

Banach-Steinhaus theorem

Dominique Unruh José Manuel Rodríguez Caballero

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Abstract

We formalize in Isabelle/HOL a result [2] due to S. Banach and H. Steinhaus [1] known as Banach-Steinhaus theorem or Uniform boundedness principle: a pointwise-bounded family of continuous linear operators from a Banach space to a normed space is uniformly bounded. Our approach is an adaptation to Isabelle/HOL of a proof due to A. Sokal [3].

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1 Missing results for the proof of Banach-Steinhaus theorem

```
theory Banach-Steinhaus-Missing
  imports
    HOL-Analysis.Bounded-Linear-Function
    HOL-Analysis.Line-Segment
```

begin

1.1 Results missing for the proof of Banach-Steinhaus theorem

The results proved here are preliminaries for the proof of Banach-Steinhaus theorem using Sokal’s approach, but they do not explicitly appear in Sokal’s paper [?].

Notation for the norm

```
bundle notation-norm begin  
notation norm ( $\|-\|$ )  
end
```

```
bundle no-notation-norm begin  
no-notation norm ( $\|-\|$ )  
end
```

```
unbundle notation-norm
```

Notation for apply bilinear function

```
bundle notation-blinfun-apply begin  
notation blinfun-apply (infixr  $*_v$  70)  
end
```

```
bundle no-notation-blinfun-apply begin  
no-notation blinfun-apply (infixr  $*_v$  70)  
end
```

```
unbundle notation-blinfun-apply
```

```
lemma bdd-above-plus:
```

```
  fixes f:: $\langle 'a \Rightarrow \text{real} \rangle$ 
```

```
  assumes  $\langle \text{bdd-above } (f \text{ ' } S) \rangle$  and  $\langle \text{bdd-above } (g \text{ ' } S) \rangle$ 
```

```
  shows  $\langle \text{bdd-above } ((\lambda x. f x + g x) \text{ ' } S) \rangle$ 
```

Explanation: If the images of two real-valued functions f, g are bounded above on a set S , then the image of their sum is bounded on S .

```
proof -
```

```
  obtain M where  $\langle \bigwedge x. x \in S \implies f x \leq M \rangle$ 
```

```
    using  $\langle \text{bdd-above } (f \text{ ' } S) \rangle$  unfolding bdd-above-def by blast
```

```
  obtain N where  $\langle \bigwedge x. x \in S \implies g x \leq N \rangle$ 
```

```
    using  $\langle \text{bdd-above } (g \text{ ' } S) \rangle$  unfolding bdd-above-def by blast
```

```
  have  $\langle \bigwedge x. x \in S \implies f x + g x \leq M + N \rangle$ 
```

```
    using  $\langle \bigwedge x. x \in S \implies f x \leq M \rangle$   $\langle \bigwedge x. x \in S \implies g x \leq N \rangle$  by fastforce
```

```
  thus ?thesis unfolding bdd-above-def by blast
```

```
qed
```

The maximum of two functions

```
definition pointwise-max::  $\langle 'a \Rightarrow 'b::\text{ord} \rangle \Rightarrow \langle 'a \Rightarrow 'b \rangle \Rightarrow \langle 'a \Rightarrow 'b \rangle$  where  
 $\langle \text{pointwise-max } f g = (\lambda x. \max (f x) (g x)) \rangle$ 
```

```
lemma max-Sup-absorb-left:
```

```
  fixes f g:: $\langle 'a \Rightarrow \text{real} \rangle$ 
```

```
  assumes  $\langle X \neq \{\} \rangle$  and  $\langle \text{bdd-above } (f \text{ ' } X) \rangle$  and  $\langle \text{bdd-above } (g \text{ ' } X) \rangle$  and  $\langle \text{Sup } (f \text{ ' } X) \geq \text{Sup } (g \text{ ' } X) \rangle$ 
```

```
  shows  $\langle \text{Sup } ((\text{pointwise-max } f g) \text{ ' } X) = \text{Sup } (f \text{ ' } X) \rangle$ 
```

Explanation: For real-valued functions f and g , if the supremum of f is greater-equal the supremum of g , then the supremum of $\max f g$ equals the supremum of f . (Under some technical conditions.)

proof –

have $y\text{-Sup}$: $\langle y \in ((\lambda x. \max (f x) (g x)) \text{ ' } X) \implies y \leq \text{Sup} (f \text{ ' } X) \rangle$ **for** y

proof –

assume $\langle y \in ((\lambda x. \max (f x) (g x)) \text{ ' } X) \rangle$

then obtain x **where** $\langle y = \max (f x) (g x) \rangle$ **and** $\langle x \in X \rangle$

by *blast*

have $\langle f x \leq \text{Sup} (f \text{ ' } X) \rangle$

by (*simp add*: $\langle x \in X \rangle$ $\langle \text{bdd-above} (f \text{ ' } X) \rangle$ *cSUP-upper*)

moreover have $\langle g x \leq \text{Sup} (g \text{ ' } X) \rangle$

by (*simp add*: $\langle x \in X \rangle$ $\langle \text{bdd-above} (g \text{ ' } X) \rangle$ *cSUP-upper*)

ultimately have $\langle \max (f x) (g x) \leq \text{Sup} (f \text{ ' } X) \rangle$

using $\langle \text{Sup} (f \text{ ' } X) \geq \text{Sup} (g \text{ ' } X) \rangle$ **by** *auto*

thus *?thesis* **by** (*simp add*: $\langle y = \max (f x) (g x) \rangle$)

qed

have $y\text{-}f\text{-}X$: $\langle y \in f \text{ ' } X \implies y \leq \text{Sup} ((\lambda x. \max (f x) (g x)) \text{ ' } X) \rangle$ **for** y

proof –

assume $\langle y \in f \text{ ' } X \rangle$

then obtain x **where** $\langle x \in X \rangle$ **and** $\langle y = f x \rangle$

by *blast*

have $\langle \text{bdd-above} ((\lambda \xi. \max (f \xi) (g \xi)) \text{ ' } X) \rangle$

by (*metis* (*no-types*) $\langle \text{bdd-above} (f \text{ ' } X) \rangle$ $\langle \text{bdd-above} (g \text{ ' } X) \rangle$ *bdd-above-image-sup sup-max*)

moreover have $\langle e > 0 \implies \exists k \in (\lambda \xi. \max (f \xi) (g \xi)) \text{ ' } X. y \leq k + e \rangle$

for $e::\text{real}$

using $\langle \text{Sup} (f \text{ ' } X) \geq \text{Sup} (g \text{ ' } X) \rangle$ **by** (*smt* $\langle x \in X \rangle$ $\langle y = f x \rangle$ *image-eqI*)

ultimately show *?thesis*

using $\langle x \in X \rangle$ $\langle y = f x \rangle$ *cSUP-upper* **by** *fastforce*

qed

have $\langle \text{Sup} ((\lambda x. \max (f x) (g x)) \text{ ' } X) \leq \text{Sup} (f \text{ ' } X) \rangle$

using $y\text{-}Sup$ **by** (*simp add*: $\langle X \neq \{\} \rangle$ *cSup-least*)

moreover have $\langle \text{Sup} ((\lambda x. \max (f x) (g x)) \text{ ' } X) \geq \text{Sup} (f \text{ ' } X) \rangle$

using $y\text{-}f\text{-}X$ **by** (*metis* (*mono-tags*) *cSup-least calculation empty-is-image*)

ultimately show *?thesis* **unfolding** *pointwise-max-def* **by** *simp*

qed

lemma *max-Sup-absorb-right*:

fixes $f g::\text{'a} \Rightarrow \text{real}$

assumes $\langle X \neq \{\} \rangle$ **and** $\langle \text{bdd-above} (f \text{ ' } X) \rangle$ **and** $\langle \text{bdd-above} (g \text{ ' } X) \rangle$ **and** $\langle \text{Sup} (f \text{ ' } X) \leq \text{Sup} (g \text{ ' } X) \rangle$

shows $\langle \text{Sup} ((\text{pointwise-max } f g) \text{ ' } X) = \text{Sup} (g \text{ ' } X) \rangle$

Explanation: For real-valued functions f and g and a nonempty set X , such that the f and g are bounded above on X , if the supremum of f on X is lower-equal the supremum of g on X , then the supremum of *pointwise-max* $f g$ on X equals the supremum of g . This is the right analog of

max-Sup-absorb-left.

proof –

have $\langle \text{Sup } ((\text{pointwise-max } g \ f) \ 'X) = \text{Sup } (g \ 'X) \rangle$
using *assms* **by** (*simp add: max-Sup-absorb-left*)
moreover have $\langle \text{pointwise-max } g \ f = \text{pointwise-max } f \ g \rangle$
unfolding *pointwise-max-def* **by** *auto*
ultimately show *?thesis* **by** *simp*
qed

lemma *max-Sup*:

fixes $f \ g :: 'a \Rightarrow \text{real}$
assumes $\langle X \neq \{\} \rangle$ **and** $\langle \text{bdd-above } (f \ 'X) \rangle$ **and** $\langle \text{bdd-above } (g \ 'X) \rangle$
shows $\langle \text{Sup } ((\text{pointwise-max } f \ g) \ 'X) = \max (\text{Sup } (f \ 'X)) (\text{Sup } (g \ 'X)) \rangle$

Explanation: Let X be a nonempty set. Two supremum over X of the maximum of two real-value functions is equal to the maximum of their suprema over X , provided that the functions are bounded above on X .

proof(*cases* $\langle \text{Sup } (f \ 'X) \geq \text{Sup } (g \ 'X) \rangle$)

case *True* **thus** *?thesis* **by** (*simp add: assms(1) assms(2) assms(3) max-Sup-absorb-left*)

next

case *False*

have $f1: \neg 0 \leq \text{Sup } (f \ 'X) + - 1 * \text{Sup } (g \ 'X)$

using *False* **by** *linarith*

hence $\text{Sup } (\text{Banach-Steinhaus-Missing.pointwise-max } f \ g \ 'X) = \text{Sup } (g \ 'X)$

by (*simp add: assms(1) assms(2) assms(3) max-Sup-absorb-right*)

thus *?thesis*

using *f1* **by** *linarith*

qed

lemma *identity-telescopic*:

fixes $x :: \langle - \Rightarrow 'a :: \text{real-normed-vector} \rangle$

assumes $\langle x \longrightarrow l \rangle$

shows $\langle (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\}) \longrightarrow l - x n \rangle$

Expression of a limit as a telescopic series. Explanation: If x converges to l then the sum $\sum_{k=n..N} x (\text{Suc } k) - x k$ converges to $l - x n$ as N goes to infinity.

proof –

have $\langle (\lambda p. x (p + \text{Suc } n)) \longrightarrow l \rangle$

using $\langle x \longrightarrow l \rangle$ **by** (*rule LIMSEQ-ignore-initial-segment*)

hence $\langle (\lambda p. x (\text{Suc } n + p)) \longrightarrow l \rangle$

by (*simp add: add commute*)

hence $\langle (\lambda p. x (\text{Suc } (n + p))) \longrightarrow l \rangle$

by *simp*

hence $\langle (\lambda t. (- (x n)) + (\lambda p. x (\text{Suc } (n + p))) t) \longrightarrow (- (x n)) + l \rangle$

using *tendsto-add-const-iff* **by** *metis*

hence $f1: \langle (\lambda p. x (\text{Suc } (n + p)) - x n) \longrightarrow l - x n \rangle$

by *simp*
have $\langle \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..n+p\} = x (\text{Suc } (n+p)) - x n \rangle$ **for** p
by (*simp add: sum-Suc-diff*)
moreover have $\langle (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\}) (n + t)$
 $= (\lambda p. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..n+p\}) t \rangle$ **for** t
by *blast*
ultimately have $\langle (\lambda p. (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\}) (n + p))$
 $\longrightarrow l - x n \rangle$
using *f1 by simp*
hence $\langle (\lambda p. (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\}) (p + n)) \longrightarrow l - x$
 $n \rangle$
by (*simp add: add.commute*)
hence $\langle (\lambda p. (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\}) p) \longrightarrow l - x n \rangle$
using *Topological-Spaces.LIMSEQ-offset* [**where** $f = (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\})$
 $k) - x k) \{n..N\}$
and $a = l - x n$ **and** $k = n$] **by** *blast*
hence $\langle (\lambda M. (\lambda N. \text{sum } (\lambda k. x (\text{Suc } k) - x k) \{n..N\}) M) \longrightarrow l - x n \rangle$
by *simp*
thus *?thesis by blast*
qed

lemma *bound-Cauchy-to-lim:*

assumes $\langle y \longrightarrow x \rangle$ **and** $\langle \bigwedge n. \|y (\text{Suc } n) - y n\| \leq c \hat{\ } n \rangle$ **and** $\langle y 0 = 0 \rangle$ **and**
 $\langle c < 1 \rangle$
shows $\langle \|x - y (\text{Suc } n)\| \leq (c / (1 - c)) * c \hat{\ } n \rangle$

Inequality about a sequence of approximations assuming that the sequence of differences is bounded by a geometric progression. Explanation: Let y be a sequence converging to x . If y satisfies the inequality $\|y (\text{Suc } n) - y n\| \leq c \hat{\ } n$ for some $c < 1$ and assuming $y 0 = (0::'a)$ then the inequality $\|x - y (\text{Suc } n)\| \leq (c / (1 - c)) * c \hat{\ } n$ holds.

proof –

have $\langle c \geq 0 \rangle$
using $\langle \bigwedge n. \|y (\text{Suc } n) - y n\| \leq c \hat{\ } n \rangle$ **by** (*smt norm-imp-pos-and-ge power-Suc0-right*)
have *norm-1*: $\langle \text{norm } (\sum k = \text{Suc } n..N. y (\text{Suc } k) - y k) \leq (c \hat{\ } \text{Suc } n) / (1 -$
 $c) \rangle$ **for** N
proof (*cases* $\langle N < \text{Suc } n \rangle$)
case *True*
hence $\langle \|\text{sum } (\lambda k. y (\text{Suc } k) - y k) \{\text{Suc } n .. N\}\| = 0 \rangle$
by *auto*
thus *?thesis using* $\langle c \geq 0 \rangle \langle c < 1 \rangle$ **by** *auto*
next
case *False*
hence $\langle N \geq \text{Suc } n \rangle$
by *auto*
have $\langle c \hat{\ } (\text{Suc } N) \geq 0 \rangle$
using $\langle c \geq 0 \rangle$ **by** *auto*
have $\langle 1 - c > 0 \rangle$
by (*simp add:* $\langle c < 1 \rangle$)

hence $\langle (1 - c)/(1 - c) = 1 \rangle$
by *auto*
have $\langle \| \text{sum } (\lambda k. y (\text{Suc } k) - y k) \{ \text{Suc } n .. N \} \| \leq (\text{sum } (\lambda k. \| y (\text{Suc } k) - y k \|) \{ \text{Suc } n .. N \}) \rangle$
by *(simp add: sum-norm-le)*
hence $\langle \| \text{sum } (\lambda k. y (\text{Suc } k) - y k) \{ \text{Suc } n .. N \} \| \leq (\text{sum } (\text{power } c) \{ \text{Suc } n .. N \}) \rangle$
by *(simp add: assms(2) sum-norm-le)*
hence $\langle (1 - c) * \| \text{sum } (\lambda k. y (\text{Suc } k) - y k) \{ \text{Suc } n .. N \} \| \leq (1 - c) * (\text{sum } (\text{power } c) \{ \text{Suc } n .. N \}) \rangle$
using $\langle 0 < 1 - c \rangle$ *mult-le-cancel-iff2* **by** *blast*
also have $\langle \dots = c^{\wedge}(\text{Suc } n) - c^{\wedge}(\text{Suc } N) \rangle$
using *Set-Interval.sum-gp-multiplied* $\langle \text{Suc } n \leq N \rangle$ **by** *blast*
also have $\langle \dots \leq c^{\wedge}(\text{Suc } n) \rangle$
using $\langle c^{\wedge}(\text{Suc } N) \geq 0 \rangle$ **by** *auto*
finally have $\langle (1 - c) * \| \sum k = \text{Suc } n..N. y (\text{Suc } k) - y k \| \leq c^{\wedge} \text{Suc } n \rangle$
by *blast*
hence $\langle ((1 - c) * \| \sum k = \text{Suc } n..N. y (\text{Suc } k) - y k \|) / (1 - c) \leq (c^{\wedge} \text{Suc } n) / (1 - c) \rangle$
using $\langle 0 < 1 - c \rangle$ **by** *(smt divide-right-mono)*
thus $\langle \| \sum k = \text{Suc } n..N. y (\text{Suc } k) - y k \| \leq (c^{\wedge} \text{Suc } n) / (1 - c) \rangle$
using $\langle 0 < 1 - c \rangle$ **by** *auto*
qed
have $\langle (\lambda N. (\text{sum } (\lambda k. y (\text{Suc } k) - y k) \{ \text{Suc } n .. N \})) \longrightarrow x - y (\text{Suc } n) \rangle$
by *(metis (no-types) <y >> x identity-telescopic)*
hence $\langle (\lambda N. \| \text{sum } (\lambda k. y (\text{Suc } k) - y k) \{ \text{Suc } n .. N \} \|) \longrightarrow \| x - y (\text{Suc } n) \| \rangle$
using *tendsto-norm* **by** *blast*
hence $\langle \| x - y (\text{Suc } n) \| \leq (c^{\wedge} \text{Suc } n) / (1 - c) \rangle$
using *norm-1 Lim-bounded* **by** *blast*
hence $\langle \| x - y (\text{Suc } n) \| \leq (c^{\wedge} \text{Suc } n) / (1 - c) \rangle$
by *auto*
moreover have $\langle (c^{\wedge} \text{Suc } n) / (1 - c) = (c / (1 - c)) * (c^{\wedge} n) \rangle$
by *(simp add: divide-inverse-commute)*
ultimately show $\langle \| x - y (\text{Suc } n) \| \leq (c / (1 - c)) * (c^{\wedge} n) \rangle$ **by** *linarith*
qed

lemma *onorm-open-ball:*

includes *notation-norm*

shows $\langle \| f \| = \text{Sup } \{ \| f *_v x \| \mid x. \| x \| < 1 \} \rangle$

Explanation: Let f be a bounded linear operator. The operator norm of f is the supremum of $\| f *_v x \|$ for x such that $\| x \| < 1$.

proof *(cases <(UNIV::'a set) = 0>)*

case *True*

hence $\langle x = 0 \rangle$ **for** $x::'a$

by *auto*

hence $\langle f *_v x = 0 \rangle$ **for** x

by *(metis (full-types) blinfun.zero-right)*

```

hence ⟨||f|| = 0⟩
  by (simp add: blinfun-eqI zero-blinfun.rep-eq)
have ⟨{ ||f *v x|| | x. ||x|| < 1 } = {0}⟩
  by (smt Collect-cong ⟨∧x. f *v x = 0⟩ norm-zero singleton-conv)
hence ⟨Sup { ||f *v x|| | x. ||x|| < 1 } = 0⟩
  by simp
thus ?thesis using ⟨||f|| = 0⟩ by auto
next
case False
hence ⟨(UNIV::'a set) ≠ 0⟩
  by simp
have nonnegative: ⟨||f *v x|| ≥ 0⟩ for x
  by simp
have ⟨∃ x::'a. x ≠ 0⟩
  using ⟨UNIV ≠ 0⟩ by auto
then obtain x::'a where ⟨x ≠ 0⟩
  by blast
hence ⟨||x|| ≠ 0⟩
  by auto
define y where ⟨y = x /R ||x||⟩
have ⟨norm y = || x /R ||x|| ||⟩
  unfolding y-def by auto
also have ⟨... = ||x|| /R ||x||⟩
  by auto
also have ⟨... = 1⟩
  using ⟨||x|| ≠ 0⟩ by auto
finally have ⟨||y|| = 1⟩
  by blast
hence norm-1-non-empty: ⟨{ ||f *v x|| | x. ||x|| = 1 } ≠ {}⟩
  by blast
have norm-1-bounded: ⟨bdd-above { ||f *v x|| | x. ||x|| = 1 }⟩
  unfolding bdd-above-def apply auto
  by (metis norm-blinfun)
have norm-less-1-non-empty: ⟨{ ||f *v x|| | x. ||x|| < 1 } ≠ {}⟩
  by (metis (mono-tags, lifting) Collect-empty-eq-bot bot-empty-eq empty-iff norm-zero
    zero-less-one)
have norm-less-1-bounded: ⟨bdd-above { ||f *v x|| | x. ||x|| < 1 }⟩
proof -
  have ⟨∃ r. ||a r|| < 1 ⟶ ||f *v (a r)|| ≤ r⟩ for a :: real ⇒ 'a
  proof -
    obtain r :: ('a ⇒L 'b) ⇒ real where
      ∧f x. 0 ≤ r f ∧ (bounded-linear f ⟶ ||f *v x|| ≤ ||x|| * r f)
      using bounded-linear.nonneg-bounded by moura
    have ⟨¬ ||f|| < 0⟩
      by simp
    hence ⟨∃ r. ||f|| * ||a r|| ≤ r⟩ ∨ ⟨∃ r. ||a r|| < 1 ⟶ ||f *v a r|| ≤ r⟩
      by (meson less-eq-real-def mult-le-cancel-left2)
    thus ?thesis using dual-order.trans norm-blinfun by blast

```

qed
hence $\langle \exists M. \forall x. \|x\| < 1 \longrightarrow \|f *_v x\| \leq M \rangle$
by *metis*
thus *?thesis* **by** *auto*
qed
have *Sup-non-neg*: $\langle \text{Sup } \{\|f *_v x\| \mid x. \|x\| = 1\} \geq 0 \rangle$
by (*smt Collect-empty-eq cSup-upper mem-Collect-eq nonnegative norm-1-bounded norm-1-non-empty*)
have $\langle \{0::\text{real}\} \neq \{\} \rangle$
by *simp*
have $\langle \text{bdd-above } \{0::\text{real}\} \rangle$
by *simp*
show $\langle \|f\| = \text{Sup } \{\|f *_v x\| \mid x. \|x\| < 1\} \rangle$
proof(*cases* $\langle \forall x. f *_v x = 0 \rangle$)
case *True*
have $\langle \|f *_v x\| = 0 \rangle$ **for** x
by (*simp add: True*)
hence $\langle \{\|f *_v x\| \mid x. \|x\| < 1\} \subseteq \{0\} \rangle$
by *blast*
moreover **have** $\langle \{\|f *_v x\| \mid x. \|x\| < 1\} \supseteq \{0\} \rangle$
using *calculation norm-less-1-non-empty* **by** *fastforce*
ultimately **have** $\langle \{\|f *_v x\| \mid x. \|x\| < 1\} = \{0\} \rangle$
by *blast*
hence *Sup1*: $\langle \text{Sup } \{\|f *_v x\| \mid x. \|x\| < 1\} = 0 \rangle$
by *simp*
have $\langle \|f\| = 0 \rangle$
by (*simp add: True blinfun-eqI*)
moreover **have** $\langle \text{Sup } \{\|f *_v x\| \mid x. \|x\| < 1\} = 0 \rangle$
using *Sup1* **by** *blast*
ultimately **show** *?thesis* **by** *simp*
next
case *False*
have *norm-f-eq-leq*: $\langle y \in \{\|f *_v x\| \mid x. \|x\| = 1\} \implies y \leq \text{Sup } \{\|f *_v x\| \mid x. \|x\| < 1\} \rangle$ **for** y
proof–
assume $\langle y \in \{\|f *_v x\| \mid x. \|x\| = 1\} \rangle$
hence $\langle \exists x. y = \|f *_v x\| \wedge \|x\| = 1 \rangle$
by *blast*
then **obtain** x **where** $\langle y = \|f *_v x\| \rangle$ **and** $\langle \|x\| = 1 \rangle$
by *auto*
define y' **where** $\langle y' n = (1 - (\text{inverse } (\text{real } (\text{Suc } n)))) *_R y \rangle$ **for** n
have $\langle y' n \in \{\|f *_v x\| \mid x. \|x\| < 1\} \rangle$ **for** n
proof–
have $\langle y' n = (1 - (\text{inverse } (\text{real } (\text{Suc } n)))) *_R \|f *_v x\| \rangle$
using *y'-def* $\langle y = \|f *_v x\| \rangle$ **by** *blast*
also **have** $\langle \dots = |(1 - (\text{inverse } (\text{real } (\text{Suc } n))))| *_R \|f *_v x\| \rangle$
by (*metis (mono-tags, opaque-lifting)* $\langle y = \|f *_v x\| \rangle$ *abs-1 abs-le-self-iff*)
abs-of-nat
abs-of-nonneg add-diff-cancel-left' add-eq-if cancel-comm-monoid-add-class.diff-cancel

$\text{diff-ge-0-iff-ge eq-iff-diff-eq-0 inverse-1 inverse-le-iff-le nat.distinct}(1)$
of-nat-0
of-nat-Suc of-nat-le-0-iff zero-less-abs-iff zero-neq-one
also have $\langle \dots = \|f *_v ((1 - (\text{inverse} (\text{real} (\text{Suc } n)))) *_R x)\| \rangle$
by (*simp add: blinfun.scaleR-right*)
finally have $y'-1: \langle y' n = \|f *_v ((1 - (\text{inverse} (\text{real} (\text{Suc } n)))) *_R x)\| \rangle$
by *blast*
have $\langle \|(1 - (\text{inverse} (\text{Suc } n))) *_R x\| = (1 - (\text{inverse} (\text{real} (\text{Suc } n)))) * \|x\| \rangle$
by (*simp add: linordered-field-class.inverse-le-1-iff*)
hence $\langle \|(1 - (\text{inverse} (\text{Suc } n))) *_R x\| < 1 \rangle$
by (*simp add: \langle \|x\| = 1 \rangle*)
thus *?thesis using y'-1 by blast*
qed
have $\langle (\lambda n. (1 - (\text{inverse} (\text{real} (\text{Suc } n))))) \longrightarrow 1 \rangle$
using *Limits.LIMSEQ-inverse-real-of-nat-add-minus* **by** *simp*
hence $\langle (\lambda n. (1 - (\text{inverse} (\text{real} (\text{Suc } n)))) *_R y \longrightarrow 1 *_R y \rangle$
using *Limits.tendsto-scaleR* **by** *blast*
hence $\langle (\lambda n. (1 - (\text{inverse} (\text{real} (\text{Suc } n)))) *_R y \longrightarrow y \rangle$
by *simp*
hence $\langle (\lambda n. y' n) \longrightarrow y \rangle$
using *y'-def* **by** *simp*
hence $\langle y' \longrightarrow y \rangle$
by *simp*
have $\langle y' n \leq \text{Sup} \{ \|f *_v x\| \mid x. \|x\| < 1 \} \rangle$ **for** n
using *cSup-upper* $\langle \wedge n. y' n \in \{ \|f *_v x\| \mid x. \|x\| < 1 \} \rangle$ *norm-less-1-bounded*
by *blast*
hence $\langle y \leq \text{Sup} \{ \|f *_v x\| \mid x. \|x\| < 1 \} \rangle$
using $\langle y' \longrightarrow y \rangle$ *Topological-Spaces.Sup-lim* **by** (*meson LIMSEQ-le-const2*)
thus *?thesis by blast*
qed
hence $\langle \text{Sup} \{ \|f *_v x\| \mid x. \|x\| = 1 \} \leq \text{Sup} \{ \|f *_v x\| \mid x. \|x\| < 1 \} \rangle$
by (*metis (lifting) cSup-least norm-1-non-empty*)
have $\langle y \in \{ \|f *_v x\| \mid x. \|x\| < 1 \} \implies y \leq \text{Sup} \{ \|f *_v x\| \mid x. \|x\| = 1 \} \rangle$ **for** y
proof (*cases \langle y = 0 \rangle*)
case *True* **thus** *?thesis by (simp add: Sup-non-neg)*
next
case *False*
hence $\langle y \neq 0 \rangle$ **by** *blast*
assume $\langle y \in \{ \|f *_v x\| \mid x. \|x\| < 1 \} \rangle$
hence $\langle \exists x. y = \|f *_v x\| \wedge \|x\| < 1 \rangle$
by *blast*
then obtain x **where** $\langle y = \|f *_v x\| \rangle$ **and** $\langle \|x\| < 1 \rangle$
by *blast*
have $\langle (1/\|x\|) * y = (1/\|x\|) * \|f *_v x\| \rangle$
by (*simp add: \langle y = \|f *_v x\| \rangle*)
also have $\langle \dots = |1/\|x\|| * \|f *_v x\| \rangle$
by *simp*
also have $\langle \dots = \|(1/\|x\|) *_R (f *_v x)\| \rangle$

```

    by simp
  also have ⟨... = ||f *v ((1/||x||) *R x)||⟩
    by (simp add: blinfun.scaleR-right)
  finally have ⟨(1/||x||) * y = ||f *v ((1/||x||) *R x)||⟩
    by blast
  have ⟨x ≠ 0⟩
    using ⟨y ≠ 0⟩ ⟨y = ||f *v x||⟩ blinfun.zero-right by auto
  have ⟨|| (1/||x||) *R x || = | (1/||x||) | * ||x||⟩
    by simp
  also have ⟨... = (1/||x||) * ||x||⟩
    by simp
  finally have ⟨|| (1/||x||) *R x || = 1⟩
    using ⟨x ≠ 0⟩ by simp
  hence ⟨(1/||x||) * y ∈ { ||f *v x|| | x. ||x|| = 1 }⟩
    using ⟨1 / ||x|| * y = ||f *v (1 / ||x||) *R x||⟩ by blast
  hence ⟨(1/||x||) * y ≤ Sup { ||f *v x|| | x. ||x|| = 1 }⟩
    by (simp add: cSup-upper norm-1-bounded)
  moreover have ⟨y ≤ (1/||x||) * y⟩
    by (metis ⟨||x|| < 1⟩ ⟨y = ||f *v x||⟩ mult-le-cancel-right1 norm-not-less-zero
      order.strict-implies-order ⟨x ≠ 0⟩ less-divide-eq-1-pos zero-less-norm-iff)
  ultimately show ?thesis by linarith
qed
hence ⟨Sup { ||f *v x|| | x. ||x|| < 1 } ≤ Sup { ||f *v x|| | x. ||x|| = 1 }⟩
  by (smt cSup-least norm-less-1-non-empty)
hence ⟨Sup { ||f *v x|| | x. ||x|| = 1 } = Sup { ||f *v x|| | x. ||x|| < 1 }⟩
  using ⟨Sup { ||f *v x|| | x. norm x = 1 } ≤ Sup { ||f *v x|| | x. ||x|| < 1 }⟩ by
linarith
have f1: ⟨(SUP x. ||f *v x|| / ||x||) = Sup { ||f *v x|| / ||x|| | x. True}⟩
  by (simp add: full-SetCompr-eq)
have ⟨y ∈ { ||f *v x|| / ||x|| | x. True} ⟹ y ∈ { ||f *v x|| | x. ||x|| = 1 } ∪ {0}⟩
  for y
proof-
  assume ⟨y ∈ { ||f *v x|| / ||x|| | x. True}⟩ show ?thesis
  proof(cases ⟨y = 0⟩)
    case True thus ?thesis by simp
  next
    case False
    have ⟨∃ x. y = ||f *v x|| / ||x||⟩
      using ⟨y ∈ { ||f *v x|| / ||x|| | x. True}⟩ by auto
    then obtain x where ⟨y = ||f *v x|| / ||x||⟩
      by blast
    hence ⟨y = |(1/||x||)| * ||f *v x||⟩
      by simp
    hence ⟨y = ||(1/||x||) *R (f *v x)||⟩
      by simp
    hence ⟨y = ||f ((1/||x||) *R x)||⟩
      by (simp add: blinfun.scaleR-right)
    moreover have ⟨|| (1/||x||) *R x || = 1⟩

```

```

    using False ⟨y = ‖f *v x‖ / ‖x‖⟩ by auto
    ultimately have ⟨y ∈ {‖f *v x‖ | x. ‖x‖ = 1}⟩
      by blast
    thus ?thesis by blast
  qed
  qed
  moreover have ⟨y ∈ {‖f x‖ | x. ‖x‖ = 1} ∪ {0} ⟹ y ∈ {‖f *v x‖ / ‖x‖ | x.
True}⟩
    for y
  proof(cases ⟨y = 0⟩)
    case True thus ?thesis by auto
  next
    case False
    hence ⟨y ∉ {0}⟩
      by simp
    moreover assume ⟨y ∈ {‖f *v x‖ | x. ‖x‖ = 1} ∪ {0}⟩
    ultimately have ⟨y ∈ {‖f *v x‖ | x. ‖x‖ = 1}⟩
      by simp
    then obtain x where ⟨‖x‖ = 1⟩ and ⟨y = ‖f *v x‖⟩
      by auto
    have ⟨y = ‖f *v x‖ / ‖x‖⟩ using ⟨‖x‖ = 1⟩ ⟨y = ‖f *v x‖⟩
      by simp
    thus ?thesis by auto
  qed
  ultimately have ⟨{‖f *v x‖ / ‖x‖ | x. True} = {‖f *v x‖ | x. ‖x‖ = 1} ∪ {0}⟩
    by blast
  hence ⟨Sup {‖f *v x‖ / ‖x‖ | x. True} = Sup ({‖f *v x‖ | x. ‖x‖ = 1} ∪ {0})⟩
    by simp
  have ∧r s. ¬ (r::real) ≤ s ∨ sup r s = s
    by (metis (lifting) sup.absorb-iff1 sup-commute)
  hence ⟨Sup ({‖f *v x‖ | x. ‖x‖ = 1} ∪ {(0::real)})
    = max (Sup {‖f *v x‖ | x. ‖x‖ = 1}) (Sup {0::real})⟩
    using ⟨0 ≤ Sup {‖f *v x‖ | x. ‖x‖ = 1}⟩ ‹bdd-above {0}› ‹{0} ≠ {}›
  cSup-singleton
  cSup-union-distrib max.absorb-iff1 sup-commute norm-1-bounded norm-1-non-empty
  by (metis (no-types, lifting) )
  moreover have ⟨Sup {(0::real)} = (0::real)⟩
    by simp
  ultimately have ⟨Sup ({‖f *v x‖ | x. ‖x‖ = 1} ∪ {0}) = Sup {‖f *v x‖ | x.
‖x‖ = 1}⟩
    using Sup-non-neg by linarith
  moreover have ⟨Sup ( {‖f *v x‖ | x. ‖x‖ = 1} ∪ {0})
    = max (Sup {‖f *v x‖ | x. ‖x‖ = 1}) (Sup {0}) ⟩
    using Sup-non-neg ‹Sup ({‖f *v x‖ | x. ‖x‖ = 1} ∪ {0})
    = max (Sup {‖f *v x‖ | x. ‖x‖ = 1}) (Sup {0})›
    by auto
  ultimately have f2: ‹Sup {‖f *v x‖ / ‖x‖ | x. True} = Sup {‖f *v x‖ | x. ‖x‖
= 1}⟩
    using ‹Sup {‖f *v x‖ / ‖x‖ | x. True} = Sup ({‖f *v x‖ | x. ‖x‖ = 1} ∪ {0})›

```

by *linarith*
have $\langle (SUP\ x.\ \|f\ *v\ x\| / \|x\|) = Sup\ \{\|f\ *v\ x\| \mid x.\ \|x\| = 1\} \rangle$
using *f1 f2 by linarith*
hence $\langle (SUP\ x.\ \|f\ *v\ x\| / \|x\|) = Sup\ \{\|f\ *v\ x\| \mid x.\ \|x\| < 1\} \rangle$
by (*simp add: $\langle Sup\ \{\|f\ *v\ x\| \mid x.\ \|x\| = 1\} = Sup\ \{\|f\ *v\ x\| \mid x.\ \|x\| < 1\} \rangle$*)

thus *?thesis apply transfer by (simp add: onorm-def)*
qed
qed

lemma *onorm-r:*

includes *notation-norm*

assumes $\langle r > 0 \rangle$

shows $\langle \|f\| = Sup\ ((\lambda x.\ \|f\ *v\ x\|) \text{ ' } (ball\ 0\ r)) / r \rangle$

Explanation: The norm of f is $1 / r$ of the supremum of the norm of $f *v x$ for x in the ball of radius r centered at the origin.

proof–

have $\langle \|f\| = Sup\ \{\|f\ *v\ x\| \mid x.\ \|x\| < 1\} \rangle$

using *onorm-open-ball by blast*

moreover have $\langle \{\|f\ *v\ x\| \mid x.\ \|x\| < 1\} = (\lambda x.\ \|f\ *v\ x\|) \text{ ' } (ball\ 0\ 1) \rangle$

unfolding *ball-def by auto*

ultimately have *onorm-f:* $\langle \|f\| = Sup\ ((\lambda x.\ \|f\ *v\ x\|) \text{ ' } (ball\ 0\ 1)) \rangle$

by *simp*

have *s2:* $\langle x \in (\lambda t.\ r *R\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1 \implies x \leq r * Sup\ ((\lambda t.\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1) \rangle$ **for** x

proof–

assume $\langle x \in (\lambda t.\ r *R\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1 \rangle$

hence $\langle \exists t.\ x = r *R\ \|f\ *v\ t\| \wedge \|t\| < 1 \rangle$

by *auto*

then obtain t **where** $\langle x = r *R\ \|f\ *v\ t\| \rangle$ **and** $\langle \|t\| < 1 \rangle$

by *blast*

define y **where** $\langle y = x /R\ r \rangle$

have $\langle x = r * (inverse\ r * x) \rangle$

using $\langle x = r *R\ norm\ (f\ t) \rangle$ **by** *auto*

hence $\langle x - (r * (inverse\ r * x)) \leq 0 \rangle$

by *linarith*

hence $\langle x \leq r * (x /R\ r) \rangle$

by *auto*

have $\langle y \in (\lambda k.\ \|f\ *v\ k\|) \text{ ' } ball\ 0\ 1 \rangle$

unfolding *y-def by (smt $\langle x \in (\lambda t.\ r *R\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1 \rangle$ assms image-iff inverse-inverse-eq pos-le-divideR-eq positive-imp-inverse-positive)*

moreover have $\langle x \leq r * y \rangle$

using $\langle x \leq r * (x /R\ r) \rangle$ *y-def by blast*

ultimately have *y-norm-f:* $\langle y \in (\lambda t.\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1 \wedge x \leq r * y \rangle$

by *blast*

have $\langle (\lambda t.\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1 \neq \{\} \rangle$

by *simp*

moreover have $\langle bdd-above\ ((\lambda t.\ \|f\ *v\ t\|) \text{ ' } ball\ 0\ 1) \rangle$

by (*simp add: bounded-linear-image blinfun.bounded-linear-right bounded-imp-bdd-above*)

bounded-norm-comp)

moreover have $\langle \exists y. y \in (\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \wedge x \leq r * y \rangle$

using *y-norm-f* by *blast*

ultimately show *?thesis*

by (*smt* $\langle 0 < r \rangle$ *cSup-upper ordered-comm-semiring-class.comm-mult-left-mono*)

qed

have *s3*: $\langle (\bigwedge x. x \in (\lambda t. r * \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \implies x \leq y) \implies r * \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1) \leq y \text{ for } y \rangle$

proof–

assume $\langle \bigwedge x. x \in (\lambda t. r * \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \implies x \leq y \rangle$

have *x-leq*: $\langle x \in (\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \implies x \leq y / r \rangle$ for *x*

proof–

assume $\langle x \in (\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \rangle$

then obtain *t* where $\langle t \in \text{ball } (0::'a) \ 1 \rangle$ and $\langle x = \|f *_{\nu} t\| \rangle$

by *auto*

define *x'* where $\langle x' = r *_{\mathbb{R}} x \rangle$

have $\langle x' = r * \|f *_{\nu} t\| \rangle$

by (*simp add: x = \|f *_{\nu} t\| x'-def*)

hence $\langle x' \in (\lambda t. r * \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \rangle$

using $\langle t \in \text{ball } (0::'a) \ 1 \rangle$ by *auto*

hence $\langle x' \leq y \rangle$

using $\langle \bigwedge x. x \in (\lambda t. r * \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \implies x \leq y \rangle$ by *blast*

thus $\langle x \leq y / r \rangle$

unfolding *x'-def* using $\langle r > 0 \rangle$ by (*simp add: mult.commute pos-le-divide-eq*)

qed

have $\langle (\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1 \neq \{\} \rangle$

by *simp*

moreover have $\langle \text{bdd-above } ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1) \rangle$

by (*simp add: bounded-linear-image blinfun.bounded-linear-right bounded-imp-bdd-above*)

bounded-norm-comp)

ultimately have $\langle \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ 1) \leq y / r \rangle$

using *x-leq* by (*simp add: bdd-above ((\lambda t. \|f *_{\nu} t\|) 'ball 0 1) cSup-least*)

thus *?thesis* using $\langle r > 0 \rangle$

by (*smt divide-strict-right-mono nonzero-mult-div-cancel-left*)

qed

have *norm-scaleR*: $\langle \text{norm} \circ ((*_R) \ r) = ((*_R) \ |r|) \circ (\text{norm}::'a \Rightarrow \text{real}) \rangle$

by *auto*

have *f-x1*: $\langle f (r *_{\mathbb{R}} x) = r *_{\mathbb{R}} f x \rangle$ for *x*

by (*simp add: blinfun.scaleR-right*)

have $\langle \text{ball } (0::'a) \ r = ((*_R) \ r) \text{ 'ball } 0 \ 1 \rangle$

by (*smt assms ball-scale nonzero-mult-div-cancel-left right-inverse-eq scale-zero-right*)

hence $\langle \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 \ r) = \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \text{ '(((*_R) \ r) 'ball 0 1))) \rangle$

by *simp*

also have $\langle \dots = \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \circ ((*_R) r)) \text{ ' (ball 0 1)} \rangle$
using *Sup.SUP-image* **by** *auto*
also have $\langle \dots = \text{Sup } ((\lambda t. \|f *_{\nu} (r *_R t)\|) \text{ ' (ball 0 1)} \rangle$
using *f-x1* **by** (*simp add: comp-assoc*)
also have $\langle \dots = \text{Sup } ((\lambda t. |r| *_R \|f *_{\nu} t\|) \text{ ' (ball 0 1)} \rangle$
using *norm-scaleR f-x1* **by** *auto*
also have $\langle \dots = \text{Sup } ((\lambda t. r *_R \|f *_{\nu} t\|) \text{ ' (ball 0 1)} \rangle$
using $\langle r > 0 \rangle$ **by** *auto*
also have $\langle \dots = r * \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \text{ ' (ball 0 1)} \rangle$
apply (*rule cSup-eq-non-empty*) **apply** *simp* **using** *s2* **apply** *auto* **using** *s3*
by *auto*
also have $\langle \dots = r * \|f\| \rangle$
using *onorm-f* **by** *auto*
finally have $\langle \text{Sup } ((\lambda t. \|f *_{\nu} t\|) \text{ ' ball 0 r}) = r * \|f\| \rangle$
by *blast*
thus $\langle \|f\| = \text{Sup } ((\lambda x. \|f *_{\nu} x\|) \text{ ' (ball 0 r)}) / r \rangle$ **using** $\langle r > 0 \rangle$ **by** *simp*
qed

Pointwise convergence

definition *pointwise-convergent-to* ::

$\langle (\text{nat} \Rightarrow ('a \Rightarrow 'b::\text{topological-space})) \Rightarrow ('a \Rightarrow 'b) \Rightarrow \text{bool} \rangle$
 $\langle ((-)/ -\text{pointwise} \rightarrow (-)) \rangle$ [*60*, *60*] *60* **where**
 $\langle \text{pointwise-convergent-to } x \ l = (\forall t::'a. (\lambda n. (x \ n) \ t) \longrightarrow l \ t) \rangle$

lemma *linear-limit-linear*:

fixes $f :: \langle - \Rightarrow ('a::\text{real-vector} \Rightarrow 'b::\text{real-normed-vector}) \rangle$
assumes $\langle \bigwedge n. \text{linear } (f \ n) \rangle$ **and** $\langle f \ -\text{pointwise} \rightarrow F \rangle$
shows $\langle \text{linear } F \rangle$

Explanation: If a family of linear operators converges pointwise, then the limit is also a linear operator.

proof

show $F (x + y) = F x + F y$ **for** $x \ y$

proof –

have $\forall a. F a = \text{lim } (\lambda n. f \ n \ a)$

using $\langle f \ -\text{pointwise} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def* **by** (*metis* (*full-types*) *limI*)

moreover have $\forall f \ b \ c \ g. (\text{lim } (\lambda n. g \ n + f \ n) = (b::'b) + c \vee \neg f \longrightarrow c) \vee \neg g \longrightarrow b$

by (*metis* (*no-types*) *limI tendsto-add*)

moreover have $\bigwedge a. (\lambda n. f \ n \ a) \longrightarrow F a$

using *assms(2)* *pointwise-convergent-to-def* **by** *force*

ultimately have

$\text{lim-sum}: \langle \text{lim } (\lambda n. (f \ n) \ x + (f \ n) \ y) = \text{lim } (\lambda n. (f \ n) \ x) + \text{lim } (\lambda n. (f \ n) \ y) \rangle$

by *metis*

have $\langle (f \ n) (x + y) = (f \ n) x + (f \ n) y \rangle$ **for** n

using $\langle \bigwedge n. \text{linear } (f \ n) \rangle$ **unfolding** *linear-def* **using** *Real-Vector-Spaces.linear-iff* *assms(1)*

by *auto*
 hence $\langle \lim (\lambda n. (f n) (x + y)) = \lim (\lambda n. (f n) x + (f n) y) \rangle$
 by *simp*
 hence $\langle \lim (\lambda n. (f n) (x + y)) = \lim (\lambda n. (f n) x) + \lim (\lambda n. (f n) y) \rangle$
 using *lim-sum* by *simp*
 moreover have $\langle (\lambda n. (f n) (x + y)) \longrightarrow F (x + y) \rangle$
 using $\langle f \text{ --pointwise--} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def* by *blast*
 moreover have $\langle (\lambda n. (f n) x) \longrightarrow F x \rangle$
 using $\langle f \text{ --pointwise--} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def* by *blast*
 moreover have $\langle (\lambda n. (f n) y) \longrightarrow F y \rangle$
 using $\langle f \text{ --pointwise--} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def* by *blast*
 ultimately show *?thesis*
 by (*metis limI*)

qed

show $F (r *_R x) = r *_R F x$ for r and x

proof –

have $\langle (f n) (r *_R x) = r *_R (f n) x \rangle$ for n
 using $\langle \bigwedge n. \text{linear } (f n) \rangle$
 by (*simp add: Real-Vector-Spaces.linear-def real-vector.linear-scale*)
 hence $\langle \lim (\lambda n. (f n) (r *_R x)) = \lim (\lambda n. r *_R (f n) x) \rangle$
 by *simp*
 have $\langle \text{convergent } (\lambda n. (f n) x) \rangle$
 by (*metis assms(2) convergentI pointwise-convergent-to-def*)
 moreover have $\langle \text{isCont } (\lambda t::'b. r *_R t) tt \rangle$ for tt
 by (*simp add: bounded-linear-scaleR-right*)
 ultimately have $\langle \lim (\lambda n. r *_R ((f n) x)) = r *_R \lim (\lambda n. (f n) x) \rangle$
 using $\langle f \text{ --pointwise--} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def*
 by (*metis (mono-tags) isCont-tendsto-compose limI*)
 hence $\langle \lim (\lambda n. (f n) (r *_R x)) = r *_R \lim (\lambda n. (f n) x) \rangle$
 using $\langle \lim (\lambda n. (f n) (r *_R x)) = \lim (\lambda n. r *_R (f n) x) \rangle$ by *simp*
 moreover have $\langle (\lambda n. (f n) x) \longrightarrow F x \rangle$
 using $\langle f \text{ --pointwise--} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def* by *blast*
 moreover have $\langle (\lambda n. (f n) (r *_R x)) \longrightarrow F (r *_R x) \rangle$
 using $\langle f \text{ --pointwise--} \rightarrow F \rangle$ **unfolding** *pointwise-convergent-to-def* by *blast*
 ultimately show *?thesis*
 by (*metis limI*)

qed

qed

lemma *non-Cauchy-unbounded*:

fixes $a :: \text{real}$

assumes $\langle \bigwedge n. a n \geq 0 \rangle$ and $\langle e > 0 \rangle$

and $\langle \forall M. \exists m. \exists n. m \geq M \wedge n \geq M \wedge m > n \wedge \text{sum } a \{ \text{Suc } n..m \} \geq e \rangle$

shows $\langle (\lambda n. (\text{sum } a \{ 0..n \})) \longrightarrow \infty \rangle$

Explanation: If the sequence of partial sums of nonnegative terms is not Cauchy, then it converges to infinite.

proof –

```

define  $S::\text{ereal set}$  where  $\langle S = \text{range } (\lambda n. \text{sum } a \{0..n\}) \rangle$ 
have  $\langle \exists s \in S. k * e \leq s \rangle$  for  $k::\text{nat}$ 
proof(induction k)
  case 0
    from  $\langle \forall M. \exists m. \exists n. m \geq M \wedge n \geq M \wedge m > n \wedge \text{sum } a \{ \text{Suc } n..m \} \geq e \rangle$ 
    obtain  $m \ n$  where  $\langle m \geq 0 \rangle$  and  $\langle n \geq 0 \rangle$  and  $\langle m > n \rangle$  and  $\langle \text{sum } a \{ \text{Suc } n..m \} \geq e \rangle$  by blast
    have  $\langle n < \text{Suc } n \rangle$ 
    by simp
    hence  $\langle \{0..n\} \cup \{ \text{Suc } n..m \} = \{0..m\} \rangle$ 
    using Set-Interval.int-disj-un(7)  $\langle n < m \rangle$  by auto
    moreover have  $\langle \text{finite } \{0..n\} \rangle$ 
    by simp
    moreover have  $\langle \text{finite } \{ \text{Suc } n..m \} \rangle$ 
    by simp
    moreover have  $\langle \{0..n\} \cap \{ \text{Suc } n..m \} = \{ \} \rangle$ 
    by simp
    ultimately have  $\langle \text{sum } a \{0..n\} + \text{sum } a \{ \text{Suc } n..m \} = \text{sum } a \{0..m\} \rangle$ 
    by (metis sum.union-disjoint)
    moreover have  $\langle \text{sum } a \{ \text{Suc } n..m \} > 0 \rangle$ 
    using  $\langle e > 0 \rangle$   $\langle \text{sum } a \{ \text{Suc } n..m \} \geq e \rangle$  by linarith
    moreover have  $\langle \text{sum } a \{0..n\} \geq 0 \rangle$ 
    by (simp add: assms(1) sum-nonneg)
    ultimately have  $\langle \text{sum } a \{0..m\} > 0 \rangle$ 
    by linarith
    moreover have  $\langle \text{sum } a \{0..m\} \in S \rangle$ 
    unfolding S-def by blast
    ultimately have  $\langle \exists s \in S. 0 \leq s \rangle$ 
    using ereal-less-eq(5) by fastforce
    thus ?case
    by (simp add: zero-ereal-def)
  next
    case (Suc k)
    assume  $\langle \exists s \in S. k * e \leq s \rangle$ 
    then obtain  $s$  where  $\langle s \in S \rangle$  and  $\langle \text{ereal } (k * e) \leq s \rangle$ 
    by blast
    have  $\langle \exists N. s = \text{sum } a \{0..N\} \rangle$ 
    using  $\langle s \in S \rangle$  unfolding S-def by blast
    then obtain  $N$  where  $\langle s = \text{sum } a \{0..N\} \rangle$ 
    by blast
    from  $\langle \forall M. \exists m. \exists n. m \geq M \wedge n \geq M \wedge m > n \wedge \text{sum } a \{ \text{Suc } n..m \} \geq e \rangle$ 
    obtain  $m \ n$  where  $\langle m \geq \text{Suc } N \rangle$  and  $\langle n \geq \text{Suc } N \rangle$  and  $\langle m > n \rangle$  and  $\langle \text{sum } a \{ \text{Suc } n..m \} \geq e \rangle$ 
    by blast
    have  $\langle \text{finite } \{ \text{Suc } N..n \} \rangle$ 
    by simp
    moreover have  $\langle \text{finite } \{ \text{Suc } n..m \} \rangle$ 
    by simp
    moreover have  $\langle \{ \text{Suc } N..n \} \cup \{ \text{Suc } n..m \} = \{ \text{Suc } N..m \} \rangle$ 

```

using *Set-Interval.ivl-disj-un*
by (*smt* $\langle \text{Suc } N \leq n \rangle \langle n < m \rangle$ *atLeastSucAtMost-greaterThanAtMost less-imp-le-nat*)
moreover have $\langle \{\} = \{\text{Suc } N..n\} \cap \{\text{Suc } n..m\} \rangle$
by *simp*
ultimately have $\langle \text{sum } a \{\text{Suc } N..m\} = \text{sum } a \{\text{Suc } N..n\} + \text{sum } a \{\text{Suc } n..m\} \rangle$
n..m
by (*metis sum.union-disjoint*)
moreover have $\langle \text{sum } a \{\text{Suc } N..n\} \geq 0 \rangle$
using $\langle \bigwedge n. a \ n \geq 0 \rangle$ **by** (*simp add: sum-nonneg*)
ultimately have $\langle \text{sum } a \{\text{Suc } N..m\} \geq e \rangle$
using $\langle e \leq \text{sum } a \{\text{Suc } n..m\} \rangle$ **by** *linarith*
have $\langle \text{finite } \{0..N\} \rangle$
by *simp*
have $\langle \text{finite } \{\text{Suc } N..m\} \rangle$
by *simp*
moreover have $\langle \{0..N\} \cup \{\text{Suc } N..m\} = \{0..m\} \rangle$
using *Set-Interval.ivl-disj-un*(γ) $\langle \text{Suc } N \leq m \rangle$ **by** *auto*
moreover have $\langle \{0..N\} \cap \{\text{Suc } N..m\} = \{\} \rangle$
by *simp*
ultimately have $\langle \text{sum } a \{0..N\} + \text{sum } a \{\text{Suc } N..m\} = \text{sum } a \{0..m\} \rangle$
by (*metis* $\langle \text{finite } \{0..N\} \rangle$ *sum.union-disjoint*)
hence $\langle e + k * e \leq \text{sum } a \{0..m\} \rangle$
using $\langle \text{ereal } (\text{real } k * e) \leq s \rangle \langle s = \text{ereal } (\text{sum } a \{0..N\}) \rangle \langle e \leq \text{sum } a \{\text{Suc } N..m\} \rangle$ **by** *auto*
moreover have $\langle e + k * e = (\text{Suc } k) * e \rangle$
by (*simp add: semiring-normalization-rules*(3))
ultimately have $\langle (\text{Suc } k) * e \leq \text{sum } a \{0..m\} \rangle$
by *linarith*
hence $\langle \text{ereal } ((\text{Suc } k) * e) \leq \text{sum } a \{0..m\} \rangle$
by *auto*
moreover have $\langle \text{sum } a \{0..m\} \in S \rangle$
unfolding *S-def* **by** *blast*
ultimately show *?case* **by** *blast*
qed
hence $\langle \exists s \in S. (\text{real } n) \leq s \rangle$ **for** *n*
by (*meson assms*(2) *ereal-le-le ex-less-of-nat-mult less-le-not-le*)
hence $\langle \text{Sup } S = \infty \rangle$
using *Sup-le-iff Sup-subset-mono dual-order.strict-trans1 leD less-PInf-Ex-of-nat subsetI*
by *metis*
hence *Sup*: $\langle \text{Sup } ((\text{range } (\lambda n. (\text{sum } a \{0..n\})))::\text{ereal set}) = \infty \rangle$ **using** *S-def*
by *blast*
have $\langle \text{incseq } (\lambda n. (\text{sum } a \{..<n\})) \rangle$
using $\langle \bigwedge n. a \ n \geq 0 \rangle$ **using** *Extended-Real.incseq-sumI* **by** *auto*
hence $\langle \text{incseq } (\lambda n. (\text{sum } a \{..< \text{Suc } n\})) \rangle$
by (*meson incseq-Suc-iff*)
hence $\langle \text{incseq } (\lambda n. (\text{sum } a \{0..n\}))::\text{ereal} \rangle$
using *incseq-ereal* **by** (*simp add: atLeast0AtMost lessThan-Suc-atMost*)
hence $\langle (\lambda n. \text{sum } a \{0..n\}) \longrightarrow \text{Sup } (\text{range } (\lambda n. (\text{sum } a \{0..n\}))::\text{ereal}) \rangle$

using *LIMSEQ-SUP* **by** *auto*
thus *?thesis using Sup PInfty-neq-ereal* **by** *auto*
qed

lemma *sum-Cauchy-positive*:

fixes $a :: \langle - \Rightarrow \text{real} \rangle$

assumes $\langle \bigwedge n. a\ n \geq 0 \rangle$ **and** $\langle \exists K. \forall n. (\text{sum } a\ \{0..n\}) \leq K \rangle$

shows $\langle \text{Cauchy } (\lambda n. \text{sum } a\ \{0..n\}) \rangle$

Explanation: If a series of nonnegative reals is bounded, then the series is Cauchy.

proof (*unfold Cauchy-altdef2, rule, rule*)

fix $e :: \text{real}$

assume $\langle e > 0 \rangle$

have $\langle \exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{sum } a\ \{\text{Suc } n..m\} < e \rangle$

proof (*rule classical*)

assume $\langle \neg(\exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{sum } a\ \{\text{Suc } n..m\} < e) \rangle$

hence $\langle \forall M. \exists m. \exists n. m \geq M \wedge n \geq M \wedge m > n \wedge \neg(\text{sum } a\ \{\text{Suc } n..m\} < e) \rangle$

by *blast*

hence $\langle \forall M. \exists m. \exists n. m \geq M \wedge n \geq M \wedge m > n \wedge \text{sum } a\ \{\text{Suc } n..m\} \geq e \rangle$

by *fastforce*

hence $\langle (\lambda n. (\text{sum } a\ \{0..n\})) \longrightarrow \infty \rangle$

using *non-Cauchy-unbounded* $\langle 0 < e \rangle$ *assms(1)* **by** *blast*

from $\langle \exists K. \forall n. \text{sum } a\ \{0..n\} \leq K \rangle$

obtain K **where** $\langle \forall n. \text{sum } a\ \{0..n\} \leq K \rangle$

by *blast*

from $\langle (\lambda n. \text{sum } a\ \{0..n\}) \longrightarrow \infty \rangle$

have $\langle \forall B. \exists N. \forall n \geq N. (\lambda n. (\text{sum } a\ \{0..n\}))\ n \geq B \rangle$

using *Lim-PInfty* **by** *simp*

hence $\langle \exists n. (\text{sum } a\ \{0..n\}) \geq K+1 \rangle$

using *ereal-less-eq(3)* **by** *blast*

thus *?thesis using* $\langle \forall n. (\text{sum } a\ \{0..n\}) \leq K \rangle$ **by** *smt*

qed

have $\langle \text{sum } a\ \{\text{Suc } n..m\} = \text{sum } a\ \{0..m\} - \text{sum } a\ \{0..n\} \rangle$

if $m > n$ **for** $m\ n$

apply (*simp add: that atLeast0AtMost*) **using** *sum-up-index-split*

by (*smt less-imp-add-positive that*)

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{sum } a\ \{0..m\} - \text{sum } a\ \{0..n\} < e \rangle$

using $\langle \exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{sum } a\ \{\text{Suc } n..m\} < e \rangle$ **by** *smt*

from $\langle \exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{sum } a\ \{0..m\} - \text{sum } a\ \{0..n\} < e \rangle$

obtain M **where** $\langle \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{sum } a\ \{0..m\} - \text{sum } a\ \{0..n\} < e \rangle$

by *blast*

moreover **have** $\langle m > n \implies \text{sum } a\ \{0..m\} \geq \text{sum } a\ \{0..n\} \rangle$ **for** $m\ n$

using $\langle \bigwedge n. a\ n \geq 0 \rangle$ **by** (*simp add: sum-mono2*)

ultimately **have** $\langle \exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow |\text{sum } a\ \{0..m\} - \text{sum } a\ \{0..n\}| < e \rangle$

by *auto*

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. m \geq n \longrightarrow |\text{sum } a \{0..m\} - \text{sum } a \{0..n\}| < e \rangle$
by $(\text{metis } \langle 0 < e \rangle \text{ abs-zero cancel-comm-monoid-add-class.diff-cancel diff-is-0-eq'}$

$\text{less-irrefl-nat linorder-neqE-nat zero-less-diff})$

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. |\text{sum } a \{0..m\} - \text{sum } a \{0..n\}| < e \rangle$

by $(\text{metis abs-minus-commute nat-le-linear})$

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. \text{dist } (\text{sum } a \{0..m\}) (\text{sum } a \{0..n\}) < e \rangle$

by $(\text{simp add: dist-real-def})$

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. \text{dist } (\text{sum } a \{0..m\}) (\text{sum } a \{0..n\}) < e \rangle$ **by** *blast*

thus $\langle \exists N. \forall n \geq N. \text{dist } (\text{sum } a \{0..n\}) (\text{sum } a \{0..N\}) < e \rangle$ **by** *auto*

qed

lemma *convergent-series-Cauchy*:

fixes $a::\langle \text{nat} \Rightarrow \text{real} \rangle$ **and** $\varphi::\langle \text{nat} \Rightarrow 'a::\text{metric-space} \rangle$

assumes $\langle \exists M. \forall n. \text{sum } a \{0..n\} \leq M \rangle$ **and** $\langle \bigwedge n. \text{dist } (\varphi (\text{Suc } n)) (\varphi n) \leq a n \rangle$

shows $\langle \text{Cauchy } \varphi \rangle$

Explanation: Let a be a real-valued sequence and let φ be sequence in a metric space. If the partial sums of a are uniformly bounded and the distance between consecutive terms of φ are bounded by the sequence a , then φ is Cauchy.

proof $(\text{unfold Cauchy-altdef2, rule, rule})$

fix $e::\text{real}$

assume $\langle e > 0 \rangle$

have $\langle \bigwedge k. a k \geq 0 \rangle$

using $\langle \bigwedge n. \text{dist } (\varphi (\text{Suc } n)) (\varphi n) \leq a n \rangle$ *dual-order.trans zero-le-dist* **by** *blast*

hence $\langle \text{Cauchy } (\lambda k. \text{sum } a \{0..k\}) \rangle$

using $\langle \exists M. \forall n. \text{sum } a \{0..n\} \leq M \rangle$ *sum-Cauchy-positive* **by** *blast*

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. \text{dist } (\text{sum } a \{0..m\}) (\text{sum } a \{0..n\}) < e \rangle$

unfolding *Cauchy-def* **using** $\langle e > 0 \rangle$ **by** *blast*

hence $\langle \exists M. \forall m \geq M. \forall n \geq M. m > n \longrightarrow \text{dist } (\text{sum } a \{0..m\}) (\text{sum } a \{0..n\}) < e \rangle$

by *blast*

have $\langle \text{dist } (\text{sum } a \{0..m\}) (\text{sum } a \{0..n\}) = \text{sum } a \{\text{Suc } n..m\} \rangle$ **if** $\langle n < m \rangle$ **for** $m n$

proof –

have $\langle n < \text{Suc } n \rangle$

by *simp*

have $\langle \text{finite } \{0..n\} \rangle$

by *simp*

moreover have $\langle \text{finite } \{\text{Suc } n..m\} \rangle$

by *simp*

moreover have $\langle \{0..n\} \cup \{\text{Suc } n..m\} = \{0..m\} \rangle$

using $\langle n < \text{Suc } n \rangle \langle n < m \rangle$ **by** *auto*

moreover have $\langle \{0..n\} \cap \{\text{Suc } n..m\} = \{\} \rangle$

by *simp*

ultimately have *sum-plus*: $\langle (\text{sum } a \{0..n\}) + \text{sum } a \{\text{Suc } n..m\} = (\text{sum } a \{0..m\}) \rangle$

```

    by (metis sum.union-disjoint)
  have ‹dist (sum a {0..m}) (sum a {0..n}) = |(sum a {0..m}) - (sum a {0..n})|›
    using dist-real-def by blast
  moreover have ‹(sum a {0..m}) - (sum a {0..n}) = sum a {Suc n..m}›
    using sum-plus by linarith
  ultimately show ?thesis
    by (simp add: ‹∧k. 0 ≤ a k› sum-nonneg)
qed
hence sum-a: ‹∃ M. ∀ m ≥ M. ∀ n ≥ M. m > n → sum a {Suc n..m} < e›
  by (metis ‹∃ M. ∀ m ≥ M. ∀ n ≥ M. dist (sum a {0..m}) (sum a {0..n}) < e›)
obtain M where ‹∀ m ≥ M. ∀ n ≥ M. m > n → sum a {Suc n..m} < e›
  using sum-a ‹e > 0› by blast
hence ‹∀ m. ∀ n. Suc m ≥ Suc M ∧ Suc n ≥ Suc M ∧ Suc m > Suc n → sum
a {Suc n..Suc m - 1} < e›
  by simp
hence ‹∀ m ≥ 1. ∀ n ≥ 1. m ≥ Suc M ∧ n ≥ Suc M ∧ m > n → sum a {n..m
- 1} < e›
  by (metis Suc-le-D)
hence sum-a2: ‹∃ M. ∀ m ≥ M. ∀ n ≥ M. m > n → sum a {n..m-1} < e›
  by (meson add-leE)
have ‹dist (φ (n+p+1)) (φ n) ≤ sum a {n..n+p}› for p n :: nat
proof(induction p)
  case 0 thus ?case by (simp add: assms(2))
next
  case (Suc p) thus ?case
    by (smt Suc-eq-plus1 add-Suc-right add-less-same-cancel1 assms(2) dist-self
dist-triangle2
gr-implies-not0 sum.cl-ivl-Suc)
qed
hence ‹m > n ⇒ dist (φ m) (φ n) ≤ sum a {n..m-1}› for m n :: nat
  by (metis Suc-eq-plus1 Suc-le-D diff-Suc-1 gr0-implies-Suc less-eq-Suc-le less-imp-Suc-add

zero-less-Suc)
hence ‹∃ M. ∀ m ≥ M. ∀ n ≥ M. m > n → dist (φ m) (φ n) < e›
  using sum-a2 ‹e > 0› by smt
thus ‹∃ N. ∀ n ≥ N. dist (φ n) (φ N) < e›
  using ‹0 < e› by fastforce
qed

unbundle notation-blinfun-apply

unbundle no-notation-norm

end

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2 Banach-Steinhaus theorem

```

theory Banach-Steinhaus
  imports Banach-Steinhaus-Missing

```

begin

We formalize Banach-Steinhaus theorem as theorem *banach-steinhaus*. This theorem was originally proved in Banach-Steinhaus's paper [1]. For the proof, we follow Sokal's approach [3]. Furthermore, we prove as a corollary a result about pointwise convergent sequences of bounded operators whose domain is a Banach space.

2.1 Preliminaries for Sokal's proof of Banach-Steinhaus theorem

lemma *linear-plus-norm*:

includes *notation-norm*

assumes $\langle \text{linear } f \rangle$

shows $\langle \|f \xi\| \leq \max \|f (x + \xi)\| \|f (x - \xi)\| \rangle$

Explanation: For arbitrary x and a linear operator f , $\|f \xi\|$ is upper bounded by the maximum of the norms of the shifts of f (i.e., $f (x + \xi)$ and $f (x - \xi)$).

proof –

have $\langle \text{norm } (f \xi) = \text{norm } ((\text{inverse } (\text{of-nat } 2)) *_{\mathbb{R}} (f (x + \xi) - f (x - \xi))) \rangle$

by (*smt add-diff-cancel-left' assms diff-add-cancel diff-diff-add linear-diff midpoint-def*

midpoint-plus-self of-nat-1 of-nat-add one-add-one scaleR-half-double)

also have $\langle \dots = \text{inverse } (\text{of-nat } 2) * \text{norm } (f (x + \xi) - f (x - \xi)) \rangle$

using *Real-Vector-Spaces.real-normed-vector-class.norm-scaleR* **by** *simp*

also have $\langle \dots \leq \text{inverse } (\text{of-nat } 2) * (\text{norm } (f (x + \xi)) + \text{norm } (f (x - \xi))) \rangle$

by (*simp add: norm-triangle-ineq4*)

also have $\langle \dots \leq \max (\text{norm } (f (x + \xi))) (\text{norm } (f (x - \xi))) \rangle$

by *auto*

finally show *?thesis by blast*

qed

lemma *onorm-Sup-on-ball*:

includes *notation-norm*

assumes $\langle r > 0 \rangle$

shows $\|f\| \leq \text{Sup } ((\lambda x. \|f *_v x\|) \text{ ' } (\text{ball } x \ r)) / r$

Explanation: Let f be a bounded operator and let x be a point. For any $0 < r$, the operator norm of f is bounded above by the supremum of f applied to the open ball of radius r around x , divided by r .

proof –

have *bdd-above-3*: $\langle \text{bdd-above } ((\lambda x. \|f *_v x\|) \text{ ' } (\text{ball } 0 \ r)) \rangle$

proof –

obtain M **where** $\langle \bigwedge \xi. \|f *_v \xi\| \leq M * \text{norm } \xi \rangle$ **and** $\langle M \geq 0 \rangle$

using *norm-blinfun norm-ge-zero* **by** *blast*

hence $\langle \bigwedge \xi. \xi \in \text{ball } 0 \ r \implies \|f *_v \xi\| \leq M * r \rangle$

using $\langle r > 0 \rangle$ **by** (*smt mem-ball-0 mult-left-mono*)

thus *?thesis* **by** (*meson bdd-aboveI2*)
qed
have *bdd-above-2*: $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ ' (ball 0 r)}) \rangle$
proof–
have $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} x\|) \text{ ' (ball 0 r)}) \rangle$
by *auto*
moreover have $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} \xi\|) \text{ ' (ball 0 r)}) \rangle$
using *bdd-above-3* **by** *blast*
ultimately have $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} x\| + \|f *_{\nu} \xi\|) \text{ ' (ball 0 r)}) \rangle$
by (*rule bdd-above-plus*)
then obtain *M* **where** $\langle \bigwedge \xi. \xi \in \text{ball 0 r} \implies \|f *_{\nu} x\| + \|f *_{\nu} \xi\| \leq M \rangle$
unfolding *bdd-above-def* **by** (*meson image-eqI*)
moreover have $\langle \|f *_{\nu} (x + \xi)\| \leq \|f *_{\nu} x\| + \|f *_{\nu} \xi\| \text{ for } \xi \rangle$
by (*simp add: blinfun.add-right norm-triangle-ineq*)
ultimately have $\langle \bigwedge \xi. \xi \in \text{ball 0 r} \implies \|f *_{\nu} (x + \xi)\| \leq M \rangle$
by (*simp add: blinfun.add-right norm-triangle-le*)
thus *?thesis* **by** (*meson bdd-aboveI2*)
qed
have *bdd-above-4*: $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} (x - \xi)\|) \text{ ' (ball 0 r)}) \rangle$
proof–
obtain *K* **where** *K-def*: $\langle \bigwedge \xi. \xi \in \text{ball 0 r} \implies \|f *_{\nu} (x + \xi)\| \leq K \rangle$
using $\langle \text{bdd-above } ((\lambda \xi. \text{norm } (f (x + \xi))) \text{ ' (ball 0 r)}) \rangle$ **unfolding**
bdd-above-def
by (*meson image-eqI*)
have $\langle \xi \in \text{ball } (0::'a) \text{ r} \implies -\xi \in \text{ball 0 r} \rangle$ **for** ξ
by *auto*
thus *?thesis* **by** (*metis K-def ab-group-add-class.ab-diff-conv-add-uminus bdd-aboveI2*)
qed
have *bdd-above-1*: $\langle \text{bdd-above } ((\lambda \xi. \max \|f *_{\nu} (x + \xi)\| \|f *_{\nu} (x - \xi)\|) \text{ ' (ball 0 r)}) \rangle$
proof–
have $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ ' (ball 0 r)}) \rangle$
using *bdd-above-2* **by** *blast*
moreover have $\langle \text{bdd-above } ((\lambda \xi. \|f *_{\nu} (x - \xi)\|) \text{ ' (ball 0 r)}) \rangle$
using *bdd-above-4* **by** *blast*
ultimately show *?thesis*
unfolding *max-def* **apply** *auto* **apply** (*meson bdd-above-Int1 bdd-above-mono image-Int-subset*)
by (*meson bdd-above-Int1 bdd-above-mono image-Int-subset*)
qed
have *bdd-above-6*: $\langle \text{bdd-above } ((\lambda t. \|f *_{\nu} t\|) \text{ ' ball x r}) \rangle$
proof–
have $\langle \text{bounded } (\text{ball x r}) \rangle$
by *simp*
hence $\langle \text{bounded } ((\lambda t. \|f *_{\nu} t\|) \text{ ' ball x r}) \rangle$
by (*metis (no-types) add.left-neutral bdd-above-2 bdd-above-norm bounded-norm-comp*)

image-add-ball image-image
thus *?thesis*

by (simp add: bounded-imp-bdd-above)
 qed
 have norm-1: $\langle (\lambda \xi. \|f *_v (x + \xi)\|) \text{ 'ball } 0 r = (\lambda t. \|f *_v t\|) \text{ 'ball } x r \rangle$
 by (metis add.right-neutral ball-translation image-image)
 have bdd-above-5: $\langle \text{bdd-above } ((\lambda \xi. \text{norm } (f (x + \xi))) \text{ 'ball } 0 r) \rangle$
 by (simp add: bdd-above-2)
 have norm-2: $\langle \|\xi\| < r \implies \|f *_v (x - \xi)\| \in (\lambda \xi. \|f *_v (x + \xi)\|) \text{ 'ball } 0 r \rangle$
 for ξ
 proof-
 assume $\langle \|\xi\| < r \rangle$
 hence $\langle \xi \in \text{ball } (0::'a) r \rangle$
 by auto
 hence $\langle -\xi \in \text{ball } (0::'a) r \rangle$
 by auto
 thus ?thesis
 by (metis (no-types, lifting) ab-group-add-class.ab-diff-conv-add-uminus image-iff)
 qed
 have norm-2': $\langle \|\xi\| < r \implies \|f *_v (x + \xi)\| \in (\lambda \xi. \|f *_v (x - \xi)\|) \text{ 'ball } 0 r \rangle$
 for ξ
 proof-
 assume $\langle \text{norm } \xi < r \rangle$
 hence $\langle \xi \in \text{ball } (0::'a) r \rangle$
 by auto
 hence $\langle -\xi \in \text{ball } (0::'a) r \rangle$
 by auto
 thus ?thesis
 by (metis (no-types, lifting) diff-minus-eq-add image-iff)
 qed
 have bdd-above-6: $\langle \text{bdd-above } ((\lambda \xi. \|f *_v (x - \xi)\|) \text{ 'ball } 0 r) \rangle$
 by (simp add: bdd-above-4)
 have Sup-2: $\langle (\text{SUP } \xi \in \text{ball } 0 r. \max \|f *_v (x + \xi)\| \|f *_v (x - \xi)\|) = \max (\text{SUP } \xi \in \text{ball } 0 r. \|f *_v (x + \xi)\|) (\text{SUP } \xi \in \text{ball } 0 r. \|f *_v (x - \xi)\|) \rangle$
 for ξ
 proof-
 have $\langle \text{ball } (0::'a) r \neq \{\} \rangle$
 using $\langle r > 0 \rangle$ by auto
 moreover have $\langle \text{bdd-above } ((\lambda \xi. \|f *_v (x + \xi)\|) \text{ 'ball } 0 r) \rangle$
 using bdd-above-5 by blast
 moreover have $\langle \text{bdd-above } ((\lambda \xi. \|f *_v (x - \xi)\|) \text{ 'ball } 0 r) \rangle$
 using bdd-above-6 by blast
 ultimately show ?thesis
 using max-Sup
 by (metis (mono-tags, lifting) Banach-Steinhaus-Missing.pointwise-max-def image-cong)
 qed
 have Sup-3': $\langle \|\xi\| < r \implies \|f *_v (x + \xi)\| \in (\lambda \xi. \|f *_v (x - \xi)\|) \text{ 'ball } 0 r \rangle$ for $\xi::'a$
 by (simp add: norm-2')

have *Sup-3''*: $\langle \|\xi\| < r \implies \|f *_{\nu} (x - \xi)\| \in (\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ 'ball } 0 r \rangle$ **for**
 $\xi :: 'a$
by (*simp add: norm-2*)
have *Sup-3*: $\langle \max (SUP \xi \in \text{ball } 0 r. \|f *_{\nu} (x + \xi)\|) (SUP \xi \in \text{ball } 0 r. \|f *_{\nu} (x - \xi)\|) =$
 $(SUP \xi \in \text{ball } 0 r. \|f *_{\nu} (x + \xi)\|) \rangle$
proof–
have $\langle (\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ 'ball } 0 r = (\lambda \xi. \|f *_{\nu} (x - \xi)\|) \text{ 'ball } 0 r \rangle$
apply *auto using Sup-3' apply auto using Sup-3'' by blast*
hence $\langle Sup ((\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ 'ball } 0 r) = Sup ((\lambda \xi. \|f *_{\nu} (x - \xi)\|) \text{ 'ball } 0 r) \rangle$
by *simp*
thus *?thesis by simp*
qed
have *Sup-1*: $\langle Sup ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 r) \leq Sup ((\lambda \xi. \|f *_{\nu} \xi\|) \text{ 'ball } x r) \rangle$
proof–
have $\langle (\lambda t. \|f *_{\nu} t\|) \xi \leq \max \|f *_{\nu} (x + \xi)\| \|f *_{\nu} (x - \xi)\| \rangle$ **for** ξ
apply (*rule linear-plus-norm*) **apply** (*rule bounded-linear.linear*)
by (*simp add: blinfun.bounded-linear-right*)
moreover **have** $\langle \text{bdd-above } ((\lambda \xi. \max \|f *_{\nu} (x + \xi)\| \|f *_{\nu} (x - \xi)\|) \text{ 'ball } 0 r) \rangle$
using *bdd-above-1 by blast*
moreover **have** $\langle \text{ball } (0 :: 'a) r \neq \{\} \rangle$
using $\langle r > 0 \rangle$ **by** *auto*
ultimately **have** $\langle Sup ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 r) \leq$
 $Sup ((\lambda \xi. \max \|f *_{\nu} (x + \xi)\| \|f *_{\nu} (x - \xi)\|) \text{ 'ball } 0 r) \rangle$
using *cSUP-mono by smt*
also **have** $\langle \dots = \max (Sup ((\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ 'ball } 0 r))$
 $(Sup ((\lambda \xi. \|f *_{\nu} (x - \xi)\|) \text{ 'ball } 0 r)) \rangle$
using *Sup-2 by blast*
also **have** $\langle \dots = Sup ((\lambda \xi. \|f *_{\nu} (x + \xi)\|) \text{ 'ball } 0 r) \rangle$
using *Sup-3 by blast*
also **have** $\langle \dots = Sup ((\lambda \xi. \|f *_{\nu} \xi\|) \text{ 'ball } x r) \rangle$
by (*metis add.right-neutral ball-translation image-image*)
finally **show** *?thesis by blast*
qed
have $\langle \|f\| = (SUP x \in \text{ball } 0 r. \|f *_{\nu} x\|) / r \rangle$
using $\langle 0 < r \rangle$ *onorm-r by blast*
moreover **have** $\langle Sup ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } 0 r) / r \leq Sup ((\lambda \xi. \|f *_{\nu} \xi\|) \text{ 'ball } x r) / r \rangle$
using *Sup-1 <0 < r> divide-right-mono by fastforce*
ultimately **have** $\langle \|f\| \leq Sup ((\lambda t. \|f *_{\nu} t\|) \text{ 'ball } x r) / r \rangle$
by *simp*
thus *?thesis by simp*
qed
lemma *onorm-Sup-on-ball'*:
includes *notation-norm*

assumes $\langle r > 0 \rangle$ **and** $\langle \tau < 1 \rangle$
shows $\langle \exists \xi \in \text{ball } x \ r. \ \tau * r * \|f\| \leq \|f *_{\nu} \xi\| \rangle$

In the proof of Banach-Steinhaus theorem, we will use this variation of the lemma *onorm-Sup-on-ball*.

Explanation: Let f be a bounded operator, let x be a point and let r be a positive real number. For any real number $\tau < 1$, there is a point ξ in the open ball of radius r around x such that $\tau * r * \|f\| \leq \|f *_{\nu} \xi\|$.

proof(*cases* $\langle f = 0 \rangle$)

case *True*

thus *?thesis* **by** (*metis* *assms(1)* *centre-in-ball* *mult-zero-right* *norm-zero-order-refl* *zero-blinfun.rep-eq*)

next

case *False*

have *bdd-above-1*: $\langle \text{bdd-above } ((\lambda t. \|(*_{\nu}) f t\|) \text{ 'ball } x \ r) \rangle$ **for** $f::\langle 'a \Rightarrow_L 'b \rangle$

using *assms(1)* *bounded-linear-image* **by** (*simp* *add*: *bounded-linear-image* *blinfun.bounded-linear-right* *bounded-imp-bdd-above* *bounded-norm-comp*)

have $\langle \text{norm } f > 0 \rangle$

using $\langle f \neq 0 \rangle$ **by** *auto*

have $\langle \text{norm } f \leq \text{Sup } ((\lambda \xi. \|(*_{\nu}) f \xi\|) \text{ ' (ball } x \ r)) / r \rangle$

using $\langle r > 0 \rangle$ **by** (*simp* *add*: *onorm-Sup-on-ball*)

hence $\langle r * \text{norm } f \leq \text{Sup } ((\lambda \xi. \|(*_{\nu}) f \xi\|) \text{ ' (ball } x \ r)) \rangle$

using $\langle 0 < r \rangle$ **by** (*smt* *divide-strict-right-mono* *nonzero-mult-div-cancel-left*)

moreover **have** $\langle \tau * r * \text{norm } f < r * \text{norm } f \rangle$

using $\langle \tau < 1 \rangle$ **using** $\langle 0 < \text{norm } f \rangle$ $\langle 0 < r \rangle$ **by** *auto*

ultimately **have** $\langle \tau * r * \text{norm } f < \text{Sup } ((\text{norm } \circ ((*__{\nu}) f)) \text{ ' (ball } x \ r)) \rangle$

by *simp*

moreover **have** $\langle (\text{norm } \circ ((*__{\nu}) f)) \text{ ' (ball } x \ r) \neq \{\} \rangle$

using $\langle 0 < r \rangle$ **by** *auto*

moreover **have** $\langle \text{bdd-above } ((\text{norm } \circ ((*__{\nu}) f)) \text{ ' (ball } x \ r)) \rangle$

using *bdd-above-1* **apply** *transfer* **by** *simp*

ultimately **have** $\langle \exists t \in (\text{norm } \circ ((*__{\nu}) f)) \text{ ' (ball } x \ r). \ \tau * r * \text{norm } f < t \rangle$

by (*simp* *add*: *less-cSup-iff*)

thus *?thesis* **by** (*smt* *comp-def* *image-iff*)

qed

2.2 Banach-Steinhaus theorem

theorem *banach-steinhaus*:

fixes $f::\langle 'c \Rightarrow ('a::\text{banach} \Rightarrow_L 'b::\text{real-normed-vector}) \rangle$

assumes $\langle \bigwedge x. \text{bounded } (\text{range } (\lambda n. (f \ n) *_{\nu} x)) \rangle$

shows $\langle \text{bounded } (\text{range } f) \rangle$

This is Banach-Steinhaus Theorem.

Explanation: If a family of bounded operators on a Banach space is pointwise bounded, then it is uniformly bounded.

proof(*rule* *classical*)

assume $\langle \neg(\text{bounded } (\text{range } f)) \rangle$

```

have  $\langle \exists K. \forall n. \text{sum } (\lambda k. \text{inverse } (\text{real-of-nat } 3^k)) \{0..n\} \leq K \rangle$ 
proof -
  have  $\langle \text{summable } (\lambda n. \text{inverse } ((3::\text{real}) ^ n)) \rangle$ 
    by (simp flip: power-inverse)
  hence  $\langle \text{bounded } (\text{range } (\lambda n. \text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..<n\})) \rangle$ 
  using summable-imp-sums-bounded [where  $f = (\lambda n. \text{inverse } (\text{real-of-nat } 3^n))$ ]
    lessThan-atLeast0 by auto
  hence  $\langle \exists M. \forall h \in (\text{range } (\lambda n. \text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..<n\})). \text{norm } h \leq M \rangle$ 
    using bounded-iff by blast
  then obtain  $M$  where  $\langle h \in \text{range } (\lambda n. \text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..<n\}) \implies \text{norm } h \leq M \rangle$ 
    for  $h$ 
    by blast
  have  $\langle \text{sum } (\lambda k. \text{inverse } (\text{real-of-nat } 3^k)) \{0..n\} \leq M \rangle$  for  $n$ 
proof -
  have  $\langle \text{norm } (\text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..< \text{Suc } n\}) \leq M \rangle$ 
    using  $\langle \bigwedge h. h \in (\text{range } (\lambda n. \text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..<n\})) \implies \text{norm } h \leq M \rangle$ 
    by blast
  hence  $\langle \text{norm } (\text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..n\}) \leq M \rangle$ 
    by (simp add: atLeastLessThanSuc-atLeastAtMost)
  hence  $\langle \text{sum } (\lambda k. \text{inverse } (\text{real } 3 ^ k)) \{0..n\} \leq M \rangle$ 
    by auto
  thus ?thesis by blast
qed
  have  $\langle \text{sum } (\lambda k. \text{inverse } (\text{real-of-nat } 3^k)) \{0..n\} \leq M \rangle$  for  $n$ 
    using sum-2 by blast
  thus ?thesis by blast
qed
have  $\langle \text{of-rat } 2/3 < (1::\text{real}) \rangle$ 
  by auto
hence  $\langle \forall g::'a \Rightarrow_L 'b. \forall x. \forall r. \exists \xi. g \neq 0 \wedge r > 0 \rightarrow (\xi \in \text{ball } x r \wedge (\text{of-rat } 2/3) * r * \text{norm } g \leq \text{norm } ((*_v) g \xi)) \rangle$ 
    using onorm-Sup-on-ball' by blast
hence  $\langle \exists \xi. \forall g::'a \Rightarrow_L 'b. \forall x. \forall r. g \neq 0 \wedge r > 0 \rightarrow ((\xi g x r) \in \text{ball } x r \wedge (\text{of-rat } 2/3) * r * \text{norm } g \leq \text{norm } ((*_v) g (\xi g x r))) \rangle$ 
    by metis
then obtain  $\xi$  where  $f1: \langle \llbracket g \neq 0; r > 0 \rrbracket \implies \xi g x r \in \text{ball } x r \wedge (\text{of-rat } 2/3) * r * \text{norm } g \leq \text{norm } ((*_v) g (\xi g x r)) \rangle$ 
    for  $g::'a \Rightarrow_L 'b$  and  $x$  and  $r$ 
    by blast
have  $\langle \forall n. \exists k. \text{norm } (f k) \geq 4^n \rangle$ 
    using  $\langle \neg(\text{bounded } (\text{range } f)) \rangle$  by (metis (mono-tags, opaque-lifting) boundedI image-iff linear)
  hence  $\langle \exists k. \forall n. \text{norm } (f (k n)) \geq 4^n \rangle$ 
    by metis
  hence  $\langle \exists k. \forall n. \text{norm } ((f \circ k) n) \geq 4^n \rangle$ 

```

by *simp*
 then obtain k where $\langle \text{norm } ((f \circ k) n) \geq 4^{\wedge} n \rangle$ for n
 by *blast*
 define T where $\langle T = f \circ k \rangle$
 have $\langle T n \in \text{range } f \rangle$ for n
 unfolding T -def by *simp*
 have $\langle \text{norm } (T n) \geq \text{of-nat } (4^{\wedge} n) \rangle$ for n
 unfolding T -def using $\langle \bigwedge n. \text{norm } ((f \circ k) n) \geq 4^{\wedge} n \rangle$ by *auto*
 hence $\langle T n \neq 0 \rangle$ for n
 by (*smt* T -def $\langle \bigwedge n. 4^{\wedge} n \leq \text{norm } ((f \circ k) n) \rangle$ *norm-zero power-not-zero zero-le-power*)
 have $\langle \text{inverse } (\text{of-nat } 3^{\wedge} n) > (0::\text{real}) \rangle$ for n
 by *auto*
 define $y::\langle \text{nat} \Rightarrow 'a \rangle$ where $\langle y = \text{rec-nat } 0 (\lambda n x. \xi (T n) x (\text{inverse } (\text{of-nat } 3^{\wedge} n))) \rangle$
 have $\langle y (Suc n) \in \text{ball } (y n) (\text{inverse } (\text{of-nat } 3^{\wedge} n)) \rangle$ for n
 using $f1 \langle \bigwedge n. T n \neq 0 \rangle \langle \bigwedge n. \text{inverse } (\text{of-nat } 3^{\wedge} n) > 0 \rangle$ unfolding y -def by *auto*
 hence $\langle \text{norm } (y (Suc n) - y n) \leq \text{inverse } (\text{of-nat } 3^{\wedge} n) \rangle$ for n
 unfolding ball -def apply *auto* using dist-norm by (*smt* $\text{norm-minus-commute}$)

moreover have $\langle \exists K. \forall n. \text{sum } (\lambda k. \text{inverse } (\text{real-of-nat } 3^{\wedge} k)) \{0..n\} \leq K \rangle$
 using sum-1 by *blast*
 moreover have $\langle \text{Cauchy } y \rangle$
 using $\text{convergent-series-Cauchy}$ [where $a = \lambda n. \text{inverse } (\text{of-nat } 3^{\wedge} n)$ and $\varphi = y$] dist-norm
 by (*metis* $\text{calculation}(1)$ $\text{calculation}(2)$)
 hence $\langle \exists x. y \longrightarrow x \rangle$
 by (*simp* $\text{add: convergent-eq-Cauchy}$)
 then obtain x where $\langle y \longrightarrow x \rangle$
 by *blast*
 have $\text{norm-2: } \langle \text{norm } (x - y (Suc n)) \leq (\text{inverse } (\text{of-nat } 2)) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) \rangle$ for n
 proof –
 have $\langle \text{inverse } (\text{real-of-nat } 3) < 1 \rangle$
 by *simp*
 moreover have $\langle y 0 = 0 \rangle$
 using y -def by *auto*
 ultimately have $\langle \text{norm } (x - y (Suc n)) \leq (\text{inverse } (\text{of-nat } 3)) * \text{inverse } (1 - (\text{inverse } (\text{of-nat } 3))) * ((\text{inverse } (\text{of-nat } 3))^{\wedge} n) \rangle$
 using $\text{bound-Cauchy-to-lim}$ [where $c = \text{inverse } (\text{of-nat } 3)$ and $y = y$ and $x = x$]
 power-inverse $\text{semiring-norm}(77)$ $\langle y \longrightarrow x \rangle$
 $\langle \bigwedge n. \text{norm } (y (Suc n) - y n) \leq \text{inverse } (\text{of-nat } 3^{\wedge} n) \rangle$ by (*metis* divide-inverse)
 moreover have $\langle \text{inverse } (\text{real-of-nat } 3) * \text{inverse } (1 - (\text{inverse } (\text{of-nat } 3))) = \text{inverse } (\text{of-nat } 2) \rangle$
 by *auto*

ultimately show *?thesis*
by *(metis power-inverse)*
qed
have $\langle \text{norm } (x - y \text{ (Suc } n)) \leq (\text{inverse } (\text{of-nat } 2)) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) \rangle$ **for**
 n
using *norm-2 by blast*
have $\langle \exists M. \forall n. \text{norm } ((*_v) (T n) x) \leq M \rangle$
unfolding *T-def* **apply** *auto*
by *(metis $\langle \bigwedge x. \text{bounded } (\text{range } (\lambda n. (*_v) (f n) x)) \rangle$ bounded-iff rangeI)*
then obtain M **where** $\langle \text{norm } ((*_v) (T n) x) \leq M \rangle$ **for** n
by *blast*
have *norm-1*: $\langle \text{norm } (T n) * \text{norm } (y \text{ (Suc } n) - x) + \text{norm } ((*_v) (T n) x)$
 $\leq \text{inverse } (\text{real } 2) * \text{inverse } (\text{real } 3^{\wedge} n) * \text{norm } (T n) + \text{norm } ((*_v) (T n)$
 $x) \rangle$ **for** n
proof–
have $\langle \text{norm } (y \text{ (Suc } n) - x) \leq (\text{inverse } (\text{of-nat } 2)) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) \rangle$
using $\langle \text{norm } (x - y \text{ (Suc } n)) \leq (\text{inverse } (\text{of-nat } 2)) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) \rangle$
by *(simp add: norm-minus-commute)*
moreover have $\langle \text{norm } (T n) \geq 0 \rangle$
by *auto*
ultimately have $\langle \text{norm } (T n) * \text{norm } (y \text{ (Suc } n) - x)$
 $\leq (\text{inverse } (\text{of-nat } 2)) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) * \text{norm } (T n) \rangle$
by *(simp add: $\langle \bigwedge n. T n \neq 0 \rangle$)*
thus *?thesis by simp*
qed
have *inverse-2*: $\langle (\text{inverse } (\text{of-nat } 6)) * \text{inverse } (\text{real } 3^{\wedge} n) * \text{norm } (T n)$
 $\leq \text{norm } ((*_v) (T n) x) \rangle$ **for** n
proof–
have $\langle (\text{of-rat } 2/3) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) * \text{norm } (T n) \leq \text{norm } ((*_v) (T n) (y$
 $\text{ (Suc } n))) \rangle$
using *f1* $\langle \bigwedge n. T n \neq 0 \rangle$ $\langle \bigwedge n. \text{inverse } (\text{of-nat } 3^{\wedge} n) > 0 \rangle$ **unfolding** *y-def*
by *auto*
also have $\langle \dots = \text{norm } ((*_v) (T n) ((y \text{ (Suc } n) - x) + x)) \rangle$
by *auto*
also have $\langle \dots = \text{norm } ((*_v) (T n) (y \text{ (Suc } n) - x) + (*_v) (T n) x) \rangle$
apply *transfer apply auto by (metis diff-add-cancel linear-simps(1))*
also have $\langle \dots \leq \text{norm } ((*_v) (T n) (y \text{ (Suc } n) - x)) + \text{norm } ((*_v) (T n) x) \rangle$
by *(simp add: norm-triangle-ineq)*
also have $\langle \dots \leq \text{norm } (T n) * \text{norm } (y \text{ (Suc } n) - x) + \text{norm } ((*_v) (T n) x) \rangle$
apply *transfer apply auto using onorm by auto*
also have $\langle \dots \leq (\text{inverse } (\text{of-nat } 2)) * (\text{inverse } (\text{of-nat } 3^{\wedge} n)) * \text{norm } (T n)$
 $+ \text{norm } ((*_v) (T n) x) \rangle$
using *norm-1 by blast*
finally have $\langle (\text{of-rat } 2/3) * \text{inverse } (\text{real } 3^{\wedge} n) * \text{norm } (T n)$
 $\leq \text{inverse } (\text{real } 2) * \text{inverse } (\text{real } 3^{\wedge} n) * \text{norm } (T n)$
 $+ \text{norm } ((*_v) (T n) x) \rangle$
by *blast*
hence $\langle (\text{of-rat } 2/3) * \text{inverse } (\text{real } 3^{\wedge} n) * \text{norm } (T n)$
 $- \text{inverse } (\text{real } 2) * \text{inverse } (\text{real } 3^{\wedge} n) * \text{norm } (T n) \leq \text{norm } ((*_v) (T$

$n) x\rangle$
 by *linarith*
moreover have $\langle (of-rat\ 2/3) * inverse\ (real\ 3\ ^n) * norm\ (T\ n)$
 $\quad - inverse\ (real\ 2) * inverse\ (real\ 3\ ^n) * norm\ (T\ n)$
 $\quad = (inverse\ (of-nat\ 6)) * inverse\ (real\ 3\ ^n) * norm\ (T\ n)\rangle$
 by *fastforce*
ultimately show $\langle (inverse\ (of-nat\ 6)) * inverse\ (real\ 3\ ^n) * norm\ (T\ n) \leq$
 $norm\ ((*_v)\ (T\ n)\ x)\rangle$
 by *linarith*
qed
have *inverse-3*: $\langle (inverse\ (of-nat\ 6)) * (of-rat\ (4/3)\ ^n)$
 $\quad \leq (inverse\ (of-nat\ 6)) * inverse\ (real\ 3\ ^n) * norm\ (T\ n)\rangle$ **for** n
proof-
have $\langle of-rat\ (4/3)\ ^n = inverse\ (real\ 3\ ^n) * (of-nat\ 4\ ^n)\rangle$
apply auto by (*metis divide-inverse-commute of-rat-divide power-divide*
of-rat-numeral-eq)
also have $\langle \dots \leq inverse\ (real\ 3\ ^n) * norm\ (T\ n)\rangle$
using $\langle \wedge n. norm\ (T\ n) \geq of-nat\ (4\ ^n)\rangle$ **by** *simp*
finally have $\langle of-rat\ (4/3)\ ^n \leq inverse\ (real\ 3\ ^n) * norm\ (T\ n)\rangle$
by *blast*
moreover have $\langle inverse\ (of-nat\ 6) > (0::real)\rangle$
by *auto*
ultimately show *?thesis* **by** *auto*
qed
have *inverse-1*: $\langle (inverse\ (of-nat\ 6)) * (of-rat\ (4/3)\ ^n) \leq M\rangle$ **for** n
proof-
have $\langle (inverse\ (of-nat\ 6)) * (of-rat\ (4/3)\ ^n)$
 $\quad \leq (inverse\ (of-nat\ 6)) * inverse\ (real\ 3\ ^n) * norm\ (T\ n)\rangle$
using *inverse-3* **by** *blast*
also have $\langle \dots \leq norm\ ((*_v)\ (T\ n)\ x)\rangle$
using *inverse-2* **by** *blast*
finally have $\langle (inverse\ (of-nat\ 6)) * (of-rat\ (4/3)\ ^n) \leq norm\ ((*_v)\ (T\ n)\ x)\rangle$
by *auto*
thus *?thesis* **using** $\langle \wedge n. norm\ ((*_v)\ (T\ n)\ x) \leq M\rangle$ **by** *smt*
qed
have $\langle \exists n. M < (inverse\ (of-nat\ 6)) * (of-rat\ (4/3)\ ^n)\rangle$
using *Real.real-arch-pow* **by** *auto*
moreover have $\langle (inverse\ (of-nat\ 6)) * (of-rat\ (4/3)\ ^n) \leq M\rangle$ **for** n
using *inverse-1* **by** *blast*
ultimately show *?thesis* **by** *smt*
qed

2.3 A consequence of Banach-Steinhaus theorem

corollary *bounded-linear-limit-bounded-linear*:

fixes $f::\langle nat \Rightarrow ('a::banach \Rightarrow_L 'b::real-normed-vector)\rangle$

assumes $\langle \wedge x. convergent\ (\lambda n. (f\ n)\ *_v\ x)\rangle$

shows $\langle \exists g. (\lambda n. (*_v)\ (f\ n)) -pointwise \rightarrow (*_v)\ g\rangle$

Explanation: If a sequence of bounded operators on a Banach space

converges pointwise, then the limit is also a bounded operator.

proof –

have $\langle \exists l. (\lambda n. (*_v) (f n) x) \longrightarrow l \rangle$ **for** x
by (*simp add: $\langle \bigwedge x. \text{convergent } (\lambda n. (*_v) (f n) x) \rangle \text{convergentD}$*)
hence $\langle \exists F. (\lambda n. (*_v) (f n)) \text{--pointwise--}\rightarrow F \rangle$
unfolding *pointwise-convergent-to-def* **by** *metis*
obtain F **where** $\langle (\lambda n. (*_v) (f n)) \text{--pointwise--}\rightarrow F \rangle$
using $\langle \exists F. (\lambda n. (*_v) (f n)) \text{--pointwise--}\rightarrow F \rangle$ **by** *auto*
have $\langle \bigwedge x. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle$
using $\langle (\lambda n. (*_v) (f n)) \text{--pointwise--}\rightarrow F \rangle$ **apply** *transfer*
by (*simp add: pointwise-convergent-to-def*)
have $\langle \text{bounded } (\text{range } f) \rangle$
using $\langle \bigwedge x. \text{convergent } (\lambda n. (*_v) (f n) x) \rangle$ *banach-steinhaus*
 $\langle \bigwedge x. \exists l. (\lambda n. (*_v) (f n) x) \longrightarrow l \rangle$ *convergent-imp-bounded* **by** *blast*
have *norm-f-n*: $\langle \exists M. \forall n. \text{norm } (f n) \leq M \rangle$
unfolding *bounded-def*
by (*meson UNIV-I $\langle \text{bounded } (\text{range } f) \rangle \text{bounded-iff image-eqI}$*)
have $\langle \text{isCont } (\lambda t::'b. \text{norm } t) \ y \rangle$ **for** $y::'b$
using *Limits.isCont-norm* **by** *simp*
hence $\langle (\lambda n. \text{norm } ((*_v) (f n) x)) \longrightarrow (\text{norm } (F x)) \rangle$ **for** x
using $\langle \bigwedge x::'a. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle$ **by** (*simp add: tendsto-norm*)
hence *norm-f-n-x*: $\langle \exists M. \forall n. \text{norm } ((*_v) (f n) x) \leq M \rangle$ **for** x
using *Elementary-Metric-Spaces.convergent-imp-bounded*
by (*metis UNIV-I $\langle \bigwedge x::'a. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle \text{bounded-iff image-eqI}$*)
have *norm-f*: $\langle \exists K. \forall n. \forall x. \text{norm } ((*_v) (f n) x) \leq \text{norm } x * K \rangle$
proof –
have $\langle \exists M. \forall n. \text{norm } ((*_v) (f n) x) \leq M \rangle$ **for** x
using *norm-f-n-x* $\langle \bigwedge x. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle$ **by** *blast*
hence $\langle \exists M. \forall n. \text{norm } (f n) \leq M \rangle$
using *norm-f-n* **by** *simp*
then obtain $M::\text{real}$ **where** $\langle \exists M. \forall n. \text{norm } (f n) \leq M \rangle$
by *blast*
have $\langle \forall n. \forall x. \text{norm } ((*_v) (f n) x) \leq \text{norm } x * \text{norm } (f n) \rangle$
apply *transfer* **apply** *auto* **by** (*metis mult.commute onorm*)
thus *?thesis* **using** $\langle \exists M. \forall n. \text{norm } (f n) \leq M \rangle$
by (*metis (no-types, opaque-lifting) dual-order.trans norm-eq-zero order-refl mult-le-cancel-iff2 vector-space-over-itself.scale-zero-left zero-less-norm-iff*)
qed
have *norm-F-x*: $\langle \exists K. \forall x. \text{norm } (F x) \leq \text{norm } x * K \rangle$
proof –
have $\langle \exists K. \forall n. \forall x. \text{norm } ((*_v) (f n) x) \leq \text{norm } x * K \rangle$
using *norm-f* $\langle \bigwedge x. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle$ **by** *auto*
thus *?thesis*
using $\langle \bigwedge x::'a. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle$ **apply** *transfer*
by (*metis Lim-bounded tendsto-norm*)
qed
have $\langle \text{linear } F \rangle$
proof(*rule linear-limit-linear*)

```

show  $\langle \text{linear } ((*_v) (f n)) \rangle$  for  $n$ 
  apply transfer apply auto by (simp add: bounded-linear.linear)
show  $\langle f \text{ -pointwise} \rightarrow F \rangle$ 
  using  $\langle (\lambda n. (*_v) (f n)) \text{ -pointwise} \rightarrow F \rangle$  by auto
qed
moreover have  $\langle \text{bounded-linear-axioms } F \rangle$ 
using norm-F-x by (simp add:  $\langle \bigwedge x. (\lambda n. (*_v) (f n) x) \longrightarrow F x \rangle$  bounded-linear-axioms-def)

ultimately have  $\langle \text{bounded-linear } F \rangle$ 
  unfolding bounded-linear-def by blast
hence  $\langle \exists g. (*_v) g = F \rangle$ 
  using bounded-linear-Blinfun-apply by auto
thus ?thesis using  $\langle (\lambda n. (*_v) (f n)) \text{ -pointwise} \rightarrow F \rangle$  apply transfer by auto
qed

end

```

References

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