ddvd-trans*[*of - d * b*]; *force*)
apply (*rule ddvd-trans*[*of - d * a*]; *force*)
done
qed
qed

lemma *gcd-greatest-mult*: **assumes** *cad*: *c dvd a * d* **and** *cbd*: *c dvd b * d*
shows *c dvd gcd a b * d*
proof−
from *gcd-greatest*[*OF assms*] **have** *c*: *c dvd gcd (d * a) (d * b)* **by** (*auto simp*: *ac-simps*)
note *gcd-mult-ddvd*[*of d a b*]
then have *gcd (d * a) (d * b) dvd gcd a b * d* **by** (*auto simp*: *ac-simps*)
from *dvd-trans*[*OF c this*] **show** *?thesis* .
qed

lemma *listgcd-greatest-mult*: (⋀ *x* :: *'a*. *x ∈ set xs ⟹ y dvd x * z*) ⟹ *y dvd listgcd xs * z*
proof (*induct xs*)
case (*Cons x xs*)

```
  from Cons have y dvd x * z y dvd listgcd xs * z by auto
  thus ?case unfolding listgcd-simps by (rule gcd-greatest-mult)
qed (simp)

lemma dvd-factor-mult-gcd:
  assumes dvd: k dvd p * q k dvd p * r
    and q0: q ≠ 0 and r0: r ≠ 0
  shows k dvd p * gcd q r
proof −
  from dvd gcd-greatest[of k p * q p * r]
  have k dvd gcd (p * q) (p * r) by simp
  also from gcd-mult-ddvd[of p q r]
  have ... dvd (p * gcd q r) by auto
  finally show ?thesis .
qed

lemma coprime-mult-cross-dvd:
  assumes coprime: coprime p q and eq: p' * p = q' * q
  shows p dvd q' (is ?g1) and q dvd p' (is ?g2)
proof (atomize(full), cases p = 0 ∨ q = 0)
  case True
  then show ?g1 ∧ ?g2
  proof
    assume p0: p = 0 with coprime have q dvd 1 by auto
    with eq p0 show ?thesis by auto
  next
    assume q0: q = 0 with coprime have p dvd 1 by auto
    with eq q0 show ?thesis by auto
  qed
next
  case False
  {
    fix p q r p' q' :: 'a
   assume cop: coprime p q and eq: p' * p = q' * q and p: p ≠ 0 and q: q ≠ 0
      and r: r dvd p r dvd q
    let ?gcd = gcd q p
    from eq have p' * p dvd q' * q by auto
    hence d1: p dvd q' * q by (rule dvd-mult-right)
    have d2: p dvd q' * p by auto
    from dvd-factor-mult-gcd[OF d1 d2 q p] have 1: p dvd q' * ?gcd .
    from q p have 2: ?gcd dvd q by auto
    from q p have 3: ?gcd dvd p by auto
    from cop[unfolded coprime-def', rule-format, OF 3 2] have ?gcd dvd 1 .
    from 1 dvd-mult-unit-iff[OF this] have p dvd q' by auto
  } note main = this
   from main[OF coprime eq,of 1] False coprime coprime-commute main[OF -
eq[symmetric], of 1]
  show ?g1 ∧ ?g2 by auto
qed
```

end

subclass (**in** *ring-gcd*) *idom-gcd* **by** (*unfold-locales*, *auto*)

lemma *coprime-rewrites*: *comm-monoid-mult.coprime* ((∗)) *1* = *coprime*
apply (*intro ext*)
apply (*subst comm-monoid-mult.coprime-def'*)
apply (*unfold-locales*)
apply (*unfold dvd-rewrites*)
apply (*fold coprime-def'*) **..**

locale *gcd-condition* =
fixes *ty* :: *'a* :: *idom itself*
assumes *gcd-exists*: ⋀*a b* :: *'a*. ∃ *x*. *is-gcd x a b*
begin
sublocale *idom-gcd* (∗) *1* :: *'a* (+) *0* (−) *uminus some-gcd*
rewrites *dvd.dvd* ((∗)) = (*dvd*)
and *comm-monoid-mult.coprime* ((∗)) *1* = *Unique-Factorization.coprime*
proof−
have *is-gcd* (*some-gcd a b*) *a b* **for** *a b* :: *'a* **by** (*intro is-gcd-some-gcdI gcd-exists*)
from *this*[*unfolded is-gcd-def*]
show *class.idom-gcd* (∗) (*1* :: *'a*) (+) *0* (−) *uminus some-gcd* **by** (*unfold-locales*, *auto simp*: *dvd-rewrites*)
qed (*simp-all add*: *dvd-rewrites coprime-rewrites*)
end

instance *semiring-gcd* ⊆ *comm-monoid-gcd* **by** (*intro-classes*, *auto*)

lemma *listgcd-connect*: *listgcd* = *gcd-list*
proof (*intro ext*)
fix *xs* :: *'a list*
show *listgcd xs* = *gcd-list xs* **by**(*induct xs*, *auto*)
qed

interpretation *some-gcd*: *gcd-condition TYPE*(*'a*::*ufd*)
proof(*unfold-locales*, *intro exI*)
interpret *factorial-monoid mk-monoid* :: *'a monoid* **by** (*fact factorial-monoid*)
note *d* = *dvd.dvd-def some-gcd-def carrier-0*
fix *a b* :: *'a*
show *is-gcd* (*some-gcd a b*) *a b*
proof (*cases a* = *0* ∨ *b* = *0*)
case *True*
thus *?thesis* **using** *some-gcd-0* **by** *auto*
next
case *False*
with *gcdof-exists*[*of a b*]
show *?thesis* **by** (*auto intro*!: *is-gcd-some-gcdI simp add*: *is-gcd-connect some-gcd-connect*)

qed
qed

lemma *some-gcd-listgcd-dvd-listgcd*: *some-gcd.listgcd xs dvd listgcd xs*
by (*induct xs*, *auto simp*:*some-gcd-dvd intro*:*dvd-imp-gcd-dvd-gcd*)

lemma *listgcd-dvd-some-gcd-listgcd*: *listgcd xs dvd some-gcd.listgcd xs*
by (*induct xs*, *auto simp*:*some-gcd-dvd intro*:*dvd-imp-gcd-dvd-gcd*)

context *factorial-ring-gcd* **begin**

Do not declare the following as subclass, to avoid conflict in *field* $\subseteq$ *gcd-condition* vs. *factorial-ring-gcd* $\subseteq$ *gcd-condition*.

sublocale *as-ufd*: *ufd*
proof(*unfold-locales*, *goal-cases*)
case (*1 x*)
from *prime-factorization-exists*[*OF* ‹$x \neq 0$›]
obtain *F* **where** *f*: $\bigwedge f.\ f \in\# F \Longrightarrow$ *prime-elem f*
and *Fx*: *normalize* (*prod-mset F*) = *normalize x* **by** *auto*
from *associatedE2*[*OF Fx*] **obtain** *u* **where** *u*: *is-unit u x* = *u* $*$ *prod-mset F*
by *blast*
from ‹$\neg$ *is-unit x*› *Fx* **have** *F* $\neq$ {#} **by** *auto*
then obtain *g G* **where** *F*: *F* = *add-mset g G* **by** (*cases F*, *auto*)
then have *g* $\in\#$ *F* **by** *auto*
with *f*[*OF this*]*prime-elem-iff-irreducible*
irreducible-mult-unit-left[*OF unit-factor-is-unit*[*OF* ‹$x \neq 0$›]]
have *g*: *irreducible* (*u* $*$ *g*) **using** *u*(*1*)
by (*subst irreducible-mult-unit-left*) *simp-all*
show *?case*
proof (*intro exI conjI mset-factorsI*)
show *prod-mset* (*add-mset* (*u* $*$ *g*) *G*) = *x*
using ‹$x \neq 0$› **by** (*simp add*: *F ac-simps u*)
fix *f* **assume** *f* $\in\#$ *add-mset* (*u* $*$ *g*) *G*
with *f*[*unfolded F*] *g prime-elem-iff-irreducible*
show *irreducible f* **by** *auto*
qed *auto*
next
case (*2 x F G*)
note *transpD*[*OF multiset.rel-transp*[*OF ddvd-transp*],*trans*]
obtain *fs* **where** *F*: *F* = *mset fs* **by** (*metis ex-mset*)
have *list-all2* (*ddvd*) *fs* (*map normalize fs*) **by** (*intro list-all2-all-nthI*, *auto*)
then have *FH*: *rel-mset* (*ddvd*) *F* (*image-mset normalize F*) **by** (*unfold rel-mset-def F*, *force*)
also
have *FG*: *image-mset normalize F* = *image-mset normalize G*
proof (*intro prime-factorization-unique''*)
from *2* **have** *xF*: *x* = *prod-mset F* **and** *xG*: *x* = *prod-mset G* **by** *auto*
from *xF* **have** *normalize x* = *normalize* (*prod-mset* (*image-mset normalize F*))
by (*simp add*: *normalize-prod-mset-normalize*)

with *xG* **have** *nFG*: ... = *normalize* (*prod-mset* (*image-mset normalize G*))
by (*simp-all add*: *normalize-prod-mset-normalize*)
then show *normalize* ($\prod$ *i*∈#*image-mset normalize F*. *i*) =
normalize ($\prod$ *i*∈#*image-mset normalize G*. *i*) **by** *auto*
next
from *2 prime-elem-iff-irreducible* **have** *f* ∈# *F* ⟹ *prime-elem f g* ∈# *G* ⟹ *prime-elem g* **for** *f g*
by (*auto intro*: *prime-elemI*)
then show *Multiset.Ball* (*image-mset normalize F*) *prime*
Multiset.Ball (*image-mset normalize G*) *prime* **by** *auto*
qed
also
obtain *gs* **where** *G*: *G* = *mset gs* **by** (*metis ex-mset*)
have *list-all2* ((*ddvd*)$^{-1-1}$) *gs* (*map normalize gs*) **by** (*intro list-all2-all-nthI*, *auto*)
then have *rel-mset* (*ddvd*) (*image-mset normalize G*) *G*
by (*subst multiset.rel-flip*[*symmetric*], *unfold rel-mset-def G*, *force*)
finally show *?case*.
qed

end

instance *int* :: *ufd* **by** (*intro ufd.intro-of-class as-ufd.ufd-axioms*)
instance *int* :: *idom-gcd* **by** (*intro-classes*, *auto*)

instance *field* ⊆ *ufd* **by** (*intro-classes*, *auto simp*: *dvd-field-iff*)

end

4 Unique Factorization Domain for Polynomials

In this theory we prove that the polynomials over a unique factorization domain (UFD) form a UFD.

theory *Unique-Factorization-Poly*
imports
Unique-Factorization
Polynomial-Factorization.Missing-Polynomial-Factorial
Subresultants.More-Homomorphisms
HOL−Computational-Algebra.Field-as-Ring
begin

hide-const (**open**) *module.smult*
hide-const (**open**) *Divisibility.irreducible*

instantiation *fract* :: (*idom*) {*normalization-euclidean-semiring*, *euclidean-ring*}
begin

definition [*simp*]: *normalize-fract* ≡ (*normalize-field* :: *'a fract* ⇒ -)

definition [*simp*]: *unit-factor-fract* = (*unit-factor-field* :: *'a fract* ⇒ -)
definition [*simp*]: *euclidean-size-fract* = (*euclidean-size-field* :: *'a fract* ⇒ -)
definition [*simp*]: *modulo-fract* = (*mod-field* :: *'a fract* ⇒ -)

instance by *standard* (*simp-all add*: *dvd-field-iff divide-simps*)

end

instantiation *fract* :: (*idom*) *euclidean-ring-gcd*
begin

definition *gcd-fract* :: *'a fract* ⇒ *'a fract* ⇒ *'a fract* **where**
gcd-fract ≡ *Euclidean-Algorithm.gcd*
definition *lcm-fract* :: *'a fract* ⇒ *'a fract* ⇒ *'a fract* **where**
lcm-fract ≡ *Euclidean-Algorithm.lcm*
definition *Gcd-fract* :: *'a fract set* ⇒ *'a fract* **where**
Gcd-fract ≡ *Euclidean-Algorithm.Gcd*
definition *Lcm-fract* :: *'a fract set* ⇒ *'a fract* **where**
Lcm-fract ≡ *Euclidean-Algorithm.Lcm*

instance
by (*standard*, *simp-all add*: *gcd-fract-def lcm-fract-def Gcd-fract-def Lcm-fract-def*)

end

instantiation *fract* :: (*idom*) *unique-euclidean-ring*
begin

definition [*simp*]: *division-segment-fract* (*x* :: *'a fract*) = (*1* :: *'a fract*)

instance by *standard* (*auto split*: *if-splits*)
end

instance *fract* :: (*idom*) *field-gcd* **by** *standard auto*

definition *divides-ff* :: *'a*::*idom fract* ⇒ *'a fract* ⇒ *bool*
where *divides-ff x y* ≡ ∃ *r*. *y* = *x* ∗ *to-fract r*

lemma *ff-list-pairs*:
∃ *xs*. *X* = *map* (λ (*x*,*y*). *Fraction-Field.Fract x y*) *xs* ∧ *0* ∉ *snd* ‘ *set xs*
proof (*induct X*)
case (*Cons a X*)
from *Cons*(*1*) **obtain** *xs* **where** *X*: *X* = *map* (λ (*x*,*y*). *Fraction-Field.Fract x y*) *xs* **and** *xs*: *0* ∉ *snd* ‘ *set xs*
by *auto*
obtain *x y* **where** *a*: *a* = *Fraction-Field.Fract x y* **and** *y*: *y* ≠ *0* **by** (*cases a*, *auto*)

```
  show ?case unfolding X a using xs y
    by (intro exI[of - (x,y) # xs], auto)
qed auto

lemma divides-ff-to-fract[simp]: divides-ff (to-fract x) (to-fract y) ⟷ x dvd y
  unfolding divides-ff-def dvd-def
  by (simp add: to-fract-def eq-fract(1) mult.commute)

lemma
  shows divides-ff-mult-cancel-left[simp]: divides-ff (z * x) (z * y) ⟷ z = 0 ∨
divides-ff x y
    and divides-ff-mult-cancel-right[simp]: divides-ff (x * z) (y * z) ⟷ z = 0 ∨
divides-ff x y
  unfolding divides-ff-def by auto

definition gcd-ff-list :: 'a::ufd fract list ⇒ 'a fract ⇒ bool where
  gcd-ff-list X g = (
    (∀ x ∈ set X. divides-ff g x) ∧
    (∀ d. (∀ x ∈ set X. divides-ff d x) ⟶ divides-ff d g))

lemma gcd-ff-list-exists: ∃ g. gcd-ff-list (X :: 'a::ufd fract list) g
proof −
  interpret some-gcd: idom-gcd (*) 1 :: 'a (+) 0 (−) uminus some-gcd
    rewrites dvd.dvd ((*)) = (dvd) by (unfold-locales, auto simp: dvd-rewrites)
  from ff-list-pairs[of X] obtain xs where X: X = map (λ (x,y). Fraction-Field.Fract
x y) xs
    and xs: 0 ∉ snd ‘ set xs by auto
  define r where r ≡ prod-list (map snd xs)
  have r: r ≠ 0 unfolding r-def prod-list-zero-iff using xs by auto
  define ys where ys ≡ map (λ (x,y). x * prod-list (remove1 y (map snd xs))) xs
  {
    fix i
    assume i < length X
    hence i: i < length xs unfolding X by auto
    obtain x y where xsi: xs ! i = (x,y) by force
    with i have (x,y) ∈ set xs unfolding set-conv-nth by force
    hence y-mem: y ∈ set (map snd xs) by force
    with xs have y: y ≠ 0 by force
    from i have id1: ys ! i = x * prod-list (remove1 y (map snd xs)) unfolding
ys-def using xsi by auto
    from i xsi have id2: X ! i = Fraction-Field.Fract x y unfolding X by auto
    have lp: prod-list (remove1 y (map snd xs)) * y = r unfolding r-def
      by (rule prod-list-remove1[OF y-mem])
    have ys ! i ∈ set ys using i unfolding ys-def by auto
    moreover have to-fract (ys ! i) = to-fract r * (X ! i)
      unfolding id1 id2 to-fract-def mult-fract
      by (subst eq-fract(1), force, force simp: y, simp add: lp)
    ultimately have ys ! i ∈ set ys to-fract (ys ! i) = to-fract r * (X ! i) .
  } note ys = this
```

define G **where** $G \equiv$ *some-gcd.listgcd ys*
define g **where** $g \equiv$ *to-fract G* $*$ *Fraction-Field.Fract 1 r*
have *len*: *length X* $=$ *length ys* **unfolding** *X ys-def* **by** *auto*
show *?thesis*
proof (*rule exI*[*of - g*], *unfold gcd-ff-list-def*, *intro ballI conjI impI allI*)
fix x
assume $x \in$ *set X*
then obtain i **where** i: $i <$ *length X* **and** x: $x =$ *X* ! i **unfolding** *set-conv-nth*
by *auto*
from *ys*[*OF i*] **have** *id*: *to-fract* (*ys* ! i) $=$ *to-fract r* $*$ x
and *ysi*: *ys* ! $i \in$ *set ys* **unfolding** x **by** *auto*
from *some-gcd.listgcd*[*OF ysi*] **have** *G dvd ys* ! i **unfolding** *G-def* **.**
then obtain d **where** *ysi*: *ys* ! $i =$ *G* $*$ d **unfolding** *dvd-def* **by** *auto*
have *to-fract d* $*$ (*to-fract G* $*$ *Fraction-Field.Fract 1 r*) $=$ x $*$ (*to-fract r* $*$
Fraction-Field.Fract 1 r)
using *id*[*unfolded ysi*]
by (*simp add*: *ac-simps*)
also have $\ldots = x$ **using** *r* **unfolding** *to-fract-def* **by** (*simp add*: *eq-fract*
One-fract-def)
finally have *to-fract d* $*$ (*to-fract G* $*$ *Fraction-Field.Fract 1 r*) $= x$ **by** *simp*
thus *divides-ff g x* **unfolding** *divides-ff-def g-def*
by (*intro exI*[*of - d*], *auto*)
next
fix d
assume $\forall x \in$ *set X*. *divides-ff d x*
hence *Ball* ((λ *x*. *to-fract r* $*$ *x*) ‘ *set X*) (*divides-ff* (*to-fract r* $*$ *d*)) **by** *simp*
also have (λ *x*. *to-fract r* $*$ *x*) ‘ *set X* $=$ *to-fract* ‘ *set ys*
unfolding *set-conv-nth* **using** *ys len* **by** *force*
finally have *dvd*: *Ball* (*set ys*) (λ *y*. *divides-ff* (*to-fract r* $*$ *d*) (*to-fract y*)) **by**
auto
obtain *nd dd* **where** *d*: $d =$ *Fraction-Field.Fract nd dd* **and** *dd*: $dd \neq 0$ **by**
(*cases d*, *auto*)
{
fix y
assume $y \in$ *set ys*
hence *divides-ff* (*to-fract r* $*$ *d*) (*to-fract y*) **using** *dvd* **by** *auto*
from *this*[*unfolded divides-ff-def d to-fract-def mult-fract*]
obtain *ra* **where** *Fraction-Field.Fract y 1* $=$ *Fraction-Field.Fract* (*r* $*$ *nd* $*$
ra) *dd* **by** *auto*
hence *y* $*$ *dd* $=$ *ra* $*$ (*r* $*$ *nd*) **by** (*simp add*: *eq-fract dd*)
hence *r* $*$ *nd dvd y* $*$ *dd* **by** *auto*
}
hence *r* $*$ *nd dvd some-gcd.listgcd ys* $*$ *dd* **by** (*rule some-gcd.listgcd-greatest-mult*)
hence *divides-ff* (*to-fract r* $*$ *d*) (*to-fract G*) **unfolding** *to-fract-def d mult-fract*
G-def divides-ff-def **by** (*auto simp add*: *eq-fract dd dvd-def*)
also have *to-fract G* $=$ *to-fract r* $*$ *g* **unfolding** *g-def* **using** *r*
by (*auto simp*: *to-fract-def eq-fract*)
finally show *divides-ff d g* **using** *r* **by** *simp*
qed

qed

definition *some-gcd-ff-list* :: *'a* :: *ufd fract list* ⇒ *'a fract* **where**
some-gcd-ff-list xs = (*SOME g*. *gcd-ff-list xs g*)

lemma *some-gcd-ff-list*: *gcd-ff-list xs* (*some-gcd-ff-list xs*)
unfolding *some-gcd-ff-list-def* **using** *gcd-ff-list-exists*[*of xs*]
by (*rule someI-ex*)

lemma *some-gcd-ff-list-divides*: *x* ∈ *set xs* ⟹ *divides-ff* (*some-gcd-ff-list xs*) *x*
using *some-gcd-ff-list*[*of xs*] **unfolding** *gcd-ff-list-def* **by** *auto*

lemma *some-gcd-ff-list-greatest*: (∀ *x* ∈ *set xs*. *divides-ff d x*) ⟹ *divides-ff d* (*some-gcd-ff-list xs*)
using *some-gcd-ff-list*[*of xs*] **unfolding** *gcd-ff-list-def* **by** *auto*

lemma *divides-ff-refl*[*simp*]: *divides-ff x x*
unfolding *divides-ff-def*
by (*rule exI*[*of - 1*], *auto simp*: *to-fract-def One-fract-def*)

lemma *divides-ff-trans*:
divides-ff x y ⟹ *divides-ff y z* ⟹ *divides-ff x z*
unfolding *divides-ff-def*
by (*auto simp del*: *to-fract-hom.hom-mult simp add*: *to-fract-hom.hom-mult*[*symmetric*])

lemma *divides-ff-mult-right*: *a* ≠ *0* ⟹ *divides-ff* (*x* ∗ *inverse a*) *y* ⟹ *divides-ff x* (*a* ∗ *y*)
unfolding *divides-ff-def divide-inverse*[*symmetric*] **by** *auto*

definition *eq-dff* :: *'a* :: *ufd fract* ⇒ *'a fract* ⇒ *bool* (**infix** ‹=*dff*› *50*) **where**
x =*dff y* ⟷ *divides-ff x y* ∧ *divides-ff y x*

lemma *eq-dffI*[*intro*]: *divides-ff x y* ⟹ *divides-ff y x* ⟹ *x* =*dff y*
unfolding *eq-dff-def* **by** *auto*

lemma *eq-dff-refl*[*simp*]: *x* =*dff x*
by (*intro eq-dffI*, *auto*)

lemma *eq-dff-sym*: *x* =*dff y* ⟹ *y* =*dff x* **unfolding** *eq-dff-def* **by** *auto*

lemma *eq-dff-trans*[*trans*]: *x* =*dff y* ⟹ *y* =*dff z* ⟹ *x* =*dff z*
unfolding *eq-dff-def* **using** *divides-ff-trans* **by** *auto*

lemma *eq-dff-cancel-right*[*simp*]: *x* ∗ *y* =*dff x* ∗ *z* ⟷ *x* = *0* ∨ *y* =*dff z*
unfolding *eq-dff-def* **by** *auto*

lemma *eq-dff-mult-right-trans*[*trans*]: *x* =*dff y* ∗ *z* ⟹ *z* =*dff u* ⟹ *x* =*dff y* ∗ *u*
using *eq-dff-trans* **by** *force*

lemma *some-gcd-ff-list-smult*: $a \neq 0 \Longrightarrow$ *some-gcd-ff-list* (*map* ((∗) *a*) *xs*) =*dff a* ∗ *some-gcd-ff-list xs*
proof
let *?g* = *some-gcd-ff-list* (*map* ((∗) *a*) *xs*)
show *divides-ff* (*a* ∗ *some-gcd-ff-list xs*) *?g*
by (*rule some-gcd-ff-list-greatest*, *insert some-gcd-ff-list-divides*[*of* - *xs*], *auto simp*: *divides-ff-def*)
assume *a*: $a \neq 0$
show *divides-ff ?g* (*a* ∗ *some-gcd-ff-list xs*)
proof (*rule divides-ff-mult-right*[*OF a some-gcd-ff-list-greatest*], *intro ballI*)
fix *x*
assume *x*: $x \in$ *set xs*
have *divides-ff* (*?g* ∗ *inverse a*) *x* = *divides-ff* (*inverse a* ∗ *?g*) (*inverse a* ∗ (*a* ∗ *x*))
using *a* **by** (*simp add*: *field-simps*)
also have ... **using** *a x* **by** (*auto intro*: *some-gcd-ff-list-divides*)
finally show *divides-ff* (*?g* ∗ *inverse a*) *x* **.**
qed
qed

definition *content-ff* :: *'a*::*ufd fract poly* ⇒ *'a fract* **where**
content-ff p = *some-gcd-ff-list* (*coeffs p*)

lemma *content-ff-iff*: *divides-ff x* (*content-ff p*) ⟷ (∀ *c* ∈ *set* (*coeffs p*). *divides-ff x c*) (**is** *?l* = *?r*)
proof
assume *?l*
from *divides-ff-trans*[*OF this*, *unfolded content-ff-def*, *OF some-gcd-ff-list-divides*]
show *?r* **..**
next
assume *?r*
thus *?l* **unfolding** *content-ff-def* **by** (*intro some-gcd-ff-list-greatest*, *auto*)
qed

lemma *content-ff-divides-ff*: $x \in$ *set* (*coeffs p*) $\Longrightarrow$ *divides-ff* (*content-ff p*) *x*
unfolding *content-ff-def* **by** (*rule some-gcd-ff-list-divides*)

lemma *content-ff-0*[*simp*]: *content-ff 0* = *0*
using *content-ff-iff*[*of 0 0*] **by** (*auto simp*: *divides-ff-def*)

lemma *content-ff-0-iff*[*simp*]: (*content-ff p* = *0*) = (*p* = *0*)
proof (*cases p* = *0*)
case *False*
define *a* **where** *a* ≡ *last* (*coeffs p*)
define *xs* **where** *xs* ≡ *coeffs p*
from *False*
have *mem*: $a \in$ *set* (*coeffs p*) **and** *a*: $a \neq 0$
unfolding *a-def last-coeffs-eq-coeff-degree*[*OF False*] *coeffs-def* **by** *auto*
from *content-ff-divides-ff*[*OF mem*] **have** *divides-ff* (*content-ff p*) *a* **.**

with *a* **have** *content-ff p* ≠ *0* **unfolding** *divides-ff-def* **by** *auto*
with *False* **show** *?thesis* **by** *auto*
qed *auto*

lemma *content-ff-eq-dff-nonzero*: *content-ff p* =*dff x* ⟹ *x* ≠ *0* ⟹ *p* ≠ *0*
using *divides-ff-def eq-dff-def* **by** *force*

lemma *content-ff-smult*: *content-ff* (*smult* (*a*::′*a*::*ufd fract*) *p*) =*dff a* ∗ *content-ff p*
proof (*cases a* = *0*)
case *False* **note** *a* = *this*
have *id*: *coeffs* (*smult a p*) = *map* ((∗) *a*) (*coeffs p*)
unfolding *coeffs-smult* **using** *a* **by** (*simp add*: *Polynomial.coeffs-smult*)
show *?thesis* **unfolding** *content-ff-def id* **using** *some-gcd-ff-list-smult*[*OF a*] **.**
qed *simp*

definition *normalize-content-ff*
where *normalize-content-ff* (*p*::′*a*::*ufd fract poly*) ≡ *smult* (*inverse* (*content-ff p*)) *p*

lemma *smult-normalize-content-ff*: *smult* (*content-ff p*) (*normalize-content-ff p*) = *p*
unfolding *normalize-content-ff-def*
by (*cases p* = *0*, *auto*)

lemma *content-ff-normalize-content-ff-1*: **assumes** *p0*: *p* ≠ *0*
shows *content-ff* (*normalize-content-ff p*) =*dff 1*
proof −
have *content-ff p* = *content-ff* (*smult* (*content-ff p*) (*normalize-content-ff p*))
unfolding *smult-normalize-content-ff* **..**
also have . . . =*dff content-ff p* ∗ *content-ff* (*normalize-content-ff p*) **by** (*rule content-ff-smult*)
finally show *?thesis* **unfolding** *eq-dff-def divides-ff-def* **using** *p0* **by** *auto*
qed

lemma *content-ff-to-fract*: **assumes** *set* (*coeffs p*) ⊆ *range to-fract*
shows *content-ff p* ∈ *range to-fract*
proof −
have *divides-ff 1* (*content-ff p*) **using** *assms*
unfolding *content-ff-iff* **unfolding** *divides-ff-def*[*abs-def*] **by** *auto*
thus *?thesis* **unfolding** *divides-ff-def* **by** *auto*
qed

lemma *content-ff-map-poly-to-fract*: *content-ff* (*map-poly to-fract* (*p* :: ′*a* :: *ufd poly*)) ∈ *range to-fract*
by (*rule content-ff-to-fract*, *subst coeffs-map-poly*, *auto*)

lemma *range-coeffs-to-fract*: **assumes** *set* (*coeffs p*) ⊆ *range to-fract*
shows ∃ *m*. *coeff p i* = *to-fract m*

proof −
from *assms*(*1*) *to-fract-0* **have** *coeff p i* ∈ *range to-fract* **using** *range-coeff* [*of p*]
by *auto* (*metis contra-subsetD to-fract-hom.hom-zero insertE range-eqI*)
thus *?thesis* **by** *auto*
qed

lemma *divides-ff-coeff*: **assumes** *set* (*coeffs p*) ⊆ *range to-fract* **and** *divides-ff* (*to-fract n*) (*coeff p i*)
shows ∃ *m*. *coeff p i* = *to-fract n* ∗ *to-fract m*
proof −
from *range-coeffs-to-fract*[*OF assms*(*1*)] **obtain** *k* **where** *pi*: *coeff p i* = *to-fract k* **by** *auto*
from *assms*(*2*)[*unfolded this*] **have** *n dvd k* **by** *simp*
then obtain *j* **where** *k*: *k* = *n* ∗ *j* **unfolding** *Rings.dvd-def* **by** *auto*
show *?thesis* **unfolding** *pi k* **by** *auto*
qed

definition *inv-embed* :: *'a* :: *ufd fract* ⇒ *'a* **where**
inv-embed = *the-inv to-fract*

lemma *inv-embed*[*simp*]: *inv-embed* (*to-fract x*) = *x*
unfolding *inv-embed-def*
by (*rule the-inv-f-f*, *auto simp*: *inj-on-def*)

lemma *inv-embed-0*[*simp*]: *inv-embed 0* = *0* **unfolding** *to-fract-0*[*symmetric*] *inv-embed*
by *simp*

lemma *range-to-fract-embed-poly*: **assumes** *set* (*coeffs p*) ⊆ *range to-fract*
shows *p* = *map-poly to-fract* (*map-poly inv-embed p*)
proof −
have *p* = *map-poly* (*to-fract o inv-embed*) *p*
by (*rule sym*, *rule map-poly-idI*, *insert assms*, *auto*)
also have ... = *map-poly to-fract* (*map-poly inv-embed p*)
by (*subst map-poly-map-poly*, *auto*)
finally show *?thesis* **.**
qed

lemma *content-ff-to-fract-coeffs-to-fract*: **assumes** *content-ff p* ∈ *range to-fract*
shows *set* (*coeffs p*) ⊆ *range to-fract*
proof
fix *x*
assume *x* ∈ *set* (*coeffs p*)
from *content-ff-divides-ff*[*OF this*] *assms*[*unfolded eq-dff-def*] **show** *x* ∈ *range to-fract*
unfolding *divides-ff-def* **by** (*auto simp del*: *to-fract-hom.hom-mult simp*: *to-fract-hom.hom-mult*[*symmetric*]
qed

lemma *content-ff-1-coeffs-to-fract*: **assumes** *content-ff p* =*dff 1*

```
  shows set (coeffs p) ⊆ range to-fract
proof
  fix x
  assume x ∈ set (coeffs p)
   from content-ff-divides-ff[OF this] assms[unfolded eq-dff-def] show x ∈ range
to-fract
   unfolding divides-ff-def by (auto simp del: to-fract-hom.hom-mult simp: to-fract-hom.hom-mult[symmetric]
qed

lemma gauss-lemma:
  fixes p q :: 'a :: ufd fract poly
  shows content-ff (p * q) =dff content-ff p * content-ff q
proof (cases p = 0 ∨ q = 0)
  case False
  hence p: p ≠ 0 and q: q ≠ 0 by auto
  let ?c = content-ff :: 'a fract poly ⇒ 'a fract
  {
    fix p q :: 'a fract poly
    assume cp1: ?c p =dff 1 and cq1: ?c q =dff 1
    define ip where ip ≡ map-poly inv-embed p
    define iq where iq ≡ map-poly inv-embed q
    interpret map-poly-hom: map-poly-comm-ring-hom to-fract..
   from content-ff-1-coeffs-to-fract[OF cp1] have cp: set (coeffs p) ⊆ range to-fract
.
   from content-ff-1-coeffs-to-fract[OF cq1] have cq: set (coeffs q) ⊆ range to-fract
.
    have ip: p = map-poly to-fract ip unfolding ip-def
      by (rule range-to-fract-embed-poly[OF cp])
    have iq: q = map-poly to-fract iq unfolding iq-def
      by (rule range-to-fract-embed-poly[OF cq])
    have cpq0: ?c (p * q) ≠ 0
      unfolding content-ff-0-iff using cp1 cq1 content-ff-eq-dff-nonzero[of - 1] by
auto
    have cpq: set (coeffs (p * q)) ⊆ range to-fract unfolding ip iq
   unfolding map-poly-hom.hom-mult[symmetric] to-fract-hom.coeffs-map-poly-hom
by auto
    have ctnt: ?c (p * q) ∈ range to-fract using content-ff-to-fract[OF cpq] .
    then obtain cpq where id: ?c (p * q) = to-fract cpq by auto
   have dvd: divides-ff 1 (?c (p * q)) using ctnt unfolding divides-ff-def by auto
   from cpq0[unfolded id] have cpq0: cpq ≠ 0 unfolding to-fract-def Zero-fract-def
by auto
    hence cpqM: cpq ∈ carrier mk-monoid by auto
    have ?c (p * q) =dff 1
    proof (rule ccontr)
      assume ¬ ?c (p * q) =dff 1
      with dvd have ¬ divides-ff (?c (p * q)) 1
        unfolding eq-dff-def by auto
      from this[unfolded id divides-ff-def] have cpq: ⋀ r. cpq * r ≠ 1
        by (auto simp: to-fract-def One-fract-def eq-fract)
```

then have *cpq1*: ¬ *cpq dvd 1* **by** (*auto elim*:*dvdE simp*:*ac-simps*)
from *mset-factors-exist*[*OF cpq0 cpq1*]
obtain *F* **where** *F*: *mset-factors F cpq* **by** *auto*
have *F* ≠ {#} **using** *F* **by** *auto*
then obtain *f* **where** *f*: *f* ∈# *F* **by** *auto*
with *F* **have** *irrf*: *irreducible f* **and** *f0*: *f* ≠ *0* **by** (*auto dest*: *mset-factorsD*)
from *irrf* **have** *pf*: *prime-elem f* **by** *simp*
note ∗ = *this*[*unfolded prime-elem-def*]
from ∗ **have** *no-unit*: ¬ *f dvd 1* **by** *auto*
from ∗ *f0* **have** *prime*: ⋀ *a b*. *f dvd a* ∗ *b* ⟹ *f dvd a* ∨ *f dvd b* **unfolding** *dvd-def* **by** *force*
let *?f* = *to-fract f*
from *F f*
have *fdvd*: *f dvd cpq* **by** (*auto intro*:*mset-factors-imp-dvd*)
hence *divides-ff ?f* (*to-fract cpq*) **by** *simp*
from *divides-ff-trans*[*OF this*, *folded id*, *OF content-ff-divides-ff*]
have *dvd*: ⋀ *z*. *z* ∈ *set* (*coeffs* (*p* ∗ *q*)) ⟹ *divides-ff ?f z* **.**
{
fix *p* :: *'a fract poly*
assume *cp*: *?c p* =*dff 1*
let *?P* = λ *i*. ¬ *divides-ff ?f* (*coeff p i*)
{
assume ∀ *c* ∈ *set* (*coeffs p*). *divides-ff ?f c*
hence *n*: *divides-ff ?f* (*?c p*) **unfolding** *content-ff-iff* **by** *auto*
from *divides-ff-trans*[*OF this*] *cp*[*unfolded eq-dff-def*] **have** *divides-ff ?f 1* **by** *auto*
also have *1* = *to-fract 1* **by** *simp*
finally have *f dvd 1* **by** (*unfold divides-ff-to-fract*)
hence *False* **using** *no-unit* **unfolding** *dvd-def* **by** (*auto simp*: *ac-simps*)
}
then obtain *cp* **where** *cp*: *cp* ∈ *set* (*coeffs p*) **and** *ncp*: ¬ *divides-ff ?f cp* **by** *auto*
hence *cp* ∈ *range* (*coeff p*) **unfolding** *range-coeff* **by** *auto*
with *ncp* **have** ∃ *i*. *?P i* **by** *auto*
from *LeastI-ex*[*OF this*] *not-less-Least*[*of - ?P*]
have ∃ *i*. *?P i* ∧ (∀ *j*. *j* < *i* ⟶ *divides-ff ?f* (*coeff p j*)) **by** *blast*
} note *cont* = *this*
from *cont*[*OF cp1*] **obtain** *r* **where**
r: ¬ *divides-ff ?f* (*coeff p r*) **and** *r'*: ⋀ *i*. *i* < *r* ⟹ *divides-ff ?f* (*coeff p i*) **by** *auto*
have ∀ *i*. ∃ *k*. *i* < *r* ⟶ *coeff p i* = *?f* ∗ *to-fract k* **using** *divides-ff-coeff*[*OF cp r'*] **by** *blast*
from *choice*[*OF this*] **obtain** *rr* **where** *r'*: ⋀ *i*. *i* < *r* ⟹ *coeff p i* = *?f* ∗ *to-fract* (*rr i*) **by** *blast*
let *?r* = *coeff p r*
from *cont*[*OF cq1*] **obtain** *s* **where**
s: ¬ *divides-ff ?f* (*coeff q s*) **and** *s'*: ⋀ *i*. *i* < *s* ⟹ *divides-ff ?f* (*coeff q i*) **by** *auto*
have ∀ *i*. ∃ *k*. *i* < *s* ⟶ *coeff q i* = *?f* ∗ *to-fract k* **using** *divides-ff-coeff*[*OF*

cq s'] **by** *blast*
from *choice*[*OF this*] **obtain** *ss* **where** *s'*: $\bigwedge$ *i*. *i* < *s* $\Longrightarrow$ *coeff q i* = *?f* ∗ *to-fract* (*ss i*) **by** *blast*
from *range-coeffs-to-fract*[*OF cp*] **have** ∀ *i*. ∃ *m*. *coeff p i* = *to-fract m* **..**
from *choice*[*OF this*] **obtain** *pi* **where** *pi*: $\bigwedge$ *i*. *coeff p i* = *to-fract* (*pi i*) **by** *blast*
from *range-coeffs-to-fract*[*OF cq*] **have** ∀ *i*. ∃ *m*. *coeff q i* = *to-fract m* **..**
from *choice*[*OF this*] **obtain** *qi* **where** *qi*: $\bigwedge$ *i*. *coeff q i* = *to-fract* (*qi i*) **by** *blast*
let *?s* = *coeff q s*
let *?g* = λ *i*. *coeff p i* ∗ *coeff q* (*r* + *s* − *i*)
define *a* **where** *a* = ($\sum$ *i*∈{..<*r*}. (*rr i* ∗ *qi* (*r* + *s* − *i*)))
define *b* **where** *b* = ($\sum$ *i* ∈ {*Suc r*..*r* + *s*}. *pi i* ∗ (*ss* (*r* + *s* − *i*)))
have *coeff* (*p* ∗ *q*) (*r* + *s*) = ($\sum$ *i*≤*r* + *s*. *?g i*) **unfolding** *coeff-mult* **..**
also have {..*r*+*s*} = {..< *r*} ∪ {*r* .. *r*+*s*} **by** *auto*
also have ($\sum$ *i*∈{..<*r*} ∪ {*r*..*r* + *s*}. *?g i*)
= ($\sum$ *i*∈{..<*r*}. *?g i*) + ($\sum$ *i* ∈ {*r*..*r* + *s*}. *?g i*)
by (*rule sum.union-disjoint*, *auto*)
also have ($\sum$ *i*∈{..<*r*}. *?g i*) = ($\sum$ *i*∈{..<*r*}. *?f* ∗ (*to-fract* (*rr i*) ∗ *to-fract* (*qi* (*r* + *s* − *i*))))
by (*rule sum.cong*[*OF refl*], *insert r' qi*, *auto*)
also have ... = *to-fract* (*f* ∗ *a*) **by** (*simp add*: *a-def sum-distrib-left*)
also have ($\sum$ *i* ∈ {*r*..*r* + *s*}. *?g i*) = *?g r* + ($\sum$ *i* ∈ {*Suc r*..*r* + *s*}. *?g i*)
by (*subst sum.remove*[*of* - *r*], *auto intro*: *sum.cong*)
also have ($\sum$ *i* ∈ {*Suc r*..*r* + *s*}. *?g i*) = ($\sum$ *i* ∈ {*Suc r*..*r* + *s*}. *?f* ∗ (*to-fract* (*pi i*) ∗ *to-fract* (*ss* (*r* + *s* − *i*))))
by (*rule sum.cong*[*OF refl*], *insert s' pi*, *auto*)
also have ... = *to-fract* (*f* ∗ *b*) **by** (*simp add*: *sum-distrib-left b-def*)
finally have *cpq*: *coeff* (*p* ∗ *q*) (*r* + *s*) = *to-fract* (*f* ∗ (*a* + *b*)) + *?r* ∗ *?s* **by** (*simp add*: *field-simps*)
{
fix *i*
from *dvd*[*of coeff* (*p* ∗ *q*) *i*] **have** *divides-ff ?f* (*coeff* (*p* ∗ *q*) *i*) **using** *range-coeff*[*of p* ∗ *q*]
by (*cases coeff* (*p* ∗ *q*) *i* = *0*, *auto simp*: *divides-ff-def*)
}
from *this*[*of r* + *s*, *unfolded cpq*] **have** *divides-ff ?f* (*to-fract* (*f* ∗ (*a* + *b*) + *pi r* ∗ *qi s*))
unfolding *pi qi* **by** *simp*
from *this*[*unfolded divides-ff-to-fract*] **have** *f dvd pi r* ∗ *qi s*
by (*metis dvd-add-times-triv-left-iff mult.commute*)
from *prime*[*OF this*] **have** *f dvd pi r* ∨ *f dvd qi s* **by** *auto*
with *r s* **show** *False* **unfolding** *pi qi* **by** *auto*
qed
} note *main* = *this*
define *n* **where** *n* ≡ *normalize-content-ff* :: *'a fract poly* ⇒ *'a fract poly*
let *?s* = λ *p*. *smult* (*content-ff p*) (*n p*)
have *?c* (*p* ∗ *q*) = *?c* (*?s p* ∗ *?s q*) **unfolding** *smult-normalize-content-ff n-def* **by** *simp*

```
  also have ?s p * ?s q = smult (?c p * ?c q) (n p * n q) by (simp add:
mult.commute)
  also have ?c (...) =dff (?c p * ?c q) * ?c (n p * n q) by (rule content-ff-smult)
  also have ?c (n p * n q) =dff 1 unfolding n-def
    by (rule main, insert p q, auto simp: content-ff-normalize-content-ff-1)
  finally show ?thesis by simp
qed auto

abbreviation (input) content-ff-ff p ≡ content-ff (map-poly to-fract p)

lemma factorization-to-fract:
  assumes q: q ≠ 0 and factor: map-poly to-fract (p :: 'a :: ufd poly) = q * r
  shows ∃ q' r' c. c ≠ 0 ∧ q = smult c (map-poly to-fract q') ∧
    r = smult (inverse c) (map-poly to-fract r') ∧
    content-ff-ff q' =dff 1 ∧ p = q' * r'
proof −
  let ?c = content-ff
  let ?p = map-poly to-fract p
  interpret map-poly-inj-comm-ring-hom to-fract :: 'a ⇒ -..
  define cq where cq ≡ normalize-content-ff q
  define cr where cr ≡ smult (content-ff q) r
  define q' where q' ≡ map-poly inv-embed cq
  define r' where r' ≡ map-poly inv-embed cr
  from content-ff-map-poly-to-fract have cp-ff: ?c ?p ∈ range to-fract by auto
  from smult-normalize-content-ff[of q] have cqs: q = smult (content-ff q) cq un-
folding cq-def ..
   from content-ff-normalize-content-ff-1[OF q] have c-cq: content-ff cq =dff 1
unfolding cq-def .
   from content-ff-1-coeffs-to-fract[OF this] have cq-ff: set (coeffs cq) ⊆ range
to-fract .
  have factor: ?p = cq * cr unfolding factor cr-def using cqs
    by (metis mult-smult-left mult-smult-right)
  from gauss-lemma[of cq cr] have cp: ?c ?p =dff ?c cq * ?c cr unfolding factor
.
  with c-cq have ?c ?p =dff ?c cr
    by (metis eq-dff-mult-right-trans mult.commute mult.right-neutral)
  with cp-ff have ?c cr ∈ range to-fract
    by (metis divides-ff-def to-fract-hom.hom-mult eq-dff-def image-iff range-eqI)
 from content-ff-to-fract-coeffs-to-fract[OF this] have cr-ff: set (coeffs cr) ⊆ range
to-fract by auto
  have cq: cq = map-poly to-fract q' unfolding q'-def
    by (rule range-to-fract-embed-poly[OF cq-ff])
  have cr: cr = map-poly to-fract r' unfolding r'-def
    by (rule range-to-fract-embed-poly[OF cr-ff])
  from factor[unfolded cq cr]
  have p: p = q' * r' by (simp add: injectivity)
  from c-cq have ctnt: content-ff-ff q' =dff 1 using cq q'-def by force
  from cqs have idq: q = smult (?c q) (map-poly to-fract q') unfolding cq .
  with q have cq: ?c q ≠ 0 by auto
```

```
  have r = smult (inverse (?c q)) cr unfolding cr-def using cq by auto
  also have cr = map-poly to-fract r' by (rule cr)
  finally have idr: r = smult (inverse (?c q)) (map-poly to-fract r') by auto
  from cq p ctnt idq idr show ?thesis by blast
qed

lemma irreducible-PM-M-PFM:
  assumes irr: irreducible p
  shows degree p = 0 ∧ irreducible (coeff p 0) ∨
  degree p ≠ 0 ∧ irreducible (map-poly to-fract p) ∧ content-ff-ff p =dff 1
proof−
  interpret map-poly-inj-idom-hom to-fract..
  from irr[unfolded irreducible-altdef]
  have p0: p ≠ 0 and irr: ¬ p dvd 1 ⋀ b. b dvd p ⟹ ¬ p dvd b ⟹ b dvd 1 by
auto
  show ?thesis
  proof (cases degree p = 0)
    case True
    from degree0-coeffs[OF True] obtain a where p: p = [:a:] by auto
    note irr = irr[unfolded p]
    from p p0 have a0: a ≠ 0 by auto
    moreover have ¬ a dvd 1 using irr(1) by simp
    moreover {
      fix b
      assume b dvd a ¬ a dvd b
      hence [:b:] dvd [:a:] ¬ [:a:] dvd [:b:] unfolding const-poly-dvd .
      from irr(2)[OF this] have b dvd 1 unfolding const-poly-dvd-1 .
    }
    ultimately have irreducible a unfolding irreducible-altdef by auto
    with True show ?thesis unfolding p by auto
  next
    case False
    let ?E = map-poly to-fract
    let ?p = ?E p
    have dp: degree ?p ≠ 0 using False by simp
    from p0 have p': ?p ≠ 0 by simp
    moreover have ¬ ?p dvd 1
      proof
        assume ?p dvd 1 then obtain q where id: ?p * q = 1 unfolding dvd-def
by auto
        have deg: degree (?p * q) = degree ?p + degree q
          by (rule degree-mult-eq, insert id, auto)
        from arg-cong[OF id, of degree, unfolded deg] dp show False by auto
      qed
    moreover {
      fix q
      assume q dvd ?p and ndvd: ¬ ?p dvd q
      then obtain r where fact: ?p = q * r unfolding dvd-def by auto
      with p' have q0: q ≠ 0 by auto
```

```
    from factorization-to-fract[OF this fact] obtain q' r' c where *: c ≠ 0 q =
smult c (?E q')
      r = smult (inverse c) (?E r') content-ff-ff q' =dff 1
      p = q' * r' by auto
    hence q' dvd p unfolding dvd-def by auto
    note irr = irr(2)[OF this]
    have ¬ p dvd q'
    proof
      assume p dvd q'
      then obtain u where q': q' = p * u unfolding dvd-def by auto
      from arg-cong[OF this, of λ x. smult c (?E x), unfolded *(2)[symmetric]]
      have q = ?p * smult c (?E u) by simp
      hence ?p dvd q unfolding dvd-def by blast
      with ndvd show False ..
    qed
    from irr[OF this] have q' dvd 1 .
    from divides-degree[OF this] have degree q' = 0 by auto
    from degree0-coeffs[OF this] obtain a' where q' = [:a':] by auto
    from *(2)[unfolded this] obtain a where q: q = [:a:]
      by (simp add: to-fract-hom.map-poly-pCons-hom)
    with q0 have a: a ≠ 0 by auto
    have q dvd 1 unfolding q const-poly-dvd-1 using a unfolding dvd-def
      by (intro exI[of - inverse a], auto)
  }
  ultimately have irr-p': irreducible ?p unfolding irreducible-altdef by auto
  let ?c = content-ff
  have ?c ?p ∈ range to-fract
    by (rule content-ff-to-fract, unfold to-fract-hom.coeffs-map-poly-hom, auto)
  then obtain c where cp: ?c ?p = to-fract c by auto
  from p' cp have c: c ≠ 0 by auto
  have ?c ?p =dff 1 unfolding cp
  proof (rule ccontr)
    define cp where cp = normalize-content-ff ?p
    from smult-normalize-content-ff[of ?p] have cps: ?p = smult (to-fract c) cp
unfolding cp-def cp ..
   from content-ff-normalize-content-ff-1[OF p'] have c-cp: content-ff cp =dff 1
unfolding cp-def .
    from range-to-fract-embed-poly[OF content-ff-1-coeffs-to-fract[OF c-cp]] ob-
tain cp' where cp = ?E cp' by auto
    from cps[unfolded this] have p = smult c cp' by (simp add: injectivity)
    hence dvd: [: c :] dvd p unfolding dvd-def by auto
    have ¬ p dvd [: c :] using divides-degree[of p [: c :]] c False by auto
    from irr(2)[OF dvd this] have c dvd 1 by simp
    assume ¬ to-fract c =dff 1
  from this[unfolded eq-dff-def One-fract-def to-fract-def[symmetric] divides-ff-def
to-fract-mult]
  have c1: ⋀ r. 1 ≠ c * r by (auto simp: ac-simps simp del: to-fract-hom.hom-mult
simp: to-fract-hom.hom-mult[symmetric])
    with ‹c dvd 1› show False unfolding dvd-def by blast
```

```
    qed
    with False irr-p' show ?thesis by auto
  qed
qed

lemma irreducible-M-PM:
  fixes p :: 'a :: ufd poly assumes 0: degree p = 0 and irr: irreducible (coeff p 0)
  shows irreducible p
proof (cases p = 0)
  case True
  thus ?thesis using assms by auto
next
  case False
  from degree0-coeffs[OF 0] obtain a where p: p = [:a:] by auto
  with False have a0: a ≠ 0 by auto
  from p irr have irreducible a by auto
  from this[unfolded irreducible-altdef]
  have a1: ¬ a dvd 1 and irr: ⋀ b. b dvd a ⟹ ¬ a dvd b ⟹ b dvd 1 by auto
  {
    fix b
    assume *: b dvd [:a:] ¬ [:a:] dvd b
    from divides-degree[OF this(1)] a0 have degree b = 0 by auto
    from degree0-coeffs[OF this] obtain bb where b: b = [: bb :] by auto
    from * irr[of bb] have b dvd 1 unfolding b const-poly-dvd by auto
  }
  with a0 a1 show ?thesis by (auto simp: irreducible-altdef p)
qed

lemma primitive-irreducible-imp-degree:
 primitive (p::'a::{semiring-gcd,idom} poly) ⟹ irreducible p ⟹ degree p > 0
  by (unfold irreducible-primitive-connect[symmetric], auto)

lemma irreducible-degree-field:
  fixes p :: 'a :: field poly assumes irreducible p
  shows degree p > 0
proof−
  {
    assume degree p = 0
    from degree0-coeffs[OF this] assms obtain a where p: p = [:a:] and a: a ≠ 0
by auto
    hence 1 = p * [:inverse a:] by auto
    hence p dvd 1 ..
    hence p ∈ Units mk-monoid by simp
    with assms have False unfolding irreducible-def by auto
  } then show ?thesis by auto
qed

lemma irreducible-PFM-PM: assumes
  irr: irreducible (map-poly to-fract p) and ct: content-ff-ff p =dff 1
```

```
  shows irreducible p
proof −
  let ?E = map-poly to-fract
  let ?p = ?E p
  from ct have p0: p ≠ 0 by (auto simp: eq-dff-def divides-ff-def)
  moreover
    from irreducible-degree-field[OF irr] have deg: degree p ≠ 0 by simp
    from irr[unfolded irreducible-altdef]
    have irr: ⋀ b. b dvd ?p ⟹ ¬ ?p dvd b ⟹ b dvd 1 by auto
    have ¬ p dvd 1 using deg divides-degree[of p 1] by auto
  moreover {
    fix q :: 'a poly
    assume dvd: q dvd p and ndvd: ¬ p dvd q
    from dvd obtain r where pqr: p = q * r ..
    from arg-cong[OF this, of ?E] have pqr': ?p = ?E q * ?E r by simp
    from p0 pqr have q: q ≠ 0 and r: r ≠ 0 by auto
    have dp: degree p = degree q + degree r unfolding pqr
      by (subst degree-mult-eq, insert q r, auto)
    from eq-dff-trans[OF eq-dff-sym[OF gauss-lemma[of ?E q ?E r, folded pqr']] ct]
    have ct: content-ff (?E q) * content-ff (?E r) =dff 1 .
     from content-ff-map-poly-to-fract obtain cq where cq: content-ff (?E q) =
to-fract cq by auto
     from content-ff-map-poly-to-fract obtain cr where cr: content-ff (?E r) =
to-fract cr by auto
    note ct[unfolded cq cr to-fract-mult eq-dff-def divides-ff-def]
    from this[folded hom-distribs]
    obtain c where c: cq * cr * c = 1 by (auto simp del: to-fract-hom.hom-mult
simp: to-fract-hom.hom-mult[symmetric])
    hence one: 1 = cq * (c * cr) 1 = cr * (c * cq) by (auto simp: ac-simps)
    {
      assume *: degree q ≠ 0 ∧ degree r ≠ 0
      with dp have degree q < degree p by auto
      hence degree (?E q) < degree (?E p) by simp
      hence ndvd: ¬ ?p dvd ?E q using divides-degree[of ?p ?E q] q by auto
      have ?E q dvd ?p unfolding pqr' by auto
      from irr[OF this ndvd] have ?E q dvd 1 .
      from divides-degree[OF this] * have False by auto
    }
    hence degree q = 0 ∨ degree r = 0 by blast
    then have q dvd 1
    proof
      assume degree q = 0
     from degree0-coeffs[OF this] q obtain a where q: q = [:a:] and a: a ≠ 0 by
auto
      hence id: set (coeffs (?E q)) = {to-fract a} by auto
     have divides-ff (to-fract a) (content-ff (?E q)) unfolding content-ff-iff id by
auto
      from this[unfolded cq divides-ff-def, folded hom-distribs]
      obtain rr where cq: cq = a * rr by (auto simp del: to-fract-hom.hom-mult
```

simp: *to-fract-hom.hom-mult*[*symmetric*])
with *one*(*1*) **have** *1* = *a* ∗ (*rr* ∗ *c* ∗ *cr*) **by** (*auto simp*: *ac-simps*)
hence *a dvd 1* **..**
thus *?thesis* **by** (*simp add*: *q*)
next
assume *degree r* = *0*
from *degree0-coeffs*[*OF this*] *r* **obtain** *a* **where** *r*: *r* = [:*a*:] **and** *a*: *a* ≠ *0* **by** *auto*
hence *id*: *set* (*coeffs* (*?E r*)) = {*to-fract a*} **by** *auto*
have *divides-ff* (*to-fract a*) (*content-ff* (*?E r*)) **unfolding** *content-ff-iff id* **by** *auto*
note *this*[*unfolded cr divides-ff-def to-fract-mult*]
note *this*[*folded hom-distribs*]
then obtain *rr* **where** *cr*: *cr* = *a* ∗ *rr* **by** (*auto simp del*: *to-fract-hom.hom-mult simp*: *to-fract-hom.hom-mult*[*symmetric*])
with *one*(*2*) **have** *one*: *1* = *a* ∗ (*rr* ∗ *c* ∗ *cq*) **by** (*auto simp*: *ac-simps*)
from *arg-cong*[*OF pqr*[*unfolded r*], *of* λ *p*. *p* ∗ [:*rr* ∗ *c* ∗ *cq*:]]
have *p* ∗ [:*rr* ∗ *c* ∗ *cq*:] = *q* ∗ [:*a* ∗ (*rr* ∗ *c* ∗ *cq*):] **by** (*simp add*: *ac-simps*)
also have ... = *q* **unfolding** *one*[*symmetric*] **by** *auto*
finally obtain *r* **where** *q* = *p* ∗ *r* **by** *blast*
hence *p dvd q* **..**
with *ndvd* **show** *?thesis* **by** *auto*
qed
}
ultimately show *?thesis* **by** (*auto simp*:*irreducible-altdef*)
qed

lemma *irreducible-cases*: *irreducible p* ⟷
degree p = *0* ∧ *irreducible* (*coeff p 0*) ∨
degree p ≠ *0* ∧ *irreducible* (*map-poly to-fract p*) ∧ *content-ff-ff p* =*dff 1*
using *irreducible-PM-M-PFM irreducible-M-PM irreducible-PFM-PM*
by *blast*

lemma *dvd-PM-iff*: *p dvd q* ⟷ *divides-ff* (*content-ff-ff p*) (*content-ff-ff q*) ∧
map-poly to-fract p dvd map-poly to-fract q
proof −
interpret *map-poly-inj-idom-hom to-fract***..**
let *?E* = *map-poly to-fract*
show *?thesis* (**is** *?l* = *?r*)
proof
assume *p dvd q*
then obtain *r* **where** *qpr*: *q* = *p* ∗ *r* **..**
from *arg-cong*[*OF this*, *of ?E*]
have *dvd*: *?E p dvd ?E q* **by** *auto*
from *content-ff-map-poly-to-fract* **obtain** *cq* **where** *cq*: *content-ff-ff q* = *to-fract cq* **by** *auto*
from *content-ff-map-poly-to-fract* **obtain** *cp* **where** *cp*: *content-ff-ff p* = *to-fract cp* **by** *auto*
from *content-ff-map-poly-to-fract* **obtain** *cr* **where** *cr*: *content-ff-ff r* = *to-fract*

cr **by** *auto*
from *gauss-lemma*[*of ?E p ?E r*, *folded hom-distribs qpr*, *unfolded cq cp cr*]
have *divides-ff* (*content-ff-ff p*) (*content-ff-ff q*) **unfolding** *cq cp eq-dff-def*
by (*metis divides-ff-def divides-ff-trans*)
with *dvd* **show** *?r* **by** *blast*
next
assume *?r*
show *?l*
proof (*cases q = 0*)
case *True*
with ‹*?r*› **show** *?l* **by** *auto*
next
case *False* **note** *q = this*
hence *q'*: *?E q ≠ 0* **by** *auto*
from ‹*?r*› **obtain** *rr* **where** *qpr*: *?E q = ?E p * rr* **unfolding** *dvd-def* **by** *auto*
with *q* **have** *p*: *p ≠ 0* **and** *Ep*: *?E p ≠ 0* **and** *rr*: *rr ≠ 0* **by** *auto*
from *gauss-lemma*[*of ?E p rr*, *folded qpr*]
have *ct*: *content-ff-ff q =dff content-ff-ff p * content-ff rr*
by *auto*
from *content-ff-map-poly-to-fract*[*of p*] **obtain** *cp* **where** *cp*: *content-ff-ff p = to-fract cp* **by** *auto*
from *content-ff-map-poly-to-fract*[*of q*] **obtain** *cq* **where** *cq*: *content-ff-ff q = to-fract cq* **by** *auto*
from ‹*?r*›[*unfolded cp cq*] **have** *divides-ff* (*to-fract cp*) (*to-fract cq*) **..**
with *ct*[*unfolded cp cq eq-dff-def*] **have** *content-ff rr ∈ range to-fract*
by (*metis* (*no-types*, *lifting*) *Ep content-ff-0-iff cp divides-ff-def*
divides-ff-trans mult.commute mult-right-cancel range-eqI)
from *range-to-fract-embed-poly*[*OF content-ff-to-fract-coeffs-to-fract*[*OF this*]]
obtain *r*
where *rr*: *rr = ?E r* **by** *auto*
from *qpr*[*unfolded rr*, *folded hom-distribs*]
have *q = p * r* **by** (*rule injectivity*)
thus *p dvd q* **..**
qed
qed
qed

lemma *factorial-monoid-poly*: *factorial-monoid* (*mk-monoid* :: *'a* :: *ufd poly monoid*)
proof (*fold factorial-condition-one*, *intro conjI*)
interpret *M*: *factorial-monoid mk-monoid* :: *'a monoid* **by** (*fact factorial-monoid*)
interpret *PFM*: *factorial-monoid mk-monoid* :: *'a fract poly monoid*
by (*rule as-ufd.factorial-monoid*)
interpret *PM*: *comm-monoid-cancel mk-monoid* :: *'a poly monoid* **by** (*unfold-locales*, *auto*)
let *?E = map-poly to-fract*
show *divisor-chain-condition-monoid* (*mk-monoid*::*'a poly monoid*)
proof (*unfold-locales*, *unfold mk-monoid-simps*)
let *?rel'* = {(*x*::*'a poly*, *y*). *x ≠ 0 ∧ y ≠ 0 ∧ properfactor x y*}

let *?rel''* = {(*x*::'*a*, *y*). *x* ≠ *0* ∧ *y* ≠ *0* ∧ *properfactor x y*}
let *?relPM* = {(*x*, *y*). *x* ≠ *0* ∧ *y* ≠ *0* ∧ *x dvd y* ∧ ¬ *y dvd* (*x* :: '*a poly*)}
let *?relM* = {(*x*, *y*). *x* ≠ *0* ∧ *y* ≠ *0* ∧ *x dvd y* ∧ ¬ *y dvd* (*x* :: '*a*)}
have *id*: *?rel'* = *?relPM* **using** *properfactor-nz* **by** *auto*
have *id'*: *?rel''* = *?relM* **using** *properfactor-nz* **by** *auto*
have *wf ?rel''* **using** *M.division-wellfounded* **by** *auto*
hence *wfM*: *wf ?relM* **using** *id'* **by** *auto*
let *?c* = λ *p*. *inv-embed* (*content-ff-ff p*)
let *?f* = λ *p*. (*degree p*, *?c p*)
note *wf* = *wf-inv-image*[*OF wf-lex-prod*[*OF wf-less wfM*], *of ?f*]
show *wf ?rel'* **unfolding** *id*
proof (*rule wf-subset*[*OF wf*], *clarify*)
fix *p q* :: '*a poly*
assume *p*: *p* ≠ *0* **and** *q*: *q* ≠ *0* **and** *dvd*: *p dvd q* **and** *ndvd*: ¬ *q dvd p*
from *dvd* **obtain** *r* **where** *qpr*: *q* = *p* ∗ *r* **..**
from *degree-mult-eq*[*of p r*, *folded qpr*] *q qpr* **have** *r*: *r* ≠ *0*
and *deg*: *degree q* = *degree p* + *degree r* **by** *auto*
show (*p*,*q*) ∈ *inv-image* ({(*x*, *y*). *x* < *y*} <∗*lex*∗> *?relM*) *?f*
proof (*cases degree p* = *degree q*)
case *False*
with *deg* **have** *degree p* < *degree q* **by** *auto*
thus *?thesis* **by** *auto*
next
case *True*
with *deg* **have** *degree r* = *0* **by** *simp*
from *degree0-coeffs*[*OF this*] *r* **obtain** *a* **where** *ra*: *r* = [:*a*:] **and** *a*: *a* ≠ *0*
by *auto*
from *arg-cong*[*OF qpr*, *of* λ *p*. *?E p* ∗ [:*inverse* (*to-fract a*):]] *a*
have *?E p* = *?E q* ∗ [:*inverse* (*to-fract a*):]
by (*auto simp*: *ac-simps ra*)
hence *?E q dvd ?E p* **..**
with *ndvd dvd-PM-iff* **have** *ndvd*: ¬ *divides-ff* (*content-ff-ff q*) (*content-ff-ff*
p) **by** *auto*
from *content-ff-map-poly-to-fract* **obtain** *cq* **where** *cq*: *content-ff-ff q* =
to-fract cq **by** *auto*
from *content-ff-map-poly-to-fract* **obtain** *cp* **where** *cp*: *content-ff-ff p* =
to-fract cp **by** *auto*
from *ndvd*[*unfolded cp cq*] **have** *ndvd*: ¬ *cq dvd cp* **by** *simp*
from *iffD1*[*OF dvd-PM-iff*,*OF dvd*,*unfolded cq cp*]
have *dvd*: *cp dvd cq* **by** *simp*
have *c-p*: *?c p* = *cp* **unfolding** *cp* **by** *simp*
have *c-q*: *?c q* = *cq* **unfolding** *cq* **by** *simp*
from *q cq* **have** *cq0*: *cq* ≠ *0* **by** *auto*
from *p cp* **have** *cp0*: *cp* ≠ *0* **by** *auto*
from *ndvd cq0 cp0 dvd* **have** (*?c p*, *?c q*) ∈ *?relM* **unfolding** *c-p c-q* **by**
auto
with *True* **show** *?thesis* **by** *auto*
qed
qed

```
  qed
  show primeness-condition-monoid (mk-monoid::'a poly monoid)
  proof (unfold-locales, unfold mk-monoid-simps)
    fix p :: 'a poly
    assume p: p ≠ 0 and irred p
    then have irr: irreducible p by auto
    from p have p': ?E p ≠ 0 by auto
    from irreducible-PM-M-PFM[OF irr] have choice: degree p = 0 ∧ irred (coeff
p 0)
      ∨ degree p ≠ 0 ∧ irred (?E p) ∧ content-ff-ff p =dff 1 by auto
    show Divisibility.prime mk-monoid p
    proof (rule Divisibility.primeI, unfold mk-monoid-simps mem-Units)
      show ¬ p dvd 1
      proof
        assume p dvd 1
        from divides-degree[OF this] have dp: degree p = 0 by auto
        from degree0-coeffs[OF this] p obtain a where p: p = [:a:] and a: a ≠ 0
by auto
        with choice have irr: irreducible a by auto
        from ‹p dvd 1›[unfolded p] have a dvd 1 by auto
        with irr show False unfolding irreducible-def by auto
      qed
      fix q r :: 'a poly
      assume q: q ≠ 0 and r: r ≠ 0 and factor p (q * r)
      from this[unfolded factor-idom] have p dvd q * r by auto
        from iffD1[OF dvd-PM-iff this] have dvd-ct: divides-ff (content-ff-ff p)
(content-ff (?E (q * r)))
        and dvd-E: ?E p dvd ?E q * ?E r by auto
      from gauss-lemma[of ?E q ?E r] divides-ff-trans[OF dvd-ct, of content-ff-ff q
* content-ff-ff r]
      have dvd-ct: divides-ff (content-ff-ff p) (content-ff-ff q * content-ff-ff r)
        unfolding eq-dff-def by auto
      from choice
      have p dvd q ∨ p dvd r
      proof
        assume degree p ≠ 0 ∧ irred (?E p) ∧ content-ff-ff p =dff 1
        hence deg: degree p ≠ 0 and irr: irred (?E p) and ct: content-ff-ff p =dff
1 by auto
          from PFM.irreducible-prime[OF irr] p have prime: Divisibility.prime
mk-monoid (?E p) by auto
        from q r have Eq: ?E q ∈ carrier mk-monoid and Er: ?E r ∈ carrier
mk-monoid
          and q': ?E q ≠ 0 and r': ?E r ≠ 0 and qr': ?E q * ?E r ≠ 0 by auto
        from PFM.prime-divides[OF Eq Er prime] q' r' qr' dvd-E
        have dvd-E: ?E p dvd ?E q ∨ ?E p dvd ?E r by simp
       from ct have ct: divides-ff (content-ff-ff p) 1 unfolding eq-dff-def by auto
      moreover have ⋀ q. divides-ff 1 (content-ff-ff q) using content-ff-map-poly-to-fract
          unfolding divides-ff-def by auto
        from divides-ff-trans[OF ct this] have ct: ⋀ q. divides-ff (content-ff-ff p)
```

```
(content-ff-ff q) .
      with dvd-E show ?thesis using dvd-PM-iff by blast
    next
      assume degree p = 0 ∧ irred (coeff p 0)
      hence deg: degree p = 0 and irr: irred (coeff p 0) by auto
      from degree0-coeffs[OF deg] p obtain a where p: p = [:a:] and a: a ≠ 0
by auto
      with irr have irr: irred a and aM: a ∈ carrier mk-monoid by auto
        from M.irreducible-prime[OF irr aM] have prime: Divisibility.prime
mk-monoid a .
       from content-ff-map-poly-to-fract obtain cq where cq: content-ff-ff q =
to-fract cq by auto
       from content-ff-map-poly-to-fract obtain cp where cp: content-ff-ff p =
to-fract cp by auto
       from content-ff-map-poly-to-fract obtain cr where cr: content-ff-ff r =
to-fract cr by auto
      have divides-ff (to-fract a) (content-ff-ff p) unfolding p content-ff-iff using
a by auto
       from divides-ff-trans[OF this[unfolded cp] dvd-ct[unfolded cp cq cr]]
       have divides-ff (to-fract a) (to-fract (cq * cr)) by simp
        hence dvd: a dvd cq * cr by (auto simp add: divides-ff-def simp del:
to-fract-hom.hom-mult simp: to-fract-hom.hom-mult[symmetric])
       from content-ff-divides-ff[of to-fract a ?E p] have divides-ff (to-fract cp)
(to-fract a)
        using cp a p by auto
       hence cpa: cp dvd a by simp
       from a q r cq cr have aM: a ∈ carrier mk-monoid and qM: cq ∈ carrier
mk-monoid and rM: cr ∈ carrier mk-monoid
        and q': cq ≠ 0 and r': cr ≠ 0 and qr': cq * cr ≠ 0
        by auto
       from M.prime-divides[OF qM rM prime] q' r' qr' dvd
       have a dvd cq ∨ a dvd cr by simp
       with dvd-trans[OF cpa] have dvd: cp dvd cq ∨ cp dvd cr by auto
       have ⋀ q. ?E p * (smult (inverse (to-fract a)) q) = q unfolding p using
a by (auto simp: one-poly-def)
       hence Edvd: ⋀ q. ?E p dvd q unfolding dvd-def by metis
       from dvd Edvd show ?thesis apply (subst(1 2) dvd-PM-iff) unfolding cp
cq cr by auto
     qed
     thus factor p q ∨ factor p r unfolding factor-idom using p q r by auto
   qed
 qed
qed

instance poly :: (ufd) ufd
 by (intro ufd.intro-of-class factorial-monoid-imp-ufd factorial-monoid-poly)


lemma primitive-iff-some-content-dvd-1:
```

fixes $f :: 'a :: ufd\ poly$
shows $primitive\ f \longleftrightarrow some\text{-}gcd.listgcd\ (coeffs\ f)\ dvd\ 1$ (**is** - $\longleftrightarrow$ $?c\ dvd\ 1$)
proof(*intro iffI primitiveI*)
fix x
assume $(\bigwedge y.\ y \in set\ (coeffs\ f) \Longrightarrow x\ dvd\ y)$
from *some-gcd.listgcd-greatest*[*of coeffs f*, *OF this*]
have $x\ dvd\ ?c$ **by** *simp*
also assume $?c\ dvd\ 1$
finally show $x\ dvd\ 1$.
next
assume $primitive\ f$
from *primitiveD*[*OF this some-gcd.listgcd*[*of - coeffs f*]]
show $?c\ dvd\ 1$ **by** *auto*
qed

end

5 Polynomials in Rings and Fields

5.1 Polynomials in Rings

We use a locale to work with polynomials in some integer-modulo ring.

theory *Poly-Mod*
imports
HOL−Computational-Algebra.Primes
Polynomial-Factorization.Square-Free-Factorization
Unique-Factorization-Poly
begin

locale *poly-mod* = **fixes** $m :: int$
begin

definition $M :: int \Rightarrow int$ **where** $M\ x = x\ mod\ m$

lemma *M-0*[*simp*]: $M\ 0 = 0$
by (*auto simp add*: *M-def*)

lemma *M-M*[*simp*]: $M\ (M\ x) = M\ x$
by (*auto simp add*: *M-def*)

lemma *M-plus*[*simp*]: $M\ (M\ x + y) = M\ (x + y)\ M\ (x + M\ y) = M\ (x + y)$
by (*auto simp add*: *M-def mod-simps*)

lemma *M-minus*[*simp*]: $M\ (M\ x - y) = M\ (x - y)\ M\ (x - M\ y) = M\ (x - y)$
by (*auto simp add*: *M-def mod-simps*)

lemma *M-times*[*simp*]: $M\ (M\ x * y) = M\ (x * y)\ M\ (x * M\ y) = M\ (x * y)$
by (*auto simp add*: *M-def mod-simps*)

lemma *M-sum*: *M* (*sum* (λ *x*. *M* (*f x*)) *A*) = *M* (*sum f A*)
proof (*induct A rule*: *infinite-finite-induct*)
case (*insert x A*)
from *insert*(*1*−*2*) **have** *M* (∑ *x*∈*insert x A*. *M* (*f x*)) = *M* (*f x* + *M* ((∑ *x*∈*A*. *M* (*f x*)))) **by** *simp*
also have *M* ((∑ *x*∈*A*. *M* (*f x*))) = *M* ((∑ *x*∈*A*. *f x*)) **using** *insert* **by** *simp*
finally show *?case* **using** *insert* **by** *simp*
qed *auto*

definition *inv-M* :: *int* ⇒ *int* **where**
inv-M = (λ *x*. *if x* + *x* ≤ *m then x else x* − *m*)

lemma *M-inv-M-id*[*simp*]: *M* (*inv-M x*) = *M x*
unfolding *inv-M-def M-def* **by** *simp*

definition *Mp* :: *int poly* ⇒ *int poly* **where** *Mp* = *map-poly M*

lemma *Mp-0*[*simp*]: *Mp 0* = *0* **unfolding** *Mp-def* **by** *auto*

lemma *Mp-coeff*: *coeff* (*Mp f*) *i* = *M* (*coeff f i*) **unfolding** *Mp-def*
by (*simp add*: *M-def coeff-map-poly*)

abbreviation *eq-m* :: *int poly* ⇒ *int poly* ⇒ *bool* (**infixl** ‹=*m*› *50*) **where**
f =*m g* ≡ (*Mp f* = *Mp g*)

notation *eq-m* (**infixl** ‹=*m*› *50*)

abbreviation *degree-m* :: *int poly* ⇒ *nat* **where**
degree-m f ≡ *degree* (*Mp f*)

lemma *mult-Mp*[*simp*]: *Mp* (*Mp f* ∗ *g*) = *Mp* (*f* ∗ *g*) *Mp* (*f* ∗ *Mp g*) = *Mp* (*f* ∗ *g*)
proof −
{
fix *f g*
have *Mp* (*Mp f* ∗ *g*) = *Mp* (*f* ∗ *g*)
unfolding *poly-eq-iff Mp-coeff* **unfolding** *coeff-mult Mp-coeff*
proof
fix *n*
show *M* (∑ *i*≤*n*. *M* (*coeff f i*) ∗ *coeff g* (*n* − *i*)) = *M* (∑ *i*≤*n*. *coeff f i* ∗ *coeff g* (*n* − *i*))
by (*subst M-sum*[*symmetric*], *rule sym*, *subst M-sum*[*symmetric*], *unfold M-times*, *simp*)
qed
}
from *this*[*of f g*] *this*[*of g f*] **show** *Mp* (*Mp f* ∗ *g*) = *Mp* (*f* ∗ *g*) *Mp* (*f* ∗ *Mp g*) = *Mp* (*f* ∗ *g*)
by (*auto simp*: *ac-simps*)

qed

lemma *plus-Mp*[*simp*]: *Mp* (*Mp f* + *g*) = *Mp* (*f* + *g*) *Mp* (*f* + *Mp g*) = *Mp* (*f* + *g*)
unfolding *poly-eq-iff Mp-coeff* **unfolding** *coeff-mult Mp-coeff* **by** (*auto simp add*: *Mp-coeff*)

lemma *minus-Mp*[*simp*]: *Mp* (*Mp f* − *g*) = *Mp* (*f* − *g*) *Mp* (*f* − *Mp g*) = *Mp* (*f* − *g*)
unfolding *poly-eq-iff Mp-coeff* **unfolding** *coeff-mult Mp-coeff* **by** (*auto simp add*: *Mp-coeff*)

lemma *Mp-smult*[*simp*]: *Mp* (*smult* (*M a*) *f*) = *Mp* (*smult a f*) *Mp* (*smult a* (*Mp f*)) = *Mp* (*smult a f*)
unfolding *Mp-def smult-as-map-poly*
by (*rule poly-eqI*, *auto simp*: *coeff-map-poly*)+

lemma *Mp-Mp*[*simp*]: *Mp* (*Mp f*) = *Mp f* **unfolding** *Mp-def*
by (*intro poly-eqI*, *auto simp*: *coeff-map-poly*)

lemma *Mp-smult-m-0*[*simp*]: *Mp* (*smult m f*) = *0*
by (*intro poly-eqI*, *auto simp*: *Mp-coeff*, *auto simp*: *M-def*)

definition *dvdm* :: *int poly* ⇒ *int poly* ⇒ *bool* (**infix** ‹*dvdm*› *50*) **where**
f dvdm g = (∃ *h*. *g* =*m f* ∗ *h*)
notation *dvdm* (**infix** ‹*dvdm*› *50*)

lemma *dvdmE*:
assumes *fg*: *f dvdm g*
and *main*: ⋀*h*. *g* =*m f* ∗ *h* ⟹ *Mp h* = *h* ⟹ *thesis*
shows *thesis*
proof−
from *fg* **obtain** *h* **where** *g* =*m f* ∗ *h* **by** (*auto simp*: *dvdm-def*)
then have *g* =*m f* ∗ *Mp h* **by** *auto*
from *main*[*OF this*] **show** *thesis* **by** *auto*
qed

lemma *Mp-dvdm*[*simp*]: *Mp f dvdm g* ⟷ *f dvdm g*
and *dvdm-Mp*[*simp*]: *f dvdm Mp g* ⟷ *f dvdm g* **by** (*auto simp*: *dvdm-def*)

definition *irreducible-m*
where *irreducible-m f* = (¬*f* =*m 0* ∧ ¬ *f dvdm 1* ∧ (∀ *a b*. *f* =*m a* ∗ *b* ⟶ *a dvdm 1* ∨ *b dvdm 1*))

definition *irreducible*$_d$*-m* :: *int poly* ⇒ *bool* **where** *irreducible*$_d$*-m f* ≡
degree-m f > *0* ∧
(∀ *g h*. *degree-m g* < *degree-m f* ⟶ *degree-m h* < *degree-m f* ⟶ ¬ *f* =*m g* ∗ *h*)

definition *prime-elem-m*
where *prime-elem-m f* ≡ ¬ *f* =*m* *0* ∧ ¬ *f dvdm 1* ∧ (∀ *g h*. *f dvdm g* ∗ *h* ⟶ *f dvdm g* ∨ *f dvdm h*)

lemma *degree-m-le-degree* [*intro*!]: *degree-m f* ≤ *degree f*
by (*simp add*: *Mp-def degree-map-poly-le*)

lemma *irreducible*$_d$*-mI*:
assumes *f0*: *degree-m f* > *0*
and *main*: ⋀*g h*. *Mp g* = *g* ⟹ *Mp h* = *h* ⟹ *degree g* > *0* ⟹ *degree g* < *degree-m f* ⟹ *degree h* > *0* ⟹ *degree h* < *degree-m f* ⟹ *f* =*m* *g* ∗ *h* ⟹ *False*
shows *irreducible*$_d$*-m f*
proof (*unfold irreducible*$_d$*-m-def*, *intro conjI allI impI f0 notI*)
fix *g h*
assume *deg*: *degree-m g* < *degree-m f degree-m h* < *degree-m f* **and** *f* =*m* *g* ∗ *h*
then have *f*: *f* =*m* *Mp g* ∗ *Mp h* **by** *simp*
have *degree-m f* ≤ *degree-m g* + *degree-m h*
unfolding *f* **using** *degree-mult-le order.trans* **by** *blast*
with *main*[*of Mp g Mp h*] *deg f* **show** *False* **by** *auto*
qed

lemma *irreducible*$_d$*-mE*:
assumes *irreducible*$_d$*-m f*
and *degree-m f* > *0* ⟹ (⋀*g h*. *degree-m g* < *degree-m f* ⟹ *degree-m h* < *degree-m f* ⟹ ¬ *f* =*m* *g* ∗ *h*) ⟹ *thesis*
shows *thesis*
using *assms* **by** (*unfold irreducible*$_d$*-m-def*, *auto*)

lemma *irreducible*$_d$*-mD*:
assumes *irreducible*$_d$*-m f*
shows *degree-m f* > *0* **and** ⋀*g h*. *degree-m g* < *degree-m f* ⟹ *degree-m h* < *degree-m f* ⟹ ¬ *f* =*m* *g* ∗ *h*
using *assms* **by** (*auto elim*: *irreducible*$_d$*-mE*)

definition *square-free-m* :: *int poly* ⇒ *bool* **where**
square-free-m f = (¬ *f* =*m* *0* ∧ (∀ *g*. *degree-m g* ≠ *0* ⟶ ¬ (*g* ∗ *g dvdm f*)))

definition *coprime-m* :: *int poly* ⇒ *int poly* ⇒ *bool* **where**
coprime-m f g = (∀ *h*. *h dvdm f* ⟶ *h dvdm g* ⟶ *h dvdm 1*)

lemma *Mp-square-free-m*[*simp*]: *square-free-m* (*Mp f*) = *square-free-m f*
unfolding *square-free-m-def dvdm-def* **by** *simp*

lemma *square-free-m-cong*: *square-free-m f* ⟹ *Mp f* = *Mp g* ⟹ *square-free-m g*
unfolding *square-free-m-def dvdm-def* **by** *simp*

lemma *Mp-prod-mset*[*simp*]: *Mp* (*prod-mset* (*image-mset Mp b*)) = *Mp* (*prod-mset*

b)
proof (*induct b*)
case (*add x b*)
have *Mp* (*prod-mset* (*image-mset Mp* ({#*x*#}+*b*))) = *Mp* (*Mp x* ∗ *prod-mset* (*image-mset Mp b*)) **by** *simp*
also have ... = *Mp* (*Mp x* ∗ *Mp* (*prod-mset* (*image-mset Mp b*))) **by** *simp*
also have ... = *Mp* (*Mp x* ∗ *Mp* (*prod-mset b*)) **unfolding** *add* **by** *simp*
finally show *?case* **by** *simp*
qed *simp*

lemma *Mp-prod-list*: *Mp* (*prod-list* (*map Mp b*)) = *Mp* (*prod-list b*)
proof (*induct b*)
case (*Cons b xs*)
have *Mp* (*prod-list* (*map Mp* (*b* # *xs*))) = *Mp* (*Mp b* ∗ *prod-list* (*map Mp xs*)) **by** *simp*
also have ... = *Mp* (*Mp b* ∗ *Mp* (*prod-list* (*map Mp xs*))) **by** *simp*
also have ... = *Mp* (*Mp b* ∗ *Mp* (*prod-list xs*)) **unfolding** *Cons* **by** *simp*
finally show *?case* **by** *simp*
qed *simp*

Polynomial evaluation modulo

definition *M-poly p x* ≡ *M* (*poly p x*)

lemma *M-poly-Mp*[*simp*]: *M-poly* (*Mp p*) = *M-poly p*
proof(*intro ext*, *induct p*)
case *0* **show** *?case* **by** *auto*
next
case *IH*: (*pCons a p*)
from *IH*(*1*) **have** *M-poly* (*Mp* (*pCons a p*)) *x* = *M* (*a* + *M*(*x* ∗ *M-poly* (*Mp p*) *x*))
by (*simp add*: *M-poly-def Mp-def*)
also note *IH*(*2*)[*of x*]
finally show *?case* **by** (*simp add*: *M-poly-def*)
qed

lemma *Mp-lift-modulus*: **assumes** *f* =*m g*
shows *poly-mod.eq-m* (*m* ∗ *k*) (*smult k f*) (*smult k g*)
using *assms* **unfolding** *poly-eq-iff poly-mod.Mp-coeff coeff-smult*
unfolding *poly-mod.M-def* **by** *simp*

lemma *Mp-ident-product*: *n* > *0* ⟹ *Mp f* = *f* ⟹ *poly-mod.Mp* (*m* ∗ *n*) *f* = *f*
unfolding *poly-eq-iff poly-mod.Mp-coeff poly-mod.M-def*
by (*auto simp add*: *zmod-zmult2-eq*) (*metis mod-div-trivial mod-0*)

lemma *Mp-shrink-modulus*: **assumes** *poly-mod.eq-m* (*m* ∗ *k*) *f g k* ≠ *0*
shows *f* =*m g*
proof −
from *assms* **have** *a*: ⋀ *n*. *coeff f n mod* (*m* ∗ *k*) = *coeff g n mod* (*m* ∗ *k*)
unfolding *poly-eq-iff poly-mod.Mp-coeff* **unfolding** *poly-mod.M-def* **by** *auto*

show *?thesis* **unfolding** *poly-eq-iff poly-mod.Mp-coeff* **unfolding** *poly-mod.M-def*
proof
fix *n*
show *coeff f n mod m = coeff g n mod m* **using** *a*[*of n*] ‹*k* ≠ *0*›
by (*metis mod-mult-right-eq mult.commute mult-cancel-left mult-mod-right*)
qed
qed

lemma *degree-m-le*: *degree-m f* ≤ *degree f* **unfolding** *Mp-def* **by** (*rule degree-map-poly-le*)

lemma *degree-m-eq*: *coeff f* (*degree f*) *mod m* ≠ *0* ⟹ *m* > *1* ⟹ *degree-m f* = *degree f*
using *degree-m-le*[*of f*] **unfolding** *Mp-def*
by (*auto intro*: *degree-map-poly simp*: *Mp-def poly-mod.M-def*)

lemma *degree-m-mult-le*:
assumes *eq*: *f* =*m g* ∗ *h*
shows *degree-m f* ≤ *degree-m g* + *degree-m h*
proof −
have *degree-m f* = *degree-m* (*Mp g* ∗ *Mp h*) **using** *eq* **by** *simp*
also have ... ≤ *degree* (*Mp g* ∗ *Mp h*) **by** (*rule degree-m-le*)
also have ... ≤ *degree-m g* + *degree-m h* **by** (*rule degree-mult-le*)
finally show *?thesis* **by** *auto*
qed

lemma *degree-m-smult-le*: *degree-m* (*smult c f*) ≤ *degree-m f*
by (*metis Mp-0 coeff-0 degree-le degree-m-le degree-smult-eq poly-mod.Mp-smult*(*2*) *smult-eq-0-iff*)

lemma *irreducible-m-Mp*[*simp*]: *irreducible-m* (*Mp f*) ⟷ *irreducible-m f* **by** (*simp add*: *irreducible-m-def*)

lemma *eq-m-irreducible-m*: *f* =*m g* ⟹ *irreducible-m f* ⟷ *irreducible-m g*
using *irreducible-m-Mp* **by** *metis*

definition *mset-factors-m* **where** *mset-factors-m F p* ≡
F ≠ {#} ∧ (∀*f*. *f* ∈# *F* ⟶ *irreducible-m f*) ∧ *p* =*m prod-mset F*

end

declare *poly-mod.M-def*[*code*]
declare *poly-mod.Mp-def*[*code*]
declare *poly-mod.inv-M-def*[*code*]

definition *Irr-Mon* :: *'a* :: *comm-semiring-1 poly set*
where *Irr-Mon* = {*x*. *irreducible x* ∧ *monic x*}

definition *factorization* :: *'a* :: *comm-semiring-1 poly set* ⇒ *'a poly* ⇒ (*'a* × *'a*

poly multiset) ⇒ *bool* **where**
factorization Factors f cfs ≡ (*case cfs of* (*c*,*fs*) ⇒ *f* = (*smult c* (*prod-mset fs*)) ∧ (*set-mset fs* ⊆ *Factors*))

definition *unique-factorization* :: *'a* :: *comm-semiring-1 poly set* ⇒ *'a poly* ⇒ (*'a* × *'a poly multiset*) ⇒ *bool* **where**
unique-factorization Factors f cfs = (*Collect* (*factorization Factors f*) = {*cfs*})

lemma *irreducible-multD*:
assumes *l*: *irreducible* (*a*∗*b*)
shows *a dvd 1* ∧ *irreducible b* ∨ *b dvd 1* ∧ *irreducible a*
proof−
from *l* **have** *a dvd 1* ∨ *b dvd 1* **by** *auto*
then show *?thesis*
proof(*elim disjE*)
assume *a*: *a dvd 1*
with *l* **have** *irreducible b*
unfolding *irreducible-def*
by (*meson is-unit-mult-iff mult.left-commute mult-not-zero*)
with *a* **show** *?thesis* **by** *auto*
next
assume *a*: *b dvd 1*
with *l* **have** *irreducible a*
unfolding *irreducible-def*
by (*meson is-unit-mult-iff mult-not-zero semiring-normalization-rules*(*16*))
with *a* **show** *?thesis* **by** *auto*
qed
qed

lemma *irreducible-dvd-prod-mset*:
fixes *p* :: *'a* :: *field poly*
assumes *irr*: *irreducible p* **and** *dvd*: *p dvd prod-mset as*
shows ∃ *a* ∈# *as*. *p dvd a*
proof −
from *irr*[*unfolded irreducible-def*] **have** *deg*: *degree p* ≠ *0* **by** *auto*
hence *p1*: ¬ *p dvd 1* **unfolding** *dvd-def*
by (*metis degree-1 nonzero-mult-div-cancel-left div-poly-less linorder-neqE-nat mult-not-zero not-less0 zero-neq-one*)
from *dvd* **show** *?thesis*
proof (*induct as*)
case (*add a as*)
hence *prod-mset* (*add-mset a as*) = *a* ∗ *prod-mset as* **by** *auto*
from *add*(*2*)[*unfolded this*] *add*(*1*) *irr*
show *?case* **by** *auto*
qed (*insert p1*, *auto*)
qed

lemma *monic-factorization-unique-mset*:
fixes *P*::*'a*::*field poly multiset*

```
  assumes eq: prod-mset P = prod-mset Q
    and P: set-mset P ⊆ {q. irreducible q ∧ monic q}
    and Q: set-mset Q ⊆ {q. irreducible q ∧ monic q}
  shows P = Q
proof −
  {
    fix P Q :: 'a poly multiset
    assume id: prod-mset P = prod-mset Q
    and P: set-mset P ⊆ {q. irreducible q ∧ monic q}
    and Q: set-mset Q ⊆ {q. irreducible q ∧ monic q}
    hence P ⊆# Q
    proof (induct P arbitrary: Q)
      case (add x P Q')
      from add(3) have irr: irreducible x and mon: monic x by auto
      have ∃ a ∈# Q'. x dvd a
      proof (rule irreducible-dvd-prod-mset[OF irr])
        show x dvd prod-mset Q' unfolding add(2)[symmetric] by simp
      qed
      then obtain y Q where Q': Q' = add-mset y Q and xy: x dvd y by (meson
mset-add)
      from add(4) Q' have irr': irreducible y and mon': monic y by auto
      have x = y  using irr irr' xy mon mon'
        by (metis irreducibleD' irreducible-not-unit poly-dvd-antisym)
      hence Q': Q' = Q + {#x#} using Q' by auto
      from mon have x0: x ≠ 0 by auto
      from arg-cong[OF add(2)[unfolded Q'], of λ z. z div x]
      have eq: prod-mset P = prod-mset Q using x0 by auto
      from add(3−4)[unfolded Q']
      have set-mset P ⊆ {q. irreducible q ∧ monic q} set-mset Q ⊆ {q. irreducible
q ∧ monic q}
        by auto
      from add(1)[OF eq this] show ?case unfolding Q' by auto
    qed auto
  }
  from this[OF eq P Q] this[OF eq[symmetric] Q P]
  show ?thesis by auto
qed


lemma exactly-one-monic-factorization:
  assumes mon: monic (f :: 'a :: field poly)
  shows ∃! fs. f = prod-mset fs ∧ set-mset fs ⊆ {q. irreducible q ∧ monic q}
proof −
  from monic-irreducible-factorization[OF mon]
  obtain gs g where fin: finite gs and f: f = (∏ a∈gs. a ^ Suc (g a))
    and gs: gs ⊆ {q. irreducible q ∧ monic q}
    by blast
  from fin
  have ∃ fs. set-mset fs ⊆ gs ∧ prod-mset fs = (∏ a∈gs. a ^ Suc (g a))
```

```
  proof (induct gs)
    case (insert a gs)
    from insert(3) obtain fs where *: set-mset fs ⊆ gs prod-mset fs = (∏ a∈gs.
a ^ Suc (g a)) by auto
    let ?fs = fs + replicate-mset (Suc (g a)) a
    show ?case
    proof (rule exI[of - fs + replicate-mset (Suc (g a)) a], intro conjI)
      show set-mset ?fs ⊆ insert a gs using *(1) by auto
      show prod-mset ?fs = (∏ a∈insert a gs. a ^ Suc (g a))
        by (subst prod.insert[OF insert(1–2)], auto simp: *(2))
    qed
  qed simp
  then obtain fs where set-mset fs ⊆ gs prod-mset fs = (∏ a∈gs. a ^ Suc (g a))
by auto
  with gs f have ex: ∃fs. f = prod-mset fs ∧ set-mset fs ⊆ {q. irreducible q ∧
monic q}
    by (intro exI[of - fs], auto)
  thus ?thesis using monic-factorization-unique-mset by blast
qed

lemma monic-prod-mset:
  fixes as :: 'a :: idom poly multiset
  assumes ⋀ a. a ∈ set-mset as ⟹ monic a
  shows monic (prod-mset as) using assms
  by (induct as, auto intro: monic-mult)

lemma exactly-one-factorization:
  assumes f: f ≠ (0 :: 'a :: field poly)
  shows ∃! cfs. factorization Irr-Mon f cfs
proof −
  let ?a = coeff f (degree f)
  let ?b = inverse ?a
  let ?g = smult ?b f
  define g where g = ?g
  from f have a: ?a ≠ 0 ?b ≠ 0 by (auto simp: field-simps)
  hence monic g unfolding g-def by simp
  note ex1 = exactly-one-monic-factorization[OF this, folded Irr-Mon-def]
  then obtain fs where g: g = prod-mset fs set-mset fs ⊆ Irr-Mon by auto
  let ?cfs = (?a,fs)
 have cfs: factorization Irr-Mon f ?cfs unfolding factorization-def split g(1)[symmetric]
    using g(2) unfolding g-def by (simp add: a field-simps)
  show ?thesis
  proof (rule, rule cfs)
    fix dgs
    assume fact: factorization Irr-Mon f dgs
    obtain d gs where dgs: dgs = (d,gs) by force
    from fact[unfolded factorization-def dgs split]
    have fd: f = smult d (prod-mset gs) and gs: set-mset gs ⊆ Irr-Mon by auto
   have monic (prod-mset gs) by (rule monic-prod-mset, insert gs[unfolded Irr-Mon-def],
```

auto)
hence *d*: *d* = *?a* **unfolding** *fd* **by** *auto*
from *arg-cong*[*OF fd*, *of* λ *x*. *smult ?b x*, *unfolded d g-def*[*symmetric*]]
have *g* = *prod-mset gs* **using** *a* **by** (*simp add*: *field-simps*)
with *ex1 g gs* **have** *gs* = *fs* **by** *auto*
thus *dgs* = *?cfs* **unfolding** *dgs d* **by** *auto*
qed
qed

lemma *mod-ident-iff*:
‹(*x* :: *int*) *mod m* = *x* ⟷ *x* ∈ {*0* ..< *m*}›
if ‹*m* > *0*›
proof −
from *that pos-mod-bound* [*of m x*] *pos-mod-sign* [*of m x*] **have** ‹*0* ≤ *x mod m*›
‹*x mod m* < *m*›
by *simp-all*
with *that* **show** *?thesis* **by** *auto*
qed

declare *prod-mset-prod-list*[*simp*]

lemma *mult-1-is-id*[*simp*]: (∗) (*1* :: *'a* :: *ring-1*) = *id* **by** *auto*

context *poly-mod*
begin

lemma *degree-m-eq-monic*: *monic f* ⟹ *m* > *1* ⟹ *degree-m f* = *degree f*
by (*rule degree-m-eq*) *auto*

lemma *monic-degree-m-lift*: **assumes** *monic f k* > *1 m* > *1*
shows *monic* (*poly-mod.Mp* (*m* ∗ *k*) *f*)
proof −
have *deg*: *degree* (*poly-mod.Mp* (*m* ∗ *k*) *f*) = *degree f*
by (*rule poly-mod.degree-m-eq-monic*[*of f m* ∗ *k*], *insert assms*, *auto simp*: *less-1-mult*)
show *?thesis* **unfolding** *poly-mod.Mp-coeff deg assms poly-mod.M-def* **using** *assms*(*2*−)
by (*simp add*: *less-1-mult*)
qed

end

locale *poly-mod-2* = *poly-mod m* **for** *m* +
assumes *m1*: *m* > *1*
begin

lemma *M-1*[*simp*]: *M 1* = *1* **unfolding** *M-def* **using** *m1*
by *auto*

lemma *Mp-1*[*simp*]: *Mp 1 = 1* **unfolding** *Mp-def* **by** *simp*

lemma *monic-degree-m*[*simp*]: *monic f* $\Longrightarrow$ *degree-m f = degree f*
using *degree-m-eq-monic*[*of f*] **using** *m1* **by** *auto*

lemma *monic-Mp*: *monic f* $\Longrightarrow$ *monic* (*Mp f*)
by (*auto simp*: *Mp-coeff*)

lemma *Mp-0-smult-sdiv-poly*: **assumes** *Mp f = 0*
shows *smult m* (*sdiv-poly f m*) = *f*
proof (*intro poly-eqI*, *unfold Mp-coeff coeff-smult sdiv-poly-def*, *subst coeff-map-poly*, *force*)
fix *n*
from *assms* **have** *coeff* (*Mp f*) *n = 0* **by** *simp*
hence *0*: *coeff f n mod m = 0* **unfolding** *Mp-coeff M-def* **.**
thus *m* $*$ (*coeff f n div m*) = *coeff f n* **by** *auto*
qed

lemma *Mp-product-modulus*: *m′ = m* $*$ *k* $\Longrightarrow$ *k > 0* $\Longrightarrow$ *Mp* (*poly-mod.Mp m′ f*)
= Mp f
by (*intro poly-eqI*, *unfold poly-mod.Mp-coeff poly-mod.M-def*, *auto simp*: *mod-mod-cancel*)

lemma *inv-M-rev*: **assumes** *bnd*: *2* $*$ *abs c < m*
shows *inv-M* (*M c*) = *c*
proof (*cases c* $\geq$ *0*)
case *True*
with *bnd* **show** *?thesis* **unfolding** *M-def inv-M-def* **by** *auto*
next
case *False*
have *2*: $\bigwedge$ *v* :: *int*. *2* $*$ *v = v + v* **by** *auto*
from *False* **have** *c*: *c < 0* **by** *auto*
from *bnd c* **have** *c + m > 0 c + m < m* **by** *auto*
with *c* **have** *cm*: *c mod m = c + m*
by (*metis le-less mod-add-self2 mod-pos-pos-trivial*)
from *c bnd* **have** *2* $*$ (*c mod m*) > *m* **unfolding** *cm* **by** *auto*
with *bnd c* **show** *?thesis* **unfolding** *M-def inv-M-def cm* **by** *auto*
qed

end

lemma (**in** *poly-mod*) *degree-m-eq-prime*:
assumes *f0*: *Mp f* $\neq$ *0*
and *deg*: *degree-m f = degree f*
and *eq*: *f* $=_m$ *g* $*$ *h*
and *p*: *prime m*
shows *degree-m f = degree-m g + degree-m h*
proof −

interpret *poly-mod-2 m* **using** *prime-ge-2-int*[*OF p*] **unfolding** *poly-mod-2-def*
by *simp*
from *f0 eq* **have** *Mp* (*Mp g* ∗ *Mp h*) ≠ *0* **by** *auto*
hence *Mp g* ∗ *Mp h* ≠ *0* **using** *Mp-0* **by** (*cases Mp g* ∗ *Mp h*, *auto*)
hence *g0*: *Mp g* ≠ *0* **and** *h0*: *Mp h* ≠ *0* **by** *auto*
have *degree* (*Mp* (*g* ∗ *h*)) = *degree-m* (*Mp g* ∗ *Mp h*) **by** *simp*
also have ... = *degree* (*Mp g* ∗ *Mp h*)
proof (*rule degree-m-eq*[*OF - m1*], *rule*)
have *id*: ⋀ *g*. *coeff* (*Mp g*) (*degree* (*Mp g*)) *mod m* = *coeff* (*Mp g*) (*degree* (*Mp g*))
unfolding *M-def*[*symmetric*] *Mp-coeff* **by** *simp*
from *p* **have** *p'*: *prime m* **unfolding** *prime-int-nat-transfer* **unfolding** *prime-nat-iff*
by *auto*
assume *coeff* (*Mp g* ∗ *Mp h*) (*degree* (*Mp g* ∗ *Mp h*)) *mod m* = *0*
from *this*[*unfolded coeff-degree-mult*]
have *coeff* (*Mp g*) (*degree* (*Mp g*)) *mod m* = *0* ∨ *coeff* (*Mp h*) (*degree* (*Mp h*)) *mod m* = *0*
unfolding *dvd-eq-mod-eq-0*[*symmetric*] **using** *m1 prime-dvd-mult-int*[*OF p'*]
by *auto*
with *g0 h0* **show** *False* **unfolding** *id* **by** *auto*
qed
also have ... = *degree* (*Mp g*) + *degree* (*Mp h*)
by (*rule degree-mult-eq*[*OF g0 h0*])
finally show *?thesis* **using** *eq* **by** *simp*
qed

lemma *monic-smult-add-small*: **assumes** *f* = *0* ∨ *degree f* < *degree g* **and** *mon*: *monic g*
shows *monic* (*g* + *smult q f*)
proof (*cases f* = *0*)
case *True*
thus *?thesis* **using** *mon* **by** *auto*
next
case *False*
with *assms* **have** *degree f* < *degree g* **by** *auto*
hence *degree* (*smult q f*) < *degree g* **by** (*meson degree-smult-le not-less order-trans*)
thus *?thesis* **using** *mon* **using** *coeff-eq-0 degree-add-eq-left* **by** *fastforce*
qed

context *poly-mod*
begin

definition *factorization-m* :: *int poly* ⇒ (*int* × *int poly multiset*) ⇒ *bool* **where**
factorization-m f cfs ≡ (*case cfs of* (*c*,*fs*) ⇒ *f* =*m* (*smult c* (*prod-mset fs*)) ∧
(∀ *f* ∈ *set-mset fs*. *irreducible*$_d$-*m f* ∧ *monic* (*Mp f*)))

definition *Mf* :: *int* × *int poly multiset* ⇒ *int* × *int poly multiset* **where**
Mf cfs ≡ *case cfs of* (*c*,*fs*) ⇒ (*M c*, *image-mset Mp fs*)

lemma *Mf-Mf*[*simp*]: *Mf* (*Mf x*) = *Mf x*
proof (*cases x*, *auto simp*: *Mf-def*, *goal-cases*)
  case (*1 c fs*)
  show *?case* **by** (*induct fs*, *auto*)
qed

definition *equivalent-fact-m* :: *int* × *int poly multiset* ⇒ *int* × *int poly multiset* ⇒ *bool* **where**
  equivalent-fact-m cfs dgs = (*Mf cfs* = *Mf dgs*)

definition *unique-factorization-m* :: *int poly* ⇒ (*int* × *int poly multiset*) ⇒ *bool* **where**
  unique-factorization-m f cfs = (*Mf* ' *Collect* (*factorization-m f*) = {*Mf cfs*})

lemma *Mp-irreducible*$_d$-*m*[*simp*]: *irreducible*$_d$-*m* (*Mp f*) = *irreducible*$_d$-*m f*
  unfolding *irreducible*$_d$-*m-def dvdm-def* **by** *simp*

lemma *Mf-factorization-m*[*simp*]: *factorization-m f* (*Mf cfs*) = *factorization-m f cfs*
  unfolding *factorization-m-def Mf-def*
proof (*cases cfs*, *simp*, *goal-cases*)
  case (*1 c fs*)
  have *Mp* (*smult c* (*prod-mset fs*)) = *Mp* (*smult* (*M c*) (*Mp* (*prod-mset fs*))) **by** *simp*
  also have ... = *Mp* (*smult* (*M c*) (*Mp* (*prod-mset* (*image-mset Mp fs*))))
    unfolding *Mp-prod-mset* **by** *simp*
  also have ... = *Mp* (*smult* (*M c*) (*prod-mset* (*image-mset Mp fs*))) **unfolding** *Mp-smult* **..**
  finally show *?case* **by** *auto*
qed

lemma *unique-factorization-m-imp-factorization*: **assumes** *unique-factorization-m f cfs*
  shows *factorization-m f cfs*
proof −
  from *assms*[*unfolded unique-factorization-m-def*] **obtain** *dfs* **where**
    fact: *factorization-m f dfs* **and** *id*: *Mf cfs* = *Mf dfs* **by** *blast*
  from *fact* **have** *factorization-m f* (*Mf dfs*) **by** *simp*
  from *this*[*folded id*] **show** *?thesis* **by** *simp*
qed

lemma *unique-factorization-m-alt-def*: *unique-factorization-m f cfs* = (*factorization-m f cfs*
  ∧ (∀ *dgs*. *factorization-m f dgs* ⟶ *Mf dgs* = *Mf cfs*))
  using *unique-factorization-m-imp-factorization*[*of f cfs*]
  unfolding *unique-factorization-m-def* **by** *auto*

end

context *poly-mod-2*
begin

lemma *factorization-m-lead-coeff*: **assumes** *factorization-m f* (*c*,*fs*)
shows *lead-coeff* (*Mp f*) = *M c*
proof −
note ∗ = *assms*[*unfolded factorization-m-def split*]
have *monic* (*prod-mset* (*image-mset Mp fs*)) **by** (*rule monic-prod-mset*, *insert* ∗, *auto*)
hence *monic* (*Mp* (*prod-mset* (*image-mset Mp fs*))) **by** (*rule monic-Mp*)
from *this*[*unfolded Mp-prod-mset*] **have** *monic*: *monic* (*Mp* (*prod-mset fs*)) **by** *simp*
from ∗ **have** *lead-coeff* (*Mp f*) = *lead-coeff* (*Mp* (*smult c* (*prod-mset fs*))) **by** *simp*
also have *Mp* (*smult c* (*prod-mset fs*)) = *Mp* (*smult* (*M c*) (*Mp* (*prod-mset fs*))) **by** *simp*
finally show *?thesis*
using *monic* ‹*smult c* (*prod-mset fs*) =*m smult* (*M c*) (*Mp* (*prod-mset fs*))›
by (*metis M-M M-def Mp-0 Mp-coeff lead-coeff-smult m1 mult-cancel-left2 poly-mod.degree-m-eq smult-eq-0-iff*)
qed

lemma *factorization-m-smult*: **assumes** *factorization-m f* (*c*,*fs*)
shows *factorization-m* (*smult d f*) (*c* ∗ *d*,*fs*)
proof −
note ∗ = *assms*[*unfolded factorization-m-def split*]
from ∗ **have** *f*: *Mp f* = *Mp* (*smult c* (*prod-mset fs*)) **by** *simp*
have *Mp* (*smult d f*) = *Mp* (*smult d* (*Mp f*)) **by** *simp*
also have ... = *Mp* (*smult* (*c* ∗ *d*) (*prod-mset fs*)) **unfolding** *f* **by** (*simp add*: *ac-simps*)
finally show *?thesis* **using** *assms*
unfolding *factorization-m-def split* **by** *auto*
qed

lemma *factorization-m-prod*: **assumes** *factorization-m f* (*c*,*fs*) *factorization-m g* (*d*,*gs*)
shows *factorization-m* (*f* ∗ *g*) (*c* ∗ *d*, *fs* + *gs*)
proof −
note ∗ = *assms*[*unfolded factorization-m-def split*]
have *Mp* (*f* ∗ *g*) = *Mp* (*Mp f* ∗ *Mp g*) **by** *simp*
also have *Mp f* = *Mp* (*smult c* (*prod-mset fs*)) **using** ∗ **by** *simp*
also have *Mp g* = *Mp* (*smult d* (*prod-mset gs*)) **using** ∗ **by** *simp*
finally have *Mp* (*f* ∗ *g*) = *Mp* (*smult* (*c* ∗ *d*) (*prod-mset* (*fs* + *gs*))) **unfolding** *mult-Mp*
by (*simp add*: *ac-simps*)
with ∗ **show** *?thesis* **unfolding** *factorization-m-def split* **by** *auto*
qed

lemma *Mp-factorization-m*[*simp*]: *factorization-m* (*Mp f*) *cfs* = *factorization-m f cfs*
unfolding *factorization-m-def* **by** *simp*

lemma *Mp-unique-factorization-m*[*simp*]:
unique-factorization-m (*Mp f*) *cfs* = *unique-factorization-m f cfs*
unfolding *unique-factorization-m-alt-def* **by** *simp*

lemma *unique-factorization-m-cong*: *unique-factorization-m f cfs* ⟹ *Mp f* = *Mp g*
⟹ *unique-factorization-m g cfs*
unfolding *Mp-unique-factorization-m*[*of f*, *symmetric*] **by** *simp*

lemma *unique-factorization-mI*: **assumes** *factorization-m f* (*c*,*fs*)
and ⋀ *d gs*. *factorization-m f* (*d*,*gs*) ⟹ *Mf* (*d*,*gs*) = *Mf* (*c*,*fs*)
shows *unique-factorization-m f* (*c*,*fs*)
unfolding *unique-factorization-m-alt-def*
by (*intro conjI*[*OF assms*(*1*)] *allI impI*, *insert assms*(*2*), *auto*)

lemma *unique-factorization-m-smult*: **assumes** *uf*: *unique-factorization-m f* (*c*,*fs*)
and *d*: *M* (*di* ∗ *d*) = *1*
shows *unique-factorization-m* (*smult d f*) (*c* ∗ *d*,*fs*)
proof (*rule unique-factorization-mI*[*OF factorization-m-smult*])
show *factorization-m f* (*c*, *fs*) **using** *uf*[*unfolded unique-factorization-m-alt-def*]
by *auto*
fix *e gs*
assume *fact*: *factorization-m* (*smult d f*) (*e*,*gs*)
from *factorization-m-smult*[*OF this*, *of di*]
have *factorization-m* (*Mp* (*smult di* (*smult d f*))) (*e* ∗ *di*, *gs*) **by** *simp*
also have *Mp* (*smult di* (*smult d f*)) = *Mp* (*smult* (*M* (*di* ∗ *d*)) *f*) **by** *simp*
also have ... = *Mp f* **unfolding** *d* **by** *simp*
finally have *fact*: *factorization-m f* (*e* ∗ *di*, *gs*) **by** *simp*
with *uf*[*unfolded unique-factorization-m-alt-def*] **have** *eq*: *Mf* (*e* ∗ *di*, *gs*) = *Mf* (*c*, *fs*) **by** *blast*
from *eq*[*unfolded Mf-def*] **have** *M* (*e* ∗ *di*) = *M c* **by** *simp*
from *arg-cong*[*OF this*, *of* λ *x*. *M* (*x* ∗ *d*)]
have *M* (*e* ∗ *M* (*di* ∗ *d*)) = *M* (*c* ∗ *d*) **by** (*simp add*: *ac-simps*)
from *this*[*unfolded d*] **have** *e*: *M e* = *M* (*c* ∗ *d*) **by** *simp*
with *eq*
show *Mf* (*e*,*gs*) = *Mf* (*c* ∗ *d*, *fs*) **unfolding** *Mf-def split* **by** *simp*
qed

lemma *unique-factorization-m-smultD*: **assumes** *uf*: *unique-factorization-m* (*smult d f*) (*c*,*fs*)
and *d*: *M* (*di* ∗ *d*) = *1*
shows *unique-factorization-m f* (*c* ∗ *di*,*fs*)
proof −
from *d* **have** *d'*: *M* (*d* ∗ *di*) = *1* **by** (*simp add*: *ac-simps*)
show *?thesis*

proof (*rule unique-factorization-m-cong*[*OF unique-factorization-m-smult*[*OF uf d'*]],
rule poly-eqI, *unfold Mp-coeff coeff-smult*)
fix *n*
have *M* (*di* ∗ (*d* ∗ *coeff f n*)) = *M* (*M* (*di* ∗ *d*) ∗ *coeff f n*) **by** (*auto simp*: *ac-simps*)
from *this*[*unfolded d*] **show** *M* (*di* ∗ (*d* ∗ *coeff f n*)) = *M* (*coeff f n*) **by** *simp*
qed
qed

lemma *degree-m-eq-lead-coeff*: *degree-m f* = *degree f* ⟹ *lead-coeff* (*Mp f*) = *M* (*lead-coeff f*)
by (*simp add*: *Mp-coeff*)

lemma *unique-factorization-m-zero*: **assumes** *unique-factorization-m f* (*c*,*fs*)
shows *M c* ≠ *0*
proof
assume *c*: *M c* = *0*
from *unique-factorization-m-imp-factorization*[*OF assms*]
have *Mp f* = *Mp* (*smult* (*M c*) (*prod-mset fs*)) **unfolding** *factorization-m-def split*
by *simp*
from *this*[*unfolded c*] **have** *f*: *Mp f* = *0* **by** *simp*
have *factorization-m f* (*0*,{#})
unfolding *factorization-m-def split f* **by** *auto*
moreover have *Mf* (*0*,{#}) = (*0*,{#}) **unfolding** *Mf-def* **by** *auto*
ultimately have *fact1*: (*0*, {#}) ∈ *Mf* ‘ *Collect* (*factorization-m f*) **by** *force*
define *g* :: *int poly* **where** *g* = [:*0*,*1*:]
have *mpg*: *Mp g* = [:*0*,*1*:] **unfolding** *Mp-def*
by (*auto simp*: *g-def*)
{
fix *g h*
assume ∗: *degree* (*Mp g*) = *0 degree* (*Mp h*) = *0* [:*0*, *1*:] = *Mp* (*g* ∗ *h*)
from *arg-cong*[*OF* ∗(*3*), *of degree*] **have** *1* = *degree-m* (*Mp g* ∗ *Mp h*) **by** *simp*
also have ... ≤ *degree* (*Mp g* ∗ *Mp h*) **by** (*rule degree-m-le*)
also have ... ≤ *degree* (*Mp g*) + *degree* (*Mp h*) **by** (*rule degree-mult-le*)
also have ... ≤ *0* **using** ∗ **by** *simp*
finally have *False* **by** *simp*
} **note** *irr* = *this*
have *factorization-m f* (*0*,{# *g* #})
unfolding *factorization-m-def split* **using** *irr*
by (*auto simp*: *irreducible*$_d$*-m-def f mpg*)
moreover have *Mf* (*0*,{# *g* #}) = (*0*,{# *g* #}) **unfolding** *Mf-def* **by** (*auto simp*: *mpg*, *simp add*: *g-def*)
ultimately have *fact2*: (*0*, {#*g*#}) ∈ *Mf* ‘ *Collect* (*factorization-m f*) **by** *force*
note [*simp*] = *assms*[*unfolded unique-factorization-m-def*]
from *fact1*[*simplified*, *folded fact2*[*simplified*]] **show** *False* **by** *auto*
qed

end

context *poly-mod*
begin

lemma *dvdm-smult*: **assumes** *f dvdm g*
shows *f dvdm smult c g*
proof −
from *assms*[*unfolded dvdm-def*] **obtain** *h* **where** *g*: *g* =*m f* ∗ *h* **by** *auto*
show *?thesis* **unfolding** *dvdm-def*
proof (*intro exI*[*of - smult c h*])
have *Mp* (*smult c g*) = *Mp* (*smult c* (*Mp g*)) **by** *simp*
also have *Mp g* = *Mp* (*f* ∗ *h*) **using** *g* **by** *simp*
finally show *Mp* (*smult c g*) = *Mp* (*f* ∗ *smult c h*) **by** *simp*
qed
qed

lemma *dvdm-factor*: **assumes** *f dvdm g*
shows *f dvdm g* ∗ *h*
proof −
from *assms*[*unfolded dvdm-def*] **obtain** *k* **where** *g*: *g* =*m f* ∗ *k* **by** *auto*
show *?thesis* **unfolding** *dvdm-def*
proof (*intro exI*[*of - h* ∗ *k*])
have *Mp* (*g* ∗ *h*) = *Mp* (*Mp g* ∗ *h*) **by** *simp*
also have *Mp g* = *Mp* (*f* ∗ *k*) **using** *g* **by** *simp*
finally show *Mp* (*g* ∗ *h*) = *Mp* (*f* ∗ (*h* ∗ *k*)) **by** (*simp add*: *ac-simps*)
qed
qed

lemma *square-free-m-smultD*: **assumes** *square-free-m* (*smult c f*)
shows *square-free-m f*
unfolding *square-free-m-def*
proof (*intro conjI allI impI*)
fix *g*
assume *degree-m g* ≠ *0*
with *assms*[*unfolded square-free-m-def*] **have** ¬ *g* ∗ *g dvdm smult c f* **by** *auto*
thus ¬ *g* ∗ *g dvdm f* **using** *dvdm-smult*[*of g* ∗ *g f c*] **by** *blast*
next
from *assms*[*unfolded square-free-m-def*] **have** ¬ *smult c f* =*m 0* **by** *simp*
thus ¬ *f* =*m 0*
by (*metis Mp-smult*(*2*) *smult-0-right*)
qed

lemma *square-free-m-smultI*: **assumes** *sf*: *square-free-m f*
and *inv*: *M* (*ci* ∗ *c*) = *1*
shows *square-free-m* (*smult c f*)
proof −
have *square-free-m* (*smult ci* (*smult c f*))

```
  proof (rule square-free-m-cong[OF sf], rule poly-eqI, unfold Mp-coeff coeff-smult)
    fix n
     have M (ci * (c * coeff f n)) = M ( M (ci * c) * coeff f n) by (simp add:
ac-simps)
    from this[unfolded inv] show M (coeff f n) = M (ci * (c * coeff f n)) by simp
  qed
  from square-free-m-smultD[OF this] show ?thesis .
qed


lemma square-free-m-factor: assumes square-free-m (f * g)
  shows square-free-m f square-free-m g
proof −
  {
    fix f g
    assume sf: square-free-m (f * g)
    have square-free-m f
      unfolding square-free-m-def
    proof (intro conjI allI impI)
      fix h
      assume degree-m h ≠ 0
      with sf[unfolded square-free-m-def] have ¬ h * h dvdm f * g by auto
      thus ¬ h * h dvdm f using dvdm-factor[of h * h f g] by blast
    next
      from sf[unfolded square-free-m-def] have ¬ f * g =m 0 by simp
      thus ¬ f =m 0
        by (metis mult.commute mult-zero-right poly-mod.mult-Mp(2))
    qed
  }
  from this[of f g] this[of g f] assms
  show square-free-m f square-free-m g by (auto simp: ac-simps)
qed

end

context poly-mod-2
begin

lemma Mp-ident-iff: Mp f = f ⟷ (∀ n. coeff f n ∈ {0 ..< m})
proof −
  have m0: m > 0 using m1 by simp
 show ?thesis unfolding poly-eq-iff Mp-coeff M-def mod-ident-iff[OF m0] by simp
qed

lemma Mp-ident-iff': Mp f = f ⟷ (set (coeffs f) ⊆ {0 ..< m})
proof −
  have 0: 0 ∈ {0 ..< m} using m1 by auto
  have ran: (∀ n. coeff f n ∈ {0..<m}) ⟷ range (coeff f) ⊆ {0 ..< m} by blast
  show ?thesis unfolding Mp-ident-iff ran using range-coeff[of f] 0 by auto
```

qed
end

lemma *Mp-Mp-pow-is-Mp*: $n \neq 0 \Longrightarrow p > 1 \Longrightarrow$ *poly-mod.Mp p* (*poly-mod.Mp* ($p \hat{} n$) *f*)
= *poly-mod.Mp p f*
using *poly-mod-2.Mp-product-modulus poly-mod-2-def* **by**(*subst power-eq-if*, *auto*)

lemma *M-M-pow-is-M*: $n \neq 0 \Longrightarrow p > 1 \Longrightarrow$ *poly-mod.M p* (*poly-mod.M* ($p \hat{} n$) *f*)
= *poly-mod.M p f* **using** *Mp-Mp-pow-is-Mp*[*of n p* [:*f*:]]
by (*metis coeff-pCons-0 poly-mod.Mp-coeff*)

definition *inverse-mod* :: *int* ⇒ *int* ⇒ *int* **where**
inverse-mod x m = *fst* (*bezout-coefficients x m*)

lemma *inverse-mod*:
(*inverse-mod x m* ∗ *x*) *mod m* = *1*
if *coprime x m m* > *1*
proof −
from *bezout-coefficients* [*of x m inverse-mod x m snd* (*bezout-coefficients x m*)]
have *inverse-mod x m* ∗ *x* + *snd* (*bezout-coefficients x m*) ∗ *m* = *gcd x m*
by (*simp add*: *inverse-mod-def*)
with *that* **have** *inverse-mod x m* ∗ *x* + *snd* (*bezout-coefficients x m*) ∗ *m* = *1*
by *simp*
then have (*inverse-mod x m* ∗ *x* + *snd* (*bezout-coefficients x m*) ∗ *m*) *mod m* = *1 mod m*
by *simp*
with ‹*m* > *1*› **show** *?thesis*
by *simp*
qed

lemma *inverse-mod-pow*:
(*inverse-mod x* (*p* ^ *n*) ∗ *x*) *mod* (*p* ^ *n*) = *1*
if *coprime x p p* > *1 n* ≠ *0*
using *that* **by** (*auto intro*: *inverse-mod*)

lemma (**in** *poly-mod*) *inverse-mod-coprime*:
assumes *p*: *prime m*
and *cop*: *coprime x m* **shows** *M* (*inverse-mod x m* ∗ *x*) = *1*
unfolding *M-def* **using** *inverse-mod-pow*[*OF cop*, *of 1*] *p*
by (*auto simp*: *prime-int-iff*)

lemma (**in** *poly-mod*) *inverse-mod-coprime-exp*:
assumes *m*: *m* = *p*^*n* **and** *p*: *prime p*
and *n*: *n* ≠ *0* **and** *cop*: *coprime x p*
shows *M* (*inverse-mod x m* ∗ *x*) = *1*
unfolding *M-def* **unfolding** *m* **using** *inverse-mod-pow*[*OF cop - n*] *p*
by (*auto simp*: *prime-int-iff*)

```
locale poly-mod-prime = poly-mod p for p :: int +
  assumes prime: prime p
begin

sublocale poly-mod-2 p using prime unfolding poly-mod-2-def
  using prime-gt-1-int by force

lemma square-free-m-prod-imp-coprime-m: assumes sf: square-free-m (A * B)
  shows coprime-m A B
  unfolding coprime-m-def
proof (intro allI impI)
  fix h
  assume dvd: h dvdm A h dvdm B
  then obtain ha hb where *: Mp A = Mp (h * ha) Mp B = Mp (h * hb)
    unfolding dvdm-def by auto
  have AB: Mp (A * B) = Mp (Mp A * Mp B) by simp
  from this[unfolded *, simplified]
  have eq: Mp (A * B) = Mp (h * h * (ha * hb)) by (simp add: ac-simps)
  hence dvd-hh: (h * h) dvdm (A * B) unfolding dvdm-def by auto
  {
    assume degree-m h ≠ 0
    from sf[unfolded square-free-m-def, THEN conjunct2, rule-format, OF this]
    have ¬ h * h dvdm A * B .
    with dvd-hh have False by simp
  }
  hence degree (Mp h) = 0 by auto
  then obtain c where hc: Mp h = [: c :] by (rule degree-eq-zeroE)
  {
    assume c = 0
    hence Mp h = 0 unfolding hc by auto
    with *(1) have Mp A = 0
      by (metis Mp-0 mult-zero-left poly-mod.mult-Mp(1))
    with sf[unfolded square-free-m-def, THEN conjunct1] have False
      by (simp add: AB)
  }
  hence c0: c ≠ 0 by auto
  with arg-cong[OF hc[symmetric], of λ f. coeff f 0, unfolded Mp-coeff M-def] m1
  have c ≥ 0 c < p by auto
  with c0 have c-props:c > 0 c < p by auto
  with prime have prime p by simp
  with c-props have coprime p c
    by (auto intro: prime-imp-coprime dest: zdvd-not-zless)
  then have coprime c p
    by (simp add: ac-simps)
  from inverse-mod-coprime[OF prime this]
  obtain d where d: M (c * d) = 1 by (auto simp: ac-simps)
  show h dvdm 1 unfolding dvdm-def
  proof (intro exI[of - [:d:]])
```

```
    have Mp (h * [: d :]) = Mp (Mp h * [: d :]) by simp
    also have ... = Mp ([: c * d :]) unfolding hc by (auto simp: ac-simps)
    also have ... = [: M (c * d) :] unfolding Mp-def
      by (metis (no-types) M-0 map-poly-pCons Mp-0 Mp-def d zero-neq-one)
    also have ... = 1 unfolding d by simp
    finally show Mp 1 = Mp (h * [:d:]) by simp
  qed
qed

lemma coprime-exp-mod: coprime lu p ⟹ n ≠ 0 ⟹ lu mod p ^ n ≠ 0
  using prime by fastforce

end

context poly-mod
begin

definition Dp :: int poly ⇒ int poly where
  Dp f = map-poly (λ a. a div m) f

lemma Dp-Mp-eq: f = Mp f + smult m (Dp f)
  by (rule poly-eqI, auto simp: Mp-coeff M-def Dp-def coeff-map-poly)

lemma dvd-imp-dvdm:
  assumes a dvd b shows a dvdm b
  by (metis assms dvd-def dvdm-def)

lemma dvdm-add:
  assumes a: u dvdm a
  and b: u dvdm b
  shows u dvdm (a+b)
proof −
  obtain a' where a: a =m u*a' using a unfolding dvdm-def by auto
  obtain b' where b: b =m u*b' using b unfolding dvdm-def by auto
  have Mp (a + b) = Mp (u*a'+u*b') using a b
    by (metis poly-mod.plus-Mp(1) poly-mod.plus-Mp(2))
  also have ... = Mp (u * (a'+ b'))
    by (simp add: distrib-left)
  finally show ?thesis unfolding dvdm-def by auto
qed


lemma monic-dvdm-constant:
  assumes uk: u dvdm [:k:]
  and u1: monic u and u2: degree u > 0
  shows k mod m = 0
proof −
  have d1: degree-m [:k:] = degree [:k:]
    by (metis degree-pCons-0 le-zero-eq poly-mod.degree-m-le)
```

```
  obtain h where h: Mp [:k:] = Mp (u * h)
    using uk unfolding dvdm-def by auto
  have d2: degree-m [:k:] = degree-m (u*h) using h by metis
  have d3: degree (map-poly M (u * map-poly M h)) = degree (u * map-poly M h)

    by (rule degree-map-poly)
       (metis coeff-degree-mult leading-coeff-0-iff mult.right-neutral M-M Mp-coeff
Mp-def u1)
  thus ?thesis using assms d1 d2 d3
    by (auto, metis M-def map-poly-pCons degree-mult-right-le h leD map-poly-0
       mult-poly-0-right pCons-eq-0-iff M-0 Mp-def mult-Mp(2))
qed

lemma div-mod-imp-dvdm:
  assumes ∃ q r. b = q * a + Polynomial.smult m r
  shows a dvdm b
proof −
  from assms  obtain q r where b:b = a * q + smult m r
    by (metis mult.commute)
  have a: Mp (Polynomial.smult m r) = 0 by auto
  show ?thesis
  proof (unfold dvdm-def, rule exI[of - q])
    have Mp (a * q + smult m r) = Mp (a * q + Mp (smult m r))
      using plus-Mp(2)[of a*q smult m r] by auto
    also have ... = Mp (a*q) by auto
    finally show eq-m b (a * q) using b by auto
  qed
qed

lemma lead-coeff-monic-mult:
  fixes p :: 'a :: {comm-semiring-1,semiring-no-zero-divisors} poly
  assumes monic p shows lead-coeff (p * q) = lead-coeff q
  using assms by (simp add: lead-coeff-mult)

lemma degree-m-mult-eq:
  assumes p: monic p and q: lead-coeff q mod m ≠ 0 and m1: m > 1
  shows degree (Mp (p * q)) = degree p + degree q
proof−
  have lead-coeff (p * q) mod m ≠ 0
    using q p by (auto simp: lead-coeff-monic-mult)
  with m1 show ?thesis
    by (auto simp: degree-m-eq intro!: degree-mult-eq)
qed

lemma dvdm-imp-degree-le:
  assumes pq: p dvdm q and p: monic p and q0: Mp q ≠ 0 and m1: m > 1
  shows degree p ≤ degree q
proof−
  from q0
```

have *q*: *lead-coeff* (*Mp q*) *mod m* ≠ *0*
by (*metis Mp-Mp Mp-coeff leading-coeff-neq-0 M-def*)
from *pq* **obtain** *r* **where** *Mpq*: *Mp q* = *Mp* (*p* ∗ *Mp r*) **by** (*auto elim*: *dvdmE*)
with *p q* **have** *lead-coeff* (*Mp r*) *mod m* ≠ *0*
by (*metis Mp-Mp Mp-coeff leading-coeff-0-iff mult-poly-0-right M-def*)
from *degree-m-mult-eq*[*OF p this m1*] *Mpq*
have *degree p* ≤ *degree-m q* **by** *simp*
thus *?thesis* **using** *degree-m-le le-trans* **by** *blast*
qed

lemma *dvdm-uminus* [*simp*]:
p dvdm −*q* ⟷ *p dvdm q*
by (*metis add.inverse-inverse dvdm-smult smult-1-left smult-minus-left*)

lemma *Mp-const-poly*: *Mp* [:*a*:] = [:*a mod m*:]
by (*simp add*: *Mp-def M-def Polynomial.map-poly-pCons*)

lemma *dvdm-imp-div-mod*:
assumes *u dvdm g*
shows ∃ *q r*. *g* = *q*∗*u* + *smult m r*
proof −
obtain *q* **where** *q*: *Mp g* = *Mp* (*u*∗*q*)
using *assms* **unfolding** *dvdm-def* **by** *fast*
have (*u*∗*q*) = *Mp* (*u*∗*q*) + *smult m* (*Dp* (*u*∗*q*))
by (*simp add*: *poly-mod.Dp-Mp-eq*[*of u*∗*q*])
hence *uq*: *Mp* (*u*∗*q*) = (*u*∗*q*) − *smult m* (*Dp* (*u*∗*q*))
by *auto*
have *g*: *g* = *Mp g* + *smult m* (*Dp g*)
by (*simp add*: *poly-mod.Dp-Mp-eq*[*of g*])
also have ... = *poly-mod.Mp m* (*u*∗*q*) + *smult m* (*Dp g*) **using** *q* **by** *simp*
also have ... = *u* ∗ *q* − *smult m* (*Dp* (*u* ∗ *q*)) + *smult m* (*Dp g*)
unfolding *uq* **by** *auto*
also have ... = *u* ∗ *q* + *smult m* (−*Dp* (*u*∗*q*)) + *smult m* (*Dp g*) **by** *auto*
also have ... = *u* ∗ *q* + *smult m* (−*Dp* (*u*∗*q*) + *Dp g*)
unfolding *smult-add-right* **by** *auto*
also have ... = *q* ∗ *u* + *smult m* (−*Dp* (*u*∗*q*) + *Dp g*) **by** *auto*
finally show *?thesis* **by** *auto*
qed

corollary *div-mod-iff-dvdm*:
shows *a dvdm b* = (∃ *q r*. *b* = *q* ∗ *a* + *Polynomial.smult m r*)
using *div-mod-imp-dvdm dvdm-imp-div-mod* **by** *blast*

lemma *dvdmE′*:
assumes *p dvdm q* **and** ⋀*r*. *q* =*m p* ∗ *Mp r* ⟹ *thesis*
shows *thesis*
using *assms* **by** (*auto simp*: *dvdm-def*)

end

context *poly-mod-2*
begin
lemma *factorization-m-mem-dvdm*: **assumes** *fact*: *factorization-m f* (*c*,*fs*)
and *mem*: *Mp g* ∈# *image-mset Mp fs*
shows *g dvdm f*
proof −
from *fact* **have** *factorization-m f* (*Mf* (*c*, *fs*)) **by** *auto*
then obtain *l* **where** *f*: *factorization-m f* (*l*, *image-mset Mp fs*) **by** (*auto simp*: *Mf-def*)
from *multi-member-split*[*OF mem*] **obtain** *ls* **where**
fs: *image-mset Mp fs* = {# *Mp g* #} + *ls* **by** *auto*
from *f*[*unfolded fs split factorization-m-def*] **show** *g dvdm f*
unfolding *dvdm-def*
by (*intro exI*[*of* - *smult l* (*prod-mset ls*)], *auto simp del*: *Mp-smult*
simp add: *Mp-smult*(*2*)[*of* - *Mp g* ∗ *prod-mset ls*, *symmetric*], *simp*)
qed

lemma *dvdm-degree*: *monic u* ⟹ *u dvdm f* ⟹ *Mp f* ≠ *0* ⟹ *degree u* ≤ *degree f*
using *dvdm-imp-degree-le m1* **by** *blast*

end

lemma (**in** *poly-mod-prime*) *pl-dvdm-imp-p-dvdm*:
assumes *l0*: *l* ≠ *0*
and *pl-dvdm*: *poly-mod.dvdm* (*p*^*l*) *a b*
shows *a dvdm b*
proof −
from *l0* **have** *l-gt-0*: *l* > *0* **by** *auto*
with *m1* **interpret** *pl*: *poly-mod-2 p*^*l* **by** (*unfold-locales*, *auto*)
from *l-gt-0* **have** *p-rw*: *p* ∗ *p* ^ (*l* − *1*) = *p* ^ *l*
by (*cases l*) *simp-all*
obtain *q r* **where** *b*: *b* = *q* ∗ *a* + *smult* (*p*^*l*) *r* **using** *pl.dvdm-imp-div-mod*[*OF pl-dvdm*] **by** *auto*
have *smult* (*p*^*l*) *r* = *smult p* (*smult* (*p* ^ (*l* − *1*)) *r*) **unfolding** *smult-smult p-rw* **..**
hence *b2*: *b* = *q* ∗ *a* + *smult p* (*smult* (*p* ^ (*l* − *1*)) *r*) **using** *b* **by** *auto*
show *?thesis*
by (*rule div-mod-imp-dvdm*, *rule exI*[*of* - *q*],
rule exI[*of* - (*smult* (*p* ^ (*l* − *1*)) *r*)], *auto simp add*: *b2*)
qed

end

5.2 Polynomials in a Finite Field

We connect polynomials in a prime field with integer polynomials modulo some prime.

theory *Poly-Mod-Finite-Field*
imports
Finite-Field
Polynomial-Interpolation.Ring-Hom-Poly
HOL−Types-To-Sets.Types-To-Sets
More-Missing-Multiset
Poly-Mod
begin

declare *rel-mset-Zero*[*transfer-rule*]

lemma *mset-transfer*[*transfer-rule*]: (*list-all2 rel* ===> *rel-mset rel*) *mset mset*
proof (*intro rel-funI*)
show *list-all2 rel xs ys* ⟹ *rel-mset rel* (*mset xs*) (*mset ys*) **for** *xs ys*
proof (*induct xs arbitrary*: *ys*)
case *Nil*
then show *?case* **by** *auto*
next
case *IH*: (*Cons x xs*)
then show *?case* **by** (*auto dest*!:*msed-rel-invL simp*: *list-all2-Cons1 intro*!:*rel-mset-Plus*)
qed
qed

abbreviation *to-int-poly* :: *'a* :: *finite mod-ring poly* ⇒ *int poly* **where**
to-int-poly ≡ *map-poly to-int-mod-ring*

interpretation *to-int-poly-hom*: *map-poly-inj-zero-hom to-int-mod-ring* **..**

lemma *irreducible*$_d$*-def-0*:
fixes *f* :: *'a* :: {*comm-semiring-1*,*semiring-no-zero-divisors*} *poly*
shows *irreducible*$_d$ *f* = (*degree f* ≠ *0* ∧
(∀ *g h*. *degree g* ≠ *0* ⟶ *degree h* ≠ *0* ⟶ *f* ≠ *g* ∗ *h*))
proof−
have *degree g* ≠ *0* ⟹ *g* ≠ *0* **for** *g* :: *'a poly* **by** *auto*
note *1* = *degree-mult-eq*[*OF this this*, *simplified*]
then show *?thesis* **by** (*force elim*!: *irreducible*$_d$*E*)
qed

5.3 Transferring to class-based mod-ring

locale *poly-mod-type* = *poly-mod m*
for *m* **and** *ty* :: *'a* :: *nontriv itself* +
assumes *m*: *m* = *CARD*(*'a*)

begin

lemma *m1*: $m > 1$ **using** *nontriv*[**where** $'a = 'a$] **by** (*auto simp*:*m*)

sublocale *poly-mod-2* **using** *m1* **by** *unfold-locales*

definition *MP-Rel* :: *int poly* ⇒ $'a$ *mod-ring poly* ⇒ *bool*
  where *MP-Rel f f'* ≡ (*Mp f* = *to-int-poly f'*)

definition *M-Rel* :: *int* ⇒ $'a$ *mod-ring* ⇒ *bool*
  where *M-Rel x x'* ≡ (*M x* = *to-int-mod-ring x'*)

definition *MF-Rel* ≡ *rel-prod M-Rel* (*rel-mset MP-Rel*)

lemma *to-int-mod-ring-plus*: *to-int-mod-ring* ((*x* :: $'a$ *mod-ring*) + *y*) = *M* (*to-int-mod-ring x* + *to-int-mod-ring y*)
  unfolding *M-def* **using** *m* **by** (*transfer*, *auto*)

lemma *to-int-mod-ring-times*: *to-int-mod-ring* ((*x* :: $'a$ *mod-ring*) ∗ *y*) = *M* (*to-int-mod-ring x* ∗ *to-int-mod-ring y*)
  unfolding *M-def* **using** *m* **by** (*transfer*, *auto*)

lemma *degree-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> (=)) *degree-m degree*
  unfolding *MP-Rel-def rel-fun-def*
  by (*auto intro*!: *degree-map-poly*)

lemma *eq-M-Rel*[*transfer-rule*]: (*M-Rel* ===> *M-Rel* ===> (=)) (λ *x y*. *M x* = *M y*) (=)
  unfolding *M-Rel-def rel-fun-def* **by** *auto*

interpretation *to-int-mod-ring-hom*: *map-poly-inj-zero-hom to-int-mod-ring*..

lemma *eq-MP-Rel*[*transfer-rule*]: (*MP-Rel* ===> *MP-Rel* ===> (=)) (=*m*) (=)
  unfolding *MP-Rel-def rel-fun-def* **by** *auto*

lemma *eq-Mf-Rel*[*transfer-rule*]: (*MF-Rel* ===> *MF-Rel* ===> (=)) (λ *x y*. *Mf x* = *Mf y*) (=)
proof (*intro rel-funI*, *goal-cases*)
  case (*1 cfs Cfs dgs Dgs*)
  obtain *c fs* **where** *cfs*: *cfs* = (*c*,*fs*) **by** *force*
  obtain *C Fs* **where** *Cfs*: *Cfs* = (*C*,*Fs*) **by** *force*
  obtain *d gs* **where** *dgs*: *dgs* = (*d*,*gs*) **by** *force*
  obtain *D Gs* **where** *Dgs*: *Dgs* = (*D*,*Gs*) **by** *force*
  note *pairs* = *cfs Cfs dgs Dgs*
  from *1*[*unfolded pairs MF-Rel-def rel-prod.simps*]
 have ∗[*transfer-rule*]: *M-Rel c C M-Rel d D rel-mset MP-Rel fs Fs rel-mset MP-Rel gs Gs*
    by *auto*
  have *eq1*: (*M c* = *M d*) = (*C* = *D*) **by** *transfer-prover*

```
  from ∗(3)[unfolded rel-mset-def] obtain fs′ Fs′ where fs-eq: mset fs′ = fs mset
Fs′ = Fs
    and rel-f: list-all2 MP-Rel fs′ Fs′ by auto
  from ∗(4)[unfolded rel-mset-def] obtain gs′ Gs′ where gs-eq: mset gs′ = gs mset
Gs′ = Gs
    and rel-g: list-all2 MP-Rel gs′ Gs′ by auto
  have eq2: (image-mset Mp fs = image-mset Mp gs) = (Fs = Gs)
    using ∗(3−4)
  proof (induct fs arbitrary: Fs gs Gs)
    case (empty Fs gs Gs)
    from empty(1) have Fs: Fs = {#} unfolding rel-mset-def by auto
    with empty show ?case by (cases gs; cases Gs; auto simp: rel-mset-def)
  next
    case (add f fs Fs′ gs′ Gs′)
    note [transfer-rule] = add(3)
    from msed-rel-invL[OF add(2)]
    obtain Fs F where Fs′: Fs′ = Fs + {#F#} and rel[transfer-rule]:
      MP-Rel f F rel-mset MP-Rel fs Fs by auto
    note IH = add(1)[OF rel(2)]
    {
      from add(3)[unfolded rel-mset-def] obtain gs Gs where id: mset gs = gs′
mset Gs = Gs′
        and rel: list-all2 MP-Rel gs Gs by auto
      have Mp f ∈# image-mset Mp gs′ ⟷ F ∈# Gs′
      proof −
        have ?thesis = ((Mp f ∈ Mp ‘ set gs) = (F ∈ set Gs))
          unfolding id[symmetric] by simp
        also have ... using rel
        proof (induct gs Gs rule: list-all2-induct)
          case (Cons g gs G Gs)
          note [transfer-rule] = Cons(1−2)
          have id: (Mp g = Mp f) = (F = G) by (transfer, auto)
          show ?case using id Cons(3) by auto
        qed auto
        finally show ?thesis by simp
      qed
    } note id = this
    show ?case
    proof (cases Mp f ∈# image-mset Mp gs′)
      case False
      have Mp f ∈# image-mset Mp (fs + {#f#}) by auto
      with False have F: image-mset Mp (fs + {#f#}) ≠ image-mset Mp gs′ by
metis
      with False[unfolded id] show ?thesis unfolding Fs′ by auto
    next
      case True
      then obtain g where fg: Mp f = Mp g and g: g ∈# gs′ by auto
      from g obtain gs where gs′: gs′ = add-mset g gs by (rule mset-add)
      from msed-rel-invL[OF add(3)[unfolded gs′]]
```

```
      obtain Gs G where Gs': Gs' = Gs + {# G #} and gG[transfer-rule]:
MP-Rel g G and
        gsGs: rel-mset MP-Rel gs Gs by auto
      have FG: F = G by (transfer, simp add: fg)
      note IH = IH[OF gsGs]
      show ?thesis unfolding gs' Fs' Gs' by (simp add: fg IH FG)
    qed
  qed
  show (Mf cfs = Mf dgs) = (Cfs = Dgs) unfolding pairs Mf-def split
    by (simp add: eq1 eq2)
qed

lemmas coeff-map-poly-of-int = coeff-map-poly[of of-int, OF of-int-0]

lemma plus-MP-Rel[transfer-rule]: (MP-Rel ===> MP-Rel ===> MP-Rel) (+)
(+)
  unfolding MP-Rel-def
proof (intro rel-funI, goal-cases)
  case (1 x f y g)
  have Mp (x + y) = Mp (Mp x + Mp y) by simp
  also have ... = Mp (map-poly to-int-mod-ring f + map-poly to-int-mod-ring g)
unfolding 1 ..
  also have ... = map-poly to-int-mod-ring (f + g) unfolding poly-eq-iff Mp-coeff

      by (auto simp: to-int-mod-ring-plus)
  finally show ?case .
qed

lemma times-MP-Rel[transfer-rule]: (MP-Rel ===> MP-Rel ===> MP-Rel)
((*)) ((*))
  unfolding MP-Rel-def
proof (intro rel-funI, goal-cases)
  case (1 x f y g)
  have Mp (x * y) = Mp (Mp x * Mp y) by simp
  also have ... = Mp (map-poly to-int-mod-ring f * map-poly to-int-mod-ring g)
unfolding 1 ..
  also have ... = map-poly to-int-mod-ring (f * g)
  proof -
    { fix n :: nat
      define A where A = {.. n}
      have finite A unfolding A-def by auto
      then have M (∑ i≤n. to-int-mod-ring (coeff f i) * to-int-mod-ring (coeff g
(n - i))) =
          to-int-mod-ring (∑ i≤n. coeff f i * coeff g (n - i))
        unfolding A-def[symmetric]
      proof (induct A)
        case (insert a A)
        have ?case = ?case (is (?l = ?r) = -) by simp
        have ?r = to-int-mod-ring (coeff f a * coeff g (n - a) + (∑ i∈ A. coeff f i
```

```
∗ coeff g (n − i)))
        using insert(1−2) by auto
      note r = this[unfolded to-int-mod-ring-plus to-int-mod-ring-times]
    from insert(1−2) have ?l = M (to-int-mod-ring (coeff f a) ∗ to-int-mod-ring
(coeff g (n − a))
        + M (∑ i∈A. to-int-mod-ring (coeff f i) ∗ to-int-mod-ring (coeff g (n −
i))))
        by simp
      also have M (∑ i∈A. to-int-mod-ring (coeff f i) ∗ to-int-mod-ring (coeff g
(n − i))) = to-int-mod-ring (∑ i∈A. coeff f i ∗ coeff g (n − i))
        unfolding insert ..
      finally
      show ?case unfolding r by simp
    qed auto
  }
  then show ?thesis by (auto intro!:poly-eqI simp: coeff-mult  Mp-coeff)
 qed
 finally show ?case .
qed

lemma smult-MP-Rel[transfer-rule]: (M-Rel ===> MP-Rel ===> MP-Rel) smult
smult
 unfolding MP-Rel-def M-Rel-def
proof (intro rel-funI, goal-cases)
 case (1 x x' f f')
 thus ?case unfolding poly-eq-iff coeff Mp-coeff
   coeff-smult M-def
 proof (intro allI, goal-cases)
   case (1 n)
   have x ∗ coeff f n mod m = (x mod m) ∗ (coeff f n mod m) mod m
     by (simp add: mod-simps)
   also have ... = to-int-mod-ring x' ∗ (to-int-mod-ring (coeff f' n)) mod m
     using 1 by auto
   also have  ... = to-int-mod-ring (x' ∗ coeff f' n)
     unfolding to-int-mod-ring-times M-def by simp
   finally show ?case by auto
 qed
qed

lemma one-M-Rel[transfer-rule]: M-Rel 1 1
 unfolding M-Rel-def M-def
 unfolding m by auto

lemma one-MP-Rel[transfer-rule]: MP-Rel 1 1
 unfolding MP-Rel-def poly-eq-iff Mp-coeff M-def
 unfolding m by auto

lemma zero-M-Rel[transfer-rule]: M-Rel 0 0
 unfolding M-Rel-def M-def
```

unfolding *m* **by** *auto*

lemma *zero-MP-Rel*[*transfer-rule*]: *MP-Rel 0 0*
unfolding *MP-Rel-def poly-eq-iff Mp-coeff M-def*
unfolding *m* **by** *auto*

lemma *listprod-MP-Rel*[*transfer-rule*]: (*list-all2 MP-Rel* ===> *MP-Rel*) *prod-list prod-list*
proof (*intro rel-funI*, *goal-cases*)
case (*1 xs ys*)
thus *?case*
proof (*induct xs ys rule*: *list-all2-induct*)
case (*Cons x xs y ys*)
note [*transfer-rule*] = *this*
show *?case* **by** *simp transfer-prover*
qed (*simp add*: *one-MP-Rel*)
qed

lemma *prod-mset-MP-Rel*[*transfer-rule*]: (*rel-mset MP-Rel* ===> *MP-Rel*) *prod-mset prod-mset*
proof (*intro rel-funI*, *goal-cases*)
case (*1 xs ys*)
have (*MP-Rel* ===> *MP-Rel* ===> *MP-Rel*) ((∗)) ((∗)) *MP-Rel 1 1* **by** *transfer-prover*+
from *1 this* **show** *?case*
proof (*induct xs ys rule*: *rel-mset-induct*)
case (*add R x xs y ys*)
note [*transfer-rule*] = *this*
show *?case* **by** *simp transfer-prover*
qed (*simp add*: *one-MP-Rel*)
qed

lemma *right-unique-MP-Rel*[*transfer-rule*]: *right-unique MP-Rel*
unfolding *right-unique-def MP-Rel-def* **by** *auto*

lemma *M-to-int-mod-ring*: *M* (*to-int-mod-ring* (*x* :: *'a mod-ring*)) = *to-int-mod-ring x*
unfolding *M-def* **unfolding** *m* **by** (*transfer*, *auto*)

lemma *Mp-to-int-poly*: *Mp* (*to-int-poly* (*f* :: *'a mod-ring poly*)) = *to-int-poly f*
by (*auto simp*: *poly-eq-iff Mp-coeff M-to-int-mod-ring*)

lemma *right-total-M-Rel*[*transfer-rule*]: *right-total M-Rel*
unfolding *right-total-def M-Rel-def* **using** *M-to-int-mod-ring* **by** *blast*

lemma *left-total-M-Rel*[*transfer-rule*]: *left-total M-Rel*
unfolding *left-total-def M-Rel-def*[*abs-def*]
proof
fix *x*

show $\exists\ x' :: 'a\ mod\text{-}ring.\ M\ x = to\text{-}int\text{-}mod\text{-}ring\ x'$ **unfolding** *M-def* **unfolding** *m*
by (*rule exI*[*of - of-int x*], *transfer*, *simp*)
qed

lemma *bi-total-M-Rel*[*transfer-rule*]: *bi-total M-Rel*
using *right-total-M-Rel left-total-M-Rel* **by** (*metis bi-totalI*)

lemma *right-total-MP-Rel*[*transfer-rule*]: *right-total MP-Rel*
unfolding *right-total-def MP-Rel-def*
proof
fix *f* :: *'a mod-ring poly*
show $\exists x.\ Mp\ x = to\text{-}int\text{-}poly\ f$
by (*intro exI*[*of - to-int-poly f*], *simp add*: *Mp-to-int-poly*)
qed

lemma *to-int-mod-ring-of-int-M*: *to-int-mod-ring* (*of-int x* :: *'a mod-ring*) = *M x*
unfolding *M-def*
unfolding *m* **by** *transfer auto*

lemma *Mp-f-representative*: *Mp f* = *to-int-poly* (*map-poly of-int f* :: *'a mod-ring poly*)
unfolding *Mp-def* **by** (*auto intro*: *poly-eqI simp*: *coeff-map-poly to-int-mod-ring-of-int-M*)

lemma *left-total-MP-Rel*[*transfer-rule*]: *left-total MP-Rel*
unfolding *left-total-def MP-Rel-def*[*abs-def*] **using** *Mp-f-representative* **by** *blast*

lemma *bi-total-MP-Rel*[*transfer-rule*]: *bi-total MP-Rel*
using *right-total-MP-Rel left-total-MP-Rel* **by** (*metis bi-totalI*)

lemma *bi-total-MF-Rel*[*transfer-rule*]: *bi-total MF-Rel*
unfolding *MF-Rel-def*[*abs-def*]
by (*intro prod.bi-total-rel multiset.bi-total-rel bi-total-MP-Rel bi-total-M-Rel*)

lemma *right-total-MF-Rel*[*transfer-rule*]: *right-total MF-Rel*
using *bi-total-MF-Rel* **unfolding** *bi-total-alt-def* **by** *auto*

lemma *left-total-MF-Rel*[*transfer-rule*]: *left-total MF-Rel*
using *bi-total-MF-Rel* **unfolding** *bi-total-alt-def* **by** *auto*

lemma *domain-RT-rel*[*transfer-domain-rule*]: *Domainp MP-Rel* = (*λ f. True*)
proof
fix *f* :: *int poly*
show *Domainp MP-Rel f* = *True* **unfolding** *MP-Rel-def*[*abs-def*] *Domainp.simps*
by (*auto simp*: *Mp-f-representative*)
qed

lemma *mem-MP-Rel*[*transfer-rule*]: (*MP-Rel* ===> *rel-set MP-Rel* ===> (=))
(*λ x Y. ∃ y ∈ Y. eq-m x y*) (∈)

proof (*intro rel-funI iffI*)
fix *x y X Y* **assume** *xy*: *MP-Rel x y* **and** *XY*: *rel-set MP-Rel X Y*
{ assume $\exists x' \in X.\ x =m\ x'$
then obtain *x'* **where** *x'X*: $x' \in X$ **and** *xx'*: $x =m\ x'$ **by** *auto*
with *xy* **have** *x'y*: *MP-Rel x' y* **by** (*auto simp*: *MP-Rel-def*)
from *rel-setD1*[*OF XY x'X*] **obtain** *y'* **where** *MP-Rel x' y'* **and** $y' \in Y$ **by** *auto*
with *x'y*
show $y \in Y$ **by** (*auto simp*: *MP-Rel-def*)
}
assume $y \in Y$
from *rel-setD2*[*OF XY this*] **obtain** *x'* **where** *x'X*: $x' \in X$ **and** *x'y*: *MP-Rel x' y* **by** *auto*
from *xy x'y* **have** $x =m\ x'$ **by** (*auto simp*: *MP-Rel-def*)
with *x'X* **show** $\exists x' \in X.\ x =m\ x'$ **by** *auto*
qed

lemma *conversep-MP-Rel-OO-MP-Rel* [*simp*]: $MP\text{-}Rel^{-1-1}\ OO\ MP\text{-}Rel = (=)$
using *Mp-to-int-poly* **by** (*intro ext*, *auto simp*: *OO-def MP-Rel-def*)

lemma *MP-Rel-OO-conversep-MP-Rel* [*simp*]: $MP\text{-}Rel\ OO\ MP\text{-}Rel^{-1-1} = eq\text{-}m$
by (*intro ext*, *auto simp*: *OO-def MP-Rel-def Mp-f-representative*)

lemma *conversep-MP-Rel-OO-eq-m* [*simp*]: $MP\text{-}Rel^{-1-1}\ OO\ eq\text{-}m = MP\text{-}Rel^{-1-1}$
by (*intro ext*, *auto simp*: *OO-def MP-Rel-def*)

lemma *eq-m-OO-MP-Rel* [*simp*]: *eq-m OO MP-Rel* = *MP-Rel*
by (*intro ext*, *auto simp*: *OO-def MP-Rel-def*)

lemma *eq-mset-MP-Rel* [*transfer-rule*]: (*rel-mset MP-Rel* ===> *rel-mset MP-Rel* ===> (=)) (*rel-mset eq-m*) (=)
proof (*intro rel-funI iffI*)
fix *A B X Y*
assume *AX*: *rel-mset MP-Rel A X* **and** *BY*: *rel-mset MP-Rel B Y*
{
assume *AB*: *rel-mset eq-m A B*
from *AX* **have** *rel-mset* $MP\text{-}Rel^{-1-1}$ *X A* **by** (*simp add*: *multiset.rel-flip*)
note *rel-mset-OO*[*OF this AB*]
note *rel-mset-OO*[*OF this BY*]
then show *X* = *Y* **by** (*simp add*: *multiset.rel-eq*)
}
assume *X* = *Y*
with *BY* **have** *rel-mset* $MP\text{-}Rel^{-1-1}$ *X B* **by** (*simp add*: *multiset.rel-flip*)
from *rel-mset-OO*[*OF AX this*]
show *rel-mset eq-m A B* **by** *simp*
qed

lemma *dvd-MP-Rel*[*transfer-rule*]: (*MP-Rel* ===> *MP-Rel* ===> (=)) (*dvdm*) (*dvd*)

unfolding *dvdm-def*[*abs-def*] *dvd-def*[*abs-def*]
by *transfer-prover*

lemma *irreducible-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> (=)) *irreducible-m irreducible*
unfolding *irreducible-m-def irreducible-def*
by *transfer-prover*

lemma *irreducible$_d$-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> (=)) *irreducible$_d$-m irreducible$_d$*
unfolding *irreducible$_d$-m-def*[*abs-def*] *irreducible$_d$-def*[*abs-def*]
by *transfer-prover*

lemma *UNIV-M-Rel*[*transfer-rule*]: *rel-set M-Rel* {*0..<m*} *UNIV*
unfolding *rel-set-def M-Rel-def*[*abs-def*] *M-def*
by (*auto simp*: *M-def m*, *goal-cases*, *metis to-int-mod-ring-of-int-mod-ring*, (*transfer*, *auto*)+)

lemma *coeff-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> (=) ===> *M-Rel*) *coeff coeff*
unfolding *rel-fun-def M-Rel-def MP-Rel-def Mp-coeff*[*symmetric*] **by** *auto*

lemma *M-1-1*: *M 1 = 1* **unfolding** *M-def* **unfolding** *m* **by** *simp*

lemma *square-free-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> (=)) *square-free-m square-free*
unfolding *square-free-m-def*[*abs-def*] *square-free-def*[*abs-def*]
by (*transfer-prover-start*, *transfer-step*+, *auto*)

lemma *mset-factors-m-MP-Rel* [*transfer-rule*]: (*rel-mset MP-Rel* ===> *MP-Rel* ===> (=)) *mset-factors-m mset-factors*
unfolding *mset-factors-def mset-factors-m-def*
by (*transfer-prover-start*, *transfer-step*+, *auto dest*:*eq-m-irreducible-m*)

lemma *coprime-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> *MP-Rel* ===> (=)) *coprime-m coprime*
unfolding *coprime-m-def*[*abs-def*] *coprime-def'* [*abs-def*]
by (*transfer-prover-start*, *transfer-step*+, *auto*)

lemma *prime-elem-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> (=)) *prime-elem-m prime-elem*
unfolding *prime-elem-m-def prime-elem-def* **by** *transfer-prover*

end

context *poly-mod-2* **begin**

lemma *non-empty*: {*0..<m*} ≠ {} **using** *m1* **by** *auto*

lemma *type-to-set*:

assumes *type-def*: $\exists (Rep :: 'b \Rightarrow int)\ Abs.\ type\text{-}definition\ Rep\ Abs\ \{0\ ..<\ m :: int\}$
shows *class.nontriv* (*TYPE*('*b*)) (**is** *?a*) **and** $m = int\ CARD('b)$ (**is** *?b*)
proof −
from *type-def* **obtain** *rep* :: $'b \Rightarrow int$ **and** *abs* :: $int \Rightarrow 'b$ **where** *t*: *type-definition rep abs* $\{0\ ..<\ m\}$ **by** *auto*
have *card* (*UNIV* :: '*b set*) = *card* $\{0\ ..<\ m\}$ **using** *t* **by** (*rule type-definition.card*)
also have ... = *m* **using** *m1* **by** *auto*
finally show *?b* **..**
then show *?a* **unfolding** *class.nontriv-def* **using** *m1* **by** *auto*
qed

end

locale *poly-mod-prime-type* = *poly-mod-type m ty* **for** *m* :: *int* **and**
ty :: '*a* :: *prime-card itself*
begin

lemma *factorization-MP-Rel* [*transfer-rule*]:
(*MP-Rel* ===> *MF-Rel* ===> (=)) *factorization-m* (*factorization Irr-Mon*)
unfolding *rel-fun-def*
proof (*intro allI impI*, *goal-cases*)
case (*1 f F cfs Cfs*)
note [*transfer-rule*] = *1*(*1*)
obtain *c fs* **where** *cfs*: *cfs* = (*c*,*fs*) **by** *force*
obtain *C Fs* **where** *Cfs*: *Cfs* = (*C*,*Fs*) **by** *force*
from *1*(*2*)[*unfolded rel-prod.simps cfs Cfs MF-Rel-def*]
have *tr*[*transfer-rule*]: *M-Rel c C rel-mset MP-Rel fs Fs* **by** *auto*
have *eq*: (*f* =*m smult c* (*prod-mset fs*) = (*F* = *smult C* (*prod-mset Fs*)))
by *transfer-prover*
have *set-mset Fs* ⊆ *Irr-Mon* = (∀ *x* ∈# *Fs*. $irreducible_d\ x \wedge monic\ x$) **unfolding** *Irr-Mon-def* **by** *auto*
also have ... = (∀ *f*∈#*fs*. $irreducible_d$-*m f* ∧ *monic* (*Mp f*))
proof (*rule sym*, *transfer-prover-start*, *transfer-step*+)
{
fix *f*
assume *f* ∈# *fs*
have *monic* (*Mp f*) ⟷ *M* (*coeff f* (*degree-m f*)) = *M 1*
unfolding *Mp-coeff*[*symmetric*] **by** *simp*
}
thus (∀ *f*∈#*fs*. $irreducible_d$-*m f* ∧ *monic* (*Mp f*)) =
(∀ *x*∈#*fs*. $irreducible_d$-*m x* ∧ *M* (*coeff x* (*degree-m x*)) = *M 1*) **by** *auto*
qed
finally
show *factorization-m f cfs* = *factorization Irr-Mon F Cfs* **unfolding** *cfs Cfs factorization-m-def factorization-def split eq* **by** *simp*
qed

lemma *unique-factorization-MP-Rel* [*transfer-rule*]: (*MP-Rel* ===> *MF-Rel* ===>

```
(=))
  unique-factorization-m (unique-factorization Irr-Mon)
  unfolding rel-fun-def
proof (intro allI impI, goal-cases)
  case (1 f F cfs Cfs)
  note [transfer-rule] = 1(1,2)
  let ?F = factorization Irr-Mon F
  let ?f = factorization-m f
  let ?R = Collect ?F
  let ?L = Mf ' Collect ?f
  note X-to-x = right-total-MF-Rel[unfolded right-total-def, rule-format]
  {
    fix X
    assume X ∈ ?R
    hence F: ?F X by simp
    from X-to-x[of X] obtain x where rel[transfer-rule]: MF-Rel x X by blast
    from F[untransferred] have Mf x ∈ ?L by blast
    with rel have ∃ x. Mf x ∈ ?L ∧ MF-Rel x X by blast
  } note R-to-L = this
  show unique-factorization-m f cfs = unique-factorization Irr-Mon F Cfs unfolding
    unique-factorization-m-def unique-factorization-def
  proof −
    have fF: ?F Cfs = ?f cfs by transfer simp
    have (?L = {Mf cfs}) = (?L ⊆ {Mf cfs} ∧ Mf cfs ∈ ?L) by blast
    also have ?L ⊆ {Mf cfs} = (∀ dfs. ?f dfs ⟶ Mf dfs = Mf cfs) by blast
    also have ... = (∀ y. ?F y ⟶ y = Cfs) (is ?left = ?right)
    proof (rule; intro allI impI)
      fix Dfs
      assume ∗: ?left and F: ?F Dfs
       from X-to-x[of Dfs] obtain dfs where [transfer-rule]: MF-Rel dfs Dfs by
auto
      from F[untransferred] have f: ?f dfs .
      from ∗[rule-format, OF f] have eq: Mf dfs = Mf cfs by simp
    have (Mf dfs = Mf cfs) = (Dfs = Cfs) by (transfer-prover-start, transfer-step+,
simp)
      thus Dfs = Cfs using eq by simp
    next
      fix dfs
      assume ∗: ?right and f: ?f dfs
      from left-total-MF-Rel obtain Dfs where
        rel[transfer-rule]: MF-Rel dfs Dfs unfolding left-total-def by blast
      have ?F Dfs by (transfer, rule f)
      from ∗[rule-format, OF this] have eq: Dfs = Cfs .
    have (Mf dfs = Mf cfs) = (Dfs = Cfs) by (transfer-prover-start, transfer-step+,
simp)
      thus Mf dfs = Mf cfs using eq by simp
    qed
    also have Mf cfs ∈ ?L = (∃ dfs. ?f dfs ∧ Mf cfs = Mf dfs) by auto
```

```
    also have ... = ?F Cfs  unfolding fF
    proof
      assume ∃ dfs. ?f dfs ∧ Mf cfs  = Mf dfs
      then obtain dfs where f: ?f dfs and id: Mf dfs = Mf cfs by auto
      from f have ?f (Mf dfs) by simp
      from this[unfolded id] show ?f cfs by simp
    qed blast
    finally show (?L = {Mf cfs}) = (?R = {Cfs}) by auto
  qed
qed

end

context begin
private lemma 1: poly-mod-type TYPE('a :: nontriv) m = (m = int CARD('a))
  and 2: class.nontriv TYPE('a) = (CARD('a) ≥ 2)
  unfolding poly-mod-type-def class.prime-card-def class.nontriv-def poly-mod-prime-type-def
by auto

private lemma 3: poly-mod-prime-type TYPE('b) m = (m = int CARD('b))
  and 4: class.prime-card TYPE('b :: prime-card) = prime CARD('b :: prime-card)

  unfolding poly-mod-type-def class.prime-card-def class.nontriv-def poly-mod-prime-type-def
by auto

lemmas poly-mod-type-simps = 1 2 3 4
end


lemma remove-duplicate-premise: (PROP P ⟹ PROP P ⟹ PROP Q) ≡ (PROP
P ⟹ PROP Q) (is ?l ≡ ?r)
proof (intro Pure.equal-intr-rule)
  assume p: PROP P and ppq: PROP ?l
  from ppq[OF p p] show PROP Q.
next
  assume p: PROP P and pq: PROP ?r
  from pq[OF p] show PROP Q.
qed

context poly-mod-prime begin

lemma type-to-set:
  assumes type-def: ∃(Rep :: 'b ⇒ int) Abs. type-definition Rep Abs {0 ..< p ::
int}
  shows class.prime-card (TYPE('b)) (is ?a) and p = int CARD('b) (is ?b)
proof −
  from prime have p2: p ≥ 2 by (rule prime-ge-2-int)
  from type-def obtain rep :: 'b ⇒ int and abs :: int ⇒ 'b where t: type-definition
rep abs {0 ..< p} by auto
```

have *card* (*UNIV* :: *'b set*) = *card* {*0* ..< *p*} **using** *t* **by** (*rule type-definition.card*)
also have ... = *p* **using** *p2* **by** *auto*
finally show *?b* **..**
then show *?a* **unfolding** *class.prime-card-def* **using** *prime p2* **by** *auto*
qed
end

lemmas (**in** *poly-mod-type*) *prime-elem-m-dvdm-multD* = *prime-elem-dvd-multD*
[**where** *'a* = *'a mod-ring poly,untransferred*]
lemmas (**in** *poly-mod-2*) *prime-elem-m-dvdm-multD* = *poly-mod-type.prime-elem-m-dvdm-multD*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *nontriv*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*, *OF non-empty*]

lemmas(**in** *poly-mod-prime-type*) *degree-m-mult-eq* = *degree-mult-eq*
[**where** *'a* = *'a mod-ring*, *untransferred*]
lemmas(**in** *poly-mod-prime*) *degree-m-mult-eq* = *poly-mod-prime-type.degree-m-mult-eq*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*, *OF non-empty*]

lemma(**in** *poly-mod-prime*) *irreducible$_d$-lifting*:
assumes *n*: *n* ≠ *0*
and *deg*: *poly-mod.degree-m* (*p*^*n*) *f* = *degree-m f*
and *irr*: *irreducible$_d$-m f*
shows *poly-mod.irreducible$_d$-m* (*p*^*n*) *f*
proof −
interpret *q*: *poly-mod-2 p*^*n* **unfolding** *poly-mod-2-def* **using** *n m1* **by** *auto*
show *q.irreducible$_d$-m f*
proof (*rule q.irreducible$_d$-mI*)
from *deg irr* **show** *q.degree-m f* > *0* **by** (*auto elim*: *irreducible$_d$-mE*)
then have *pdeg-f*: *degree-m f* ≠ *0* **by** (*simp add*: *deg*)
note *pMp-Mp* = *Mp-Mp-pow-is-Mp*[*OF n m1*]
fix *g h*
assume *deg-g*: *degree g* < *q.degree-m f* **and** *deg-h*: *degree h* < *q.degree-m f*
and *eq*: *q.eq-m f* (*g* ∗ *h*)
from *eq* **have** *p-f*: *f* =*m* (*g* ∗ *h*) **using** *pMp-Mp* **by** *metis*
have ¬*g* =*m 0* **and** ¬*h* =*m 0*
apply (*metis degree-0 mult-zero-left Mp-0 p-f pdeg-f poly-mod.mult-Mp*(*1*))
by (*metis degree-0 mult-eq-0-iff Mp-0 mult-Mp*(*2*) *p-f pdeg-f*)
note [*simp*] = *degree-m-mult-eq*[*OF this*]
from *degree-m-le*[*of g*] *deg-g*
have *2*: *degree-m g* < *degree-m f* **by** (*fold deg*, *auto*)
from *degree-m-le*[*of h*] *deg-h*
have *3*: *degree-m h* < *degree-m f* **by** (*fold deg*, *auto*)
from *irreducible$_d$-mD*(*2*)[*OF irr 2 3*] *p-f*
show *False* **by** *auto*
qed
qed

lemmas (**in** *poly-mod-prime-type*) *mset-factors-exist* =
mset-factors-exist[**where** *'a* = *'a mod-ring poly*,*untransferred*]
lemmas (**in** *poly-mod-prime*) *mset-factors-exist* = *poly-mod-prime-type.mset-factors-exist*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*,
unfolded remove-duplicate-premise, *cancel-type-definition*, *OF non-empty*]

lemmas (**in** *poly-mod-prime-type*) *mset-factors-unique* =
mset-factors-unique[**where** *'a* = *'a mod-ring poly*,*untransferred*]
lemmas (**in** *poly-mod-prime*) *mset-factors-unique* = *poly-mod-prime-type.mset-factors-unique*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*,
unfolded remove-duplicate-premise, *cancel-type-definition*, *OF non-empty*]

lemmas (**in** *poly-mod-prime-type*) *prime-elem-iff-irreducible* =
prime-elem-iff-irreducible[**where** *'a* = *'a mod-ring poly*,*untransferred*]
lemmas (**in** *poly-mod-prime*) *prime-elem-iff-irreducible*[*simp*] = *poly-mod-prime-type.prime-elem-iff-irreducible*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*,
unfolded remove-duplicate-premise, *cancel-type-definition*, *OF non-empty*]

lemmas (**in** *poly-mod-prime-type*) *irreducible-connect* =
irreducible-connect-field[**where** *'a* = *'a mod-ring*, *untransferred*]
lemmas (**in** *poly-mod-prime*) *irreducible-connect*[*simp*] = *poly-mod-prime-type.irreducible-connect*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*,
unfolded remove-duplicate-premise, *cancel-type-definition*, *OF non-empty*]

lemmas (**in** *poly-mod-prime-type*) *irreducible-degree* =
irreducible-degree-field[**where** *'a* = *'a mod-ring*, *untransferred*]
lemmas (**in** *poly-mod-prime*) *irreducible-degree* = *poly-mod-prime-type.irreducible-degree*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*,
unfolded remove-duplicate-premise, *cancel-type-definition*, *OF non-empty*]

end

5.4 Karatsuba's Multiplication Algorithm for Polynomials

theory *Karatsuba-Multiplication*
imports
Polynomial-Interpolation.Missing-Polynomial
begin

lemma *karatsuba-main-step*: **fixes** *f* :: *'a* :: *comm-ring-1 poly*
assumes *f*: *f* = *monom-mult n f1* + *f0* **and** *g*: *g* = *monom-mult n g1* + *g0*
shows
monom-mult (*n* + *n*) (*f1* ∗ *g1*) + (*monom-mult n* (*f1* ∗ *g1* − (*f1* − *f0*) ∗ (*g1*
− *g0*) + *f0* ∗ *g0*) + *f0* ∗ *g0*) = *f* ∗ *g*
unfolding *assms*
by (*auto simp*: *field-simps mult-monom monom-mult-def*)

lemma *karatsuba-single-sided*: **fixes** *f* :: *'a* :: *comm-ring-1 poly*
assumes *f* = *monom-mult n f1* + *f0*
shows *monom-mult n* (*f1* ∗ *g*) + *f0* ∗ *g* = *f* ∗ *g*
unfolding *assms* **by** (*auto simp*: *field-simps mult-monom monom-mult-def*)

definition *split-at* :: *nat* ⇒ *'a list* ⇒ *'a list* × *'a list* **where**
[*code del*]: *split-at n xs* = (*take n xs*, *drop n xs*)

lemma *split-at-code*[*code*]:
split-at n [] = ([],[])
split-at n (*x* # *xs*) = (*if n* = *0 then* ([], *x* # *xs*) *else case split-at* (*n*−*1*) *xs of*
(*bef*,*aft*)
⇒ (*x* # *bef*, *aft*))
unfolding *split-at-def* **by** (*force*, *cases n*, *auto*)

fun *coeffs-minus* :: *'a* :: *ab-group-add list* ⇒ *'a list* ⇒ *'a list* **where**
coeffs-minus (*x* # *xs*) (*y* # *ys*) = ((*x* − *y*) # *coeffs-minus xs ys*)
| *coeffs-minus xs* [] = *xs*
| *coeffs-minus* [] *ys* = *map uminus ys*

The following constant determines at which size we will switch to the standard multiplication algorithm.

definition *karatsuba-lower-bound* **where** [*termination-simp*]: *karatsuba-lower-bound* = (*7* :: *nat*)

fun *karatsuba-main* :: *'a* :: *comm-ring-1 list* ⇒ *nat* ⇒ *'a list* ⇒ *nat* ⇒ *'a poly* **where**
karatsuba-main f n g m = (*if n* ≤ *karatsuba-lower-bound* ∨ *m* ≤ *karatsuba-lower-bound then*
let ff = *poly-of-list f in foldr* (λ*a p*. *smult a ff* + *pCons 0 p*) *g 0*
else let n2 = *n div 2 in*
if m > *n2 then* (*case split-at n2 f of*
(*f0*,*f1*) ⇒ *case split-at n2 g of*
(*g0*,*g1*) ⇒ *let*
p1 = *karatsuba-main f1* (*n* − *n2*) *g1* (*m* − *n2*);
p2 = *karatsuba-main* (*coeffs-minus f1 f0*) *n2* (*coeffs-minus g1 g0*) *n2*;
p3 = *karatsuba-main f0 n2 g0 n2*
in monom-mult (*n2* + *n2*) *p1* + (*monom-mult n2* (*p1* − *p2* + *p3*) + *p3*))
else case split-at n2 f of
(*f0*,*f1*) ⇒ *let*
p1 = *karatsuba-main f1* (*n* − *n2*) *g m*;
p2 = *karatsuba-main f0 n2 g m*
in monom-mult n2 p1 + *p2*)

declare *karatsuba-main.simps*[*simp del*]

lemma *poly-of-list-split-at*: **assumes** *split-at n f* = (*f0*,*f1*)

shows *poly-of-list f = monom-mult n (poly-of-list f1) + poly-of-list f0*
proof −
from *assms* **have** *id*: *f1 = drop n f f0 = take n f* **unfolding** *split-at-def* **by** *auto*
show *?thesis* **unfolding** *id*
proof (*rule poly-eqI*)
fix *i*
show *coeff (poly-of-list f) i =*
coeff (monom-mult n (poly-of-list (drop n f)) + poly-of-list (take n f)) i
unfolding *monom-mult-def coeff-monom-mult coeff-add poly-of-list-def coeff-Poly*
by (*cases n ≤ i*; *cases i ≥ length f*, *auto simp*: *nth-default-nth nth-default-beyond*)
qed
qed

lemma *coeffs-minus*: *poly-of-list (coeffs-minus f1 f0) = poly-of-list f1 − poly-of-list f0*
proof (*rule poly-eqI*, *unfold poly-of-list-def coeff-diff coeff-Poly*)
fix *i*
show *nth-default 0 (coeffs-minus f1 f0) i = nth-default 0 f1 i − nth-default 0 f0 i*
proof (*induct f1 f0 arbitrary*: *i rule*: *coeffs-minus.induct*)
case (*1 x xs y ys*)
thus *?case* **by** (*cases i*, *auto*)
next
case (*3 x xs*)
thus *?case* **unfolding** *coeffs-minus.simps*
by (*subst nth-default-map-eq*[*of uminus 0 0*], *auto*)
qed *auto*
qed

lemma *karatsuba-main*: *karatsuba-main f n g m = poly-of-list f ∗ poly-of-list g*
proof (*induct n arbitrary*: *f g m rule*: *less-induct*)
case (*less n f g m*)
note *simp*[*simp*] = *karatsuba-main.simps*[*of f n g m*]
show *?case* (**is** *?lhs = ?rhs*)
proof (*cases* (*n ≤ karatsuba-lower-bound ∨ m ≤ karatsuba-lower-bound*) = *False*)
case *False*
hence *lhs*: *?lhs = foldr* (*λa p. smult a* (*poly-of-list f*) + *pCons 0 p*) *g 0* **by** *simp*
have *rhs*: *?rhs = poly-of-list g ∗ poly-of-list f* **by** *simp*
also have *. . . = foldr* (*λa p. smult a* (*poly-of-list f*) + *pCons 0 p*) (*strip-while* ((=) *0*) *g*) *0*
unfolding *times-poly-def fold-coeffs-def poly-of-list-impl* **..**
also have *. . . = ?lhs* **unfolding** *lhs*
proof (*induct g*)
case (*Cons x xs*)
have ∀ *x*∈*set xs. x = 0 ⟹ foldr* (*λa p. smult a* (*Poly f*) + *pCons 0 p*) *xs 0 = 0*
by (*induct xs*, *auto*)

```
    thus ?case using Cons by (auto simp: cCons-def Cons)
   qed auto
   finally show ?thesis by simp
  next
   case True
   let ?n2 = n div 2
   have ?n2 < n n - ?n2 < n using True unfolding karatsuba-lower-bound-def
by auto
   note IH = less[OF this(1)] less[OF this(2)]
   obtain f1 f0 where f: split-at ?n2 f = (f0,f1) by force
   obtain g1 g0 where g: split-at ?n2 g = (g0,g1) by force
   note fsplit = poly-of-list-split-at[OF f]
   note gsplit = poly-of-list-split-at[OF g]
   show ?lhs = ?rhs unfolding simp Let-def f g split IH True if-False coeffs-minus
     karatsuba-single-sided[OF fsplit] karatsuba-main-step[OF fsplit gsplit] by auto
  qed
qed


definition karatsuba-mult-poly :: 'a :: comm-ring-1 poly ⇒ 'a poly ⇒ 'a poly where
  karatsuba-mult-poly f g = (let ff = coeffs f; gg = coeffs g; n = length ff; m =
length gg
   in (if n ≤ karatsuba-lower-bound ∨ m ≤ karatsuba-lower-bound then if n ≤ m
   then foldr (λa p. smult a g + pCons 0 p) ff 0
   else foldr (λa p. smult a f + pCons 0 p) gg 0
   else if n ≤ m
   then karatsuba-main gg m ff n
   else karatsuba-main ff n gg m))

lemma karatsuba-mult-poly: karatsuba-mult-poly f g = f * g
proof -
  note d = karatsuba-mult-poly-def Let-def
  let ?len = length (coeffs f) ≤ length (coeffs g)
  show ?thesis (is ?lhs = ?rhs)
   proof (cases length (coeffs f) ≤ karatsuba-lower-bound ∨ length (coeffs g) ≤
karatsuba-lower-bound)
   case True note outer = this
   show ?thesis
   proof (cases ?len)
    case True
     with outer have ?lhs = foldr (λa p. smult a g + pCons 0 p) (coeffs f) 0
unfolding d by auto
    also have ... = ?rhs unfolding times-poly-def fold-coeffs-def by auto
    finally show ?thesis .
   next
    case False
     with outer have ?lhs = foldr (λa p. smult a f + pCons 0 p) (coeffs g) 0
unfolding d by auto
    also have ... = g * f unfolding times-poly-def fold-coeffs-def by auto
```

also have ... = *?rhs* **by** *simp*
finally show *?thesis* **.**
qed
next
case *False* **note** *outer* = *this*
show *?thesis*
proof (*cases ?len*)
case *True*
with *outer* **have** *?lhs* = *karatsuba-main* (*coeffs g*) (*length* (*coeffs g*)) (*coeffs f*) (*length* (*coeffs f*))
unfolding *d* **by** *auto*
also have ... = *g* ∗ *f* **unfolding** *karatsuba-main* **by** *auto*
also have ... = *?rhs* **by** *auto*
finally show *?thesis* **.**
next
case *False*
with *outer* **have** *?lhs* = *karatsuba-main* (*coeffs f*) (*length* (*coeffs f*)) (*coeffs g*) (*length* (*coeffs g*))
unfolding *d* **by** *auto*
also have ... = *?rhs* **unfolding** *karatsuba-main* **by** *auto*
finally show *?thesis* **.**
qed
qed
qed

lemma *karatsuba-mult-poly-code-unfold*[*code-unfold*]: (∗) = *karatsuba-mult-poly*
by (*intro ext*, *unfold karatsuba-mult-poly*, *auto*)

The following declaration will resolve a race-conflict between (∗) = *karatsuba-mult-poly* and *monom 1 ?n* ∗ *?f* = *monom-mult ?n ?f*
?f ∗ *monom 1 ?n* = *monom-mult ?n ?f*.

lemmas *karatsuba-monom-mult-code-unfold*[*code-unfold*] =
monom-mult-unfold[**where** *f* = *f* :: *'a* :: *comm-ring-1 poly* **for** *f*, *unfolded karatsuba-mult-poly-code-unfold*]

end

5.5 Record Based Version

We provide an implementation for polynomials which may be parametrized by the ring- or field-operations. These don't have to be type-based!

5.5.1 Definitions

theory *Polynomial-Record-Based*
imports
Arithmetic-Record-Based
Karatsuba-Multiplication
begin

context
  fixes *ops* :: *'i arith-ops-record* (**structure**)
begin
private abbreviation (*input*) *zero* **where** *zero* ≡ *arith-ops-record.zero ops*
private abbreviation (*input*) *one* **where** *one* ≡ *arith-ops-record.one ops*
private abbreviation (*input*) *plus* **where** *plus* ≡ *arith-ops-record.plus ops*
private abbreviation (*input*) *times* **where** *times* ≡ *arith-ops-record.times ops*
private abbreviation (*input*) *minus* **where** *minus* ≡ *arith-ops-record.minus ops*
private abbreviation (*input*) *uminus* **where** *uminus* ≡ *arith-ops-record.uminus ops*
private abbreviation (*input*) *divide* **where** *divide* ≡ *arith-ops-record.divide ops*
private abbreviation (*input*) *inverse* **where** *inverse* ≡ *arith-ops-record.inverse ops*
private abbreviation (*input*) *modulo* **where** *modulo* ≡ *arith-ops-record.modulo ops*
private abbreviation (*input*) *normalize* **where** *normalize* ≡ *arith-ops-record.normalize ops*
private abbreviation (*input*) *unit-factor* **where** *unit-factor* ≡ *arith-ops-record.unit-factor ops*
private abbreviation (*input*) *DP* **where** *DP* ≡ *arith-ops-record.DP ops*

definition *is-poly* :: *'i list* ⇒ *bool* **where**
  is-poly xs ⟷ *list-all DP xs* ∧ *no-trailing* (*HOL.eq zero*) *xs*

definition *cCons-i* :: *'i* ⇒ *'i list* ⇒ *'i list*
where
  cCons-i x xs = (*if xs* = [] ∧ *x* = *zero then* [] *else x* # *xs*)

fun *plus-poly-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* **where**
  plus-poly-i (*x* # *xs*) (*y* # *ys*) = *cCons-i* (*plus x y*) (*plus-poly-i xs ys*)
| *plus-poly-i xs* [] = *xs*
| *plus-poly-i* [] *ys* = *ys*

definition *uminus-poly-i* :: *'i list* ⇒ *'i list* **where**
  [*code-unfold*]: *uminus-poly-i* = *map uminus*

fun *minus-poly-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* **where**
  minus-poly-i (*x* # *xs*) (*y* # *ys*) = *cCons-i* (*minus x y*) (*minus-poly-i xs ys*)
| *minus-poly-i xs* [] = *xs*
| *minus-poly-i* [] *ys* = *uminus-poly-i ys*

abbreviation (*input*) *zero-poly-i* :: *'i list* **where**
  zero-poly-i ≡ []

definition *one-poly-i* :: *'i list* **where**
  [*code-unfold*]: *one-poly-i* = [*one*]

definition *smult-i* :: *'i* ⇒ *'i list* ⇒ *'i list* **where**
smult-i a pp = (*if a* = *zero then* [] *else strip-while* ((=) *zero*) (*map* (*times a*) *pp*))

definition *sdiv-i* :: *'i list* ⇒ *'i* ⇒ *'i list* **where**
sdiv-i pp a = (*strip-while* ((=) *zero*) (*map* (λ *c*. *divide c a*) *pp*))

definition *poly-of-list-i* :: *'i list* ⇒ *'i list* **where**
poly-of-list-i = *strip-while* ((=) *zero*)

fun *coeffs-minus-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* **where**
coeffs-minus-i (*x* # *xs*) (*y* # *ys*) = (*minus x y* # *coeffs-minus-i xs ys*)
| *coeffs-minus-i xs* [] = *xs*
| *coeffs-minus-i* [] *ys* = *map uminus ys*

definition *monom-mult-i* :: *nat* ⇒ *'i list* ⇒ *'i list* **where**
monom-mult-i n xs = (*if xs* = [] *then xs else replicate n zero* @ *xs*)

fun *karatsuba-main-i* :: *'i list* ⇒ *nat* ⇒ *'i list* ⇒ *nat* ⇒ *'i list* **where**
karatsuba-main-i f n g m = (*if n* ≤ *karatsuba-lower-bound* ∨ *m* ≤ *karatsuba-lower-bound*
then
let ff = *poly-of-list-i f in foldr* (λ*a p*. *plus-poly-i* (*smult-i a ff*) (*cCons-i zero p*))
g zero-poly-i
else let n2 = *n div 2 in*
if m > *n2 then* (*case split-at n2 f of*
(*f0*,*f1*) ⇒ *case split-at n2 g of*
(*g0*,*g1*) ⇒ *let*
p1 = *karatsuba-main-i f1* (*n* − *n2*) *g1* (*m* − *n2*);
p2 = *karatsuba-main-i* (*coeffs-minus-i f1 f0*) *n2* (*coeffs-minus-i g1 g0*) *n2*;
p3 = *karatsuba-main-i f0 n2 g0 n2*
in plus-poly-i (*monom-mult-i* (*n2* + *n2*) *p1*)
(*plus-poly-i* (*monom-mult-i n2* (*plus-poly-i* (*minus-poly-i p1 p2*) *p3*)) *p3*))
else case split-at n2 f of
(*f0*,*f1*) ⇒ *let*
p1 = *karatsuba-main-i f1* (*n* − *n2*) *g m*;
p2 = *karatsuba-main-i f0 n2 g m*
in plus-poly-i (*monom-mult-i n2 p1*) *p2*)

definition *times-poly-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* **where**
times-poly-i f g ≡ (*let n* = *length f*; *m* = *length g*
in (*if n* ≤ *karatsuba-lower-bound* ∨ *m* ≤ *karatsuba-lower-bound then if n* ≤ *m*
then
foldr (λ*a p*. *plus-poly-i* (*smult-i a g*) (*cCons-i zero p*)) *f zero-poly-i else*
foldr (λ*a p*. *plus-poly-i* (*smult-i a f*) (*cCons-i zero p*)) *g zero-poly-i else*
if n ≤ *m then karatsuba-main-i g m f n else karatsuba-main-i f n g m*))

definition *coeff-i* :: *'i list* ⇒ *nat* ⇒ *'i* **where**
coeff-i = *nth-default zero*

definition *degree-i* :: *'i list* ⇒ *nat* **where**

degree-i pp ≡ *length pp* − *1*

definition *lead-coeff-i* :: *'i list* ⇒ *'i* **where**
lead-coeff-i pp = (*case pp of* [] ⇒ *zero* | - ⇒ *last pp*)

definition *monic-i* :: *'i list* ⇒ *bool* **where**
monic-i pp = (*lead-coeff-i pp* = *one*)

fun *minus-poly-rev-list-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* **where**
minus-poly-rev-list-i (*x* # *xs*) (*y* # *ys*) = (*minus x y*) # (*minus-poly-rev-list-i xs ys*)
| *minus-poly-rev-list-i xs* [] = *xs*
| *minus-poly-rev-list-i* [] (*y* # *ys*) = []

fun *divmod-poly-one-main-i* :: *'i list* ⇒ *'i list* ⇒ *'i list*
⇒ *nat* ⇒ *'i list* × *'i list* **where**
divmod-poly-one-main-i q r d (*Suc n*) = (*let*
a = *hd r*;
qqq = *cCons-i a q*;
rr = *tl* (*if a* = *zero then r else minus-poly-rev-list-i r* (*map* (*times a*) *d*))
in divmod-poly-one-main-i qqq rr d n)
| *divmod-poly-one-main-i q r d 0* = (*q*,*r*)

fun *mod-poly-one-main-i* :: *'i list* ⇒ *'i list*
⇒ *nat* ⇒ *'i list* **where**
mod-poly-one-main-i r d (*Suc n*) = (*let*
a = *hd r*;
rr = *tl* (*if a* = *zero then r else minus-poly-rev-list-i r* (*map* (*times a*) *d*))
in mod-poly-one-main-i rr d n)
| *mod-poly-one-main-i r d 0* = *r*

definition *pdivmod-monic-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* × *'i list* **where**
pdivmod-monic-i cf cg ≡ *case*
divmod-poly-one-main-i [] (*rev cf*) (*rev cg*) (*1* + *length cf* − *length cg*)
of (*q*,*r*) ⇒ (*poly-of-list-i q*, *poly-of-list-i* (*rev r*))

definition *dupe-monic-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* ⇒ *'i list* ⇒ *'i list* × *'i list* **where**
dupe-monic-i D H S T U = (*case pdivmod-monic-i* (*times-poly-i T U*) *D of* (*Q*,*R*) ⇒
(*plus-poly-i* (*times-poly-i S U*) (*times-poly-i H Q*), *R*))

definition *of-int-poly-i* :: *int poly* ⇒ *'i list* **where**
of-int-poly-i f = *map* (*arith-ops-record.of-int ops*) (*coeffs f*)

definition *to-int-poly-i* :: *'i list* ⇒ *int poly* **where**
to-int-poly-i f = *poly-of-list* (*map* (*arith-ops-record.to-int ops*) *f*)

definition *dupe-monic-i-int* :: *int poly* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* ⇒ *int*

```
poly ⇒ int poly × int poly where
  dupe-monic-i-int D H S T = (let
      d = of-int-poly-i D;
      h = of-int-poly-i H;
      s = of-int-poly-i S;
      t = of-int-poly-i T
    in (λ U. case dupe-monic-i d h s t (of-int-poly-i U) of
      (D',H') ⇒ (to-int-poly-i D', to-int-poly-i H')))

definition div-field-poly-i :: 'i list ⇒ 'i list ⇒ 'i list where
  div-field-poly-i cf cg = (
      if cg = [] then zero-poly-i
        else let ilc = inverse (last cg); ch = map (times ilc) cg;
                 q = fst (divmod-poly-one-main-i [] (rev cf) (rev ch) (1 + length cf
− length cg))
            in poly-of-list-i ((map (times ilc) q)))

definition mod-field-poly-i :: 'i list ⇒ 'i list ⇒ 'i list where
  mod-field-poly-i cf cg = (
      if cg = [] then cf
        else let ilc = inverse (last cg); ch = map (times ilc) cg;
                 r = mod-poly-one-main-i (rev cf) (rev ch) (1 + length cf − length
cg)
            in poly-of-list-i (rev r))

definition normalize-poly-i :: 'i list ⇒ 'i list where
  normalize-poly-i xs = smult-i (inverse (unit-factor (lead-coeff-i xs))) xs

definition unit-factor-poly-i :: 'i list ⇒ 'i list where
  unit-factor-poly-i xs = cCons-i (unit-factor (lead-coeff-i xs)) []

fun pderiv-main-i :: 'i ⇒ 'i list ⇒ 'i list where
  pderiv-main-i f (x # xs) = cCons-i (times f x) (pderiv-main-i (plus f one) xs)
| pderiv-main-i f [] = []

definition pderiv-i :: 'i list ⇒ 'i list where
  pderiv-i xs = pderiv-main-i one (tl xs)

definition dvd-poly-i :: 'i list ⇒ 'i list ⇒ bool where
  dvd-poly-i xs ys = (∃ zs. is-poly zs ∧ ys = times-poly-i xs zs)

definition irreducible-i :: 'i list ⇒ bool where
  irreducible-i xs = (degree-i xs ≠ 0 ∧
  (∀ q r. is-poly q ⟶ is-poly r ⟶ degree-i q < degree-i xs ⟶ degree-i r < degree-i
xs
    ⟶ xs ≠ times-poly-i q r))

definition poly-ops :: 'i list arith-ops-record where
  poly-ops ≡ Arith-Ops-Record
```

zero-poly-i
one-poly-i
plus-poly-i
times-poly-i
minus-poly-i
uminus-poly-i
div-field-poly-i
(λ -. []) — not defined
mod-field-poly-i
normalize-poly-i
unit-factor-poly-i
(λ *i*. *if i = 0 then* [] *else* [*arith-ops-record.of-int ops i*])
(λ -. *0*) — not defined
is-poly

definition *gcd-poly-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* **where**
gcd-poly-i = arith-ops.gcd-eucl-i poly-ops

definition *euclid-ext-poly-i* :: *'i list* ⇒ *'i list* ⇒ (*'i list* × *'i list*) × *'i list* **where**
euclid-ext-poly-i = arith-ops.euclid-ext-i poly-ops

definition *separable-i* :: *'i list* ⇒ *bool* **where**
separable-i xs ≡ *gcd-poly-i xs* (*pderiv-i xs*) = *one-poly-i*

end

5.5.2 Properties

definition *pdivmod-monic* :: *'a*::*comm-ring-1 poly* ⇒ *'a poly* ⇒ *'a poly* × *'a poly* **where**
pdivmod-monic f g ≡ *let cg = coeffs g*; *cf = coeffs f*;
(*q*, *r*) = *divmod-poly-one-main-list* [] (*rev cf*) (*rev cg*) (*1 + length cf − length cg*)
in (*poly-of-list q*, *poly-of-list* (*rev r*))

lemma *coeffs-smult'*: *coeffs* (*smult a p*) = (*if a = 0 then* [] *else strip-while* ((=) *0*) (*map* (*Groups.times a*) (*coeffs p*)))
by (*simp add*: *coeffs-map-poly smult-conv-map-poly*)

lemma *coeffs-sdiv*: *coeffs* (*sdiv-poly p a*) = (*strip-while* ((=) *0*) (*map* (λ *x*. *x div a*) (*coeffs p*)))
unfolding *sdiv-poly-def* **by** (*rule coeffs-map-poly*)

lifting-forget *poly.lifting*

context *ring-ops*
begin

definition *poly-rel* :: *'i list* ⇒ *'a poly* ⇒ *bool* **where**
poly-rel x x' ⟷ *list-all2 R x* (*coeffs x'*)

lemma *right-total-poly-rel*[*transfer-rule*]:
right-total poly-rel
using *list.right-total-rel*[*of R*] *right-total* **unfolding** *poly-rel-def right-total-def* **by** *auto*

lemma *poly-rel-inj*: *poly-rel x y* ⟹ *poly-rel x z* ⟹ *y* = *z*
using *list.bi-unique-rel*[*OF bi-unique*] **unfolding** *poly-rel-def coeffs-eq-iff bi-unique-def* **by** *auto*

lemma *bi-unique-poly-rel*[*transfer-rule*]: *bi-unique poly-rel*
using *list.bi-unique-rel*[*OF bi-unique*] **unfolding** *poly-rel-def bi-unique-def coeffs-eq-iff* **by** *auto*

lemma *Domainp-is-poly* [*transfer-domain-rule*]:
Domainp poly-rel = *is-poly ops*
unfolding *poly-rel-def* [*abs-def*] *is-poly-def* [*abs-def*]
proof (*intro ext iffI*, *unfold Domainp-iff*)
note *DPR* = *fun-cong* [*OF list.Domainp-rel* [*of R*, *unfolded DPR*],
unfolded Domainp-iff]
let *?no-trailing* = *no-trailing* (*HOL.eq zero*)
fix *xs*
have *no-trailing*: *no-trailing* (*HOL.eq 0*) *xs'* ⟷ *?no-trailing xs*
if *list-all2 R xs xs'* **for** *xs'*
proof (*cases xs rule*: *rev-cases*)
case *Nil*
with *that* **show** *?thesis*
by *simp*
next
case (*snoc ys y*)
with *that* **have** *xs'* ≠ []
by *auto*
then obtain *ys' y'* **where** *xs'* = *ys'* @ [*y'*]
by (*cases xs' rule*: *rev-cases*) *simp-all*
with *that snoc* **show** *?thesis*
by *simp* (*meson bi-unique bi-unique-def zero*)
qed
let *?DPR* = *arith-ops-record.DP ops*
{
assume ∃ *x'*. *list-all2 R xs* (*coeffs x'*)
then obtain *xs'* **where** ∗: *list-all2 R xs* (*coeffs xs'*) **by** *auto*
with *DPR* [*of xs*] **have** *list-all ?DPR xs* **by** *auto*
then show *list-all ?DPR xs* ∧ *?no-trailing xs*
using *no-trailing* [*OF* ∗] **by** *simp*
}
{
assume *list-all ?DPR xs* ∧ *?no-trailing xs*

```
    with DPR [of xs] obtain xs' where *: list-all2 R xs xs' and ?no-trailing xs
      by auto
    from no-trailing [OF *] this(2) have no-trailing (HOL.eq 0) xs'
      by simp
    hence coeffs (poly-of-list xs') = xs' unfolding poly-of-list-impl by auto
    with * show ∃x'. list-all2 R xs (coeffs x') by metis
  }
qed


lemma poly-rel-zero[transfer-rule]: poly-rel zero-poly-i 0
  unfolding poly-rel-def by auto


lemma poly-rel-one[transfer-rule]: poly-rel (one-poly-i ops) 1
  unfolding poly-rel-def one-poly-i-def by (simp add: one)


lemma poly-rel-cCons[transfer-rule]: (R ===> list-all2 R ===> list-all2 R)
(cCons-i ops) cCons
  unfolding cCons-i-def[abs-def] cCons-def[abs-def]
  by transfer-prover


lemma poly-rel-pCons[transfer-rule]: (R ===> poly-rel ===> poly-rel) (cCons-i
ops) pCons
  unfolding rel-fun-def poly-rel-def coeffs-pCons-eq-cCons cCons-def[symmetric]
  using poly-rel-cCons[unfolded rel-fun-def] by auto


lemma poly-rel-eq[transfer-rule]: (poly-rel ===> poly-rel ===> (=)) (=) (=)
  unfolding poly-rel-def[abs-def] coeffs-eq-iff[abs-def] rel-fun-def
  by (metis bi-unique bi-uniqueDl bi-uniqueDr list.bi-unique-rel)


lemma poly-rel-plus[transfer-rule]: (poly-rel ===> poly-rel ===> poly-rel) (plus-poly-i
ops) (+)
proof (intro rel-funI)
  fix x1 y1 x2 y2
  assume poly-rel x1 x2 and poly-rel y1 y2
  thus poly-rel (plus-poly-i ops x1 y1) (x2 + y2)
    unfolding poly-rel-def coeffs-eq-iff coeffs-plus-eq-plus-coeffs
  proof (induct x1 y1 arbitrary: x2 y2 rule: plus-poly-i.induct)
    case (1 x1 xs1 y1 ys1 X2 Y2)
    from 1(2) obtain x2 xs2 where X2: coeffs X2 = x2 # coeffs xs2
      by (cases X2, auto simp: cCons-def split: if-splits)
    from 1(3) obtain y2 ys2 where Y2: coeffs Y2 = y2 # coeffs ys2
      by (cases Y2, auto simp: cCons-def split: if-splits)
    from 1(2) 1(3) have [transfer-rule]: R x1 x2 R y1 y2
```

```
    and ∗: list-all2 R xs1 (coeffs xs2) list-all2 R ys1 (coeffs ys2) unfolding X2
Y2 by auto
   note [transfer-rule] = 1(1)[OF ∗]
   show ?case unfolding X2 Y2 by simp transfer-prover
 next
   case (2 xs1 xs2 ys2)
   thus ?case by (cases coeffs xs2, auto)
 next
   case (3 xs2 y1 ys1 Y2)
   thus ?case by (cases Y2, auto simp: cCons-def)
 qed
qed


lemma poly-rel-uminus[transfer-rule]: (poly-rel ===> poly-rel) (uminus-poly-i ops)
Groups.uminus
proof (intro rel-funI)
 fix x y
 assume poly-rel x y
 hence [transfer-rule]: list-all2 R x (coeffs y) unfolding poly-rel-def .
 show poly-rel (uminus-poly-i ops x) (−y)
   unfolding poly-rel-def coeffs-uminus uminus-poly-i-def by transfer-prover
qed


lemma poly-rel-minus[transfer-rule]: (poly-rel ===> poly-rel ===> poly-rel) (minus-poly-i
ops) (−)
proof (intro rel-funI)
 fix x1 y1 x2 y2
 assume poly-rel x1 x2 and poly-rel y1 y2
 thus poly-rel (minus-poly-i ops x1 y1) (x2 − y2)
   unfolding diff-conv-add-uminus
   unfolding poly-rel-def coeffs-eq-iff coeffs-plus-eq-plus-coeffs coeffs-uminus
 proof (induct x1 y1 arbitrary: x2 y2 rule: minus-poly-i.induct)
   case (1 x1 xs1 y1 ys1 X2 Y2)
   from 1(2) obtain x2 xs2 where X2: coeffs X2 = x2 # coeffs xs2
     by (cases X2, auto simp: cCons-def split: if-splits)
   from 1(3) obtain y2 ys2 where Y2: coeffs Y2 = y2 # coeffs ys2
     by (cases Y2, auto simp: cCons-def split: if-splits)
   from 1(2) 1(3) have [transfer-rule]: R x1 x2 R y1 y2
     and ∗: list-all2 R xs1 (coeffs xs2) list-all2 R ys1 (coeffs ys2) unfolding X2
Y2 by auto
   note [transfer-rule] = 1(1)[OF ∗]
   show ?case unfolding X2 Y2 by simp transfer-prover
 next
   case (2 xs1 xs2 ys2)
   thus ?case by (cases coeffs xs2, auto)
 next
   case (3 xs2 y1 ys1 Y2)
```

from *3(1)* **have** *id0*: *coeffs ys1 = coeffs 0* **by** (*cases ys1*, *auto*)
have *id1*: *minus-poly-i ops* [] (*xs2 # y1*) = *uminus-poly-i ops* (*xs2 # y1*) **by** *simp*
from *3(2)* **have** [*transfer-rule*]: *poly-rel* (*xs2 # y1*) *Y2* **unfolding** *poly-rel-def* **by** *simp*
show *?case* **unfolding** *id0 id1 coeffs-uminus*[*symmetric*] *coeffs-plus-eq-plus-coeffs*[*symmetric*] *poly-rel-def*[*symmetric*] **by** *simp transfer-prover*
qed
qed

lemma *poly-rel-smult*[*transfer-rule*]: (*R* ===> *poly-rel* ===> *poly-rel*) (*smult-i ops*) *smult*
unfolding *rel-fun-def poly-rel-def coeffs-smult' smult-i-def*
proof (*intro allI impI*, *goal-cases*)
case (*1 x y xs ys*)
note [*transfer-rule*] = *1*
show *?case* **by** *transfer-prover*
qed

lemma *poly-rel-coeffs*[*transfer-rule*]: (*poly-rel* ===> *list-all2 R*) (λ *x*. *x*) *coeffs*
unfolding *rel-fun-def poly-rel-def* **by** *auto*

lemma *poly-rel-poly-of-list*[*transfer-rule*]: (*list-all2 R* ===> *poly-rel*) (*poly-of-list-i ops*) *poly-of-list*
unfolding *rel-fun-def poly-of-list-i-def poly-rel-def poly-of-list-impl*
proof (*intro allI impI*, *goal-cases*)
case (*1 x y*)
note [*transfer-rule*] = *this*
show *?case* **by** *transfer-prover*
qed

lemma *poly-rel-monom-mult*[*transfer-rule*]:
((=) ===> *poly-rel* ===> *poly-rel*) (*monom-mult-i ops*) *monom-mult*
unfolding *rel-fun-def monom-mult-i-def poly-rel-def monom-mult-code Let-def*
proof (*auto*, *goal-cases*)
case (*1 x xs y*)
show *?case* **by** (*induct x*, *auto simp*: *1(3) zero*)
qed

declare *karatsuba-main-i.simps*[*simp del*]

lemma *list-rel-coeffs-minus-i*: **assumes** *list-all2 R x1 x2 list-all2 R y1 y2*
shows *list-all2 R* (*coeffs-minus-i ops x1 y1*) (*coeffs-minus x2 y2*)
proof −
note *simps* = *coeffs-minus-i.simps coeffs-minus.simps*
show *?thesis* **using** *assms*

```
  proof (induct x1 y1 arbitrary: x2 y2 rule: coeffs-minus-i.induct)
    case (1 x xs y ys)
   from 1(2−) obtain Y Ys where y2: y2 = Y # Ys unfolding list-all2-conv-all-nth
by (cases y2, auto)
    with 1(2−) have y: R y Y list-all2 R ys Ys by auto
   from 1(2−) obtain X Xs where x2: x2 = X # Xs unfolding list-all2-conv-all-nth
by (cases x2, auto)
    with 1(2−) have x: R x X list-all2 R xs Xs by auto
    from 1(1)[OF x(2) y(2)] x(1) y(1)
    show ?case unfolding x2 y2 simps using minus[unfolded rel-fun-def] by auto
  next
    case (3 y ys)
    from 3 have x2: x2 = [] by auto
   from 3 obtain Y Ys where y2: y2 = Y # Ys unfolding list-all2-conv-all-nth
by (cases y2, auto)
    obtain y1 where y1: y # ys = y1 by auto
     show ?case unfolding y2 simps x2 unfolding y2[symmetric] list-all2-map2
list-all2-map1
      using 3(2) unfolding y1 using uminus[unfolded rel-fun-def]
      unfolding list-all2-conv-all-nth by auto
  qed auto
qed


lemma poly-rel-karatsuba-main: list-all2 R x1 x2 ⟹ list-all2 R y1 y2 ⟹
  poly-rel (karatsuba-main-i ops x1 n y1 m) (karatsuba-main x2 n y2 m)
proof (induct n arbitrary: x1 y1 x2 y2 m rule: less-induct)
  case (less n f g F G m)
  note simp[simp] = karatsuba-main.simps[of F n G m] karatsuba-main-i.simps[of
ops f n g m]
  note IH = less(1)
  note rel[transfer-rule] = less(2−3)
  show ?case (is poly-rel ?lhs ?rhs)
 proof (cases (n ≤ karatsuba-lower-bound ∨ m ≤ karatsuba-lower-bound) = False)
    case False
    from False
   have lhs: ?lhs = foldr (λa p. plus-poly-i ops (smult-i ops a (poly-of-list-i ops f))
        (cCons-i ops zero p)) g [] by simp
    from False have rhs: ?rhs = foldr (λa p. smult a (poly-of-list F) + pCons 0
p) G 0 by simp
    show ?thesis unfolding lhs rhs by transfer-prover
  next
    case True note ∗ = this
    let ?n2 = n div 2
    have ?n2 < n n − ?n2 < n using True unfolding karatsuba-lower-bound-def
by auto
    note IH = IH[OF this(1)] IH[OF this(2)]
    obtain f1 f0 where f: split-at ?n2 f = (f0,f1) by force
    obtain g1 g0 where g: split-at ?n2 g = (g0,g1) by force
```

obtain *F1 F0* **where** *F*: *split-at ?n2 F* = (*F0*,*F1*) **by** *force*
obtain *G1 G0* **where** *G*: *split-at ?n2 G* = (*G0*,*G1*) **by** *force*
from *rel f F* **have** *relf*[*transfer-rule*]: *list-all2 R f0 F0 list-all2 R f1 F1*
unfolding *split-at-def* **by** *auto*
from *rel g G* **have** *relg*[*transfer-rule*]: *list-all2 R g0 G0 list-all2 R g1 G1*
unfolding *split-at-def* **by** *auto*
show *?thesis*
proof (*cases ?n2* < *m*)
case *True*
obtain *p1 P1* **where** *p1*: *p1* = *karatsuba-main-i ops f1* (*n* − *n div 2*) *g1* (*m* − *n div 2*)
P1 = *karatsuba-main F1* (*n* − *n div 2*) *G1* (*m* − *n div 2*) **by** *auto*
obtain *p2 P2* **where** *p2*: *p2* = *karatsuba-main-i ops* (*coeffs-minus-i ops f1 f0*) (*n div 2*)
(*coeffs-minus-i ops g1 g0*) (*n div 2*)
P2 = *karatsuba-main* (*coeffs-minus F1 F0*) (*n div 2*)
(*coeffs-minus G1 G0*) (*n div 2*) **by** *auto*
obtain *p3 P3* **where** *p3*: *p3* = *karatsuba-main-i ops f0* (*n div 2*) *g0* (*n div 2*)
P3 = *karatsuba-main F0* (*n div 2*) *G0* (*n div 2*) **by** *auto*
from ∗ *True* **have** *lhs*: *?lhs* = *plus-poly-i ops* (*monom-mult-i ops* (*n div 2* + *n div 2*) *p1*)
(*plus-poly-i ops*
(*monom-mult-i ops* (*n div 2*)
(*plus-poly-i ops* (*minus-poly-i ops p1 p2*) *p3*)) *p3*)
unfolding *simp Let-def f g split p1 p2 p3* **by** *auto*
have [*transfer-rule*]: *poly-rel p1 P1* **using** *IH*(*2*)[*OF relf*(*2*) *relg*(*2*)] **unfolding** *p1* .
have [*transfer-rule*]: *poly-rel p3 P3* **using** *IH*(*1*)[*OF relf*(*1*) *relg*(*1*)] **unfolding** *p3* .
have [*transfer-rule*]: *poly-rel p2 P2* **unfolding** *p2*
by (*rule IH*(*1*)[*OF list-rel-coeffs-minus-i list-rel-coeffs-minus-i*], *insert relf relg*)
from *True* ∗ **have** *rhs*: *?rhs* = *monom-mult* (*n div 2* + *n div 2*) *P1* +
(*monom-mult* (*n div 2*) (*P1* − *P2* + *P3*) + *P3*)
unfolding *simp Let-def F G split p1 p2 p3* **by** *auto*
show *?thesis* **unfolding** *lhs rhs* **by** *transfer-prover*
next
case *False*
obtain *p1 P1* **where** *p1*: *p1* = *karatsuba-main-i ops f1* (*n* − *n div 2*) *g m*
P1 = *karatsuba-main F1* (*n* − *n div 2*) *G m* **by** *auto*
obtain *p2 P2* **where** *p2*: *p2* = *karatsuba-main-i ops f0* (*n div 2*) *g m*
P2 = *karatsuba-main F0* (*n div 2*) *G m* **by** *auto*
from ∗ *False* **have** *lhs*: *?lhs* = *plus-poly-i ops* (*monom-mult-i ops* (*n div 2*) *p1*) *p2*
unfolding *simp Let-def f split p1 p2* **by** *auto*
from ∗ *False* **have** *rhs*: *?rhs* = *monom-mult* (*n div 2*) *P1* + *P2*
unfolding *simp Let-def F split p1 p2* **by** *auto*
have [*transfer-rule*]: *poly-rel p1 P1* **using** *IH*(*2*)[*OF relf*(*2*) *rel*(*2*)] **unfolding**

```
p1 .
    have [transfer-rule]: poly-rel p2 P2 using IH(1)[OF relf(1) rel(2)] unfolding
p2 .
     show ?thesis unfolding lhs rhs by transfer-prover
   qed
  qed
qed


lemma poly-rel-times[transfer-rule]: (poly-rel ===> poly-rel ===> poly-rel) (times-poly-i
ops) ((*))
proof (intro rel-funI)
  fix x1 y1 x2 y2
  assume x12[transfer-rule]: poly-rel x1 x2 and y12 [transfer-rule]: poly-rel y1 y2
 hence X12[transfer-rule]: list-all2 R x1 (coeffs x2) and Y12[transfer-rule]: list-all2
R y1 (coeffs y2)
    unfolding poly-rel-def by auto
  hence len: length (coeffs x2) = length x1 length (coeffs y2) = length y1
    unfolding list-all2-conv-all-nth by auto
 let ?cond1 = length x1 ≤ karatsuba-lower-bound ∨ length y1 ≤ karatsuba-lower-bound

  let ?cond2 = length x1 ≤ length y1
  note d = karatsuba-mult-poly[symmetric] karatsuba-mult-poly-def Let-def
      times-poly-i-def len if-True if-False
  consider (TT) ?cond1 = True ?cond2 = True | (TF) ?cond1 = True ?cond2
= False
    | (FT) ?cond1 = False ?cond2 = True | (FF) ?cond1 = False ?cond2 = False
by auto
  thus poly-rel (times-poly-i ops x1 y1) (x2 * y2)
  proof (cases)
    case TT
    show ?thesis unfolding d TT
     unfolding poly-rel-def coeffs-eq-iff times-poly-def times-poly-i-def fold-coeffs-def
      by transfer-prover
  next
    case TF
    show ?thesis unfolding d TF
     unfolding poly-rel-def coeffs-eq-iff times-poly-def times-poly-i-def fold-coeffs-def
      by transfer-prover
  next
    case FT
    show ?thesis unfolding d FT
      by (rule poly-rel-karatsuba-main[OF Y12 X12])
  next
    case FF
    show ?thesis unfolding d FF
      by (rule poly-rel-karatsuba-main[OF X12 Y12])
  qed
qed
```

lemma *poly-rel-coeff*[*transfer-rule*]: (*poly-rel* ===> (=) ===> *R*) (*coeff-i ops*) *coeff*
unfolding *poly-rel-def rel-fun-def coeff-i-def nth-default-coeffs-eq*[*symmetric*]
proof (*intro allI impI*, *clarify*)
fix *x y n*
assume [*transfer-rule*]: *list-all2 R x* (*coeffs y*)
show *R* (*nth-default zero x n*) (*nth-default 0* (*coeffs y*) *n*) **by** *transfer-prover*
qed

lemma *poly-rel-degree*[*transfer-rule*]: (*poly-rel* ===> (=)) *degree-i degree*
unfolding *poly-rel-def rel-fun-def degree-i-def degree-eq-length-coeffs*
by (*simp add*: *list-all2-lengthD*)

lemma *lead-coeff-i-def'*: *lead-coeff-i ops x* = (*coeff-i ops*) *x* (*degree-i x*)
unfolding *lead-coeff-i-def degree-i-def coeff-i-def*
proof (*cases x*, *auto*, *goal-cases*)
case (*1 a xs*)
hence *id*: *last xs* = *last* (*a* # *xs*) **by** *auto*
show *?case* **unfolding** *id* **by** (*subst last-conv-nth-default*, *auto*)
qed

lemma *poly-rel-lead-coeff*[*transfer-rule*]: (*poly-rel* ===> *R*) (*lead-coeff-i ops*) *lead-coeff*
unfolding *lead-coeff-i-def'* [*abs-def*] **by** *transfer-prover*

lemma *poly-rel-minus-poly-rev-list*[*transfer-rule*]:
(*list-all2 R* ===> *list-all2 R* ===> *list-all2 R*) (*minus-poly-rev-list-i ops*) *minus-poly-rev-list*
proof (*intro rel-funI*, *goal-cases*)
case (*1 x1 x2 y1 y2*)
thus *?case*
proof (*induct x1 y1 arbitrary*: *x2 y2 rule*: *minus-poly-rev-list-i.induct*)
case (*1 x1 xs1 y1 ys1 X2 Y2*)
from *1*(*2*) **obtain** *x2 xs2* **where** *X2*: *X2* = *x2* # *xs2* **by** (*cases X2*, *auto*)
from *1*(*3*) **obtain** *y2 ys2* **where** *Y2*: *Y2* = *y2* # *ys2* **by** (*cases Y2*, *auto*)
from *1*(*2*) *1*(*3*) **have** [*transfer-rule*]: *R x1 x2 R y1 y2*
and ∗: *list-all2 R xs1 xs2 list-all2 R ys1 ys2* **unfolding** *X2 Y2* **by** *auto*
note [*transfer-rule*] = *1*(*1*)[*OF* ∗]
show *?case* **unfolding** *X2 Y2* **by** (*simp*, *intro conjI*, *transfer-prover*+)
next
case (*2 xs1 xs2 ys2*)
thus *?case* **by** (*cases xs2*, *auto*)
next
case (*3 xs2 y1 ys1 Y2*)
thus *?case* **by** (*cases Y2*, *auto*)

```
  qed
qed


lemma divmod-poly-one-main-i: assumes len: n ≤ length Y and rel: list-all2 R
x X list-all2 R y Y
    list-all2 R z Z and n: n = N
 shows rel-prod (list-all2 R) (list-all2 R) (divmod-poly-one-main-i ops x y z n)
    (divmod-poly-one-main-list X Y Z N)
   using len rel unfolding n
proof (induct N arbitrary: x X y Y z Z)
  case (Suc n x X y Y z Z)
  from Suc(2,4) have [transfer-rule]: R (hd y) (hd Y) by (cases y; cases Y, auto)

  note [transfer-rule] = Suc(3−5)
  have id: ?case = (rel-prod (list-all2 R) (list-all2 R)
     (divmod-poly-one-main-i ops (cCons-i ops (hd y) x)
       (tl (if hd y = zero then y else minus-poly-rev-list-i ops y (map (times (hd y))
z))) z n)
     (divmod-poly-one-main-list (cCons (hd Y) X)
       (tl (if hd Y = 0 then Y else minus-poly-rev-list Y (map ((∗) (hd Y)) Z))) Z
n))
     by (simp add: Let-def)
  show ?case unfolding id
  proof (rule Suc(1), goal-cases)
    case 1
    show ?case using Suc(2) by simp
  qed (transfer-prover+)
qed simp


lemma mod-poly-one-main-i: assumes len: n ≤ length X and rel: list-all2 R x X
list-all2 R y Y
    and n: n = N
 shows list-all2 R (mod-poly-one-main-i ops x y n)
    (mod-poly-one-main-list X Y N)
   using len rel unfolding n
proof (induct N arbitrary: x X y Y)
  case (Suc n y Y z Z)
  from Suc(2,3) have [transfer-rule]: R (hd y) (hd Y) by (cases y; cases Y, auto)

  note [transfer-rule] = Suc(3−4)
  have id: ?case = (list-all2 R
     (mod-poly-one-main-i ops
       (tl (if hd y = zero then y else minus-poly-rev-list-i ops y (map (times (hd y))
z))) z n)
     (mod-poly-one-main-list
       (tl (if hd Y = 0 then Y else minus-poly-rev-list Y (map ((∗) (hd Y)) Z))) Z
n))
```

```
    by (simp add: Let-def)
  show ?case unfolding id
  proof (rule Suc(1), goal-cases)
    case 1
    show ?case using Suc(2) by simp
  qed (transfer-prover+)
qed simp

lemma poly-rel-dvd[transfer-rule]: (poly-rel ===> poly-rel ===> (=)) (dvd-poly-i
ops) (dvd)
  unfolding dvd-poly-i-def[abs-def] dvd-def[abs-def]
  by (transfer-prover-start, transfer-step+, auto)

lemma poly-rel-monic[transfer-rule]: (poly-rel ===> (=)) (monic-i ops) monic
  unfolding monic-i-def lead-coeff-i-def' by transfer-prover

lemma poly-rel-pdivmod-monic: assumes mon: monic Y
  and x: poly-rel x X and y: poly-rel y Y
  shows rel-prod poly-rel poly-rel (pdivmod-monic-i ops x y) (pdivmod-monic X Y)
proof −
  note [transfer-rule] = x y
  note listall = this[unfolded poly-rel-def]
  note defs = pdivmod-monic-def pdivmod-monic-i-def Let-def
  from mon obtain k where len: length (coeffs Y) = Suc k unfolding poly-rel-def
list-all2-iff
    by (cases coeffs Y, auto)
  have [transfer-rule]:
    rel-prod (list-all2 R) (list-all2 R)
      (divmod-poly-one-main-i ops [] (rev x) (rev y) (1 + length x − length y))
       (divmod-poly-one-main-list [] (rev (coeffs X)) (rev (coeffs Y)) (1 + length
(coeffs X) − length (coeffs Y)))
    by (rule divmod-poly-one-main-i, insert x y listall, auto, auto simp: poly-rel-def
list-all2-iff len)
  show ?thesis unfolding defs by transfer-prover
qed

lemma ring-ops-poly: ring-ops (poly-ops ops) poly-rel
  by (unfold-locales, auto simp: poly-ops-def
  bi-unique-poly-rel
  right-total-poly-rel
  poly-rel-times
  poly-rel-zero
  poly-rel-one
  poly-rel-minus
  poly-rel-uminus
  poly-rel-plus
  poly-rel-eq
  Domainp-is-poly)
end
```

context *idom-ops*
begin

lemma *poly-rel-pderiv* [*transfer-rule*]: (*poly-rel* ===> *poly-rel*) (*pderiv-i ops*) *pderiv*
proof (*intro rel-funI*, *unfold poly-rel-def coeffs-pderiv-code pderiv-i-def pderiv-coeffs-def*)
fix *xs xs′*
assume *list-all2 R xs* (*coeffs xs′*)
then obtain *ys ys′ y y′* **where** *id*: *tl xs* = *ys tl* (*coeffs xs′*) = *ys′ one* = *y 1* = *y′* **and**
R: *list-all2 R ys ys′ R y y′*
by (*cases xs*; *cases coeffs xs′*; *auto simp*: *one*)
show *list-all2 R* (*pderiv-main-i ops one* (*tl xs*))
(*pderiv-coeffs-code 1* (*tl* (*coeffs xs′*)))
unfolding *id* **using** *R*
proof (*induct ys ys′ arbitrary*: *y y′ rule*: *list-all2-induct*)
case (*Cons x xs x′ xs′ y y′*)
note [*transfer-rule*] = *Cons*(*1*,*2*,*4*)
have *R* (*plus y one*) (*y′* + *1*) **by** *transfer-prover*
note [*transfer-rule*] = *Cons*(*3*)[*OF this*]
show *?case* **by** (*simp*, *transfer-prover*)
qed *simp*
qed

lemma *poly-rel-irreducible*[*transfer-rule*]: (*poly-rel* ===> (=)) (*irreducible-i ops*) *irreducible*$_d$
unfolding *irreducible-i-def*[*abs-def*] *irreducible*$_d$*-def*[*abs-def*]
by (*transfer-prover-start*, *transfer-step*+, *auto*)

lemma *idom-ops-poly*: *idom-ops* (*poly-ops ops*) *poly-rel*
using *ring-ops-poly* **unfolding** *ring-ops-def idom-ops-def* **by** *auto*
end

context *idom-divide-ops*
begin

lemma *poly-rel-sdiv*[*transfer-rule*]: (*poly-rel* ===> *R* ===> *poly-rel*) (*sdiv-i ops*) *sdiv-poly*
unfolding *rel-fun-def poly-rel-def coeffs-sdiv sdiv-i-def*
proof (*intro allI impI*, *goal-cases*)
case (*1 x y xs ys*)
note [*transfer-rule*] = *1*
show *?case* **by** *transfer-prover*
qed
end

context *field-ops*
begin

```
lemma poly-rel-div[transfer-rule]: (poly-rel ===> poly-rel ===> poly-rel)
  (div-field-poly-i ops) (div)
proof (intro rel-funI, goal-cases)
  case (1 x X y Y)
  note [transfer-rule] = this
  note listall = this[unfolded poly-rel-def]
  note defs = div-field-poly-impl div-field-poly-impl-def div-field-poly-i-def Let-def
  show ?case
  proof (cases y = [])
    case True
    with 1(2) have nil: coeffs Y = [] unfolding poly-rel-def by auto
    show ?thesis unfolding defs True nil poly-rel-def by auto
  next
    case False
    from append-butlast-last-id[OF False] obtain ys yl where y: y = ys @ [yl] by
metis
    from False listall have coeffs Y ≠ [] by auto
    from append-butlast-last-id[OF this] obtain Ys Yl where Y: coeffs Y = Ys
@ [Yl] by metis
    from listall have [transfer-rule]: R yl Yl by (simp add: y Y)
    have id: last (coeffs Y) = Yl last (y) = yl
      ⋀ t e. (if y = [] then t else e) = e
      ⋀ t e. (if coeffs Y = [] then t else e) = e unfolding y Y by auto
    have [transfer-rule]: (rel-prod (list-all2 R) (list-all2 R))
      (divmod-poly-one-main-i ops [] (rev x) (rev (map (times (inverse yl)) y))
        (1 + length x − length y))
      (divmod-poly-one-main-list [] (rev (coeffs X))
              (rev (map ((*) (Fields.inverse Yl)) (coeffs Y)))
              (1 + length (coeffs X) − length (coeffs Y)))
    proof (rule divmod-poly-one-main-i, goal-cases)
      case 5
      from listall show ?case by (simp add: list-all2-lengthD)
    next
      case 1
      from listall show ?case by (simp add: list-all2-lengthD Y)
    qed transfer-prover+
    show ?thesis unfolding defs id by transfer-prover
  qed
qed


lemma poly-rel-mod[transfer-rule]: (poly-rel ===> poly-rel ===> poly-rel)
  (mod-field-poly-i ops) (mod)
proof (intro rel-funI, goal-cases)
  case (1 x X y Y)
  note [transfer-rule] = this
  note listall = this[unfolded poly-rel-def]
  note defs = mod-poly-code mod-field-poly-i-def Let-def
```

```
  show ?case
  proof (cases y = [])
    case True
    with 1(2) have nil: coeffs Y = [] unfolding poly-rel-def by auto
    show ?thesis unfolding defs True nil poly-rel-def by (simp add: listall)
  next
    case False
    from append-butlast-last-id[OF False] obtain ys yl where y: y = ys @ [yl] by
metis
    from False listall have coeffs Y ≠ [] by auto
    from append-butlast-last-id[OF this] obtain Ys Yl where Y: coeffs Y = Ys
@ [Yl] by metis
    from listall have [transfer-rule]: R yl Yl by (simp add: y Y)
    have id: last (coeffs Y) = Yl last (y) = yl
      ⋀ t e. (if y = [] then t else e) = e
      ⋀ t e. (if coeffs Y = [] then t else e) = e unfolding y Y by auto
    have [transfer-rule]: list-all2 R
      (mod-poly-one-main-i ops (rev x) (rev (map (times (inverse yl)) y))
        (1 + length x − length y))
      (mod-poly-one-main-list (rev (coeffs X))
              (rev (map ((*) (Fields.inverse Yl)) (coeffs Y)))
              (1 + length (coeffs X) − length (coeffs Y)))
    proof (rule mod-poly-one-main-i, goal-cases)
      case 4
      from listall show ?case by (simp add: list-all2-lengthD)
    next
      case 1
      from listall show ?case by (simp add: list-all2-lengthD Y)
    qed transfer-prover+
    show ?thesis unfolding defs id by transfer-prover
  qed
qed


lemma poly-rel-normalize [transfer-rule]: (poly-rel ===> poly-rel)
  (normalize-poly-i ops) Rings.normalize
  unfolding normalize-poly-old-def normalize-poly-i-def lead-coeff-i-def'
  by transfer-prover


lemma poly-rel-unit-factor [transfer-rule]: (poly-rel ===> poly-rel)
  (unit-factor-poly-i ops) Rings.unit-factor
  unfolding unit-factor-poly-def unit-factor-poly-i-def lead-coeff-i-def'
  unfolding monom-0 by transfer-prover

lemma idom-divide-ops-poly: idom-divide-ops (poly-ops ops) poly-rel
proof −
  interpret poly: idom-ops poly-ops ops poly-rel by (rule idom-ops-poly)
  show ?thesis
```

```
    by (unfold-locales, simp add: poly-rel-div poly-ops-def)
qed

lemma euclidean-ring-ops-poly: euclidean-ring-ops (poly-ops ops) poly-rel
proof −
  interpret poly: idom-ops poly-ops ops poly-rel by (rule idom-ops-poly)
  have id: arith-ops-record.normalize (poly-ops ops) = normalize-poly-i ops
    arith-ops-record.unit-factor (poly-ops ops) = unit-factor-poly-i ops
    unfolding poly-ops-def by simp-all
  show ?thesis
    by (unfold-locales, simp add: poly-rel-mod poly-ops-def, unfold id,
      simp add: poly-rel-normalize, insert poly-rel-div poly-rel-unit-factor,
      auto simp: poly-ops-def)
qed


lemma poly-rel-gcd [transfer-rule]: (poly-rel ===> poly-rel ===> poly-rel) (gcd-poly-i
ops) gcd
proof −
 interpret poly: euclidean-ring-ops poly-ops ops poly-rel by (rule euclidean-ring-ops-poly)
  show ?thesis using poly.gcd-eucl-i unfolding gcd-poly-i-def gcd-eucl .
qed


lemma poly-rel-euclid-ext [transfer-rule]: (poly-rel ===> poly-rel ===>
  rel-prod (rel-prod poly-rel poly-rel) poly-rel) (euclid-ext-poly-i ops) euclid-ext
proof −
 interpret poly: euclidean-ring-ops poly-ops ops poly-rel by (rule euclidean-ring-ops-poly)
  show ?thesis using poly.euclid-ext-i unfolding euclid-ext-poly-i-def .
qed

end



context ring-ops
begin
notepad
begin
  fix xs x ys y
  assume [transfer-rule]: poly-rel xs x poly-rel ys y
  have x * y = y * x by simp
  from this[untransferred]
  have times-poly-i ops xs ys = times-poly-i ops ys xs .
end
end
end
```

5.5.3 Over a Finite Field

theory *Poly-Mod-Finite-Field-Record-Based*
imports
Poly-Mod-Finite-Field
Finite-Field-Record-Based
Polynomial-Record-Based
begin

locale *arith-ops-record* = *arith-ops ops* + *poly-mod m* **for** *ops* :: *'i arith-ops-record*
and *m* :: *int*
begin
definition *M-rel-i* :: *'i* ⇒ *int* ⇒ *bool* **where**
M-rel-i f F = (*arith-ops-record.to-int ops f* = *M F*)

definition *Mp-rel-i* :: *'i list* ⇒ *int poly* ⇒ *bool* **where**
Mp-rel-i f F = (*map* (*arith-ops-record.to-int ops*) *f* = *coeffs* (*Mp F*))

lemma *Mp-rel-i-Mp*[*simp*]: *Mp-rel-i f* (*Mp F*) = *Mp-rel-i f F* **unfolding** *Mp-rel-i-def*
by *auto*

lemma *Mp-rel-i-Mp-to-int-poly-i*: *Mp-rel-i f F* ⟹ *Mp* (*to-int-poly-i ops f*) = *to-int-poly-i ops f*
unfolding *Mp-rel-i-def to-int-poly-i-def* **by** *simp*
end

locale *mod-ring-gen* = *ring-ops ff-ops R* **for** *ff-ops* :: *'i arith-ops-record* **and**
R :: *'i* ⇒ *'a* :: *nontriv mod-ring* ⇒ *bool* +
fixes *p* :: *int*
assumes *p*: *p* = *int CARD*(*'a*)
and *of-int*: 0 ≤ *x* ⟹ *x* < *p* ⟹ *R* (*arith-ops-record.of-int ff-ops x*) (*of-int x*)
and *to-int*: *R y z* ⟹ *arith-ops-record.to-int ff-ops y* = *to-int-mod-ring z*
and *to-int'*: 0 ≤ *arith-ops-record.to-int ff-ops y* ⟹ *arith-ops-record.to-int ff-ops y* < *p* ⟹
R y (*of-int* (*arith-ops-record.to-int ff-ops y*))
begin

lemma *nat-p*: *nat p* = *CARD*(*'a*) **unfolding** *p* **by** *simp*

sublocale *poly-mod-type p TYPE*(*'a*)
by (*unfold-locales*, *rule p*)

lemma *coeffs-to-int-poly*: *coeffs* (*to-int-poly* (*x* :: *'a mod-ring poly*)) = *map to-int-mod-ring* (*coeffs x*)
by (*rule coeffs-map-poly*, *auto*)

lemma *coeffs-of-int-poly*: *coeffs* (*of-int-poly* (*Mp x*) :: *'a mod-ring poly*) = *map of-int* (*coeffs* (*Mp x*))
apply (*rule coeffs-map-poly*)
by (*metis M-0 M-M Mp-coeff leading-coeff-0-iff of-int-hom.hom-zero to-int-mod-ring-of-int-M*)

lemma *to-int-poly-i*: **assumes** *poly-rel f g* **shows** *to-int-poly-i ff-ops f = to-int-poly g*
proof −
have ∗: *map* (*arith-ops-record.to-int ff-ops*) *f* = *coeffs* (*to-int-poly g*)
unfolding *coeffs-to-int-poly*
by (*rule nth-equalityI*, *insert assms*, *auto simp*: *list-all2-conv-all-nth poly-rel-def to-int*)
show *?thesis* **unfolding** *coeffs-eq-iff to-int-poly-i-def poly-of-list-def coeffs-Poly* ∗ *strip-while-coeffs*..
qed

lemma *poly-rel-of-int-poly*: **assumes** *id*: *f′* = *of-int-poly-i ff-ops* (*Mp f*) *f″* = *of-int-poly* (*Mp f*)
shows *poly-rel f′ f″* **unfolding** *id poly-rel-def*
unfolding *list-all2-conv-all-nth coeffs-of-int-poly of-int-poly-i-def length-map*
by (*rule conjI*[*OF refl*], *intro allI impI*, *simp add*: *nth-coeffs-coeff Mp-coeff M-def*, *rule of-int*,
insert p, *auto*)

sublocale *arith-ops-record ff-ops p* .

lemma *Mp-rel-iI*: *poly-rel f1 f2* ⟹ *MP-Rel f3 f2* ⟹ *Mp-rel-i f1 f3*
unfolding *Mp-rel-i-def MP-Rel-def poly-rel-def*
by (*auto simp add*: *list-all2-conv-all-nth to-int intro*: *nth-equalityI*)

lemma *M-rel-iI*: *R f1 f2* ⟹ *M-Rel f3 f2* ⟹ *M-rel-i f1 f3*
unfolding *M-rel-i-def M-Rel-def* **by** (*simp add*: *to-int*)

lemma *M-rel-iI′*: **assumes** *R f1 f2*
shows *M-rel-i f1* (*arith-ops-record.to-int ff-ops f1*)
by (*rule M-rel-iI*[*OF assms*], *simp add*: *to-int*[*OF assms*] *M-Rel-def M-to-int-mod-ring*)

lemma *Mp-rel-iI′*: **assumes** *poly-rel f1 f2*
shows *Mp-rel-i f1* (*to-int-poly-i ff-ops f1*)
proof (*rule Mp-rel-iI*[*OF assms*], *unfold to-int-poly-i*[*OF assms*])
show *MP-Rel* (*to-int-poly f2*) *f2* **unfolding** *MP-Rel-def* **by** (*simp add*: *Mp-to-int-poly*)
qed

lemma *M-rel-iD*: **assumes** *M-rel-i f1 f3*
shows
R f1 (*of-int* (*M f3*))
M-Rel f3 (*of-int* (*M f3*))
proof −
show *M-Rel f3* (*of-int* (*M f3*))
using *M-Rel-def to-int-mod-ring-of-int-M* **by** *auto*
from *assms* **show** *R f1* (*of-int* (*M f3*))
unfolding *M-rel-i-def*
by (*metis int-one-le-iff-zero-less leD linear m1 poly-mod.M-def pos-mod-sign*

pos-mod-bound to-int′)
qed

lemma *Mp-rel-iD*: **assumes** *Mp-rel-i f1 f3*
shows
poly-rel f1 (*of-int-poly* (*Mp f3*))
MP-Rel f3 (*of-int-poly* (*Mp f3*))
proof −
show *Rel*: *MP-Rel f3* (*of-int-poly* (*Mp f3*))
using *MP-Rel-def Mp-Mp Mp-f-representative* **by** *auto*
let *?ti* = *arith-ops-record.to-int ff-ops*
from *assms*[*unfolded Mp-rel-i-def*] **have**
∗: *coeffs* (*Mp f3*) = *map ?ti f1* **by** *auto*
{
fix *x*
assume *x* ∈ *set f1*
hence *?ti x* ∈ *set* (*map ?ti f1*) **by** *auto*
from *this*[*folded* ∗] **have** *?ti x* ∈ *range M*
by (*metis* (*no-types*, *lifting*) *MP-Rel-def M-to-int-mod-ring Rel coeffs-to-int-poly ex-map-conv range-eqI*)
hence *?ti x* ≥ *0 ?ti x* < *p* **unfolding** *M-def* **using** *m1* **by** *auto*
hence *R x* (*of-int* (*?ti x*))
by (*rule to-int′*)
}
thus *poly-rel f1* (*of-int-poly* (*Mp f3*)) **using** ∗
unfolding *poly-rel-def coeffs-of-int-poly*
by (*auto simp*: *list-all2-map2 list-all2-same*)
qed
end

locale *prime-field-gen* = *field-ops ff-ops R* + *mod-ring-gen ff-ops R p* **for** *ff-ops* :: *′i arith-ops-record* **and**
R :: *′i* ⇒ *′a* :: *prime-card mod-ring* ⇒ *bool* **and** *p* :: *int*
begin

sublocale *poly-mod-prime-type p TYPE*(*′a*)
by (*unfold-locales*, *rule p*)

end

lemma (**in** *mod-ring-locale*) *mod-ring-rel-of-int*:
0 ≤ *x* ⟹ *x* < *p* ⟹ *mod-ring-rel x* (*of-int x*)
unfolding *mod-ring-rel-def*
by (*transfer*, *auto simp*: *p*)

context *prime-field*
begin

lemma *prime-field-finite-field-ops-int*: *prime-field-gen* (*finite-field-ops-int p*) *mod-ring-rel p*
proof −
interpret *field-ops finite-field-ops-int p mod-ring-rel* **by** (*rule finite-field-ops-int*)
show *?thesis*
by (*unfold-locales*, *rule p*,
auto simp: *finite-field-ops-int-def p mod-ring-rel-def of-int-of-int-mod-ring*)
qed

lemma *prime-field-finite-field-ops-integer*: *prime-field-gen* (*finite-field-ops-integer* (*integer-of-int p*)) *mod-ring-rel-integer p*
proof −
interpret *field-ops finite-field-ops-integer* (*integer-of-int p*) *mod-ring-rel-integer* **by** (*rule finite-field-ops-integer*, *simp*)
have *pp*: *p* = *int-of-integer* (*integer-of-int p*) **by** *auto*
interpret *int*: *prime-field-gen finite-field-ops-int p mod-ring-rel*
by (*rule prime-field-finite-field-ops-int*)
show *?thesis*
by (*unfold-locales*, *rule p*, *auto simp*: *finite-field-ops-integer-def*
mod-ring-rel-integer-def[*OF pp*] *urel-integer-def*[*OF pp*] *mod-ring-rel-of-int*
int.to-int[*symmetric*] *finite-field-ops-int-def*)
qed

lemma *prime-field-finite-field-ops32*: **assumes** *small*: *p* ≤ *65535*
shows *prime-field-gen* (*finite-field-ops32* (*uint32-of-int p*)) *mod-ring-rel32 p*
proof −
let *?pp* = *uint32-of-int p*
have *ppp*: *p* = *int-of-uint32 ?pp*
by (*subst int-of-uint32-inv*, *insert small p2*, *auto*)
note ∗ = *ppp small*
interpret *field-ops finite-field-ops32 ?pp mod-ring-rel32*
by (*rule finite-field-ops32*, *insert* ∗)
interpret *int*: *prime-field-gen finite-field-ops-int p mod-ring-rel*
by (*rule prime-field-finite-field-ops-int*)
show *?thesis*
proof (*unfold-locales*, *rule p*, *auto simp*: *finite-field-ops32-def*)
fix *x*
assume *x*: *0* ≤ *x x* < *p*
from *int.of-int*[*OF this*] **have** *mod-ring-rel x* (*of-int x*) **by** (*simp add*: *finite-field-ops-int-def*)
thus *mod-ring-rel32* (*uint32-of-int x*) (*of-int x*) **unfolding** *mod-ring-rel32-def*[*OF* ∗]
by (*intro exI*[*of - x*], *auto simp*: *urel32-def*[*OF* ∗], *subst int-of-uint32-inv*, *insert* ∗ *x*, *auto*)
next
fix *y z*
assume *mod-ring-rel32 y z*
from *this*[*unfolded mod-ring-rel32-def*[*OF* ∗]] **obtain** *x* **where** *yx*: *urel32 y x* **and** *xz*: *mod-ring-rel x z* **by** *auto*

from *int.to-int*[*OF xz*] **have** *zx*: *to-int-mod-ring z* = *x* **by** (*simp add*: *finite-field-ops-int-def*)
show *int-of-uint32 y* = *to-int-mod-ring z* **unfolding** *zx* **using** *yx* **unfolding** *urel32-def*[*OF* ∗] **by** *simp*
next
fix *y*
show *0* ≤ *int-of-uint32 y* ⟹ *int-of-uint32 y* < *p* ⟹ *mod-ring-rel32 y* (*of-int* (*int-of-uint32 y*))
unfolding *mod-ring-rel32-def*[*OF* ∗] *urel32-def*[*OF* ∗]
by (*intro exI*[*of - int-of-uint32 y*], *auto simp*: *mod-ring-rel-of-int*)
qed
qed

lemma *prime-field-finite-field-ops64*: **assumes** *small*: *p* ≤ *4294967295*
shows *prime-field-gen* (*finite-field-ops64* (*uint64-of-int p*)) *mod-ring-rel64 p*
proof −
let *?pp* = *uint64-of-int p*
have *ppp*: *p* = *int-of-uint64 ?pp*
by (*subst int-of-uint64-inv*, *insert small p2*, *auto*)
note ∗ = *ppp small*
interpret *field-ops finite-field-ops64 ?pp mod-ring-rel64*
by (*rule finite-field-ops64*, *insert* ∗)
interpret *int*: *prime-field-gen finite-field-ops-int p mod-ring-rel*
by (*rule prime-field-finite-field-ops-int*)
show *?thesis*
proof (*unfold-locales*, *rule p*, *auto simp*: *finite-field-ops64-def*)
fix *x*
assume *x*: *0* ≤ *x x* < *p*
from *int.of-int*[*OF this*] **have** *mod-ring-rel x* (*of-int x*) **by** (*simp add*: *finite-field-ops-int-def*)
thus *mod-ring-rel64* (*uint64-of-int x*) (*of-int x*) **unfolding** *mod-ring-rel64-def*[*OF* ∗]
by (*intro exI*[*of - x*], *auto simp*: *urel64-def*[*OF* ∗], *subst int-of-uint64-inv*, *insert* ∗ *x*, *auto*)
next
fix *y z*
assume *mod-ring-rel64 y z*
from *this*[*unfolded mod-ring-rel64-def*[*OF* ∗]] **obtain** *x* **where** *yx*: *urel64 y x* **and** *xz*: *mod-ring-rel x z* **by** *auto*
from *int.to-int*[*OF xz*] **have** *zx*: *to-int-mod-ring z* = *x* **by** (*simp add*: *finite-field-ops-int-def*)
show *int-of-uint64 y* = *to-int-mod-ring z* **unfolding** *zx* **using** *yx* **unfolding** *urel64-def*[*OF* ∗] **by** *simp*
next
fix *y*
show *0* ≤ *int-of-uint64 y* ⟹ *int-of-uint64 y* < *p* ⟹ *mod-ring-rel64 y* (*of-int* (*int-of-uint64 y*))
unfolding *mod-ring-rel64-def*[*OF* ∗] *urel64-def*[*OF* ∗]
by (*intro exI*[*of - int-of-uint64 y*], *auto simp*: *mod-ring-rel-of-int*)

```
  qed
qed
end

context mod-ring-locale
begin
lemma mod-ring-finite-field-ops-int: mod-ring-gen (finite-field-ops-int p) mod-ring-rel
p
proof −
 interpret ring-ops finite-field-ops-int p mod-ring-rel by (rule ring-finite-field-ops-int)
  show ?thesis
    by (unfold-locales, rule p,
      auto simp: finite-field-ops-int-def p mod-ring-rel-def of-int-of-int-mod-ring)
qed

lemma mod-ring-finite-field-ops-integer: mod-ring-gen (finite-field-ops-integer (integer-of-int
p)) mod-ring-rel-integer p
proof −
  interpret ring-ops finite-field-ops-integer (integer-of-int p) mod-ring-rel-integer
by (rule ring-finite-field-ops-integer, simp)
  have pp: p = int-of-integer (integer-of-int p) by auto
  interpret int: mod-ring-gen finite-field-ops-int p mod-ring-rel
    by (rule mod-ring-finite-field-ops-int)
  show ?thesis
    by (unfold-locales, rule p, auto simp: finite-field-ops-integer-def
      mod-ring-rel-integer-def[OF pp] urel-integer-def[OF pp] mod-ring-rel-of-int
      int.to-int[symmetric] finite-field-ops-int-def)
qed


lemma mod-ring-finite-field-ops32: assumes small: p ≤ 65535
  shows mod-ring-gen (finite-field-ops32 (uint32-of-int p)) mod-ring-rel32 p
proof −
  let ?pp = uint32-of-int p
  have ppp: p = int-of-uint32 ?pp
    by (subst int-of-uint32-inv, insert small p2, auto)
  note ∗ = ppp small
  interpret ring-ops finite-field-ops32 ?pp mod-ring-rel32
    by (rule ring-finite-field-ops32, insert ∗)
  interpret int: mod-ring-gen finite-field-ops-int p mod-ring-rel
    by (rule mod-ring-finite-field-ops-int)
  show ?thesis
  proof (unfold-locales, rule p, auto simp: finite-field-ops32-def)
    fix x
    assume x: 0 ≤ x x < p
      from int.of-int[OF this] have mod-ring-rel x (of-int x) by (simp add: fi-
nite-field-ops-int-def)
   thus mod-ring-rel32 (uint32-of-int x) (of-int x) unfolding mod-ring-rel32-def[OF
∗]
```

by (*intro exI*[*of* - *x*], *auto simp*: *urel32-def*[*OF* ∗], *subst int-of-uint32-inv*, *insert* ∗ *x*, *auto*)
next
fix *y z*
assume *mod-ring-rel32 y z*
from *this*[*unfolded mod-ring-rel32-def*[*OF* ∗]] **obtain** *x* **where** *yx*: *urel32 y x* **and** *xz*: *mod-ring-rel x z* **by** *auto*
from *int.to-int*[*OF xz*] **have** *zx*: *to-int-mod-ring z* = *x* **by** (*simp add*: *finite-field-ops-int-def*)
show *int-of-uint32 y* = *to-int-mod-ring z* **unfolding** *zx* **using** *yx* **unfolding** *urel32-def*[*OF* ∗] **by** *simp*
next
fix *y*
show *0* ≤ *int-of-uint32 y* ⟹ *int-of-uint32 y* < *p* ⟹ *mod-ring-rel32 y* (*of-int* (*int-of-uint32 y*))
unfolding *mod-ring-rel32-def*[*OF* ∗] *urel32-def*[*OF* ∗]
by (*intro exI*[*of* - *int-of-uint32 y*], *auto simp*: *mod-ring-rel-of-int*)
qed
qed

lemma *mod-ring-finite-field-ops64*: **assumes** *small*: *p* ≤ *4294967295*
shows *mod-ring-gen* (*finite-field-ops64* (*uint64-of-int p*)) *mod-ring-rel64 p*
proof −
let *?pp* = *uint64-of-int p*
have *ppp*: *p* = *int-of-uint64 ?pp*
by (*subst int-of-uint64-inv*, *insert small p2*, *auto*)
note ∗ = *ppp small*
interpret *ring-ops finite-field-ops64 ?pp mod-ring-rel64*
by (*rule ring-finite-field-ops64*, *insert* ∗)
interpret *int*: *mod-ring-gen finite-field-ops-int p mod-ring-rel*
by (*rule mod-ring-finite-field-ops-int*)
show *?thesis*
proof (*unfold-locales*, *rule p*, *auto simp*: *finite-field-ops64-def*)
fix *x*
assume *x*: *0* ≤ *x x* < *p*
from *int.of-int*[*OF this*] **have** *mod-ring-rel x* (*of-int x*) **by** (*simp add*: *finite-field-ops-int-def*)
thus *mod-ring-rel64* (*uint64-of-int x*) (*of-int x*) **unfolding** *mod-ring-rel64-def*[*OF* ∗]
by (*intro exI*[*of* - *x*], *auto simp*: *urel64-def*[*OF* ∗], *subst int-of-uint64-inv*, *insert* ∗ *x*, *auto*)
next
fix *y z*
assume *mod-ring-rel64 y z*
from *this*[*unfolded mod-ring-rel64-def*[*OF* ∗]] **obtain** *x* **where** *yx*: *urel64 y x* **and** *xz*: *mod-ring-rel x z* **by** *auto*
from *int.to-int*[*OF xz*] **have** *zx*: *to-int-mod-ring z* = *x* **by** (*simp add*: *finite-field-ops-int-def*)
show *int-of-uint64 y* = *to-int-mod-ring z* **unfolding** *zx* **using** *yx* **unfolding**

urel64-def[*OF* ∗] **by** *simp*
next
fix *y*
show *0* ≤ *int-of-uint64 y* ⟹ *int-of-uint64 y* < *p* ⟹ *mod-ring-rel64 y* (*of-int* (*int-of-uint64 y*))
unfolding *mod-ring-rel64-def*[*OF* ∗] *urel64-def*[*OF* ∗]
by (*intro exI*[*of* - *int-of-uint64 y*], *auto simp*: *mod-ring-rel-of-int*)
qed
qed
end

end

5.6 Chinese Remainder Theorem for Polynomials

We prove the Chinese Remainder Theorem, and strengthen it by showing uniqueness

theory *Chinese-Remainder-Poly*
imports
HOL−Number-Theory.Residues
Polynomial-Factorization.Polynomial-Irreducibility
Polynomial-Interpolation.Missing-Polynomial
begin

lemma *cong-add-poly*:
[(*a*::*'b*::{*field-gcd*} *poly*) = *b*] (*mod m*) ⟹ [*c* = *d*] (*mod m*) ⟹ [*a* + *c* = *b* + *d*] (*mod m*)
by (*fact cong-add*)

lemma *cong-mult-poly*:
[(*a*::*'b*::{*field-gcd*} *poly*) = *b*] (*mod m*) ⟹ [*c* = *d*] (*mod m*) ⟹ [*a* ∗ *c* = *b* ∗ *d*] (*mod m*)
by (*fact cong-mult*)

lemma *cong-mult-self-poly*: [(*a*::*'b*::{*field-gcd*} *poly*) ∗ *m* = *0*] (*mod m*)
by (*fact cong-mult-self-right*)

lemma *cong-scalar2-poly*: [(*a*::*'b*::{*field-gcd*} *poly*)= *b*] (*mod m*) ⟹ [*k* ∗ *a* = *k* ∗ *b*] (*mod m*)
by (*fact cong-scalar-left*)

lemma *cong-sum-poly*:
(⋀*x*. *x* ∈ *A* ⟹ [((*f x*)::*'b*::{*field-gcd*} *poly*) = *g x*] (*mod m*)) ⟹
[(∑ *x*∈*A*. *f x*) = (∑ *x*∈*A*. *g x*)] (*mod m*)
by (*rule cong-sum*)

lemma *cong-iff-lin-poly*: ([(*a*::*'b*::{*field-gcd*} *poly*) = *b*] (*mod m*)) = (∃ *k*. *b* = *a* + *m* ∗ *k*)
using *cong-diff-iff-cong-0* [*of b a m*] **by** (*auto simp add*: *cong-0-iff dvd-def alge-*

bra-simps dest: *cong-sym*)

lemma *cong-solve-poly*: (*a*::'*b*::{*field-gcd*} *poly*) ≠ *0* ⟹ ∃ *x*. [*a* ∗ *x* = *gcd a n*] (*mod n*)
proof (*cases n* = *0*)
  case *True*
  note *n0*=*True*
  show *?thesis*
  proof (*cases monic a*)
    case *True*
    have *n*: *normalize a* = *a* **by** (*rule normalize-monic*[*OF True*])
    show *?thesis*
    by (*rule exI*[*of - 1*], *auto simp add*: *n0 n cong-def*)
  next
    case *False*
    show *?thesis*
    by (*auto simp add*: *True cong-def normalize-poly-old-def map-div-is-smult-inverse*)
        (*metis mult.right-neutral mult-smult-right*)
  qed
next
 case *False*
 note *n-not-0* = *False*
 show *?thesis*
   using *bezout-coefficients-fst-snd* [*of a n*, *symmetric*]
   by (*auto simp add*: *cong-iff-lin-poly mult.commute* [*of a*] *mult.commute* [*of n*])
qed

lemma *cong-solve-coprime-poly*:
assumes *coprime-an*:*coprime* (*a*::'*b*::{*field-gcd*} *poly*) *n*
shows ∃ *x*. [*a* ∗ *x* = *1*] (*mod n*)
proof (*cases a* = *0*)
  case *True*
  show *?thesis* **unfolding** *cong-def*
    using *True coprime-an* **by** *auto*
next
  case *False*
  show *?thesis*
    using *coprime-an cong-solve-poly*[*OF False*, *of n*]
    unfolding *cong-def*
    by *presburger*
qed

lemma *cong-dvd-modulus-poly*:
  [*x* = *y*] (*mod m*) ⟹ *n dvd m* ⟹ [*x* = *y*] (*mod n*) **for** *x y* :: '*b*::{*field-gcd*} *poly*
  by (*auto simp add*: *cong-iff-lin-poly elim*!: *dvdE*)

lemma *chinese-remainder-aux-poly*:
  fixes *A* :: '*a set*

```
    and m :: 'a ⇒ 'b::{field-gcd} poly
  assumes fin: finite A
    and cop: ∀ i ∈ A. (∀ j ∈ A. i ≠ j ⟶ coprime (m i) (m j))
  shows ∃ b. (∀ i ∈ A. [b i = 1] (mod m i) ∧ [b i = 0] (mod (∏ j ∈ A − {i}. m
j)))
proof (rule finite-set-choice, rule fin, rule ballI)
  fix i
  assume i : A
  with cop have coprime (∏ j ∈ A − {i}. m j) (m i)
    by (auto intro: prod-coprime-left)
  then have ∃ x. [(∏ j ∈ A − {i}. m j) * x = 1] (mod m i)
    by (elim cong-solve-coprime-poly)
  then obtain x where [(∏ j ∈ A − {i}. m j) * x = 1] (mod m i)
    by auto
  moreover have [(∏ j ∈ A − {i}. m j) * x = 0]
    (mod (∏ j ∈ A − {i}. m j))
    by (subst mult.commute, rule cong-mult-self-poly)
  ultimately show ∃ a. [a = 1] (mod m i) ∧ [a = 0]
      (mod prod m (A − {i}))
    by blast
qed



lemma chinese-remainder-poly:
  fixes A :: 'a set
    and m :: 'a ⇒ 'b::{field-gcd} poly
    and u :: 'a ⇒ 'b poly
  assumes fin: finite A
    and cop: ∀ i∈A. (∀ j∈A. i ≠ j ⟶ coprime (m i) (m j))
  shows ∃ x. (∀ i∈A. [x = u i] (mod m i))
proof −
  from chinese-remainder-aux-poly [OF fin cop] obtain b where
    bprop: ∀ i∈A. [b i = 1] (mod m i) ∧
      [b i = 0] (mod (∏ j ∈ A − {i}. m j))
    by blast
  let ?x = ∑ i∈A. (u i) * (b i)
  show ?thesis
  proof (rule exI, clarify)
    fix i
    assume a: i : A
    show [?x = u i] (mod m i)
    proof −
      from fin a have ?x = (∑ j ∈ {i}. u j * b j) +
          (∑ j ∈ A − {i}. u j * b j)
        by (subst sum.union-disjoint [symmetric], auto intro: sum.cong)
      then have [?x = u i * b i + (∑ j ∈ A − {i}. u j * b j)] (mod m i)
        unfolding cong-def
        by auto
```

also have $[u\ i * b\ i + (\sum j \in A - \{i\}.\ u\ j * b\ j) =$
$u\ i * 1 + (\sum j \in A - \{i\}.\ u\ j * 0)]\ (mod\ m\ i)$
apply (*rule cong-add-poly*)
apply (*rule cong-scalar2-poly*)
using *bprop a* **apply** *blast*
apply (*rule cong-sum*)
apply (*rule cong-scalar2-poly*)
using *bprop* **apply** *auto*
apply (*rule cong-dvd-modulus-poly*)
apply (*drule* (*1*) *bspec*)
apply (*erule conjE*)
apply *assumption*
apply *rule*
using *fin a* **apply** *auto*
done
thus *?thesis*
by (*metis* (*no-types*, *lifting*) *a add.right-neutral fin mult-cancel-left1 mult-cancel-right1*

sum.not-neutral-contains-not-neutral sum.remove)
qed
qed
qed

lemma *cong-trans-poly*:
$[(a::'b::\{field\text{-}gcd\}\ poly) = b]\ (mod\ m) \Longrightarrow [b = c]\ (mod\ m) \Longrightarrow [a = c]\ (mod\ m)$
by (*fact cong-trans*)

lemma *cong-mod-poly*: $(n::'b::\{field\text{-}gcd\}\ poly) \sim= 0 \Longrightarrow [a\ mod\ n = a]\ (mod\ n)$
by *auto*

lemma *cong-sym-poly*: $[(a::'b::\{field\text{-}gcd\}\ poly) = b]\ (mod\ m) \Longrightarrow [b = a]\ (mod\ m)$
by (*fact cong-sym*)

lemma *cong-1-poly*: $[(a::'b::\{field\text{-}gcd\}\ poly) = b]\ (mod\ 1)$
by (*fact cong-1*)

lemma *coprime-cong-mult-poly*:
assumes $[(a::'b::\{field\text{-}gcd\}\ poly) = b]\ (mod\ m)$ **and** $[a = b]\ (mod\ n)$ **and** *coprime m n*
shows $[a = b]\ (mod\ m * n)$
using *divides-mult assms*
by (*metis* (*no-types*, *opaque-lifting*) *cong-dvd-modulus-poly cong-iff-lin-poly dvd-mult2 dvd-refl minus-add-cancel mult.right-neutral*)

lemma *coprime-cong-prod-poly*:

$(\forall\, i{\in}A.\ (\forall\, j{\in}A.\ i \neq j \longrightarrow \mathit{coprime}\ (m\ i)\ (m\ j))) \Longrightarrow$
$(\forall\, i{\in}A.\ [(x::'b::\{\mathit{field\text{-}gcd}\}\ \mathit{poly}) = y]\ (\mathit{mod}\ m\ i)) \Longrightarrow$
$[x = y]\ (\mathit{mod}\ (\prod i{\in}A.\ m\ i))$
apply (*induct A rule*: *infinite-finite-induct*)
apply *auto*
apply (*metis coprime-cong-mult-poly prod-coprime-right*)
done

lemma *cong-less-modulus-unique-poly*:
$[(x::'b::\{\mathit{field\text{-}gcd}\}\ \mathit{poly}) = y]\ (\mathit{mod}\ m) \Longrightarrow \mathit{degree}\ x < \mathit{degree}\ m \Longrightarrow \mathit{degree}\ y < \mathit{degree}\ m \Longrightarrow x = y$
by (*simp add*: *cong-def mod-poly-less*)

lemma *chinese-remainder-unique-poly*:
fixes $A :: 'a\ \mathit{set}$
and $m :: 'a \Rightarrow 'b::\{\mathit{field\text{-}gcd}\}\ \mathit{poly}$
and $u :: 'a \Rightarrow 'b\ \mathit{poly}$
assumes *nz*: $\forall\, i{\in}A.\ (m\ i) \neq 0$
and *cop*: $\forall\, i{\in}A.\ (\forall\, j{\in}A.\ i \neq j \longrightarrow \mathit{coprime}\ (m\ i)\ (m\ j))$

and *not-constant*: $0 < \mathit{degree}\ (\mathit{prod}\ m\ A)$
shows $\exists!x.\ \mathit{degree}\ x < (\sum i{\in}A.\ \mathit{degree}\ (m\ i)) \wedge (\forall\, i{\in}A.\ [x = u\ i]\ (\mathit{mod}\ m\ i))$
proof −
from *not-constant* **have** *fin*: *finite A*
by (*metis degree-1 gr-implies-not0 prod.infinite*)
from *chinese-remainder-poly* [*OF fin cop*]
obtain *y* **where** *one*: $(\forall\, i{\in}A.\ [y = u\ i]\ (\mathit{mod}\ m\ i))$
by *blast*
let *?x* = $y\ \mathit{mod}\ (\prod i{\in}A.\ m\ i)$
have *degree-prod-sum*: $\mathit{degree}\ (\mathit{prod}\ m\ A) = (\sum i{\in}A.\ \mathit{degree}\ (m\ i))$
by (*rule degree-prod-eq-sum-degree*[*OF nz*])
from *fin nz* **have** *prodnz*: $(\prod i{\in}A.\ (m\ i)) \neq 0$
by *auto*

have *less*: $\mathit{degree}\ ?x < (\sum i{\in}A.\ \mathit{degree}\ (m\ i))$
unfolding *degree-prod-sum*[*symmetric*]
using *degree-mod-less*[*OF prodnz*, *of y*]
using *not-constant*
by *auto*
have *cong*: $\forall\, i{\in}A.\ [?x = u\ i]\ (\mathit{mod}\ m\ i)$
apply *auto*
apply (*rule cong-trans-poly*)
prefer *2*
using *one* **apply** *auto*
apply (*rule cong-dvd-modulus-poly*)
apply (*rule cong-mod-poly*)
using *prodnz* **apply** *auto*
apply *rule*

apply (*rule fin*)
apply *assumption*
done
have *unique*: $\forall z.\ degree\ z < (\sum i{\in}A.\ degree\ (m\ i)) \wedge (\forall i{\in}A.\ [z = u\ i]\ (mod\ m\ i)) \longrightarrow z = ?x$
proof (*clarify*)
fix $z::'b\ poly$
assume *zless*: $degree\ z < (\sum i{\in}A.\ degree\ (m\ i))$
assume *zcong*: $(\forall i{\in}A.\ [z = u\ i]\ (mod\ m\ i))$
have *deg1*: $degree\ z < degree\ (prod\ m\ A)$
using *degree-prod-sum zless* **by** *simp*
have *deg2*: $degree\ ?x < degree\ (prod\ m\ A)$
by (*metis deg1 degree-0 degree-mod-less gr0I gr-implies-not0*)
have $\forall i{\in}A.\ [?x = z]\ (mod\ m\ i)$
apply *clarify*
apply (*rule cong-trans-poly*)
using *cong* **apply** (*erule bspec*)
apply (*rule cong-sym-poly*)
using *zcong* **by** *auto*
with *fin cop* **have** $[?x = z]\ (mod\ (\prod i{\in}A.\ m\ i))$
by (*intro coprime-cong-prod-poly*) *auto*
with *zless* **show** $z = ?x$
apply (*intro cong-less-modulus-unique-poly*)
apply (*erule cong-sym-poly*)
apply (*auto simp add*: *deg1 deg2*)
done
qed
from *less cong unique* **show** *?thesis* **by** *blast*
qed

end

6 The Berlekamp Algorithm

theory *Berlekamp-Type-Based*
imports
Jordan-Normal-Form.Matrix-Kernel
Jordan-Normal-Form.Gauss-Jordan-Elimination
Jordan-Normal-Form.Missing-VectorSpace
Polynomial-Factorization.Square-Free-Factorization
Polynomial-Factorization.Missing-Multiset
Finite-Field
Chinese-Remainder-Poly
Poly-Mod-Finite-Field
HOL−Computational-Algebra.Field-as-Ring
begin

hide-const (**open**) *up-ring.coeff up-ring.monom Modules.module subspace Modules.module-hom*

6.1 Auxiliary lemmas

context
fixes $g :: 'b \Rightarrow 'a :: comm\text{-}monoid\text{-}mult$
begin
lemma *prod-list-map-filter*: *prod-list* (*map g* (*filter f xs*)) $*$ *prod-list* (*map g* (*filter* (λ *x*. $\neg$ *f x*) *xs*))
= *prod-list* (*map g xs*)
by (*induct xs*, *auto simp*: *ac-simps*)

lemma *prod-list-map-partition*:
assumes *List.partition f xs* = (*ys*, *zs*)
shows *prod-list* (*map g xs*) = *prod-list* (*map g ys*) $*$ *prod-list* (*map g zs*)
using *assms* **by** (*subst prod-list-map-filter*[*symmetric*, *of* - *f*], *auto simp*: *o-def*)
end

lemma *coprime-id-is-unit*:
fixes *a*::'*b*::*semiring-gcd*
shows *coprime a a* $\longleftrightarrow$ *is-unit a*
using *dvd-unit-imp-unit* **by** *auto*

lemma *dim-vec-of-list*[*simp*]: *dim-vec* (*vec-of-list x*) = *length x*
by (*transfer*, *auto*)

lemma *length-list-of-vec*[*simp*]: *length* (*list-of-vec A*) = *dim-vec A*
by (*transfer'*, *auto*)

lemma *list-of-vec-vec-of-list*[*simp*]: *list-of-vec* (*vec-of-list a*) = *a*
proof −
{
fix *aa* :: '*a list*
have *map* (λ*n*. *if n* < *length aa then aa* ! *n else undef-vec* (*n* − *length aa*)) [*0*..<*length aa*]
= *map* ((!) *aa*) [*0*..<*length aa*]
by *simp*
hence *map* (λ*n*. *if n* < *length aa then aa* ! *n else undef-vec* (*n* − *length aa*)) [*0*..<*length aa*] = *aa*
by (*simp add*: *map-nth*)
}
thus *?thesis* **by** (*transfer*, *simp add*: *mk-vec-def*)
qed

context
assumes *SORT-CONSTRAINT*('*a*::*finite*)
begin

lemma *inj-Poly-list-of-vec'*: *inj-on* (*Poly* $\circ$ *list-of-vec*) {*v*. *dim-vec v* = *n*}
proof (*rule comp-inj-on*)
show *inj-on list-of-vec* {*v*. *dim-vec v* = *n*}

```
    by (auto simp add: inj-on-def, transfer, auto simp add: mk-vec-def)
  show inj-on Poly (list-of-vec ‘ {v. dim-vec v = n})
  proof (auto simp add: inj-on-def)
    fix x y::'c vec assume n = dim-vec x and dim-xy: dim-vec y = dim-vec x
    and Poly-eq: Poly (list-of-vec x) = Poly (list-of-vec y)
    note [simp del] = nth-list-of-vec
    show list-of-vec x = list-of-vec y
    proof (rule nth-equalityI, auto simp: dim-xy)
      have length-eq: length (list-of-vec x ) = length (list-of-vec y)
        using dim-xy by (transfer, auto)
      fix i assume i < dim-vec x
       thus list-of-vec x ! i = list-of-vec y ! i using Poly-eq unfolding poly-eq-iff
coeff-Poly-eq
        using dim-xy unfolding nth-default-def by (auto, presburger)
     qed
  qed
qed

corollary inj-Poly-list-of-vec: inj-on (Poly ∘ list-of-vec) (carrier-vec n)
  using inj-Poly-list-of-vec' unfolding carrier-vec-def .

lemma list-of-vec-rw-map: list-of-vec m = map (λn. m $ n) [0..<dim-vec m]
    by (transfer, auto simp add: mk-vec-def)

lemma degree-Poly':
assumes xs: xs ≠ []
shows degree (Poly xs) < length xs
using xs
by (induct xs, auto intro: Poly.simps(1))

lemma vec-of-list-list-of-vec[simp]: vec-of-list (list-of-vec a) = a
by (transfer, auto simp add: mk-vec-def)

lemma row-mat-of-rows-list:
assumes b: b < length A
and nc: ∀ i. i < length A ⟶ length (A ! i) = nc
shows (row (mat-of-rows-list nc A) b) = vec-of-list (A ! b)
proof (auto simp add: vec-eq-iff)
  show dim-col (mat-of-rows-list nc A) = length (A ! b)
    unfolding mat-of-rows-list-def using b nc by auto
  fix i assume i: i < length (A ! b)
  show row (mat-of-rows-list nc A) b $ i = vec-of-list (A ! b) $ i
    using i b nc
    unfolding mat-of-rows-list-def row-def
    by (transfer, auto simp add: mk-vec-def mk-mat-def)
qed

lemma degree-Poly-list-of-vec:
assumes n: x ∈ carrier-vec n
```

and *n0*: $n > 0$
shows *degree* (*Poly* (*list-of-vec x*)) < *n*
proof −
have *x-dim*: *dim-vec x* = *n* **using** *n* **by** *auto*
have *l*: (*list-of-vec x*) ≠ []
by (*auto simp add*: *list-of-vec-rw-map vec-of-dim-0*[*symmetric*] *n0 n x-dim*)
have *degree* (*Poly* (*list-of-vec x*)) < *length* (*list-of-vec x*) **by** (*rule degree-Poly'*[*OF l*])
also have ... = *n* **using** *x-dim* **by** *auto*
finally show *?thesis* **.**
qed

lemma *list-of-vec-nth*:
assumes *i*: *i* < *dim-vec x*
shows *list-of-vec x* ! *i* = *x* $ *i*
using *i*
by (*transfer*, *auto simp add*: *mk-vec-def*)

lemma *coeff-Poly-list-of-vec-nth'*:
assumes *i*: *i* < *dim-vec x*
shows *coeff* (*Poly* (*list-of-vec x*)) *i* = *x* $ *i*
using *i*
by (*auto simp add*: *list-of-vec-nth nth-default-def*)

lemma *list-of-vec-row-nth*:
assumes *x*: *x* < *dim-col A*
shows *list-of-vec* (*row A i*) ! *x* = *A* $$ (*i*, *x*)
using *x* **unfolding** *row-def* **by** (*transfer'*, *auto simp add*: *mk-vec-def*)

lemma *coeff-Poly-list-of-vec-nth*:
assumes *x*: *x* < *dim-col A*
shows *coeff* (*Poly* (*list-of-vec* (*row A i*))) *x* = *A* $$ (*i*, *x*)
proof −
have *coeff* (*Poly* (*list-of-vec* (*row A i*))) *x* = *nth-default 0* (*list-of-vec* (*row A i*)) *x*
unfolding *coeff-Poly-eq* **by** *simp*
also have ... = *A* $$ (*i*, *x*) **using** *x list-of-vec-row-nth*
unfolding *nth-default-def* **by** (*auto simp del*: *nth-list-of-vec*)
finally show *?thesis* **.**
qed

lemma *inj-on-list-of-vec*: *inj-on list-of-vec* (*carrier-vec n*)
unfolding *inj-on-def* **unfolding** *list-of-vec-rw-map* **by** *auto*

lemma *vec-of-list-carrier*[*simp*]: *vec-of-list x* ∈ *carrier-vec* (*length x*)
unfolding *carrier-vec-def* **by** *simp*

lemma *card-carrier-vec*: *card* (*carrier-vec n*:: *'b*::*finite vec set*) = *CARD*(*'b*) ^ *n*
proof −

```
  let ?A = UNIV::'b set
  let ?B = {xs. set xs ⊆ ?A ∧ length xs = n}
  let ?C = (carrier-vec n:: 'b::finite vec set)
  have card ?C = card ?B
  proof −
    have bij-betw (list-of-vec) ?C ?B
    proof (unfold bij-betw-def, auto)
      show inj-on list-of-vec (carrier-vec n) by (rule inj-on-list-of-vec)
      fix x::'b list
      assume n: n = length x
      thus x ∈ list-of-vec ‘ carrier-vec (length x)
        unfolding image-def
        by auto (rule bexI[of - vec-of-list x], auto)
    qed
    thus ?thesis using bij-betw-same-card by blast
  qed
  also have ... = card ?A ^ n
    by (rule card-lists-length-eq, simp)
  finally show ?thesis .
qed


lemma finite-carrier-vec[simp]: finite (carrier-vec n:: 'b::finite vec set)
  by (rule card-ge-0-finite, unfold card-carrier-vec, auto)


lemma row-echelon-form-dim0-row:
assumes A ∈ carrier-mat 0 n
shows row-echelon-form A
using assms
unfolding row-echelon-form-def pivot-fun-def Let-def by auto

lemma row-echelon-form-dim0-col:
assumes A ∈ carrier-mat n 0
shows row-echelon-form A
using assms
unfolding row-echelon-form-def pivot-fun-def Let-def by auto

lemma row-echelon-form-one-dim0[simp]: row-echelon-form (1_m 0)
  unfolding row-echelon-form-def pivot-fun-def Let-def by auto

lemma Poly-list-of-vec-0[simp]: Poly (list-of-vec (0_v 0)) = [:0:]
  by (simp add: poly-eq-iff nth-default-def)

lemma monic-normalize:
assumes (p :: 'b :: {field,euclidean-ring-gcd} poly) ≠ 0 shows monic (normalize
p)
by (simp add: assms normalize-poly-old-def)
```

lemma *exists-factorization-prod-list*:
fixes *P*::'*b*::*field poly list*
assumes *degree* (*prod-list P*) > *0*
  and ⋀ *u*. *u* ∈ *set P* ⟹ *degree u* > *0* ∧ *monic u*
  and *square-free* (*prod-list P*)
shows ∃ *Q*. *prod-list Q* = *prod-list P* ∧ *length P* ≤ *length Q*
        ∧ (∀ *u*. *u* ∈ *set Q* ⟶ *irreducible u* ∧ *monic u*)
using *assms*
proof (*induct P*)
  case *Nil*
  thus *?case* **by** *auto*
next
  case (*Cons x P*)
  have *sf-P*: *square-free* (*prod-list P*)
   by (*metis Cons.prems*(*3*) *dvd-triv-left prod-list.Cons mult.commute square-free-factor*)
  have *deg-x*: *degree x* > *0* **using** *Cons.prems* **by** *auto*
  have *distinct-P*: *distinct P*
   by (*meson Cons.prems*(*2*) *Cons.prems*(*3*) *distinct.simps*(*2*) *square-free-prod-list-distinct*)
  have ∃ *A*. *finite A* ∧ *x* = ∏ *A* ∧ *A* ⊆ {*q*. *irreducible q* ∧ *monic q*}
    proof (*rule monic-square-free-irreducible-factorization*)
      show *monic x* **by** (*simp add*: *Cons.prems*(*2*))
      show *square-free x*
        by (*metis Cons.prems*(*3*) *dvd-triv-left prod-list.Cons square-free-factor*)
    qed
    from *this* **obtain** *A* **where** *fin-A*: *finite A*
    and *xA*: *x* = ∏ *A*
    and *A*: *A* ⊆ {*q*. *irreducible*$_d$ *q* ∧ *monic q*}
    by *auto*
    obtain *A'* **where** *s*: *set A'* = *A* **and** *length-A'*: *length A'* = *card A*
      using ‹*finite A*› *distinct-card finite-distinct-list* **by** *force*
  have *A-not-empty*: *A* ≠ {} **using** *xA deg-x* **by** *auto*
  have *x-prod-list-A'*: *x* = *prod-list A'*
  proof −
    have *x* = ∏ *A* **using** *xA* **by** *simp*
    also have ... = *prod id A* **by** *simp*
    also have ... = *prod id* (*set A'*) **unfolding** *s* **by** *simp*
    also have ... = *prod-list* (*map id A'*)
      by (*rule prod.distinct-set-conv-list*, *simp add*: *card-distinct length-A' s*)
    also have ... = *prod-list A'* **by** *auto*
    finally show *?thesis* **.**
  qed
  show *?case*
  proof (*cases P* = [])
    case *True*
    show *?thesis*
    proof (*rule exI*[*of* - *A'*], *auto simp add*: *True*)
      show *prod-list A'* = *x* **using** *x-prod-list-A'* **by** *simp*
      show *Suc 0* ≤ *length A'* **using** *A-not-empty* **using** *s length-A'*

```
      by (simp add: Suc-leI card-gt-0-iff fin-A)
    show ⋀u. u ∈ set A' ⟹ irreducible u using s A by auto
    show ⋀u. u ∈ set A' ⟹ monic u using s A by auto
  qed
 next
  case False
  have hyp: ∃ Q. prod-list Q = prod-list P
   ∧ length P ≤ length Q ∧ (∀ u. u ∈ set Q ⟶ irreducible u ∧ monic u)
  proof (rule Cons.hyps[OF - - sf-P])
   have set-P: set P ≠ {} using False by auto
   have prod-list P = prod-list (map id P) by simp
   also have ... = prod id (set P)
    using prod.distinct-set-conv-list[OF distinct-P, of id] by simp
   also have ... = ∏ (set P) by simp
   finally have prod-list P = ∏ (set P) .
   hence degree (prod-list P) = degree (∏ (set P)) by simp
   also have ... = degree (prod id (set P)) by simp
   also have ... = (∑ i∈(set P). degree (id i))
   proof (rule degree-prod-eq-sum-degree)
    show ∀ i∈set P. id i ≠ 0 using Cons.prems(2) by force
   qed
   also have ... > 0
    by (metis Cons.prems(2) List.finite-set set-P gr0I id-apply insert-iff list.set(2)
sum-pos)
   finally show degree (prod-list P) > 0 by simp
   show ⋀u. u ∈ set P ⟹ degree u > 0 ∧ monic u using Cons.prems by auto
  qed
  from this obtain Q where QP: prod-list Q = prod-list P and length-PQ: length
P ≤ length Q
  and monic-irr-Q: (∀ u. u ∈ set Q ⟶ irreducible u ∧ monic u) by blast
  show ?thesis
  proof (rule exI[of - A' @ Q], auto simp add: monic-irr-Q)
   show prod-list A' * prod-list Q = x * prod-list P unfolding QP x-prod-list-A'
by auto
   have length A' ≠ 0 using A-not-empty using s length-A' by auto
   thus Suc (length P) ≤ length A' + length Q using QP length-PQ by linarith
   show ⋀u. u ∈ set A' ⟹ irreducible u using s A by auto
   show ⋀u. u ∈ set A' ⟹ monic u using s A by auto
  qed
qed
qed

lemma normalize-eq-imp-smult:
  fixes p :: 'b :: {euclidean-ring-gcd} poly
  assumes n: normalize p = normalize q
  shows ∃ c. c ≠ 0 ∧ q = smult c p
proof(cases p = 0)
  case True with n show ?thesis by (auto intro:exI[of - 1])
next
```

case *p0*: *False*
have *degree-eq*: *degree p* = *degree q* **using** *n degree-normalize* **by** *metis*
hence *q0*: *q* ≠ *0* **using** *p0 n* **by** *auto*
have *p-dvd-q*: *p dvd q* **using** *n* **by** (*simp add*: *associatedD1*)
from *p-dvd-q* **obtain** *k* **where** *q*: *q* = *k* ∗ *p* **unfolding** *dvd-def* **by** (*auto simp*: *ac-simps*)
with *q0* **have** *k* ≠ *0* **by** *auto*
then have *degree k* = *0*
using *degree-eq degree-mult-eq p0 q* **by** *fastforce*
then obtain *c* **where** *k*: *k* = [: *c* :] **by** (*metis degree-0-id*)
with ‹*k* ≠ *0*› **have** *c* ≠ *0* **by** *auto*
have *q* = *smult c p* **unfolding** *q k* **by** *simp*
with ‹*c* ≠ *0*› **show** *?thesis* **by** *auto*
qed

lemma *prod-list-normalize*:
fixes *P* :: *'b* :: {*idom-divide*,*normalization-semidom-multiplicative*} *poly list*
shows *normalize* (*prod-list P*) = *prod-list* (*map normalize P*)
proof (*induct P*)
case *Nil*
show *?case* **by** *auto*
next
case (*Cons p P*)
have *normalize* (*prod-list* (*p* # *P*)) = *normalize p* ∗ *normalize* (*prod-list P*)
using *normalize-mult* **by** *auto*
also have ... = *normalize p* ∗ *prod-list* (*map normalize P*) **using** *Cons.hyps* **by** *auto*
also have ... = *prod-list* (*normalize p* # (*map normalize P*)) **by** *auto*
also have ... = *prod-list* (*map normalize* (*p* # *P*)) **by** *auto*
finally show *?case* **.**
qed

lemma *prod-list-dvd-prod-list-subset*:
fixes *A*::*'b*::*comm-monoid-mult list*
assumes *dA*: *distinct A*
and *dB*: *distinct B*
and *s*: *set A* ⊆ *set B*
shows *prod-list A dvd prod-list B*
proof −
have *prod-list A* = *prod-list* (*map id A*) **by** *auto*
also have ... = *prod id* (*set A*)
by (*rule prod.distinct-set-conv-list*[*symmetric*, *OF dA*])
also have ... *dvd prod id* (*set B*)
by (*rule prod-dvd-prod-subset*[*OF - s*], *auto*)
also have ... = *prod-list* (*map id B*)
by (*rule prod.distinct-set-conv-list*[*OF dB*])
also have ... = *prod-list B* **by** *simp*
finally show *?thesis* **.**

qed

end

lemma *gcd-monic-constant*:
  gcd f g ∈ {*1*, *f*} **if** *monic f* **and** *degree g* = *0*
    for *f g* :: *'a* :: {*field-gcd*} *poly*
proof (*cases g* = *0*)
  case *True*
  moreover from ‹*monic f*› **have** *normalize f* = *f*
    by (*rule normalize-monic*)
  ultimately show *?thesis*
    by *simp*
next
  case *False*
  with ‹*degree g* = *0*› **have** *is-unit g*
    by *simp*
  then have *Rings.coprime f g*
    by (*rule is-unit-right-imp-coprime*)
  then show *?thesis*
    by *simp*
qed

lemma *distinct-find-base-vectors*:
fixes *A*::*'a*::*field mat*
assumes *ref*: *row-echelon-form A*
  and *A*: *A* ∈ *carrier-mat nr nc*
shows *distinct* (*find-base-vectors A*)
proof −
  note *non-pivot-base* = *non-pivot-base*[*OF ref A*]
  let *?pp* = *set* (*pivot-positions A*)
  from *A* **have** *dim*: *dim-row A* = *nr dim-col A* = *nc* **by** *auto*
  {
    fix *j j'*
    assume *j*: *j* < *nc j* ∉ *snd* ‘ *?pp* **and** *j'*: *j'* < *nc j'* ∉ *snd* ‘ *?pp* **and** *neq*: *j'* ≠ *j*
    from *non-pivot-base*(*2*)[*OF j*] *non-pivot-base*(*4*)[*OF j' j neq*]
   have *non-pivot-base A* (*pivot-positions A*) *j* ≠ *non-pivot-base A* (*pivot-positions A*) *j'* **by** *auto*
  }
  hence *inj*: *inj-on* (*non-pivot-base A* (*pivot-positions A*))
    (*set* [*j*←[*0*..<*nc*] . *j* ∉ *snd* ‘ *?pp*]) **unfolding** *inj-on-def* **by** *auto*
  thus *?thesis* **unfolding** *find-base-vectors-def Let-def* **unfolding** *distinct-map dim* **by** *auto*
qed

lemma *length-find-base-vectors*:
fixes *A*::*'a*::*field mat*
assumes *ref*: *row-echelon-form A*
  and *A*: *A* ∈ *carrier-mat nr nc*

shows *length* (*find-base-vectors A*) = *card* (*set* (*find-base-vectors A*))
using *distinct-card*[*OF distinct-find-base-vectors*[*OF ref A*]] **by** *auto*

6.2 Previous Results

definition *power-poly-f-mod* :: *'a*::*field poly* ⇒ *'a poly* ⇒ *nat* ⇒ *'a poly* **where**
power-poly-f-mod modulus = (λ*a n*. *a* ^ *n mod modulus*)

lemma *power-poly-f-mod-binary*: *power-poly-f-mod m a n* = (*if n* = *0 then 1 mod m*
else let (*d*, *r*) = *Euclidean-Rings.divmod-nat n 2*;
rec = *power-poly-f-mod m* ((*a* ∗ *a*) *mod m*) *d in*
if r = *0 then rec else* (*rec* ∗ *a*) *mod m*)
for *m a* :: *'a* :: {*field-gcd*} *poly*
proof −
note *d* = *power-poly-f-mod-def*
show *?thesis*
proof (*cases n* = *0*)
case *True*
thus *?thesis* **unfolding** *d* **by** *simp*
next
case *False*
obtain *q r* **where** *div*: *Euclidean-Rings.divmod-nat n 2* = (*q*,*r*) **by** *force*
hence *n*: *n* = *2* ∗ *q* + *r* **and** *r*: *r* = *0* ∨ *r* = *1* **unfolding** *Euclidean-Rings.divmod-nat-def*
by *auto*
have *id*: *a* ^ (*2* ∗ *q*) = (*a* ∗ *a*) ^ *q*
by (*simp add*: *power-mult-distrib semiring-normalization-rules*)
show *?thesis*
proof (*cases r* = *0*)
case *True*
show *?thesis*
using *power-mod* [*of a* ∗ *a m q*]
by (*auto simp add*: *Euclidean-Rings.divmod-nat-def Let-def True n d div id*)
next
case *False*
with *r* **have** *r*: *r* = *1* **by** *simp*
show *?thesis*
by (*auto simp add*: *d r div Let-def mod-simps*)
(*simp add*: *n r mod-simps ac-simps power-mult-distrib power-mult power2-eq-square*)
qed
qed
qed

fun *power-polys* **where**
power-polys mul-p u curr-p (*Suc i*) = *curr-p* #
power-polys mul-p u ((*curr-p* ∗ *mul-p*) *mod u*) *i*
| *power-polys mul-p u curr-p 0* = []

context
assumes *SORT-CONSTRAINT('a::prime-card)*
begin

lemma *fermat-theorem-mod-ring* [*simp*]:
  fixes *a::'a mod-ring*
  shows *a ^ CARD('a) = a*
proof (*cases a = 0*)
  case *True*
  then show *?thesis* **by** *auto*
next
  case *False*
  then show *?thesis*
  proof *transfer*
    fix *a*
    assume *a ∈ {0..<int CARD('a)}* **and** *a ≠ 0*
    then have *a: 1 ≤ a a < int CARD('a)*
      by *simp-all*
    then have [*simp*]: *a mod int CARD('a) = a*
      by *simp*
    from *a* **have** *¬ int CARD('a) dvd a*
      by (*auto simp add: zdvd-not-zless*)
    then have *¬ CARD('a) dvd nat |a|*
      by *simp*
    with *a* **have** *¬ CARD('a) dvd nat a*
      by *simp*
    with *prime-card* **have** [*nat a ^ (CARD('a) − 1) = 1*] (*mod CARD('a)*)
      by (*rule fermat-theorem*)
    with *a* **have** *int (nat a ^ (CARD('a) − 1) mod CARD('a)) = 1*
      by (*simp add: cong-def*)
    with *a* **have** *a ^ (CARD('a) − 1) mod CARD('a) = 1*
      by (*simp add: of-nat-mod*)
    then have *a * (a ^ (CARD('a) − 1) mod int CARD('a)) = a*
      by *simp*
   then have (*a * (a ^ (CARD('a) − 1) mod int CARD('a))) mod int CARD('a)*
= a mod int CARD('a)
      by (*simp only:*)
    then show *a ^ CARD('a) mod int CARD('a) = a*
      by (*simp add: mod-simps semiring-normalization-rules(27)*)
  qed
qed

lemma *mod-eq-dvd-iff-poly*: ((*x::'a mod-ring poly*) *mod n = y mod n*) = (*n dvd x*
− y)
proof
  assume *H*: *x mod n = y mod n*
  hence *x mod n − y mod n = 0* **by** *simp*
  hence (*x mod n − y mod n*) *mod n = 0* **by** *simp*

hence $(x - y)$ *mod* $n = 0$ **by** (*simp add*: *mod-diff-eq*)
thus n *dvd* $x - y$ **by** (*simp add*: *dvd-eq-mod-eq-0*)
next
assume *H*: n *dvd* $x - y$
then obtain k **where** *k*: $x-y = n*k$ **unfolding** *dvd-def* **by** *blast*
hence $x = n*k + y$ **using** *diff-eq-eq* **by** *blast*
hence x *mod* $n = (n*k + y)$ *mod* n **by** *simp*
thus x *mod* $n = y$ *mod* n **by** (*simp add*: *mod-add-left-eq*)
qed

lemma *cong-gcd-eq-poly*:
gcd a m = *gcd b m* **if** [(*a*::'*a mod-ring poly*) = *b*] (*mod m*)
using *that* **by** (*simp add*: *cong-def*) (*metis gcd-mod-left mod-by-0*)

lemma *coprime-h-c-poly*:
fixes *h*::'*a mod-ring poly*
assumes *c1* ≠ *c2*
shows *coprime* (*h* − [:*c1*:]) (*h* − [:*c2*:])
proof (*intro coprimeI*)
fix *d* **assume** *d dvd h* − [:*c1*:]
and *d dvd h* − [:*c2*:]
hence *h mod d* = [:*c1*:] *mod d* **and** *h mod d* = [:*c2*:] *mod d*
using *mod-eq-dvd-iff-poly* **by** *simp*+
hence [:*c1*:] *mod d* = [:*c2*:] *mod d* **by** *simp*
hence *d dvd* [:*c2* − *c1*:]
by (*metis* (*no-types*) *mod-eq-dvd-iff-poly diff-pCons right-minus-eq*)
thus *is-unit d*
by (*metis* (*no-types*) *assms dvd-trans is-unit-monom-0 monom-0 right-minus-eq*)
qed

lemma *coprime-h-c-poly2*:
fixes *h*::'*a mod-ring poly*
assumes *coprime* (*h* − [:*c1*:]) (*h* − [:*c2*:])
and ¬ *is-unit* (*h* − [:*c1*:])
shows *c1* ≠ *c2*
using *assms coprime-id-is-unit* **by** *blast*

lemma *degree-minus-eq-right*:
fixes *p*::'*b*::*ab-group-add poly*
shows *degree q* < *degree p* ⟹ *degree* (*p* − *q*) = *degree p*
using *degree-add-eq-left*[*of* −*q p*] *degree-minus* **by** *auto*

lemma *coprime-prod*:
fixes *A*::'*a mod-ring set* **and** *g*::'*a mod-ring* ⇒ '*a mod-ring poly*
assumes ∀ *x*∈*A*. *coprime* (*g a*) (*g x*)
shows *coprime* (*g a*) (*prod* (λ*x*. *g x*) *A*)
proof −

```
  have f: finite A by simp
  show ?thesis
  using f using assms
  proof (induct A)
    case (insert x A)
    have (∏ c∈insert x A. g c) = (g x) * (∏ c∈A. g c)
      by (simp add: insert.hyps(2))
    with insert.prems show ?case
      by (auto simp: insert.hyps(3) prod-coprime-right)
  qed auto
qed


lemma coprime-prod2:
  fixes A::'b::semiring-gcd set
  assumes ∀ x∈A. coprime (a) (x) and f: finite A
  shows coprime (a) (prod (λx. x) A)
  using f using assms
proof (induct A)
  case (insert x A)
  have (∏ c∈insert x A. c) = (x) * (∏ c∈A. c)
    by (simp add: insert.hyps)
  with insert.prems show ?case
    by (simp add: insert.hyps prod-coprime-right)
qed auto


lemma divides-prod:
  fixes g::'a mod-ring ⇒ 'a mod-ring poly
  assumes ∀ c1 c2. c1 ∈ A ∧ c2 ∈ A ∧ c1 ≠ c2 ⟶ coprime (g c1) (g c2)
  assumes ∀ c∈ A. g c dvd f
  shows (∏ c∈A. g c) dvd f
proof −
  have finite-A: finite A using finite[of A] .
  thus ?thesis using assms
  proof (induct A)
    case (insert x A)
    have (∏ c∈insert x A. g c) =  g x * (∏ c∈ A. g c)
      by (simp add: insert.hyps(2))
    also have ... dvd f
    proof (rule divides-mult)
      show g x dvd f using insert.prems by auto
      show prod g A dvd f using insert.hyps(3) insert.prems by auto
      from insert show Rings.coprime (g x) (prod g A)
        by (auto intro: prod-coprime-right)
    qed
    finally show ?case .
  qed auto
```

qed

lemma *poly-monom-identity-mod-p*:
monom (*1*::*'a mod-ring*) (*CARD*(*'a*)) − *monom 1 1* = *prod* (λ*x*. [:*0*,*1*:] − [:*x*:]) (*UNIV*::*'a mod-ring set*)
(**is** *?lhs* = *?rhs*)
proof −
let *?f*=(λ*x*::*'a mod-ring*. [:*0*,*1*:] − [:*x*:])
have *?rhs dvd ?lhs*
proof (*rule divides-prod*)
{
fix *a*::*'a mod-ring*
have *poly ?lhs a* = *0*
by (*simp add*: *poly-monom*)
hence ([:*0*,*1*:] − [:*a*:]) *dvd ?lhs*
using *poly-eq-0-iff-dvd* **by** *fastforce*
}
thus ∀ *x*∈*UNIV*::*'a mod-ring set*. [:*0*, *1*:] − [:*x*:] *dvd monom 1 CARD*(*'a*) − *monom 1 1* **by** *fast*
show ∀ *c1 c2*. *c1* ∈ *UNIV* ∧ *c2* ∈ *UNIV* ∧ *c1* ≠ (*c2* :: *'a mod-ring*) ⟶ *coprime* ([:*0*, *1*:] − [:*c1*:]) ([:*0*, *1*:] − [:*c2*:])
by (*auto dest*!: *coprime-h-c-poly*[*of* - - [:*0*,*1*:]])
qed
from *this* **obtain** *g* **where** *g*: *?lhs* = *?rhs* ∗ *g* **using** *dvdE* **by** *blast*
have *degree-lhs-card*: *degree ?lhs* = *CARD*(*'a*)
proof −
have *degree* (*monom* (*1*::*'a mod-ring*) *1*) = *1* **by** (*simp add*: *degree-monom-eq*)
moreover have *d-c*: *degree* (*monom* (*1*::*'a mod-ring*) *CARD*(*'a*)) = *CARD*(*'a*)
by (*simp add*: *degree-monom-eq*)
ultimately have *degree* (*monom* (*1*::*'a mod-ring*) *1*) < *degree* (*monom* (*1*::*'a mod-ring*) *CARD*(*'a*))
using *prime-card* **unfolding** *prime-nat-iff* **by** *auto*
hence *degree ?lhs* = *degree* (*monom* (*1*::*'a mod-ring*) *CARD*(*'a*))
by (*rule degree-minus-eq-right*)
thus *?thesis* **unfolding** *d-c* **.**
qed
have *degree-rhs-card*: *degree ?rhs* = *CARD*(*'a*)
proof −
have *degree* (*prod ?f UNIV*) = *sum* (*degree* ∘ *?f*) *UNIV*
∧ *coeff* (*prod ?f UNIV*) (*sum* (*degree* ∘ *?f*) *UNIV*) = *1*
by (*rule degree-prod-sum-monic*, *auto*)
moreover have *sum* (*degree* ∘ *?f*) *UNIV* = *CARD*(*'a*) **by** *auto*
ultimately show *?thesis* **by** *presburger*
qed
have *monic-lhs*: *monic ?lhs* **using** *degree-lhs-card* **by** *auto*
have *monic-rhs*: *monic ?rhs* **by** (*rule monic-prod*, *simp*)
have *degree-eq*: *degree ?rhs* = *degree ?lhs* **unfolding** *degree-lhs-card degree-rhs-card*

..
  have *g-not-0*: $g \neq 0$ **using** *g monic-lhs* **by** *auto*
  have *degree-g0*: *degree* $g = 0$
  proof −
    have *degree* (*?rhs* ∗ *g*) = *degree ?rhs* + *degree g*
      by (*rule degree-monic-mult*[*OF monic-rhs g-not-0*])
    thus *?thesis* **using** *degree-eq g* **by** *simp*
  qed
  have *monic-g*: *monic g* **using** *monic-factor g monic-lhs monic-rhs* **by** *auto*
  have $g = 1$ **using** *monic-degree-0*[*OF monic-g*] *degree-g0* **by** *simp*
  thus *?thesis* **using** *g* **by** *auto*
qed

lemma *poly-identity-mod-p*:
  $v\hat{}(CARD('a)) - v = prod\ (\lambda x.\ v - [:x:])\ (UNIV::'a\ mod\text{-}ring\ set)$
 proof −
  have *id*: *monom 1 1* $\circ_p$ $v = v$ [:*0*, *1*:] $\circ_p$ $v = v$ **unfolding** *pcompose-def*
    apply (*auto*)
    by (*simp add*: *fold-coeffs-def*)
  have *id2*: *monom 1* (*CARD*('*a*)) $\circ_p$ $v = v\ \hat{}\ (CARD('a))$ **by** (*metis id*(*1*) *pcompose-hom.hom-power x-pow-n*)
  show *?thesis* **using** *arg-cong*[*OF poly-monom-identity-mod-p*, *of* $\lambda\ f.\ f \circ_p v$]
   unfolding *pcompose-hom.hom-minus pcompose-hom.hom-prod id pcompose-const id2* **.**
qed

lemma *coprime-gcd*:
  fixes *h*::'*a mod-ring poly*
  assumes *Rings.coprime* (*h*−[:*c1*:]) (*h*−[:*c2*:])
  shows *Rings.coprime* (*gcd f*(*h*−[:*c1*:])) (*gcd f* (*h*−[:*c2*:]))
  using *assms coprime-divisors* **by** *blast*

lemma *divides-prod-gcd*:
  fixes *h*::'*a mod-ring poly*
  assumes $\forall$ *c1 c2*. *c1* ∈ *A* ∧ *c2* ∈ *A* ∧ *c1* ≠ *c2* ⟶ *coprime* (*h*−[:*c1*:]) (*h*−[:*c2*:])
  shows ($\prod$ *c*∈*A*. *gcd f* (*h* − [:*c*:])) *dvd f*
proof −
  have *finite-A*: *finite A* **using** *finite*[*of A*] **.**
  thus *?thesis* **using** *assms*
  proof (*induct A*)
    case (*insert x A*)
    have ($\prod$ *c*∈*insert x A*. *gcd f* (*h* − [:*c*:])) = (*gcd f* (*h* − [:*x*:])) ∗ ($\prod$ *c*∈ *A*. *gcd*

f (*h* − [:*c*:]))
by (*simp add*: *insert.hyps*(*2*))
also have ... *dvd f*
proof (*rule divides-mult*)
show *gcd f* (*h* − [:*x*:]) *dvd f* **by** *simp*
show (∏ *c*∈*A*. *gcd f* (*h* − [:*c*:])) *dvd f* **using** *insert.hyps*(*3*) *insert.prems* **by** *auto*
show *Rings.coprime* (*gcd f* (*h* − [:*x*:])) (∏ *c*∈*A*. *gcd f* (*h* − [:*c*:]))
by (*rule prod-coprime-right*)
(*metis Berlekamp-Type-Based.coprime-h-c-poly coprime-gcd coprime-iff-coprime insert.hyps*(*2*))
qed
finally show *?case* **.**
qed *auto*
qed

lemma *monic-prod-gcd*:
assumes *f*: *finite A* **and** *f0*: (*f* :: *'b* :: {*field-gcd*} *poly*) ≠ *0*
shows *monic* (∏ *c*∈*A*. *gcd f* (*h* − [:*c*:]))
using *f*
proof (*induct A*)
case (*insert x A*)
have *rw*: (∏ *c*∈*insert x A*. *gcd f* (*h* − [:*c*:]))
= (*gcd f* (*h* − [:*x*:])) ∗ (∏ *c*∈ *A*. *gcd f* (*h* − [:*c*:]))
by (*simp add*: *insert.hyps*)
show *?case*
proof (*unfold rw*, *rule monic-mult*)
show *monic* (*gcd f* (*h* − [:*x*:]))
using *poly-gcd-monic*[*of f*] *f0*
using *insert.prems insert-iff* **by** *blast*
show *monic* (∏ *c*∈*A*. *gcd f* (*h* − [:*c*:]))
using *insert.hyps*(*3*) *insert.prems* **by** *blast*
qed
qed *auto*

lemma *coprime-not-unit-not-dvd*:
fixes *a*::*'b*::*semiring-gcd*
assumes *a dvd b*
and *coprime b c*
and ¬ *is-unit a*
shows ¬ *a dvd c*
using *assms coprime-divisors coprime-id-is-unit* **by** *fastforce*

lemma *divides-prod2*:
fixes *A*::*'b*::*semiring-gcd set*
assumes *f*: *finite A*
and ∀ *a*∈*A*. *a dvd c*
and ∀ *a1 a2*. *a1* ∈ *A* ∧ *a2* ∈ *A* ∧ *a1* ≠ *a2* ⟶ *coprime a1 a2*
shows ∏ *A dvd c*

using *assms*
proof (*induct A*)
case (*insert x A*)
have $\prod$ (*insert x A*) = *x* ∗ $\prod$ *A* **by** (*simp add*: *insert.hyps*(*1*) *insert.hyps*(*2*))
also have ... *dvd c*
proof (*rule divides-mult*)
show *x dvd c* **by** (*simp add*: *insert.prems*)
show $\prod$ *A dvd c* **using** *insert* **by** *auto*
from *insert* **show** *Rings.coprime x* ($\prod$ *A*)
by (*auto intro*: *prod-coprime-right*)
qed
finally show *?case* **.**
qed *auto*

lemma *coprime-polynomial-factorization*:
fixes *a1* :: *'b* :: {*field-gcd*} *poly*
assumes *irr*: *as* ⊆ {*q*. *irreducible q* ∧ *monic q*}
and *finite as* **and** *a1*: *a1* ∈ *as* **and** *a2*: *a2* ∈ *as* **and** *a1-not-a2*: *a1* ≠ *a2*
shows *coprime a1 a2*
proof (*rule ccontr*)
assume *not-coprime*: ¬ *coprime a1 a2*
let *?b*= *gcd a1 a2*
have *b-dvd-a1*: *?b dvd a1* **and** *b-dvd-a2*: *?b dvd a2* **by** *simp*+
have *irr-a1*: *irreducible a1* **using** *a1 irr* **by** *blast*
have *irr-a2*: *irreducible a2* **using** *a2 irr* **by** *blast*
have *a2-not0*: *a2* ≠ *0* **using** *a2 irr* **by** *auto*
have *degree-a1*: *degree a1* ≠ *0* **using** *irr-a1* **by** *auto*
have *degree-a2*: *degree a2* ≠ *0* **using** *irr-a2* **by** *auto*
have *not-a2-dvd-a1*: ¬ *a2 dvd a1*
proof (*rule ccontr*, *simp*)
assume *a2-dvd-a1*: *a2 dvd a1*
from *this* **obtain** *k* **where** *k*: *a1* = *a2* ∗ *k* **unfolding** *dvd-def* **by** *auto*
have *k-not0*: *k* ≠ *0* **using** *degree-a1 k* **by** *auto*
show *False*
proof (*cases degree a2* = *degree a1*)
case *False*
have *degree a2* < *degree a1*
using *False a2-dvd-a1 degree-a1 divides-degree*
by *fastforce*
hence ¬ *irreducible a1*
using *degree-a2 a2-dvd-a1 degree-a2*
by (*metis degree-a1 irreducible*$_d$*D*(*2*) *irreducible*$_d$*-multD irreducible-connect-field k neq0-conv*)
thus *False* **using** *irr-a1* **by** *contradiction*
next
case *True*
have *degree a1* = *degree a2* + *degree k*
unfolding *k* **using** *degree-mult-eq*[*OF a2-not0 k-not0*] **by** *simp*

```
    hence degree k = 0 using True by simp
    hence k = 1 using monic-factor a1 a2 irr k monic-degree-0 by auto
    hence a1 = a2 using k by simp
    thus False using a1-not-a2 by contradiction
   qed
  qed
  have b-not0: ?b ≠ 0 by (simp add: a2-not0)
  have degree-b: degree ?b > 0
   using not-coprime[simplified] b-not0 is-unit-gcd is-unit-iff-degree by blast
  have degree ?b < degree a2
   by (meson b-dvd-a1 b-dvd-a2 irreducibleD' dvd-trans gcd-dvd-1 irr-a2 not-a2-dvd-a1
not-coprime)
  hence ¬ irreducible_d a2 using degree-a2 b-dvd-a2 degree-b
   by (metis degree-smult-eq irreducible_d-dvd-smult less-not-refl3)
  thus False using irr-a2 by auto
qed


theorem Berlekamp-gcd-step:
fixes f::'a mod-ring poly and h::'a mod-ring poly
assumes hq-mod-f: [h^(CARD('a)) = h] (mod f) and monic-f: monic f and sf-f:
square-free f
shows f = prod (λc. gcd f (h − [:c:])) (UNIV::'a mod-ring set)  (is ?lhs = ?rhs)
proof (cases f=0)
  case True
  thus ?thesis using coeff-0 monic-f zero-neq-one by auto
  next
  case False note f-not-0 = False
  show ?thesis
  proof (rule poly-dvd-antisym)
   show ?rhs dvd f
    using coprime-h-c-poly by (intro divides-prod-gcd, auto)
   have monic ?rhs by (rule monic-prod-gcd[OF - f-not-0], simp)
   thus coeff f (degree f) = coeff ?rhs (degree ?rhs)
    using monic-f by auto
   next
   show f dvd ?rhs
   proof −
    let ?p = CARD('a)
    obtain P  where finite-P: finite P
    and f-desc-square-free: f = (∏ a∈P. a)
    and P: P ⊆ {q. irreducible q ∧ monic q}
     using monic-square-free-irreducible-factorization[OF monic-f sf-f] by auto
    have f-dvd-hqh: f dvd (h^?p − h) using hq-mod-f unfolding cong-def
     using mod-eq-dvd-iff-poly by blast
    also have hq-h-rw: ... = prod (λc. h − [:c:]) (UNIV::'a mod-ring set)
     by (rule poly-identity-mod-p)
    finally have f-dvd-hc: f dvd prod (λc. h − [:c:]) (UNIV::'a mod-ring set) by
simp
```

```
      have f = ∏ P using f-desc-square-free by simp
      also have ... dvd ?rhs
      proof (rule divides-prod2[OF finite-P])
        show ∀ a1 a2. a1 ∈ P ∧ a2 ∈ P ∧ a1 ≠ a2 ⟶ coprime a1 a2
          using coprime-polynomial-factorization[OF P finite-P] by simp
        show ∀ a∈P. a dvd (∏ c∈UNIV. gcd f (h − [:c:]))
        proof
          fix fi assume fi-P: fi ∈ P
          show fi dvd ?rhs
          proof (rule dvd-prod, auto)
            show fi dvd f using f-desc-square-free fi-P
             using dvd-prod-eqI finite-P by blast
            hence fi dvd (h^?p − h) using dvd-trans f-dvd-hqh by auto
            also have ... = prod (λc. h − [:c:]) (UNIV::'a mod-ring set)
              unfolding hq-h-rw by simp
          finally have fi-dvd-prod-hc: fi dvd prod (λc. h − [:c:]) (UNIV::'a mod-ring
set) .
            have irr-fi: irreducible (fi) using fi-P P by blast
           have fi-not-unit: ¬ is-unit fi using irr-fi by (simp add: irreducible_dD(1)
poly-dvd-1)
            have fi-dvd-hc: ∃ c∈UNIV::'a mod-ring set. fi dvd (h−[:c:])
              by (rule irreducible-dvd-prod[OF - fi-dvd-prod-hc], simp add: irr-fi)
            thus ∃ c. fi dvd h − [:c:] by simp
          qed
        qed
      qed
      finally show f dvd ?rhs .
    qed
  qed
qed
```

6.3 Definitions

```
definition berlekamp-mat :: 'a mod-ring poly ⇒ 'a mod-ring mat where
  berlekamp-mat u = (let n = degree u;
    mul-p = power-poly-f-mod u [:0,1:] (CARD('a));
    xks = power-polys mul-p u 1 n
   in
    mat-of-rows-list n (map (λ cs. let coeffs-cs = (coeffs cs);
                                       k = n − length (coeffs cs)
                                   in (coeffs cs) @ replicate k 0) xks))


definition berlekamp-resulting-mat :: ('a mod-ring) poly ⇒ 'a mod-ring mat where
berlekamp-resulting-mat u = (let Q = berlekamp-mat u;
    n = dim-row Q;
    QI = mat n n (λ (i,j). if i = j then Q $$ (i,j) − 1 else Q $$ (i,j))
    in (gauss-jordan-single (transpose-mat QI)))
```

definition *berlekamp-basis* :: *'a mod-ring poly ⇒ 'a mod-ring poly list* **where**
berlekamp-basis u = (*map* (*Poly o list-of-vec*) (*find-base-vectors* (*berlekamp-resulting-mat u*)))

lemma *berlekamp-basis-code*[*code*]: *berlekamp-basis u* =
(*map* (*poly-of-list o list-of-vec*) (*find-base-vectors* (*berlekamp-resulting-mat u*)))
unfolding *berlekamp-basis-def poly-of-list-def* **..**

primrec *berlekamp-factorization-main* :: *nat ⇒ 'a mod-ring poly list ⇒ 'a mod-ring poly list ⇒ nat ⇒ 'a mod-ring poly list* **where**
berlekamp-factorization-main i divs (*v # vs*) *n* = (*if v = 1 then berlekamp-factorization-main i divs vs n else*
if length divs = n then divs else
let facts = [*w . u ← divs, s ← [0 ..< CARD('a)], w ← [gcd u* (*v* − [:*of-int s*:])], *w ≠ 1*];
(*lin*,*nonlin*) = *List.partition* (λ *q. degree q* = *i*) *facts*
in lin @ *berlekamp-factorization-main i nonlin vs* (*n* − *length lin*))
| *berlekamp-factorization-main i divs* [] *n* = *divs*

definition *berlekamp-monic-factorization* :: *nat ⇒ 'a mod-ring poly ⇒ 'a mod-ring poly list* **where**
berlekamp-monic-factorization d f = (*let*
vs = *berlekamp-basis f*;
n = *length vs*;
fs = *berlekamp-factorization-main d* [*f*] *vs n*
in fs)

6.4 Properties

lemma *power-polys-works*:
fixes *u*::*'b*::*unique-euclidean-semiring*
assumes *i*: *i* < *n* **and** *c*: *curr-p* = *curr-p mod u*
shows *power-polys mult-p u curr-p n* ! *i* = *curr-p* ∗ *mult-p* ^ *i mod u*
using *i c*
proof (*induct n arbitrary*: *curr-p i*)
case *0* **thus** *?case* **by** *simp*
next
case (*Suc n*)
have *p-rw*: *power-polys mult-p u curr-p* (*Suc n*) ! *i*
= (*curr-p # power-polys mult-p u* (*curr-p* ∗ *mult-p mod u*) *n*) ! *i*
by *simp*
show *?case*
proof (*cases i=0*)
case *True*
show *?thesis* **using** *Suc.prems* **unfolding** *p-rw True* **by** *auto*
next
case *False* **note** *i-not-0* = *False*
show *?thesis*
proof (*cases i* < *n*)

```
      case True note i-less-n = True
      have power-polys mult-p u curr-p (Suc n) ! i = power-polys mult-p u (curr-p
* mult-p mod u) n ! (i − 1)
        unfolding p-rw using nth-Cons-pos False by auto
      also have ... = (curr-p * mult-p mod u) * mult-p ^(i−1) mod u
        by (rule Suc.hyps) (auto simp add: i-less-n less-imp-diff-less)
      also have ... = curr-p * mult-p ^i mod u
        using False by (cases i) (simp-all add: algebra-simps mod-simps)
      finally show ?thesis .
    next
      case False
      hence i-n: i = n using Suc.prems by auto
      have power-polys mult-p u curr-p (Suc n) ! i = power-polys mult-p u (curr-p
* mult-p mod u) n ! (n − 1)
          unfolding p-rw using nth-Cons-pos i-n i-not-0 by auto
      also have ... = (curr-p * mult-p mod u) * mult-p ^(n−1) mod u
      proof (rule Suc.hyps)
        show n − 1 < n using i-n i-not-0 by linarith
        show curr-p * mult-p mod u = curr-p * mult-p mod u mod u by simp
      qed
      also have ... = curr-p * mult-p ^i mod u
          using i-n [symmetric] i-not-0 by (cases i) (simp-all add: algebra-simps
mod-simps)
      finally show ?thesis .
    qed
  qed
qed


lemma length-power-polys[simp]: length (power-polys mult-p u curr-p n) = n
  by (induct n arbitrary: curr-p, auto)



lemma Poly-berlekamp-mat:
assumes k: k < degree u
shows Poly (list-of-vec (row (berlekamp-mat u) k)) = [:0,1:]^(CARD('a) * k) mod
u
proof −
  let ?map =(map (λcs. coeffs cs @ replicate (degree u − length (coeffs cs)) 0)
              (power-polys (power-poly-f-mod u [:0, 1:] (nat (int CARD('a)))) u 1
(degree u)))
  have row (berlekamp-mat u) k = row (mat-of-rows-list (degree u) ?map) k
    by (simp add: berlekamp-mat-def Let-def)
  also have ... = vec-of-list (?map ! k)
  proof−
    {
      fix i assume i: i < degree u
```

```
      then have ‹u ≠ 0›
        by auto
      let ?c= power-polys (power-poly-f-mod u [:0, 1:] CARD('a)) u 1 (degree u) !
i
      let ?coeffs-c=(coeffs ?c)
      have ?c = 1*([:0, 1:] ^ CARD('a) mod u)^i mod u
      proof (unfold power-poly-f-mod-def, rule power-polys-works[OF i])
        show 1 = 1 mod u using k mod-poly-less by force
      qed
       also have ... = [:0, 1:] ^ (CARD('a) * i) mod u by (simp add: power-mod
power-mult)
      finally have c-rw: ?c = [:0, 1:] ^ (CARD('a) * i) mod u .
      have length ?coeffs-c ≤ degree u
      proof −
        show ?thesis
        proof (cases ?c = 0)
          case True thus ?thesis by auto
          next
          case False
          have length ?coeffs-c = degree (?c) + 1 by (rule length-coeffs[OF False])
         also have ... = degree ([:0, 1:] ^ (CARD('a) * i) mod u) + 1 using c-rw
by simp
          also have ... ≤ degree u
               using ‹i < degree u› ‹u ≠ 0› degree-mod-less [of u ‹pCons 0 1 ^
(CARD('a) * i)›]
            by auto
          finally show ?thesis .
        qed
      qed
     then have length ?coeffs-c + (degree u − length ?coeffs-c) = degree u by auto
    }
    with k show ?thesis by (intro row-mat-of-rows-list, auto)
  qed
  finally have row-rw: row (berlekamp-mat u) k = vec-of-list (?map ! k) .
  have Poly (list-of-vec (row (berlekamp-mat u) k)) = Poly (list-of-vec (vec-of-list
(?map ! k)))
    unfolding row-rw ..
  also have ... = Poly (?map ! k) by simp
  also have ... = [:0,1:]^(CARD('a) * k) mod u
  proof −
    let ?cs = (power-polys (power-poly-f-mod u [:0, 1:] (nat (int CARD('a)))) u 1
(degree u)) ! k
    let ?c = coeffs ?cs @ replicate (degree u − length (coeffs ?cs)) 0
    have map-k-c: ?map ! k = ?c by (rule nth-map, simp add: k)
   have (Poly (?map ! k)) = Poly (coeffs ?cs) unfolding map-k-c Poly-append-replicate-0
..
    also have ... = ?cs by simp
    also have ... = power-polys ([:0, 1:] ^ CARD('a) mod u) u 1 (degree u) ! k
      by (simp add: power-poly-f-mod-def)
```

```
    also have ... = 1* ([:0,1:]^(CARD('a)) mod u) ^ k mod u
    proof (rule power-polys-works[OF k])
      show 1 = 1 mod u using k mod-poly-less by force
    qed
    also have ... = ([:0,1:]^(CARD('a)) mod u) ^ k mod u by auto
      also have ... = [:0,1:]^(CARD('a) * k) mod u by (simp add: power-mod
power-mult)
    finally show ?thesis .
  qed
    finally show ?thesis .
qed

corollary Poly-berlekamp-cong-mat:
assumes k: k < degree u
shows [Poly (list-of-vec (row (berlekamp-mat u) k)) = [:0,1:]^(CARD('a) * k)]
(mod u)
using Poly-berlekamp-mat[OF k] unfolding cong-def by auto

lemma mat-of-rows-list-dim[simp]:
  mat-of-rows-list n vs ∈ carrier-mat (length vs) n
  dim-row (mat-of-rows-list n vs) = length vs
  dim-col (mat-of-rows-list n vs) = n
  unfolding mat-of-rows-list-def by auto

lemma berlekamp-mat-closed[simp]:
  berlekamp-mat u ∈ carrier-mat (degree u) (degree u)
  dim-row (berlekamp-mat u) = degree u
  dim-col (berlekamp-mat u) = degree u
 unfolding carrier-mat-def berlekamp-mat-def Let-def by auto


lemma vec-of-list-coeffs-nth:
assumes i: i ∈ {..degree h} and h-not0: h ≠ 0
shows vec-of-list (coeffs h) $ i = coeff h i
proof −
  have vec-of-list (map (coeff h) [0..<degree h] @ [coeff h (degree h)]) $ i = coeff
h i
      using i
      by (transfer', auto simp add: mk-vec-def)
        (metis (no-types, lifting) Cons-eq-append-conv coeffs-def coeffs-nth degree-0
       diff-zero length-upt less-eq-nat.simps(1) list.simps(8) list.simps(9) map-append
          nth-Cons-0 upt-Suc upt-eq-Nil-conv)
  thus vec-of-list (coeffs h) $ i = coeff h i
    using i h-not0
    unfolding coeffs-def by simp
qed
```

lemma *poly-mod-sum*:
fixes *x y z* :: *'b*::*field poly*
assumes *f*: *finite A*
shows *sum f A mod z* = *sum* (λ*i*. *f i mod z*) *A*
using *f*
by (*induct*, *auto simp add*: *poly-mod-add-left*)

lemma *prime-not-dvd-fact*:
assumes *kn*: *k* < *n* **and** *prime-n*: *prime n*
shows ¬ *n dvd fact k*
using *kn*
proof (*induct k*)
case *0*
thus *?case* **using** *prime-n* **unfolding** *prime-nat-iff* **by** *auto*
next
case (*Suc k*)
show *?case*
proof (*rule ccontr*, *unfold not-not*)
assume *n dvd fact* (*Suc k*)
also have ... = *Suc k* ∗ ∏ {*1*..*k*} **unfolding** *fact-Suc* **unfolding** *fact-prod* **by** *simp*
finally have *n dvd Suc k* ∗ ∏ {*1*..*k*} **.**
hence *n dvd Suc k* ∨ *n dvd* ∏ {*1*..*k*} **using** *prime-dvd-mult-eq-nat*[*OF prime-n*] **by** *blast*
moreover have ¬ *n dvd Suc k* **by** (*simp add*: *Suc.prems*(*1*) *nat-dvd-not-less*)
moreover hence ¬ *n dvd* ∏ {*1*..*k*} **using** *Suc.hyps Suc.prems*
using *Suc-lessD fact-prod*[*of k*] **by** (*metis of-nat-id*)
ultimately show *False* **by** *simp*
qed
qed

lemma *dvd-choose-prime*:
assumes *kn*: *k* < *n* **and** *k*: *k* ≠ *0* **and** *n*: *n* ≠ *0* **and** *prime-n*: *prime n*
shows *n dvd* (*n choose k*)
proof −
have *n dvd* (*fact n*) **by** (*simp add*: *fact-num-eq-if n*)
moreover have ¬ *n dvd* (*fact k* ∗ *fact* (*n*−*k*))
proof (*rule ccontr*, *simp*)
assume *n dvd fact k* ∗ *fact* (*n* − *k*)
hence *n dvd fact k* ∨ *n dvd fact* (*n* − *k*) **using** *prime-dvd-mult-eq-nat*[*OF prime-n*] **by** *simp*
moreover have ¬ *n dvd* (*fact k*) **by** (*rule prime-not-dvd-fact*[*OF kn prime-n*])
moreover have ¬ *n dvd fact* (*n* − *k*) **using** *prime-not-dvd-fact*[*OF - prime-n*] *kn k* **by** *simp*
ultimately show *False* **by** *simp*
qed
moreover have (*fact n*::*nat*) = *fact k* ∗ *fact* (*n*−*k*) ∗ (*n choose k*)

using *binomial-fact-lemma kn* **by** *auto*
ultimately show *?thesis* **using** *prime-n*
by (*auto simp add*: *prime-dvd-mult-iff*)
qed

lemma *add-power-poly-mod-ring*:
fixes x :: *'a mod-ring poly*
shows $(x + y) \hat{} CARD('a) = x \hat{} CARD('a) + y \hat{} CARD('a)$
proof −
let *?A*=$\{0..CARD('a)\}$
let *?f*=λk. *of-nat* $(CARD('a)$ *choose* $k) * x \hat{} k * y \hat{} (CARD('a) - k)$
have *A-rw*: *?A* = *insert* $CARD('a)$ (*insert 0* (*?A* − $\{0\}$ − $\{CARD('a)\}$))
by *fastforce*
have *sum0*: *sum ?f* (*?A* − $\{0\}$ − $\{CARD('a)\}$) = *0*
proof (*rule sum.neutral*, *rule*)
fix *xa* **assume** *xa*: *xa* ∈ $\{0..CARD('a)\}$ − $\{0\}$ − $\{CARD('a)\}$
have *card-dvd-choose*: $CARD('a)$ *dvd* $(CARD('a)$ *choose xa*)
proof (*rule dvd-choose-prime*)
show *xa* < $CARD('a)$ **using** *xa* **by** *simp*
show *xa* ≠ *0* **using** *xa* **by** *simp*
show $CARD('a) \neq 0$ **by** *simp*
show *prime* $CARD('a)$ **by** (*rule prime-card*)
qed
hence *rw0*: *of-int* $(CARD('a)$ *choose xa*) = (*0* :: *'a mod-ring*)
by *transfer simp*
have *of-nat* $(CARD('a)$ *choose xa*) = [:*of-int* $(CARD('a)$ *choose xa*) :: *'a mod-ring*:]
by (*simp add*: *of-nat-poly*)
also have ... = [:*0*:] **using** *rw0* **by** *simp*
finally show *of-nat* $(CARD('a)$ *choose xa*) $* x \hat{} xa * y \hat{} (CARD('a) - xa)$ = *0* **by** *auto*
qed
have $(x + y) \hat{} CARD('a)$
= $(\sum k = 0..CARD('a).$ *of-nat* $(CARD('a)$ *choose* $k) * x \hat{} k * y \hat{} (CARD('a) - k))$
unfolding *binomial-ring* **by** (*rule sum.cong*, *auto*)
also have ... = *sum ?f* (*insert* $CARD('a)$ (*insert 0* (*?A* − $\{0\}$ − $\{CARD('a)\}$)))
using *A-rw* **by** *simp*
also have ... = *?f 0* + *?f* $CARD('a)$ + *sum ?f* (*?A* − $\{0\}$ − $\{CARD('a)\}$) **by** *auto*
also have ... = $x \hat{} CARD('a) + y \hat{} CARD('a)$ **unfolding** *sum0* **by** *auto*
finally show *?thesis* **.**
qed

lemma *power-poly-sum-mod-ring*:
fixes f :: *'b* ⇒ *'a mod-ring poly*

assumes f: *finite A*
shows $(sum\ f\ A)\ \hat{}\ CARD('a) = sum\ (\lambda i.\ (f\ i)\ \hat{}\ CARD('a))\ A$
using f **by** (*induct*, *auto simp add*: *add-power-poly-mod-ring*)

lemma *poly-power-card-as-sum-of-monoms*:
fixes $h :: 'a\ mod\text{-}ring\ poly$
shows $h\ \hat{}\ CARD('a) = (\sum i{\leq}degree\ h.\ monom\ (coeff\ h\ i)\ (CARD('a){*}i))$
proof −
have $h\ \hat{}\ CARD('a) = (\sum i{\leq}degree\ h.\ monom\ (coeff\ h\ i)\ i)\ \hat{}\ CARD('a)$
by (*simp add*: *poly-as-sum-of-monoms*)
also have $... = (\sum i{\leq}degree\ h.\ (monom\ (coeff\ h\ i)\ i)\ \hat{}\ CARD('a))$
by (*simp add*: *power-poly-sum-mod-ring*)
also have $... = (\sum i{\leq}degree\ h.\ monom\ (coeff\ h\ i)\ (CARD('a){*}i))$
proof (*rule sum.cong*, *rule*)
fix x **assume** x: $x \in \{..degree\ h\}$
show $monom\ (coeff\ h\ x)\ x\ \hat{}\ CARD('a) = monom\ (coeff\ h\ x)\ (CARD('a) * x)$
by (*unfold poly-eq-iff*, *auto simp add*: *monom-power*)
qed
finally show *?thesis* **.**
qed

lemma *degree-Poly-berlekamp-le*:
assumes i: $i < degree\ u$
shows $degree\ (Poly\ (list\text{-}of\text{-}vec\ (row\ (berlekamp\text{-}mat\ u)\ i))) < degree\ u$
by (*metis Poly-berlekamp-mat degree-0 degree-mod-less gr-implies-not0 i linorder-neqE-nat*)

lemma *monom-card-pow-mod-sum-berlekamp*:
assumes i: $i < degree\ u$
shows $monom\ 1\ (CARD('a) * i)\ mod\ u = (\sum j{<}degree\ u.\ monom\ ((berlekamp\text{-}mat\ u)\ \$\$\ (i,j))\ j)$
proof −
let *?p* = $Poly\ (list\text{-}of\text{-}vec\ (row\ (berlekamp\text{-}mat\ u)\ i))$
have *degree-not-0*: $degree\ u \neq 0$ **using** i **by** *simp*
hence *set-rw*: $\{..degree\ u - 1\} = \{..{<}degree\ u\}$ **by** *auto*
have *degree-le*: $degree\ ?p < degree\ u$
by (*rule degree-Poly-berlekamp-le*[*OF i*])
hence *degree-le2*: $degree\ ?p \leq degree\ u - 1$ **by** *auto*
have $monom\ 1\ (CARD('a) * i)\ mod\ u = [:0, 1:]\ \hat{}\ (CARD('a) * i)\ mod\ u$
using *x-as-monom x-pow-n* **by** *metis*
also have $... = ?p$
unfolding *Poly-berlekamp-mat*[*OF i*] **by** *simp*
also have $... = (\sum i{\leq}degree\ u - 1.\ monom\ (coeff\ ?p\ i)\ i)$
using *degree-le2 poly-as-sum-of-monoms'* **by** *fastforce*
also have $... = (\sum i{<}degree\ u.\ monom\ (coeff\ ?p\ i)\ i)$ **using** *set-rw* **by** *auto*

```
  also have ... = (∑ j<degree u. monom ((berlekamp-mat u) $$ (i,j)) j)
  proof (rule sum.cong, rule)
    fix x assume x: x ∈ {..<degree u}
    have coeff ?p x = berlekamp-mat u $$ (i, x)
    proof (rule coeff-Poly-list-of-vec-nth)
      show x < dim-col (berlekamp-mat u) using x by auto
    qed
    thus monom (coeff ?p x) x = monom (berlekamp-mat u $$ (i, x)) x
      by (simp add: poly-eq-iff)
  qed
  finally show ?thesis .
qed



lemma col-scalar-prod-as-sum:
assumes dim-vec v = dim-row A
shows col A j · v = (∑ i = 0..<dim-vec v. A $$ (i,j) * v $ i)
  using assms
  unfolding col-def scalar-prod-def
  by transfer' (rule sum.cong, transfer', auto simp add: mk-vec-def mk-mat-def )

lemma row-transpose-scalar-prod-as-sum:
assumes j: j < dim-col A and dim-v: dim-vec v = dim-row A
shows row (transpose-mat A) j · v = (∑ i = 0..<dim-vec v. A $$ (i,j) * v $ i)
proof −
  have row (transpose-mat A) j · v = col A j · v using j row-transpose by auto
  also have ... = (∑ i = 0..<dim-vec v. A $$ (i,j) * v $ i)
    by (rule col-scalar-prod-as-sum[OF dim-v])
  finally show ?thesis .
qed



lemma poly-as-sum-eq-monoms:
assumes ss-eq: (∑ i<n. monom (f i) i) = (∑ i<n. monom (g i) i)
and a-less-n: a<n
shows f a = g a
proof −
  let ?f=λi. if i = a then f i else 0
  let ?g=λi. if i = a then g i else 0
  have sum-f-0: sum ?f ({..<n} − {a}) = 0 by (rule sum.neutral, auto)
  have coeff (∑ i<n. monom (f i) i) a = coeff (∑ i<n. monom (g i) i) a
    using ss-eq unfolding poly-eq-iff by simp
  hence (∑ i<n. coeff (monom (f i) i) a) = (∑ i<n. coeff (monom (g i) i) a)
    by (simp add: coeff-sum)
  hence 1: (∑ i<n. if i = a then f i else 0) = (∑ i<n. if i = a then g i else 0)
    unfolding coeff-monom by auto
  have set-rw: {..<n} = (insert a ({..<n} − {a})) using a-less-n by auto
  have (∑ i<n. if i = a then f i else 0) = sum ?f (insert a ({..<n} − {a}))
```

using *set-rw* **by** *auto*
also have ... = *?f a* + *sum ?f* ({..<*n*} − {*a*})
by (*simp add*: *sum.insert-remove*)
also have ... = *?f a* **using** *sum-f-0* **by** *simp*
finally have *2*: (∑ *i*<*n*. *if i* = *a then f i else 0*) = *?f a* **.**
have *sum ?g* {..<*n*} = *sum ?g* (*insert a* ({..<*n*} − {*a*}))
using *set-rw* **by** *auto*
also have ... = *?g a* + *sum ?g* ({..<*n*} − {*a*})
by (*simp add*: *sum.insert-remove*)
also have ... = *?g a* **using** *sum-f-0* **by** *simp*
finally have *3*: (∑ *i*<*n*. *if i* = *a then g i else 0*) = *?g a* **.**
show *?thesis* **using** *1 2 3* **by** *auto*
qed

lemma *dim-vec-of-list-h*:
assumes *degree h* < *degree u*
shows *dim-vec* (*vec-of-list* ((*coeffs h*) @ *replicate* (*degree u* − *length* (*coeffs h*)) *0*))
= *degree u*
proof −
have *length* (*coeffs h*) ≤ *degree u*
by (*metis Suc-leI assms coeffs-0-eq-Nil degree-0 length-coeffs-degree*
list.size(*3*) *not-le-imp-less order.asym*)
thus *?thesis* **by** *simp*
qed

lemma *vec-of-list-coeffs-nth'*:
assumes *i*: *i* ∈ {..*degree h*} **and** *h-not0*: *h* ≠ *0*
assumes *degree h* < *degree u*
shows *vec-of-list* ((*coeffs h*) @ *replicate* (*degree u* − *length* (*coeffs h*)) *0*) $ *i* =
coeff h i
using *assms*
by (*transfer'*, *auto simp add*: *mk-vec-def coeffs-nth length-coeffs-degree nth-append*)

lemma *vec-of-list-coeffs-replicate-nth-0*:
assumes *i*: *i* ∈ {..<*degree u*}
shows *vec-of-list* (*coeffs 0* @ *replicate* (*degree u* − *length* (*coeffs 0*)) *0*) $ *i* = *coeff*
0 i
using *assms*
by (*transfer'*, *auto simp add*: *mk-vec-def*)

lemma *vec-of-list-coeffs-replicate-nth*:
assumes *i*: *i* ∈ {..<*degree u*}

assumes *degree h < degree u*
shows *vec-of-list ((coeffs h) @ replicate (degree u − length (coeffs h)) 0) $ i = coeff h i*
proof (*cases h = 0*)
case *True*
thus *?thesis* **using** *vec-of-list-coeffs-replicate-nth-0 i* **by** *auto*
next
case *False* **note** *h-not0 = False*
show *?thesis*
proof (*cases i ∈{..degree h}*)
case *True* **thus** *?thesis* **using** *assms vec-of-list-coeffs-nth′ h-not0* **by** *simp*
next
case *False*
have *c0*: *coeff h i = 0* **using** *False le-degree* **by** *auto*
thus *?thesis*
using *assms False h-not0*
by (*transfer′, auto simp add*: *mk-vec-def length-coeffs-degree nth-append c0*)
qed
qed

lemma *equation-13*:
fixes *u h*
defines *H*: *H ≡ vec-of-list ((coeffs h) @ replicate (degree u − length (coeffs h)) 0)*
assumes *deg-le*: *degree h < degree u*
shows *[h^CARD('a) = h] (mod u) ⟷ (transpose-mat (berlekamp-mat u)) ∗v H = H*
(**is** *?lhs = ?rhs*)
proof −
have *f*: *finite {..degree u}* **by** *auto*
have [*simp*]: *dim-vec H = degree u* **unfolding** *H* **using** *dim-vec-of-list-h deg-le* **by** *simp*
let *?B = (berlekamp-mat u)*
let *?f = λi. (transpose-mat ?B ∗v H) $ i*
show *?thesis*
proof
assume *rhs*: *?rhs*
have *dimv-h-dimr-B*: *dim-vec H = dim-row ?B*
by (*metis berlekamp-mat-closed(2) berlekamp-mat-closed(3) dim-mult-mat-vec index-transpose-mat(2) rhs*)
have *degree-h-less-dim-H*: *degree h < dim-vec H* **by** (*auto simp add*: *deg-le*)
have *set-rw*: *{..degree u − 1} = {..<degree u}* **using** *deg-le* **by** *auto*
have *degree h ≤ degree u − 1* **using** *deg-le* **by** *simp*
hence *h = (∑ j≤degree u − 1. monom (coeff h j) j)* **using** *poly-as-sum-of-monoms′* **by** *fastforce*
also have *... = (∑ j<degree u. monom (coeff h j) j)* **using** *set-rw* **by** *simp*

also have ... = ($\sum j$<*degree u*. *monom* (*?f j*) *j*)
proof (*rule sum.cong*, *rule*+)
fix *j* **assume** *i*: *j* ∈ {..<*degree u*}
have (*coeff h j*) = *?f j*
using *rhs vec-of-list-coeffs-replicate-nth*[*OF i deg-le*]
unfolding *H* **by** *presburger*
thus *monom* (*coeff h j*) *j* = *monom* (*?f j*) *j*
by *simp*
qed
also have ... = ($\sum j$<*degree u*. *monom* (*row* (*transpose-mat ?B*) *j* · *H*) *j*)
by (*rule sum.cong*, *auto*)
also have ... = ($\sum j$<*degree u*. *monom* ($\sum i$ = *0*..<*dim-vec H*. *?B* $$ (*i*,*j*) * *H* $ *i*) *j*)
proof (*rule sum.cong*, *rule*)
fix *x* **assume** *x*: *x* ∈ {..<*degree u*}
show *monom* (*row* (*transpose-mat ?B*) *x* · *H*) *x* =
monom ($\sum i$ = *0*..<*dim-vec H*. *?B* $$ (*i*, *x*) * *H* $ *i*) *x*
proof (*unfold monom-eq-iff*, *rule row-transpose-scalar-prod-as-sum*[*OF - dimv-h-dimr-B*])
show *x* < *dim-col ?B* **using** *x deg-le* **by** *auto*
qed
qed
also have ... = ($\sum j$<*degree u*. $\sum i$ = *0*..<*dim-vec H*. *monom* (*?B* $$ (*i*,*j*) * *H* $ *i*) *j*)
by (*auto simp add*: *monom-sum*)
also have ... = ($\sum i$ = *0*..<*dim-vec H*. $\sum j$<*degree u*. *monom* (*?B* $$ (*i*,*j*) * *H* $ *i*) *j*)
by (*rule sum.swap*)
also have ... = ($\sum i$ = *0*..<*dim-vec H*. $\sum j$<*degree u*. *monom* (*H* $ *i*) *0* * *monom* (*?B* $$ (*i*,*j*)) *j*)
proof (*rule sum.cong*, *rule*, *rule sum.cong*, *rule*)
fix *x xa*
show *monom* (*?B* $$ (*x*, *xa*) * *H* $ *x*) *xa* = *monom* (*H* $ *x*) *0* * *monom* (*?B* $$ (*x*, *xa*)) *xa*
by (*simp add*: *mult-monom*)
qed
also have ... = ($\sum i$ = *0*..<*dim-vec H*. (*monom* (*H* $ *i*) *0*) * ($\sum j$<*degree u*. *monom* (*?B* $$ (*i*,*j*)) *j*))
by (*rule sum.cong*, *auto simp*: *sum-distrib-left*)
also have ... = ($\sum i$ = *0*..<*dim-vec H*. (*monom* (*H* $ *i*) *0*) * (*monom 1* (*CARD*('*a*) * *i*) *mod u*))
proof (*rule sum.cong*, *rule*)
fix *x* **assume** *x*: *x* ∈ {*0*..<*dim-vec H*}
have ($\sum j$<*degree u*. *monom* (*?B* $$ (*x*, *j*)) *j*) = (*monom 1* (*CARD*('*a*) * *x*) *mod u*)
proof (*rule monom-card-pow-mod-sum-berlekamp*[*symmetric*])
show *x* < *degree u* **using** *x dimv-h-dimr-B* **by** *auto*
qed
thus *monom* (*H* $ *x*) *0* * ($\sum j$<*degree u*. *monom* (*?B* $$ (*x*, *j*)) *j*) =

```
      monom (H $ x) 0 * (monom 1 (CARD('a) * x) mod u) by presburger
  qed
  also have ... = (∑ i = 0..<dim-vec H. monom (H $ i) (CARD('a) * i) mod
u)
  proof (rule sum.cong, rule)
    fix x
    have h-rw: monom (H $ x) 0 mod u = monom (H $ x) 0
      by (metis deg-le degree-pCons-eq-if gr-implies-not-zero
        linorder-neqE-nat mod-poly-less monom-0)
     have monom (H $ x) (CARD('a) * x) = monom (H $ x) 0 * monom 1
(CARD('a) * x)
      unfolding mult-monom by simp
    also have ... = smult (H $ x) (monom 1 (CARD('a) * x))
      by (simp add: monom-0)
     also have ... mod u =  Polynomial.smult (H $ x) (monom 1 (CARD('a) *
x) mod u)
      using mod-smult-left by auto
    also have ... = monom (H $ x) 0 * (monom 1 (CARD('a) * x) mod u)
      by (simp add: monom-0)
    finally show monom (H $ x) 0 * (monom 1 (CARD('a) * x) mod u)
      = monom (H $ x) (CARD('a) * x) mod u ..
  qed
  also have ... = (∑ i = 0..<dim-vec H. monom (H $ i) (CARD('a) * i)) mod
u
    by (simp add: poly-mod-sum)
   also have ... = (∑ i = 0..<dim-vec H. monom (coeff h i) (CARD('a) * i))
mod u
  proof (rule arg-cong[of - - λx. x  mod u], rule sum.cong, rule)
     fix x assume x: x ∈ {0..<dim-vec H}
     have H $ x = (coeff h x)
     proof (unfold H, rule vec-of-list-coeffs-replicate-nth[OF - deg-le])
       show x ∈ {..<degree u} using x by auto
     qed
    thus monom (H $ x) (CARD('a) * x) = monom (coeff h x) (CARD('a) * x)
      by simp
  qed
  also have ... = (∑ i≤degree h. monom (coeff h i) (CARD('a) * i)) mod u
  proof (rule arg-cong[of - - λx. x mod u])
    let ?f=λi. monom (coeff h i) (CARD('a) * i)
    have ss0: (∑ i = degree h + 1 ..< dim-vec H. ?f i) = 0
      by (rule sum.neutral, simp add: coeff-eq-0)
    have set-rw: {0..< dim-vec H} = {0..degree h} ∪ {degree h + 1 ..< dim-vec
H}
      using degree-h-less-dim-H by auto
     have (∑ i = 0..<dim-vec H. ?f i) = (∑ i = 0..degree h. ?f i) + (∑ i =
degree h + 1 ..< dim-vec H. ?f i)
      unfolding set-rw by (rule sum.union-disjoint, auto)
    also have ... = (∑ i = 0..degree h. ?f i) unfolding ss0 by auto
    finally show (∑ i = 0..<dim-vec H. ?f i) = (∑ i≤degree h. ?f i)
```

```
      by (simp add: atLeast0AtMost)
    qed
    also have ... = h^CARD('a) mod u
      using poly-power-card-as-sum-of-monoms by auto
    finally show ?lhs
      unfolding cong-def
      using deg-le
      by (simp add: mod-poly-less)
next
  assume lhs: ?lhs
  have deg-le': degree h ≤ degree u − 1 using deg-le by auto
  have set-rw: {..<degree u} = {..degree u −1} using deg-le by auto
  hence (∑ i<degree u. monom (coeff h i) i) = (∑ i ≤ degree u − 1. monom
(coeff h i) i) by simp
  also have ... = (∑ i≤degree h. monom (coeff h i) i)
    unfolding poly-as-sum-of-monoms
    using poly-as-sum-of-monoms' deg-le' by auto
  also have ... = (∑ i≤degree h. monom (coeff h i) i) mod u
    by (simp add: deg-le mod-poly-less poly-as-sum-of-monoms)
  also have ... = (∑ i≤degree h. monom (coeff h i) (CARD('a)∗i)) mod u
     using lhs
    unfolding cong-def poly-as-sum-of-monoms poly-power-card-as-sum-of-monoms
     by auto
  also have ... = (∑ i≤degree h. monom (coeff h i) 0 ∗ monom 1 (CARD('a)∗i))
mod u
   by (rule arg-cong[of - - λx. x mod u], rule sum.cong, simp-all add: mult-monom)
  also have ... = (∑ i≤degree h. monom (coeff h i) 0 ∗ monom 1 (CARD('a)∗i)
mod u)
    by (simp add: poly-mod-sum)
  also have ... = (∑ i≤degree h. monom (coeff h i) 0 ∗ (monom 1 (CARD('a)∗i)
mod u))
  proof (rule sum.cong, rule)
    fix x assume x: x ∈ {..degree h}
    have h-rw: monom (coeff h x) 0 mod u = monom (coeff h x) 0
      by (metis deg-le degree-pCons-eq-if gr-implies-not-zero
        linorder-neqE-nat mod-poly-less monom-0)
     have monom (coeff h x) 0 ∗ monom 1 (CARD('a) ∗ x) = smult (coeff h x)
(monom 1 (CARD('a) ∗ x))
      by (simp add: monom-0)
     also have ... mod u =  Polynomial.smult (coeff h x) (monom 1 (CARD('a)
∗ x) mod u)
      using mod-smult-left by auto
     also have ... = monom (coeff h x) 0 ∗ (monom 1 (CARD('a) ∗ x) mod u)
      by (simp add: monom-0)
    finally show monom (coeff h x) 0 ∗ monom 1 (CARD('a) ∗ x) mod u
     = monom (coeff h x) 0 ∗ (monom 1 (CARD('a) ∗ x) mod u) .
  qed
  also have ... = (∑ i≤degree h. monom (coeff h i) 0 ∗ (∑ j<degree u. monom
(?B $$ (i, j)) j))
```

proof (*rule sum.cong*, *rule*)
fix x **assume** x: $x \in \{..degree\ h\}$
have $(monom\ 1\ (CARD('a) * x)\ mod\ u) = (\sum j{<}degree\ u.\ monom\ (?B\ \$\$\ (x, j))\ j)$
proof (*rule monom-card-pow-mod-sum-berlekamp*)
show $x < degree\ u$ **using** x *deg-le* **by** *auto*
qed
thus $monom\ (coeff\ h\ x)\ 0 * (monom\ 1\ (CARD('a) * x)\ mod\ u) =$
$monom\ (coeff\ h\ x)\ 0 * (\sum j{<}degree\ u.\ monom\ (?B\ \$\$\ (x, j))\ j)$ **by** *simp*
qed
also have $... = (\sum i{<}degree\ u.\ monom\ (coeff\ h\ i)\ 0 * (\sum j{<}degree\ u.\ monom\ (?B\ \$\$\ (i, j))\ j))$
proof −
let $?f{=}\lambda i.\ monom\ (coeff\ h\ i)\ 0 * (\sum j{<}degree\ u.\ monom\ (?B\ \$\$\ (i, j))\ j)$
have *ss0*: $(\sum i{=}degree\ h{+}1\ ..{<}\ degree\ u.\ ?f\ i) = 0$
by (*rule sum.neutral*, *simp add*: *coeff-eq-0*)
have *set-rw*: $\{0..{<}degree\ u\} = \{0..degree\ h\} \cup \{degree\ h{+}1..{<}degree\ u\}$ **using** *deg-le* **by** *auto*
have $(\sum i{=}0..{<}degree\ u.\ ?f\ i) = (\sum i{=}0..degree\ h.\ ?f\ i) + (\sum i{=}degree\ h{+}1\ ..{<}\ degree\ u.\ ?f\ i)$
unfolding *set-rw* **by** (*rule sum.union-disjoint*, *auto*)
also have $... = (\sum i{=}0..degree\ h.\ ?f\ i)$ **using** *ss0* **by** *simp*
finally show *?thesis*
by (*simp add*: *atLeast0AtMost atLeast0LessThan*)
qed
also have $... = (\sum i{<}degree\ u.\ (\sum j{<}degree\ u.\ monom\ (coeff\ h\ i)\ 0 * monom\ (?B\ \$\$\ (i, j))\ j))$
by (*simp add*: *sum-distrib-left*)
also have $... = (\sum i{<}degree\ u.\ (\sum j{<}degree\ u.\ monom\ (coeff\ h\ i * ?B\ \$\$\ (i, j))\ j))$
by (*simp add*: *mult-monom*)
also have $... = (\sum j{<}degree\ u.\ (\sum i{<}degree\ u.\ monom\ (coeff\ h\ i * ?B\ \$\$\ (i, j))\ j))$
using *sum.swap* **by** *auto*
also have $... = (\sum j{<}degree\ u.\ monom\ (\sum i{<}degree\ u.\ \ (coeff\ h\ i * ?B\ \$\$\ (i, j)))\ j)$
by (*simp add*: *monom-sum*)
finally have *ss-rw*: $(\sum i{<}degree\ u.\ monom\ (coeff\ h\ i)\ i)$
$= (\sum j{<}degree\ u.\ monom\ (\sum i{<}degree\ u.\ coeff\ h\ i * ?B\ \$\$\ (i, j))\ j)$ **.**
have *coeff-eq-sum*: $\forall i.\ i < degree\ u \longrightarrow coeff\ h\ i = (\sum j{<}degree\ u.\ coeff\ h\ j * ?B\ \$\$\ (j, i))$
using *poly-as-sum-eq-monoms*[*OF ss-rw*] **by** *fast*
have *coeff-eq-sum'*: $\forall i.\ i < degree\ u \longrightarrow H\ \$\ i = (\sum j{<}degree\ u.\ H\ \$\ j * ?B\ \$\$\ (j, i))$
proof (*rule+*)
fix i **assume** i: $i < degree\ u$
have $H\ \$\ i = coeff\ h\ i$ **by** (*simp add*: *H deg-le i vec-of-list-coeffs-replicate-nth*)
also have $... = (\sum j{<}degree\ u.\ coeff\ h\ j * ?B\ \$\$\ (j, i))$ **using** *coeff-eq-sum i* **by** *blast*

```
    also have ... = (∑j<degree u. H $ j * ?B $$ (j, i))
      by (rule sum.cong, auto simp add: H deg-le vec-of-list-coeffs-replicate-nth)
    finally show H $ i = (∑j<degree u. H $ j * ?B $$ (j, i)) .
  qed
  show (transpose-mat (?B)) *v H = H
  proof (rule eq-vecI)
    fix i
    show dim-vec (transpose-mat ?B *v H) = dim-vec (H) by auto
    assume i: i < dim-vec (H)
    have (transpose-mat ?B *v H) $ i = row (transpose-mat ?B) i · H using i by
simp
    also have ... = (∑j = 0..<dim-vec H. ?B $$ (j, i) * H $ j)
    proof (rule row-transpose-scalar-prod-as-sum)
      show i < dim-col ?B using i by simp
      show dim-vec H = dim-row ?B by simp
    qed
    also have ... = (∑j<degree u. H $ j * ?B $$ (j, i)) by (rule sum.cong, auto)
     also have ... = H $ i using coeff-eq-sum'[rule-format, symmetric, of i] i by
simp
    finally show (transpose-mat ?B *v H) $ i = H $ i .
  qed
 qed
qed


end


context
assumes SORT-CONSTRAINT('a::prime-card)
begin

lemma exists-s-factor-dvd-h-s:
fixes fi::'a mod-ring poly
assumes finite-P: finite P
      and f-desc-square-free: f = (∏ a∈P. a)
      and P: P ⊆ {q. irreducible q ∧ monic q}
      and fi-P: fi ∈ P
      and h: h ∈ {v. [v^(CARD('a)) = v] (mod f)}
      shows ∃ s. fi dvd (h − [:s:])
proof −
  let ?p = CARD('a)
       have f-dvd-hqh: f dvd (h^?p − h) using h unfolding cong-def
        using mod-eq-dvd-iff-poly by blast
      also have hq-h-rw: ... = prod (λc. h − [:c:]) (UNIV::'a mod-ring set)
        by (rule poly-identity-mod-p)
      finally have f-dvd-hc: f dvd prod (λc. h − [:c:]) (UNIV::'a mod-ring set) by
simp
          have fi dvd f using f-desc-square-free fi-P
```

using *dvd-prod-eqI finite-P* **by** *blast*
hence *fi dvd* $(h \hat{} ?p - h)$ **using** *dvd-trans f-dvd-hqh* **by** *auto*
also have ... = *prod* $(\lambda c.\ h - [:c:])$ $(UNIV::'a\ mod\text{-}ring\ set)$ **unfolding** *hq-h-rw* **by** *simp*
finally have *fi-dvd-prod-hc*: *fi dvd prod* $(\lambda c.\ h - [:c:])$ $(UNIV::'a\ mod\text{-}ring\ set)$ **.**
have *irr-fi*: *irreducible fi* **using** *fi-P P* **by** *blast*
have *fi-not-unit*: ¬ *is-unit fi* **using** *irr-fi* **by** (*simp add*: $irreducible_d D(1)$ *poly-dvd-1*)
show *?thesis* **using** *irreducible-dvd-prod*[*OF - fi-dvd-prod-hc*] *irr-fi* **by** *auto*
qed

corollary *exists-unique-s-factor-dvd-h-s*:
fixes *fi*::$'a$ *mod-ring poly*
assumes *finite-P*: *finite P*
and *f-desc-square-free*: $f = (\prod a{\in}P.\ a)$
and *P*: $P \subseteq \{q.\ irreducible\ q \wedge monic\ q\}$
and *fi-P*: $fi \in P$
and *h*: $h \in \{v.\ [v \hat{} (CARD('a)) = v]\ (mod\ f)\}$
shows $\exists!s.$ *fi dvd* $(h - [:s:])$
proof −
obtain *c* **where** *fi-dvd*: *fi dvd* $(h - [:c:])$ **using** *assms exists-s-factor-dvd-h-s* **by** *blast*
have *irr-fi*: *irreducible fi* **using** *fi-P P* **by** *blast*
have *fi-not-unit*: ¬ *is-unit fi*
by (*simp add*: *irr-fi* $irreducible_d D(1)$ *poly-dvd-1*)
show *?thesis*
proof (*rule ex1I*[*of - c*], *auto simp add*: *fi-dvd*)
fix *c2* **assume** *fi-dvd-hc2*: *fi dvd h* − [:*c2*:]
have ∗: *fi dvd* $(h - [:c:]) * (h - [:c2:])$ **using** *fi-dvd* **by** *auto*
hence *fi dvd* $(h - [:c:]) \vee$ *fi dvd* $(h - [:c2:])$
using *irr-fi* **by** *auto*
thus *c2* = *c*
using *coprime-h-c-poly coprime-not-unit-not-dvd fi-dvd fi-dvd-hc2 fi-not-unit* **by** *blast*
qed
qed

lemma *exists-two-distint*: $\exists\ a\ b::'a\ mod\text{-}ring.\ a \neq b$
by (*rule exI*[*of - 0*], *rule exI*[*of - 1*], *auto*)

lemma *coprime-cong-mult-factorization-poly*:
fixes *f*::$'b$::{*field*} *poly*
and *a b p* :: $'c$ :: {*field-gcd*} *poly*
assumes *finite-P*: *finite P*
and *P*: $P \subseteq \{q.\ irreducible\ q\}$

and p: $\forall\, p{\in}P.\ [a{=}b]\ (mod\ p)$
and *coprime-P*: $\forall\, p1\ p2.\ p1 \in P \wedge p2 \in P \wedge p1 \neq p2 \longrightarrow coprime\ p1\ p2$
shows $[a = b]\ (mod\ (\prod a{\in}P.\ a))$
using *finite-P P p coprime-P*
proof (*induct P*)
case *empty*
thus *?case* **by** *simp*
next
case (*insert p P*)
have *ab-mod-pP*: $[a{=}b]\ (mod\ (p * \prod P))$
proof (*rule coprime-cong-mult-poly*)
show $[a = b]\ (mod\ p)$ **using** *insert.prems* **by** *auto*
show $[a = b]\ (mod \prod P)$ **using** *insert.prems insert.hyps* **by** *auto*
from *insert* **show** *Rings.coprime p* $(\prod P)$
by (*auto intro*: *prod-coprime-right*)
qed
thus *?case* **by** (*simp add*: *insert.hyps*(*1*) *insert.hyps*(*2*))
qed

end

context
assumes *SORT-CONSTRAINT*('*a*::*prime-card*)
begin

lemma *W-eq-berlekamp-mat*:
fixes *u*::'*a mod-ring poly*
shows $\{v.\ [v\hat{}CARD('a) = v]\ (mod\ u) \wedge degree\ v < degree\ u\}$
$= \{h.$ *let H* = *vec-of-list* ((*coeffs h*) @ *replicate* (*degree u* − *length* (*coeffs h*)) *0*) *in*
(*transpose-mat* (*berlekamp-mat u*)) $*_v$ *H* = *H* ∧ *degree h* < *degree u*$\}$
using *equation-13* **by** (*auto simp add*: *Let-def*)

lemma *transpose-minus-1*:
assumes *dim-row Q* = *dim-col Q*
shows *transpose-mat* (*Q* − (1_m (*dim-row Q*))) = (*transpose-mat Q* − (1_m (*dim-row Q*)))
using *assms*
unfolding *mat-eq-iff* **by** *auto*

lemma *system-iff*:
fixes *v*::'*b*::*comm-ring-1 vec*
assumes *sq-Q*: *dim-row Q* = *dim-col Q* **and** *v*: *dim-row Q* = *dim-vec v*
shows (*transpose-mat Q* $*_v$ *v* = *v*) ⟷ ((*transpose-mat Q* − 1_m (*dim-row Q*)) $*_v$ *v* = 0_v (*dim-vec v*))
proof −
have *t1*:*transpose-mat Q* $*_v$ *v* − *v* = 0_v (*dim-vec v*) ⟹ (*transpose-mat Q* −

1_m (*dim-row Q*)) $*_v$ v = 0_v (*dim-vec v*)
by (*subst minus-mult-distrib-mat-vec*, *insert sq-Q*[*symmetric*] *v*, *auto*)
have *t2*:(*transpose-mat Q* − 1_m (*dim-row Q*)) $*_v$ v = 0_v (*dim-vec v*) ⟹ *transpose-mat Q* $*_v$ v − v = 0_v (*dim-vec v*)
by (*subst* (*asm*) *minus-mult-distrib-mat-vec*, *insert sq-Q*[*symmetric*] *v*, *auto*)
have *transpose-mat Q* $*_v$ v − v = v − v ⟹ *transpose-mat Q* $*_v$ v = v
proof −
assume *a1*: *transpose-mat Q* $*_v$ v − v = v − v
have *f2*: *transpose-mat Q* $*_v$ v ∈ *carrier-vec* (*dim-vec v*)
by (*metis dim-mult-mat-vec index-transpose-mat*(*2*) *sq-Q v carrier-vec-dim-vec*)
then have *f3*: 0_v (*dim-vec v*) + *transpose-mat Q* $*_v$ v = *transpose-mat Q* $*_v$ v
by (*meson left-zero-vec*)
have *f4*: 0_v (*dim-vec v*) = *transpose-mat Q* $*_v$ v − v
using *a1* **by** *auto*
have *f5*: − v ∈ *carrier-vec* (*dim-vec v*)
by *simp*
then have *f6*: − v + *transpose-mat Q* $*_v$ v = v − v
using *f2 a1* **using** *comm-add-vec minus-add-uminus-vec* **by** *fastforce*
have v − v = − v + v **by** *auto*
then have *transpose-mat Q* $*_v$ v = *transpose-mat Q* $*_v$ v − v + v
using *f6 f4 f3 f2* **by** (*metis* (*no-types*, *lifting*) *a1 assoc-add-vec comm-add-vec f5 carrier-vec-dim-vec*)
then show *?thesis*
using *a1* **by** *auto*
qed
hence (*transpose-mat Q* $*_v$ v = v) = ((*transpose-mat Q* $*_v$ v) − v = v − v) **by** *auto*
also have ... = ((*transpose-mat Q* $*_v$ v) − v = 0_v (*dim-vec v*)) **by** *auto*
also have ... = ((*transpose-mat Q* − 1_m (*dim-row Q*)) $*_v$ v = 0_v (*dim-vec v*))
using *t1 t2* **by** *auto*
finally show *?thesis*.
qed

lemma *system-if-mat-kernel*:
assumes *sq-Q*: *dim-row Q* = *dim-col Q* **and** *v*: *dim-row Q* = *dim-vec v*
shows (*transpose-mat Q* $*_v$ v = v) ⟷ v ∈ *mat-kernel* (*transpose-mat* (*Q* − (1_m (*dim-row Q*))))
proof −
have (*transpose-mat Q* $*_v$ v = v) = ((*transpose-mat Q* − 1_m (*dim-row Q*)) $*_v$ v = 0_v (*dim-vec v*))
using *assms system-iff* **by** *blast*
also have ... = (v ∈ *mat-kernel* (*transpose-mat* (*Q* − (1_m (*dim-row Q*)))))
unfolding *mat-kernel-def* **unfolding** *transpose-minus-1*[*OF sq-Q*] **unfolding** *v* **by** *auto*
finally show *?thesis* .
qed

lemma *degree-u-mod-irreducible$_d$-factor-0*:
fixes *v* **and** *u*::*'a mod-ring poly*
defines *W*: $W \equiv \{v.\ [v\ \widehat{}\ CARD('a) = v]\ (mod\ u)\}$
assumes *v*: $v \in W$
and *finite-U*: *finite U* **and** *u-U*: $u = \prod U$ **and** *U-irr-monic*: $U \subseteq \{q.\ irreducible\ q \wedge monic\ q\}$
and *fi-U*: $fi \in U$
shows *degree* $(v\ mod\ fi) = 0$
proof −
have *deg-fi*: *degree fi* > 0
using *U-irr-monic*
using *fi-U irreducible$_d$D*[*of fi*] **by** *auto*
have *fi dvd u*
using *u-U U-irr-monic finite-U dvd-prod-eqI fi-U* **by** *blast*
moreover have *u dvd* $(v\widehat{}CARD('a) - v)$
using *v* **unfolding** *W cong-def*
by (*simp add*: *mod-eq-dvd-iff-poly*)
ultimately have *fi dvd* $(v\widehat{}CARD('a) - v)$
by (*rule dvd-trans*)
then have *fi-dvd-prod-vc*: *fi dvd prod* $(\lambda c.\ v - [:c:])$ $(UNIV::'a\ mod\text{-}ring\ set)$
by (*simp add*: *poly-identity-mod-p*)
have *irr-fi*: *irreducible fi* **using** *fi-U U-irr-monic* **by** *blast*
have *fi-not-unit*: ¬ *is-unit fi*
using *irr-fi*
by (*auto simp*: *poly-dvd-1*)
have *fi-dvd-vc*: $\exists\, c.\ fi\ dvd\ v - [:c:]$
using *irreducible-dvd-prod*[*OF - fi-dvd-prod-vc*] *irr-fi* **by** *auto*
from *this* **obtain** *a* **where** *fi dvd* $v - [:a:]$ **by** *blast*
hence $v\ mod\ fi = [:a:]\ mod\ fi$ **using** *mod-eq-dvd-iff-poly* **by** *blast*
also have ... $= [:a:]$ **by** (*simp add*: *deg-fi mod-poly-less*)
finally show *?thesis* **by** *simp*
qed

definition *poly-abelian-monoid*
$= (\!|carrier = UNIV::'a\ mod\text{-}ring\ poly\ set,\ monoid.mult = ((*)),\ one = 1,\ zero = 0,\ add = (+),\ module.smult = smult|\!)$

interpretation *vector-space-poly*: *vectorspace class-ring poly-abelian-monoid*
rewrites [*simp*]: $\mathbf{0}_{poly\text{-}abelian\text{-}monoid} = 0$
and [*simp*]: $\mathbf{1}_{poly\text{-}abelian\text{-}monoid} = 1$
and [*simp*]: $(\oplus_{poly\text{-}abelian\text{-}monoid}) = (+)$
and [*simp*]: $(\otimes_{poly\text{-}abelian\text{-}monoid}) = (*)$
and [*simp*]: *carrier poly-abelian-monoid* $= UNIV$
and [*simp*]: $(\odot_{poly\text{-}abelian\text{-}monoid}) = smult$
apply *unfold-locales*

apply (*auto simp*: *poly-abelian-monoid-def class-field-def smult-add-left smult-add-right Units-def*)
by (*metis add.commute add.right-inverse*)

lemma *subspace-Berlekamp*:
assumes *f*: *degree f* ≠ *0*
shows *subspace* (*class-ring* :: *'a mod-ring ring*)
{*v*. [*v*^(*CARD*(*'a*)) = *v*] (*mod f*) ∧ (*degree v* < *degree f*)} *poly-abelian-monoid*
proof −
{ fix *v* :: *'a mod-ring poly* **and** *w* :: *'a mod-ring poly*
assume *a1*: *v* ^ *card* (*UNIV*::*'a set*) *mod f* = *v mod f*
assume *w* ^ *card* (*UNIV*::*'a set*) *mod f* = *w mod f*
then have (*v* ^ *card* (*UNIV*::*'a set*) + *w* ^ *card* (*UNIV*::*'a set*)) *mod f* = (*v* + *w*) *mod f*
using *a1* **by** (*meson mod-add-cong*)
then have (*v* + *w*) ^ *card* (*UNIV*::*'a set*) *mod f* = (*v* + *w*) *mod f*
by (*simp add*: *add-power-poly-mod-ring*)
} note *r*=*this*
thus *?thesis* **using** *f*
by (*unfold-locales*, *auto simp*: *zero-power mod-smult-left smult-power cong-def degree-add-less*)
qed

lemma *berlekamp-resulting-mat-closed*[*simp*]:
berlekamp-resulting-mat u ∈ *carrier-mat* (*degree u*) (*degree u*)
dim-row (*berlekamp-resulting-mat u*) = *degree u*
dim-col (*berlekamp-resulting-mat u*) = *degree u*
proof −
let *?A*=(*transpose-mat* (*mat* (*degree u*) (*degree u*)
(λ(*i*, *j*). *if i* = *j then berlekamp-mat u* $$ (*i*, *j*) − *1 else berlekamp-mat u* $$ (*i*, *j*))))
let *?G*=(*gauss-jordan-single ?A*)
have *?G* ∈*carrier-mat* (*degree u*) (*degree u*)
by (*rule gauss-jordan-single*(*2*)[*of ?A*], *auto*)
thus
berlekamp-resulting-mat u ∈ *carrier-mat* (*degree u*) (*degree u*)
dim-row (*berlekamp-resulting-mat u*) = *degree u*
dim-col (*berlekamp-resulting-mat u*) = *degree u*
unfolding *berlekamp-resulting-mat-def Let-def* **by** *auto*
qed

lemma *berlekamp-resulting-mat-basis*:
kernel.basis (*degree u*) (*berlekamp-resulting-mat u*) (*set* (*find-base-vectors* (*berlekamp-resulting-mat u*)))
proof (*rule find-base-vectors*(*3*))
show *berlekamp-resulting-mat u* ∈ *carrier-mat* (*degree u*) (*degree u*) **by** *simp*

```
  let ?A=(transpose-mat (mat (degree u) (degree u)
        (λ(i, j). if i = j then berlekamp-mat u $$ (i, j) − 1 else berlekamp-mat u
$$ (i, j))))
  have row-echelon-form (gauss-jordan-single ?A)
    by (rule gauss-jordan-single(3)[of ?A], auto)
  thus row-echelon-form (berlekamp-resulting-mat u)
    unfolding berlekamp-resulting-mat-def Let-def by auto
qed


lemma set-berlekamp-basis-eq: (set (berlekamp-basis u))
  = (Poly ∘ list-of-vec)‘ (set (find-base-vectors (berlekamp-resulting-mat u)))
  by (auto simp add: image-def o-def berlekamp-basis-def)


lemma berlekamp-resulting-mat-constant:
assumes deg-u: degree u = 0
shows berlekamp-resulting-mat u = 1_m 0
  by (unfold mat-eq-iff, auto simp add: deg-u)

context
  fixes u::'a::prime-card mod-ring poly
begin

lemma set-berlekamp-basis-constant:
assumes deg-u: degree u = 0
shows set (berlekamp-basis u) = {}
proof −
  have one-carrier: 1_m 0 ∈ carrier-mat 0 0 by auto
 have m: mat-kernel (1_m 0) = {(0_v 0) :: 'a mod-ring vec} unfolding mat-kernel-def
by auto
  have r: row-echelon-form (1_m 0 :: 'a mod-ring mat)
    unfolding row-echelon-form-def pivot-fun-def Let-def by auto
  have set (find-base-vectors (1_m 0)) ⊆ {0_v 0 :: 'a mod-ring vec}
    using find-base-vectors(1)[OF r one-carrier] unfolding m .
  hence set (find-base-vectors (1_m 0) :: 'a mod-ring vec list) = {}
    using find-base-vectors(2)[OF r one-carrier]
    using subset-singletonD by fastforce
  thus ?thesis
   unfolding set-berlekamp-basis-eq unfolding berlekamp-resulting-mat-constant[OF
deg-u] by auto
qed


lemma row-echelon-form-berlekamp-resulting-mat: row-echelon-form (berlekamp-resulting-mat
u)
 by (rule gauss-jordan-single(3), auto simp add: berlekamp-resulting-mat-def Let-def)

lemma mat-kernel-berlekamp-resulting-mat-degree-0:
```

assumes *d*: *degree u = 0*
shows *mat-kernel* (*berlekamp-resulting-mat u*) = {0_v *0*}
by (*auto simp add*: *mat-kernel-def mult-mat-vec-def d*)

lemma *in-mat-kernel-berlekamp-resulting-mat*:
assumes *x*: *transpose-mat* (*berlekamp-mat u*) $*_v$ *x = x*
and *x-dim*: *x* ∈ *carrier-vec* (*degree u*)
shows *x* ∈ *mat-kernel* (*berlekamp-resulting-mat u*)
proof −
let *?QI*=(*mat*(*dim-row* (*berlekamp-mat u*)) (*dim-row* (*berlekamp-mat u*))
(λ(*i*, *j*). *if i = j then berlekamp-mat u* $$ (*i*, *j*) − *1 else berlekamp-mat u* $$ (*i*, *j*)))
have ∗: (*transpose-mat* (*berlekamp-mat u*) − 1_m (*degree u*)) = *transpose-mat ?QI* **by** *auto*
have (*transpose-mat* (*berlekamp-mat u*) − 1_m (*dim-row* (*berlekamp-mat u*))) $*_v$ *x* = 0_v (*dim-vec x*)
using *system-iff*[*of berlekamp-mat u x*] *x-dim x* **by** *auto*
hence *transpose-mat ?QI* $*_v$ *x* = 0_v (*degree u*) **using** *x-dim* ∗ **by** *auto*
hence *berlekamp-resulting-mat u* $*_v$ *x* = 0_v (*degree u*)
unfolding *berlekamp-resulting-mat-def Let-def*
using *gauss-jordan-single*(*1*)[*of transpose-mat ?QI degree u degree u - x*] *x-dim*
by *auto*
thus *?thesis* **by** (*auto simp add*: *mat-kernel-def x-dim*)
qed

private abbreviation *V* ≡ *kernel.VK* (*degree u*) (*berlekamp-resulting-mat u*)
private abbreviation *W* ≡ *vector-space-poly.vs*
{*v*. [*v*^(*CARD*('*a*)) = *v*] (*mod u*) ∧ (*degree v* < *degree u*)}

interpretation *V*: *vectorspace class-ring V*
proof −
interpret *k*: *kernel* (*degree u*) (*degree u*) (*berlekamp-resulting-mat u*)
by (*unfold-locales*; *auto*)
show *vectorspace class-ring V* **by** *intro-locales*
qed

lemma *linear-Poly-list-of-vec*:
shows (*Poly* ∘ *list-of-vec*) ∈ *module-hom class-ring V* (*vector-space-poly.vs* {*v*. [*v*^(*CARD*('*a*)) = *v*] (*mod u*)})
proof (*auto simp add*: *LinearCombinations.module-hom-def Matrix.module-vec-def*)
fix *m1 m2*:: '*a mod-ring vec*
assume *m1*: *m1* ∈ *mat-kernel* (*berlekamp-resulting-mat u*)
and *m2*: *m2* ∈ *mat-kernel* (*berlekamp-resulting-mat u*)
have *m1-rw*: *list-of-vec m1 = map* (λ*n*. *m1* $ *n*) [*0*..<*dim-vec m1*]
by (*transfer*, *auto simp add*: *mk-vec-def*)
have *m2-rw*: *list-of-vec m2 = map* (λ*n*. *m2* $ *n*) [*0*..<*dim-vec m2*]
by (*transfer*, *auto simp add*: *mk-vec-def*)
have *m1* ∈ *carrier-vec* (*degree u*) **by** (*rule mat-kernelD*(*1*)[*OF - m1*], *auto*)

moreover have *m2* ∈ *carrier-vec* (*degree u*) **by** (*rule mat-kernelD(1)*[*OF - m2*], *auto*)
ultimately have *dim-eq*: *dim-vec m1* = *dim-vec m2* **by** *auto*
show *Poly* (*list-of-vec* (*m1* + *m2*)) = *Poly* (*list-of-vec m1*) + *Poly* (*list-of-vec m2*)
unfolding *poly-eq-iff m1-rw m2-rw plus-vec-def*
using *dim-eq*
by (*transfer'*, *auto simp add*: *mk-vec-def nth-default-def*)
next
fix *r m* **assume** *m*: *m* ∈ *mat-kernel* (*berlekamp-resulting-mat u*)
show *Poly* (*list-of-vec* (*r* ·$_v$ *m*)) = *smult r* (*Poly* (*list-of-vec m*))
unfolding *poly-eq-iff list-of-vec-rw-map*[*of m*] *smult-vec-def*
by (*transfer'*, *auto simp add*: *mk-vec-def nth-default-def*)
next
fix *x* **assume** *x*: *x* ∈ *mat-kernel* (*berlekamp-resulting-mat u*)
show [*Poly* (*list-of-vec x*) ^ *CARD*('*a*) = *Poly* (*list-of-vec x*)] (*mod u*)
proof (*cases degree u* = *0*)
case *True*
have *mat-kernel* (*berlekamp-resulting-mat u*) = {*0*$_v$ *0*}
by (*rule mat-kernel-berlekamp-resulting-mat-degree-0*[*OF True*])
hence *x-0*: *x* = *0*$_v$ *0* **using** *x* **by** *blast*
show *?thesis* **by** (*auto simp add*: *zero-power x-0 cong-def*)
next
case *False* **note** *deg-u* = *False*
show *?thesis*
proof −
let *?QI*=(*mat* (*degree u*) (*degree u*)
(λ(*i*, *j*). *if i* = *j then berlekamp-mat u* $$ (*i*, *j*) − *1 else berlekamp-mat u* $$ (*i*, *j*)))
let *?H*=*vec-of-list* (*coeffs* (*Poly* (*list-of-vec x*)) @ *replicate* (*degree u* − *length* (*coeffs* (*Poly* (*list-of-vec x*)))) *0*)
have *x-dim*: *dim-vec x* = *degree u* **using** *x* **unfolding** *mat-kernel-def* **by** *auto*
hence *x-carrier*[*simp*]: *x* ∈ *carrier-vec* (*degree u*) **by** (*metis carrier-vec-dim-vec*)
have *x-kernel*: *berlekamp-resulting-mat u* ∗$_v$ *x* = *0*$_v$ (*degree u*)
using *x* **unfolding** *mat-kernel-def* **by** *auto*
have *t-QI-x-0*: (*transpose-mat ?QI*) ∗$_v$ *x* = *0*$_v$ (*degree u*)
using *gauss-jordan-single*(*1*)[*of* (*transpose-mat ?QI*) *degree u degree u gauss-jordan-single* (*transpose-mat ?QI*) *x*]
using *x-kernel* **unfolding** *berlekamp-resulting-mat-def Let-def* **by** *auto*
have *l*: (*list-of-vec x*) ≠ []
by (*auto simp add*: *list-of-vec-rw-map vec-of-dim-0*[*symmetric*] *deg-u x-dim*)
have *deg-le*: *degree* (*Poly* (*list-of-vec x*)) < *degree u*
using *degree-Poly-list-of-vec*
using *x-carrier deg-u* **by** *blast*
show [*Poly* (*list-of-vec x*) ^ *CARD*('*a*) = *Poly* (*list-of-vec x*)] (*mod u*)
proof (*unfold equation-13*[*OF deg-le*])
have *QR-rw*: *?QI* = *berlekamp-mat u* − *1*$_m$ (*dim-row* (*berlekamp-mat u*)) **by** *auto*
have *dim-row* (*berlekamp-mat u*) = *dim-vec ?H*

```
          by (auto, metis le-add-diff-inverse length-list-of-vec length-strip-while-le
x-dim)
        moreover have ?H ∈ mat-kernel (transpose-mat (berlekamp-mat u − 1_m
(dim-row (berlekamp-mat u))))
        proof −
          have Hx: ?H = x
          proof (unfold vec-eq-iff, auto)
            let ?H'=vec-of-list (strip-while ((=) 0) (list-of-vec x)
             @ replicate (degree u − length (strip-while ((=) 0) (list-of-vec x))) 0)
            show length (strip-while ((=) 0) (list-of-vec x))
             + (degree u − length (strip-while ((=) 0) (list-of-vec x))) = dim-vec x
                 by (metis le-add-diff-inverse length-list-of-vec length-strip-while-le
x-dim)
            fix i assume i: i < dim-vec x
            have ?H $ i =  coeff (Poly (list-of-vec x)) i
            proof (rule vec-of-list-coeffs-replicate-nth[OF - deg-le])
             show i ∈ {..<degree u}  using x-dim i by (auto, linarith)
            qed
            also have ... = x $ i by (rule coeff-Poly-list-of-vec-nth'[OF i])
            finally show ?H' $ i = x $ i by auto
          qed
           have ?H ∈ carrier-vec (degree u) using deg-le dim-vec-of-list-h Hx by
auto
         moreover have transpose-mat (berlekamp-mat u − 1_m (degree u)) *_v ?H
= 0_v (degree u)
           using t-QI-x-0 Hx QR-rw by auto
          ultimately show ?thesis
           by (auto simp add: mat-kernel-def)
        qed
        ultimately show transpose-mat (berlekamp-mat u) *_v ?H = ?H
          using system-if-mat-kernel[of berlekamp-mat u ?H]
          by auto
        qed
      qed
    qed
qed


lemma linear-Poly-list-of-vec':
  assumes degree u > 0
  shows (Poly ∘ list-of-vec) ∈ module-hom R V W
proof (auto simp add: LinearCombinations.module-hom-def Matrix.module-vec-def)
  fix m1 m2:: 'a mod-ring vec
  assume m1: m1 ∈ mat-kernel (berlekamp-resulting-mat u)
  and m2: m2 ∈ mat-kernel (berlekamp-resulting-mat u)
  have m1-rw: list-of-vec m1 = map (λn. m1 $ n) [0..<dim-vec m1]
    by (transfer, auto simp add: mk-vec-def)
  have m2-rw: list-of-vec m2 = map (λn. m2 $ n) [0..<dim-vec m2]
    by (transfer, auto simp add: mk-vec-def)
```

```
  have m1 ∈ carrier-vec (degree u) by (rule mat-kernelD(1)[OF - m1], auto)
  moreover have m2 ∈ carrier-vec (degree u) by (rule mat-kernelD(1)[OF - m2],
auto)
  ultimately have dim-eq: dim-vec m1 = dim-vec m2 by auto
  show Poly (list-of-vec (m1 + m2)) = Poly (list-of-vec m1) + Poly (list-of-vec
m2)
    unfolding poly-eq-iff m1-rw m2-rw plus-vec-def
    using dim-eq
    by (transfer', auto simp add: mk-vec-def nth-default-def)
next
  fix r m assume m: m ∈ mat-kernel (berlekamp-resulting-mat u)
  show Poly (list-of-vec (r ·v m)) = smult r (Poly (list-of-vec m))
    unfolding poly-eq-iff list-of-vec-rw-map[of m] smult-vec-def
    by (transfer', auto simp add: mk-vec-def nth-default-def)
next
  fix x assume x: x ∈ mat-kernel (berlekamp-resulting-mat u)
  show [Poly (list-of-vec x) ^ CARD('a) = Poly (list-of-vec x)] (mod u)
  proof (cases degree u = 0)
    case True
    have mat-kernel (berlekamp-resulting-mat u) = {0v 0}
      by (rule mat-kernel-berlekamp-resulting-mat-degree-0[OF True])
    hence x-0: x = 0v 0 using x by blast
    show ?thesis by (auto simp add: zero-power x-0 cong-def)
  next
    case False note deg-u = False
    show ?thesis
    proof −
      let ?QI=(mat (degree u) (degree u)
      (λ(i, j). if i = j then berlekamp-mat u $$ (i, j) − 1 else berlekamp-mat u $$
(i, j)))
      let ?H=vec-of-list (coeffs (Poly (list-of-vec x)) @ replicate (degree u − length
(coeffs (Poly (list-of-vec x)))) 0)
      have x-dim: dim-vec x = degree u using x unfolding mat-kernel-def by auto
      hence x-carrier[simp]: x ∈ carrier-vec (degree u) by (metis carrier-vec-dim-vec)
      have x-kernel: berlekamp-resulting-mat u ∗v x = 0v (degree u)
        using x unfolding mat-kernel-def by auto
      have t-QI-x-0: (transpose-mat ?QI) ∗v x = 0v (degree u)
          using gauss-jordan-single(1)[of (transpose-mat ?QI) degree u degree u
gauss-jordan-single (transpose-mat ?QI) x]
        using x-kernel unfolding berlekamp-resulting-mat-def Let-def by auto
      have l: (list-of-vec x) ≠ []
        by (auto simp add: list-of-vec-rw-map vec-of-dim-0[symmetric] deg-u x-dim)
      have deg-le: degree (Poly (list-of-vec x)) < degree u
        using degree-Poly-list-of-vec
        using x-carrier deg-u by blast
      show [Poly (list-of-vec x) ^ CARD('a) = Poly (list-of-vec x)] (mod u)
      proof (unfold equation-13[OF deg-le])
        have QR-rw: ?QI = berlekamp-mat u − 1m (dim-row (berlekamp-mat u))
by auto
```

```
      have dim-row (berlekamp-mat u) = dim-vec ?H
        by (auto, metis le-add-diff-inverse length-list-of-vec length-strip-while-le
x-dim)
      moreover have ?H ∈ mat-kernel (transpose-mat (berlekamp-mat u − 1_m
(dim-row (berlekamp-mat u))))
      proof −
        have Hx: ?H = x
        proof (unfold vec-eq-iff, auto)
          let ?H'=vec-of-list (strip-while ((=) 0) (list-of-vec x)
           @ replicate (degree u − length (strip-while ((=) 0) (list-of-vec x))) 0)
          show length (strip-while ((=) 0) (list-of-vec x))
           + (degree u − length (strip-while ((=) 0) (list-of-vec x))) = dim-vec x
              by (metis le-add-diff-inverse length-list-of-vec length-strip-while-le
x-dim)
          fix i assume i: i < dim-vec x
          have ?H $ i =  coeff (Poly (list-of-vec x)) i
          proof (rule vec-of-list-coeffs-replicate-nth[OF - deg-le])
           show i ∈ {..<degree u}  using x-dim i by (auto, linarith)
          qed
          also have ... = x $ i by (rule coeff-Poly-list-of-vec-nth'[OF i])
          finally show ?H' $ i = x $ i by auto
        qed
         have ?H ∈ carrier-vec (degree u) using deg-le dim-vec-of-list-h Hx by
auto
        moreover have transpose-mat (berlekamp-mat u − 1_m (degree u)) *_v ?H
= 0_v (degree u)
          using t-QI-x-0 Hx QR-rw by auto
         ultimately show ?thesis
          by (auto simp add: mat-kernel-def)
      qed
      ultimately show transpose-mat (berlekamp-mat u) *_v ?H = ?H
        using system-if-mat-kernel[of berlekamp-mat u ?H]
        by auto
      qed
    qed
  qed
next
  fix x assume x: x ∈ mat-kernel (berlekamp-resulting-mat u)
  show degree (Poly (list-of-vec x)) < degree u
    by (rule degree-Poly-list-of-vec, insert assms x, auto simp: mat-kernel-def)
qed


lemma berlekamp-basis-eq-8:
  assumes v: v ∈ set (berlekamp-basis u)
  shows [v ^ CARD('a) = v] (mod u)
proof −
  {
     fix x assume x: x ∈ set (find-base-vectors (berlekamp-resulting-mat u))
```

```
    have set (find-base-vectors (berlekamp-resulting-mat u)) ⊆ mat-kernel (berlekamp-resulting-mat
u)
     proof (rule find-base-vectors(1))
      show row-echelon-form (berlekamp-resulting-mat u)
       by (rule row-echelon-form-berlekamp-resulting-mat)
     show berlekamp-resulting-mat u ∈ carrier-mat (degree u) (degree u) by simp
     qed
     hence x ∈ mat-kernel (berlekamp-resulting-mat u) using x by auto
     hence [Poly (list-of-vec x) ^ CARD('a) = Poly (list-of-vec x)] (mod u)
      using linear-Poly-list-of-vec
       unfolding LinearCombinations.module-hom-def Matrix.module-vec-def by
auto
  }
  thus [v ^ CARD('a) = v] (mod u) using v unfolding set-berlekamp-basis-eq by
auto
qed


lemma surj-Poly-list-of-vec:
  assumes deg-u: degree u > 0
  shows (Poly ∘ list-of-vec)` (carrier V) = carrier W
proof (auto simp add: image-def)
  fix xa
  assume xa: xa ∈ mat-kernel (berlekamp-resulting-mat u)
  thus [Poly (list-of-vec xa) ^ CARD('a) = Poly (list-of-vec xa)] (mod u)
    using linear-Poly-list-of-vec
   unfolding LinearCombinations.module-hom-def Matrix.module-vec-def by auto
  show degree (Poly (list-of-vec xa)) < degree u
  proof (rule degree-Poly-list-of-vec[OF - deg-u])
    show xa ∈ carrier-vec (degree u) using xa unfolding mat-kernel-def by simp
  qed
next
  fix x assume x: [x ^ CARD('a) = x] (mod u)
  and deg-x: degree x < degree u
  show ∃ xa ∈ mat-kernel (berlekamp-resulting-mat u). x = Poly (list-of-vec xa)
   proof (rule bexI[of - vec-of-list (coeffs x @ replicate (degree u − length (coeffs
x)) 0)])
    let ?X = vec-of-list (coeffs x @ replicate (degree u − length (coeffs x)) 0)
    show x = Poly (list-of-vec (vec-of-list (coeffs x @ replicate (degree u − length
(coeffs x)) 0)))
      by auto
    have X: ?X ∈ carrier-vec (degree u) unfolding carrier-vec-def
      by (auto, metis Suc-leI coeffs-0-eq-Nil deg-x degree-0 le-add-diff-inverse
          length-coeffs-degree linordered-semidom-class.add-diff-inverse list.size(3)
order.asym)
    have t: transpose-mat (berlekamp-mat u) *v ?X = ?X
      using equation-13[OF deg-x] x by auto
    show vec-of-list (coeffs x @ replicate (degree u − length (coeffs x)) 0)
    ∈ mat-kernel (berlekamp-resulting-mat u) by (rule in-mat-kernel-berlekamp-resulting-mat[OF
```

t X])
qed
qed

lemma *card-set-berlekamp-basis*: *card* (*set* (*berlekamp-basis u*)) = *length* (*berlekamp-basis u*)
proof −
have *b*: *berlekamp-resulting-mat u* ∈ *carrier-mat* (*degree u*) (*degree u*) **by** *auto*
have (*set* (*berlekamp-basis u*)) = (*Poly* ∘ *list-of-vec*) ‘ *set* (*find-base-vectors* (*berlekamp-resulting-mat u*))
unfolding *set-berlekamp-basis-eq* **..**
also have *card* ... = *card* (*set* (*find-base-vectors* (*berlekamp-resulting-mat u*)))
proof (*rule card-image*, *rule subset-inj-on*[*OF inj-Poly-list-of-vec*])
show *set* (*find-base-vectors* (*berlekamp-resulting-mat u*)) ⊆ *carrier-vec* (*degree u*)
using *find-base-vectors*(*1*)[*OF row-echelon-form-berlekamp-resulting-mat b*]
unfolding *carrier-vec-def mat-kernel-def*
by *auto*
qed
also have ... = *length* (*find-base-vectors* (*berlekamp-resulting-mat u*))
by (*rule length-find-base-vectors*[*symmetric*, *OF row-echelon-form-berlekamp-resulting-mat b*])
finally show *?thesis* **unfolding** *berlekamp-basis-def* **by** *auto*
qed

context
assumes *deg-u0*[*simp*]: *degree u* > *0*
begin

interpretation *Berlekamp-subspace*: *vectorspace class-ring W*
by (*rule vector-space-poly.subspace-is-vs*[*OF subspace-Berlekamp*], *simp*)

lemma *linear-map-Poly-list-of-vec′*: *linear-map class-ring V W* (*Poly* ∘ *list-of-vec*)
proof (*auto simp add*: *linear-map-def*)
show *vectorspace class-ring V* **by** *intro-locales*
show *vectorspace class-ring W* **by** (*rule Berlekamp-subspace.vectorspace-axioms*)
show *mod-hom class-ring V W* (*Poly* ∘ *list-of-vec*)
proof (*rule mod-hom.intro*, *unfold mod-hom-axioms-def*)
show *module class-ring V* **by** *intro-locales*
show *module class-ring W* **using** *Berlekamp-subspace.vectorspace-axioms* **by** *intro-locales*
show *Poly* ∘ *list-of-vec* ∈ *module-hom class-ring V W*
by (*rule linear-Poly-list-of-vec′*[*OF deg-u0*])
qed
qed

lemma *berlekamp-basis-basis*:
Berlekamp-subspace.basis (*set* (*berlekamp-basis u*))

proof (*unfold set-berlekamp-basis-eq*, *rule linear-map.linear-inj-image-is-basis*)
show *linear-map class-ring V W* (*Poly* ∘ *list-of-vec*)
by (*rule linear-map-Poly-list-of-vec'*)
show *inj-on* (*Poly* ∘ *list-of-vec*) (*carrier V*)
proof (*rule subset-inj-on*[*OF inj-Poly-list-of-vec*])
show *carrier V* ⊆ *carrier-vec* (*degree u*)
by (*auto simp add*: *mat-kernel-def*)
qed
show (*Poly* ∘ *list-of-vec*) ' *carrier V* = *carrier W*
using *surj-Poly-list-of-vec*[*OF deg-u0*] **by** *auto*
show *b*: *V.basis* (*set* (*find-base-vectors* (*berlekamp-resulting-mat u*)))
by (*rule berlekamp-resulting-mat-basis*)
show *V.fin-dim*
proof −
have *finite* (*set* (*find-base-vectors* (*berlekamp-resulting-mat u*))) **by** *auto*
moreover have *set* (*find-base-vectors* (*berlekamp-resulting-mat u*)) ⊆ *carrier V*
and *V.gen-set* (*set* (*find-base-vectors* (*berlekamp-resulting-mat u*)))
using *b* **unfolding** *V.basis-def* **by** *auto*
ultimately show *?thesis* **unfolding** *V.fin-dim-def* **by** *auto*
qed
qed

lemma *finsum-sum*:
fixes *f*::'*a mod-ring poly*
assumes *f*: *finite B*
and *a-Pi*: *a* ∈ *B* → *carrier R*
and *V*: *B* ⊆ *carrier W*
shows (⨁$_W$*v*∈*B*. *a v* ⊙$_W$ *v*) = *sum* (λ*v*. *smult* (*a v*) *v*) *B*
using *f a-Pi V*
proof (*induct B*)
case *empty*
thus *?case* **unfolding** *Berlekamp-subspace.module.M.finsum-empty* **by** *auto*
next
case (*insert x V*)
have *hyp*: (⨁$_W$*v* ∈ *V*. *a v* ⊙$_W$ *v*) = *sum* (λ*v*. *smult* (*a v*) *v*) *V*
proof (*rule insert.hyps*)
show *a* ∈ *V* → *carrier R*
using *insert.prems* **unfolding** *class-field-def* **by** *auto*
show *V* ⊆ *carrier W* **using** *insert.prems* **by** *simp*
qed
have (⨁$_W$*v*∈*insert x V*. *a v* ⊙$_W$ *v*) = (*a x* ⊙$_W$ *x*) ⊕$_W$ (⨁$_W$*v* ∈ *V*. *a v* ⊙$_W$ *v*)
proof (*rule abelian-monoid.finsum-insert*)
show *abelian-monoid W* **by** (*unfold-locales*)
show *finite V* **by** *fact*
show *x* ∉ *V* **by** *fact*
show (λ*v*. *a v* ⊙$_W$ *v*) ∈ *V* → *carrier W*

proof (*unfold Pi-def*, *rule*, *rule allI*, *rule impI*)
fix *v* **assume** *v*: *v*∈*V*
show *a v* ⊙$_W$ *v* ∈ *carrier W*
proof (*rule Berlekamp-subspace.smult-closed*)
show *a v* ∈ *carrier class-ring* **using** *insert.prems v* **unfolding** *Pi-def*
by (*simp add*: *class-field-def*)
show *v* ∈ *carrier W* **using** *v insert.prems* **by** *auto*
qed
qed
show *a x* ⊙$_W$ *x* ∈ *carrier W*
proof (*rule Berlekamp-subspace.smult-closed*)
show *a x* ∈ *carrier class-ring* **using** *insert.prems* **unfolding** *Pi-def*
by (*simp add*: *class-field-def*)
show *x* ∈ *carrier W* **using** *insert.prems* **by** *auto*
qed
qed
also have ... = (*a x* ⊙$_W$ *x*) + (⨁$_W$*v* ∈ *V*. *a v* ⊙$_W$ *v*) **by** *auto*
also have ... = (*a x* ⊙$_W$ *x*) + *sum* (λ*v*. *smult* (*a v*) *v*) *V* **unfolding** *hyp* **by** *simp*
also have ... = (*smult* (*a x*) *x*) + *sum* (λ*v*. *smult* (*a v*) *v*) *V* **by** *simp*
also have ... = *sum* (λ*v*. *smult* (*a v*) *v*) (*insert x V*)
by (*simp add*: *insert.hyps*(*1*) *insert.hyps*(*2*))
finally show *?case* **.**
qed

lemma *exists-vector-in-Berlekamp-subspace-dvd*:
fixes *p-i*::′*a mod-ring poly*
assumes *finite-P*: *finite P*
and *f-desc-square-free*: *u* = (∏ *a*∈*P*. *a*)
and *P*: *P* ⊆ {*q*. *irreducible q* ∧ *monic q*}
and *pi*: *p-i* ∈ *P* **and** *pj*: *p-j* ∈ *P* **and** *pi-pj*: *p-i* ≠ *p-j*
and *monic-f*: *monic u* **and** *sf-f*: *square-free u*
and *not-irr-w*: ¬ *irreducible w*
and *w-dvd-f*: *w dvd u* **and** *monic-w*: *monic w*
and *pi-dvd-w*: *p-i dvd w* **and** *pj-dvd-w*: *p-j dvd w*
shows ∃ *v*. *v* ∈ {*h*. [*h*^(*CARD*(′*a*)) = *h*] (*mod u*) ∧ *degree h* < *degree u*}
∧ *v mod p-i* ≠ *v mod p-j*
∧ *degree* (*v mod p-i*) = *0*
∧ *degree* (*v mod p-j*) = *0*
— This implies that the algorithm decreases the degree of the reducible polynomials in each step:
∧ (∃ *s*. *gcd w* (*v* − [:*s*:]) ≠ *w* ∧ *gcd w* (*v* − [:*s*:]) ≠ *1*)
proof −
have *f-not-0*: *u* ≠ *0* **using** *monic-f* **by** *auto*
have *irr-pi*: *irreducible p-i* **using** *pi P* **by** *auto*
have *irr-pj*: *irreducible p-j* **using** *pj P* **by** *auto*
obtain *m* **and** *n*::*nat* **where** *P-m*: *P* = *m* ‘ {*i*. *i* < *n*} **and** *inj-on-m*: *inj-on m* {*i*. *i* < *n*}

```
    using finite-imp-nat-seg-image-inj-on[OF finite-P] by blast
  hence n = card P by (simp add: card-image)
  have degree-prod: degree (prod m {i. i < n}) = degree u
    by (metis P-m f-desc-square-free inj-on-m prod.reindex-cong)
  have not-zero: ∀ i∈{i. i < n}. m i ≠ 0
    using P-m f-desc-square-free f-not-0 by auto
  obtain i where mi: m i = p-i and i: i < n using P-m pi by blast
  obtain j where mj: m j = p-j and j: j < n using P-m pj by blast
  have ij: i ≠ j  using  mi mj pi-pj by auto
  obtain s-i and s-j::'a mod-ring where si-sj: s-i ≠ s-j using exists-two-distint
by blast
  let ?u=λx. if x = i then [:s-i:] else if x = j then [:s-j:] else [:0:]
  have degree-si: degree [:s-i:] = 0 by auto
  have degree-sj: degree [:s-j:] = 0 by auto
  have ∃!v. degree v < (∑ i∈{i. i < n}. degree (m i)) ∧ (∀ a∈{i. i < n}. [v = ?u
a] (mod m a))
  proof (rule chinese-remainder-unique-poly)
    show ∀ a∈{i. i < n}. ∀ b∈{i. i < n}. a ≠ b ⟶ Rings.coprime (m a) (m b)
    proof (rule+)
      fix a b assume a ∈ {i. i < n} and b ∈ {i. i < n} and a ≠ b
      thus Rings.coprime (m a) (m b)
        using coprime-polynomial-factorization[OF P finite-P, simplified] P-m
        by (metis image-eqI inj-onD inj-on-m)
    qed
    show ∀ i∈{i. i < n}. m i ≠ 0 by (rule not-zero)
   show 0 < degree (prod m {i. i < n}) unfolding degree-prod using deg-u0 by
blast
  qed
  from this obtain v where v: ∀ a∈{i. i < n}. [v = ?u a] (mod m a)
  and degree-v: degree v < (∑ i∈{i. i < n}. degree (m i)) by blast
  show ?thesis
  proof (rule exI[of - v], auto)
    show vp-v-mod: [v ^ CARD('a) = v] (mod u)
    proof (unfold f-desc-square-free, rule coprime-cong-mult-factorization-poly[OF
finite-P])
      show P ⊆ {q. irreducible q} using P by blast
      show ∀ p∈P. [v ^ CARD('a) = v] (mod p)
      proof (rule ballI)
        fix p assume p: p ∈ P
        hence irr-p: irreducible_d p using P by auto
        obtain k where mk: m k = p and k: k < n using P-m p by blast
        have [v = ?u k] (mod p) using v mk k by auto
        moreover have ?u k mod p = ?u k
          apply (rule mod-poly-less) using irreducible_dD(1)[OF irr-p] by auto
        ultimately obtain s where v-mod-p: v mod p = [:s:] unfolding cong-def
by force
        hence deg-v-p: degree (v mod p) = 0 by auto
        have v mod p = [:s:] by (rule v-mod-p)
        also have ... = [:s:]^CARD('a) unfolding poly-const-pow by auto
```

also have ... = $(v \ mod \ p) \ \hat{} \ CARD('a)$ **using** *v-mod-p* **by** *auto*
also have ... = $(v \ mod \ p) \ \hat{} \ CARD('a) \ mod \ p$ **using** *calculation* **by** *auto*
also have ... = $v \hat{} CARD('a) \ mod \ p$ **using** *power-mod* **by** *blast*
finally show $[v \ \hat{} \ CARD('a) = v] \ (mod \ p)$ **unfolding** *cong-def* **..**
qed
show $\forall \ p1 \ p2. \ p1 \in P \wedge p2 \in P \wedge p1 \neq p2 \longrightarrow coprime \ p1 \ p2$
using *P coprime-polynomial-factorization finite-P* **by** *auto*
qed
have $[v = ?u \ i] \ (mod \ m \ i)$ **using** *v i* **by** *auto*
hence *v-pi-si-mod*: $v \ mod \ p\text{-}i = [:s\text{-}i:] \ mod \ p\text{-}i$ **unfolding** *cong-def mi* **by** *auto*
also have ... = $[:s\text{-}i:]$ **apply** (*rule mod-poly-less*) **using** *irr-pi* **by** *auto*
finally have *v-pi-si*: $v \ mod \ p\text{-}i = [:s\text{-}i:]$ **.**

have $[v = ?u \ j] \ (mod \ m \ j)$ **using** *v j* **by** *auto*
hence *v-pj-sj-mod*: $v \ mod \ p\text{-}j = [:s\text{-}j:] \ mod \ p\text{-}j$ **unfolding** *cong-def mj* **using** *ij* **by** *auto*
also have ... = $[:s\text{-}j:]$ **apply** (*rule mod-poly-less*) **using** *irr-pj* **by** *auto*
finally have *v-pj-sj*: $v \ mod \ p\text{-}j = [:s\text{-}j:]$ **.**
show $v \ mod \ p\text{-}i = v \ mod \ p\text{-}j \Longrightarrow False$ **using** *si-sj v-pi-si v-pj-sj* **by** *auto*
show $degree \ (v \ mod \ p\text{-}i) = 0$ **unfolding** *v-pi-si* **by** *simp*
show $degree \ (v \ mod \ p\text{-}j) = 0$ **unfolding** *v-pj-sj* **by** *simp*
show $\exists \ s. \ gcd \ w \ (v - [:s:]) \neq w \wedge gcd \ w \ (v - [:s:]) \neq 1$
proof (*rule exI*[*of - s-i*], *rule conjI*)
have *pi-dvd-v-si*: $p\text{-}i \ dvd \ v - [:s\text{-}i:]$ **using** *v-pi-si-mod mod-eq-dvd-iff-poly* **by** *blast*
have *pj-dvd-v-sj*: $p\text{-}j \ dvd \ v - [:s\text{-}j:]$ **using** *v-pj-sj-mod mod-eq-dvd-iff-poly* **by** *blast*
have *w-eq*: $w = prod \ (\lambda c. \ gcd \ w \ (v - [:c:])) \ (UNIV :: 'a \ mod\text{-}ring \ set)$
proof (*rule Berlekamp-gcd-step*)
show $[v \ \hat{} \ CARD('a) = v] \ (mod \ w)$ **using** *vp-v-mod cong-dvd-modulus-poly w-dvd-f* **by** *blast*
show *square-free w* **by** (*rule square-free-factor*[*OF w-dvd-f sf-f*])
show *monic w* **by** (*rule monic-w*)
qed
show $gcd \ w \ (v - [:s\text{-}i:]) \neq w$
proof (*rule ccontr*, *simp*)
assume *gcd-w*: $gcd \ w \ (v - [:s\text{-}i:]) = w$
show *False* **apply** (*rule* ‹$v \ mod \ p\text{-}i = v \ mod \ p\text{-}j \Longrightarrow False$›)
by (*metis irreducibleE* ‹$degree \ (v \ mod \ p\text{-}i) = 0$› *gcd-greatest-iff gcd-w irr-pj is-unit-field-poly mod-eq-dvd-iff-poly mod-poly-less neq0-conv pj-dvd-w v-pi-si*)
qed
show $gcd \ w \ (v - [:s\text{-}i:]) \neq 1$
by (*metis irreducibleE gcd-greatest-iff irr-pi pi-dvd-v-si pi-dvd-w*)
qed
show $degree \ v < degree \ u$
proof −
have $(\sum i \mid i < n. \ degree \ (m \ i)) = degree \ (prod \ m \ \{i. \ i < n\})$
by (*rule degree-prod-eq-sum-degree*[*symmetric*, *OF not-zero*])
thus *?thesis* **using** *degree-v* **unfolding** *degree-prod* **by** *auto*

qed
qed
qed

lemma *exists-vector-in-Berlekamp-basis-dvd-aux*:
assumes *basis-V*: *Berlekamp-subspace.basis B*
and *finite-V*: *finite B*
assumes *finite-P*: *finite P*
and *f-desc-square-free*: $u = (\prod a{\in}P.\ a)$
and *P*: $P \subseteq \{q.\ \textit{irreducible}\ q \wedge \textit{monic}\ q\}$
and *pi*: *p-i* ∈ *P* **and** *pj*: *p-j* ∈ *P* **and** *pi-pj*: *p-i* ≠ *p-j*
and *monic-f*: *monic u* **and** *sf-f*: *square-free u*
and *not-irr-w*: ¬ *irreducible w*
and *w-dvd-f*: *w dvd u* **and** *monic-w*: *monic w*
and *pi-dvd-w*: *p-i dvd w* **and** *pj-dvd-w*: *p-j dvd w*
shows ∃ *v* ∈ *B*. *v mod p-i* ≠ *v mod p-j*
proof (*rule ccontr*, *auto*)
have *V-in-carrier*: *B* ⊆ *carrier W*
using *basis-V* **unfolding** *Berlekamp-subspace.basis-def* **by** *auto*
assume *all-eq*: ∀ *v*∈*B*. *v mod p-i* = *v mod p-j*
obtain *x* **where** *x*: *x* ∈ {*h*. [*h* ^ *CARD*('*a*) = *h*] (*mod u*) ∧ *degree h* < *degree u*}
and *x-pi-pj*: *x mod p-i* ≠ *x mod p-j* **and** *degree* (*x mod p-i*) = *0* **and** *degree* (*x mod p-j*) = *0*
(∃ *s*. *gcd w* (*x* − [:*s*:]) ≠ *w* ∧ *gcd w* (*x* − [:*s*:]) ≠ *1*)
using *exists-vector-in-Berlekamp-subspace-dvd*[*OF - - - pi pj - - - - w-dvd-f monic-w pi-dvd-w*]
assms **by** *meson*
have *x-in*: *x* ∈ *carrier W* **using** *x* **by** *auto*
hence (∃!*a*. *a* ∈ *B* →$_E$ *carrier class-ring* ∧ *Berlekamp-subspace.lincomb a B* = *x*)
using *Berlekamp-subspace.basis-criterion*[*OF finite-V V-in-carrier*] **using** *basis-V*
by (*simp add*: *class-field-def*)
from *this* **obtain** *a* **where** *a-Pi*: *a* ∈ *B* →$_E$ *carrier class-ring*
and *lincomb-x*: *Berlekamp-subspace.lincomb a B* = *x*
by *blast*
have *fs-ss*: $(\bigoplus_W v{\in}B.\ a\ v \odot_W v)$ = *sum* (λ*v*. *smult* (*a v*) *v*) *B*
proof (*rule finsum-sum*)
show *finite B* **by** *fact*
show *a* ∈ *B* → *carrier class-ring* **using** *a-Pi* **by** *auto*
show *B* ⊆ *carrier W* **by** (*rule V-in-carrier*)
qed
have *x mod p-i* = *Berlekamp-subspace.lincomb a B mod p-i* **using** *lincomb-x* **by** *simp*
also have ... = $(\bigoplus_W v{\in}B.\ a\ v \odot_W v)$ *mod p-i* **unfolding** *Berlekamp-subspace.lincomb-def* **..**
also have ... = (*sum* (λ*v*. *smult* (*a v*) *v*) *B*) *mod p-i* **unfolding** *fs-ss* **..**

also have ... = *sum* (λ*v*. *smult* (*a v*) *v mod p-i*) *B* **using** *finite-V poly-mod-sum*
by *blast*
also have ... = *sum* (λ*v*. *smult* (*a v*) (*v mod p-i*)) *B* **by** (*meson mod-smult-left*)
also have ... = *sum* (λ*v*. *smult* (*a v*) (*v mod p-j*)) *B* **using** *all-eq* **by** *auto*
also have ... = *sum* (λ*v*. *smult* (*a v*) *v mod p-j*) *B* **by** (*metis mod-smult-left*)
also have ... = (*sum* (λ*v*. *smult* (*a v*) *v*) *B*) *mod p-j*
by (*metis* (*mono-tags*, *lifting*) *finite-V poly-mod-sum sum.cong*)
also have ... = ($\bigoplus_W v \in B$. *a v* $\odot_W$ *v*) *mod p-j* **unfolding** *fs-ss* **..**
also have ... = *Berlekamp-subspace.lincomb a B mod p-j*
unfolding *Berlekamp-subspace.lincomb-def* **..**
also have ... = *x mod p-j* **using** *lincomb-x* **by** *simp*
finally have *x mod p-i* = *x mod p-j* **.**
thus *False* **using** *x-pi-pj* **by** *contradiction*
qed

lemma *exists-vector-in-Berlekamp-basis-dvd*:
assumes *basis-V*: *Berlekamp-subspace.basis B*
and *finite-V*: *finite B*
assumes *finite-P*: *finite P*
and *f-desc-square-free*: *u* = (∏ *a*∈*P*. *a*)
and *P*: *P* ⊆ {*q*. *irreducible q* ∧ *monic q*}
and *pi*: *p-i* ∈ *P* **and** *pj*: *p-j* ∈ *P* **and** *pi-pj*: *p-i* ≠ *p-j*
and *monic-f*: *monic u* **and** *sf-f*: *square-free u*
and *not-irr-w*: ¬ *irreducible w*
and *w-dvd-f*: *w dvd u* **and** *monic-w*: *monic w*
and *pi-dvd-w*: *p-i dvd w* **and** *pj-dvd-w*: *p-j dvd w*
shows ∃ *v* ∈ *B*. *v mod p-i* ≠ *v mod p-j*
∧ *degree* (*v mod p-i*) = *0*
∧ *degree* (*v mod p-j*) = *0*
— This implies that the algorithm decreases the degree of the reducible polynomials in each step:
∧ (∃ *s*. *gcd w* (*v* − [:*s*:]) ≠ *w* ∧ ¬ *coprime w* (*v* − [:*s*:]))
proof −
have *f-not-0*: *u* ≠ *0* **using** *monic-f* **by** *auto*
have *irr-pi*: *irreducible p-i* **using** *pi P* **by** *fast*
have *irr-pj*: *irreducible p-j* **using** *pj P* **by** *fast*
obtain *v* **where** *vV*: *v* ∈ *B* **and** *v-pi-pj*: *v mod p-i* ≠ *v mod p-j*
using *assms exists-vector-in-Berlekamp-basis-dvd-aux* **by** *blast*
have *v*: *v* ∈ {*v*. [*v* ^ *CARD*('*a*) = *v*] (*mod u*)}
using *basis-V vV* **unfolding** *Berlekamp-subspace.basis-def* **by** *auto*
have *deg-v-pi*: *degree* (*v mod p-i*) = *0*
by (*rule degree-u-mod-irreducible*$_d$*-factor-0*[*OF v finite-P f-desc-square-free P pi*])
from *this* **obtain** *s-i* **where** *v-pi-si*: *v mod p-i* = [:*s-i*:] **using** *degree-eq-zeroE*
by *blast*
have *deg-v-pj*: *degree* (*v mod p-j*) = *0*
by (*rule degree-u-mod-irreducible*$_d$*-factor-0*[*OF v finite-P f-desc-square-free P pj*])

from *this* **obtain** *s-j* **where** *v-pj-sj*: *v mod p-j* = [:*s-j*:] **using** *degree-eq-zeroE*
by *blast*
have *si-sj*: *s-i* ≠ *s-j* **using** *v-pi-si v-pj-sj v-pi-pj* **by** *auto*
have (∃ *s*. *gcd w* (*v* − [:*s*:]) ≠ *w* ∧ ¬ *Rings.coprime w* (*v* − [:*s*:]))
proof (*rule exI*[*of - s-i*], *rule conjI*)
have *pi-dvd-v-si*: *p-i dvd v* − [:*s-i*:] **by** (*metis mod-eq-dvd-iff-poly mod-mod-trivial*
v-pi-si)
have *pj-dvd-v-sj*: *p-j dvd v* − [:*s-j*:] **by** (*metis mod-eq-dvd-iff-poly mod-mod-trivial*
v-pj-sj)
have *w-eq*: *w* = *prod* (λ*c*. *gcd w* (*v* − [:*c*:])) (*UNIV*::′*a mod-ring set*)
proof (*rule Berlekamp-gcd-step*)
show [*v* ^ *CARD*(′*a*) = *v*] (*mod w*) **using** *v cong-dvd-modulus-poly w-dvd-f*
by *blast*
show *square-free w* **by** (*rule square-free-factor*[*OF w-dvd-f sf-f*])
show *monic w* **by** (*rule monic-w*)
qed
show *gcd w* (*v* − [:*s-i*:]) ≠ *w*
by (*metis irreducibleE deg-v-pi gcd-greatest-iff irr-pj is-unit-field-poly mod-eq-dvd-iff-poly*
mod-poly-less neq0-conv pj-dvd-w v-pi-pj v-pi-si)
show ¬ *Rings.coprime w* (*v* − [:*s-i*:])
using *irr-pi pi-dvd-v-si pi-dvd-w*
by (*simp add*: *irreducible*$_d$*D*(*1*) *not-coprimeI*)
qed
thus *?thesis* **using** *v-pi-pj vV deg-v-pi deg-v-pj* **by** *auto*
qed

lemma *exists-bijective-linear-map-W-vec*:
assumes *finite-P*: *finite P*
and *u-desc-square-free*: *u* = (∏ *a*∈*P*. *a*)
and *P*: *P* ⊆ {*q*. *irreducible q* ∧ *monic q*}
shows ∃ *f*. *linear-map class-ring W* (*module-vec TYPE*(′*a mod-ring*) (*card P*)) *f*
∧ *bij-betw f* (*carrier W*) (*carrier-vec* (*card P*)::′*a mod-ring vec set*)
proof −
let *?B*=*carrier-vec* (*card P*)::′*a mod-ring vec set*
have *u-not-0*: *u* ≠ *0* **using** *deg-u0 degree-0* **by** *force*
obtain *m* **and** *n*::*nat* **where** *P-m*: *P* = *m* ‘ {*i*. *i* < *n*} **and** *inj-on-m*: *inj-on m*
{*i*. *i* < *n*}
using *finite-imp-nat-seg-image-inj-on*[*OF finite-P*] **by** *blast*
hence *n*: *n* = *card P* **by** (*simp add*: *card-image*)
have *degree-prod*: *degree* (*prod m* {*i*. *i* < *n*}) = *degree u*
by (*metis P-m u-desc-square-free inj-on-m prod.reindex-cong*)
have *not-zero*: ∀ *i*∈{*i*. *i* < *n*}. *m i* ≠ *0*
using *P-m u-desc-square-free u-not-0* **by** *auto*
have *deg-sum-eq*: (∑ *i*∈{*i*. *i* < *n*}. *degree* (*m i*)) = *degree u*
by (*metis degree-prod degree-prod-eq-sum-degree not-zero*)
have *coprime-mi-mj*:∀ *i*∈{*i*. *i* < *n*}. ∀ *j*∈{*i*. *i* < *n*}. *i* ≠ *j* ⟶ *coprime* (*m i*) (*m*
j)
proof (*rule*+)
fix *i j* **assume** *i*: *i* ∈ {*i*. *i* < *n*}

```
    and j: j ∈ {i. i < n} and ij: i ≠ j
    show coprime (m i) (m j)
    proof (rule coprime-polynomial-factorization[OF P finite-P])
      show m i ∈ P using i P-m by auto
      show m j ∈ P using j P-m by auto
      show m i ≠ m j using inj-on-m i ij j unfolding inj-on-def by blast
    qed
  qed
  let ?f = λv. vec n (λi. coeff (v mod (m i)) 0)
  interpret vec-VS: vectorspace class-ring (module-vec TYPE('a mod-ring) n)
    by (rule VS-Connect.vec-vs)
  interpret linear-map class-ring W (module-vec TYPE('a mod-ring) n) ?f
   by (intro-locales, unfold mod-hom-axioms-def LinearCombinations.module-hom-def,
       auto simp add: vec-eq-iff module-vec-def mod-smult-left poly-mod-add-left)
  have linear-map class-ring W (module-vec TYPE('a mod-ring) n) ?f
    by (intro-locales)
  moreover have inj-f: inj-on ?f (carrier W)
  proof (rule Ke0-imp-inj, auto simp add: mod-hom.ker-def)
    show [0 ^ CARD('a) = 0] (mod u) by (simp add: cong-def zero-power)
     show vec n (λi. 0) = 0_{module-vec TYPE('a mod-ring) n} by (auto simp add:
module-vec-def)
    fix x assume x: [x ^ CARD('a) = x] (mod u) and deg-x: degree x < degree u
    and v: vec n (λi. coeff (x mod m i) 0) = 0_{module-vec TYPE('a mod-ring) n}
    have cong-0: ∀ i∈{i. i < n}. [x = (λi. 0) i] (mod m i)
    proof (rule, unfold cong-def)
      fix i assume i: i ∈ {i. i < n}
      have deg-x-mod-mi: degree (x mod m i) = 0
     proof (rule degree-u-mod-irreducible_d-factor-0[OF - finite-P u-desc-square-free
P])
        show x ∈ {v. [v ^ CARD('a) = v] (mod u)} using x by auto
        show m i ∈ P using P-m i by auto
      qed
      thus x mod m i = 0 mod m i
       using v
       unfolding module-vec-def
     by (auto, metis i leading-coeff-neq-0 mem-Collect-eq index-vec index-zero-vec(1))
    qed
    moreover have deg-x2: degree x < (∑ i∈{i. i < n}. degree (m i))
     using deg-sum-eq deg-x by simp
    moreover have ∀ i∈{i. i < n}. [0 = (λi. 0) i] (mod m i)
     by (auto simp add: cong-def)
    moreover have degree 0 < (∑ i∈{i. i < n}. degree (m i))
     using degree-prod deg-sum-eq deg-u0 by force
    moreover have ∃!x. degree x < (∑ i∈{i. i < n}. degree (m i))
      ∧ (∀ i∈{i. i < n}. [x = (λi. 0) i] (mod m i))
    proof (rule chinese-remainder-unique-poly[OF not-zero])
     show 0 < degree (prod m {i. i < n})
      using deg-u0 degree-prod by linarith
    qed (insert coprime-mi-mj, auto)
```

```
    ultimately show x = 0 by blast
  qed
  moreover have ?f ‘ (carrier W) = ?B
  proof (auto simp add: image-def)
    fix xa
    show n = card P by (auto simp add: n)
    next
    fix x::'a mod-ring vec assume x: x ∈ carrier-vec (card P)
    have  ∃!v. degree v < (∑ i∈{i. i < n}. degree (m i)) ∧ (∀ i∈{i. i < n}. [v =
(λi. [:x $ i:]) i] (mod m i))
    proof (rule chinese-remainder-unique-poly[OF not-zero])
      show 0 < degree (prod m {i. i < n})
        using deg-u0 degree-prod by linarith
    qed (insert coprime-mi-mj, auto)
    from this obtain v where deg-v: degree v < (∑ i∈{i. i < n}. degree (m i))
      and v-x-cong: (∀ i ∈ {i. i < n}. [v = (λi. [:x $ i:]) i] (mod m i)) by auto
    show ∃ xa. [xa ^ CARD('a) = xa] (mod u) ∧ degree xa < degree u
      ∧ x = vec n (λi. coeff (xa mod m i) 0)
    proof (rule exI[of - v], auto)
      show v: [v ^ CARD('a) = v] (mod u)
     proof (unfold u-desc-square-free, rule coprime-cong-mult-factorization-poly[OF
finite-P], auto)
        fix y assume y: y ∈ P thus irreducible y using P by blast
        obtain i where i: i ∈ {i. i < n} and mi: y = m i using P-m y by blast
        have irreducible (m i) using i P-m P by auto
        moreover have [v = [:x $ i:]] (mod m i) using v-x-cong i by auto
        ultimately have v-mi-eq-xi: v mod m i = [:x $ i:]
          by (auto simp: cong-def intro!: mod-poly-less)
       have xi-pow-xi: [:x $ i:]^CARD('a) = [:x $ i:] by (simp add: poly-const-pow)
        hence (v mod m i)^CARD('a) = v mod m i using v-mi-eq-xi by auto
        hence (v mod m i)^CARD('a) = (v^CARD('a) mod m i)
          by (metis mod-mod-trivial power-mod)
      thus [v ^ CARD('a) = v] (mod y) unfolding mi cong-def v-mi-eq-xi xi-pow-xi
by simp
      next
        fix p1 p2 assume p1 ∈ P and p2 ∈ P and p1 ≠ p2
        then show Rings.coprime p1 p2
          using coprime-polynomial-factorization[OF P finite-P] by auto
      qed
      show degree v < degree u using deg-v deg-sum-eq degree-prod by presburger
      show x = vec n (λi. coeff (v mod m i) 0)
      proof (unfold vec-eq-iff, rule conjI)
        show dim-vec x = dim-vec (vec n (λi. coeff (v mod m i) 0)) using x n by
simp
          show ∀ i<dim-vec (vec n (λi. coeff (v mod m i) 0)). x $ i = vec n (λi.
coeff (v mod m i) 0) $ i
        proof (auto)
          fix i assume i: i < n
          have deg-mi: irreducible (m i) using i P-m P by auto
```

```
      have deg-v-mi: degree (v mod m i) = 0
    proof (rule degree-u-mod-irreducible_d-factor-0[OF - finite-P u-desc-square-free
P])
        show v ∈ {v. [v ^ CARD('a) = v] (mod u)} using v by fast
        show m i ∈ P using P-m i by auto
      qed
      have v mod m i = [:x $ i:] mod m i using v-x-cong i unfolding cong-def
by auto
      also have ... = [:x $ i:] using deg-mi by (auto intro!: mod-poly-less)
      finally show x $ i = coeff (v mod m i) 0 by simp
    qed
  qed
 qed
 ultimately show ?thesis unfolding bij-betw-def n by auto
qed

lemma fin-dim-kernel-berlekamp: V.fin-dim
proof −
 have finite (set (find-base-vectors (berlekamp-resulting-mat u))) by auto
 moreover have set (find-base-vectors (berlekamp-resulting-mat u)) ⊆ carrier V
 and V.gen-set (set (find-base-vectors (berlekamp-resulting-mat u)))
  using berlekamp-resulting-mat-basis[of u] unfolding V.basis-def by auto
 ultimately show ?thesis unfolding V.fin-dim-def by auto
qed

lemma Berlekamp-subspace-fin-dim: Berlekamp-subspace.fin-dim
proof (rule linear-map.surj-fin-dim[OF linear-map-Poly-list-of-vec'])
 show (Poly ∘ list-of-vec) ' carrier V = carrier W
  using surj-Poly-list-of-vec[OF deg-u0] by auto
 show V.fin-dim by (rule fin-dim-kernel-berlekamp)
qed

context
 fixes P
 assumes finite-P: finite P
 and u-desc-square-free: u = (∏ a∈P. a)
 and P: P ⊆ {q. irreducible q ∧ monic q}
begin

interpretation RV: vec-space TYPE('a mod-ring) card P .

lemma Berlekamp-subspace-eq-dim-vec: Berlekamp-subspace.dim = RV.dim
proof −
 obtain f where lm-f: linear-map class-ring W (module-vec TYPE('a mod-ring)
(card P)) f
 and bij-f: bij-betw f (carrier W) (carrier-vec (card P)::'a mod-ring vec set)
   using exists-bijective-linear-map-W-vec[OF finite-P u-desc-square-free P] by
blast
```

```
  show ?thesis
  proof (rule linear-map.dim-eq[OF lm-f Berlekamp-subspace-fin-dim])
    show inj-on f (carrier W) by (rule bij-betw-imp-inj-on[OF bij-f])
     show  f ‘ carrier W = carrier RV.V using bij-f unfolding bij-betw-def by
auto
  qed
qed


lemma Berlekamp-subspace-dim: Berlekamp-subspace.dim = card P
  using Berlekamp-subspace-eq-dim-vec RV.dim-is-n by simp

corollary card-berlekamp-basis-number-factors: card (set (berlekamp-basis u)) =
card P
  unfolding Berlekamp-subspace-dim[symmetric]
 by (rule Berlekamp-subspace.dim-basis[symmetric], auto simp add: berlekamp-basis-basis)


lemma length-berlekamp-basis-numbers-factors: length (berlekamp-basis u) = card
P
  using card-set-berlekamp-basis card-berlekamp-basis-number-factors by auto


end
end
end
end

context
  assumes SORT-CONSTRAINT('a :: prime-card)
begin

context
  fixes f :: 'a mod-ring poly and n
  assumes sf: square-free f
  and n: n = length (berlekamp-basis f)
  and monic-f: monic f
begin
lemma berlekamp-basis-length-factorization: assumes f: f = prod-list us
  and d: ⋀ u. u ∈ set us ⟹ degree u > 0
  shows length us ≤ n
proof (cases degree f = 0)
  case True
  have us = []
  proof (rule ccontr)
    assume us ≠ []
    from this obtain u where u: u ∈ set us by fastforce
    hence deg-u: degree u > 0 using d by auto
    have degree f = degree (prod-list us) unfolding f ..
```

```
    also have ... = sum-list (map degree us)
    proof (rule degree-prod-list-eq)
      fix p assume p: p ∈ set us
      show p ≠ 0 using d[OF p] degree-0 by auto
    qed
    also have  ... ≥ degree u by (simp add: member-le-sum-list u)
    finally have degree f > 0 using deg-u by auto
    thus False using True by auto
  qed
  thus ?thesis by simp
next
  case False
  hence f-not-0: f ≠ 0 using degree-0 by fastforce
  obtain P where fin-P: finite P and f-P: f = ∏ P and P: P ⊆ {p. irreducible
p ∧ monic p}
    using monic-square-free-irreducible-factorization[OF monic-f sf] by auto
  have n-card-P: n = card P
    using P False f-P fin-P length-berlekamp-basis-numbers-factors n by blast
  have distinct-us: distinct us using d f sf square-free-prod-list-distinct by blast
  let ?us'=(map normalize us)
  have distinct-us': distinct ?us'
  proof (auto simp add: distinct-map)
    show distinct us by (rule distinct-us)
    show inj-on normalize (set us)
    proof (auto simp add: inj-on-def, rule ccontr)
       fix x y assume x: x ∈ set us and y: y ∈ set us and n: normalize x =
normalize y
      and x-not-y: x ≠ y
      from normalize-eq-imp-smult[OF n]
      obtain c where c0: c ≠ 0 and y-smult: y = smult c x by blast
      have sf-xy: square-free (x*y)
      proof (rule square-free-factor[OF - sf])
        have x*y = prod-list [x,y] by simp
        also have ... dvd prod-list us
       by (rule prod-list-dvd-prod-list-subset, auto simp add: x y x-not-y distinct-us)
        also have ... = f unfolding f ..
        finally show x * y dvd f .
      qed
      have x * y = smult c (x*x) using y-smult mult-smult-right by auto
      hence sf-smult: square-free (smult c (x*x)) using sf-xy by auto
      have x*x dvd (smult c (x*x)) by (simp add: dvd-smult)
      hence ¬ square-free (smult c (x*x))
       by (metis d square-free-def x)
      thus False using sf-smult by contradiction
    qed
  qed
  have length-us-us': length us = length ?us' by simp
  have f-us': f = prod-list ?us'
  proof −
```

```
    have f = normalize f using monic-f f-not-0 by (simp add: normalize-monic)
    also have ... = prod-list ?us' by (unfold f, rule prod-list-normalize[of us])
    finally show ?thesis .
  qed
  have ∃ Q. prod-list Q = prod-list ?us' ∧ length ?us' ≤ length Q
        ∧ (∀ u. u ∈ set Q ⟶ irreducible u ∧ monic u)
  proof (rule exists-factorization-prod-list)
    show degree (prod-list ?us') > 0 using False f-us' by auto
    show square-free (prod-list ?us') using f-us' sf by auto
    fix u assume u: u ∈ set ?us'
    have u-not0: u ≠ 0 using d u degree-0 by fastforce
    have degree u > 0 using d u by auto
    moreover have monic u using u monic-normalize[OF u-not0] by auto
    ultimately show degree u > 0 ∧ monic u by simp
  qed
  from this obtain Q
  where Q-us': prod-list Q = prod-list ?us'
  and length-us'-Q: length ?us' ≤ length Q
  and Q: (∀ u. u ∈ set Q ⟶ irreducible u ∧ monic u)
  by blast
  have distinct-Q: distinct Q
  proof (rule square-free-prod-list-distinct)
    show square-free (prod-list Q) using Q-us' f-us' sf by auto
   show ⋀u. u ∈ set Q ⟹ degree u > 0 using Q irreducible-degree-field by auto
  qed
  have set-Q-P: set Q = P
  proof (rule monic-factorization-uniqueness)
    show ∏ (set Q) = ∏ P using Q-us'
      by (metis distinct-Q f-P f-us' list.map-ident prod.distinct-set-conv-list)
  qed (insert P Q fin-P, auto)
  hence length Q = card P using distinct-Q distinct-card by fastforce
  have length us = length ?us' by (rule length-us-us')
  also have ... ≤ length Q using length-us'-Q by auto
  also have ... = card (set Q) using distinct-card[OF distinct-Q] by simp
  also have ... = card P using set-Q-P by simp
  finally show ?thesis using n-card-P by simp
qed

lemma berlekamp-basis-irreducible: assumes f: f = prod-list us
  and n-us: length us = n
  and us: ⋀ u. u ∈ set us ⟹ degree u > 0
  and u: u ∈ set us
  shows irreducible u
proof (fold irreducible-connect-field, intro irreducible_dI[OF us[OF u]])
  fix q r :: 'a mod-ring poly
  assume dq: degree q > 0 and qu: degree q < degree u and dr: degree r > 0 and
uqr: u = q * r
  with us[OF u] have q: q ≠ 0 and r: r ≠ 0 by auto
  from split-list[OF u] obtain xs ys where id: us = xs @ u # ys by auto
```

```
  let ?us = xs @ q # r # ys
  have f: f = prod-list ?us unfolding f id uqr by simp
  {
    fix x
    assume x ∈ set ?us
    with us[unfolded id] dr dq have degree x > 0 by auto
  }
  from berlekamp-basis-length-factorization[OF f this]
  have length ?us ≤ n by simp
  also have ... = length us unfolding n-us by simp
  also have ... < length ?us unfolding id by simp
  finally show False by simp
qed
end

lemma not-irreducible-factor-yields-prime-factors:
  assumes uf: u dvd (f :: 'b :: {field-gcd} poly) and fin: finite P
      and fP: f = ∏ P and P: P ⊆ {q. irreducible q ∧ monic q}
    and u: degree u > 0 ¬ irreducible u
  shows ∃ pi pj. pi ∈ P ∧ pj ∈ P ∧ pi ≠ pj ∧ pi dvd u ∧ pj dvd u
proof −
  from finite-distinct-list[OF fin] obtain ps where Pps: P = set ps and dist:
distinct ps by auto
  have fP: f = prod-list ps unfolding fP Pps using dist
    by (simp add: prod.distinct-set-conv-list)
  note P = P[unfolded Pps]
  have set ps ⊆ P unfolding Pps by auto
  from uf[unfolded fP] P dist this
  show ?thesis
  proof (induct ps)
    case Nil
    with u show ?case using divides-degree[of u 1] by auto
  next
    case (Cons p ps)
    from Cons(3) have ps: set ps ⊆ {q. irreducible q ∧ monic q} by auto
    from Cons(2) have dvd: u dvd p * prod-list ps by simp
    obtain k where gcd: u = gcd p u * k by (meson dvd-def gcd-dvd2)
    from Cons(3) have *: monic p irreducible p p ≠ 0 by auto
    from monic-irreducible-gcd[OF *(1), of u] *(2)
    have gcd p u = 1 ∨ gcd p u = p by auto
    thus ?case
    proof
      assume gcd p u = 1
      then have Rings.coprime p u
        by (rule gcd-eq-1-imp-coprime)
      with dvd have u dvd prod-list ps
        using coprime-dvd-mult-right-iff coprime-imp-coprime by blast
      from Cons(1)[OF this ps] Cons(4−5) show ?thesis by auto
    next
```

```
      assume gcd p u = p
      with gcd have upk: u = p * k by auto
      hence p: p dvd u by auto
      from dvd[unfolded upk] *(3) have kps: k dvd prod-list ps by auto
      from dvd u * have dk: degree k > 0
        by (metis gr0I irreducible-mult-unit-right is-unit-iff-degree mult-zero-right
upk)
      from ps kps have ∃ q ∈ set ps. q dvd k
      proof (induct ps)
        case Nil
        with dk show ?case using divides-degree[of k 1] by auto
      next
        case (Cons p ps)
        from Cons(3) have dvd: k dvd p * prod-list ps by simp
        obtain l where gcd: k = gcd p k * l by (meson dvd-def gcd-dvd2)
        from Cons(2) have *: monic p irreducible p p ≠ 0 by auto
        from monic-irreducible-gcd[OF *(1), of k] *(2)
        have gcd p k = 1 ∨ gcd p k = p by auto
        thus ?case
        proof
          assume gcd p k = 1
          with dvd have k dvd prod-list ps
            by (metis dvd-triv-left gcd-greatest-mult mult.left-neutral)
          from Cons(1)[OF - this] Cons(2) show ?thesis by auto
        next
          assume gcd p k = p
          with gcd have upk: k = p * l by auto
          hence p: p dvd k by auto
          thus ?thesis by auto
        qed
      qed
      then obtain q where q: q ∈ set ps and dvd: q dvd k by auto
      from dvd upk have qu: q dvd u by auto
      from Cons(4) q have p ≠ q by auto
      thus ?thesis using q p qu Cons(5) by auto
    qed
  qed
qed

lemma berlekamp-factorization-main:
  fixes f::'a mod-ring poly
  assumes sf-f: square-free f
    and vs: vs = vs1 @ vs2
    and vsf: vs = berlekamp-basis f
    and n-bb: n = length (berlekamp-basis f)
    and n: n = length us1 + n2
    and us: us = us1 @ berlekamp-factorization-main d divs vs2 n2
    and us1: ⋀ u. u ∈ set us1 ⟹ monic u ∧ irreducible u
    and divs: ⋀ d. d ∈ set divs ⟹ monic d ∧ degree d > 0
```

```
    and vs1: ⋀ u v i. v ∈ set vs1 ⟹ u ∈ set us1 ∪ set divs
      ⟹ i < CARD('a) ⟹ gcd u (v − [:of-nat i:]) ∈ {1,u}
    and f: f = prod-list (us1 @ divs)
    and deg-f: degree f > 0
    and d: ⋀ g. g dvd f ⟹ degree g = d ⟹ irreducible g
  shows f = prod-list us ∧ (∀ u ∈ set us. monic u ∧ irreducible u)
proof −
  have mon-f: monic f unfolding f
    by (rule monic-prod-list, insert divs us1, auto)
  from monic-square-free-irreducible-factorization[OF mon-f sf-f] obtain P where
    P: finite P f = ∏ P P ⊆ {q. irreducible q ∧ monic q} by auto
  hence f0: f ≠ 0 by auto
  show ?thesis
    using vs n us divs f us1 vs1
  proof (induct vs2 arbitrary: divs n2 us1 vs1)
    case (Cons v vs2)
    show ?case
    proof (cases v = 1)
      case False
      from Cons(2) vsf have v: v ∈ set (berlekamp-basis f) by auto
      from berlekamp-basis-eq-8[OF this] have vf: [v ^ CARD('a) = v] (mod f) .
      let ?gcd = λ u i. gcd u (v − [:of-int i:])
      let ?gcdn = λ u i. gcd u (v − [:of-nat i:])
      let ?map = λ u. (map (λ i. ?gcd u i) [0 ..< CARD('a)])
      define udivs where udivs ≡ λ u. filter (λ w. w ≠ 1) (?map u)
      {
        obtain xs where xs: [0..<CARD('a)] = xs by auto
        have udivs = (λ u. [w. i ← [0 ..< CARD('a)], w ← [?gcd u i], w ≠ 1])
          unfolding udivs-def xs
          by (intro ext, auto simp: o-def, induct xs, auto)
      } note udivs-def' = this
      define facts where facts ≡ [ w . u ← divs, w ← udivs u]
      {
        fix u
        assume u: u ∈ set divs
       then obtain bef aft where divs: divs = bef @ u # aft by (meson split-list)
        from Cons(5)[OF u] have mon-u: monic u by simp
        have uf: u dvd f unfolding Cons(6) divs by auto
      from vf uf have vu: [v ^ CARD('a) = v] (mod u) by (rule cong-dvd-modulus-poly)
        from square-free-factor[OF uf sf-f] have sf-u: square-free u .
        let ?g = ?gcd u
        from mon-u have u0: u ≠ 0 by auto
        have u = (∏ c∈UNIV. gcd u (v − [:c:]))
          using Berlekamp-gcd-step[OF vu mon-u sf-u] .
        also have ... = (∏ i ∈ {0..< int CARD('a)}. ?g i)
       by (rule sym, rule prod.reindex-cong[OF to-int-mod-ring-hom.inj-f range-to-int-mod-ring[symmetric]],
          simp add: of-int-of-int-mod-ring)
        finally have u-prod: u = (∏ i ∈ {0..< int CARD('a)}. ?g i) .
        let ?S = {0..<int CARD('a)} − {i. ?g i = 1}
```

```
      {
        fix i
        assume i ∈ ?S
        hence ?g i ≠ 1 by auto
          moreover have mgi: monic (?g i) by (rule poly-gcd-monic, insert u0,
auto)
        ultimately have degree (?g i) > 0
          using monic-degree-0 by blast
        note this mgi
      } note gS = this

      have int-set: int ' set [0..<CARD('a)] = {0 ..< int CARD('a)}
        by (simp add: image-int-atLeastLessThan)

      have inj: inj-on ?g ?S unfolding inj-on-def
      proof (intro ballI impI)
        fix i j
        assume i: i ∈ ?S and j: j ∈ ?S and gij: ?g i = ?g j
        show i = j
        proof (rule ccontr)
          define S where S = {0..<int CARD('a)} - {i,j}
          have id: {0..<int CARD('a)} = (insert i (insert j S)) and S: i ∉ S j ∉
S finite S
            using i j unfolding S-def by auto
          assume ij: i ≠ j
          have u = (∏ i ∈ {0..< int CARD('a)}. ?g i) by fact
          also have ... = ?g i * ?g j * (∏ i ∈ S. ?g i)
            unfolding id using S ij by auto
          also have ... = ?g i * ?g i * (∏ i ∈ S. ?g i) unfolding gij by simp
          finally have dvd: ?g i * ?g i dvd u unfolding dvd-def by auto
            with sf-u[unfolded square-free-def, THEN conjunct2, rule-format, OF
gS(1)[OF i]]
          show False by simp
        qed
      qed

      have u = (∏ i ∈ {0..< int CARD('a)}. ?g i) by fact
      also have ... = (∏ i ∈ ?S. ?g i)
        by (rule sym, rule prod.setdiff-irrelevant, auto)
      also have ... = ∏ (set (udivs u)) unfolding udivs-def set-filter set-map
          by (rule sym, rule prod.reindex-cong[of ?g, OF inj - refl], auto simp:
int-set[symmetric])
      finally have u-udivs: u = ∏ (set (udivs u)) .
      {
        fix w
        assume mem: w ∈ set (udivs u)
        then obtain i where w: w = ?g i and i: i ∈ ?S
          unfolding udivs-def set-filter set-map int-set by auto
        have wu: w dvd u by (simp add: w)
```

```
      let ?v = λ j. v − [:of-nat j:]
      define j where j = nat i
    from i have j: of-int i = (of-nat j :: 'a mod-ring) j < CARD('a) unfolding
j-def by auto
      from gS[OF i, folded w] have ∗: degree w > 0 monic w w ≠ 0 by auto
      from w have w dvd ?v j using j by simp
      hence gcdj: ?gcdn w j = w by (metis gcd.commute gcd-left-idem j(1) w)
      {
        fix j'
        assume j': j' < CARD('a)
        have ?gcdn w j' ∈ {1,w}
        proof (rule ccontr)
          assume not: ?gcdn w j' ∉ {1,w}
          with gcdj have neq: int j' ≠ int j by auto

          let ?h = ?gcdn w j'
          from ∗(3) not have deg: degree ?h > 0
            using monic-degree-0 poly-gcd-monic by auto
          have hw: ?h dvd w by auto
          have ?h dvd ?gcdn u j' using wu using dvd-trans by auto
          also have ?gcdn u j' = ?g j' by simp
          finally have hj': ?h dvd ?g j' by auto
          from divides-degree[OF this] deg u0 have degj': degree (?g j') > 0 by
auto
          hence j'1: ?g j' ≠ 1 by auto
         with j' have mem': ?g j' ∈ set (udivs u) unfolding udivs-def by auto
          from degj' j' have j'S: int j' ∈ ?S by auto
          from i j have jS: int j ∈ ?S by auto
          from inj-on-contraD[OF inj neq j'S jS]
          have neq: w ≠ ?g j' using w j by auto
          have cop: ¬ coprime w (?g j') using hj' hw deg
            by (metis coprime-not-unit-not-dvd poly-dvd-1 Nat.neq0-conv)
          obtain w' where w': ?g j' = w' by auto
          from u-udivs sf-u have square-free (∏(set (udivs u))) by simp
          from square-free-prodD[OF this finite-set mem mem'] cop neq
          show False by simp
        qed
      }
      from gS[OF i, folded w] i this
       have degree w > 0 monic w ⋀ j. j < CARD('a) ⟹ ?gcdn w j ∈ {1,w}
by auto
    } note udivs = this
    let ?is = filter (λ i. ?g i ≠ 1) (map int [0 ..< CARD('a)])
    have id: udivs u = map ?g ?is
      unfolding udivs-def filter-map o-def ..
    have dist: distinct (udivs u) unfolding id distinct-map
    proof (rule conjI[OF distinct-filter], unfold distinct-map)
      have ?S = set ?is unfolding int-set[symmetric] by auto
      thus inj-on ?g (set ?is) using inj by auto
```

```
    qed (auto simp: inj-on-def)
    from u-udivs prod.distinct-set-conv-list[OF dist, of id]
    have prod-list (udivs u) = u by auto
    note udivs this dist
  } note udivs = this
  have facts: facts = concat (map udivs divs)
    unfolding facts-def by auto
   obtain lin nonlin where part: List.partition (λ q. degree q = d) facts =
(lin,nonlin)
    by force
  from Cons(6) have f = prod-list us1 * prod-list divs by auto
  also have prod-list divs = prod-list facts unfolding facts using udivs(4)
    by (induct divs, auto)
  finally have f: f = prod-list us1 * prod-list facts .
  note facts' = facts
  {
    fix u
    assume u: u ∈ set facts
     from u[unfolded facts] obtain u' where u': u' ∈ set divs and u: u ∈ set
(udivs u') by auto
    from u' udivs(1-2)[OF u' u] prod-list-dvd[OF u, unfolded udivs(4)[OF u']]
    have degree u > 0 monic u ∃ u' ∈ set divs. u dvd u' by auto
  } note facts = this
  have not1: (v = 1) = False using False by auto
  have us = us1 @ (if length divs = n2 then divs
     else let (lin, nonlin) = List.partition (λq. degree q = d) facts
         in lin @ berlekamp-factorization-main d nonlin vs2 (n2 − length lin))
    unfolding Cons(4) facts-def udivs-def' berlekamp-factorization-main.simps
Let-def not1 if-False
    by (rule arg-cong[where f = λ x. us1 @ x], rule if-cong, simp-all)
  hence res: us = us1 @ (if length divs = n2 then divs else
        lin @ berlekamp-factorization-main d nonlin vs2 (n2 − length lin))
    unfolding part by auto
  show ?thesis
  proof (cases length divs = n2)
    case False
   with res have us: us = (us1 @ lin) @ berlekamp-factorization-main d nonlin
vs2 (n2 − length lin)
      by auto
    from Cons(2) have vs: vs = (vs1 @ [v]) @ vs2 by auto
    have f: f = prod-list ((us1 @ lin) @ nonlin)
      unfolding f using prod-list-partition[OF part] by simp
    {
      fix u
      assume u ∈ set ((us1 @ lin) @ nonlin)
      with part have u ∈ set facts ∪ set us1 by auto
   with facts Cons(7) have degree u > 0 by (auto simp: irreducible-degree-field)
    } note deg = this
    from berlekamp-basis-length-factorization[OF sf-f n-bb mon-f f deg, unfolded
```

```
Cons(3)]
      have n2 ≥ length lin by auto
      hence n: n = length (us1 @ lin) + (n2 − length lin)
        unfolding Cons(3) by auto
      show ?thesis
      proof (rule Cons(1)[OF vs n us - f])
        fix u
        assume u ∈ set nonlin
        with part have u ∈ set facts by auto
        from facts[OF this] show monic u ∧ degree u > 0 by auto
      next
        fix u
        assume u: u ∈ set (us1 @ lin)
        {
          assume ∗: ¬ (monic u ∧ irreducible_d u)
          with Cons(7) u have u ∈ set lin by auto
          with part have uf: u ∈ set facts and deg: degree u = d by auto
          from facts[OF uf] obtain u' where u' ∈ set divs and uu': u dvd u' by
auto
           from this(1) have u' dvd f unfolding Cons(6) using prod-list-dvd[of
u'] by auto
          with uu' have u dvd f by (rule dvd-trans)
          from facts[OF uf] d[OF this deg] ∗ have False by auto
        }
        thus monic u ∧ irreducible u by auto
      next
        fix w u i
        assume w: w ∈ set (vs1 @ [v])
          and u: u ∈ set (us1 @ lin) ∪ set nonlin
          and i: i < CARD('a)
        from u part have u: u ∈ set us1 ∪ set facts by auto
        show gcd u (w − [:of-nat i:]) ∈ {1, u}
        proof (cases u ∈ set us1)
          case True
          from Cons(7)[OF this] have monic u irreducible u by auto
          thus ?thesis by (rule monic-irreducible-gcd)
        next
          case False
          with u have u: u ∈ set facts by auto
          show ?thesis
          proof (cases w = v)
            case True
            from u[unfolded facts'] obtain u' where u: u ∈ set (udivs u')
              and u': u' ∈ set divs by auto
            from udivs(3)[OF u' u i] show ?thesis unfolding True .
          next
            case False
            with w have w: w ∈ set vs1 by auto
            from u obtain u' where u': u' ∈ set divs and dvd: u dvd u'
```

```
            using facts(3)[of u] dvd-refl[of u] by blast
          from w have w ∈ set vs1 ∨ w = v by auto
          from facts(1−2)[OF u] have u: monic u by auto
          from Cons(8)[OF w - i] u'
          have gcd u' (w − [:of-nat i:]) ∈ {1, u'} by auto
          with dvd u show ?thesis by (rule monic-gcd-dvd)
        qed
      qed
    qed
  next
    case True
    with res have us: us = us1 @ divs by auto
    from Cons(3) True have n: n = length us unfolding us by auto
    show ?thesis unfolding us[symmetric]
    proof (intro conjI ballI)
      show f: f = prod-list us unfolding us using Cons(6) by simp
      {
        fix u
        assume u ∈ set us
        hence degree u > 0 using Cons(5) Cons(7)[unfolded irreducible_d-def]
          unfolding us by (auto simp: irreducible-degree-field)
      } note deg = this
      fix u
      assume u: u ∈ set us
      thus monic u unfolding us using Cons(5) Cons(7) by auto
      show irreducible u
        by (rule berlekamp-basis-irreducible[OF sf-f n-bb mon-f f n[symmetric]
deg u])
    qed
  qed
next
  case True
  with Cons(4) have us: us = us1 @ berlekamp-factorization-main d divs vs2
n2 by simp
  from Cons(2) True have vs: vs = (vs1 @ [1]) @ vs2 by auto
  show ?thesis
  proof (rule Cons(1)[OF vs Cons(3) us Cons(5−7)], goal-cases)
    case (3 v u i)
    show ?case
    proof (cases v = 1)
      case False
      with 3 Cons(8)[of v u i] show ?thesis by auto
    next
      case True
      hence deg: degree (v − [: of-nat i :]) = 0
        by (metis (no-types, opaque-lifting) degree-pCons-0 diff-pCons diff-zero
pCons-one)
      from 3(2) Cons(5,7)[of u] have monic u by auto
      from gcd-monic-constant[OF this deg] show ?thesis .
```

```
      qed
    qed
  qed
 next
  case Nil
  with vsf have vs1: vs1 = berlekamp-basis f by auto
  from Nil(3) have us: us = us1 @ divs by auto
  from Nil(4,6) have md: ⋀ u. u ∈ set us ⟹ monic u ∧ degree u > 0
   unfolding us by (auto simp: irreducible-degree-field)
  from Nil(7)[unfolded vs1] us
  have no-further-splitting-possible:
   ⋀ u v i. v ∈ set (berlekamp-basis f) ⟹ u ∈ set us
   ⟹ i < CARD('a) ⟹ gcd u (v − [:of-nat i:]) ∈ {1, u} by auto
  from Nil(5) us have prod: f = prod-list us by simp
  show ?case
  proof (intro conjI ballI)
   fix u
   assume u: u ∈ set us
   from md[OF this] have mon-u: monic u and deg-u: degree u > 0 by auto
   from prod u have uf: u dvd f by (simp add: prod-list-dvd)
    from monic-square-free-irreducible-factorization[OF mon-f sf-f] obtain P
where
     P: finite P f = ∏ P P ⊆ {q. irreducible q ∧ monic q} by auto
   show irreducible u
   proof (rule ccontr)
    assume irr-u: ¬ irreducible u
    from not-irreducible-factor-yields-prime-factors[OF uf P deg-u this]
    obtain pi pj where pij: pi ∈ P pj ∈ P pi ≠ pj pi dvd u pj dvd u by blast
    from exists-vector-in-Berlekamp-basis-dvd[OF
      deg-f berlekamp-basis-basis[OF deg-f, folded vs1] finite-set
      P pij(1−3) mon-f sf-f irr-u uf mon-u pij(4−5), unfolded vs1]
    obtain v s where v: v ∈ set (berlekamp-basis f)
     and gcd: gcd u (v − [:s:]) ∉ {1,u} using is-unit-gcd by auto
    from surj-of-nat-mod-ring[of s] obtain i where i: i < CARD('a) and s: s
= of-nat i by auto
    from no-further-splitting-possible[OF v u i] gcd[unfolded s]
    show False by auto
   qed
  qed (insert prod md, auto)
 qed
qed

lemma berlekamp-monic-factorization:
 fixes f::'a mod-ring poly
 assumes sf-f: square-free f
  and us: berlekamp-monic-factorization d f = us
  and d: ⋀ g. g dvd f ⟹ degree g = d ⟹ irreducible g
  and deg: degree f > 0
  and mon: monic f
```

shows $f = \mathit{prod\text{-}list}\ us \wedge (\forall\ u \in \mathit{set}\ us.\ \mathit{monic}\ u \wedge \mathit{irreducible}\ u)$
proof −
from $us[\mathit{unfolded}\ \mathit{berlekamp\text{-}monic\text{-}factorization\text{-}def}\ \mathit{Let\text{-}def}]\ \mathit{deg}$
have us: $us = [] @ \mathit{berlekamp\text{-}factorization\text{-}main}\ d\ [f]\ (\mathit{berlekamp\text{-}basis}\ f)\ (\mathit{length}\ (\mathit{berlekamp\text{-}basis}\ f))$
by (auto)
have id: $\mathit{berlekamp\text{-}basis}\ f = [] @ \mathit{berlekamp\text{-}basis}\ f$
$\mathit{length}\ (\mathit{berlekamp\text{-}basis}\ f) = \mathit{length}\ [] + \mathit{length}\ (\mathit{berlekamp\text{-}basis}\ f)$
$f = \mathit{prod\text{-}list}\ ([] @ [f])$
by auto
show $f = \mathit{prod\text{-}list}\ us \wedge (\forall\ u \in \mathit{set}\ us.\ \mathit{monic}\ u \wedge \mathit{irreducible}\ u)$
by $(\mathit{rule}\ \mathit{berlekamp\text{-}factorization\text{-}main}[\mathit{OF}\ \mathit{sf\text{-}f}\ \mathit{id}(1)\ \mathit{refl}\ \mathit{refl}\ \mathit{id}(2)\ us\ \text{-}\ \text{-}\ \text{-}\ \mathit{id}(3)],$
$\mathit{insert}\ \mathit{mon}\ \mathit{deg}\ d,\ \mathit{auto})$
qed
end

end

7 Distinct Degree Factorization

theory *Distinct-Degree-Factorization*
imports
Finite-Field
Polynomial-Factorization.Square-Free-Factorization
Berlekamp-Type-Based
begin

definition $\mathit{factors\text{-}of\text{-}same\text{-}degree} :: \mathit{nat} \Rightarrow 'a :: \mathit{field}\ \mathit{poly} \Rightarrow \mathit{bool}$ **where**
$\mathit{factors\text{-}of\text{-}same\text{-}degree}\ i\ f = (i \neq 0 \wedge \mathit{degree}\ f \neq 0 \wedge \mathit{monic}\ f \wedge (\forall\ g.\ \mathit{irreducible}\ g \longrightarrow g\ \mathit{dvd}\ f \longrightarrow \mathit{degree}\ g = i))$

lemma $\mathit{factors\text{-}of\text{-}same\text{-}degreeD}$: **assumes** $\mathit{factors\text{-}of\text{-}same\text{-}degree}\ i\ f$
shows $i \neq 0\ \mathit{degree}\ f \neq 0\ \mathit{monic}\ f\ g\ \mathit{dvd}\ f \Longrightarrow \mathit{irreducible}\ g = (\mathit{degree}\ g = i)$
proof −
note $* = \mathit{assms}[\mathit{unfolded}\ \mathit{factors\text{-}of\text{-}same\text{-}degree\text{-}def}]$
show i: $i \neq 0$ **and** f: $\mathit{degree}\ f \neq 0\ \mathit{monic}\ f$ **using** $*$ **by** auto
assume gf: $g\ \mathit{dvd}\ f$
with $*$ **have** $\mathit{irreducible}\ g \Longrightarrow \mathit{degree}\ g = i$ **by** auto
moreover
{
assume $**$: $\mathit{degree}\ g = i\ \neg\ \mathit{irreducible}\ g$
with $\mathit{irreducible}_d\text{-}\mathit{factor}[\mathit{of}\ g]\ i$ **obtain** $h1\ h2$ **where** irr: $\mathit{irreducible}\ h1$ **and** gh: $g = h1 * h2$
and $\mathit{deg\text{-}h2}$: $\mathit{degree}\ h2 < \mathit{degree}\ g$ **by** auto
from $**\ i$ **have** $g0$: $g \neq 0$ **by** auto
from $gf\ gh\ g0$ **have** $h1\ \mathit{dvd}\ f$ **using** $\mathit{dvd\text{-}mult\text{-}left}$ **by** blast
from $*\ f\ \mathit{this}\ \mathit{irr}$ **have** $\mathit{deg\text{-}h}$: $\mathit{degree}\ h1 = i$ **by** auto
from $\mathit{arg\text{-}cong}[\mathit{OF}\ gh,\ \mathit{of}\ \mathit{degree}]\ g0$ **have** $\mathit{degree}\ g = \mathit{degree}\ h1 + \mathit{degree}\ h2$
by $(\mathit{simp}\ \mathit{add}$: $\mathit{degree\text{-}mult\text{-}eq}\ gh)$

```
    with **(1) deg-h have degree h2 = 0 by auto
    from degree0-coeffs[OF this] obtain c where h2: h2 = [:c:] by auto
    with gh g0 have g: g = smult c h1 c ≠ 0 by auto
    with irr **(2) irreducible-smult-field[of c h1] have False by auto
  }
  ultimately show irreducible g = (degree g = i) by auto
qed



hide-const order
hide-const up-ring.monom


theorem (in field) finite-field-mult-group-has-gen2:
  assumes finite:finite (carrier R)
  shows ∃ a ∈ carrier (mult-of R). group.ord (mult-of R) a = order (mult-of R)
  ∧ carrier (mult-of R) = {a[^]i | i::nat . i ∈ UNIV}
proof −
  note mult-of-simps[simp]
  have finite': finite (carrier (mult-of R)) using finite by (rule finite-mult-of)

  interpret G: group mult-of R rewrites
     ([^]mult-of R) = (([^]) :: - ⇒ nat ⇒ -) and 1mult-of R = 1
    by (rule field-mult-group) (simp-all add: fun-eq-iff nat-pow-def)

  let ?N = λ x . card {a ∈ carrier (mult-of R). group.ord (mult-of R) a  = x}
  have 0 < order R − 1 unfolding Coset.order-def using card-mono[OF finite,
of {0, 1}] by simp
  then have *: 0 < order (mult-of R) using assms by (simp add: order-mult-of)
   have fin: finite {d. d dvd order (mult-of R) } using dvd-nat-bounds[OF *] by
force

  have (∑ d | d dvd order (mult-of R). ?N d)
       = card (UN d:{d . d dvd order (mult-of R) }. {a ∈ carrier (mult-of R).
group.ord (mult-of R) a  = d})
     (is - = card ?U)
    using fin finite by (subst card-UN-disjoint) auto
  also have ?U = carrier (mult-of R)
  proof
    { fix x assume x:x ∈ carrier (mult-of R)
      hence x':x∈carrier (mult-of R) by simp
      then have group.ord (mult-of R) x dvd order (mult-of R)
          using finite' G.ord-dvd-group-order[OF x'] by (simp add: order-mult-of)
      hence x ∈ ?U using dvd-nat-bounds[of order (mult-of R) group.ord (mult-of
R) x] x by blast
    } thus carrier (mult-of R) ⊆ ?U by blast
  qed auto
  also have card ... = Coset.order (mult-of R)
```

using *order-mult-of finite'* **by** (*simp add*: *Coset.order-def*)
finally have *sum-Ns-eq*: $(\sum d \mid d\ dvd\ order\ (mult\text{-}of\ R).\ ?N\ d) = order\ (mult\text{-}of\ R)$ **.**

{ fix *d* **assume** *d*:*d dvd order* (*mult-of R*)
have *card* $\{a \in carrier\ (mult\text{-}of\ R).\ group.ord\ (mult\text{-}of\ R)\ a = d\} \leq phi'\ d$
proof *cases*
assume *card* $\{a \in carrier\ (mult\text{-}of\ R).\ group.ord\ (mult\text{-}of\ R)\ a = d\} = 0$
thus *?thesis* **by** *presburger*
next
assume *card* $\{a \in carrier\ (mult\text{-}of\ R).\ group.ord\ (mult\text{-}of\ R)\ a = d\} \neq 0$
hence $\exists\, a \in carrier\ (mult\text{-}of\ R).\ group.ord\ (mult\text{-}of\ R)\ a = d$ **by** (*auto simp*: *card-eq-0-iff*)
thus *?thesis* **using** *num-elems-of-ord-eq-phi'*[*OF finite d*] **by** *auto*
qed
}
hence *all-le*:$\bigwedge i.\ i \in \{d.\ d\ dvd\ order\ (mult\text{-}of\ R)\ \} \Longrightarrow (\lambda i.\ card\ \{a \in carrier\ (mult\text{-}of\ R).\ group.ord\ (mult\text{-}of\ R)\ a = i\})\ i \leq (\lambda i.\ phi'\ i)\ i$ **by** *fast*
hence *le*:$(\sum i \mid i\ dvd\ order\ (mult\text{-}of\ R).\ ?N\ i) \leq (\sum i \mid i\ dvd\ order\ (mult\text{-}of\ R).\ phi'\ i)$
using *sum-mono*[*of* $\{d\ .\ d\ dvd\ order\ (mult\text{-}of\ R)\}$ $\lambda i.\ card\ \{a \in carrier\ (mult\text{-}of\ R).\ group.ord\ (mult\text{-}of\ R)\ a = i\}$] **by** *presburger*
have $order\ (mult\text{-}of\ R) = (\sum d \mid d\ dvd\ order\ (mult\text{-}of\ R).\ phi'\ d)$ **using** $*$
by (*simp add*: *sum-phi'-factors*)
hence *eq*:$(\sum i \mid i\ dvd\ order\ (mult\text{-}of\ R).\ ?N\ i) = (\sum i \mid i\ dvd\ order\ (mult\text{-}of\ R).\ phi'\ i)$ **using** *le sum-Ns-eq* **by** *presburger*
have $\bigwedge i.\ i \in \{d.\ d\ dvd\ order\ (mult\text{-}of\ R)\ \} \Longrightarrow ?N\ i = (\lambda i.\ phi'\ i)\ i$
proof (*rule ccontr*)
fix *i*
assume *i1*:$i \in \{d.\ d\ dvd\ order\ (mult\text{-}of\ R)\}$ **and** $?N\ i \neq phi'\ i$
hence $?N\ i = 0$
using *num-elems-of-ord-eq-phi'*[*OF finite, of i*] **by** (*auto simp*: *card-eq-0-iff*)
moreover have $0 < i$ **using** $*$ *i1* **by** (*simp add*: *dvd-nat-bounds*[*of order* (*mult-of R*) *i*])
ultimately have $?N\ i < phi'\ i$ **using** *phi'-nonzero* **by** *presburger*
hence $(\sum i \mid i\ dvd\ order\ (mult\text{-}of\ R).\ ?N\ i) < (\sum i \mid i\ dvd\ order\ (mult\text{-}of\ R).\ phi'\ i)$
using *sum-strict-mono-ex1*[*OF fin, of ?N* $\lambda\ i\ .\ phi'\ i$] *i1 all-le* **by** *auto*
thus *False* **using** *eq* **by** *force*
qed
hence $?N\ (order\ (mult\text{-}of\ R)) > 0$ **using** $*$ **by** (*simp add*: *phi'-nonzero*)
then obtain *a* **where** *a*:$a \in carrier\ (mult\text{-}of\ R)$ **and** *a-ord*:$group.ord\ (mult\text{-}of\ R)\ a = order\ (mult\text{-}of\ R)$
by (*auto simp add*: *card-gt-0-iff*)
hence *set-eq*:$\{a[\hat{}]i \mid i::nat.\ i \in UNIV\} = (\lambda x.\ a[\hat{}]x)\ `\{0\ ..\ group.ord\ (mult\text{-}of\ R)\ a - 1\}$

```
    using G.ord-elems[OF finite'] by auto
  have card-eq:card ((λx. a[^]x) ` {0 .. group.ord (mult-of R) a − 1}) = card {0
.. group.ord (mult-of R) a − 1}
    by (intro card-image G.ord-inj finite' a)
  hence card ((λ x . a[^]x) ` {0 .. group.ord (mult-of R) a − 1}) = card {0 ..order
(mult-of R) − 1}
    using assms by (simp add: card-eq a-ord)
  hence card-R-minus-1:card {a[^]i | i::nat. i ∈ UNIV} =  order (mult-of R)
    using ∗ by (subst set-eq) auto
  have ∗∗:{a[^]i | i::nat. i ∈ UNIV} ⊆ carrier (mult-of R)
    using G.nat-pow-closed[OF a] by auto
  with - have carrier (mult-of R) = {a[^]i|i::nat. i ∈ UNIV}
   by (rule card-seteq[symmetric]) (simp-all add: card-R-minus-1 finite Coset.order-def
del: UNIV-I)
  thus ?thesis using a a-ord by blast
qed



lemma add-power-prime-poly-mod-ring[simp]:
fixes x :: 'a::{prime-card} mod-ring poly
shows (x + y) ^ CARD('a)^n = x ^ (CARD('a)^n) + y ^ CARD('a)^n
proof (induct n arbitrary: x y)
  case 0
  then show ?case by auto
next
  case (Suc n)
  define p where p: p = CARD('a)
  have (x + y) ^ p ^ Suc n =  (x + y) ^ (p ∗ p^n) by simp
  also have ... = ((x + y) ^ p) ^ (p^n)
    by (simp add: power-mult)
  also have ... = (x^p + y^p)^(p^n)
    by (simp add: add-power-poly-mod-ring p)
   also have ... = (x^p)^(p^n) + (y^p)^(p^n) using Suc.hyps unfolding p by
auto
  also have ... = x^(p^(n+1)) + y^(p^(n+1)) by (simp add: power-mult)
  finally show ?case by (simp add: p)
qed


lemma fermat-theorem-mod-ring2[simp]:
fixes a::'a::{prime-card} mod-ring
shows a ^ (CARD('a)^n) = a
proof (induct n arbitrary: a)
  case (Suc n)
  define p where p = CARD('a)
  have a ^ p ^ Suc n = a ^ (p ∗ (p ^ n)) by simp
  also have ... = (a ^ p) ^ (p ^ n) by (simp add: power-mult)
   also have ... = a^(p ^ n) using fermat-theorem-mod-ring[of a^p] unfolding
```

p-def **by** *auto*
also have ... = *a* **using** *Suc.hyps p-def* **by** *auto*
finally show *?case* **by** (*simp add*: *p-def*)
qed *auto*

lemma *fermat-theorem-power-poly*[*simp*]:
fixes *a*::'*a*::*prime-card mod-ring*
shows [:*a*:] ^ *CARD*('*a*::*prime-card*) ^ *n* = [:*a*:]
by (*auto simp add*: *Missing-Polynomial.poly-const-pow mod-poly-less*)

lemma *degree-prod-monom*: *degree* (∏ *i* = *0*..<*n*. *monom 1 1*) = *n*
by (*metis degree-monom-eq prod-pow x-pow-n zero-neq-one*)

lemma *degree-monom0*[*simp*]: *degree* (*monom a 0*) = *0* **using** *degree-monom-le*
by *auto*
lemma *degree-monom0'*[*simp*]: *degree* (*monom 0 b*) = *0* **by** *auto*

lemma *sum-monom-mod*:
assumes *b* < *degree f*
shows (∑ *i*≤*b*. *monom* (*g i*) *i*) *mod f* = (∑ *i*≤*b*. *monom* (*g i*) *i*)
using *assms*
proof (*induct b*)
case *0*
then show *?case* **by** (*auto simp add*: *mod-poly-less*)
next
case (*Suc b*)
have *hyp*: (∑ *i*≤*b*. *monom* (*g i*) *i*) *mod f* = (∑ *i*≤*b*. *monom* (*g i*) *i*)
using *Suc.prems Suc.hyps* **by** *simp*
have *rw-monom*: *monom* (*g* (*Suc b*)) (*Suc b*) *mod f* = *monom* (*g* (*Suc b*)) (*Suc b*)
by (*metis Suc.prems degree-monom-eq mod-0 mod-poly-less monom-hom.hom-0-iff*)
have *rw*: (∑ *i*≤*Suc b*. *monom* (*g i*) *i*) = (*monom* (*g* (*Suc b*)) (*Suc b*) + (∑ *i*≤*b*. *monom* (*g i*) *i*))
by *auto*
have (∑ *i*≤*Suc b*. *monom* (*g i*) *i*) *mod f*
= (*monom* (*g* (*Suc b*)) (*Suc b*) + (∑ *i*≤*b*. *monom* (*g i*) *i*)) *mod f* **using** *rw* **by** *presburger*
also have ... =((*monom* (*g* (*Suc b*)) (*Suc b*)) *mod f*) + ((∑ *i*≤*b*. *monom* (*g i*) *i*) *mod f*)
using *poly-mod-add-left* **by** *auto*
also have ... = *monom* (*g* (*Suc b*)) (*Suc b*) + (∑ *i*≤*b*. *monom* (*g i*) *i*)
using *hyp rw-monom* **by** *presburger*
also have ... = (∑ *i*≤*Suc b*. *monom* (*g i*) *i*) **using** *rw* **by** *auto*
finally show *?case* **.**
qed

lemma *x-power-aq-minus-1-rw*:
fixes *x*::*nat*

```
  assumes x: x > 1
    and a: a > 0
    and b: b > 0
  shows x ^ (a * q) - 1 = ((x^a) - 1) * sum ((^) (x^a)) {..<q}
proof -
  have xa: (x ^ a) > 0 using x by auto
  have int-rw1: int (x ^ a) - 1 = int ((x ^ a) - 1)
    using xa by linarith
  have int-rw2: sum ((^) (int (x ^ a))) {..<q} = int (sum ((^) ((x ^ a))) {..<q})
    unfolding int-sum by simp
  have int (x ^ a) ^ q = int (Suc ((x ^ a) ^ q - 1)) using xa by auto
  hence int ((x ^ a) ^ q - 1) = int (x ^ a) ^ q - 1 using xa by presburger
  also have ... = (int (x ^ a) - 1) * sum ((^) (int (x ^ a))) {..<q}
    by (rule power-diff-1-eq)
  also have ... = (int ((x ^ a) - 1)) * int (sum ((^) ( (x ^ a))) {..<q})
    unfolding int-rw1 int-rw2 by simp
  also have ... = int (((x ^ a) - 1) * (sum ((^) ( (x ^ a))) {..<q})) by auto
  finally have aux: int ((x ^ a) ^ q - 1) = int (((x ^ a) - 1) * sum ((^) (x ^
a)) {..<q}) .
  have x ^ (a * q) - 1 = (x^a)^q - 1
    by (simp add: power-mult)
  also have ... = ((x^a) - 1) * sum ((^) (x^a)) {..<q}
    using aux unfolding int-int-eq .
  finally show ?thesis .
qed

lemma dvd-power-minus-1-conv1:
  fixes x::nat
  assumes x: x > 1
    and a: a > 0
    and xa-dvd: x ^ a - 1 dvd x^b - 1
    and b0: b > 0
  shows a dvd b
proof -
  define r where r[simp]: r = b mod a
  define q where q[simp]: q = b div a
  have b: b = a * q + r by auto
  have ra: r < a by (simp add: a)
  hence xr-less-xa: x ^ r - 1 < x ^ a - 1
    using x power-strict-increasing-iff diff-less-mono x by simp
  have dvd: x ^ a - 1 dvd x ^ (a * q) - 1
    using x-power-aq-minus-1-rw[OF x a b0] unfolding dvd-def by auto
  have x^b - 1 = x^b - x^r + x^r - 1
    using assms(1) assms(4) by auto
  also have ... = x^r * (x^(a*q) - 1) + x^r - 1
    by (metis (no-types, lifting) b diff-mult-distrib2 mult.commute nat-mult-1-right
power-add)
  finally have x^b - 1 = x^r * (x^(a*q) - 1) + x^r - 1 .
  hence x ^ a - 1 dvd x ^ r * (x ^ (a * q) - 1) + x ^ r - 1 using xa-dvd by
```

presburger
hence *x^a − 1 dvd x^r − 1*
by (*metis* (*no-types*) *diff-add-inverse diff-commute dvd dvd-diff-nat dvd-trans dvd-triv-right*)
hence *r = 0*
using *xr-less-xa*
by (*meson nat-dvd-not-less neq0-conv one-less-power x zero-less-diff*)
thus *?thesis* **by** *auto*
qed

lemma *dvd-power-minus-1-conv2*:
fixes *x*::*nat*
assumes *x*: *x > 1*
and *a*: *a > 0*
and *a-dvd-b*: *a dvd b*
and *b0*: *b > 0*
shows *x ^ a − 1 dvd x^b − 1*
proof −
define *q* **where** *q*[*simp*]: *q = b div a*
have *b*: *b = a * q* **using** *a-dvd-b* **by** *auto*
have *x^b − 1 = ((x ^ a) − 1) * sum ((^) (x ^ a)) {..<q}*
unfolding *b* **by** (*rule x-power-aq-minus-1-rw*[*OF x a b0*])
thus *?thesis* **unfolding** *dvd-def* **by** *auto*
qed

corollary *dvd-power-minus-1-conv*:
fixes *x*::*nat*
assumes *x*: *x > 1*
and *a*: *a > 0*
and *b0*: *b > 0*
shows *a dvd b = (x ^ a − 1 dvd x^b − 1)*
using *assms dvd-power-minus-1-conv1 dvd-power-minus-1-conv2* **by** *blast*

locale *poly-mod-type-irr = poly-mod-type m TYPE('a::prime-card)* **for** *m* +
fixes *f*::*'a*::{*prime-card*} *mod-ring poly*
assumes *irr-f*: *irreducible$_d$ f*
begin

definition *plus-irr* :: *'a mod-ring poly ⇒'a mod-ring poly ⇒ 'a mod-ring poly*
where *plus-irr a b = (a + b) mod f*

definition *minus-irr* :: *'a mod-ring poly ⇒'a mod-ring poly ⇒ 'a mod-ring poly*
where *minus-irr x y ≡ (x − y) mod f*

definition *uminus-irr* :: *'a mod-ring poly* ⇒*'a mod-ring poly*
where *uminus-irr x* = −*x*

definition *mult-irr* :: *'a mod-ring poly* ⇒*'a mod-ring poly* ⇒ *'a mod-ring poly*
where *mult-irr x y* = ((*x*∗*y*) *mod f*)

definition *carrier-irr* :: *'a mod-ring poly set*
where *carrier-irr* = {*x*. *degree x* < *degree f*}

definition *power-irr* :: *'a mod-ring poly* ⇒ *nat* ⇒ *'a mod-ring poly*
where *power-irr p n* = ((*p*^*n*) *mod f*)

definition *R* = (|*carrier* = *carrier-irr*, *monoid.mult* = *mult-irr*, *one* = *1*, *zero* = *0*, *add* = *plus-irr*|)

lemma *degree-f*[*simp*]: *degree f* > *0*
using *irr-f irreducible*$_d$*D*(*1*) **by** *blast*

lemma *element-in-carrier*: (*a* ∈ *carrier R*) = (*degree a* < *degree f*)
unfolding *R-def carrier-irr-def* **by** *auto*

lemma *f-dvd-ab*:
a = *0* ∨ *b* = *0* **if** *f dvd a* ∗ *b*
and *a*: *degree a* < *degree f*
and *b*: *degree b* < *degree f*
proof (*rule ccontr*)
assume ¬ (*a* = *0* ∨ *b* = *0*)
then have *a* ≠ *0* **and** *b* ≠ *0*
by *simp-all*
with *a b* **have** ¬ *f dvd a* **and** ¬ *f dvd b*
by (*auto simp add*: *mod-poly-less dvd-eq-mod-eq-0*)
moreover from ‹*f dvd a* ∗ *b*› *irr-f* **have** *f dvd a* ∨ *f dvd b*
by *auto*
ultimately show *False*
by *simp*
qed

lemma *ab-mod-f0*:
a = *0* ∨ *b* = *0* **if** *a* ∗ *b mod f* = *0*
and *a*: *degree a* < *degree f*
and *b*: *degree b* < *degree f*
using *that f-dvd-ab* **by** *auto*

lemma *irreducible*$_d$*D2*:
fixes *p q* :: *'b*::{*comm-semiring-1*,*semiring-no-zero-divisors*} *poly*
assumes *irreducible*$_d$ *p*
and *degree q* < *degree p* **and** *degree q* ≠ *0*
shows ¬ *q dvd p*

using *assms irreducible$_d$-dvd-smult* **by** *force*

lemma *times-mod-f-1-imp-0*:
assumes *x*: *degree x < degree f*
and *x2*: $\forall$ *xa. x * xa mod f = 1* $\longrightarrow$ $\neg$ *degree xa < degree f*
shows *x = 0*
proof (*rule ccontr*)
assume *x3*: *x ≠ 0*
let *?u = fst* (*bezout-coefficients f x*)
let *?v = snd* (*bezout-coefficients f x*)
have *?u * f + ?v * x = gcd f x* **using** *bezout-coefficients-fst-snd* **by** *auto*
also have *... = 1*
proof (*rule ccontr*)
assume *g*: *gcd f x ≠ 1*
have *degree* (*gcd f x*) *< degree f*
by (*metis degree-0 dvd-eq-mod-eq-0 gcd-dvd1 gcd-dvd2 irr-f*
irreducible$_d$D(*1*) *mod-poly-less nat-neq-iff x x3*)
have $\neg$ *gcd f x dvd f*
proof (*rule irreducible$_d$D2*[*OF irr-f*])
show *degree* (*gcd f x*) *< degree f*
by (*metis degree-0 dvd-eq-mod-eq-0 gcd-dvd1 gcd-dvd2 irr-f*
irreducible$_d$D(*1*) *mod-poly-less nat-neq-iff x x3*)
show *degree* (*gcd f x*) *≠ 0*
by (*metis* (*no-types*, *opaque-lifting*) *g degree-mod-less′ gcd.bottom-left-bottom gcd-eq-0-iff*
gcd-left-idem gcd-mod-left gr-implies-not0 x)
qed
moreover have *gcd f x dvd f* **by** *auto*
ultimately show *False* **by** *contradiction*
qed
finally have *?v*x mod f = 1*
by (*metis degree-1 degree-f mod-mult-self3 mod-poly-less*)
hence (*x*(*?v mod f*)) *mod f = 1*
by (*simp add*: *mod-mult-right-eq mult.commute*)
moreover have *degree* (*?v mod f*) *< degree f*
by (*metis degree-0 degree-f degree-mod-less′ not-gr-zero*)
ultimately show *False* **using** *x2* **by** *auto*
qed

sublocale *field-R*: *field R*
proof −
have *∗*: $\exists$ *y. degree y < degree f* $\wedge$ *f dvd x + y* **if** *degree x < degree f*
for *x* :: *'a mod-ring poly*
proof −
from *that* **have** *degree* (− *x*) *< degree f*
by *simp*
moreover have *f dvd* (*x* + − *x*)
by *simp*

```
    ultimately show ?thesis
      by blast
  qed
  have **: degree (x * y mod f) < degree f
    if degree x < degree f and degree y < degree f
    for x y :: 'a mod-ring poly
    using that by (cases x = 0 ∨ y = 0)
      (auto intro: degree-mod-less' dest: f-dvd-ab)
  show field R
      by standard (auto simp add: R-def carrier-irr-def plus-irr-def mult-irr-def
Units-def algebra-simps degree-add-less mod-poly-less mod-add-eq mult-poly-add-left
mod-mult-left-eq mod-mult-right-eq mod-eq-0-iff-dvd ab-mod-f0 * ** dest: times-mod-f-1-imp-0)
qed

lemma zero-in-carrier[simp]: 0 ∈ carrier-irr unfolding carrier-irr-def by auto

lemma card-carrier-irr[simp]: card carrier-irr = CARD('a)^(degree f)
proof −
  let ?A = (carrier-vec (degree f):: 'a mod-ring vec set)
  have bij-A-carrier: bij-betw (Poly ∘ list-of-vec) ?A carrier-irr
  proof (unfold bij-betw-def, rule conjI)
    show inj-on (Poly ∘ list-of-vec) ?A by (rule inj-Poly-list-of-vec)
    show (Poly ∘ list-of-vec) ` ?A = carrier-irr
    proof (unfold image-def o-def carrier-irr-def, auto)
      fix xa assume xa ∈ ?A thus degree (Poly (list-of-vec xa)) < degree f
        using degree-Poly-list-of-vec irr-f by blast
    next
      fix x::'a mod-ring poly
      assume deg-x: degree x < degree f
      let ?xa = vec-of-list (coeffs x @ replicate (degree f − length (coeffs x)) 0)
      show ∃ xa∈carrier-vec (degree f). x = Poly (list-of-vec xa)
        by (rule bexI[of - ?xa], unfold carrier-vec-def, insert deg-x)
          (auto simp add: degree-eq-length-coeffs)
    qed
  qed
  have CARD('a)^(degree f) = card ?A
    by (simp add: card-carrier-vec)
  also have ... = card carrier-irr using bij-A-carrier bij-betw-same-card by blast
  finally show ?thesis ..
qed

lemma finite-carrier-irr[simp]: finite (carrier-irr)
proof −
  have degree f > degree 0 using degree-0 by auto
  hence carrier-irr ≠ {} using degree-0 unfolding carrier-irr-def
    by blast
  moreover have card carrier-irr ≠ 0 by auto
  ultimately show ?thesis using card-eq-0-iff by metis
qed
```

lemma *finite-carrier-R*[*simp*]: *finite* (*carrier R*) **unfolding** *R-def* **by** *simp*

lemma *finite-carrier-mult-of*[*simp*]: *finite* (*carrier* (*mult-of R*))
unfolding *carrier-mult-of* **by** *auto*

lemma *constant-in-carrier*[*simp*]: [:*a*:] ∈ *carrier R*
unfolding *R-def carrier-irr-def* **by** *auto*

lemma *mod-in-carrier*[*simp*]: *a mod f* ∈ *carrier R*
unfolding *R-def carrier-irr-def*
by (*auto*, *metis degree-0 degree-f degree-mod-less' less-not-refl*)

lemma *order-irr*: *Coset.order* (*mult-of R*) = *CARD*('*a*)^*degree f* − *1*
by (*simp add*: *card-Diff-singleton Coset.order-def carrier-mult-of R-def*)

lemma *element-power-order-eq-1*:
assumes *x*: *x* ∈ *carrier* (*mult-of R*)
shows *x* [^]$_{(mult\text{-}of\ R)}$ *Coset.order* (*mult-of R*) = **1**$_{(mult\text{-}of\ R)}$
by (*meson field-R.field-mult-group finite-carrier-mult-of group.pow-order-eq-1 x*)

corollary *element-power-order-eq-1'*:
assumes *x*: *x* ∈ *carrier* (*mult-of R*)
shows*x* [^]$_{(mult\text{-}of\ R)}$ *CARD*('*a*)^*degree f* = *x*
proof −
have *x* [^]$_{(mult\text{-}of\ R)}$ *CARD*('*a*)^*degree f*
= *x* ⊗$_{(mult\text{-}of\ R)}$ *x* [^]$_{(mult\text{-}of\ R)}$ (*CARD*('*a*)^*degree f* − *1*)
by (*metis Diff-iff One-nat-def Suc-pred field-R.m-comm field-R.nat-pow-Suc field-R.nat-pow-closed mult-of-simps*(*1*) *mult-of-simps*(*2*) *nat-pow-mult-of neq0-conv power-eq-0-iff x zero-less-card-finite*)
also have *x* ⊗$_{(mult\text{-}of\ R)}$ *x* [^]$_{(mult\text{-}of\ R)}$ (*CARD*('*a*)^*degree f* − *1*) = *x*
by (*metis carrier-mult-of element-power-order-eq-1 field-R.Units-closed field-R.field-Units field-R.r-one monoid.simps*(*2*) *mult-mult-of mult-of-def order-irr x*)
finally show *?thesis* **.**
qed

lemma *pow-irr*[*simp*]: *x* [^]$_{(R)}$ *n*= *x*^*n mod f*
by (*induct n*, *auto simp add*: *mod-poly-less nat-pow-def R-def mult-of-def mult-irr-def carrier-irr-def mod-mult-right-eq mult.commute*)

lemma *pow-irr-mult-of*[*simp*]: *x* [^]$_{(mult\text{-}of\ R)}$ *n*= *x*^*n mod f*
by (*induct n*, *auto simp add*: *mod-poly-less nat-pow-def R-def mult-of-def mult-irr-def carrier-irr-def mod-mult-right-eq mult.commute*)

lemma *fermat-theorem-power-poly-R*[*simp*]: [:*a*:] [^]$_R$ *CARD*('*a*) ^ *n* = [:*a*:]
by (*auto simp add*: *Missing-Polynomial.poly-const-pow mod-poly-less*)

lemma *times-mod-expand*:
($a \otimes_{(R)} b$) = ((a *mod* f) $\otimes_{(R)}$ (b *mod* f))
by (*simp add*: *mod-mult-eq R-def mult-irr-def*)

lemma *mult-closed-power*:
assumes *x*: $x \in$ *carrier* R **and** *y*: $y \in$ *carrier* R
and x [^]$_{(R)}$ *CARD*('*a*) ^ $m' = x$
and y [^]$_{(R)}$ *CARD*('*a*) ^ $m' = y$
shows ($x \otimes_{(R)} y$) [^]$_{(R)}$ *CARD*('*a*) ^ $m' = (x \otimes_{(R)} y)$
using *assms assms field-R.nat-pow-distrib* **by** *auto*

lemma *add-closed-power*:
assumes *x1*: x [^]$_{(R)}$ *CARD*('*a*) ^ $m' = x$
and *y1*: y [^]$_{(R)}$ *CARD*('*a*) ^ $m' = y$
shows ($x \oplus_{(R)} y$) [^]$_{(R)}$ *CARD*('*a*) ^ $m' = (x \oplus_{(R)} y)$
proof −
have ($x + y$) ^ *CARD*('*a*) ^ $m' = x$^(*CARD*('*a*) ^ m') + y ^ (*CARD*('*a*) ^ m')
by *auto*
hence ($x + y$) ^ *CARD*('*a*) ^ m' *mod* f = (x^(*CARD*('*a*) ^ m') + y ^ (*CARD*('*a*)
^ m')) *mod* f **by** *auto*
hence ($x \oplus_{(R)} y$) [^]$_{(R)}$ *CARD*('*a*) ^ m'
= (x [^]$_{(R)}$ *CARD*('*a*)^m') $\oplus_{(R)}$ (y [^]$_{(R)}$ *CARD*('*a*)^m')
by (*auto*, *unfold R-def plus-irr-def*, *auto simp add*: *mod-add-eq power-mod*)
also have ... = $x \oplus_{(R)} y$ **unfolding** *x1 y1* **by** *simp*
finally show *?thesis* **.**
qed

lemma *x-power-pm-minus-1*:
assumes *x*: $x \in$ *carrier* (*mult-of R*)
and x [^]$_{(R)}$ *CARD*('*a*) ^ $m' = x$
shows x [^]$_{(R)}$ (*CARD*('*a*) ^ $m' - 1$) = $\mathbf{1}_{(R)}$
by (*metis* (*no-types*, *lifting*) *One-nat-def Suc-pred assms*(*2*) *carrier-mult-of field-R.Units-closed*
field-R.Units-l-cancel field-R.field-Units field-R.l-one field-R.m-rcancel field-R.nat-pow-Suc
field-R.nat-pow-closed field-R.one-closed field-R.r-null field-R.r-one x zero-less-card-finite
zero-less-power)

context
begin

private lemma *monom-a-1-P*:

assumes *m*: *monom 1 1* ∈ *carrier R*
and *eq*: *monom 1 1* $[\hat{}]_{(R)}$ (*CARD*('*a*) ^ *m'*) = *monom 1 1*
shows *monom a 1* $[\hat{}]_{(R)}$ (*CARD*('*a*) ^ *m'*) = *monom a 1*
proof −
have *monom a 1* = [:*a*:] ∗ (*monom 1 1*)
by (*metis One-nat-def monom-0 monom-Suc mult.commute pCons-0-as-mult*)
also have ... = [:*a*:] $\otimes_{(R)}$ (*monom 1 1*)
by (*auto simp add*: *R-def mult-irr-def*)
(*metis One-nat-def assms*(*2*) *mod-mod-trivial mod-smult-left pow-irr*)
finally have *eq2*: *monom a 1* = [:*a*:] $\otimes_R$ *monom 1 1* **.**
show *?thesis* **unfolding** *eq2*
by (*rule mult-closed-power*[*OF - m - eq*], *insert fermat-theorem-power-poly-R*, *auto*)
qed

private lemma *prod-monom-1-1*:
defines *P* == (λ *x n*. (*x*[$\hat{}$]$_{(R)}$ (*CARD*('*a*) ^ *n*) = *x*))
assumes *m*: *monom 1 1* ∈ *carrier R*
and *eq*: *P* (*monom 1 1*) *n*
shows *P* ((∏ *i* = *0*..<*b*::*nat*. *monom 1 1*) *mod f*) *n*
proof (*induct b*)
case *0*
then show *?case* **unfolding** *P-def*
by (*simp add*: *power-mod*)
next
case (*Suc b*)
let *?N* = (∏ *i* = *0*..<*b*. *monom 1 1*)
have *eq2*: (∏ *i* = *0*..<*Suc b*. *monom 1 1*) *mod f* = *monom 1 1* $\otimes_{(R)}$ (∏ *i* = *0*..<*b*. *monom 1 1*)
by (*metis field-R.m-comm field-R.nat-pow-Suc mod-in-carrier mod-mod-trivial pow-irr prod-pow times-mod-expand*)
also have ... = (*monom 1 1 mod f*) $\otimes_{(R)}$ ((∏ *i* = *0*..<*b*. *monom 1 1*) *mod f*)
by (*rule times-mod-expand*)
finally have *eq2*: (∏ *i* = *0*..<*Suc b*. *monom 1 1*) *mod f*
= (*monom 1 1 mod f*) $\otimes_{(R)}$ ((∏ *i* = *0*..<*b*. *monom 1 1*) *mod f*) **.**
show *?case*
unfolding *eq2 P-def*
proof (*rule mult-closed-power*)
show (*monom 1 1 mod f*) $[\hat{}]_R$ *CARD*('*a*) ^ *n* = *monom 1 1 mod f*
using *P-def element-in-carrier eq m mod-poly-less* **by** *force*
show ((∏ *i* = *0*..<*b*. *monom 1 1*) *mod f*) $[\hat{}]_R$ *CARD*('*a*) ^ *n* = (∏ *i* = *0*..<*b*. *monom 1 1*) *mod f*
using *P-def Suc.hyps* **by** *blast*
qed (*auto*)
qed

private lemma *monom-1-b*:

defines $P == (\lambda\ x\ n.\ (x[\widehat{\ }]_{(R)}\ (CARD('a)\ \widehat{\ }\ n) = x))$
assumes *m*: $monom\ 1\ 1 \in carrier\ R$
and *monom-1-1*: $P\ (monom\ 1\ 1)\ m'$
and *b*: $b < degree\ f$
shows $P\ (monom\ 1\ b)\ m'$
proof −
have $monom\ 1\ b = (\prod i = 0..<b.\ monom\ 1\ 1)$
by (*metis prod-pow x-pow-n*)
also have $... = (\prod i = 0..<b.\ monom\ 1\ 1)\ mod\ f$
by (*rule mod-poly-less*[*symmetric*], *auto*)
(*metis One-nat-def b degree-linear-power x-as-monom*)
finally have *eq2*: $monom\ 1\ b = (\prod i = 0..<b.\ monom\ 1\ 1)\ mod\ f$ **.**
show *?thesis* **unfolding** *eq2 P-def*
by (*rule prod-monom-1-1*[*OF m monom-1-1*[*unfolded P-def*]])
qed

private lemma *monom-a-b*:
defines $P == (\lambda\ x\ n.\ (x[\widehat{\ }]_{(R)}\ (CARD('a)\ \widehat{\ }\ n) = x))$
assumes *m*: $monom\ 1\ 1 \in carrier\ R$
and *m1*: $P\ (monom\ 1\ 1)\ m'$
and *b*: $b < degree\ f$
shows $P\ (monom\ a\ b)\ m'$
proof −
have $monom\ a\ b = smult\ a\ (monom\ 1\ b)$
by (*simp add*: *smult-monom*)
also have $... = [:a:] * (monom\ 1\ b)$ **by** *auto*
also have $... = [:a:] \otimes_{(R)} (monom\ 1\ b)$
unfolding *R-def mult-irr-def*
by (*simp add*: *b degree-monom-eq mod-poly-less*)
finally have *eq*: $monom\ a\ b = [:a:] \otimes_{(R)} (monom\ 1\ b)$ **.**
show *?thesis* **unfolding** *eq P-def*
proof (*rule mult-closed-power*)
show $[:a:]\ [\widehat{\ }]_R\ CARD('a)\ \widehat{\ }\ m' = [:a:]$ **by** (*rule fermat-theorem-power-poly-R*)
show $monom\ 1\ b\ [\widehat{\ }]_R\ CARD('a)\ \widehat{\ }\ m' = monom\ 1\ b$
unfolding *P-def* **by** (*rule monom-1-b*[*OF m m1*[*unfolded P-def*] *b*])
show $monom\ 1\ b \in carrier\ R$ **unfolding** *element-in-carrier* **using** *b*
by (*simp add*: *degree-monom-eq*)
qed (*auto*)
qed

private lemma *sum-monoms-P*:
defines $P == (\lambda\ x\ n.\ (x[\widehat{\ }]_{(R)}\ (CARD('a)\ \widehat{\ }\ n) = x))$
assumes *m*: $monom\ 1\ 1 \in carrier\ R$
and *monom-1-1*: $P\ (monom\ 1\ 1)\ n$
and *b*: $b < degree\ f$
shows $P\ ((\sum i{\le}b.\ monom\ (g\ i)\ i))\ n$

```
  using b
proof (induct b)
  case 0
  then show ?case unfolding P-def
    by (simp add: poly-const-pow mod-poly-less monom-0)
next
  case (Suc b)
  have b: b < degree f using Suc.prems by auto
  have rw: (∑ i≤b. monom (g i) i) mod f = (∑ i≤b. monom (g i) i) by (rule
sum-monom-mod[OF b])
  have rw2: (monom (g (Suc b)) (Suc b) mod f) = monom (g (Suc b)) (Suc b)
   by (metis Suc.prems field-R.nat-pow-eone m monom-a-b pow-irr power-0 power-one-right)
  have hyp: P (∑ i≤b. monom (g i) i) n using Suc.prems Suc.hyps by auto
  have (∑ i≤Suc b. monom (g i) i) = monom (g (Suc b)) (Suc b) + (∑ i≤b.
monom (g i) i)
    by simp
  also have ... = (monom (g (Suc b)) (Suc b) mod f) + ((∑ i≤b. monom (g i) i)
mod f)
    using rw rw2 by argo
  also have ... = monom (g (Suc b)) (Suc b) ⊕_R (∑ i≤b. monom (g i) i)
    unfolding R-def plus-irr-def
    by (simp add: poly-mod-add-left)
  finally have eq: (∑ i≤Suc b. monom (g i) i)
    = monom (g (Suc b)) (Suc b) ⊕_R (∑ i≤b. monom (g i) i) .
  show ?case unfolding eq P-def
  proof (rule add-closed-power)
    show monom (g (Suc b)) (Suc b) [^]_R CARD('a) ^ n = monom (g (Suc b))
(Suc b)
      by (rule monom-a-b[OF m monom-1-1[unfolded P-def] Suc.prems])
    show (∑ i≤b. monom (g i) i) [^]_R CARD('a) ^ n = (∑ i≤b. monom (g i) i)
      using hyp unfolding P-def by simp
  qed
qed

lemma element-carrier-P:
  defines P ≡ (λ x n. (x[^]_(R) (CARD('a) ^ n) = x))
  assumes m: monom 1 1 ∈ carrier R
  and monom-1-1: P (monom 1 1) m'
  and a: a ∈ carrier R
shows P a m'
proof −
  have degree-a: degree a < degree f using a element-in-carrier by simp
  have P (∑ i≤degree a. monom (poly.coeff a i) i) m'
    unfolding P-def
    by (rule sum-monoms-P[OF m monom-1-1[unfolded P-def] degree-a])
  thus ?thesis unfolding poly-as-sum-of-monoms by simp
qed
end
```

end

lemma *degree-divisor1*:
assumes *f*: *irreducible* (*f* :: *'a* :: *prime-card mod-ring poly*)
and *d*: *degree f* = *d*
shows *f dvd* (*monom 1 1*)^(*CARD*(*'a*)^*d*) − *monom 1 1*
proof −
interpret *poly-mod-type-irr CARD*(*'a*) *f* **by** (*unfold-locales*, *auto simp add*: *f*)
show *?thesis*
proof (*cases d* = *1*)
case *True*
show *?thesis*
proof (*cases monom 1 1 mod f* = *0*)
case *True*
then show *?thesis*
by (*metis Suc-pred dvd-diff dvd-mult2 mod-eq-0-iff-dvd power.simps*(*2*)
zero-less-card-finite zero-less-power)
next
case *False* **note** *mod-f-not0* = *False*
have *monom 1* (*CARD*(*'a*)) *mod f* = *monom 1 1 mod f*
proof −
let *?g1* = (*monom 1* (*CARD*(*'a*))) *mod f*
let *?g2* = (*monom 1 1*) *mod f*
have *deg-g1*: *degree ?g1* < *degree f* **and** *deg-g2*: *degree ?g2* < *degree f*
by (*metis True card-UNIV-unit d degree-0 degree-mod-less' zero-less-card-finite zero-neq-one*)+
have *g2*: *?g2* [^]$_{(mult\text{-}of\ R)}$ *CARD*(*'a*)^*degree f* = *?g2* ^ (*CARD*(*'a*)^*degree f*) *mod f*
by (*rule pow-irr-mult-of*)
have *?g2* [^]$_{(mult\text{-}of\ R)}$ *CARD*(*'a*)^*degree f* = *?g2*
by (*rule element-power-order-eq-1'*, *insert mod-f-not0 deg-g2*,
auto simp add: *carrier-mult-of R-def carrier-irr-def*)
hence *?g2* ^ *CARD*(*'a*) *mod f* = *?g2 mod f* **using** *True d* **by** *auto*
hence *?g1 mod f* = *?g2 mod f* **by** (*metis mod-mod-trivial power-mod x-pow-n*)
thus *?thesis* **by** *simp*
qed
thus *?thesis* **by** (*metis True mod-eq-dvd-iff-poly power-one-right x-pow-n*)
qed
next
case *False*
have *deg-f1*: *1* < *degree f*
using *False d degree-f* **by** *linarith*
have *monom 1 1* [^]$_{(mult\text{-}of\ R)}$ *CARD*(*'a*)^*degree f* = *monom 1 1*
by (*rule element-power-order-eq-1'*, *insert deg-f1*)
(*auto simp add*: *carrier-mult-of R-def carrier-irr-def degree-monom-eq*)
hence *monom 1 1*^*CARD*(*'a*)^*degree f mod f* = *monom 1 1 mod f*
using *deg-f1* **by** (*auto*, *metis mod-mod-trivial*)
thus *?thesis* **using** *d mod-eq-dvd-iff-poly* **by** *blast*

```
  qed
qed


lemma degree-divisor2:
  assumes f: irreducible (f :: 'a :: prime-card mod-ring poly)
  and d: degree f = d
  and c-ge-1: 1 ≤ c and cd: c < d
shows ¬ f dvd monom 1 1 ^ CARD('a) ^ c − monom 1 1
proof (rule ccontr)
  interpret poly-mod-type-irr CARD('a) f by (unfold-locales, auto simp add: f)
  have field-R: field R
    by (simp add: field-R.field-axioms)
  assume ¬ ¬ f dvd monom 1 1 ^ CARD('a) ^ c − monom 1 1
  hence f-dvd: f dvd monom 1 1 ^ CARD('a) ^ c − monom 1 1 by simp
  obtain a where a-R: a ∈ carrier (mult-of R)
    and ord-a: group.ord (mult-of R) a = order (mult-of R)
    and gen: carrier (mult-of R) = {a [^]_R i |i. i ∈ (UNIV::nat set)}
    using field.finite-field-mult-group-has-gen2[OF field-R] by auto
  have d-not1: d>1 using c-ge-1 cd by auto
  have monom-in-carrier: monom 1 1 ∈ carrier (mult-of R)
    using d-not1 unfolding carrier-mult-of R-def carrier-irr-def
    by (simp add: d degree-monom-eq)
  then have monom 1 1 ∉ {0_R}
    by auto
  then obtain k where monom 1 1 = a ^ k mod f
    using gen monom-in-carrier by auto
  then have k: a [^]_R k = monom 1 1
    by simp
  have a-m-1: a [^]_R (CARD('a)^c − 1) = 1_R
  proof (rule x-power-pm-minus-1[OF a-R])
    let ?x = monom 1 1::'a mod-ring poly
    show a [^]_R CARD('a) ^ c = a
    proof (rule element-carrier-P)
      show ?x ∈ carrier R
        by (metis k mod-in-carrier pow-irr)
      have ?x ^ CARD('a)^ c mod f = ?x mod f using f-dvd
        using mod-eq-dvd-iff-poly by blast
      thus ?x [^]_R CARD('a)^ c = ?x
        by (metis d d-not1 degree-monom-eq mod-poly-less one-neq-zero pow-irr)
      show a ∈ carrier R using a-R unfolding carrier-mult-of by auto
    qed
  qed
  have Group.group (mult-of R)
    by (simp add: field-R.field-mult-group)
  moreover have finite (carrier (mult-of R)) by auto
  moreover have a ∈ carrier (mult-of R) by (rule a-R )
  moreover have a [^]_mult-of R (CARD('a) ^ c − 1) = 1_mult-of R
    using a-m-1 unfolding mult-of-def
```

```
    by (auto, metis mult-of-def pow-irr-mult-of nat-pow-mult-of)
  ultimately have ord-dvd: group.ord (mult-of R) a dvd (CARD('a)^c - 1)
    by (meson group.pow-eq-id)
  have d dvd c
  proof (rule dvd-power-minus-1-conv1[OF nontriv])
    show 0 < d using cd by auto
    show CARD('a) ^ d - 1 dvd CARD('a) ^ c - 1
      using ord-dvd by (simp add: d ord-a order-irr)
    show 0 < c using c-ge-1 by auto
  qed
  thus False using c-ge-1 cd
    using nat-dvd-not-less by auto
qed

lemma degree-divisor: assumes irreducible (f :: 'a :: prime-card mod-ring poly)
degree f = d
  shows f dvd (monom 1 1)^(CARD('a)^d) - monom 1 1
  and 1 ≤ c ⟹ c < d ⟹ ¬ f dvd (monom 1 1)^(CARD('a)^c) - monom 1 1
    using assms degree-divisor1 degree-divisor2 by blast+

context
  assumes SORT-CONSTRAINT('a :: prime-card)
begin

function dist-degree-factorize-main ::
  'a mod-ring poly ⇒ 'a mod-ring poly ⇒ nat ⇒ (nat × 'a mod-ring poly) list
  ⇒ (nat × 'a mod-ring poly) list where
  dist-degree-factorize-main v w d res = (if v = 1 then res else if d + d > degree v
    then (degree v, v) # res else let
      w = w^(CARD('a)) mod v;
      d = Suc d;
      gd = gcd (w - monom 1 1) v
      in if gd = 1 then dist-degree-factorize-main v w d res else
      let v' = v div gd in
      dist-degree-factorize-main v' (w mod v') d ((d,gd) # res))
  by pat-completeness auto


termination
proof (relation measure (λ (v,w,d,res). Suc (degree v) - d), goal-cases)
  case (3 v w d res x xa xb xc)
  have xb dvd v unfolding 3 by auto
  hence xc dvd v unfolding 3 by (metis dvd-def dvd-div-mult-self)
  from divides-degree[OF this] 3
  show ?case by auto
qed auto

declare dist-degree-factorize-main.simps[simp del]
```

lemma *dist-degree-factorize-main*: **assumes**
dist: *dist-degree-factorize-main v w d res = facts* **and**
w: *w = (monom 1 1)^(CARD('a)^d) mod v* **and**
sf: *square-free u* **and**
mon: *monic u* **and**
prod: *u = v * prod-list (map snd res)* **and**
deg: ⋀ *f*. *irreducible f ⟹ f dvd v ⟹ degree f > d* **and**
res: ⋀ *i f*. *(i,f) ∈ set res ⟹ i ≠ 0 ∧ degree f ≠ 0 ∧ monic f ∧ (∀ g*. *irreducible g ⟶ g dvd f ⟶ degree g = i)*
shows *u = prod-list (map snd facts) ∧ (∀ i f*. *(i,f) ∈ set facts ⟶ factors-of-same-degree i f)*
using *dist w prod res deg* **unfolding** *factors-of-same-degree-def*
proof *(induct v w d res rule*: *dist-degree-factorize-main.induct)*
case *(1 v w d res)*
note *IH = 1(1−2)*
note *result = 1(3)*
note *w = 1(4)*
note *u = 1(5)*
note *res = 1(6)*
note *fact = 1(7)*
note *[simp] = dist-degree-factorize-main.simps[of - - d]*
let *?x = monom 1 1* :: *'a mod-ring poly*
show *?case*
proof *(cases v = 1)*
case *True*
thus *?thesis* **using** *result u mon res* **by** *auto*
next
case *False* **note** *v = this*
note *IH = IH[OF this]*
have *mon-prod*: *monic (prod-list (map snd res))* **by** *(rule monic-prod-list, insert res, auto)*
with *mon[unfolded u]* **have** *mon-v*: *monic v* **by** *(simp add*: *coeff-degree-mult)*
with *False* **have** *deg-v*: *degree v ≠ 0* **by** *(simp add*: *monic-degree-0)*
show *?thesis*
proof *(cases degree v < d + d)*
case *True*
with *result False* **have** *facts*: *facts = (degree v, v) # res* **by** *simp*
show *?thesis*
proof *(intro allI conjI impI)*
fix *i f g*
assume *∗*: *(i,f) ∈ set facts irreducible g g dvd f*
show *degree g = i*
proof *(cases (i,f) ∈ set res)*
case *True*
from *res[OF this] ∗* **show** *?thesis* **by** *auto*
next
case *False*
with *∗ facts* **have** *id*: *i = degree v f = v* **by** *auto*
note *∗ = ∗(2−3)[unfolded id]*

```
        from fact[OF *] have dg: d < degree g by auto
        from divides-degree[OF *(2)] mon-v have deg-gv: degree g ≤ degree v by
auto
        from *(2) obtain h where vgh: v = g * h unfolding dvd-def by auto
        from arg-cong[OF this, of degree] mon-v have dvgh: degree v = degree g
+ degree h
          by (metis deg-v degree-mult-eq degree-mult-eq-0)
        with dg deg-gv dg True have deg-h: degree h < d by auto
        {
          assume degree h = 0
          with dvgh have degree g = degree v by simp
        }
        moreover
        {
          assume deg-h0: degree h ≠ 0
          hence ∃ k. irreducible_d k ∧ k dvd h
            using dvd-triv-left irreducible_d-factor by blast
          then obtain k where irr: irreducible k and k dvd h by auto
          from dvd-trans[OF this(2), of v] vgh have k dvd v by auto
          from fact[OF irr this] have dk: d < degree k .
          from divides-degree[OF ‹k dvd h›] deg-h0 have degree k ≤ degree h by
auto
          with deg-h have degree k < d by auto
          with dk have False by auto
        }
        ultimately have degree g = degree v by auto
        thus ?thesis unfolding id by auto
      qed
    qed (insert v mon-v deg-v u facts res, force+)
  next
    case False
    note IH = IH[OF this refl refl refl]
    let ?p = CARD('a)
    let ?w = w ^ ?p mod v
    let ?g = gcd (?w - ?x) v
    let ?v = v div ?g
    let ?d = Suc d
    from result[simplified] v False
    have result: (if ?g = 1 then dist-degree-factorize-main v ?w ?d res
            else dist-degree-factorize-main ?v (?w mod ?v) ?d ((?d, ?g) # res))
= facts
      by (auto simp: Let-def)
    from mon-v have mon-g: monic ?g by (metis deg-v degree-0 poly-gcd-monic)
    have ww: ?w = ?x ^ ?p ^ ?d mod v unfolding w
        by simp (metis (mono-tags, opaque-lifting) One-nat-def mult.commute
power-Suc power-mod power-mult x-pow-n)
    have gv: ?g dvd v by auto
    hence gv': v div ?g dvd v
      by (metis dvd-def dvd-div-mult-self)
```

```
    {
      fix f
      assume irr: irreducible f and fv: f dvd v and degree f = ?d
      from degree-divisor(1)[OF this(1,3)]
      have f dvd ?x ^ ?p ^ ?d - ?x by auto
      hence f dvd (?x ^ ?p ^ ?d - ?x) mod v using fv by (rule dvd-mod)
       also have (?x ^ ?p ^ ?d - ?x) mod v = ?x ^ ?p ^ ?d mod v - ?x mod v
by (rule poly-mod-diff-left)
      also have ?x ^ ?p ^ ?d mod v = ?w mod v unfolding ww by auto
        also have ... - ?x mod v = (w ^ ?p mod v - ?x) mod v by (metis
poly-mod-diff-left)
      finally have f dvd (w^?p mod v - ?x) using fv by (rule dvd-mod-imp-dvd)
      with fv have f dvd ?g by auto
    } note deg-d-dvd-g = this
    show ?thesis
    proof (cases ?g = 1)
      case True
       with result have dist: dist-degree-factorize-main v ?w ?d res = facts by
auto
      show ?thesis
      proof (rule IH(1)[OF True dist ww u res])
        fix f
        assume irr: irreducible f and fv: f dvd v
        from fact[OF this] have d < degree f .
        moreover have degree f ≠ ?d
        proof
          assume degree f = ?d
          from divides-degree[OF deg-d-dvd-g[OF irr fv this]] mon-v
          have degree f ≤ degree ?g by auto
          with irr have degree ?g ≠ 0 unfolding irreducible_d-def by auto
          with True show False by auto
        qed
        ultimately show ?d < degree f by auto
      qed
    next
      case False
      with result
      have result: dist-degree-factorize-main ?v (?w mod ?v) ?d ((?d, ?g) # res)
= facts
        by auto
     from False mon-g have deg-g: degree ?g ≠ 0 by (simp add: monic-degree-0)
      have www: ?w mod ?v = monom 1 1 ^ ?p ^ ?d mod ?v using gv'
        by (simp add: mod-mod-cancel ww)
      from square-free-factor[OF - sf, of v] u have sfv: square-free v by auto
      have u: u = ?v * prod-list (map snd ((?d, ?g) # res))
        unfolding u by simp
      show ?thesis
      proof (rule IH(2)[OF False refl result www u], goal-cases)
        case (1 i f)
```

```
      show ?case
      proof (cases (i,f) ∈ set res)
        case True
        from res[OF this] show ?thesis by auto
      next
        case False
        with 1 have id: i = ?d f = ?g by auto
        show ?thesis unfolding id
        proof (intro conjI impI allI)
          fix g
          assume ∗: irreducible g g dvd ?g
          hence gv: g dvd v using dvd-trans[of g ?g v] by simp
          from fact[OF ∗(1) this] have dg: d < degree g .
          {
            assume degree g > ?d
            from degree-divisor(2)[OF ∗(1) refl - this]
            have ndvd: ¬ g dvd ?x ^ ?p ^ ?d − ?x by auto
            from ∗(2) have g dvd ?w − ?x by simp
            from this[unfolded ww]
            have g dvd ?x ^ ?p ^ ?d mod v − ?x .
              with gv have g dvd (?x ^ ?p ^ ?d mod v − ?x) mod v by (metis
dvd-mod)
            also have (?x ^ ?p ^ ?d mod v − ?x) mod v = (?x ^ ?p ^ ?d − ?x)
mod v
              by (metis mod-diff-left-eq)
                finally have g dvd ?x ^ ?p ^ ?d − ?x using gv by (rule
dvd-mod-imp-dvd)
            with ndvd have False by auto
          }
          with dg show degree g = ?d by presburger
        qed (insert mon-g deg-g, auto)
      qed
    next
      case (2 f)
      note irr = 2(1)
      from dvd-trans[OF 2(2) gv'] have fv: f dvd v .
      from fact[OF irr fv] have df: d < degree f degree f ≠ 0 by auto
      {
        assume degree f = ?d
        from deg-d-dvd-g[OF irr fv this] have fg: f dvd ?g .
        from gv have id: v = (v div ?g) ∗ ?g by simp
        from sfv id have square-free (v div ?g ∗ ?g) by simp
        from square-free-multD(1)[OF this 2(2) fg] have degree f = 0 .
        with df have False by auto
      }
      with df show ?d < degree f by presburger
    qed
  qed
qed
```

```
  qed
qed

definition distinct-degree-factorization
  :: 'a mod-ring poly ⇒ (nat × 'a mod-ring poly) list where
  distinct-degree-factorization f =
     (if degree f = 1 then [(1,f)] else dist-degree-factorize-main f (monom 1 1) 0
[])

lemma distinct-degree-factorization: assumes
  dist: distinct-degree-factorization f = facts and
  u: square-free f and
  mon: monic f
shows f = prod-list (map snd facts) ∧ (∀ i f. (i,f) ∈ set facts ⟶ factors-of-same-degree
i f)
proof −
  note dist = dist[unfolded distinct-degree-factorization-def]
  show ?thesis
  proof (cases degree f ≤ 1)
    case False
    hence degree f > 1 and dist: dist-degree-factorize-main f (monom 1 1) 0 [] =
facts
      using dist by auto
    hence ∗: monom 1 (Suc 0) = monom 1 (Suc 0) mod f
      by (simp add: degree-monom-eq mod-poly-less)
    show ?thesis
       by (rule dist-degree-factorize-main[OF dist - u mon], insert ∗, auto simp:
irreducible_d-def)
  next
    case True
    hence degree f = 0 ∨ degree f = 1 by auto
    thus ?thesis
    proof
      assume degree f = 0
      with mon have f: f = 1 using monic-degree-0 by blast
       hence facts = [] using dist unfolding dist-degree-factorize-main.simps[of -
- 0]
        by auto
      thus ?thesis using f by auto
    next
      assume deg: degree f = 1
      hence facts: facts = [(1,f)] using dist by auto
      show ?thesis unfolding facts factors-of-same-degree-def
      proof (intro conjI allI impI; clarsimp)
        fix g
        assume irreducible g g dvd f
          thus degree g = Suc 0 using deg divides-degree[of g f] by (auto simp:
irreducible_d-def)
      qed (insert mon deg, auto)
```

qed
qed
qed
end

end

8 A Combined Factorization Algorithm for Polynomials over GF(p)

8.1 Type Based Version

We combine Berlekamp's algorithm with the distinct degree factorization to obtain an efficient factorization algorithm for square-free polynomials in GF(p).

theory *Finite-Field-Factorization*
imports *Berlekamp-Type-Based*
Distinct-Degree-Factorization
begin

We prove soundness of the finite field factorization, indepedendent on whether distinct-degree-factorization is applied as preprocessing or not.

consts *use-distinct-degree-factorization* :: *bool*

context
assumes *SORT-CONSTRAINT*('*a*::*prime-card*)
begin

definition *finite-field-factorization* :: '*a mod-ring poly* ⇒ '*a mod-ring* × '*a mod-ring poly list* **where**
finite-field-factorization f = (*if degree f* = *0 then* (*lead-coeff f*,[]) *else let*
a = *lead-coeff f*;
u = *smult* (*inverse a*) *f*;
gs = (*if use-distinct-degree-factorization then distinct-degree-factorization u else* [(*1*,*u*)]);
(*irr*,*hs*) = *List.partition* (λ (*i*,*f*). *degree f* = *i*) *gs*
in (*a*,*map snd irr* @ *concat* (*map* (λ (*i*,*g*). *berlekamp-monic-factorization i g*) *hs*)))

lemma *finite-field-factorization-explicit*:
fixes *f*::'*a mod-ring poly*
assumes *sf-f*: *square-free f*
and *us*: *finite-field-factorization f* = (*c*,*us*)
shows *f* = *smult c* (*prod-list us*) ∧ (∀ *u* ∈ *set us*. *monic u* ∧ *irreducible u*)
proof (*cases degree f* = *0*)
case *False* **note** *f* = *this*
define *g* **where** *g* = *smult* (*inverse c*) *f*

obtain *gs* **where** *dist*: (*if use-distinct-degree-factorization then distinct-degree-factorization g else* [(*1*,*g*)]) = *gs* **by** *auto*
note *us* = *us*[*unfolded finite-field-factorization-def Let-def*]
from *us f* **have** *c*: *c* = *lead-coeff f* **by** *auto*
obtain *irr hs* **where** *part*: *List.partition* (λ (*i*, *f*). *degree f* = *i*) *gs* = (*irr*,*hs*) **by** *force*
from *arg-cong*[*OF this*, *of fst*] **have** *irr*: *irr* = *filter* (λ (*i*, *f*). *degree f* = *i*) *gs* **by** *auto*
from *us*[*folded c*, *folded g-def*, *unfolded dist part split*] *f*
have *us*: *us* = *map snd irr* @ *concat* (*map* (λ(*x*, *y*). *berlekamp-monic-factorization x y*) *hs*) **by** *auto*
from *f c* **have** *c0*: *c* ≠ *0* **by** *auto*
from *False c0* **have** *deg-g*: *degree g* ≠ *0* **unfolding** *g-def* **by** *auto*
have *mon-g*: *monic g* **unfolding** *g-def*
by (*metis c c0 field-class.field-inverse lead-coeff-smult*)
from *sf-f* **have** *sf-g*: *square-free g* **unfolding** *g-def* **by** (*simp add*: *c0*)
from *c0* **have** *f*: *f* = *smult c g* **unfolding** *g-def* **by** *auto*
have *g* = *prod-list* (*map snd gs*) ∧ (∀ (*i*,*f*) ∈ *set gs*. *degree f* > *0* ∧ *monic f* ∧ (∀ *h*. *h dvd f* ⟶ *degree h* = *i* ⟶ *irreducible h*))
proof (*cases use-distinct-degree-factorization*)
case *True*
with *dist* **have** *distinct-degree-factorization g* = *gs* **by** *auto*
note *dist* = *distinct-degree-factorization*[*OF this sf-g mon-g*]
from *dist* **have** *g*: *g* = *prod-list* (*map snd gs*) **by** *auto*
show *?thesis*
proof (*intro conjI*[*OF g*] *ballI*, *clarify*)
fix *i f*
assume (*i*,*f*) ∈ *set gs*
with *dist* **have** *factors-of-same-degree i f* **by** *auto*
from *factors-of-same-degreeD*[*OF this*]
show *degree f* > *0* ∧ *monic f* ∧ (∀ *h*. *h dvd f* ⟶ *degree h* = *i* ⟶ *irreducible h*) **by** *auto*
qed
next
case *False*
with *dist* **have** *gs*: *gs* = [(*1*,*g*)] **by** *auto*
show *?thesis* **unfolding** *gs* **using** *deg-g mon-g linear-irreducible*$_d$[**where** *'a* = *'a mod-ring*] **by** *auto*
qed
hence *g-gs*: *g* = *prod-list* (*map snd gs*)
and *mon-gs*: ⋀ *i f*. (*i*, *f*) ∈ *set gs* ⟹ *monic f* ∧ *degree f* > *0*
and *irrI*: ⋀ *i f h* . (*i*, *f*) ∈ *set gs* ⟹ *h dvd f* ⟹ *degree h* = *i* ⟹ *irreducible h* **by** *auto*
have *g*: *g* = *prod-list* (*map snd irr*) ∗ *prod-list* (*map snd hs*) **unfolding** *g-gs*
using *prod-list-map-partition*[*OF part*] **.**
{
fix *f*
assume *f* ∈ *snd* ‘ *set irr*
from *this*[*unfolded irr*] **obtain** *i* **where** ∗: (*i*,*f*) ∈ *set gs degree f* = *i* **by** *auto*

have *f dvd f* **by** *auto*
from *irrI*[*OF* ∗(*1*) *this* ∗(*2*)] *mon-gs*[*OF* ∗(*1*)] **have** *monic f irreducible f* **by** *auto*
} note *irr* = *this*
let *?berl* = λ *hs*. *concat* (*map* (λ(*x*, *y*). *berlekamp-monic-factorization x y*) *hs*)
have *set hs* ⊆ *set gs* **using** *part* **by** *auto*
hence *prod-list* (*map snd hs*) = *prod-list* (*?berl hs*)
∧ (∀ *f* ∈ *set* (*?berl hs*). *monic f* ∧ *irreducible*$_d$ *f*)
proof (*induct hs*)
case (*Cons ih hs*)
obtain *i h* **where** *ih*: *ih* = (*i*,*h*) **by** *force*
have *?berl* (*Cons ih hs*) = *berlekamp-monic-factorization i h* @ *?berl hs* **unfolding** *ih* **by** *auto*
from *Cons*(*2*)[*unfolded ih*] **have** *mem*: (*i*,*h*) ∈ *set gs* **and** *sub*: *set hs* ⊆ *set gs* **by** *auto*
note *IH* = *Cons*(*1*)[*OF sub*]
from *mem* **have** *h* ∈ *set* (*map snd gs*) **by** *force*
from *square-free-factor*[*OF prod-list-dvd*[*OF this*], *folded g-gs*, *OF sf-g*] **have** *sf*: *square-free h* **.**
from *mon-gs*[*OF mem*] *irrI*[*OF mem*] **have** ∗: *degree h* > *0 monic h*
⋀ *g*. *g dvd h* ⟹ *degree g* = *i* ⟹ *irreducible g* **by** *auto*
from *berlekamp-monic-factorization*[*OF sf refl* ∗(*3*) ∗(*1*−*2*), *of i*]
have *berl*: *prod-list* (*berlekamp-monic-factorization i h*) = *h*
and *irr*: ⋀ *f*. *f* ∈ *set* (*berlekamp-monic-factorization i h*) ⟹ *monic f* ∧ *irreducible f* **by** *auto*
have *prod-list* (*map snd* (*Cons ih hs*)) = *h* ∗ *prod-list* (*map snd hs*) **unfolding** *ih* **by** *simp*
also have *prod-list* (*map snd hs*) = *prod-list* (*?berl hs*) **using** *IH* **by** *auto*
finally have *prod-list* (*map snd* (*Cons ih hs*)) = *prod-list* (*?berl* (*Cons ih hs*))
unfolding *ih* **using** *berl* **by** *auto*
thus *?case* **using** *IH irr* **unfolding** *ih* **by** *auto*
qed *auto*
with *g irr* **have** *main*: *g* = *prod-list us* ∧ (∀ *u* ∈ *set us*. *monic u* ∧ *irreducible*$_d$ *u*) **unfolding** *us*
by *auto*
thus *?thesis* **unfolding** *f* **using** *sf-g* **by** *auto*
next
case *True*
with *us*[*unfolded finite-field-factorization-def*] **have** *c* = *lead-coeff f* **and** *us*: *us* = [] **by** *auto*
with *degree0-coeffs*[*OF True*] **have** *f*: *f* = [:*c*:] **by** *auto*
show *?thesis* **unfolding** *us f* **by** (*auto simp*: *normalize-poly-def*)
qed

lemma *finite-field-factorization*:
fixes *f*::*'a mod-ring poly*
assumes *sf-f*: *square-free f*
and *us*: *finite-field-factorization f* = (*c*,*us*)
shows *unique-factorization Irr-Mon f* (*c*, *mset us*)

proof −
from *finite-field-factorization-explicit*[*OF sf-f us*]
have *fact*: *factorization Irr-Mon f* (*c*, *mset us*)
unfolding *factorization-def split Irr-Mon-def* **by** (*auto simp*: *prod-mset-prod-list*)
from *sf-f*[*unfolded square-free-def*] **have** $f \neq 0$ **by** *auto*
from *exactly-one-factorization*[*OF this*] *fact*
show *?thesis* **unfolding** *unique-factorization-def* **by** *auto*
qed
end

Experiments revealed that preprocessing via distinct-degree-factorization slows down the factorization algorithm (statement for implementation in AFP 2017)

overloading *use-distinct-degree-factorization* ≡ *use-distinct-degree-factorization*
begin
definition *use-distinct-degree-factorization*
where [*code-unfold*]: *use-distinct-degree-factorization* = *False*
end
end

8.2 Record Based Version

theory *Finite-Field-Factorization-Record-Based*
imports
Finite-Field-Factorization
Matrix-Record-Based
Poly-Mod-Finite-Field-Record-Based
HOL−Types-To-Sets.Types-To-Sets
Jordan-Normal-Form.Matrix-IArray-Impl
Jordan-Normal-Form.Gauss-Jordan-IArray-Impl
Polynomial-Interpolation.Improved-Code-Equations
Polynomial-Factorization.Missing-List
begin

hide-const(**open**) *monom coeff*

Whereas ⟦*square-free ?f*; *finite-field-factorization ?f* = (*?c*, *?us*)⟧ ⟹ *unique-factorization Irr-Mon ?f* (*?c*, *mset ?us*) provides a result for a polynomials over GF(p), we now develop a theorem which speaks about integer polynomials modulo p.

lemma (**in** *poly-mod-prime-type*) *finite-field-factorization-modulo-ring*:
assumes *g*: (*g* :: *'a mod-ring poly*) = *of-int-poly f*
and *sf*: *square-free-m f*
and *fact*: *finite-field-factorization g* = (*d*,*gs*)
and *c*: *c* = *to-int-mod-ring d*
and *fs*: *fs* = *map to-int-poly gs*
shows *unique-factorization-m f* (*c*, *mset fs*)
proof −
have [*transfer-rule*]: *MP-Rel f g* **unfolding** *g MP-Rel-def* **by** (*simp add*: *Mp-f-representative*)

have *sg*: *square-free g* **by** (*transfer*, *rule sf*)
have [*transfer-rule*]: *M-Rel c d* **unfolding** *M-Rel-def c* **by** (*rule M-to-int-mod-ring*)
have *fs-gs*[*transfer-rule*]: *list-all2 MP-Rel fs gs*
unfolding *fs list-all2-map1 MP-Rel-def*[*abs-def*] *Mp-to-int-poly* **by** (*simp add*: *list.rel-refl*)
have [*transfer-rule*]: *rel-mset MP-Rel* (*mset fs*) (*mset gs*)
using *fs-gs* **using** *rel-mset-def* **by** *blast*
have [*transfer-rule*]: *MF-Rel* (*c*,*mset fs*) (*d*,*mset gs*) **unfolding** *MF-Rel-def* **by** *transfer-prover*
from *finite-field-factorization*[*OF sg fact*]
have *uf*: *unique-factorization Irr-Mon g* (*d*,*mset gs*) **by** *auto*
from *uf*[*untransferred*] **show** *unique-factorization-m f* (*c*, *mset fs*) **.**
qed

We now have to implement *finite-field-factorization*.

context
fixes *p* :: *int*
and *ff-ops* :: *'i arith-ops-record*
begin

fun *power-poly-f-mod-i* :: (*'i list* ⇒ *'i list*) ⇒ *'i list* ⇒ *nat* ⇒ *'i list* **where**
power-poly-f-mod-i modulus a n = (*if n* = *0 then modulus* (*one-poly-i ff-ops*)
else let (*d*,*r*) = *Euclidean-Rings.divmod-nat n 2*;
rec = *power-poly-f-mod-i modulus* (*modulus* (*times-poly-i ff-ops a a*)) *d in*
if r = *0 then rec else modulus* (*times-poly-i ff-ops rec a*))

declare *power-poly-f-mod-i.simps*[*simp del*]

fun *power-polys-i* :: *'i list* ⇒ *'i list* ⇒ *'i list* ⇒ *nat* ⇒ *'i list list* **where**
power-polys-i mul-p u curr-p (*Suc i*) = *curr-p* #
power-polys-i mul-p u (*mod-field-poly-i ff-ops* (*times-poly-i ff-ops curr-p mul-p*) *u*) *i*
| *power-polys-i mul-p u curr-p 0* = []

lemma *length-power-polys-i*[*simp*]: *length* (*power-polys-i x y z n*) = *n*
by (*induct n arbitrary*: *x y z*, *auto*)

definition *berlekamp-mat-i* :: *'i list* ⇒ *'i mat* **where**
berlekamp-mat-i u = (*let n* = *degree-i u*;
ze = *arith-ops-record.zero ff-ops*; *on* = *arith-ops-record.one ff-ops*;
mul-p = *power-poly-f-mod-i* (λ *v*. *mod-field-poly-i ff-ops v u*)
[*ze*, *on*] (*nat p*);
xks = *power-polys-i mul-p u* [*on*] *n*
in mat-of-rows-list n (*map* (λ *cs*. *cs* @ *replicate* (*n* − *length cs*) *ze*) *xks*))

definition *berlekamp-resulting-mat-i* :: *'i list* ⇒ *'i mat* **where**
berlekamp-resulting-mat-i u = (*let Q* = *berlekamp-mat-i u*;
n = *dim-row Q*;
QI = *mat n n* (λ (*i*,*j*). *if i* = *j then arith-ops-record.minus ff-ops* (*Q* $$ (*i*,*j*))

```
(arith-ops-record.one ff-ops) else Q $$ (i,j))
    in (gauss-jordan-single-i ff-ops (transpose-mat QI)))

definition berlekamp-basis-i :: 'i list ⇒ 'i list list where
  berlekamp-basis-i u = (map (poly-of-list-i ff-ops o list-of-vec)
    (find-base-vectors-i ff-ops (berlekamp-resulting-mat-i u)))

primrec berlekamp-factorization-main-i :: 'i ⇒ 'i ⇒ nat ⇒ 'i list list ⇒ 'i list list
⇒ nat ⇒ 'i list list where
  berlekamp-factorization-main-i ze on d divs (v # vs) n = (
    if v = [on] then berlekamp-factorization-main-i ze on d divs vs n else
    if length divs = n then divs else
    let of-int = arith-ops-record.of-int ff-ops;
        facts = filter (λ w. w ≠ [on])
          [ gcd-poly-i ff-ops u (minus-poly-i ff-ops v (if s = 0 then [] else [of-int (int
s)])) .
            u ← divs, s ← [0 ..< nat p]];
      (lin,nonlin) = List.partition (λ q. degree-i q = d) facts
      in lin @ berlekamp-factorization-main-i ze on d nonlin vs (n − length lin))
| berlekamp-factorization-main-i ze on d divs [] n = divs

definition berlekamp-monic-factorization-i :: nat ⇒ 'i list ⇒ 'i list list where
  berlekamp-monic-factorization-i d f = (let
     vs = berlekamp-basis-i f
   in berlekamp-factorization-main-i (arith-ops-record.zero ff-ops) (arith-ops-record.one
ff-ops) d [f] vs (length vs))

partial-function (tailrec) dist-degree-factorize-main-i ::
  'i ⇒ 'i ⇒ nat ⇒ 'i list ⇒ 'i list ⇒ nat ⇒ (nat × 'i list) list
  ⇒ (nat × 'i list) list where
  [code]: dist-degree-factorize-main-i ze on dv v w d res = (if v = [on] then res else
if d + d > dv
    then (dv, v) # res else let
      w = power-poly-f-mod-i (λ f. mod-field-poly-i ff-ops f v) w (nat p);
      d = Suc d;
      gd = gcd-poly-i ff-ops (minus-poly-i ff-ops w [ze,on]) v
      in if gd = [on] then dist-degree-factorize-main-i ze on dv v w d res else
      let v' = div-field-poly-i ff-ops v gd
      in dist-degree-factorize-main-i ze on (degree-i v') v' (mod-field-poly-i ff-ops w
v') d ((d,gd) # res))

definition distinct-degree-factorization-i
  :: 'i list ⇒ (nat × 'i list) list where
  distinct-degree-factorization-i f = (let ze = arith-ops-record.zero ff-ops;
     on = arith-ops-record.one ff-ops in if degree-i f = 1 then [(1,f)] else
     dist-degree-factorize-main-i ze on (degree-i f) f [ze,on] 0 [])

definition finite-field-factorization-i :: 'i list ⇒ 'i × 'i list list where
  finite-field-factorization-i f = (if degree-i f = 0 then (lead-coeff-i ff-ops f,[]) else
```

```
let
    a = lead-coeff-i ff-ops f;
    u = smult-i ff-ops (arith-ops-record.inverse ff-ops a) f;
    gs = (if use-distinct-degree-factorization then distinct-degree-factorization-i u
else [(1,u)]);
    (irr,hs) = List.partition (λ (i,f). degree-i f = i) gs
    in (a,map snd irr @ concat (map (λ (i,g). berlekamp-monic-factorization-i i g)
hs)))
end

context prime-field-gen
begin

lemma power-polys-i: assumes i: i < n and [transfer-rule]: poly-rel f f' poly-rel
g g'
  and h: poly-rel h h'
  shows poly-rel (power-polys-i ff-ops g f h n ! i) (power-polys g' f' h' n ! i)
  using i h
proof (induct n arbitrary: h h' i)
  case (Suc n h h' i) note * = this
  note [transfer-rule] = *(3)
  show ?case
  proof (cases i)
    case 0
    with Suc show ?thesis by auto
  next
    case (Suc j)
    with *(2−) have j < n by auto
    note IH = *(1)[OF this]
    show ?thesis unfolding Suc by (simp, rule IH, transfer-prover)
  qed
qed simp

lemma power-poly-f-mod-i: assumes m: (poly-rel ===> poly-rel) m (λ x'. x' mod
m')
 shows poly-rel f f' ⟹ poly-rel (power-poly-f-mod-i ff-ops m f n) (power-poly-f-mod
m' f' n)
proof −
  from m have m: ⋀ x x'. poly-rel x x' ⟹ poly-rel (m x) (x' mod m')
    unfolding rel-fun-def by auto
 show poly-rel f f' ⟹ poly-rel (power-poly-f-mod-i ff-ops m f n) (power-poly-f-mod
m' f' n)
  proof (induct n arbitrary: f f' rule: less-induct)
    case (less n f f')
    note f[transfer-rule] = less(2)
    show ?case
    proof (cases n = 0)
      case True
      show ?thesis
```

```
      by (simp add: True power-poly-f-mod-i.simps power-poly-f-mod-binary,
        rule m[OF poly-rel-one])
    next
      case False
      hence n: (n = 0) = False by simp
      obtain q r where div: Euclidean-Rings.divmod-nat n 2 = (q,r) by force
      from this[unfolded Euclidean-Rings.divmod-nat-def] n have q < n by auto
      note IH = less(1)[OF this]
      have rec: poly-rel (power-poly-f-mod-i ff-ops m (m (times-poly-i ff-ops f f)) q)

        (power-poly-f-mod m' (f' * f' mod m') q)
        by (rule IH, rule m, transfer-prover)
      have other: poly-rel
        (m (times-poly-i ff-ops (power-poly-f-mod-i ff-ops m (m (times-poly-i ff-ops
f f)) q) f))
        (power-poly-f-mod m' (f' * f' mod m') q * f' mod m')
        by (rule m, rule poly-rel-times[unfolded rel-fun-def, rule-format, OF rec f])
      show ?thesis unfolding power-poly-f-mod-i.simps[of - - - n] Let-def
        power-poly-f-mod-binary[of - - n] div split n if-False using rec other by auto
    qed
  qed
qed

lemma berlekamp-mat-i[transfer-rule]: (poly-rel ===> mat-rel R)
  (berlekamp-mat-i p ff-ops) berlekamp-mat
proof (intro rel-funI)
  fix f f'
  let ?ze = arith-ops-record.zero ff-ops
  let ?on = arith-ops-record.one ff-ops
  assume f[transfer-rule]: poly-rel f f'
  have deg: degree-i f = degree f' by transfer-prover
  {
    fix i j
    assume i: i < degree f' and j: j < degree f'
    define cs where cs = (λcs :: 'i list. cs @ replicate (degree f' − length cs) ?ze)
    define cs' where cs' = (λcs :: 'a mod-ring poly. coeffs cs @ replicate (degree
f' − length (coeffs cs)) 0)
    define poly where poly = power-polys-i ff-ops
        (power-poly-f-mod-i ff-ops (λv. mod-field-poly-i ff-ops v f) [?ze, ?on] (nat
p)) f [?on]
        (degree f')
    define poly' where poly' = (power-polys (power-poly-f-mod f' [:0, 1:] (nat p))
f' 1 (degree f'))
    have *: poly-rel (power-poly-f-mod-i ff-ops (λv. mod-field-poly-i ff-ops v f) [?ze,
?on] (nat p))
      (power-poly-f-mod f' [:0, 1:] (nat p))
      by (rule power-poly-f-mod-i, transfer-prover, simp add: poly-rel-def one zero)
    have [transfer-rule]: poly-rel (poly ! i) (poly' ! i)
      unfolding poly-def poly'-def
```

```
    by (rule power-polys-i[OF i f *], simp add: poly-rel-def one)
  have *: list-all2 R (cs (poly ! i)) (cs' (poly' ! i))
    unfolding cs-def cs'-def by transfer-prover
  from list-all2-nthD[OF *[unfolded poly-rel-def], of j] j
  have R (cs (poly ! i) ! j) (cs' (poly' ! i) ! j) unfolding cs-def by auto
  hence R
        (mat-of-rows-list (degree f')
          (map (λcs. cs @ replicate (degree f' − length cs) ?ze)
            (power-polys-i ff-ops
              (power-poly-f-mod-i ff-ops (λv. mod-field-poly-i ff-ops v f) [?ze, ?on]
(nat p)) f [?on]
                (degree f'))) $$
          (i, j))
        (mat-of-rows-list (degree f')
          (map (λcs. coeffs cs @ replicate (degree f' − length (coeffs cs)) 0)
            (power-polys (power-poly-f-mod f' [:0, 1:] (nat p)) f' 1 (degree f'))) $$
          (i, j))
    unfolding mat-of-rows-list-def length-map length-power-polys-i power-polys-works
      length-power-polys index-mat[OF i j] split
      unfolding poly-def cs-def poly'-def cs'-def using i
      by auto
 } note main = this
 show mat-rel R (berlekamp-mat-i p ff-ops f) (berlekamp-mat f')
  unfolding berlekamp-mat-i-def berlekamp-mat-def Let-def nat-p[symmetric] deg
  unfolding mat-rel-def
  by (intro conjI allI impI, insert main, auto)
qed

lemma berlekamp-resulting-mat-i[transfer-rule]: (poly-rel ===> mat-rel R)
 (berlekamp-resulting-mat-i p ff-ops) berlekamp-resulting-mat
proof (intro rel-funI)
 fix f f'
 assume poly-rel f f'
 from berlekamp-mat-i[unfolded rel-fun-def, rule-format, OF this]
 have bmi: mat-rel R (berlekamp-mat-i p ff-ops f) (berlekamp-mat f') .
 show mat-rel R (berlekamp-resulting-mat-i p ff-ops f) (berlekamp-resulting-mat
f')
  unfolding berlekamp-resulting-mat-def Let-def berlekamp-resulting-mat-i-def
  by (rule gauss-jordan-i[unfolded rel-fun-def, rule-format],
   insert bmi, auto simp: mat-rel-def one intro!: minus[unfolded rel-fun-def, rule-format])
qed

lemma berlekamp-basis-i[transfer-rule]: (poly-rel ===> list-all2 poly-rel)
 (berlekamp-basis-i p ff-ops) berlekamp-basis
 unfolding berlekamp-basis-i-def[abs-def] berlekamp-basis-code[abs-def] o-def
 by transfer-prover

lemma berlekamp-factorization-main-i[transfer-rule]:
 ((=) ===> list-all2 poly-rel ===> list-all2 poly-rel ===> (=) ===> list-all2
```

```
poly-rel)
    (berlekamp-factorization-main-i p ff-ops (arith-ops-record.zero ff-ops)
      (arith-ops-record.one ff-ops))
    berlekamp-factorization-main
proof (intro rel-funI, clarify, goal-cases)
  case (1 - d xs xs' ys ys' - n)
  let ?ze = arith-ops-record.zero ff-ops
  let ?on = arith-ops-record.one ff-ops
  let ?of-int = arith-ops-record.of-int ff-ops
  from 1(2) 1(1) show ?case
  proof (induct ys ys' arbitrary: xs xs' n rule: list-all2-induct)
    case (Cons y ys y' ys' xs xs' n)
    note trans[transfer-rule] = Cons(1,2,4)
    obtain clar0 clar1 clar2 where clarify: ⋀ s u. gcd-poly-i ff-ops u
                      (minus-poly-i ff-ops y
                     (if s = 0 then [] else [?of-int (int s)])) = clar0 s u
        [0..<nat p] = clar1
        [?on] = clar2 by auto
    define facts where facts = concat (map (λu. concat
                     (map (λs. if gcd-poly-i ff-ops u
                                  (minus-poly-i ff-ops y (if s = 0 then [] else [?of-int
(int s)])) ≠
                                [?on]
                               then [gcd-poly-i ff-ops u
                                  (minus-poly-i ff-ops y (if s = 0 then [] else [?of-int
(int s)]))]
                               else [])
                       [0..<nat p])) xs)
    define Facts where Facts = [w←concat
                     (map (λu. map (λs. gcd-poly-i ff-ops u
                                    (minus-poly-i ff-ops y
                                      (if s = 0 then [] else [?of-int (int s)])))
                                  [0..<nat p])
                       xs) .  w ≠ [?on]]
    have Facts: Facts = facts
      unfolding Facts-def facts-def clarify
    proof (induct xs)
      case (Cons x xs)
      show ?case by (simp add: Cons, induct clar1, auto)
    qed simp
    define facts' where facts' = concat
          (map (λu. concat
                    (map (λx. if gcd u (y' − [:of-nat x:]) ≠ 1
                              then [gcd u (y' − [:of-int (int x):])] else [])
                      [0..<nat p]))
            xs')
    have id: ⋀ x. of-int (int x) = of-nat x [?on] = one-poly-i ff-ops
      by (auto simp: one-poly-i-def)
    have facts[transfer-rule]: list-all2 poly-rel facts facts'
```

```
    unfolding facts-def facts'-def
  apply (rule concat-transfer[unfolded rel-fun-def, rule-format])
  apply (rule list.map-transfer[unfolded rel-fun-def, rule-format, OF - trans(3)])
  apply (rule concat-transfer[unfolded rel-fun-def, rule-format])
  apply (rule list-all2-map-map)
  proof (unfold id)
    fix f f' x
    assume [transfer-rule]: poly-rel f f' and x: x ∈ set [0..<nat p]
    hence *: 0 ≤ int x int x < p by auto
   from of-int[OF this] have rel[transfer-rule]: R (?of-int (int x)) (of-nat x) by
auto
    {
      assume 0 < x
      with * have *: 0 < int x int x < p by auto
      have (of-nat x :: 'a mod-ring) = of-int (int x) by simp
      also have ... ≠ 0 unfolding of-int-of-int-mod-ring using * unfolding p
        by (transfer', auto)
    }
    with rel have [transfer-rule]: poly-rel (if x = 0 then [] else [?of-int (int x)])
[:of-nat x:]
      unfolding poly-rel-def by (auto simp add: cCons-def p)
    show list-all2 poly-rel
        (if gcd-poly-i ff-ops f (minus-poly-i ff-ops y (if x = 0 then [] else [?of-int
(int x)])) ≠ one-poly-i ff-ops
        then [gcd-poly-i ff-ops f (minus-poly-i ff-ops y (if x = 0 then [] else [?of-int
(int x)]))]
         else [])
        (if gcd f' (y' − [:of-nat x:]) ≠ 1 then [gcd f' (y' − [:of-nat x:])] else [])
      by transfer-prover
  qed
  have id1: berlekamp-factorization-main-i p ff-ops ?ze ?on d xs (y # ys) n = (
   if y = [?on] then berlekamp-factorization-main-i p ff-ops ?ze ?on d xs ys n else
    if length xs = n then xs else
    (let fac = facts;
        (lin, nonlin) = List.partition (λq. degree-i q = d) fac
          in lin @ berlekamp-factorization-main-i p ff-ops ?ze ?on d nonlin ys (n
− length lin)))
    unfolding berlekamp-factorization-main-i.simps Facts[symmetric]
    by (simp add: o-def Facts-def Let-def)
  have id2: berlekamp-factorization-main d xs' (y' # ys') n = (
    if y' = 1 then berlekamp-factorization-main d xs' ys' n
    else if length xs' = n then xs' else
    (let fac = facts';
        (lin, nonlin) = List.partition (λq. degree q = d) fac
          in lin @ berlekamp-factorization-main d nonlin ys' (n − length lin)))
    by (simp add: o-def facts'-def nat-p)
  have len: length xs = length xs' by transfer-prover
  have id3: (y = [?on]) = (y' = 1)
  by (transfer-prover-start, transfer-step+, simp add: one-poly-i-def finite-field-ops-int-def)
```

```
    show ?case
    proof (cases y' = 1)
      case True
      hence id4: (y' = 1) = True by simp
      show ?thesis unfolding id1 id2 id3 id4 if-True
        by (rule Cons(3), transfer-prover)
    next
      case False
      hence id4: (y' = 1) = False by simp
      note id1 = id1[unfolded id3 id4 if-False]
      note id2 = id2[unfolded id4 if-False]
      show ?thesis
      proof (cases length xs' = n)
        case True
        thus ?thesis unfolding id1 id2 Let-def len using trans by simp
      next
        case False
        hence id: (length xs' = n) = False by simp
        have id': length [q←facts . degree-i q = d] = length [q←facts'. degree q =
d]
          by transfer-prover
       have [transfer-rule]: list-all2 poly-rel (berlekamp-factorization-main-i p ff-ops
?ze ?on d [x←facts . degree-i x ≠ d] ys
         (n − length [q←facts . degree-i q = d]))
         (berlekamp-factorization-main d [x←facts' . degree x ≠ d] ys'
         (n − length [q←facts' . degree q = d]))
          unfolding id'
          by (rule Cons(3), transfer-prover)
        show ?thesis unfolding id1 id2 Let-def len id if-False
          unfolding partition-filter-conv o-def split by transfer-prover
      qed
    qed
  qed simp
qed

lemma berlekamp-monic-factorization-i[transfer-rule]:
  ((=) ===> poly-rel ===> list-all2 poly-rel)
     (berlekamp-monic-factorization-i p ff-ops) berlekamp-monic-factorization
  unfolding berlekamp-monic-factorization-i-def[abs-def] berlekamp-monic-factorization-def[abs-def]
Let-def
  by transfer-prover

lemma dist-degree-factorize-main-i:
  poly-rel F f ⟹ poly-rel G g ⟹ list-all2 (rel-prod (=) poly-rel) Res res
   ⟹ list-all2 (rel-prod (=) poly-rel)
      (dist-degree-factorize-main-i p ff-ops
         (arith-ops-record.zero ff-ops) (arith-ops-record.one ff-ops) (degree-i F) F G
d Res)
      (dist-degree-factorize-main f g d res)
```

proof (*induct f g d res arbitrary*: *F G Res rule*: *dist-degree-factorize-main.induct*)
case (*1 v w d res V W Res*)
let *?ze* = *arith-ops-record.zero ff-ops*
let *?on* = *arith-ops-record.one ff-ops*
note *simp* = *dist-degree-factorize-main.simps*[*of v w d*]
dist-degree-factorize-main-i.simps[*of p ff-ops ?ze ?on degree-i V V W d*]
have *v*[*transfer-rule*]: *poly-rel V v* **by** (*rule 1*)
have *w*[*transfer-rule*]: *poly-rel W w* **by** (*rule 1*)
have *res*[*transfer-rule*]: *list-all2* (*rel-prod* (=) *poly-rel*) *Res res* **by** (*rule 1*)
have [*transfer-rule*]: *poly-rel* [*?on*] *1*
by (*simp add*: *one poly-rel-def*)
have *id1*: (*V* = [*?on*]) = (*v* = *1*) **unfolding** *finite-field-ops-int-def* **by** *transfer-prover*
have *id2*: *degree-i V* = *degree v* **by** *transfer-prover*
note *simp* = *simp*[*unfolded id1 id2*]
note *IH* = *1*(*1*,*2*)
show *?case*
proof (*cases v* = *1*)
case *True*
with *res* **show** *?thesis* **unfolding** *id2 simp* **by** *simp*
next
case *False*
with *id1* **have** (*v* = *1*) = *False* **by** *auto*
note *simp* = *simp*[*unfolded this if-False*]
note *IH* = *IH*[*OF False*]
show *?thesis*
proof (*cases degree v* < *d* + *d*)
case *True*
thus *?thesis* **unfolding** *id2 simp* **using** *res v* **by** *auto*
next
case *False*
hence (*degree v* < *d* + *d*) = *False* **by** *auto*
note *simp* = *simp*[*unfolded this if-False*]
let *?P* = *power-poly-f-mod-i ff-ops* (λ*f*. *mod-field-poly-i ff-ops f V*) *W* (*nat p*)
let *?G* = *gcd-poly-i ff-ops* (*minus-poly-i ff-ops ?P* [*?ze*, *?on*]) *V*
let *?g* = *gcd* (*w* ^ *CARD*('*a*) *mod v* − *monom 1 1*) *v*
define *G* **where** *G* = *?G*
define *g* **where** *g* = *?g*
note *simp* = *simp*[*unfolded Let-def*, *folded G-def g-def*]
note *IH* = *IH*[*OF False refl refl refl*]
have [*transfer-rule*]: *poly-rel* [*?ze*,*?on*] (*monom 1 1*) **unfolding** *poly-rel-def*
by (*auto simp*: *coeffs-monom one zero*)
have *id*: *w* ^ *CARD*('*a*) *mod v* = *power-poly-f-mod v w* (*nat p*)
unfolding *power-poly-f-mod-def* **by** (*simp add*: *p*)
have *P*[*transfer-rule*]: *poly-rel ?P* (*w* ^ *CARD*('*a*) *mod v*) **unfolding** *id*
by (*rule power-poly-f-mod-i*[*OF - w*], *transfer-prover*)
have *g*[*transfer-rule*]: *poly-rel G g* **unfolding** *G-def g-def* **by** *transfer-prover*
have *id3*: (*G* = [*?on*]) = (*g* = *1*) **by** *transfer-prover*

```
      note simp = simp[unfolded id3]
      show ?thesis
      proof (cases g = 1)
        case True
        from IH(1)[OF this[unfolded g-def] v P res] True
        show ?thesis unfolding id2 simp by simp
      next
        case False
        have vg: poly-rel (div-field-poly-i ff-ops V G) (v div g) by transfer-prover
        have poly-rel (mod-field-poly-i ff-ops ?P
            (div-field-poly-i ff-ops V G)) (w ^ CARD('a) mod v mod (v div g)) by
transfer-prover
          note IH = IH(2)[OF False[unfolded g-def] refl vg[unfolded G-def g-def]
this[unfolded G-def g-def],
            folded g-def G-def]
        have list-all2 (rel-prod (=) poly-rel) ((Suc d, G) # Res) ((Suc d, g) # res)
          using g res by auto
        note IH = IH[OF this]
        from False have (g = 1) = False by simp
        note simp = simp[unfolded this if-False]
        show ?thesis unfolding id2 simp using IH by simp
      qed
    qed
  qed
qed

lemma distinct-degree-factorization-i[transfer-rule]: (poly-rel ===> list-all2 (rel-prod
(=) poly-rel))
  (distinct-degree-factorization-i p ff-ops) distinct-degree-factorization
proof
  fix F f
  assume f[transfer-rule]: poly-rel F f
  have id: (degree-i F = 1) = (degree f = 1) by transfer-prover
  note d = distinct-degree-factorization-i-def distinct-degree-factorization-def
  let ?ze = arith-ops-record.zero ff-ops
  let ?on = arith-ops-record.one ff-ops
  show list-all2 (rel-prod (=) poly-rel) (distinct-degree-factorization-i p ff-ops F)
          (distinct-degree-factorization f)
  proof (cases degree f = 1)
    case True
    with id f show ?thesis unfolding d by auto
  next
    case False
    from False id have ?thesis = (list-all2 (rel-prod (=) poly-rel)
      (dist-degree-factorize-main-i p ff-ops ?ze ?on (degree-i F) F [?ze, ?on] 0 [])
     (dist-degree-factorize-main f (monom 1 1) 0 [])) unfolding d Let-def by simp

    also have ...
      by (rule dist-degree-factorize-main-i[OF f], auto simp: poly-rel-def
```

```
      coeffs-monom one zero)
    finally show ?thesis .
  qed
qed

lemma finite-field-factorization-i[transfer-rule]:
  (poly-rel ===> rel-prod R (list-all2 poly-rel))
    (finite-field-factorization-i p ff-ops) finite-field-factorization
  unfolding finite-field-factorization-i-def finite-field-factorization-def Let-def lead-coeff-i-def'
  by transfer-prover
```

Since the implementation is sound, we can now combine it with the soundness result of the finite field factorization.

```
lemma finite-field-i-sound:
  assumes f': f' = of-int-poly-i ff-ops (Mp f)
  and berl-i: finite-field-factorization-i p ff-ops f' = (c',fs')
  and sq: square-free-m f
  and fs: fs = map (to-int-poly-i ff-ops) fs'
  and c: c = arith-ops-record.to-int ff-ops c'
  shows unique-factorization-m f (c, mset fs)
    ∧ c ∈ {0 ..< p}
    ∧ (∀ fi ∈ set fs. set (coeffs fi) ⊆ {0 ..< p})
proof −
  define f'' :: 'a mod-ring poly where f'' = of-int-poly (Mp f)
  have rel-f[transfer-rule]: poly-rel f' f''
    by (rule poly-rel-of-int-poly[OF f'], simp add: f''-def)
  interpret pff: idom-ops poly-ops ff-ops poly-rel
    by (rule idom-ops-poly)
  obtain c'' fs'' where berl: finite-field-factorization f'' = (c'',fs'') by force
  from rel-funD[OF finite-field-factorization-i rel-f, unfolded rel-prod-conv assms(2)
split berl]
  have rel[transfer-rule]: R c' c'' list-all2 poly-rel fs' fs'' by auto
  from to-int[OF rel(1)] have cc': c = to-int-mod-ring c'' unfolding c by simp
  from m1 have ‹M c ∈ {0 ..< p}›
    by (simp add: M-def cc')
  then have c: ‹c ∈ {0 ..< p}›
    by (simp add: M-to-int-mod-ring cc')
  {
    fix f
    assume f ∈ set fs'
     with rel(2) obtain f' where poly-rel f f'  unfolding list-all2-conv-all-nth
set-conv-nth
      by auto
    hence is-poly ff-ops f using fun-cong[OF Domainp-is-poly, of f]
      unfolding Domainp-iff[abs-def] by auto
  }
  hence fs': Ball (set fs') (is-poly ff-ops) by auto
  define mon :: 'a mod-ring poly ⇒ bool where mon = monic
  have [transfer-rule]: (poly-rel ===> (=)) (monic-i ff-ops) mon unfolding mon-def
```

```
    by (rule poly-rel-monic)
  have len: length fs' = length fs'' by transfer-prover
  have fs': fs = map to-int-poly fs'' unfolding fs
  proof (rule nth-map-conv[OF len], intro allI impI)
    fix i
    assume i: i < length fs'
    obtain f g where id: fs' ! i = f fs'' ! i = g by auto
    from i rel(2)[unfolded list-all2-conv-all-nth[of - fs' fs'']] id
    have poly-rel f g by auto
    from to-int-poly-i[OF this] have to-int-poly-i ff-ops f = to-int-poly g .
    thus to-int-poly-i ff-ops (fs' ! i) = to-int-poly (fs'' ! i) unfolding id .
  qed
  have f: f'' = of-int-poly f unfolding poly-eq-iff f''-def
   by (simp add: to-int-mod-ring-hom.injectivity to-int-mod-ring-of-int-M Mp-coeff)
  have *: unique-factorization-m f (c, mset fs)
    using finite-field-factorization-modulo-ring[OF f sq berl cc' fs'] by auto
  have fs': (∀ fi∈set fs. set (coeffs fi) ⊆ {0..<p}) unfolding fs'
    using range-to-int-mod-ring[where 'a = 'a]
    by (auto simp: coeffs-to-int-poly p)
  with c fs *
  show ?thesis by blast
qed
end

definition finite-field-factorization-main :: int ⇒ 'i arith-ops-record ⇒ int poly ⇒
int × int poly list where
  finite-field-factorization-main p ff-ops f ≡
    let (c',fs') = finite-field-factorization-i p ff-ops (of-int-poly-i ff-ops (poly-mod.Mp
p f))
      in (arith-ops-record.to-int ff-ops c', map (to-int-poly-i ff-ops) fs')

lemma(in prime-field-gen) finite-field-factorization-main:
  assumes res: finite-field-factorization-main p ff-ops f = (c,fs)
  and sq: square-free-m f
  shows unique-factorization-m f (c, mset fs)
    ∧ c ∈ {0 ..< p}
    ∧ (∀ fi ∈ set fs. set (coeffs fi) ⊆ {0 ..< p})
proof −
  obtain c' fs' where
    res': finite-field-factorization-i p ff-ops (of-int-poly-i ff-ops (Mp f)) = (c', fs')
by force
  show ?thesis
    by (rule finite-field-i-sound[OF refl res' sq],
      insert res[unfolded finite-field-factorization-main-def res'], auto)
qed

definition finite-field-factorization-int :: int ⇒ int poly ⇒ int × int poly list
where
```

```
  finite-field-factorization-int p = (
    if p ≤ 65535
    then finite-field-factorization-main p (finite-field-ops32 (uint32-of-int p))
    else if p ≤ 4294967295
    then finite-field-factorization-main p (finite-field-ops64 (uint64-of-int p))
    else finite-field-factorization-main p (finite-field-ops-integer (integer-of-int p)))

context poly-mod-prime begin
lemmas finite-field-factorization-main-integer = prime-field-gen.finite-field-factorization-main
 [OF prime-field.prime-field-finite-field-ops-integer, unfolded prime-field-def mod-ring-locale-def,
   unfolded poly-mod-type-simps, internalize-sort 'a :: prime-card, OF type-to-set,
unfolded remove-duplicate-premise, cancel-type-definition, OF non-empty]

lemmas finite-field-factorization-main-uint32 = prime-field-gen.finite-field-factorization-main
 [OF prime-field.prime-field-finite-field-ops32, unfolded prime-field-def mod-ring-locale-def,
   unfolded poly-mod-type-simps, internalize-sort 'a :: prime-card, OF type-to-set,
unfolded remove-duplicate-premise, cancel-type-definition, OF non-empty]

lemmas finite-field-factorization-main-uint64 = prime-field-gen.finite-field-factorization-main
 [OF prime-field.prime-field-finite-field-ops64, unfolded prime-field-def mod-ring-locale-def,
   unfolded poly-mod-type-simps, internalize-sort 'a :: prime-card, OF type-to-set,
unfolded remove-duplicate-premise, cancel-type-definition, OF non-empty]

lemma finite-field-factorization-int:
  assumes sq: poly-mod.square-free-m p f
  and result: finite-field-factorization-int p f = (c,fs)
  shows poly-mod.unique-factorization-m p f (c, mset fs)
    ∧ c ∈ {0 ..< p}
    ∧ (∀ fi ∈ set fs. set (coeffs fi) ⊆ {0 ..< p})
  using finite-field-factorization-main-integer[OF  - sq, of c fs]
    finite-field-factorization-main-uint32[OF - - sq, of c fs]
    finite-field-factorization-main-uint64[OF - - sq, of c fs]
    result[unfolded finite-field-factorization-int-def]
  by (auto split: if-splits)

end

end
```

9 Hensel Lifting

9.1 Properties about Factors

We define and prove properties of Hensel-lifting. Here, we show the result that Hensel-lifting can lift a factorization mod p to a factorization mod p^n. For the lifting we have proofs for both versions, the original linear Hensel-lifting or the quadratic approach from Zassenhaus. Via the linear version, we also show a uniqueness result, however only in the binary case,

i.e., where $f = g \cdot h$. Uniqueness of the general case will later be shown in theory Berlekamp-Hensel by incorporating the factorization algorithm for finite fields algorithm.

theory *Hensel-Lifting*
imports
HOL−Computational-Algebra.Euclidean-Algorithm
Poly-Mod-Finite-Field-Record-Based
Polynomial-Factorization.Square-Free-Factorization
begin

lemma *uniqueness-poly-equality*:
fixes *f g* :: *'a* :: {*factorial-ring-gcd*,*semiring-gcd-mult-normalize*} *poly*
assumes *cop*: *coprime f g*
and *deg*: $B = 0 \vee$ *degree* $B <$ *degree f* $B' = 0 \vee$ *degree* $B' <$ *degree f*
and *f*: $f \neq 0$ **and** *eq*: $A * f + B * g = A' * f + B' * g$
shows $A = A'$ $B = B'$
proof −
from *eq* **have** ∗: $(A - A') * f = (B' - B) * g$ **by** (*simp add*: *field-simps*)
hence *f dvd* $(B' - B) * g$ **unfolding** *dvd-def* **by** (*intro exI*[*of* - $A - A'$], *auto simp*: *field-simps*)
with *cop*[*simplified*] **have** *dvd*: *f dvd* $(B' - B)$
by (*simp add*: *coprime-dvd-mult-right-iff ac-simps*)
from *divides-degree*[*OF this*] **have** *degree f* ≤ *degree* $(B' - B) \vee B = B'$ **by** *auto*
with *degree-diff-le-max*[*of B' B*] *deg*
show $B = B'$ **by** *auto*
with ∗ *f* **show** $A = A'$ **by** *auto*
qed

lemmas (**in** *poly-mod-prime-type*) *uniqueness-poly-equality* =
uniqueness-poly-equality[**where** *'a*=*'a mod-ring*, *untransferred*]
lemmas (**in** *poly-mod-prime*) *uniqueness-poly-equality* = *poly-mod-prime-type.uniqueness-poly-equality*
[*unfolded poly-mod-type-simps*, *internalize-sort 'a* :: *prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*, *OF non-empty*]

lemma *pseudo-divmod-main-list-1-is-divmod-poly-one-main-list*:
pseudo-divmod-main-list (*1* :: *'a* :: *comm-ring-1*) *q f g n* = *divmod-poly-one-main-list q f g n*
by (*induct n arbitrary*: *q f g*, *auto simp*: *Let-def*)

lemma *pdivmod-monic-pseudo-divmod*: **assumes** *g*: *monic g* **shows** *pdivmod-monic f g* = *pseudo-divmod f g*
proof −
from *g* **have** *id*: (*coeffs g* = []) = *False* **by** *auto*
from *g* **have** *mon*: *hd* (*rev* (*coeffs g*)) = *1* **by** (*metis coeffs-eq-Nil hd-rev id last-coeffs-eq-coeff-degree*)
show *?thesis*
unfolding *pseudo-divmod-impl pseudo-divmod-list-def id if-False pdivmod-monic-def Let-def mon*

pseudo-divmod-main-list-1-is-divmod-poly-one-main-list **by** (*auto split*: *prod.splits*)
qed

lemma *pdivmod-monic*: **assumes** *g*: *monic g* **and** *res*: *pdivmod-monic f g* = (*q*, *r*)
shows *f* = *g* * *q* + *r r* = *0* ∨ *degree r* < *degree g*
proof −
from *g* **have** *g0*: *g* ≠ *0* **by** *auto*
from *pseudo-divmod*[*OF g0 res*[*unfolded pdivmod-monic-pseudo-divmod*[*OF g*]], *unfolded g*]
show *f* = *g* * *q* + *r r* = *0* ∨ *degree r* < *degree g* **by** *auto*
qed

definition *dupe-monic* :: *'a* :: *comm-ring-1 poly* ⇒ *'a poly* ⇒ *'a poly* ⇒ *'a poly* ⇒ *'a poly* ⇒
'a poly * *'a poly* **where**
dupe-monic D H S T U = (*case pdivmod-monic* (*T* * *U*) *D of* (*q*,*r*) ⇒
(*S* * *U* + *H* * *q*, *r*))

lemma *dupe-monic*: **assumes** *1*: *D***S* + *H***T* = *1*
and *mon*: *monic D*
and *dupe*: *dupe-monic D H S T U* = (*A*,*B*)
shows *A* * *D* + *B* * *H* = *U B* = *0* ∨ *degree B* < *degree D*
proof −
obtain *Q R* **where** *div*: *pdivmod-monic* ((*T* * *U*)) *D* = (*Q*,*R*) **by** *force*
from *dupe*[*unfolded dupe-monic-def div split*]
have *A*: *A* = (*S* * *U* + *H* * *Q*) **and** *B*: *B* = *R* **by** *auto*
from *pdivmod-monic*[*OF mon div*] **have** *TU*: *T* * *U* = *D* * *Q* + *R* **and**
deg: *R* = *0* ∨ *degree R* < *degree D* **by** *auto*
hence *R*: *R* = *T* * *U* − *D* * *Q* **by** *simp*
have *A* * *D* + *B* * *H* = (*D* * *S* + *H* * *T*) * *U* **unfolding** *A B R* **by** (*simp add*: *field-simps*)
also have ... = *U* **unfolding** *1* **by** *simp*
finally show *eq*: *A* * *D* + *B* * *H* = *U* **.**
show *B* = *0* ∨ *degree B* < *degree D* **using** *deg* **unfolding** *B* **.**
qed

lemma *dupe-monic-unique*: **fixes** *D* :: *'a* :: {*factorial-ring-gcd*,*semiring-gcd-mult-normalize*} *poly*
assumes *1*: *D***S* + *H***T* = *1*
and *mon*: *monic D*
and *dupe*: *dupe-monic D H S T U* = (*A*,*B*)
and *cop*: *coprime D H*
and *other*: *A'* * *D* + *B'* * *H* = *U B'* = *0* ∨ *degree B'* < *degree D*
shows *A'* = *A B'* = *B*
proof −
from *dupe-monic*[*OF 1 mon dupe*] **have** *one*: *A* * *D* + *B* * *H* = *U B* = *0* ∨ *degree B* < *degree D* **by** *auto*
from *mon* **have** *D0*: *D* ≠ *0* **by** *auto*
from *uniqueness-poly-equality*[*OF cop one*(*2*) *other*(*2*) *D0*, *of A A'*, *unfolded*

other, *OF one*(*1*)]
show $A' = A$ $B' = B$ **by** *auto*
qed

context *ring-ops*
begin
lemma *poly-rel-dupe-monic-i*: **assumes** *mon*: *monic D*
and *rel*: *poly-rel d D poly-rel h H poly-rel s S poly-rel t T poly-rel u U*
shows *rel-prod poly-rel poly-rel* (*dupe-monic-i ops d h s t u*) (*dupe-monic D H S T U*)
proof −
note *defs* = *dupe-monic-i-def dupe-monic-def*
note [*transfer-rule*] = *rel*
have [*transfer-rule*]: *rel-prod poly-rel poly-rel*
(*pdivmod-monic-i ops* (*times-poly-i ops t u*) *d*)
(*pdivmod-monic* (*T* ∗ *U*) *D*)
by (*rule poly-rel-pdivmod-monic*[*OF mon*], *transfer-prover*+)
show *?thesis* **unfolding** *defs* **by** *transfer-prover*
qed
end

context *mod-ring-gen*
begin

lemma *monic-of-int-poly*: *monic D* ⟹ *monic* (*of-int-poly* (*Mp D*) :: *'a mod-ring poly*)
using *Mp-f-representative Mp-to-int-poly monic-Mp* **by** *auto*

lemma *dupe-monic-i*: **assumes** *dupe-i*: *dupe-monic-i ff-ops d h s t u* = (*a*,*b*)
and *1*: *D*∗*S* + *H*∗*T* =*m* *1*
and *mon*: *monic D*
and *A*: *A* = *to-int-poly-i ff-ops a*
and *B*: *B* = *to-int-poly-i ff-ops b*
and *d*: *Mp-rel-i d D*
and *h*: *Mp-rel-i h H*
and *s*: *Mp-rel-i s S*
and *t*: *Mp-rel-i t T*
and *u*: *Mp-rel-i u U*
shows
A ∗ *D* + *B* ∗ *H* =*m* *U*
B = *0* ∨ *degree B* < *degree D*
Mp-rel-i a A
Mp-rel-i b B
proof −
let *?I* = λ *f*. *of-int-poly* (*Mp f*) :: *'a mod-ring poly*
let *?i* = *to-int-poly-i ff-ops*
note *dd* = *Mp-rel-iD*[*OF d*]
note *hh* = *Mp-rel-iD*[*OF h*]
note *ss* = *Mp-rel-iD*[*OF s*]

note *tt* = *Mp-rel-iD*[*OF t*]
note *uu* = *Mp-rel-iD*[*OF u*]
obtain *A′ B′* **where** *dupe*: *dupe-monic* (*?I D*) (*?I H*) (*?I S*) (*?I T*) (*?I U*) = (*A′*,*B′*) **by** *force*
from *poly-rel-dupe-monic-i*[*OF monic-of-int-poly*[*OF mon*] *dd*(*1*) *hh*(*1*) *ss*(*1*) *tt*(*1*) *uu*(*1*), *unfolded dupe-i dupe*]
have *a*: *poly-rel a A′* **and** *b*: *poly-rel b B′* **by** *auto*
show *aa*: *Mp-rel-i a A* **by** (*rule Mp-rel-iI′*[*OF a*, *folded A*])
show *bb*: *Mp-rel-i b B* **by** (*rule Mp-rel-iI′*[*OF b*, *folded B*])
note *Aa* = *Mp-rel-iD*[*OF aa*]
note *Bb* = *Mp-rel-iD*[*OF bb*]
from *poly-rel-inj*[*OF a Aa*(*1*)] *A* **have** *A*: *A′* = *?I A* **by** *simp*
from *poly-rel-inj*[*OF b Bb*(*1*)] *B* **have** *B*: *B′* = *?I B* **by** *simp*
note *Mp* = *dd*(*2*) *hh*(*2*) *ss*(*2*) *tt*(*2*) *uu*(*2*)
note [*transfer-rule*] = *Mp*
have (=) (*D* ∗ *S* + *H* ∗ *T* =*m* *1*) (*?I D* ∗ *?I S* + *?I H* ∗ *?I T* = *1*) **by** *transfer-prover*
with *1* **have** *11*: *?I D* ∗ *?I S* + *?I H* ∗ *?I T* = *1* **by** *simp*
from *dupe-monic*[*OF 11 monic-of-int-poly*[*OF mon*] *dupe*, *unfolded A B*]
have *res*: *?I A* ∗ *?I D* + *?I B* ∗ *?I H* = *?I U ?I B* = *0* ∨ *degree* (*?I B*) < *degree* (*?I D*) **by** *auto*
note [*transfer-rule*] = *Aa*(*2*) *Bb*(*2*)
have (=) (*A* ∗ *D* + *B* ∗ *H* =*m* *U*) (*?I A* ∗ *?I D* + *?I B* ∗ *?I H* = *?I U*)
(=) (*B* =*m* *0* ∨ *degree-m B* < *degree-m D*) (*?I B* = *0* ∨ *degree* (*?I B*) < *degree* (*?I D*)) **by** *transfer-prover+*
with *res* **have** ∗: *A* ∗ *D* + *B* ∗ *H* =*m* *U B* =*m* *0* ∨ *degree-m B* < *degree-m D* **by** *auto*
show *A* ∗ *D* + *B* ∗ *H* =*m* *U* **by** *fact*
have *B*: *Mp B* = *B* **using** *Mp-rel-i-Mp-to-int-poly-i assms*(*5*) *bb* **by** *blast*
from ∗(*2*) **show** *B* = *0* ∨ *degree B* < *degree D* **unfolding** *B* **using** *degree-m-le*[*of D*] **by** *auto*
qed

lemma *Mp-rel-i-of-int-poly-i*: **assumes** *Mp F* = *F*
shows *Mp-rel-i* (*of-int-poly-i ff-ops F*) *F*
by (*metis Mp-f-representative Mp-rel-iI′ assms poly-rel-of-int-poly to-int-poly-i*)

lemma *dupe-monic-i-int*: **assumes** *dupe-i*: *dupe-monic-i-int ff-ops D H S T U* = (*A*,*B*)
and *1*: *D*∗*S* + *H*∗*T* =*m* *1*
and *mon*: *monic D*
and *norm*: *Mp D* = *D Mp H* = *H Mp S* = *S Mp T* = *T Mp U* = *U*
shows
A ∗ *D* + *B* ∗ *H* =*m* *U*
B = *0* ∨ *degree B* < *degree D*
Mp A = *A*
Mp B = *B*
proof −
let *?oi* = *of-int-poly-i ff-ops*

```
  let ?ti = to-int-poly-i ff-ops
  note rel = norm[THEN Mp-rel-i-of-int-poly-i]
  obtain a b where dupe: dupe-monic-i ff-ops (?oi D) (?oi H) (?oi S) (?oi T)
(?oi U) = (a,b) by force
  from dupe-i[unfolded dupe-monic-i-int-def this Let-def] have AB: A = ?ti a B
= ?ti b by auto
  from dupe-monic-i[OF dupe 1 mon AB rel] Mp-rel-i-Mp-to-int-poly-i
  show A * D + B * H =m U
    B = 0 ∨ degree B < degree D
    Mp A = A
    Mp B = B
    unfolding AB by auto
qed

end

definition dupe-monic-dynamic
  :: int ⇒ int poly ⇒ int poly ⇒ int poly ⇒ int poly ⇒ int poly ⇒ int poly × int
poly where
  dupe-monic-dynamic p = (
    if p ≤ 65535
    then dupe-monic-i-int (finite-field-ops32 (uint32-of-int p))
    else if p ≤ 4294967295
    then dupe-monic-i-int (finite-field-ops64 (uint64-of-int p))
    else dupe-monic-i-int (finite-field-ops-integer (integer-of-int p)))

context poly-mod-2
begin

lemma dupe-monic-i-int-finite-field-ops-integer: assumes
      dupe-i: dupe-monic-i-int (finite-field-ops-integer (integer-of-int m)) D H S T
U = (A,B)
  and 1: D*S + H*T =m 1
  and mon: monic D
  and norm: Mp D = D Mp H = H Mp S = S Mp T = T Mp U = U
shows
  A * D + B * H =m U
  B = 0 ∨ degree B < degree D
  Mp A = A
  Mp B = B
  using m1 mod-ring-gen.dupe-monic-i-int[OF
      mod-ring-locale.mod-ring-finite-field-ops-integer[unfolded mod-ring-locale-def],

      internalize-sort 'a :: nontriv, OF type-to-set, unfolded remove-duplicate-premise,

        cancel-type-definition, OF - assms] by auto

lemma dupe-monic-i-int-finite-field-ops32: assumes
      m: m ≤ 65535
```

and *dupe-i*: *dupe-monic-i-int* (*finite-field-ops32* (*uint32-of-int m*)) *D H S T U* = (*A*,*B*)
and *1*: *D***S* + *H***T* =*m* *1*
and *mon*: *monic D*
and *norm*: *Mp D* = *D Mp H* = *H Mp S* = *S Mp T* = *T Mp U* = *U*
shows
A * *D* + *B* * *H* =*m* *U*
B = *0* ∨ *degree B* < *degree D*
Mp A = *A*
Mp B = *B*
using *m1 mod-ring-gen.dupe-monic-i-int*[*OF*
mod-ring-locale.mod-ring-finite-field-ops32[*unfolded mod-ring-locale-def*],
internalize-sort *'a* :: *nontriv*, *OF type-to-set*, *unfolded remove-duplicate-premise*,
cancel-type-definition, *OF - assms*] **by** *auto*

lemma *dupe-monic-i-int-finite-field-ops64*: **assumes**
m: *m* ≤ *4294967295*
and *dupe-i*: *dupe-monic-i-int* (*finite-field-ops64* (*uint64-of-int m*)) *D H S T U* = (*A*,*B*)
and *1*: *D***S* + *H***T* =*m* *1*
and *mon*: *monic D*
and *norm*: *Mp D* = *D Mp H* = *H Mp S* = *S Mp T* = *T Mp U* = *U*
shows
A * *D* + *B* * *H* =*m* *U*
B = *0* ∨ *degree B* < *degree D*
Mp A = *A*
Mp B = *B*
using *m1 mod-ring-gen.dupe-monic-i-int*[*OF*
mod-ring-locale.mod-ring-finite-field-ops64[*unfolded mod-ring-locale-def*],
internalize-sort *'a* :: *nontriv*, *OF type-to-set*, *unfolded remove-duplicate-premise*,
cancel-type-definition, *OF - assms*] **by** *auto*

lemma *dupe-monic-dynamic*: **assumes** *dupe*: *dupe-monic-dynamic m D H S T U* = (*A*,*B*)
and *1*: *D***S* + *H***T* =*m* *1*
and *mon*: *monic D*
and *norm*: *Mp D* = *D Mp H* = *H Mp S* = *S Mp T* = *T Mp U* = *U*
shows
A * *D* + *B* * *H* =*m* *U*
B = *0* ∨ *degree B* < *degree D*
Mp A = *A*
Mp B = *B*
using *dupe*
dupe-monic-i-int-finite-field-ops32[*OF - - 1 mon norm*, *of A B*]
dupe-monic-i-int-finite-field-ops64[*OF - - 1 mon norm*, *of A B*]
dupe-monic-i-int-finite-field-ops-integer[*OF - 1 mon norm*, *of A B*]
unfolding *dupe-monic-dynamic-def* **by** (*auto split*: *if-splits*)

end

context *poly-mod*
begin

definition *dupe-monic-int* :: *int poly* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* ⇒
int poly ∗ *int poly* **where**
dupe-monic-int D H S T U = (*case pdivmod-monic* (*Mp* (*T* ∗ *U*)) *D of* (*q*,*r*) ⇒
(*Mp* (*S* ∗ *U* + *H* ∗ *q*), *Mp r*))

end

declare *poly-mod.dupe-monic-int-def*[*code*]

Old direct proof on int poly. It does not permit to change implementation. This proof is still present, since we did not export the uniqueness part from the type-based uniqueness result ⟦*?D* ∗ *?S* + *?H* ∗ *?T* = *1*; *monic ?D*; *dupe-monic ?D ?H ?S ?T ?U* = (*?A*, *?B*); *comm-monoid-mult-class.coprime ?D ?H*; *?A′* ∗ *?D* + *?B′* ∗ *?H* = *?U*; *?B′* = *0* ∨ *degree ?B′* < *degree ?D*⟧ ⟹ *?A′* = *?A*

⟦*?D* ∗ *?S* + *?H* ∗ *?T* = *1*; *monic ?D*; *dupe-monic ?D ?H ?S ?T ?U* = (*?A*, *?B*); *comm-monoid-mult-class.coprime ?D ?H*; *?A′* ∗ *?D* + *?B′* ∗ *?H* = *?U*; *?B′* = *0* ∨ *degree ?B′* < *degree ?D*⟧ ⟹ *?B′* = *?B* via the various relations.

lemma (**in** *poly-mod-2*) *dupe-monic-int*: **assumes** *1*: *D*∗*S* + *H*∗*T* =*m* *1*
and *mon*: *monic D*
and *dupe*: *dupe-monic-int D H S T U* = (*A*,*B*)
shows *A* ∗ *D* + *B* ∗ *H* =*m* *U B* = *0* ∨ *degree B* < *degree D Mp A* = *A Mp B* = *B*
coprime-m D H ⟹ *A′* ∗ *D* + *B′* ∗ *H* =*m* *U* ⟹ *B′* = *0* ∨ *degree B′* < *degree D* ⟹ *Mp D* = *D*
⟹ *Mp A′* = *A′* ⟹ *Mp B′* = *B′* ⟹ *prime m*
⟹ *A′* = *A* ∧ *B′* = *B*
proof −
obtain *Q R* **where** *div*: *pdivmod-monic* (*Mp* (*T* ∗ *U*)) *D* = (*Q*,*R*) **by** *force*
from *dupe*[*unfolded dupe-monic-int-def div split*]
have *A*: *A* = *Mp* (*S* ∗ *U* + *H* ∗ *Q*) **and** *B*: *B* = *Mp R* **by** *auto*
from *pdivmod-monic*[*OF mon div*] **have** *TU*: *Mp* (*T* ∗ *U*) = *D* ∗ *Q* + *R* **and**
deg: *R* = *0* ∨ *degree R* < *degree D* **by** *auto*
hence *Mp R* = *Mp* (*Mp* (*T* ∗ *U*) − *D* ∗ *Q*) **by** *simp*
also have ... = *Mp* (*T* ∗ *U* − *Mp* (*Mp* (*Mp D* ∗ *Q*))) **unfolding** *Mp-Mp*
unfolding *minus-Mp*
using *minus-Mp mult-Mp* **by** *metis*
also have ... = *Mp* (*T* ∗ *U* − *D* ∗ *Q*) **by** *simp*
finally have *r*: *Mp R* = *Mp* (*T* ∗ *U* − *D* ∗ *Q*) **by** *simp*
have *Mp* (*A* ∗ *D* + *B* ∗ *H*) = *Mp* (*Mp* (*A* ∗ *D*) + *Mp* (*B* ∗ *H*)) **by** *simp*

also have $Mp\ (A * D) = Mp\ ((S * U + H * Q) * D)$ **unfolding** A **by** *simp*
also have $Mp\ (B * H) = Mp\ (Mp\ R * Mp\ H)$ **unfolding** B **by** *simp*
also have $\ldots = Mp\ ((T * U - D * Q) * H)$ **unfolding** r **by** *simp*
also have $Mp\ (Mp\ ((S * U + H * Q) * D) + Mp\ ((T * U - D * Q) * H)) =$
$Mp\ ((S * U + H * Q) * D + (T * U - D * Q) * H)$ **by** *simp*
also have $(S * U + H * Q) * D + (T * U - D * Q) * H = (D * S + H * T) * U$
by (*simp add*: *field-simps*)
also have $Mp \ldots = Mp\ (Mp\ (D * S + H * T) * U)$ **by** *simp*
also have $Mp\ (D * S + H * T) = 1$ **using** *1* **by** *simp*
finally show *eq*: $A * D + B * H =_m U$ **by** *simp*
have *id*: *degree-m* $(Mp\ R) =$ *degree-m* R **by** *simp*
have *id'*: *degree* $D =$ *degree-m* D **using** *mon* **by** *simp*
show *degB*: $B = 0 \vee$ *degree* $B <$ *degree* D **using** *deg* **unfolding** B *id id'*
using *degree-m-le*[*of* R] **by** (*cases* $R = 0$, *auto*)
show *Mp*: $Mp\ A = A\ Mp\ B = B$ **unfolding** $A\ B$ **by** *auto*
assume *another*: $A' * D + B' * H =_m U$ **and** *degB'*: $B' = 0 \vee$ *degree* $B' <$ *degree* D
and *norm*: $Mp\ A' = A'\ Mp\ B' = B'$ **and** *cop*: *coprime-m* $D\ H$ **and** *D*: $Mp\ D = D$
and *prime*: *prime* m
from *degB Mp D* **have** *degB*: $B =_m 0 \vee$ *degree-m* $B <$ *degree-m* D **by** *auto*
from *degB' Mp D norm* **have** *degB'*: $B' =_m 0 \vee$ *degree-m* $B' <$ *degree-m* D **by** *auto*
from *mon D* **have** *D0*: $\neg\ (D =_m 0)$ **by** *auto*
from *prime* **interpret** *poly-mod-prime* m **by** *unfold-locales*
from *another eq* **have** $A' * D + B' * H =_m A * D + B * H$ **by** *simp*
from *uniqueness-poly-equality*[*OF cop degB' degB D0 this*]
show $A' = A \wedge B' = B$ **unfolding** *norm Mp* **by** *auto*
qed

lemma *coprime-bezout-coefficients*:
assumes *cop*: *coprime* $f\ g$
and *ext*: *bezout-coefficients* $f\ g = (a,\ b)$
shows $a * f + b * g = 1$
using *assms bezout-coefficients* [*of* $f\ g\ a\ b$]
by *simp*

lemma (**in** *poly-mod-prime-type*) *bezout-coefficients-mod-int*: **assumes** f: $(F :: 'a\ mod\text{-}ring\ poly) =$ *of-int-poly* f
and g: $(G :: 'a\ mod\text{-}ring\ poly) =$ *of-int-poly* g
and *cop*: *coprime-m* $f\ g$
and *fact*: *bezout-coefficients* $F\ G = (A,B)$
and a: $a =$ *to-int-poly* A
and b: $b =$ *to-int-poly* B
shows $f * a + g * b =_m 1$
proof −
have f[*transfer-rule*]: *MP-Rel f F* **unfolding** *f MP-Rel-def* **by** (*simp add*: *Mp-f-representative*)

have *g*[*transfer-rule*]: *MP-Rel g G* **unfolding** *g MP-Rel-def* **by** (*simp add*: *Mp-f-representative*)
have [*transfer-rule*]: *MP-Rel a A* **unfolding** *a MP-Rel-def* **by** (*rule Mp-to-int-poly*)
have [*transfer-rule*]: *MP-Rel b B* **unfolding** *b MP-Rel-def* **by** (*rule Mp-to-int-poly*)
from *cop* **have** *coprime F G* **using** *coprime-MP-Rel*[*unfolded rel-fun-def*] *f g* **by** *auto*
from *coprime-bezout-coefficients* [*OF this fact*]
have $A * F + B * G = 1$ **.**
from *this* [*untransferred*]
show *?thesis* **by** (*simp add*: *ac-simps*)
qed

definition *bezout-coefficients-i* :: *'i arith-ops-record* ⇒ *'i list* ⇒ *'i list* ⇒ *'i list* × *'i list* **where**
bezout-coefficients-i ff-ops f g = *fst* (*euclid-ext-poly-i ff-ops f g*)

definition *euclid-ext-poly-mod-main* :: *int* ⇒ *'a arith-ops-record* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* × *int poly* **where**
euclid-ext-poly-mod-main p ff-ops f g = (*case bezout-coefficients-i ff-ops* (*of-int-poly-i ff-ops f*) (*of-int-poly-i ff-ops g*) *of*
(*a*,*b*) ⇒ (*to-int-poly-i ff-ops a*, *to-int-poly-i ff-ops b*))

definition *euclid-ext-poly-dynamic* :: *int* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* × *int poly* **where**
euclid-ext-poly-dynamic p = (
if p ≤ *65535*
then euclid-ext-poly-mod-main p (*finite-field-ops32* (*uint32-of-int p*))
else if p ≤ *4294967295*
then euclid-ext-poly-mod-main p (*finite-field-ops64* (*uint64-of-int p*))
else euclid-ext-poly-mod-main p (*finite-field-ops-integer* (*integer-of-int p*)))

context *prime-field-gen*
begin
lemma *bezout-coefficients-i*[*transfer-rule*]:
(*poly-rel* ===> *poly-rel* ===> *rel-prod poly-rel poly-rel*)
(*bezout-coefficients-i ff-ops*) *bezout-coefficients*
unfolding *bezout-coefficients-i-def bezout-coefficients-def*
by *transfer-prover*

lemma *bezout-coefficients-i-sound*: **assumes** *f*: *f'* = *of-int-poly-i ff-ops f Mp f* = *f*
and *g*: *g'* = *of-int-poly-i ff-ops g Mp g* = *g*
and *cop*: *coprime-m f g*
and *res*: *bezout-coefficients-i ff-ops f' g'* = (*a'*,*b'*)
and *a*: *a* = *to-int-poly-i ff-ops a'*
and *b*: *b* = *to-int-poly-i ff-ops b'*
shows *f* ∗ *a* + *g* ∗ *b* =*m* *1*
Mp a = *a Mp b* = *b*
proof −
from *f* **have** *f'*: *f'* = *of-int-poly-i ff-ops* (*Mp f*) **by** *simp*

define f'' **where** $f'' \equiv$ *of-int-poly* (*Mp f*) :: *'a mod-ring poly*
have f'': $f'' =$ *of-int-poly f* **unfolding** f''*-def f* **by** *simp*
have *rel-f*[*transfer-rule*]: *poly-rel f' f''*
by (*rule poly-rel-of-int-poly*[*OF f'*], *simp add*: f'' *f*)
from *g* **have** g': $g' =$ *of-int-poly-i ff-ops* (*Mp g*) **by** *simp*
define g'' **where** $g'' \equiv$ *of-int-poly* (*Mp g*) :: *'a mod-ring poly*
have g'': $g'' =$ *of-int-poly g* **unfolding** g''*-def g* **by** *simp*
have *rel-g*[*transfer-rule*]: *poly-rel g' g''*
by (*rule poly-rel-of-int-poly*[*OF g'*], *simp add*: g'' *g*)
obtain a'' b'' **where** *eucl*: *bezout-coefficients f'' g''* = (a'',b'') **by** *force*
from *bezout-coefficients-i*[*unfolded rel-fun-def rel-prod-conv*, *rule-format*, *OF rel-f rel-g*,
unfolded res split eucl]
have *rel*[*transfer-rule*]: *poly-rel a' a'' poly-rel b' b''* **by** *auto*
with *to-int-poly-i* **have** *a*: *a* = *to-int-poly a''*
and *b*: *b* = *to-int-poly b''* **unfolding** *a b* **by** *auto*
from *bezout-coefficients-mod-int* [*OF f'' g'' cop eucl a b*]
show $f * a + g * b =_m 1$ **.**
show *Mp a* = *a Mp b* = *b* **unfolding** *a b* **by** (*auto simp*: *Mp-to-int-poly*)
qed

lemma *euclid-ext-poly-mod-main*: **assumes** *cop*: *coprime-m f g*
and *f*: *Mp f* = *f* **and** *g*: *Mp g* = *g*
and *res*: *euclid-ext-poly-mod-main m ff-ops f g* = (*a*,*b*)
shows $f * a + g * b =_m 1$
Mp a = *a Mp b* = *b*
proof −
obtain a' b' **where** res': *bezout-coefficients-i ff-ops* (*of-int-poly-i ff-ops f*)
(*of-int-poly-i ff-ops g*) = (a', b') **by** *force*
show $f * a + g * b =_m 1$
Mp a = *a Mp b* = *b*
by (*insert bezout-coefficients-i-sound*[*OF refl f refl g cop res'*]
res [*unfolded euclid-ext-poly-mod-main-def res'*], *auto*)
qed

end

context *poly-mod-prime* **begin**

lemmas *euclid-ext-poly-mod-integer* = *prime-field-gen.euclid-ext-poly-mod-main*
[*OF prime-field.prime-field-finite-field-ops-integer*,
unfolded prime-field-def mod-ring-locale-def poly-mod-type-simps, *internalize-sort*
'a :: *prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*,
OF non-empty]

lemmas *euclid-ext-poly-mod-uint32* = *prime-field-gen.euclid-ext-poly-mod-main*
[*OF prime-field.prime-field-finite-field-ops32*,
unfolded prime-field-def mod-ring-locale-def poly-mod-type-simps, *internalize-sort*
'a :: *prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*,

OF non-empty]

lemmas *euclid-ext-poly-mod-uint64* = *prime-field-gen.euclid-ext-poly-mod-main*[*OF prime-field.prime-field-finite-field-ops64*,
unfolded prime-field-def mod-ring-locale-def poly-mod-type-simps, *internalize-sort* 'a :: *prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*, *OF non-empty*]

lemma *euclid-ext-poly-dynamic*:
assumes *cop*: *coprime-m f g* **and** *f*: *Mp f* = *f* **and** *g*: *Mp g* = *g*
and *res*: *euclid-ext-poly-dynamic p f g* = (*a*,*b*)
shows *f* ∗ *a* + *g* ∗ *b* =*m* *1*
Mp a = *a Mp b* = *b*
using *euclid-ext-poly-mod-integer*[*OF cop f g*, *of p a b*]
euclid-ext-poly-mod-uint32[*OF - cop f g*, *of p a b*]
euclid-ext-poly-mod-uint64[*OF - cop f g*, *of p a b*]
res[*unfolded euclid-ext-poly-dynamic-def*] **by** (*auto split*: *if-splits*)

end

lemma *range-sum-prod*: **assumes** *xy*: *x* ∈ {*0*..<*q*} (*y* :: *int*) ∈ {*0*..<*p*}
shows *x* + *q* ∗ *y* ∈ {*0*..<*p* ∗ *q*}
proof −
{
fix *x q* :: *int*
have *x* ∈ {*0* ..< *q*} ⟷ *0* ≤ *x* ∧ *x* < *q* **by** *auto*
} **note** *id* = *this*
from *xy* **have** *0*: *0* ≤ *x* + *q* ∗ *y* **by** *auto*
have *x* + *q* ∗ *y* ≤ *q* − *1* + *q* ∗ *y* **using** *xy* **by** *simp*
also have *q* ∗ *y* ≤ *q* ∗ (*p* − *1*) **using** *xy* **by** *auto*
finally have *x* + *q* ∗ *y* ≤ *q* − *1* + *q* ∗ (*p* − *1*) **by** *auto*
also have ... = *p* ∗ *q* − *1* **by** (*simp add*: *field-simps*)
finally show *?thesis* **using** *0* **by** *auto*
qed

context
fixes *C* :: *int poly*
begin

context
fixes *p* :: *int* **and** *S T D1 H1* :: *int poly*
begin

fun *linear-hensel-main* **where**
linear-hensel-main (*Suc 0*) = (*D1*,*H1*)
| *linear-hensel-main* (*Suc n*) = (
let (*D*,*H*) = *linear-hensel-main n*;
q = *p* ^ *n*;
U = *poly-mod.Mp p* (*sdiv-poly* (*C* − *D* ∗ *H*) *q*); — *H2* + *H3*

```
      (A,B) = poly-mod.dupe-monic-int p D1 H1 S T U
    in (D + smult q B, H + smult q A)) — H4
  | linear-hensel-main 0 = (D1,H1)

lemma linear-hensel-main: assumes 1: poly-mod.eq-m p (D1 * S + H1 * T) 1
  and equiv: poly-mod.eq-m p (D1 * H1) C
  and monD1: monic D1
  and normDH1: poly-mod.Mp p D1 = D1 poly-mod.Mp p H1 = H1
  and res: linear-hensel-main n = (D,H)
  and n: n ≠ 0
  and prime: prime p — p > 1 suffices if one does not need uniqueness
  and cop: poly-mod.coprime-m p D1 H1
  shows poly-mod.eq-m (p^n) (D * H) C
    ∧ monic D
    ∧ poly-mod.eq-m p D D1 ∧ poly-mod.eq-m p H H1
    ∧ poly-mod.Mp (p^n) D = D
    ∧ poly-mod.Mp (p^n) H = H ∧
    (poly-mod.eq-m (p^n) (D' * H') C ⟶
     poly-mod.eq-m p D' D1 ⟶
     poly-mod.eq-m p H' H1 ⟶
     poly-mod.Mp (p^n) D' = D' ⟶
     poly-mod.Mp (p^n) H' = H' ⟶ monic D' ⟶ D' = D ∧ H' = H)

  using res n
proof (induct n arbitrary: D H D' H')
  case (Suc n D' H' D'' H'')
  show ?case
  proof (cases n = 0)
    case True
    with Suc equiv monD1 normDH1 show ?thesis by auto
  next
    case False
    hence n: n ≠ 0 by auto
    let ?q = p^n
    let ?pq = p * p^n
    from prime have p: p > 1 using prime-gt-1-int by force
    from n p have q: ?q > 1 by auto
    from n p have pq: ?pq > 1 by (metis power-gt1-lemma)
    interpret p: poly-mod-2 p using p unfolding poly-mod-2-def .
    interpret q: poly-mod-2 ?q using q unfolding poly-mod-2-def .
    interpret pq: poly-mod-2 ?pq using pq unfolding poly-mod-2-def .
    obtain D H where rec: linear-hensel-main n = (D,H) by force
    obtain V where V: sdiv-poly (C - D * H) ?q = V by force
    obtain U where U: p.Mp (sdiv-poly (C - D * H) ?q) = U by auto
    obtain A B where dupe: p.dupe-monic-int D1 H1 S T U = (A,B) by force
    note IH = Suc(1)[OF rec n]
    from IH
    have CDH: q.eq-m (D * H) C
      and monD: monic D
```

and *p-eq*: *p.eq-m D D1 p.eq-m H H1*
and *norm*: *q.Mp D = D q.Mp H = H* **by** *auto*
from *n* **obtain** *k* **where** *n*: *n = Suc k* **by** (*cases n*, *auto*)
have *qq*: *?q * ?q = ?pq * p^k* **unfolding** *n* **by** *simp*
from *Suc(2)*[*unfolded n linear-hensel-main.simps*, *folded n*, *unfolded rec split Let-def U dupe*]
have *D'*: *D' = D + smult ?q B* **and** *H'*: *H' = H + smult ?q A* **by** *auto*
note *dupe = p.dupe-monic-int*[*OF 1 monD1 dupe*]
from *CDH* **have** *q.Mp C − q.Mp (D * H) = 0* **by** *simp*
hence *q.Mp (q.Mp C − q.Mp (D * H)) = 0* **by** *simp*
hence *q.Mp (C − D*H) = 0* **by** *simp*
from *q.Mp-0-smult-sdiv-poly*[*OF this*] **have** *CDHq*: *smult ?q (sdiv-poly (C − D * H) ?q) = C − D * H* **.**
have *ADBHU*: *p.eq-m (A * D + B * H) U* **using** *p-eq dupe(1)*
by (*metis* (*mono-tags*, *lifting*) *p.mult-Mp(2) poly-mod.plus-Mp*)
have *pq.Mp (D' * H') = pq.Mp ((D + smult ?q B) * (H + smult ?q A))*
unfolding *D' H'* **by** *simp*
also have *(D + smult ?q B) * (H + smult ?q A) = (D * H + smult ?q (A * D + B * H)) + smult (?q * ?q) (A * B)*
by (*simp add*: *field-simps smult-distribs*)
also have *pq.Mp ... = pq.Mp (D * H + pq.Mp (smult ?q (A * D + B * H)) + pq.Mp (smult (?q * ?q) (A * B)))*
using *pq.plus-Mp* **by** *metis*
also have *pq.Mp (smult (?q * ?q) (A * B)) = 0* **unfolding** *qq*
by (*metis pq.Mp-smult-m-0 smult-smult*)
finally have *DH'*: *pq.Mp (D' * H') = pq.Mp (D * H + pq.Mp (smult ?q (A * D + B * H)))* **by** *simp*
also have *pq.Mp (smult ?q (A * D + B * H)) = pq.Mp (smult ?q U)*
using *p.Mp-lift-modulus*[*OF ADBHU*, *of ?q*] **by** *simp*
also have *... = pq.Mp (C − D * H)*
unfolding *arg-cong*[*OF CDHq*, *of pq.Mp*, *symmetric*] *U*[*symmetric*] *V*
by (*rule p.Mp-lift-modulus*[*of - - ?q*], *auto*)
also have *pq.Mp (D * H + pq.Mp (C − D * H)) = pq.Mp C* **by** *simp*
finally have *CDH*: *pq.eq-m C (D' * H')* **by** *simp*

have *deg*: *degree D1 = degree D* **using** *p-eq(1) monD1 monD*
by (*metis p.monic-degree-m*)
have *mon*: *monic D'* **unfolding** *D'* **using** *dupe(2) monD* **unfolding** *deg* **by** (*rule monic-smult-add-small*)
have *normD'*: *pq.Mp D' = D'*
unfolding *D' pq.Mp-ident-iff poly-mod.Mp-coeff plus-poly.rep-eq coeff-smult*
proof
fix *i*
from *norm(1) dupe(4)* **have** *coeff D i ∈ {0..<?q} coeff B i ∈ {0..<p}*
unfolding *p.Mp-ident-iff q.Mp-ident-iff* **by** *auto*
thus *coeff D i + ?q * coeff B i ∈ {0..< ?pq}* **by** (*rule range-sum-prod*)
qed
have *normH'*: *pq.Mp H' = H'*
unfolding *H' pq.Mp-ident-iff poly-mod.Mp-coeff plus-poly.rep-eq coeff-smult*

```
  proof
    fix i
    from norm(2) dupe(3) have coeff H i ∈ {0..<?q} coeff A i ∈ {0..<p}
      unfolding p.Mp-ident-iff q.Mp-ident-iff by auto
    thus coeff H i + ?q * coeff A i ∈ {0..< ?pq} by (rule range-sum-prod)
  qed
  have eq: p.eq-m D D' p.eq-m H H' unfolding D' H' n
     poly-eq-iff p.Mp-coeff p.M-def by (auto simp: field-simps)
  with p-eq have eq: p.eq-m D' D1 p.eq-m H' H1 by auto
  {
    assume CDH'': pq.eq-m C (D'' * H'')
      and DH1'': p.eq-m D1 D'' p.eq-m H1 H''
      and norm'': pq.Mp D'' = D'' pq.Mp H'' = H''
      and monD'': monic D''
    from q.Dp-Mp-eq[of D''] obtain d B' where D'': D'' = q.Mp d + smult ?q
B' by auto
    from q.Dp-Mp-eq[of H''] obtain h A' where H'': H'' = q.Mp h + smult ?q
A' by auto
    {
      fix A B
      assume *: pq.Mp (q.Mp A + smult ?q B) = q.Mp A + smult ?q B
      have p.Mp B = B unfolding p.Mp-ident-iff
      proof
        fix i
        from arg-cong[OF *, of λ f. coeff f i, unfolded pq.Mp-coeff pq.M-def]
       have coeff (q.Mp A + smult ?q B) i ∈ {0 ..< ?pq} using * pq.Mp-ident-iff
by blast
        hence sum: coeff (q.Mp A) i + ?q * coeff B i ∈ {0 ..< ?pq} by auto
        have q.Mp (q.Mp A) = q.Mp A by auto
       from this[unfolded q.Mp-ident-iff] have A: coeff (q.Mp A) i ∈ {0 ..< p^n}
by auto
        {
          assume coeff B i < 0 hence coeff B i ≤ -1 by auto
           from mult-left-mono[OF this, of ?q] q.m1 have ?q * coeff B i ≤ -?q
by simp
          with A sum have False by auto
        } hence coeff B i ≥ 0 by force
        moreover
        {
          assume coeff B i ≥ p
         from mult-left-mono[OF this, of ?q] q.m1 have ?q * coeff B i ≥ ?pq by
simp
          with A sum have False by auto
        } hence coeff B i < p by force
        ultimately show coeff B i ∈ {0 ..< p} by auto
      qed
    } note norm-convert = this
   from norm-convert[OF norm''(1)[unfolded D'']] have normB': p.Mp B' = B'
.
```

from *norm-convert*[*OF norm''*(*2*)[*unfolded H''*]] **have** *normA'*: *p.Mp A' = A'*
.
let *?d = q.Mp d*
let *?h = q.Mp h*
{
assume *lt*: *degree ?d < degree B'*
hence *eq*: *degree D'' = degree B'* **unfolding** *D''* **using** *q.m1 p.m1*
by (*subst degree-add-eq-right*, *auto*)
from *lt* **have** [*simp*]: *coeff ?d* (*degree B'*) *= 0* **by** (*rule coeff-eq-0*)
from *monD''*[*unfolded eq*, *unfolded D''*, *simplified*] *False q.m1 lt* **have** *False*
by (*metis mod-mult-self1-is-0 poly-mod.M-def q.M-1 zero-neq-one*)
}
hence *deg-dB'*: *degree ?d ≥ degree B'* **by** *presburger*
{
assume *eq*: *degree ?d = degree B'* **and** *B'*: *B' ≠ 0*
let *?B = coeff B'* (*degree B'*)
from *normB'*[*unfolded p.Mp-ident-iff*, *rule-format*, *of degree B'*] *B'*
have *?B* ∈ {*0..<p*} − {*0*} **by** *simp*
hence *bnds*: *?B > 0 ?B < p* **by** *auto*
have *degD''*: *degree D'' ≤ degree ?d* **unfolding** *D''* **using** *eq* **by** (*simp add*: *degree-add-le*)
have *?q * ?B ≥ 1 * 1* **by** (*rule mult-mono*, *insert q.m1 bnds*, *auto*)
moreover have *coeff D''* (*degree ?d*) *= 1 + ?q * ?B* **using** *monD''*
unfolding *D''* **using** *eq*
by (*metis D'' coeff-smult monD'' plus-poly.rep-eq poly-mod.Dp-Mp-eq*
poly-mod.degree-m-eq-monic poly-mod.plus-Mp(*1*)
q.Mp-smult-m-0 q.m1 q.monic-Mp q.plus-Mp(*2*))
ultimately have *gt*: *coeff D''* (*degree ?d*) *> 1* **by** *auto*
hence *coeff D''* (*degree ?d*) *≠ 0* **by** *auto*
hence *degree D'' ≥ degree ?d* **by** (*rule le-degree*)
with *degree-add-le-max*[*of ?d smult ?q B'*, *folded D'*] *eq*
have *deg*: *degree D'' = degree ?d* **using** *degD''* **by** *linarith*
from *gt*[*folded this*] **have** ¬ *monic D''* **by** *auto*
with *monD''* **have** *False* **by** *auto*
}
with *deg-dB'* **have** *deg-dB2*: *B' = 0 ∨ degree B' < degree ?d* **by** *fastforce*
have *d*: *q.Mp D'' = ?d* **unfolding** *D''*
by (*metis add.right-neutral poly-mod.Mp-smult-m-0 poly-mod.plus-Mp*)
have *h*: *q.Mp H'' = ?h* **unfolding** *H''*
by (*metis add.right-neutral poly-mod.Mp-smult-m-0 poly-mod.plus-Mp*)
from *CDH''* **have** *pq.Mp C = pq.Mp* (*D'' * H''*) **by** *simp*
from *arg-cong*[*OF this*, *of q.Mp*]
have *q.Mp C = q.Mp* (*D'' * H''*)
using *p.m1 q.Mp-product-modulus* **by** *auto*
also have *... = q.Mp* (*q.Mp D'' * q.Mp H''*) **by** *simp*
also have *... = q.Mp* (*?d * ?h*) **unfolding** *d h* **by** *simp*
finally have *eqC*: *q.eq-m* (*?d * ?h*) *C* **by** *auto*
have *d1*: *p.eq-m ?d D1* **unfolding** *d*[*symmetric*] **using** *DH1''*
using *assms*(*4*) *n p.Mp-product-modulus p.m1* **by** *auto*

```
    have h1: p.eq-m ?h H1 unfolding h[symmetric] using DH1''
      using assms(5) n p.Mp-product-modulus p.m1 by auto
    have mond: monic (q.Mp d) using monD'' deg-dB2 unfolding D''
      using d q.monic-Mp[OF monD''] by simp
    from eqC d1 h1 mond IH[of q.Mp d q.Mp h] have IH: ?d = D ?h = H by
auto
    from deg-dB2[unfolded IH] have degB': B' = 0 ∨ degree B' < degree D by
auto
   from IH have D'': D'' = D + smult ?q B' and H'': H'' = H + smult ?q A'
      unfolding D'' H'' by auto
    have pq.Mp (D'' * H'') = pq.Mp (D' * H') using CDH'' CDH by simp
    also have pq.Mp (D'' * H'') = pq.Mp ((D + smult ?q B') * (H + smult ?q
A'))
      unfolding D'' H'' by simp
    also have (D + smult ?q B') * (H + smult ?q A') = (D * H + smult ?q (A'
* D + B' * H)) + smult (?q * ?q) (A' * B')
      by (simp add: field-simps smult-distribs)
    also have pq.Mp ... = pq.Mp (D * H + pq.Mp (smult ?q (A' * D + B' *
H)) + pq.Mp (smult (?q * ?q) (A' * B')))
      using pq.plus-Mp by metis
    also have pq.Mp (smult (?q * ?q) (A' * B')) = 0 unfolding qq
      by (metis pq.Mp-smult-m-0 smult-smult)
    finally have pq.Mp (D * H + pq.Mp (smult ?q (A' * D + B' * H)))
      = pq.Mp (D * H + pq.Mp (smult ?q (A * D + B * H))) unfolding DH'
by simp
    hence pq.Mp (smult ?q (A' * D + B' * H)) = pq.Mp (smult ?q (A * D + B
* H))
        by (metis (no-types, lifting) add-diff-cancel-left' poly-mod.minus-Mp(1)
poly-mod.plus-Mp(2))
      hence p.Mp (A' * D + B' * H) = p.Mp (A * D + B * H) unfolding
poly-eq-iff p.Mp-coeff pq.Mp-coeff coeff-smult
      by (insert p, auto simp: p.M-def pq.M-def)
    hence p.Mp (A' * D1 + B' * H1) = p.Mp (A * D1 + B * H1) using p-eq
      by (metis p.mult-Mp(2) poly-mod.plus-Mp)
    hence eq: p.eq-m (A' * D1 + B' * H1) U using dupe(1) by auto
    have degree D = degree D1 using monD monD1
       arg-cong[OF p-eq(1), of degree]
       p.degree-m-eq-monic[OF - p.m1] by auto
    hence B' = 0 ∨ degree B' < degree D1 using degB' by simp
    from dupe(5)[OF cop eq this normDH1(1) normA' normB' prime] have A'
= A B' = B by auto
    hence D'' = D' H'' = H' unfolding D'' H'' D' H' by auto
   }
   thus ?thesis using normD' normH' CDH mon eq by simp
  qed
qed simp
end
end
```

definition *linear-hensel-binary* :: *int* ⇒ *nat* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* ⇒ *int poly* × *int poly* **where**
linear-hensel-binary p n C D H = (*let*
(*S*,*T*) = *euclid-ext-poly-dynamic p D H*
in linear-hensel-main C p S T D H n)

lemma (**in** *poly-mod-prime*) *unique-hensel-binary*:
assumes *prime*: *prime p*
and *cop*: *coprime-m D H* **and** *eq*: *eq-m* (*D* ∗ *H*) *C*
and *normalized-input*: *Mp D* = *D Mp H* = *H*
and *monic-input*: *monic D*
and *n*: *n* ≠ *0*
shows ∃! (*D'*,*H'*). — *D'*, *H'* are computed via *linear-hensel-binary*
poly-mod.eq-m (*p*^*n*) (*D'* ∗ *H'*) *C* — the main result: equivalence mod *p*^*n*
∧ *monic D'* — monic output
∧ *eq-m D D'* ∧ *eq-m H H'* — apply ‘*mod p*‘ on *D'* and *H'* yields *D* and *H* again
∧ *poly-mod.Mp* (*p*^*n*) *D'* = *D'* ∧ *poly-mod.Mp* (*p*^*n*) *H'* = *H'* — output is normalized
proof −
obtain *D' H'* **where** *hensel-result*: *linear-hensel-binary p n C D H* = (*D'*,*H'*) **by** *force*
from *m1* **have** *p*: *p* > *1* **.**
obtain *S T* **where** *ext*: *euclid-ext-poly-dynamic p D H* = (*S*,*T*) **by** *force*
obtain *D1 H1* **where** *main*: *linear-hensel-main C p S T D H n* = (*D1*,*H1*) **by** *force*
from *hensel-result*[*unfolded linear-hensel-binary-def ext split Let-def main*]
have *id*: *D1* = *D' H1* = *H'* **by** *auto*
note *eucl* = *euclid-ext-poly-dynamic* [*OF cop normalized-input ext*]
from *linear-hensel-main* [*OF eucl*(*1*)
eq monic-input normalized-input main [*unfolded id*] *n prime cop*]
show *?thesis* **by** (*intro ex1I*, *auto*)
qed

context
fixes *C* :: *int poly*
begin

lemma *hensel-step-main*: **assumes**
one-q: *poly-mod.eq-m q* (*D* ∗ *S* + *H* ∗ *T*) *1*
and *one-p*: *poly-mod.eq-m p* (*D1* ∗ *S1* + *H1* ∗ *T1*) *1*
and *CDHq*: *poly-mod.eq-m q C* (*D* ∗ *H*)
and *D1D*: *poly-mod.eq-m p D1 D*
and *H1H*: *poly-mod.eq-m p H1 H*
and *S1S*: *poly-mod.eq-m p S1 S*
and *T1T*: *poly-mod.eq-m p T1 T*
and *mon*: *monic D*
and *mon1*: *monic D1*
and *q*: *q* > *1*

and *p*: *p* > *1*
and *D1*: *poly-mod.Mp p D1* = *D1*
and *H1*: *poly-mod.Mp p H1* = *H1*
and *S1*: *poly-mod.Mp p S1* = *S1*
and *T1*: *poly-mod.Mp p T1* = *T1*
and *D*: *poly-mod.Mp q D* = *D*
and *H*: *poly-mod.Mp q H* = *H*
and *S*: *poly-mod.Mp q S* = *S*
and *T*: *poly-mod.Mp q T* = *T*
and *U1*: *U1* = *poly-mod.Mp p* (*sdiv-poly* (*C* − *D* ∗ *H*) *q*)
and *dupe1*: *dupe-monic-dynamic p D1 H1 S1 T1 U1* = (*A*,*B*)
and *D'*: *D'* = *D* + *smult q B*
and *H'*: *H'* = *H* + *smult q A*
and *U2*: *U2* = *poly-mod.Mp q* (*sdiv-poly* (*S*∗*D'* + *T*∗*H'* − *1*) *p*)
and *dupe2*: *dupe-monic-dynamic q D H S T U2* = (*A'*,*B'*)
and *rq*: *r* = *p* ∗ *q*
and *pq*: *p dvd q*
and *S'*: *S'* = *poly-mod.Mp r* (*S* − *smult p A'*)
and *T'*: *T'* = *poly-mod.Mp r* (*T* − *smult p B'*)
shows *poly-mod.eq-m r C* (*D'* ∗ *H'*)
poly-mod.Mp r D' = *D'*
poly-mod.Mp r H' = *H'*
poly-mod.Mp r S' = *S'*
poly-mod.Mp r T' = *T'*
poly-mod.eq-m r (*D'* ∗ *S'* + *H'* ∗ *T'*) *1*
monic D'
unfolding *rq*
proof −
from *pq* **obtain** *k* **where** *qp*: *q* = *p* ∗ *k* **unfolding** *dvd-def* **by** *auto*
from *arg-cong*[*OF qp*, *of sgn*] *q p* **have** *k0*: *k* > *0* **unfolding** *sgn-mult* **by** (*auto simp*: *sgn-1-pos*)
from *qp* **have** *qq*: *q* ∗ *q* = *p* ∗ *q* ∗ *k* **by** *auto*
let *?r* = *p* ∗ *q*
interpret *poly-mod-2 p* **by** (*standard*, *insert p*, *auto*)
interpret *q*: *poly-mod-2 q* **by** (*standard*, *insert q*, *auto*)
from *p q* **have** *r*: *?r* > *1* **by** (*simp add*: *less-1-mult*)
interpret *r*: *poly-mod-2 ?r* **using** *r* **unfolding** *poly-mod-2-def* **.**
have *Mp-conv*: *Mp* (*q.Mp x*) = *Mp x* **for** *x* **unfolding** *qp*
by (*rule Mp-product-modulus*[*OF refl k0*])
from *arg-cong*[*OF CDHq*, *of Mp*, *unfolded Mp-conv*] **have** *Mp C* = *Mp* (*Mp D* ∗ *Mp H*)
by *simp*
also have *Mp D* = *Mp D1* **using** *D1D* **by** *simp*
also have *Mp H* = *Mp H1* **using** *H1H* **by** *simp*
finally have *CDHp*: *eq-m C* (*D1* ∗ *H1*) **by** *simp*
have *Mp U1* = *U1* **unfolding** *U1* **by** *simp*
note *dupe1* = *dupe-monic-dynamic*[*OF dupe1 one-p mon1 D1 H1 S1 T1 this*]
have *q.Mp U2* = *U2* **unfolding** *U2* **by** *simp*
note *dupe2* = *q.dupe-monic-dynamic*[*OF dupe2 one-q mon D H S T this*]

from *CDHq* **have** *q.Mp C − q.Mp (D ∗ H) = 0* **by** *simp*
hence *q.Mp (q.Mp C − q.Mp (D ∗ H)) = 0* **by** *simp*
hence *q.Mp (C − D∗H) = 0* **by** *simp*
from *q.Mp-0-smult-sdiv-poly*[*OF this*] **have** *CDHq*: *smult q (sdiv-poly (C − D ∗ H) q) = C − D ∗ H* **.**
{
fix *A B*
have *Mp (A ∗ D1 + B ∗ H1) = Mp (Mp (A ∗ D1) + Mp (B ∗ H1))* **by** *simp*
also have *Mp (A ∗ D1) = Mp (A ∗ Mp D1)* **by** *simp*
also have *... = Mp (A ∗ D)* **unfolding** *D1D* **by** *simp*
also have *Mp (B ∗ H1) = Mp (B ∗ Mp H1)* **by** *simp*
also have *... = Mp (B ∗ H)* **unfolding** *H1H* **by** *simp*
finally have *Mp (A ∗ D1 + B ∗ H1) = Mp (A ∗ D + B ∗ H)* **by** *simp*
} note *D1H1 = this*
have *r.Mp (D′ ∗ H′) = r.Mp ((D + smult q B) ∗ (H + smult q A))*
unfolding *D′ H′* **by** *simp*
also have *(D + smult q B) ∗ (H + smult q A) = (D ∗ H + smult q (A ∗ D + B ∗ H)) + smult (q ∗ q) (A ∗ B)*
by (*simp add*: *field-simps smult-distribs*)
also have *r.Mp ... = r.Mp (D ∗ H + r.Mp (smult q (A ∗ D + B ∗ H)) + r.Mp (smult (q ∗ q) (A ∗ B)))*
using *r.plus-Mp* **by** *metis*
also have *r.Mp (smult (q ∗ q) (A ∗ B)) = 0* **unfolding** *qq*
by (*metis r.Mp-smult-m-0 smult-smult*)
also have *r.Mp (smult q (A ∗ D + B ∗ H)) = r.Mp (smult q U1)*
proof (*rule Mp-lift-modulus*[*of - - q*])
show *Mp (A ∗ D + B ∗ H) = Mp U1* **using** *dupe1*(*1*) **unfolding** *D1H1* **by** *simp*
qed
also have *... = r.Mp (C − D ∗ H)*
unfolding *arg-cong*[*OF CDHq*, *of r.Mp*, *symmetric*]
using *Mp-lift-modulus*[*of U1 sdiv-poly (C − D ∗ H) q q*] **unfolding** *U1*
by *simp*
also have *r.Mp (D ∗ H + r.Mp (C − D ∗ H) + 0) = r.Mp C* **by** *simp*
finally show *CDH*: *r.eq-m C (D′ ∗ H′)* **by** *simp*
have *degree D1 = degree (Mp D1)* **using** *mon1* **by** *simp*
also have *... = degree D* **unfolding** *D1D* **using** *mon* **by** *simp*
finally have *deg-eq*: *degree D1 = degree D* **by** *simp*
show *mon*: *monic D′* **unfolding** *D′* **using** *dupe1*(*2*) *mon* **unfolding** *deg-eq* **by** (*rule monic-smult-add-small*)
have *Mp (S ∗ D′ + T ∗ H′ − 1) = Mp (Mp (D ∗ S + H ∗ T) + (smult q (S ∗ B + T ∗ A) − 1))*
unfolding *D′ H′ plus-Mp* **by** (*simp add*: *field-simps smult-distribs*)
also have *Mp (D ∗ S + H ∗ T) = Mp (Mp (D1 ∗ Mp S) + Mp (H1 ∗ Mp T))* **using** *D1H1*[*of S T*] **by** (*simp add*: *ac-simps*)
also have *... = 1* **using** *one-p* **unfolding** *S1S*[*symmetric*] *T1T*[*symmetric*] **by** *simp*
also have *Mp (1 + (smult q (S ∗ B + T ∗ A) − 1)) = Mp (smult q (S ∗ B + T ∗ A))* **by** *simp*

also have ... = *0* **unfolding** *qp* **by** (*metis Mp-smult-m-0 smult-smult*)
finally have *Mp* (*S* ∗ *D'* + *T* ∗ *H'* − *1*) = *0* **.**
from *Mp-0-smult-sdiv-poly*[*OF this*]
have *SDTH*: *smult p* (*sdiv-poly* (*S* ∗ *D'* + *T* ∗ *H'* − *1*) *p*) = *S* ∗ *D'* + *T* ∗ *H'* − *1* **.**
have *swap*: *q* ∗ *p* = *p* ∗ *q* **by** *simp*
have *r.Mp* (*D'* ∗ *S'* + *H'* ∗ *T'*) =
r.Mp ((*D* + *smult q B*) ∗ (*S* − *smult p A'*) + (*H* + *smult q A*) ∗ (*T* − *smult p B'*))
unfolding *D' S' H' T' rq* **using** *r.plus-Mp r.mult-Mp* **by** *metis*
also have ... = *r.Mp* ((*D* ∗ *S* + *H* ∗ *T* +
smult q (*B* ∗ *S* + *A* ∗ *T*)) − *smult p* (*A'* ∗ *D* + *B'* ∗ *H*) − *smult ?r* (*A* ∗ *B'* + *B* ∗ *A'*))
by (*simp add*: *field-simps smult-distribs*)
also have ... = *r.Mp* ((*D* ∗ *S* + *H* ∗ *T* +
smult q (*B* ∗ *S* + *A* ∗ *T*)) − *r.Mp* (*smult p* (*A'* ∗ *D* + *B'* ∗ *H*)) − *r.Mp* (*smult ?r* (*A* ∗ *B'* + *B* ∗ *A'*)))
using *r.plus-Mp r.minus-Mp* **by** *metis*
also have *r.Mp* (*smult ?r* (*A* ∗ *B'* + *B* ∗ *A'*)) = *0* **by** *simp*
also have *r.Mp* (*smult p* (*A'* ∗ *D* + *B'* ∗ *H*)) = *r.Mp* (*smult p U2*)
using *q.Mp-lift-modulus*[*OF dupe2*(*1*), *of p*] **unfolding** *swap* **.**
also have ... = *r.Mp* (*S* ∗ *D'* + *T* ∗ *H'* − *1*)
unfolding *arg-cong*[*OF SDTH*, *of r.Mp*, *symmetric*]
using *q.Mp-lift-modulus*[*of U2 sdiv-poly* (*S* ∗ *D'* + *T* ∗ *H'* − *1*) *p p*]
unfolding *U2 swap* **by** *simp*
also have *S* ∗ *D'* + *T* ∗ *H'* − *1* = *S* ∗ *D* + *T* ∗ *H* + *smult q* (*B* ∗ *S* + *A* ∗ *T*) − *1*
unfolding *D' H'* **by** (*simp add*: *field-simps smult-distribs*)
also have *r.Mp* (*D* ∗ *S* + *H* ∗ *T* + *smult q* (*B* ∗ *S* + *A* ∗ *T*) −
r.Mp (*S* ∗ *D* + *T* ∗ *H* + *smult q* (*B* ∗ *S* + *A* ∗ *T*) − *1*) − *0*)
= *1* **by** *simp*
finally show *1*: *r.eq-m* (*D'* ∗ *S'* + *H'* ∗ *T'*) *1* **by** *simp*
show *D'*: *r.Mp D'* = *D'* **unfolding** *D' r.Mp-ident-iff poly-mod.Mp-coeff plus-poly.rep-eq coeff-smult*
proof
fix *n*
from *D dupe1*(*4*) **have** *coeff D n* ∈ {*0*..<*q*} *coeff B n* ∈ {*0*..<*p*}
unfolding *q.Mp-ident-iff Mp-ident-iff* **by** *auto*
thus *coeff D n* + *q* ∗ *coeff B n* ∈ {*0*..<*?r*} **by** (*metis range-sum-prod*)
qed
show *H'*: *r.Mp H'* = *H'* **unfolding** *H' r.Mp-ident-iff poly-mod.Mp-coeff plus-poly.rep-eq coeff-smult*
proof
fix *n*
from *H dupe1*(*3*) **have** *coeff H n* ∈ {*0*..<*q*} *coeff A n* ∈ {*0*..<*p*}
unfolding *q.Mp-ident-iff Mp-ident-iff* **by** *auto*
thus *coeff H n* + *q* ∗ *coeff A n* ∈ {*0*..<*?r*} **by** (*metis range-sum-prod*)
qed
show *poly-mod.Mp ?r S'* = *S' poly-mod.Mp ?r T'* = *T'*

unfolding S' T' *rq* **by** *auto*
qed

definition *hensel-step* **where**
hensel-step p q S1 T1 D1 H1 S T D H = (
let U = *poly-mod.Mp p* (*sdiv-poly* (*C* − *D* ∗ *H*) *q*); — *Z2 and Z3*
(*A*,*B*) = *dupe-monic-dynamic p D1 H1 S1 T1 U*;
D' = *D* + *smult q B*; — *Z4*
H' = *H* + *smult q A*;
U' = *poly-mod.Mp q* (*sdiv-poly* (*S*∗D' + *T*∗H' − *1*) *p*); — *Z5 + Z6*
(*A*′,*B*′) = *dupe-monic-dynamic q D H S T U*′;
q' = *p* ∗ *q*;
S' = *poly-mod.Mp q*′ (*S* − *smult p A*′); — *Z7*
T' = *poly-mod.Mp q*′ (*T* − *smult p B*′)
in (S',T',D',H'))

definition *quadratic-hensel-step q S T D H* = *hensel-step q q S T D H S T D H*

lemma *quadratic-hensel-step-code*[*code*]:
quadratic-hensel-step q S T D H =
(*let dupe* = *dupe-monic-dynamic q D H S T*; — this will share the conversions of *D H S T*
U = *poly-mod.Mp q* (*sdiv-poly* (*C* − *D* ∗ *H*) *q*);
(*A*, *B*) = *dupe U*;
D' = *D* + *Polynomial.smult q B*;
H' = *H* + *Polynomial.smult q A*;
U' = *poly-mod.Mp q* (*sdiv-poly* (*S* ∗ D' + *T* ∗ H' − *1*) *q*);
(*A*′, *B*′) = *dupe U*′;
q' = *q* ∗ *q*;
S' = *poly-mod.Mp q*′ (*S* − *Polynomial.smult q A*′);
T' = *poly-mod.Mp q*′ (*T* − *Polynomial.smult q B*′)
in (S', T', D', H'))
unfolding *quadratic-hensel-step-def*[*unfolded hensel-step-def*] *Let-def* **..**

definition *simple-quadratic-hensel-step* **where** — do not compute new values S' and T'
simple-quadratic-hensel-step q S T D H = (
let U = *poly-mod.Mp q* (*sdiv-poly* (*C* − *D* ∗ *H*) *q*); — *Z2 + Z3*
(*A*,*B*) = *dupe-monic-dynamic q D H S T U*;
D' = *D* + *smult q B*; — *Z4*
H' = *H* + *smult q A*
in (D',H'))

lemma *hensel-step*: **assumes** *step*: *hensel-step p q S1 T1 D1 H1 S T D H* = (S', T', D', H')
and *one-p*: *poly-mod.eq-m p* (*D1* ∗ *S1* + *H1* ∗ *T1*) *1*
and *mon1*: *monic D1*
and *p*: *p* > *1*
and *CDHq*: *poly-mod.eq-m q C* (*D* ∗ *H*)

and *one-q*: *poly-mod.eq-m q* (*D* ∗ *S* + *H* ∗ *T*) *1*
and *D1D*: *poly-mod.eq-m p D1 D*
and *H1H*: *poly-mod.eq-m p H1 H*
and *S1S*: *poly-mod.eq-m p S1 S*
and *T1T*: *poly-mod.eq-m p T1 T*
and *mon*: *monic D*
and *q*: *q* > *1*
and *D1*: *poly-mod.Mp p D1* = *D1*
and *H1*: *poly-mod.Mp p H1* = *H1*
and *S1*: *poly-mod.Mp p S1* = *S1*
and *T1*: *poly-mod.Mp p T1* = *T1*
and *D*: *poly-mod.Mp q D* = *D*
and *H*: *poly-mod.Mp q H* = *H*
and *S*: *poly-mod.Mp q S* = *S*
and *T*: *poly-mod.Mp q T* = *T*
and *rq*: *r* = *p* ∗ *q*
and *pq*: *p dvd q*
shows
poly-mod.eq-m r C (*D*′ ∗ *H*′)
poly-mod.eq-m r (*D*′ ∗ *S*′ + *H*′ ∗ *T*′) *1*
poly-mod.Mp r D′ = *D*′
poly-mod.Mp r H′ = *H*′
poly-mod.Mp r S′ = *S*′
poly-mod.Mp r T′ = *T*′
poly-mod.Mp p D1 = *poly-mod.Mp p D*′
poly-mod.Mp p H1 = *poly-mod.Mp p H*′
poly-mod.Mp p S1 = *poly-mod.Mp p S*′
poly-mod.Mp p T1 = *poly-mod.Mp p T*′
monic D′
proof −
define *U* **where** *U*: *U* = *poly-mod.Mp p* (*sdiv-poly* (*C* − *D* ∗ *H*) *q*)
note *step* = *step*[*unfolded hensel-step-def Let-def*, *folded U*]
obtain *A B* **where** *dupe1*: *dupe-monic-dynamic p D1 H1 S1 T1 U* = (*A*,*B*) **by** *force*
note *step* = *step*[*unfolded dupe1 split*]
from *step* **have** *D*′: *D*′ = *D* + *smult q B* **and** *H*′: *H*′ = *H* + *smult q A*
by (*auto split*: *prod.splits*)
define *U*′ **where** *U*′: *U*′ = *poly-mod.Mp q* (*sdiv-poly* (*S* ∗ *D*′ + *T* ∗ *H*′ − *1*) *p*)
obtain *A*′ *B*′ **where** *dupe2*: *dupe-monic-dynamic q D H S T U*′ = (*A*′,*B*′) **by** *force*
from *step*[*folded D*′ *H*′, *folded U*′, *unfolded dupe2 split*, *folded rq*]
have *S*′: *S*′ = *poly-mod.Mp r* (*S* − *Polynomial.smult p A*′) **and**
T′: *T*′ = *poly-mod.Mp r* (*T* − *Polynomial.smult p B*′) **by** *auto*
from *hensel-step-main*[*OF one-q one-p CDHq D1D H1H S1S T1T mon mon1 q p D1 H1 S1 T1 D H S T U*
dupe1 D′ *H*′ *U*′ *dupe2 rq pq S*′ *T*′]
show *poly-mod.eq-m r* (*D*′ ∗ *S*′ + *H*′ ∗ *T*′) *1*
poly-mod.eq-m r C (*D*′ ∗ *H*′)

poly-mod.Mp r D′ = D′
poly-mod.Mp r H′ = H′
poly-mod.Mp r S′ = S′
poly-mod.Mp r T′ = T′
monic D′ **by** *auto*
from *pq* **obtain** *s* **where** *q*: *q = p * s* **by** (*metis dvdE*)
show *poly-mod.Mp p D1 = poly-mod.Mp p D′*
poly-mod.Mp p H1 = poly-mod.Mp p H′
unfolding *q D′ D1D H′ H1H*
by (*metis add.right-neutral poly-mod.Mp-smult-m-0 poly-mod.plus-Mp(2) smult-smult*)+

from ‹*q > 1*› **have** *q0*: *q > 0* **by** *auto*
show *poly-mod.Mp p S1 = poly-mod.Mp p S′*
poly-mod.Mp p T1 = poly-mod.Mp p T′
unfolding *S′ S1S T′ T1T poly-mod-2.Mp-product-modulus*[*OF poly-mod-2.intro*[*OF* ‹*p > 1*›] *rq q0*]
by (*metis group-add-class.diff-0-right poly-mod.Mp-smult-m-0 poly-mod.minus-Mp(2)*)+

qed

lemma *quadratic-hensel-step*: **assumes** *step*: *quadratic-hensel-step q S T D H = (S′, T′, D′, H′)*
and *CDH*: *poly-mod.eq-m q C (D * H)*
and *one*: *poly-mod.eq-m q (D * S + H * T) 1*
and *D*: *poly-mod.Mp q D = D*
and *H*: *poly-mod.Mp q H = H*
and *S*: *poly-mod.Mp q S = S*
and *T*: *poly-mod.Mp q T = T*
and *mon*: *monic D*
and *q*: *q > 1*
and *rq*: *r = q * q*
shows
*poly-mod.eq-m r C (D′ * H′)*
*poly-mod.eq-m r (D′ * S′ + H′ * T′) 1*
poly-mod.Mp r D′ = D′
poly-mod.Mp r H′ = H′
poly-mod.Mp r S′ = S′
poly-mod.Mp r T′ = T′
poly-mod.Mp q D = poly-mod.Mp q D′
poly-mod.Mp q H = poly-mod.Mp q H′
poly-mod.Mp q S = poly-mod.Mp q S′
poly-mod.Mp q T = poly-mod.Mp q T′
monic D′
proof (*atomize*(*full*), *goal-cases*)
case *1*
from *hensel-step*[*OF step*[*unfolded quadratic-hensel-step-def*] *one mon q CDH one refl refl refl refl mon q D H S T D H S T rq*]
show *?case* **by** *auto*
qed

```
context
  fixes p :: int and S1 T1 D1 H1 :: int poly
begin
private lemma decrease[termination-simp]: ¬ j ≤ 1 ⟹ odd j ⟹ Suc (j div 2)
< j by presburger

fun quadratic-hensel-loop where
  quadratic-hensel-loop (j :: nat) = (
      if j ≤ 1 then (p, S1, T1, D1, H1) else
      if even j then
          (case quadratic-hensel-loop (j div 2) of
            (q, S, T, D, H) ⇒
          let qq = q * q in
          (case quadratic-hensel-step q S T D H of — quadratic step
            (S', T', D', H') ⇒ (qq, S', T', D', H')))
      else — odd j
        (case quadratic-hensel-loop (j div 2 + 1) of
          (q, S, T, D, H) ⇒
          (case quadratic-hensel-step q S T D H of — quadratic step
            (S', T', D', H') ⇒
                let qq = q * q; pj = qq div p; down = poly-mod.Mp pj in
                  (pj, down S', down T', down D', down H'))))

definition quadratic-hensel-main j = (case quadratic-hensel-loop j of
    (qq, S, T, D, H) ⇒ (D, H))

declare quadratic-hensel-loop.simps[simp del]

— unroll the definition of hensel-loop so that in outermost iteration we can use
simple-hensel-step
lemma quadratic-hensel-main-code[code]: quadratic-hensel-main j = (
   if j ≤ 1 then (D1, H1)
      else if even j
      then (case quadratic-hensel-loop (j div 2) of
            (q, S, T, D, H) ⇒
               simple-quadratic-hensel-step q S T D H)
       else (case quadratic-hensel-loop (j div 2 + 1) of
            (q, S, T, D, H) ⇒
              (case simple-quadratic-hensel-step q S T D H of
               (D', H') ⇒ let down = poly-mod.Mp (q * q div p) in (down D', down
H'))))
  unfolding quadratic-hensel-loop.simps[of j] quadratic-hensel-main-def Let-def
  by (simp split: if-splits prod.splits option.splits sum.splits
      add: quadratic-hensel-step-code simple-quadratic-hensel-step-def Let-def)


context
  fixes j :: nat
```

```
  assumes 1: poly-mod.eq-m p (D1 * S1 + H1 * T1) 1
  and CDH1: poly-mod.eq-m p C (D1 * H1)
  and mon1: monic D1
  and p: p > 1
  and D1: poly-mod.Mp p D1 = D1
  and H1: poly-mod.Mp p H1 = H1
  and S1: poly-mod.Mp p S1 = S1
  and T1: poly-mod.Mp p T1 = T1
  and j: j ≥ 1
begin

lemma quadratic-hensel-loop:
  assumes quadratic-hensel-loop j = (q, S, T, D, H)
  shows (poly-mod.eq-m q C (D * H) ∧ monic D
    ∧ poly-mod.eq-m p D1 D ∧ poly-mod.eq-m p H1 H
    ∧ poly-mod.eq-m q (D * S + H * T) 1
    ∧ poly-mod.Mp q D = D ∧ poly-mod.Mp q H = H
    ∧ poly-mod.Mp q S = S ∧ poly-mod.Mp q T = T
    ∧ q = p^j)
  using j assms
proof (induct j arbitrary: q S T D H rule: less-induct)
  case (less j q' S' T' D' H')
  note res = less(3)
  interpret poly-mod-2 p using p by (rule poly-mod-2.intro)
  let ?hens = quadratic-hensel-loop
  note simp[simp] = quadratic-hensel-loop.simps[of j]
  show ?case
  proof (cases j = 1)
    case True
    show ?thesis using res simp unfolding True using CDH1 1 mon1 D1 H1 S1
T1 by auto
  next
    case False
    with less(2) have False: (j ≤ 1) = False by auto
     have mod-2: k ≥ 1 ⟹ poly-mod-2 (p^k) for k by (intro poly-mod-2.intro,
insert p, auto)
    {
      fix k D
      assume *: k ≥ 1 k ≤ j poly-mod.Mp (p ^ k) D = D
      from *(2) have {0..<p ^ k} ⊆ {0..<p ^ j} using p by auto
      hence poly-mod.Mp (p ^ j) D = D
        unfolding poly-mod-2.Mp-ident-iff[OF mod-2[OF less(2)]]
        using *(3)[unfolded poly-mod-2.Mp-ident-iff[OF mod-2[OF *(1)]]] by blast
    } note lift-norm = this
    show ?thesis
    proof (cases even j)
      case True
      let ?j2 = j div 2
      from False have lt: ?j2 < j 1 ≤ ?j2 by auto
```

```
    obtain q S T D H where rec: ?hens ?j2 = (q, S, T, D, H) by (cases ?hens
?j2, auto)
    note IH = less(1)[OF lt rec]
    from IH
    have ∗: poly-mod.eq-m q C (D ∗ H)
      poly-mod.eq-m q (D ∗ S + H ∗ T) 1
      monic D
      eq-m D1 D
      eq-m H1 H
      poly-mod.Mp q D = D
      poly-mod.Mp q H = H
      poly-mod.Mp q S = S
      poly-mod.Mp q T = T
      q = p ^ ?j2
      by auto
    hence norm: poly-mod.Mp (p ^ j) D = D poly-mod.Mp (p ^ j) H = H
      poly-mod.Mp (p ^ j) S = S poly-mod.Mp (p ^ j) T = T
      using lift-norm[OF lt(2)] by auto
    from lt p have q: q > 1 unfolding ∗ by simp
    let ?step = quadratic-hensel-step q S T D H
    obtain S2 T2 D2 H2 where step-res: ?step = (S2, T2, D2, H2) by (cases
?step, auto)
    note step = quadratic-hensel-step[OF step-res ∗(1,2,6−9,3) q refl]
    let ?qq = q ∗ q
    {
      fix D D2
      assume poly-mod.Mp q D = poly-mod.Mp q D2
        from arg-cong[OF this, of Mp] Mp-Mp-pow-is-Mp[of ?j2, OF - p, folded
∗(10)] lt
      have Mp D = Mp D2 by simp
    } note shrink = this
    have ∗∗: poly-mod.eq-m ?qq C (D2 ∗ H2)
      poly-mod.eq-m ?qq (D2 ∗ S2 + H2 ∗ T2) 1
      monic D2
      eq-m D1 D2
      eq-m H1 H2
      poly-mod.Mp ?qq D2 = D2
      poly-mod.Mp ?qq H2 = H2
      poly-mod.Mp ?qq S2 = S2
      poly-mod.Mp ?qq T2 = T2
      using step shrink[of H H2] shrink[of D D2] ∗(4−7) by auto
    note simp = simp False if-False rec split Let-def step-res option.simps
    from True have j: p ^ j = p ^ (2 ∗ ?j2) by auto
    with ∗(10) have qq: q ∗ q = p ^ j
      by (simp add: power-mult-distrib semiring-normalization-rules(30−))
    from res[unfolded simp] True have id': q' = ?qq S' = S2 T' = T2 D' = D2
H' = H2 by auto
    show ?thesis unfolding id' using ∗∗ by (auto simp: qq)
  next
```

case *odd*: *False*
hence *False'*: (*even j*) = *False* **by** *auto*
let *?j2* = *j div 2 + 1*
from *False odd* **have** *lt*: *?j2* < *j 1* ≤ *?j2* **by** *presburger+*
obtain *q S T D H* **where** *rec*: *?hens ?j2* = (*q*, *S*, *T*, *D*, *H*) **by** (*cases ?hens ?j2*, *auto*)
note *IH* = *less(1)*[*OF lt rec*]
note *simp* = *simp False if-False rec sum.simps split Let-def False' option.simps*
from *IH* **have** ∗: *poly-mod.eq-m q C* (*D* ∗ *H*)
poly-mod.eq-m q (*D* ∗ *S* + *H* ∗ *T*) *1*
monic D
eq-m D1 D
eq-m H1 H
poly-mod.Mp q D = *D*
poly-mod.Mp q H = *H*
poly-mod.Mp q S = *S*
poly-mod.Mp q T = *T*
q = *p* ^ *?j2*
by *auto*
hence *norm*: *poly-mod.Mp* (*p* ^ *j*) *D* = *D poly-mod.Mp* (*p* ^ *j*) *H* = *H*
using *lift-norm*[*OF lt(2)*] *lt* **by** *auto*
from *lt p* **have** *q*: *q* > *1* **unfolding** ∗
using *mod-2 poly-mod-2.m1* **by** *blast*
let *?step* = *quadratic-hensel-step q S T D H*
obtain *S2 T2 D2 H2* **where** *step-res*: *?step* = (*S2*, *T2*, *D2*, *H2*) **by** (*cases ?step*, *auto*)
have *dvd*: *q dvd q* **by** *auto*
note *step* = *quadratic-hensel-step*[*OF step-res* ∗(*1*,*2*,*6*−*9*,*3*) *q refl*]
let *?qq* = *q* ∗ *q*
{
fix *D D2*
assume *poly-mod.Mp q D* = *poly-mod.Mp q D2*
from *arg-cong*[*OF this*, *of Mp*] *Mp-Mp-pow-is-Mp*[*of ?j2*, *OF - p*, *folded* ∗(*10*)] *lt*
have *Mp D* = *Mp D2* **by** *simp*
} note *shrink* = *this*
have ∗∗: *poly-mod.eq-m ?qq C* (*D2* ∗ *H2*)
poly-mod.eq-m ?qq (*D2* ∗ *S2* + *H2* ∗ *T2*) *1*
monic D2
eq-m D1 D2
eq-m H1 H2
poly-mod.Mp ?qq D2 = *D2*
poly-mod.Mp ?qq H2 = *H2*
poly-mod.Mp ?qq S2 = *S2*
poly-mod.Mp ?qq T2 = *T2*
using *step shrink*[*of H H2*] *shrink*[*of D D2*] ∗(*4*−*7*) **by** *auto*
note *simp* = *simp False if-False rec split Let-def step-res option.simps*
from *odd* **have** *j*: *Suc j* = *2* ∗ *?j2* **by** *auto*
from *arg-cong*[*OF this*, *of* λ *j*. *p* ^ *j div p*]

```
      have pj: p ^ j = q * q div p and qq: q * q = p ^ j * p unfolding *(10)
using p
       by (simp add: power-mult-distrib semiring-normalization-rules(30-))+
     let ?pj = p ^ j
     from res[unfolded simp] pj
     have id:
       q' = p^j
       S' = poly-mod.Mp ?pj S2
       T' = poly-mod.Mp ?pj T2
       D' = poly-mod.Mp ?pj D2
       H' = poly-mod.Mp ?pj H2
       by auto
     interpret pj: poly-mod-2 ?pj by (rule mod-2[OF ‹1 ≤ j›])
     have norm: pj.Mp D' = D' pj.Mp H' = H'
       unfolding id by (auto simp: poly-mod.Mp-Mp)
     have mon: monic D' using pj.monic-Mp[OF step(11)] unfolding id .
     have id': Mp (pj.Mp D) = Mp D for D using ‹1 ≤ j›
       by (simp add: Mp-Mp-pow-is-Mp p)
     have eq: eq-m D1 D2 ⟹ eq-m D1 (pj.Mp D2) for D1 D2
       unfolding id' by auto
     have id'': pj.Mp (poly-mod.Mp (q * q) D) = pj.Mp D for D
       unfolding qq by (rule pj.Mp-product-modulus[OF refl], insert p, auto)
     {
       fix D1 D2
       assume poly-mod.eq-m (q * q) D1 D2
       hence poly-mod.Mp (q * q) D1 = poly-mod.Mp (q * q) D2 by simp
       from arg-cong[OF this, of pj.Mp]
       have pj.Mp D1 = pj.Mp D2 unfolding id'' .
     } note eq' = this
     from eq'[OF step(1)] have eq1: pj.eq-m C (D' * H') unfolding id by simp
     from eq'[OF step(2)] have eq2: pj.eq-m (D' * S' + H' * T') 1
       unfolding id by (metis pj.mult-Mp pj.plus-Mp)
     from **(4-5) have eq3: eq-m D1 D' eq-m H1 H'
       unfolding id by (auto intro: eq)
     from norm mon eq1 eq2 eq3
     show ?thesis unfolding id by simp
   qed
 qed
qed

lemma quadratic-hensel-main: assumes res: quadratic-hensel-main j = (D,H)
 shows poly-mod.eq-m (p^j) C (D * H)
 monic D
 poly-mod.eq-m p D1 D
 poly-mod.eq-m p H1 H
 poly-mod.Mp (p^j) D = D
 poly-mod.Mp (p^j) H = H
proof (atomize(full), goal-cases)
 case 1
```

```
  let ?hen = quadratic-hensel-loop j
  from res obtain q S T where hen: ?hen = (q, S, T, D, H)
    by (cases ?hen, auto simp: quadratic-hensel-main-def)
  from quadratic-hensel-loop[OF hen] show ?case by auto
qed
end
end
end

datatype 'a factor-tree = Factor-Leaf 'a int poly | Factor-Node 'a 'a factor-tree
'a factor-tree

fun factor-node-info :: 'a factor-tree ⇒ 'a where
  factor-node-info (Factor-Leaf i x) = i
| factor-node-info (Factor-Node i l r) = i

fun factors-of-factor-tree :: 'a factor-tree ⇒ int poly multiset where
  factors-of-factor-tree (Factor-Leaf i x) = {#x#}
| factors-of-factor-tree (Factor-Node i l r) = factors-of-factor-tree l + factors-of-factor-tree
r

fun product-factor-tree :: int ⇒ 'a factor-tree ⇒ int poly factor-tree where
  product-factor-tree p (Factor-Leaf i x) = (Factor-Leaf x x)
| product-factor-tree p (Factor-Node i l r) = (let
    L = product-factor-tree p l;
    R = product-factor-tree p r;
    f = factor-node-info L;
    g = factor-node-info R;
    fg = poly-mod.Mp p (f * g)
   in Factor-Node fg L R)

fun sub-trees :: 'a factor-tree ⇒ 'a factor-tree set where
  sub-trees (Factor-Leaf i x) = {Factor-Leaf i x}
| sub-trees (Factor-Node i l r) = insert (Factor-Node i l r) (sub-trees l ∪ sub-trees
r)

lemma sub-trees-refl[simp]: t ∈ sub-trees t by (cases t, auto)

lemma product-factor-tree: assumes ⋀ x. x ∈# factors-of-factor-tree t ⟹ poly-mod.Mp
p x = x
  shows u ∈ sub-trees (product-factor-tree p t) ⟹ factor-node-info u = f ⟹
  poly-mod.Mp p f = f ∧ f = poly-mod.Mp p (prod-mset (factors-of-factor-tree u))
∧
  factors-of-factor-tree (product-factor-tree p t) = factors-of-factor-tree t
  using assms
proof (induct t arbitrary: u f)
  case (Factor-Node i l r u f)
  interpret poly-mod p .
  let ?L = product-factor-tree p l
```

```
  let ?R = product-factor-tree p r
  let ?f = factor-node-info ?L
  let ?g = factor-node-info ?R
  let ?fg = Mp (?f * ?g)
  have Mp ?f = ?f ∧ ?f = Mp (prod-mset (factors-of-factor-tree ?L)) ∧
      (factors-of-factor-tree ?L) = (factors-of-factor-tree l)
      by (rule Factor-Node(1)[OF sub-trees-refl refl], insert Factor-Node(5), auto)
  hence IH1: ?f = Mp (prod-mset (factors-of-factor-tree ?L))
      (factors-of-factor-tree ?L) = (factors-of-factor-tree l) by blast+
  have Mp ?g = ?g ∧ ?g = Mp (prod-mset (factors-of-factor-tree ?R)) ∧
      (factors-of-factor-tree ?R) = (factors-of-factor-tree r)
      by (rule Factor-Node(2)[OF sub-trees-refl refl], insert Factor-Node(5), auto)
  hence IH2: ?g = Mp (prod-mset (factors-of-factor-tree ?R))
      (factors-of-factor-tree ?R) = (factors-of-factor-tree r) by blast+
  have id: (factors-of-factor-tree (product-factor-tree p (Factor-Node i l r))) =
    (factors-of-factor-tree (Factor-Node i l r)) by (simp add: Let-def IH1 IH2)
  from Factor-Node(3) consider (root) u = Factor-Node ?fg ?L ?R
    | (l) u ∈ sub-trees ?L | (r) u ∈ sub-trees ?R
    by (auto simp: Let-def)
  thus ?case
  proof cases
    case root
    with Factor-Node have f: f = ?fg by auto
    show ?thesis unfolding f root id by (simp add: Let-def ac-simps IH1 IH2)
  next
    case l
    have Mp f = f ∧ f = Mp (prod-mset (factors-of-factor-tree u))
      using Factor-Node(1)[OF l Factor-Node(4)] Factor-Node(5) by auto
    thus ?thesis unfolding id by blast
  next
    case r
    have Mp f = f ∧ f = Mp (prod-mset (factors-of-factor-tree u))
      using Factor-Node(2)[OF r Factor-Node(4)] Factor-Node(5) by auto
    thus ?thesis unfolding id by blast
  qed
qed auto

fun create-factor-tree-simple :: int poly list ⇒ unit factor-tree where
  create-factor-tree-simple xs = (let n = length xs in if n ≤ 1 then Factor-Leaf ()
(hd xs)
    else let i = n div 2;
      xs1 = take i xs;
      xs2 = drop i xs
     in Factor-Node () (create-factor-tree-simple xs1) (create-factor-tree-simple xs2)
      )

declare create-factor-tree-simple.simps[simp del]

lemma create-factor-tree-simple: xs ≠ [] ⟹ factors-of-factor-tree (create-factor-tree-simple
```

```
xs) = mset xs
proof (induct xs rule: wf-induct[OF wf-measure[of length]])
  case (1 xs)
  from 1(2) have xs: length xs ≠ 0 by auto
  then consider (base) length xs = 1 | (step) length xs > 1 by linarith
  thus ?case
  proof cases
    case base
    then obtain x where xs: xs = [x] by (cases xs; cases tl xs; auto)
    thus ?thesis by (auto simp: create-factor-tree-simple.simps)
  next
    case step
    let ?i = length xs div 2
    let ?xs1 = take ?i xs
    let ?xs2 = drop ?i xs
    from step have xs1: (?xs1, xs) ∈ measure length ?xs1 ≠ [] by auto
    from step have xs2: (?xs2, xs) ∈ measure length ?xs2 ≠ [] by auto
   from step have id: create-factor-tree-simple xs = Factor-Node () (create-factor-tree-simple
(take ?i xs))
        (create-factor-tree-simple (drop ?i xs)) unfolding create-factor-tree-simple.simps[of
xs] Let-def by auto
    have xs: xs = ?xs1 @ ?xs2 by auto
    show ?thesis unfolding id arg-cong[OF xs, of mset] mset-append
      using 1(1)[rule-format, OF xs1] 1(1)[rule-format, OF xs2]
      by auto
  qed
qed
```

We define a better factorization tree which balances the trees according to their degree., cf. Modern Computer Algebra, Chapter 15.5 on Multifactor Hensel lifting.

```
fun partition-factors-main :: nat ⇒ ('a × nat) list ⇒ ('a × nat) list × ('a × nat)
list where
  partition-factors-main s [] = ([], [])
| partition-factors-main s ((f,d) # xs) = (if d ≤ s then case partition-factors-main
(s − d) xs of
    (l,r) ⇒ ((f,d) # l, r) else case partition-factors-main d xs of
    (l,r) ⇒ (l, (f,d) # r))

lemma partition-factors-main: partition-factors-main s xs = (a,b) ⟹ mset xs =
mset a + mset b
  by (induct s xs arbitrary: a b rule: partition-factors-main.induct, auto split:
if-splits prod.splits)

definition partition-factors :: ('a × nat) list ⇒ ('a × nat) list × ('a × nat) list
where
  partition-factors xs = (let n = sum-list (map snd xs) div 2 in
    case partition-factors-main n xs of
    ([], x # y # ys) ⇒ ([x], y # ys)
```

```
  | (x # y # ys, []) ⇒ ([x], y # ys)
  | pair ⇒ pair)

lemma partition-factors: partition-factors xs = (a,b) ⟹ mset xs = mset a + mset
b
  unfolding partition-factors-def Let-def
   by (cases partition-factors-main (sum-list (map snd xs) div 2) xs, auto split:
list.splits
    simp: partition-factors-main)

lemma partition-factors-length: assumes ¬ length xs ≤ 1 (a,b) = partition-factors
xs
  shows [termination-simp]: length a < length xs length b < length xs and a ≠ []
b ≠ []
proof −
  obtain ys zs where main: partition-factors-main (sum-list (map snd xs) div 2)
xs = (ys,zs) by force
  note res = assms(2)[unfolded partition-factors-def Let-def main split]
  from arg-cong[OF partition-factors-main[OF main], of size] have len: length xs
= length ys + length zs by auto
  with assms(1) have len2: length ys + length zs ≥ 2 by auto
  from res len2 have length a < length xs ∧ length b < length xs ∧ a ≠ [] ∧ b ≠
[] unfolding len
    by (cases ys; cases zs; cases tl ys; cases tl zs; auto)
  thus length a < length xs length b < length xs a ≠ [] b ≠ [] by blast+
qed

fun create-factor-tree-balanced :: (int poly × nat)list ⇒ unit factor-tree where
  create-factor-tree-balanced xs = (if length xs ≤ 1 then Factor-Leaf () (fst (hd xs))
else
     case partition-factors xs of (l,r) ⇒ Factor-Node ()
      (create-factor-tree-balanced l)
      (create-factor-tree-balanced r))

definition create-factor-tree :: int poly list ⇒ unit factor-tree where
  create-factor-tree xs = (let ys = map (λ f. (f, degree f)) xs;
     zs = rev (sort-key snd ys)
     in create-factor-tree-balanced zs)

lemma create-factor-tree-balanced: xs ≠ [] ⟹ factors-of-factor-tree (create-factor-tree-balanced
xs) = mset (map fst xs)
proof (induct xs rule: create-factor-tree-balanced.induct)
  case (1 xs)
  show ?case
  proof (cases length xs ≤ 1)
    case True
    with 1(3) obtain x where xs: xs = [x] by (cases xs; cases tl xs, auto)
    show ?thesis unfolding xs by auto
  next
```

```
    case False
    obtain a b where part: partition-factors xs = (a,b) by force
    note abp = this[symmetric]
    note nonempty = partition-factors-length(3-4)[OF False abp]
    note IH = 1(1)[OF False abp nonempty(1)] 1(2)[OF False abp nonempty(2)]
   show ?thesis unfolding create-factor-tree-balanced.simps[of xs] part split using

      False IH partition-factors[OF part] by auto
  qed
qed

lemma create-factor-tree: assumes xs ≠ []
  shows factors-of-factor-tree (create-factor-tree xs) = mset xs
proof -
  let ?xs = rev (sort-key snd (map (λf. (f, degree f)) xs))
  from assms have set xs ≠ {} by auto
  hence set ?xs ≠ {} by auto
  hence xs: ?xs ≠ [] by blast
 show ?thesis unfolding create-factor-tree-def Let-def create-factor-tree-balanced[OF
xs]
    by (auto, induct xs, auto)
qed

context
  fixes p :: int and n :: nat
begin

definition quadratic-hensel-binary :: int poly ⇒ int poly ⇒ int poly ⇒ int poly ×
int poly where
  quadratic-hensel-binary C D H = (
     case euclid-ext-poly-dynamic p D H of
      (S,T) ⇒ quadratic-hensel-main C p S T D H n)

fun hensel-lifting-main :: int poly ⇒ int poly factor-tree ⇒ int poly list where
  hensel-lifting-main U (Factor-Leaf - -) = [U]
| hensel-lifting-main U (Factor-Node - l r) = (let
    v = factor-node-info l;
    w = factor-node-info r;
    (V,W) = quadratic-hensel-binary U v w
    in hensel-lifting-main V l @ hensel-lifting-main W r)

definition hensel-lifting-monic :: int poly ⇒ int poly list ⇒ int poly list where
  hensel-lifting-monic u vs = (if vs = [] then [] else let
     pn = p^n;
     C = poly-mod.Mp pn u;
     tree = product-factor-tree p (create-factor-tree vs)
     in hensel-lifting-main C tree)

definition hensel-lifting :: int poly ⇒ int poly list ⇒ int poly list where
```

hensel-lifting f gs = (*let lc* = *lead-coeff f*;
ilc = *inverse-mod lc* (*p*^*n*);
g = *smult ilc f*
in hensel-lifting-monic g gs)

end

context *poly-mod-prime* **begin**

context
fixes *n* :: *nat*
assumes *n*: *n* ≠ *0*
begin

abbreviation *hensel-binary* ≡ *quadratic-hensel-binary p n*

abbreviation *hensel-main* ≡ *hensel-lifting-main p n*

lemma *hensel-binary*:
assumes *cop*: *coprime-m D H* **and** *eq*: *eq-m C* (*D* ∗ *H*)
and *normalized-input*: *Mp D* = *D Mp H* = *H*
and *monic-input*: *monic D*
and *hensel-result*: *hensel-binary C D H* = (*D*′,*H*′)
shows *poly-mod.eq-m* (*p*^*n*) *C* (*D*′ ∗ *H*′) — the main result: equivalence mod *p*^*n*
∧ *monic D*′ — monic output
∧ *eq-m D D*′ ∧ *eq-m H H*′ — apply ‘*mod p*‘ on *D*′ and *H*′ yields *D* and *H* again
∧ *poly-mod.Mp* (*p*^*n*) *D*′ = *D*′ ∧ *poly-mod.Mp* (*p*^*n*) *H*′ = *H*′ — output is normalized
proof −
from *m1* **have** *p*: *p* > *1* **.**
obtain *S T* **where** *ext*: *euclid-ext-poly-dynamic p D H* = (*S*,*T*) **by** *force*
obtain *D1 H1* **where** *main*: *quadratic-hensel-main C p S T D H n* = (*D1*,*H1*) **by** *force*
note *hen* = *hensel-result*[*unfolded quadratic-hensel-binary-def ext split Let-def main*]
from *n* **have** *n*: *n* ≥ *1* **by** *simp*
note *eucl* = *euclid-ext-poly-dynamic*[*OF cop normalized-input ext*]
note *main* = *quadratic-hensel-main*[*OF eucl*(*1*) *eq monic-input p normalized-input eucl*(*2*−) *n main*]
show *?thesis* **using** *hen main* **by** *auto*
qed

lemma *hensel-main*:
assumes *eq*: *eq-m C* (*prod-mset* (*factors-of-factor-tree Fs*))
and ⋀ *F*. *F* ∈# *factors-of-factor-tree Fs* ⟹ *Mp F* = *F* ∧ *monic F*
and *hensel-result*: *hensel-main C Fs* = *Gs*
and *C*: *monic C poly-mod.Mp* (*p*^*n*) *C* = *C*

and *sf*: *square-free-m C*
and $\bigwedge f\ t.\ t \in$ *sub-trees Fs* $\Longrightarrow$ *factor-node-info t* $= f \Longrightarrow f =$ *Mp* (*prod-mset* (*factors-of-factor-tree t*))
shows *poly-mod.eq-m* (*p*^*n*) *C* (*prod-list Gs*) — the main result: equivalence mod *p*^*n*
$\wedge$ *factors-of-factor-tree Fs* = *mset* (*map Mp Gs*)
$\wedge$ ($\forall$ *G*. *G* $\in$ *set Gs* $\longrightarrow$ *monic G* $\wedge$ *poly-mod.Mp* (*p*^*n*) *G* = *G*)
using *assms*
proof (*induct Fs arbitrary*: *C Gs*)
case (*Factor-Leaf f fs C Gs*)
thus *?case* **by** *auto*
next
case (*Factor-Node f l r C Gs*) **note** $*$ = *this*
note *simps* = *hensel-lifting-main.simps*
note *IH1* = $*(1)$[*rule-format*]
note *IH2* = $*(2)$[*rule-format*]
note *res* = $*(5)$[*unfolded simps Let-def*]
note *eq* = $*(3)$
note *Fs* = $*(4)$
note *C* = $*(6,7)$
note *sf* = $*(8)$
note *inv* = $*(9)$
interpret *pn*: *poly-mod-2 p*^*n* **apply** (*unfold-locales*) **using** *m1 n* **by** *auto*
let *?Mp* = *pn.Mp*
define *D* **where** *D* $\equiv$ *prod-mset* (*factors-of-factor-tree l*)
define *H* **where** *H* $\equiv$ *prod-mset* (*factors-of-factor-tree r*)
let *?D* = *Mp D*
let *?H* = *Mp H*
let *?D'* = *factor-node-info l*
let *?H'* = *factor-node-info r*
obtain *A B* **where** *hen*: *hensel-binary C ?D' ?H'* = (*A*,*B*) **by** *force*
note *res* = *res*[*unfolded hen split*]
obtain *AD* **where** *AD'*: *AD* = *hensel-main A l* **by** *auto*
obtain *BH* **where** *BH'*: *BH* = *hensel-main B r* **by** *auto*
from *inv*[*of l*, *OF - refl*] **have** *D'*: *?D'* = *?D* **unfolding** *D-def* **by** *auto*
from *inv*[*of r*, *OF - refl*] **have** *H'*: *?H'* = *?H* **unfolding** *H-def* **by** *auto*
from *eq*[*simplified*]
have *eq'*: *Mp C* = *Mp* (*?D* $*$ *?H*) **unfolding** *D-def H-def* **by** *simp*
from *square-free-m-cong*[*OF sf*, *of ?D* $*$ *?H*, *OF eq'*]
have *sf'*: *square-free-m* (*?D* $*$ *?H*) **.**
from *poly-mod-prime.square-free-m-prod-imp-coprime-m*[*OF - this*]
have *cop'*: *coprime-m ?D ?H* **unfolding** *poly-mod-prime-def* **using** *prime* **.**
from *eq'* **have** *eq'*: *eq-m C* (*?D* $*$ *?H*) **by** *simp*
have *monD*: *monic D* **unfolding** *D-def* **by** (*rule monic-prod-mset*, *insert Fs*, *auto*)
from *hensel-binary*[*OF - - - - - hen*, *unfolded D' H'*, *OF cop' eq' Mp-Mp Mp-Mp monic-Mp*[*OF monD*]]
have *step*: *poly-mod.eq-m* (*p* ^ *n*) *C* (*A* $*$ *B*) $\wedge$ *monic A* $\wedge$ *eq-m ?D A* $\wedge$
eq-m ?H B $\wedge$ *?Mp A* = *A* $\wedge$ *?Mp B* = *B* **.**

from *res* **have** *Gs*: $Gs = AD$ @ BH **by** (*simp add*: AD' BH')
have *AD*: *eq-m A ?D ?Mp A* = *A eq-m A* (*prod-mset* (*factors-of-factor-tree l*))
and *monA*: *monic A*
using *step* **by** (*auto simp*: *D-def*)
note *sf-fact* = *square-free-m-factor*[*OF sf'*]
from *square-free-m-cong*[*OF sf-fact*(*1*)] *AD* **have** *sfA*: *square-free-m A* **by** *auto*
have *IH1*: *poly-mod.eq-m* (*p* ^ *n*) *A* (*prod-list AD*) ∧
factors-of-factor-tree l = *mset* (*map Mp AD*) ∧
(∀ *G*. *G* ∈ *set AD* ⟶ *monic G* ∧ *?Mp G* = *G*)
by (*rule IH1*[*OF AD*(*3*) *Fs AD'*[*symmetric*] *monA AD*(*2*) *sfA inv*], *auto*)
have *BH*: *eq-m B ?H pn.Mp B* = *B eq-m B* (*prod-mset* (*factors-of-factor-tree r*))
using *step* **by** (*auto simp*: *H-def*)
from *step* **have** *pn.eq-m C* (*A* ∗ *B*) **by** *simp*
hence *?Mp C* = *?Mp* (*A* ∗ *B*) **by** *simp*
with *C AD*(*2*) **have** *pn.Mp C* = *pn.Mp* (*A* ∗ *pn.Mp B*) **by** *simp*
from *arg-cong*[*OF this*, *of lead-coeff*] *C*
have *monic* (*pn.Mp* (*A* ∗ *B*)) **by** *simp*
then have *lead-coeff* (*pn.Mp A*) ∗ *lead-coeff* (*pn.Mp B*) = *1*
by (*metis lead-coeff-mult leading-coeff-neq-0 local.step mult-cancel-right2 pn.degree-m-eq pn.m1 poly-mod.M-def poly-mod.Mp-coeff*)
with *monA AD*(*2*) *BH*(*2*) **have** *monB*: *monic B* **by** *simp*
from *square-free-m-cong*[*OF sf-fact*(*2*)] *BH* **have** *sfB*: *square-free-m B* **by** *auto*
have *IH2*: *poly-mod.eq-m* (*p* ^ *n*) *B* (*prod-list BH*) ∧
factors-of-factor-tree r = *mset* (*map Mp BH*) ∧
(∀ *G*. *G* ∈ *set BH* ⟶ *monic G* ∧ *?Mp G* = *G*)
by (*rule IH2*[*OF BH*(*3*) *Fs BH'*[*symmetric*] *monB BH*(*2*) *sfB inv*], *auto*)
from *step* **have** *?Mp C* = *?Mp* (*?Mp A* ∗ *?Mp B*) **by** *auto*
also have *?Mp A* = *?Mp* (*prod-list AD*) **using** *IH1* **by** *auto*
also have *?Mp B* = *?Mp* (*prod-list BH*) **using** *IH2* **by** *auto*
finally have *poly-mod.eq-m* (*p* ^ *n*) *C* (*prod-list AD* ∗ *prod-list BH*)
by (*auto simp*: *poly-mod.mult-Mp*)
thus *?case* **unfolding** *Gs* **using** *IH1 IH2* **by** *auto*
qed

lemma *hensel-lifting-monic*:
assumes *eq*: *poly-mod.eq-m p C* (*prod-list Fs*)
and *Fs*: ⋀ *F*. *F* ∈ *set Fs* ⟹ *poly-mod.Mp p F* = *F* ∧ *monic F*
and *res*: *hensel-lifting-monic p n C Fs* = *Gs*
and *mon*: *monic* (*poly-mod.Mp* (*p*^*n*) *C*)
and *sf*: *poly-mod.square-free-m p C*
shows *poly-mod.eq-m* (*p*^*n*) *C* (*prod-list Gs*)
mset (*map* (*poly-mod.Mp p*) *Gs*) = *mset Fs*
G ∈ *set Gs* ⟹ *monic G* ∧ *poly-mod.Mp* (*p*^*n*) *G* = *G*
proof −
note *res* = *res*[*unfolded hensel-lifting-monic-def Let-def*]
let *?Mp* = *poly-mod.Mp* (*p* ^ *n*)
let *?C* = *?Mp C*
interpret *poly-mod-prime p*
by (*unfold-locales*, *insert n prime*, *auto*)

interpret *pn*: *poly-mod-2 p^n* **using** *m1 n poly-mod-2.intro* **by** *auto*
from *eq n* **have** *eq*: *eq-m* (*?Mp C*) (*prod-list Fs*)
using *Mp-Mp-pow-is-Mp eq m1 n* **by** *force*
have *poly-mod.eq-m* (*p^n*) *C* (*prod-list Gs*) ∧ *mset* (*map* (*poly-mod.Mp p*) *Gs*) = *mset Fs*
∧ (*G* ∈ *set Gs* ⟶ *monic G* ∧ *poly-mod.Mp* (*p^n*) *G* = *G*)
proof (*cases Fs* = [])
case *True*
with *res* **have** *Gs*: *Gs* = [] **by** *auto*
from *eq* **have** *Mp ?C* = *1* **unfolding** *True* **by** *simp*
hence *degree* (*Mp ?C*) = *0* **by** *simp*
with *degree-m-eq-monic*[*OF mon m1*] **have** *degree ?C* = *0* **by** *simp*
with *mon* **have** *?C* = *1* **using** *monic-degree-0* **by** *blast*
thus *?thesis* **unfolding** *True Gs* **by** *auto*
next
case *False*
let *?t* = *create-factor-tree Fs*
note *tree* = *create-factor-tree*[*OF False*]
from *False res* **have** *hen*: *hensel-main ?C* (*product-factor-tree p ?t*) = *Gs* **by** *auto*
have *tree1*: *x* ∈# *factors-of-factor-tree ?t* ⟹ *Mp x* = *x* **for** *x* **unfolding** *tree* **using** *Fs* **by** *auto*
from *product-factor-tree*[*OF tree1 sub-trees-refl refl*, *of ?t*]
have *id*: (*factors-of-factor-tree* (*product-factor-tree p ?t*)) =
(*factors-of-factor-tree ?t*) **by** *auto*
have *eq*: *eq-m ?C* (*prod-mset* (*factors-of-factor-tree* (*product-factor-tree p ?t*)))
unfolding *id tree* **using** *eq* **by** *auto*
have *id'*: *Mp C* = *Mp ?C* **using** *n* **by** (*simp add*: *Mp-Mp-pow-is-Mp m1*)
have *pn.eq-m ?C* (*prod-list Gs*) ∧ *mset Fs* = *mset* (*map Mp Gs*) ∧ (∀ *G*. *G* ∈ *set Gs* ⟶ *monic G* ∧ *pn.Mp G* = *G*)
by (*rule hensel-main*[*OF eq Fs hen mon pn.Mp-Mp square-free-m-cong*[*OF sf id'*], *unfolded id tree*],
insert product-factor-tree[*OF tree1*], *auto*)
thus *?thesis* **by** *auto*
qed
thus *poly-mod.eq-m* (*p^n*) *C* (*prod-list Gs*)
mset (*map* (*poly-mod.Mp p*) *Gs*) = *mset Fs*
G ∈ *set Gs* ⟹ *monic G* ∧ *poly-mod.Mp* (*p^n*) *G* = *G* **by** *blast*+
qed

lemma *hensel-lifting*:
assumes *res*: *hensel-lifting p n f fs* = *gs* — result of hensel is fact. *gs*
and *cop*: *coprime* (*lead-coeff f*) *p*
and *sf*: *poly-mod.square-free-m p f*
and *fact*: *poly-mod.factorization-m p f* (*c*, *mset fs*) — input is fact. *fs mod p*
and *c*: *c* ∈ {*0*..<*p*}
and *norm*: (∀ *fi*∈*set fs*. *set* (*coeffs fi*) ⊆ {*0*..<*p*})

shows *poly-mod.factorization-m* $(p\hat{}n)$ f $(lead\text{-}coeff\ f,\ mset\ gs)$ — factorization mod $p\hat{}n$
$sort\ (map\ degree\ fs) = sort\ (map\ degree\ gs)$ — degrees stay the same
$\bigwedge g.\ g \in set\ gs \Longrightarrow monic\ g \wedge poly\text{-}mod.Mp\ (p\hat{}n)\ g = g \wedge$ — monic and normalized
$irreducible\text{-}m\ g \wedge$ — irreducibility even mod p
$degree\text{-}m\ g = degree\ g$ — mod p does not change degree of g
proof −
interpret *poly-mod-prime p* **using** *prime* **by** *unfold-locales*
interpret *q*: *poly-mod-2* $p\hat{}n$ **using** *m1 n* **unfolding** *poly-mod-2-def* **by** *auto*
from *fact* **have** *eq*: *eq-m f* (*smult c* (*prod-list fs*))
and *mon-fs*: $(\forall fi \in set\ fs.\ monic\ (Mp\ fi) \wedge irreducible_d\text{-}m\ fi)$
unfolding *factorization-m-def* **by** *auto*
{
fix *f*
assume $f \in set\ fs$
with *mon-fs norm* **have** $set\ (coeffs\ f) \subseteq \{0..<p\}$ **and** *monic* (*Mp f*) **by** *auto*
hence *monic f* **using** *Mp-ident-iff'* **by** *force*
} note $mon\text{-}fs' = this$
have *Mp-id*: $\bigwedge f.\ Mp\ (q.Mp\ f) = Mp\ f$ **by** (*simp add*: *Mp-Mp-pow-is-Mp m1 n*)
let *?lc* = *lead-coeff f*
let $?q = p\ \hat{}\ n$
define *ilc* **where** $ilc \equiv inverse\text{-}mod\ ?lc\ ?q$
define *F* **where** $F \equiv smult\ ilc\ f$
from *res*[*unfolded hensel-lifting-def Let-def*]
have *hen*: $hensel\text{-}lifting\text{-}monic\ p\ n\ F\ fs = gs$
unfolding *ilc-def F-def* **.**
from *m1 n cop* **have** *inv*: $q.M\ (ilc * ?lc) = 1$
by (*auto simp add*: *q.M-def inverse-mod-pow ilc-def*)
hence *ilc0*: $ilc \neq 0$ **by** (*cases ilc* = *0*, *auto*)
{
fix *q*
assume $ilc * ?lc = ?q * q$
from *arg-cong*[*OF this*, *of q.M*] **have** $q.M\ (ilc * ?lc) = 0$
unfolding *q.M-def* **by** *auto*
with *inv* **have** *False* **by** *auto*
} note *not-dvd* = *this*
have *mon*: *monic* (*q.Mp F*) **unfolding** *F-def q.Mp-coeff coeff-smult*
by (*subst q.degree-m-eq* [*OF - q.m1*]) (*auto simp*: *inv ilc0* [*symmetric*] *intro*: *not-dvd*)
have $q.Mp\ f = q.Mp\ (smult\ (q.M\ (?lc * ilc))\ f)$ **using** *inv* **by** (*simp add*: *ac-simps*)
also have ... = $q.Mp\ (smult\ ?lc\ F)$ **by** (*simp add*: *F-def*)
finally have *f*: $q.Mp\ f = q.Mp\ (smult\ ?lc\ F)$ **.**
from *arg-cong*[*OF f*, *of Mp*]
have *f-p*: $Mp\ f = Mp\ (smult\ ?lc\ F)$
by (*simp add*: *Mp-Mp-pow-is-Mp n m1*)
from *arg-cong*[*OF this*, *of square-free-m*, *unfolded Mp-square-free-m*] *sf*

have *square-free-m* (*smult ?lc F*) **by** *simp*
from *square-free-m-smultD*[*OF this*] **have** *sf*: *square-free-m F* **.**

define c' **where** $c' \equiv M\ (c * ilc)$
from *factorization-m-smult*[*OF fact*, *of ilc*, *folded F-def*]
have *fact*: *factorization-m F* (c', *mset fs*) **unfolding** c'*-def factorization-m-def*
by *auto*
hence *eq*: *eq-m F* (*smult* c' (*prod-list fs*)) **unfolding** *factorization-m-def* **by** *auto*
from *factorization-m-lead-coeff*[*OF fact*] *monic-Mp*[*OF mon*, *unfolded Mp-id*]
have $M\ c' = 1$
by *auto*
hence c': $c' = 1$ **unfolding** c'*-def* **by** *auto*
with *eq* **have** *eq*: *eq-m F* (*prod-list fs*) **by** *auto*
{
fix *f*
assume $f \in$ *set fs*
with *mon-fs' norm* **have** *Mp f* $= f \wedge$ *monic f* **unfolding** *Mp-ident-iff'*
by *auto*
} note *fs* $=$ *this*
note *hen* $=$ *hensel-lifting-monic*[*OF eq fs hen mon sf*]
from *hen*(*2*) **have** *gs-fs*: *mset* (*map Mp gs*) $=$ *mset fs* **by** *auto*
have *eq*: *q.eq-m f* (*smult ?lc* (*prod-list gs*))
unfolding *f* **using** *arg-cong*[*OF hen*(*1*), *of* λ *f*. *q.Mp* (*smult ?lc f*)] **by** *simp*
{
fix *g*
assume *g*: $g \in$ *set gs*
from *hen*(*3*)[*OF - g*] **have** *mon-g*: *monic g* **and** *Mp-g*: *q.Mp g* $= g$ **by** *auto*
from *g* **have** *Mp g* $\in\#$ *mset* (*map Mp gs*) **by** *auto*
from *this*[*unfolded gs-fs*] **obtain** *f* **where** *f*: $f \in$ *set fs* **and** *fg*: *eq-m f g* **by**
auto
from *mon-fs f fs* **have** *irr-f*: *irreducible*$_d$*-m f* **and** *mon-f*: *monic f* **and** *Mp-f*:
Mp f $= f$ **by** *auto*
have *deg*: *degree-m g* $=$ *degree g*
by (*rule degree-m-eq-monic*[*OF mon-g m1*])
from *irr-f fg* **have** *irr-g*: *irreducible*$_d$*-m g*
unfolding *irreducible*$_d$*-m-def dvdm-def* **by** *simp*
have *q.irreducible*$_d$*-m g*
by (*rule irreducible*$_d$*-lifting*[*OF n - irr-g*], *unfold deg*, *rule q.degree-m-eq-monic*[*OF*
mon-g q.m1])
note *mon-g Mp-g deg irr-g this*
} note *g* $=$ *this*
{
fix *g*
assume *g* $\in$ *set gs*
from *g*[*OF this*]
show *monic g* $\wedge$ *q.Mp g* $= g \wedge$ *irreducible-m g* $\wedge$ *degree-m g* $=$ *degree g* **by**
auto
}
show *sort* (*map degree fs*) $=$ *sort* (*map degree gs*)

proof (*rule sort-key-eq-sort-key*)
have *mset* (*map degree fs*) = *image-mset degree* (*mset fs*) **by** *auto*
also have ... = *image-mset degree* (*mset* (*map Mp gs*)) **unfolding** *gs-fs* **..**
also have ... = *mset* (*map degree* (*map Mp gs*)) **unfolding** *mset-map* **..**
also have *map degree* (*map Mp gs*) = *map degree-m gs* **by** *auto*
also have ... = *map degree gs* **using** *g*(*3*) **by** *auto*
finally show *mset* (*map degree fs*) = *mset* (*map degree gs*) **.**
qed *auto*
show *q.factorization-m f* (*lead-coeff f*, *mset gs*)
using *eq g* **unfolding** *q.factorization-m-def* **by** *auto*
qed

end

end
end

theory *Hensel-Lifting-Type-Based*
imports *Hensel-Lifting*
begin

9.2 Hensel Lifting in a Type-Based Setting

lemma *degree-smult-eq-iff*:
degree (*smult a p*) = *degree p* ⟷ *degree p* = *0* ∨ *a* ∗ *lead-coeff p* ≠ *0*
by (*metis* (*no-types*, *lifting*) *coeff-smult degree-0 degree-smult-le le-antisym le-degree le-zero-eq leading-coeff-0-iff*)

lemma *degree-smult-eqI*[*intro*!]:
assumes *degree p* ≠ *0* ⟹ *a* ∗ *lead-coeff p* ≠ *0*
shows *degree* (*smult a p*) = *degree p*
using *assms degree-smult-eq-iff* **by** *auto*

lemma *degree-mult-eq2*:
assumes *lc*: *lead-coeff p* ∗ *lead-coeff q* ≠ *0*
shows *degree* (*p* ∗ *q*) = *degree p* + *degree q* (**is** - = *?r*)
proof(*intro antisym*[*OF degree-mult-le*] *le-degree*, *unfold coeff-mult*)
let *?f* = λ*i*. *coeff p i* ∗ *coeff q* (*?r* − *i*)
have (∑ *i*≤*?r*. *?f i*) = *sum ?f* {*..degree p*} + *sum ?f* {*Suc* (*degree p*)*..?r*}
by (*rule sum-up-index-split*)
also have *sum ?f* {*Suc* (*degree p*)*..?r*} = *0*
proof−
{ fix *x* **assume** *x* > *degree p*
then have *coeff p x* = *0* **by** (*rule coeff-eq-0*)
then have *?f x* = *0* **by** *auto*
}
then show *?thesis* **by** (*intro sum.neutral*, *auto*)
qed
also have *sum ?f* {*..degree p*} = *sum ?f* {*..*<*degree p*} + *?f* (*degree p*)

```
    by(fold lessThan-Suc-atMost, unfold sum.lessThan-Suc, auto)
  also have sum ?f {..<degree p} = 0
    proof−
      {fix x assume x < degree p
        then have coeff q (?r − x) = 0 by (intro coeff-eq-0, auto)
        then have ?f x = 0 by auto
      }
      then show ?thesis by (intro sum.neutral, auto)
    qed
  finally show (∑ i≤?r. ?f i) ≠ 0 using assms by (auto simp:)
qed

lemma degree-mult-eq-left-unit:
  fixes p q :: 'a :: comm-semiring-1 poly
  assumes unit: lead-coeff p dvd 1 and q0: q ≠ 0
  shows degree (p * q) = degree p + degree q
proof(intro degree-mult-eq2 notI)
  from unit obtain c where lead-coeff p * c = 1 by (elim dvdE,auto)
  then have c * lead-coeff p = 1 by (auto simp: ac-simps)
  moreover assume lead-coeff p * lead-coeff q = 0
    then have c * lead-coeff p * lead-coeff q = 0 by (auto simp: ac-simps)
  ultimately have lead-coeff q = 0 by auto
  with q0 show False by auto
qed

context ring-hom begin
lemma monic-degree-map-poly-hom: monic p ⟹ degree (map-poly hom p) = degree p
  by (auto intro: degree-map-poly)

lemma monic-map-poly-hom: monic p ⟹ monic (map-poly hom p)
  by (simp add: monic-degree-map-poly-hom)

end

lemma of-nat-zero:
  assumes CARD('a::nontriv) dvd n
  shows (of-nat n :: 'a mod-ring) = 0
  apply (transfer fixing: n) using assms by (presburger)

abbreviation rebase :: 'a :: nontriv mod-ring ⇒ 'b :: nontriv mod-ring (‹@-› [100]100)
  where @x ≡ of-int (to-int-mod-ring x)

abbreviation rebase-poly :: 'a :: nontriv mod-ring poly ⇒ 'b :: nontriv mod-ring poly (‹#-› [100]100)
  where #x ≡ of-int-poly (to-int-poly x)

lemma rebase-self [simp]:
```

$@x = x$
by (*simp add*: *of-int-of-int-mod-ring*)

lemma *map-poly-rebase* [*simp*]:
map-poly rebase p = *#p*
by (*induct p*) *simp-all*

lemma *rebase-poly-0*: *#0* = *0*
by *simp*

lemma *rebase-poly-1*: *#1* = *1*
by *simp*

lemma *rebase-poly-pCons*[*simp*]: *#pCons a p* = *pCons* (*@a*) (*#p*)
by(*cases a* = *0* ∧ *p* = *0*, *simp*, *fold map-poly-rebase*, *subst map-poly-pCons*, *auto*)

lemma *rebase-poly-self*[*simp*]: *#p* = *p* **by** (*induct p*, *auto*)

lemma *degree-rebase-poly-le*: *degree* (*#p*) ≤ *degree p*
by (*fold map-poly-rebase*, *subst degree-map-poly-le*, *auto*)

lemma(**in** *comm-ring-hom*) *degree-map-poly-unit*: **assumes** *lead-coeff p dvd 1*
shows *degree* (*map-poly hom p*) = *degree p*
using *hom-dvd-1*[*OF assms*] **by** (*auto intro*: *degree-map-poly*)

lemma *rebase-poly-eq-0-iff*:
(*#p* :: *'a* :: *nontriv mod-ring poly*) = *0* ⟷ (∀ *i*. (*@coeff p i* :: *'a mod-ring*) = *0*) (**is** *?l* ⟷ *?r*)
proof(*intro iffI*)
assume *?l*
then have *coeff* (*#p* :: *'a mod-ring poly*) *i* = *0* **for** *i* **by** *auto*
then show *?r* **by** *auto*
next
assume *?r*
then have *coeff* (*#p* :: *'a mod-ring poly*) *i* = *0* **for** *i* **by** *auto*
then show *?l* **by** (*intro poly-eqI*, *auto*)
qed

lemma *mod-mod-le*:
assumes *ab*: (*a*::*int*) ≤ *b* **and** *a0*: *0* < *a* **and** *c0*: *c* ≥ *0* **shows** (*c mod a*) *mod b* = *c mod a*
by (*meson pos-mod-bound pos-mod-sign a0 ab less-le-trans mod-pos-pos-trivial*)

locale *rebase-ge* =
fixes *ty1* :: *'a* :: *nontriv itself* **and** *ty2* :: *'b* :: *nontriv itself*
assumes *card*: *CARD*(*'a*) ≤ *CARD*(*'b*)
begin

lemma *ab*: *int CARD*(*'a*) ≤ *CARD*(*'b*) **using** *card* **by** *auto*

lemma *rebase-eq-0*[*simp*]:
shows (@(*x* :: *'a mod-ring*) :: *'b mod-ring*) = *0* ⟷ *x* = *0*
using *card* **by** (*transfer*, *auto*)

lemma *degree-rebase-poly-eq*[*simp*]:
shows *degree* (#(*p* :: *'a mod-ring poly*) :: *'b mod-ring poly*) = *degree p*
by (*subst degree-map-poly*; *simp*)

lemma *lead-coeff-rebase-poly*[*simp*]:
lead-coeff (#(*p*::*'a mod-ring poly*) :: *'b mod-ring poly*) = @*lead-coeff p*
by *simp*

lemma *to-int-mod-ring-rebase*: *to-int-mod-ring*(@(*x* :: *'a mod-ring*)::*'b mod-ring*) = *to-int-mod-ring x*
using *card* **by** (*transfer*, *auto*)

lemma *rebase-id*[*simp*]: @(@(*x*::*'a mod-ring*) :: *'b mod-ring*) = @*x*
using *card* **by** (*transfer*, *auto*)

lemma *rebase-poly-id*[*simp*]: #(#(*p*::*'a mod-ring poly*) :: *'b mod-ring poly*) = #*p*
by (*induct p*, *auto*)

end

locale *rebase-dvd* =
fixes *ty1* :: *'a* :: *nontriv itself* **and** *ty2* :: *'b* :: *nontriv itself*
assumes *dvd*: *CARD*(*'b*) *dvd CARD*(*'a*)
begin

lemma *ab*: *CARD*(*'a*) ≥ *CARD*(*'b*) **by** (*rule dvd-imp-le*[*OF dvd*], *auto*)

lemma *rebase-id*[*simp*]: @(@(*x*::*'b mod-ring*) :: *'a mod-ring*) = *x* **using** *ab* **by** (*transfer*, *auto*)

lemma *rebase-poly-id*[*simp*]: #(#(*p*::*'b mod-ring poly*) :: *'a mod-ring poly*) = *p* **by** (*induct p*, *auto*)

lemma *rebase-of-nat*[*simp*]: (@(*of-nat n* :: *'a mod-ring*) :: *'b mod-ring*) = *of-nat n*
apply *transfer* **apply** (*rule mod-mod-cancel*) **using** *dvd* **by** *presburger*

lemma *mod-1-lift-nat*:
assumes (*of-int* (*int x*) :: *'a mod-ring*) = *1*
shows (*of-int* (*int x*) :: *'b mod-ring*) = *1*
proof −
from *assms* **have** *int x mod CARD*(*'a*) = *1*
by *transfer*
then have *x mod CARD*(*'a*) = *1*

by (*simp add*: *of-nat-mod* [*symmetric*])
then have *x mod CARD('b) = 1*
by (*metis dvd mod-mod-cancel one-mod-card*)
then have *int x mod CARD('b) = 1*
by (*simp add*: *of-nat-mod* [*symmetric*])
then show *?thesis*
by *transfer*
qed

sublocale *comm-ring-hom rebase* :: *'a mod-ring* ⇒ *'b mod-ring*
proof
fix *x y* :: *'a mod-ring*
show *hom-add*: (@(*x*+*y*) :: *'b mod-ring*) = @*x* + @*y*
by *transfer* (*simp add*: *mod-simps dvd mod-mod-cancel*)
show (@(*x*∗*y*) :: *'b mod-ring*) = @*x* ∗ @*y*
by *transfer* (*simp add*: *mod-simps dvd mod-mod-cancel*)
qed *auto*

lemma *of-nat-CARD-eq-0*[*simp*]: (*of-nat CARD('a)* :: *'b mod-ring*) = *0*
using *dvd* **by** (*transfer*, *presburger*)

interpretation *map-poly-hom*: *map-poly-comm-ring-hom rebase* :: *'a mod-ring* ⇒ *'b mod-ring*..

sublocale *poly*: *comm-ring-hom rebase-poly* :: *'a mod-ring poly* ⇒ *'b mod-ring poly*
by (*fold map-poly-rebase*, *unfold-locales*)

lemma *poly-rebase*[*simp*]: @*poly p x* = *poly* (#(*p* :: *'a mod-ring poly*) :: *'b mod-ring poly*) (@(*x*::*'a mod-ring*) :: *'b mod-ring*)
by (*fold map-poly-rebase poly-map-poly*, *rule*)

lemma *rebase-poly-smult*[*simp*]: (#(*smult a p* :: *'a mod-ring poly*) :: *'b mod-ring poly*) = *smult* (@*a*) (#*p*)
by(*induct p*, *auto simp*: *hom-distribs*)

end

locale *rebase-mult* =
fixes *ty1* :: *'a* :: *nontriv itself*
and *ty2* :: *'b* :: *nontriv itself*
and *ty3* :: *'d* :: *nontriv itself*
assumes *d*: *CARD('a)* = *CARD('b)* ∗ *CARD('d)*
begin

sublocale *rebase-dvd ty1 ty2* **using** *d* **by** (*unfold-locales*, *auto*)

lemma *rebase-mult-eq*[*simp*]: (*of-nat CARD('d)* ∗ *a* :: *'a mod-ring*) = *of-nat CARD('d)* ∗ *a'* ⟷ (@*a* :: *'b mod-ring*) = @*a'*
proof−

from *dvd* **obtain** d' **where** $CARD('a) = d' * CARD('b)$ **by** (*elim dvdE*, *auto*)
then show *?thesis* **by** (*transfer*, *auto simp*:*d*)
qed

lemma *rebase-poly-smult-eq*[*simp*]:
fixes a a' :: $'a$ *mod-ring poly*
defines $d \equiv$ *of-nat* $CARD('d)$ :: $'a$ *mod-ring*
shows *smult* d a = *smult* d $a' \longleftrightarrow$ $(\#a :: 'b$ *mod-ring poly*$) = \#a'$ (**is** *?l* $\longleftrightarrow$ *?r*)
proof (*intro iffI*)
assume *l*: *?l* **show** *?r*
proof (*intro poly-eqI*)
fix n
from l **have** *coeff* (*smult* d a) n = *coeff* (*smult* d a') n **by** *auto*
then have $d * coeff\ a\ n = d * coeff\ a'\ n$ **by** *auto*
from *this*[*unfolded d-def rebase-mult-eq*]
show *coeff* $(\#a :: 'b$ *mod-ring poly*$)$ n = *coeff* $(\#a')$ n **by** *auto*
qed
next
assume *r*: *?r* **show** *?l*
proof(*intro poly-eqI*)
fix n
from r **have** *coeff* $(\#a :: 'b$ *mod-ring poly*$)$ n = *coeff* $(\#a')$ n **by** *auto*
then have $(@coeff\ a\ n :: 'b$ *mod-ring*$) = @coeff\ a'\ n$ **by** *auto*
from *this*[*folded d-def rebase-mult-eq*]
show *coeff* (*smult* d a) n = *coeff* (*smult* d a') n **by** *auto*
qed
qed

lemma *rebase-eq-0-imp-ex-mult*:
$(@(a :: 'a$ *mod-ring*$) :: 'b$ *mod-ring*$) = 0 \Longrightarrow (\exists\, c :: 'd$ *mod-ring*. $a =$ *of-nat* $CARD('b) * @c)$ (**is** *?l* $\Longrightarrow$ *?r*)
proof(*cases* $CARD('a) = CARD('b)$)
case *True* **then show** *?l* $\Longrightarrow$ *?r*
by (*transfer*, *auto*)
next
case *False*
have [*simp*]: *int* $CARD('b)$ *mod int* $CARD('a)$ = *int* $CARD('b)$
by(*rule mod-pos-pos-trivial*, *insert ab False*, *auto*)
{
fix a
assume a: $0 \le a$ $a <$ *int* $CARD('a)$ **and** *mod*: a *mod int* $CARD('b) = 0$
from *mod* **have** *int* $CARD('b)$ *dvd* a **by** *auto*
then obtain i **where** $*$: $a =$ *int* $CARD('b) * i$ **by** (*elim dvdE*, *auto*)
from $* a$ **have** $i <$ *int* $CARD('d)$ **by** (*simp add*:*d*)
moreover
hence $(i$ *mod int* $CARD('a)) = i$
by (*metis dual-order.order-iff-strict less-le-trans not-le of-nat-less-iff* $* a(1)$ $a(2)$

```
        mod-pos-pos-trivial mult-less-cancel-right1 nat-neq-iff nontriv of-nat-1)
    with * a have a = int CARD('b) * (i mod int CARD('a)) mod int CARD('a)
      by (auto simp:d)
   moreover from * a have 0 ≤ i
    using linordered-semiring-strict-class.mult-pos-neg of-nat-0-less-iff zero-less-card-finite
     by (simp add: zero-le-mult-iff)
   ultimately have ∃ i≥0. i < int CARD('d) ∧ a = int CARD('b) * (i mod int
CARD('a)) mod int CARD('a)
     by (auto intro: exI[of - i])
 }
 then show ?l ⟹ ?r by (transfer, auto simp:d)
qed

lemma rebase-poly-eq-0-imp-ex-smult:
 (#(p :: 'a mod-ring poly) :: 'b mod-ring poly) = 0 ⟹
  (∃ p' :: 'd mod-ring poly. (p = 0 ⟷ p' = 0) ∧ degree p' ≤ degree p ∧ p = smult
(of-nat CARD('b)) (#p'))
 (is ?l ⟹ ?r)
proof(induct p)
 case 0
 then show ?case by (intro exI[of - 0],auto)
next
 case IH: (pCons a p)
 from IH(3) have (#p :: 'b mod-ring poly) = 0 by auto
 from IH(2)[OF this] obtain p' :: 'd mod-ring poly
 where *: p = 0 ⟷ p' = 0 degree p' ≤ degree p p = smult (of-nat CARD('b))
(#p') by (elim exE conjE)
 from IH have (@a :: 'b mod-ring) = 0 by auto
 from rebase-eq-0-imp-ex-mult[OF this]
 obtain a' :: 'd mod-ring where a': of-nat CARD('b) * (@a') = a by auto
 from IH(1) have pCons a p ≠ 0 by auto
 moreover from *(1,2) have degree (pCons a' p') ≤ degree (pCons a p) by auto
 moreover from a' *(3)
 have pCons a p = smult (of-nat CARD('b)) (#pCons a' p') by auto
 ultimately show ?case by (intro exI[of - pCons a' p'], auto)
qed

end



lemma mod-mod-nat[simp]: a mod b mod (b * c :: nat) = a mod b
 by (simp add: mod-mult2-eq)

locale Knuth-ex-4-6-2-22-base =
 fixes ty-p :: 'p :: nontriv itself
   and ty-q :: 'q :: nontriv itself
   and ty-pq :: 'pq :: nontriv itself
 assumes pq: CARD('pq) = CARD('p) * CARD('q)
```

and *p-dvd-q*: *CARD*(*'p*) *dvd CARD*(*'q*)
begin

sublocale *rebase-q-to-p*: *rebase-dvd TYPE*(*'q*) *TYPE*(*'p*) **using** *p-dvd-q* **by** (*unfold-locales*, *auto*)
sublocale *rebase-pq-to-p*: *rebase-mult TYPE*(*'pq*) *TYPE*(*'p*) *TYPE*(*'q*) **using** *pq* **by** (*unfold-locales*, *auto*)
sublocale *rebase-pq-to-q*: *rebase-mult TYPE*(*'pq*) *TYPE*(*'q*) *TYPE*(*'p*) **using** *pq* **by** (*unfold-locales*, *auto*)

sublocale *rebase-p-to-q*: *rebase-ge TYPE*(*'p*) *TYPE* (*'q*) **by** (*unfold-locales*, *insert p-dvd-q*, *simp add*: *dvd-imp-le*)
sublocale *rebase-p-to-pq*: *rebase-ge TYPE*(*'p*) *TYPE* (*'pq*) **by** (*unfold-locales*, *simp add*: *pq*)
sublocale *rebase-q-to-pq*: *rebase-ge TYPE*(*'q*) *TYPE* (*'pq*) **by** (*unfold-locales*, *simp add*: *pq*)

definition *p* ≡ *if* (*ty-p* :: *'p itself*) = *ty-p then CARD*(*'p*) *else undefined*
lemma *p*[*simp*]: *p* ≡ *CARD*(*'p*) **unfolding** *p-def* **by** *auto*

definition *q* ≡ *if* (*ty-q* :: *'q itself*) = *ty-q then CARD*(*'q*) *else undefined*
lemma *q*[*simp*]: *q* = *CARD*(*'q*) **unfolding** *q-def* **by** *auto*

lemma *p1*: *int p* > *1*
  using *nontriv* [**where** *?'a* = *'p*] *p* **by** *simp*
lemma *q1*: *int q* > *1*
  using *nontriv* [**where** *?'a* = *'q*] *q* **by** *simp*
lemma *q0*: *int q* > *0*
  using *q1* **by** *auto*

lemma *pq2*[*simp*]: *CARD*(*'pq*) = *p* ∗ *q* **using** *pq* **by** *simp*

lemma *qq-eq-0*[*simp*]: (*of-nat CARD*(*'q*) ∗ *of-nat CARD*(*'q*) :: *'pq mod-ring*) = *0*
proof−
  have (*of-nat* (*q* ∗ *q*) :: *'pq mod-ring*) = *0* **by** (*rule of-nat-zero*, *auto simp*: *p-dvd-q*)
  then show *?thesis* **by** *auto*
qed

lemma *of-nat-q*[*simp*]: *of-nat q* :: *'q mod-ring* ≡ *0* **by** (*fold of-nat-card-eq-0*, *auto*)

lemma *rebase-rebase*[*simp*]: (@(@(*x*::*'pq mod-ring*) :: *'q mod-ring*) :: *'p mod-ring*) = @*x*
  using *p-dvd-q* **by** (*transfer*) (*simp add*: *mod-mod-cancel*)

lemma *rebase-rebase-poly*[*simp*]: (#(#(*f*::*'pq mod-ring poly*) :: *'q mod-ring poly*) :: *'p mod-ring poly*) = #*f*
  by (*induct f*, *auto*)

end

definition *dupe-monic* **where**
dupe-monic D H S T U = (*case pdivmod-monic* (*T* * *U*) *D of* (*q*,*r*) ⇒ (*S* * *U* + *H* * *q*, *r*))

lemma *dupe-monic*:
fixes *D* :: *'a* :: *prime-card mod-ring poly*
assumes *1*: *D**S* + *H***T* = *1*
and *mon*: *monic D*
and *dupe*: *dupe-monic D H S T U* = (*A*,*B*)
shows *A* * *D* + *B* * *H* = *U B* = *0* ∨ *degree B* < *degree D*
coprime D H ⟹ *A'* * *D* + *B'* * *H* = *U* ⟹ *B'* = *0* ∨ *degree B'* < *degree D* ⟹ *A'* = *A* ∧ *B'* = *B*
proof −
obtain *q r* **where** *div*: *pdivmod-monic* (*T* * *U*) *D* = (*q*,*r*) **by** *force*
from *dupe*[*unfolded dupe-monic-def div split*]
have *A*: *A* = (*S* * *U* + *H* * *q*) **and** *B*: *B* = *r* **by** *auto*
from *pdivmod-monic*[*OF mon div*] **have** *TU*: *T* * *U* = *D* * *q* + *r* **and**
deg: *r* = *0* ∨ *degree r* < *degree D* **by** *auto*
hence *r*: *r* = *T* * *U* − *D* * *q* **by** *simp*
have *A* * *D* + *B* * *H* = (*S* * *U* + *H* * *q*) * *D* + (*T* * *U* − *D* * *q*) * *H* **unfolding** *A B r* **by** *simp*
also have ... = (*D* * *S* + *H* * *T*) * *U* **by** (*simp add*: *field-simps*)
also have *D* * *S* + *H* * *T* = *1* **using** *1* **by** *simp*
finally show *eq*: *A* * *D* + *B* * *H* = *U* **by** *simp*
show *degB*: *B* = *0* ∨ *degree B* < *degree D* **using** *deg* **unfolding** *B* **by** (*cases r* = *0*, *auto*)
assume *another*: *A'* * *D* + *B'* * *H* = *U* **and** *degB'*: *B'* = *0* ∨ *degree B'* < *degree D*
and *cop*: *coprime D H*
from *degB* **have** *degB*: *B* = *0* ∨ *degree B* < *degree D* **by** *auto*
from *degB'* **have** *degB'*: *B'* = *0* ∨ *degree B'* < *degree D* **by** *auto*
from *mon* **have** *D0*: *D* ≠ *0* **by** *auto*
from *another eq* **have** *A'* * *D* + *B'* * *H* = *A* * *D* + *B* * *H* **by** *simp*
from *uniqueness-poly-equality*[*OF cop degB' degB D0 this*]
show *A'* = *A* ∧ *B'* = *B* **by** *auto*
qed

locale *Knuth-ex-4-6-2-22-main* = *Knuth-ex-4-6-2-22-base p-ty q-ty pq-ty*
for *p-ty* :: *'p*::*nontriv itself*
and *q-ty* :: *'q*::*nontriv itself*
and *pq-ty* :: *'pq*::*nontriv itself* +
fixes *a b* :: *'p mod-ring poly* **and** *u* :: *'pq mod-ring poly* **and** *v w* :: *'q mod-ring poly*
assumes *uvw*: (#*u* :: *'q mod-ring poly*) = *v* * *w*
and *degu*: *degree u* = *degree v* + *degree w*

```
    and avbw: (a * #v + b * #w :: 'p mod-ring poly) = 1
    and monic-v: monic v

    and bv: degree b < degree v
begin

lemma deg-v: degree (#v :: 'p mod-ring poly) = degree v
  using monic-v by (simp add: of-int-hom.monic-degree-map-poly-hom)

lemma u0: u ≠ 0 using degu bv by auto

lemma ex-f: ∃ f :: 'p mod-ring poly. u = #v * #w + smult (of-nat q) (#f)
proof−
 from uvw have (#(u − #v * #w) :: 'q mod-ring poly) = 0 by (auto simp:hom-distribs)
  from rebase-pq-to-q.rebase-poly-eq-0-imp-ex-smult[OF this]
  obtain f :: 'p mod-ring poly where u − #v * #w = smult (of-nat q) (#f) by
force
  then have u = #v * #w + smult (of-nat q) (#f) by (metis add-diff-cancel-left'
add-diff-eq)
  then show ?thesis by (intro exI[of - f], auto)
qed

definition f :: 'p mod-ring poly ≡ SOME f. u = #v * #w + smult (of-nat q)
(#f)

lemma u: u = #v * #w + smult (of-nat q) (#f)
  using ex-f[folded some-eq-ex] f-def by auto

lemma t-ex: ∃ t :: 'p mod-ring poly. degree (b * f − t * #v) < degree v
proof−
  define v' where v' ≡ #v :: 'p mod-ring poly
  from monic-v
  have 1: lead-coeff v' = 1 by (simp add: v'-def deg-v)
  then have 4: v' ≠ 0 by auto
  obtain t rem :: 'p mod-ring poly
  where pseudo-divmod (b * f) v' = (t,rem) by force
  from pseudo-divmod[OF 4 this, folded, unfolded 1]
  have b * f = v' * t + rem and deg: rem = 0 ∨ degree rem < degree v' by auto
  then have rem = b * f − t * v' by(auto simp: ac-simps)
  also have ... = b * f − #(#t :: 'p mod-ring poly) * v' (is - = - − ?t * v') by
simp
  also have ... = b * f − ?t * #v
    by (unfold v'-def, rule)
  finally have degree rem = degree ... by auto
  with deg bv have degree (b * f − ?t * #v :: 'p mod-ring poly) < degree v by
(auto simp: v'-def deg-v)
  then show ?thesis by (rule exI)
qed
```

definition t **where** $t \equiv SOME\ t :: 'p\ mod\text{-}ring\ poly.\ degree\ (b * f - t * \#v) < degree\ v$

definition $v' \equiv b * f - t * \#v$
definition $w' \equiv a * f + t * \#w$

lemma f: $w' * \#v + v' * \#w = f$ (**is** $?l = -$)
proof−
  have $?l = f * (a * \#v + b * \#w :: 'p\ mod\text{-}ring\ poly)$ **by** (*simp add*: v'-*def* w'-*def ring-distribs ac-simps*)
  also
    from *avbw* **have** $(\#(a * \#v + b * \#w) :: 'p\ mod\text{-}ring\ poly) = 1$ **by** *auto*
    then have $(a * \#v + b * \#w :: 'p\ mod\text{-}ring\ poly) = 1$ **by** *auto*
  finally show *?thesis* **by** *auto*
qed

lemma *degv'*: $degree\ v' < degree\ v$ **by** (*unfold* v'-*def t-def*, *rule someI-ex*, *rule t-ex*)

lemma *degqf*[*simp*]: $degree\ (smult\ (of\text{-}nat\ CARD('q))\ (\#f :: 'pq\ mod\text{-}ring\ poly)) = degree\ (\#f :: 'pq\ mod\text{-}ring\ poly)$
proof (*intro degree-smult-eqI*)
  assume $degree\ (\#f :: 'pq\ mod\text{-}ring\ poly) \neq 0$
  then have *f0*: $degree\ f \neq 0$ **by** *simp*
  moreover define l **where** $l \equiv lead\text{-}coeff\ f$
  ultimately have *l0*: $l \neq 0$ **by** *auto*
  then show $of\text{-}nat\ CARD('q) * lead\text{-}coeff\ (\#f::'pq\ mod\text{-}ring\ poly) \neq 0$
  apply (*unfold rebase-p-to-pq.lead-coeff-rebase-poly*, *fold l-def*)
  apply (*transfer*)
  using *q1* **by** (*simp add*: *pq mod-mod-cancel*)
qed

lemma *degw'*: $degree\ w' \leq degree\ w$
proof(*rule ccontr*)
  let $?f = \#f :: 'pq\ mod\text{-}ring\ poly$
  let $?qf = smult\ (of\text{-}nat\ q)\ (\#f) :: 'pq\ mod\text{-}ring\ poly$

  have $degree\ (\#w::'p\ mod\text{-}ring\ poly) \leq degree\ w$ **by** (*rule degree-rebase-poly-le*)
  also assume $\neg\ degree\ w' \leq degree\ w$
  then have *1*: $degree\ w < degree\ w'$ **by** *auto*
  finally have *2*: $degree\ (\#w :: 'p\ mod\text{-}ring\ poly) < degree\ w'$ **by** *auto*
  then have *w'0*: $w' \neq 0$ **by** *auto*

  have *3*: $degree\ (\#v * w') = degree\ (\#v :: 'p\ mod\text{-}ring\ poly) + degree\ w'$
      using *monic-v*[*unfolded*] **by** (*intro degree-monic-mult*[*OF - w'0*], *auto simp*: *deg-v*)

  have $degree\ f \leq degree\ u$
  proof(*rule ccontr*)
    assume ¬*?thesis*

then have $*$: *degree u* $<$ *degree f* **by** *auto*
with *degu* **have** *1*: *degree v* $+$ *degree w* $<$ *degree f* **by** *auto*
define *lcf* **where** *lcf* $\equiv$ *lead-coeff f*
with *1* **have** *lcf0*: *lcf* $\neq$ *0* **by** (*unfold*, *auto*)
have *degree f* $=$ *degree ?qf* **by** *simp*
also have ... $=$ *degree* (#*v* $*$ #*w* $+$ *?qf*)
proof(*rule sym*, *rule degree-add-eq-right*)
from *1 degree-mult-le*[*of* #*v*::*'pq mod-ring poly* #*w*]
show *degree* (#*v* $*$ #*w* :: *'pq mod-ring poly*) $<$ *degree ?qf* **by** *simp*
qed
also have ... $<$ *degree f* **using** $*$ *u* **by** *auto*
finally show *False* **by** *auto*
qed
with *degu* **have** *degree f* $\leq$ *degree v* $+$ *degree w* **by** *auto*
also note *f*[*symmetric*]
finally have *degree* (*w'* $*$ #*v* $+$ *v'* $*$ #*w*) $\leq$ *degree v* $+$ *degree w*.
moreover have *degree* (*w'* $*$ #*v* $+$ *v'* $*$ #*w*) $=$ *degree* (*w'* $*$ #*v*)
proof(*rule degree-add-eq-left*)
have *degree* (*v'* $*$ #*w*) $\leq$ *degree v'* $+$ *degree* (#*w* :: *'p mod-ring poly*)
by(*rule degree-mult-le*)
also have ... $<$ *degree v* $+$ *degree* (#*w* :: *'p mod-ring poly*) **using** *degv'* **by** *auto*
also have ... $<$ *degree* (#*v* :: *'p mod-ring poly*) $+$ *degree w'* **using** *2* **by** (*auto simp*: *deg-v*)
also have ... $=$ *degree* (#*v* $*$ *w'*) **using** *3* **by** *auto*
finally show *degree* (*v'* $*$ #*w*) $<$ *degree* (*w'* $*$ #*v*) **by** (*auto simp*: *ac-simps*)
qed
ultimately have *degree* (*w'* $*$ #*v*) $\leq$ *degree v* $+$ *degree w* **by** *auto*
moreover
from *3* **have** *degree* (*w'* $*$ #*v*) $=$ *degree w'* $+$ *degree v* **by** (*auto simp*: *ac-simps deg-v*)
with *1* **have** *degree w* $+$ *degree v* $<$ *degree* (*w'* $*$ #*v*) **by** *auto*
ultimately show *False* **by** *auto*
qed

abbreviation *qv'* $\equiv$ *smult* (*of-nat q*) (#*v'*) :: *'pq mod-ring poly*
abbreviation *qw'* $\equiv$ *smult* (*of-nat q*) (#*w'*) :: *'pq mod-ring poly*

abbreviation *V* $\equiv$ #*v* $+$ *qv'*
abbreviation *W* $\equiv$ #*w* $+$ *qw'*

lemma *vV*: *v* $=$ #*V* **by** (*auto simp*: *v'-def hom-distribs*)

lemma *wW*: *w* $=$ #*W* **by** (*auto simp*: *w'-def hom-distribs*)

lemma *uVW*: *u* $=$ *V* $*$ *W*
by (*subst u*, *fold f*, *simp add*: *ring-distribs add.left-cancel smult-add-right*[*symmetric*] *hom-distribs*)

lemma *degV*: *degree V* $=$ *degree v*

```
  and lcV: lead-coeff V = @lead-coeff v
  and degW: degree W = degree w
proof−
  from p1 q1 have int p < int p * int q by auto
  from less-trans[OF - this]
  have 1: l < int p ⟹ l < int p * int q for l by auto
  have degree qv' = degree (#v' :: 'pq mod-ring poly)
  proof (rule degree-smult-eqI, safe, unfold rebase-p-to-pq.degree-rebase-poly-eq)
    define l where l ≡ lead-coeff v'
    assume degree v' > 0
    then have lead-coeff v' ≠ 0 by auto
    then have (@l :: 'pq mod-ring) ≠ 0 by (simp add: l-def)
    then have (of-nat q * @l :: 'pq mod-ring) ≠ 0
      apply (transfer fixing:q-ty) using p-dvd-q p1 q1 1 by auto
    moreover assume  of-nat q * coeff (#v') (degree v') = (0 :: 'pq mod-ring)
    ultimately show False by (auto simp: l-def)
  qed
  also from degv' have ... < degree (#v::'pq mod-ring poly) by simp
  finally have *: degree qv' < degree (#v :: 'pq mod-ring poly).
  from degree-add-eq-left[OF *]
  show **: degree V = degree v by (simp add: v'-def)

  from * have coeff qv' (degree v) = 0 by (intro coeff-eq-0, auto)
  then show lead-coeff V = @lead-coeff v by (unfold **, auto simp: v'-def)

  with u0 uVW have degree (V * W) = degree V + degree W
    by (intro degree-mult-eq-left-unit, auto simp: monic-v)
  from this[folded uVW, unfolded degu **] show degree W = degree w by auto
qed

end

locale Knuth-ex-4-6-2-22-prime = Knuth-ex-4-6-2-22-main ty-p ty-q ty-pq a b u v
w
  for ty-p :: 'p :: prime-card itself
  and ty-q :: 'q :: nontriv itself
  and ty-pq :: 'pq :: nontriv itself
  and a b u v w +
  assumes coprime: coprime (#v :: 'p mod-ring poly) (#w)

begin

lemma coprime-preserves: coprime (#V :: 'p mod-ring poly) (#W)
 apply (intro coprimeI,simp add: rebase-q-to-p.of-nat-CARD-eq-0[simplified] hom-distribs)
  using coprime by (elim coprimeE, auto)

lemma pre-unique:
  assumes f2: w'' * #v + v'' * #w = f
      and degv'': degree v'' < degree v
```

shows $v'' = v' \wedge w'' = w'$
proof(*intro conjI*)
from *f f2*
have $w' * \#v + v' * \#w = w'' * \#v + v'' * \#w$ **by** *auto*
also have $... - w'' * \#v = v'' * \#w$ **by** *auto*
also have $... - v' * \#w = (v''- v') * \#w$ **by** (*auto simp*: *left-diff-distrib*)
finally have $*$: $(w' - w'') * \#v = (v''- v') * \#w$ **by** (*auto simp*: *left-diff-distrib*)
then have $\#v$ *dvd* $(v'' - v') * \#w$ **by** (*auto intro*: *dvdI*[*of - - w' − w''*] *simp*: *ac-simps*)
with *coprime* **have** $\#v$ *dvd* $v'' - v'$
by (*simp add*: *coprime-dvd-mult-left-iff*)
moreover have *degree* $(v'' - v') <$ *degree v* **by** (*rule degree-diff-less*[*OF degv'' degv'*])
ultimately have $v'' - v' = 0$
by (*metis deg-v degree-0 gr-implies-not-zero poly-divides-conv0*)
then show $v'' = v'$ **by** *auto*
with $*$ **have** $(w' - w'') * \#v = 0$ **by** *auto*
with *bv* **have** $w' - w'' = 0$
by (*metis deg-v degree-0 gr-implies-not-zero mult-eq-0-iff*)
then show $w'' = w'$ **by** *auto*
qed

lemma *unique*:
assumes *vV2*: $v = \#V2$ **and** *wW2*: $w = \#W2$ **and** *uVW2*: $u = V2 * W2$
and *degV2*: *degree V2* = *degree v* **and** *degW2*: *degree W2* = *degree w*
and *lc*: *lead-coeff V2* = @*lead-coeff v*
shows $V2 = V$ $W2 = W$
proof−
from *vV2* **have** $(\#(V2 - \#v) :: 'q$ *mod-ring poly*$) = 0$ **by** (*auto simp*: *hom-distribs*)
from *rebase-pq-to-q.rebase-poly-eq-0-imp-ex-smult*[*OF this*]
obtain $v'' :: 'p$ *mod-ring poly*
where *deg*: *degree* $v'' \leq$ *degree* $(V2 - \#v)$
and v'': $V2 - \#v =$ *smult* (*of-nat CARD*($'q$)) $(\#v'')$ **by** (*elim exE conjE*)
then have *V2*: $V2 = \#v + ...$ **by** (*metis add-diff-cancel-left' diff-add-cancel*)

from *lc*[*unfolded degV2*, *unfolded V2*]
have *of-nat q* $*$ (@*coeff* v'' (*degree v*) $:: 'pq$ *mod-ring*) = *of-nat q* $*$ *0* **by** *auto*
from *this*[*unfolded q rebase-pq-to-p.rebase-mult-eq*]
have *coeff* v'' (*degree v*) = *0* **by** *simp*
moreover have *degree* $v'' \leq$ *degree v* **using** *deg degV2*
by (*metis degree-diff-le le-antisym nat-le-linear rebase-q-to-pq.degree-rebase-poly-eq*)
ultimately have *degv''*: *degree* $v'' <$ *degree v*
using *bv eq-zero-or-degree-less* **by** *fastforce*

from *wW2* **have** $(\#(W2 - \#w) :: 'q$ *mod-ring poly*$) = 0$ **by** (*auto simp*: *hom-distribs*)
from *rebase-pq-to-q.rebase-poly-eq-0-imp-ex-smult*[*OF this*] *pq*
obtain $w'' :: 'p$ *mod-ring poly* **where** w'': $W2 - \#w =$ *smult* (*of-nat q*) $(\#w'')$
by *force*

then have *W2*: *W2* = #*w* + ... **by** (*metis add-diff-cancel-left' diff-add-cancel*)

have *u* = #*v* * #*w* + *smult* (*of-nat q*) (#*w''* * #*v* + #*v''* * #*w*) + *smult* (*of-nat* (*q* * *q*)) (#*v''* * #*w''*)
by(*simp add*: *uVW2 V2 W2 ring-distribs smult-add-right ac-simps*)
also have *smult* (*of-nat* (*q* * *q*)) (#*v''* * #*w''* :: *'pq mod-ring poly*) = *0* **by** *simp*
finally have *u* − #*v* * #*w* = *smult* (*of-nat q*) (#*w''* * #*v* + #*v''* * #*w*) **by** *auto*
also have *u* − #*v* * #*w* = *smult* (*of-nat q*) (#*f*) **by** (*subst u*, *simp*)
finally have *w''* * #*v* + *v''* * #*w* = *f* **by** (*simp add*: *hom-distribs*)
from *pre-unique*[*OF this degv''*]
have *pre*: *v''* = *v'* *w''* = *w'* **by** *auto*
with *V2 W2* **show** *V2* = *V W2* = *W* **by** *auto*
qed

end

definition
hensel-1 (*ty* ::*'p* :: *prime-card itself*)
(*u* :: *'pq* :: *nontriv mod-ring poly*) (*v* :: *'q* :: *nontriv mod-ring poly*) (*w* :: *'q mod-ring poly*) ≡
if v = *1 then* (*1*,*u*) *else*
let (*s*, *t*) = *bezout-coefficients* (#*v* :: *'p mod-ring poly*) (#*w*) *in*
let (*a*, *b*) = *dupe-monic* (#*v*::*'p mod-ring poly*) (#*w*) *s t 1 in*
(*Knuth-ex-4-6-2-22-main.V TYPE*(*'q*) *b u v w*, *Knuth-ex-4-6-2-22-main.W TYPE*(*'q*) *a b u v w*)

lemma *hensel-1*:
fixes *u* :: *'pq* :: *nontriv mod-ring poly*
and *v w* :: *'q* :: *nontriv mod-ring poly*
assumes *CARD*(*'pq*) = *CARD*(*'p* :: *prime-card*) * *CARD*(*'q*)
and *CARD*(*'p*) *dvd CARD*(*'q*)
and *uvw*: #*u* = *v* * *w*
and *degu*: *degree u* = *degree v* + *degree w*
and *monic*: *monic v*
and *coprime*: *coprime* (#*v* :: *'p mod-ring poly*) (#*w*)
and *out*: *hensel-1 TYPE*(*'p*) *u v w* = (*V'*,*W'*)
shows *u* = *V'* * *W'* ∧ *v* = #*V'* ∧ *w* = #*W'* ∧ *degree V'* = *degree v* ∧ *degree W'* = *degree w* ∧
monic V' ∧ *coprime* (#*V'* :: *'p mod-ring poly*) (#*W'*) (**is** *?main*)
and (∀ *V'' W''*. *u* = *V''* * *W''* ⟶ *v* = #*V''* ⟶ *w* = #*W''* ⟶
degree V'' = *degree v* ⟶ *degree W''* = *degree w* ⟶ *lead-coeff V''* = @*lead-coeff v* ⟶
V'' = *V'* ∧ *W''* = *W'*) (**is** *?unique*)
proof−
from *monic*
have *degv*: *degree* (#*v* :: *'p mod-ring poly*) = *degree v*
by (*simp add*: *of-int-hom.monic-degree-map-poly-hom*)
from *monic*

```
  have monic2: monic (#v :: 'p mod-ring poly)
    by (auto simp: degv)
  obtain s t where bezout: bezout-coefficients (#v :: 'p mod-ring poly) (#w) = (s,
t)
    by (auto simp add: prod-eq-iff)
  then have s * #v + t * #w = gcd (#v :: 'p mod-ring poly) (#w)
    by (rule bezout-coefficients)
  with coprime have vswt: #v * s + #w * t = 1
    by (simp add: ac-simps)
  obtain a b where dupe: dupe-monic (#v) (#w) s t 1 = (a, b) by force
  from dupe-monic(1,2)[OF vswt monic2, where U=1, unfolded this]
  have avbw: a * #v + b * #w = 1 and degb: b = 0 ∨ degree b < degree (#v::'p
mod-ring poly) by auto
  have ?main ∧ ?unique
  proof (cases b = 0)
    case b0: True
    with avbw have a * #v = 1 by auto
    then have degree (#v :: 'p mod-ring poly) = 0
      by (metis degree-1 degree-mult-eq-0 mult-zero-left one-neq-zero)
    from this[unfolded degv] monic-degree-0[OF monic[unfolded]]
    have 1: v = 1 by auto
    with b0 out uvw have 2: V' = 1 W' = u
      by (unfold split hensel-1-def Let-def dupe) auto
   have 3: ?unique apply (simp add: 1 2) by (metis monic-degree-0 mult.left-neutral)
    with uvw degu show ?thesis unfolding 1 2 by auto
  next
    case b0: False
    with degb degv have degb: degree b < degree v by auto
    then have v1: v ≠ 1 by auto
    interpret Knuth-ex-4-6-2-22-prime TYPE('p) TYPE('q) TYPE('pq) a b
      by (unfold-locales; fact assms degb avbw)
    show ?thesis
    proof (intro conjI)
      from out [unfolded hensel-1-def] v1
      have 1 [simp]: V' = V W' = W by (auto simp: bezout dupe)
      from uVW show u = V' * W' by auto
      from degV show [simp]: degree V' = degree v by simp
      from degW show [simp]: degree W' = degree w by simp
      from lcV have lead-coeff V' = @lead-coeff v by simp
      with monic-v show monic V' by (simp add:)
      from vV show v = #V' by simp
      from wW show w = #W' by simp
      from coprime-preserves show coprime (#V' :: 'p mod-ring poly) (#W') by
simp
      show 9: ?unique by (unfold 1, intro allI conjI impI; rule unique)
    qed
  qed
  then show ?main ?unique by (fact conjunct1, fact conjunct2)
qed
```

end

9.3 Result is Unique

We combine the finite field factorization algorithm with Hensel-lifting to obtain factorizations mod p^n. Moreover, we prove results on unique-factorizations in mod p^n which admit to extend the uniqueness result for binary Hensel-lifting to the general case. As a consequence, our factorization algorithm will produce unique factorizations mod p^n.

theory *Berlekamp-Hensel*
imports
Finite-Field-Factorization-Record-Based
Hensel-Lifting
begin

hide-const *coeff monom*

definition *berlekamp-hensel* :: *int* ⇒ *nat* ⇒ *int poly* ⇒ *int poly list* **where**
berlekamp-hensel p n f = (*case finite-field-factorization-int p f of*
(-,*fs*) ⇒ *hensel-lifting p n f fs*)

Finite field factorization in combination with Hensel-lifting delivers factorization modulo p^k where factors are irreducible modulo p. Assumptions: input polynomial is square-free modulo p.

context *poly-mod-prime* **begin**

lemma *berlekamp-hensel-main*:
assumes *n*: *n* ≠ *0*
and *res*: *berlekamp-hensel p n f* = *gs*
and *cop*: *coprime* (*lead-coeff f*) *p*
and *sf*: *square-free-m f*
and *berl*: *finite-field-factorization-int p f* = (*c*,*fs*)
shows *poly-mod.factorization-m* (*p* ^ *n*) *f* (*lead-coeff f*, *mset gs*) — factorization mod *p*^*n*
and *sort* (*map degree fs*) = *sort* (*map degree gs*)
and ⋀ *g*. *g* ∈ *set gs* ⟹ *monic g* ∧ *poly-mod.Mp* (*p*^*n*) *g* = *g* ∧ — monic and normalized
poly-mod.irreducible-m p g ∧ — irreducibility even mod *p*
poly-mod.degree-m p g = *degree g* — mod *p* does not change degree of *g*
proof −
from *res*[*unfolded berlekamp-hensel-def berl split*]
have *hen*: *hensel-lifting p n f fs* = *gs* **.**
note *bh* = *finite-field-factorization-int*[*OF sf berl*]
from *bh* **have** *poly-mod.factorization-m p f* (*c*, *mset fs*) *c* ∈ {*0*..<*p*} (∀ *fi*∈*set fs*. *set* (*coeffs fi*) ⊆ {*0*..<*p*})
by (*auto simp*: *poly-mod.unique-factorization-m-alt-def*)
note *hen* = *hensel-lifting*[*OF n hen cop sf*, *OF this*]

show *poly-mod.factorization-m* (*p* ^ *n*) *f* (*lead-coeff f*, *mset gs*)
sort (*map degree fs*) = *sort* (*map degree gs*)
⋀ *g*. *g* ∈ *set gs* ⟹ *monic g* ∧ *poly-mod.Mp* (*p*^*n*) *g* = *g* ∧
poly-mod.irreducible-m p g ∧
poly-mod.degree-m p g = *degree g* **using** *hen* **by** *auto*
qed

theorem *berlekamp-hensel*:
assumes *cop*: *coprime* (*lead-coeff f*) *p*
and *sf*: *square-free-m f*
and *res*: *berlekamp-hensel p n f* = *gs*
and *n*: *n* ≠ *0*
shows *poly-mod.factorization-m* (*p*^*n*) *f* (*lead-coeff f*, *mset gs*) — factorization mod *p*^*n*
and ⋀ *g*. *g* ∈ *set gs* ⟹ *poly-mod.Mp* (*p*^*n*) *g* = *g* ∧ *poly-mod.irreducible-m p g*
— normalized and *irreducible* even mod *p*
proof −
obtain *c fs* **where** *finite-field-factorization-int p f* = (*c*,*fs*) **by** *force*
from *berlekamp-hensel-main*[*OF n res cop sf this*]
show *poly-mod.factorization-m* (*p*^*n*) *f* (*lead-coeff f*, *mset gs*)
⋀ *g*. *g* ∈ *set gs* ⟹ *poly-mod.Mp* (*p*^*n*) *g* = *g* ∧ *poly-mod.irreducible-m p g* **by** *auto*
qed

lemma *berlekamp-and-hensel-separated*:
assumes *cop*: *coprime* (*lead-coeff f*) *p*
and *sf*: *square-free-m f*
and *res*: *hensel-lifting p n f fs* = *gs*
and *berl*: *finite-field-factorization-int p f* = (*c*,*fs*)
and *n*: *n* ≠ *0*
shows *berlekamp-hensel p n f* = *gs*
and *sort* (*map degree fs*) = *sort* (*map degree gs*)
proof −
show *berlekamp-hensel p n f* = *gs* **unfolding** *res*[*symmetric*]
berlekamp-hensel-def hensel-lifting-def berl split Let-def **..**
from *berlekamp-hensel-main*[*OF n this cop sf berl*] **show** *sort* (*map degree fs*) = *sort* (*map degree gs*)
by *auto*
qed

end

lemma *prime-cop-exp-poly-mod*:
assumes *prime*: *prime p* **and** *cop*: *coprime c p* **and** *n*: *n* ≠ *0*
shows *poly-mod.M* (*p*^*n*) *c* ∈ {*1* ..< *p*^*n*}
proof −
from *prime* **have** *p1*: *p* > *1* **by** (*simp add*: *prime-int-iff*)
interpret *poly-mod-2 p*^*n* **unfolding** *poly-mod-2-def* **using** *p1 n* **by** *simp*

```
  from cop p1 m1 have M c ≠ 0
    by (auto simp add: M-def)
  moreover have M c < p^n M c ≥ 0 unfolding M-def using m1 by auto
  ultimately show ?thesis by auto
qed

context poly-mod-2
begin

context
  fixes p :: int
  assumes prime: prime p
begin

interpretation p: poly-mod-prime p using prime by unfold-locales

lemma coprime-lead-coeff-factor: assumes coprime (lead-coeff (f * g)) p
  shows coprime (lead-coeff f) p coprime (lead-coeff g) p
proof −
  {
    fix f g
    assume cop: coprime (lead-coeff (f * g)) p
    from this[unfolded lead-coeff-mult]
    have coprime (lead-coeff f) p using prime
      by simp
  }
  from this[OF assms] this[of g f] assms
  show coprime (lead-coeff f) p coprime (lead-coeff g) p by (auto simp: ac-simps)
qed

lemma unique-factorization-m-factor: assumes uf: unique-factorization-m (f * g)
(c,hs)
  and cop: coprime (lead-coeff (f * g)) p
  and sf: p.square-free-m (f * g)
  and n: n ≠ 0
  and m: m = p^n
  shows ∃ fs gs. unique-factorization-m f (lead-coeff f,fs)
  ∧ unique-factorization-m g (lead-coeff g,gs)
  ∧ Mf (c,hs) = Mf (lead-coeff f * lead-coeff g, fs + gs)
  ∧ image-mset Mp fs = fs ∧ image-mset Mp gs = gs
proof −
  from prime have p1: 1 < p by (simp add: prime-int-iff)
  interpret p: poly-mod-2 p by (standard, rule p1)
  note sf = p.square-free-m-factor[OF sf]
  note cop = coprime-lead-coeff-factor[OF cop]
  from cop have copm: coprime (lead-coeff f) m coprime (lead-coeff g) m
    by (simp-all add: m)
  have df: degree-m f = degree f
    by (rule degree-m-eq[OF - m1], insert copm(1) m1, auto)
```

have *dg*: *degree-m g* = *degree g*
by (*rule degree-m-eq*[*OF - m1*], *insert copm*(*2*) *m1*, *auto*)
define *fs* **where** *fs* ≡ *mset* (*berlekamp-hensel p n f*)
define *gs* **where** *gs* ≡ *mset* (*berlekamp-hensel p n g*)
from *p.berlekamp-hensel*[*OF cop*(*1*) *sf*(*1*) *refl n*, *folded m*]
have *f*: *factorization-m f* (*lead-coeff f*,*fs*)
and *f-id*: ⋀ *f*. *f* ∈# *fs* ⟹ *Mp f* = *f* **unfolding** *fs-def* **by** *auto*
from *p.berlekamp-hensel*[*OF cop*(*2*) *sf*(*2*) *refl n*, *folded m*]
have *g*: *factorization-m g* (*lead-coeff g*,*gs*)
and *g-id*: ⋀ *f*. *f* ∈# *gs* ⟹ *Mp f* = *f* **unfolding** *gs-def* **by** *auto*
from *factorization-m-prod*[*OF f g*] *uf*[*unfolded unique-factorization-m-alt-def*]
have *eq*: *Mf* (*lead-coeff f* ∗ *lead-coeff g*, *fs* + *gs*) = *Mf* (*c*,*hs*) **by** *blast*
have *uff*: *unique-factorization-m f* (*lead-coeff f*,*fs*)
proof (*rule unique-factorization-mI*[*OF f*])
fix *e ks*
assume *factorization-m f* (*e*,*ks*)
from *factorization-m-prod*[*OF this g*] *uf*[*unfolded unique-factorization-m-alt-def*]
factorization-m-lead-coeff[*OF this*, *unfolded degree-m-eq-lead-coeff*[*OF df*]]
have *Mf* (*e* ∗ *lead-coeff g*, *ks* + *gs*) = *Mf* (*c*,*hs*) **and** *e*: *M* (*lead-coeff f*) = *M* *e* **by** *blast*+
from *this*[*folded eq*, *unfolded Mf-def split*]
have *ks*: *image-mset Mp ks* = *image-mset Mp fs* **by** *auto*
show *Mf* (*e*, *ks*) = *Mf* (*lead-coeff f*, *fs*) **unfolding** *Mf-def split ks e* **by** *simp*
qed
have *idf*: *image-mset Mp fs* = *fs* **using** *f-id* **by** (*induct fs*, *auto*)
have *idg*: *image-mset Mp gs* = *gs* **using** *g-id* **by** (*induct gs*, *auto*)
have *ufg*: *unique-factorization-m g* (*lead-coeff g*,*gs*)
proof (*rule unique-factorization-mI*[*OF g*])
fix *e ks*
assume *factorization-m g* (*e*,*ks*)
from *factorization-m-prod*[*OF f this*] *uf*[*unfolded unique-factorization-m-alt-def*]
factorization-m-lead-coeff[*OF this*, *unfolded degree-m-eq-lead-coeff*[*OF dg*]]
have *Mf* (*lead-coeff f* ∗ *e*, *fs* + *ks*) = *Mf* (*c*,*hs*) **and** *e*: *M* (*lead-coeff g*) = *M* *e* **by** *blast*+
from *this*[*folded eq*, *unfolded Mf-def split*]
have *ks*: *image-mset Mp ks* = *image-mset Mp gs* **by** *auto*
show *Mf* (*e*, *ks*) = *Mf* (*lead-coeff g*, *gs*) **unfolding** *Mf-def split ks e* **by** *simp*
qed
from *uff ufg eq*[*symmetric*] *idf idg* **show** *?thesis* **by** *auto*
qed

lemma *unique-factorization-factorI*:
assumes *ufact*: *unique-factorization-m* (*f* ∗ *g*) *FG*
and *cop*: *coprime* (*lead-coeff* (*f* ∗ *g*)) *p*
and *sf*: *poly-mod.square-free-m p* (*f* ∗ *g*)
and *n*: *n* ≠ *0*
and *m*: *m* = *p*^*n*
shows *factorization-m f F* ⟹ *unique-factorization-m f F*
and *factorization-m g G* ⟹ *unique-factorization-m g G*

proof −
obtain *c fg* **where** *FG*: *FG* = (*c*,*fg*) **by** *force*
from *unique-factorization-m-factor*[*OF ufact*[*unfolded FG*] *cop sf n m*]
obtain *fs gs* **where** *ufact*: *unique-factorization-m f* (*lead-coeff f*, *fs*)
unique-factorization-m g (*lead-coeff g*, *gs*) **by** *auto*
from *ufact*(*1*) **show** *factorization-m f F* ⟹ *unique-factorization-m f F*
by (*metis unique-factorization-m-alt-def*)
from *ufact*(*2*) **show** *factorization-m g G* ⟹ *unique-factorization-m g G*
by (*metis unique-factorization-m-alt-def*)
qed

end

lemma *monic-Mp-prod-mset*: **assumes** *fs*: ⋀ *f*. *f* ∈# *fs* ⟹ *monic* (*Mp f*)
shows *monic* (*Mp* (*prod-mset fs*))
proof −
have *monic* (*prod-mset* (*image-mset Mp fs*))
by (*rule monic-prod-mset*, *insert fs*, *auto*)
from *monic-Mp*[*OF this*] **have** *monic* (*Mp* (*prod-mset* (*image-mset Mp fs*))) **.**
also have *Mp* (*prod-mset* (*image-mset Mp fs*)) = *Mp* (*prod-mset fs*) **by** (*rule Mp-prod-mset*)
finally show *?thesis* **.**
qed

lemma *degree-Mp-mult-monic*: **assumes** *monic f monic g*
shows *degree* (*Mp* (*f* ∗ *g*)) = *degree f* + *degree g*
by (*metis zero-neq-one assms degree-monic-mult leading-coeff-0-iff monic-degree-m monic-mult*)

lemma *factorization-m-degree*: **assumes** *factorization-m f* (*c*,*fs*)
and *0*: *Mp f* ≠ *0*
shows *degree-m f* = *sum-mset* (*image-mset degree-m fs*)
proof −
note *a* = *assms*[*unfolded factorization-m-def split*]
hence *deg*: *degree-m f* = *degree-m* (*smult c* (*prod-mset fs*))
and *fs*: ⋀ *f*. *f* ∈# *fs* ⟹ *monic* (*Mp f*) **by** *auto*
define *gs* **where** *gs* ≡ *Mp* (*prod-mset fs*)
from *monic-Mp-prod-mset*[*OF fs*] **have** *mon-gs*: *monic gs* **unfolding** *gs-def* **.**
have *d*:*degree* (*Mp* (*Polynomial.smult c gs*)) = *degree gs*
proof −
have *f1*: *0* ≠ *c* **by** (*metis 0 Mp-0 a*(*1*) *smult-eq-0-iff*)
then have *M c* ≠ *0* **by** (*metis* (*no-types*) *0 assms*(*1*) *factorization-m-lead-coeff leading-coeff-0-iff*)
then show *degree* (*Mp* (*Polynomial.smult c gs*)) = *degree gs*
unfolding *monic-degree-m*[*OF mon-gs*,*symmetric*]
using *f1* **by** (*metis coeff-smult degree-m-eq degree-smult-eq m1 mon-gs monic-degree-m mult-cancel-left1 poly-mod.M-def*)
qed
note *deg*

also have *degree-m* (*smult c* (*prod-mset fs*)) = *degree-m* (*smult c gs*)
unfolding *gs-def* **by** *simp*
also have ... = *degree gs* **using** *d*.
also have ... = *sum-mset* (*image-mset degree-m fs*) **unfolding** *gs-def*
using *fs*
proof (*induct fs*)
case (*add f fs*)
have *mon*: *monic* (*Mp f*) *monic* (*Mp* (*prod-mset fs*)) **using** *monic-Mp-prod-mset*[*of fs*]
add(*2*) **by** *auto*
have *degree* (*Mp* (*prod-mset* (*add-mset f fs*))) = *degree* (*Mp* (*Mp f* ∗ *Mp* (*prod-mset fs*)))
by (*auto simp*: *ac-simps*)
also have ... = *degree* (*Mp f*) + *degree* (*Mp* (*prod-mset fs*))
by (*rule degree-Mp-mult-monic*[*OF mon*])
also have *degree* (*Mp* (*prod-mset fs*)) = *sum-mset* (*image-mset degree-m fs*)
by (*rule add*(*1*), *insert add*(*2*), *auto*)
finally show *?case* **by** (*simp add*: *ac-simps*)
qed *simp*
finally show *?thesis* .
qed

lemma *degree-m-mult-le*: *degree-m* (*f* ∗ *g*) ≤ *degree-m f* + *degree-m g*
using *degree-m-mult-le* **by** *auto*

lemma *degree-m-prod-mset-le*: *degree-m* (*prod-mset fs*) ≤ *sum-mset* (*image-mset degree-m fs*)
proof (*induct fs*)
case *empty*
show *?case* **by** *simp*
next
case (*add f fs*)
then show *?case* **using** *degree-m-mult-le*[*of f prod-mset fs*] **by** *auto*
qed

end

context *poly-mod-prime*
begin

lemma *unique-factorization-m-factor-partition*: **assumes** *l0*: *l* ≠ *0*
and *uf*: *poly-mod.unique-factorization-m* (*p*^*l*) *f* (*lead-coeff f*, *mset gs*)
and *f*: *f* = *f1* ∗ *f2*
and *cop*: *coprime* (*lead-coeff f*) *p*
and *sf*: *square-free-m f*
and *part*: *List.partition* (λ*gi*. *gi dvdm f1*) *gs* = (*gs1*, *gs2*)
shows *poly-mod.unique-factorization-m* (*p*^*l*) *f1* (*lead-coeff f1*, *mset gs1*)
poly-mod.unique-factorization-m (*p*^*l*) *f2* (*lead-coeff f2*, *mset gs2*)

proof −
interpret *pl*: *poly-mod-2 p^l* **by** (*standard*, *insert m1 l0*, *auto*)
let *?I* = *image-mset pl.Mp*
note *Mp-pow* [*simp*] = *Mp-Mp-pow-is-Mp*[*OF l0 m1*]
have [*simp*]: *pl.Mp x dvdm u* = (*x dvdm u*) **for** *x u* **unfolding** *dvdm-def* **using** *Mp-pow*[*of x*]
by (*metis poly-mod.mult-Mp*(*1*))
have *gs-split*: *set gs* = *set gs1* ∪ *set gs2* **using** *part* **by** *auto*
from *pl.unique-factorization-m-factor*[*OF prime uf*[*unfolded f*] - - *l0 refl*, *folded f*, *OF cop sf*]
obtain *hs1 hs2* **where** *uf'*: *pl.unique-factorization-m f1* (*lead-coeff f1*, *hs1*)
pl.unique-factorization-m f2 (*lead-coeff f2*, *hs2*)
and *gs-hs*: *?I* (*mset gs*) = *hs1* + *hs2*
unfolding *pl.Mf-def split* **by** *auto*
have *gs-gs*: *?I* (*mset gs*) = *?I* (*mset gs1*) + *?I* (*mset gs2*) **using** *part*
by (*auto*, *induct gs arbitrary*: *gs1 gs2*, *auto*)
with *gs-hs* **have** *gs-hs12*: *?I* (*mset gs1*) + *?I* (*mset gs2*) = *hs1* + *hs2* **by** *auto*
note *pl-dvdm-imp-p-dvdm* = *pl-dvdm-imp-p-dvdm*[*OF l0*]
note *fact* = *pl.unique-factorization-m-imp-factorization*[*OF uf*]
have *gs1*: *?I* (*mset gs1*) = {#*x* ∈# *?I* (*mset gs*). *x dvdm f1*#}
using *part* **by** (*auto*, *induct gs arbitrary*: *gs1 gs2*, *auto*)
also have ... = {#*x* ∈# *hs1*. *x dvdm f1*#} + {#*x* ∈# *hs2*. *x dvdm f1*#}
unfolding *gs-hs* **by** *simp*
also have {#*x* ∈# *hs2*. *x dvdm f1*#} = {#}
proof (*rule ccontr*)
assume ¬ *?thesis*
then obtain *x* **where** *x*: *x* ∈# *hs2* **and** *dvd*: *x dvdm f1* **by** *fastforce*
from *x gs-hs* **have** *x* ∈# *?I* (*mset gs*) **by** *auto*
with *fact*[*unfolded pl.factorization-m-def*]
have *xx*: *pl.irreducible*$_d$*-m x monic x* **by** *auto*
from *square-free-m-prod-imp-coprime-m*[*OF sf*[*unfolded f*]]
have *cop-h-f*: *coprime-m f1 f2* **by** *auto*
from *pl.factorization-m-mem-dvdm*[*OF pl.unique-factorization-m-imp-factorization*[*OF uf'*(*2*)], *of x*] *x*
have *pl.dvdm x f2* **by** *auto*
hence *x dvdm f2* **by** (*rule pl-dvdm-imp-p-dvdm*)
from *cop-h-f*[*unfolded coprime-m-def*, *rule-format*, *OF dvd this*]
have *x dvdm 1* **by** *auto*
from *dvdm-imp-degree-le*[*OF this xx*(*2*) - *m1*] **have** *degree x* = *0* **by** *auto*
with *xx* **show** *False* **unfolding** *pl.irreducible*$_d$*-m-def* **by** *auto*
qed
also have {#*x* ∈# *hs1*. *x dvdm f1*#} = *hs1*
proof (*rule ccontr*)
assume ¬ *?thesis*
from *filter-mset-inequality*[*OF this*]
obtain *x* **where** *x*: *x* ∈# *hs1* **and** *dvd*: ¬ *x dvdm f1* **by** *blast*
from *pl.factorization-m-mem-dvdm*[*OF pl.unique-factorization-m-imp-factorization*[*OF uf'*(*1*)],
of x] *x dvd*

```
    have pl.dvdm x f1 by auto
    from pl-dvdm-imp-p-dvdm[OF this] dvd show False by auto
  qed
  finally have gs-hs1: ?I (mset gs1) = hs1 by simp
  with gs-hs12 have ?I (mset gs2) = hs2 by auto
  with uf' gs-hs1 have pl.unique-factorization-m f1 (lead-coeff f1, ?I (mset gs1))
     pl.unique-factorization-m f2 (lead-coeff f2, ?I (mset gs2)) by auto
  thus pl.unique-factorization-m f1 (lead-coeff f1, mset gs1)
     pl.unique-factorization-m f2 (lead-coeff f2, mset gs2)
    unfolding pl.unique-factorization-m-def
    by (auto simp: pl.Mf-def image-mset.compositionality o-def)
qed

lemma factorization-pn-to-factorization-p: assumes fact: poly-mod.factorization-m
(p^n) C (c,fs)
  and sf: square-free-m C
  and n: n ≠ 0
shows factorization-m C (c,fs)
proof −
  let ?q = p^n
  from n m1 have q: ?q > 1 by simp
  interpret q: poly-mod-2 ?q by (standard, insert q, auto)
  from fact[unfolded q.factorization-m-def]
  have eq: q.Mp C = q.Mp (Polynomial.smult c (prod-mset fs))
    and irr: ⋀ f. f ∈# fs ⟹ q.irreducible_d-m f
    and mon: ⋀ f. f ∈# fs ⟹ monic (q.Mp f)
    by auto
  from arg-cong[OF eq, of Mp]
  have eq: eq-m C (smult c (prod-mset fs))
    by (simp add: Mp-Mp-pow-is-Mp m1 n)
  show ?thesis unfolding factorization-m-def split
  proof (rule conjI[OF eq], intro ballI conjI)
    fix f
    assume f: f ∈# fs
    from mon[OF this] have mon-qf: monic (q.Mp f) .
    hence lc: lead-coeff (q.Mp f) = 1 by auto
    from mon-qf show mon-f: monic (Mp f)
      by (metis Mp-Mp-pow-is-Mp m1 monic-Mp n)
    from irr[OF f] have irr: q.irreducible_d-m f .
    hence q.degree-m f ≠ 0 unfolding q.irreducible_d-m-def by auto
    also have q.degree-m f = degree-m f using mon[OF f]
      by (metis Mp-Mp-pow-is-Mp m1 monic-degree-m n)
    finally have deg: degree-m f ≠ 0 by auto
    from f obtain gs where fs: fs = {#f#} + gs
      by (metis mset-subset-eq-single subset-mset.add-diff-inverse)
    from eq[unfolded fs] have Mp C = Mp (f * smult c (prod-mset gs)) by auto
    from square-free-m-factor[OF square-free-m-cong[OF sf this]]
    have sf-f: square-free-m f by simp
    have sf-Mf: square-free-m (q.Mp f)
```

```
    by (rule square-free-m-cong[OF sf-f], auto simp: Mp-Mp-pow-is-Mp n m1)
  have coprime (lead-coeff (q.Mp f)) p using mon[OF f] prime by simp
  from berlekamp-hensel[OF this sf-Mf refl n, unfolded lc] obtain gs where
    qfact: q.factorization-m (q.Mp f) (1, mset gs)
    and ⋀ g. g ∈ set gs ⟹ irreducible-m g by blast
  hence fact: q.Mp f = q.Mp (prod-list gs)
     and gs: ⋀ g. g∈ set gs ⟹ irreducible_d-m g ∧ q.irreducible_d-m g ∧ monic
(q.Mp g)
    unfolding q.factorization-m-def by auto
  from q.factorization-m-degree[OF qfact]
  have deg: q.degree-m (q.Mp f) = sum-mset (image-mset q.degree-m (mset gs))
    using mon-qf by fastforce
  from irr[unfolded q.irreducible_d-m-def]
  have sum-mset (image-mset q.degree-m (mset gs)) ≠ 0 by (fold deg, auto)
  then obtain g gs' where gs1: gs = g # gs' by (cases gs, auto)
  {
    assume gs' ≠ []
    then obtain h hs where gs2: gs' = h # hs by (cases gs', auto)
    from deg gs[unfolded q.irreducible_d-m-def]
    have small: q.degree-m g < q.degree-m f
      q.degree-m h + sum-mset (image-mset q.degree-m (mset hs)) < q.degree-m
f
      unfolding gs1 gs2 by auto
    have q.eq-m f (g * (h * prod-list hs))
      using fact unfolding gs1 gs2 by simp
    with irr[unfolded q.irreducible_d-m-def, THEN conjunct2, rule-format, of g h
* prod-list hs]
      small(1) have ¬ q.degree-m (h * prod-list hs) < q.degree-m f by auto
    hence q.degree-m f ≤ q.degree-m (h * prod-list hs) by simp
    also have ... = q.degree-m (prod-mset ({#h#} + mset hs)) by simp
    also have ... ≤ sum-mset (image-mset q.degree-m ({#h#} + mset hs))
      by (rule q.degree-m-prod-mset-le)
    also have ... < q.degree-m f using small(2) by simp
    finally have False by simp
  }
  hence gs1: gs = [g] unfolding gs1 by (cases gs', auto)
  with fact have q.Mp f = q.Mp g by auto
  from arg-cong[OF this, of Mp] have eq: Mp f = Mp g
    by (simp add: Mp-Mp-pow-is-Mp m1 n)
  from gs[unfolded gs1] have g: irreducible_d-m g by auto
  with eq show irreducible_d-m f unfolding irreducible_d-m-def by auto
 qed
qed

lemma unique-monic-hensel-factorization:
 assumes ufact: unique-factorization-m C (1,Fs)
 and C: monic C square-free-m C
 and n: n ≠ 0
 shows ∃ Gs. poly-mod.unique-factorization-m (p^n) C (1, Gs)
```

```
  using ufact C
proof (induct Fs arbitrary: C rule: wf-induct[OF wf-measure[of size]])
  case (1 Fs C)
  let ?q = p^n
  from n m1 have q: ?q > 1 by simp
  interpret q: poly-mod-2 ?q by (standard, insert q, auto)
  note [simp] = Mp-Mp-pow-is-Mp[OF n m1]
  note IH = 1(1)[rule-format]
  note ufact = 1(2)
  hence fact: factorization-m C (1, Fs) unfolding unique-factorization-m-alt-def
by auto
  note monC = 1(3)
  note sf = 1(4)
  let ?n = size Fs
  {
    fix d gs
    assume qfact: q.factorization-m C (d,gs)
    from q.factorization-m-lead-coeff[OF this] q.monic-Mp[OF monC]
    have d1: q.M d = 1 by auto

    from factorization-pn-to-factorization-p[OF qfact sf n]
    have factorization-m C (d,gs) .
    with ufact d1 have q.M d = 1 M d = 1 image-mset Mp gs = image-mset Mp
Fs
      unfolding unique-factorization-m-alt-def Mf-def by auto
  } note pre-unique = this
  show ?case
  proof (cases Fs)
    case empty
    with fact C have Mp C = 1 unfolding factorization-m-def by auto
    hence degree (Mp C) = 0 by simp
    with degree-m-eq-monic[OF monC m1] have degree C = 0 by simp
    with monC have C1: C = 1 using monic-degree-0 by blast
    with fact have fact: q.factorization-m C (1,{#})
      by (auto simp: q.factorization-m-def)
    show ?thesis
    proof (rule exI, rule q.unique-factorization-mI[OF fact])
      fix d gs
      assume fact: q.factorization-m C (d,gs)
      from pre-unique[OF this, unfolded empty]
      show q.Mf (d, gs) = q.Mf (1, {#}) by (auto simp: q.Mf-def)
    qed
  next
    case (add D H) note FDH = this
    let ?D = Mp D
    let ?H = Mp (prod-mset H)
    from fact have monFs: ⋀ F. F ∈# Fs ⟹ monic (Mp F)
      and prod: eq-m C (prod-mset Fs) unfolding factorization-m-def by auto
    hence monD: monic ?D unfolding FDH by auto
```

from *square-free-m-cong*[*OF sf*, *of D* ∗ *prod-mset H*] *prod*[*unfolded FDH*]
have *square-free-m* (*D* ∗ *prod-mset H*) **by** (*auto simp*: *ac-simps*)
from *square-free-m-prod-imp-coprime-m*[*OF this*]
have *coprime-m D* (*prod-mset H*) **.**
hence *cop'*: *coprime-m ?D ?H* **unfolding** *coprime-m-def dvdm-def Mp-Mp* **by** *simp*
from *fact* **have** *eq'*: *eq-m* (*?D* ∗ *?H*) *C*
unfolding *FDH* **by** (*simp add*: *factorization-m-def ac-simps*)
note *unique-hensel-binary*[*OF prime cop' eq' Mp-Mp Mp-Mp monD n*]
from *ex1-implies-ex*[*OF this*] *this*
obtain *A B* **where** *CAB*: *q.eq-m* (*A* ∗ *B*) *C* **and** *monA*: *monic A* **and** *DA*: *eq-m ?D A*
and *HB*: *eq-m ?H B* **and** *norm*: *q.Mp A* = *A q.Mp B* = *B*
and *unique*: ⋀ *D' H'*. *q.eq-m* (*D'* ∗ *H'*) *C* ⟹
monic D' ⟹
eq-m (*Mp D*) *D'* ⟹ *eq-m* (*Mp* (*prod-mset H*)) *H'* ⟹ *q.Mp D'* = *D'* ⟹ *q.Mp H'* = *H'*
⟹ *D'* = *A* ∧ *H'* = *B* **by** *blast*
note *hensel-bin-wit* = *CAB monA DA HB norm*
from *monA* **have** *monA'*: *monic* (*q.Mp A*) **by** (*rule q.monic-Mp*)
from *q.monic-Mp*[*OF monC*] *CAB* **have** *monicP*:*monic* (*q.Mp* (*A* ∗ *B*)) **by** *auto*
have *f4*: ⋀*p*. *coeff* (*A* ∗ *p*) (*degree* (*A* ∗ *p*)) = *coeff p* (*degree p*)
by (*simp add*: *coeff-degree-mult monA*)
have *f2*: ⋀*p n i*. *coeff p n mod i* = *coeff* (*poly-mod.Mp i p*) *n*
using *poly-mod.M-def poly-mod.Mp-coeff* **by** *presburger*
hence *coeff B* (*degree B*) = *0* ∨ *monic B*
using *monicP f4* **by** (*metis* (*no-types*) *norm*(*2*) *q.degree-m-eq q.m1*)
hence *monB*: *monic B*
using *f4 monicP* **by** (*metis norm*(*2*) *leading-coeff-0-iff*)
from *monA monB* **have** *lcAB*: *lead-coeff* (*A* ∗ *B*) = *1* **by** (*rule monic-mult*)
hence *copAB*: *coprime* (*lead-coeff* (*A* ∗ *B*)) *p* **by** *auto*
from *arg-cong*[*OF CAB*, *of Mp*]
have *CAB'*: *eq-m C* (*A* ∗ *B*) **by** *auto*
from *sf CAB'* **have** *sfAB*: *square-free-m* (*A* ∗ *B*) **using** *square-free-m-cong* **by** *blast*
from *CAB' ufact* **have** *ufact*: *unique-factorization-m* (*A* ∗ *B*) (*1*, *Fs*)
using *unique-factorization-m-cong* **by** *blast*
have (*1* :: *nat*) ≠ *0 p* = *p* ^ *1* **by** *auto*
note *u-factor* = *unique-factorization-factorI*[*OF prime ufact copAB sfAB this*]
from *fact DA* **have** *irreducible*$_d$*-m D eq-m A D* **unfolding** *add factorization-m-def* **by** *auto*
hence *irreducible*$_d$*-m A* **using** *Mp-irreducible*$_d$*-m* **by** *fastforce*
from *irreducible*$_d$*-lifting*[*OF n - this*] **have** *irrA*: *q.irreducible*$_d$*-m A* **using** *monA*
by (*simp add*: *m1 poly-mod.degree-m-eq-monic q.m1*)

from *add* **have** *lenH*: (*H*,*Fs*) ∈ *measure size* **by** *auto*
from *HB fact* **have** *factB*: *factorization-m B* (*1*, *H*)
unfolding *FDH factorization-m-def* **by** *auto*

from *u-factor*(*2*)[*OF factB*] **have** *ufactB*: *unique-factorization-m B* (*1*, *H*) **.**

from *sfAB* **have** *sfB*: *square-free-m B* **by** (*rule square-free-m-factor*)
from *IH*[*OF lenH ufactB monB sfB*] **obtain** *Bs* **where**
IH2: *q.unique-factorization-m B* (*1*, *Bs*) **by** *auto*

from *CAB* **have** *q.Mp C* = *q.Mp* (*q.Mp A* ∗ *q.Mp B*) **by** *simp*
also have *q.Mp A* ∗ *q.Mp B* = *q.Mp A* ∗ *q.Mp* (*prod-mset Bs*)
using *IH2* **unfolding** *q.unique-factorization-m-alt-def q.factorization-m-def*
by *auto*
also have *q.Mp* ... = *q.Mp* (*A* ∗ *prod-mset Bs*) **by** *simp*
finally have *factC*: *q.factorization-m C* (*1*, {# *A* #} + *Bs*) **using** *IH2 monA′ irrA*
by (*auto simp*: *q.unique-factorization-m-alt-def q.factorization-m-def*)
show *?thesis*
proof (*rule exI*, *rule q.unique-factorization-mI*[*OF factC*])
fix *d gs*
assume *dgs*: *q.factorization-m C* (*d*,*gs*)
from *pre-unique*[*OF dgs*, *unfolded add*] **have** *d1*: *q.M d* = *1* **and**
gs-fs: *image-mset Mp gs* = {# *Mp D* #} + *image-mset Mp H* **by** (*auto simp*: *ac-simps*)
have ∀ *f m p ma*. *image-mset f m* ≠ *add-mset* (*p*::*int poly*) *ma* ∨
(∃ *mb pa*. *m* = *add-mset* (*pa*::*int poly*) *mb* ∧ *f pa* = *p* ∧ *image-mset f mb* = *ma*)
by (*simp add*: *msed-map-invR*)
then obtain *g hs* **where** *gs*: *gs* = {# *g* #} + *hs* **and** *gD*: *Mp g* = *Mp D*
and *hsH*: *image-mset Mp hs* = *image-mset Mp H*
using *gs-fs* **by** (*metis add-mset-add-single union-commute*)
from *dgs*[*unfolded q.factorization-m-def split*]
have *eq*: *q.Mp C* = *q.Mp* (*smult d* (*prod-mset gs*))
and *irr-mon*: ⋀ *g*. *g*∈#*gs* ⟹ *q.irreducible*$_d$*-m g* ∧ *monic* (*q.Mp g*)
using *d1* **by** *auto*
note *eq*
also have *q.Mp* (*smult d* (*prod-mset gs*)) = *q.Mp* (*smult* (*q.M d*) (*prod-mset gs*))
by *simp*
also have ... = *q.Mp* (*prod-mset gs*) **unfolding** *d1* **by** *simp*
finally have *eq*: *q.eq-m* (*q.Mp g* ∗ *q.Mp* (*prod-mset hs*)) *C* **unfolding** *gs* **by** *simp*
from *gD* **have** *Dg*: *eq-m* (*Mp D*) (*q.Mp g*) **by** *simp*
have *Mp* (*prod-mset H*) = *Mp* (*prod-mset* (*image-mset Mp H*)) **by** *simp*
also have ... = *Mp* (*prod-mset hs*) **unfolding** *hsH*[*symmetric*] **by** *simp*
finally have *Hhs*: *eq-m* (*Mp* (*prod-mset H*)) (*q.Mp* (*prod-mset hs*)) **by** *simp*
from *irr-mon*[*of g*, *unfolded gs*] **have** *mon-g*: *monic* (*q.Mp g*) **by** *auto*
from *unique*[*OF eq mon-g Dg Hhs q.Mp-Mp q.Mp-Mp*]
have *gA*: *q.Mp g* = *A* **and** *hsB*: *q.Mp* (*prod-mset hs*) = *B* **by** *auto*
have *q.factorization-m B* (*1*, *hs*) **unfolding** *q.factorization-m-def split*
by (*simp add*: *hsB norm irr-mon*[*unfolded gs*])
with *IH2* **have** *hsBs*: *q.Mf* (*1*,*hs*) = *q.Mf* (*1*,*Bs*) **unfolding** *q.unique-factorization-m-alt-def*

```
by blast
      show q.Mf (d, gs) = q.Mf (1, {# A #} + Bs)
        using gA hsBs d1 unfolding gs q.Mf-def by auto
    qed
  qed
qed

theorem berlekamp-hensel-unique:
  assumes cop: coprime (lead-coeff f) p
  and sf: poly-mod.square-free-m p f
  and res: berlekamp-hensel p n f = gs
  and n: n ≠ 0
  shows poly-mod.unique-factorization-m (p^n) f (lead-coeff f, mset gs) — unique
factorization mod p^n
    ⋀ g. g ∈ set gs ⟹ poly-mod.Mp (p^n) g = g   — normalized
proof −
  let ?q = p^n
  interpret q: poly-mod-2 ?q unfolding poly-mod-2-def using m1 n by simp
  from berlekamp-hensel[OF assms]
  have bh-fact: q.factorization-m f (lead-coeff f, mset gs) by auto
  from berlekamp-hensel[OF assms]
  show ⋀ g. g ∈ set gs ⟹ poly-mod.Mp (p^n) g = g by blast
    from prime have p1: p > 1 by (simp add: prime-int-iff)
  let ?lc = coeff f (degree f)
  define ilc where ilc ≡ inverse-mod ?lc (p ^ n)
  from cop p1 n have inv: q.M (ilc * ?lc) = 1
    by (auto simp add: q.M-def ilc-def inverse-mod-pow)
  hence ilc0: ilc ≠ 0 by (cases ilc = 0, auto)
  {
    fix q
    assume ilc * ?lc = ?q * q
    from arg-cong[OF this, of q.M] have q.M (ilc * ?lc) = 0
      unfolding q.M-def by auto
    with inv have False by auto
  } note not-dvd = this
  let ?in = q.Mp (smult ilc f)
  have mon: monic ?in unfolding q.Mp-coeff coeff-smult
    by (subst q.degree-m-eq[OF - q.m1], insert not-dvd, auto simp: inv ilc0)
   have q.Mp f = q.Mp (smult (q.M (?lc * ilc)) f) using inv by (simp add:
ac-simps)
  also have ... = q.Mp (smult ?lc (smult ilc f)) by simp
  finally have f: q.Mp f = q.Mp (smult ?lc (smult ilc f)) .
  from arg-cong[OF f, of Mp]
  have Mp f = Mp (smult ?lc (smult ilc f))
    by (simp add: Mp-Mp-pow-is-Mp n p1)
  from arg-cong[OF this, of square-free-m, unfolded Mp-square-free-m] sf
  have square-free-m (smult (coeff f (degree f)) (smult ilc f)) by simp
  from square-free-m-smultD[OF this] have sf: square-free-m (smult ilc f) .
  have Mp-in: Mp ?in = Mp (smult ilc f)
```

by (*simp add*: *Mp-Mp-pow-is-Mp n p1*)
from *Mp-square-free-m*[*of ?in*, *unfolded Mp-in*] *sf* **have** *sf*: *square-free-m ?in*
unfolding *Mp-square-free-m* **by** *simp*
obtain *a b* **where** *finite-field-factorization-int p ?in* = (*a*,*b*) **by** *force*
from *finite-field-factorization-int*[*OF sf this*]
have *ufact*: *unique-factorization-m ?in* (*a*, *mset b*) **by** *auto*
from *unique-factorization-m-imp-factorization*[*OF this*]
have *fact*: *factorization-m ?in* (*a*, *mset b*) **.**
from *factorization-m-lead-coeff*[*OF this*] *monic-Mp*[*OF mon*]
have *M a* = *1* **by** *auto*
with *ufact* **have** *unique-factorization-m ?in* (*1*, *mset b*)
unfolding *unique-factorization-m-def Mf-def* **by** *auto*
from *unique-monic-hensel-factorization*[*OF this mon sf n*]
obtain *hs* **where** *q.unique-factorization-m ?in* (*1*, *hs*) **by** *auto*
hence *unique*: *q.unique-factorization-m* (*smult ilc f*) (*1*, *hs*)
unfolding *unique-factorization-m-def Mf-def* **by** *auto*
from *q.factorization-m-smult*[*OF q.unique-factorization-m-imp-factorization*[*OF unique*], *of ?lc*]
have *q.factorization-m* (*smult* (*ilc* ∗ *?lc*) *f*) (*?lc*, *hs*) **by** (*simp add*: *ac-simps*)
moreover have *q.Mp* (*smult* (*q.M* (*ilc* ∗ *?lc*)) *f*) = *q.Mp f* **unfolding** *inv* **by** *simp*
ultimately have *fact*: *q.factorization-m f* (*?lc*, *hs*)
unfolding *q.factorization-m-def* **by** *auto*
have *q.unique-factorization-m f* (*?lc*, *hs*)
proof (*rule q.unique-factorization-mI*[*OF fact*])
fix *d us*
assume *other-fact*: *q.factorization-m f* (*d*,*us*)
from *q.factorization-m-lead-coeff*[*OF this*] **have** *lc*: *q.M d* = *lead-coeff* (*q.Mp f*) **..**
have *lc*: *q.M d* = *q.M ?lc* **unfolding** *lc*
by (*metis bh-fact q.factorization-m-lead-coeff*)
from *q.factorization-m-smult*[*OF other-fact*, *of ilc*] *unique*
have *eq*: *q.Mf* (*d* ∗ *ilc*, *us*) = *q.Mf* (*1*, *hs*) **unfolding** *q.unique-factorization-m-def* **by** *auto*
thus *q.Mf* (*d*, *us*) = *q.Mf* (*?lc*, *hs*) **using** *lc* **unfolding** *q.Mf-def* **by** *auto*
qed
with *bh-fact* **show** *q.unique-factorization-m f* (*lead-coeff f*, *mset gs*)
unfolding *q.unique-factorization-m-alt-def* **by** *metis*
qed

lemma *hensel-lifting-unique*:
assumes *n*: *n* ≠ *0*
and *res*: *hensel-lifting p n f fs* = *gs* — result of hensel is fact. *gs*
and *cop*: *coprime* (*lead-coeff f*) *p*
and *sf*: *poly-mod.square-free-m p f*
and *fact*: *poly-mod.factorization-m p f* (*c*, *mset fs*) — input is fact. *fs mod p*
and *c*: *c* ∈ {*0*..<*p*}
and *norm*: (∀ *fi*∈*set fs*. *set* (*coeffs fi*) ⊆ {*0*..<*p*})

shows *poly-mod.unique-factorization-m* $(p\hat{}n)$ f (*lead-coeff* f, *mset gs*) — unique factorization mod $p\hat{}n$
$\quad$ *sort* (*map degree fs*) = *sort* (*map degree gs*) — degrees stay the same
$\quad$ $\bigwedge g.\ g \in set\ gs \Longrightarrow monic\ g \wedge poly\text{-}mod.Mp\ (p\hat{}n)\ g = g \wedge$ — monic and normalized
$\quad$ *poly-mod.irreducible-m p g* $\wedge$ — irreducibility even mod p
$\quad$ *poly-mod.degree-m p g* = *degree g* — mod p does not change degree of g
proof −
note *hensel* = *hensel-lifting*[*OF assms*]
show *sort* (*map degree fs*) = *sort* (*map degree gs*)
$\bigwedge g.\ g \in set\ gs \Longrightarrow monic\ g \wedge poly\text{-}mod.Mp\ (p\hat{}n)\ g = g \wedge$
poly-mod.irreducible-m p g $\wedge$
poly-mod.degree-m p g = *degree g* **using** *hensel* **by** *auto*
from *berlekamp-hensel-unique*[*OF cop sf refl n*]
have *poly-mod.unique-factorization-m* $(p\ \hat{}n)$ f (*lead-coeff* f, *mset* (*berlekamp-hensel p n f*)) **by** *auto*
with *hensel*(*1*) **show** *poly-mod.unique-factorization-m* $(p\hat{}n)$ f (*lead-coeff* f, *mset gs*)
by (*metis poly-mod.unique-factorization-m-alt-def*)
qed

end

end

10 Reconstructing Factors of Integer Polynomials

10.1 Square-Free Polynomials over Finite Fields and Integers

theory *Square-Free-Int-To-Square-Free-GFp*
imports
Subresultants.Subresultant-Gcd
Polynomial-Factorization.Rational-Factorization
Finite-Field
Polynomial-Factorization.Square-Free-Factorization
begin

lemma *square-free-int-rat*: **assumes** *sf*: *square-free f*
shows *square-free* (*map-poly rat-of-int f*)
proof −
let *?r* = *map-poly rat-of-int*
from *sf*[*unfolded square-free-def*] **have** *f0*: $f \neq 0 \bigwedge q.\ degree\ q \neq 0 \Longrightarrow \neg\ q * q\ dvd\ f$ **by** *auto*
show *?thesis*
proof (*rule square-freeI*)
show *?r f* $\neq$ *0* **using** *f0* **by** *auto*

```
    fix q
    assume dq: degree q > 0 and dvd: q * q dvd ?r f
    hence q0: q ≠ 0 by auto
    obtain q' c where norm: rat-to-normalized-int-poly q = (c,q') by force
    from rat-to-normalized-int-poly[OF norm] have c0: c ≠ 0 by auto
    note q = rat-to-normalized-int-poly(1)[OF norm]
    from dvd obtain k where rf: ?r f = q * (q * k) unfolding dvd-def by (auto
simp: ac-simps)
    from rat-to-int-factor-explicit[OF this norm] obtain s where
      f: f = q' * smult (content f) s by auto
    let ?s = smult (content f) s
    from arg-cong[OF f, of ?r] c0
    have ?r f = q * (smult (inverse c) (?r ?s))
      by (simp add: field-simps q hom-distribs)
    from arg-cong[OF this[unfolded rf], of λ f. f div q] q0
    have q * k = smult (inverse c) (?r ?s)
      by (metis nonzero-mult-div-cancel-left)
    from arg-cong[OF this, of smult c] have ?r ?s = q * smult c k using c0
      by (auto simp: field-simps)
    from rat-to-int-factor-explicit[OF this norm] obtain t where ?s = q' * t by
blast
     from f[unfolded this] sf[unfolded square-free-def] f0 have degree q' = 0 by
auto
    with rat-to-normalized-int-poly(4)[OF norm] dq show False by auto
  qed
qed

lemma content-free-unit:
  assumes content (p::'a::{idom,semiring-gcd} poly) = 1
  shows p dvd 1 ⟷ degree p = 0
  by (insert assms, auto dest!:degree0-coeffs simp: normalize-1-iff poly-dvd-1)

lemma square-free-imp-resultant-non-zero: assumes sf: square-free (f :: int poly)
  shows resultant f (pderiv f) ≠ 0
proof (cases degree f = 0)
  case True
  from degree0-coeffs[OF this] obtain c where f: f = [:c:] by auto
  with sf have c: c ≠ 0 unfolding square-free-def by auto
  show ?thesis unfolding f by simp
next
  case False note deg = this
  define pp where pp = primitive-part f
  define c where c = content f
  from sf have f0: f ≠ 0 unfolding square-free-def by auto
  hence c0: c ≠ 0 unfolding c-def by auto
 have f: f = smult c pp unfolding c-def pp-def unfolding content-times-primitive-part[of
f] ..
  from sf[unfolded f] c0 have sf': square-free pp by (metis dvd-smult smult-0-right
square-free-def)
```

```
  from deg[unfolded f] c0 have deg': ⋀ x. degree pp > 0 ∨ x by auto
  from content-primitive-part[OF f0] have cp: content pp = 1 unfolding pp-def
.
  let ?p' = pderiv pp
  {
    assume resultant pp ?p' = 0
    from this[unfolded resultant-0-gcd] have ¬ coprime pp ?p' by auto
    then obtain r where r: r dvd pp r dvd ?p' ¬ r dvd 1
      by (blast elim: not-coprimeE)
    from r(1) obtain k where pp = r * k ..
    from pos-zmult-eq-1-iff-lemma[OF arg-cong[OF this,
      of content, unfolded content-mult cp, symmetric]] content-ge-0-int[of r]
    have cr: content r = 1 by auto
    with r(3) content-free-unit have dr: degree r ≠ 0 by auto
    let ?r = map-poly rat-of-int
      from r(1) have dvd: ?r r dvd ?r pp unfolding dvd-def by (auto simp:
hom-distribs)
   from r(2) have ?r r dvd ?r ?p' apply (intro of-int-poly-hom.hom-dvd) by auto
    also have ?r ?p' = pderiv (?r pp) unfolding of-int-hom.map-poly-pderiv ..
    finally have dvd': ?r r dvd pderiv (?r pp) by auto
    from dr have dr': degree (?r r) ≠ 0 by simp
    from square-free-imp-separable[OF square-free-int-rat[OF sf']]
    have separable (?r pp) .
    hence cop: coprime (?r pp) (pderiv (?r pp)) unfolding separable-def .
    from f0 f have pp0: pp ≠ 0 by auto
    from dvd dvd' have ?r r dvd gcd (?r pp) (pderiv (?r pp)) by auto
     from divides-degree[OF this] pp0 have degree (?r r) ≤ degree (gcd (?r pp)
(pderiv (?r pp)))
      by auto
    with dr' have degree (gcd (?r pp) (pderiv (?r pp))) ≠ 0 by auto
    hence ¬ coprime (?r pp) (pderiv (?r pp)) by auto
    with cop have False by auto
  }
  hence resultant pp ?p' ≠ 0 by auto
  with resultant-smult-left[OF c0, of pp ?p', folded f] c0
  have resultant f ?p' ≠ 0 by auto
  with resultant-smult-right[OF c0, of f ?p', folded pderiv-smult f] c0
  show resultant f (pderiv f) ≠ 0 by auto
qed

lemma large-mod-0: assumes (n :: int) > 1 |k| < n k mod n = 0 shows k = 0
proof −
  from ‹k mod n = 0› have n dvd k
    by auto
  then obtain m where k = n * m ..
  with ‹n > 1› ‹|k| < n› show ?thesis
    by (auto simp add: abs-mult)
qed
```

definition *separable-bound* :: *int poly* ⇒ *int* **where**
separable-bound f = *max* (*abs* (*resultant f* (*pderiv f*)))
(*max* (*abs* (*lead-coeff f*)) (*abs* (*lead-coeff* (*pderiv f*))))

lemma *square-free-int-imp-resultant-non-zero-mod-ring*: **assumes** *sf*: *square-free f*

and *large*: *int CARD*('*a*) > *separable-bound f*
shows *resultant* (*map-poly of-int f* :: '*a* :: *prime-card mod-ring poly*) (*pderiv* (*map-poly of-int f*)) ≠ *0*
∧ *map-poly of-int f* ≠ (*0* :: '*a mod-ring poly*)
proof (*intro conjI*, *rule notI*)
let *?i* = *of-int* :: *int* ⇒ '*a mod-ring*
let *?m* = *of-int-poly* :: - ⇒ '*a mod-ring poly*
let *?f* = *?m f*
from *sf*[*unfolded square-free-def*] **have** *f0*: *f* ≠ *0* **by** *auto*
hence *lf*: *lead-coeff f* ≠ *0* **by** *auto*
{
fix *k* :: *int*
have *C1*: *int CARD*('*a*) > *1* **using** *prime-card*[**where** '*a* = '*a*] **by** (*auto simp*: *prime-nat-iff*)
assume *abs k* < *CARD*('*a*) **and** *?i k* = *0*
hence *k* = *0* **unfolding** *of-int-of-int-mod-ring*
by (*transfer*) (*rule large-mod-0*[*OF C1*])
} note *of-int-0* = *this*
from *square-free-imp-resultant-non-zero*[*OF sf*]
have *non0*: *resultant f* (*pderiv f*) ≠ *0* **.**
{
fix *g* :: *int poly*
assume *abs*: *abs* (*lead-coeff g*) < *CARD*('*a*)
have *degree* (*?m g*) = *degree g* **by** (*rule degree-map-poly*, *insert of-int-0*[*OF abs*], *auto*)
} note *deg* = *this*
note *large* = *large*[*unfolded separable-bound-def*]
from *of-int-0*[*of lead-coeff f*] *large lf* **have** *?i* (*lead-coeff f*) ≠ *0* **by** *auto*
thus *f0*: *?f* ≠ *0* **unfolding** *poly-eq-iff* **by** *auto*
assume *0*: *resultant ?f* (*pderiv ?f*) = *0*
have *resultant ?f* (*pderiv ?f*) = *?i* (*resultant f* (*pderiv f*))
unfolding *of-int-hom.map-poly-pderiv*[*symmetric*]
by (*subst of-int-hom.resultant-map-poly*(*1*)[*OF deg deg*], *insert large*, *auto simp*: *hom-distribs*)
from *of-int-0*[*OF - this*[*symmetric*, *unfolded 0*]] *non0*
show *False* **using** *large* **by** *auto*
qed

lemma *square-free-int-imp-separable-mod-ring*: **assumes** *sf*: *square-free f*
and *large*: *int CARD*('*a*) > *separable-bound f*
shows *separable* (*map-poly of-int f* :: '*a* :: *prime-card mod-ring poly*)
proof −
define *g* **where** *g* = *map-poly* (*of-int* :: *int* ⇒ '*a mod-ring*) *f*

from *square-free-int-imp-resultant-non-zero-mod-ring*[*OF sf large*]
have *res*: *resultant g* (*pderiv g*) ≠ *0* **and** *g*: *g* ≠ *0* **unfolding** *g-def* **by** *auto*
from *res*[*unfolded resultant-0-gcd*] **have** *degree* (*gcd g* (*pderiv g*)) = *0* **by** *auto*
from *degree0-coeffs*[*OF this*]
have *separable g* **unfolding** *separable-def*
by (*metis degree-pCons-0 g gcd-eq-0-iff is-unit-gcd is-unit-iff-degree*)
thus *?thesis* **unfolding** *g-def* **.**
qed

lemma *square-free-int-imp-square-free-mod-ring*: **assumes** *sf*: *square-free f*
and *large*: *int CARD*('*a*) > *separable-bound f*
shows *square-free* (*map-poly of-int f* :: '*a* :: *prime-card mod-ring poly*)
using *separable-imp-square-free*[*OF square-free-int-imp-separable-mod-ring*[*OF assms*]]
.

end

10.2 Finding a Suitable Prime

The Berlekamp-Zassenhaus algorithm demands for an input polynomial f to determine a prime p such that f is square-free mod p and such that p and the leading coefficient of f are coprime. To this end, we first prove that such a prime always exists, provided that f is square-free over the integers. Second, we provide a generic algorithm which searches for primes have a certain property P. Combining both results gives us the suitable prime for the Berlekamp-Zassenhaus algorithm.

theory *Suitable-Prime*
imports
Poly-Mod
Finite-Field-Record-Based
HOL−Types-To-Sets.Types-To-Sets
Poly-Mod-Finite-Field-Record-Based
Polynomial-Record-Based
Square-Free-Int-To-Square-Free-GFp
begin

lemma *square-free-separable-GFp*: **fixes** *f* :: '*a* :: *prime-card mod-ring poly*
assumes *card*: *CARD*('*a*) > *degree f*
and *sf*: *square-free f*
shows *separable f*
proof (*rule ccontr*)
assume ¬ *separable f*
with *square-free-separable-main*[*OF sf*]
obtain *g k* **where** ∗: *f* = *g* ∗ *k degree g* ≠ *0* **and** *g0*: *pderiv g* = *0* **by** *auto*
from *assms* **have** *f*: *f* ≠ *0* **unfolding** *square-free-def* **by** *auto*
have *degree f* = *degree g* + *degree k* **using** *f* **unfolding** ∗(*1*)
by (*subst degree-mult-eq*, *auto*)
with *card* **have** *card*: *degree g* < *CARD*('*a*) **by** *auto*

from $*(2)$ **obtain** n **where** *deg*: *degree* $g = Suc\ n$ **by** (*cases degree g*, *auto*)
from $*(2)$ **have** *cg*: *coeff g* (*degree g*) $\neq 0$ **by** *auto*
from *g0* **have** *coeff* (*pderiv g*) $n = 0$ **by** *auto*
from *this*[*unfolded coeff-pderiv*, *folded deg*] *cg*
have *of-nat* (*degree g*) = (*0* :: *'a mod-ring*) **by** *auto*
from *of-nat-0-mod-ring-dvd*[*OF this*] **have** *CARD*(*'a*) *dvd degree g* **.**
with *card* **show** *False* **by** (*simp add*: *deg nat-dvd-not-less*)
qed

lemma *square-free-iff-separable-GFp*: **assumes** *degree f* < *CARD*(*'a*)
shows *square-free* (*f* :: *'a* :: *prime-card mod-ring poly*) = *separable f*
using *separable-imp-square-free*[*of f*] *square-free-separable-GFp*[*OF assms*] **by** *auto*

definition *separable-impl-main* :: *int* ⇒ *'i arith-ops-record* ⇒ *int poly* ⇒ *bool* **where**
separable-impl-main p ff-ops f = *separable-i ff-ops* (*of-int-poly-i ff-ops* (*poly-mod.Mp p f*))

lemma (**in** *prime-field-gen*) *separable-impl*:
shows *separable-impl-main p ff-ops f* ⟹ *square-free-m f*
p > *degree-m f* ⟹ *p* > *separable-bound f* ⟹ *square-free f*
⟹ *separable-impl-main p ff-ops f* **unfolding** *separable-impl-main-def*
proof −
define *F* **where** *F*: (*F* :: *'a mod-ring poly*) = *of-int-poly* (*Mp f*)
let *?f'* = *of-int-poly-i ff-ops* (*Mp f*)
define *f''* **where** *f''* ≡ *of-int-poly* (*Mp f*) :: *'a mod-ring poly*
have *rel-f*[*transfer-rule*]: *poly-rel ?f' f''*
by (*rule poly-rel-of-int-poly*[*OF refl*], *simp add*: *f''-def*)
have *separable-i ff-ops ?f'* ⟷ *gcd f''* (*pderiv f''*) = *1*
unfolding *separable-i-def* **by** *transfer-prover*
also have ... ⟷ *coprime f''* (*pderiv f''*)
by (*auto simp add*: *gcd-eq-1-imp-coprime*)
finally have *id*: *separable-i ff-ops ?f'* ⟷ *separable f''* **unfolding** *separable-def coprime-iff-coprime* **.**
have *Mprel* [*transfer-rule*]: *MP-Rel* (*Mp f*) *F* **unfolding** *F MP-Rel-def*
by (*simp add*: *Mp-f-representative*)
have *square-free f''* = *square-free F* **unfolding** *f''-def F* **by** *simp*
also have ... = *square-free-m* (*Mp f*)
by (*transfer*, *simp*)
also have ... = *square-free-m f* **by** *simp*
finally have *id2*: *square-free f''* = *square-free-m f* **.**
from *separable-imp-square-free*[*of f''*]
show *separable-i ff-ops ?f'* ⟹ *square-free-m f*
unfolding *id id2* **by** *auto*
let *?m* = *map-poly* (*of-int* :: *int* ⇒ *'a mod-ring*)
let *?f* = *?m f*
assume *p* > *degree-m f* **and** *bnd*: *p* > *separable-bound f* **and** *sf*: *square-free f*
with *rel-funD*[*OF degree-MP-Rel Mprel*, *folded p*]

have $p > degree\ F$ **by** *simp*
hence $CARD('a) > degree\ f''$ **unfolding** f''*-def* $F\ p$ **by** *simp*
from *square-free-iff-separable-GFp*[*OF this*]
have *separable-i ff-ops ?f'* $=$ *square-free* f'' **unfolding** *id id2* **by** *simp*
also have $\ldots =$ *square-free* F **unfolding** f''*-def* F **by** *simp*
also have $F =$ *?f* **unfolding** F
by (*rule poly-eqI*, (*subst coeff-map-poly*, *force*)+, *unfold Mp-coeff*,
auto simp: *M-def*, *transfer*, *auto simp*: *p*)
also have *square-free ?f* **using** *square-free-int-imp-square-free-mod-ring*[**where**
$'a = 'a$, *OF sf*] *bnd m* **by** *auto*
finally
show *separable-i ff-ops ?f'* **.**
qed

context *poly-mod-prime* **begin**

lemmas *separable-impl-integer* $=$ *prime-field-gen.separable-impl*
[*OF prime-field.prime-field-finite-field-ops-integer*, *unfolded prime-field-def mod-ring-locale-def*,
unfolded poly-mod-type-simps, *internalize-sort* $'a :: prime\text{-}card$, *OF type-to-set*,
unfolded remove-duplicate-premise,cancel-type-definition, *OF non-empty*]

lemmas *separable-impl-uint32* $=$ *prime-field-gen.separable-impl*
[*OF prime-field.prime-field-finite-field-ops32*, *unfolded prime-field-def mod-ring-locale-def*,
unfolded poly-mod-type-simps, *internalize-sort* $'a :: prime\text{-}card$, *OF type-to-set*,
unfolded remove-duplicate-premise,cancel-type-definition, *OF non-empty*]

lemmas *separable-impl-uint64* $=$ *prime-field-gen.separable-impl*
[*OF prime-field.prime-field-finite-field-ops64* , *unfolded prime-field-def mod-ring-locale-def*,
unfolded poly-mod-type-simps, *internalize-sort* $'a :: prime\text{-}card$, *OF type-to-set*,
unfolded remove-duplicate-premise,cancel-type-definition, *OF non-empty*]

end

definition *separable-impl* :: *int* $\Rightarrow$ *int poly* $\Rightarrow$ *bool* **where**
separable-impl p $=$ (
if $p \leq 65535$
then separable-impl-main p (*finite-field-ops32* (*uint32-of-int p*))
else if $p \leq 4294967295$
then separable-impl-main p (*finite-field-ops64* (*uint64-of-int p*))
else separable-impl-main p (*finite-field-ops-integer* (*integer-of-int p*)))

lemma *square-free-mod-imp-square-free*: **assumes**
p: *prime p* **and** *sf*: *poly-mod.square-free-m p f*
and *cop*: *coprime* (*lead-coeff f*) *p*
shows *square-free f*
proof −
interpret *poly-mod p* **.**
from *sf*[*unfolded square-free-m-def*] **have** *f0*: *Mp f* $\neq$ *0* **and** *ndvd*: $\bigwedge$ *g. degree-m*
$g > 0 \Longrightarrow \neg\ (g * g)\ dvdm\ f$

```
    by auto
  from f0 have ff0: f ≠ 0 by auto
  show square-free f unfolding square-free-def
  proof (intro conjI[OF ff0] allI impI notI)
    fix g
    assume deg: degree g > 0 and dvd: g * g dvd f
    then obtain h where f: f = g * g * h unfolding dvd-def by auto
    from arg-cong[OF this, of Mp] have (g * g) dvdm f unfolding dvdm-def by
auto
    with ndvd[of g] have deg0: degree-m g = 0 by auto
    hence g0: M (lead-coeff g) = 0 unfolding Mp-def using deg
      by (metis M-def deg0 p poly-mod.degree-m-eq prime-gt-1-int neq0-conv)
    from p have p0: p ≠ 0 by auto
    from arg-cong[OF f, of lead-coeff] have lead-coeff f = lead-coeff g * lead-coeff
g * lead-coeff h
      by (auto simp: lead-coeff-mult)
    hence lead-coeff g dvd lead-coeff f by auto
    with cop have cop: coprime (lead-coeff g) p
      by (auto elim: coprime-imp-coprime intro: dvd-trans)
    with p0 have coprime (lead-coeff g mod p) p by simp
    also have lead-coeff g mod p = 0
      using M-def g0 by simp
    finally show False using p
      unfolding prime-int-iff
      by (simp add: prime-int-iff)
  qed
qed

lemma(in poly-mod-prime) separable-impl:
  shows separable-impl p f ⟹ square-free-m f
    nat p > degree-m f ⟹ nat p > separable-bound f ⟹ square-free f
    ⟹ separable-impl p f
  using
    separable-impl-integer[of f]
    separable-impl-uint32[of f]
    separable-impl-uint64[of f]
  unfolding separable-impl-def by (auto split: if-splits)

lemma coprime-lead-coeff-large-prime: assumes prime: prime (p :: int)
  and large: p > abs (lead-coeff f)
  and f: f ≠ 0
  shows coprime (lead-coeff f) p
proof −
  {
    fix lc
    assume 0 < lc lc < p
    then have ¬ p dvd lc
      by (simp add: zdvd-not-zless)
    with ‹prime p› have coprime p lc
```

```
    by (auto intro: prime-imp-coprime)
  then have coprime lc p
    by (simp add: ac-simps)
 } note main = this
 define lc where lc = lead-coeff f
 from f have lc0: lc ≠ 0 unfolding lc-def by auto
 from large have large: p > abs lc unfolding lc-def by auto
 have coprime lc p
 proof (cases lc > 0)
  case True
  from large have p > lc by auto
  from main[OF True this] show ?thesis .
 next
  case False
  let ?mlc = - lc
  from large False lc0 have ?mlc > 0 p > ?mlc by auto
  from main[OF this] show ?thesis by simp
 qed
 thus ?thesis unfolding lc-def by auto
qed

lemma prime-for-berlekamp-zassenhaus-exists: assumes sf: square-free f
 shows ∃ p. prime p ∧ (coprime (lead-coeff f) p ∧ separable-impl p f)
proof (rule ccontr)
 from assms have f0: f ≠ 0 unfolding square-free-def by auto
 define n where n = max (max (abs (lead-coeff f)) (degree f)) (separable-bound
f)
 assume contr: ¬ ?thesis
 {
  fix p :: int
  assume prime: prime p and n: p > n
  then interpret poly-mod-prime p by unfold-locales
   from n have large: p > abs (lead-coeff f) nat p > degree f nat p > separa-
ble-bound f
    unfolding n-def by auto
  from coprime-lead-coeff-large-prime[OF prime large(1) f0]
  have cop: coprime (lead-coeff f) p by auto
  with prime contr have nsf: ¬ separable-impl p f by auto
  from large(2) have nat p > degree-m f using degree-m-le[of f] by auto
  from separable-impl(2)[OF this large(3) sf] nsf have False by auto
 }
 hence no-large-prime: ⋀ p. prime p ⟹ p > n ⟹ False by auto
 from bigger-prime[of nat n] obtain p where ∗: prime p p > nat n by auto
 define q where q ≡ int p
 from ∗ have prime q q > n unfolding q-def by auto
 from no-large-prime[OF this]
 show False .
qed
```

```
definition next-primes :: nat ⇒ nat × nat list where
  next-primes n = (if n = 0 then next-candidates 0 else
    let (m,ps) = next-candidates n in (m,filter prime ps))

partial-function (tailrec) find-prime-main ::
  (nat ⇒ bool) ⇒ nat ⇒ nat list ⇒ nat where
  [code]: find-prime-main f np ps = (case ps of [] ⇒
    let (np',ps') = next-primes np
      in find-prime-main f np' ps'
    | (p # ps) ⇒ if f p then p else find-prime-main f np ps)

definition find-prime :: (nat ⇒ bool) ⇒ nat where
  find-prime f = find-prime-main f 0 []


lemma next-primes: assumes res: next-primes n = (m,ps)
  and n: candidate-invariant n
  shows candidate-invariant m sorted ps distinct ps n < m
  set ps = {i. prime i ∧ n ≤ i ∧ i < m}
proof −
  have candidate-invariant m ∧ sorted ps ∧ distinct ps ∧ n < m ∧
    set ps = {i. prime i ∧ n ≤ i ∧ i < m}
  proof (cases n = 0)
    case True
     with res[unfolded next-primes-def] have nc: next-candidates 0 = (m,ps) by
auto
    from this[unfolded next-candidates-def] have ps: ps = primes-1000 and m: m
= 1001 by auto
    have ps: set ps = {i. prime i ∧ n ≤ i ∧ i < m} unfolding m True ps
      by (auto simp: primes-1000)
    with next-candidates[OF nc n[unfolded True]] True
    show ?thesis by simp
  next
    case False
   with res[unfolded next-primes-def Let-def] obtain qs where nc: next-candidates
n = (m, qs)
      and ps: ps = filter prime qs by (cases next-candidates n, auto)
    have sorted qs ⟹ sorted ps unfolding ps using sorted-filter[of id qs prime]
by auto
    with next-candidates[OF nc n] show ?thesis unfolding ps by auto
  qed
  thus candidate-invariant m sorted ps distinct ps n < m
    set ps = {i. prime i ∧ n ≤ i ∧ i < m} by auto
qed

lemma find-prime: assumes ∃ n. prime n ∧ f n
  shows prime (find-prime f) ∧ f (find-prime f)
proof −
  from assms obtain n where fn: prime n f n by auto
```

```
{
  fix i ps m
  assume candidate-invariant i
    and n ∈ set ps ∨ n ≥ i
    and m = (Suc n − i, length ps)
    and ⋀ p. p ∈ set ps ⟹ prime p
  hence prime (find-prime-main f i ps) ∧ f (find-prime-main f i ps)
  proof (induct m arbitrary: i ps rule: wf-induct[OF wf-measures[of [fst, snd]]])
    case (1 m i ps)
    note IH = 1(1)[rule-format]
    note can = 1(2)
    note n = 1(3)
    note m = 1(4)
    note ps = 1(5)
    note simps [simp] = find-prime-main.simps[of f i ps]
    show ?case
    proof (cases ps)
      case Nil
      with n have i-n: i ≤ n by auto
      obtain j qs where np: next-primes i = (j,qs) by force
      note j = next-primes[OF np can]
       from j(4) i-n m have meas: ((Suc n − j, length qs), m) ∈ measures [fst,
snd] by auto
      have n: n ∈ set qs ∨ j ≤ n unfolding j(5) using i-n fn by auto
     show ?thesis unfolding simps using IH[OF meas j(1) n refl] j(5) by (simp
add: Nil np)
    next
      case (Cons p qs)
      show ?thesis
      proof (cases f p)
        case True
        thus ?thesis unfolding simps using ps unfolding Cons by simp
      next
        case False
          have m: ((Suc n − i, length qs), m) ∈ measures [fst, snd] using m
unfolding Cons by simp
        have n: n ∈ set qs ∨ i ≤ n using False n fn by (auto simp: Cons)
        from IH[OF m can n refl ps]
        show ?thesis unfolding simps using Cons False by simp
      qed
    qed
  qed
 } note main = this
 have candidate-invariant 0 by (simp add: candidate-invariant-def)
  from main[OF this - refl, of Nil] show ?thesis unfolding find-prime-def by
auto
qed

definition suitable-prime-bz :: int poly ⇒ int where
```

```
  suitable-prime-bz f ≡ let lc = lead-coeff f in int (find-prime (λ n. let p = int n
in
      coprime lc p ∧ separable-impl p f))

lemma suitable-prime-bz: assumes sf: square-free f and p: p = suitable-prime-bz
f
  shows prime p coprime (lead-coeff f) p poly-mod.square-free-m p f
proof −
  let ?lc = lead-coeff f
  from prime-for-berlekamp-zassenhaus-exists[OF sf, unfolded Let-def]
  obtain P where ∗: prime P ∧ coprime ?lc P ∧ separable-impl P f
    by auto
  hence prime (nat P) using prime-int-nat-transfer by blast
  with ∗ have ∃ p. prime p ∧ coprime ?lc (int p) ∧ separable-impl p f
    by (intro exI [of - nat P]) (auto dest: prime-gt-0-int)
  from find-prime[OF this]
  have prime: prime p and cop: coprime ?lc p and sf: separable-impl p f
    unfolding p suitable-prime-bz-def Let-def by auto
  then interpret poly-mod-prime p by unfold-locales
  from prime cop separable-impl(1)[OF sf]
  show prime p coprime ?lc p square-free-m f by auto
qed

definition square-free-heuristic :: int poly ⇒ int option where
  square-free-heuristic f = (let lc = lead-coeff f in
    find (λ p. coprime lc p ∧ separable-impl p f) [2, 3, 5, 7, 11, 13, 17, 19, 23])

lemma find-Some-D: find f xs = Some y ⟹ y ∈ set xs ∧ f y unfolding find-Some-iff
by auto

lemma square-free-heuristic: assumes square-free-heuristic f = Some p
  shows coprime (lead-coeff f) p ∧ separable-impl p f ∧ prime p
proof −
  from find-Some-D[OF assms[unfolded square-free-heuristic-def Let-def]]
  show ?thesis by auto
qed

end
```

10.3 Maximal Degree during Reconstruction

We define a function which computes an upper bound on the degree of a factor for which we have to reconstruct the integer values of the coefficients. This degree will determine how large the second parameter of the factor-bound will be.

In essence, if the Berlekamp-factorization will produce n factors with degrees $d_1, \ldots, d_n$, then our bound will be the sum of the $\frac{n}{2}$ largest degrees. The reason is that we will combine at most $\frac{n}{2}$ factors before reconstruction.

Soundness of the bound is proven, as well as a monotonicity property.

```
theory Degree-Bound
  imports Containers.Set-Impl
  HOL−Library.Multiset
  Polynomial-Interpolation.Missing-Polynomial
  Efficient−Mergesort.Efficient-Sort
begin

definition max-factor-degree :: nat list ⇒ nat where
  max-factor-degree degs = (let
     ds = sort degs
     in sum-list (drop (length ds div 2) ds))

definition degree-bound where degree-bound vs = max-factor-degree (map degree
vs)

lemma insort-middle: sort (xs @ x # ys) = insort x (sort (xs @ ys))
  by (metis append.assoc sort-append-Cons-swap sort-snoc)

lemma sum-list-insort[simp]:
  sum-list (insort (d :: 'a :: {comm-monoid-add,linorder}) xs) = d + sum-list xs
proof (induct xs)
  case (Cons x xs)
  thus ?case by (cases d ≤ x, auto simp: ac-simps)
qed simp

lemma half-largest-elements-mono: sum-list (drop (length ds div 2) (sort ds))
    ≤ sum-list (drop (Suc (length ds) div 2) (insort (d :: nat) (sort ds)))
proof −
  define n  where n  = length ds div 2
  define m  where m  = Suc (length ds) div 2
  define xs where xs = sort ds
  have xs: sorted xs unfolding xs-def by auto
  have nm: m ∈ {n, Suc n} unfolding n-def m-def by auto
  show ?thesis unfolding n-def[symmetric] m-def[symmetric] xs-def[symmetric]
    using nm xs
  proof (induct xs arbitrary: n m d)
    case (Cons x xs n m d)
    show ?case
    proof (cases n)
      case 0
      with Cons(2) have m: m = 0 ∨ m = 1 by auto
      show ?thesis
      proof (cases d ≤ x)
        case True
        hence ins: insort d (x # xs) = d # x # xs by auto
        show ?thesis unfolding ins 0 using True m by auto
      next
        case False
```

```
      hence ins: insort d (x # xs) = x # insort d xs by auto
      show ?thesis unfolding ins 0 using False m by auto
    qed
  next
    case (Suc nn)
     with Cons(2) obtain mm where m: m = Suc mm and mm: mm ∈ {nn,
Suc nn} by auto
    from Cons(3) have sort: sorted xs by (simp)
    note IH = Cons(1)[OF mm]
    show ?thesis
    proof (cases d ≤ x)
      case True
      with Cons(3) have ins: insort d (x # xs) = d # insort x xs
        by (cases xs, auto)
      show ?thesis unfolding ins Suc m using IH[OF sort] by auto
    next
      case False
      hence ins: insort d (x # xs) = x # insort d xs by auto
      show ?thesis unfolding ins Suc m using IH[OF sort] Cons(3) by auto
    qed
  qed
 qed auto
qed

lemma max-factor-degree-mono:
  max-factor-degree (map degree (fold remove1 ws vs)) ≤ max-factor-degree (map
degree vs)
 unfolding max-factor-degree-def Let-def length-sort length-map
proof (induct ws arbitrary: vs)
 case (Cons w ws vs)
 show ?case
 proof (cases w ∈ set vs)
   case False
   hence remove1 w vs = vs by (rule remove1-idem)
   thus ?thesis using Cons[of vs] by auto
 next
   case True
   then obtain bef aft where vs: vs = bef @ w # aft and rem1: remove1 w vs
= bef @ aft
     by (metis remove1.simps(2) remove1-append split-list-first)
   let ?exp = λ ws vs. sum-list (drop (length (fold remove1 ws vs) div 2)
     (sort (map degree (fold remove1 ws vs))))
   let ?bnd = λ vs. sum-list (drop (length vs div 2) (sort (map degree vs)))
   let ?bd = λ vs. sum-list (drop (length vs div 2) (sort vs))
   define ba where ba = bef @ aft
   define ds where ds = map degree ba
   define d  where d  = degree w
   have ?exp (w # ws) vs = ?exp ws (bef @ aft) by (auto simp: rem1)
   also have ... ≤ ?bnd ba unfolding ba-def by (rule Cons)
```

```
    also have ... = ?bd ds unfolding ds-def by simp
    also have ... ≤ sum-list (drop (Suc (length ds) div 2) (insort d (sort ds)))
      by (rule half-largest-elements-mono)
     also have ... = ?bnd vs unfolding vs ds-def d-def by (simp add: ba-def
insort-middle)
    finally show ?exp (w # ws) vs ≤ ?bnd vs by simp
  qed
qed auto

lemma mset-sub-decompose: mset ds ⊆# mset bs + as ⟹ length ds < length bs
⟹ ∃ b1 b b2.
   bs = b1 @ b # b2 ∧ mset ds ⊆# mset (b1 @ b2) + as
proof (induct ds arbitrary: bs as)
  case Nil
  hence bs = [] @ hd bs # tl bs by auto
  thus ?case by fastforce
next
  case (Cons d ds bs as)
  have d ∈# mset (d # ds) by auto
  with Cons(2) have d: d ∈# mset bs + as by (rule mset-subset-eqD)
  hence d ∈ set bs ∨ d ∈# as by auto
  thus ?case
  proof
    assume d ∈ set bs
    from this[unfolded in-set-conv-decomp] obtain b1 b2 where bs: bs = b1 @ d
# b2 by auto
    from Cons(2) Cons(3)
    have mset ds ⊆# mset (b1 @ b2) + as length ds < length (b1 @ b2) by (auto
simp: ac-simps bs)
    from Cons(1)[OF this] obtain b1' b b2' where split: b1 @ b2 = b1' @ b #
b2'
      and sub: mset ds ⊆# mset (b1' @ b2') + as by auto
    from split[unfolded append-eq-append-conv2]
    obtain us where b1 = b1' @ us ∧ us @ b2 = b # b2' ∨ b1 @ us = b1' ∧ b2
= us @ b # b2' ..
    thus ?thesis
    proof
      assume b1 @ us = b1' ∧ b2 = us @ b # b2'
      hence ∗: b1 @ us = b1' b2 = us @ b # b2' by auto
      hence bs: bs = (b1 @ d # us) @ b # b2' unfolding bs by auto
      show ?thesis
        by (intro exI conjI, rule bs, insert ∗ sub, auto simp: ac-simps)
    next
      assume b1 = b1' @ us ∧ us @ b2 = b # b2'
      hence ∗: b1 = b1' @ us us @ b2 = b # b2' by auto
      show ?thesis
      proof (cases us)
        case Nil
        with ∗ have ∗: b1 = b1' b2 = b # b2' by auto
```

```
      hence bs: bs = (b1' @ [d]) @ b # b2' unfolding bs by simp
      show ?thesis
        by (intro exI conjI, rule bs, insert * sub, auto simp: ac-simps)
    next
      case (Cons u vs)
      with * have *: b1 = b1' @ b # vs vs @ b2 = b2' by auto
      hence bs: bs = b1' @ b # (vs @ d # b2) unfolding bs by auto
      show ?thesis
        by (intro exI conjI, rule bs, insert * sub, auto simp: ac-simps)
    qed
  qed
next
  define as' where as' = as - {#d#}
  assume d ∈# as
  hence as': as = {#d#} + as' unfolding as'-def by auto
  from Cons(2)[unfolded as'] Cons(3) have mset ds ⊆# mset bs + as' length ds
< length bs
    by (auto simp: ac-simps)
  from Cons(1)[OF this] obtain b1 b b2 where bs: bs = b1 @ b # b2 and
    sub: mset ds ⊆# mset (b1 @ b2) + as' by auto
  show ?thesis
    by (intro exI conjI, rule bs, insert sub, auto simp: as' ac-simps)
qed
qed


lemma max-factor-degree-aux: fixes es :: nat list
  assumes sub: mset ds ⊆# mset es
    and len: length ds + length ds ≤ length es and sort: sorted es
  shows sum-list ds ≤ sum-list (drop (length es div 2) es)
proof -
  define bef where bef = take (length es div 2) es
  define aft where aft = drop (length es div 2) es
  have es: es = bef @ aft unfolding bef-def aft-def by auto
  from len have len: length ds ≤ length bef length ds ≤ length aft unfolding
bef-def aft-def
    by auto
  from sub have sub: mset ds ⊆# mset bef + mset aft unfolding es by auto
  from sort have sort: sorted (bef @ aft) unfolding es .
  show ?thesis unfolding aft-def[symmetric] using sub len sort
  proof (induct ds arbitrary: bef aft)
    case (Cons d ds bef aft)
    have d ∈# mset (d # ds) by auto
    with Cons(2) have d ∈# mset bef + mset aft by (rule mset-subset-eqD)
    hence d ∈ set bef ∨ d ∈ set aft by auto
    thus ?case
    proof
      assume d ∈ set aft
      from this[unfolded in-set-conv-decomp] obtain a1 a2 where aft: aft = a1 @
```

```
d # a2 by auto
      from Cons(4) have len-a: length ds ≤ length (a1 @ a2) unfolding aft by
auto
      from Cons(2)[unfolded aft] Cons(3)
       have mset ds ⊆# mset bef + (mset (a1 @ a2)) length ds < length bef by
auto
      from mset-sub-decompose[OF this]
      obtain b b1 b2
        where bef: bef = b1 @ b # b2 and sub: mset ds ⊆# (mset (b1 @ b2) +
mset (a1 @ a2)) by auto
      from Cons(3) have len-b: length ds ≤ length (b1 @ b2) unfolding bef by
auto
      from Cons(5)[unfolded bef aft] have sort: sorted ( (b1 @ b2) @ (a1 @ a2))
        unfolding sorted-append by auto
      note IH = Cons(1)[OF sub len-b len-a sort]
      show ?thesis using IH unfolding aft by simp
    next
      assume d ∈ set bef
      from this[unfolded in-set-conv-decomp] obtain b1 b2 where bef: bef = b1 @
d # b2 by auto
      from Cons(3) have len-b: length ds ≤ length (b1 @ b2) unfolding bef by
auto
      from Cons(2)[unfolded bef] Cons(4)
       have mset ds ⊆# mset aft + (mset (b1 @ b2)) length ds < length aft by
(auto simp: ac-simps)
      from mset-sub-decompose[OF this]
      obtain a a1 a2
        where aft: aft = a1 @ a # a2 and sub: mset ds ⊆# (mset (b1 @ b2) +
mset (a1 @ a2))
        by (auto simp: ac-simps)
      from Cons(4) have len-a: length ds ≤ length (a1 @ a2) unfolding aft by
auto
      from Cons(5)[unfolded bef aft] have sort: sorted ( (b1 @ b2) @ (a1 @ a2))
and ad: d ≤ a
        unfolding sorted-append by auto
      note IH = Cons(1)[OF sub len-b len-a sort]
      show ?thesis using IH ad unfolding aft by simp
    qed
  qed auto
qed

lemma max-factor-degree: assumes sub: mset ws ⊆# mset vs
  and len: length ws + length ws ≤ length vs
  shows degree (prod-list ws) ≤ max-factor-degree (map degree vs)
proof −
  define ds where ds ≡ map degree ws
  define es where es ≡ sort (map degree vs)
  from sub len have sub: mset ds ⊆# mset es and len: length ds + length ds ≤
length es
```

and *es*: *sorted es*
unfolding *ds-def es-def*
by (*auto simp*: *image-mset-subseteq-mono*)
have *degree* (*prod-list ws*) ≤ *sum-list* (*map degree ws*) **by** (*rule degree-prod-list-le*)
also have ... ≤ *max-factor-degree* (*map degree vs*)
unfolding *max-factor-degree-def Let-def ds-def*[*symmetric*] *es-def*[*symmetric*]
using *sub len es* **by** (*rule max-factor-degree-aux*)
finally show *?thesis* **.**
qed

lemma *degree-bound*: **assumes** *sub*: *mset ws* ⊆# *mset vs*
and *len*: *length ws* + *length ws* ≤ *length vs*
shows *degree* (*prod-list ws*) ≤ *degree-bound vs*
using *max-factor-degree*[*OF sub len*] **unfolding** *degree-bound-def* **by** *auto*

end

10.4 Mahler Measure

This part contains a definition of the Mahler measure, it contains Landau's inequality and the Graeffe-transformation. We also assemble a heuristic to approximate the Mahler's measure.

theory *Mahler-Measure*
imports
Sqrt-Babylonian.Sqrt-Babylonian
Poly-Mod-Finite-Field-Record-Based
Polynomial-Factorization.Fundamental-Theorem-Algebra-Factorized
Polynomial-Factorization.Missing-Multiset
begin

context *comm-monoid-list* **begin**
lemma *induct-gen-abs*:
assumes ⋀ *a r*. *a*∈*set lst* ⟹ *P* (*f* (*h a*) *r*) (*f* (*g a*) *r*)
⋀ *x y z*. *P x y* ⟹ *P y z* ⟹ *P x z*
P (*F* (*map g lst*)) (*F* (*map g lst*))
shows *P* (*F* (*map h lst*)) (*F* (*map g lst*))
using *assms* **proof**(*induct lst arbitrary*:*P*)
case (*Cons a as P*)
have *inl*:*a*∈*set* (*a*#*as*) **by** *auto*
let *?uf* = λ *v w*. *P* (*f* (*g a*) *v*) (*f* (*g a*) *w*)
have *p-suc*:*?uf* (*F* (*map g as*)) (*F* (*map g as*))
using *Cons.prems*(*3*) **by** *auto*
{ fix *r aa* **assume** *aa* ∈ *set as* **hence** *ins*:*aa* ∈ *set* (*a*#*as*) **by** *auto*
have *P* (*f* (*g a*) (*f* (*h aa*) *r*)) (*f* (*g a*) (*f* (*g aa*) *r*))
using *Cons.prems*(*1*)[*of aa f r* (*g a*),*OF ins*]
by (*auto simp*: *assoc commute left-commute*)
} note *h* = *this*
from *Cons.hyps*(*1*)[*of ?uf*, *OF h Cons.prems*(*2*)[*simplified*] *p-suc*]
have *e1*:*P* (*f* (*g a*) (*F* (*map h as*))) (*f* (*g a*) (*F* (*map g as*))) **by** *simp*

have *e2*:*P* (*f* (*h a*) (*F* (*map h as*))) (*f* (*g a*) (*F* (*map h as*)))
using *Cons.prems*(*1*)[*OF inl*] **by** *blast*
from *Cons*(*3*)[*OF e2 e1*] **show** *?case* **by** *auto* **next**
qed *auto*
end

lemma *prod-induct-gen*:
assumes $\bigwedge$ *a r*. *f* (*h a* ∗ *r* :: *'a* :: {*comm-monoid-mult*}) = *f* (*g a* ∗ *r*)
shows *f* (∏ *v*←*lst*. *h v*) = *f* (∏ *v*←*lst*. *g v*)
proof − **let** *?P x y* = *f x* = *f y*
show *?thesis* **using** *comm-monoid-mult-class.prod-list.induct-gen-abs*[*of - ?P*,*OF assms*] **by** *auto*
qed

abbreviation *complex-of-int*::*int* ⇒ *complex* **where**
complex-of-int ≡ *of-int*

definition *l2norm-list* :: *int list* ⇒ *int* **where**
l2norm-list lst = ⌊*sqrt* (*sum-list* (*map* (λ *a*. *a* ∗ *a*) *lst*))⌋

abbreviation *l2norm* :: *int poly* ⇒ *int* **where**
l2norm p ≡ *l2norm-list* (*coeffs p*)

abbreviation *norm2 p* ≡ ∑ *a*←*coeffs p*. (*cmod a*)^2

abbreviation *l2norm-complex* **where**
l2norm-complex p ≡ *sqrt* (*norm2 p*)

abbreviation *height* :: *int poly* ⇒ *int* **where**
height p ≡ *max-list* (*map* (*nat* ∘ *abs*) (*coeffs p*))

definition *complex-roots-complex* **where**
complex-roots-complex (*p*::*complex poly*) = (*SOME as*. *smult* (*coeff p* (*degree p*)) (∏ *a*←*as*. [:− *a*, *1*:]) = *p* ∧ *length as* = *degree p*)

lemma *complex-roots*:
smult (*lead-coeff p*) (∏ *a*←*complex-roots-complex p*. [:− *a*, *1*:]) = *p*
length (*complex-roots-complex p*) = *degree p*
using *someI-ex*[*OF fundamental-theorem-algebra-factorized*]
unfolding *complex-roots-complex-def* **by** *simp-all*

lemma *complex-roots-c* [*simp*]:
complex-roots-complex [:*c*:] = []
using *complex-roots*(*2*) [*of* [:*c*:]] **by** *simp*

declare *complex-roots*(*2*)[*simp*]

lemma *complex-roots-1* [*simp*]:
complex-roots-complex 1 = []

using *complex-roots-c* [*of 1*] **by** (*simp add*: *pCons-one*)

lemma *linear-term-irreducible*$_d$[*simp*]: *irreducible*$_d$ [: *a*, *1*:]
by (*rule linear-irreducible*$_d$, *simp*)

definition *complex-roots-int* **where**
complex-roots-int (*p*::*int poly*) = *complex-roots-complex* (*map-poly of-int p*)

lemma *complex-roots-int*:
smult (*lead-coeff p*) (∏ *a*←*complex-roots-int p*. [:− *a*, *1*:]) = *map-poly of-int p*
length (*complex-roots-int p*) = *degree p*
proof −
show *smult* (*lead-coeff p*) (∏ *a*←*complex-roots-int p*. [:− *a*, *1*:]) = *map-poly of-int p*
length (*complex-roots-int p*) = *degree p*
using *complex-roots*[*of map-poly of-int p*] **unfolding** *complex-roots-int-def* **by** *auto*
qed

The measure for polynomials, after K. Mahler

definition *mahler-measure-poly* **where**
mahler-measure-poly p = *cmod* (*lead-coeff p*) ∗ (∏ *a*←*complex-roots-complex p*. (*max 1* (*cmod a*)))

definition *mahler-measure* **where**
mahler-measure p = *mahler-measure-poly* (*map-poly complex-of-int p*)

definition *mahler-measure-monic* **where**
mahler-measure-monic p = (∏ *a*←*complex-roots-complex p*. (*max 1* (*cmod a*)))

lemma *mahler-measure-poly-via-monic* :
mahler-measure-poly p = *cmod* (*lead-coeff p*) ∗ *mahler-measure-monic p*
unfolding *mahler-measure-poly-def mahler-measure-monic-def* **by** *simp*

lemma *smult-inj*[*simp*]: **assumes** (*a*::*'a*::*idom*) ≠ *0* **shows** *inj* (*smult a*)
proof−
interpret *map-poly-inj-zero-hom* (∗) *a* **using** *assms* **by** (*unfold-locales*, *auto*)
show *?thesis* **unfolding** *smult-as-map-poly* **by** (*rule inj-f*)
qed

definition *reconstruct-poly*::*'a*::*idom* ⇒ *'a list* ⇒ *'a poly* **where**
reconstruct-poly c roots = *smult c* (∏ *a*←*roots*. [:− *a*, *1*:])

lemma *reconstruct-is-original-poly*:
reconstruct-poly (*lead-coeff p*) (*complex-roots-complex p*) = *p*
using *complex-roots*(*1*) **by** (*simp add*: *reconstruct-poly-def*)

lemma *reconstruct-with-type-conversion*:
smult (*lead-coeff* (*map-poly of-int f*)) (*prod-list* (*map* (λ *a*. [:− *a*, *1*:]) (*complex-roots-int*

f)))
  = *map-poly of-int f*
unfolding *complex-roots-int-def complex-roots*(*1*) **by** *simp*

lemma *reconstruct-prod*:
  shows *reconstruct-poly* (*a*::*complex*) *as* ∗ *reconstruct-poly b bs*
      = *reconstruct-poly* (*a* ∗ *b*) (*as* @ *bs*)
unfolding *reconstruct-poly-def* **by** *auto*

lemma *linear-term-inj*[*simplified*,*simp*]: *inj* (λ *a*. [:− *a*, *1*::′*a*::*idom*:])
  unfolding *inj-on-def* **by** *simp*

lemma *reconstruct-poly-monic-defines-mset*:
  assumes (∏ *a*←*as*. [:− *a*, *1*:]) = (∏ *a*←*bs*. [:− *a*, *1*::′*a*::*field*:])
  shows *mset as* = *mset bs*
proof −
  let *?as* = *mset* (*map* (λ *a*. [:− *a*, *1*:]) *as*)
  let *?bs* = *mset* (*map* (λ *a*. [:− *a*, *1*:]) *bs*)
 have *eq-smult*:*prod-mset ?as* = *prod-mset ?bs* **using** *assms* **by** (*metis prod-mset-prod-list*)
  have *irr*:⋀ *as*::′*a list*. *set-mset* (*mset* (*map* (λ *a*. [:− *a*, *1*:]) *as*)) ⊆ {*q*. *irreducible*
q ∧ *monic q*}
    by (*auto intro*!: *linear-term-irreducible*$_d$[*of* −-::′*a*, *simplified*])
  from *monic-factorization-unique-mset*[*OF eq-smult irr irr*]
  show *?thesis* **apply** (*subst inj-eq*[*OF multiset.inj-map*,*symmetric*]) **by** *auto*
qed

lemma *reconstruct-poly-defines-mset-of-argument*:
  assumes (*a*::′*a*::*field*) ≠ *0*
        reconstruct-poly a as = *reconstruct-poly a bs*
  shows *mset as* = *mset bs*
proof −
  have *eq-smult*:*smult a* (∏ *a*←*as*. [:− *a*, *1*:]) = *smult a* (∏ *a*←*bs*. [:− *a*, *1*:])
    using *assms*(*2*) **by** (*auto simp*:*reconstruct-poly-def*)
 from *reconstruct-poly-monic-defines-mset*[*OF Fun.injD*[*OF smult-inj*[*OF assms*(*1*)]
eq-smult]]
  show *?thesis* **by** *simp*
qed

lemma *complex-roots-complex-prod* [*simp*]:
  assumes *f* ≠ *0 g* ≠ *0*
  shows *mset* (*complex-roots-complex* (*f* ∗ *g*))
      = *mset* (*complex-roots-complex f*) + *mset* (*complex-roots-complex g*)
proof −
  let *?p* = *f* ∗ *g*
  let *?lc v* = (*lead-coeff* (*v*:: *complex poly*))
  have *nonzero-prod*:*?lc ?p* ≠ *0* **using** *assms* **by** *auto*
 from *reconstruct-prod*[*of ?lc f complex-roots-complex f ?lc g complex-roots-complex*
g]
  have *reconstruct-poly* (*?lc ?p*) (*complex-roots-complex ?p*)

```
    = reconstruct-poly (?lc ?p) (complex-roots-complex f @ complex-roots-complex
g)
   unfolding lead-coeff-mult[symmetric] reconstruct-is-original-poly by auto
  from reconstruct-poly-defines-mset-of-argument[OF nonzero-prod this]
  show ?thesis by simp
qed

lemma mset-mult-add:
  assumes mset (a::'a::field list) = mset b + mset c
  shows prod-list a = prod-list b * prod-list c
  unfolding prod-mset-prod-list[symmetric]
  using prod-mset-Un[of mset b mset c,unfolded assms[symmetric]].

lemma mset-mult-add-2:
  assumes mset a = mset b + mset c
  shows prod-list (map i a::'b::field list) = prod-list (map i b) * prod-list (map i c)
proof −
  have r:mset (map i a) = mset (map i b) + mset (map i c)  using assms
    by (metis map-append mset-append mset-map)
  show ?thesis using mset-mult-add[OF r] by auto
qed

lemma measure-mono-eq-prod:
  assumes f ≠ 0 g ≠ 0
 shows mahler-measure-monic (f * g) = mahler-measure-monic f * mahler-measure-monic
g
  unfolding mahler-measure-monic-def
  using mset-mult-add-2[OF complex-roots-complex-prod[OF assms],of λ a. max 1
(cmod a)] by simp

lemma mahler-measure-poly-0[simp]: mahler-measure-poly 0 = 0 unfolding mahler-measure-poly-via-monic
by auto

lemma measure-eq-prod:
  mahler-measure-poly (f * g) = mahler-measure-poly f * mahler-measure-poly g
proof −
  consider f = 0 | g = 0 | (both) f ≠ 0 g ≠ 0 by auto
  thus ?thesis proof(cases)
    case both show ?thesis unfolding mahler-measure-poly-via-monic norm-mult
lead-coeff-mult
      by (auto simp: measure-mono-eq-prod[OF both])
  qed (simp-all)
qed

lemma prod-cmod[simp]:
  cmod (∏ a←lst. f a) = (∏ a←lst. cmod (f a))
  by(induct lst,auto simp:real-normed-div-algebra-class.norm-mult)

lemma lead-coeff-of-prod[simp]:
```

lead-coeff ($\prod$ *a*←*lst*. *f a*::'*a*::*idom poly*) = ($\prod$ *a*←*lst*. *lead-coeff* (*f a*))
by(*induct lst*,*auto simp*:*lead-coeff-mult*)

lemma *ineq-about-squares*:**assumes** $x \leq (y::real)$ **shows** $x \leq c\hat{}2 + y$ **using** *assms*
by (*simp add*: *add.commute add-increasing2*)

lemma *first-coeff-le-tail*:(*cmod* (*lead-coeff g*))$\hat{}2 \leq$ ($\sum$ *a*←*coeffs g*. (*cmod a*)$\hat{}2$)
proof(*induct g*)
case (*pCons a p*)
thus *?case* **proof**(*cases p* = *0*) **case** *False*
show *?thesis* **using** *pCons* **unfolding** *lead-coeff-pCons*(*1*)[*OF False*]
by(*cases a* = *0*,*simp-all add*:*ineq-about-squares*)
qed *simp*
qed *simp*

lemma *square-prod-cmod*[*simp*]:
(*cmod* (*a* ∗ *b*))$\hat{}2$ = *cmod a* $\hat{}$ *2* ∗ *cmod b* $\hat{}$ *2*
by (*simp add*: *norm-mult power-mult-distrib*)

lemma *sum-coeffs-smult-cmod*:
($\sum$ *a*←*coeffs* (*smult v p*). (*cmod a*)$\hat{}2$) = (*cmod v*)$\hat{}2$ ∗ ($\sum$ *a*←*coeffs p*. (*cmod a*)$\hat{}2$)
(**is** *?l* = *?r*)
proof −
have *?l* = ($\sum$ *a*←*coeffs p*. (*cmod v*)$\hat{}2$ ∗ (*cmod a*)$\hat{}2$) **by**(*cases v*=*0*;*induct p*,*auto*)
thus *?thesis* **by** (*auto simp*:*sum-list-const-mult*)
qed

abbreviation *linH a* ≡ *if* (*cmod a* > *1*) *then* [:− *1*,*cnj a*:] *else* [:− *a*,*1*:]

lemma *coeffs-cong-1*[*simp*]: *cCons a v* = *cCons b v* ⟷ *a* = *b* **unfolding** *cCons-def*
by *auto*

lemma *strip-while-singleton*[*simp*]:
strip-while ((=) *0*) [*v* ∗ *a*] = *cCons* (*v* ∗ *a*) [] **unfolding** *cCons-def strip-while-def*
by *auto*

lemma *coeffs-times-linterm*:
shows *coeffs* (*pCons 0* (*smult a p*) + *smult b p*) = *strip-while* (*HOL.eq* (*0*::'*a*::{*comm-ring-1*}))
(*map* (λ(*c*,*d*).*b*∗*d*+*c*∗*a*) (*zip* (*0* # *coeffs p*) (*coeffs p* @ [*0*]))) **proof** −
{fix *v*
have *coeffs* (*smult b p* + *pCons* (*a*∗ *v*) (*smult a p*)) = *strip-while* (*HOL.eq 0*) (*map* (λ(*c*,*d*).*b*∗*d*+*c*∗*a*) (*zip* ([*v*] @ *coeffs p*) (*coeffs p* @ [*0*])))
proof(*induct p arbitrary*:*v*) **case** (*pCons pa ps*) **thus** *?case* **by** *auto* **qed** *auto*
}
from *this*[*of 0*] **show** *?thesis* **by** (*simp add*: *add.commute*)
qed

lemma *filter-distr-rev*[*simp*]:
shows *filter f* (*rev lst*) = *rev* (*filter f lst*)
by(*induct lst*;*auto*)

lemma *strip-while-filter*:
shows *filter* ((≠) *0*) (*strip-while* ((=) *0*) (*lst*::'*a*::*zero list*)) = *filter* ((≠) *0*) *lst*
proof − {**fix** *lst*::'*a list*
have *filter* ((≠) *0*) (*dropWhile* ((=) *0*) *lst*) = *filter* ((≠) *0*) *lst* **by** (*induct lst*;*auto*)
hence (*filter* ((≠) *0*) (*strip-while* ((=) *0*) (*rev lst*))) = *filter* ((≠) *0*) (*rev lst*)
unfolding *strip-while-def* **by**(*simp*)}
from *this*[*of rev lst*] **show** *?thesis* **by** *simp*
qed

lemma *sum-stripwhile*[*simp*]:
assumes *f 0* = *0*
shows (∑ *a*←*strip-while* ((=) *0*) *lst*. *f a*) = (∑ *a*←*lst*. *f a*)
proof −
{**fix** *lst*
have (∑ *a*←*filter* ((≠) *0*) *lst*. *f a*) = (∑ *a*←*lst*. *f a*) **by**(*induct lst*,*auto simp*:*assms*)}
note *f*=*this*
have *sum-list* (*map f* (*filter* ((≠) *0*) (*strip-while* ((=) *0*) *lst*)))
= *sum-list* (*map f* (*filter* ((≠) *0*) *lst*))
using *strip-while-filter*[*of lst*] **by**(*simp*)
thus *?thesis* **unfolding** *f*.
qed

lemma *complex-split* : *Complex a b* = *c* ⟷ (*a* = *Re c* ∧ *b* = *Im c*)
using *complex-surj* **by** *auto*

lemma *norm-times-const*:(∑ *y*←*lst*. (*cmod* (*a* ∗ *y*))^2) = (*cmod a*)^2 ∗ (∑ *y*←*lst*. (*cmod y*)^2)
by(*induct lst*,*auto simp*:*ring-distribs*)

fun *bisumTail* **where**
bisumTail f (*Cons a* (*Cons b bs*)) = *f a b* + *bisumTail f* (*Cons b bs*) |
bisumTail f (*Cons a Nil*) = *f a 0* |
bisumTail f Nil = *f 1 0*
fun *bisum* **where**
bisum f (*Cons a as*) = *f 0 a* + *bisumTail f* (*Cons a as*) |
bisum f Nil = *f 0 0*

lemma *bisumTail-is-map-zip*:
(∑ *x*←*zip* (*v* # *l1*) (*l1* @ [*0*]). *f x*) = *bisumTail* (λ*x y* .*f* (*x*,*y*)) (*v*#*l1*)
by(*induct l1 arbitrary*:*v*,*auto*)

lemma *bisum-is-map-zip*:

$(\sum x{\leftarrow}zip\ (0\ \#\ l1)\ (l1\ @\ [0]).\ f\ x) = bisum\ (\lambda x\ y.\ f\ (x,y))\ l1$
using *bisumTail-is-map-zip*[*of f hd l1 tl l1*] **by**(*cases l1*,*auto*)
lemma *map-zip-is-bisum*:
$bisum\ f\ l1 = (\sum (x,y){\leftarrow}zip\ (0\ \#\ l1)\ (l1\ @\ [0]).\ f\ x\ y)$
using *bisum-is-map-zip*[*of* $\lambda(x,y).\ f\ x\ y$] **by** *auto*

lemma *bisum-outside* :
$(bisum\ (\lambda\ x\ y.\ f1\ x - f2\ x\ y + f3\ y)\ lst :: 'a :: field)$
$= sum\text{-}list\ (map\ f1\ lst) + f1\ 0 - bisum\ f2\ lst + sum\text{-}list\ (map\ f3\ lst) + f3\ 0$
proof(*cases lst*)
case (*Cons a lst*) **show** *?thesis* **unfolding** *map-zip-is-bisum Cons* **by**(*induct lst arbitrary*:*a*,*auto*)
qed *auto*

lemma *Landau-lemma*:
$(\sum a{\leftarrow}coeffs\ (\prod a{\leftarrow}lst.\ [:-\ a,\ 1:]).\ (cmod\ a)^2) = (\sum a{\leftarrow}coeffs\ (\prod a{\leftarrow}lst.\ linH\ a).\ (cmod\ a)^2)$
(**is** *norm2 ?l* = *norm2 ?r*)
proof −
have *a*:$\bigwedge$ *a*. $(cmod\ a)^2 = Re\ (a * cnj\ a)$ **using** *complex-norm-square*
unfolding *complex-split complex-of-real-def* **by** *simp*
have *b*:$\bigwedge$ *x a y*. $(cmod\ (x - a * y))\hat{}2$
$= (cmod\ x)^2 - Re\ (a * y * cnj\ x + x * cnj\ (a * y)) + (cmod\ (a * y))\hat{}2$
unfolding *left-diff-distrib right-diff-distrib a complex-cnj-diff* **by** *simp*
have *c*:$\bigwedge$ *y a x*. $(cmod\ (cnj\ a * x - y))^2$
$= (cmod\ (a * x))^2 - Re\ (a * y * cnj\ x + x * cnj\ (a * y)) + (cmod\ y)\hat{}2$
unfolding *left-diff-distrib right-diff-distrib a complex-cnj-diff*
by (*simp add*: *mult.assoc mult.left-commute*)
{ fix *f1 a*
have *norm2* $([:-\ a,\ 1\ :] * f1) = bisum\ (\lambda x\ y.\ cmod\ (x - a * y)\hat{}2)\ (coeffs\ f1)$
by(*simp add*: *bisum-is-map-zip*[*of - coeffs f1*] *coeffs-times-linterm*[*of 1 - −a*,*simplified*])
also have $\ldots = norm2\ f1 + cmod\ a\hat{}2 * norm2\ f1$
$- bisum\ (\lambda x\ y.\ Re\ (a * y * cnj\ x + x * cnj\ (a * y)))\ (coeffs\ f1)$
unfolding *b bisum-outside norm-times-const* **by** *simp*
also have $\ldots = bisum\ (\lambda x\ y.\ cmod\ (cnj\ a * x - y)\hat{}2)\ (coeffs\ f1)$
unfolding *c bisum-outside norm-times-const* **by** *auto*
also have $\ldots = norm2\ ([:-\ 1,\ cnj\ a\ :] * f1)$
using *coeffs-times-linterm*[*of cnj a - −1*]
by(*simp add*: *bisum-is-map-zip*[*of - coeffs f1*] *mult.commute*)
finally have *norm2* $([:-\ a,\ 1\ :] * f1) = \ldots$.**}**
hence *h*:$\bigwedge$ *a f1*. *norm2* $([:-\ a,\ 1\ :] * f1) = norm2\ (linH\ a * f1)$ **by** *auto*
show *?thesis* **by**(*rule prod-induct-gen*[*OF h*])
qed

lemma *Landau-inequality*:
$mahler\text{-}measure\text{-}poly\ f \leq l2norm\text{-}complex\ f$

proof −
let *?f* = *reconstruct-poly* (*lead-coeff f*) (*complex-roots-complex f*)
let *?roots* = (*complex-roots-complex f*)
let *?g* = ∏ *a*←*?roots*. *linH a*

have *max*:⋀*a*. *cmod* (*if* 1 < *cmod a then cnj a else* 1) = *max* 1 (*cmod a*)
by *simp*
have ⋀*a*. 1 < *cmod a* ⟹ *a* ≠ 0 **by** *auto*
hence ⋀*a*. *lead-coeff* (*linH a*) = (*if* (*cmod a* > 1) *then cnj a else* 1) **by**(*auto simp*:*if-split*)
hence *lead-coeff-g*:*cmod* (*lead-coeff ?g*) = (∏ *a*←*?roots*. *max* 1 (*cmod a*)) **by**(*auto simp*:*max*)

have *norm2 f* = (∑ *a*←*coeffs ?f*. (*cmod a*)^2) **unfolding** *reconstruct-is-original-poly*..
also have ... = *cmod* (*lead-coeff f*)^2 ∗ (∑ *a*←*coeffs* (∏ *a*←*?roots*. [:− *a*, 1:]). (*cmod a*)^2)
unfolding *reconstruct-poly-def* **using** *sum-coeffs-smult-cmod*.
finally have *fg-norm*:*norm2 f* = *cmod* (*lead-coeff f*)^2 ∗ (∑ *a*←*coeffs ?g*. (*cmod a*)^2)
unfolding *Landau-lemma* **by** *auto*

have (*cmod* (*lead-coeff ?g*))^2 ≤ (∑ *a*←*coeffs ?g*. (*cmod a*)^2)
using *first-coeff-le-tail* **by** *blast*
from *ordered-comm-semiring-class.comm-mult-left-mono*[*OF this*]
have (*cmod* (*lead-coeff f*) ∗ *cmod* (*lead-coeff ?g*))^2 ≤ (∑ *a*←*coeffs f*. (*cmod a*)^2)
unfolding *fg-norm* **by** (*simp add*:*power-mult-distrib*)
hence *cmod* (*lead-coeff f*) ∗ (∏ *a*←*?roots*. *max* 1 (*cmod a*)) ≤ *sqrt* (*norm2 f*)
using *NthRoot.real-le-rsqrt lead-coeff-g* **by** *auto*
thus *mahler-measure-poly f* ≤ *sqrt* (*norm2 f*)
using *reconstruct-with-type-conversion*[*unfolded complex-roots-int-def*]
by (*simp add*: *mahler-measure-poly-via-monic mahler-measure-monic-def complex-roots-int-def*)
qed

lemma *prod-list-ge1*:
assumes *Ball* (*set x*) (λ (*a*::*real*). *a* ≥ 1)
shows *prod-list x* ≥ 1
using *assms* **proof**(*induct x*)
case (*Cons a as*)
have ∀ *a*∈*set as*. 1 ≤ *a* 1 ≤ *a* **using** *Cons*(2) **by** *auto*
thus *?case* **using** *Cons.hyps mult-mono'* **by** *fastforce*
qed *auto*

lemma *mahler-measure-monic-ge-1*: *mahler-measure-monic p* ≥ 1
unfolding *mahler-measure-monic-def* **by**(*rule prod-list-ge1*,*simp*)

lemma *mahler-measure-monic-ge-0*: *mahler-measure-monic p* ≥ 0
using *mahler-measure-monic-ge-1 le-numeral-extra*(1) *order-trans* **by** *blast*

lemma *mahler-measure-ge-0*: $0 \leq$ *mahler-measure h* **unfolding** *mahler-measure-def mahler-measure-poly-via-monic*
by (*simp add*: *mahler-measure-monic-ge-0*)

lemma *mahler-measure-constant*[*simp*]: *mahler-measure-poly* [:*c*:] = *cmod c*
proof −
have *main*: *complex-roots-complex* [:*c*:] = [] **unfolding** *complex-roots-complex-def*
by (*rule some-equality*, *auto*)
show *?thesis* **unfolding** *mahler-measure-poly-def main* **by** *auto*
qed

lemma *mahler-measure-factor*[*simplified*,*simp*]: *mahler-measure-poly* [:− *a*, *1*:] = *max 1* (*cmod a*)
proof −
have *main*: *complex-roots-complex* [:− *a*, *1*:] = [*a*] **unfolding** *complex-roots-complex-def*
proof (*rule some-equality*, *auto*, *goal-cases*)
case (*1 as*)
thus *?case* **by** (*cases as*, *auto*)
qed
show *?thesis* **unfolding** *mahler-measure-poly-def main* **by** *auto*
qed

lemma *mahler-measure-poly-explicit*: *mahler-measure-poly* (*smult c* (∏ *a*←*as*. [:− *a*, *1*:]))
= *cmod c* ∗ (∏ *a*←*as*. (*max 1* (*cmod a*)))
proof (*cases c* = *0*)
case *True*
thus *?thesis* **by** *auto*
next
case *False* **note** *c* = *this*
show *?thesis*
proof (*induct as*)
case (*Cons a as*)
have *mahler-measure-poly* (*smult c* (∏ *a*←*a* # *as*. [:− *a*, *1*:]))
= *mahler-measure-poly* (*smult c* (∏ *a*←*as*. [:− *a*, *1*:]) ∗ [: −*a*, *1* :])
by (*rule arg-cong*[*of - - mahler-measure-poly*], *unfold list.simps prod-list.Cons mult-smult-left*, *simp*)
also have . . . = *mahler-measure-poly* (*smult c* (∏ *a*←*as*. [:− *a*, *1*:])) ∗ *mahler-measure-poly* ([:− *a*, *1*:])
(**is** - = *?l* ∗ *?r*) **by** (*rule measure-eq-prod*)
also have *?l* = *cmod c* ∗ (∏ *a*←*as*. *max 1* (*cmod a*)) **unfolding** *Cons* **by** *simp*
also have *?r* = *max 1* (*cmod a*) **by** *simp*
finally show *?case* **by** *simp*
next
case *Nil*
show *?case* **by** *simp*
qed
qed

lemma *mahler-measure-poly-ge-1*:
assumes $h \neq 0$
shows $(1::real) \leq$ *mahler-measure h*
proof −
have *rc*: $|$*real-of-int i*$|$ = *of-int* $|i|$ **for** *i* **by** *simp*
from *assms* **have** *cmod* (*lead-coeff* (*map-poly complex-of-int h*)) > 0 **by** *simp*
hence *cmod* (*lead-coeff* (*map-poly complex-of-int h*)) ≥ 1
by(*cases lead-coeff h* = 0, *auto simp del*: *leading-coeff-0-iff*)
from *mult-mono*[*OF this mahler-measure-monic-ge-1 norm-ge-zero*]
show *?thesis* **unfolding** *mahler-measure-def mahler-measure-poly-via-monic*
by *auto*
qed

lemma *mahler-measure-dvd*: **assumes** $f \neq 0$ **and** *h dvd f*
shows *mahler-measure h* ≤ *mahler-measure f*
proof −
from *assms* **obtain** *g* **where** *f*: *f* = *g* ∗ *h* **unfolding** *dvd-def* **by** *auto*
from *f assms* **have** *g0*: *g* ≠ 0 **by** *auto*
hence *mg*: *mahler-measure g* ≥ 1 **by** (*rule mahler-measure-poly-ge-1*)
have *1* ∗ *mahler-measure h* ≤ *mahler-measure f*
unfolding *mahler-measure-def f measure-eq-prod*
of-int-poly-hom.hom-mult **unfolding** *mahler-measure-def*[*symmetric*]
by (*rule mult-right-mono*[*OF mg mahler-measure-ge-0*])
thus *?thesis* **by** *simp*
qed

definition *graeffe-poly* :: *'a* ⇒ *'a* :: *comm-ring-1 list* ⇒ *nat* ⇒ *'a poly* **where**
graeffe-poly c as m = *smult* (*c* ^ (2^*m*)) (∏ *a*←*as*. [:− (*a* ^ (2^*m*)), *1*:])

context
fixes *f* :: *complex poly* **and** *c as*
assumes *f*: *f* = *smult c* (∏ *a*←*as*. [:− *a*, *1*:])
begin
lemma *mahler-graeffe*: *mahler-measure-poly* (*graeffe-poly c as m*) = (*mahler-measure-poly f*)^(2^*m*)
proof −
have *graeffe*: *graeffe-poly c as m* = *smult* (*c* ^ 2 ^ *m*) (∏ *a*←(*map* (λ *a*. *a* ^ 2 ^ *m*) *as*). [:− *a*, *1*:])
unfolding *graeffe-poly-def*
by (*rule arg-cong*[*of - - smult* (*c* ^ 2 ^ *m*)], *induct as*, *auto*)
{
fix *n* :: *nat*
assume *n*: *n* > 0
have *id*: *max 1* (*cmod a* ^ *n*) = *max 1* (*cmod a*) ^ *n* **for** *a*
proof (*cases cmod a* ≤ *1*)
case *True*

```
    hence cmod a ^ n ≤ 1 by (simp add: power-le-one)
    with True show ?thesis by (simp add: max-def)
  qed (auto simp: max-def)
  have (∏ x←as. max 1 (cmod x ^ n)) = (∏ a←as. max 1 (cmod a)) ^ n
    by (induct as, auto simp: field-simps n id)
 }
 thus ?thesis unfolding f mahler-measure-poly-explicit graeffe
  by (auto simp: o-def field-simps norm-power)
qed
end

fun drop-half :: 'a list ⇒ 'a list where
  drop-half (x # y # ys) = x # drop-half ys
| drop-half xs = xs

fun alternate :: 'a list ⇒ 'a list × 'a list where
  alternate (x # y # ys) = (case alternate ys of (evn, od) ⇒ (x # evn, y # od))
| alternate xs = (xs,[])

definition poly-square-subst :: 'a :: comm-ring-1 poly ⇒ 'a poly where
  poly-square-subst f = poly-of-list (drop-half (coeffs f))

definition poly-even-odd :: 'a :: comm-ring-1 poly ⇒ 'a poly × 'a poly where
   poly-even-odd f = (case alternate (coeffs f) of (evn,od) ⇒ (poly-of-list evn,
poly-of-list od))


lemma poly-square-subst-coeff: coeff (poly-square-subst f) i = coeff f (2 * i)
proof −
  have id: coeff f (2 * i) = coeff (Poly (coeffs f)) (2 * i) by simp
  obtain xs where xs: coeffs f = xs by auto
  show ?thesis unfolding poly-square-subst-def poly-of-list-def coeff-Poly-eq id xs
  proof (induct xs arbitrary: i rule: drop-half.induct)
    case (1 x y ys i) thus ?case by (cases i, auto)
  next
    case (2-2 x i) thus ?case by (cases i, auto)
  qed auto
qed

lemma poly-even-odd-coeff: assumes poly-even-odd f = (ev,od)
  shows coeff ev i = coeff f (2 * i) coeff od i = coeff f (2 * i + 1)
proof −
  have id: ⋀ i. coeff f i = coeff (Poly (coeffs f)) i by simp
  obtain xs where xs: coeffs f = xs by auto
  from assms[unfolded poly-even-odd-def]
  have ev-od: ev = Poly (fst (alternate xs)) od = Poly (snd (alternate xs))
    by (auto simp: xs split: prod.splits)
  have coeff ev i = coeff f (2 * i) ∧ coeff od i = coeff f (2 * i + 1)
    unfolding poly-of-list-def coeff-Poly-eq id xs ev-od
```

```
  proof (induct xs arbitrary: i rule: alternate.induct)
    case (1 x y ys i) thus ?case by (cases alternate ys; cases i, auto)
  next
    case (2-2 x i) thus ?case by (cases i, auto)
  qed auto
  thus coeff ev i = coeff f (2 * i) coeff od i = coeff f (2 * i + 1) by auto
qed

lemma poly-square-subst: poly-square-subst (f ∘p (monom 1 2)) = f
  by (rule poly-eqI, unfold poly-square-subst-coeff, subst coeff-pcompose-x-pow-n,
auto)

lemma poly-even-odd: assumes poly-even-odd f = (g,h)
  shows f = g ∘p monom 1 2 + monom 1 1 * (h ∘p monom 1 2)
proof −
  note id = poly-even-odd-coeff[OF assms]
  show ?thesis
  proof (rule poly-eqI, unfold coeff-add coeff-monom-mult)
    fix n :: nat
    obtain m i where mi: m = n div 2 i = n mod 2 by auto
    have nmi: n = 2 * m + i i < 2 0 < (2 :: nat) 1 < (2 :: nat) unfolding mi
by auto
    have (2 :: nat) ≠ 0 by auto
    show coeff f n = coeff (g ∘p monom 1 2) n + (if 1 ≤ n then 1 * coeff (h ∘p
monom 1 2) (n − 1) else 0)
    proof (cases i = 1)
      case True
      hence id1: 2 * m + i − 1 = 2 * m + 0 by auto
       show ?thesis unfolding nmi id id1 coeff-pcompose-monom[OF nmi(2)] co-
eff-pcompose-monom[OF nmi(3)]
        unfolding True by auto
    next
      case False
      with nmi have i0: i = 0 by auto
      show ?thesis
      proof (cases m)
        case (Suc k)
        hence id1: 2 * m + i − 1 = 2 * k + 1 using i0 by auto
        show ?thesis unfolding nmi id coeff-pcompose-monom[OF nmi(2)]
          coeff-pcompose-monom[OF nmi(4)] id1 unfolding Suc i0 by auto
      next
        case 0
       show ?thesis unfolding nmi id coeff-pcompose-monom[OF nmi(2)] unfold-
ing i0 0 by auto
      qed
    qed
  qed
qed
```

context
fixes $f :: 'a :: idom\ poly$
begin

lemma *graeffe-0*: $f = smult\ c\ (\prod a{\leftarrow}as.\ [:-\ a,\ 1:]) \Longrightarrow graeffe\text{-}poly\ c\ as\ 0 = f$
unfolding *graeffe-poly-def* **by** *auto*

lemma *graeffe-recursion*: **assumes** *graeffe-poly c as m = f*
shows *graeffe-poly c as (Suc m) = smult ((−1) ^ (degree f)) (poly-square-subst (f * f ∘p [:0,−1:]))*
proof −
let *?g = graeffe-poly c as m*
have $f * f \circ_p [:0,-1:] = ?g * ?g \circ_p [:0,-1:]$ **unfolding** *assms* **by** *simp*
also have $?g \circ_p [:0,-1:] = smult\ ((-\ 1)\ \hat{}\ length\ as)\ (smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:]))$
unfolding *graeffe-poly-def*
proof (*induct as*)
case (*Cons a as*)
have *?case* = $((smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ ([:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:] \circ_p [:0,\ -\ 1:]) * (\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:]) \circ_p [:0,\ -\ 1:]) =$
$smult\ (-1 * (-\ 1)\ \hat{}\ length\ as)\ (smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ ([:\ a\ \hat{}\ 2\ \hat{}\ m,\ 1:] * (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:])))$
unfolding *list.simps prod-list.Cons pcompose-smult pcompose-mult* **by** *simp*
also have $smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ ([:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:] \circ_p [:0,\ -\ 1:] * (\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:]) \circ_p [:0,\ -\ 1:])$
$= smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ ((\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:]) \circ_p [:0,\ -\ 1:]) * [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:] \circ_p [:0,\ -\ 1:]$
unfolding *mult-smult-left* **by** *simp*
also have $smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ ((\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:]) \circ_p [:0,\ -\ 1:]) =$
$smult\ ((-\ 1)\ \hat{}\ length\ as)\ (smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:]))$
unfolding *pcompose-smult*[*symmetric*] *Cons* **..**
also have $[:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:] \circ_p [:0,\ -\ 1:] = smult\ (-1)\ [:\ a\hat{}2\hat{}m,\ 1:]$ **by** *simp*
finally have *id*: $?case = (smult\ ((-\ 1)\ \hat{}\ length\ as)\ (smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:])) * smult\ (-\ 1)\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:] =$
$smult\ (-\ 1 * (-\ 1)\ \hat{}\ length\ as)\ (smult\ (c\ \hat{}\ 2\ \hat{}\ m)\ ([:a\ \hat{}\ 2\ \hat{}\ m,\ 1:] * (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:])))$ **by** *simp*
obtain *c d* **where** *id′*: $(\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:]) = c\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:] = d$ **by** *auto*
show *?case* **unfolding** *id* **unfolding** *id′* **by** (*simp add*: *ac-simps*)
qed *simp*
finally have $f * f \circ_p [:0,\ -\ 1:] =$
$smult\ ((-\ 1)\ \hat{}\ length\ as * (c\ \hat{}\ 2\ \hat{}\ m * c\ \hat{}\ 2\ \hat{}\ m))\ ((\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:]) * (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:]))$
unfolding *graeffe-poly-def* **by** (*simp add*: *ac-simps*)
also have $c\ \hat{}\ 2\ \hat{}\ m * c\ \hat{}\ 2\ \hat{}\ m = c\ \hat{}\ 2\ \hat{}\ (Suc\ m)$ **by** (*simp add*: *semiring-normalization-rules*(*36*))
also have $(\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ m),\ 1:]) * (\prod a{\leftarrow}as.\ [:a\ \hat{}\ 2\ \hat{}\ m,\ 1:]) =$
$(\prod a{\leftarrow}as.\ [:-\ (a\ \hat{}\ 2\ \hat{}\ (Suc\ m)),\ 1:]) \circ_p monom\ 1\ 2$

```
  proof (induct as)
    case (Cons a as)
    have id: (monom 1 2 :: 'a poly) = [:0,0,1:]
      by (metis monom-altdef pCons-0-as-mult power2-eq-square smult-1-left)
    have (∏ a←a # as. [:− (a ^ 2 ^ m), 1:]) * (∏ a←a # as. [:a ^ 2 ^ m, 1:])
      = ([:− (a ^ 2 ^ m), 1:] * [: a ^ 2 ^ m, 1:]) * ((∏ a← as. [:− (a ^ 2 ^ m),
1:]) * (∏ a← as. [:a ^ 2 ^ m, 1:]))
        (is - = ?a * ?b)
      unfolding list.simps prod-list.Cons by (simp only: ac-simps)
    also have ?b = (∏ a←as. [:− (a ^ 2 ^ Suc m), 1:]) ∘p monom 1 2 unfolding
Cons by simp
      also have ?a = [: − (a ^ 2 ^ (Suc m)), 0 , 1:] by (simp add: semir-
ing-normalization-rules(36))
    also have ... = [: − (a ^ 2 ^ (Suc m)), 1:] ∘p monom 1 2 by (simp add: id)
    also have [: − (a ^ 2 ^ (Suc m)), 1:] ∘p monom 1 2 * (∏ a←as. [:− (a ^ 2 ^
Suc m), 1:]) ∘p monom 1 2 =
       (∏ a←a # as. [:− (a ^ 2 ^ Suc m), 1:]) ∘p monom 1 2 unfolding pcom-
pose-mult[symmetric] by simp
    finally show ?case .
  qed simp
  finally have f * f ∘p [:0, − 1:] = (smult ((− 1) ^ length as) (graeffe-poly c as
(Suc m)) ∘p monom 1 2)
    unfolding graeffe-poly-def pcompose-smult by simp
  from arg-cong[OF this, of λ f. smult ((− 1) ^ length as) (poly-square-subst f),
unfolded poly-square-subst]
  have graeffe-poly c as (Suc m) = smult ((− 1) ^ length as) (poly-square-subst (f
* f ∘p [:0, − 1:])) by simp
  also have ... = smult ((− 1) ^ degree f) (poly-square-subst (f * f ∘p [:0, − 1:]))

  proof (cases f = 0)
    case True
    thus ?thesis by (auto simp: poly-square-subst-def)
  next
    case False
    with assms have c0: c ≠ 0 unfolding graeffe-poly-def by auto
    from arg-cong[OF assms, of degree]
     have degree f = degree (smult (c ^ 2 ^ m) (∏ a←as. [:− (a ^ 2 ^ m), 1:]))
unfolding graeffe-poly-def by auto
      also have ... = degree (∏ a←as. [:− (a ^ 2 ^ m), 1:]) unfolding de-
gree-smult-eq using c0 by auto
    also have ... = length as unfolding degree-linear-factors by simp
    finally show ?thesis by simp
  qed
  finally show ?thesis .
qed
end

definition graeffe-one-step :: 'a ⇒ 'a :: idom poly ⇒ 'a poly where
  graeffe-one-step c f = smult c (poly-square-subst (f * f ∘p [:0,−1:]))
```

lemma *graeffe-one-step-code*[*code*]: **fixes** $c :: 'a :: idom$
shows *graeffe-one-step c f* = (*case poly-even-odd f of* (*g*,*h*)
⇒ *smult c* (*g* ∗ *g* − *monom 1 1* ∗ *h* ∗ *h*))
proof −
obtain *g h* **where** *eo*: *poly-even-odd f* = (*g*,*h*) **by** *force*
from *poly-even-odd*[*OF eo*] **have** *fgh*: *f* = *g* ∘$_p$ *monom 1 2* + *monom 1 1* ∗ *h* ∘$_p$ *monom 1 2* **by** *auto*
have *m2*: *monom* (*1* :: *'a*) *2* = [:*0*,*0*,*1*:] *monom* (*1* :: *'a*) *1* = [:*0*,*1*:]
unfolding *coeffs-eq-iff coeffs-monom*
by (*auto simp add*: *numeral-2-eq-2*)
show *?thesis* **unfolding** *eo split graeffe-one-step-def*
proof (*rule arg-cong*[*of - - smult c*])
let *?g* = *g* ∘$_p$ *monom 1 2*
let *?h* = *h* ∘$_p$ *monom 1 2*
let *?x* = *monom* (*1* :: *'a*) *1*
have *2*: *2* = *Suc* (*Suc 0*) **by** *simp*
have *f* ∗ *f* ∘$_p$ [:*0*, − *1*:] = (*g* ∘$_p$ *monom 1 2* + *monom 1 1* ∗ *h* ∘$_p$ *monom 1 2*) ∗
(*g* ∘$_p$ *monom 1 2* + *monom 1 1* ∗ *h* ∘$_p$ *monom 1 2*) ∘$_p$ [:*0*, − *1*:] **unfolding** *fgh* **by** *simp*
also have (*g* ∘$_p$ *monom 1 2* + *monom 1 1* ∗ *h* ∘$_p$ *monom 1 2*) ∘$_p$ [:*0*, − *1*:]
= *g* ∘$_p$ (*monom 1 2* ∘$_p$ [:*0*, − *1*:]) + *monom 1 1* ∘$_p$ [:*0*, − *1*:] ∗ *h* ∘$_p$ (*monom 1 2* ∘$_p$ [:*0*, − *1*:])
unfolding *pcompose-add pcompose-mult pcompose-assoc* **by** *simp*
also have *monom* (*1* :: *'a*) *2* ∘$_p$ [:*0*, − *1*:] = *monom 1 2* **unfolding** *m2* **by** *auto*
also have *?x* ∘$_p$ [:*0*, − *1*:] = [:*0*, −*1*:] **unfolding** *m2* **by** *auto*
also have [:*0*, − *1*:] ∗ *h* ∘$_p$ *monom 1 2* = (−*?x*) ∗ *?h* **unfolding** *m2* **by** *simp*
also have (*?g* + *?x* ∗ *?h*) ∗ (*?g* + (− *?x*) ∗ *?h*) = (*?g* ∗ *?g* − (*?x* ∗ *?x*) ∗ *?h* ∗ *?h*)
by (*auto simp*: *field-simps*)
also have *?x* ∗ *?x* = *?x* ∘$_p$ *monom 1 2* **unfolding** *mult-monom* **by** (*insert m2*, *simp add*: *2*)
also have (*?g* ∗ *?g* − ... ∗ *?h* ∗ *?h*) = (*g* ∗ *g* − *?x* ∗ *h* ∗ *h*) ∘$_p$ *monom 1 2*
unfolding *pcompose-diff pcompose-mult* **by** *auto*
finally have *poly-square-subst* (*f* ∗ *f* ∘$_p$ [:*0*, − *1*:])
= *poly-square-subst* ((*g* ∗ *g* − *?x* ∗ *h* ∗ *h*) ∘$_p$ *monom 1 2*) **by** *simp*
also have ... = *g* ∗ *g* − *?x* ∗ *h* ∗ *h* **unfolding** *poly-square-subst* **by** *simp*
finally show *poly-square-subst* (*f* ∗ *f* ∘$_p$ [:*0*, − *1*:]) = *g* ∗ *g* − *?x* ∗ *h* ∗ *h* **.**
qed
qed

fun *graeffe-poly-impl-main* :: *'a* ⇒ *'a* :: *idom poly* ⇒ *nat* ⇒ *'a poly* **where**
graeffe-poly-impl-main c f 0 = *f*
| *graeffe-poly-impl-main c f* (*Suc m*) = *graeffe-one-step c* (*graeffe-poly-impl-main c f m*)

lemma *graeffe-poly-impl-main*: **assumes** *f* = *smult c* (∏ *a*←*as*. [:− *a*, *1*:])

shows *graeffe-poly-impl-main* $((-1)\hat{}$*degree f*$)$ *f m* = *graeffe-poly c as m*
proof (*induct m*)
case *0*
show *?case* **using** *graeffe-0*[*OF assms*] **by** *simp*
next
case (*Suc m*)
have [*simp*]: *degree* (*graeffe-poly c as m*) = *degree f* **unfolding** *graeffe-poly-def degree-smult-eq assms*
degree-linear-factors **by** *auto*
from *arg-cong*[*OF Suc*, *of degree*]
show *?case* **unfolding** *graeffe-recursion*[*OF Suc*[*symmetric*]]
by (*simp add*: *graeffe-one-step-def*)
qed

definition *graeffe-poly-impl* :: *'a* :: *idom poly* ⇒ *nat* ⇒ *'a poly* **where**
graeffe-poly-impl f = *graeffe-poly-impl-main* $((-1)\hat{}$(*degree f*)) *f*

lemma *graeffe-poly-impl*: **assumes** *f* = *smult c* ($\prod$ *a*←*as*. [:− *a*, *1*:])
shows *graeffe-poly-impl f m* = *graeffe-poly c as m*
using *graeffe-poly-impl-main*[*OF assms*] **unfolding** *graeffe-poly-impl-def* **.**

lemma *drop-half-map*: *drop-half* (*map f xs*) = *map f* (*drop-half xs*)
by (*induct xs rule*: *drop-half.induct*, *auto*)

lemma (**in** *inj-comm-ring-hom*) *map-poly-poly-square-subst*:
map-poly hom (*poly-square-subst f*) = *poly-square-subst* (*map-poly hom f*)
unfolding *poly-square-subst-def coeffs-map-poly-hom drop-half-map poly-of-list-def*
by (*rule poly-eqI*, *auto simp*: *nth-default-map-eq*)

context *inj-idom-hom*
begin

lemma *graeffe-poly-impl-hom*:
map-poly hom (*graeffe-poly-impl f m*) = *graeffe-poly-impl* (*map-poly hom f*) *m*
proof −
interpret *mh*: *map-poly-inj-idom-hom*..
obtain *c* **where** *c*: $(((-\ 1)\ \hat{}$ *degree f*$)$ :: *'a*$)$ = *c* **by** *auto*
have *c'*: $(((-\ 1)\ \hat{}$ *degree f*$)$ :: *'b*$)$ = *hom c* **unfolding** *c*[*symmetric*] **by** (*simp add*:*hom-distribs*)
show *?thesis* **unfolding** *graeffe-poly-impl-def degree-map-poly-hom c c'*
apply (*induct m arbitrary*: *f*; *simp*)
by (*unfold graeffe-one-step-def hom-distribs map-poly-poly-square-subst map-poly-pcompose*,*simp*)
qed
end

lemma *graeffe-poly-impl-mahler*: *mahler-measure* (*graeffe-poly-impl f m*) = *mahler-measure* $f\ \hat{}\ 2\ \hat{}\ m$
proof −
let *?c* = *complex-of-int*

```
  let ?cc = map-poly ?c
  let ?f = ?cc f
  note eq = complex-roots(1)[of ?f]
  interpret inj-idom-hom complex-of-int by (standard, auto)
  show ?thesis
    unfolding mahler-measure-def mahler-graeffe[OF eq[symmetric], symmetric]
    graeffe-poly-impl[OF eq[symmetric], symmetric] by (simp add: of-int-hom.graeffe-poly-impl-hom)
qed

definition mahler-landau-graeffe-approximation :: nat ⇒ nat ⇒ int poly ⇒ int
where
  mahler-landau-graeffe-approximation kk dd f = (let
     no = sum-list (map (λ a. a * a) (coeffs f))
    in root-int-floor kk (dd * no))

lemma mahler-landau-graeffe-approximation-core:
  assumes g: g = graeffe-poly-impl f k
  shows mahler-measure f ≤ root (2 ^ Suc k) (real-of-int (∑ a←coeffs g. a * a))
proof −
  have mahler-measure f = root (2^k) (mahler-measure f ^ (2^k))
    by (simp add: real-root-power-cancel mahler-measure-ge-0)
  also have ... = root (2^k) (mahler-measure g)
    unfolding graeffe-poly-impl-mahler g by simp
  also have ... = root (2^k) (root 2 (((mahler-measure g)^2)))
    by (simp add: real-root-power-cancel mahler-measure-ge-0)
  also have ... = root (2^Suc k) (((mahler-measure g)^2))
    by (metis power-Suc2 real-root-mult-exp)
  also have ... ≤ root (2 ^ Suc k) (real-of-int (∑ a←coeffs g. a * a))
  proof (rule real-root-le-mono, force)
    have square-mono: 0 ≤ (x :: real) ⟹ x ≤ y ⟹ x * x ≤ y * y for x y
      by (simp add: mult-mono')
    obtain gs where gs: coeffs g = gs by auto
    have (mahler-measure g)^2 ≤ real-of-int |∑ a←coeffs g. a * a|
      using square-mono[OF mahler-measure-ge-0 Landau-inequality[of of-int-poly
g, folded mahler-measure-def]]
    by (auto simp: power2-eq-square coeffs-map-poly o-def of-int-hom.hom-sum-list)
    also have |∑ a←coeffs g. a * a| = (∑ a←coeffs g. a * a) unfolding gs
      by (induct gs, auto)
    finally show (mahler-measure g)^2 ≤ real-of-int (∑ a←coeffs g. a * a) .
  qed
  finally show mahler-measure f ≤ root (2 ^ Suc k) (real-of-int (∑ a←coeffs g. a
* a)) .
qed

lemma Landau-inequality-mahler-measure: mahler-measure f ≤ sqrt (real-of-int
(∑ a←coeffs f. a * a))
  by (rule order.trans[OF mahler-landau-graeffe-approximation-core[OF refl, of -
0]],
  auto simp: graeffe-poly-impl-def sqrt-def)
```

lemma *mahler-landau-graeffe-approximation*:
assumes *g*: *g* = *graeffe-poly-impl f k dd* = *d*^(2^(*Suc k*)) *kk* = 2^(*Suc k*)
shows ⌊*real d* ∗ *mahler-measure f*⌋ ≤ *mahler-landau-graeffe-approximation kk dd g*
proof −
have *id1*: *real-of-int* (*int* (*d* ^ 2 ^ *Suc k*)) = (*real d*) ^ 2 ^ *Suc k* **by** *simp*
have *id2*: *root* (2 ^ *Suc k*) (*real d* ^ 2 ^ *Suc k*) = *real d*
by (*simp add*: *real-root-power-cancel*)
show *?thesis* **unfolding** *mahler-landau-graeffe-approximation-def Let-def root-int-floor of-int-mult g*(2−3)
by (*rule floor-mono*, *unfold real-root-mult id1 id2*, *rule mult-left-mono*, *rule mahler-landau-graeffe-approximation-core*[*OF g*(1)], *auto*)
qed

context
fixes *bnd* :: *nat*
begin

function *mahler-approximation-main* :: *nat* ⇒ *int* ⇒ *int poly* ⇒ *int* ⇒ *nat* ⇒ *nat* ⇒ *int* **where**
mahler-approximation-main dd c g mm k kk = (*let mmm* = *mahler-landau-graeffe-approximation kk dd g*;
new-mm = (*if k* = 0 *then mmm else min mm mmm*)
in (*if k* ≥ *bnd then new-mm else*
— abort after *bnd* iterations of Graeffe transformation
mahler-approximation-main (*dd* ∗ *dd*) *c* (*graeffe-one-step c g*) *new-mm* (*Suc k*) (2 ∗ *kk*)))
by *pat-completeness auto*

termination by (*relation measure* (λ (*dd*,*c*,*f*,*mm*,*k*,*kk*). *Suc bnd* − *k*), *auto*)
declare *mahler-approximation-main.simps*[*simp del*]

lemma *mahler-approximation-main*: **assumes** *k* ≠ 0 ⟹ ⌊*real d* ∗ *mahler-measure f*⌋ ≤ *mm*
and *c* = (−1)^(*degree f*)
and *g* = *graeffe-poly-impl-main c f k dd* = *d*^(2^(*Suc k*)) *kk* = 2^(*Suc k*)
shows ⌊*real d* ∗ *mahler-measure f*⌋ ≤ *mahler-approximation-main dd c g mm k kk*
using *assms*
proof (*induct c g mm k kk rule*: *mahler-approximation-main.induct*)
case (*1 dd c g mm k kk*)
let *?df* = ⌊*real d* ∗ *mahler-measure f*⌋
note *dd* = *1*(*5*)
note *kk* = *1*(*6*)
note *g* = *1*(*4*)
note *c* = *1*(*3*)
note *mm* = *1*(*2*)
note *IH* = *1*(*1*)

note *mahl* = *mahler-approximation-main.simps*[*of dd c g mm k kk*]
define *mmm* **where** *mmm* = *mahler-landau-graeffe-approximation kk dd g*
define *new-mm* **where** *new-mm* = (*if k* = *0 then mmm else min mm mmm*)
let *?cond* = *bnd* ≤ *k*
have *id*: *mahler-approximation-main dd c g mm k kk* = (*if ?cond then new-mm*
else mahler-approximation-main (*dd* ∗ *dd*) *c* (*graeffe-one-step c g*) *new-mm*
(*Suc k*) (*2* ∗ *kk*))
unfolding *mahl mmm-def*[*symmetric*] *Let-def new-mm-def*[*symmetric*] **by** *simp*
have *gg*: *g* = (*graeffe-poly-impl f k*) **unfolding** *g graeffe-poly-impl-def c* **..**
from *mahler-landau-graeffe-approximation*[*OF gg dd kk*, *folded mmm-def*]
have *mmm*: *?df* ≤ *mmm* **.**
with *mm* **have** *new-mm*: *?df* ≤ *new-mm* **unfolding** *new-mm-def* **by** *auto*
show *?case*
proof (*cases ?cond*)
case *True*
show *?thesis* **unfolding** *id* **using** *True new-mm* **by** *auto*
next
case *False*
hence *id*: *mahler-approximation-main dd c g mm k kk* =
mahler-approximation-main (*dd* ∗ *dd*) *c* (*graeffe-one-step c g*) *new-mm* (*Suc*
k) (*2* ∗ *kk*)
unfolding *id* **by** *auto*
have *id′*: *graeffe-one-step c g* = *graeffe-poly-impl-main c f* (*Suc k*)
unfolding *g* **by** *simp*
have *dd* ∗ *dd* = *d* ^ *2* ^ *Suc* (*Suc k*) *2* ∗ *kk* = *2* ^ *Suc* (*Suc k*) **unfolding** *dd*
kk
semiring-normalization-rules(*26*) **by** *auto*
from *IH*[*OF mmm-def new-mm-def False new-mm c id′ this*]
show *?thesis* **unfolding** *id* **.**
qed
qed

definition *mahler-approximation* :: *nat* ⇒ *int poly* ⇒ *int* **where**
mahler-approximation d f = *mahler-approximation-main* (*d* ∗ *d*) ((−*1*)^(*degree*
f)) *f* (−*1*) *0 2*

lemma *mahler-approximation*: ⌊*real d* ∗ *mahler-measure f*⌋ ≤ *mahler-approximation*
d f
unfolding *mahler-approximation-def*
by (*rule mahler-approximation-main*, *auto simp*: *semiring-normalization-rules*(*29*))

end

end

10.5 The Mignotte Bound

theory *Factor-Bound*
imports

```
  Mahler-Measure
  Polynomial-Factorization.Gauss-Lemma
  Subresultants.Coeff-Int
begin

lemma binomial-mono-left: n ≤ N ⟹ n choose k ≤ N choose k
proof (induct n arbitrary: k N)
  case (0 k N)
  thus ?case by (cases k, auto)
next
  case (Suc n k N) note IH = this
  show ?case
  proof (cases k)
    case (Suc kk)
    from IH obtain NN where N: N = Suc NN and le: n ≤ NN by (cases N,
auto)
    show ?thesis unfolding N Suc using IH(1)[OF le]
      by (simp add: add-le-mono)
  qed auto
qed

definition choose-int where choose-int m n = (if n < 0 then 0 else m choose (nat
n))

lemma choose-int-suc[simp]:
  choose-int (Suc n) i = choose-int n (i−1) + choose-int n i
proof(cases nat i)
  case 0 thus ?thesis by (simp add:choose-int-def) next
  case (Suc v) hence nat (i − 1) = v i≠0 by simp-all
    thus ?thesis unfolding choose-int-def Suc by simp
qed

lemma sum-le-1-prod: assumes d: 1 ≤ d and c: 1 ≤ c
  shows c + d ≤ 1 + c ∗ (d :: real)
proof −
  from d c have (c − 1) ∗ (d − 1) ≥ 0 by auto
  thus ?thesis by (auto simp: field-simps)
qed

lemma mignotte-helper-coeff-int: cmod (coeff-int (∏ a←lst. [:− a, 1:]) i)
    ≤ choose-int (length lst − 1) i ∗ (∏ a←lst. (max 1 (cmod a)))
    + choose-int (length lst − 1) (i − 1)
proof(induct lst arbitrary:i)
  case Nil thus ?case by (auto simp:coeff-int-def choose-int-def)
  case (Cons v xs i)
  show ?case
  proof (cases xs = [])
    case True
    show ?thesis unfolding True
```

```
    by (cases nat i, cases nat (i − 1), auto simp: coeff-int-def choose-int-def)
  next
    case False
    hence id: length (v # xs) − 1 = Suc (length xs − 1) by auto
    have id': choose-int (length xs) i = choose-int (Suc (length xs − 1)) i for i
      using False by (cases xs, auto)
    let ?r = (∏ a←xs. [:− a, 1:])
    let ?mv = (∏ a←xs. (max 1 (cmod a)))
    let ?c1 = real (choose-int (length xs − 1) (i − 1 − 1))
    let ?c2 = real (choose-int (length (v # xs) − 1) i − choose-int (length xs −
1) i)
    let ?m xs n = choose-int (length xs − 1) n ∗ (∏ a←xs. (max 1 (cmod a)))
    have le1:1 ≤ max 1 (cmod v) by auto
    have le2:cmod v ≤ max 1 (cmod v) by auto
    have mv-ge-1:1 ≤ ?mv by (rule prod-list-ge1, auto)
    obtain a b c d where abcd :
      a = real (choose-int (length xs − 1) i)
      b = real (choose-int (length xs − 1) (i − 1))
      c = (∏ a←xs. max 1 (cmod a))
      d = cmod v by auto
    {
      have c1: c ≥ 1 unfolding abcd by (rule mv-ge-1)
      have b: b = 0 ∨ b ≥ 1 unfolding abcd by auto
      have a: a = 0 ∨ a ≥ 1 unfolding abcd by auto
      hence a0: a ≥ 0 by auto
      have acd: a ∗ (c ∗ d) ≤ a ∗ (c ∗ max 1 d) using a0 c1
        by (simp add: mult-left-mono)
      from b have b ∗ (c + d) ≤ b ∗ (1  + (c ∗ max 1 d))
      proof
        assume b ≥ 1
        hence ?thesis = (c + d ≤ 1 + c ∗ max 1 d) by simp
        also have ...
        proof (cases d ≥ 1)
          case False
          hence id: max 1 d = 1 by simp
          show ?thesis using False unfolding id by simp
        next
          case True
          hence id: max 1 d = d by simp
          show ?thesis using True c1 unfolding id by (rule sum-le-1-prod)
        qed
        finally show ?thesis .
      qed auto
      with acd have b ∗ c + (b ∗ d + a ∗ (c ∗ d)) ≤ b + (a ∗ (c ∗ max 1 d) + b
∗ (c ∗ max 1 d))
        by (auto simp: field-simps)
    } note abcd-main = this
    have cmod (coeff-int ([:− v, 1:] ∗ ?r) i) ≤ cmod (coeff-int ?r (i − 1)) + cmod
(coeff-int (smult v ?r) i)
```

```
    using norm-triangle-ineq4 by auto
  also have ... ≤ ?m xs (i − 1) + (choose-int (length xs − 1) (i − 1 − 1)) +
cmod (coeff-int (smult v ?r) i)
    using Cons[of i−1] by auto
  also have choose-int (length xs − 1) (i − 1) = choose-int (length (v # xs) −
1) i − choose-int (length xs − 1) i
    unfolding id choose-int-suc by auto
  also have ?c2 ∗ (∏ a←xs. max 1 (cmod a)) + ?c1 +
    cmod (coeff-int (smult v (∏ a←xs. [:− a, 1:])) i) ≤
    ?c2 ∗ (∏ a←xs. max 1 (cmod a)) + ?c1 + cmod v ∗ (
      real (choose-int (length xs − 1) i) ∗ (∏ a←xs. max 1 (cmod a)) +
      real (choose-int (length xs − 1) (i − 1)))
   using mult-mono'[OF order-refl Cons, of cmod v i, simplified] by (auto simp:
norm-mult)
   also have ... ≤ ?m (v # xs) i + (choose-int (length xs) (i − 1)) using
abcd-main[unfolded abcd]
    by (simp add: field-simps id')
  finally show ?thesis by simp
 qed
qed

lemma mignotte-helper-coeff-int': cmod (coeff-int (∏ a←lst. [:− a, 1:]) i)
  ≤ ((length lst − 1) choose i) ∗ (∏ a←lst. (max 1 (cmod a)))
  + min i 1 ∗ ((length lst − 1) choose (nat (i − 1)))
 by (rule order.trans[OF mignotte-helper-coeff-int], auto simp: choose-int-def min-def)

lemma mignotte-helper-coeff:
  cmod (coeff h i) ≤ (degree h − 1 choose i) ∗ mahler-measure-poly h
    + min i 1 ∗ (degree h − 1 choose (i − 1)) ∗ cmod (lead-coeff h)
proof −
 let ?r = complex-roots-complex h
 have cmod (coeff h i) = cmod (coeff (smult (lead-coeff h) (∏ a←?r. [:− a, 1:]))
i)
   unfolding complex-roots by auto
  also have ... = cmod (lead-coeff h) ∗ cmod (coeff (∏ a←?r. [:− a, 1:]) i)
by(simp add:norm-mult)
 also have ... ≤ cmod (lead-coeff h) ∗ ((degree h − 1 choose i) ∗ mahler-measure-monic
h +
   (min i 1 ∗ ((degree h − 1) choose nat (int i − 1))))
   unfolding mahler-measure-monic-def
   by (rule mult-left-mono, insert mignotte-helper-coeff-int'[of ?r i], auto)
  also have ... = (degree h − 1 choose i) ∗ mahler-measure-poly h + cmod
(lead-coeff h) ∗ (
   min i 1 ∗ ((degree h − 1) choose nat (int i − 1)))
   unfolding mahler-measure-poly-via-monic by (simp add: field-simps)
 also have nat (int i − 1) = i − 1 by (cases i, auto)
 finally show ?thesis by (simp add: ac-simps split: if-splits)
qed
```

lemma *mignotte-coeff-helper*:
abs (*coeff h i*) ≤
(*degree h − 1 choose i*) ∗ *mahler-measure h* +
(*min i 1* ∗ (*degree h − 1 choose* (*i − 1*)) ∗ *abs* (*lead-coeff h*))
unfolding *mahler-measure-def*
using *mignotte-helper-coeff*[*of of-int-poly h i*]
by *auto*

lemma *cmod-through-lead-coeff*[*simp*]:
cmod (*lead-coeff* (*of-int-poly h*)) = *abs* (*lead-coeff h*)
by *simp*

lemma *choose-approx*: *n* ≤ *N* ⟹ *n choose k* ≤ *N choose* (*N div 2*)
by (*rule order.trans*[*OF binomial-mono-left binomial-maximum*])

For Mignotte's factor bound, we currently do not support queries for individual coefficients, as we do not have a combined factor bound algorithm.

definition *mignotte-bound* :: *int poly* ⇒ *nat* ⇒ *int* **where**
mignotte-bound f d = (*let d′* = *d − 1*; *d2* = *d′ div 2*; *binom* = (*d′ choose d2*) *in*
(*mahler-approximation 2 binom f* + *binom* ∗ *abs* (*lead-coeff f*)))

lemma *mignotte-bound-main*:
assumes *f* ≠ *0 g dvd f degree g* ≤ *n*
shows |*coeff g k*| ≤ ⌊*real* (*n − 1 choose k*) ∗ *mahler-measure f*⌋ +
int (*min k 1* ∗ (*n − 1 choose* (*k − 1*))) ∗ |*lead-coeff f*|
proof−
let *?bnd* = *2*
let *?n* = (*n − 1*) *choose k*
let *?n′* = *min k 1* ∗ ((*n − 1*) *choose* (*k − 1*))
let *?approx* = *mahler-approximation ?bnd ?n f*
obtain *h* **where** *gh*:*g* ∗ *h* = *f* **using** *assms* **by** (*metis dvdE*)
have *nz*:*g*≠*0 h*≠*0* **using** *gh assms*(*1*) **by** *auto*
have *g1*:(*1*::*real*) ≤ *mahler-measure h* **using** *mahler-measure-poly-ge-1 gh assms*(*1*)
by *auto*
note *g0* = *mahler-measure-ge-0*
have *to-n*: (*degree g − 1 choose k*) ≤ *real ?n*
using *binomial-mono-left*[*of degree g − 1 n − 1 k*] *assms*(*3*) **by** *auto*
have *to-n′*: *min k 1* ∗ (*degree g − 1 choose* (*k − 1*)) ≤ *real ?n′*
using *binomial-mono-left*[*of degree g − 1 n − 1 k − 1*] *assms*(*3*)
by (*simp add*: *min-def*)
have |*coeff g k*| ≤ (*degree g − 1 choose k*) ∗ *mahler-measure g*
+ (*real* (*min k 1* ∗ (*degree g − 1 choose* (*k − 1*))) ∗ |*lead-coeff g*|)
using *mignotte-coeff-helper*[*of g k*] **by** *simp*
also have ... ≤ *?n* ∗ *mahler-measure f* + *real ?n′* ∗ |*lead-coeff f*|
proof (*rule add-mono*[*OF mult-mono*[*OF to-n*] *mult-mono*[*OF to-n′*]])
have *mahler-measure g* ≤ *mahler-measure g* ∗ *mahler-measure h* **using** *g1*
g0[*of g*]
using *mahler-measure-poly-ge-1 nz*(*1*) **by** *force*
thus *mahler-measure g* ≤ *mahler-measure f*

```
    using measure-eq-prod[of of-int-poly g of-int-poly h]
    unfolding mahler-measure-def gh[symmetric] by (auto simp: hom-distribs)
  have *: lead-coeff f = lead-coeff g * lead-coeff h
   unfolding arg-cong[OF gh, of lead-coeff, symmetric] by (rule lead-coeff-mult)
  have |lead-coeff h| ≠ 0 using nz(2) by auto
  hence lh: |lead-coeff h| ≥ 1 by linarith
 have |lead-coeff f| = |lead-coeff g| * |lead-coeff h| unfolding * by (rule abs-mult)
  also have ... ≥ |lead-coeff g| * 1
    by (rule mult-mono, insert lh, auto)
  finally have |lead-coeff g| ≤ |lead-coeff f| by simp
  thus real-of-int |lead-coeff g| ≤ real-of-int |lead-coeff f| by simp
qed (auto simp: g0)
finally have |coeff g k| ≤ ?n * mahler-measure f + real-of-int (?n' * |lead-coeff
f|) by simp
from floor-mono[OF this, folded floor-add-int]
have |coeff g k| ≤ floor (?n * mahler-measure f) + ?n' * |lead-coeff f| by linarith
thus ?thesis unfolding mignotte-bound-def Let-def using mahler-approximation[of
?n f ?bnd] by auto
qed

lemma Mignotte-bound:
  shows of-int |coeff g k| ≤ (degree g choose k) * mahler-measure g
proof (cases k ≤ degree g ∧ g ≠ 0)
  case False
  hence coeff g k = 0 using le-degree by (cases g = 0, auto)
  thus ?thesis using mahler-measure-ge-0[of g] by auto
next
  case kg: True
  hence g: g ≠ 0 g dvd g by auto
  from mignotte-bound-main[OF g le-refl, of k]
  have real-of-int |coeff g k|
    ≤ of-int ⌊real (degree g − 1 choose k) * mahler-measure g⌋ +
      of-int (int (min k 1 * (degree g − 1 choose (k − 1))) * |lead-coeff g|) by
linarith
  also have ... ≤ real (degree g − 1 choose k) * mahler-measure g
    + real (min k 1 * (degree g − 1 choose (k − 1))) * (of-int |lead-coeff g| * 1)
    by (rule add-mono, force, auto)
  also have ... ≤ real (degree g − 1 choose k) * mahler-measure g
    + real (min k 1 * (degree g − 1 choose (k − 1))) * mahler-measure g
    by (rule add-left-mono[OF mult-left-mono],
    unfold mahler-measure-def mahler-measure-poly-def,
    rule mult-mono, auto intro!: prod-list-ge1)
  also have ... =
    (real ((degree g − 1 choose k) + (min k 1 * (degree g − 1 choose (k − 1)))))
* mahler-measure g
    by (auto simp: field-simps)
  also have (degree g − 1 choose k) + (min k 1 * (degree g − 1 choose (k − 1)))
= degree g choose k
  proof (cases k = 0)
```

```
    case False
    then obtain kk where k: k = Suc kk by (cases k, auto)
    with kg obtain gg where g: degree g = Suc gg by (cases degree g, auto)
    show ?thesis unfolding k g by auto
  qed auto
  finally show ?thesis .
qed

lemma mignotte-bound:
  assumes f ≠ 0 g dvd f degree g ≤ n
  shows |coeff g k| ≤ mignotte-bound f n
proof −
  let ?bnd = 2
  let ?n = (n − 1) choose ((n − 1) div 2)
  have to-n: (n − 1 choose k) ≤ real ?n for k
    using choose-approx[OF le-refl] by auto
  from mignotte-bound-main[OF assms, of k]
  have |coeff g k| ≤
    ⌊real (n − 1 choose k) ∗ mahler-measure f⌋ +
    int (min k 1 ∗ (n − 1 choose (k − 1))) ∗ |lead-coeff f| .
  also have ... ≤ ⌊real (n − 1 choose k) ∗ mahler-measure f⌋ +
    int ((n − 1 choose (k − 1))) ∗ |lead-coeff f|
    by (rule add-left-mono[OF mult-right-mono], cases k, auto)
  also have ... ≤ mignotte-bound f n
    unfolding mignotte-bound-def Let-def
    by (rule add-mono[OF order.trans[OF floor-mono[OF mult-right-mono]
   mahler-approximation[of ?n f ?bnd]] mult-right-mono], insert to-n mahler-measure-ge-0,
auto)
  finally show ?thesis .
qed
```

As indicated before, at the moment the only available factor bound is Mignotte's one. As future work one might use a combined bound.

```
definition factor-bound :: int poly ⇒ nat ⇒ int where
  factor-bound = mignotte-bound

lemma factor-bound: assumes f ≠ 0 g dvd f degree g ≤ n
  shows |coeff g k| ≤ factor-bound f n
  unfolding factor-bound-def by (rule mignotte-bound[OF assms])
```

We further prove a result for factor bounds and scalar multiplication.

```
lemma factor-bound-ge-0: f ≠ 0 ⟹ factor-bound f n ≥ 0
  using factor-bound[of f 1 n 0] by auto

lemma factor-bound-smult: assumes f: f ≠ 0 and d: d ≠ 0
  and dvd: g dvd smult d f and deg: degree g ≤ n
  shows |coeff g k| ≤ |d| ∗ factor-bound f n
proof −
  let ?nf = primitive-part f let ?cf = content f
```

let *?ng* = *primitive-part g* **let** *?cg* = *content g*
from *content-dvd-contentI*[*OF dvd*] **have** *?cg dvd abs d* ∗ *?cf*
unfolding *content-smult-int* **.**
hence *dvd-c*: *?cg dvd d* ∗ *?cf* **using** *d*
by (*metis abs-content-int abs-mult dvd-abs-iff*)
from *primitive-part-dvd-primitive-partI*[*OF dvd*] **have** *?ng dvd smult* (*sgn d*) *?nf*
unfolding *primitive-part-smult-int* **.**
hence *dvd-n*: *?ng dvd ?nf* **using** *d*
by (*metis content-eq-zero-iff dvd dvd-smult-int f mult-eq-0-iff content-times-primitive-part smult-smult*)
define *gc* **where** *gc* = *gcd ?cf ?cg*
define *cg* **where** *cg* = *?cg div gc*
from *dvd d f* **have** *g*: *g* ≠ *0* **by** *auto*
from *f* **have** *cf*: *?cf* ≠ *0* **by** *auto*
from *g* **have** *cg*: *?cg* ≠ *0* **by** *auto*
hence *gc*: *gc* ≠ *0* **unfolding** *gc-def* **by** *auto*
have *cg-dvd*: *cg dvd ?cg* **unfolding** *cg-def gc-def* **using** *g* **by** (*simp add*: *div-dvd-iff-mult*)
have *cg-id*: *?cg* = *cg* ∗ *gc* **unfolding** *gc-def cg-def* **using** *g cf* **by** *simp*
from *dvd-smult-int*[*OF d dvd*] **have** *ngf*: *?ng dvd f* **.**
have *gcf*: |*gc*| *dvd content f* **unfolding** *gc-def* **by** *auto*
have *dvd-f*: *smult gc ?ng dvd f*
proof (*rule dvd-content-dvd*,
unfold content-smult-int content-primitive-part[*OF g*]
primitive-part-smult-int primitive-part-idemp)
show |*gc*| ∗ *1 dvd content f* **using** *gcf* **by** *auto*
show *smult* (*sgn gc*) (*primitive-part g*) *dvd primitive-part f*
using *dvd-n cf gc* **using** *zsgn-def* **by** *force*
qed
have *cg dvd d* **using** *dvd-c* **unfolding** *gc-def cg-def* **using** *cf cg d*
by (*simp add*: *div-dvd-iff-mult dvd-gcd-mult*)
then obtain *h* **where** *dcg*: *d* = *cg* ∗ *h* **unfolding** *dvd-def* **by** *auto*
with *d* **have** *h* ≠ *0* **by** *auto*
hence *h1*: |*h*| ≥ *1* **by** *simp*
have *degree* (*smult gc* (*primitive-part g*)) = *degree g*
using *gc* **by** *auto*
from *factor-bound*[*OF f dvd-f*, *unfolded this*, *OF deg*, *of k*, *unfolded coeff-smult*]
have *le*: |*gc* ∗ *coeff ?ng k*| ≤ *factor-bound f n* **.**
note *f0* = *factor-bound-ge-0*[*OF f*, *of n*]
from *mult-left-mono*[*OF le*, *of abs cg*]
have |*cg* ∗ *gc* ∗ *coeff ?ng k*| ≤ |*cg*| ∗ *factor-bound f n*
unfolding *abs-mult*[*symmetric*] **by** *simp*
also have *cg* ∗ *gc* ∗ *coeff ?ng k* = *coeff* (*smult ?cg ?ng*) *k* **unfolding** *cg-id* **by** *simp*
also have ... = *coeff g k* **unfolding** *content-times-primitive-part* **by** *simp*
finally have |*coeff g k*| ≤ *1* ∗ (|*cg*| ∗ *factor-bound f n*) **by** *simp*
also have ... ≤ |*h*| ∗ (|*cg*| ∗ *factor-bound f n*)
by (*rule mult-right-mono*[*OF h1*], *insert f0*, *auto*)
also have ... = (|*cg* ∗ *h*|) ∗ *factor-bound f n* **by** (*simp add*: *abs-mult*)
finally show *?thesis* **unfolding** *dcg* **.**

qed

end

10.6 Iteration of Subsets of Factors

theory *Sublist-Iteration*
imports
Polynomial-Factorization.Missing-Multiset
Polynomial-Factorization.Missing-List
HOL−Library.IArray
begin

Misc lemmas **lemma** *mem-snd-map*: $(\exists x.\ (x,\ y) \in S) \longleftrightarrow y \in snd\ `\ S$ **by** *force*

lemma *filter-upt*: **assumes** $l \leq m\ m < n$ **shows** $filter\ ((\leq)\ m)\ [l..<n] = [m..<n]$
proof(*insert assms*, *induct n*)
case *0* **then show** *?case* **by** *auto*
next
case (*Suc n*) **then show** *?case* **by** (*cases* $m = n$, *auto*)
qed

lemma *upt-append*: $i < j \Longrightarrow j < k \Longrightarrow [i..<j]@[j..<k] = [i..<k]$
proof(*induct k arbitrary*: *j*)
case *0* **then show** *?case* **by** *auto*
next
case (*Suc k*) **then show** *?case* **by** (*cases* $j = k$, *auto*)
qed

lemma *IArray-sub*[*simp*]: (!!) *as* = (!) (*IArray.list-of as*) **by** *auto*
declare *IArray.sub-def*[*simp del*]

Following lemmas in this section are for *subseqs*

lemma *subseqs-Cons*[*simp*]: *subseqs* (*x*#*xs*) = *map* (*Cons x*) (*subseqs xs*) @ *subseqs xs*
by (*simp add*: *Let-def*)

declare *subseqs.simps*(*2*) [*simp del*]

lemma *singleton-mem-set-subseqs* [*simp*]: $[x] \in set\ (subseqs\ xs) \longleftrightarrow x \in set\ xs$ **by** (*induct xs*, *auto*)

lemma *Cons-mem-set-subseqsD*: $y\#ys \in set\ (subseqs\ xs) \Longrightarrow y \in set\ xs$ **by** (*induct xs*, *auto*)

lemma *subseqs-subset*: $ys \in set\ (subseqs\ xs) \Longrightarrow set\ ys \subseteq set\ xs$
by (*metis Pow-iff image-eqI subseqs-powset*)

lemma *Cons-mem-set-subseqs-Cons*:
y#ys ∈ *set* (*subseqs* (*x#xs*)) ⟷ (*y* = *x* ∧ *ys* ∈ *set* (*subseqs xs*)) ∨ *y#ys* ∈ *set* (*subseqs xs*)
by *auto*

lemma *sorted-subseqs-sorted*:
sorted xs ⟹ *ys* ∈ *set* (*subseqs xs*) ⟹ *sorted ys*
proof(*induct xs arbitrary*: *ys*)
case *Nil* **thus** *?case* **by** *simp*
next
case *Cons* **thus** *?case* **using** *subseqs-subset* **by** *fastforce*
qed

lemma *subseqs-of-subseq*: *ys* ∈ *set* (*subseqs xs*) ⟹ *set* (*subseqs ys*) ⊆ *set* (*subseqs xs*)
proof(*induct xs arbitrary*: *ys*)
case *Nil* **then show** *?case* **by** *auto*
next
case *IHx*: (*Cons x xs*)
from *IHx.prems* **show** *?case*
proof(*induct ys*)
case *Nil* **then show** *?case* **by** *auto*
next
case *IHy*: (*Cons y ys*)
from *IHy.prems*[*unfolded subseqs-Cons*]
consider *y* = *x ys* ∈ *set* (*subseqs xs*) | *y* # *ys* ∈ *set* (*subseqs xs*) **by** *auto*
then show *?case*
proof(*cases*)
case *1* **with** *IHx.hyps* **show** *?thesis* **by** *auto*
next
case *2* **from** *IHx.hyps*[*OF this*] **show** *?thesis* **by** *auto*
qed
qed
qed

lemma *mem-set-subseqs-append*: *xs* ∈ *set* (*subseqs ys*) ⟹ *xs* ∈ *set* (*subseqs* (*zs* @ *ys*))
by (*induct zs*, *auto*)

lemma *Cons-mem-set-subseqs-append*:
x ∈ *set ys* ⟹ *xs* ∈ *set* (*subseqs zs*) ⟹ *x#xs* ∈ *set* (*subseqs* (*ys*@*zs*))
proof(*induct ys*)
case *Nil* **then show** *?case* **by** *auto*
next
case *IH*: (*Cons y ys*)
then consider *x* = *y* | *x* ∈ *set ys* **by** *auto*
then show *?case*
proof(*cases*)

case *1* **with** *IH* **show** *?thesis* **by** (*auto intro*: *mem-set-subseqs-append*)
next
case *2* **from** *IH.hyps*[*OF this IH.prems*(*2*)] **show** *?thesis* **by** *auto*
qed
qed

lemma *Cons-mem-set-subseqs-sorted*:
sorted xs $\Longrightarrow$ *y*#*ys* $\in$ *set* (*subseqs xs*) $\Longrightarrow$ *y*#*ys* $\in$ *set* (*subseqs* (*filter* (λ*x*. *y* $\leq$ *x*) *xs*))
by (*induct xs*) (*auto simp*: *Let-def*)

lemma *subseqs-map*[*simp*]: *subseqs* (*map f xs*) = *map* (*map f*) (*subseqs xs*) **by** (*induct xs*, *auto*)

lemma *subseqs-of-indices*: *map* (*map* (*nth xs*)) (*subseqs* [*0*..<*length xs*]) = *subseqs xs*
proof (*induct xs*)
case *Nil* **then show** *?case* **by** *auto*
next
case (*Cons x xs*)
from *this*[*symmetric*]
have *subseqs xs* = *map* (*map* ((!) (*x*#*xs*))) (*subseqs* [*Suc 0*..<*Suc* (*length xs*)])
by (*fold map-Suc-upt*, *simp*)
then show *?case* **by** (*unfold length-Cons upt-conv-Cons*[*OF zero-less-Suc*], *simp*)
qed

Specification **definition** *subseq-of-length n xs ys* $\equiv$ *ys* $\in$ *set* (*subseqs xs*) $\wedge$ *length ys* = *n*

lemma *subseq-of-lengthI*[*intro*]:
assumes *ys* $\in$ *set* (*subseqs xs*) *length ys* = *n*
shows *subseq-of-length n xs ys*
by (*insert assms*, *unfold subseq-of-length-def*, *auto*)

lemma *subseq-of-lengthD*[*dest*]:
assumes *subseq-of-length n xs ys*
shows *ys* $\in$ *set* (*subseqs xs*) *length ys* = *n*
by (*insert assms*, *unfold subseq-of-length-def*, *auto*)

lemma *subseq-of-length0*[*simp*]: *subseq-of-length 0 xs ys* $\longleftrightarrow$ *ys* = [] **by** *auto*

lemma *subseq-of-length-Nil*[*simp*]: *subseq-of-length n* [] *ys* $\longleftrightarrow$ *n* = *0* $\wedge$ *ys* = []
by (*auto simp*: *subseq-of-length-def*)

lemma *subseq-of-length-Suc-upt*:
subseq-of-length (*Suc n*) [*0*..<*m*] *xs* $\longleftrightarrow$
(*if n* = *0 then length xs* = *Suc 0* $\wedge$ *hd xs* < *m*
else hd xs < *hd* (*tl xs*) $\wedge$ *subseq-of-length n* [*0*..<*m*] (*tl xs*)) (**is** *?l* $\longleftrightarrow$ *?r*)
proof(*cases n*)

```
  case 0
  show ?thesis
  proof(intro iffI)
    assume l: ?l
    with 0 have 1: length xs = Suc 0 by auto
    then have xs: xs = [hd xs] by (metis length-0-conv length-Suc-conv list.sel(1))
    with l have [hd xs] ∈ set (subseqs [0..<m]) by auto
    with 1 show ?r by (unfold 0, auto)
  next
    assume ?r
    with 0 have 1: length xs = Suc 0 and 2: hd xs < m by auto
    then have xs: xs = [hd xs] by (metis length-0-conv length-Suc-conv list.sel(1))
    from 2 show ?l by (subst xs, auto simp: 0)
  qed
next
  case n: (Suc n')
  show ?thesis
  proof (intro iffI)
    assume ?l
    with n have 1: length xs = Suc (Suc n') and 2: xs ∈ set (subseqs [0..<m])
by auto
    from 1[unfolded length-Suc-conv]
    obtain x y ys where xs: xs = x#y#ys and n': length ys = n' by auto
    have sorted xs by(rule sorted-subseqs-sorted[OF - 2], auto)
    from this[unfolded xs] have x ≤ y by auto
    moreover
      from 2 have distinct xs by (rule subseqs-distinctD, auto)
      from this[unfolded xs] have x ≠ y by auto
    ultimately have x < y by auto
    moreover
        from 2 have y#ys ∈ set (subseqs [0..<m]) by (unfold xs, auto dest:
Cons-in-subseqsD)
      with n n' have subseq-of-length n [0..<m] (y#ys) by auto
    ultimately show ?r by (auto simp: xs)
  next
    assume r: ?r
    with n have len: length xs = Suc (Suc n')
     and ∗: hd xs < hd (tl xs) tl xs ∈ set (subseqs [0..<m]) by auto
    from len[unfolded length-Suc-conv] obtain x y ys
    where xs: xs = x#y#ys and n': length ys = n' by auto
    with ∗ have xy: x < y and yys: y#ys ∈ set (subseqs [0..<m]) by auto
    from Cons-mem-set-subseqs-sorted[OF - yys]
    have y#ys ∈ set (subseqs (filter ((≤) y) [0..<m])) by auto
    also from Cons-mem-set-subseqsD[OF yys] have ym: y < m by auto
      then have filter ((≤) y) [0..<m] = [y..<m] by (auto intro: filter-upt)
    finally have y#ys ∈ set (subseqs [y..<m]) by auto
    with xy have x#y#ys ∈ set (subseqs (x#[y..<m])) by auto
    also from xy have ... ⊆ set (subseqs ([0..<y] @ [y..<m]))
       by (intro subseqs-of-subseq Cons-mem-set-subseqs-append, auto intro: sub-
```

seqs-refl)
  also from *xy ym* **have** [*0..<y*] @ [*y..<m*] = [*0..<m*] **by** (*auto intro*: *upt-append*)
  finally have *xs* ∈ *set* (*subseqs* [*0..<m*]) **by** (*unfold xs*)
  with *len*[*folded n*] **show** *?l* **by** *auto*
 qed
qed

lemma *subseqs-of-length-of-indices*:
  { *ys. subseq-of-length n xs ys* } = { *map* (*nth xs*) *is* | *is. subseq-of-length n* [*0..<length xs*] *is* }
 by(*unfold subseq-of-length-def*, *subst subseqs-of-indices*[*symmetric*], *auto*)

lemma *subseqs-of-length-Suc-Cons*:
 { *ys. subseq-of-length* (*Suc n*) (*x*#*xs*) *ys* } =
  Cons x ‘ { *ys. subseq-of-length n xs ys* } ∪ { *ys. subseq-of-length* (*Suc n*) *xs ys* }
 by (*unfold subseq-of-length-def*, *auto*)

datatype (*'a*,*'b*,*'state*)*subseqs-impl* = *Sublists-Impl*
 (*create-subseqs*: *'b* ⇒ *'a list* ⇒ *nat* ⇒ (*'b* × *'a list*)*list* × *'state*)
 (*next-subseqs*: *'state* ⇒ (*'b* × *'a list*)*list* × *'state*)

locale *subseqs-impl* =
 fixes *f* :: *'a* ⇒ *'b* ⇒ *'b*
 and *sl-impl* :: (*'a*,*'b*,*'state*)*subseqs-impl*
begin

definition *S* :: *'b* ⇒ *'a list* ⇒ *nat* ⇒ (*'b* × *'a list*)*set* **where**
 S base elements n = { (*foldr f ys base*, *ys*) | *ys. subseq-of-length n elements ys* }

end

locale *correct-subseqs-impl* = *subseqs-impl f sl-impl*
 for *f* :: *'a* ⇒ *'b* ⇒ *'b*
 and *sl-impl* :: (*'a*,*'b*,*'state*)*subseqs-impl* +
 fixes *invariant* :: *'b* ⇒ *'a list* ⇒ *nat* ⇒ *'state* ⇒ *bool*
 assumes *create-subseqs*: *create-subseqs sl-impl base elements n* = (*out*, *state*) ⟹ *invariant base elements n state* ∧ *set out* = *S base elements n*
 and *next-subseqs*:
  invariant base elements n state ⟹
   next-subseqs sl-impl state = (*out*, *state'*) ⟹
   invariant base elements (*Suc n*) *state'* ∧ *set out* = *S base elements* (*Suc n*)

Basic Implementation **fun** *subseqs-i-n-main* :: (*'a* ⇒ *'b* ⇒ *'b*) ⇒ *'b* ⇒ *'a list* ⇒ *nat* ⇒ *nat* ⇒ (*'b* × *'a list*) *list* **where**
 subseqs-i-n-main f b xs i n = (*if i* = *0 then* [(*b*,[])] *else if i* = *n then* [(*foldr f xs b*, *xs*)]
  else case xs of
   (*y* # *ys*) ⇒ *map* (λ (*c*,*zs*) ⇒ (*c*,*y* # *zs*)) (*subseqs-i-n-main f* (*f y b*) *ys* (*i* −

```
1) (n − 1))
      @ subseqs-i-n-main f b ys i (n − 1))
declare subseqs-i-n-main.simps[simp del]

definition subseqs-length :: ('a ⇒ 'b ⇒ 'b) ⇒ 'b ⇒ nat ⇒ 'a list ⇒ ('b × 'a list)
list where
  subseqs-length f b i xs = (
    let n = length xs in if i > n then [] else subseqs-i-n-main f b xs i n)

lemma subseqs-length: assumes f-ac: ⋀ x y z. f x (f y z) = f y (f x z)
  shows set (subseqs-length f a n xs) =
  { (foldr f ys a, ys) | ys. ys ∈ set (subseqs xs) ∧ length ys = n}
proof −
  show ?thesis
  proof (cases length xs < n)
    case True
    thus ?thesis unfolding subseqs-length-def Let-def
      using length-subseqs[of xs] subseqs-length-simple-False by auto
  next
    case False
    hence id: (length xs < n) = False and n ≤ length xs by auto
    from this(2) show ?thesis unfolding subseqs-length-def Let-def id if-False
    proof (induct xs arbitrary: n a rule: length-induct[rule-format])
      case (1 xs n a)
      note n = 1(2)
      note IH = 1(1)
      note simp[simp] = subseqs-i-n-main.simps[of f - xs n]
      show ?case
      proof (cases n = 0)
        case True
        thus ?thesis unfolding simp by simp
      next
        case False note 0 = this
        show ?thesis
        proof (cases n = length xs)
          case True
          have ?thesis = ({(foldr f xs a, xs)} = (λ ys. (foldr f ys a, ys)) ‘ {ys. ys ∈
set (subseqs xs) ∧ length ys = length xs})
            unfolding simp using 0 True by auto
          from this[unfolded full-list-subseqs] show ?thesis by auto
        next
          case False
          with n have n: n < length xs by auto
          from 0 obtain m where m: n = Suc m by (cases n, auto)
          from n 0 obtain y ys where xs: xs = y # ys by (cases xs, auto)
          from n m xs have le: m ≤ length ys n ≤ length ys by auto
          from xs have lt: length ys < length xs by auto
          have sub: set (subseqs-i-n-main f a xs n (length xs)) =
            (λ(c, zs). (c, y # zs)) ‘ set (subseqs-i-n-main f (f y a) ys m (length ys))
```

```
∪
          set (subseqs-i-n-main f a ys n (length ys))
          unfolding simp using 0 False by (simp add: xs m)
        have fold: ⋀ ys. foldr f ys (f y a) = f y (foldr f ys a)
          by (induct-tac ys, auto simp: f-ac)
        show ?thesis unfolding sub IH[OF lt le(1)] IH[OF lt le(2)]
          unfolding m xs by (auto simp: Let-def fold)
      qed
    qed
  qed
 qed
qed

definition basic-subseqs-impl :: ('a ⇒ 'b ⇒ 'b) ⇒ ('a, 'b, 'b × 'a list × nat)subseqs-impl
where
 basic-subseqs-impl f = Sublists-Impl
   (λ a xs n. (subseqs-length f a n xs, (a,xs,n)))
   (λ (a,xs,n). (subseqs-length f a (Suc n) xs, (a,xs,Suc n)))

lemma basic-subseqs-impl: assumes f-ac: ⋀ x y z. f x (f y z) = f y (f x z)
 shows correct-subseqs-impl f (basic-subseqs-impl f)
   (λ a xs n triple. (a,xs,n) = triple)
 by (unfold-locales; unfold subseqs-impl.S-def basic-subseqs-impl-def subseq-of-length-def,
     insert subseqs-length[of f, OF f-ac], auto)

Improved Implementation  datatype ('a,'b,'state) subseqs-foldr-impl = Sub-
lists-Foldr-Impl
 (subseqs-foldr: 'b ⇒ 'a list ⇒ nat ⇒ 'b list × 'state)
 (next-subseqs-foldr: 'state ⇒ 'b list × 'state)

locale subseqs-foldr-impl =
 fixes f :: 'a ⇒ 'b ⇒ 'b
 and impl :: ('a,'b,'state) subseqs-foldr-impl
begin
definition S where S base elements n ≡ { foldr f ys base | ys. subseq-of-length n
elements ys }
end

locale correct-subseqs-foldr-impl = subseqs-foldr-impl f impl
 for f and impl :: ('a,'b,'state) subseqs-foldr-impl +
 fixes invariant :: 'b ⇒ 'a list ⇒ nat ⇒ 'state ⇒ bool
 assumes subseqs-foldr:
   subseqs-foldr impl base elements n = (out, state) ⟹
    invariant base elements n state ∧ set out = S base elements n
 and next-subseqs-foldr:
   next-subseqs-foldr impl state = (out, state') ⟹ invariant base elements n state
⟹
    invariant base elements (Suc n) state' ∧ set out = S base elements (Suc n)
```

locale *my-subseqs* =
  fixes $f :: 'a \Rightarrow 'b \Rightarrow 'b$
begin

context fixes *head* :: *'a* **and** *tail* :: *'a iarray*
begin

fun *next-subseqs1* **and** *next-subseqs2*
where *next-subseqs1 ret0 ret1 []* = (*ret0*, (*head*, *tail*, *ret1*))
  | *next-subseqs1 ret0 ret1* ((*i*,*v*)#*prevs*) = *next-subseqs2* (*f head v* # *ret0*) *ret1 prevs v* [*0*..<*i*]
  | *next-subseqs2 ret0 ret1 prevs v []* = *next-subseqs1 ret0 ret1 prevs*
  | *next-subseqs2 ret0 ret1 prevs v* (*j*#*js*) =
    (*let v'* = *f* (*tail* !! *j*) *v in next-subseqs2* (*v'* # *ret0*) ((*j*,*v'*) # *ret1*) *prevs v js*)

definition *next-subseqs2-set v js* ≡ { (*j*, *f* (*tail* !! *j*) *v*) | *j*. *j* ∈ *set js* }

definition *out-subseqs2-set v js* ≡ { *f* (*tail* !! *j*) *v* | *j*. *j* ∈ *set js* }

definition *next-subseqs1-set prevs* ≡ ⋃ { *next-subseqs2-set v* [*0*..<*i*] | *v i*. (*i*,*v*) ∈ *set prevs* }

definition *out-subseqs1-set prevs* ≡
  (*f head* ∘ *snd*) ‘ *set prevs* ∪ (⋃ { *out-subseqs2-set v* [*0*..<*i*] | *v i*. (*i*,*v*) ∈ *set prevs* })

fun *next-subseqs1-spec* **where**
  next-subseqs1-spec out nexts prevs (*out'*, (*head'*,*tail'*,*nexts'*)) ⟷
   set nexts' = *set nexts* ∪ *next-subseqs1-set prevs* ∧
   set out' = *set out* ∪ *out-subseqs1-set prevs*

fun *next-subseqs2-spec* **where**
  next-subseqs2-spec out nexts prevs v js (*out'*, (*head'*,*tail'*,*nexts'*)) ⟷
   set nexts' = *set nexts* ∪ *next-subseqs1-set prevs* ∪ *next-subseqs2-set v js* ∧
   set out' = *set out* ∪ *out-subseqs1-set prevs* ∪ *out-subseqs2-set v js*

lemma *next-subseqs2-Cons*:
  next-subseqs2-set v (*j*#*js*) = *insert* (*j*, *f* (*tail*!!*j*) *v*) (*next-subseqs2-set v js*)
  by (*auto simp*: *next-subseqs2-set-def*)

lemma *out-subseqs2-Cons*:
  out-subseqs2-set v (*j*#*js*) = *insert* (*f* (*tail*!!*j*) *v*) (*out-subseqs2-set v js*)
  by (*auto simp*: *out-subseqs2-set-def*)

lemma *next-subseqs1-set-as-next-subseqs2-set*:
  next-subseqs1-set ((*i*,*v*) # *prevs*) = *next-subseqs1-set prevs* ∪ *next-subseqs2-set v* [*0*..<*i*]
  by (*auto simp*: *next-subseqs1-set-def*)

```
lemma out-subseqs1-set-as-out-subseqs2-set:
  out-subseqs1-set ((i,v) # prevs) =
   { f head v } ∪ out-subseqs1-set prevs ∪ out-subseqs2-set v [0..<i]
  by (auto simp: out-subseqs1-set-def)

lemma next-subseqs1-spec:
   shows ⋀out nexts. next-subseqs1-spec out nexts prevs (next-subseqs1 out nexts
prevs)
    and ⋀out nexts. next-subseqs2-spec out nexts prevs v js (next-subseqs2 out nexts
prevs v js)
proof(induct rule: next-subseqs1-next-subseqs2.induct)
  case (1 ret0 ret1)
  then show ?case by (simp add: next-subseqs1-set-def out-subseqs1-set-def)
next
  case (2 ret0 ret1 i v prevs)
  show ?case
  proof(cases next-subseqs1 out nexts ((i, v) # prevs))
    case split: (fields out' head' tail' nexts')
   have next-subseqs2-spec (f head v # out) nexts prevs v [0..<i] (out', (head',tail',nexts'))
      by (fold split, unfold next-subseqs1.simps, rule 2)
    then show ?thesis
      apply (unfold next-subseqs2-spec.simps split)
    by (auto simp: next-subseqs1-set-as-next-subseqs2-set out-subseqs1-set-as-out-subseqs2-set)
  qed
next
  case (3 ret0 ret1 prevs v)
  show ?case
  proof (cases next-subseqs1 out nexts prevs)
    case split: (fields out' head' tail' nexts')
     from 3[of out nexts] show ?thesis by(simp add: split next-subseqs2-set-def
out-subseqs2-set-def)
  qed
next
  case (4 ret0 ret1 prevs v j js)
  define tj where tj = tail !! j
  define nexts'' where nexts'' = (j, f tj v) # nexts
  define out'' where out'' = (f tj v) # out
  let ?n = next-subseqs2 out'' nexts'' prevs v js
  show ?case
  proof (cases ?n)
    case split: (fields out' head' tail' nexts')
    show ?thesis
      apply (unfold next-subseqs2.simps Let-def)
      apply (fold tj-def)
      apply (fold out''-def nexts''-def)
    apply (unfold split next-subseqs2-spec.simps next-subseqs2-Cons out-subseqs2-Cons)
      using 4[OF refl, of out'' nexts'', unfolded split]
      apply (auto simp: tj-def nexts''-def out''-def)
      done
```

```
  qed
qed

end

fun next-subseqs where next-subseqs (head,tail,prevs) = next-subseqs1 head tail []
[] prevs

fun create-subseqs
where create-subseqs base elements 0 = (
      if elements = [] then ([base],(undefined, IArray [], []))
      else let head = hd elements; tail = IArray (tl elements) in
        ([base], (head, tail, [(IArray.length tail, base)])))
  |   create-subseqs base elements (Suc n) =
      next-subseqs (snd (create-subseqs base elements n))

definition impl where impl = Sublists-Foldr-Impl create-subseqs next-subseqs

sublocale subseqs-foldr-impl f impl .

definition set-prevs where set-prevs base tail n ≡
  { (i, foldr f (map ((!) tail) is) base) | i is.
   subseq-of-length n [0..<length tail] is ∧ i = (if n = 0 then length tail else hd is)
}

lemma snd-set-prevs:
  snd ‘ (set-prevs base tail n) = (λas. foldr f as base) ‘ { as. subseq-of-length n tail
as }
  by (subst subseqs-of-length-of-indices, auto simp: set-prevs-def image-Collect)


fun invariant where invariant base elements n (head,tail,prevs) =
  (if elements = [] then prevs = []
   else head = hd elements ∧ tail = IArray (tl elements) ∧ set prevs = set-prevs
base (tl elements) n)


lemma next-subseq-preserve:
  assumes next-subseqs (head,tail,prevs) = (out, (head',tail',prevs'))
  shows head' = head tail' = tail
proof−
  define P :: 'b list × - × - × (nat × 'b) list ⇒ bool
  where P ≡ λ (out, (head',tail',prevs')). head' = head ∧ tail' = tail
  { fix ret0 ret1 v js
    have ∗: P (next-subseqs1 head tail ret0 ret1 prevs)
     and  P (next-subseqs2 head tail ret0 ret1 prevs v js)
   by(induct rule: next-subseqs1-next-subseqs2.induct, simp add: P-def, auto simp:
Let-def)
  }
```

from *this*(*1*)[*unfolded P-def*, *of* [] [], *folded next-subseqs.simps*] *assms*
show *head′* = *head tail′* = *tail* **by** *auto*
qed

lemma *next-subseqs-spec*:
assumes *nxt*: *next-subseqs* (*head*,*tail*,*prevs*) = (*out*, (*head′*,*tail′*,*prevs′*))
shows *set prevs′* = { (*j*, *f* (*tail* !! *j*) *v*) | *v i j*. (*i*,*v*) ∈ *set prevs* ∧ *j* < *i* } (**is** *?g1*)
and *set out* = (*f head* ∘ *snd*) ‘ *set prevs* ∪ *snd* ‘ *set prevs′* (**is** *?g2*)
proof−
note *next-subseqs1-spec*(*1*)[*of head tail Nil Nil prevs*]
note *this*[*unfolded nxt*[*simplified*]]
note *this*[*unfolded next-subseqs1-spec.simps*]
note *this*[*unfolded next-subseqs1-set-def out-subseqs1-set-def*]
note ∗ = *this*[*unfolded next-subseqs2-set-def out-subseqs2-set-def*]
then show *g1*: *?g1* **by** *auto*
also have *snd* ‘ ... = (⋃ {{(*f* (*tail* !! *j*) *v*) | *j*. *j* < *i*} | *v i*. (*i*, *v*) ∈ *set prevs*})
by (*unfold image-Collect*, *auto*)
finally have ∗∗: *snd* ‘ *set prevs′* = ...
with *conjunct2*[*OF* ∗] **show** *?g2* **by** *simp*
qed

lemma *next-subseq-prevs*:
assumes *nxt*: *next-subseqs* (*head*,*tail*,*prevs*) = (*out*, (*head′*,*tail′*,*prevs′*))
and *inv-prevs*: *set prevs* = *set-prevs base* (*IArray.list-of tail*) *n*
shows *set prevs′* = *set-prevs base* (*IArray.list-of tail*) (*Suc n*) (**is** *?l* = *?r*)
proof(*intro equalityI subsetI*)
fix *t*
assume *r*: *t* ∈ *?r*
from *this*[*unfolded set-prevs-def*] **obtain** *iis*
where *t*: *t* = (*hd iis*, *foldr f* (*map* ((!!) *tail*) *iis*) *base*)
and *sl*: *subseq-of-length* (*Suc n*) [*0*..<*IArray.length tail*] *iis* **by** *auto*
from *sl* **have** *length iis* > *0* **by** *auto*
then obtain *i is* **where** *iis*: *iis* = *i*#*is* **by** (*meson list.set-cases nth-mem*)
define *v* **where** *v* = *foldr f* (*map* ((!!) *tail*) *is*) *base*
note *sl*[*unfolded subseq-of-length-Suc-upt*]
note *nxt* = *next-subseqs-spec*[*OF nxt*]
show *t* ∈ *?l*
proof(*cases n* = *0*)
case *True*
from *sl*[*unfolded subseq-of-length-Suc-upt*] *t*
show *?thesis* **by** (*unfold nxt*[*unfolded inv-prevs*] *True set-prevs-def length-Suc-conv*, *auto*)
next
case [*simp*]: *False*
from *sl*[*unfolded subseq-of-length-Suc-upt iis*,*simplified*]
have *i*: *i* < *hd is* **and** *is*: *subseq-of-length n* [*0*..<*IArray.length tail*] *is* **by** *auto*
then have ∗: (*hd is*, *v*) ∈ *set-prevs base* (*IArray.list-of tail*) *n*
by (*unfold set-prevs-def*, *auto intro*!: *exI*[*of - is*] *simp*: *v-def*)
with *i* **have** (*i*, *f* (*tail* !! *i*) *v*) ∈ {(*j*, *f* (*tail* !! *j*) *v*) | *j*. *j* < *hd is*} **by** *auto*

```
    with t[unfolded iis] have t ∈ ... by (auto simp: v-def)
    with ∗ show ?thesis by (unfold nxt[unfolded inv-prevs], auto)
  qed
next
  fix t
  assume l: t ∈ ?l
  from l[unfolded next-subseqs-spec(1)[OF nxt]]
  obtain j v i
  where t: t = (j, f (tail!!j) v)
    and j: j < i
    and iv: (i,v) ∈ set prevs by auto
  from iv[unfolded inv-prevs set-prevs-def, simplified]
  obtain is
  where v: v = foldr f (map ((!!) tail) is) base
    and is: subseq-of-length n [0..<IArray.length tail] is
    and i: if n = 0 then i = IArray.length tail else i = hd is by auto
  from is j i have jis: subseq-of-length (Suc n) [0..<IArray.length tail] (j#is)
    by (unfold subseq-of-length-Suc-upt, auto)
  then show t ∈ ?r by (auto intro!: exI[of - j#is] simp: set-prevs-def t v)
qed

lemma invariant-next-subseqs:
  assumes inv: invariant base elements n state
      and nxt: next-subseqs state = (out, state')
  shows invariant base elements (Suc n) state'
proof(cases elements = [])
  case True with inv nxt show ?thesis by(cases state, auto)
next
  case False with inv nxt show ?thesis
  proof (cases state)
    case state: (fields head tail prevs)
    note inv = inv[unfolded state]
    show ?thesis
    proof (cases state')
      case state': (fields head' tail' prevs')
      note nxt = nxt[unfolded state state']
      note [simp] = next-subseq-preserve[OF nxt]
      from False inv
      have set prevs = set-prevs base (IArray.list-of tail) n by auto
      from False next-subseq-prevs[OF nxt this] inv
      show ?thesis by(auto simp: state')
    qed
  qed
qed

lemma out-next-subseqs:
  assumes inv: invariant base elements n state
      and nxt: next-subseqs state = (out, state')
  shows set out = S base elements (Suc n)
```

```
proof (cases state)
  case state: (fields head tail prevs)
  show ?thesis
  proof(cases elements = [])
    case True
    with inv nxt show ?thesis by (auto simp: state S-def)
  next
    case elements: False
    show ?thesis
    proof(cases state')
      case state': (fields head' tail' prevs')
      from elements inv[unfolded state,simplified]
      have head = hd elements
       and tail = IArray (tl elements)
       and prevs: set prevs = set-prevs base (tl elements) n by auto
       with elements have elements2: elements = head # IArray.list-of tail  by
auto
      let ?f = λas. (foldr f as base)
      have set out = ?f ‘ {ys. subseq-of-length (Suc n) elements ys}
      proof−
          from invariant-next-subseqs[OF inv nxt, unfolded state' invariant.simps
if-not-P[OF elements]]
        have tail': tail' = IArray (tl elements)
         and prevs': set prevs' = set-prevs base (tl elements) (Suc n) by auto
        note next-subseqs-spec(2)[OF nxt[unfolded state state'], unfolded this]
        note this[folded image-comp, unfolded snd-set-prevs]
        also note prevs
        also note snd-set-prevs
        also have f head ‘ ?f ‘ { as. subseq-of-length n (tl elements) as } =
          ?f ‘ Cons head ‘ { as. subseq-of-length n (tl elements) as } by (auto simp:
image-def)
        also note image-Un[symmetric]
        also have
          ((#) head ‘ {as. subseq-of-length n (tl elements) as} ∪
           {as. subseq-of-length (Suc n) (tl elements) as}) =
           {as. subseq-of-length (Suc n) elements as}
        by (unfold subseqs-of-length-Suc-Cons elements2, auto)
        finally show ?thesis.
      qed
      then show ?thesis by (auto simp: S-def)
    qed
  qed
qed

lemma create-subseqs:
  create-subseqs base elements n = (out, state) ⟹
   invariant base elements n state ∧ set out = S base elements n
proof(induct n arbitrary: out state)
 case 0 then show ?case by (cases elements, cases state, auto simp: S-def Let-def
```

set-prevs-def)
next
case (*Suc n*) **show** *?case*
proof (*cases create-subseqs base elements n*)
case *1*: (*fields out'' head tail prevs*)
show *?thesis*
proof (*cases next-subseqs* (*head*, *tail*, *prevs*))
case (*fields out' head' tail' prevs'*)
note *2* = *this*[*unfolded next-subseq-preserve*[*OF this*]]
from *Suc*(*2*)[*unfolded create-subseqs.simps 1 snd-conv 2*]
have *3*: *out'* = *out state* = (*head*,*tail*,*prevs'*) **by** *auto*
from *Suc*(*1*)[*OF 1*]
have *inv*: *invariant base elements n* (*head*, *tail*, *prevs*) **by** *auto*
from *out-next-subseqs*[*OF inv 2*] *invariant-next-subseqs*[*OF inv 2*]
show *?thesis* **by** (*auto simp*: *3*)
qed
qed
qed

sublocale *correct-subseqs-foldr-impl f impl invariant*
by (*unfold-locales*; *auto simp*: *impl-def invariant-next-subseqs out-next-subseqs create-subseqs*)

lemma *impl-correct*: *correct-subseqs-foldr-impl f impl invariant* **..**
end

lemmas [*code*] =
my-subseqs.next-subseqs.simps
my-subseqs.next-subseqs1.simps
my-subseqs.next-subseqs2.simps
my-subseqs.create-subseqs.simps
my-subseqs.impl-def

end

10.7 Reconstruction of Integer Factorization

We implemented Zassenhaus reconstruction-algorithm, i.e., given a factorization of f mod p^n, the aim is to reconstruct a factorization of f over the integers.

theory *Reconstruction*
imports
Berlekamp-Hensel
Polynomial-Factorization.Gauss-Lemma
Polynomial-Factorization.Dvd-Int-Poly
Polynomial-Factorization.Gcd-Rat-Poly
Degree-Bound
Factor-Bound
Sublist-Iteration

Poly-Mod
begin

hide-const *coeff monom*

Misc lemmas **lemma** *foldr-of-Cons*[*simp*]: *foldr Cons xs ys* = *xs* @ *ys* **by** (*induct xs*, *auto*)

lemma *foldr-map-prod*[*simp*]:
foldr (λ*x*. *map-prod* (*f x*) (*g x*)) *xs base* = (*foldr f xs* (*fst base*), *foldr g xs* (*snd base*))
by (*induct xs*, *auto*)

The main part **context** *poly-mod*
begin

definition *inv-Mp* :: *int poly* ⇒ *int poly* **where**
inv-Mp = *map-poly inv-M*

definition *mul-const* :: *int poly* ⇒ *int* ⇒ *int* **where**
mul-const p c = (*coeff p 0* ∗ *c*) *mod m*

fun *prod-list-m* :: *int poly list* ⇒ *int poly* **where**
prod-list-m (*f* # *fs*) = *Mp* (*f* ∗ *prod-list-m fs*)
| *prod-list-m* [] = *1*

context
fixes *sl-impl* :: (*int poly*, *int* × *int poly list*, *'state*)*subseqs-foldr-impl*
and *m2* :: *int*
begin
definition *inv-M2* :: *int* ⇒ *int* **where**
inv-M2 = (λ *x*. *if x* ≤ *m2 then x else x* − *m*)

definition *inv-Mp2* :: *int poly* ⇒ *int poly* **where**
inv-Mp2 = *map-poly inv-M2*

partial-function (*tailrec*) *reconstruction* :: *'state* ⇒ *int poly* ⇒ *int poly*
⇒ *int* ⇒ *nat* ⇒ *nat* ⇒ *int poly list* ⇒ *int poly list*
⇒ (*int* × (*int poly list*)) *list* ⇒ *int poly list* **where**
reconstruction state u luu lu d r vs res cands = (*case cands of Nil*
⇒ *let d'* = *Suc d*
in if d' + *d'* > *r then* (*u* # *res*) *else*
(*case next-subseqs-foldr sl-impl state of* (*cands*,*state'*) ⇒
reconstruction state' u luu lu d' r vs res cands)
| (*lv'*,*ws*) # *cands'* ⇒ *let*
lv = *inv-M2 lv'* — *lv* is last coefficient of *vb* below
in if lv dvd coeff luu 0 then let
vb = *inv-Mp2* (*Mp* (*smult lu* (*prod-list-m ws*)))
in if vb dvd luu then

```
      let pp-vb = primitive-part vb;
          u' = u div pp-vb;
          r' = r − length ws;
          res' = pp-vb # res
        in if d + d > r'
          then u' # res'
          else let
              lu' = lead-coeff u';
              vs' = fold remove1 ws vs;
              (cands'', state') = subseqs-foldr sl-impl (lu',[]) vs' d
            in reconstruction state' u' (smult lu' u') lu' d r' vs' res' cands''
      else reconstruction state u luu lu d r vs res cands'
      else reconstruction state u luu lu d r vs res cands')
  end
end


declare poly-mod.reconstruction.simps[code]
declare poly-mod.prod-list-m.simps[code]
declare poly-mod.mul-const-def[code]
declare poly-mod.inv-M2-def[code]
declare poly-mod.inv-Mp2-def[code-unfold]
declare poly-mod.inv-Mp-def[code-unfold]

definition zassenhaus-reconstruction-generic ::
  (int poly, int × int poly list, 'state) subseqs-foldr-impl
  ⇒ int poly list ⇒ int ⇒ nat ⇒ int poly ⇒ int poly list where
  zassenhaus-reconstruction-generic sl-impl vs p n f = (let
     lf = lead-coeff f;
     pn = p^n;
     (-, state) = subseqs-foldr sl-impl (lf,[]) vs 0
   in
     poly-mod.reconstruction pn sl-impl (pn div 2) state f (smult lf f) lf 0 (length
vs) vs [] [])

lemma coeff-mult-0: coeff (f * g) 0 = coeff f 0 * coeff g 0
  by (metis poly-0-coeff-0 poly-mult)

lemma lead-coeff-factor: assumes u: u = v * (w :: 'a ::idom poly)
  shows smult (lead-coeff u) u = (smult (lead-coeff w) v) * (smult (lead-coeff v)
w)
  lead-coeff (smult (lead-coeff w) v) = lead-coeff u lead-coeff (smult (lead-coeff v)
w) = lead-coeff u
  unfolding u by (auto simp: lead-coeff-mult lead-coeff-smult)

lemma not-irreducible_d-lead-coeff-factors: assumes ¬ irreducible_d (u :: 'a :: idom
poly) degree u ≠ 0
  shows ∃ f g. smult (lead-coeff u) u = f * g ∧ lead-coeff f = lead-coeff u ∧
lead-coeff g = lead-coeff u
```

```
  ∧ degree f < degree u ∧ degree g < degree u
proof −
  from assms[unfolded irreducible_d-def, simplified]
  obtain v w where deg: degree v < degree u degree w < degree u and u: u = v
* w by auto
  define f where f = smult (lead-coeff w) v
  define g where g = smult (lead-coeff v) w
  note lf = lead-coeff-factor[OF u, folded f-def g-def]
  show ?thesis
  proof (intro exI conjI, (rule lf)+)
    show degree f < degree u degree g < degree u unfolding f-def g-def using deg
u by auto
  qed
qed

lemma mset-subseqs-size: mset ‘ {ys. ys ∈ set (subseqs xs) ∧ length ys = n} =
  {ws. ws ⊆# mset xs ∧ size ws = n}
proof (induct xs arbitrary: n)
  case (Cons x xs n)
  show ?case (is ?l = ?r)
  proof (cases n)
    case 0
    thus ?thesis by (auto simp: Let-def)
  next
    case (Suc m)
    have ?r = {ws. ws ⊆# mset (x # xs)} ∩ {ps. size ps = n} by auto
    also have {ws. ws ⊆# mset (x # xs)} = {ps. ps ⊆# mset xs} ∪ ((λ ps. ps +
{#x#}) ‘ {ps. ps ⊆# mset xs})
      by (simp add: multiset-subset-insert)
    also have ... ∩ {ps. size ps = n} = {ps. ps ⊆# mset xs ∧ size ps = n}
     ∪ ((λ ps. ps + {#x#}) ‘ {ps. ps ⊆# mset xs ∧ size ps = m}) unfolding Suc
by auto
    finally have id: ?r =
      {ps. ps ⊆# mset xs ∧ size ps = n} ∪ (λps. ps + {#x#}) ‘ {ps. ps ⊆# mset
xs ∧ size ps = m} .
    have ?l = mset ‘ {ys ∈ set (subseqs xs). length ys = Suc m}
     ∪ mset ‘ {ys ∈ (#) x ‘ set (subseqs xs). length ys = Suc m}
      unfolding Suc by (auto simp: Let-def)
    also have mset ‘ {ys ∈ (#) x ‘ set (subseqs xs). length ys = Suc m}
     = (λps. ps + {#x#}) ‘ mset ‘ {ys ∈ set (subseqs xs). length ys = m} by force
    finally have id': ?l = mset ‘ {ys ∈ set (subseqs xs). length ys = Suc m} ∪
      (λps. ps + {#x#}) ‘ mset ‘ {ys ∈ set (subseqs xs). length ys = m} .
    show ?thesis unfolding id id' Cons[symmetric] unfolding Suc by simp
  qed
qed auto

context poly-mod-2
begin
lemma prod-list-m[simp]: prod-list-m fs = Mp (prod-list fs)
```

by (*induct fs*, *auto*)

lemma *inv-Mp-coeff*: *coeff* (*inv-Mp f*) *n* = *inv-M* (*coeff f n*)
unfolding *inv-Mp-def*
by (*rule coeff-map-poly*, *insert m1*, *auto simp*: *inv-M-def*)

lemma *Mp-inv-Mp-id*[*simp*]: *Mp* (*inv-Mp f*) = *Mp f*
unfolding *poly-eq-iff Mp-coeff inv-Mp-coeff* **by** *simp*

lemma *inv-Mp-rev*: **assumes** *bnd*: ⋀ *n*. *2* ∗ *abs* (*coeff f n*) < *m*
shows *inv-Mp* (*Mp f*) = *f*
proof (*rule poly-eqI*)
fix *n*
define *c* **where** *c* = *coeff f n*
from *bnd*[*of n*, *folded c-def*] **have** *bnd*: *2* ∗ *abs c* < *m* **by** *auto*
show *coeff* (*inv-Mp* (*Mp f*)) *n* = *coeff f n* **unfolding** *inv-Mp-coeff Mp-coeff c-def*[*symmetric*]
using *inv-M-rev*[*OF bnd*] **.**
qed

lemma *mul-const-commute-below*: *mul-const x* (*mul-const y z*) = *mul-const y* (*mul-const x z*)
unfolding *mul-const-def* **by** (*metis mod-mult-right-eq mult.left-commute*)

context
fixes *p n*
and *sl-impl* :: (*int poly*, *int* × *int poly list*, *'state*)*subseqs-foldr-impl*
and *sli* :: *int* × *int poly list* ⇒ *int poly list* ⇒ *nat* ⇒ *'state* ⇒ *bool*
assumes *prime*: *prime p*
and *m*: *m* = *p*^*n*
and *n*: *n* ≠ *0*
and *sl-impl*: *correct-subseqs-foldr-impl* (λ*x*. *map-prod* (*mul-const x*) (*Cons x*)) *sl-impl sli*
begin
private definition *test-dvd-exec lu u ws* = (¬ *inv-Mp* (*Mp* (*smult lu* (*prod-mset ws*))) *dvd smult lu u*)

private definition *test-dvd u ws* = (∀ *v l*. *v dvd u* ⟶ *0* < *degree v* ⟶ *degree v* < *degree u*
⟶ ¬ *v* =*m smult l* (*prod-mset ws*))

private definition *large-m u vs* = (∀ *v n*. *v dvd u* ⟶ *degree v* ≤ *degree-bound vs* ⟶ *2* ∗ *abs* (*coeff v n*) < *m*)

lemma *large-m-factor*: *large-m u vs* ⟹ *v dvd u* ⟹ *large-m v vs*
unfolding *large-m-def* **using** *dvd-trans* **by** *auto*

lemma *test-dvd-factor*: **assumes** *u*: *u* ≠ *0* **and** *test*: *test-dvd u ws* **and** *vu*: *v dvd*

u
shows *test-dvd v ws*
proof −
from *vu* **obtain** *w* **where** *uv*: $u = v * w$ **unfolding** *dvd-def* **by** *auto*
from *u* **have** *deg*: *degree u* = *degree v* + *degree w* **unfolding** *uv*
by (*subst degree-mult-eq*, *auto*)
show *?thesis* **unfolding** *test-dvd-def*
proof (*intro allI impI*, *goal-cases*)
case (*1 f l*)
from *1*(*1*) **have** *fu*: *f dvd u* **unfolding** *uv* **by** *auto*
from *1*(*3*) **have** *deg*: *degree f* < *degree u* **unfolding** *deg* **by** *auto*
from *test*[*unfolded test-dvd-def*, *rule-format*, *OF fu 1*(*2*) *deg*]
show *?case* **.**
qed
qed

lemma *coprime-exp-mod*: *coprime lu p* ⟹ *prime p* ⟹ *n* ≠ *0* ⟹ *lu mod p* ^ *n* ≠ *0*
by (*auto simp add*: *abs-of-pos prime-gt-0-int*)

interpretation *correct-subseqs-foldr-impl* λ*x*. *map-prod* (*mul-const x*) (*Cons x*) *sl-impl sli* **by** *fact*

lemma *reconstruction*: **assumes**
res: *reconstruction sl-impl m2 state u* (*smult lu u*) *lu d r vs res cands* = *fs*
and *f*: *f* = *u* ∗ *prod-list res*
and *meas*: *meas* = (*r* − *d*, *cands*)
and *dr*: *d* + *d* ≤ *r*
and *r*: *r* = *length vs*
and *cands*: *set cands* ⊆ *S* (*lu*,[]) *vs d*
and *d0*: *d* = *0* ⟹ *cands* = []
and *lu*: *lu* = *lead-coeff u*
and *factors*: *unique-factorization-m u* (*lu*,*mset vs*)
and *sf*: *poly-mod.square-free-m p u*
and *cop*: *coprime lu p*
and *norm*: ⋀ *v*. *v* ∈ *set vs* ⟹ *Mp v* = *v*
and *tests*: ⋀ *ws*. *ws* ⊆# *mset vs* ⟹ *ws* ≠ {#} ⟹
size ws < *d* ∨ *size ws* = *d* ∧ *ws* ∉ (*mset o snd*) ‘ *set cands*
⟹ *test-dvd u ws*
and *irr*: ⋀ *f*. *f* ∈ *set res* ⟹ *irreducible*$_d$ *f*
and *deg*: *degree u* > *0*
and *cands-ne*: *cands* ≠ [] ⟹ *d* < *r*
and *large*: ∀ *v n*. *v dvd smult lu u* ⟶ *degree v* ≤ *degree-bound vs*
⟶ *2* ∗ *abs* (*coeff v n*) < *m*
and *f0*: *f* ≠ *0*
and *state*: *sli* (*lu*,[]) *vs d state*
and *m2*: *m2* = *m div 2*
shows *f* = *prod-list fs* ∧ (∀ *fi* ∈ *set fs*. *irreducible*$_d$ *fi*)
proof −

from *large* **have** *large*: *large-m* (*smult lu u*) *vs* **unfolding** *large-m-def* **by** *auto*
interpret *p*: *poly-mod-prime p* **using** *prime* **by** *unfold-locales*
define *R* **where** *R* ≡ *measures* [
λ (*n* :: *nat*,*cds* :: (*int* × *int poly list*) *list*). *n*,
λ (*n*,*cds*). *length cds*]
have *wf*: *wf R* **unfolding** *R-def* **by** *simp*
have *mset-snd-S*: ⋀ *vs lu d*. (*mset* ∘ *snd*) ‘ *S* (*lu*,[]) *vs d* =
{ *ws*. *ws* ⊆# *mset vs* ∧ *size ws* = *d*}
by (*fold mset-subseqs-size image-comp*, *unfold S-def image-Collect*, *auto*)
have *inv-M2*[*simp*]: *inv-M2 m2* = *inv-M* **unfolding** *inv-M2-def m2 inv-M-def*
by (*intro ext*, *auto*)
have *inv-Mp2*[*simp*]: *inv-Mp2 m2* = *inv-Mp* **unfolding** *inv-Mp2-def inv-Mp-def*
by *simp*
have *p-Mp*[*simp*]: ⋀ *f*. *p.Mp* (*Mp f*) = *p.Mp f* **using** *m p.m1 n Mp-Mp-pow-is-Mp*
by *blast*
{
fix *u lu vs*
assume *eq*: *Mp u* = *Mp* (*smult lu* (*prod-mset vs*)) **and** *cop*: *coprime lu p* **and**
size: *size vs* ≠ *0*
and *mi*: ⋀ *v*. *v* ∈# *vs* ⟹ *irreducible*$_d$*-m v* ∧ *monic v*
from *cop p.m1* **have** *lu0*: *lu* ≠ *0* **by** *auto*
from *size* **have** *vs* ≠ {#} **by** *auto*
then obtain *v vs*′ **where** *vs-v*: *vs* = *vs*′ + {#*v*#} **by** (*cases vs*, *auto*)
have *mon*: *monic* (*prod-mset vs*)
by (*rule monic-prod-mset*, *insert mi*, *auto*)
hence *vs0*: *prod-mset vs* ≠ *0* **by** (*metis coeff-0 zero-neq-one*)
from *mon* **have** *l-vs*: *lead-coeff* (*prod-mset vs*) = *1* **.**
have *deg-ws*: *degree-m* (*smult lu* (*prod-mset vs*)) = *degree* (*smult lu* (*prod-mset*
vs))
by (*rule degree-m-eq*[*OF - m1*], *unfold lead-coeff-smult*,
insert cop n p.m1 l-vs, *auto simp*: *m*)
with *eq* **have** *degree-m u* = *degree* (*smult lu* (*prod-mset vs*)) **by** *auto*
also have ... = *degree* (*prod-mset vs*′ ∗ *v*) **unfolding** *degree-smult-eq vs-v*
using *lu0* **by** (*simp add*:*ac-simps*)
also have ... = *degree* (*prod-mset vs*′) + *degree v*
by (*rule degree-mult-eq*, *insert vs0*[*unfolded vs-v*], *auto*)
also have ... ≥ *degree v* **by** *simp*
finally have *deg-v*: *degree v* ≤ *degree-m u* **.**
from *mi*[*unfolded vs-v*, *of v*] **have** *irreducible*$_d$*-m v* **by** *auto*
hence *0* < *degree-m v* **unfolding** *irreducible*$_d$*-m-def* **by** *auto*
also have ... ≤ *degree v* **by** (*rule degree-m-le*)
also have ... ≤ *degree-m u* **by** (*rule deg-v*)
also have ... ≤ *degree u* **by** (*rule degree-m-le*)
finally have *degree u* > *0* **by** *auto*
} **note** *deg-non-zero* = *this*
{
fix *u* :: *int poly* **and** *vs* :: *int poly list* **and** *d* :: *nat*
assume *deg-u*: *degree u* > *0*
and *cop*: *coprime* (*lead-coeff u*) *p*

and *uf*: *unique-factorization-m u* (*lead-coeff u*, *mset vs*)
and *sf*: *p.square-free-m u*
and *norm*: $\bigwedge$ *v*. *v* ∈ *set vs* ⟹ *Mp v* = *v*
and *d*: *size* (*mset vs*) < *d* + *d*
and *tests*: $\bigwedge$ *ws*. *ws* ⊆# *mset vs* ⟹ *ws* ≠ {#} ⟹ *size ws* < *d* ⟹ *test-dvd u ws*
from *deg-u* **have** *u0*: *u* ≠ *0* **by** *auto*
have *irreducible*$_d$ *u*
proof (*rule irreducible*$_d$*I*[*OF deg-u*])
fix *q q*′ :: *int poly*
assume *deg*: *degree q* > *0 degree q* < *degree u degree q*′ > *0 degree q*′ < *degree u*
and *uq*: *u* = *q* ∗ *q*′
then have *qu*: *q dvd u* **and** *q*′*u*: *q*′ *dvd u* **by** *auto*
from *u0* **have** *deg-u*: *degree u* = *degree q* + *degree q*′ **unfolding** *uq*
by (*subst degree-mult-eq*, *auto*)
from *coprime-lead-coeff-factor*[*OF prime cop*[*unfolded uq*]]
have *cop-q*: *coprime* (*lead-coeff q*) *p coprime* (*lead-coeff q*′) *p* **by** *auto*
from *unique-factorization-m-factor*[*OF prime uf*[*unfolded uq*] - - *n m*, *folded uq*,
OF cop sf]
obtain *fs gs l* **where** *uf-q*: *unique-factorization-m q* (*lead-coeff q*, *fs*)
and *uf-q*′: *unique-factorization-m q*′ (*lead-coeff q*′, *gs*)
and *Mf-eq*: *Mf* (*l*, *mset vs*) = *Mf* (*lead-coeff q* ∗ *lead-coeff q*′, *fs* + *gs*)
and *fs-id*: *image-mset Mp fs* = *fs*
and *gs-id*: *image-mset Mp gs* = *gs* **by** *auto*
from *Mf-eq fs-id gs-id* **have** *image-mset Mp* (*mset vs*) = *fs* + *gs*
unfolding *Mf-def* **by** *auto*
also have *image-mset Mp* (*mset vs*) = *mset vs* **using** *norm* **by** (*induct vs*, *auto*)
finally have *eq*: *mset vs* = *fs* + *gs* **by** *simp*
from *uf-q*[*unfolded unique-factorization-m-alt-def factorization-m-def split*]
have *q-eq*: *q* =*m smult* (*lead-coeff q*) (*prod-mset fs*) **by** *auto*
have *degree-m q* = *degree q*
by (*rule degree-m-eq*[*OF - m1*], *insert cop-q*(*1*) *n p.m1*, *unfold m*,
auto simp:)
with *q-eq* **have** *degm-q*: *degree q* = *degree* (*Mp* (*smult* (*lead-coeff q*) (*prod-mset fs*))) **by** *auto*
with *deg* **have** *fs-nempty*: *fs* ≠ {#}
by (*cases fs*; *cases lead-coeff q* = *0*; *auto simp*: *Mp-def*)
from *uf-q*′[*unfolded unique-factorization-m-alt-def factorization-m-def split*]
have *q*′*-eq*: *q*′ =*m smult* (*lead-coeff q*′) (*prod-mset gs*) **by** *auto*
have *degree-m q*′ = *degree q*′
by (*rule degree-m-eq*[*OF - m1*], *insert cop-q*(*2*) *n p.m1*, *unfold m*,
auto simp:)
with *q*′*-eq* **have** *degm-q*′: *degree q*′ = *degree* (*Mp* (*smult* (*lead-coeff q*′) (*prod-mset gs*))) **by** *auto*
with *deg* **have** *gs-nempty*: *gs* ≠ {#}
by (*cases gs*; *cases lead-coeff q*′ = *0*; *auto simp*: *Mp-def*)

from *eq* **have** *size*: *size fs* + *size gs* = *size* (*mset vs*) **by** *auto*
with *d* **have** *choice*: *size fs* < *d* ∨ *size gs* < *d* **by** *auto*
from *choice* **show** *False*
proof
assume *fs*: *size fs* < *d*
from *eq* **have** *sub*: *fs* ⊆# *mset vs* **using** *mset-subset-eq-add-left*[*of fs gs*] **by** *auto*
have *test-dvd u fs*
by (*rule tests*[*OF sub fs-nempty*, *unfolded Nil*], *insert fs*, *auto*)
from *this*[*unfolded test-dvd-def*] *uq deg q-eq* **show** *False* **by** *auto*
next
assume *gs*: *size gs* < *d*
from *eq* **have** *sub*: *gs* ⊆# *mset vs* **using** *mset-subset-eq-add-left*[*of fs gs*] **by** *auto*
have *test-dvd u gs*
by (*rule tests*[*OF sub gs-nempty*, *unfolded Nil*], *insert gs*, *auto*)
from *this*[*unfolded test-dvd-def*] *uq deg q′-eq* **show** *False* **unfolding** *uq* **by** *auto*
qed
qed
} **note** *irreducible*$_d$*-via-tests* = *this*
show *?thesis* **using** *assms*(1−16) *large state*
proof (*induct meas arbitrary*: *u lu d r vs res cands state rule*: *wf-induct*[*OF wf*])
case (*1 meas u lu d r vs res cands state*)
note *IH* = *1*(*1*)[*rule-format*]
note *res* = *1*(*2*)[*unfolded reconstruction.simps*[**where** *cands* = *cands*]]
note *f* = *1*(*3*)
note *meas* = *1*(*4*)
note *dr* = *1*(*5*)
note *r* = *1*(*6*)
note *cands* = *1*(*7*)
note *d0* = *1*(*8*)
note *lu* = *1*(*9*)
note *factors* = *1*(*10*)
note *sf* = *1*(*11*)
note *cop* = *1*(*12*)
note *norm* = *1*(*13*)
note *tests* = *1*(*14*)
note *irr* = *1*(*15*)
note *deg-u* = *1*(*16*)
note *cands-empty* = *1*(*17*)
note *large* = *1*(*18*)
note *state* = *1*(*19*)
from *unique-factorization-m-zero*[*OF factors*]
have *Mlu0*: *M lu* ≠ *0* **by** *auto*
from *Mlu0* **have** *lu0*: *lu* ≠ *0* **by** *auto*
from *this*[*unfolded lu*] **have** *u0*: *u* ≠ *0* **by** *auto*
from *unique-factorization-m-imp-factorization*[*OF factors*]

```
    have fact: factorization-m u (lu,mset vs) by auto
    from this[unfolded factorization-m-def split] norm
    have vs: u =m smult lu (prod-list vs) and
      vs-mi: ⋀ f. f∈#mset vs ⟹ irreducible_d-m f ∧ monic f by auto
    let ?luu = smult lu u
    show ?case
    proof (cases cands)
      case Nil
      note res = res[unfolded this]
      let ?d' = Suc d
      show ?thesis
      proof (cases r < ?d' + ?d')
        case True
        with res have fs: fs = u # res by (simp add: Let-def)
        from True[unfolded r] have size: size (mset vs) < ?d' + ?d' by auto
        have irreducible_d u
         by (rule irreducible_d-via-tests[OF deg-u cop[unfolded lu] factors(1)[unfolded
lu]
          sf norm size tests], auto simp: Nil)
        with fs f irr show ?thesis by simp
      next
        case False
        with dr have dr: ?d' + ?d' ≤ r and dr': ?d' < r by auto
      obtain state' cands' where sln: next-subseqs-foldr sl-impl state = (cands',state')
by force
       from next-subseqs-foldr[OF sln state] have state': sli (lu,[]) vs (Suc d) state'
          and cands': set cands' = S (lu,[]) vs (Suc d) by auto
        let ?new = subseqs-length mul-const lu ?d' vs
         have R: ((r − Suc d, cands'), meas) ∈ R unfolding meas R-def using
False by auto
        from res False sln
       have fact: reconstruction sl-impl m2 state' u ?luu lu ?d' r vs res cands' = fs
by auto
        show ?thesis
         proof (rule IH[OF R fact f refl dr r - - lu factors sf cop norm - irr deg-u
dr' large state'], goal-cases)
          case (4 ws)
          show ?case
          proof (cases size ws = Suc d)
            case False
            with 4 have size ws < Suc d by auto
            thus ?thesis by (intro tests[OF 4(1−2)], unfold Nil, auto)
          next
            case True
           from 4(3)[unfolded cands' mset-snd-S] True 4(1) show ?thesis by auto
          qed
        qed (auto simp: cands')
      qed
    next
```

```
      case (Cons c cds)
      with d0 have d0: d > 0 by auto
      obtain lv' ws where c: c = (lv',ws) by force
      let ?lv = inv-M lv'
      define vb where vb ≡ inv-Mp (Mp (smult lu (prod-list ws)))
      note res = res[unfolded Cons c list.simps split]
      from cands[unfolded Cons c S-def] have ws: ws ∈ set (subseqs vs) length ws
= d
        and lv'': lv' = foldr mul-const ws lu by auto
      from subseqs-sub-mset[OF ws(1)] have ws-vs: mset ws ⊆# mset vs set ws ⊆
set vs
        using set-mset-mono subseqs-length-simple-False by auto fastforce
      have mon-ws: monic (prod-mset (mset ws))
        by (rule monic-prod-mset, insert ws-vs vs-mi, auto)
      have l-ws: lead-coeff (prod-mset (mset ws)) = 1 using mon-ws .
      have lv': M lv' = M (coeff (smult lu (prod-list ws)) 0)
        unfolding lv'' coeff-smult
        by (induct ws arbitrary: lu, auto simp: mul-const-def M-def coeff-mult-0)
           (metis mod-mult-right-eq mult.left-commute)
      show ?thesis
      proof (cases ?lv dvd coeff ?luu 0 ∧ vb dvd ?luu)
        case False
        have ndvd: ¬ vb dvd ?luu
        proof
          assume dvd: vb dvd ?luu
          hence coeff vb 0 dvd coeff ?luu 0 by (metis coeff-mult-0 dvd-def)
          with dvd False have ?lv ≠ coeff vb 0 by auto
          also have lv' = M lv' using ws(2) d0 unfolding lv''
            by (cases ws, force, simp add: M-def mul-const-def)
           also have inv-M (M lv') = coeff vb 0 unfolding vb-def inv-Mp-coeff
Mp-coeff lv' by simp
          finally show False by simp
        qed
        from False res
        have res: reconstruction sl-impl m2 state u ?luu lu d r vs res cds = fs
          unfolding vb-def Let-def by auto
        have R: ((r − d, cds), meas) ∈ R unfolding meas Cons R-def by auto
        from cands have cands: set cds ⊆ S (lu,[]) vs d
          unfolding Cons by auto
        show ?thesis
       proof (rule IH[OF R res f refl dr r cands - lu factors sf cop norm - irr deg-u
- large state], goal-cases)
          case (3 ws')
          show ?case
          proof (cases ws' = mset ws)
            case False
            show ?thesis
              by (rule tests[OF 3(1−2)], insert 3(3) False, force simp: Cons c)
          next
```

```
        case True
        have test: test-dvd-exec lu u ws'
          unfolding True test-dvd-exec-def using ndvd unfolding vb-def by
simp
        show ?thesis unfolding test-dvd-def
        proof (intro allI impI notI, goal-cases)
          case (1 v l)
          note deg-v = 1(2-3)
         from 1(1) obtain w where u: u = v * w unfolding dvd-def by auto
          from u0 have deg: degree u = degree v + degree w unfolding u
            by (subst degree-mult-eq, auto)
          define v' where v' = smult (lead-coeff w) v
          define w' where w' = smult (lead-coeff v) w
          let ?ws = smult (lead-coeff w * l) (prod-mset ws')
          from arg-cong[OF 1(4), of λ f. Mp (smult (lead-coeff w) f)]
          have v'-ws': Mp v' = Mp ?ws unfolding v'-def
            by simp
          from lead-coeff-factor[OF u, folded v'-def w'-def]
          have prod: ?luu = v' * w' and lc: lead-coeff v' = lu and lead-coeff w'
= lu
            unfolding lu by auto
         with lu0 have lc0: lead-coeff v ≠ 0 lead-coeff w ≠ 0 unfolding v'-def
w'-def by auto
          from deg-v have deg-w: 0 < degree w degree w < degree u unfolding
deg by auto
          from deg-v deg-w lc0
          have deg: 0 < degree v' degree v' < degree u 0 < degree w' degree w'
< degree u
            unfolding v'-def w'-def by auto
          from prod have v-dvd: v' dvd ?luu by auto
          with test[unfolded test-dvd-exec-def]
          have neq: v' ≠ inv-Mp (Mp (smult lu (prod-mset ws'))) by auto
          have deg-m-v': degree-m v' = degree v'
            by (rule degree-m-eq[OF - m1], unfold lc m,
            insert cop prime n coprime-exp-mod, auto)
          with v'-ws' have degree v' = degree-m ?ws by simp
          also have ... ≤ degree-m (prod-mset ws') by (rule degree-m-smult-le)
          also have ... = degree-m (prod-list ws) unfolding True by simp
          also have ... ≤ degree (prod-list ws) by (rule degree-m-le)
          also have ... ≤ degree-bound vs
            using ws-vs(1) ws(2) dr[unfolded r] degree-bound by auto
          finally have degree v' ≤ degree-bound vs .
         from inv-Mp-rev[OF large[unfolded large-m-def, rule-format, OF v-dvd
this]]
          have inv: inv-Mp (Mp v') = v' by simp
          from arg-cong[OF v'-ws', of inv-Mp, unfolded inv]
          have v': v' = inv-Mp (Mp ?ws) by auto
          have deg-ws: degree-m ?ws = degree ?ws
          proof (rule degree-m-eq[OF - m1],
```

unfold lead-coeff-smult True l-ws, *rule*)
assume *lead-coeff w ∗ l ∗ 1 mod m = 0*
hence *0*: *M* (*lead-coeff w ∗ l*) = *0* **unfolding** *M-def* **by** *simp*
have *Mp ?ws* = *Mp* (*smult* (*M* (*lead-coeff w ∗ l*)) (*prod-mset ws′*)) **by** *simp*
also have ... = *0* **unfolding** *0* **by** *simp*
finally have *Mp ?ws* = *0* **by** *simp*
hence *v′* = *0* **unfolding** *v′* **by** (*simp add*: *inv-Mp-def*)
with *deg* **show** *False* **by** *auto*
qed
from *arg-cong*[*OF v′*, *of* λ *f*. *lead-coeff* (*Mp f*), *simplified*]
have *M lu* = *M* (*lead-coeff v′*) **using** *lc* **by** *simp*
also have ... = *lead-coeff* (*Mp v′*)
by (*rule degree-m-eq-lead-coeff*[*OF deg-m-v′*, *symmetric*])
also have ... = *lead-coeff* (*Mp ?ws*)
using *arg-cong*[*OF v′*, *of* λ *f*. *lead-coeff* (*Mp f*)] **by** *simp*
also have ... = *M* (*lead-coeff ?ws*)
by (*rule degree-m-eq-lead-coeff*[*OF deg-ws*])
also have ... = *M* (*lead-coeff w ∗ l*) **unfolding** *lead-coeff-smult True l-ws* **by** *simp*
finally have *id*: *M lu* = *M* (*lead-coeff w ∗ l*) **.**
note *v′*
also have *Mp ?ws* = *Mp* (*smult* (*M* (*lead-coeff w ∗ l*)) (*prod-mset ws′*)) **by** *simp*
also have ... = *Mp* (*smult lu* (*prod-mset ws′*)) **unfolding** *id*[*symmetric*] **by** *simp*
finally show *False* **using** *neq* **by** *simp*
qed
qed
qed (*insert d0 Cons cands-empty*, *auto*)
next
case *True*
define *pp-vb* **where** *pp-vb* ≡ *primitive-part vb*
define *u′* **where** *u′* ≡ *u div pp-vb*
define *lu′* **where** *lu′* ≡ *lead-coeff u′*
let *?luu′* = *smult lu′ u′*
define *vs′* **where** *vs′* ≡ *fold remove1 ws vs*
obtain *state′ cands′* **where** *slc*: *subseqs-foldr sl-impl* (*lu′*,[]) *vs′ d* = (*cands′*, *state′*) **by** *force*
from *subseqs-foldr*[*OF slc*] **have** *state′*: *sli* (*lu′*,[]) *vs′ d state′*
and *cands′*: *set cands′* = *S* (*lu′*,[]) *vs′ d* **by** *auto*
let *?res′* = *pp-vb* # *res*
let *?r′* = *r* − *length ws*
note *defs* = *vb-def pp-vb-def u′-def lu′-def vs′-def slc*
from *fold-remove1-mset*[*OF subseqs-sub-mset*[*OF ws*(*1*)]]
have *vs-split*: *mset vs* = *mset vs′* + *mset ws* **unfolding** *vs′-def* **by** *auto*
hence *vs′-diff*: *mset vs′* = *mset vs* − *mset ws* **and** *ws-sub*: *mset ws* ⊆# *mset vs* **by** *auto*
from *arg-cong*[*OF vs-split*, *of size*]

have r': $?r' = length\ vs'$ **unfolding** $defs\ r$ **by** $simp$
from $arg\text{-}cong[OF\ vs\text{-}split,\ of\ prod\text{-}mset]$
have $prod\text{-}vs$: $prod\text{-}list\ vs = prod\text{-}list\ vs' * prod\text{-}list\ ws$ **by** $simp$
from $arg\text{-}cong[OF\ vs\text{-}split,\ of\ set\text{-}mset]$ **have** $set\text{-}vs$: $set\ vs = set\ vs' \cup set\ ws$ **by** $auto$
note $inv = inverse\text{-}mod\text{-}coprime\text{-}exp[OF\ m\ prime\ n]$
note $p\text{-}inv = p.inverse\text{-}mod\text{-}coprime[OF\ prime]$
from $True\ res\ slc$
have res: $(if\ ?r' < d + d\ then\ u' \#\ ?res'\ else\ reconstruction\ sl\text{-}impl\ m2\ state'$
$u'\ ?luu'\ lu'\ d\ ?r'\ vs'\ ?res'\ cands') = fs$
unfolding $Let\text{-}def\ defs$ **by** $auto$
from $True$ **have** dvd: $vb\ dvd\ ?luu$ **by** $simp$
from $dvd\text{-}smult\text{-}int[OF\ lu0\ this]$ **have** ppu: $pp\text{-}vb\ dvd\ u$ **unfolding** $defs$ **by** $simp$
hence u: $u = pp\text{-}vb * u'$ **unfolding** $u'\text{-}def$
by $(metis\ dvdE\ mult\text{-}eq\text{-}0\text{-}iff\ nonzero\text{-}mult\text{-}div\text{-}cancel\text{-}left)$
hence uu': $u'\ dvd\ u$ **unfolding** $dvd\text{-}def$ **by** $auto$
have f: $f = u' * prod\text{-}list\ ?res'$ **using** $f\ u$ **by** $auto$
let $?fact = smult\ lu\ (prod\text{-}mset\ (mset\ ws))$
have $Mp\text{-}vb$: $Mp\ vb = Mp\ (smult\ lu\ (prod\text{-}list\ ws))$ **unfolding** $vb\text{-}def$ **by** $simp$
have $pp\text{-}vb\text{-}vb$: $smult\ (content\ vb)\ pp\text{-}vb = vb$ **unfolding** $pp\text{-}vb\text{-}def$ **by** $(rule\ content\text{-}times\text{-}primitive\text{-}part)$
{
have $smult\ (content\ vb)\ u = (smult\ (content\ vb)\ pp\text{-}vb) * u'$ **unfolding** u **by** $simp$
also have $smult\ (content\ vb)\ pp\text{-}vb = vb$ **by** $fact$
finally have $smult\ (content\ vb)\ u = vb * u'$ **by** $simp$
from $arg\text{-}cong[OF\ this,\ of\ Mp]$
have $Mp\ (Mp\ vb * u') = Mp\ (smult\ (content\ vb)\ u)$ **by** $simp$
hence $Mp\ (smult\ (content\ vb)\ u) = Mp\ (?fact * u')$ **unfolding** $Mp\text{-}vb$ **by** $simp$
} **note** $prod = this$
from $arg\text{-}cong[OF\ this,\ of\ p.Mp]$
have $prod'$: $p.Mp\ (smult\ (content\ vb)\ u) = p.Mp\ (?fact * u')$ **by** $simp$
from dvd **have** $lead\text{-}coeff\ vb\ dvd\ lead\text{-}coeff\ (smult\ lu\ u)$
by $(metis\ dvd\text{-}def\ lead\text{-}coeff\text{-}mult)$
hence $ldvd$: $lead\text{-}coeff\ vb\ dvd\ lu * lu$ **unfolding** $lead\text{-}coeff\text{-}smult\ lu$ **by** $simp$
from cop **have** $cop\text{-}lu$: $coprime\ (lu * lu)\ p$
by $simp$
from $coprime\text{-}divisors\ [OF\ ldvd\ dvd\text{-}refl]\ cop\text{-}lu$
have $cop\text{-}lvb$: $coprime\ (lead\text{-}coeff\ vb)\ p$ **by** $simp$
then have $cop\text{-}vb$: $coprime\ (content\ vb)\ p$
by $(auto\ intro: coprime\text{-}divisors[OF\ content\text{-}dvd\text{-}coeff\ dvd\text{-}refl])$
from u **have** $u'\ dvd\ u$ **unfolding** $dvd\text{-}def$ **by** $auto$
hence $lead\text{-}coeff\ u'\ dvd\ lu$ **unfolding** lu **by** $(metis\ dvd\text{-}def\ lead\text{-}coeff\text{-}mult)$
from $coprime\text{-}divisors[OF\ this\ dvd\text{-}refl]\ cop$
have $coprime\ (lead\text{-}coeff\ u')\ p$ **by** $simp$

hence *coprime* (*lu* ∗ *lead-coeff u′*) *p* **and** *cop-lu′*: *coprime lu′ p*
using *cop* **by** (*auto simp*: *lu′-def*)
hence *cop′*: *coprime* (*lead-coeff* (*?fact* ∗ *u′*)) *p*
unfolding *lead-coeff-mult lead-coeff-smult l-ws* **by** *simp*
have *p.square-free-m* (*smult* (*content vb*) *u*) **using** *cop-vb sf p-inv*
by (*auto intro*!: *p.square-free-m-smultI*)
from *p.square-free-m-cong*[*OF this prod′*]
have *sf′*: *p.square-free-m* (*?fact* ∗ *u′*) **by** *simp*
from *p.square-free-m-factor*[*OF this*]
have *sf-u′*: *p.square-free-m u′* **by** *simp*
have *unique-factorization-m* (*smult* (*content vb*) *u*) (*lu* ∗ *content vb*, *mset vs*)
using *cop-vb factors inv* **by** (*auto intro*: *unique-factorization-m-smult*)
from *unique-factorization-m-cong*[*OF this prod*]
have *uf*: *unique-factorization-m* (*?fact* ∗ *u′*) (*lu* ∗ *content vb*, *mset vs*) **.**
{
from *unique-factorization-m-factor*[*OF prime uf cop′ sf′ n m*]
obtain *fs gs* **where** *uf1*: *unique-factorization-m ?fact* (*lu*, *fs*)
and *uf2*: *unique-factorization-m u′* (*lu′*, *gs*)
and *eq*: *Mf* (*lu* ∗ *content vb*, *mset vs*) = *Mf* (*lu* ∗ *lead-coeff u′*, *fs* + *gs*)
unfolding *lead-coeff-smult l-ws lu′-def*
by *auto*
have *factorization-m ?fact* (*lu*, *mset ws*)
unfolding *factorization-m-def split* **using** *set-vs vs-mi norm* **by** *auto*
with *uf1*[*unfolded unique-factorization-m-alt-def*] **have** *Mf* (*lu*,*mset ws*) = *Mf* (*lu*, *fs*)
by *blast*
hence *fs-ws*: *image-mset Mp fs* = *image-mset Mp* (*mset ws*) **unfolding** *Mf-def split* **by** *auto*
from *eq*[*unfolded Mf-def split*]
have *image-mset Mp* (*mset vs*) = *image-mset Mp fs* + *image-mset Mp gs* **by** *auto*
from *this*[*unfolded fs-ws vs-split*] **have** *gs*: *image-mset Mp gs* = *image-mset Mp* (*mset vs′*)
by (*simp add*: *ac-simps*)
from *uf1* **have** *uf1*: *unique-factorization-m ?fact* (*lu*, *mset ws*)
unfolding *unique-factorization-m-def Mf-def split fs-ws* **by** *simp*
from *uf2* **have** *uf2*: *unique-factorization-m u′* (*lu′*, *mset vs′*)
unfolding *unique-factorization-m-def Mf-def split gs* **by** *simp*
note *uf1 uf2*
}
hence *factors*: *unique-factorization-m u′* (*lu′*, *mset vs′*)
unique-factorization-m ?fact (*lu*, *mset ws*) **by** *auto*
have *lu′*: *lu′* = *lead-coeff u′* **unfolding** *lu′-def* **by** *simp*
have *vb0*: *vb* ≠ *0* **using** *dvd lu0 u0* **by** *auto*
from *ws*(*2*) **have** *size-ws*: *size* (*mset ws*) = *d* **by** *auto*
with *d0* **have** *size-ws0*: *size* (*mset ws*) ≠ *0* **by** *auto*
then obtain *w ws′* **where** *ws-w*: *ws* = *w* # *ws′* **by** (*cases ws*, *auto*)
from *Mp-vb* **have** *Mp-vb′*: *Mp vb* = *Mp* (*smult lu* (*prod-mset* (*mset ws*)))

by *auto*
have *deg-vb*: *degree vb* > *0*
by (*rule deg-non-zero*[*OF Mp-vb' cop size-ws0 vs-mi*], *insert vs-split*, *auto*)
also have *degree vb* = *degree pp-vb* **using** *arg-cong*[*OF pp-vb-vb*, *of degree*]
unfolding *degree-smult-eq* **using** *vb0* **by** *auto*
finally have *deg-pp*: *degree pp-vb* > *0* **by** *auto*
hence *pp-vb0*: *pp-vb* ≠ *0* **by** *auto*
from *factors*(*1*)[*unfolded unique-factorization-m-alt-def factorization-m-def*]
have *eq-u'*: *Mp u'* = *Mp* (*smult lu'* (*prod-mset* (*mset vs'*))) **by** *auto*
from *r'*[*unfolded ws*(*2*)] *dr* **have** *length vs'* + *d* = *r* **by** *auto*
from *this cands-empty*[*unfolded Cons*] **have** *size* (*mset vs'*) ≠ *0* **by** *auto*
from *deg-non-zero*[*OF eq-u' cop-lu' this vs-mi*]
have *deg-u'*: *degree u'* > *0* **unfolding** *vs-split* **by** *auto*
have *irr-pp*: *irreducible*$_d$ *pp-vb*
proof (*rule irreducible*$_d$*I*[*OF deg-pp*])
fix *q r* :: *int poly*
assume *deg-q*: *degree q* > *0 degree q* < *degree pp-vb*
and *deg-r*: *degree r* > *0 degree r* < *degree pp-vb*
and *pp-qr*: *pp-vb* = *q* ∗ *r*
then have *qvb*: *q dvd pp-vb* **by** *auto*
from *dvd-trans*[*OF qvb ppu*] **have** *qu*: *q dvd u* **.**
have *degree pp-vb* = *degree q* + *degree r* **unfolding** *pp-qr*
by (*subst degree-mult-eq*, *insert pp-qr pp-vb0*, *auto*)
have *uf*: *unique-factorization-m* (*smult* (*content vb*) *pp-vb*) (*lu*, *mset ws*)
unfolding *pp-vb-vb*
by (*rule unique-factorization-m-cong*[*OF factors*(*2*)], *insert Mp-vb*, *auto*)
from *unique-factorization-m-smultD*[*OF uf inv*] *cop-vb*
have *uf*: *unique-factorization-m pp-vb* (*lu* ∗ *inverse-mod* (*content vb*) *m*, *mset ws*) **by** *auto*
from *ppu* **have** *lead-coeff pp-vb dvd lu* **unfolding** *lu* **by** (*metis dvd-def lead-coeff-mult*)
from *coprime-divisors*[*OF this dvd-refl*] *cop*
have *cop-pp*: *coprime* (*lead-coeff pp-vb*) *p* **by** *simp*
from *coprime-lead-coeff-factor*[*OF prime cop-pp*[*unfolded pp-qr*]]
have *cop-qr*: *coprime* (*lead-coeff q*) *p coprime* (*lead-coeff r*) *p* **by** *auto*
from *p.square-free-m-factor*[*OF sf*[*unfolded u*]]
have *sf-pp*: *p.square-free-m pp-vb* **by** *simp*
from *unique-factorization-m-factor*[*OF prime uf*[*unfolded pp-qr*] - - *n m*, *folded pp-qr*, *OF cop-pp sf-pp*]
obtain *fs gs l* **where** *uf-q*: *unique-factorization-m q* (*lead-coeff q*, *fs*)
and *uf-r*: *unique-factorization-m r* (*lead-coeff r*, *gs*)
and *Mf-eq*: *Mf* (*l*, *mset ws*) = *Mf* (*lead-coeff q* ∗ *lead-coeff r*, *fs* + *gs*)
and *fs-id*: *image-mset Mp fs* = *fs*
and *gs-id*: *image-mset Mp gs* = *gs* **by** *auto*
from *Mf-eq* **have** *image-mset Mp* (*mset ws*) = *image-mset Mp fs* + *image-mset Mp gs*
unfolding *Mf-def* **by** *auto*
also have *image-mset Mp* (*mset ws*) = *mset ws* **using** *norm ws-vs*(*2*) **by** (*induct ws*, *auto*)

finally have *eq*: *mset ws = image-mset Mp fs + image-mset Mp gs* **by** *simp*
from *arg-cong*[*OF this*, *of size*, *unfolded size-ws*] **have** *size*: *size fs + size gs = d* **by** *auto*
from *uf-q*[*unfolded unique-factorization-m-alt-def factorization-m-def split*]
have *q-eq*: *q =m smult* (*lead-coeff q*) (*prod-mset fs*) **by** *auto*
have *degree-m q = degree q*
by (*rule degree-m-eq*[*OF - m1*], *insert cop-qr*(*1*) *n p.m1*, *unfold m*, *auto simp*:)
with *q-eq* **have** *degm-q*: *degree q = degree* (*Mp* (*smult* (*lead-coeff q*) (*prod-mset fs*))) **by** *auto*
with *deg-q* **have** *fs-nempty*: *fs ≠* {#}
by (*cases fs*; *cases lead-coeff q = 0*; *auto simp*: *Mp-def*)
from *uf-r*[*unfolded unique-factorization-m-alt-def factorization-m-def split*]
have *r-eq*: *r =m smult* (*lead-coeff r*) (*prod-mset gs*) **by** *auto*
have *degree-m r = degree r*
by (*rule degree-m-eq*[*OF - m1*], *insert cop-qr*(*2*) *n p.m1*, *unfold m*, *auto simp*:)
with *r-eq* **have** *degm-r*: *degree r = degree* (*Mp* (*smult* (*lead-coeff r*) (*prod-mset gs*))) **by** *auto*
with *deg-r* **have** *gs-nempty*: *gs ≠* {#}
by (*cases gs*; *cases lead-coeff r = 0*; *auto simp*: *Mp-def*)
from *gs-nempty* **have** *size gs ≠ 0* **by** *auto*
with *size* **have** *size-fs*: *size fs < d* **by** *linarith*
note ∗ = *tests*[*unfolded test-dvd-def*, *rule-format*, *OF - fs-nempty - qu*, *of lead-coeff q*]
from *ppu* **have** *degree pp-vb ≤ degree u*
using *dvd-imp-degree-le u0* **by** *blast*
with *deg-q q-eq size-fs*
have ¬ *fs ⊆*# *mset vs* **by** (*auto dest*!:∗)
thus *False* **unfolding** *vs-split eq fs-id gs-id* **using** *mset-subset-eq-add-left*[*of fs mset vs′ + gs*]
by (*auto simp*: *ac-simps*)
qed
{
fix *ws′*
assume ∗: *ws′ ⊆*# *mset vs′ ws′ ≠* {#}
size ws′ < d ∨ size ws′ = d ∧ ws′ ∉ (*mset ∘ snd*) ‘ *set cands′*
from ∗(*1*) **have** *ws′ ⊆*# *mset vs* **unfolding** *vs-split*
by (*simp add*: *subset-mset.add-increasing2*)
from *tests*[*OF this* ∗(*2*)] ∗(*3*)[*unfolded cands′ mset-snd-S*] ∗(*1*)
have *test-dvd u ws′* **by** *auto*
from *test-dvd-factor*[*OF u0 this*[*unfolded lu*] *uu′*]
have *test-dvd u′ ws′* **.**
} **note** *tests′ = this*
show *?thesis*
proof (*cases ?r′ < d + d*)
case *True*
with *res* **have** *res*: *fs = u′ # ?res′* **by** *auto*

```
        from True r' have size: size (mset vs') < d + d by auto
        have irreducible_d u'
              by (rule irreducible_d-via-tests[OF deg-u' cop-lu'[unfolded lu'] fac-
tors(1)[unfolded lu']
          sf-u' norm size tests'], insert set-vs, auto)
        with f res irr irr-pp show ?thesis by auto
      next
        case False
       have res: reconstruction sl-impl m2 state' u' ?luu' lu' d ?r' vs' ?res' cands'
= fs
          using False res by auto
        from False have dr: d + d ≤ ?r' by auto
         from False dr r r' d0 ws Cons have le: ?r' − d < r − d by (cases ws,
auto)
        hence R: ((?r' − d, cands'), meas) ∈ R unfolding meas R-def by simp
        have dr': d < ?r' using le False ws(2) by linarith
        have luu': lu' dvd lu using ‹lead-coeff u' dvd lu› unfolding lu' .
        have large-m (smult lu' u') vs
          by (rule large-m-factor[OF large dvd-dvd-smult], insert uu' luu')
        moreover have degree-bound vs' ≤ degree-bound vs
          unfolding vs'-def degree-bound-def by (rule max-factor-degree-mono)
        ultimately have large': large-m (smult lu' u') vs' unfolding large-m-def
by auto
        show ?thesis
          by (rule IH[OF R res f refl dr r' - - lu' factors(1) sf-u' cop-lu' norm
tests' - deg-u'
          dr' large' state'], insert irr irr-pp d0 Cons set-vs, auto simp: cands')
      qed
    qed
  qed
 qed
qed
end
end


definition zassenhaus-reconstruction ::
  int poly list ⇒ int ⇒ nat ⇒ int poly ⇒ int poly list where
  zassenhaus-reconstruction vs p n f = (let
    mul = poly-mod.mul-const (p^n);
    sl-impl = my-subseqs.impl (λx. map-prod (mul x) (Cons x))
    in zassenhaus-reconstruction-generic sl-impl vs p n f)

context
  fixes p n f hs
  assumes prime: prime p
  and cop: coprime (lead-coeff f) p
  and sf: poly-mod.square-free-m p f
  and deg: degree f > 0
```

```
  and bh: berlekamp-hensel p n f = hs
  and bnd: 2 * |lead-coeff f| * factor-bound f (degree-bound hs) < p ^ n
begin

private lemma n: n ≠ 0
proof
  assume n: n = 0
  hence pn: p^n = 1 by auto
  let ?f = smult (lead-coeff f) f
  let ?d = degree-bound hs
  have f: f ≠ 0 using deg by auto
  hence lead-coeff f ≠ 0 by auto
  hence lf: abs (lead-coeff f) > 0 by auto
  obtain c d where c: factor-bound f (degree-bound hs) = c abs (lead-coeff f) = d
by auto
  {
    assume *: 1 ≤ c 2 * d * c < 1 0 < d
    hence 1 ≤ d by auto
    from mult-mono[OF this *(1)] * have 1 ≤ d * c by auto
    hence 2 * d * c ≥ 2 by auto
    with * have False by auto
  } note tedious = this
  have 1 ≤ factor-bound f ?d
    using factor-bound[OF f, of 1 ?d 0] by auto
  also have ... = 0 using bnd unfolding pn
    using factor-bound-ge-0[of f degree-bound hs, OF f] lf unfolding c
    by (cases c ≥ 1; insert tedious, auto)
  finally show False by simp
qed

interpretation p: poly-mod-prime p using prime by unfold-locales

lemma zassenhaus-reconstruction-generic:
  assumes sl-impl: correct-subseqs-foldr-impl (λv. map-prod (poly-mod.mul-const
(p^n) v) (Cons v)) sl-impl sli
  and res: zassenhaus-reconstruction-generic sl-impl hs p n f = fs
  shows f = prod-list fs ∧ (∀ fi ∈ set fs. irreducible_d fi)
proof −
  let ?lc = lead-coeff f
  let ?ff = smult ?lc f
  let ?q = p^n
  have p1: p > 1 using prime unfolding prime-int-iff by simp
  interpret poly-mod-2 p^n using p1 n unfolding poly-mod-2-def by simp
 obtain cands state where slc: subseqs-foldr sl-impl (lead-coeff f, []) hs 0 = (cands,
state) by force
  interpret correct-subseqs-foldr-impl λx. map-prod (mul-const x) (Cons x) sl-impl
sli by fact
  from subseqs-foldr[OF slc] have state: sli (lead-coeff f, []) hs 0 state by auto
 from res[unfolded zassenhaus-reconstruction-generic-def bh split Let-def slc fst-conv]
```

have *res*: *reconstruction sl-impl* (*?q div 2*) *state f ?ff ?lc 0* (*length hs*) *hs* [] [] = *fs* **by** *auto*
from *p.berlekamp-hensel-unique*[*OF cop sf bh n*]
have *ufact*: *unique-factorization-m f* (*?lc*, *mset hs*) **by** *simp*
note *bh* = *p.berlekamp-hensel*[*OF cop sf bh n*]
from *deg* **have** *f0*: *f* ≠ *0* **and** *lf0*: *?lc* ≠ *0* **by** *auto*
hence *ff0*: *?ff* ≠ *0* **by** *auto*
have *bnd*: ∀ *g k*. *g dvd ?ff* ⟶ *degree g* ≤ *degree-bound hs* ⟶ *2* ∗ |*coeff g k*| < *p* ^ *n*
proof (*intro allI impI*, *goal-cases*)
case (*1 g k*)
from *factor-bound-smult*[*OF f0 lf0 1*, *of k*]
have |*coeff g k*| ≤ |*?lc*| ∗ *factor-bound f* (*degree-bound hs*) **.**
hence *2* ∗ |*coeff g k*| ≤ *2* ∗ |*?lc*| ∗ *factor-bound f* (*degree-bound hs*) **by** *auto*
also have ... < *p*^*n* **using** *bnd* **.**
finally show *?case* **.**
qed
note *bh'* = *bh*[*unfolded factorization-m-def split*]
have *deg-f*: *degree-m f* = *degree f*
using *cop unique-factorization-m-zero* [*OF ufact*] *n*
by (*auto simp add*: *M-def intro*: *degree-m-eq* [*OF - m1*])
have *mon-hs*: *monic* (*prod-list hs*) **using** *bh'* **by** (*auto intro*: *monic-prod-list*)
have *Mlc*: *M ?lc* ∈ {*1* ..< *p*^*n*}
by (*rule prime-cop-exp-poly-mod*[*OF prime cop n*])
hence *?lc* ≠ *0* **by** *auto*
hence *f0*: *f* ≠ *0* **by** *auto*
have *degm*: *degree-m* (*smult ?lc* (*prod-list hs*)) = *degree* (*smult ?lc* (*prod-list hs*))

by (*rule degree-m-eq*[*OF - m1*], *insert n bh mon-hs Mlc*, *auto simp*: *M-def*)
from *reconstruction*[*OF prime refl n sl-impl res - refl - refl - refl refl ufact sf cop - - - deg - bnd f0*] *bh*(*2*) *state*
show *?thesis* **by** *simp*
qed

lemma *zassenhaus-reconstruction-irreducible*$_d$:
assumes *res*: *zassenhaus-reconstruction hs p n f* = *fs*
shows *f* = *prod-list fs* ∧ (∀ *fi* ∈ *set fs*. *irreducible*$_d$ *fi*)
by (*rule zassenhaus-reconstruction-generic*[*OF my-subseqs.impl-correct res*[*unfolded zassenhaus-reconstruction-def Let-def*]])

corollary *zassenhaus-reconstruction*:
assumes *pr*: *primitive f*
assumes *res*: *zassenhaus-reconstruction hs p n f* = *fs*
shows *f* = *prod-list fs* ∧ (∀ *fi* ∈ *set fs*. *irreducible fi*)
using *zassenhaus-reconstruction-irreducible*$_d$[*OF res*] *pr*
irreducible-primitive-connect[*OF primitive-prod-list*]
by *auto*
end

end

theory *Code-Abort-Gcd*
imports
HOL−Computational-Algebra.Polynomial-Factorial
begin

Dummy code-setup for *Gcd* and *Lcm* in the presence of Container.

definition *dummy-Gcd* **where** *dummy-Gcd x* = *Gcd x*
definition *dummy-Lcm* **where** *dummy-Lcm x* = *Lcm x*
declare [[*code abort*: *dummy-Gcd*]]

lemma *dummy-Gcd-Lcm*: *Gcd x* = *dummy-Gcd x Lcm x* = *dummy-Lcm x*
unfolding *dummy-Gcd-def dummy-Lcm-def* **by** *auto*

lemmas *dummy-Gcd-Lcm-poly* [*code*] = *dummy-Gcd-Lcm*
[**where** *?'a* = *'a* :: {*factorial-ring-gcd*,*semiring-gcd-mult-normalize*} *poly*]
lemmas *dummy-Gcd-Lcm-int* [*code*] = *dummy-Gcd-Lcm* [**where** *?'a* = *int*]
lemmas *dummy-Gcd-Lcm-nat* [*code*] = *dummy-Gcd-Lcm* [**where** *?'a* = *nat*]

declare [[*code abort*: *Euclidean-Algorithm.Gcd Euclidean-Algorithm.Lcm*]]

end

11 The Polynomial Factorization Algorithm

11.1 Factoring Square-Free Integer Polynomials

We combine all previous results, i.e., Berlekamp's algorithm, Hensel-lifting, the reconstruction of Zassenhaus, Mignotte-bounds, etc., to eventually assemble the factorization algorithm for integer polynomials.

theory *Berlekamp-Zassenhaus*
imports
Berlekamp-Hensel
Polynomial-Factorization.Gauss-Lemma
Polynomial-Factorization.Dvd-Int-Poly
Reconstruction
Suitable-Prime
Degree-Bound
Code-Abort-Gcd
begin

context
begin
private partial-function (*tailrec*) *find-exponent-main* :: *int* ⇒ *int* ⇒ *nat* ⇒ *int* ⇒ *nat* **where**
[*code*]: *find-exponent-main p pm m bnd* = (*if pm* > *bnd then m*
else find-exponent-main p (*pm* ∗ *p*) (*Suc m*) *bnd*)

```
definition find-exponent :: int ⇒ int ⇒ nat where
  find-exponent p bnd = find-exponent-main p p 1 bnd

lemma find-exponent: assumes p: p > 1
  shows p ^ find-exponent p bnd > bnd find-exponent p bnd ≠ 0
proof −
  {
    fix m and n
    assume n = nat (1 + bnd − p^m) and m ≥ 1
    hence bnd < p ^ find-exponent-main p (p^m) m bnd ∧ find-exponent-main p
(p^m) m bnd ≥ 1
    proof (induct n arbitrary: m rule: less-induct)
      case (less n m)
      note simp = find-exponent-main.simps[of p p^m]
      show ?case
      proof (cases bnd < p ^ m)
        case True
        thus ?thesis using less unfolding simp by simp
      next
        case False
        hence id: find-exponent-main p (p ^ m) m bnd = find-exponent-main p (p
^ Suc m) (Suc m) bnd
          unfolding simp by (simp add: ac-simps)
        show ?thesis unfolding id
          by (rule less(1)[OF - refl], unfold less(2), insert False p, auto)
      qed
    qed
  }
  from this[OF refl, of 1]
  show p ^ find-exponent p bnd > bnd find-exponent p bnd ≠ 0
    unfolding find-exponent-def by auto
qed

end

definition berlekamp-zassenhaus-factorization :: int poly ⇒ int poly list where
  berlekamp-zassenhaus-factorization f = (let
    — find suitable prime
    p = suitable-prime-bz f;
    — compute finite field factorization
    (-, fs) = finite-field-factorization-int p f;
    — determine maximal degree that we can build by multiplying at most half of
the factors
    max-deg = degree-bound fs;
    — determine a number large enough to represent all coefficients of every
    — factor of lc ∗ f that has at most degree most max-deg
    bnd = 2 ∗ |lead-coeff f| ∗ factor-bound f max-deg;
    — determine k such that p^k > bnd
```

```
    k = find-exponent p bnd;
    — perform hensel lifting to lift factorization to mod p^k
    vs = hensel-lifting p k f fs
    — reconstruct integer factors
  in zassenhaus-reconstruction vs p k f)

theorem berlekamp-zassenhaus-factorization-irreducible_d:
  assumes res: berlekamp-zassenhaus-factorization f = fs
  and sf: square-free f
  and deg: degree f > 0
  shows f = prod-list fs ∧ (∀ fi ∈ set fs. irreducible_d fi)
proof −
  let ?lc = lead-coeff f
  define p where p ≡ suitable-prime-bz f
  obtain c gs where berl: finite-field-factorization-int p f = (c,gs) by force
  let ?degs = map degree gs
  note res = res[unfolded berlekamp-zassenhaus-factorization-def Let-def, folded
p-def,
    unfolded berl split, folded]
  from suitable-prime-bz[OF sf refl]
  have prime: prime p and cop: coprime ?lc p and sf: poly-mod.square-free-m p f
    unfolding p-def by auto
  from prime interpret poly-mod-prime p by unfold-locales
  define n where n = find-exponent p (2 * abs ?lc * factor-bound f (degree-bound
gs))
  note n = find-exponent[OF m1, of 2 * abs ?lc * factor-bound f (degree-bound
gs),
    folded n-def]
  note bh = berlekamp-and-hensel-separated[OF cop sf refl berl n(2)]
  have db: degree-bound (berlekamp-hensel p n f) = degree-bound gs unfolding bh
    degree-bound-def max-factor-degree-def by simp
  note res = res[folded n-def bh(1)]
  show ?thesis
    by (rule zassenhaus-reconstruction-irreducible_d[OF prime cop sf deg refl - res],
insert n db, auto)
qed

corollary berlekamp-zassenhaus-factorization-irreducible:
  assumes res: berlekamp-zassenhaus-factorization f = fs
    and sf: square-free f
    and pr: primitive f
    and deg: degree f > 0
  shows f = prod-list fs ∧ (∀ fi ∈ set fs. irreducible fi)
  using pr irreducible-primitive-connect[OF primitive-prod-list]
    berlekamp-zassenhaus-factorization-irreducible_d[OF res sf deg] by auto

end
```

11.2 A fast coprimality approximation

We adapt the integer polynomial gcd algorithm so that it first tests whether f and g are coprime modulo a few primes. If so, we are immediately done.

theory *Gcd-Finite-Field-Impl*
imports
Suitable-Prime
Code-Abort-Gcd
HOL−Library.Code-Target-Int
begin

definition *coprime-approx-main* :: *int* ⇒ *'i arith-ops-record* ⇒ *int poly* ⇒ *int poly* ⇒ *bool* **where**
coprime-approx-main p ff-ops f g = (*gcd-poly-i ff-ops* (*of-int-poly-i ff-ops* (*poly-mod.Mp p f*))
(*of-int-poly-i ff-ops* (*poly-mod.Mp p g*)) = *one-poly-i ff-ops*)

lemma (**in** *prime-field-gen*) *coprime-approx-main*:
shows *coprime-approx-main p ff-ops f g* ⟹ *coprime-m f g*
proof −
define *F* **where** *F*: (*F* :: *'a mod-ring poly*) = *of-int-poly* (*Mp f*)
define *G* **where** *G*: (*G* :: *'a mod-ring poly*) = *of-int-poly* (*Mp g*) **let** *?f'* = *of-int-poly-i ff-ops* (*Mp f*)
let *?g'* = *of-int-poly-i ff-ops* (*Mp g*)
define *f''* **where** *f''* ≡ *of-int-poly* (*Mp f*) :: *'a mod-ring poly*
define *g''* **where** *g''* ≡ *of-int-poly* (*Mp g*) :: *'a mod-ring poly*
have *rel-f*[*transfer-rule*]: *poly-rel ?f' f''*
by (*rule poly-rel-of-int-poly*[*OF refl*], *simp add*: *f''-def*)
have *rel-f*[*transfer-rule*]: *poly-rel ?g' g''*
by (*rule poly-rel-of-int-poly*[*OF refl*], *simp add*: *g''-def*)
have *id*: (*gcd-poly-i ff-ops* (*of-int-poly-i ff-ops* (*Mp f*)) (*of-int-poly-i ff-ops* (*Mp g*)) = *one-poly-i ff-ops*)
= *coprime f'' g''* (**is** *?P* ⟷ *?Q*)
proof −
have *?P* ⟷ *gcd f'' g'' = 1*
unfolding *separable-i-def* **by** *transfer-prover*
also have ... ⟷ *?Q*
by (*simp add*: *coprime-iff-gcd-eq-1*)
finally show *?thesis* **.**
qed
have *fF*: *MP-Rel* (*Mp f*) *F* **unfolding** *F MP-Rel-def*
by (*simp add*: *Mp-f-representative*)
have *gG*: *MP-Rel* (*Mp g*) *G* **unfolding** *G MP-Rel-def*
by (*simp add*: *Mp-f-representative*)
have *coprime f'' g''* = *coprime F G* **unfolding** *f''-def F g''-def G* **by** *simp*
also have ... = *coprime-m* (*Mp f*) (*Mp g*)
using *coprime-MP-Rel*[*unfolded rel-fun-def*, *rule-format*, *OF fF gG*] **by** *simp*
also have ... = *coprime-m f g* **unfolding** *coprime-m-def dvdm-def* **by** *simp*
finally have *id2*: *coprime f'' g''* = *coprime-m f g* **.**

show *coprime-approx-main p ff-ops f g* $\Longrightarrow$ *coprime-m f g* **unfolding** *coprime-approx-main-def id id2* **by** *auto*
qed

context *poly-mod-prime* **begin**

lemmas *coprime-approx-main-uint32 = prime-field-gen.coprime-approx-main*[*OF*

prime-field.prime-field-finite-field-ops32, *unfolded prime-field-def mod-ring-locale-def poly-mod-type-simps*, *internalize-sort 'a :: prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*, *OF non-empty*]

lemmas *coprime-approx-main-uint64 = prime-field-gen.coprime-approx-main*[*OF*

prime-field.prime-field-finite-field-ops64, *unfolded prime-field-def mod-ring-locale-def poly-mod-type-simps*, *internalize-sort 'a :: prime-card*, *OF type-to-set*, *unfolded remove-duplicate-premise*, *cancel-type-definition*, *OF non-empty*]

end

lemma *coprime-mod-imp-coprime*: **assumes**
p: *prime p* **and**
cop-m: *poly-mod.coprime-m p f g* **and**
cop: *coprime* (*lead-coeff f*) *p* $\vee$ *coprime* (*lead-coeff g*) *p* **and**
cnt: *content f = 1* $\vee$ *content g = 1*
shows *coprime f g*
proof −
interpret *poly-mod-prime p* **by** (*standard*, *rule p*)
from *cop-m*[*unfolded coprime-m-def*] **have** *cop-m*: $\bigwedge$ *h. h dvdm f* $\Longrightarrow$ *h dvdm g* $\Longrightarrow$ *h dvdm 1* **by** *auto*
show *?thesis*
proof (*rule coprimeI*)
fix *h*
assume *dvd*: *h dvd f h dvd g*
hence *h dvdm f h dvdm g* **unfolding** *dvdm-def dvd-def* **by** *auto*
from *cop-m*[*OF this*] **obtain** *k* **where** *unit*: *Mp* (*h* ∗ *Mp k*) = *1* **unfolding** *dvdm-def* **by** *auto*
from *content-dvd-contentI*[*OF dvd*(*1*)] *content-dvd-contentI*[*OF dvd*(*2*)] *cnt*
have *cnt*: *content h = 1* **by** *auto*
let *?k = Mp k*
from *unit* **have** *h0*: *h* ≠ *0* **by** *auto*
from *unit* **have** *k0*: *?k* ≠ *0* **by** *fastforce*
from *p* **have** *p0*: *p* ≠ *0* **by** *auto*
from *dvd* **have** *lead-coeff h dvd lead-coeff f lead-coeff h dvd lead-coeff g*
by (*metis dvd-def lead-coeff-mult*)+
with *cop* **have** *coph*: *coprime* (*lead-coeff h*) *p*
by (*meson dvd-trans not-coprime-iff-common-factor*)
let *?k = Mp k*
from *arg-cong*[*OF unit*, *of degree*] **have** *degm0*: *degree-m* (*h* ∗ *?k*) = *0* **by** *simp*

```
    have lead-coeff ?k ∈ {0 ..< p} unfolding Mp-coeff M-def using m1 by simp
    with k0 have lk: lead-coeff ?k ≥ 1 lead-coeff ?k < p
      by (auto simp add: int-one-le-iff-zero-less order.not-eq-order-implies-strict)
   have id: lead-coeff (h * ?k) = lead-coeff h * lead-coeff ?k unfolding lead-coeff-mult
..
    from coph prime lk have coprime (lead-coeff h * lead-coeff ?k) p
      by (simp add: ac-simps prime-imp-coprime zdvd-not-zless)
    with id have cop-prod: coprime (lead-coeff (h * ?k)) p by simp
    from h0 k0 have lc0: lead-coeff (h * ?k) ≠ 0
      unfolding lead-coeff-mult by auto
    from p have lcp: lead-coeff (h * ?k) mod p ≠ 0
      using M-1 M-def cop-prod by auto
    have deg-eq: degree-m (h * ?k) = degree (h * Mp k)
      by (rule degree-m-eq[OF - m1], insert lcp)
    from this[unfolded degm0] have degree (h * Mp k) = 0 by simp
    with degree-mult-eq[OF h0 k0] have deg0: degree h = 0 by auto
    from degree0-coeffs[OF this] obtain h0 where h: h = [:h0:] by auto
    have content h = abs h0 unfolding content-def h by (cases h0 = 0, auto)
    hence abs h0 = 1 using cnt by auto
    hence h0 ∈ {−1,1} by auto
    hence h = 1 ∨ h = −1 unfolding h by (auto)
    thus is-unit h by auto
  qed
qed
```

We did not try to optimize the set of chosen primes. They have just been picked randomly from a list of primes.

```
definition gcd-primes32 :: int list where
  gcd-primes32 = [383, 1409, 19213, 22003, 41999]

lemma gcd-primes32: p ∈ set gcd-primes32 ⟹ prime p ∧ p ≤ 65535
proof −
  have list-all (λ p. prime p ∧ p ≤ 65535) gcd-primes32 by eval
  thus p ∈ set gcd-primes32 ⟹ prime p ∧ p ≤ 65535 by (auto simp: list-all-iff)
qed

definition gcd-primes64 :: int list where
  gcd-primes64 = [383, 21984191, 50329901, 80329901, 219849193]

lemma gcd-primes64: p ∈ set gcd-primes64 ⟹ prime p ∧ p ≤ 4294967295
proof −
  have list-all (λ p. prime p ∧ p ≤ 4294967295) gcd-primes64 by eval
   thus p ∈ set gcd-primes64 ⟹ prime p ∧ p ≤ 4294967295 by (auto simp:
list-all-iff)
qed

definition coprime-heuristic :: int poly ⇒ int poly ⇒ bool where
  coprime-heuristic f g = (let lcf = lead-coeff f; lcg = lead-coeff g in
   find (λ p. (coprime lcf p ∨ coprime lcg p) ∧ coprime-approx-main p (finite-field-ops64
```

```
(uint64-of-int p)) f g)
    gcd-primes64 ≠ None)

lemma coprime-heuristic: assumes coprime-heuristic f g
  and content f = 1 ∨ content g = 1
  shows coprime f g
proof (cases find (λp. (coprime (lead-coeff f) p ∨ coprime (lead-coeff g) p) ∧
           coprime-approx-main p (finite-field-ops64 (uint64-of-int p)) f g)
   gcd-primes64)
  case (Some p)
   from find-Some-D[OF Some] gcd-primes64 have p: prime p and small: p ≤
4294967295
    and cop: coprime (lead-coeff f) p ∨ coprime (lead-coeff g) p
    and copp: coprime-approx-main p (finite-field-ops64 (uint64-of-int p)) f g by
auto
  interpret poly-mod-prime p using p by unfold-locales
  from coprime-approx-main-uint64[OF small copp] have poly-mod.coprime-m p f
g by auto
  from coprime-mod-imp-coprime[OF p this cop assms(2)] show coprime f g .
qed (insert assms(1)[unfolded coprime-heuristic-def], auto simp: Let-def)

definition gcd-int-poly :: int poly ⇒ int poly ⇒ int poly where
  gcd-int-poly f g =
    (if f = 0 then normalize g
     else if g = 0 then normalize f
          else let
            cf = Polynomial.content f;
            cg = Polynomial.content g;
            ct = gcd cf cg;
            ff = map-poly (λ x. x div cf) f;
            gg = map-poly (λ x. x div cg) g
           in if coprime-heuristic ff gg then [:ct:] else smult ct (gcd-poly-code-aux ff
gg))

lemma gcd-int-poly-code[code-unfold]: gcd = gcd-int-poly
proof (intro ext)
  fix f g :: int poly
  let ?ff = primitive-part f
  let ?gg = primitive-part g
  note d = gcd-int-poly-def gcd-poly-code gcd-poly-code-def
  show gcd f g = gcd-int-poly f g
  proof (cases f = 0 ∨ g = 0 ∨ ¬ coprime-heuristic ?ff ?gg)
    case True
    thus ?thesis unfolding d by (auto simp: Let-def primitive-part-def)
  next
    case False
    hence cop: coprime-heuristic ?ff ?gg by simp
    from False have f ≠ 0 by auto
    from content-primitive-part[OF this] coprime-heuristic[OF cop]
```

```
    have id: gcd ?ff ?gg = 1 by auto
    show ?thesis unfolding gcd-poly-decompose[of f g] unfolding gcd-int-poly-def
Let-def id
      using False by (auto simp: primitive-part-def)
  qed
qed

end

theory Square-Free-Factorization-Int
imports
  Square-Free-Int-To-Square-Free-GFp
  Suitable-Prime
  Code-Abort-Gcd
  Gcd-Finite-Field-Impl
begin

definition yun-wrel :: int poly ⇒ rat ⇒ rat poly ⇒ bool where
  yun-wrel F c f = (map-poly rat-of-int F = smult c f)

definition yun-rel :: int poly ⇒ rat ⇒ rat poly ⇒ bool where
  yun-rel F c f = (yun-wrel F c f
    ∧ content F = 1 ∧ lead-coeff F > 0 ∧ monic f)

definition yun-erel :: int poly ⇒ rat poly ⇒ bool where
  yun-erel F f = (∃ c. yun-rel F c f)

lemma yun-wrelD: assumes yun-wrel F c f
  shows map-poly rat-of-int F = smult c f
  using assms unfolding yun-wrel-def by auto

lemma yun-relD: assumes yun-rel F c f
  shows yun-wrel F c f map-poly rat-of-int F = smult c f
    degree F = degree f F ≠ 0 lead-coeff F > 0 monic f
    f = 1 ⟷ F = 1 content F = 1
proof −
  note ∗ = assms[unfolded yun-rel-def yun-wrel-def, simplified]
  then have degree (map-poly rat-of-int F) = degree f by auto
  then show deg: degree F = degree f by simp
  show F ≠ 0 lead-coeff F > 0 monic f content F = 1
    map-poly rat-of-int F = smult c f
    yun-wrel F c f using ∗ by (auto simp: yun-wrel-def)
  {
    assume f = 1
    with deg have degree F = 0 by auto
   from degree0-coeffs[OF this] obtain c where F: F = [:c:] and c: c = lead-coeff
F by auto
    from c ∗ have c0: c > 0 by auto
    hence cF: content F = c unfolding F content-def by auto
```

```
    with ∗ have c = 1 by auto
    with F have F = 1 by simp
  }
  moreover
  {
    assume F = 1
    with deg have degree f = 0 by auto
    with ‹monic f› have f = 1
      using monic-degree-0 by blast
  }
  ultimately show (f = 1) ⟷ (F = 1) by auto
qed

lemma yun-erel-1-eq: assumes yun-erel F f
  shows (F = 1) ⟷ (f = 1)
proof −
  from assms[unfolded yun-erel-def] obtain c where yun-rel F c f by auto
  from yun-relD[OF this] show ?thesis by simp
qed

lemma yun-rel-1[simp]: yun-rel 1 1 1
  by (auto simp: yun-rel-def yun-wrel-def content-def)

lemma yun-erel-1[simp]: yun-erel 1 1 unfolding yun-erel-def using yun-rel-1 by
blast

lemma yun-rel-mult: yun-rel F c f ⟹ yun-rel G d g ⟹ yun-rel (F ∗ G) (c ∗ d)
(f ∗ g)
  unfolding yun-rel-def yun-wrel-def content-mult lead-coeff-mult
  by (auto simp: monic-mult hom-distribs)

lemma yun-erel-mult: yun-erel F f ⟹ yun-erel G g ⟹ yun-erel (F ∗ G) (f ∗ g)

  unfolding yun-erel-def using yun-rel-mult[of F - f G - g] by blast

lemma yun-rel-pow: assumes yun-rel F c f
  shows yun-rel (F^n) (c^n) (f^n)
  by (induct n, insert assms yun-rel-mult, auto)

lemma yun-erel-pow: yun-erel F f ⟹ yun-erel (F^n) (f^n)
  using yun-rel-pow unfolding yun-erel-def by blast

lemma yun-wrel-pderiv: assumes yun-wrel F c f
  shows yun-wrel (pderiv F) c (pderiv f)
 by (unfold yun-wrel-def, simp add: yun-wrelD[OF assms] pderiv-smult hom-distribs)

lemma yun-wrel-minus: assumes yun-wrel F c f yun-wrel G c g
  shows yun-wrel (F − G) c (f − g)
 using assms unfolding yun-wrel-def by (auto simp: smult-diff-right hom-distribs)
```

lemma *yun-wrel-div*: **assumes** *f*: *yun-wrel F c f* **and** *g*: *yun-wrel G d g*
and *dvd*: *G dvd F g dvd f*
and *G0*: *G* ≠ *0*
shows *yun-wrel* (*F div G*) (*c* / *d*) (*f div g*)
proof −
let *?r* = *rat-of-int*
let *?rp* = *map-poly ?r*
from *dvd* **obtain** *H h* **where** *fgh*: *F* = *G* ∗ *H f* = *g* ∗ *h* **unfolding** *dvd-def* **by** *auto*
from *G0 yun-wrelD*[*OF g*] **have** *g0*: *g* ≠ *0* **and** *d0*: *d* ≠ *0* **by** *auto*
from *arg-cong*[*OF fgh*(*1*), *of* λ *x. x div G*] **have** *H*: *H* = *F div G* **using** *G0* **by** *simp*
from *arg-cong*[*OF fgh*(*1*), *of ?rp*] **have** *?rp F* = *?rp G* ∗ *?rp H* **by** (*auto simp*: *hom-distribs*)
from *arg-cong*[*OF this*, *of* λ *x. x div ?rp G*] *G0* **have** *id*: *?rp H* = *?rp F div ?rp G* **by** *auto*
have *?rp* (*F div G*) = *?rp F div ?rp G* **unfolding** *H*[*symmetric*] *id* **by** *simp*
also have ... = *smult c f div smult d g* **using** *f g* **unfolding** *yun-wrel-def* **by** *auto*
also have ... = *smult* (*c* / *d*) (*f div g*) **unfolding** *div-smult-right div-smult-left*
by (*simp add*: *field-simps*)
finally show *?thesis* **unfolding** *yun-wrel-def* **by** *simp*
qed

lemma *yun-rel-div*: **assumes** *f*: *yun-rel F c f* **and** *g*: *yun-rel G d g*
and *dvd*: *G dvd F g dvd f*
shows *yun-rel* (*F div G*) (*c* / *d*) (*f div g*)
proof −
note *ff* = *yun-relD*[*OF f*]
note *gg* = *yun-relD*[*OF g*]
show *?thesis* **unfolding** *yun-rel-def*
proof (*intro conjI*)
from *yun-wrel-div*[*OF ff*(*1*) *gg*(*1*) *dvd gg*(*4*)]
show *yun-wrel* (*F div G*) (*c* / *d*) (*f div g*) **by** *auto*
from *dvd* **have** *fg*: *f* = *g* ∗ (*f div g*) **by** *auto*
from *arg-cong*[*OF fg*, *of monic*] *ff*(*6*) *gg*(*6*)
show *monic* (*f div g*) **using** *monic-factor* **by** *blast*
from *dvd* **have** *FG*: *F* = *G* ∗ (*F div G*) **by** *auto*
from *arg-cong*[*OF FG*, *of content*, *unfolded content-mult*] *ff*(*8*) *gg*(*8*)
show *content* (*F div G*) = *1* **by** *simp*
from *arg-cong*[*OF FG*, *of lead-coeff*, *unfolded lead-coeff-mult*] *ff*(*5*) *gg*(*5*)
show *lead-coeff* (*F div G*) > *0* **by** (*simp add*: *zero-less-mult-iff*)
qed
qed

lemma *yun-wrel-gcd*: **assumes** *yun-wrel F c′ f yun-wrel G c g* **and** *c*: *c′* ≠ *0 c* ≠

0
and *d*: *d* = *rat-of-int* (*lead-coeff* (*gcd F G*)) *d* ≠ *0*
shows *yun-wrel* (*gcd F G*) *d* (*gcd f g*)
proof −
let *?r* = *rat-of-int*
let *?rp* = *map-poly ?r*
have *smult d* (*gcd f g*) = *smult d* (*gcd* (*smult c' f*) (*smult c g*))
by (*simp add*: *c gcd-smult-left gcd-smult-right*)
also have ... = *smult d* (*gcd* (*?rp F*) (*?rp G*)) **using** *assms*(*1*−*2*)[*unfolded yun-wrel-def*] **by** *simp*
also have ... = *smult* (*d* ∗ *inverse d*) (*?rp* (*gcd F G*))
unfolding *gcd-rat-to-gcd-int d* **by** *simp*
also have *d* ∗ *inverse d* = *1* **using** *d* **by** *auto*
finally show *?thesis* **unfolding** *yun-wrel-def* **by** *simp*
qed

lemma *yun-rel-gcd*: **assumes** *f*: *yun-rel F c f* **and** *g*: *yun-wrel G c' g* **and** *c'*: *c'* ≠ *0*
and *d*: *d* = *rat-of-int* (*lead-coeff* (*gcd F G*))
shows *yun-rel* (*gcd F G*) *d* (*gcd f g*)
unfolding *yun-rel-def*
proof (*intro conjI*)
note *ff* = *yun-relD*[*OF f*]
from *ff* **have** *c0*: *c* ≠ *0* **by** *auto*
from *ff d* **have** *d0*: *d* ≠ *0* **by** *auto*
from *yun-wrel-gcd*[*OF ff*(*1*) *g c0 c' d d0*]
show *yun-wrel* (*gcd F G*) *d* (*gcd f g*) **by** *auto*
from *ff* **have** *gcd f g* ≠ *0* **by** *auto*
thus *monic* (*gcd f g*) **by** (*simp add*: *poly-gcd-monic*)
obtain *H* **where** *H*: *gcd F G* = *H* **by** *auto*
obtain *lc* **where** *lc*: *coeff H* (*degree H*) = *lc* **by** *auto*
from *ff* **have** *gcd F G* ≠ *0* **by** *auto*
hence *H* ≠ *0 lc* ≠ *0* **unfolding** *H*[*symmetric*] *lc*[*symmetric*] **by** *auto*
thus *0* < *lead-coeff* (*gcd F G*) **unfolding**
arg-cong[*OF normalize-gcd*[*of F G*], *of lead-coeff*, *symmetric*]
unfolding *normalize-poly-eq-map-poly H*
by (*auto*, *subst Polynomial.coeff-map-poly*, *auto*,
subst Polynomial.degree-map-poly, *auto simp*: *sgn-if*)
have *H dvd F* **unfolding** *H*[*symmetric*] **by** *auto*
then obtain *K* **where** *F*: *F* = *H* ∗ *K* **unfolding** *dvd-def* **by** *auto*
from *arg-cong*[*OF this*, *of content*, *unfolded content-mult ff*(*8*)]
content-ge-0-int[*of H*] **have** *content H* = *1*
by (*auto simp add*: *zmult-eq-1-iff*)
thus *content* (*gcd F G*) = *1* **unfolding** *H* **.**
qed

lemma *yun-factorization-main-int*: **assumes** *f*: *f* = *p div gcd p* (*pderiv p*)

```
    and g = pderiv p div gcd p (pderiv p) monic p
    and yun-gcd.yun-factorization-main gcd f g i hs = res
    and yun-gcd.yun-factorization-main gcd F G i Hs = Res
    and yun-rel F c f yun-wrel G c g list-all2 (rel-prod yun-erel (=)) Hs hs
  shows list-all2 (rel-prod yun-erel (=)) Res res
proof −
  let ?P = λ f g. ∀ i hs res F G Hs Res c.
    yun-gcd.yun-factorization-main gcd f g i hs = res
    ⟶ yun-gcd.yun-factorization-main gcd F G i Hs = Res
    ⟶ yun-rel F c f ⟶ yun-wrel G c g ⟶ list-all2 (rel-prod yun-erel (=)) Hs hs
    ⟶ list-all2 (rel-prod yun-erel (=)) Res res
  note simps = yun-gcd.yun-factorization-main.simps
  note rel = yun-relD
  let ?rel = λ F f. map-poly rat-of-int F = smult (rat-of-int (lead-coeff F)) f
  show ?thesis
  proof (induct rule: yun-factorization-induct[of ?P, rule-format, OF - - assms])
    case (1 f g i hs res F G Hs Res c)
    from rel[OF 1(4)] 1(1) have f = 1 F = 1 by auto
    from 1(2−3)[unfolded simps[of - 1] this] have res = hs Res = Hs by auto
    with 1(6) show ?case by simp
  next
    case (2 f g i hs res F G Hs Res c)
    define d where d = g − pderiv f
    define a where a = gcd f d
    define D where D = G − pderiv F
    define A where A = gcd F D
    note f = 2(5)
    note g = 2(6)
    note hs = 2(7)
    note f1 = 2(1)
    from f1 rel[OF f] have ∗: (f = 1) = False (F = 1) = False and c: c ≠ 0 by
auto
    note res = 2(3)[unfolded simps[of - f] ∗ if-False Let-def, folded d-def a-def]
    note Res = 2(4)[unfolded simps[of - F] ∗ if-False Let-def, folded D-def A-def]
    note IH = 2(2)[folded d-def a-def, OF res Res]
    obtain c′ where c′: c′ = rat-of-int (lead-coeff (gcd F D)) by auto
    show ?case
    proof (rule IH)
      from yun-wrel-minus[OF g yun-wrel-pderiv[OF rel(1)[OF f]]]
      have d: yun-wrel D c d unfolding D-def d-def .
      have a: yun-rel A c′ a unfolding A-def a-def
        by (rule yun-rel-gcd[OF f d c c′])
      hence yun-erel A a unfolding yun-erel-def by auto
      thus list-all2 (rel-prod yun-erel (=)) ((A, Suc i) # Hs) ((a, Suc i) # hs)
        using hs by auto
      have A0: A ≠ 0 by (rule rel(4)[OF a])
      have A dvd D a dvd d unfolding A-def a-def by auto
      from yun-wrel-div[OF d rel(1)[OF a] this A0]
      show yun-wrel (D div A) (c / c′) (d div a) .
```

```
    have A dvd F a dvd f unfolding A-def a-def by auto
    from yun-rel-div[OF f a this]
    show yun-rel (F div A) (c / c') (f div a) .
   qed
  qed
qed

lemma yun-monic-factorization-int-yun-rel: assumes
    res: yun-gcd.yun-monic-factorization gcd f = res
    and Res: yun-gcd.yun-monic-factorization gcd F = Res
    and f: yun-rel F c f
  shows list-all2 (rel-prod yun-erel (=)) Res res
proof −
  note ff = yun-relD[OF f]
  let ?g = gcd f (pderiv f)
  let ?yf = yun-gcd.yun-factorization-main gcd (f div ?g) (pderiv f div ?g) 0 []
  let ?G = gcd F (pderiv F)
  let ?yF = yun-gcd.yun-factorization-main gcd (F div ?G) (pderiv F div ?G) 0 []

  obtain r R where r: ?yf = r and R: ?yF = R by blast
  from res[unfolded yun-gcd.yun-monic-factorization-def Let-def r]
  have res: res = [(a, i)←r . a ≠ 1] by simp
  from Res[unfolded yun-gcd.yun-monic-factorization-def Let-def R]
  have Res: Res = [(A, i)←R . A ≠ 1] by simp
  from yun-wrel-pderiv[OF ff(1)] have f': yun-wrel (pderiv F) c (pderiv f) .
  from ff have c: c ≠ 0 by auto
  from yun-rel-gcd[OF f f' c refl] obtain d where g: yun-rel ?G d ?g ..
  from yun-rel-div[OF f g] have 1: yun-rel (F div ?G) (c / d) (f div ?g) by auto
  from yun-wrel-div[OF f' yun-relD(1)[OF g] - - yun-relD(4)[OF g]]
  have 2: yun-wrel (pderiv F div ?G) (c / d) (pderiv f div ?g) by auto
  from yun-factorization-main-int[OF refl refl ff(6) r R 1 2]
  have list-all2 (rel-prod yun-erel (=)) R r by simp
  thus ?thesis unfolding res Res
    by (induct R r rule: list-all2-induct, auto dest: yun-erel-1-eq)
qed

lemma yun-rel-same-right: assumes yun-rel f c G yun-rel g d G
  shows f = g
proof −
  note f = yun-relD[OF assms(1)]
  note g = yun-relD[OF assms(2)]
  let ?r = rat-of-int
  let ?rp = map-poly ?r
  from g have d: d ≠ 0 by auto
  obtain a b where quot: quotient-of (c / d) = (a,b) by force
  from quotient-of-nonzero[of c/d, unfolded quot] have b: b ≠ 0 by simp
  note f(2)
   also have smult c G = smult (c / d) (smult d G) using d by (auto simp:
field-simps)
```

```
  also have smult d G = ?rp g using g(2) by simp
  also have cd: c / d = (?r a / ?r b) using quotient-of-div[OF quot] .
  finally have fg: ?rp f = smult (?r a / ?r b) (?rp g) by simp
  from f have c ≠ 0 by auto
  with cd d have a: a ≠ 0 by auto
  from arg-cong[OF fg, of λ x. smult (?r b) x]
  have smult (?r b) (?rp f) = smult (?r a) (?rp g) using b by auto
  hence ?rp (smult b f) = ?rp (smult a g) by (auto simp: hom-distribs)
  then have fg: [:b:] * f = [:a:] * g by auto
  from arg-cong[OF this, of content, unfolded content-mult f(8) g(8)]
  have content [: b :] = content [: a :] by simp
  hence abs: abs a = abs b unfolding content-def using b a by auto
 from arg-cong[OF fg, of λ x. lead-coeff x > 0, unfolded lead-coeff-mult] f(5) g(5)
a b
  have (a > 0) = (b > 0) by (simp add: zero-less-mult-iff)
  with a b abs have a = b by auto
  with arg-cong[OF fg, of λ x. x div [:b:]] b show ?thesis
    by (metis nonzero-mult-div-cancel-left pCons-eq-0-iff)
qed



definition square-free-factorization-int-main :: int poly ⇒ (int poly × nat) list
where
  square-free-factorization-int-main f = (case square-free-heuristic f of None ⇒
    yun-gcd.yun-monic-factorization gcd f | Some p ⇒ [(f,1)])

lemma square-free-factorization-int-main: assumes res: square-free-factorization-int-main
f = fs
  and ct: content f = 1 and lc: lead-coeff f > 0
  and deg: degree f ≠ 0
shows square-free-factorization f (1,fs) ∧ (∀ fi i. (fi, i) ∈ set fs ⟶ content fi =
1 ∧ lead-coeff fi > 0) ∧
  distinct (map snd fs)
proof (cases square-free-heuristic f)
  case None
  from lc have f0: f ≠ 0 by auto
  from res None have fs: yun-gcd.yun-monic-factorization gcd f = fs
    unfolding square-free-factorization-int-main-def by auto
  let ?r = rat-of-int
  let ?rp = map-poly ?r
  define G where G = smult (inverse (lead-coeff (?rp f))) (?rp f)
  have ?rp f ≠ 0 using f0 by auto
  hence mon: monic G unfolding G-def coeff-smult by simp
  obtain Fs where Fs: yun-gcd.yun-monic-factorization gcd G = Fs by blast
  from lc have lg: lead-coeff (?rp f) ≠ 0 by auto
  let ?c = lead-coeff (?rp f)
  define c where c = ?c
  have rp: ?rp f = smult c G unfolding G-def c-def by (simp add: field-simps)
```

have *in-rel*: *yun-rel f c G* **unfolding** *yun-rel-def yun-wrel-def*
using *rp mon lc ct* **by** *auto*
from *yun-monic-factorization-int-yun-rel*[*OF Fs fs in-rel*]
have *out-rel*: *list-all2* (*rel-prod yun-erel* (=)) *fs Fs* **by** *auto*
from *yun-monic-factorization*[*OF Fs mon*]
have *square-free-factorization G* (*1*, *Fs*) **and** *dist*: *distinct* (*map snd Fs*) **by** *auto*
note *sff* = *square-free-factorizationD*[*OF this*(*1*)]
from *out-rel* **have** *map snd fs* = *map snd Fs* **by** (*induct fs Fs rule*: *list-all2-induct*, *auto*)
with *dist* **have** *dist'*: *distinct* (*map snd fs*) **by** *auto*
have *main*: *square-free-factorization f* (*1*, *fs*) ∧ (∀ *fi i*. (*fi*, *i*) ∈ *set fs* ⟶ *content fi* = *1* ∧ *lead-coeff fi* > *0*)
unfolding *square-free-factorization-def split*
proof (*intro conjI allI impI*)
from *ct* **have** *f* ≠ *0* **by** *auto*
thus *f* = *0* ⟹ *1* = *0 f* = *0* ⟹ *fs* = [] **by** *auto*
from *dist'* **show** *distinct fs* **by** (*simp add*: *distinct-map*)
{
fix *a i*
assume *a*: (*a*,*i*) ∈ *set fs*
with *out-rel* **obtain** *bj* **where** *bj* ∈ *set Fs* **and** *rel-prod yun-erel* (=) (*a*,*i*) *bj*
unfolding *list-all2-conv-all-nth set-conv-nth* **by** *fastforce*
then obtain *b* **where** *b*: (*b*,*i*) ∈ *set Fs* **and** *ab*: *yun-erel a b* **by** (*cases bj*, *auto simp*: *rel-prod.simps*)
from *sff*(*2*)[*OF b*] **have** *b'*: *square-free b degree b* ≠ *0* **and** *i*: *i* > *0* **by** *auto*
from *ab* **obtain** *c* **where** *rel*: *yun-rel a c b* **unfolding** *yun-erel-def* **by** *auto*
note *aa* = *yun-relD*[*OF this*]
from *aa* **have** *c0*: *c* ≠ *0* **by** *auto*
from *b' aa*(*3*) **show** *degree a* > *0* **by** *simp*
from *square-free-smult*[*OF c0 b'*(*1*), *folded aa*(*2*)]
show *square-free a* **unfolding** *square-free-def* **by** (*force simp*: *dvd-def hom-distribs*)
show *i* > *0* **by** *fact*
show *cnt*: *content a* = *1* **and** *lc*: *lead-coeff a* > *0* **using** *aa* **by** *auto*
fix *A I*
assume *A*: (*A*,*I*) ∈ *set fs* **and** *diff*: (*a*,*i*) ≠ (*A*,*I*)
from *a*[*unfolded set-conv-nth*] **obtain** *k* **where** *k*: *fs* ! *k* = (*a*,*i*) *k* < *length fs* **by** *auto*
from *A*[*unfolded set-conv-nth*] **obtain** *K* **where** *K*: *fs* ! *K* = (*A*,*I*) *K* < *length fs* **by** *auto*
from *diff k K* **have** *kK*: *k* ≠ *K* **by** *auto*
from *dist'*[*unfolded distinct-conv-nth length-map*, *rule-format*, *OF k*(*2*) *K*(*2*) *kK*]
have *iI*: *i* ≠ *I* **using** *k K* **by** *simp*
from *A out-rel* **obtain** *Bj* **where** *Bj* ∈ *set Fs* **and** *rel-prod yun-erel* (=) (*A*,*I*) *Bj*
unfolding *list-all2-conv-all-nth set-conv-nth* **by** *fastforce*
then obtain *B* **where** *B*: (*B*,*I*) ∈ *set Fs* **and** *AB*: *yun-erel A B* **by** (*cases Bj*, *auto simp*: *rel-prod.simps*)
then obtain *C* **where** *Rel*: *yun-rel A C B* **unfolding** *yun-erel-def* **by** *auto*

```
      note AA = yun-relD[OF this]
      from iI have (b,i) ≠ (B,I) by auto
      from sff(3)[OF b B this] have cop: coprime b B by simp
      from AA have C: C ≠ 0 by auto
      from yun-rel-gcd[OF rel AA(1) C refl] obtain c where yun-rel (gcd a A) c
(gcd b B) by auto
      note rel = yun-relD[OF this]
      from rel(2) cop have ?rp (gcd a A) = [: c :] by simp
      from arg-cong[OF this, of degree] have degree (gcd a A) = 0 by simp
      from degree0-coeffs[OF this] obtain c where gcd: gcd a A = [: c :] by auto
      from rel(8) rel(5) show Rings.coprime a A
        by (auto intro!: gcd-eq-1-imp-coprime simp add: gcd)
    }
    let ?prod = λ fs. (∏(a, i)∈set fs. a ^ i)
    let ?pr = λ fs. (∏(a, i)←fs. a ^ i)
    define pr where pr = ?prod fs
   from ‹distinct fs› have pfs: ?prod fs = ?pr fs by (rule prod.distinct-set-conv-list)
   from ‹distinct Fs› have pFs: ?prod Fs = ?pr Fs by (rule prod.distinct-set-conv-list)
    from out-rel have yun-erel (?prod fs) (?prod Fs) unfolding pfs pFs
    proof (induct fs Fs rule: list-all2-induct)
      case (Cons ai fs Ai Fs)
      obtain a i where ai: ai = (a,i) by force
      from Cons(1) ai obtain A where Ai: Ai = (A,i)
        and rel: yun-erel a A by (cases Ai, auto simp: rel-prod.simps)
     show ?case unfolding ai Ai using yun-erel-mult[OF yun-erel-pow[OF rel, of
i] Cons(3)]
        by auto
    qed simp
    also have ?prod Fs = G using sff(1) by simp
    finally obtain d where rel: yun-rel pr d G unfolding yun-erel-def pr-def by
auto
    with in-rel have f = pr by (rule yun-rel-same-right)
    thus f = smult 1 (?prod fs) unfolding pr-def by simp
  qed
  from main dist' show ?thesis by auto
next
  case (Some p)
   from res[unfolded square-free-factorization-int-main-def Some] have fs: fs =
[(f,1)] by auto
  from lc have f0: f ≠ 0 by auto
   from square-free-heuristic[OF Some] poly-mod-prime.separable-impl(1)[of p f]
square-free-mod-imp-square-free[of p f] deg
  show ?thesis unfolding fs
    by (auto simp: ct lc square-free-factorization-def f0 poly-mod-prime-def)
qed

definition square-free-factorization-int' :: int poly ⇒ int × (int poly × nat)list
where
  square-free-factorization-int' f = (if degree f = 0
```

```
    then (lead-coeff f,[]) else (let — content factorization
      c = content f;
      d = (sgn (lead-coeff f) * c);
      g = sdiv-poly f d
      — and square-free factorization
    in (d, square-free-factorization-int-main g)))


lemma square-free-factorization-int': assumes res: square-free-factorization-int' f
= (d, fs)
  shows square-free-factorization f (d,fs)
    (fi, i) ∈ set fs ⟹ content fi = 1 ∧ lead-coeff fi > 0
    distinct (map snd fs)
proof −
  note res = res[unfolded square-free-factorization-int'-def Let-def]
  have square-free-factorization f (d,fs)
    ∧ ((fi, i) ∈ set fs ⟶ content fi = 1 ∧ lead-coeff fi > 0)
    ∧ distinct (map snd fs)
  proof (cases degree f = 0)
    case True
    from degree0-coeffs[OF True] obtain c where f: f = [: c :] by auto
    thus ?thesis using res by (simp add: square-free-factorization-def)
  next
    case False
    let ?s = sgn (lead-coeff f)
    have s: ?s ∈ {−1,1} using False unfolding sgn-if by auto
    define g where g = smult ?s f
    let ?d = ?s * content f
    have content g = content ([:?s:] * f) unfolding g-def by simp
    also have ... = content [:?s:] * content f unfolding content-mult by simp
    also have content [:?s:] = 1 using s by (auto simp: content-def)
    finally have cg: content g = content f by simp
    from False res
    have d: d = ?d and fs: fs = square-free-factorization-int-main (sdiv-poly f ?d)
by auto
    let ?g = primitive-part g
    define ng where ng = primitive-part g
    note fs
     also have sdiv-poly f ?d = sdiv-poly g (content g) unfolding cg unfolding
g-def
        by (rule poly-eqI, unfold coeff-sdiv-poly coeff-smult, insert s, auto simp:
div-minus-right)
    finally have fs: square-free-factorization-int-main ng = fs
      unfolding primitive-part-alt-def ng-def by simp
    have lead-coeff f ≠ 0 using False by auto
    hence lg: lead-coeff g > 0 unfolding g-def lead-coeff-smult
    by (meson linorder-neqE-linordered-idom sgn-greater sgn-less zero-less-mult-iff)
    hence g0: g ≠ 0 by auto
    from g0 have content g ≠ 0 by simp
```

```
    from arg-cong[OF content-times-primitive-part[of g], of lead-coeff, unfolded
lead-coeff-smult]
    lg content-ge-0-int[of g] have lg': lead-coeff ng > 0 unfolding ng-def
    by (metis ‹content g ≠ 0› dual-order.antisym dual-order.strict-implies-order
zero-less-mult-iff)
   from content-primitive-part[OF g0] have c-ng: content ng = 1 unfolding
ng-def .
  have degree ng = degree f using ‹content [:sgn (lead-coeff f):] = 1› g-def ng-def
    by (auto simp add: sgn-eq-0-iff)
   with False have degree ng ≠ 0 by auto
   note main = square-free-factorization-int-main[OF fs c-ng lg' this]
   show ?thesis
   proof (intro conjI impI)
    {
     assume (fi, i) ∈ set fs
     with main show content fi = 1 0 < lead-coeff fi by auto
    }
    have d0: d ≠ 0 using ‹content [:?s:] = 1› d by (auto simp:sgn-eq-0-iff)
    have smult d ng = smult ?s (smult (content g) (primitive-part g))
     unfolding ng-def d cg by simp
   also have smult (content g) (primitive-part g) = g using content-times-primitive-part
.
    also have smult ?s g = f unfolding g-def using s by auto
    finally have id: smult d ng = f .
    from main have square-free-factorization ng (1, fs) by auto
    from square-free-factorization-smult[OF d0 this]
    show square-free-factorization f (d,fs) unfolding id by simp
    show distinct (map snd fs) using main by auto
   qed
  qed
  thus  square-free-factorization f (d,fs)
   (fi, i) ∈ set fs ⟹ content fi = 1 ∧ lead-coeff fi > 0 distinct (map snd fs) by
auto
qed


definition x-split :: 'a :: semiring-0 poly ⇒ nat × 'a poly where
  x-split f = (let fs = coeffs f; zs = takeWhile ((=) 0) fs
    in case zs of [] ⇒ (0,f) | - ⇒ (length zs, poly-of-list (dropWhile ((=) 0) fs)))

lemma x-split: assumes x-split f = (n, g)
  shows f = monom 1 n * g n ≠ 0 ∨ f ≠ 0 ⟹ ¬ monom 1 1 dvd g
proof −
  define zs where zs = takeWhile ((=) 0) (coeffs f)
  note res = assms[unfolded zs-def[symmetric] x-split-def Let-def]
  have f = monom 1 n * g ∧ ((n ≠ 0 ∨ f ≠ 0) ⟶ ¬ (monom 1 1 dvd g)) (is -
∧ (- ⟶ ¬ (?x dvd -)))
  proof (cases f = 0)
    case True
```

```
    with res have n = 0 g = 0 unfolding zs-def by auto
    thus ?thesis using True by auto
  next
    case False note f = this
    show ?thesis
    proof (cases zs = [])
      case True
    hence choice: coeff f 0 ≠ 0 using f unfolding zs-def coeff-f-0-code poly-compare-0-code
        by (cases coeffs f, auto)
     have dvd: ?x dvd h ⟷ coeff h 0 = 0 for h by (simp add: monom-1-dvd-iff')
      from True choice res f show ?thesis unfolding dvd by auto
    next
      case False
      define ys where ys = dropWhile ((=) 0) (coeffs f)
     have dvd: ?x dvd h ⟷ coeff h 0 = 0 for h by (simp add: monom-1-dvd-iff')
      from res False have n: n = length zs and g: g = poly-of-list ys unfolding
ys-def
        by (cases zs, auto)+
      obtain xx where xx: coeffs f = xx by auto
      have coeffs f = zs @ ys unfolding zs-def ys-def by auto
      also have zs = replicate n 0 unfolding zs-def n xx by (induct xx, auto)
      finally have ff: coeffs f = replicate n 0 @ ys by auto
      from f have lead-coeff f ≠ 0 by auto
      then have nz: coeffs f ≠ [] last (coeffs f) ≠ 0
        by (simp-all add: last-coeffs-eq-coeff-degree)
      have ys: ys ≠ [] using nz[unfolded ff] by auto
      with ys-def have hd: hd ys ≠ 0 by (metis (full-types) hd-dropWhile)
     hence coeff (poly-of-list ys) 0 ≠ 0 unfolding poly-of-list-def coeff-Poly using
ys by (cases ys, auto)
      moreover have coeffs (Poly ys) = ys
        by (simp add: ys-def strip-while-dropWhile-commute)
      then have coeffs (monom-mult n (Poly ys)) = replicate n 0 @ ys
      by (simp add: coeffs-eq-iff monom-mult-def [symmetric] ff ys monom-mult-code)
      ultimately show ?thesis unfolding dvd g
        by (auto simp add: coeffs-eq-iff monom-mult-def [symmetric] ff)
    qed
  qed
  thus f = monom 1 n * g n ≠ 0 ∨ f ≠ 0 ⟹ ¬ monom 1 1 dvd g by auto
qed


definition square-free-factorization-int :: int poly ⇒ int × (int poly × nat)list
where
  square-free-factorization-int f = (case x-split f of (n,g) — extract x^n
    ⇒ case square-free-factorization-int' g of (d,fs)
    ⇒ if n = 0 then (d,fs) else (d, (monom 1 1, n) # fs))

lemma square-free-factorization-int: assumes res: square-free-factorization-int f
= (d, fs)
```

```
  shows square-free-factorization f (d,fs)
    (fi, i) ∈ set fs ⟹ primitive fi ∧ lead-coeff fi > 0
proof −
  obtain n g where xs: x-split f = (n,g) by force
  obtain c hs where sf: square-free-factorization-int' g = (c,hs) by force
  from res[unfolded square-free-factorization-int-def xs sf split]
  have d: d = c and fs: fs = (if n = 0 then hs else (monom 1 1, n) # hs) by
(cases n, auto)
  note sff = square-free-factorization-int'(1−2)[OF sf]
  note xs = x-split[OF xs]
  let ?x = monom 1 1 :: int poly
  have x: primitive ?x ∧ lead-coeff ?x = 1 ∧ degree ?x = 1
    by (auto simp add: degree-monom-eq content-def monom-Suc)
  thus (fi, i) ∈ set fs ⟹ primitive fi ∧ lead-coeff fi > 0 using sff(2) unfolding
fs
    by (cases n, auto)
  show square-free-factorization f (d,fs)
  proof (cases n)
    case 0
    with d fs sff xs show ?thesis by auto
  next
    case (Suc m)
    with xs have fg: f = monom 1 (Suc m) ∗ g and dvd: ¬ ?x dvd g by auto
    from Suc have fs: fs = (?x,Suc m) # hs unfolding fs by auto
    have degx: degree ?x = 1 by code-simp
   from irreducible_d-square-free[OF linear-irreducible_d[OF this]] have sfx: square-free
?x by auto
    have fg: f = ?x ^ n ∗ g unfolding fg Suc by (metis x-pow-n)
    have eq0: ?x ^ n ∗ g = 0 ⟷ g = 0 by simp
    note sf = square-free-factorizationD[OF sff(1)]
    {
      fix a i
      assume ai: (a,i) ∈ set hs
      with sf(4) sf(2) have g0: g ≠ 0 and i > 0 by auto
      from split-list[OF ai] obtain ys zs where hs: hs = ys @ (a,i) # zs by auto
      have a dvd g unfolding square-free-factorization-prod-list[OF sff(1)] hs
        by (rule dvd-smult, insert ‹i > 0›, cases i, auto simp add: ac-simps)
      moreover have ¬ ?x dvd g using xs[unfolded Suc] by auto
      ultimately have dvd: ¬ ?x dvd a using dvd-trans by blast
      from sf(2)[OF ai] have a ≠ 0 by auto
      have 1 = gcd ?x a
      proof (rule gcdI)
        fix d
        assume d: d dvd ?x d dvd a
       from content-dvd-contentI[OF d(1)] x have cnt: is-unit (content d) by auto
        show is-unit d
        proof (cases degree d = 1)
          case False
          with divides-degree[OF d(1), unfolded degx] have degree d = 0 by auto
```

from *degree0-coeffs*[*OF this*] **obtain** *c* **where** *dc*: *d* = [:*c*:] **by** *auto*
from *cnt*[*unfolded dc*] **have** *is-unit c* **by** (*auto simp*: *content-def*, *cases c = 0*, *auto*)
hence *d* ∗ *d* = *1* **unfolding** *dc* **by** (*cases c* = −*1*; *cases c* = *1*, *auto*)
thus *is-unit d* **by** (*metis dvd-triv-right*)
next
case *True*
from *d*(*1*) **obtain** *e* **where** *xde*: *?x* = *d* ∗ *e* **unfolding** *dvd-def* **by** *auto*
from *arg-cong*[*OF this*, *of degree*] *degx* **have** *degree d* + *degree e* = *1*
by (*metis True add.right-neutral degree-0 degree-mult-eq one-neq-zero*)
with *True* **have** *degree e* = *0* **by** *auto*
from *degree0-coeffs*[*OF this*] *xde* **obtain** *e* **where** *xde*: *?x* = [:*e*:] ∗ *d* **by** *auto*
from *arg-cong*[*OF this*, *of content*, *unfolded content-mult*] *x*
have *content* [:*e*:] ∗ *content d* = *1* **by** *auto*
also have *content* [:*e* :] = *abs e* **by** (*auto simp*: *content-def*, *cases e* = *0*, *auto*)
finally have |*e*| ∗ *content d* = *1* **.**
from *pos-zmult-eq-1-iff-lemma*[*OF this*] **have** *e* ∗ *e* = *1* **by** (*cases e* = *1*; *cases e* = −*1*, *auto*)
with *arg-cong*[*OF xde*, *of smult e*] **have** *d* = *?x* ∗ [:*e*:] **by** *auto*
hence *?x dvd d* **unfolding** *dvd-def* **by** *blast*
with *d*(*2*) **have** *?x dvd a* **by** (*metis dvd-trans*)
with *dvd* **show** *?thesis* **by** *auto*
qed
qed *auto*
hence *coprime ?x a*
by (*simp add*: *gcd-eq-1-imp-coprime*)
note *this dvd*
} **note** *hs-dvd-x* = *this*
from *hs-dvd-x*[*of ?x Suc m*]
have *nmem*: (*?x*,*Suc m*) ∉ *set hs* **by** *auto*
hence *eq*: *?x* ^ *n* ∗ *g* = *smult c* (∏(*a*, *i*)∈*set fs*. *a* ^ *i*)
unfolding *sf*(*1*) **unfolding** *fs Suc* **by** *simp*
show *?thesis* **unfolding** *fg d* **unfolding** *square-free-factorization-def split eq0* **unfolding** *eq*
proof (*intro conjI allI impI*, *rule refl*)
fix *a i*
assume *ai*: (*a*,*i*) ∈ *set fs*
thus *square-free a degree a* > *0 i* > *0* **using** *sf*(*2*) *sfx degx* **unfolding** *fs* **by** *auto*
fix *b j*
assume *bj*: (*b*,*j*) ∈ *set fs* **and** *diff*: (*a*,*i*) ≠ (*b*,*j*)
consider (*hs-hs*) (*a*,*i*) ∈ *set hs* (*b*,*j*) ∈ *set hs*
| (*hs-x*) (*a*,*i*) ∈ *set hs b* = *?x*
| (*x-hs*) (*b*,*j*) ∈ *set hs a* = *?x*
using *ai bj diff* **unfolding** *fs* **by** *auto*
then show *Rings.coprime a b*
proof *cases*

case *hs-hs*
from *sf*(*3*)[*OF this diff*] **show** *?thesis* **.**
next
case *hs-x*
from *hs-dvd-x*(*1*)[*OF hs-x*(*1*)] **show** *?thesis* **unfolding** *hs-x*(*2*) **by** (*simp add*: *ac-simps*)
next
case *x-hs*
from *hs-dvd-x*(*1*)[*OF x-hs*(*1*)] **show** *?thesis* **unfolding** *x-hs*(*2*) **by** *simp*
qed
next
show $g = 0 \Longrightarrow c = 0$ **using** *sf*(*4*) **by** *auto*
show $g = 0 \Longrightarrow fs = []$ **using** *sf*(*4*) *xs Suc* **by** *auto*
show *distinct fs* **using** *sf*(*5*) *nmem* **unfolding** *fs* **by** *auto*
qed
qed
qed

end

11.3 Factoring Arbitrary Integer Polynomials

We combine the factorization algorithm for square-free integer polynomials with a square-free factorization algorithm to a factorization algorithm for integer polynomials which does not make any assumptions.

theory *Factorize-Int-Poly*
imports
Berlekamp-Zassenhaus
Square-Free-Factorization-Int
begin

hide-const *coeff monom*
lifting-forget *poly.lifting*

typedef *int-poly-factorization-algorithm* = {*alg*.
$\forall$ (*f* :: *int poly*) *fs*. *square-free f* $\longrightarrow$ *degree f* > *0* $\longrightarrow$ *alg f* = *fs* $\longrightarrow$
(*f* = *prod-list fs* $\wedge$ ($\forall$ *fi* $\in$ *set fs*. *irreducible*$_d$ *fi*))}
by (*rule exI*[*of* - *berlekamp-zassenhaus-factorization*],
insert berlekamp-zassenhaus-factorization-irreducible$_d$, *auto*)

setup-lifting *type-definition-int-poly-factorization-algorithm*

lift-definition *int-poly-factorization-algorithm* :: *int-poly-factorization-algorithm* $\Rightarrow$
(*int poly* $\Rightarrow$ *int poly list*) **is** λ *x*. *x* **.**

lemma *int-poly-factorization-algorithm-irreducible*$_d$:
assumes *int-poly-factorization-algorithm alg f* = *fs*
and *square-free f*

and *degree f > 0*
shows *f = prod-list fs ∧ (∀ fi ∈ set fs. irreducible$_d$ fi)*
using *assms* **by** (*transfer*, *auto*)

corollary *int-poly-factorization-algorithm-irreducible*:
assumes *res*: *int-poly-factorization-algorithm alg f = fs*
and *sf*: *square-free f*
and *deg*: *degree f > 0*
and *pr*: *primitive f*
shows *f = prod-list fs ∧ (∀ fi ∈ set fs. irreducible fi ∧ degree fi > 0 ∧ primitive fi)*
proof (*intro conjI ballI*)
note ∗ = *int-poly-factorization-algorithm-irreducible$_d$*[*OF res sf deg*]
from ∗ **show** *f*: *f = prod-list fs* **by** *auto*
fix *fi* **assume** *fi*: *fi ∈ set fs*
with *primitive-prod-list*[*OF pr*[*unfolded f*]] **show** *primitive fi* **by** *auto*
from *irreducible-primitive-connect*[*OF this*] ∗ *pr*[*unfolded f*] *fi*
show *irreducible fi* **by** *auto*
from ∗ *fi* **show** *degree fi > 0* **by** (*auto*)
qed

lemma *irreducible-imp-square-free*:
assumes *irr*: *irreducible* (*p*::*'a*::*idom poly*) **shows** *square-free p*
proof(*intro square-freeI*)
from *irr* **show** *p0*: *p ≠ 0* **by** *auto*
fix *a* **assume** *a ∗ a dvd p*
then obtain *b* **where** *paab*: *p = a ∗ (a ∗ b)* **by** (*elim dvdE*, *auto*)
assume *degree a > 0*
then have *a1*: *¬ a dvd 1* **by** (*auto simp*: *poly-dvd-1*)
then have *ab1*: *¬ a ∗ b dvd 1* **using** *dvd-mult-left* **by** *auto*
from *paab irr a1 ab1* **show** *False* **by** *force*
qed

lemma *not-mem-set-dropWhileD*: *x ∉ set* (*dropWhile P xs*) ⟹ *x ∈ set xs* ⟹ *P x*
by (*metis dropWhile-append3 in-set-conv-decomp*)

lemma *primitive-reflect-poly*:
fixes *f* :: *'a* :: *comm-semiring-1 poly*
shows *primitive* (*reflect-poly f*) = *primitive f*
proof−
have (∀ *a* ∈ *set* (*coeffs f*). *x dvd a*) ⟷ (∀ *a* ∈ *set* (*dropWhile* ((=) *0*) (*coeffs f*)). *x dvd a*) **for** *x*
by (*auto dest*: *not-mem-set-dropWhileD set-dropWhileD*)
then show *?thesis* **by** (*auto simp*: *primitive-def coeffs-reflect-poly*)
qed

lemma *gcd-list-sub*:
assumes *set xs* ⊆ *set ys* **shows** *gcd-list ys dvd gcd-list xs*
by (*metis Gcd-fin.subset assms semiring-gcd-class.gcd-dvd1*)

lemma *content-reflect-poly*:
content (*reflect-poly f*) = *content f* (**is** *?l* = *?r*)
proof−
have *l*: *?l* = *gcd-list* (*dropWhile* ((=) *0*) (*coeffs f*)) (**is** - = *gcd-list ?xs*)
by (*simp add*: *content-def reflect-poly-def*)
have *set ?xs* ⊆ *set* (*coeffs f*) **by** (*auto dest*: *set-dropWhileD*)
from *gcd-list-sub*[*OF this*]
have *?r dvd gcd-list ?xs* **by** (*simp add*: *content-def*)
with *l* **have** *rl*: *?r dvd ?l* **by** *auto*
have *set* (*coeffs f*) ⊆ *set* (*0* # *?xs*) **by** (*auto dest*: *not-mem-set-dropWhileD*)
from *gcd-list-sub*[*OF this*]
have *gcd-list ?xs dvd ?r* **by** (*simp add*: *content-def*)
with *l* **have** *lr*: *?l dvd ?r* **by** *auto*
from *rl lr* **show** *?l* = *?r* **by** (*simp add*: *associated-eqI*)
qed

lemma *coeff-primitive-part*: *content f* ∗ *coeff* (*primitive-part f*) *i* = *coeff f i*
using *arg-cong*[*OF content-times-primitive-part*[*of f*], *of* λ*f*. *coeff f* -, *unfolded coeff-smult*].

lemma *smult-cancel*[*simp*]:
fixes *c* :: *'a* :: *idom*
shows *smult c f* = *smult c g* ⟷ *c* = *0* ∨ *f* = *g*
proof−
have *l*: *smult c f* = [:*c*:] ∗ *f* **by** *simp*
have *r*: *smult c g* = [:*c*:] ∗ *g* **by** *simp*
show *?thesis* **unfolding** *l r mult-cancel-left* **by** *simp*
qed

lemma *primitive-part-reflect-poly*:
fixes *f* :: *'a* :: {*semiring-gcd*,*idom*} *poly*
shows *primitive-part* (*reflect-poly f*) = *reflect-poly* (*primitive-part f*) (**is** *?l* = *?r*)
using *content-times-primitive-part*[*of reflect-poly f*]
proof−
note *content-reflect-poly*[*of f*, *symmetric*]
also have *smult* (*content* (*reflect-poly f*)) *?l* = *reflect-poly f* **by** *simp*
also have ... = *reflect-poly* (*smult* (*content f*) (*primitive-part f*)) **by** *simp*
finally show *?thesis* **unfolding** *reflect-poly-smult smult-cancel* **by** *auto*
qed

lemma *reflect-poly-eq-zero*[*simp*]:
reflect-poly f = *0* ⟷ *f* = *0*
proof

```
  assume reflect-poly f = 0
  then have coeff (reflect-poly f) 0 = 0 by simp
  then have lead-coeff f = 0 by simp
  then show f = 0 by simp
qed simp

lemma irreducible_d-reflect-poly-main:
  fixes f :: 'a :: {idom, semiring-gcd} poly
  assumes nz: coeff f 0 ≠ 0
    and irr: irreducible_d (reflect-poly f)
  shows irreducible_d f
proof
  let ?r = reflect-poly
  from irr degree-reflect-poly-eq[OF nz] show degree f > 0 by auto
  fix g h
  assume deg: degree g < degree f degree h < degree f and fgh: f = g * h
  from arg-cong[OF fgh, of λ f. coeff f 0] nz
  have nz': coeff g 0 ≠ 0 by (auto simp: coeff-mult-0)
  note rfgh = arg-cong[OF fgh, of reflect-poly, unfolded reflect-poly-mult[of g h]]
 from deg degree-reflect-poly-le[of g] degree-reflect-poly-le[of h] degree-reflect-poly-eq[OF
nz]
  have degree (?r h) < degree (?r f) degree (?r g) < degree (?r f) by auto
  with irr rfgh show False by auto
qed

lemma irreducible_d-reflect-poly:
  fixes f :: 'a :: {idom, semiring-gcd} poly
  assumes nz: coeff f 0 ≠ 0
  shows irreducible_d (reflect-poly f) = irreducible_d f
proof
  assume irreducible_d (reflect-poly f)
  from irreducible_d-reflect-poly-main[OF nz this] show irreducible_d f .
next
  from nz have nzr: coeff (reflect-poly f) 0 ≠ 0 by auto
  assume irreducible_d f
  with nz have irreducible_d (reflect-poly (reflect-poly f)) by simp
  from irreducible_d-reflect-poly-main[OF nzr this]
  show irreducible_d (reflect-poly f) .
qed

lemma irreducible-reflect-poly:
  fixes f :: 'a :: {idom,semiring-gcd} poly
  assumes nz: coeff f 0 ≠ 0
  shows irreducible (reflect-poly f) = irreducible f (is ?l = ?r)
proof (cases degree f = 0)
  case True then obtain f0 where f = [:f0:] by (auto dest: degree0-coeffs)
  then show ?thesis by simp
next
  case deg: False
```

```
  show ?thesis
  proof (cases primitive f)
    case False
    with deg irreducible-imp-primitive[of f] irreducible-imp-primitive[of reflect-poly
f] nz
    show ?thesis unfolding primitive-reflect-poly by auto
  next
    case cf: True
    let ?r = reflect-poly
    from nz have nz': coeff (?r f) 0 ≠ 0 by auto
    let ?ir = irreducible_d
    from irreducible_d-reflect-poly[OF nz] irreducible_d-reflect-poly[OF nz'] nz
    have ?ir f ⟷ ?ir (reflect-poly f) by auto
    also have ... ⟷ irreducible (reflect-poly f)
      by (rule irreducible-primitive-connect, unfold primitive-reflect-poly, fact cf)
    finally show ?thesis
      by (unfold irreducible-primitive-connect[OF cf], auto)
  qed
qed


lemma reflect-poly-dvd: (f :: 'a :: idom poly) dvd g ⟹ reflect-poly f dvd reflect-poly
g
  unfolding dvd-def by (auto simp: reflect-poly-mult)

lemma square-free-reflect-poly: fixes f :: 'a :: idom poly
  assumes sf: square-free f
  and nz: coeff f 0 ≠ 0
shows square-free (reflect-poly f) unfolding square-free-def
proof (intro allI conjI impI notI)
  let ?r = reflect-poly
  from sf[unfolded square-free-def]
  have f0: f ≠ 0 and sf: ⋀ q. 0 < degree q ⟹ q * q dvd f ⟹ False by auto
  from f0 nz show ?r f = 0 ⟹ False by auto
  fix q
  assume 0: 0 < degree q and dvd: q * q dvd ?r f
  from dvd have q dvd ?r f by auto
  then obtain x where id: ?r f = q * x by fastforce
  {
    assume coeff q 0 = 0
    hence coeff (?r f) 0 = 0 using id by (auto simp: coeff-mult)
    with nz have False by auto
  }
  hence nzq: coeff q 0 ≠ 0 by auto
  from dvd have ?r (q * q) dvd ?r (?r f) by (rule reflect-poly-dvd)
  also have ?r (?r f) = f using nz by auto
  also have ?r (q * q) = ?r q * ?r q by (rule reflect-poly-mult)
  finally have ?r q * ?r q dvd f .
  from sf[OF - this] 0 nzq show False by simp
```

qed

lemma *gcd-reflect-poly*: **fixes** *f* :: *'a* :: {*factorial-ring-gcd*, *semiring-gcd-mult-normalize*} *poly*
assumes *nz*: *coeff f 0* ≠ *0 coeff g 0* ≠ *0*
shows *gcd* (*reflect-poly f*) (*reflect-poly g*) = *normalize* (*reflect-poly* (*gcd f g*))
proof (*rule sym*, *rule gcdI*)
have *gcd f g dvd f* **by** *auto*
from *reflect-poly-dvd*[*OF this*]
show *normalize* (*reflect-poly* (*gcd f g*)) *dvd reflect-poly f* **by** *simp*
have *gcd f g dvd g* **by** *auto*
from *reflect-poly-dvd*[*OF this*]
show *normalize* (*reflect-poly* (*gcd f g*)) *dvd reflect-poly g* **by** *simp*
show *normalize* (*normalize* (*reflect-poly* (*gcd f g*))) = *normalize* (*reflect-poly* (*gcd f g*)) **by** *auto*
fix *h*
assume *hf*: *h dvd reflect-poly f* **and** *hg*: *h dvd reflect-poly g*
from *hf* **obtain** *k* **where** *reflect-poly f* = *h* ∗ *k* **unfolding** *dvd-def* **by** *auto*
from *arg-cong*[*OF this*, *of* λ *f*. *coeff f 0*, *unfolded coeff-mult-0*] *nz*(*1*) **have** *h*: *coeff h 0* ≠ *0* **by** *auto*
from *reflect-poly-dvd*[*OF hf*] *reflect-poly-dvd*[*OF hg*]
have *reflect-poly h dvd f reflect-poly h dvd g* **using** *nz* **by** *auto*
hence *reflect-poly h dvd gcd f g* **by** *auto*
from *reflect-poly-dvd*[*OF this*] *h* **have** *h dvd reflect-poly* (*gcd f g*) **by** *auto*
thus *h dvd normalize* (*reflect-poly* (*gcd f g*)) **by** *auto*
qed

lemma *linear-primitive-irreducible*:
fixes *f* :: *'a* :: {*comm-semiring-1*,*semiring-no-zero-divisors*} *poly*
assumes *deg*: *degree f* = *1* **and** *cf*: *primitive f*
shows *irreducible f*
proof (*intro irreducibleI*)
fix *a b* **assume** *fab*: *f* = *a* ∗ *b*
with *deg* **have** *a0*: *a* ≠ *0* **and** *b0*: *b* ≠ *0* **by** *auto*
from *deg*[*unfolded fab*] *degree-mult-eq*[*OF this*] **have** *degree a* = *0* ∨ *degree b* = *0* **by** *auto*
then show *a dvd 1* ∨ *b dvd 1*
proof
assume *degree a* = *0*
then obtain *a0* **where** *a*: *a* = [:*a0*:] **by** (*auto dest*:*degree0-coeffs*)
with *fab* **have** *c* ∈ *set* (*coeffs f*) ⟹ *a0 dvd c* **for** *c* **by** (*cases a0* = *0*, *auto simp*: *coeffs-smult*)
with *cf* **show** *?thesis* **by** (*auto dest*: *primitiveD simp*: *a*)
next
assume *degree b* = *0*
then obtain *b0* **where** *b*: *b* = [:*b0*:] **by** (*auto dest*:*degree0-coeffs*)
with *fab* **have** *c* ∈ *set* (*coeffs f*) ⟹ *b0 dvd c* **for** *c* **by** (*cases b0* = *0*, *auto simp*: *coeffs-smult*)
with *cf* **show** *?thesis* **by** (*auto dest*: *primitiveD simp*: *b*)

```
  qed
qed (insert deg, auto simp: poly-dvd-1)

lemma square-free-factorization-last-coeff-nz:
  assumes sff: square-free-factorization f (a, fs)
  and mem: (fi,i) ∈ set fs
  and nz: coeff f 0 ≠ 0
shows coeff fi 0 ≠ 0
proof
  assume fi: coeff fi 0 = 0
  note sff-list = square-free-factorization-prod-list[OF sff]
  note sff = square-free-factorizationD[OF sff]
  from sff-list have coeff f 0 = a * coeff (∏ (a, i)←fs. a ^ i) 0 by simp
  with split-list[OF mem] fi sff(2)[OF mem] have coeff f 0 = 0
    by (cases i, auto simp: coeff-mult)
  with nz show False by simp
qed



context
  fixes alg :: int-poly-factorization-algorithm
begin

definition main-int-poly-factorization :: int poly ⇒ int poly list where
  main-int-poly-factorization f = (let df = degree f
    in if df = 1 then [f] else
    if abs (coeff f 0) < abs (coeff f df) — take reciprocal polynomial, if f(0) < lc(f)
     then map reflect-poly (int-poly-factorization-algorithm alg (reflect-poly f))
     else int-poly-factorization-algorithm alg f)


definition internal-int-poly-factorization :: int poly ⇒ int × (int poly × nat) list
where
  internal-int-poly-factorization f = (
    case square-free-factorization-int f of
     (a,gis) ⇒ (a, [ (h,i) . (g,i) ← gis, h ← main-int-poly-factorization g ])
  )

lemma internal-int-poly-factorization-code[code]: internal-int-poly-factorization f =
(
    case square-free-factorization-int f of (a,gis) ⇒
  (a, concat (map (λ (g,i). (map (λ f. (f,i)) (main-int-poly-factorization g))) gis)))
  unfolding internal-int-poly-factorization-def by auto


definition factorize-int-last-nz-poly :: int poly ⇒ int × (int poly × nat) list where
  factorize-int-last-nz-poly f = (let df = degree f
    in if df = 0 then (coeff f 0, []) else if df = 1 then (content f,[(primitive-part
```

```
f,1)]) else
    internal-int-poly-factorization f)


definition factorize-int-poly-generic :: int poly ⇒ int × (int poly × nat) list where
  factorize-int-poly-generic f = (case x-split f of (n,g) — extract x^n
    ⇒ if g = 0 then (0,[]) else case factorize-int-last-nz-poly g of (a,fs)
    ⇒ if n = 0 then (a,fs) else (a, (monom 1 1, n) # fs))


lemma factorize-int-poly-0[simp]: factorize-int-poly-generic 0 = (0,[])
  unfolding factorize-int-poly-generic-def x-split-def by simp

lemma main-int-poly-factorization:
  assumes res: main-int-poly-factorization f = fs
  and sf: square-free f
  and df: degree f > 0
  and nz: coeff f 0 ≠ 0
shows f = prod-list fs ∧ (∀ fi ∈ set fs. irreducible_d fi)
proof (cases degree f = 1)
  case True
  with res[unfolded main-int-poly-factorization-def Let-def]
  have fs = [f] by auto
  with True show ?thesis by auto
next
  case False
  hence *: (if degree f = 1 then t :: int poly list else e) = e for t e by auto
  note res = res[unfolded main-int-poly-factorization-def Let-def *]
  show ?thesis
  proof (cases abs (coeff f 0) < abs (coeff f (degree f)))
    case False
    with res have int-poly-factorization-algorithm alg f = fs by auto
    from int-poly-factorization-algorithm-irreducible_d[OF this sf df] show ?thesis .
  next
    case True
    let ?f = reflect-poly f
    from square-free-reflect-poly[OF sf nz] have sf: square-free ?f .
    from nz df have df: degree ?f > 0 by simp
    from True res obtain gs where fs: fs = map reflect-poly gs
      and gs: int-poly-factorization-algorithm alg (reflect-poly f) = gs
      by auto
    from int-poly-factorization-algorithm-irreducible_d[OF gs sf df]
    have id: reflect-poly ?f = reflect-poly (prod-list gs) ?f = prod-list gs
      and irr: ⋀ gi. gi ∈ set gs ⟹ irreducible_d gi by auto
    from id(1) have f-fs: f = prod-list fs unfolding fs using nz
      by (simp add: reflect-poly-prod-list)
    {
      fix fi
      assume fi ∈ set fs
```

```
    from this[unfolded fs] obtain gi where gi: gi ∈ set gs and fi: fi = reflect-poly
gi by auto
    {
      assume coeff gi 0 = 0
      with id(2) split-list[OF gi] have coeff ?f 0 = 0
        by (auto simp: coeff-mult)
      with nz have False by auto
    }
    hence nzg: coeff gi 0 ≠ 0 by auto
    from irreducible_d-reflect-poly[OF nzg] irr[OF gi] have irreducible_d fi unfold-
ing fi by simp
  }
  with f-fs show ?thesis by auto
 qed
qed

lemma internal-int-poly-factorization-mem:
 assumes f: coeff f 0 ≠ 0
 and res: internal-int-poly-factorization f = (c,fs)
 and mem: (fi,i) ∈ set fs
 shows irreducible fi irreducible_d fi and primitive fi and degree fi ≠ 0 i ≠ 0
proof −
 obtain a psi where a-psi: square-free-factorization-int f = (a, psi)
   by force
 from square-free-factorization-int[OF this]
 have sff: square-free-factorization f (a, psi)
   and cnt: ⋀ fi i. (fi, i) ∈ set psi ⟹ primitive fi by blast+
 from square-free-factorization-last-coeff-nz[OF sff - f]
 have nz-fi: ⋀ fi i. (fi, i) ∈ set psi ⟹ coeff fi 0 ≠ 0 by auto
 note res = res[unfolded internal-int-poly-factorization-def a-psi Let-def split]
 obtain fact where fact: fact = (λ (q,i :: nat). (map (λ f. (f,i)) (main-int-poly-factorization
q))) by auto
 from res[unfolded split Let-def]
 have c: c = a and fs: fs = concat (map fact psi)
   unfolding fact by auto
 note sff' = square-free-factorizationD[OF sff]
 from mem[unfolded fs, simplified] obtain d j where psi: (d,j) ∈ set psi
   and fi: (fi, i) ∈ set (fact (d,j)) by auto
 from square-free-factorizationD(2)[OF sff psi] have j > 0 by auto
 obtain hs where d: main-int-poly-factorization d = hs by force
 from fi[unfolded split fact] have j = i by auto
 with ‹j > 0› show i ≠ 0 by auto
 from fi[unfolded d split fact] have fi: fi ∈ set hs by auto
 from main-int-poly-factorization[OF d - - nz-fi[OF psi]] sff'(2)[OF psi] cnt[OF
psi]
 have main: d = prod-list hs ⋀ fi. fi ∈ set hs ⟹ irreducible_d fi by auto
 from main split-list[OF fi] have content fi dvd content d by auto
 with cnt[OF psi] show cnt: primitive fi by simp
 from main(2)[OF fi] show irr: irreducible_d fi .
```

```
  show irreducible fi
    using irreducible-primitive-connect[OF cnt] irr by blast
  from irr show degree fi ≠ 0 by auto
qed

lemma internal-int-poly-factorization:
  assumes f: coeff f 0 ≠ 0
  and res: internal-int-poly-factorization f = (c,fs)
  shows square-free-factorization f (c,fs)
proof −
  obtain a psi where a-psi: square-free-factorization-int f = (a, psi)
    by force
  from square-free-factorization-int[OF this]
  have sff: square-free-factorization f (a, psi)
    and pr: ⋀ fi i. (fi, i) ∈ set psi ⟹ primitive fi by blast+
 obtain fact where fact: fact = (λ (q,i :: nat). (map (λ f. (f,i)) (main-int-poly-factorization
q))) by auto
  from res[unfolded split Let-def]
  have c: c = a and fs: fs = concat (map fact psi)
    unfolding fact internal-int-poly-factorization-def a-psi by auto
  note sff' = square-free-factorizationD[OF sff]
  show ?thesis unfolding square-free-factorization-def split
  proof (intro conjI impI allI)
   show f = 0 ⟹ c = 0 f = 0 ⟹ fs = [] using sff'(4) unfolding c fs by auto
   {
     fix a i
     assume (a,i) ∈ set fs
     from irreducible-imp-square-free internal-int-poly-factorization-mem[OF f res
this]
     show square-free a degree a > 0 i > 0 by auto
   }
   from square-free-factorization-last-coeff-nz[OF sff - f]
   have nz: ⋀ fi i. (fi, i) ∈ set psi ⟹ coeff fi 0 ≠ 0 by auto
   have eq: f = smult c (∏ (a, i)←fs. a ^ i) unfolding
     prod.distinct-set-conv-list[OF sff'(5)]
     sff'(1) c
    proof (rule arg-cong[where f = smult a], unfold fs, insert sff'(2) nz, induct
psi)
     case (Cons pi psi)
     obtain p i where pi: pi = (p,i) by force
     obtain gs where gs: main-int-poly-factorization p = gs by auto
      from Cons(2)[of p i] have p: square-free p degree p > 0 unfolding pi by
auto
     from Cons(3)[of p i] have nz: coeff p 0 ≠ 0 unfolding pi by auto
      from main-int-poly-factorization[OF gs p nz] have pgs: p = prod-list gs by
auto
     have fact: fact (p,i) = map (λ g. (g,i)) gs unfolding fact split gs by auto
     have cong: ⋀ x y X Y. x = X ⟹ y = Y ⟹ x * y = X * Y by auto
       show ?case unfolding pi list.simps prod-list.Cons split fact concat.simps
```

```
prod-list.append
      map-append
    proof (rule cong)
     show p ^ i = (∏(a, i)←map (λg. (g, i)) gs. a ^ i) unfolding pgs
      by (induct gs, auto simp: ac-simps power-mult-distrib)
     show (∏(a, i)←psi. a ^ i) = (∏(a, i)←concat (map fact psi). a ^ i)
      by (rule Cons(1), insert Cons(2−3), auto)
    qed
   qed simp
   {
    fix i j l fi
    assume ∗: j < length psi l < length (fact (psi ! j)) fact (psi ! j) ! l = (fi, i)
    from ∗ have psi: psi ! j ∈ set psi by auto
    obtain d k where dk: psi ! j = (d,k) by force
    with ∗ have psij: psi ! j = (d,i) unfolding fact split by auto
     from sff'(2)[OF psi[unfolded psij]] have d: square-free d degree d > 0 by
auto
    from nz[OF psi[unfolded psij]] have d0: coeff d 0 ≠ 0 .
    from ∗ psij fact
     have bz: main-int-poly-factorization d = map fst (fact (psi ! j)) by (auto
simp: o-def)
    from main-int-poly-factorization[OF bz d d0] pr[OF psi[unfolded dk]]
    have dhs: d = prod-list (map fst (fact (psi ! j))) by auto
    from ∗ have mem: fi ∈ set (map fst (fact (psi ! j)))
     by (metis fst-conv image-eqI nth-mem set-map)
    from mem dhs psij d have ∃ d. fi ∈ set (map fst (fact (psi ! j))) ∧
     d = prod-list (map fst (fact (psi ! j))) ∧
     psi ! j = (d, i) ∧
     square-free d by blast
   } note deconstruct = this
   {
    fix k K fi i Fi I
    assume k: k < length fs K < length fs and f: fs ! k = (fi, i) fs ! K = (Fi, I)
    and diff: k ≠ K
    from nth-concat-diff[OF k[unfolded fs] diff, folded fs, unfolded length-map]
     obtain j l J L where diff: (j, l) ≠ (J, L)
      and j: j < length psi J < length psi
      and l: l < length (map fact psi ! j) L < length (map fact psi ! J)
      and fs: fs ! k = map fact psi ! j ! l fs ! K = map fact psi ! J ! L by blast+
    hence psij: psi ! j ∈ set psi by auto
    from j have id: map fact psi ! j = fact (psi ! j) map fact psi ! J = fact (psi
! J) by auto
    note l = l[unfolded id] note fs = fs[unfolded id]
    from j have psi: psi ! j ∈ set psi psi ! J ∈ set psi by auto
    from deconstruct[OF j(1) l(1) fs(1)[unfolded f, symmetric]]
    obtain d where mem: fi ∈ set (map fst (fact (psi ! j)))
     and d: d = prod-list (map fst (fact (psi ! j))) psi ! j = (d, i) square-free d
by blast
    from deconstruct[OF j(2) l(2) fs(2)[unfolded f, symmetric]]
```

obtain D **where** *Mem*: $Fi \in set\ (map\ fst\ (fact\ (psi\ !\ J)))$
and D: $D = prod\text{-}list\ (map\ fst\ (fact\ (psi\ !\ J)))\ psi\ !\ J = (D,\ I)$ *square-free* D **by** *blast*
from *pr*[*OF psij*[*unfolded d*(2)]] **have** *cnt*: *primitive d* **.**
have *coprime fi Fi*
proof (*cases* $J = j$)
case *False*
from *sff*′(5) *False j* **have** $(d,i) \neq (D,I)$
unfolding *distinct-conv-nth d*(2)[*symmetric*] *D*(2)[*symmetric*] **by** *auto*
from *sff*′(3)[*OF psi*[*unfolded d*(2) *D*(2)] *this*]
have *cop*: *coprime d D* **by** *auto*
from *prod-list-dvd*[*OF mem*, *folded d*(1)] **have** *fid*: *fi dvd d* **by** *auto*
from *prod-list-dvd*[*OF Mem*, *folded D*(1)] **have** *FiD*: *Fi dvd D* **by** *auto*
from *coprime-divisors*[*OF fid FiD*] *cop* **show** *?thesis* **by** *simp*
next
case *True* **note** $id = this$
from *id diff* **have** *diff*: $l \neq L$ **by** *auto*
obtain *bz* **where** *bz*: $bz = map\ fst\ (fact\ (psi\ !\ j))$ **by** *auto*
from *fs*[*unfolded f*] *l*
have *fi*: $fi = bz\ !\ l\ Fi = bz\ !\ L$
unfolding *id bz* **by** (*metis fst-conv nth-map*)+
from *d*[*folded bz*] **have** *sf*: *square-free* (*prod-list bz*) **by** *auto*
from *d*[*folded bz*] *cnt* **have** *cnt*: *content* (*prod-list bz*) = *1* **by** *auto*
from *l* **have** *l*: $l < length\ bz\ L < length\ bz$ **unfolding** *bz id* **by** *auto*
from *l fi* **have** *fi* $\in$ *set bz* **by** *auto*
from *content-dvd-1*[*OF cnt prod-list-dvd*[*OF this*]] **have** *cnt*: *content fi* = *1*
.
obtain *g* **where** *g*: $g = gcd\ fi\ Fi$ **by** *auto*
have *g*′: *g dvd fi g dvd Fi* **unfolding** *g* **by** *auto*
define *bef* **where** $bef = take\ l\ bz$
define *aft* **where** $aft = drop\ (Suc\ l)\ bz$
from *id-take-nth-drop*[*OF l*(1)] *l* **have** *bz*: $bz = bef\ @\ fi\ \#\ aft$ **and** *bef*: *length bef* = *l*
unfolding *bef-def aft-def fi* **by** *auto*
with *l diff* **have** *mem*: $Fi \in set\ (bef\ @\ aft)$ **unfolding** *fi*(2) **by** (*auto simp*: *nth-append*)
from *split-list*[*OF this*] **obtain** *Bef Aft* **where** *ba*: $bef\ @\ aft = Bef\ @\ Fi\ \#\ Aft$ **by** *auto*
have $prod\text{-}list\ bz = fi * prod\text{-}list\ (bef\ @\ aft)$ **unfolding** *bz* **by** *simp*
also have $prod\text{-}list\ (bef\ @\ aft) = Fi * prod\text{-}list\ (Bef\ @\ Aft)$ **unfolding** *ba* **by** *auto*
finally have $fi * Fi\ dvd\ prod\text{-}list\ bz$ **by** *auto*
with *g*′ **have** $g * g\ dvd\ prod\text{-}list\ bz$ **by** (*meson dvd-trans mult-dvd-mono*)
with *sf*[*unfolded square-free-def*] **have** *deg*: *degree g* = *0* **by** *auto*
from *content-dvd-1*[*OF cnt g*′(1)] **have** *cnt*: *content g* = *1* **.**
from *degree0-coeffs*[*OF deg*] **obtain** *c* **where** *gc*: $g = [:\ c\ :]$ **by** *auto*
from *cnt*[*unfolded gc content-def*, *simplified*] **have** *abs c* = *1*
by (*cases c* = *0*, *auto*)
with *g gc* **have** *gcd fi Fi* $\in \{1,-1\}$ **by** *fastforce*

```
      thus coprime fi Fi
        by (auto intro!: gcd-eq-1-imp-coprime)
          (metis dvd-minus-iff dvd-refl is-unit-gcd-iff one-neq-neg-one)
    qed
  } note cop = this
  show dist: distinct fs unfolding distinct-conv-nth
  proof (intro impI allI)
    fix k K
    assume k: k < length fs K < length fs and diff: k ≠ K
    obtain fi i Fi I where f: fs ! k = (fi,i) fs ! K = (Fi,I) by force+
    from cop[OF k f diff] have cop: coprime fi Fi .
    from k(1) f(1) have (fi,i) ∈ set fs unfolding set-conv-nth by force
    from internal-int-poly-factorization-mem[OF assms(1) res this] have degree
fi > 0 by auto
    hence ¬ is-unit fi by (simp add: poly-dvd-1)
    with cop coprime-id-is-unit[of fi] have fi ≠ Fi by auto
    thus fs ! k ≠ fs ! K unfolding f by auto
  qed
  show f = smult c (∏(a, i)∈set fs. a ^ i) unfolding eq
    prod.distinct-set-conv-list[OF dist] by simp
  fix fi i Fi I
  assume mem: (fi, i) ∈ set fs (Fi,I) ∈ set fs and diff: (fi, i) ≠ (Fi, I)
  then obtain k K where k: k < length fs K < length fs
    and f: fs ! k = (fi, i) fs ! K = (Fi, I) unfolding set-conv-nth by auto
  with diff have diff: k ≠ K by auto
  from cop[OF k f diff] show Rings.coprime fi Fi by auto
 qed
qed

lemma factorize-int-last-nz-poly: assumes res: factorize-int-last-nz-poly f = (c,fs)
  and nz: coeff f 0 ≠ 0
shows square-free-factorization f (c,fs)
 (fi,i) ∈ set fs ⟹ irreducible fi
 (fi,i) ∈ set fs ⟹ degree fi ≠ 0
proof (atomize(full))
 from nz have lz: lead-coeff f ≠ 0 by auto
 note res = res[unfolded factorize-int-last-nz-poly-def Let-def]
 consider (0) degree f = 0
   | (1) degree f = 1
   | (2) degree f > 1 by linarith
 then show square-free-factorization f (c,fs) ∧ ((fi,i) ∈ set fs ⟶ irreducible fi)
∧ ((fi,i) ∈ set fs ⟶ degree fi ≠ 0)
 proof cases
   case 0
   from degree0-coeffs[OF 0] obtain a where f: f = [:a:] by auto
   from res show ?thesis unfolding square-free-factorization-def f by auto
 next
   case 1
   then have irr: irreducible (primitive-part f)
```

```
    by (auto intro!: linear-primitive-irreducible content-primitive-part)
  from irreducible-imp-square-free[OF irr] have sf: square-free (primitive-part f)
.
   from 1 have f0: f ≠ 0 by auto
    from res irr sf f0 show ?thesis unfolding square-free-factorization-def by
(auto simp: 1)
 next
   case 2
   with res have internal-int-poly-factorization f = (c,fs) by auto
  from internal-int-poly-factorization[OF nz this] internal-int-poly-factorization-mem[OF
nz this]
   show ?thesis by auto
 qed
qed

lemma factorize-int-poly: assumes res: factorize-int-poly-generic f = (c,fs)
shows square-free-factorization f (c,fs)
  (fi,i) ∈ set fs ⟹ irreducible fi
  (fi,i) ∈ set fs ⟹ degree fi ≠ 0
proof (atomize(full))
  obtain n g where xs: x-split f = (n,g) by force
  obtain d hs where fact: factorize-int-last-nz-poly g = (d,hs) by force
  from res[unfolded factorize-int-poly-generic-def xs split fact]
  have res: (if g = 0 then (0, []) else if n = 0 then (d, hs) else (d, (monom 1 1,
n) # hs)) = (c, fs) .
  note xs = x-split[OF xs]
 show square-free-factorization f (c,fs) ∧ ((fi,i) ∈ set fs ⟶ irreducible fi) ∧ ((fi,i)
∈ set fs ⟶ degree fi ≠ 0)
  proof (cases g = 0)
   case True
   hence f = 0 c = 0 fs = [] using res xs by auto
   thus ?thesis unfolding square-free-factorization-def by auto
  next
   case False
   with xs have ¬ monom 1 1 dvd g by auto
   hence coeff g 0 ≠ 0 by (simp add: monom-1-dvd-iff')
   note fact = factorize-int-last-nz-poly[OF fact this]
   let ?x = monom 1 1 :: int poly
   have x: content ?x = 1 ∧ lead-coeff ?x = 1 ∧ degree ?x = 1
    by (auto simp add: degree-monom-eq monom-Suc content-def)
   from res False have res: (if n = 0 then (d, hs) else (d, (?x, n) # hs)) = (c,
fs) by auto
   show ?thesis
   proof (cases n)
    case 0
    with res xs have id: fs = hs c = d f = g by auto
    from fact show ?thesis unfolding id by auto
   next
    case (Suc m)
```

with *res* **have** *id*: $c = d\ fs = (?x,Suc\ m) \#\ hs$ **by** *auto*
from *Suc xs* **have** *fg*: $f = monom\ 1\ (Suc\ m) * g$ **and** *dvd*: $\neg\ ?x\ dvd\ g$ **by** *auto*
from *x linear-primitive-irreducible*[*of ?x*] **have** *irr*: *irreducible ?x* **by** *auto*
from *irreducible-imp-square-free*[*OF this*] **have** *sfx*: *square-free ?x* **.**
from *irr fact* **have** *one*: $(fi,\ i) \in set\ fs \longrightarrow irreducible\ fi \wedge degree\ fi \neq 0$ **unfolding** *id* **by** (*auto simp*: *degree-monom-eq*)
have *fg*: $f = ?x\ \hat{}\ n * g$ **unfolding** *fg Suc* **by** (*metis x-pow-n*)
from *x* **have** *degx*: *degree ?x* = *1* **by** *simp*
note *sf* = *square-free-factorizationD*[*OF fact*(*1*)]
{
fix *a i*
assume *ai*: $(a,i) \in set\ hs$
with *sf*(*4*) **have** *g0*: $g \neq 0$ **by** *auto*
from *sf*(*2*)[*OF ai*] **have** *i*: $i \neq 0$ **by** *auto*
from *split-list*[*OF ai*] **obtain** *ys zs* **where** *hs*: $hs = ys\ @\ (a,i) \#\ zs$ **by** *auto*
have *a dvd g* **unfolding** *square-free-factorization-prod-list*[*OF fact*(*1*)] *hs*
by (*rule dvd-smult*, *insert i*, *cases i*, *auto simp add*: *ac-simps*)
moreover have $\neg\ ?x\ dvd\ g$ **using** *xs*[*unfolded Suc*] **by** *auto*
ultimately have *dvd*: $\neg\ ?x\ dvd\ a$ **using** *dvd-trans* **by** *blast*
from *sf*(*2*)[*OF ai*] **have** $a \neq 0$ **by** *auto*
have *1* = *gcd ?x a*
proof (*rule gcdI*)
fix *d*
assume *d*: *d dvd ?x d dvd a*
from *content-dvd-contentI*[*OF d*(*1*)] *x* **have** *cnt*: *is-unit* (*content d*) **by** *auto*
show *is-unit d*
proof (*cases degree d* = *1*)
case *False*
with *divides-degree*[*OF d*(*1*), *unfolded degx*] **have** *degree d* = *0* **by** *auto*
from *degree0-coeffs*[*OF this*] **obtain** *c* **where** *dc*: $d = [:c:]$ **by** *auto*
from *cnt*[*unfolded dc*] **have** *is-unit c* **by** (*auto simp*: *content-def*, *cases c* = *0*, *auto*)
hence $d * d = 1$ **unfolding** *dc* **by** (*auto*, *cases c* = -1; *cases c* = *1*, *auto*)
thus *is-unit d* **by** (*metis dvd-triv-right*)
next
case *True*
from *d*(*1*) **obtain** *e* **where** *xde*: $?x = d * e$ **unfolding** *dvd-def* **by** *auto*
from *arg-cong*[*OF this*, *of degree*] *degx* **have** *degree d* + *degree e* = *1*
by (*metis True add.right-neutral degree-0 degree-mult-eq one-neq-zero*)
with *True* **have** *degree e* = *0* **by** *auto*
from *degree0-coeffs*[*OF this*] *xde* **obtain** *e* **where** *xde*: $?x = [:e:] * d$ **by** *auto*
from *arg-cong*[*OF this*, *of content*, *unfolded content-mult*] *x*
have *content* $[:e:] *$ *content d* = *1* **by** *auto*
also have *content* $[:e\ :] = abs\ e$ **by** (*auto simp*: *content-def*, *cases e* = *0*, *auto*)

```
          finally have |e| * content d = 1 .
           from pos-zmult-eq-1-iff-lemma[OF this] have e * e = 1 by (cases e =
1; cases e = −1, auto)
          with arg-cong[OF xde, of smult e] have d = ?x * [:e:] by auto
          hence ?x dvd d unfolding dvd-def by blast
          with d(2) have ?x dvd a by (metis dvd-trans)
          with dvd show ?thesis by auto
        qed
      qed auto
      hence coprime ?x a
        by (simp add: gcd-eq-1-imp-coprime)
      note this dvd
    } note hs-dvd-x = this
    from hs-dvd-x[of ?x Suc m]
    have nmem: (?x,Suc m) ∉ set hs by auto
    hence eq: ?x ^ n * g = smult d (∏(a, i)∈set fs. a ^ i)
      unfolding sf(1) unfolding id Suc by simp
    have eq0: ?x ^ n * g = 0 ⟷ g = 0 by simp
  have square-free-factorization f (d,fs) unfolding fg id(1) square-free-factorization-def
split eq0 unfolding eq
    proof (intro conjI allI impI, rule refl)
      fix a i
      assume ai: (a,i) ∈ set fs
       thus square-free a degree a > 0 i > 0 using sf(2) sfx degx unfolding id
by auto
      fix b j
      assume bj: (b,j) ∈ set fs and diff: (a,i) ≠ (b,j)
      consider (hs-hs) (a,i) ∈ set hs (b,j) ∈ set hs
        | (hs-x) (a,i) ∈ set hs b = ?x
        | (x-hs) (b,j) ∈ set hs a = ?x
        using ai bj diff unfolding id by auto
      thus Rings.coprime a b
      proof cases
        case hs-hs
        from sf(3)[OF this diff] show ?thesis .
      next
        case hs-x
        from hs-dvd-x(1)[OF hs-x(1)] show ?thesis unfolding hs-x(2)
          by (simp add: ac-simps)
      next
        case x-hs
        from hs-dvd-x(1)[OF x-hs(1)] show ?thesis unfolding x-hs(2)
          by simp
      qed
    next
      show g = 0 ⟹ d = 0 using sf(4) by auto
      show g = 0 ⟹ fs = [] using sf(4) xs Suc by auto
      show distinct fs using sf(5) nmem unfolding id by auto
    qed
```

thus *?thesis* **using** *one* **unfolding** *id* **by** *auto*
qed
qed
qed
end

lift-definition *berlekamp-zassenhaus-factorization-algorithm* :: *int-poly-factorization-algorithm*
is *berlekamp-zassenhaus-factorization*
using *berlekamp-zassenhaus-factorization-irreducible*$_d$ **by** *blast*

abbreviation *factorize-int-poly* **where**
factorize-int-poly ≡ *factorize-int-poly-generic berlekamp-zassenhaus-factorization-algorithm*

end

11.4 Factoring Rational Polynomials

We combine the factorization algorithm for integer polynomials with Gauss Lemma to a factorization algorithm for rational polynomials.

theory *Factorize-Rat-Poly*
imports
Factorize-Int-Poly
begin

interpretation *content-hom*: *monoid-mult-hom*
content::′*a*::{*factorial-semiring*, *semiring-gcd*, *normalization-semidom-multiplicative*}
poly ⇒ -
by (*unfold-locales*, *auto simp*: *content-mult*)

lemma *prod-dvd-1-imp-all-dvd-1*:
assumes *finite X* **and** *prod f X dvd 1* **and** *x* ∈ *X* **shows** *f x dvd 1*
proof (*insert assms*, *induct rule*:*finite-induct*)
case *IH*: (*insert x′ X*)
show *?case*
proof (*cases x* = *x′*)
case *True*
with *IH* **show** *?thesis* **using** *dvd-trans*[*of f x′ f x′* ∗ - *1*]
by (*metis dvd-triv-left prod.insert*)
next
case *False*
then show *?thesis* **using** *IH* **by** (*auto intro*!: *IH*(*3*) *dvd-trans*[*of prod f X* - ∗
prod f X 1])
qed
qed *simp*

context
fixes *alg* :: *int-poly-factorization-algorithm*
begin

definition *factorize-rat-poly-generic* :: *rat poly* ⇒ *rat* × (*rat poly* × *nat*) *list* **where**
factorize-rat-poly-generic f = (*case rat-to-normalized-int-poly f of*
(*c*,*g*) ⇒ *case factorize-int-poly-generic alg g of* (*d*,*fs*) ⇒ (*c* ∗ *rat-of-int d*,
map (λ (*fi*,*i*). (*map-poly rat-of-int fi*, *i*)) *fs*))

lemma *factorize-rat-poly-0*[*simp*]: *factorize-rat-poly-generic 0* = (*0*,[])
unfolding *factorize-rat-poly-generic-def rat-to-normalized-int-poly-def* **by** *simp*

lemma *factorize-rat-poly*:
assumes *res*: *factorize-rat-poly-generic f* = (*c*,*fs*)
shows *square-free-factorization f* (*c*,*fs*)
and (*fi*,*i*) ∈ *set fs* ⟹ *irreducible fi*
proof(*atomize*(*full*), *cases f=0*, *goal-cases*)
case *1* **with** *res* **show** *?case* **by** (*auto simp*: *square-free-factorization-def*)
next
case *2* **show** *?case*
proof (*unfold square-free-factorization-def split*, *intro conjI impI allI*)
let *?r* = *rat-of-int*
let *?rp* = *map-poly ?r*
obtain *d g* **where** *ri*: *rat-to-normalized-int-poly f* = (*d*,*g*) **by** *force*
obtain *e gs* **where** *fi*: *factorize-int-poly-generic alg g* = (*e*,*gs*) **by** *force*
from *res*[*unfolded factorize-rat-poly-generic-def ri fi split*]
have *c*: *c* = *d* ∗ *?r e* **and** *fs*: *fs* = *map* (λ (*fi*,*i*). (*?rp fi*, *i*)) *gs* **by** *auto*
from *factorize-int-poly*[*OF fi*]
have *irr*: (*fi*, *i*) ∈ *set gs* ⟹ *irreducible fi* ∧ *content fi* = *1* **for** *fi i*
using *irreducible-imp-primitive*[*of fi*] **by** *auto*
note *sff* = *factorize-int-poly*(*1*)[*OF fi*]
note *sff′* = *square-free-factorizationD*[*OF sff*]
{
fix *n f*
have *?rp* (*f* ^ *n*) = (*?rp f*) ^ *n*
by (*induct n*, *auto simp*: *hom-distribs*)
} note *exp* = *this*
show *dist*: *distinct fs* **using** *sff′*(*5*) **unfolding** *fs distinct-map inj-on-def* **by** *auto*
interpret *mh*: *map-poly-inj-idom-hom rat-of-int*..
have *f* = *smult d* (*?rp g*) **using** *rat-to-normalized-int-poly*[*OF ri*] **by** *auto*
also have ... = *smult d* (*?rp* (*smult e* (∏ (*a*, *i*)∈*set gs*. *a* ^ *i*))) **using** *sff′*(*1*) **by** *simp*
also have ... = *smult c* (*?rp* (∏ (*a*, *i*)∈*set gs*. *a* ^ *i*)) **unfolding** *c* **by** (*simp add*: *hom-distribs*)
also have *?rp* (∏ (*a*, *i*)∈*set gs*. *a* ^ *i*) = (∏ (*a*, *i*)∈*set fs*. *a* ^ *i*)
unfolding *prod.distinct-set-conv-list*[*OF sff′*(*5*)] *prod.distinct-set-conv-list*[*OF dist*]
unfolding *fs*
by (*insert exp*, *auto intro*!: *arg-cong*[*of* - - λ*x*. *prod-list* (*map x gs*)] *simp*: *hom-distribs of-int-poly-hom.hom-prod-list*)
finally show *f*: *f* = *smult c* (∏ (*a*, *i*)∈*set fs*. *a* ^ *i*) **by** *auto*
{

fix *a i*
assume *ai*: (*a*,*i*) ∈ *set fs*
from *ai* **obtain** *A* **where** *a*: *a* = *?rp A* **and** *A*: (*A*,*i*) ∈ *set gs* **unfolding** *fs*
by *auto*
fix *b j*
assume (*b*,*j*) ∈ *set fs* **and** *diff*: (*a*,*i*) ≠ (*b*,*j*)
from *this*(*1*) **obtain** *B* **where** *b*: *b* = *?rp B* **and** *B*: (*B*,*j*) ∈ *set gs* **unfolding**
fs **by** *auto*
from *diff*[*unfolded a b*] **have** (*A*,*i*) ≠ (*B*,*j*) **by** *auto*
from *sff'*(*3*)[*OF A B this*]
show *Rings.coprime a b*
by (*auto simp add*: *coprime-iff-gcd-eq-1 gcd-rat-to-gcd-int a b*)
}
{
fix *fi i*
assume (*fi*,*i*) ∈ *set fs*
then obtain *gi* **where** *fi*: *fi* = *?rp gi* **and** *gi*: (*gi*,*i*) ∈ *set gs* **unfolding** *fs*
by *auto*
from *sff'*(*2*)[*OF gi*] **show** *0* < *i* **by** *auto*
from *irr*[*OF gi*] **have** *cf-gi*: *primitive gi* **by** *auto*
then have *primitive* (*?rp gi*) **by** (*auto simp*: *content-field-poly*)
note [*simp*] = *irreducible-primitive-connect*[*OF cf-gi*] *irreducible-primitive-connect*[*OF*
this]
show *irreducible fi*
using *irr*[*OF gi*] *fi irreducible*$_d$*-int-rat*[*of gi*,*simplified*] **by** *auto*
then show *degree fi* > *0 square-free fi* **unfolding** *fi*
by (*auto intro*: *irreducible-imp-square-free*)
}
{
assume *f* = *0* **with** *ri* **have** ∗: *d* = *1 g* = *0* **unfolding** *rat-to-normalized-int-poly-def*
by *auto*
with *sff'*(*4*)[*OF* ∗(*2*)] **show** *c* = *0 fs* = [] **unfolding** *c fs* **by** *auto*
}
qed
qed

end

abbreviation *factorize-rat-poly* **where**
factorize-rat-poly ≡ *factorize-rat-poly-generic berlekamp-zassenhaus-factorization-algorithm*

end

12 External Interface

We provide two functions for external usage that work on lists and integers only, so that they can easily be accessed via these primitive datatypes.

theory *Factorization-External-Interface*
imports
Factorize-Rat-Poly
Factorize-Int-Poly
begin

declare *Lcm-fin.set-eq-fold*[*code-unfold*]

definition *factor-int-poly* :: *integer list* ⇒ *integer* × (*integer list* × *integer*) *list*
where
factor-int-poly p = *map-prod integer-of-int* (*map* (*map-prod* (*map integer-of-int o coeffs*) *integer-of-nat*))
(*factorize-int-poly* (*poly-of-list* (*map int-of-integer p*)))

Just for clarifying the representation, we present a part of the soundness statement of the factorization algorithm with conversions included

lemma *factor-int-poly*: **assumes** *factor-int-poly p* = (*c*, *qes*)
shows *poly-of-list* (*map int-of-integer p*) = *smult* (*int-of-integer c*)
(∏ (*q*, *e*)←*qes*. *poly-of-list* (*map int-of-integer q*) ^ *nat-of-integer e*)
(**is** *?p* = *?prod*)
proof −
obtain *C Qes* **where** *fact*: *factorize-int-poly ?p* = (*C*,*Qes*) **by** *force*
from *square-free-factorization-prod-list*[*OF factorize-int-poly*(*1*)[*OF this*]]
have *?p* = *smult C* (∏ (*x*, *y*)←*Qes*. *x* ^ *y*) **.**
also have ... = *?prod* **using** *assms*[*unfolded factor-int-poly-def fact*, *symmetric*]
by (*intro arg-cong2*[*of* - - - - λ *x y*. *smult x* (*prod-list y*)], *auto simp*: *o-def*)
finally show *?thesis* **.**
qed

Note that coefficients are listed with lowest coefficient as head of list

value *coeffs* (*monom 1 3*) :: *int list*
value *factor-int-poly* [*0*,*0*,*0*,*5*]
value *factor-int-poly* [*0*,*1*,−*2*,*1*]

definition *integers-of-rat* **where** *integers-of-rat x* = *map-prod integer-of-int integer-of-int* (*quotient-of x*)
fun *rat-of-integers* **where** *rat-of-integers* (*n*,*d*) = (*rat-of-int* (*int-of-integer n*) / *rat-of-int* (*int-of-integer d*))

definition *integer-of-rat* **where** *integer-of-rat x* = *integer-of-int* (*fst* (*quotient-of x*))
definition *rat-of-integer* **where** *rat-of-integer x* = *rat-of-int* (*int-of-integer x*)

lemma *integers-of-rat*[*simp*]: *rat-of-integers* (*integers-of-rat x*) = *x*
proof −
obtain *n d* **where** *id*: *quotient-of x* = (*n*,*d*) **by** *force*
from *quotient-of-div*[*OF id*]
show *?thesis* **unfolding** *integers-of-rat-def id* **by** *auto*

qed

lemma *integer-of-rat*[*simp*]: **assumes** $x \in \mathbb{Z}$
shows *rat-of-integer* (*integer-of-rat x*) = *x*
proof −
from *assms* **obtain** *y* **where** *x*: *x* = *of-int y* **unfolding** *Ints-def* **by** *auto*
hence *id*: *quotient-of x* = (*y*,*1*) **by** *simp*
from *quotient-of-div*[*OF id*]
show *?thesis* **unfolding** *integer-of-rat-def rat-of-integer-def id* **by** *auto*
qed

definition *factor-rat-poly* :: (*integer* × *integer*) *list* ⇒ (*integer* × *integer*) × (*integer list* × *integer*) *list* **where**
factor-rat-poly p = *map-prod integers-of-rat* (*map* (*map-prod* (*map integer-of-rat o coeffs*) *integer-of-nat*))
(*factorize-rat-poly* (*poly-of-list* (*map rat-of-integers p*)))

lemma *factor-rat-poly*: **assumes** *factor-rat-poly p* = (*c*, *qes*)
shows *poly-of-list* (*map rat-of-integers p*) = *smult* (*rat-of-integers c*)
(∏ (*q*, *e*)←*qes*. *poly-of-list* (*map rat-of-integer q*) ^ *nat-of-integer e*)
(**is** *?p* = *?prod*)
proof −
obtain *C Qes* **where** *fact*: *factorize-rat-poly ?p* = (*C*,*Qes*) **by** *force*
{
fix *a b*
assume *ab*: (*a*,*b*) ∈ *set Qes*
with *fact*[*unfolded factorize-rat-poly-generic-def*] **have** *a*: *a* ∈ *range of-int-poly*
by (*auto split*: *prod.splits*)
have *map* (λ*x*. *rat-of-integer* (*integer-of-rat x*)) (*coeffs a*) = *map* (λ *x*. *x*) (*coeffs a*)
by (*intro map-cong*[*OF refl integer-of-rat*], *insert a*, *force*)
hence *Poly* (*map* (λ*x*. *rat-of-integer* (*integer-of-rat x*)) (*coeffs a*)) = *a* **by** *simp*
} note *eq* = *this*
from *square-free-factorization-prod-list*[*OF factorize-rat-poly*(*1*)[*OF fact*]]
have *?p* = *smult C* (∏ (*x*, *y*)←*Qes*. *x* ^ *y*) **.**
also have ... = *?prod* **using** *assms*[*unfolded factor-rat-poly-def fact*, *symmetric*]
apply (*intro arg-cong2*[*of* - - - - λ *x y*. *smult x* (*prod-list y*)])
subgoal by *simp*
subgoal using *eq* **by** (*auto simp add*: *o-def intro*!: *arg-cong*[*of* - - λ *x*. *x* ^ -])
done
finally show *?thesis* **.**
qed

Note that rational numbers in the input are encoded as pairs, whereas the polynomials in the output are just integer polynomials, i.e., only the constant factor is a rational number

value *factor-rat-poly* [(*1*,*6*),(−*1*,*3*),(*1*,*6*)]

end

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