

Actuarial Mathematics

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Abstract

Actuarial Mathematics is a theory in applied mathematics, which is mainly used for determining the prices of insurance products and evaluating the liability of a company associating with insurance contracts. It is related to calculus, probability theory and financial theory, etc.

In this entry, I formalize the very basic part of Actuarial Mathematics in Isabelle/HOL. It includes the theory of interest, survival model, and life table. The theory of interest deals with interest rates, present value factors, an annuity certain, etc. The survival model is a probabilistic model that represents the human mortality. The life table is based on the survival model and used for practical calculations.

I have already formalized the basic part of Actuarial Mathematics in Coq (<https://github.com/Yosuke-Ito-345/Actuary>) in a purely axiomatic manner. In contrast, Isabelle formalization is based on the probability theory and the survival model is developed as generally as possible. Such rigorous and general formulation seems very rare; at least I cannot find any similar documentation on the web.

This formalization in Isabelle is still at an early stage, and I cannot guarantee the backward compatibility in the future development. If you heavily depend on the “Actuarial Mathematics” library, please let me know.

Contents

1	Preliminary Lemmas	1
1.1	Lemmas on <i>indicator</i> for a Linearly Ordered Type	12
2	Additional Lemmas for the <i>HOL</i>-Analysis Library	15
2.1	Set Lebesgue Integrability on Affine Transformation	30
2.2	Set Lebesgue Integral on Affine Transformation	32
2.3	Alternative Integral Test	34
2.4	Interchange of Differentiation and Lebesgue Integration . . .	36

3 Additional Lemmas for the <i>HOL-Probability</i> Library	45
3.1 More Properties of <i>cdf</i> 's	46
3.2 Conditional Probability Space	54
3.3 Complementary Cumulative Distribution Function	59
3.4 Hazard Rate	63
4 Theory of Interest	70
5 Survival Model	88
5.1 General Theory of Survival Model	88
5.1.1 Introduction of Survival Function for X	88
5.1.2 Introduction of Future-Lifetime Random Variable $T(x)$	91
5.1.3 Actuarial Notations on the Survival Model	92
5.1.4 Properties of Survival Function for $T(x)$	93
5.1.5 Properties of $\$p\{-t&x\}$	98
5.1.6 Properties of Survival Function for X	99
5.1.7 Introduction of Cumulative Distributive Function for X	99
5.1.8 Properties of Cumulative Distributive Function for $T(x)$	100
5.1.9 Properties of $\$q\{-t&x\}$	102
5.1.10 Properties of Cumulative Distributive Function for X	104
5.1.11 Relations between $\$p\{-t&x\}$ and $\$q\{-t&x\}$	105
5.1.12 Properties of Life Expectation	106
5.2 Piecewise Differentiable Survival Function	120
5.2.1 Properties of Survival Function for X	120
5.2.2 Properties of Cumulative Distributive Function for X	121
5.2.3 Introduction of Probability Density Functions of X and $T(x)$	122
5.2.4 Properties of Survival Function for $T(x)$	123
5.2.5 Properties of Cumulative Distributive Function for $T(x)$	125
5.2.6 Properties of Probability Density Function of $T(x)$	126
5.2.7 Properties of Probability Density Function of X	132
5.2.8 Relations between Life Expectation and Probability Density Function	133
5.2.9 Introduction of Force of Mortality	136
5.2.10 Properties of Force of Mortality	138
5.2.11 Properties of Curtate Life Expectation	150
5.3 Limited Survival Function	151
6 Life Table	153
6.1 Basic Properties of Life Table	153
6.2 Construction of Survival Model from Life Table	158
6.2.1 Relations between Life Table and Survival Function for X	159

6.2.2	Relations between Life Table and Cumulative Distributive Function for X	160
6.2.3	Relations between Life Table and Survival Function for $T(x)$	161
6.2.4	Relations between Life Table and Cumulative Distributive Function for $T(x)$	163
6.2.5	Life Table and Actuarial Notations	163
6.3	Piecewise Differentiable Life Table	167
6.4	Interpolations	173
6.5	Limited Life Table	189
7	Examples	190
theory Preliminaries		
imports HOL-Library.Rewrite HOL-Library.Extended-Nonnegative-Real HOL-Library.Extended-Real HOL-Probability.Probability		
begin		
declare [[show-types]]		
notation powr (infixr $\cdot\wedge\cdot$ 80)		
1 Preliminary Lemmas		
lemma	Collect-conj-eq2: $\{x \in A. P x \wedge Q x\} = \{x \in A. P x\} \cap \{x \in A. Q x\}$	
by	blast	
lemma	vimage-compl-atMost:	
fixes	$f :: 'a \Rightarrow 'b::linorder$	
shows	$(f -` \{..y\}) = f -` \{y <..\}$	
by	fastforce	
context	linorder	
begin		
lemma	Icc-minus-Ico:	
fixes	$a b$	
assumes	$a \leq b$	
shows	$\{a..b\} - \{a..<b\} = \{b\}$	
proof		
{	fix x assume $x \in \{a..b\} - \{a..<b\}$	
hence	$x \in \{b\}$ by force }	
thus	$\{a..b\} - \{a..<b\} \subseteq \{b\}$ by blast	
next		
show	$\{b\} \subseteq \{a..b\} - \{a..<b\}$ using assms by simp	
qed		
lemma	Icc-minus-Ioc:	
fixes	$a b$	
assumes	$a \leq b$	

```

shows  $\{a..b\} - \{a<..b\} = \{a\}$ 
proof
  { fix x assume  $x \in \{a..b\} - \{a<..b\}$ 
    hence  $x \in \{a\}$  by force }
    thus  $\{a..b\} - \{a<..b\} \subseteq \{a\}$  by blast
next
  show  $\{a\} \subseteq \{a..b\} - \{a<..b\}$  using assms by simp
qed

lemma Int-atLeastAtMost-Unbounded[simp]:  $\{a..\} \text{ Int } \{..b\} = \{a..b\}$ 
  by auto

lemma Int-greaterThanAtMost-Unbounded[simp]:  $\{a<..\} \text{ Int } \{..b\} = \{a<..b\}$ 
  by auto

lemma Int-atLeastLessThan-Unbounded[simp]:  $\{a..\} \text{ Int } \{..<b\} = \{a..<b\}$ 
  by auto

lemma Int-greaterThanLessThan-Unbounded[simp]:  $\{a<..\} \text{ Int } \{..<b\} = \{a<..<b\}$ 
  by auto

end

lemma Ico-real-nat-disjoint:
  disjoint-family  $(\lambda n::nat. \{a + real n .. < a + real n + 1\})$  for  $a::real$ 
proof -
  { fix m n :: nat
    assume  $\{a + real m .. < a + real m + 1\} \cap \{a + real n .. < a + real n + 1\} \neq \{\}$ 
    then obtain x::real
      where  $x \in \{a + real m .. < a + real m + 1\} \cap \{a + real n .. < a + real n + 1\}$  by force
      hence  $m = n$  by simp }
    thus ?thesis unfolding disjoint-family-on-def by blast
qed

corollary Ico-nat-disjoint: disjoint-family  $(\lambda n::nat. \{real n .. < real n + 1\})$ 
  using Ico-real-nat-disjoint[of 0] by simp

lemma Ico-real-nat-union:  $(\bigcup n::nat. \{a + real n .. < a + real n + 1\}) = \{a..\}$ 
for  $a::real$ 
proof
  show  $(\bigcup n::nat. \{a + real n .. < a + real n + 1\}) \subseteq \{a..\}$  by auto
next
  show  $\{a..\} \subseteq (\bigcup n::nat. \{a + real n .. < a + real n + 1\})$ 
proof
  fix x assume  $x \in \{a..\}$ 
  hence  $a \leq x$  by simp

```

```

hence nat ⌊x-a⌋ ≤ x-a ∧ x-a < nat ⌊x-a⌋ + 1 by linarith
hence a + nat ⌊x-a⌋ ≤ x ∧ x < a + nat ⌊x-a⌋ + 1 by auto
thus x ∈ (⋃ n::nat. {a + real n ..< a + real n + 1}) by auto
qed
qed

```

```

corollary Ico-nat-union: (⋃ n::nat. {real n ..< real n + 1}) = {0..}
using Ico-real-nat-union[of 0] by simp

```

```

lemma Ico-nat-union-finite: (⋃ (n::nat)<m. {real n ..< real n + 1}) = {0..<m}
proof
  show (⋃ (n::nat)<m. {real n ..< real n + 1}) ⊆ {0..<m} by auto
next
  show {0..<m} ⊆ (⋃ (n::nat)<m. {real n ..< real n + 1})
proof
  fix x::real
  assume ∗: x ∈ {0..<m}
  hence nat ⌊x⌋ < m using of-nat-floor by fastforce
  moreover with ∗ have nat ⌊x⌋ ≤ x ∧ x < nat ⌊x⌋ + 1
    by (metis Nat.add-0-right atLeastLessThan-iff le-nat-floor
      less-one linorder-not-le nat-add-left-cancel-le of-nat-floor)
  ultimately show x ∈ (⋃ (n::nat)<m. {real n ..< real n + 1})
    unfolding atLeastLessThan-def by force
qed
qed

```

```

lemma seq-part-multiple: fixes m n :: nat assumes m ≠ 0 defines A ≡ λi::nat.
{i*m ..< (i+1)*m}
shows ∀ i j. i ≠ j → A i ∩ A j = {} and (⋃ i<n. A i) = {..< n*m}
proof –
  { fix i j :: nat
    have i ≠ j → A i ∩ A j = {}
    proof (erule contrapos-np)
      assume A i ∩ A j ≠ {}
      then obtain k where k ∈ A i ∩ A j by blast
      hence i*m < (j+1)*m ∧ j*m < (i+1)*m unfolding A-def by force
      hence i < j+1 ∧ j < i+1 using mult-less-cancel2 by blast
      thus i = j by force
    qed }
    thus ∀ i j. i ≠ j → A i ∩ A j = {} by blast
next
  show (⋃ i<n. A i) = {..< n*m}
  proof
    show (⋃ i<n. A i) ⊆ {..< n*m}
    proof
      fix x::nat
      assume x ∈ (⋃ i<n. A i)
      then obtain i where i-n: i < n and i-x: x < (i+1)*m unfolding A-def by
      force
    
```

```

hence  $i+1 \leq n$  by linarith
hence  $x < n*m$  by (meson less-le-trans mult-le-cancel2 i-x)
thus  $x \in \{.. < n*m\}$ 
      using diff-mult-distrib mult-1 i-n by auto
qed
next
show  $\{.. < n*m\} \subseteq (\bigcup_{i < n} A_i)$ 
proof
fix  $x::nat$ 
let  $?i = x \text{ div } m$ 
assume  $x \in \{.. < n*m\}$ 
hence  $?i < n$  by (simp add: less-mult-imp-div-less)
moreover have  $?i*m \leq x \wedge x < (?i+1)*m$ 
      using assms div-times-less-eq-dividend dividend-less-div-times by auto
ultimately show  $x \in (\bigcup_{i < n} A_i)$  unfolding A-def by force
qed
qed
qed

lemma frontier-Icc-real: frontier {a..b} = {a, b} if a ≤ b for a b :: real
unfolding frontier-def using that by force

lemma(in field) divide-mult-cancel[simp]: fixes a b assumes b ≠ 0
shows a / b * b = a
by (simp add: assms)

lemma inverse-powr: (1/a).^b = a.^b - b if a > 0 for a b :: real
by (smt (verit) that powr-divide powr-minus-divide powr-one-eq-one)

lemma powr-eq-one-iff-gen[simp]: a.^x = 1 ↔ x = 0 if a > 0 a ≠ 1 for a x :: real
by (metis powr-eq-0-iff powr-inj powr-zero-eq-one that)

lemma powr-less-cancel2: 0 < a ⇒ 0 < x ⇒ 0 < y ⇒ x.^a < y.^a ⇒ x < y
for a x y ::real
proof -
assume a-pos: 0 < a and x-pos: 0 < x and y-pos: 0 < y
show x.^a < y.^a ⇒ x < y
proof (erule contrapos-pp)
assume ¬ x < y
hence x ≥ y by fastforce
hence x.^a ≥ y.^a
proof (cases x = y)
case True
thus ?thesis by simp
next
case False
hence x.^a > y.^a

```

```

using ‹x ≥ y› powr-less-mono2 a-pos y-pos by auto
thus ?thesis by auto
qed
thus ¬ x. ^a < y. ^a by fastforce
qed
qed

lemma geometric-increasing-sum-aux: (1 - r) ^ 2 * (∑ k < n. (k + 1) * r ^ k) = 1 - (n + 1) * r ^ n + n * r ^ (n + 1)
  for n::nat and r::real
proof (induct n)
  case 0
  thus ?case by simp
next
  case (Suc n)
  thus ?case
    apply (simp only: sum.lessThan-Suc)
    apply (subst distrib-left)
    apply (subst Suc.hyps)
    by (subst power2-diff, simp add: field-simps power2-eq-square)
qed

lemma geometric-increasing-sum: (∑ k < n. (k + 1) * r ^ k) = (1 - (n + 1) * r ^ n + n * r ^ (n + 1)) / (1 - r) ^ 2
  if r ≠ 1 for n::nat and r::real
  by (subst geometric-increasing-sum-aux[THEN sym], simp add: that)

lemma Reals-UNIV[simp]: ℝ = {x::real. True}
  unfolding Reals-def by auto

lemma Lim-cong:
  assumes ∀ F x in F. f x = g x
  shows Lim F f = Lim F g
  unfolding t2-space-class.Lim-def using tendsto-cong assms by fastforce

lemma LIM-zero-iff': ((λx. l - f x) —→ 0) F = (f —→ l) F
  for f :: 'a ⇒ 'b::real-normed-vector
  unfolding tendsto-iff dist-norm
  by (rewrite minus-diff-eq[THEN sym], rewrite norm-minus-cancel) simp

lemma antimono-onI:
  (∀r s. r ∈ A ⇒ s ∈ A ⇒ r ≤ s ⇒ f r ≥ f s) ⇒ antimono-on A f
  by (rule monotone-onI)

lemma antimono-onD:
  [antimono-on A f; r ∈ A; s ∈ A; r ≤ s] ⇒ f r ≥ f s
  by (rule monotone-onD)

lemma antimono-imp-mono-on: antimono f ⇒ antimono-on A f

```

```

by (simp add: antimonoD antimono-onI)

lemma antimono-on-subset: antimono-on A f ==> B ⊆ A ==> antimono-on B f
  by (rule monotone-on-subset)

lemma mono-on-antimono-on:
  fixes f :: 'a::order => 'b::ordered-ab-group-add
  shows mono-on A f <=> antimono-on A (λr. - f r)
  by (simp add: monotone-on-def)

corollary mono-antimono:
  fixes f :: 'a::order => 'b::ordered-ab-group-add
  shows mono f <=> antimono (λr. - f r)
  by (rule mono-on-antimono-on)

lemma mono-on-at-top-le:
  fixes a :: 'a::linorder and b :: 'b::{order-topology, linordered-ab-group-add}
  and f :: 'a => 'b
  assumes f-mono: mono-on {a..} f and f-to-l: (f —> l) at-top
  shows ∀x. x ∈ {a..} ==> f x ≤ l
  proof (unfold atomize-ball)
    { fix b assume b-a: b ≥ a and fb-l: ¬ f b ≤ l
      let ?eps = f b - l
      have lim-top: ∀S. open S ==> l ∈ S ==> eventually (λx. f x ∈ S) at-top
        using assms tendsto-def by auto
      have eventually (λx. f x ∈ {l - ?eps <..< l + ?eps}) at-top
        using fb-l by (intro lim-top; force)
      then obtain N where fn-in: ∀n. n ≥ N ==> f n ∈ {l - ?eps <..< l + ?eps}
        using eventually-at-top-linorder by metis
      let ?n = max b N
      have f ?n < f b using fn-in by force
      moreover have f ?n ≥ f b using f-mono b-a by (simp add: le-max-iff-disj
mono-on-def)
        ultimately have False by simp }
      thus ∀x∈{a..}. f x ≤ l
        apply -
        by (rule notnotD, rewrite Set.ball-simps) auto
    qed

corollary mono-at-top-le:
  fixes b :: 'b::{order-topology, linordered-ab-group-add} and f :: 'a::linorder => 'b
  assumes mono f and (f —> b) at-top
  shows ∀x. f x ≤ b
  using mono-on-at-top-le assms by (metis atLeast-iff mono-imp-mono-on nle-le)

lemma mono-on-at-bot-ge:
  fixes a :: 'a::linorder and b :: 'b::{order-topology, linordered-ab-group-add}
  and f :: 'a => 'b
  assumes f-mono: mono-on {..a} f and f-to-l: (f —> l) at-bot

```

```

shows  $\bigwedge x. x \in \{..a\} \implies f x \geq l$ 
proof (unfold atomize-ball)
{ fix b assume b-a:  $b \leq a$  and fb-l:  $\neg f b \geq l$ 
let ?eps =  $l - f b$ 
have lim-bot:  $\bigwedge S. \text{open } S \implies l \in S \implies \text{eventually } (\lambda x. f x \in S) \text{ at-bot}$ 
using assms tendsto-def by auto
have eventually ( $\lambda x. f x \in \{l - ?\text{eps} <.. < l + ?\text{eps}\}$ ) at-bot
using fb-l by (intro lim-bot; force)
then obtain N where fn-in:  $\bigwedge n. n \leq N \implies f n \in \{l - ?\text{eps} <.. < l + ?\text{eps}\}$ 
using eventually-at-bot-linorder by metis
let ?n = min b N
have  $f ?n > f b$  using fn-in by force
moreover have  $f ?n \leq f b$  using f-mono b-a by (simp add: min.coboundedI1
mono-onD)
ultimately have False by simp }
thus  $\forall x \in \{..a\}. f x \geq l$ 
apply –
by (rule notnotD, rewrite Set.ball-simps) auto
qed

```

```

corollary mono-at-bot-ge:
fixes b :: 'b::{'order-topology, linordered-ab-group-add} and f :: 'a::linorder  $\Rightarrow$  'b
assumes mono f and (f —> b) at-bot
shows  $\bigwedge x. f x \geq b$ 
using mono-on-at-bot-ge assms by (metis atMost-iff mono-imp-mono-on nle-le)

```

```

lemma antimono-on-at-top-ge:
fixes a :: 'a::linorder and b :: 'b::{'order-topology, linordered-ab-group-add}
and f :: 'a  $\Rightarrow$  'b
assumes f-antimono: antimono-on {a..} f and f-to-l: (f —> l) at-top
shows  $\bigwedge x. x \in \{a..\} \implies f x \geq l$ 
proof (unfold atomize-ball)
{ fix b assume b-a:  $b \geq a$  and fb-l:  $\neg f b \geq l$ 
let ?eps =  $l - f b$ 
have lim-top:  $\bigwedge S. \text{open } S \implies l \in S \implies \text{eventually } (\lambda x. f x \in S) \text{ at-top}$ 
using assms tendsto-def by auto
have eventually ( $\lambda x. f x \in \{l - ?\text{eps} <.. < l + ?\text{eps}\}$ ) at-top
using fb-l by (intro lim-top; force)
then obtain N where fn-in:  $\bigwedge n. n \geq N \implies f n \in \{l - ?\text{eps} <.. < l + ?\text{eps}\}$ 
using eventually-at-top-linorder by metis
let ?n = max b N
have  $f ?n > f b$  using fn-in by force
moreover have  $f ?n \leq f b$  using f-antimono b-a
by (simp add: antimono-onD le-max-iff-disj)
ultimately have False by simp }
thus  $\forall x \in \{a..\}. f x \geq l$ 
apply –
by (rule notnotD, rewrite Set.ball-simps) auto
qed

```

corollary *antimono-at-top-le*:

fixes $b :: 'b : \{order-topology, linordered-ab-group-add\}$ and $f :: 'a :: linorder \Rightarrow 'b$
assumes *antimono f* and $(f \longrightarrow b)$ at-top
shows $\bigwedge x. f x \geq b$
using *antimono-on-at-top-ge assms antimono-imp-mono-on* by *blast*

lemma *antimono-on-at-bot-ge*:

fixes $a :: 'a :: linorder$ and $b :: 'b : \{order-topology, linordered-ab-group-add\}$
and $f :: 'a \Rightarrow 'b$
assumes *f-antimono: antimono-on {..a} f* and *f-to-l: (f → l) at-bot*
shows $\bigwedge x. x \in \{..a\} \implies f x \leq l$
proof (*unfold atomize-ball*)
{ fix b assume $b-a: b \leq a$ and $fb-l: \neg f b \leq l$
let $?eps = f b - l$
have *lim-bot: ∏S. open S → l ∈ S → eventually (λx. f x ∈ S) at-bot*
using *assms tendsto-def* by *auto*
have *eventually (λx. f x ∈ {l - ?eps <..< l + ?eps}) at-bot*
using *fb-l* by (*intro lim-bot; force*)
then obtain N where *fn-in: ∏n. n ≤ N → f n ∈ {l - ?eps <..< l + ?eps}*
using *eventually-at-bot-linorder* by *metis*
let $?n = min b N$
have $f ?n < f b$ using *fn-in* by *force*
moreover have $f ?n \geq f b$ using *f-antimono b-a* by (*simp add: min.coboundedI1 antimono-onD*)
ultimately have *False* by *simp* }
thus $\forall x \in \{..a\}. f x \leq l$
apply –
by (*rule notnotD, rewrite Set.ball-simps*) *auto*
qed

corollary *antimono-at-bot-ge*:

fixes $b :: 'b : \{order-topology, linordered-ab-group-add\}$ and $f :: 'a :: linorder \Rightarrow 'b$
assumes *antimono f* and $(f \longrightarrow b)$ at-bot
shows $\bigwedge x. f x \leq b$
using *antimono-on-at-bot-ge assms antimono-imp-mono-on* by *blast*

lemma *continuous-cdivide*:

fixes $c :: 'a :: real-normed-field$
assumes $c \neq 0$ continuous $F f$
shows continuous $F (\lambda x. f x / c)$
using *assms continuous-mult-right* by (*rewrite field-class.field-divide-inverse*)
auto

lemma *continuous-mult-left-iff*:

fixes $c :: 'a :: real-normed-field$
assumes $c \neq 0$
shows continuous $F f \longleftrightarrow$ continuous $F (\lambda x. c * f x)$
using *continuous-mult-left continuous-cdivide assms* by *force*

```

lemma continuous-mult-right-iff:
  fixes c::'a::real-normed-field
  assumes c ≠ 0
  shows continuous F f ↔ continuous F (λx. f x * c)
  using continuous-mult-right continuous-cdivide assms by force

lemma continuous-cdivide-iff:
  fixes c::'a::real-normed-field
  assumes c ≠ 0
  shows continuous F f ↔ continuous F (λx. f x / c)
proof
  assume continuous F f
  thus continuous F (λx. f x / c)
    by (intro continuous-cdivide) (simp add: assms)
next
  assume continuous F (λx. f x / c)
  hence continuous F (λx. f x / c * c) using continuous-mult-right by fastforce
  thus continuous F f using assms by force
qed

lemma continuous-cong:
  assumes eventually (λx. f x = g x) F f (Lim F (λx. x)) = g (Lim F (λx. x))
  shows continuous F f ↔ continuous F g
  unfolding continuous-def using assms filterlim-cong by force

lemma continuous-at-within-cong:
  assumes f x = g x eventually (λx. f x = g x) (at x within s)
  shows continuous (at x within s) f ↔ continuous (at x within s) g
proof (cases ⟨x ∈ closure (s - {x})⟩)
  case True
  thus ?thesis
    using assms apply (intro continuous-cong, simp)
    by (rewrite Lim-ident-at, simp add: at-within-eq-bot-iff)+ simp
next
  case False
  hence at x within s = bot using not-trivial-limit-within by blast
  thus ?thesis by simp
qed

lemma continuous-within-shift:
  fixes a x :: 'a :: {topological-ab-group-add, t2-space}
  and s :: 'a set
  and f :: 'a ⇒ 'b::topological-space
  shows continuous (at x within s) (λx. f (x+a)) ↔ continuous (at (x+a) within plus a ` s) f
proof
  assume continuous (at x within s) (λx. f (x+a))
  moreover have continuous (at (x+a) within plus a ` s) (plus (-a))

```

```

    by (simp add: continuous-at-imp-continuous-at-within)
moreover have plus (-a) ` plus a ` s = s by force
ultimately show continuous (at (x+a) within plus a ` s) f
    using continuous-within-compose unfolding comp-def by force
next
assume continuous (at (x+a) within plus a ` s) f
moreover have continuous (at x within s) (plus a)
    by (simp add: continuous-at-imp-continuous-at-within)
ultimately show continuous (at x within s) (λx. f (x+a))
    using continuous-within-compose unfolding comp-def by (force simp add:
add.commute)
qed

lemma isCont-shift:
fixes a x :: 'a :: {topological-ab-group-add, t2-space}
and f :: 'a ⇒ 'b::topological-space
shows isCont (λx. f (x+a)) x ↔ isCont f (x+a)
using continuous-within-shift by force

lemma has-real-derivative-at-split:
(f has-real-derivative D) (at x) ↔
(f has-real-derivative D) (at-left x) ∧ (f has-real-derivative D) (at-right x)
proof –
have at x = at x within ({..

```

```

by (rewrite DERIV-cmult-iff[of c], simp-all add: assms mult-ac)

lemma DERIV-cdivide-iff:
assumes c ≠ 0
shows (f has-field-derivative D) (at x within s) ↔
((λx. f x / c) has-field-derivative D / c) (at x within s)
apply (rewrite field-class.field-divide-inverse)++
using DERIV-cmult-right-iff assms inverse-nonzero-iff-nonzero by blast

lemma DERIV-ln-divide-chain:
fixes f :: real ⇒ real
assumes f x > 0 and (f has-real-derivative D) (at x within s)
shows ((λx. ln (f x)) has-real-derivative (D / f x)) (at x within s)
proof –
have DERIV ln (f x) :> 1 / f x using assms by (intro DERIV-ln-divide) blast
thus ?thesis using DERIV-chain2 assms by fastforce
qed

lemma inverse-fun-has-integral-ln:
fixes f :: real ⇒ real and f' :: real ⇒ real
assumes a ≤ b and
      ∀x. x ∈ {a..b} ⇒ f x > 0 and
      continuous-on {a..b} f and
      ∀x. x ∈ {a <.. < b} ⇒ (f has-real-derivative f' x) (at x)
shows ((λx. f' x / f x) has-integral (ln (f b) - ln (f a))) {a..b}
proof –
have continuous-on {a..b} (λx. ln (f x)) using assms by (intro continuous-intros; force)
moreover have ∀x. x ∈ {a <.. < b} ⇒ ((λx. ln (f x)) has-vector-derivative f' x / f x) (at x)
apply (rewrite has-real-derivative-iff-has-vector-derivative[THEN sym])
using assms by (intro DERIV-ln-divide-chain; simp)
ultimately show ?thesis using assms by (intro fundamental-theorem-of-calculus-interior; simp)
qed

lemma DERIV-fun-powr2:
fixes a::real
assumes a-pos: a > 0
and f: DERIV f x :> r
shows DERIV (λx. a.^(f x)) x :> a.^(f x) * r * ln a
proof –
let ?g = (λx. a)
have g: DERIV ?g x :> 0 by simp
have pos: ?g x > 0 by (simp add: a-pos)
show ?thesis
using DERIV-powr[OF g pos f] a-pos by (auto simp add: field-simps)
qed

```

```

lemma has-real-derivative-powr2:
  assumes a-pos: a > 0
  shows (( $\lambda x. a \cdot x$ ) has-real-derivative a.  $\hat{x} * \ln a$ ) (at x)
proof -
  let ?f = ( $\lambda x. x :: \text{real}$ )
  have f: DERIV ?f x :> 1 by simp
  thus ?thesis using DERIV-fun-powr2[OF a-pos f] by simp
qed

lemma field-differentiable-shift:
  ( $f$  field-differentiable (at  $(x + z)$ )) = (( $\lambda x. f(x + z)$ ) field-differentiable (at x))
  unfolding field-differentiable-def using DERIV-shift by force

```

1.1 Lemmas on indicator for a Linearly Ordered Type

```

lemma indicator-Icc-shift:
  fixes a b t x :: 'a::linordered-ab-group-add
  shows indicator {a..b} x = indicator {t+a..t+b} (t+x)
  by (simp add: indicator-def)

lemma indicator-Icc-shift-inverse:
  fixes a b t x :: 'a::linordered-ab-group-add
  shows indicator {a-t..b-t} x = indicator {a..b} (t+x)
  by (metis add.commute diff-add-cancel indicator-Icc-shift)

lemma indicator-Ici-shift:
  fixes a t x :: 'a::linordered-ab-group-add
  shows indicator {a..} x = indicator {t+a..} (t+x)
  by (simp add: indicator-def)

lemma indicator-Ici-shift-inverse:
  fixes a t x :: 'a::linordered-ab-group-add
  shows indicator {a-t..} x = indicator {a..} (t+x)
  by (metis add.commute diff-add-cancel indicator-Ici-shift)

lemma indicator-Iic-shift:
  fixes b t x :: 'a::linordered-ab-group-add
  shows indicator {..b} x = indicator {..t+b} (t+x)
  by (simp add: indicator-def)

lemma indicator-Iic-shift-inverse:
  fixes b t x :: 'a::linordered-ab-group-add
  shows indicator {..b-t} x = indicator {..b} (t+x)
  by (metis add.commute diff-add-cancel indicator-Iic-shift)

lemma indicator-Icc-reverse:
  fixes a b t x :: 'a::linordered-ab-group-add
  shows indicator {a..b} x = indicator {t-b..t-a} (t-x)

```

```

by (metis add-le-cancel-left atLeastAtMost-iff diff-add-cancel indicator-simps le-diff-eq)

lemma indicator-Icc-reverse-inverse:
  fixes a b t x :: 'a::linordered-ab-group-add
  shows indicator {t-b..t-a} x = indicator {a..b} (t-x)
  by (metis add-diff-cancel-left' diff-add-cancel indicator-Icc-reverse)

lemma indicator-Ici-reverse:
  fixes a t x :: 'a::linordered-ab-group-add
  shows indicator {a..} x = indicator {..t-a} (t-x)
  by (simp add: indicator-def)

lemma indicator-Ici-reverse-inverse:
  fixes b t x :: 'a::linordered-ab-group-add
  shows indicator {t-b..} x = indicator {..b} (t-x)
  by (metis add-diff-cancel-left' diff-add-cancel indicator-Ici-reverse)

lemma indicator-Iic-reverse:
  fixes b t x :: 'a::linordered-ab-group-add
  shows indicator {..b} x = indicator {t-b..} (t-x)
  by (simp add: indicator-def)

lemma indicator-Iic-reverse-inverse:
  fixes a t x :: 'a::linordered-field
  shows indicator {..t-a} x = indicator {a..} (t-x)
  by (simp add: indicator-def)

lemma indicator-Icc-affine-pos:
  fixes a b c t x :: 'a::linordered-field
  assumes c > 0
  shows indicator {a..b} x = indicator {t+c*a..t+c*b} (t + c*x)
  unfolding indicator-def using assms by simp

lemma indicator-Icc-affine-pos-inverse:
  fixes a b c t x :: 'a::linordered-field
  assumes c > 0
  shows indicator {(a-t)/c..(b-t)/c} x = indicator {a..b} (t + c*x)
  using indicator-Icc-affine-pos[where a=(a-t)/c and b=(b-t)/c and c=c and
  t=t] assms by simp

lemma indicator-Ici-affine-pos:
  fixes a c t x :: 'a::linordered-field
  assumes c > 0
  shows indicator {a..} x = indicator {t+c*a..} (t + c*x)
  unfolding indicator-def using assms by simp

lemma indicator-Ici-affine-pos-inverse:
  fixes a c t x :: 'a::linordered-field
  assumes c > 0

```

```

shows indicator {(a-t)/c..} x = indicator {a..} (t + c*x)
using indicator-Ici-affine-pos[where a=(a-t)/c and c=c and t=t] assms by
simp

lemma indicator-Iic-affine-pos:
  fixes b c t x :: 'a::linordered-field
  assumes c > 0
  shows indicator {..b} x = indicator {..t+c*b} (t + c*x)
  unfolding indicator-def using assms by simp

lemma indicator-Iic-affine-pos-inverse:
  fixes b c t x :: 'a::linordered-field
  assumes c > 0
  shows indicator {..(b-t)/c} x = indicator {..b} (t + c*x)
  using indicator-Iic-affine-pos[where b=(b-t)/c and c=c and t=t] assms by
simp

lemma indicator-Icc-affine-neg:
  fixes a b c t x :: 'a::linordered-field
  assumes c < 0
  shows indicator {a..b} x = indicator {t+c*b..t+c*a} (t + c*x)
  unfolding indicator-def using assms by auto

lemma indicator-Icc-affine-neg-inverse:
  fixes a b c t x :: 'a::linorderered-field
  assumes c < 0
  shows indicator {((b-t)/c..(a-t)/c)} x = indicator {a..b} (t + c*x)
  using indicator-Icc-affine-neg[where a=(b-t)/c and b=(a-t)/c and c=c and
t=t] assms by simp

lemma indicator-Ici-affine-neg:
  fixes a c t x :: 'a::linordered-field
  assumes c < 0
  shows indicator {a..} x = indicator {..t+c*a} (t + c*x)
  unfolding indicator-def using assms by simp

lemma indicator-Ici-affine-neg-inverse:
  fixes b c t x :: 'a::linorderded-field
  assumes c < 0
  shows indicator {((b-t)/c..)} x = indicator {..b} (t + c*x)
  using indicator-Ici-affine-neg[where a=(b-t)/c and c=c and t=t] assms by
simp

lemma indicator-Iic-affine-neg:
  fixes b c t x :: 'a::linorderded-field
  assumes c < 0
  shows indicator {..b} x = indicator {t+c*b..} (t + c*x)
  unfolding indicator-def using assms by simp

```

```

lemma indicator-Iic-affine-neg-inverse:
  fixes a c t x :: 'a::linordered-field
  assumes c < 0
  shows indicator {..(a-t)/c} x = indicator {a..} (t + c*x)
  using indicator-Iic-affine-neg[where b=(a-t)/c and c=c and t=t] assms by
  simp

```

2 Additional Lemmas for the HOL–Analysis Library

```

lemma differentiable-eq-field-differentiable-real:
  fixes f :: real  $\Rightarrow$  real
  shows f differentiable F  $\longleftrightarrow$  f field-differentiable F
  unfolding field-differentiable-def differentiable-def has-real-derivative
  using has-real-derivative-iff by presburger

```

```

lemma differentiable-on-eq-field-differentiable-real:
  fixes f :: real  $\Rightarrow$  real
  shows f differentiable-on s  $\longleftrightarrow$  ( $\forall x \in s$ . f field-differentiable (at x within s))
  unfolding differentiable-on-def using differentiable-eq-field-differentiable-real by
  simp

```

```

lemma differentiable-on-cong :
  assumes  $\bigwedge x$ .  $x \in s \implies f x = g x$  and f differentiable-on s
  shows g differentiable-on s
  proof –
    { fix x assume  $x \in s$ 
      hence f differentiable at x within s using assms unfolding differentiable-on-def
      by simp
      from this obtain D where (f has-derivative D) (at x within s)
      unfolding differentiable-def by blast
      hence (g has-derivative D) (at x within s)
      using has-derivative-transform assms  $\langle x \in s \rangle$  by metis
      hence g differentiable at x within s unfolding differentiable-def by blast }
      hence  $\forall x \in s$ . g differentiable at x within s by simp
      thus ?thesis unfolding differentiable-on-def by simp
  qed

```

```

lemma C1-differentiable-imp-deriv-continuous-on:
  f C1-differentiable-on S  $\implies$  continuous-on S (deriv f)
  using C1-differentiable-on-eq field-derivative-eq-vector-derivative by auto

```

```

lemma deriv-shift:
  assumes f field-differentiable at (x+a)
  shows deriv f (x+a) = deriv ( $\lambda s$ . f (x+s)) a
  proof –
    have (f has-field-derivative deriv f (x+a)) (at (x+a))
    using DERIV-deriv-iff-field-differentiable assms
    by force
    hence (( $\lambda s$ . f (x+s)) has-field-derivative deriv f (x+a)) (at a)

```

```

using DERIV-at-within-shift has-field-derivative-at-within by blast
moreover have (( $\lambda s. f(x+s)$ ) has-field-derivative deriv ( $\lambda s. f(x+s)$ ) a) (at a)
  using DERIV-imp-deriv calculation by fastforce
ultimately show ?thesis using DERIV-unique by force
qed

lemma piecewise-differentiable-on-cong:
assumes f piecewise-differentiable-on i
  and  $\bigwedge x. x \in i \implies f x = g x$ 
shows g piecewise-differentiable-on i
proof -
  have continuous-on i g
  using continuous-on-cong-simp assms piecewise-differentiable-on-imp-continuous-on
by force
moreover have  $\exists S. \text{finite } S \wedge (\forall x \in i - S. g \text{ differentiable (at } x \text{ within } i))$ 
proof -
  from assms piecewise-differentiable-on-def
  obtain S where fin: finite S and  $\forall x \in i - S. f \text{ differentiable (at } x \text{ within } i)$ 
by metis
  hence  $\bigwedge x. x \in i - S \implies f \text{ differentiable (at } x \text{ within } i)$  by simp
  hence  $\bigwedge x. x \in i - S \implies g \text{ differentiable (at } x \text{ within } i)$ 
  using has-derivative-transform assms by (metis DiffD1 differentiable-def)
  with fin show ?thesis by blast
qed
ultimately show ?thesis unfolding piecewise-differentiable-on-def by simp
qed

lemma differentiable-piecewise:
assumes continuous-on i f
  and f differentiable-on i
shows f piecewise-differentiable-on i
unfolding piecewise-differentiable-on-def using assms differentiable-onD by auto

lemma piecewise-differentiable-scaleR:
assumes f piecewise-differentiable-on S
shows  $(\lambda x. a *_R f x)$  piecewise-differentiable-on S
proof -
  from assms obtain T where fin: finite T  $\bigwedge x. x \in S - T \implies f \text{ differentiable at } x \text{ within } S$ 
  unfolding piecewise-differentiable-on-def by blast
  hence  $\bigwedge x. x \in S - T \implies (\lambda x. a *_R f x) \text{ differentiable at } x \text{ within } S$ 
  using differentiable-scaleR by simp
  moreover have continuous-on S  $(\lambda x. a *_R f x)$ 
  using assms continuous-on-scaleR continuous-on-const piecewise-differentiable-on-def
by blast
ultimately show  $(\lambda x. a *_R f x)$  piecewise-differentiable-on S
  using fin piecewise-differentiable-on-def by blast
qed

```

```

lemma differentiable-on-piecewise-compose:
  assumes f piecewise-differentiable-on S
    and g differentiable-on f ` S
  shows g ∘ f piecewise-differentiable-on S
proof –
  from assms obtain T where fin: finite T ∧ x. x ∈ S − T ⇒ f differentiable
  at x within S
    unfolding piecewise-differentiable-on-def by blast
    hence ∀x. x ∈ S − T ⇒ g ∘ f differentiable at x within S
      by (meson DiffD1 assms differentiable-chain-within differentiable-onD im-
      age-eqI)
    hence ∃T. finite T ∧ (∀x∈S−T. g ∘ f differentiable at x within S) using fin
      by blast
    moreover have continuous-on S (g ∘ f)
      using assms continuous-on-compose differentiable-imp-continuous-on
      unfolding piecewise-differentiable-on-def by blast
      ultimately show ?thesis
      unfolding piecewise-differentiable-on-def by force
qed

lemma MVT-order-free:
  fixes r h :: real
  defines I ≡ {r..r+h} ∪ {r+h..r}
  assumes continuous-on I f and f differentiable-on interior I
  obtains t where t ∈ {0<..<1} and f (r+h) − f r = h * deriv f (r + t*h)
proof –
  consider (h-pos) h > 0 | (h-0) h = 0 | (h-neg) h < 0 by force
  thus ?thesis
  proof cases
    case h-pos
    hence r < r+h I = {r..r+h} unfolding I-def by simp-all
    moreover hence interior I = {r<..<r+h} by simp
    moreover hence ∀x. [r < x; x < r+h] ⇒ f differentiable (at x)
      using assms by (simp add: differentiable-on-eq-differentiable-at)
    ultimately obtain z where r < z ∧ z < r+h ∧ f (r+h) − f r = h * deriv f z
      using MVT assms by (smt (verit) DERIV-imp-deriv)
    moreover hence (z-r) / h ∈ {0<..<1} by simp
    moreover have z = r + (z-r)/h * h using h-pos by simp
    ultimately show ?thesis using that by presburger
  next
    case h-0
    have 1/2 ∈ {0::real<..<1} by simp
    moreover have f (r+h) − f r = 0 using h-0 by simp
    moreover have h * deriv f (r + (1/2)*h) = 0 using h-0 by simp
    ultimately show ?thesis using that by presburger
  next case h-neg
    hence r+h < r I = {r+h..r} unfolding I-def by simp-all
    moreover hence interior I = {r+h<..<r} by simp
    moreover hence ∀x. [r+h < x; x < r] ⇒ f differentiable (at x)

```

```

using assms by (simp add: differentiable-on-eq-differentiable-at)
ultimately obtain z where r+h < z ∧ z < r ∧ f r - f (r+h) = -h * deriv
f z
  using MVT assms by (smt (verit) DERIV-imp-deriv)
  moreover hence (z-r) / h ∈ {0<..<1} by (simp add: neg-less-divide-eq)
  moreover have z = r + (z-r)/h * h using h-neg by simp
  ultimately show ?thesis using that mult-minus-left by fastforce
qed
qed

lemma integral-combine2:
fixes f :: real ⇒ 'a::banach
assumes a ≤ c c ≤ b
  and f integrable-on {a..c} f integrable-on {c..b}
shows integral {a..c} f + integral {c..b} f = integral {a..b} f
apply (rule integral-unique[THEN sym])
apply (rule has-integral-combine[of a c b], simp-all add: assms)
using has-integral-integral assms by auto

lemma has-integral-null-interval: fixes a b :: real and f::real ⇒ real assumes a
≥ b
shows (f has-integral 0) {a..b}
using assms content-real-eq-0 by blast

lemma has-integral-interval-reverse: fixes f :: real ⇒ real and a b :: real
assumes a ≤ b
  and continuous-on {a..b} f
shows ((λx. f (a+b-x)) has-integral (integral {a..b} f)) {a..b}
proof -
let ?g = λx. a + b - x
let ?g' = λx. -1
have g-C0: continuous-on {a..b} ?g using continuous-on-op-minus by simp
have Dg-g': ∀x. x ∈ {a..b} ⇒ (?g has-field-derivative ?g' x) (at x within {a..b})
  by (auto intro!: derivative-eq-intros)
show ?thesis
using has-integral-substitution-general
  [of {} a b ?g a b f, simplified, OF assms g-C0 Dg-g', simplified]
apply (simp add: has-integral-null-interval[OF assms(1), THEN integral-unique])
  by (simp add: has-integral-neg-iff)
qed

lemma FTC-real-deriv-has-integral:
fixes F :: real ⇒ real
assumes a ≤ b
  and F piecewise-differentiable-on {a<..<b}
  and continuous-on {a..b} F
shows (deriv F has-integral F b - F a) {a..b}
proof -
obtain S where fin: finite S and

```

```

diff:  $\bigwedge x. x \in \{a < .. < b\} - S \implies F$  differentiable at  $x$  within  $\{a < .. < b\} - S$ 
using assms unfolding piecewise-differentiable-on-def
by (meson Diff-subset differentiable-within-subset)
hence  $\bigwedge x. x \in \{a < .. < b\} - S \implies (F \text{ has-real-derivative deriv } F x) \text{ (at } x)$ 
proof -
fix x assume x-in:  $x \in \{a < .. < b\} - S$ 
have open ( $\{a < .. < b\} - S$ )
using fin finite-imp-closed by (metis open-Diff open-greaterThanLessThan)
hence at x within  $\{a < .. < b\} - S = \text{at } x$  by (meson x-in at-within-open)
hence  $F$  differentiable at  $x$  using diff x-in by (smt (verit))
thus ( $F$  has-real-derivative deriv  $F x$ ) (at  $x$ )
using DERIV-deriv-iff-real-differentiable by simp
qed
thus ?thesis
by (intro fundamental-theorem-of-calculus-interior-strong[where S=S];
simp add: assms fin has-real-derivative-iff-has-vector-derivative)
qed

lemma integrable-spike-cong:
assumes negligible S  $\bigwedge x. x \in T - S \implies g x = f x$ 
shows  $f$  integrable-on  $T \longleftrightarrow g$  integrable-on  $T$ 
using integrable-spike assms by force

lemma has-integral-powr2-from-0:
fixes a c :: real
assumes a-pos:  $a > 0$  and a-neq-1:  $a \neq 1$  and c-nneg:  $c \geq 0$ 
shows  $((\lambda x. a. \hat{x}) \text{ has-integral } ((a. \hat{c} - 1) / (\ln a))) \{0..c\}$ 
proof -
have  $((\lambda x. a. \hat{x}) \text{ has-integral } ((a. \hat{c}) / (\ln a) - (a. \hat{0}) / (\ln a))) \{0..c\}$ 
proof (rule fundamental-theorem-of-calculus[OF c-nneg])
fix x::real
assume x ∈ {0..c}
show  $((\lambda y. a. \hat{y} / \ln a) \text{ has-vector-derivative } a. \hat{x}) \text{ (at } x \text{ within } \{0..c\})$ 
apply (insert has-real-derivative-powr2[OF a-pos, of x])
apply (drule DERIV-cdivide[where c = ln a], simp add: assms)
apply (rule has-vector-derivative-within-subset[where S=UNIV and T={0..c}], auto)
by (rule iffD1[OF has-real-derivative-iff-has-vector-derivative])
qed
thus ?thesis
using assms powr-zero-eq-one by (simp add: field-simps)
qed

lemma integrable-on-powr2-from-0:
fixes a c :: real
assumes a-pos:  $a > 0$  and a-neq-1:  $a \neq 1$  and c-nneg:  $c \geq 0$ 
shows  $(\lambda x. a. \hat{x}) \text{ integrable-on } \{0..c\}$ 
using has-integral-powr2-from-0[OF assms] unfolding integrable-on-def by blast

```

```

lemma integrable-on-powr2-from-0-general:
  fixes a c :: real
  assumes a-pos: a > 0 and c-nneg: c ≥ 0
  shows (λx. a. ^x) integrable-on {0..c}
  proof (cases a = 1)
    case True
    thus ?thesis
      using has-integral-const-real by auto
  next
    case False
    thus ?thesis
      using has-integral-powr2-from-0 False assms by auto
  qed

lemma has-bochner-integral-power:
  fixes a b :: real and k :: nat
  assumes a ≤ b
  shows has-bochner-integral lborel (λx. x^k * indicator {a..b} x) ((b^(k+1) − a^(k+1)) / (k+1))
  proof −
    have ∀x. ((λx. x^(k+1) / (k+1)) has-real-derivative x^k) (at x)
    using DERIV-pow by (intro derivative-eq-intros) auto
    hence has-bochner-integral lborel (λx. x^k * indicator {a..b} x) (b^(k+1)/(k+1) − a^(k+1)/(k+1))
    by (intro has-bochner-integral-FTC-Icc-real; simp add: assms)
    thus ?thesis by (simp add: diff-divide-distrib)
  qed

corollary integrable-power: (a::real) ≤ b ⇒ integrable lborel (λx. x^k * indicator {a..b} x)
  using has-bochner-integral-power integrable.intros by blast

lemma has-integral-set-integral-real:
  fixes f::'a::euclidean-space ⇒ real and A :: 'a set
  assumes f: set-integrable lborel A f
  shows (f has-integral (set-lebesgue-integral lborel A f)) A
  using assms has-integral-integral-real[where f=λx. indicat-real A x * f x]
  unfolding set-integrable-def set-lebesgue-integral-def
  by simp (smt (verit, ccfv-SIG) has-integral-cong has-integral-restrict-UNIV indicator-times-eq-if)

lemma set-borel-measurable-lborel:
  set-borel-measurable lborel A f ←→ set-borel-measurable borel A f
  unfolding set-borel-measurable-def by simp

lemma restrict-space-whole[simp]: restrict-space M (space M) = M
  unfolding restrict-space-def by (simp add: measure-of-of-measure)

```

```

lemma deriv-measurable-real:
  fixes f :: real  $\Rightarrow$  real
  assumes f differentiable-on S open S f  $\in$  borel-measurable borel
  shows set-borel-measurable borel S (deriv f)
proof -
  have  $\bigwedge x. x \in S \implies \text{deriv } f x = \lim (\lambda i. (f(x + 1 / \text{Suc } i) - f x) / (1 / \text{Suc } i))$ 
  proof -
    fix x assume x-S: x  $\in$  S
    hence f field-differentiable (at x within S)
    using differentiable-on-eq-field-differentiable-real assms by simp
    hence (f has-field-derivative deriv f x) (at x)
    using assms DERIV-deriv-iff-field-differentiable x-S at-within-open by force
    hence ( $\lambda h. (f(x+h) - f x) / h$ )  $\rightarrow$  deriv f x using DERIV-def by auto
    hence  $\forall d. (\forall i. d i \in \text{UNIV} - \{0::\text{real}\}) \rightarrow d \rightarrow 0 \rightarrow$ 
       $((\lambda h. (f(x+h) - f x) / h) \circ d) \rightarrow \text{deriv } f x$ 
    using tendsto-at-iff-sequentially by blast
    moreover have  $\forall i. 1 / \text{Suc } i \in \text{UNIV} - \{0::\text{real}\}$  by simp
    moreover have  $(\lambda i. 1 / \text{Suc } i) \rightarrow 0$  using LIMSEQ-Suc lim-const-over-n
    by blast
    ultimately have  $((\lambda h. (f(x+h) - f x) / h) \circ (\lambda i. 1 / \text{Suc } i)) \rightarrow \text{deriv } f x$  by auto
    thus deriv f x =  $\lim (\lambda i. (f(x + 1 / \text{Suc } i) - f x) / (1 / \text{Suc } i))$ 
    unfolding comp-def by (simp add: limI)
  qed
  moreover have  $(\lambda x. \text{indicator } S x * \lim (\lambda i. (f(x + 1 / \text{Suc } i) - f x) / (1 / \text{Suc } i)))$ 
     $\in$  borel-measurable borel
  using assms by (measurable, simp, measurable)
  ultimately show ?thesis
  unfolding set-borel-measurable-def measurable-cong
  by simp (smt (verit) indicator-simps(2) measurable-cong mult-eq-0-iff)
qed

```

```

lemma piecewise-differentiable-on-deriv-measurable-real:
  fixes f :: real  $\Rightarrow$  real
  assumes f piecewise-differentiable-on S open S f  $\in$  borel-measurable borel
  shows set-borel-measurable borel S (deriv f)
proof -
  from assms obtain T where fin: finite T and
    diff:  $\bigwedge x. x \in S - T \implies f$  differentiable at x within S
  unfolding piecewise-differentiable-on-def by blast
  with assms have open (S - T) using finite-imp-closed by blast
  moreover hence f differentiable-on (S - T)
  unfolding differentiable-on-def using assms by (metis Diff-iff at-within-open
  diff)
  ultimately have set-borel-measurable borel (S - T) (deriv f)
  by (intro deriv-measurable-real; simp add: assms)

```

```

thus ?thesis
  unfolding set-borel-measurable-def apply simp
  apply (rule measurable-discrete-difference
    [where X=T and f=λx. indicat-real (S - T) x * deriv f x], simp-all)
  using fin uncountable-infinite apply blast
  by (simp add: indicator-diff)
qed

lemma borel-measurable-antimono:
  fixes f :: real ⇒ real
  shows antimono f ⟹ f ∈ borel-measurable borel
  using borel-measurable-mono by (smt (verit, del-insts) borel-measurable-uminus-eq
monotone-on-def)

lemma set-borel-measurable-restrict-space-iff:
  fixes f :: 'a ⇒ 'b::real-normed-vector
  assumes Ω[measurable, simp]: Ω ∩ space M ∈ sets M
  shows f ∈ borel-measurable (restrict-space M Ω) ⟷ set-borel-measurable M Ω
f
  using assms borel-measurable-restrict-space-iff set-borel-measurable-def by blast

lemma set-integrable-restrict-space-iff:
  fixes f :: 'a ⇒ 'b:{banach, second-countable-topology}
  assumes A ∈ sets M
  shows set-integrable M A f ⟷ integrable (restrict-space M A) f
  unfolding set-integrable-def using assms
  by (rewrite integrable-restrict-space; simp)

lemma set-lebesgue-integral-restrict-space:
  fixes f :: 'a ⇒ 'b:{banach, second-countable-topology}
  assumes A ∈ sets M
  shows set-lebesgue-integral M A f = integralL (restrict-space M A) f
  unfolding set-lebesgue-integral-def using assms integral-restrict-space
  by (metis (mono-tags) sets.Int-space-eq2)

lemma distr-borel-lborel: distr M borel f = distr M lborel f
  by (metis distr-cong sets-lborel)

lemma AE-translation:
  assumes AE x in lborel. P x shows AE x in lborel. P (a+x)
proof -
  from assms obtain N where P: ∀x. x ∈ space lborel - N ⟹ P x and null:
N ∈ null-sets lborel
  using AE-E3 by blast
  hence {y. a+y ∈ N} ∈ null-sets lborel
  using null-sets-translation[of N - a, simplified] by (simp add: add.commute)
  moreover have ∀y. y ∈ space lborel - {y. a+y ∈ N} ⟹ P (a+y) using P
by force
  ultimately show AE y in lborel. P (a+y)

```

by (smt (verit, del-insts) Diff-iff eventually-ae-filter mem-Collect-eq subsetI)
qed

lemma set-AE-translation:

assumes AE $x \in S$ in lborel. $P x$ **shows** AE $x \in plus(-a) \cup S$ in lborel. $P(a+x)$
proof –
have AE x in lborel. $a+x \in S \rightarrow P(a+x)$ **using assms by** (rule AE-translation)
moreover have $\bigwedge x. a+x \in S \leftrightarrow x \in plus(-a) \cup S$ **by force**
ultimately show ?thesis **by simp**
qed

lemma AE-scale-measure-iff:

assumes $r > 0$
shows (AE x in (scale-measure $r M$). $P x$) \leftrightarrow (AE x in M . $P x$)
unfolding ae-filter-def null-sets-def
apply (rewrite space-scale-measure, simp)
using assms by (smt (verit) Collect-cong not-gr-zero)

lemma nn-set-integral-cong2:

assumes AE $x \in A$ in M . $f x = g x$
shows $(\int^+ x \in A. f x \partial M) = (\int^+ x \in A. g x \partial M)$
proof –
{ fix x
assume $x \in space M$
have $(x \in A \rightarrow f x = g x) \rightarrow f x * indicator A x = g x * indicator A x$ **by force**
hence AE x in M . $(x \in A \rightarrow f x = g x) \rightarrow f x * indicator A x = g x * indicator A x$
by (rule AE-I2)
hence AE x in M . $f x * indicator A x = g x * indicator A x$ **using assms AE-mp**
by auto
thus ?thesis **by** (rule nn-integral-cong-AE)
qed

lemma set-lebesgue-integral-cong-AE2:

assumes [measurable]: $A \in sets M$ set-borel-measurable M A set-borel-measurable M A g
assumes AE $x \in A$ in M . $f x = g x$
shows $(LINT x:A|M. f x) = (LINT x:A|M. g x)$
proof –
let $?fA = \lambda x. indicator A x *_R f x$ **and** $?gA = \lambda x. indicator A x *_R g x$
have $?fA \in borel-measurable M$ $?gA \in borel-measurable M$
using assms unfolding set-borel-measurable-def **by simp-all**
moreover have AE $x \in A$ in M . $?fA x = ?gA x$ **using assms by simp**
ultimately have $(LINT x:A|M. ?fA x) = (LINT x:A|M. ?gA x)$
by (intro set-lebesgue-integral-cong-AE; simp)
moreover have $(LINT x:A|M. f x) = (LINT x:A|M. ?fA x)$ $(LINT x:A|M. g x) = (LINT x:A|M. ?gA x)$
unfolding set-lebesgue-integral-def

```

by (metis indicator-scaleR-eq-if) +
ultimately show ?thesis by simp
qed

proposition set-nn-integral-eq-set-integral:
assumes AE x:A in M. 0 ≤ f x set-integrable M A f
shows (ʃ+x∈A. f x ∂M) = (ʃ x∈A. f x ∂M)
proof –
have (ʃ+x∈A. f x ∂M) = ʃ+x. ennreal (f x * indicator A x) ∂M
  using nn-integral-set-ennreal by blast
also have ... = ʃ x. f x * indicator A x ∂M
  using assms unfolding set-integrable-def
  by (rewrite nn-integral-eq-integral; force simp add: mult.commute)
also have ... = (ʃ x∈A. f x ∂M) unfolding set-lebesgue-integral-def by (simp
add: mult.commute)
  finally show ?thesis .
qed

proposition nn-integral-disjoint-family-on-finite:
assumes [measurable]: f ∈ borel-measurable M ∧(n::nat). n ∈ S ⇒ B n ∈ sets
M
  and disjoint-family-on B S finite S
shows (ʃ+x ∈ (ʃ n∈S. B n). f x ∂M) = (ʃ n∈S. (ʃ+x ∈ B n. f x ∂M))
proof –
let ?A = λn::nat. if n ∈ S then B n else {}
have ∧n::nat. ?A n ∈ sets M by simp
moreover have disjoint-family ?A
  unfolding disjoint-family-on-def
proof –
{ fix m n :: nat
  assume m ≠ n
  hence (if m ∈ S then B m else {}) ∩ (if n ∈ S then B n else {}) = {}
    apply simp
    using assms unfolding disjoint-family-on-def by blast }
thus ∀ m::nat∈UNIV. ∀ n::nat∈UNIV. m ≠ n →
  (if m ∈ S then B m else {}) ∩ (if n ∈ S then B n else {}) = {}
  by blast
qed
ultimately have set-nn-integral M (ʃ (range ?A)) f = (ʃ n. set-nn-integral M
(?A n) f)
  by (rewrite nn-integral-disjoint-family; simp)
moreover have set-nn-integral M (ʃ (range ?A)) f = (ʃ+x ∈ (ʃ n∈S. B n).
f x ∂M)
proof –
have ∪ (range ?A) = (ʃ n∈S. B n) by simp
thus ?thesis by simp
qed
moreover have (ʃ n. set-nn-integral M (?A n) f) = (ʃ n∈S. set-nn-integral
M (B n) f)

```

```

    by (rewrite suminf-finite[of S]; simp add: assms)
ultimately show ?thesis by simp
qed

lemma nn-integral-distr-set:
assumes T ∈ measurable M M' and f ∈ borel-measurable (distr M M' T)
and A ∈ sets M' and ⋀x. x ∈ space M ⟹ T x ∈ A
shows integralN (distr M M' T) f = set-nn-integral (distr M M' T) A f
proof -
have integralN (distr M M' T) f = (ʃ+x ∈ (space M'). f x ∂(distr M M' T))
by (rewrite nn-set-integral-space[THEN sym], simp)
also have ... = (ʃ+x ∈ A. f x ∂(distr M M' T))
proof -
have [simp]: sym-diff (space M') A = space M' - A
using assms by (metis Diff-mono sets.sets-into-space sup.orderE)
show ?thesis
apply (rule nn-integral-null-delta; simp add: assms)
unfolding null-sets-def using assms
apply (simp, rewrite emeasure-distr; simp)
unfolding vimage-def using emeasure-empty
by (metis (no-types, lifting) Diff-disjoint disjoint-iff-not-equal mem-Collect-eq)
qed
finally show ?thesis .
qed

```

```

lemma measure-eqI-Ioc:
fixes M N :: real measure
assumes sets: sets M = sets borel sets N = borel
assumes fin: ⋀a b. a ≤ b ⟹ emeasure M {a <.. b} < ∞
assumes eq: ⋀a b. a ≤ b ⟹ emeasure M {a <.. b} = emeasure N {a <.. b}
shows M = N
proof (rule measure-eqI-generator-eq-countable)
let ?Ioc = λ(a::real,b::real). {a <.. b} let ?E = range ?Ioc
show Int-stable ?E using Int-stable-def Int-greaterThanAtMost by force
show ?E ⊆ Pow UNIV sets M = sigma-sets UNIV ?E sets N = sigma-sets UNIV
?E
unfolding sets by (auto simp add: borel-sigma-sets-Ioc)
show ⋀I. I ∈ ?E ⟹ emeasure M I = emeasure N I
proof -
fix I assume I ∈ ?E
then obtain a b where I = {a <.. b} by auto
thus emeasure M I = emeasure N I by (smt (verit, best) eq greaterThanAt-
Most-empty)
qed
show ?Ioc ` (Rats × Rats) ⊆ ?E (∪ i ∈ (Rats × Rats). ?Ioc i) = UNIV
using Rats-no-bot-less Rats-no-top-le by auto
show countable (?Ioc ` (Rats × Rats)) using countable-rat by blast
show ⋀I. I ∈ ?Ioc ` (Rats × Rats) ⟹ emeasure M I ≠ ∞

```

proof –
fix I **assume** $I \in ?Ioc`(\text{Rats} \times \text{Rats})$
then obtain $a b$ **where** $(a,b) \in (\text{Rats} \times \text{Rats})$ $I = \{a <.. b\}$ **by** blast
thus $\text{emeasure } M I \neq \infty$ **by** (smt (verit, best) $Ioc\text{-inj fin order.strict-implies-not-eq}$)
qed
qed

lemma (in finite-measure) distributed-measure:
assumes distributed $M N X f$
and $\bigwedge x. x \in \text{space } N \implies f x \geq 0$
and $A \in \text{sets } N$
shows measure $M (X - ` A \cap \text{space } M) = (\int x. \text{indicator } A x * f x \partial N)$

proof –
have [simp]: $(\lambda x. \text{indicat-real } A x * f x) \in \text{borel-measurable } N$
using assms **apply** (measurable; simp?)
using distributed-real-measurable assms **by** force
have $\text{emeasure } M (X - ` A \cap \text{space } M) = (\int^+ x \in A. \text{ennreal } (f x) \partial N)$
by (rule distributed-emeasure; simp add: assms)
moreover have enn2real $(\int^+ x \in A. \text{ennreal } (f x) \partial N) = \int x. \text{indicator } A x * f x \partial N$
apply (rewrite enn2real-nn-integral-eq-integral
[where $f = \lambda x. \text{ennreal } (\text{indicator } A x * f x)$, THEN sym]; (simp add: assms)?)
using distributed-emeasure assms
by (smt (verit) emeasure-finite indicator-mult-ennreal mult.commute
nn-integral-cong top.not-eq-extremum)
ultimately show ?thesis **using** measure-def **by** metis
qed

lemma set-integrable-const[simp]:
 $A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies \text{set-integrable } M A (\lambda-. c)$
using has-bochner-integral-indicator **unfolding** set-integrable-def **by** simp

lemma set-integral-const[simp]:
 $A \in \text{sets } M \implies \text{emeasure } M A < \infty \implies \text{set-lebesgue-integral } M A (\lambda-. c) = \text{measure } M A *_R c$
unfolding set-lebesgue-integral-def **using** has-bochner-integral-indicator **by** force

lemma set-integral-empty-0[simp]: $\text{set-lebesgue-integral } M \{\} f = 0$
unfolding set-lebesgue-integral-def **by** simp

lemma set-integral-nonneg[simp]:
fixes $f :: 'a \Rightarrow \text{real}$ **and** $A :: 'a \text{ set}$
shows $(\bigwedge x. x \in A \implies 0 \leq f x) \implies 0 \leq \text{set-lebesgue-integral } M A f$
unfolding set-lebesgue-integral-def **by** (simp add: indicator-times-eq-if(1))

lemma
fixes $f :: 'a \Rightarrow 'b :: \{\text{banach}, \text{second-countable-topology}\}$ **and** $w :: 'a \Rightarrow \text{real}$
assumes $A \in \text{sets } M$ set-borel-measurable $M A f$

```

 $\bigwedge i. \text{set-borel-measurable } M A (s i) \text{ set-integrable } M A w$ 
assumes lim:  $\text{AE } x \in A \text{ in } M. (\lambda i. s i x) \longrightarrow f x$ 
assumes bound:  $\bigwedge i::\text{nat}. \text{AE } x \in A \text{ in } M. \text{norm } (s i x) \leq w x$ 
shows set-integrable-dominated-convergence: set-integrable  $M A f$ 
and set-integrable-dominated-convergence2:  $\bigwedge i. \text{set-integrable } M A (s i)$ 
and set-integral-dominated-convergence:
 $(\lambda i. \text{set-lebesgue-integral } M A (s i)) \longrightarrow \text{set-lebesgue-integral } M A f$ 
proof -
  have  $(\lambda x. \text{indicator } A x *_R f x) \in \text{borel-measurable } M$  and
     $\bigwedge i. (\lambda x. \text{indicator } A x *_R s i x) \in \text{borel-measurable } M$  and
      integrable  $M (\lambda x. \text{indicator } A x *_R w x)$ 
    using assms unfolding set-borel-measurable-def set-integrable-def by simp-all
    moreover have  $\text{AE } x \text{ in } M. (\lambda i. \text{indicator } A x *_R s i x) \longrightarrow \text{indicator } A x$ 
     $*_R f x$ 
  proof -
    obtain N where N-null:  $N \in \text{null-sets } M$  and
      si-f:  $\bigwedge x. x \in \text{space } M - N \implies x \in A \longrightarrow (\lambda i. s i x) \longrightarrow f x$ 
      using lim AE-E3 by (smt (verit))
    hence  $\bigwedge x. x \in \text{space } M - N \implies (\lambda i. \text{indicator } A x *_R s i x) \longrightarrow \text{indicator}$ 
     $A x *_R f x$ 
    proof -
      fix x assume asm:  $x \in \text{space } M - N$ 
      thus  $(\lambda i. \text{indicator } A x *_R s i x) \longrightarrow \text{indicator } A x *_R f x$ 
      proof (cases  $\langle x \in A \rangle$ )
        case True
          with si-f asm show ?thesis by simp
        next
          case False
          thus ?thesis by simp
        qed
      qed
      thus ?thesis by (smt (verit) AE-I' DiffI N-null mem-Collect-eq subsetI)
    qed
    moreover have  $\bigwedge i. \text{AE } x \text{ in } M. \text{norm } (\text{indicator } A x *_R s i x) \leq \text{indicator } A x$ 
     $*_R w x$ 
  proof -
    fix i
    from bound obtain N where N-null:  $N \in \text{null-sets } M$  and
       $\bigwedge x. x \in \text{space } M - N \implies x \in A \longrightarrow \text{norm } (s i x) \leq w x$ 
      using AE-E3 by (smt (verit))
    hence  $\bigwedge x. x \in \text{space } M - N \implies \text{norm } (\text{indicator } A x *_R s i x) \leq \text{indicator}$ 
     $A x *_R w x$ 
    by (simp add: indicator-scaleR-eq-if)
    with N-null show  $\text{AE } x \text{ in } M. \text{norm } (\text{indicator } A x *_R s i x) \leq \text{indicator } A x$ 
     $*_R w x$ 
    by (smt (verit) DiffI eventually-ae-filter mem-Collect-eq subsetI)
  qed
  ultimately show set-integrable  $M A f \bigwedge i. \text{set-integrable } M A (s i)$ 
   $(\lambda i. \text{set-lebesgue-integral } M A (s i)) \longrightarrow \text{set-lebesgue-integral } M A f$ 

```

```

unfolding set-integrable-def set-lebesgue-integral-def
by (rule integrable-dominated-convergence, rule integrable-dominated-convergence2,
      rule integral-dominated-convergence)
qed

lemma absolutely-integrable-on-iff-set-integrable:
  fixes f :: 'a::euclidean-space  $\Rightarrow$  real
  assumes f  $\in$  borel-measurable lborel
  and S  $\in$  sets lborel
  shows set-integrable lborel S f  $\longleftrightarrow$  f absolutely-integrable-on S
  unfolding set-integrable-def apply (simp, rewrite integrable-completion[THEN
  sym])
  apply measurable using assms by simp-all

corollary integrable-on-iff-set-integrable-nonneg:
  fixes f :: 'a::euclidean-space  $\Rightarrow$  real
  assumes  $\bigwedge x. x \in S \implies f x \geq 0$  f  $\in$  borel-measurable lborel
  and S  $\in$  sets lborel
  shows set-integrable lborel S f  $\longleftrightarrow$  f integrable-on S
  using absolutely-integrable-on-iff-set-integrable assms
  by (metis absolutely-integrable-on-iff-nonneg)

lemma integrable-on-iff-set-integrable-nonneg-AE:
  fixes f :: 'a::euclidean-space  $\Rightarrow$  real
  assumes AE x $\in$ S in lborel. f x  $\geq 0$  f  $\in$  borel-measurable lborel
  and S  $\in$  sets lborel
  shows set-integrable lborel S f  $\longleftrightarrow$  f integrable-on S
proof –
  from assms obtain N where nonneg:  $\bigwedge x. x \in S - N \implies f x \geq 0$  and null:
  N  $\in$  null-sets lborel
  by (smt (verit, ccfv-threshold) AE-E3 Diff-iff UNIV-I space-borel space-lborel)
  let ?g =  $\lambda x. \text{if } x \in N \text{ then } 0 \text{ else } f x$ 
  have [simp]: negligible N using null negligible-iff-null-sets null-sets-completionI
  by blast
  have N  $\in$  sets lborel using null by auto
  hence [simp]: ?g  $\in$  borel-measurable borel using assms by force
  have set-integrable lborel S f  $\longleftrightarrow$  set-integrable lborel S ?g
proof –
  have AE x $\in$ S in lborel. f x = ?g x by (rule AE-I'[of N], simp-all add: null,
  blast)
  thus ?thesis using assms by (intro set-integrable-cong-AE[of f - ?g S]; simp)
  qed
  also have ...  $\longleftrightarrow$  ?g integrable-on S
  using assms by (intro integrable-on-iff-set-integrable-nonneg; simp add: nonneg)
  also have ...  $\longleftrightarrow$  f integrable-on S by (rule integrable-spike-cong[of N]; simp)
  finally show ?thesis .
qed

lemma set-borel-integral-eq-integral-nonneg:

```

```

fixes f :: 'a::euclidean-space  $\Rightarrow$  real
assumes  $\bigwedge x. x \in S \implies f x \geq 0$  f  $\in$  borel-measurable borel S  $\in$  sets borel
shows ( $\text{LINT } x : S \mid \text{lborel. } f x$ ) = integral S f
    — Note that  $0 = 0$  holds when the integral diverges.
proof (cases ⟨set-integrable lborel S f⟩)
  case True
    thus ?thesis using set-borel-integral-eq-integral by force
  next
    case False
    hence ( $\text{LINT } x : S \mid \text{lborel. } f x$ ) = 0
      unfolding set-lebesgue-integral-def set-integrable-def
      by (rewrite not-integrable-integral-eq; simp)
    moreover have integral S f = 0
      apply (rule not-integrable-integral)
      using False assms by (rewrite in asm integrable-on-iff-set-integrable-nonneg;
      simp)
      ultimately show ?thesis ..
  qed

lemma set-borel-integral-eq-integral-nonneg-AE:
fixes f :: 'a::euclidean-space  $\Rightarrow$  real
assumes AE  $x \in S$  in lborel.  $f x \geq 0$  f  $\in$  borel-measurable borel S  $\in$  sets borel
shows ( $\text{LINT } x : S \mid \text{lborel. } f x$ ) = integral S f
    — Note that  $0 = 0$  holds when the integral diverges.
proof (cases ⟨set-integrable lborel S f⟩)
  case True
    thus ?thesis using set-borel-integral-eq-integral by force
  next
    case False
    hence ( $\text{LINT } x : S \mid \text{lborel. } f x$ ) = 0
      unfolding set-lebesgue-integral-def set-integrable-def
      by (rewrite not-integrable-integral-eq; simp)
    moreover have integral S f = 0
      apply (rule not-integrable-integral)
      using False assms by (rewrite in asm integrable-on-iff-set-integrable-nonneg-AE;
      simp)
      ultimately show ?thesis ..
  qed

```

2.1 Set Lebesgue Integrability on Affine Transformation

```

lemma set-integrable-Icc-affine-pos-iff:
fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a b c t :: real
assumes c > 0
shows set-integrable lborel  $\{(a-t)/c..(b-t)/c\}$  ( $\lambda x. f(t + c*x)$ )
     $\longleftrightarrow$  set-integrable lborel {a..b} f
unfolding set-integrable-def using assms
apply (rewrite indicator-Icc-affine-pos-inverse, simp)
by (rule lborel-integrable-real-affine-iff) simp

```

```

corollary set-integrable-Icc-shift:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a b t :: real
  shows set-integrable lborel {a-t..b-t} ( $\lambda x. f(t+x)$ )  $\longleftrightarrow$  set-integrable lborel
  {a..b} f
  using set-integrable-Icc-affine-pos-iff[where c=1] by simp

lemma set-integrable-Ici-affine-pos-iff:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a c t :: real
  assumes c > 0
  shows set-integrable lborel {(a-t)/c..} ( $\lambda x. f(t + c*x)$ )
   $\longleftrightarrow$  set-integrable lborel {a..} f
  unfolding set-integrable-def using assms
  apply (rewrite indicator-Ici-affine-pos-inverse, simp)
  by (rule lborel-integrable-real-affine-iff) simp

corollary set-integrable-Ici-shift:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a t :: real
  shows set-integrable lborel {a-t..} ( $\lambda x. f(t+x)$ )  $\longleftrightarrow$  set-integrable lborel {a..} f
  using set-integrable-Ici-affine-pos-iff[where c=1] by simp

lemma set-integrable-Iic-affine-pos-iff:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and b c t :: real
  assumes c > 0
  shows set-integrable lborel {..(b-t)/c} ( $\lambda x. f(t + c*x)$ )
   $\longleftrightarrow$  set-integrable lborel {..b} f
  unfolding set-integrable-def using assms
  apply (rewrite indicator-Iic-affine-pos-inverse, simp)
  by (rule lborel-integrable-real-affine-iff) simp

corollary set-integrable-Iic-shift:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and b t :: real
  shows set-integrable lborel {..b-t} ( $\lambda x. f(t+x)$ )  $\longleftrightarrow$  set-integrable lborel {..b} f
  using set-integrable-Iic-affine-pos-iff[where c=1] by simp

lemma set-integrable-Icc-affine-neg-iff:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a b c t :: real
  assumes c < 0
  shows set-integrable lborel {(b-t)/c..(a-t)/c} ( $\lambda x. f(t + c*x)$ )
   $\longleftrightarrow$  set-integrable lborel {a..b} f
  unfolding set-integrable-def using assms
  apply (rewrite indicator-Icc-affine-neg-inverse, simp)
  by (rule lborel-integrable-real-affine-iff) simp

corollary set-integrable-Icc-reverse:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a b t :: real
  shows set-integrable lborel {t-b..t-a} ( $\lambda x. f(t-x)$ )  $\longleftrightarrow$  set-integrable lborel
  {a..b} f
  using set-integrable-Icc-affine-neg-iff[where c=-1] by simp

```

```

lemma set-integrable-Ici-affine-neg-iff:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and b c t :: real
  assumes c < 0
  shows set-integrable lborel  $\{(b-t)/c..\}$  ( $\lambda x. f(t + c*x)$ )
     $\longleftrightarrow$  set-integrable lborel  $\{..b\}$  f
  unfolding set-integrable-def using assms
  apply (rewrite indicator-Ici-affine-neg-inverse, simp)
  by (rule lborel-integrable-real-affine-iff) simp

corollary set-integrable-Ici-reverse:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and b t :: real
  shows set-integrable lborel  $\{t-b..\}$  ( $\lambda x. f(t-x)$ )  $\longleftrightarrow$  set-integrable lborel  $\{..b\}$  f
  using set-integrable-Ici-affine-neg-iff[where c=-1] by simp

lemma set-integrable-Iic-affine-neg-iff:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a c t :: real
  assumes c < 0
  shows set-integrable lborel  $\{..(a-t)/c\}$  ( $\lambda x. f(t + c*x)$ )
     $\longleftrightarrow$  set-integrable lborel  $\{a..\}$  f
  unfolding set-integrable-def using assms
  apply (rewrite indicator-Iic-affine-neg-inverse, simp)
  by (rule lborel-integrable-real-affine-iff) simp

corollary set-integrable-Iic-reverse:
  fixes f :: real  $\Rightarrow$  'a:{banach, second-countable-topology} and a t :: real
  shows set-integrable lborel  $\{..t-a\}$  ( $\lambda x. f(t-x)$ )  $\longleftrightarrow$  set-integrable lborel  $\{a..\}$  f
  using set-integrable-Iic-affine-neg-iff[where c=-1] by simp

```

2.2 Set Lebesgue Integral on Affine Transformation

```

lemma lborel-set-integral-Icc-affine-pos:
  fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a b c :: real
  assumes c > 0
  shows  $(\int x \in \{a..b\}. f x \partial\text{lborel}) = c *_R (\int x \in \{(a-t)/c..(b-t)/c\}. f(t + c*x) \partial\text{lborel})$ 
  unfolding set-lebesgue-integral-def using assms
  apply (rewrite indicator-Icc-affine-pos-inverse, simp)
  using lborel-integral-real-affine[where c=c] by force

corollary lborel-set-integral-Icc-shift:
  fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a b :: real
  shows  $(\int x \in \{a..b\}. f x \partial\text{lborel}) = (\int x \in \{a-t..b-t\}. f(t+x) \partial\text{lborel})$ 
  using lborel-set-integral-Icc-affine-pos[where c=1] by simp

lemma lborel-set-integral-Ici-affine-pos:
  fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a c :: real
  assumes c > 0
  shows  $(\int x \in \{a..\}. f x \partial\text{lborel}) = c *_R (\int x \in \{(a-t)/c..\}. f(t + c*x) \partial\text{lborel})$ 

```

```

unfolding set-lebesgue-integral-def using assms
apply (rewrite indicator-Ici-affine-pos-inverse, simp)
using lborel-integral-real-affine[where c=c] by force

corollary lborel-set-integral-Ici-shift:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a::real
shows  $(\int x \in \{a..\}. f x) \partial borel = (\int x \in \{a-t..\}. f (t+x)) \partial borel$ 
using lborel-set-integral-Ici-affine-pos[where c=1] by simp

lemma lborel-set-integral-Iic-affine-pos:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and b c :: real
assumes c > 0
shows  $(\int x \in \{..b\}. f x) \partial borel = c * R (\int x \in \{(b-t)/c..\}. f (t + c*x)) \partial borel$ 
unfolding set-lebesgue-integral-def using assms
apply (rewrite indicator-Iic-affine-pos-inverse, simp)
using lborel-integral-real-affine[where c=c] by force

corollary lborel-set-integral-Iic-shift:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and b::real
shows  $(\int x \in \{..b\}. f x) \partial borel = (\int x \in \{..b-t\}. f (t+x)) \partial borel$ 
using lborel-set-integral-Iic-affine-pos[where c=1] by simp

lemma lborel-set-integral-Icc-affine-neg:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a b c :: real
assumes c < 0
shows  $(\int x \in \{a..b\}. f x) \partial borel = -c * R (\int x \in \{(b-t)/c..(a-t)/c\}. f (t + c*x)) \partial borel$ 
unfolding set-lebesgue-integral-def using assms
apply (rewrite indicator-Icc-affine-neg-inverse, simp)
using lborel-integral-real-affine[where c=c] by force

corollary lborel-set-integral-Icc-reverse:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a b :: real
shows  $(\int x \in \{a..b\}. f x) \partial borel = (\int x \in \{t-b..t-a\}. f (t-x)) \partial borel$ 
using lborel-set-integral-Icc-affine-neg[where c=-1] by simp

lemma lborel-set-integral-Ici-affine-neg:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and b c :: real
assumes c < 0
shows  $(\int x \in \{..b\}. f x) \partial borel = -c * R (\int x \in \{(b-t)/c..\}. f (t + c*x)) \partial borel$ 
unfolding set-lebesgue-integral-def using assms
apply (rewrite indicator-Ici-affine-neg-inverse, simp)
using lborel-integral-real-affine[where c=c] by force

corollary lborel-set-integral-Ici-reverse:
fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and b::real
shows  $(\int x \in \{..b\}. f x) \partial borel = (\int x \in \{t-b..\}. f (t-x)) \partial borel$ 
using lborel-set-integral-Ici-affine-neg[where c=-1] by simp

```

```

lemma lborel-set-integral-Iic-affine-neg:
  fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a c :: real
  assumes c < 0
  shows ( $\int_{x \in \{a..\}} f x \, d\text{lborel}$ ) = -c *R ( $\int_{x \in \{(a-t)/c..(a-t)/c\}} f (t + c*x) \, d\text{lborel}$ )
  unfolding set-lebesgue-integral-def using assms
  apply (rewrite indicator-Iic-affine-neg-inverse, simp)
  using lborel-integral-real-affine[where c=c] by force

corollary lborel-set-integral-Iic-reverse:
  fixes f :: real  $\Rightarrow$  'a :: {banach, second-countable-topology} and a::real
  shows ( $\int_{x \in \{a..\}} f x \, d\text{lborel}$ ) = ( $\int_{x \in \{..t-a\}} f (t-x) \, d\text{lborel}$ )
  using lborel-set-integral-Iic-affine-neg[where c=-1] by simp

lemma set-integrable-Ici-equiv-aux:
  fixes f :: real  $\Rightarrow$  'a::banach, second-countable-topology} and a b :: real
  assumes  $\bigwedge c d. \text{set-integrable lborel } \{c..d\} f a \leq b$ 
  shows set-integrable lborel {a..} f  $\longleftrightarrow$  set-integrable lborel {b..} f
proof
  assume set-integrable lborel {a..} f
  thus set-integrable lborel {b..} f by (rule set-integrable-subset; simp add: assms)
next
  assume set-integrable lborel {b..} f
  moreover have set-integrable lborel {a..b} f using assms by blast
  moreover have {a..} = {a..b}  $\cup$  {b..} using assms by auto
  ultimately show set-integrable lborel {a..} f using set-integrable-Un by force
qed

corollary set-integrable-Ici-equiv:
  fixes f :: real  $\Rightarrow$  'a::banach, second-countable-topology} and a b :: real
  assumes  $\bigwedge c d. \text{set-integrable lborel } \{c..d\} f$ 
  shows set-integrable lborel {a..} f  $\longleftrightarrow$  set-integrable lborel {b..} f
  using set-integrable-Ici-equiv-aux assms by (smt (verit))

lemma set-integrable-Iic-equiv:
  fixes f :: real  $\Rightarrow$  real and a b :: real
  assumes  $\bigwedge c d. \text{set-integrable lborel } \{c..d\} f$ 
  shows set-integrable lborel {..a} f  $\longleftrightarrow$  set-integrable lborel {..b} f (is ?LHS  $\longleftrightarrow$  ?RHS)
proof -
  have ?LHS  $\longleftrightarrow$  set-integrable lborel {-a..} ( $\lambda x. f (-x)$ )
  using set-integrable-Ici-reverse[where t=0] by force
  also have ...  $\longleftrightarrow$  set-integrable lborel {-b..} ( $\lambda x. f (-x)$ )
proof -
  have  $\bigwedge c d. \text{set-integrable lborel } \{c..d\} (\lambda x. f (-x))$ 
  apply (rewrite at {..} minus-minus[THEN sym])
  apply (rewrite at {..} minus-minus[THEN sym])
  using assms set-integrable-Icc-reverse[where t=0] by force
  thus ?thesis by (rule set-integrable-Ici-equiv)
qed

```

also have ... $\longleftrightarrow ?RHS$ **using** *set-integrable-Ici-reverse[where t=0]* **by force**
finally show *?thesis* .

qed

2.3 Alternative Integral Test

lemma *nn-integral-suminf-Ico-real-nat*:

fixes $a::real$ **and** $f :: real \Rightarrow ennreal$
assumes $f \in borel-measurable lborel$
shows $(\int^+ x \in \{a..\}. f x \partial borel) = (\sum k. \int^+ x \in \{a+k..<a+k+1\}. f x \partial borel)$
apply (*rewrite Ico-real-nat-union[THEN sym]*)
using *Ico-real-nat-disjoint assms by (intro nn-integral-disjoint-family; simp)*

lemma *set-integrable-iff-bounded*:

fixes $f :: 'a \Rightarrow 'b :: \{banach, second-countable-topology\}$
assumes $A \in sets M$
shows *set-integrable M A f \longleftrightarrow set-borel-measurable M A f \wedge $(\int^+ x \in A. norm(f x) \partial M) < \infty$*
unfolding *set-integrable-def set-borel-measurable-def using integrable-iff-bounded*
by (*smt (verit, ccfv-threshold) indicator-mult-ennreal indicator-pos-le*
mult.commute nn-integral-cong norm-scaleR)

theorem *set-integrable-iff-summable*:

fixes $a::real$ **and** $f :: real \Rightarrow real$
assumes *antimono-on {a..} f \wedge $x. a \leq x \implies f x \geq 0$ $f \in borel-measurable lborel$*
shows *set-integrable lborel {a..} f \longleftrightarrow summable ($\lambda k. f(a+k)$)*
proof
assume *asm: set-integrable lborel {a..} f*
have [*measurable*]: $(\lambda x. ennreal(f x)) \in borel-measurable lborel$ **using** *assms by simp*
have $\forall k \geq 0. norm(f(a+(k+1::nat))) \leq (\int x \in \{a+k..<a+k+1\}. f x \partial borel)$
proof –
{ **fix** $k::nat$
have *norm(f(a+(k+1::nat))) = f(a+k+1)*
using *assms by (smt (verit) of-nat-0-le-iff of-nat-1 of-nat-add real-norm-def)*
also have ... $= (\int x \in \{a+k..<a+k+1\}. f(a+k+1) \partial borel)$
unfolding *set-lebesgue-integral-def by simp*
also have ... $\leq (\int x \in \{a+k..<a+k+1\}. f x \partial borel)$
apply (*rule set-integral-mono, simp*)
apply (*rule set-integrable-restrict-space[of lborel {a..}], simp add: asm*)
apply (*rewrite sets-restrict-space, force*)
using *assms unfolding mono-on-def monotone-on-def by simp*
finally have *norm(f(a+(k+1::nat))) \leq (\int x \in \{a+k..<a+k+1\}. f x \partial borel)*
. }
thus *?thesis by simp*
qed
moreover have *summable ($\lambda k. \int x \in \{a+k..<a+k+1\}. f x \partial borel$)*
proof –

```

have ( $\int^+_{x \in \{a..\}} ennreal (f x) \partial borel \neq \infty$ 
  using asm unfolding set-integrable-def apply simp
  by (smt (verit) indicator-mult-ennreal infinity-ennreal-def mult.commute
       nn-integral-cong real-integrable-def)
thus ?thesis
apply (rewrite in asm nn-integral-suminf-Ico-real-nat, simp)
apply (rule summable-suminf-not-top)
using assms apply (intro set-integral-nonneg, force)
apply (rewrite set-nn-integral-eq-set-integral[THEN sym], simp add: assms)
by (rule set-integrable-subset[of lborel {a..}], simp-all add: asm) force
qed
ultimately have summable ( $\lambda k. f (a+(k+1::nat))$ )
  using summable-comparison-test by (smt (verit, del-insts))
thus summable ( $\lambda k. f (a+k)$ ) using summable-iff-shift by blast
next
assume asm: summable ( $\lambda k. f (a+k)$ )
hence ( $\int^+_{x \in \{a..\}} ennreal |f x| \partial borel < \infty$ )
proof -
  have ( $\int^+_{x \in \{a..\}} ennreal |f x| \partial borel = (\int^+_{x \in \{a..\}} ennreal (f x) \partial borel)$ 
    using assms by (metis abs-of-nonneg atLeast-iff indicator-simps(2) mult-eq-0-iff)
  also have ... = ( $\sum k. \int^+_{x \in \{a+k..<a+k+1\}} ennreal (f x) \partial borel$ )
    using assms by (rewrite nn-integral-suminf-Ico-real-nat; simp)
  also have ...  $\leq (\sum k. \int^+_{x \in \{a+k..<a+k+1\}} ennreal (f (a+k)) \partial borel)$ 
  proof -
    have  $\bigwedge (k::nat) x. x \in \{a+k..<a+k+1\} \implies f x \leq f (a+k)$ 
      using assms unfolding monotone-on-def by auto
    thus ?thesis
      apply (intro suminf-le, simp-all)
      by (rule nn-integral-mono)
        (metis (no-types, opaque-lifting) atLeastLessThan-iff dual-order.refl ennreal-leI
         indicator-simps(2) mult-eq-0-iff mult-mono zero-le)
    qed
    also have ... = ( $\sum k. ennreal (f (a+k))$ )
      apply (rule suminf-cong)
      by (rewrite nn-integral-cmult-indicator; simp)
    also have ...  $< \infty$ 
      unfolding infinity-ennreal-def apply (rewrite less-top[THEN sym])
      using asm assms by (smt (verit) of-nat-0-le-iff suminf-cong suminf-ennreal2
           top-neq-ennreal)
      finally show ?thesis .
    qed
    moreover have set-borel-measurable lborel {a..} f
      using assms unfolding set-borel-measurable-def by simp
    ultimately show set-integrable lborel {a..} f by (rewrite set-integrable-iff-bounded)
  auto
qed

```

2.4 Interchange of Differentiation and Lebesgue Integration

definition measurable-extension :: '*a measure* \Rightarrow '*b measure* \Rightarrow ('*a* \Rightarrow '*b*) \Rightarrow '*a* \Rightarrow '*b* **where**

measurable-extension *M N f* =

(SOME *g*. *g* \in *M* \rightarrow_M *N* \wedge ($\exists S \in$ (null-sets *M*). {*x* \in space *M*. *f x* \neq *g x*} \subseteq *S*))

— The term *measurable-extension* is proposed by Reynald Affeldt.

— This function is used to make an almost-everywhere-defined function measurable.

lemma

fixes *f g*

assumes *g* \in *M* \rightarrow_M *N* *S* \in null-sets *M* {*x* \in space *M*. *f x* \neq *g x*} \subseteq *S*

shows measurable-extensionI: AE *x* in *M*. *f x* = measurable-extension *M N f x* and

measurable-extensionI2: AE *x* in *M*. *g x* = measurable-extension *M N f x* **and**

measurable-extension-measurable: measurable-extension *M N f* \in measurable *M N*

proof —

let ?G = λg . *g* \in *M* \rightarrow_M *N* **and** ?S = λg . $\exists S \in$ null-sets *M*. {*x* \in space *M*. *f x* \neq *g x*} \subseteq *S*

show AE *x* in *M*. *f x* = measurable-extension *M N f x*

unfolding measurable-extension-def

apply (rule someI2[of λg . ?G *g* \wedge ?S *g* *g*])

using assms apply blast

using AE-I' by auto

moreover have AE *x* in *M*. *g x* = *f x*

using assms by (smt (verit, best) AE-I' Collect-cong)

ultimately show AE *x* in *M*. *g x* = measurable-extension *M N f x* **by force**

show measurable-extension *M N f* \in measurable *M N*

unfolding measurable-extension-def

apply (rule conjE[of ?G *g* ?S *g*])

using assms apply auto[1]

using someI-ex[of λg . ?G *g* \wedge ?S *g*] **by auto**

qed

corollary measurable-measurable-extension-AE:

fixes *f*

assumes *f* \in *M* \rightarrow_M *N*

shows AE *x* in *M*. *f x* = measurable-extension *M N f x*

by (rule measurable-extensionI[**where** *g=f* **and** *S={}*]); simp add: assms)

definition borel-measurable-extension ::

'*a measure* \Rightarrow ('*a* \Rightarrow '*b*::topological-space) \Rightarrow '*a* \Rightarrow '*b* **where**

borel-measurable-extension *M f* = measurable-extension *M borel f*

lemma

fixes *f g*

assumes *g* \in borel-measurable *M* *S* \in null-sets *M* {*x* \in space *M*. *f x* \neq *g x*} \subseteq *S*

```

shows borel-measurable-extensionI: AE x in M. f x = borel-measurable-extension
M f x and
  borel-measurable-extensionI2: AE x in M. g x = borel-measurable-extension M
f x and
  borel-measurable-extension-measurable: borel-measurable-extension M f ∈ borel-measurable
M
unfolding borel-measurable-extension-def using assms
apply -
  using measurable-extensionI apply blast
  using measurable-extensionI2 apply blast
  using measurable-extension-measurable by blast

corollary borel-measurable-measurable-extension-AE:
  fixes f
  assumes f ∈ borel-measurable M
  shows AE x in M. f x = borel-measurable-extension M f x
  using assms measurable-measurable-extension-AE unfolding borel-measurable-extension-def
  by auto

definition set-borel-measurable-extension :: 
  'a measure ⇒ 'a set ⇒ ('a ⇒ 'b::topological-space) ⇒ 'a ⇒ 'b
  where set-borel-measurable-extension M A f = borel-measurable-extension (restrict-space
M A) f

lemma
  fixes f g :: 'a ⇒ 'b::real-normed-vector and A
  assumes A ∈ sets M set-borel-measurable M A g S ∈ null-sets M {x ∈ A. f x ≠
g x} ⊆ S
  shows set-borel-measurable-extensionI:
    AE x∈A in M. f x = set-borel-measurable-extension M A f x and
    set-borel-measurable-extensionI2:
      AE x∈A in M. g x = set-borel-measurable-extension M A f x and
      set-borel-measurable-extension-measurable:
        set-borel-measurable M A (set-borel-measurable-extension M A f)

proof -
  have g ∈ borel-measurable (restrict-space M A)
  using assms by (rewrite set-borel-measurable-restrict-space-iff; simp)
  moreover have S ∩ A ∈ null-sets (restrict-space M A)
  using assms null-sets-restrict-space by (metis Int-lower2 null-set-Int2)
  moreover have {x ∈ space (restrict-space M A). f x ≠ g x} ⊆ S ∩ A
  using assms by (rewrite space-restrict-space2; simp)
  ultimately show AE x∈A in M. f x = set-borel-measurable-extension M A f x
  and
    AE x∈A in M. g x = set-borel-measurable-extension M A f x and
    set-borel-measurable M A (set-borel-measurable-extension M A f)
    unfolding set-borel-measurable-extension-def using assms
    apply -
    apply (rewrite AE-restrict-space-iff[THEN sym], simp)
    apply (rule borel-measurable-extensionI[of g - S ∩ A]; simp)

```

```

apply (rewrite AE-restrict-space-iff[THEN sym], simp)
apply (rule borel-measurable-extensionI2[of g - S ∩ A]; simp)
apply (rewrite set-borel-measurable-restrict-space-iff[THEN sym], simp)
  by (rule borel-measurable-extension-measurable[of g - S ∩ A]; simp)
qed

```

corollary *set-borel-measurable-measurable-extension-AE*:

```

fixes f::'a ⇒ 'b::real-normed-vector and A
assumes set-borel-measurable M A f A ∈ sets M
shows AE x∈A in M. f x = set-borel-measurable-extension M A f x
using set-borel-measurable-restrict-space-iff
  borel-measurable-measurable-extension-AE AE-restrict-space-iff
unfolding set-borel-measurable-extension-def
  by (smt (verit) AE-cong sets.Int-space-eq2 assms)

```

proposition *interchange-deriv-LINT-general*:

```

fixes a b :: real and f :: real ⇒ 'a ⇒ real and g :: 'a ⇒ real
assumes f-integ: ∀r. r ∈ {a <.. < b} ⇒ integrable M (f r) and
  f-diff: AE x in M. (λr. f r x) differentiable-on {a <.. < b} and
  Df-bound: AE x in M. ∀r ∈ {a <.. < b}. |deriv (λr. f r x) r| ≤ g x integrable M g
shows ∀r. r ∈ {a <.. < b} ⇒ ((λr. ∫ x. f r x ∂M) has-real-derivative
  ∫ x. borel-measurable-extension M (λx. deriv (λr. f r x) r) x ∂M) (at r)
proof –

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Preparation

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have f-msr: ∀r. r ∈ {a <.. < b} ⇒ f r ∈ borel-measurable M using f-integ by
  auto
from f-diff obtain N1 where N1-null: N1 ∈ null-sets M and
  ∀x. x ∈ space M - N1 ⇒ (λs. f s x) differentiable-on {a <.. < b}
  by (smt (verit) AE-E3)
hence f-diffN1: ∀x. x ∈ space M - N1 ⇒ (λs. f s x) differentiable-on {a <.. < b}
  by (meson Diff-iff sets.sets-into-space subset-eq)
from Df-bound obtain N2 where N2-null: N2 ∈ null-sets M and
  ∀x. x ∈ space M - N2 ⇒ ∀r ∈ {a <.. < b}. |deriv (λs. f s x) r| ≤ g x
  by (smt (verit) AE-E3)
hence Df-boundN2: ∀x. x ∈ space M - N2 ⇒ ∀r ∈ {a <.. < b}. |deriv (λs. f s x) r| ≤ g x
  by (meson Diff-iff sets.sets-into-space subset-eq)
define N where N ≡ N1 ∪ N2
let ?CN = space M - N
have N-null: N ∈ null-sets M and N-msr: N ∈ sets M
  unfolding N-def using N1-null N2-null by auto
have f-diffCN: ∀x. x ∈ ?CN ⇒ (λs. f s x) differentiable-on {a <.. < b}
  unfolding N-def using f-diffN1 by simp
define Df :: real ⇒ 'a ⇒ real where
  Df r x ≡ indicator ({a <.. < b} × ?CN) (r, x) * deriv (λs. f s x) r for r x
have Df-boundCN: ∀x. x ∈ ?CN ⇒ ∀r ∈ {a <.. < b}. |Df r x| ≤ g x
  unfolding Df-def N-def using Df-boundN2 by simp

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Main Part of the Proof

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fix r assume r-ab:  $r \in \{a < .. < b\}$ 
then obtain e where e-pos:  $e > 0$  and ball-ab:  $\text{ball } r e \subseteq \{a < .. < b\}$ 
  by (meson openE open-greaterThanLessThan)
have  $\bigwedge d:\text{nat} \Rightarrow \text{real}. \llbracket \forall i. d \ i \in \text{UNIV} - \{0\}; d \longrightarrow 0 \rrbracket \implies$ 
   $((\lambda h. ((\int x. f(r+h) x \partial M) - \int x. f r x \partial M) / h) \circ d) \longrightarrow$ 
   $\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x \partial M$ 
proof -
  fix d::nat $\Rightarrow$ real assume d-neq0:  $\forall i. d \ i \in \text{UNIV} - \{0\}$  and d-to0:  $d \longrightarrow 0$ 
  then obtain m where  $\forall i \geq m. |d \ i - 0| < e$  using LIMSEQ-def e-pos
  dist-real-def by metis
  hence rd-ab:  $\bigwedge n. r + d(n+m) \in \{a < .. < b\}$  using dist-real-def ball-ab by (simp
  add: subset-eq)
  hence fd-msr:  $\bigwedge n. (\lambda x. (f(r+d(n+m)) x - f r x) / d(n+m)) \in \text{borel-measurable}$ 
  M
  using r-ab by (measurable; (intro f-msr)?; simp)
  hence limf-msr:  $(\lambda x. \lim (\lambda n. (f(r+d(n+m)) x - f r x) / d(n+m))) \in$ 
  borel-measurable M
  by measurable
  moreover have limf-Df:  $\bigwedge x. x \in ?CN \implies (\lambda n. (f(r+d(n+m)) x - f r x) / d(n+m)) \longrightarrow Df r x$ 
  / d(n+m)
  proof -
    fix x assume x-CN:  $x \in ?CN$ 
    hence ( $\lambda s. f s x$ ) field-differentiable (at r)
      using f-diffCN r-ab
      by (metis at-within-open differentiable-on-eq-field-differentiable-real
      open-greaterThanLessThan)
    hence  $((\lambda h. (f(r+h) x - f r x) / h) \longrightarrow Df r x)$  (at 0)
      apply (rewrite in asm DERIV-deriv-iff-field-differentiable[THEN sym])
      unfolding Df-def using r-ab x-CN by (simp add: DERIV-def)
    hence  $(\lambda i. (f(r+d i) x - f r x) / d i) \longrightarrow Df r x$ 
      apply (rewrite in asm tends-to-at-iff-sequentially)
      apply (rule allE'[where x=d], simp)
      unfolding comp-def using d-neq0 d-to0 by simp
    thus  $(\lambda n. (f(r+d(n+m)) x - f r x) / d(n+m)) \longrightarrow Df r x$ 
      by (rule LIMSEQ-ignore-initial-segment[where k=m])
  qed
  ultimately have Df-eq:
     $\bigwedge x. Df r x = \text{indicator } ?CN x * \lim (\lambda n. (f(r+d(n+m)) x - f r x) / d(n+m))$ 
  proof -
    fix x
    show Df r x = indicator ?CN x * lim ( $\lambda n. (f(r+d(n+m)) x - f r x) / d(n+m)$ )
    proof (cases  $x \in ?CN$ )
      case True
      hence lim ( $\lambda n. (f(r+d(n+m)) x - f r x) / d(n+m)$ ) = Df r x
        by (intro limI, rule limf-Df)
      thus ?thesis using True by simp
    next

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case False
thus ?thesis unfolding Df-def by simp
qed
qed
hence Df-msr: Df r ∈ borel-measurable M
apply (rewrite in λx. □ Df-eq)
apply (measurable; (rule limf-msr) ?)
using N-null unfolding null-sets-def by force
have ((λh. ((∫ x. f (r+h) x ∂M) − ∫ x. f r x ∂M) / h) ∘ d) —→
    ∫ x. lim (λn. (f (r + d (n+m)) x − f r x) / d (n+m)) ∂M
proof −
    have (λn. ∫ x. (f (r + d (n+m)) x − f r x) / d (n+m) ∂M) —→
        ∫ x. lim (λn. (f (r + d (n+m)) x − f r x) / d (n+m)) ∂M
    proof — by Lebesgue's Dominated Convergence Theorem
        have AE x in M. (λn. (f (r + d (n+m)) x − f r x) / d (n+m)) —→
            lim (λn. (f (r + d (n+m)) x − f r x) / d (n+m))
        using limf-Df Df-eq N-null by (smt (verit) DiffI AE-I' limI mem-Collect-eq
subset-eq)
        moreover have ∀n. AE x in M. norm ((f (r + d (n+m)) x − f r x) / d
(n+m)) ≤ g x
        proof −
            fix n
            { fix x assume x-CN: x ∈ ?CN
                let ?I = {r..(r + d (n+m))} ∪ {(r + d (n+m))..r}
                have f-diffI: (λs. f s x) differentiable-on ?I
                apply (rule differentiable-on-subset[where t={a<..<b}], rule f-diffCN,
rule x-CN)
                using r-ab rd-ab[of n] by (rewrite Un-subset-iff, auto)
                hence continuous-on ?I (λs. f s x) (λs. f s x) differentiable-on interior ?I
                apply −
                using differentiable-imp-continuous-on apply blast
                by (metis differentiable-on-subset interior-subset)
                then obtain t where t-01: t ∈ {0 <.. < 1} and
                    f-MVT: f (r + d (n+m)) x − f r x = d (n+m) * deriv (λs. f s x) (r
+ t * (d (n+m)))
                by (rule MVT-order-free)
                hence 0 < t t < 1 by simp-all
                hence rtd-ab: r + t * (d (n+m)) ∈ {a <.. < b}
                using r-ab rd-ab[of n]
            by simp (smt (verit, ccfv-threshold) mult-less-cancel-left mult-less-cancel-right2)
            have d (n+m) * deriv (λs. f s x) (r + t * (d (n+m))) =
                d (n+m) * Df (r + t * (d (n+m))) x
            proof −
                have r + t * (d (n+m)) ∈ {a <.. < b}
                using r-ab rd-ab[of n] t-01
                by (smt (verit) ball-eq-greaterThanLessThan dist-real-def
                    greaterThanLessThan-eq-iff greaterThanLessThan-eq-ball mem-ball
                    mult-le-cancel-right1 mult-minus-right mult-pos-neg)
            thus ?thesis unfolding Df-def using x-CN by simp

```

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qed
with f-MVT have  $(f(r + d(n+m))x - f r x) / d(n+m) = Df(r + t * (d(n+m)))x$ 
using d-neq0 by simp
moreover have  $|Df(r + t * (d(n+m)))x| \leq g x$  using Df-boundCN
x-CN rtd-ab by simp
ultimately have  $|(f(r + d(n+m))x - f r x) / d(n+m)| \leq g x$  by
simp }
thus AE x in M. norm  $((f(r + d(n+m))x - f r x) / d(n+m)) \leq g x$ 
unfolding real-norm-def using AE-I' N-null
by (smt (verit, ccfv-threshold) Diff-iff mem-Collect-eq subsetI)
qed
ultimately show  $((\lambda n. \int x. (f(r + d(n+m))x - f r x) / d(n+m)) \partial M)$ 
→
 $\int x. \lim (\lambda n. (f(r + d(n+m))x - f r x) / d(n+m)) \partial M$ 
using limf-msr fd-msr Df-bound
by (intro integral-dominated-convergence[where w=g], simp-all)
qed
moreover have  $\wedge n. ((\int x. f(r + d(n+m))x \partial M) - \int x. f r x \partial M) / d(n+m) =$ 
 $\int x. (f(r + d(n+m))x - f r x) / d(n+m) \partial M$ 
using d-neq0 apply simp
by (rewrite Bochner-Integration.integral-diff;
(rule f-integ | simp); (rule rd-ab | rule r-ab))
ultimately show ?thesis
unfolding comp-def using d-neq0
apply -
by (rule LIMSEQ-offset[where k=m]) simp
qed
moreover have  $(\int x. \lim (\lambda n. (f(r + d(n+m))x - f r x) / d(n+m)) \partial M)$ 
=
 $\int x. borel-measurable-extension M (\lambda y. deriv (\lambda s. f s y) r) x \partial M$ 
proof -
have  $(\int x. \lim (\lambda n. (f(r + d(n+m))x - f r x) / d(n+m)) \partial M) = \int x. Df_{r x} \partial M$ 
proof -
have AE x in M. lim  $(\lambda n. (f(r + d(n+m))x - f r x) / d(n+m)) = Df_{r x}$ 
proof -
{ fix x assume x-CN:  $x \in ?CN$ 
hence  $\lim (\lambda n. (f(r + d(n+m))x - f r x) / d(n+m)) = Df_{r x}$  by
(simp add: Df-eq) }
thus ?thesis using AE-I' N-null by (smt (verit, del-insts) DiffI mem-Collect-eq
subsetI)
qed
thus ?thesis using limf-msr Df-msr by (intro integral-cong-AE; simp)
qed
also have ... =  $\int x. borel-measurable-extension M (\lambda y. deriv (\lambda s. f s y) r)$ 
x ∂M

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proof -
  have  $\text{AE } x \text{ in } M. Df r x = \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x$  and
     $\text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) \in \text{borel-measurable } M$ 
proof -
  have  $\{x \in \text{space } M. \text{deriv} (\lambda s. f s x) r \neq Df r x\} \subseteq N$ 
proof -
  { fix  $x$  assume  $x \in ?CN$ 
    hence  $\text{deriv} (\lambda s. f s x) r = Df r x$  unfolding  $Df\text{-def}$  using  $r\text{-ab}$  by
       $\text{simp }$ 
    thus  $?thesis$  by  $\text{blast}$ 
    qed
    thus  $\text{AE } x \text{ in } M. Df r x = \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x$  and
       $\text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) \in \text{borel-measurable } M$ 
    using  $Df\text{-msr } N\text{-null}$ 
    apply -
    apply (rule  $\text{borel-measurable-extensionI2[where } S=N]; \text{simp}$ )
    by (rule  $\text{borel-measurable-extension-measurable[where } g=Df r]; \text{simp}$ )
    qed
    thus  $?thesis$  using  $Df\text{-msr}$  by (intro  $\text{integral-cong-AE}; \text{simp}$ )
    qed
    finally show  $?thesis$ .
    qed
    ultimately show  $((\lambda h. ((\int x. f (r+h) x \partial M) - \int x. f r x \partial M) / h) \circ d)$ 
  →
     $\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x \partial M$ 
    using  $\text{tendsto-cong-limit}$  by  $\text{simp}$ 
    qed
    thus  $((\lambda s. \int x. f s x \partial M) \text{ has-real-derivative}$ 
       $\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x \partial M) \text{ (at } r)$ 
      by (rewrite  $\text{DERIV-def}$ , rewrite  $\text{tendsto-at-iff-sequentially}$ )  $\text{simp}$ 
    qed

```

proposition *interchange-deriv-LINT*:

```

fixes  $a b :: \text{real}$  and  $f :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$  and  $g :: 'a \Rightarrow \text{real}$ 
assumes  $\bigwedge r. r \in \{a <.. < b\} \implies \text{integrable } M (f r)$  and
   $\text{AE } x \text{ in } M. (\lambda r. f r x) \text{ differentiable-on } \{a <.. < b\}$  and
   $\bigwedge r. r \in \{a <.. < b\} \implies (\lambda x. (\text{deriv} (\lambda r. f r x) r)) \in \text{borel-measurable } M$  and
   $\text{AE } x \text{ in } M. \forall r \in \{a <.. < b\}. |\text{deriv} (\lambda r. f r x) r| \leq g x \text{ integrable } M g$ 
  shows  $\bigwedge r. r \in \{a <.. < b\} \implies ((\lambda r. \int x. f r x \partial M) \text{ has-real-derivative}$ 
     $\int x. \text{deriv} (\lambda r. f r x) r \partial M) \text{ (at } r)$ 
proof -
  fix  $r$  assume  $r\text{-ab}: r \in \{a <.. < b\}$ 
  hence  $Df\text{-msr}: (\lambda x. \text{deriv} (\lambda s. f s x) r) \in \text{borel-measurable } M$  using  $\text{assms}$  by
   $\text{simp}$ 
  have  $((\lambda s. \int x. f s x \partial M) \text{ has-real-derivative}$ 

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 $\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x \partial M) \text{ (at } r)$ 
using assms r-ab by (intro interchange-deriv-LINT-general; simp)
moreover have ( $\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv} (\lambda s. f s y) r) x \partial M) =$ 
 $\int x. \text{deriv} (\lambda s. f s x) r \partial M$ 
apply (rule integral-cong-AE)
apply (rule borel-measurable-extension-measurable
    [where  $g = \lambda y. \text{deriv} (\lambda s. f s y) r$  and  $S = \{\}$ , simp-all add: Df-msr)
using borel-measurable-measurable-extension-AE Df-msr by (smt (verit) AE-cong)
ultimately show (( $\lambda r. \int x. f r x \partial M$ ) has-real-derivative  $\int x. \text{deriv} (\lambda r. f r x) r \partial M$ ) (at r)
    by simp
qed

```

proposition *interchange-deriv-LINT-set-general*:

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fixes  $a b :: \text{real}$  and  $f :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$  and  $g :: 'a \Rightarrow \text{real}$  and  $A :: 'a \text{ set}$ 
assumes  $A\text{-msr}: A \in \text{sets } M$  and
     $f\text{-integ}: \bigwedge r. r \in \{a < .. < b\} \implies \text{set-integrable } M A (f r)$  and
     $f\text{-diff}: \text{AE } x \in A \text{ in } M. (\lambda r. f r x) \text{ differentiable-on } \{a < .. < b\}$  and
     $Df\text{-bound}: \text{AE } x \in A \text{ in } M. \forall r \in \{a < .. < b\}. |\text{deriv} (\lambda r. f r x) r| \leq g x \text{ set-integrable } M A g$ 
shows  $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative}$ 
 $(\int x \in A. \text{set-borel-measurable-extension } M A (\lambda x. \text{deriv} (\lambda r. f r x) r) x \partial M))$ 
(at r)
proof –
    let  $?M\text{-}A = \text{restrict-space } M A$ 
    have  $\bigwedge r. r \in \{a < .. < b\} \implies \text{integrable } ?M\text{-}A (f r)$ 
        using A-msr f-integ set-integrable-restrict-space-iff by auto
    moreover have  $\text{AE } x \text{ in } ?M\text{-}A. (\lambda r. f r x) \text{ differentiable-on } \{a < .. < b\}$ 
        using AE-restrict-space-iff A-msr f-diff by (metis sets.Int-space-eq2)
    moreover have  $\text{AE } x \text{ in } ?M\text{-}A. \forall r \in \{a < .. < b\}. |\text{deriv} (\lambda r. f r x) r| \leq g x$  and
         $\text{integrable } ?M\text{-}A g$ 
        using A-msr Df-bound set-integrable-restrict-space-iff
            apply –
            by (simp add: AE-restrict-space-iff, auto)
    ultimately have  $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \text{integral}^L ?M\text{-}A (f r)) \text{ has-real-derivative}$ 
 $\text{integral}^L ?M\text{-}A (\text{borel-measurable-extension } ?M\text{-}A (\lambda x. \text{deriv} (\lambda r. f r x) r)))$  (at r)
        by (rule interchange-deriv-LINT-general[where M=restrict-space M A]) auto
        thus  $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative}$ 
 $(\int x \in A. \text{set-borel-measurable-extension } M A (\lambda x. \text{deriv} (\lambda r. f r x) r) x \partial M))$ 
(at r)
        unfolding set-borel-measurable-extension-def using assms
        by (rewrite set-lebesgue-integral-restrict-space, simp)+
qed

```

proposition *interchange-deriv-LINT-set*:

```

fixes  $a b :: \text{real}$  and  $f :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$  and  $g :: 'a \Rightarrow \text{real}$  and  $A :: 'a \text{ set}$ 
assumes  $A \in \text{sets } M$  and

```

$\bigwedge r. r \in \{a < .. < b\} \implies$ set-integrable $M A (f r)$ **and**
 $\forall E x \in A \text{ in } M. (\lambda r. f r x) \text{ differentiable-on } \{a < .. < b\}$ **and**
 $\bigwedge r. r \in \{a < .. < b\} \implies$ set-borel-measurable $M A (\lambda x. (\text{deriv } (\lambda r. f r x) r))$ **and**
 $\forall E x \in A \text{ in } M. \forall r \in \{a < .. < b\}. |\text{deriv } (\lambda r. f r x) r| \leq g x$ set-integrable $M A g$
shows $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative}$
 $(\int x \in A. \text{deriv } (\lambda r. f r x) r \partial M))$ (at r)
proof –
fix r **assume** $r-ab: r \in \{a < .. < b\}$
hence $Df-msr: \text{set-borel-measurable } M A (\lambda x. \text{deriv } (\lambda s. f s x) r)$ **using** assms
by *simp*
have $((\lambda s. \int x \in A. f s x \partial M) \text{ has-real-derivative}$
 $(\int x \in A. \text{set-borel-measurable-extension } M A (\lambda y. \text{deriv } (\lambda s. f s y) r) x \partial M))$
(at r)
using $\text{assms } r-ab$ **by** (*intro interchange-deriv-LINT-set-general*; *simp*)
moreover have $(\int x \in A. \text{set-borel-measurable-extension } M A (\lambda y. \text{deriv } (\lambda s. f s y) r) x \partial M) =$
 $(\int x \in A. \text{deriv } (\lambda s. f s x) r \partial M)$
apply (*rule set-lebesgue-integral-cong-AE2*, *simp add: assms*)
apply (*rule set-borel-measurable-extension-measurable*
[where $g = \lambda y. \text{deriv } (\lambda s. f s y) r$ **and** $S = \{\}$ **], simp-all add: Df-msr assms**)
using *set-borel-measurable-measurable-extension-AE* *Df-msr assms* **by** (*smt (verit) AE-cong*)
ultimately show
 $((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative } (\int x \in A. \text{deriv } (\lambda r. f r x) r \partial M))$ (at r)
by *simp*
qed

3 Additional Lemmas for the HOL–Probability Library

lemma (in finite-borel-measure)
fixes $F :: \text{real} \Rightarrow \text{real}$
assumes $\text{nondec}F : \bigwedge x y. x \leq y \implies F x \leq F y$ **and**
 $\text{right-cont-}F : \bigwedge a. \text{continuous } (\text{at-right } a) F$ **and**
 $\text{lim-}F\text{-at-bot} : (F \xrightarrow{} 0) \text{ at-bot}$ **and**
 $\text{lim-}F\text{-at-top} : (F \xrightarrow{} m) \text{ at-top}$ **and**
 $m : 0 \leq m$
shows $\text{emeasure-interval-measure-Ioi} : \text{emeasure } (\text{interval-measure } F) \{x <..\} = m - F x$
and $\text{measure-interval-measure-Ioi} : \text{measure } (\text{interval-measure } F) \{x <..\} = m - F x$
proof –
interpret $F\text{-FM} : \text{finite-measure interval-measure } F$
using *finite-borel-measure.axioms(1)* *finite-borel-measure-interval-measure lim-}F\text{-at-bot}*
lim-}F\text{-at-top m nondec}F right-cont-}F **by** *blast*
have $\text{UNIV} = \{..x\} \cup \{x <..\}$ **by** *auto*
moreover have $\{..x\} \cap \{x <..\} = \{\}$ **by** *auto*

```

ultimately have emeasure (interval-measure F) UNIV =
  emeasure (interval-measure F) {..x} + emeasure (interval-measure F) {x<..}
  by (simp add: plus-emeasure)
moreover have emeasure (interval-measure F) UNIV = m
  using assms interval-measure-UNIV by presburger
ultimately show ∃: emeasure (interval-measure F) {x<..} = m - F x
  using assms emeasure-interval-measure-Iic
  by (metis ennreal-add-diff-cancel-left ennreal-minus measure-interval-measure-Iic
       measure-nonneg top-neq-ennreal)
hence ennreal (measure (interval-measure F) {x<..}) = m - F x
  using emeasure-eq-measure by (metis emeasure-eq-ennreal-measure top-neq-ennreal)
moreover have ∀x. F x ≤ m
  using lim-F-at-top nondecF by (intro mono-at-top-le[where f=F]; simp add:
mono-def)
ultimately show measure (interval-measure F) {x<..} = m - F x
  using ennreal-inj F-FM.emeasure-eq-measure by force
qed

lemma (in prob-space) cond-prob-nonneg[simp]: cond-prob M P Q ≥ 0
  by (auto simp: cond-prob-def)

lemma (in prob-space) cond-prob-whole-1: cond-prob M P P = 1 if prob {ω ∈ space M. P ω} ≠ 0
  unfolding cond-prob-def using that by simp

lemma (in prob-space) cond-prob-0-null: cond-prob M P Q = 0 if prob {ω ∈ space M. Q ω} = 0
  unfolding cond-prob-def using that by simp

lemma (in prob-space) cond-prob-AE-prob:
  assumes {ω ∈ space M. P ω} ∈ events {ω ∈ space M. Q ω} ∈ events
  and AE ω in M. Q ω
  shows cond-prob M P Q = prob {ω ∈ space M. P ω}
proof -
  let ?setP = {ω ∈ space M. P ω}
  let ?setQ = {ω ∈ space M. Q ω}
  have [simp]: prob ?setQ = 1 using assms prob-Collect-eq-1 by simp
  hence cond-prob M P Q = prob (?setP ∩ ?setQ)
    unfolding cond-prob-def by (simp add: Collect-conj-eq2)
  also have ... = prob ?setP
  proof (rule antisym)
    show prob (?setP ∩ ?setQ) ≤ prob ?setP
      using assms finite-measure-mono inf-sup-ord(1) by blast
  next
    show prob ?setP ≤ prob (?setP ∩ ?setQ)
    proof -
      have prob (?setP ∩ ?setQ) = prob ?setP + prob ?setQ - prob (?setP ∪
        ?setQ)
        using assms by (smt (verit) finite-measure-Diff' finite-measure-Union'

```

```

sup-commute)
also have ... = prob ?setP + (1 - prob (?setP ∪ ?setQ)) by simp
also have ... ≥ prob ?setP by simp
finally show ?thesis .
qed
qed
finally show ?thesis .
qed

```

3.1 More Properties of *cdf*'s

```

context finite-borel-measure
begin

```

```

lemma cdf-diff-eq2:
assumes x ≤ y
shows cdf M y - cdf M x = measure M {x <.. y}
proof (cases «x = y»)
  case True
  thus ?thesis by force
next
  case False
  hence x < y using assms by simp
  thus ?thesis by (rule cdf-diff-eq)
qed

```

```

end

```

```

context prob-space
begin

```

```

lemma cdf-distr-measurable [measurable]:
assumes [measurable]: random-variable borel X
shows cdf (distr M borel X) ∈ borel-measurable borel
proof (rule borel-measurable-mono)
show mono (cdf (distr M borel X))
unfolding mono-def
using finite-borel-measure.cdf-nondecreasing
by (simp add: real-distribution.finite-borel-measure-M)
qed

```

```

lemma cdf-distr-P:
assumes random-variable borel X
shows cdf (distr M borel X) x = P(ω in M. X ω ≤ x)
unfolding cdf-def apply (rewrite measure-distr; (simp add: assms) ?)
unfolding vimage-def by (rule arg-cong[where f=prob], force)

```

```

lemma cdf-continuous-distr-P-lt:
assumes random-variable borel X isCont (cdf (distr M borel X)) x

```

```

shows cdf (distr M borel X) x = P(ω in M. X ω < x)
proof -
  have P(ω in M. X ω < x) = measure (distr M borel X) {..<x}
    apply (rewrite measure-distr, simp-all add: assms)
  unfolding vimage-def by simp (smt (verit) Collect-cong Int-def mem-Collect-eq)
  also have ... = measure (distr M borel X) ({..<x} ∪ {x})
  apply (rewrite finite-measure.measure-zero-union, simp-all add: assms finite-measure-distr)
    using finite-borel-measure.isCont-cdf real-distribution.finite-borel-measure-M
  assms by blast
  also have ... = measure (distr M borel X) {..x} by (metis ivl-disj-un-singleton(2))
  also have ... = cdf (distr M borel X) x unfolding cdf-def by simp
  finally show ?thesis by simp
qed

lemma cdf-distr-diff-P:
  assumes x ≤ y
  and random-variable borel X
  shows cdf (distr M borel X) y - cdf (distr M borel X) x = P(ω in M. x < X ω
  ∧ X ω ≤ y)
proof -
  interpret distrX-FBM: finite-borel-measure distr M borel X
    using real-distribution.finite-borel-measure-M real-distribution-distr assms by
  simp
  have cdf (distr M borel X) y - cdf (distr M borel X) x = measure (distr M borel
  X) {x<..y}
    by (rewrite distrX-FBM.cdf-diff-eq2; simp add: assms)
  also have ... = P(ω in M. x < X ω ∧ X ω ≤ y)
    apply (rewrite measure-distr; (simp add: assms)?)
    unfolding vimage-def by (rule arg-cong[where f=prob], force)
  finally show ?thesis .
qed

lemma cdf-distr-max:
  fixes c::real
  assumes [measurable]: random-variable borel X
  shows cdf (distr M borel (λx. max (X x) c)) u = cdf (distr M borel X) u *
  indicator {c..} u
proof (cases `c ≤ u`)
  case True
  thus ?thesis
    unfolding cdf-def
    apply (rewrite measure-distr; simp?)+
    by (smt (verit) Collect-cong atMost-iff vimage-def)
next
  case False
  thus ?thesis
    unfolding cdf-def
    apply (rewrite measure-distr; simp?)+
    by (smt (verit, best) Int-emptyI atMost-iff measure-empty vimage-eq)

```

qed

```

lemma cdf-distr-min:
  fixes c::real
  assumes [measurable]: random-variable borel X
  shows cdf (distr M borel (λx. min (X x) c)) u =
    1 - (1 - cdf (distr M borel X) u) * indicator {..} u
  proof (cases `c ≤ u`)
    case True
    thus ?thesis
      unfolding cdf-def
      using real-distribution.finite-borel-measure-M real-distribution-distr
      apply (rewrite measure-distr; simp?)
      by (smt (verit, del-insts) Int-absorb1 atMost-iff prob-space subset-vimage-iff)
  next
    case False
    thus ?thesis
      unfolding cdf-def
      using real-distribution.finite-borel-measure-M real-distribution-distr
      apply (rewrite measure-distr; simp?)+
      using prob-space-axioms assms
      by (smt (verit) Collect-cong Int-def atMost-iff prob-space prob-space.cdf-distr-P
vimage-eq)
  qed

lemma cdf-distr-floor-P:
  fixes X :: 'a ⇒ real
  assumes [measurable]: random-variable borel X
  shows cdf (distr M borel (λx. ⌊X x⌋)) u = ℙ(x in M. X x < ⌊u⌋ + 1)
  unfolding cdf-def
  apply (rewrite measure-distr; simp?)
  apply (rule arg-cong[where f=prob])
  unfolding vimage-def using floor-le-iff le-floor-iff by blast

lemma expectation-nonneg-tail:
  assumes [measurable]: random-variable borel X
  and X-nonneg: ∀x. x ∈ space M ⇒ X x ≥ 0
  defines F u ≡ cdf (distr M borel X) u
  shows (∫⁺x. ennreal (X x) ∂M) = (∫⁺u∈{0..}. ennreal (1 - F u) ∂lborel)
  proof -
    let ?distrX = distr M borel X
    have (∫⁺x. ennreal (X x) ∂M) = (∫⁺u. ennreal u ∂?distrX)
      by (rewrite nn-integral-distr; simp)
    also have ... = (∫⁺u∈{0..}. ennreal u ∂?distrX)
      by (rule nn-integral-distr-set; simp add: X-nonneg)
    also have ... = (∫⁺u∈{0..}. (∫⁺v∈{0..}. indicator {..

```

```

apply (rewrite nn-integral-indicator, measurable, simp)
  by (metis atLeastLessThan-def diff-zero emeasure-lborel-Ico inf.commute)
thus ?thesis by (metis (no-types, lifting) indicator-eq-0-iff mult-eq-0-iff)
qed
also have ... = ( $\int^+ v \in \{0..\} . (\int^+ u \in \{0..\} . indicator \{.. < u\} v \partial?distrX) \partial borel$ )
proof -
  have ( $\int^+ u \in \{0..\} . (\int^+ v \in \{0..\} . indicator \{.. < u\} v \partial?distrX) \partial?distrX$ ) =
     $\int^+ u . (\int^+ v. indicator \{.. < u\} v * indicator \{0..\} v * indicator \{0..\} u \partial borel)$ 
 $\partial?distrX$ 
  by (rewrite nn-integral-multc; simp)
  also have ... =
     $\int^+ v . (\int^+ u. indicator \{.. < u\} v * indicator \{0..\} v * indicator \{0..\} u \partial?distrX) \partial borel$ 
  apply (rewrite pair-sigma-finite.Fubini'; simp?)
  using pair-sigma-finite.intro assms
  prob-space-distr prob-space-imp-sigma-finite sigma-finite-lborel
  apply blast
  by measurable auto
  also have ... = ( $\int^+ v \in \{0..\} . (\int^+ u \in \{0..\} . indicator \{.. < u\} v \partial?distrX) \partial borel$ )
  apply (rewrite nn-integral-multc[THEN sym]; measurable; simp?)
  apply (rule nn-integral-cong)+
  using mult.assoc mult.commute by metis
  finally show ?thesis by simp
qed
also have ... = ( $\int^+ v \in \{0..\} . (\int^+ u. indicator \{v < ..\} u \partial?distrX) \partial borel$ )
  apply (rule nn-integral-cong)
  apply (rewrite nn-integral-multc[THEN sym], measurable; (simp del: nn-integral-indicator)?)+
  apply (rule nn-integral-cong)
  using lessThan-iff greaterThan-iff atLeast-iff indicator-simps
  by (smt (verit, del-insts) mult-1 mult-eq-0-iff)
also have ... = ( $\int^+ v \in \{0..\} . ennreal (1 - F v) \partial borel$ )
  apply (rule nn-integral-cong, simp)
  apply (rewrite emeasure-distr; simp?)
  apply (rewrite vimage-compl-atMost[THEN sym])
  unfolding F-def cdf-def
  apply (rewrite measure-distr; simp?)
  apply (rewrite prob-compl[THEN sym], simp)
  by (metis (no-types, lifting) Diff-Compl Diff-Diff-Int Int-commute emeasure-eq-measure)
  finally show ?thesis .
qed

lemma expectation-nonpos-tail:
assumes [measurable]: random-variable borel X
  and X-nonpos:  $\bigwedge x. x \in space M \implies X x \leq 0$ 
defines F u ≡ cdf (distr M borel X) u
shows ( $\int^+ x. ennreal (- X x) \partial M$ ) = ( $\int^+ u \in \{..0\}. ennreal (F u) \partial borel$ )
proof -
  let ?distrX = distr M borel X

```

```

have ( $\int^+ x. \text{ennreal} (-X x) \partial M$ ) = ( $\int^+ u. \text{ennreal} (-u) \partial ?\text{distr} X$ )
  by (rewrite nn-integral-distr; simp)
also have ... = ( $\int^+ u \in \{..0\}. \text{ennreal} (-u) \partial ?\text{distr} X$ )
proof -
  have [simp]:  $\{..0::\text{real}\} \cup \{0 <..\} = \text{UNIV}$  by force
  have ( $\int^+ u. \text{ennreal} (-u) \partial ?\text{distr} X$ ) =
    ( $\int^+ u \in \{..0\}. \text{ennreal} (-u) \partial ?\text{distr} X$ ) + ( $\int^+ u \in \{0 <..\}. \text{ennreal} (-u) \partial ?\text{distr} X$ )
    by (rewrite nn-integral-disjoint-pair[THEN sym], simp-all, force)
  also have ... = ( $\int^+ u \in \{..0\}. \text{ennreal} (-u) \partial ?\text{distr} X$ )
    apply (rewrite nn-integral-zero'[of  $\lambda u. \text{ennreal} (-u) * \text{indicator} \{0 <..\} u$ ; simp?])
    unfolding indicator-def using always-eventually ennreal-lt-0 by force
    finally show ?thesis .
qed
also have ... = ( $\int^+ u \in \{..0\}. (\int^+ v \in \{..0\}. \text{indicator} \{u..\} v \partial \text{lborel}) \partial ?\text{distr} X$ )
proof -
  have  $\bigwedge u::\text{real}. u \in \{..0\} \implies \text{ennreal} (-u) = (\int^+ v \in \{..0\}. \text{indicator} \{u..\} v \partial \text{lborel})$ 
    by (rewrite indicator-inter-arith[THEN sym]) simp
  thus ?thesis by (metis (no-types, lifting) indicator-eq-0-iff mult-eq-0-iff)
qed
also have ... = ( $\int^+ v \in \{..0\}. (\int^+ u \in \{..0\}. \text{indicator} \{u..\} v \partial ?\text{distr} X) \partial \text{lborel}$ )
proof -
  have ( $\int^+ u \in \{..0\}. (\int^+ v \in \{..0\}. \text{indicator} \{u..\} v \partial \text{lborel}) \partial ?\text{distr} X$ ) =
     $\int^+ u. (\int^+ v. \text{indicator} \{u..\} v * \text{indicator} \{..0\} v * \text{indicator} \{..0\} u \partial \text{lborel})$ 
     $\partial ?\text{distr} X$ 
    by (rewrite nn-integral-multc; simp)
  also have ... =
     $\int^+ v. (\int^+ u. \text{indicator} \{u..\} v * \text{indicator} \{..0\} v * \text{indicator} \{..0\} u \partial ?\text{distr} X)$ 
     $\partial \text{lborel}$ 
    apply (rewrite pair-sigma-finite.Fubini'; simp?)
    using pair-sigma-finite.intro assms
    prob-space-distr prob-space-imp-sigma-finite sigma-finite-lborel
    apply blast
    by measurable auto
  also have ... = ( $\int^+ v \in \{..0\}. (\int^+ u \in \{..0\}. \text{indicator} \{u..\} v \partial ?\text{distr} X) \partial \text{lborel}$ )
    apply (rewrite nn-integral-multc[THEN sym]; measurable; simp?)
    apply (rule nn-integral-cong)+
    using mult.assoc mult.commute by metis
    finally show ?thesis by simp
qed
also have ... = ( $\int^+ v \in \{..0\}. (\int^+ u. \text{indicator} \{..v\} u \partial ?\text{distr} X) \partial \text{lborel}$ )
  apply (rule nn-integral-cong)
  apply (rewrite nn-integral-multc[THEN sym], measurable; (simp del: nn-integral-indicator)?)+
  apply (rule nn-integral-cong)
  using atMost-iff atLeast-iff indicator-simps by (smt (verit, del-insts) mult-1
mult-eq-0-iff)
also have ... = ( $\int^+ v \in \{..0\}. \text{ennreal} (F v) \partial \text{lborel}$ )
  apply (rule nn-integral-cong, simp)

```

```

apply (rewrite emeasure-distr; simp?)
unfolding F-def cdf-def
by (rewrite measure-distr; simp add: emeasure-eq-measure)
finally show ?thesis .
qed

proposition expectation-tail:
assumes [measurable]: integrable M X
defines F u ≡ cdf (distr M borel X) u
shows expectation X = (LBINT u:{0..}. 1 - F u) - (LBINT u:{..0}. F u)
proof -
have expectation X = expectation (λx. max (X x) 0) + expectation (λx. min (X x) 0)
using integrable-max integrable-min
apply (rewrite Bochner-Integration.integral-add[THEN sym], measurable)
by (rule Bochner-Integration.integral-cong; simp)
also have ... = expectation (λx. max (X x) 0) - expectation (λx. - min (X x))
0 by force
also have ... = (LBINT u:{0..}. 1 - F u) - (LBINT u:{..0}. F u)
proof -
have expectation (λx. max (X x) 0) = (LBINT u:{0..}. 1 - F u)
proof -
have expectation (λx. max (X x) 0) = enn2real (ʃ+x. ennreal (max (X x) 0) ∂M)
by (rule integral-eq-nn-integral; simp)
also have ... = enn2real (ʃ+u∈{0..}. ennreal (1 - F u) ∂lborel)
apply (rewrite expectation-nonneg-tail; simp?)
apply (rewrite cdf-distr-max, simp)
unfolding F-def
by (metis (opaque-lifting) indicator-simps mult.commute mult-1 mult-eq-0-iff)
also have ... = enn2real (ʃ+u. ennreal ((1 - F u) * indicator {0..} u)
∂lborel)
by (simp add: indicator-mult-ennreal mult.commute)
also have ... = (LBINT u:{0..}. 1 - F u)
apply (rewrite integral-eq-nn-integral[THEN sym], simp add: F-def)
unfolding F-def using real-distribution.cdf-bounded-prob apply force
unfolding set-lebesgue-integral-def by (rule Bochner-Integration.integral-cong;
simp)
finally show ?thesis .
qed
moreover have expectation (λx. - min (X x) 0) = (LBINT u:{..0}. F u)
proof -
have expectation (λx. - min (X x) 0) = enn2real (ʃ+x. ennreal (- min (X x) 0) ∂M)
by (rule integral-eq-nn-integral; simp)
also have ... = enn2real (ʃ+u∈{..0}. ennreal (F u) ∂lborel)
proof -
let ?distrminX = distr M borel (λx. min (X x) 0)
have [simp]: sym-diff {..0} {..<0} = {0::real} by force

```

```

have enn2real ( $\int^+ x. \text{ennreal} (- \min (X x) 0) \partial M) =$ 
enn2real ( $\int^+ u \in \{..0\}. \text{ennreal} (\text{cdf} ?\text{distrmin}_X u) \partial borel)$ 
  by (rewrite expectation-nonpos-tail; simp)
also have ... = enn2real ( $\int^+ u \in \{..<0\}. \text{ennreal} (\text{cdf} ?\text{distrmin}_X u) \partial borel)$ 
  by (rewrite nn-integral-null-delta, auto)
also have ... = enn2real ( $\int^+ u \in \{..<0\}. \text{ennreal} (F u) \partial borel)$ 
  apply (rewrite cdf-distr-min, simp)
  apply (rule arg-cong[where f=enn2real], rule nn-integral-cong)
unfolding F-def by (smt (verit) indicator-simps mult-cancel-left1 mult-eq-0-iff)
also have ... = enn2real ( $\int^+ u \in \{..0\}. \text{ennreal} (F u) \partial borel)$ 
  by (rewrite nn-integral-null-delta, auto simp add: sup-commute)
finally show ?thesis .
qed
also have ... = enn2real ( $\int^+ u. \text{ennreal} (F u * \text{indicator} \{..0\} u) \partial borel)$ 
  using mult.commute indicator-mult-ennreal by metis
also have ... = (LBINT u:{..0}. F u)
  apply (rewrite integral-eq-nn-integral[THEN sym], simp add: F-def)
  unfolding F-def
  using finite-borel-measure.cdf-nonneg real-distribution.finite-borel-measure-M
apply simp
unfolding set-lebesgue-integral-def by (rule Bochner-Integration.integral-cong;
simp)
  finally show ?thesis .
qed
ultimately show ?thesis by simp
qed
finally show ?thesis .
qed

proposition distributed-deriv-cdf:
assumes [measurable]: random-variable borel X
defines F u ≡ cdf (distr M borel X) u
assumes finite S ∧ x. x ∉ S ⇒ (F has-real-derivative f x) (at x)
  and continuous-on UNIV F f ∈ borel-measurable lborel
shows distributed M lborel X f
proof –
have FBM: finite-borel-measure (distr M borel X)
  using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
then interpret distrX-FBM: finite-borel-measure distr M borel X .
have FBML: finite-borel-measure (distr M lborel X) using FBM distr-borel-lborel
by (smt (verit))
then interpret distrlX-FBM: finite-borel-measure distr M lborel X .
have [simp]: ( $\lambda x. \text{ennreal} (f x)) \in \text{borel-measurable borel}$  using assms by simp
moreover have distr M lborel X = density lborel f
proof –
have  $\bigwedge a b. a \leq b \Rightarrow \text{emeasure} (\text{distr} M \text{lborel} X) \{a <.. b\} < \top$ 
  using distrlX-FBM.emeasure-real less-top-ennreal by blast
moreover have  $\bigwedge a b. a \leq b \Rightarrow$ 

```

```

emeasure (distr M lborel X) {a<..b} = emeasure (density lborel f) {a<..b}
proof -
  fix a b :: real assume a ≤ b
  hence [simp]: sym-diff {a<..b} {a..b} = {a} by force
    have emeasure (density lborel f) {a<..b} = (ʃ+x∈{a<..b}. ennreal (f x)
      ∂lborel)
      by (rewrite emeasure-density; simp)
    also have ... = (ʃ+x∈{a..b}. ennreal (f x) ∂lborel) by (rewrite nn-integral-null-delta,
      auto)
    also have ... = ∫+x. ennreal (indicat-real {a..b} x * f x) ∂lborel
      by (metis indicator-mult-ennreal mult.commute)
    also have ... = ennreal (F b - F a)
    proof -
      define g where g x = (if x ∈ S then 0 else f x) for x :: real
      have [simp]: ∀x. g x ≥ 0
        unfolding g-def
        apply (split if-split, auto)
        apply (rule mono-on-imp-deriv-nonneg[of UNIV F], auto)
        unfolding F-def mono-on-def using distrX-FBM.cdf-nondecreasing apply
      blast
      using assms unfolding F-def by force
      have (ʃ+x. ennreal (indicat-real {a..b} x * f x) ∂lborel)
        = ∫+x. ennreal (indicat-real {a..b} x * g x) ∂lborel
        apply (rule nn-integral-cong-AE)
        apply (rule AE-mp[where P=λx. x ∉ S])
        using assms finite-imp-null-set-lborel AE-not-in apply blast
        unfolding g-def by simp
      also have ... = ennreal (F b - F a)
        apply (rewrite nn-integral-has-integral-lebesgue, simp)
        apply (rule fundamental-theorem-of-calculus-strong[of S], auto simp: ‹a ≤ b› g-def assms)
        using has-real-derivative-iff-has-vector-derivative assms apply presburger
        using assms continuous-on-subset subsetI by fastforce
        finally show ?thesis .
    qed
    also have ... = emeasure (distr M lborel X) {a <.. b}
      apply (rewrite distrlX-FBM.emeasure-Ioc, simp add: ‹a ≤ b›)
      unfolding F-def cdf-def
      apply (rewrite ennreal-minus[THEN sym], simp)+
      by (metis distr-borel-lborel)
    finally show emeasure (distr M lborel X) {a<..b} = emeasure (density lborel
      f) {a<..b}
      by simp
    qed
    ultimately show ?thesis by (intro measure-eqI-Ioc; simp)
  qed
  ultimately show ?thesis unfolding distributed-def by auto
qed

```

end

Define the conditional probability space. This is obtained by rescaling the restricted space of a probability space.

3.2 Conditional Probability Space

```
lemma restrict-prob-space:
  assumes measure-space Ω A μ a ∈ A
    and 0 < μ a μ a < ∞
  shows prob-space (scale-measure (1 / μ a) (restrict-space (measure-of Ω A μ)
a))
proof
  let ?M = measure-of Ω A μ
  let ?Ma = restrict-space ?M a
  let ?rMa = scale-measure (1 / μ a) ?Ma
  have emeasure ?rMa (space ?rMa) = 1 / μ a * emeasure ?Ma (space ?rMa) by
simp
  also have ... = 1 / μ a * emeasure ?M (space ?rMa)
  using assms
  apply (rewrite emeasure-restrict-space)
  apply (simp add: measure-space-def sigma-algebra.sets-measure-of-eq)
  by (simp add: space-restrict-space space-scale-measure) simp
  also have ... = 1 / μ a * emeasure ?M (space ?Ma) by (rewrite space-scale-measure)
simp
  also have ... = 1 / μ a * emeasure ?M a
  using assms
  apply (rewrite space-restrict-space2)
  by (simp add: measure-space-closed) +
  also have ... = 1
  using assms measure-space-def
  apply (rewrite emeasure-measure-of-sigma, blast+)
  by (simp add: ennreal-divide-times)
  finally show emeasure ?rMa (space ?rMa) = 1 .
qed

definition cond-prob-space :: 'a measure ⇒ 'a set ⇒ 'a measure (infix `|` 200)
  where M|A ≡ scale-measure (1 / emeasure M A) (restrict-space M A)

context prob-space
begin

lemma cond-prob-space-whole[simp]: M `|` space M = M
  unfolding cond-prob-space-def using emeasure-space-1 by simp

lemma cond-prob-space-correct:
  assumes A ∈ events prob A > 0
  shows prob-space (M|A)
  unfolding cond-prob-space-def
```

```

apply (subst(2) measure-of-of-measure[of M, THEN sym])
using assms
by (intro restrict-prob-space; (simp add: measure-space)?; simp-all add: emeasure-eq-measure)

lemma space-cond-prob-space:
assumes A ∈ events
shows space (M|A) = A
unfolding cond-prob-space-def using assms by (simp add: space-scale-measure)

lemma sets-cond-prob-space: sets (M|A) = (∩) A ` events
unfolding cond-prob-space-def by (metis sets-restrict-space sets-scale-measure)

lemma measure-cond-prob-space:
assumes A ∈ events B ∈ events
and prob A > 0
shows measure (M|A) (B ∩ A) = prob (B ∩ A) / prob A
proof -
have 1 / emeasure M A = ennreal (1 / prob A)
using assms by (smt (verit) divide-ennreal emeasure-eq-measure ennreal-1)
hence measure (M|A) (B ∩ A) = (1 / prob A) * measure (restrict-space M A)
(B ∩ A)
unfolding cond-prob-space-def using measure-scale-measure by force
also have ... = (1 / prob A) * prob (B ∩ A)
using measure-restrict-space assms by (metis inf.cobounded2 sets.Int-space-eq2)
also have ... = prob (B ∩ A) / prob A by simp
finally show ?thesis .
qed

corollary measure-cond-prob-space-subset:
assumes A ∈ events B ∈ events B ⊆ A
and prob A > 0
shows measure (M|A) B = prob B / prob A
proof -
have B = B ∩ A using assms by auto
moreover have measure (M|A) (B ∩ A) = prob (B ∩ A) / prob A
using assms measure-cond-prob-space by simp
ultimately show ?thesis by auto
qed

lemma cond-cond-prob-space:
assumes A ∈ events B ∈ events B ⊆ A prob B > 0
shows (M|A)|B = M|B
proof (rule measure-eqI)
have pA-pos[simp]: prob A > 0 using assms by (smt (verit, ccfv-SIG) finite-measure-mono)
interpret MA-PS: prob-space M|A using cond-prob-space-correct assms by simp
interpret MB-PS: prob-space M|B using cond-prob-space-correct assms by simp
have 1 / emeasure M A = ennreal (1 / prob A)

```

```

using pA-pos by (smt (verit, ccfv-SIG) divide-ennreal emeasure-eq-measure
ennreal-1)
hence [simp]:  $0 < MA\text{-PS}.prob B$ 
using assms pA-pos
by (metis divide-eq-0-iff measure-cond-prob-space-subset zero-less-measure-iff)
have [simp]:  $B \in MA\text{-PS}.events$ 
using assms by (rewrite sets-cond-prob-space, unfold image-def) blast
have [simp]: finite-measure  $((M\setminus A)\setminus B)$ 
by (rule prob-space.finite-measure, rule prob-space.cond-prob-space-correct,
simp-all add: MA-PS.prob-space-axioms)
show sets-MAB: sets  $((M\setminus A)\setminus B) = sets (M\setminus B)$ 
apply (rewrite prob-space.sets-cond-prob-space)
using MA-PS.prob-space-axioms apply presburger
apply (rewrite sets-cond-prob-space, unfold image-def) +
using assms by blast
show  $\bigwedge C. C \in sets ((M\setminus A)\setminus B) \implies emeasure ((M\setminus A)\setminus B) C = emeasure (M\setminus B)$ 
C
proof -
fix C assume C ∈ sets  $((M\setminus A)\setminus B)$ 
hence C ∈ sets  $(M\setminus B)$  using sets-MAB by simp
from this obtain D where D ∈ events  $C = B \cap D$ 
by (rewrite in asm sets-cond-prob-space, auto)
hence [simp]: C ∈ events and [simp]:  $C \subseteq B$  and [simp]:  $C \subseteq A$  using assms
by auto
hence [simp]: C ∈ MA-PS.events
using assms by (rewrite sets-cond-prob-space, unfold image-def) blast
show emeasure  $((M\setminus A)\setminus B) C = emeasure (M\setminus B) C$ 
apply (rewrite finite-measure.emeasure-eq-measure, simp) +
apply (rewrite ennreal-inj, simp-all)
apply (rewrite prob-space.measure-cond-prob-space-subset,
simp-all add: assms prob-space-axioms MA-PS.prob-space-axioms) +
using pA-pos by fastforce
qed
qed

lemma cond-prob-space-prob:
assumes[measurable]: Measurable.pred M P Measurable.pred M Q
and  $\mathcal{P}(x \text{ in } M. Q x) > 0$ 
shows measure  $(M \downarrow \{x \in space M. Q x\}) \{x \in space M. P x \wedge Q x\} = \mathcal{P}(x \text{ in } M. P x \mid Q x)$ 
proof -
let ?SetP =  $\{x \in space M. P x\}$ 
let ?SetQ =  $\{x \in space M. Q x\}$ 
have measure  $(M \downarrow ?SetQ) \{x \in space M. P x \wedge Q x\} = measure (M \downarrow ?SetQ)$ 
(?SetP ∩ ?SetQ)
by (smt (verit, ccfv-SIG) Collect-cong Int-def mem-Collect-eq)
also have ... = prob (?SetP ∩ ?SetQ) / prob ?SetQ
using assms by (rewrite measure-cond-prob-space; simp?)
also have ... =  $\mathcal{P}(x \text{ in } M. P x \mid Q x)$ 

```

unfolding cond-prob-def assms by (smt (verit) Collect-cong Int-def mem-Collect-eq)
finally show ?thesis .

qed

lemma cond-prob-space-cond-prob:

assumes [measurable]: Measurable.pred M P Measurable.pred M Q
and $\mathcal{P}(x \text{ in } M. Q x) > 0$
shows $\mathcal{P}(x \text{ in } M. P x \mid Q x) = \mathcal{P}(x \text{ in } (M \setminus \{x \in \text{space } M. Q x\}). P x)$

proof –

let ?setQ = $\{x \in \text{space } M. Q x\}$

have $\mathcal{P}(x \text{ in } M. P x \mid Q x) = \text{measure } (M \setminus ?setQ) \{x \in \text{space } M. P x \wedge Q x\}$

using cond-prob-space-prob assms by simp

also have ... = $\mathcal{P}(x \text{ in } (M \setminus ?setQ). P x)$

proof –

have $\{x \in \text{space } M. P x \wedge Q x\} = \{x \in \text{space } (M \setminus ?setQ). P x\}$

using space-cond-prob-space assms by force

thus ?thesis by simp

qed

finally show ?thesis .

qed

lemma cond-prob-neg:

assumes[measurable]: Measurable.pred M P Measurable.pred M Q

and $\mathcal{P}(x \text{ in } M. Q x) > 0$

shows $\mathcal{P}(x \text{ in } M. \neg P x \mid Q x) = 1 - \mathcal{P}(x \text{ in } M. P x \mid Q x)$

proof –

let ?setP = $\{x \in \text{space } M. P x\}$

let ?setQ = $\{x \in \text{space } M. Q x\}$

interpret setQ-PS: prob-space $M \setminus ?setQ$ **using** cond-prob-space-correct assms by simp

have [simp]: $\{x \in \text{space } (M \setminus ?setQ). P x\} \in \text{setQ-PS.events}$

proof –

have $\{x \in \text{space } (M \setminus ?setQ). P x\} = ?setQ \cap ?setP$ **using** space-cond-prob-space by force

thus ?thesis **using** sets-cond-prob-space by simp

qed

have $\mathcal{P}(x \text{ in } M. \neg P x \mid Q x) = \mathcal{P}(x \text{ in } M \setminus ?setQ. \neg P x)$

by (rewrite cond-prob-space-cond-prob; simp add: assms)

also have ... = $1 - \mathcal{P}(x \text{ in } M \setminus ?setQ. P x)$ **by** (rewrite setQ-PS.prob-neg, simp-all add: assms)

also have ... = $1 - \mathcal{P}(x \text{ in } M. P x \mid Q x)$

by (rewrite cond-prob-space-cond-prob; simp add: assms)

finally show ?thesis .

qed

lemma random-variable-cond-prob-space:

assumes $A \in \text{events prob } A > 0$

and [measurable]: random-variable borel X

shows $X \in \text{borel-measurable } (M \setminus A)$

```

proof (rule borel-measurableI)
  fix  $S :: \text{`b set}$ 
  assume [measurable]: open S
  show  $X -` S \cap \text{space } (M \downharpoonright A) \in \text{sets } (M \downharpoonright A)$ 
    apply (rewrite space-cond-prob-space, simp add: assms)
    apply (rewrite sets-cond-prob-space, simp add: image-def)
    apply (rule bexI[of - ] X -` S \cap \text{space } M; measurable)
    using sets.Int-space-eq2 Int-commute assms by auto
  qed

lemma AE-cond-prob-space-iff:
  assumes  $A \in \text{events prob } A > 0$ 
  shows  $(\text{AE } x \text{ in } M \downharpoonright A. P x) \longleftrightarrow (\text{AE } x \text{ in } M. x \in A \longrightarrow P x)$ 
proof –
  have [simp]:  $1 / \text{emeasure } M A > 0$ 
  using assms divide-ennreal emeasure-eq-measure ennreal-1
  by (smt (verit) divide-pos-pos ennreal-eq-0-iff not-gr-zero)
  show ?thesis
    unfolding cond-prob-space-def
    apply (rewrite AE-scale-measure-iff, simp)
    by (rewrite AE-restrict-space-iff; simp add: assms)
  qed

lemma integral-cond-prob-space-nn:
  assumes  $A \in \text{events prob } A > 0$ 
  and [measurable]: random-variable borel X
  and [nonneg]:  $\text{AE } x \text{ in } M. x \in A \longrightarrow 0 \leq X x$ 
  shows  $\text{integral}^L (M \downharpoonright A) X = \text{expectation } (\lambda x. \text{indicator } A x * X x) / \text{prob } A$ 
proof –
  have [simp]:  $X \in \text{borel-measurable } (M \downharpoonright A)$ 
  by (rule random-variable-cond-prob-space; (simp add: assms))
  have [simp]:  $\text{AE } x \text{ in } (M \downharpoonright A). 0 \leq X x$ 
  by (rewrite AE-cond-prob-space-iff; simp add: assms)
  have [simp]: random-variable borel  $(\lambda x. \text{indicat-real } A x * X x)$ 
  using borel-measurable-indicator assms by force
  have [simp]:  $\text{AE } x \text{ in } M. 0 \leq \text{indicat-real } A x * X x$  using nonneg by fastforce
  have  $\text{integral}^L (M \downharpoonright A) X = \text{enn2real } (\int^+ x. \text{ennreal } (X x) \partial(M \downharpoonright A))$ 
  by (rewrite integral-eq-nn-integral; simp)
  also have ...  $= \text{enn2real } (1 / \text{prob } A * \int^+ x. \text{ennreal } (X x) \partial(\text{restrict-space } M A))$ 
  unfolding cond-prob-space-def
  apply (rewrite nn-integral-scale-measure, simp add: measurable-restrict-space1)
  using divide-ennreal emeasure-eq-measure ennreal-1 assms by (smt (verit))
  also have ...  $= \text{enn2real } (1 / \text{prob } A * (\int^+ x. \text{ennreal } (\text{indicator } A x * X x) \partial M))$ 
  apply (rewrite nn-integral-restrict-space, simp add: assms)
  by (metis indicator-mult-ennreal mult.commute)
  also have ...  $= \text{integral}^L M (\lambda x. \text{indicator } A x * X x) / \text{prob } A$ 
  apply (rewrite integral-eq-nn-integral; simp?)

```

```

    by (simp add: divide-nonneg-pos enn2real-mult)
  finally show ?thesis by simp
qed

end

```

Define the complementary cumulative distribution function, also known as tail distribution. The theory presented below is a slight modification of the subsection "Properties of cdf's" in the theory *Distribution-Functions*.

3.3 Complementary Cumulative Distribution Function

```

definition ccdf :: real measure ⇒ real ⇒ real
  where ccdf M ≡ λx. measure M {x<..}
    — complementary cumulative distribution function (tail distribution)

```

```

lemma ccdf-def2: ccdf M x = measure M {x<..}
  by (simp add: ccdf-def)

```

```

context finite-borel-measure
begin

```

```

lemma add-cdf-ccdf: cdf M x + ccdf M x = measure M (space M)

```

```

proof -

```

```

  have {..x} ∪ {x<..} = UNIV by auto
  moreover have {..x} ∩ {x<..} = {} by auto

```

```

  ultimately have emeasure M {..x} + emeasure M {x<..} = emeasure M UNIV

```

```

    using plus-emeasure M-is-borel atMost-borel greaterThan-borel by metis

```

```

  hence measure M {..x} + measure M {x<..} = measure M UNIV

```

```

    using finite-emeasure-space emeasure-eq-measure ennreal-inj

```

```

    by (smt (verit, ccfv-SIG) ennreal-plus measure-nonneg)

```

```

  thus ?thesis unfolding cdf-def ccdf-def using borel-UNIV by simp

```

```

qed

```

```

lemma ccdf-cdf: ccdf M = (λx. measure M (space M) - cdf M x)

```

```

  by (rule ext) (smt (verit) add-cdf-ccdf)

```

```

lemma cdf-ccdf: cdf M = (λx. measure M (space M) - ccdf M x)

```

```

  by (rule ext) (smt (verit) add-cdf-ccdf)

```

```

lemma isCont-cdf-ccdf: isCont (cdf M) x ↔ isCont (ccdf M) x

```

```

proof

```

```

  show isCont (cdf M) x ==> isCont (ccdf M) x by (rewrite ccdf-cdf) auto

```

```

next

```

```

  show isCont (ccdf M) x ==> isCont (cdf M) x by (rewrite cdf-ccdf) auto

```

```

qed

```

```

lemma isCont-ccdf: isCont (ccdf M) x ↔ measure M {x} = 0

```

```

  using isCont-cdf-ccdf isCont-cdf by simp

```

```

lemma continuous-cdf-ccdf:
  shows continuous F (cdf M)  $\longleftrightarrow$  continuous F (ccdf M)
    (is ?LHS  $\longleftrightarrow$  ?RHS)
proof
  assume ?LHS
  thus ?RHS using continuous-diff continuous-const by (rewrite ccdf-cdf) blast
next
  assume ?RHS
  thus ?LHS using continuous-diff continuous-const by (rewrite cdf-ccdf) blast
qed

lemma has-real-derivative-cdf-ccdf:
  ( $(\text{cdf } M \text{ has-real-derivative } D)$  F  $\longleftrightarrow$  ( $(\text{ccdf } M \text{ has-real-derivative } -D)$  F)
proof
  assume ( $(\text{cdf } M \text{ has-real-derivative } D)$  F
  thus ( $(\text{ccdf } M \text{ has-real-derivative } -D)$  F
    using ccdf-cdf DERIV-const Deriv.field-differentiable-diff by fastforce
next
  assume ( $(\text{ccdf } M \text{ has-real-derivative } -D)$  F
  thus ( $(\text{cdf } M \text{ has-real-derivative } D)$  F
    using cdf-ccdf DERIV-const Deriv.field-differentiable-diff by fastforce
qed

lemma differentiable-cdf-ccdf: ( $(\text{cdf } M \text{ differentiable } F)$   $\longleftrightarrow$  ( $(\text{ccdf } M \text{ differentiable } F)$ )
unfolding differentiable-def
apply (rewrite has-real-derivative-iff[THEN sym])+  

apply (rewrite has-real-derivative-cdf-ccdf)
by (metis verit-minus-simplify(4))

lemma deriv-cdf-ccdf:
assumes cdf M differentiable at x
shows deriv (cdf M) x = - deriv (ccdf M) x
using has-real-derivative-cdf-ccdf differentiable-cdf-ccdf assms
by (simp add: DERIV-deriv-iff-real-differentiable DERIV-imp-deriv)

lemma ccdf-diff-eq2:
assumes x ≤ y
shows ccdf M x - ccdf M y = measure M {x <.. y}
proof –
  have ccdf M x - ccdf M y = cdf M y - cdf M x using add-cdf-ccdf by (smt (verit))
  also have ... = measure M {x <.. y} using cdf-diff-eq2 assms by simp
  finally show ?thesis .
qed

lemma ccdf-nonincreasing: x ≤ y  $\implies$  ccdf M x ≥ ccdf M y
using add-cdf-ccdf cdf-nondecreasing by (smt (verit))

```

```

lemma ccdf-nonneg:  $\text{ccdf } M \ x \geq 0$ 
  using add-cdf-ccdf cdf-bounded by (smt (verit))

lemma ccdf-bounded:  $\text{ccdf } M \ x \leq \text{measure } M \ (\text{space } M)$ 
  using add-cdf-ccdf cdf-nonneg by (smt (verit))

lemma ccdf-lim-at-top:  $(\text{ccdf } M \longrightarrow 0) \text{ at-top}$ 
proof -
  have  $((\lambda x. \text{measure } M \ (\text{space } M) - \text{cdf } M \ x) \longrightarrow \text{measure } M \ (\text{space } M) - \text{measure } M \ (\text{space } M)) \text{ at-top}$ 
    apply (intro tendsto-intros)
    by (rule cdf-lim-at-top)
    thus ?thesis
      by (rewrite ccdf-cdf) simp
  qed

lemma ccdf-lim-at-bot:  $(\text{ccdf } M \longrightarrow \text{measure } M \ (\text{space } M)) \text{ at-bot}$ 
proof -
  have  $((\lambda x. \text{measure } M \ (\text{space } M) - \text{cdf } M \ x) \longrightarrow \text{measure } M \ (\text{space } M) - 0) \text{ at-bot}$ 
    apply (intro tendsto-intros)
    by (rule cdf-lim-at-bot)
    thus ?thesis
      by (rewrite ccdf-cdf) simp
  qed

lemma ccdf-is-right-cont: continuous (at-right a) (ccdf M)
proof -
  have continuous (at-right a) ( $\lambda x. \text{measure } M \ (\text{space } M) - \text{cdf } M \ x$ )
    apply (intro continuous-intros)
    by (rule cdf-is-right-cont)
    thus ?thesis by (rewrite ccdf-cdf) simp
  qed

end

context prob-space
begin

lemma ccdf-distr-measurable [measurable]:
  assumes [measurable]: random-variable borel X
  shows ccdf (distr M borel X) ∈ borel-measurable borel
  using real-distribution.finite-borel-measure-M by (rewrite finite-borel-measure.ccdf-cdf; simp)

lemma ccdf-distr-P:
  assumes random-variable borel X
  shows ccdf (distr M borel X) x = P(ω in M. X ω > x)

```

```

unfolding ccdf-def apply (rewrite measure-distr; (simp add: assms) ?)
unfolding vimage-def by (rule arg-cong[where f=prob]) force

lemma ccdf-continuous-distr-P-ge:
assumes random-variable borel X isCont (ccdf (distr M borel X)) x
shows ccdf (distr M borel X) x = P(ω in M. X ω ≥ x)
proof -
have ccdf (distr M borel X) x = measure (distr M borel X) {x<..} unfolding
ccdf-def by simp
also have ... = measure (distr M borel X) ({x<..} ∪ {x})
apply (rewrite finite-measure.measure-zero-union, simp-all add: assms finite-measure-distr)
using finite-borel-measure.isCont-ccdf real-distribution.finite-borel-measure-M
assms by blast
also have ... = measure (distr M borel X) {x..} by (metis Un-commute ivl-disj-un-singleton(1))
also have ... = P(ω in M. X ω ≥ x)
apply (rewrite measure-distr, simp-all add: assms)
unfolding vimage-def by simp (smt (verit) Collect-cong Int-def mem-Collect-eq)
finally show ?thesis .
qed

lemma ccdf-distr-diff-P:
assumes x ≤ y
and random-variable borel X
shows ccdf (distr M borel X) x - ccdf (distr M borel X) y = P(ω in M. x < X
ω ∧ X ω ≤ y)
proof -
interpret distrX-FBM: finite-borel-measure distr M borel X
using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
have ccdf (distr M borel X) x - ccdf (distr M borel X) y = measure (distr M
borel X) {x<..y}
by (rewrite distrX-FBM.ccdf-diff-eq2; simp add: assms)
also have ... = P(ω in M. x < X ω ∧ X ω ≤ y)
apply (rewrite measure-distr; (simp add: assms) ?)
unfolding vimage-def by (rule arg-cong[where f=prob], force)
finally show ?thesis .
qed

end

context real-distribution
begin

lemma ccdf-bounded-prob: ∀x. ccdf M x ≤ 1
by (subst prob-space[THEN sym], rule ccdf-bounded)

lemma ccdf-lim-at-bot-prob: (ccdf M —→ 1) at-bot
by (subst prob-space[THEN sym], rule ccdf-lim-at-bot)

```

end

Introduce the hazard rate. This notion will be used to define the force of mortality.

3.4 Hazard Rate

context *prob-space*
begin

definition *hazard-rate* :: $('a \Rightarrow \text{real}) \Rightarrow \text{real} \Rightarrow \text{real}$
where *hazard-rate* $X t \equiv$
 $\text{Lim} (\text{at-right } 0) (\lambda dt. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) / dt)$
— Here, X is supposed to be a random variable.

lemma *hazard-rate-0-ccdf-0*:
assumes [*measurable*]: *random-variable borel X*
and *ccdf (distr M borel X) t = 0*
shows *hazard-rate X t = 0*
— Note that division by 0 is calculated as 0 in Isabelle/HOL.
proof —
have $\bigwedge dt. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) = 0$
unfolding *cond-prob-def* **using** *ccdf-distr-P assms by simp*
hence *hazard-rate X t = Lim (at-right 0) (λdt::real. 0)*
unfolding *hazard-rate-def* **by** (*rewrite Lim-cong; simp*)
also have ... = 0 **by** (*rule tendsto-Lim; simp*)
finally show ?thesis .
qed

lemma *hazard-rate-deriv-cdf*:
assumes [*measurable*]: *random-variable borel X*
and (*cdf (distr M borel X)*) *differentiable at t*
shows *hazard-rate X t = deriv (cdf (distr M borel X)) t / ccdf (distr M borel X) t*
proof (*cases <ccdf (distr M borel X) t = 0*)
case *True*
with *hazard-rate-0-ccdf-0* **show** ?thesis **by simp**
next
case *False*
let $?F = \text{cdf} (\text{distr } M \text{ borel } X)$
have $\forall_F dt \text{ in at-right } 0. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) / dt =$
 $(?F(t + dt) - ?F t) / dt / \text{ccdf} (\text{distr } M \text{ borel } X) t$
apply (*rule eventually-at-rightI[where b=1]; simp*)
unfolding *cond-prob-def*
apply (*rewrite cdf-distr-diff-P; simp*)
apply (*rewrite ccdf-distr-P[THEN sym]; simp*)
by (*smt (verit) Collect-cong mult.commute*)
moreover have $((\lambda dt. (?F(t + dt) - ?F t) / dt) / \text{ccdf} (\text{distr } M \text{ borel } X) t$

```

deriv ?F t / ccdf (distr M borel X) t) (at-right 0)
apply (rule tendsto-intros, simp-all add: False)
apply (rule Lim-at-imp-Lim-at-within)
using DERIV-deriv-iff-real-differentiable assms DERIV-def by blast
ultimately show ?thesis
  unfolding hazard-rate-def using tendsto-cong by (intro tendsto-Lim; force)
qed

lemma deriv-cdf-hazard-rate:
assumes [measurable]: random-variable borel X
  and (cdf (distr M borel X)) differentiable at t
shows deriv (cdf (distr M borel X)) t = ccdf (distr M borel X) t * hazard-rate
X t
proof -
  interpret distrX-FBM: finite-borel-measure distr M borel X
    using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
  show ?thesis
  proof (cases `ccdf (distr M borel X) t = 0`)
    case True
    hence cdf (distr M borel X) t = 1
      using distrX-FBM.cdf-ccdf
      by simp (metis assms(1) distrX-FBM.borel-UNIV prob-space.prob-space
prob-space-distr)
    moreover obtain D where (cdf (distr M borel X) has-real-derivative D) (at
t)
      using assms real-differentiable-def by atomize-elim blast
    ultimately have (cdf (distr M borel X) has-real-derivative 0) (at t)
      using assms
      by (smt (verit) DERIV-local-max real-distribution.cdf-bounded-prob real-distribution-distr)
      thus ?thesis using True by (simp add: DERIV-imp-deriv)
    next
    case False
    thus ?thesis using hazard-rate-deriv-cdf assms by simp
  qed
qed

lemma hazard-rate-deriv-ccdf:
assumes [measurable]: random-variable borel X
  and (ccdf (distr M borel X)) differentiable at t
shows hazard-rate X t = - deriv (ccdf (distr M borel X)) t / ccdf (distr M borel
X) t
proof -
  interpret distrX-FBM: finite-borel-measure distr M borel X
    using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
  show ?thesis
    using hazard-rate-deriv-cdf distrX-FBM.deriv-cdf-ccdf assms distrX-FBM.differentiable-cdf-ccdf
    by presburger

```

qed

```
lemma deriv-ccdf-hazard-rate:
  assumes [measurable]: random-variable borel X
    and (ccdf (distr M borel X)) differentiable at t
  shows deriv (ccdf (distr M borel X)) t = - ccdf (distr M borel X) t * hazard-rate
X t
proof -
  interpret distrX-FBM: finite-borel-measure distr M borel X
    using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
  show ?thesis
    using deriv-cdf-hazard-rate distrX-FBM.deriv-cdf-ccdf assms distrX-FBM.differentiable-cdf-ccdf
      by simp
qed

lemma hazard-rate-deriv-ln-ccdf:
  assumes [measurable]: random-variable borel X
    and (ccdf (distr M borel X)) differentiable at t
    and ccdf (distr M borel X) t ≠ 0
  shows hazard-rate X t = - deriv (λt. ln (ccdf (distr M borel X) t)) t
proof -
  interpret distrX-FBM: finite-borel-measure distr M borel X
    using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
  let ?srvl = ccdf (distr M borel X)
  have ?srvl t > 0 using distrX-FBM.ccdf-nonneg assms by (smt (verit))
  moreover have (?srvl has-real-derivative (deriv ?srvl t)) (at t)
    using DERIV-deriv-iff-real-differentiable assms by blast
  ultimately have ((λt. ln (?srvl t)) has-real-derivative 1 / ?srvl t * deriv ?srvl
t) (at t)
    by (rule derivative-intros)
  hence deriv (λt. ln (?srvl t)) t = deriv ?srvl t / ?srvl t by (simp add: DE-
RIV-imp-deriv)
  also have ... = - hazard-rate X t using hazard-rate-deriv-ccdf assms by simp
  finally show ?thesis by simp
qed

lemma hazard-rate-has-integral:
  assumes [measurable]: random-variable borel X
    and t ≤ u
    and (ccdf (distr M borel X)) piecewise-differentiable-on {t<..<u}
    and continuous-on {t..u} (ccdf (distr M borel X))
    and ∀s. s ∈ {t..u} ==> ccdf (distr M borel X) s ≠ 0
  shows
    (hazard-rate X has-integral ln (ccdf (distr M borel X) t / ccdf (distr M borel X)
u)) {t..u}
proof -
  interpret distrX-FBM: finite-borel-measure distr M borel X
```

```

using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
let ?srvl = ccdf (distr M borel X)
have [simp]:  $\bigwedge s. t \leq s \wedge s \leq u \implies ?srvl s > 0$ 
  using distrX-FBM.ccdf-nonneg assms by (smt (verit) atLeastAtMost-iff)
have (deriv ( $\lambda s. -\ln (?srvl s)$ ) has-integral  $= -\ln (?srvl u) - -\ln (?srvl t)$ )
{t..u}
proof -
  have continuous-on {t..u} ( $\lambda s. -\ln (?srvl s)$ )
    by (rule continuous-intros, rule continuous-on-ln, auto simp add: assms)
  moreover hence ( $\lambda s. -\ln (?srvl s)$ ) piecewise-differentiable-on {t<..<u}
  proof -
    have ?srvl ` {t<..<u}  $\subseteq \{0 <..\}$ 
    proof -
      { fix s assume s ∈ {t<..<u}
        hence ?srvl s ≠ 0 using assms by simp
        moreover have ?srvl s ≥ 0 using distrX-FBM.ccdf-nonneg by simp
        ultimately have ?srvl s > 0 by simp }
      thus ?thesis by auto
    qed
    hence ( $\lambda r. -\ln r$ ) o ?srvl piecewise-differentiable-on {t<..<u}
      apply (intro differentiable-on-piecewise-compose, simp add: assms)
      apply (rule derivative-intros)
      apply (rule differentiable-on-subset[of ln {0 <..}], simp-all)
      apply (rewrite differentiable-on-eq-field-differentiable-real, auto)
      unfolding field-differentiable-def using DERIV-ln by (metis has-field-derivative-at-within)
      thus ?thesis unfolding comp-def by simp
    qed
    ultimately show ?thesis by (intro FTC-real-deriv-has-integral; simp add:
assms)
  qed
  hence ln: (deriv ( $\lambda s. -\ln (?srvl s)$ ) has-integral ln (?srvl t / ?srvl u)) {t..u}
    by simp (rewrite ln-div; force simp: assms)
  thus ((hazard-rate X) has-integral ln (?srvl t / ?srvl u)) {t..u}
  proof -
    from assms obtain S0 where finS0: finite S0 and
      diffS0:  $\bigwedge s. s \in \{t <.. < u\} - S0 \implies ?srvl$  differentiable at s within {t<..<u}
      unfolding piecewise-differentiable-on-def by blast
    from this obtain S where finite S and  $\bigwedge s. s \in \{t..u\} - S \implies ?srvl$  differentiable at s
    proof (atomize-elim)
      let ?S = S0 ∪ {t, u}
      have finite ?S using finS0 by simp
      moreover have  $\forall s. s \in \{t..u\} - ?S \longrightarrow$  ccdf (distr M borel X) differentiable
      at s
      proof -
        { fix s assume s-in: s ∈ {t..u} - ?S
          hence ?srvl differentiable at s within {t <.. < u} using diffS0 by simp
          hence ?srvl differentiable at s
        }
    qed
  qed

```

```

    using s-in by (rewrite at-within-open[THEN sym], simp-all) }
    thus ?thesis by blast
qed
ultimately show
   $\exists S. \text{finite } S \wedge (\forall s. s \in \{t..u\} - S \longrightarrow \text{ccdf}(\text{distr } M \text{ borel } X) \text{ differentiable}$ 
at s)
  by blast
qed
thus ?thesis
apply (rewrite has-integral-spike-finite-eq [of S - deriv (\lambda s. - ln (?srvl s))], simp)
apply (rewrite hazard-rate-deriv-ln-ccdf, simp-all add: assms)
apply (rewrite deriv-minus, simp-all)
apply (rewrite in asm differentiable-eq-field-differentiable-real)
apply (rewrite comp-def[THEN sym], rule field-differentiable-compose[of ?srvl], simp-all)
unfolding field-differentiable-def apply (rule exI, rule DERIV-ln, simp)
using ln by simp
qed
qed

corollary hazard-rate-integrable:
assumes [measurable]: random-variable borel X
and t ≤ u
and (ccdf (distr M borel X)) piecewise-differentiable-on {t < .. < u}
and continuous-on {t..u} (ccdf (distr M borel X))
and  $\bigwedge s. s \in \{t..u\} \implies \text{ccdf}(\text{distr } M \text{ borel } X) s \neq 0$ 
shows hazard-rate X integrable-on {t..u}
using has-integral-integrable hazard-rate-has-integral assms by blast

lemma ccdf-exp-cumulative-hazard:
assumes [measurable]: random-variable borel X
and t ≤ u
and (ccdf (distr M borel X)) piecewise-differentiable-on {t < .. < u}
and continuous-on {t..u} (ccdf (distr M borel X))
and  $\bigwedge s. s \in \{t..u\} \implies \text{ccdf}(\text{distr } M \text{ borel } X) s \neq 0$ 
shows ccdf (distr M borel X) u / ccdf (distr M borel X) t =
exp (- integral {t..u} (hazard-rate X))
proof -
interpret distrX-FBM: finite-borel-measure distr M borel X
  using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
let ?srvl = ccdf (distr M borel X)
have [simp]:  $\bigwedge s. t \leq s \wedge s \leq u \implies ?\text{srvl } s > 0$ 
  using distrX-FBM.ccdf-nonneg assms by (smt (verit) atLeastAtMost-iff)
have integral {t..u} (hazard-rate X) = ln (?srvl t / ?srvl u)
  using hazard-rate-has-integral has-integral-integrable-integral assms by auto
also have ... = - ln (?srvl u / ?srvl t) using ln-div assms by simp
finally have - integral {t..u} (hazard-rate X) = ln (?srvl u / ?srvl t) by simp

```

```

thus ?thesis using assms by simp
qed

lemma hazard-rate-density-ccdf:
assumes distributed M lborel X f
  and  $\bigwedge s. f s \geq 0$   $t < u$  continuous-on  $\{t..u\} f$ 
shows hazard-rate X t =  $f t / \text{ccdf} (\text{distr } M \text{ borel } X) t$ 
proof (cases `ccdf (distr M borel X) t = 0`)
  case True
  thus ?thesis
    apply (rewrite hazard-rate-0-ccdf-0, simp-all)
    using assms(1) distributed-measurable by force
next
  case False
  have [simp]:  $t \leq u$  using assms by simp
  have [measurable]: random-variable borel X
    using assms distributed-measurable measurable-lborel1 by blast
  have [measurable]:  $f \in \text{borel-measurable lborel}$ 
    using assms distributed-real-measurable by metis
  have [simp]: integrable lborel f
  proof -
    have prob ( $X -` \text{UNIV} \cap \text{space } M$ ) = LINT x|lborel. indicat-real UNIV x * f x
      by (rule distributed-measure; simp add: assms)
    thus ?thesis
      using prob-space not-integrable-integral-eq by fastforce
  qed
  have  $((\lambda dt. (\text{LBINT } s:\{t..t+dt\}. f s) / dt) \longrightarrow f t)$  (at-right 0)
  proof -
    have  $\bigwedge dt. (\int^+ x. \text{ennreal} (\text{indicat-real } \{t..t+dt\} x * f x) \partial \text{lborel}) < \infty$ 
    proof -
      fix dt :: real
      have  $(\int^+ x. \text{ennreal} (\text{indicat-real } \{t..t+dt\} x * f x) \partial \text{lborel}) =$ 
        set-nn-integral lborel  $\{t..t+dt\} f$ 
        by (metis indicator-mult-ennreal mult.commute)
      moreover have emeasure M ( $X -` \{t..t+dt\} \cap \text{space } M$ ) = set-nn-integral
        lborel  $\{t..t+dt\} f$ 
        by (rule distributed-emeasure; simp add: assms)
      moreover have emeasure M ( $X -` \{t..t+dt\} \cap \text{space } M$ ) <  $\infty$ 
        using emeasure-eq-measure ennreal-less-top infinity-ennreal-def by presburger
        ultimately show  $(\int^+ x. \text{ennreal} (\text{indicat-real } \{t..t+dt\} x * f x) \partial \text{lborel}) <$ 
           $\infty$  by simp
    qed
    hence  $\bigwedge dt. (\text{LBINT } s:\{t..t+dt\}. f s) = \text{integral } \{t..t+dt\} f$ 
      apply (intro set-borel-integral-eq-integral)
      unfolding set-integrable-def
      apply (rule integrableI-nonneg; simp)
      by (rule AE-I2, simp add: assms)
    moreover have  $((\lambda dt. (\text{integral } \{t..t+dt\} f) / dt) \longrightarrow f t)$  (at-right 0)
    proof -

```

```

have (( $\lambda x. \text{integral } \{t..x\} f$ ) has-real-derivative  $f t$ ) (at  $t$  within  $\{t..u\}$ )
  by (rule integral-has-real-derivative; simp add: assms)
  moreover have (at  $t$  within  $\{t..u\}$ ) = (at  $(t+0)$  within  $(+)t` \{0..u-t\}$ ) by
    simp
    ultimately have
      (( $\lambda x. \text{integral } \{t..x\} f$ )  $\circ$   $(+)t$  has-real-derivative  $f t$ ) (at  $0$  within  $\{0..u-t\}$ )
      by (metis DERIV-at-within-shift-lemma)
      hence (( $\lambda dt. (\text{integral } \{t..t+dt\} f) / dt$ ) —→  $f t$ ) (at  $0$  within  $\{0..u-t\}$ )
        using has-field-derivative-iff by force
      thus ?thesis using at-within-Icc-at-right assms by (smt (verit))
    qed
    ultimately show ?thesis by simp
  qed
  moreover have  $\bigwedge dt. dt > 0 \implies \mathcal{P}(x \text{ in } M. X x \in \{t <.. t+dt\}) = (\text{LBINT}_{s:\{t..t+dt\}. f s})$ 
  proof -
    fix  $dt :: \text{real}$  assume  $dt > 0$ 
    hence [simp]:  $\text{sym-diff } \{t <.. t + dt\} \{t .. t + dt\} = \{t\}$  by force
    have prob ( $X -` \{t <.. t+dt\} \cap \text{space } M$ ) =  $\int s. \text{indicator } \{t <.. t+dt\} s * f s$ 
     $\partial\text{borel}$ 
    by (rule distributed-measure; simp add: assms)
    hence  $\mathcal{P}(x \text{ in } M. X x \in \{t <.. t+dt\}) = (\text{LBINT}_{s:\{t..t+dt\}. f s})$ 
    unfolding set-lebesgue-integral-def vimage-def Int-def by simp (smt (verit)
    Collect-cong)
    moreover have  $(\text{LBINT}_{s:\{t..t+dt\}. f s}) = (\text{LBINT}_{s:\{t..t+dt\}. f s})$ 
    by (rule set-integral-null-delta; force)
    ultimately show  $\mathcal{P}(x \text{ in } M. X x \in \{t <.. t+dt\}) = (\text{LBINT}_{s:\{t..t+dt\}. f s})$ 
  by simp
  qed
  ultimately have  $((\lambda dt. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt) / dt) —\rightarrow f t)$ 
  (at-right 0)
  by simp (smt (verit) Lim-cong-within greaterThan-iff)
  hence  $((\lambda dt. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) / dt) —\rightarrow$ 
     $f t / \text{ccdf } (\text{distr } M \text{ borel } X) t)$  (at-right 0)
  unfolding cond-prob-def
  apply (rewrite ccdf-distr-P[THEN sym]; simp)
  apply (rewrite mult.commute, rewrite divide-divide-eq-left[THEN sym])
  by (rule tendsto-intros; (simp add: False)?) (smt (verit) Collect-cong Lim-cong-within)
  thus ?thesis unfolding hazard-rate-def by (intro tendsto-Lim; simp)
qed

end

end
theory Interest
  imports Preliminaries
begin

```

4 Theory of Interest

```

locale interest =
  fixes i :: real — i stands for an interest rate.
  assumes v-futr-pos: 1 + i > 0 — Assume that the future value is positive.
  begin

    definition i-nom :: nat ⇒ real (⟨$i{-}⟩ [0] 200)
    where $i{-}m} ≡ m * ((1+i).⟨(1/m) - 1⟩) — nominal interest rate

    definition i-force :: real (⟨$δ⟩ 200)
    where $δ ≡ ln (1+i) — force of interest

    definition d-nom :: nat ⇒ real (⟨$d{-}⟩ [0] 200)
    where $d{-}m} ≡ $i{-}m} / (1 + $i{-}m}/m) — discount rate

    abbreviation d-nom-yr :: real (⟨$d⟩ 200)
    where $d ≡ $d{-}1} — Post-fix yr stands for "year".

    definition v-pres :: real (⟨$v⟩ 200)
    where $v ≡ 1 / (1+i) — present value factor

    definition ann :: nat ⇒ nat ⇒ real (⟨$a{-}⟩'--> [0,101] 200)
    where $a{-}m}-n ≡ ∑ k< n*m. $v.⟨((k+1::nat)/m) / m
    — present value of an immediate annuity

    abbreviation ann-yr :: nat ⇒ real (⟨$a'--> [101] 200)
    where $a{-}n ≡ $a{-}1}-n

    definition acc :: nat ⇒ nat ⇒ real (⟨$s{-}⟩'--> [0,101] 200)
    where $s{-}m}-n ≡ ∑ k< n*m. (1+i).⟨((k::nat)/m) / m
    — future value of an immediate annuity
    — The name acc stands for "accumulation".

    abbreviation acc-yr :: nat ⇒ real (⟨$s'--> 200)
    where $s{-}n ≡ $s{-}1}-n

    definition ann-due :: nat ⇒ nat ⇒ real (⟨$a''{-}⟩'--> [0,101] 200)
    where $a''{-}m}-n ≡ ∑ k< n*m. $v.⟨((k::nat)/m) / m
    — present value of an annuity-due

    abbreviation ann-due-yr :: nat ⇒ real (⟨$a''{-}⟩'--> [101] 200)
    where $a''{-}n ≡ $a''{-}1}-n

    definition acc-due :: nat ⇒ nat ⇒ real (⟨$s''{-}⟩'--> [0,101] 200)
    where $s''{-}m}-n ≡ ∑ k< n*m. (1+i).⟨((k+1::nat)/m) / m
    — future value of an annuity-due

    abbreviation acc-due-yr :: nat ⇒ real (⟨$s''{-}⟩'--> [101] 200)

```

where $\$s''\text{-}n \equiv \$s''\hat{\wedge}\{1\}\text{-}n$

definition $ann\text{-}cont :: real \Rightarrow real (\langle\$a'''\text{-} \rangle [101] 200)$
where $\$a'\text{-}n \equiv integral \{0..n\} (\lambda t::real. \$v.\hat{t})$
— present value of a continuous annuity

definition $acc\text{-}cont :: real \Rightarrow real (\langle\$s'''\text{-} \rangle [101] 200)$
where $\$s'\text{-}n \equiv integral \{0..n\} (\lambda t::real. (1+i).\hat{t})$
— future value of a continuous annuity

definition $perp :: nat \Rightarrow real (\langle\$a\hat{\wedge}\{-\}'\text{-}\infty \rangle [0] 200)$
where $\$a\hat{\wedge}\{m\}\text{-}\infty \equiv 1 / \$i\hat{\wedge}\{m\}$
— present value of a perpetual annuity

abbreviation $perp\text{-}yr :: real (\langle\$a'\text{-}\infty \rangle 200)$
where $\$a\text{-}\infty \equiv \$a\hat{\wedge}\{1\}\text{-}\infty$

definition $perp\text{-}due :: nat \Rightarrow real (\langle\$a''''\hat{\wedge}\{-\}'\text{-}\infty \rangle [0] 200)$
where $\$a''\hat{\wedge}\{m\}\text{-}\infty \equiv 1 / \$d\hat{\wedge}\{m\}$
— present value of a perpetual annuity-due

abbreviation $perp\text{-}due\text{-}yr :: real (\langle\$a''''\text{-}\infty \rangle 200)$
where $\$a''\text{-}\infty \equiv \$a''\hat{\wedge}\{1\}\text{-}\infty$

definition $ann\text{-}incr :: nat \Rightarrow nat \Rightarrow nat \Rightarrow real (\langle\$'(I\hat{\wedge}\{-\}a)\hat{\wedge}\{-\}'\text{-} \rangle [0,0,101] 200)$
where $\$(I\hat{\wedge}\{l\}a)\hat{\wedge}\{m\}\text{-}n \equiv \sum k < n*m. \$v.\hat{((k+1::nat)/m)} * \lceil l*(k+1::nat)/m \rceil$
 $/ (l*m)$
— present value of an increasing annuity
— This is my original definition.
— Here, l represents the number of increments per unit time.

abbreviation $ann\text{-}incr\text{-}lvl :: nat \Rightarrow nat \Rightarrow real (\langle\$'(Ia')\hat{\wedge}\{-\}'\text{-} \rangle [0,101] 200)$
where $\$(Ia)\hat{\wedge}\{m\}\text{-}n \equiv \$(I\hat{\wedge}\{1\}a)\hat{\wedge}\{m\}\text{-}n$
— The post-fix *lvl* stands for "level".

abbreviation $ann\text{-}incr\text{-}yr :: nat \Rightarrow real (\langle\$'(Ia')'\text{-} \rangle [101] 200)$
where $\$(Ia)\text{-}n \equiv \$(Ia)\hat{\wedge}\{1\}\text{-}n$

definition $acc\text{-}incr :: nat \Rightarrow nat \Rightarrow nat \Rightarrow real (\langle\$'(I\hat{\wedge}\{-\}s')\hat{\wedge}\{-\}'\text{-} \rangle [0,0,101] 200)$
where $\$(I\hat{\wedge}\{l\}s)\hat{\wedge}\{m\}\text{-}n \equiv \sum k < n*m. (1+i).\hat{((n-(k+1::nat)/m) * \lceil l*(k+1::nat)/m \rceil)}$
 $/ (l*m)$
— future value of an increasing annuity

abbreviation $acc\text{-}incr\text{-}lvl :: nat \Rightarrow nat \Rightarrow real (\langle\$'(Is')\hat{\wedge}\{-\}'\text{-} \rangle [0,101] 200)$
where $\$(Is)\hat{\wedge}\{m\}\text{-}n \equiv \$(I\hat{\wedge}\{1\}s)\hat{\wedge}\{m\}\text{-}n$

abbreviation $acc\text{-}incr\text{-}yr :: nat \Rightarrow real (\langle\$'(Is')'\text{-} \rangle [101] 200)$
where $\$(Is)\text{-}n \equiv \$(Is)\hat{\wedge}\{1\}\text{-}n$

definition *ann-due-incr* :: *nat* \Rightarrow *nat* \Rightarrow *real* ($\langle \$'(I\hat{\{}}\{-\}a''')\hat{\{}}\{-\}' \rightarrow [0,0,101] 200 \rangle$)
where $\$(I\hat{\{}}\{l\}a'')\hat{\{}}\{m\}-n \equiv \sum k < n * m. \$v.\hat{\{}}((k::nat)/m) * \lceil l*(k+1::nat)/m \rceil / (l*m)$

abbreviation *ann-due-incr-lvl* :: *nat* \Rightarrow *nat* \Rightarrow *real* ($\langle \$'(Ia''')\hat{\{}}\{-\}' \rightarrow [0,101] 200 \rangle$)
where $\$(Ia'')\hat{\{}}\{m\}-n \equiv \$(I\hat{\{}}\{1\}a'')\hat{\{}}\{m\}-n$

abbreviation *ann-due-incr-yr* :: *nat* \Rightarrow *real* ($\langle \$'(Ia''')\hat{\{}}\{-\}' \rightarrow [101] 200 \rangle$)
where $\$(Ia'')\hat{\{}}\{1\}-n \equiv \$(Ia'')\hat{\{}}\{1\}-n$

definition *acc-due-incr* :: *nat* \Rightarrow *nat* \Rightarrow *real* ($\langle \$'(I\hat{\{}}\{-\}s''')\hat{\{}}\{-\}' \rightarrow [0,0,101] 200 \rangle$)
where $\$(I\hat{\{}}\{l\}s'')\hat{\{}}\{m\}-n \equiv \sum k < n * m. (1+i).\hat{\{}}(n-(k::nat)/m) * \lceil l*(k+1::nat)/m \rceil / (l*m)$

abbreviation *acc-due-incr-lvl* :: *nat* \Rightarrow *nat* \Rightarrow *real* ($\langle \$'(Is''')\hat{\{}}\{-\}' \rightarrow [0,101] 200 \rangle$)
where $\$(Is'')\hat{\{}}\{m\}-n \equiv \$(I\hat{\{}}\{1\}s'')\hat{\{}}\{m\}-n$

abbreviation *acc-due-incr-yr* :: *nat* \Rightarrow *real* ($\langle \$'(Is''')\hat{\{}}\{-\}' \rightarrow [101] 200 \rangle$)
where $\$(Is'')\hat{\{}}\{1\}-n \equiv \$(Is'')\hat{\{}}\{1\}-n$

definition *perp-incr* :: *nat* \Rightarrow *nat* \Rightarrow *real* ($\langle \$'(I\hat{\{}}\{-\}a')\hat{\{}}\{-\}' \rightarrow [0,0] 200 \rangle$)
where $\$(I\hat{\{}}\{l\}a)\hat{\{}}\{m\}-\infty \equiv \lim (\lambda n. \$I\hat{\{}}\{l\}a)\hat{\{}}\{m\}-n)$

abbreviation *perp-incr-lvl* :: *nat* \Rightarrow *real* ($\langle \$'(Ia')\hat{\{}}\{-\}' \rightarrow [0] 200 \rangle$)
where $\$(Ia)\hat{\{}}\{m\}-\infty \equiv \$(I\hat{\{}}\{1\}a)\hat{\{}}\{m\}-\infty$

abbreviation *perp-incr-yr* :: *real* ($\langle \$'(Ia')\hat{\{}}\{-\}' \rightarrow [0] 200 \rangle$)
where $\$(Ia)-\infty \equiv \$(Ia)\hat{\{}}\{1\}-\infty$

definition *perp-due-incr* :: *nat* \Rightarrow *nat* \Rightarrow *real* ($\langle \$'(I\hat{\{}}\{-\}a''')\hat{\{}}\{-\}' \rightarrow [0,0] 200 \rangle$)
where $\$(I\hat{\{}}\{l\}a'')\hat{\{}}\{m\}-\infty \equiv \lim (\lambda n. \$I\hat{\{}}\{l\}a'')\hat{\{}}\{m\}-n)$

abbreviation *perp-due-incr-lvl* :: *nat* \Rightarrow *real* ($\langle \$'(Ia''')\hat{\{}}\{-\}' \rightarrow [0] 200 \rangle$)
where $\$(Ia'')\hat{\{}}\{m\}-\infty \equiv \$(I\hat{\{}}\{1\}a'')\hat{\{}}\{m\}-\infty$

abbreviation *perp-due-incr-yr* :: *real* ($\langle \$'(Ia''')\hat{\{}}\{-\}' \rightarrow [0] 200 \rangle$)
where $\$(Ia'')\hat{\{}}\{1\}-\infty \equiv \$(Ia'')\hat{\{}}\{1\}-\infty$

lemma *v-futr-m-pos*: $1 + \$i\hat{\{}}\{m\}/m > 0$ if $m \neq 0$ **for** $m::nat$
using *v-futr-pos i-nom-def* **by** *force*

lemma *i-nom-1[simp]*: $\$i\hat{\{}}\{1\} = i$
using *v-futr-pos i-nom-def* **by** *force*

lemma *i-nom-eff*: $(1 + \$i\hat{\{}}\{m\}/m)\hat{\{}}\{m\} = 1 + i$ if $m \neq 0$ **for** $m::nat$

**unfolding $i\text{-nom-def}$ using $\text{less-imp-neq } v\text{-futr-pos}$ that
apply ($\text{simp, subst powr-realpow[THEN sym], simp}$)
by ($\text{subst powr-powr, simp}$)**

**lemma $i\text{-nom-i}: 1 + \$i\hat{\wedge}\{m\}/m = (1+i).\hat{\wedge}(1/m)$ if $m \neq 0$ for $m::nat$
unfolding $i\text{-nom-def}$ by (simp add: that)**

**lemma $i\text{-nom-0-iff-i-0}: \$i\hat{\wedge}\{m\} = 0 \longleftrightarrow i = 0$ if $m \neq 0$ for $m::nat$
proof**

**assume $\$i\hat{\wedge}\{m\} = 0$
hence $\star: (1+i).\hat{\wedge}(1/m) = (1+i).\hat{\wedge}0$
unfolding $i\text{-nom-def}$ using $v\text{-futr-pos}$ that by simp
show $i = 0$
proof (rule $ccontr$)
assume $i \neq 0$
hence $1/m = 0$ using $\text{powr-inj} \star v\text{-futr-pos}$ by (smt (verit))
thus False using that by simp**

qed

next

**assume $i = 0$
thus $\$i\hat{\wedge}\{m\} = 0$
unfolding $i\text{-nom-def}$ by simp**

qed

lemma $i\text{-nom-pos-iff-i-pos}: \$i\hat{\wedge}\{m\} > 0 \longleftrightarrow i > 0$ if $m \neq 0$ for $m::nat$

proof

**assume $\$i\hat{\wedge}\{m\} > 0$
hence $\star: (1+i).\hat{\wedge}(1/m) > 1.\hat{\wedge}(1/m)$
unfolding $i\text{-nom-def}$ using $v\text{-futr-pos}$ that by ($\text{simp add: zero-less-mult-iff}$)
thus $i > 0$
using $\text{powr-less-cancel2}[of 1/m 1 1+i]$ $v\text{-futr-pos}$ that by simp**

next

**assume $i > 0$
hence $(1+i).\hat{\wedge}(1/m) > 1.\hat{\wedge}(1/m)$
using powr-less-mono2 $v\text{-futr-pos}$ that by simp
thus $\$i\hat{\wedge}\{m\} > 0$
unfolding $i\text{-nom-def}$ using that by ($\text{simp add: zero-less-mult-iff}$)**

qed

lemma $e\text{-delta}: \exp(\$\delta) = 1 + i$

unfolding $i\text{-force-def}$ by ($\text{simp add: v-futr-pos}$)

lemma $\text{delta-0-iff-i-0}: \$\delta = 0 \longleftrightarrow i = 0$

proof

**assume $\$\delta = 0$
thus $i = 0$
using $e\text{-delta}$ by auto**

next

assume $i = 0$

```

thus $δ = 0
  unfolding i-force-def by simp
qed

lemma lim-i-nom: (λm. $i^{m}) —→ $δ
proof -
  let ?f = λh. ((1+i).^h - 1) / h
  have D1ipwr: DERIV (λh. (1+i).^h) 0 :> $δ
    unfolding i-force-def
    using has-real-derivative-powr2[OF v-futr-pos, where x=0] v-futr-pos by simp
    hence limf: (?f —→ $δ) (at 0)
      unfolding DERIV-def using v-futr-pos by auto
      hence (λm. $i^{Suc m}) —→ $δ
        unfolding i-nom-def using tends-to-at-iff-sequentially[of ?f $δ 0 ℝ, THEN
        iffD1]
          apply simp
          apply (drule-tac x=λm. 1 / Suc m in spec, simp, drule mp)
          subgoal using lim-1-over-n LIMSEQ-Suc by force
            by (simp add: o-def mult.commute)
          thus ?thesis
            by (simp add: LIMSEQ-imp-Suc)
qed

lemma d-nom-0-iff-i-0: $d^{m} = 0 ↔ i = 0 if m ≠ 0 for m::nat
proof -
  have $d^{m} = 0 ↔ $i^{m} = 0
    unfolding d-nom-def using v-futr-m-pos by (smt (verit) divide-eq-0-iff of-nat-0)
    thus ?thesis
      using i-nom-0-iff-i-0 that by auto
qed

lemma d-nom-pos-iff-i-pos: $d^{m} > 0 ↔ i > 0 if m ≠ 0 for m::nat
proof -
  have $d^{m} > 0 ↔ $i^{m} > 0
    unfolding d-nom-def using zero-less-divide-iff i-nom-pos-iff-i-pos v-futr-m-pos
    that by (smt (verit))
    thus ?thesis
      using i-nom-pos-iff-i-pos that by auto
qed

lemma d-nom-i-nom: 1 - $d^{m}/m = 1 / (1 + $i^{m}/m) if m ≠ 0 for m::nat
proof -
  have 1 - $d^{m}/m = 1 - ($i^{m}/m) / (1 + $i^{m}/m)
    by (simp add: d-nom-def)
  also have ... = 1 / (1 + $i^{m}/m)
    using v-futr-m-pos
    by (smt (verit, ccfv-SIG) add-divide-distrib that div-self)
  finally show ?thesis .
qed

```

```

lemma lim-d-nom: ( $\lambda m. \delta^{\wedge}\{m\}$ )  $\longrightarrow$   $\delta$ 
proof -
  have ( $\lambda m. i^{\wedge}\{m\}/m$ )  $\longrightarrow$  0
    using lim-i-nom tends-to-divide-0 tends-to-of-nat by blast
  hence ( $\lambda m. 1 + i^{\wedge}\{m\}/m$ )  $\longrightarrow$  1
    by (metis add.right-neutral tends-to-add-const-iff)
  thus ?thesis
  unfolding d-nom-def using lim-i-nom tends-to-divide div-by-1 by fastforce
qed

lemma v-pos:  $v > 0$ 
  unfolding v-pres-def using v-futr-pos by auto

lemma v-1-iff-i-0:  $v = 1 \longleftrightarrow i = 0$ 
proof
  assume  $v = 1$ 
  thus  $i = 0$ 
    unfolding v-pres-def by simp
next
  assume  $i = 0$ 
  thus  $v = 1$ 
    unfolding v-pres-def by simp
qed

lemma v-lt-1-iff-i-pos:  $v < 1 \longleftrightarrow i > 0$ 
proof
  assume  $v < 1$ 
  thus  $i > 0$ 
    unfolding v-pres-def by (simp add: v-futr-pos)
next
  assume  $i > 0$ 
  thus  $v < 1$ 
    unfolding v-pres-def by (simp add: v-futr-pos)
qed

lemma v-i-nom:  $v = (1 + i^{\wedge}\{m\}/m)^{\wedge}-m$  if  $m \neq 0$  for  $m::nat$ 
proof -
  have  $v = (1 + i)^{\wedge}-1$ 
  unfolding v-pres-def using v-futr-pos powr-real-def that by (simp add: powr-neg-one)
  also have ... =  $((1 + i^{\wedge}\{m\}/m)^{\wedge}m)^{\wedge}-1$ 
    using i-nom-eff that by presburger
  also have ... =  $(1 + i^{\wedge}\{m\}/m)^{\wedge}-m$ 
    using i-nom-eff i-nom-i powr-powr that by fastforce
  finally show ?thesis .
qed

lemma i-v:  $1 + i = v^{\wedge}-1$ 
  unfolding v-pres-def powr-real-def using v-futr-pos powr-neg-one by simp

```

```

lemma i-v-powr:  $(1 + i) \cdot \hat{a} = \$v \cdot \hat{-a}$  for  $a::real$   

by (subst i-v, subst powr-powr, simp)

lemma v-delta:  $\ln (\$v) = - \$\delta$   

unfolding i-force-def v-pres-def using v-futr-pos by (simp add: ln-div)

lemma is-derive-vpow: DERIV  $(\lambda t. \$v \cdot \hat{t}) t :> - \$\delta * \$v \cdot \hat{t}$   

using v-delta has-real-derivative-powr2 v-pos by (metis mult.commute)

lemma d-nom-v:  $\hat{d}\{m\} = m * (1 - \$v \cdot \hat{(1/m)})$  if  $m \neq 0$  for  $m::nat$   

proof –  

  have  $\hat{d}\{m\} = m * (1 - 1 / (1 + \hat{i}\{m\}/m))$   

    using d-nom-i-nom[THEN sym] that by force  

  also have ... =  $m * (1 - 1 / (1 + i) \cdot \hat{(1/m)})$   

    using i-nom-i that powr-minus-divide by simp  

  also have ... =  $m * (1 - \$v \cdot \hat{(1/m)})$   

    using v-pres-def v-futr-pos powr-divide by simp  

    finally show ?thesis .  

qed

lemma d-nom-i-nom-v:  $\hat{d}\{m\} = \hat{i}\{m\} * \$v \cdot \hat{(1/m)}$  if  $m \neq 0$  for  $m::nat$   

unfolding d-nom-def v-pres-def using i-nom-i powr-divide v-futr-pos that by auto

lemma a-calc:  $\hat{a}\{m\} \cdot n = (1 - \$v \cdot \hat{n}) / \hat{i}\{m\}$  if  $m \neq 0$   $i \neq 0$  for  $n m ::nat$   

proof –  

  have  $\bigwedge l::nat. l/m = (1/m) * l$  by simp  

  hence  $\star: \bigwedge l::nat. \$v \cdot \hat{(l/m)} = (\$v \cdot \hat{(1/m)}) \cdot \hat{l}$   

    using powr-powr powr-realpow v-pos by (metis powr-gt-zero)  

  hence  $\hat{a}\{m\} \cdot n = (\sum k < n * m. (\$v \cdot \hat{(1/m)}) \cdot \hat{(k+1::nat)}) / m$   

    unfolding ann-def by presburger  

  also have ... =  $\$v \cdot \hat{(1/m)} * (\sum k < n * m. (\$v \cdot \hat{(1/m)}) \cdot \hat{k}) / m$   

    by (simp, subst sum-divide-distrib[THEN sym], subst sum-distrib-left[THEN sym], simp)  

  also have ... =  $\$v \cdot \hat{(1/m)} * (((\$v \cdot \hat{(1/m)}) \cdot \hat{(n*m)} - 1) / (\$v \cdot \hat{(1/m)} - 1)) / m$   

    apply (subst geometric-sum[of \$v \cdot \hat{(1/m)} n*m]; simp?)  

    using powr-zero-eq-one[of \$v] v-pos v-1-iff-i-0 powr-inj that  

    by (smt (verit, del-insts) divide-eq-0-iff of-nat-eq-0-iff)  

  also have ... =  $((\$v \cdot \hat{(1/m)}) \cdot \hat{(n*m)} - 1) / (m * (\$v \cdot \hat{(1/m)} - 1) / \$v \cdot \hat{(1/m)})$   

    by (simp add: field-simps)  

  also have ... =  $(\$v \cdot \hat{n} - 1) / (m * (1 - 1 / \$v \cdot \hat{(1/m)}))$   

    apply (subst  $\star$ [of n*m::nat, THEN sym], simp only: of-nat-simps)  

    apply (subst nonzero-mult-mult-cancel-right[where 'a=real, of m n], simp add:  

    that)  

    apply (subst powr-realpow[OF v-pos])  

    apply (subst times-divide-eq-right[of -- \$v \cdot \hat{(1/m)}, THEN sym])  

    using v-pos by (subst diff-divide-distrib[of -- \$v \cdot \hat{(1/m)}], simp)

```

```

also have ... =  $(1 - \$v^{\wedge}n) / (m * (1 / \$v. \wedge(1/m) - 1))$ 
  using minus-divide-divide by (smt (verit) mult-minus-right)
also have ... =  $(1 - \$v^{\wedge}n) / \$i^{\wedge}\{m\}$ 
  unfolding i-nom-def v-pres-def using v-futr-pos powr-divide by auto
  finally show ?thesis .
qed

lemma a-calc-i-0:  $\$a^{\wedge}\{m\}-n = n$  if  $m \neq 0$   $i = 0$  for  $n m :: nat$ 
  unfolding ann-def v-pres-def using that by simp

lemma s-calc-i-0:  $\$s^{\wedge}\{m\}-n = n$  if  $m \neq 0$   $i = 0$  for  $n m :: nat$ 
  unfolding acc-def using that by simp

lemma s-a:  $\$s^{\wedge}\{m\}-n = (1+i)^{\wedge}n * \$a^{\wedge}\{m\}-n$  if  $m \neq 0$  for  $n m :: nat$ 
proof -
  have  $(1+i)^{\wedge}n * \$a^{\wedge}\{m\}-n = (\sum k < n*m. (1+i)^{\wedge}n * (\$v. \wedge((k+1::nat)/m) / m))$ 
    unfolding ann-def using sum-distrib-left by blast
  also have ... =  $(\sum k < n*m. (1+i). \wedge((n*m - Suc k)/m) / m)$ 
  proof -
    have  $\bigwedge k :: nat. k < n*m \implies (1+i)^{\wedge}n * (\$v. \wedge((k+1::nat)/m) / m) = (1+i). \wedge((n*m - Suc k)/m) / m$ 
    unfolding v-pres-def
    by (simp add: v-futr-pos inverse-powr diff-divide-eq-iff that flip: powr-realpow
      powr-add)
    thus ?thesis by (meson lessThan-iff sum.cong)
  qed
  also have ... =  $(\sum k < n*m. (1+i). \wedge(k/m) / m)$ 
  apply (subst atLeast0LessThan[THEN sym])+
  by (subst sum.atLeastLessThan-rev[THEN sym, of - n*m 0, simplified add-0-right],
    simp)
  also have ... =  $\$s^{\wedge}\{m\}-n$ 
  unfolding acc-def by simp
  finally show ?thesis ..
qed

lemma s-calc:  $\$s^{\wedge}\{m\}-n = ((1+i)^{\wedge}n - 1) / \$i^{\wedge}\{m\}$  if  $m \neq 0$   $i \neq 0$  for  $n m :: nat$ 
  using that v-futr-pos
  apply (subst s-a, simp, subst a-calc; simp?)
  apply (rule disjI2)
  apply (subst right-diff-distrib, simp)
  apply (rule left-right-inverse-power)
  unfolding v-pres-def by auto

lemma a''-a:  $\$a''^{\wedge}\{m\}-n = (1+i). \wedge(1/m) * \$a^{\wedge}\{m\}-n$  if  $m \neq 0$  for  $m :: nat$ 
  unfolding ann-def ann-due-def
  apply (subst sum-distrib-left, subst times-divide-eq-right, simp)
  by (subst i-v, subst powr-powr, subst powr-add[THEN sym], simp, subst add-divide-distrib,
    simp)

```

```

lemma a-a'': $a^{\wedge}\{m\}-n = $v. $\widehat{(1/m)}$  * $a'' $\widehat{\{m\}}$ -n if m ≠ 0 for m::nat
  unfolding ann-def ann-due-def
  apply (subst sum-distrib-left, subst times-divide-eq-right, simp)
  by (subst powr-add[THEN sym], subst add-divide-distrib, simp)

lemma a''-calc-i-0: $a'' $\widehat{\{m\}}$ -n = n if m ≠ 0 i = 0 for n m :: nat
  unfolding ann-due-def v-pres-def using that by simp

lemma s''-calc-i-0: $s'' $\widehat{\{m\}}$ -n = n if m ≠ 0 i = 0 for n m :: nat
  unfolding acc-due-def using that by simp

lemma a''-calc: $a'' $\widehat{\{m\}}$ -n = (1 - $v $\widehat{n}$ ) / $d $\widehat{\{m\}}$  if m ≠ 0 i ≠ 0 for n m :: nat
proof -
  have $a'' $\widehat{\{m\}}$ -n = (1+i). $\widehat{(1/m)}$  * ((1 - $v $\widehat{n}$ ) / $i $\widehat{\{m\}}$ )
    using a''-a a-calc times-divide-eq-right that by simp
  also have ... = (1 - $v $\widehat{n}$ ) / ($v. $\widehat{(1/m)}$  * $i $\widehat{\{m\}}$ )
    by (subst i-v, subst powr-powr, simp, subst powr-minus-divide, simp)
  also have ... = (1 - $v $\widehat{n}$ ) / $d $\widehat{\{m\}}$ 
    using d-nom-i-nom-v that by simp
  finally show ?thesis .
qed

lemma s''-s: $s'' $\widehat{\{m\}}$ -n = (1+i). $\widehat{(1/m)}$  * $s $\widehat{\{m\}}$ -n if m ≠ 0 for m::nat
  unfolding acc-def acc-due-def
  apply (subst sum-distrib-left, subst times-divide-eq-right)
  by (subst powr-add[THEN sym], simp, subst add-divide-distrib, simp)

lemma s-s'': $s $\widehat{\{m\}}$ -n = $v. $\widehat{(1/m)}$  * $s'' $\widehat{\{m\}}$ -n if m ≠ 0 for m::nat
  unfolding acc-def acc-due-def v-pres-def using v-futr-pos
  apply (subst sum-distrib-left, subst times-divide-eq-right, simp)
  by (subst inverse-powr, simp, subst powr-add[THEN sym], subst add-divide-distrib, simp)

lemma s''-calc: $s'' $\widehat{\{m\}}$ -n = ((1+i) $\widehat{n} - 1$ ) / $d $\widehat{\{m\}}$  if m ≠ 0 i ≠ 0 for n m :: nat
proof -
  have $s'' $\widehat{\{m\}}$ -n = (1+i). $\widehat{(1/m)}$  * ((1+i) $\widehat{n} - 1$ ) / $i $\widehat{\{m\}}$ 
    using s''-s s-calc times-divide-eq-right that by simp
  also have ... = ((1+i) $\widehat{n} - 1$ ) / ($v. $\widehat{(1/m)}$  * $i $\widehat{\{m\}}$ )
    by (subst i-v, subst powr-powr, simp, subst powr-minus-divide, simp)
  also have ... = ((1+i) $\widehat{n} - 1$ ) / $d $\widehat{\{m\}}$ 
    using d-nom-i-nom-v that by simp
  finally show ?thesis .
qed

lemma s''-a'': $s'' $\widehat{\{m\}}$ -n = (1+i) $\widehat{n}$  * $a'' $\widehat{\{m\}}$ -n if m ≠ 0 for m::nat
  using that s''-s a''-a s-a by simp

```

```

lemma a'-calc: $a'-n = (1 - $v.^n) / $δ if i ≠ 0 n ≥ 0 for n::real
  unfolding ann-cont-def
  apply (rule integral-unique)
  using has-integral-powr2-from-0[OF v-pos - that(2)] v-delta v-1-iff-i-0 that
  by (smt (verit) minus-divide-divide)

lemma a'-calc-i-0: $a'-n = n if i = 0 n ≥ 0 for n::real
  unfolding ann-cont-def
  apply (subst iffD2[OF v-1-iff-i-0], simp add: that)
  by (simp add: integral-cong that)

lemma s'-calc: $s'-n = ((1+i).^n - 1) / $δ if i ≠ 0 n ≥ 0 for n::real
  unfolding acc-cont-def
  apply (rule integral-unique)
  using has-integral-powr2-from-0[OF v-futr-pos - that(2)] i-force-def that
  by simp

lemma s'-calc-i-0: $s'-n = n if i = 0 n ≥ 0 for n::real
  unfolding acc-cont-def
  apply (subst ⟨i = 0⟩, simp)
  by (simp add: integral-cong that)

lemma s'-a': $s'-n = (1+i).^n * $a'-n if n ≥ 0 for n::real
proof -
  have (1+i).^n * $a'-n = integral {0..n} (λt. (1+i).^(n-t))
  unfolding ann-cont-def
  using integrable-on-powr2-from-0-general[of $v n] v-pos v-futr-pos that
  apply (subst integral-mult, simp)
  apply (rule integral-cong)
  unfolding v-pres-def using inverse-powr powr-add[THEN sym] by (smt (verit))
  also have ... = $s'-n
  unfolding acc-cont-def using v-futr-pos that
  apply (subst has-integral-interval-reverse[of 0 n, simplified, THEN integral-unique];
  simp?)
  by (rule continuous-on-powr; auto)
  finally show ?thesis ..
qed

lemma lim-m-a: (λm. $a^{m}-n) —→ $a'-n for n::nat
proof (rule LIMSEQ-imp-Suc)
  show (λm. $a^{m}-n) —→ $a'-n
  proof (cases i = 0)
    case True
    show ?thesis
    using a-calc-i-0 a'-calc-i-0 True by simp
  next
    case False
    show ?thesis
  qed

```

```

using False v-pos delta-0-iff-i-0
apply (subst a-calc; simp?)
apply (subst a'-calc; simp?)
apply (subst powr-realpow, simp)
apply (rule tends-to-divide; simp?)
by (rule LIMSEQ-Suc[OF lim-i-nom])
qed
qed

lemma lim-m-a'': ( $\lambda m. \$a'' \setminus \{m\} - n$ )  $\longrightarrow \$a' - n$  for  $n :: nat$ 
proof (rule LIMSEQ-imp-Suc)
show ( $\lambda m. \$a'' \setminus \{Suc m\} - n$ )  $\longrightarrow \$a' - n$ 
proof (cases i = 0)
case True
show ?thesis
using a''-calc-i-0 a'-calc-i-0 True by simp
next
case False
show ?thesis
using False v-pos delta-0-iff-i-0
apply (subst a''-calc; simp?)
apply (subst a'-calc; simp?)
apply (subst powr-realpow, simp)
apply (rule tends-to-divide; simp?)
by (rule LIMSEQ-Suc[OF lim-d-nom])
qed
qed

lemma lim-m-s: ( $\lambda m. \$s \setminus \{m\} - n$ )  $\longrightarrow \$s' - n$  for  $n :: nat$ 
proof (rule LIMSEQ-imp-Suc)
show ( $\lambda m. \$s \setminus \{Suc m\} - n$ )  $\longrightarrow \$s' - n$ 
proof (cases i = 0)
case True
show ?thesis
using s-calc-i-0 s'-calc-i-0 True by simp
next
case False
show ?thesis
using False v-futr-pos delta-0-iff-i-0
apply (subst s-calc; simp?)
apply (subst s'-calc; simp?)
apply (subst powr-realpow, simp)
apply (rule tends-to-divide; simp?)
by (rule LIMSEQ-Suc[OF lim-i-nom])
qed
qed

lemma lim-m-s'': ( $\lambda m. \$s'' \setminus \{m\} - n$ )  $\longrightarrow \$s' - n$  for  $n :: nat$ 
proof (rule LIMSEQ-imp-Suc)

```

```

show ( $\lambda m. \$s'' \wedge \{Suc m\} \cdot n$ ) —————  $\$s' \cdot n$ 
  proof (cases  $i = 0$ )
    case True
      show ?thesis
        using  $s''\text{-calc-}i\text{-}0 s'\text{-calc-}i\text{-}0$  True by simp
  next
    case False
      show ?thesis
        using False v-futr-pos delta-0-iff-i-0
        apply (subst  $s''\text{-calc}$ ; simp?)
        apply (subst  $s'\text{-calc}$ ; simp?)
        apply (subst powr-realpow; simp)
        apply (rule tendsto-divide; simp?)
        by (rule LIMSEQ-Suc[OF lim-d-nom])
  qed
qed

lemma lim-n-a: ( $\lambda n. \$a \wedge \{m\} \cdot n$ ) —————  $\$a \wedge \{m\} \cdot \infty$  if  $m \neq 0$   $i > 0$  for  $m::nat$ 
proof -
  have  $\$i \wedge \{m\} \neq 0$  using i-nom-pos-iff-i-pos that by (smt (verit))
  moreover have ( $\lambda n. \$v \wedge n$ ) ————— 0
    using LIMSEQ-realpow-zero[of $v] v-pos v-lt-1-iff-i-pos that by simp
  ultimately show ?thesis
    using that apply (subst a-calc; simp?)
    unfolding perp-def apply (rule tendsto-divide; simp?)
    using tendsto-diff[where a=1 and b=0] by auto
  qed

lemma lim-n-a'': ( $\lambda n. \$a'' \wedge \{m\} \cdot n$ ) —————  $\$a'' \wedge \{m\} \cdot \infty$  if  $m \neq 0$   $i > 0$  for  $m::nat$ 
proof -
  have  $\$d \wedge \{m\} \neq 0$  using d-nom-pos-iff-i-pos that by (smt (verit))
  moreover have ( $\lambda n. \$v \wedge n$ ) ————— 0
    using LIMSEQ-realpow-zero[of $v] v-pos v-lt-1-iff-i-pos that by simp
  ultimately show ?thesis
    using that apply (subst a''-calc; simp?)
    unfolding perp-due-def apply (rule tendsto-divide; simp?)
    using tendsto-diff[where a=1 and b=0] by auto
  qed

lemma Ilsm-Ilam:  $\$(I \wedge \{l\} s) \wedge \{m\} \cdot n = (1+i) \wedge n * \$(I \wedge \{l\} a) \wedge \{m\} \cdot n$ 
  if  $l \neq 0$   $m \neq 0$  for  $l n m :: nat$ 
  unfolding acc-incr-def ann-incr-def v-pres-def using v-futr-pos powr-realpow
  apply (subst inverse-powr; simp)
  apply (subst sum-distrib-left)
  by (subst minus-real-def, subst powr-add, subst times-divide-eq-right, subst mult.assoc,
       simp)

lemma Iam-calc:  $\$(I a) \wedge \{m\} \cdot n = (\sum j < n. (j+1)/m * (\sum k=j*m..<(j+1)*m. \$v. \wedge((k+1)/m)))$ 
  if  $m \neq 0$  for  $n m :: nat$ 

```

```

proof -
let ?I = {.. $n$ }
let ?A =  $\lambda j. \{j*m..<(j+1)*m\}$ 
let ?g =  $\lambda k. \$v. \lceil((k+1::nat)/m) * \lceil((k+1::nat)/m) / m\rceil$ 
have $(Ia) $\lceil\{m\}\cdot n = (\sum j < n. \sum k = j*m..<(j+1)*m. \$v. \lceil((k+1)/m) * \lceil((k+1)/m)\rceil / m)$ 
unfolding ann-incr-def using seq-part-multiple that
apply (simp only: mult-1)
by (subst sum.UNION-disjoint[of ?I ?A ?g, THEN sym]; simp)
also have ... =  $(\sum j < n. (j+1)/m * (\sum k = j*m..<(j+1)*m. \$v. \lceil((k+1)/m)))$ 
proof -
{ fix j k
  assume  $j*m \leq k \wedge k < (j+1)*m$ 
  hence  $j*m < k+1 \wedge k+1 \leq (j+1)*m$  by force
  hence  $j < (k+1)/m \wedge (k+1)/m \leq j+1$ 
  using pos-less-divide-eq pos-divide-le-eq of-nat-less-iff of-nat-le-iff that
  by (smt (verit) of-nat-le-0-iff of-nat-mult)
  hence  $\lceil((k+1)/m) = j+1$ 
  by (simp add: ceiling-unique) }
hence  $\bigwedge k. j*m \leq k \wedge k < (j+1)*m \implies \lceil((k+1)/m) = j+1$ 
by (metis (no-types) of-nat-1 of-nat-add)
with v-pos show ?thesis
apply -
apply (rule sum.cong, simp)
apply (subst sum-distrib-left, rule sum.cong; simp)
by (smt (verit, ccfv-SIG) of-int-1 of-int-diff of-int-of-nat-eq)
qed
finally show ?thesis .
qed

lemma Ism-calc: $(Is) $\lceil\{m\}\cdot n = (\sum j < n. (j+1)/m * (\sum k = j*m..<(j+1)*m. (1+i). \lceil(n-(k+1)/m)\rceil))$ 
if  $m \neq 0$  for  $n m :: nat$ 
using v-pos that
apply (subst Ilsm-Ilam; simp)
apply (subst Iam-calc[simplified]; simp?)
apply ((subst sum-distrib-left, rule sum.cong; simp))+
unfolding v-pres-def using v-futr-pos
apply (subst inverse-powr; simp)
apply (subst powr-realpow[THEN sym], simp)
by (subst powr-add[THEN sym]; simp)

lemma Imam-calc-aux: $(I $\lceil\{m\}a\lceil\{m\}\cdot n = (\sum k < n*m. \$v. \lceil((k+1)/m) * (k+1) / m^2)$ 
if  $m \neq 0$  for  $m :: nat$ 
unfolding ann-incr-def power-def
apply (rule sum.cong, simp)
apply (subst of-nat-mult)
using v-pos that
apply (subst nonzero-mult-div-cancel-left, simp)

```

by (subst ceiling-of-nat; simp)

lemma *Imam-calc*:

```

$(I\{m\}a)\widehat{\{m\}}_n = ($v.\widehat{(1/m)} * (1 - (n*m+1)*$v\widehat{n} + n*m*$v.\widehat{(n+1/m)})) / (m*(1-$v.\widehat{(1/m)}))^2
if i ≠ 0 m ≠ 0 for n m :: nat
proof -
  have ∗: $v.\widehat{(1/m)} > 0 using v-pos by force
  hence $(I\{m\}a)\widehat{\{m\}}_n = (∑ k < n*m. (k+1)*($v.\widehat{(1/m)})^(k+1)) / m^2
    using that
    apply (subst Imam-calc-aux, simp)
    apply (subst sum-divide-distrib[THEN sym], simp)
    apply (rule sum.cong; simp)
    using powr-realpow[THEN sym] powr-powr by (simp add: add-divide-distrib
    powr-add)
  also have ... = $v.\widehat{(1/m)} * (∑ k < n*m. (k+1)*($v.\widehat{(1/m)})^k) / m^2
    by (subst sum-distrib-left, simp add: that, rule sum.cong; simp)
  also have ... = $v.\widehat{(1/m)} *
    ((1 - (n*m+1)*($v.\widehat{(1/m)})^(n*m) + n*m*($v.\widehat{(1/m)})^(n*m+1)) / (1 -
    $v.\widehat{(1/m)})^2) / m^2
    using v-pos v-1-iff-i-0 that by (subst geometric-increasing-sum; simp?)
  also have ... = ($v.\widehat{(1/m)} * (1 - (n*m+1)*$v\widehat{n} + n*m*$v.\widehat{(n+1/m)})) / (m*(1-$v.\widehat{(1/m)}))^2
    using ∗
    apply (subst powr-realpow[of $v.\widehat{(1/m)}, THEN sym], simp) +
    apply (subst powr-powr) +
    apply (subst times-divide-eq-right[THEN sym], subst divide-divide-eq-left)
    apply (subst power-mult-distrib)
    using powr-eq-one-iff-gen v-pos v-1-iff-i-0 apply (simp add: field-simps)
    by ((subst powr-realpow, simp) +, simp)
  finally show ?thesis .
qed
```

lemma *Imam-calc-i-0*: \$(I\{m\}a)\widehat{\{m\}}_n = (n*m+1)*n / (2*m) if i = 0 m ≠ 0
for n m :: nat

proof -

```

  have $(I\{m\}a)\widehat{\{m\}}_n = (∑ k < n*m. $v.\widehat{(k+1)/m}) * (k+1) / m^2
    by (subst Imam-calc-aux, simp-all add: that)
  also have ... = (∑ k < n*m. k+1) / m^2
    apply (subst v-1-iff-i-0[THEN iffD2], simp-all add: that)
    by (subst sum-divide-distrib[THEN sym], simp)
  also have ... = (n*m*(n*m+1) div 2) / m^2
    apply (subst Suc-eq-plus1[THEN sym], subst sum-bounds-lt-plus1[of id, simplified])
    by (subst Sum-Icc-nat, simp)
  also have ... = (n*m+1)*n / (2*m)
    apply (subst real-of-nat-div, simp)
    using that by (subst power2-eq-square, simp add: field-simps)
  finally show ?thesis .
```

qed

lemma *Imsm-calc*:

$$\begin{aligned} & \$\{I^{\{m\}}s\}^{\{m\}-n} = ((1+i). \widehat{(n+1/m)} - (n*m+1)*(1+i). \widehat{(1/m)} + n*m) / \\ & (m*((1+i). \widehat{(1/m)} - 1))^{\widehat{2}} \\ & \text{if } i \neq 0 \text{ } m \neq 0 \text{ for } n \text{ } m :: \text{nat} \\ \text{proof } - & \text{have } \$\{I^{\{m\}}a\}^{\{m\}-n} = \\ & (\$v^{\widehat{n}} * ((1+i). \widehat{(n+1/m)} - (n*m+1)*(1+i). \widehat{(1/m)} + n*m)) / (m*((1+i). \widehat{(1/m)} - 1))^{\widehat{2}} \\ \text{proof } - & \text{have } \$\{I^{\{m\}}a\}^{\{m\}-n} = \\ & (\$v. \widehat{(1/m)} * (1 - (n*m+1)*\$v^{\widehat{n}} + n*m*\$v. \widehat{(n+1/m)})) / (m*(1 - \$v. \widehat{(1/m)}))^{\widehat{2}} \\ & \text{using that by (subst Imam-calc; simp)} \\ \text{also have } \dots & = (1 - (n*m+1)*\$v^{\widehat{n}} + n*m*\$v. \widehat{(n+1/m)}) / (\$v. \widehat{(1/m)} * (m*(\$v. \widehat{(-1/m)} - 1))^{\widehat{2}}) \\ & \text{apply (subgoal-tac \$v. \widehat{(-1/m)} = 1 / \$v. \widehat{(1/m)}, erule ssubst)} \\ & \text{apply ((subst power2-eq-square)+, simp add: field-simps that)} \\ & \text{by (simp add: powr-minus-divide)} \\ \text{also have } \dots & = \\ & (\$v. \widehat{(n+1/m)} * (\$v. \widehat{(-n-1/m)} - (n*m+1)*\$v. \widehat{(-1/m)} + n*m)) / \\ & (\$v. \widehat{(1/m)} * (m*(\$v. \widehat{(-1/m)} - 1))^{\widehat{2}}) \\ & \text{apply (subgoal-tac \$v. \widehat{(-n-1/m)} = 1 / \$v. \widehat{(n+1/m)} \$v. \widehat{(-1/m)} = \$v^{\widehat{n}} \\ & / \$v. \widehat{(n+1/m)}) \\ & \text{apply ((erule ssubst)+, simp-all add: field-simps)} \\ & \text{using v-pos} \\ & \text{apply (simp add: powr-diff[THEN sym] powr-realpow[THEN sym])} \\ & \text{by (smt (verit) powr-minus-divide)} \\ \text{also have } \dots & = \\ & (\$v^{\widehat{n}} * (\$v. \widehat{(-n-1/m)} - (n*m+1)*\$v. \widehat{(-1/m)} + n*m)) / ((m*(\$v. \widehat{(-1/m)} - 1))^{\widehat{2}}) \\ & \text{apply (subst powr-add[of - n 1/m])} \\ & \text{using v-pos powr-realpow by simp} \\ \text{also have } \dots & = \\ & (\$v^{\widehat{n}} * ((1+i). \widehat{(n+1/m)} - (n*m+1)*(1+i). \widehat{(1/m)} + n*m)) / ((m*((1+i). \widehat{(1/m)} - 1))^{\widehat{2}}) \\ & \text{apply (subgoal-tac } -n-1/m = -(n+1/m) - 1/m = -(1/m), (erule ssubst)+ \\ & \text{apply (subst i-v-powr[THEN sym])+} \\ & \text{by simp-all} \\ \text{finally show ?thesis .} & \\ \text{qed} & \\ \text{thus ?thesis} & \\ \text{apply } - & \\ \text{using that v-futr-pos} & \\ \text{apply (subst Ilsm-Ilam, simp)} & \\ \text{apply (erule ssubst, simp)} & \\ \text{apply (rule disjI2)} & \\ \text{by (subst power-mult-distrib[THEN sym], simp add: v-pres-def)} & \\ \text{qed} & \end{aligned}$$

lemma *Imsm-calc-i-0*: $\$(I^{\{m\}}s)^{\{m\}-n} = (n*m+1)*n / (2*m)$ if $i = 0$ $m \neq 0$
for $n \text{ } m :: \text{nat}$
using that

```

apply (subst Ilsm-Ilam, simp)
by (subst Imam-calc-i-0; simp)

lemma Ila''m-Ilam: $(I\hat{\{l\}}a'')\hat{\{m\}}-n = (1+i).\hat{(1/m)} * $(I\hat{\{l\}}a)\hat{\{m\}}-n
  if l ≠ 0 m ≠ 0 for l m n :: nat
  unfolding ann-incr-def ann-due-incr-def using that
  apply (subst i-v, subst powr-powr, simp)
  apply (subst sum-distrib-left)
  apply (rule sum.cong; simp)
  apply (rule disjI2)
  by (smt (verit) add-divide-distrib powr-add)

lemma Ia''m-calc: $(Ia'')\hat{\{m\}}-n = (∑ j < n. (j+1)/m * (∑ k=j*m..<(j+1)*m.
  $v.\hat{(k/m)}))
  if m ≠ 0 for n m :: nat
  using that
  apply (subst Ila''m-Ilam; simp del: One-nat-def)
  apply (subst Iam-calc; simp)
  apply (subst sum-distrib-left)
  apply (rule sum.cong; simp)
  apply (subst sum-distrib-left)+
  apply (rule sum.cong; simp)
  apply (subst i-v-powr)
  using powr-add[of $v, THEN sym] by (simp add: field-simps)

lemma Ima''m-calc-aux: $(I\hat{\{m\}}a'')\hat{\{m\}}-n = (∑ k < n*m. $v.\hat{(k/m)} * (k+1) /
  m\hat{\{2\}})
  if m ≠ 0 for m::nat
  using that
  apply (subst Ila''m-Ilam, simp)
  apply (subst Imam-calc-aux, simp)
  apply (subst sum-distrib-left)
  apply (rule sum.cong; simp)
  using powr-add[of $v, THEN sym] i-v-powr by (simp add: field-simps)

lemma Ima''m-calc: $(I\hat{\{m\}}a'')\hat{\{m\}}-n = (1 - (n*m+1)*$v\hat{n} + n*m*$v.\hat{(n+1/m)}) /
  / (m*(1-$v.\hat{(1/m)}))\hat{\{2\}}
  if i ≠ 0 m ≠ 0 for n m :: nat
  using that v-pos
  apply (subst Ila''m-Ilam, simp)
  apply (subst Imam-calc; simp)
  by (simp add: powr-divide v-pres-def)

lemma Ils''m-Ilsm: $(I\hat{\{l\}}s'')\hat{\{m\}}-n = (1+i).\hat{(1/m)} * $(I\hat{\{l\}}s)\hat{\{m\}}-n
  if l ≠ 0 m ≠ 0 for l m n :: nat
  unfolding acc-incr-def acc-due-incr-def using that
  apply (subst sum-distrib-left)
  apply (rule sum.cong; simp)
  apply (rule disjI2)

```

by (subst powr-add[THEN sym], subst add-divide-distrib, simp)

lemma *Ims''m-calc*:

$$\$(I^{\{m\}} s'')^{\{m\}-n} = \\ (1+i). \widehat{(1/m)} * ((1+i). \widehat{(n+1/m)} - (n*m+1)*(1+i). \widehat{(1/m)} + n*m) / \\ (m*((1+i). \widehat{(1/m)} - 1))^2$$

if $i \neq 0$ $m \neq 0$ for $n m :: nat$

using that by (simp add: *Ils''m-Ilsm Imsm-calc*)

lemma *lim-Imam*: $(\lambda n. \$ (I^{\{m\}} a)^{\{m\}-n}) \longrightarrow 1 / (\$ i^{\{m\}} * \$ d^{\{m\}})$ if $m \neq 0$ $i > 0$ for $m :: nat$

proof –

have $(\lambda n. \$ (I^{\{m\}} a)^{\{m\}-n}) =$

$$(\lambda n. \$ v. \widehat{(1/m)} * (1 - (n*m+1)*\$v^{\widehat{n}} + n*m*\$v. \widehat{(n+1/m)}) / (m*(1-\$v. \widehat{(1/m)}))^2) \\ \text{using that by (subst Imam-calc; simp)}$$

$$\text{moreover have } (\lambda n. \$ v. \widehat{(1/m)} * (1 - (n*m+1)*\$v^{\widehat{n}} + n*m*\$v. \widehat{(n+1/m)})) \\ / (m*(1-\$v. \widehat{(1/m)}))^2) \longrightarrow 1 / (\$ i^{\{m\}} * \$ d^{\{m\}})$$

proof –

have $\star: |\$v| < 1$

using *v-lt-1-iff-i-pos v-pos* that by force

hence $(\lambda n. (n*m+1)*\$v^{\widehat{n}}) \longrightarrow 0$

apply (subst tendsto-cong[of - (\lambda n. n*m*\\$v^{\widehat{n}} + \\$v^{\widehat{n}})])

apply (rule always-eventually, rule allI)

apply (simp add: distrib-right)

apply (subgoal-tac 0 = 0 + 0, erule ssubst, intro tendsto-intros; simp)

apply (subst mult.commute, subst mult.assoc)

apply (subgoal-tac 0 = real m * 0, erule ssubst, intro tendsto-intros; simp?)

by (rule powser-times-n-limit-0; simp)

moreover have $(\lambda n. n*m*\$v. \widehat{(n+1/m)}) \longrightarrow 0$

apply (subst tendsto-cong[of - (\lambda n. (m*\\$v. \widehat{(1/m)})*(n*\\$v^{\widehat{n}}))])

apply (rule always-eventually, rule allI)

apply (subst powr-add, subst powr-realpow; simp add: v-pos)

apply (subgoal-tac 0 = m*\\$v. \widehat{(1/m)} * 0, erule ssubst, intro tendsto-intros; simp?)

by (rule powser-times-n-limit-0, simp add: \star)

$$\text{ultimately have } (\lambda n. \$ v. \widehat{(1/m)} * (1 - (n*m+1)*\$v^{\widehat{n}} + n*m*\$v. \widehat{(n+1/m)})) \\ / (m*(1-\$v. \widehat{(1/m)}))^2) \longrightarrow \$v. \widehat{(1/m)} * (1 - 0 + 0) / (m*(1-\$v. \widehat{(1/m)}))^2$$

using *v-lt-1-iff-i-pos v-pos* that by (intro tendsto-intros; simp)

thus ?thesis

unfolding *i-nom-def* using *v-pos* that

apply (subst i-v-powr, subst powr-minus-divide, subst d-nom-v; simp)

by (subst(asm)(2) power2-eq-square, simp add: field-simps)

qed

ultimately show ?thesis by simp

qed

lemma *perp-incr-calc*: $\$(I^{\{m\}} a)^{\{m\}-\infty} = 1 / (\$ i^{\{m\}} * \$ d^{\{m\}})$ if $m \neq 0$ $i >$

```

0 for m::nat
  unfolding perp-incr-def by (rule limI, rule lim-Imam; simp add: that)

lemma lim-Ima''m: ( $\lambda n. \$\{I^{\wedge\{m\}}a'\}^{\wedge\{m\}-n}$ ) ————— 1 / ( $\$d^{\wedge\{m\}}$ ) $^{\wedge 2}$  if  $m \neq 0$ 
i > 0 for m::nat
  unfolding perp-due-incr-def using that
  apply (subst Ila''m-Ilam, simp, subst mult.commute, subst i-v-powr, subst powr-minus-divide)
  apply (subgoal-tac 1 / ( $\$d^{\wedge\{m\}}$ ) $^{\wedge 2}$  =  $(1 / (\$i^{\wedge\{m\}} * \$d^{\wedge\{m\}})) * (1 / \$v. \wedge(1/m))$ , erule
ssubst)
  apply (intro tendsto-intros, simp add: lim-Imam)
  by (subst power2-eq-square, subst(1) d-nom-i-nom-v; simp add: field-simps that)

lemma perp-due-incr-calc:  $\$(I^{\wedge\{m\}}a')^{\wedge\{m\}-\infty} = 1 / (\$d^{\wedge\{m\}})^{\wedge 2}$  if  $m \neq 0$   $i >$ 
0 for m::nat
  unfolding perp-due-incr-def by (rule limI, rule lim-Ima''m; simp add: that)

end

end
theory Survival-Model
imports HOL-Library.Rewrite HOL-Library.Extended-Nonnegative-Real HOL-Library.Extended-Real
HOL-Probability.Probability Preliminaries
begin

```

5 Survival Model

The survival model is built on the probability space \mathfrak{M} . Additionally, the random variable $X : \text{space } \mathfrak{M} \rightarrow \mathbb{R}$ is introduced, which represents the age at death.

locale prob-space-actuary = MM-PS: prob-space \mathfrak{M} **for** \mathfrak{M}
— Since the letter M may be used as a commutation function, adopt the letter \mathfrak{M} instead as a notation for the measure space.

```

locale survival-model = prob-space-actuary +
fixes X :: 'a ⇒ real
assumes X-RV[simp]: MM-PS.random-variable (borel :: real measure) X
and X-pos-AE[simp]: AE  $\xi$  in  $\mathfrak{M}$ .  $X \xi > 0$ 
begin

```

5.1 General Theory of Survival Model

interpretation distrX-RD: real-distribution distr \mathfrak{M} borel X
using MM-PS.real-distribution-distr **by** simp

lemma X-le-event[simp]: $\{\xi \in \text{space } \mathfrak{M}. X \xi \leq x\} \in \text{MM-PS.events}$
by measurable simp

lemma X-gt-event[simp]: $\{\xi \in \text{space } \mathfrak{M}. X \xi > x\} \in \text{MM-PS.events}$

by measurable simp

lemma *X-compl-le-gt*: space \mathfrak{M} – { $\xi \in \text{space } \mathfrak{M}$. $X \xi \leq x$ } = { $\xi \in \text{space } \mathfrak{M}$. $X \xi > x$ } for $x::\text{real}$

proof –

have space \mathfrak{M} – { $\xi \in \text{space } \mathfrak{M}$. $X \xi \leq x$ } = space \mathfrak{M} – ($X - ` \{..x\}$) by blast
also have ... = ($X - ` \{x <..\}$) \cap space \mathfrak{M} using vimage-compl-atMost by fastforce

also have ... = { $\xi \in \text{space } \mathfrak{M}$. $X \xi > x$ } by blast

finally show ?thesis .

qed

lemma *X-compl-gt-le*: space \mathfrak{M} – { $\xi \in \text{space } \mathfrak{M}$. $X \xi > x$ } = { $\xi \in \text{space } \mathfrak{M}$. $X \xi \leq x$ } for $x::\text{real}$

using *X-compl-le-gt* by blast

5.1.1 Introduction of Survival Function for X

Note that *ccdf* (*distr* \mathfrak{M} *borel* X) is the survival (distributive) function for X .

lemma *ccdfX-0-1*: *ccdf* (*distr* \mathfrak{M} *borel* X) 0 = 1

apply (rewrite MM-PS.*ccdf-distr-P*, simp)

using *X-pos-AE* MM-PS.*prob-space*

using MM-PS.*prob-Collect-eq-1* *X-gt-event* by presburger

lemma *ccdfX-unborn-1*: *ccdf* (*distr* \mathfrak{M} *borel* X) $x = 1$ if $x \leq 0$

proof (rule antisym)

show *ccdf* (*distr* \mathfrak{M} *borel* X) $x \leq 1$ using MM-PS.*ccdf-distr-P* by simp

show *ccdf* (*distr* \mathfrak{M} *borel* X) $x \geq 1$

proof –

have *ccdf* (*distr* \mathfrak{M} *borel* X) $x \geq \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) 0$

using finite-borel-measure.*ccdf-nonincreasing* distrX-RD.finite-borel-measure-M that by simp

also have *ccdf* (*distr* \mathfrak{M} *borel* X) 0 = 1 using *ccdfX-0-1* that by simp

finally show ?thesis .

qed

qed

definition *death-pt* :: ereal ($\langle \$\psi \rangle$)

where $\$\psi \equiv \text{Inf} (\text{ereal} ` \{x \in \mathbb{R}. \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = 0\})$

— This is my original notation, which is used to develop life insurance mathematics rigorously.

lemma *psi-nonneg*: $\$\psi \geq 0$

unfolding *death-pt-def*

proof (rule Inf-greatest)

fix $x'::\text{ereal}$

assume $x' \in \text{ereal} ` \{x \in \mathbb{R}. \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = 0\}$

then obtain $x::\text{real}$ where $x' = \text{ereal } x$ and *ccdf* (*distr* \mathfrak{M} *borel* X) $x = 0$ by

blast

hence $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) 0 > \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x$ **using** ccdfX-0-1
 $X\text{-pos-AE}$ **by** simp

hence $x \geq 0$
 using $\text{mono-invE finite-borel-measure.ccdf-nonincreasing distrX-RD.finite-borel-measure-M}$
 $X\text{-RV}$
 by $(\text{smt}(\text{verit}))$

thus $x' \geq 0$ **using** $\langle x' = \text{ereal } x \rangle$ **by** simp

qed

lemma $\text{ccdfX-beyond-0}: \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x = 0$ **if** $x > \$\psi$ **for** $x::\text{real}$

proof –

have $\text{ereal} \{y \in \mathbb{R}. \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) y = 0\} \neq \{\}$ **using** death-pt-def **that**
by force

hence $\exists y' \in (\text{ereal} \{y \in \mathbb{R}. \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) y = 0\}). y' < \text{ereal } x$
 using $\text{that unfolding death-pt-def by (rule cInf-lessD)}$

then obtain y'
 where $y' \in (\text{ereal} \{y \in \mathbb{R}. \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) y = 0\})$ **and** $y' < \text{ereal } x$
by blast

then obtain $y::\text{real}$
 where $y' = \text{ereal } y$ **and** $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) y = 0$ **and** $\text{ereal } y < \text{ereal } x$
by blast

hence $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) y = 0$ **and** $y < x$ **by** simp-all

hence $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x \leq 0$
 using $\text{finite-borel-measure.ccdf-nonincreasing distrX-RD.finite-borel-measure-M}$
 $X\text{-RV}$
 by $(\text{metis order-less-le})$

thus $?thesis$ **using** $\text{finite-borel-measure.ccdf-nonneg distrX-RD.finite-borel-measure-M}$
 $X\text{-RV}$ **by** $(\text{smt}(\text{verit}))$

qed

lemma $\text{ccdfX-psi-0}: \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)(\text{real-of-ereal } \$\psi) = 0$ **if** $\$psi < \infty$

proof –

have $|\$psi| \neq \infty$ **using** that psi-nonneg **by** simp

then obtain $u::\text{real}$ **where** $\$psi = \text{ereal } u$ **using** ereal-real' **by** blast

hence $\text{real-of-ereal } \$psi = u$ **by** simp

moreover have $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) u = 0$

proof –

have $\bigwedge x::\text{real}. x \neq u \implies x \in \{u <..\} \implies \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x = 0$
 by $(\text{rule ccdfX-beyond-0}, \text{simp add: } \langle \$psi = \text{ereal } u \rangle)$

hence $(\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) \longrightarrow 0)$ **(at-right** u)

apply –
 by $(\text{rule iffD2[OF Lim-cong-within[where ?g=(\lambda x.0)]], simp-all+})$

moreover have $(\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) \longrightarrow \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) u)$
(at-right u)

using $\text{finite-borel-measure.ccdf-is-right-cont distrX-RD.finite-borel-measure-M}$
 continuous-within $X\text{-RV}$ **by** blast

ultimately show $?thesis$ **using** $\text{tendsto-unique trivial-limit-at-right-real}$ **by**
blast

```

qed
ultimately show ?thesis by simp
qed

lemma ccdfX-0-equiv: ccdf (distr M borel X) x = 0  $\longleftrightarrow$  x  $\geq$  $\psi for x::real
proof
  assume ccdf (distr M borel X) x = 0
  thus ereal x  $\geq$  $\psi unfolding death-pt-def by (simp add: INF-lower)
next
  assume $\psi  $\leq$  ereal x
  hence $\psi = ereal x  $\vee$  $\psi < ereal x unfolding less-eq-ereal-def by auto
  thus ccdf (distr M borel X) x = 0
  proof
    assume *: $\psi = ereal x
    hence $\psi < \infty by simp
    moreover have real-of-ereal $\psi = x using * by simp
    ultimately show ccdf (distr M borel X) x = 0 using ccdfX-psi-0 by simp
  next
    assume $\psi < ereal x
    thus ccdf (distr M borel X) x = 0 by (rule ccdfX-beyond-0)
  qed
qed

lemma psi-pos[simp]: $\psi > 0
proof (rule not-le-imp-less, rule notI)
  show $\psi  $\leq$  (0::ereal)  $\Longrightarrow$  False
  proof -
    assume $\psi  $\leq$  (0::ereal)
    hence ccdf (distr M borel X) 0 = 0 using ccdfX-0-equiv by (simp add: zero-ereal-def)
    moreover have ccdf (distr M borel X) 0 = 1 using ccdfX-0-1 by simp
    ultimately show False by simp
  qed
qed

corollary psi-pos'[simp]: $\psi > ereal 0
using psi-pos zero-ereal-def by presburger

```

5.1.2 Introduction of Future-Lifetime Random Variable $T(x)$

definition alive :: real \Rightarrow 'a set
where alive x \equiv { $\xi \in space M$. $X \xi > x$ }

lemma alive-event[simp]: alive x \in MM-PS.events for x::real
 unfolding alive-def by simp

lemma X-alivex-measurable[measurable, simp]: X \in borel-measurable ($M \downharpoonright alive x$) for x::real
 unfolding cond-prob-space-def by (measurable; simp add: measurable-restrict-space1)

definition *futr-life* :: *real* \Rightarrow (*'a* \Rightarrow *real*) ($\langle T \rangle$)
where $T x \equiv (\lambda \xi. X \xi - x)$
— Note that $T(x) : space \mathfrak{M} \rightarrow \mathbb{R}$ represents the time until death of a person aged x .

lemma *T0-eq-X[simp]*: $T 0 = X$
unfolding *futr-life-def* **by** *simp*

lemma *Tx-measurable[measurable, simp]*: $T x \in borel-measurable \mathfrak{M}$ **for** $x::real$
unfolding *futr-life-def* **by** (*simp add: borel-measurable-diff*)

lemma *Tx-alivex-measurable[measurable, simp]*: $T x \in borel-measurable (\mathfrak{M} \downarrow alive x)$ **for** $x::real$
unfolding *futr-life-def* **by** (*simp add: borel-measurable-diff*)

lemma *alive-T*: $alive x = \{\xi \in space \mathfrak{M}. T x \xi > 0\}$ **for** $x::real$
unfolding *alive-def futr-life-def* **by** *force*

lemma *alivex-Tx-pos[simp]*: $0 < T x \xi$ **if** $\xi \in space (\mathfrak{M} \downarrow alive x)$ **for** $x::real$
using *MM-PS.space-cond-prob-space alive-T* **that by** *simp*

lemma *PT0-eq-PX-lborel*: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T 0 \xi \in A \mid T 0 \xi > 0) = \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)$
if $A \in sets lborel$ **for** $A :: real$ **set**
apply (*rewrite MM-PS.cond-prob-AE-prob, simp-all*)
using *that X-RV measurable-lborel1 predE pred-sets2* **by** *blast*

5.1.3 Actuarial Notations on the Survival Model

definition *survive* :: *real* \Rightarrow *real* ($\langle \$p' \cdot \{&\} \rangle [0,0] 200$)
where $\$p \cdot \{t \& x\} \equiv ccdf (distr (\mathfrak{M} \downarrow alive x) borel (T x)) t$
— the probability that a person aged x will survive for t years
— Note that the function $\$p \cdot \{ \cdot \& x \}$ is the survival function on $(\mathfrak{M} \downarrow alive x)$ for the random variable $T(x)$.
— The parameter t is usually nonnegative, but theoretically it can be negative.

abbreviation *survive-1* :: *real* \Rightarrow *real* ($\langle \$p' \cdot \{&\} \rangle [101] 200$)
where $\$p \cdot x \equiv \$p \cdot \{1 \& x\}$

definition *die* :: *real* \Rightarrow *real* ($\langle \$q' \cdot \{&\} \rangle [0,0] 200$)
where $\$q \cdot \{t \& x\} \equiv cdf (distr (\mathfrak{M} \downarrow alive x) borel (T x)) t$
— the probability that a person aged x will die within t years
— Note that the function $\$q \cdot \{ \cdot \& x \}$ is the cumulative distributive function on $(\mathfrak{M} \downarrow alive x)$ for the random variable $T(x)$.
— The parameter t is usually nonnegative, but theoretically it can be negative.

abbreviation *die-1* :: *real* \Rightarrow *real* ($\langle \$q' \cdot \{&\} \rangle [101] 200$)
where $\$q \cdot x \equiv \$q \cdot \{1 \& x\}$

definition *die-defer* :: *real* \Rightarrow *real* \Rightarrow *real* \Rightarrow *real* ($\langle \$q' -\{|\&-\} \rangle [0,0,0] 200$)
where $\$q\{-f|t&x\} = |\$q\{-f+t&x\} - \$q\{-f&x\}|$
— the probability that a person aged x will die within t years, deferred f years
— The parameters f and t are usually nonnegative, but theoretically they can be negative.

abbreviation *die-defer-1* :: *real* \Rightarrow *real* \Rightarrow *real* ($\langle \$q' -\{|\&-\} \rangle [0,0] 200$)
where $\$q\{-f|x\} \equiv \$q\{-f|1&x\}$

definition *life-expect* :: *real* \Rightarrow *real* ($\langle \$e'^{\circ} -\{|\&-\} \rangle [101] 200$)
where $\$e'^{\circ}-x \equiv \text{integral}^L (\mathfrak{M} \downharpoonright \text{alive } x) (T x)$
— complete life expectation
— Note that $\$e'^{\circ}-x$ is calculated as 0 when $\text{nn-integral} (\mathfrak{M} \downharpoonright \text{alive } x) (T x) = \infty$.

definition *temp-life-expect* :: *real* \Rightarrow *real* \Rightarrow *real* ($\langle \$e'^{\circ} -\{|\&-\} \rangle [0,0] 200$)
where $\$e'^{\circ}-\{x:n\} \equiv \text{integral}^L (\mathfrak{M} \downharpoonright \text{alive } x) (\lambda \xi. \min (T x \xi) n)$
— temporary complete life expectation

definition *curt-life-expect* :: *real* \Rightarrow *real* ($\langle \$e' -\{|\&-\} \rangle [101] 200$)
where $\$e-x \equiv \text{integral}^L (\mathfrak{M} \downharpoonright \text{alive } x) (\lambda \xi. \lfloor T x \xi \rfloor)$
— curtate life expectation
— Note that $\$e-x$ is calculated as 0 when $\text{nn-integral} (\mathfrak{M} \downharpoonright \text{alive } x) \lfloor T x \rfloor = \infty$.

definition *temp-curt-life-expect* :: *real* \Rightarrow *real* \Rightarrow *real* ($\langle \$e' -\{|\&-\} \rangle [0,0] 200$)
where $\$e-\{x:n\} \equiv \text{integral}^L (\mathfrak{M} \downharpoonright \text{alive } x) (\lambda \xi. \lfloor \min (T x \xi) n \rfloor)$
— temporary curtate life expectation
— In the definition n can be a real number, but in practice n is usually a natural number.

5.1.4 Properties of Survival Function for $T(x)$

context
fixes $x::\text{real}$
assumes $x < \psi$
begin

lemma *PXx-pos*[simp]: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) > 0$
proof —
have $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) = \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) x$
unfolding *alive-def* **using** *MM-PS.ccdf-distr-P* **by** *simp*
also have $\dots > 0$
using *ccdfX-0-equiv distrX-RD.ccdf-nonneg x-lt-psi* **by** (*smt (verit) linorder-not-le*)
finally show *?thesis* .

qed

lemma *PTx-pos*[simp]: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0) > 0$
apply (*rewrite alive-T[THEN sym]*)

unfolding alive-def **by** simp

interpretation alivex-PS: prob-space $\mathfrak{M} \downharpoonright \text{alive } x$
by (rule MM-PS.cond-prob-space-correct, simp-all add: alive-def)

interpretation distrTx-RD: real-distribution distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$) **by**
simp

lemma ccdfTx-cond-prob:

ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) $t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > t \mid T x \xi > 0)$ **for**
 $t : \text{real}$
apply (rewrite alivex-PS.ccdf-distr-P, simp)
unfolding alive-def
apply (rewrite MM-PS.cond-prob-space-cond-prob[THEN sym], simp-all add:
pred-def)
unfolding futr-life-def **by** simp

lemma ccdfTx-0-1: ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) $0 = 1$

apply (rewrite ccdfTx-cond-prob)
unfolding futr-life-def cond-prob-def
by (smt (verit, best) Collect-cong PXx-pos divide-eq-1-iff)

lemma ccdfTx-nonpos-1: ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) $t = 1$ **if** $t \leq 0$ **for**
 $t :: \text{real}$
using antisym ccdfTx-0-1 **that**
by (metis distrTx-RD.ccdf-bounded-prob distrTx-RD.ccdf-nonincreasing)

lemma ccdfTx-0-equiv: ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) $t = 0 \longleftrightarrow x+t \geq \ψ **for** $t :: \text{real}$

proof –
have ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) $t =$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t \wedge X \xi > x) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
apply (rewrite ccdfTx-cond-prob)
unfolding cond-prob-def futr-life-def **by** (smt (verit) Collect-cong)
hence ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) $t = 0 \longleftrightarrow$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t \wedge X \xi > x) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) = 0$
by simp
also have ... $\longleftrightarrow \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t \wedge X \xi > x) = 0$
using x-lt-psi PXx-pos **by** (smt (verit) divide-eq-0-iff)
also have ... $\longleftrightarrow x+t \geq \ψ
using ccdfX-0-equiv MM-PS.ccdf-distr-P
by (smt (verit) Collect-cong X-RV le-ereal-le linorder-not-le x-lt-psi)
finally show ?thesis .
qed

lemma ccdfTx-continuous-on-nonpos[simp]:

continuous-on $\{\dots\}$ (ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)))
by (metis atMost-iff ccdfTx-nonpos-1 continuous-on-cong continuous-on-const)

```

lemma ccdfTx-differentiable-on-nonpos[simp]:
  (ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))) differentiable-on {..0}
  by (rewrite differentiable-on-cong[where  $f = \lambda \_. 1$ ]; simp add: ccdfTx-nonpos-1)

lemma ccdfTx-has-real-derivative-0-at-neg:
  (ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))) has-real-derivative 0 (at t) if  $t < 0$  for  $t :: \text{real}$ 
  apply (rewrite has-real-derivative-iff-has-vector-derivative)
  apply (rule has-vector-derivative-transform-within-open[of  $\lambda \_. 1 - \{.. < 0\}$ ])
  using ccdfTx-nonpos-1 that by simp-all

lemma ccdfTx-integrable-Icc:
  set-integrable lborel {a..b} (ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))) for  $a b :: \text{real}$ 
proof -
  have ( $\int^+ t. \text{ennreal} (\text{indicat-real } \{a..b\} t) * \text{ccdf} (\text{distr } (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) t) \partial\text{lborel}$ 
    <  $\top$ 
proof -
  have ( $\int^+ t. \text{ennreal} (\text{indicat-real } \{a..b\} t) * \text{ccdf} (\text{distr } (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) t) \partial\text{lborel}$ 
     $\leq (\int^+ t. \text{ennreal} (\text{indicat-real } \{a..b\} t) \partial\text{lborel})$ 
  apply (rule nn-integral-mono)
  using distrTx-RD.ccdf-bounded
  by (simp add: distrTx-RD.ccdf-bounded-prob indicator-times-eq-if(1))
  also have ... = nn-integral lborel (indicator {a..b}) by (meson ennreal-indicator)
  also have ... = emeasure lborel {a..b} by (rewrite nn-integral-indicator; simp)
  also have ... <  $\top$ 
  using emeasure-lborel-Icc-eq ennreal-less-top infinity-ennreal-def by presburger
  finally show ?thesis .
qed
thus ?thesis
  unfolding set-integrable-def
  apply (intro integrableI-nonneg, simp-all)
  using distrTx-RD.ccdf-nonneg by (intro always-eventually) auto
qed

corollary ccdfTx-integrable-on-Icc:
  ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))) integrable-on {a..b} for  $a b :: \text{real}$ 
  using set-borel-integral-eq-integral ccdfTx-integrable-Icc by force

lemma ccdfTx-PX:
  ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ )))  $t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ 
  if  $t \geq 0$  for  $t :: \text{real}$ 
  apply (rewrite ccdfTx-cond-prob)
  unfolding cond-prob-def futr-life-def PXx-pos by (smt (verit) Collect-cong that)

lemma ccdfTx-ccdfX: ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ )))  $t =$ 
  ccdf (distr  $\mathfrak{M}$  borel  $X$ )  $(x + t) / \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) x$ 

```

if $t \geq 0$ **for** $t::real$
using $\text{ccdf}Tx-PX$ **that** $\text{MM-PS.ccdf-distr-P } X-RV$ **by** presburger

lemma $\text{ccdfT0-eq-ccdfX}: \text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } 0) \text{ borel } (T 0)) = \text{ccdf} (\text{distr } \mathfrak{M}$
 $\text{borel } X)$
proof
fix x
show $\text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } 0) \text{ borel } (T 0)) x = \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) x$
proof (**cases** $x \geq 0$)
case True
thus ?thesis
using $\text{survival-model.ccdf}Tx-\text{ccdf}X[\text{where } x=0]$ $\text{ccdf}X\text{-0-1 survival-model-axioms}$
by fastforce
next
case False
hence $x \leq 0$ **by** simp
thus ?thesis
apply (**rewrite** $\text{ccdf}X\text{-unborn-1}$, simp)
by (**rewrite** $\text{survival-model.ccdf}Tx\text{-nonpos-1}$; $\text{simp add: survival-model-axioms}$)

qed
qed

lemma $\text{continuous-ccdf}X-\text{ccdf}Tx$:
 $\text{continuous} (\text{at } (x+t) \text{ within } \{x..\}) (\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)) \longleftrightarrow$
 $\text{continuous} (\text{at } t \text{ within } \{0..\}) (\text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)))$
if $t \geq 0$ **for** $t::real$
proof –
let $?srvl = \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$
have [**simp**]: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) \neq 0$ **using** $PXx\text{-pos}$ **by** force
have $\star: \bigwedge u. u \geq 0 \implies \text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) u =$
 $?srvl (x + u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
using $\text{survive-def MM-PS.ccdf-distr-P } ccdfTx-PX$ **that** **by** simp
have $\text{continuous} (\text{at } t \text{ within } \{0..\}) (\text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x))) \longleftrightarrow$
 $\text{continuous} (\text{at } t \text{ within } \{0..\}) (\lambda u. ?srvl (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. x < X \xi))$
proof –
have $\forall_F u \text{ in at } t \text{ within } \{0..\}. \text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) u =$
 $?srvl (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
using \star **by** (**rewrite** $\text{eventually-at-topological}$, simp-all) **blast**
thus ?thesis
by (**intro** $\text{continuous-at-within-cong}$, $\text{simp-all add: } \star \text{ that}$)
qed
also have ... $\longleftrightarrow \text{continuous} (\text{at } t \text{ within } \{0..\}) (\lambda u. ?srvl (x+u))$
by (**rewrite** $\text{at } - = \square \text{ continuous-cdivide-iff}[\text{of } \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)]$, simp-all)
also have ... $\longleftrightarrow \text{continuous} (\text{at } (x+t) \text{ within } \{x..\}) ?srvl$
proof
let $?subx = \lambda v. v - x$
assume $LHS: \text{continuous} (\text{at } t \text{ within } \{0..\}) (\lambda u. ?srvl (x+u))$
hence $\text{continuous} (\text{at } (?subx (x+t)) \text{ within } ?subx ' \{x..\}) (\lambda u. ?srvl (x+u))$
proof –

```

have ?subx ` {x..} = {0..}
  by (metis (no-types, lifting) add.commute add-uminus-conv-diff diff-self
       image-add-atLeast image-cong)
thus ?thesis using LHS by simp
qed
moreover have continuous (at (x+t) within {x..}) ?subx by (simp add: continuous-diff)
ultimately have continuous (at (x+t) within {x..}) (λu. ?srvl (x + (?subx u)))
  using continuous-within-compose2 by force
thus continuous (at (x+t) within {x..}) ?srvl by simp
next
assume RHS: continuous (at (x+t) within {x..}) ?srvl
hence continuous (at ((plus x) t) within (plus x) ` {0..}) ?srvl by simp
moreover have continuous (at t within {0..}) (plus x) by (simp add: continuous-add)
ultimately show continuous (at t within {0..}) (λu. ?srvl (x+u))
  using continuous-within-compose2 by force
qed
finally show ?thesis by simp
qed

lemma isCont-ccdfX-ccdfTx:
isCont (ccdf (distr M borel X)) (x+t) ↔
  isCont (ccdf (distr (M ↳ alive x) borel (T x))) t
  if t > 0 for t::real
proof -
have isCont (ccdf (distr M borel X)) (x+t) ↔
  continuous (at (x+t) within {x<..}) (ccdf (distr M borel X))
  by (smt (verit) at-within-open greaterThan-iff open-greaterThan that)
also have ... ↔ continuous (at (x+t) within {x..}) (ccdf (distr M borel X))
  by (meson Ioi-le-Ico calculation continuous-within-subset top-greatest)
also have ... ↔ continuous (at t within {0..}) (ccdf (distr (M ↳ alive x) borel (T x)))
  using that continuous-ccdfX-ccdfTx by force
also have ... ↔ continuous (at t within {0<..}) (ccdf (distr (M ↳ alive x) borel (T x)))
  by (metis Ioi-le-Ico at-within-open continuous-at-imp-continuous-at-within
       continuous-within-subset greaterThan-iff open-greaterThan that)
also have ... ↔ isCont (ccdf (distr (M ↳ alive x) borel (T x))) t
  by (metis at-within-open greaterThan-iff open-greaterThan that)
finally show ?thesis .
qed

lemma has-real-derivative-ccdfX-ccdfTx:
((ccdf (distr M borel X)) has-real-derivative D) (at (x+t)) ↔
  ((ccdf (distr (M ↳ alive x) borel (T x))) has-real-derivative (D / P(ξ in M. X ξ
  > x))) (at t)
  if t > 0 for t D :: real

```

```

proof -
  have ((ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))) has-real-derivative
    ( $D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ ) (at  $t$ )  $\longleftrightarrow$ 
    (( $\lambda t.$  (ccdf (distr  $\mathfrak{M}$  borel  $X$ )) ( $x+t$ ) /  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ ) has-real-derivative
    ( $D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ ) (at  $t$ )
proof -
  let ?d =  $t/2$ 
  { fix  $u:\text{real}$  assume dist  $u t < ?d$ 
    hence  $u > 0$  by (smt (verit) dist-real-def dist-triangle-half-r)
    hence ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))  $u =$ 
      ccdf (distr  $\mathfrak{M}$  borel  $X$ ) ( $x+u$ ) / MM-PS.prob { $\xi::'a \in \text{space } \mathfrak{M}. x < X \xi$ }
      using survive-def MM-PS.ccdf-distr-P ccdfTx-PX that by simp }
    moreover have ?d > 0 using that by simp
    ultimately show ?thesis
      apply -
      apply (rule DERIV-cong-ev, simp)
      apply (rewrite eventually-nhds-metric, blast)
      by simp
    qed
    also have ...  $\longleftrightarrow$  (( $\lambda t.$  (ccdf (distr  $\mathfrak{M}$  borel  $X$ )) ( $x+t$ )) has-real-derivative  $D$ )
    (at  $t$ )
    using PXx-pos by (rewrite DERIV-cdivide-iff[of  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ , THEN
    sym]; force)
    also have ...  $\longleftrightarrow$  (ccdf (distr  $\mathfrak{M}$  borel  $X$ ) has-real-derivative  $D$ ) (at ( $x+t$ ))
    by (simp add: DERIV-shift add.commute)
    finally show ?thesis by simp
  qed

lemma differentiable-ccdfX-ccdfTx:
  (ccdf (distr  $\mathfrak{M}$  borel  $X$ )) differentiable at ( $x+t$ )  $\longleftrightarrow$ 
  (ccdf (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ))) differentiable at  $t$ 
  if  $t > 0$  for  $t:\text{real}$ 
  apply (rewrite differentiable-eq-field-differentiable-real)+
  unfolding field-differentiable-def using has-real-derivative-ccdfX-ccdfTx that
  by (smt (verit, del-insts) PXx-pos nonzero-mult-div-cancel-left)

```

5.1.5 Properties of $\$p\{-t\&x\}$

```

lemma p-0-1:  $\$p\{-0\&x\} = 1$ 
  unfolding survive-def using ccdfTx-0-1 by simp

lemma p-nonneg[simp]:  $\$p\{-t\&x\} \geq 0$  for  $t:\text{real}$ 
  unfolding survive-def using distrTx-RD.ccdf-nonneg by simp

lemma p-le-1[simp]:  $\$p\{-t\&x\} \leq 1$  for  $t:\text{real}$ 
  unfolding survive-def using distrTx-RD.ccdf-bounded-prob by auto

lemma p-0-equiv:  $\$p\{-t\&x\} = 0 \longleftrightarrow x+t \geq \$\psi$  for  $t:\text{real}$ 
  unfolding survive-def by (rule ccdfTx-0-equiv)

```

```

lemma p-PTx: $p-{t&x} = P(ξ in M. T x ξ > t | T x ξ > 0) for t::real
unfolding survive-def using ccdfTx-cond-prob by simp

lemma p-PX: $p-{t&x} = P(ξ in M. X ξ > x + t) / P(ξ in M. X ξ > x) if t ≥
0 for t::real
unfolding survive-def using ccdfTx-PX that by simp

lemma p-mult: $p-{t+t' & x} = $p-{t&x} * $p-{t' & x+t}
if t ≥ 0 t' ≥ 0 x+t < $ψ for t t' :: real
proof -
  have $p-{t+t' & x} = P(ξ in M. X ξ > x+t+t') / P(ξ in M. X ξ > x)
  apply (rewrite p-PX; (simp add: that)?)
  by (rule disjI2, smt (verit, best) Collect-cong)
  also have ... = (P(ξ in M. X ξ > x+t+t') / P(ξ in M. X ξ > x+t)) *
  (P(ξ in M. X ξ > x+t) / P(ξ in M. X ξ > x))
  using that survival-model.PXx-pos survival-model-axioms by fastforce
  also have ... = $p-{t&x} * $p-{t' & x+t}
  apply (rewrite p-PX, simp add: that)
  by (rewrite survival-model.p-PX, simp-all add: that survival-model-axioms)
  finally show ?thesis .
qed

lemma p-PTx-ge-ccdf-isCont: $p-{t&x} = P(ξ in M. T x ξ ≥ t | T x ξ > 0)
if isCont (ccdf (distr M borel X)) (x+t) t > 0 for t::real
unfolding survive-def using that isCont-ccdfX-ccdfTx
apply (rewrite alivex-PS.ccdf-continuous-distr-P-ge, simp-all)
by (rewrite MM-PS.cond-prob-space-cond-prob, simp-all add: alive-T)

end

```

5.1.6 Properties of Survival Function for X

```

lemma ccdfX-continuous-unborn[simp]: continuous-on {..0} (ccdf (distr M borel
X))
using ccdfTx-continuous-on-nonpos by (metis ccdfT0-eq-ccdfX psi-pos')

lemma ccdfX-differentiable-unborn[simp]: (ccdf (distr M borel X)) differentiable-on
{..0}
using ccdfTx-differentiable-on-nonpos by (metis ccdfT0-eq-ccdfX psi-pos')

lemma ccdfX-has-real-derivative-0-unborn:
(ccdf (distr M borel X) has-real-derivative 0) (at x) if x < 0 for x::real
using ccdfTx-has-real-derivative-0-at-neg by (metis ccdfT0-eq-ccdfX psi-pos' that)

lemma ccdfX-integrable-Icc:
set-integrable lborel {a..b} (ccdf (distr M borel X)) for a b :: real
using ccdfTx-integrable-Icc by (metis ccdfT0-eq-ccdfX psi-pos')

```

corollary *ccdfX-integrable-on-Icc*:
ccdf (distr \mathfrak{M} borel X) integrable-on { $a..b$ } for $a b :: real$
using set-borel-integral-eq-integral ccdfX-integrable-Icc by force

lemma *ccdfX-p*: *ccdf (distr \mathfrak{M} borel X) $x = \$p\{-x & 0\}$ for $x::real$*
by (metis ccdfT0-eq-ccdfX survive-def psi-pos')

5.1.7 Introduction of Cumulative Distributive Function for X

lemma *cdfX-0-0*: *cdf (distr \mathfrak{M} borel X) $0 = 0$*
using ccdfX-0-1 distrX-RD.cdf-cdf distrX-RD.prob-space by fastforce

lemma *cdfX-unborn-0*: *cdf (distr \mathfrak{M} borel X) $x = 0$ if $x \leq 0$*
using ccdfX-unborn-1 cdfX-0-0 distrX-RD.cdf-ccdf that by fastforce

lemma *cdfX-beyond-1*: *cdf (distr \mathfrak{M} borel X) $x = 1$ if $x > \$\psi$ for $x::real$*
using ccdfX-beyond-0 distrX-RD.cdf-ccdf that distrX-RD.prob-space by force

lemma *cdfX-psi-1*: *cdf (distr \mathfrak{M} borel X) ($real-of-ereal \$\psi$) = 1 if $\psi < \infty$*
using ccdfX-psi-0 distrX-RD.cdf-ccdf distrX-RD.prob-space that by fastforce

lemma *cdfX-1-equiv*: *cdf (distr \mathfrak{M} borel X) $x = 1 \longleftrightarrow x \geq \psi$ for $x::real$*
using ccdfX-0-equiv distrX-RD.cdf-ccdf distrX-RD.prob-space by force

5.1.8 Properties of Cumulative Distributive Function for $T(x)$

context

fixes $x::real$
assumes *x-lt-psi[simp]*: $x < \psi$
begin

interpretation *alivex-PS*: *prob-space $\mathfrak{M} \downharpoonright alive x$*
by (rule MM-PS.cond-prob-space-correct, simp-all add: alive-def)

interpretation *distrTx-RD*: *real-distribution distr ($\mathfrak{M} \downharpoonright alive x$) borel (T x)* by
simp

lemma *cdfTx-cond-prob*:
cdf (distr ($\mathfrak{M} \downharpoonright alive x$) borel (T x)) $t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq t \mid T x \xi > 0)$ for $t::real$
apply (rewrite distrTx-RD.cdf-ccdf, rewrite distrTx-RD.prob-space)
apply (rewrite ccdfTx-cond-prob, simp)
by (rewrite not-less[THEN sym], rewrite MM-PS.cond-prob-neg; simp)

lemma *cdfTx-0-0*: *cdf (distr ($\mathfrak{M} \downharpoonright alive x$) borel (T x)) $0 = 0$*
using ccdfTx-0-1 distrTx-RD.cdf-ccdf distrTx-RD.prob-space by force

lemma *cdfTx-nonpos-0*: *cdf (distr ($\mathfrak{M} \downharpoonright alive x$) borel (T x)) $t = 0$ if $t \leq 0$ for $t :: real$*
using ccdfTx-nonpos-1 distrTx-RD.cdf-ccdf distrTx-RD.prob-space that by force

```

lemma cdfTx-1-equiv:  $cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x)) t = 1 \longleftrightarrow x+t \geq \$\psi$ 
for  $t::real$ 
using ccdfTx-0-equiv distrTx-RD.cdf-ccdf distrTx-RD.prob-space by force

lemma cdfTx-continuous-on-nonpos[simp]:
 $continuous-on \{..0\} (cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x)))$ 
by (rewrite continuous-on-cong[where  $g=\lambda t. 0$ ]) (simp-all add: cdfTx-nonpos-0)+

lemma cdfTx-differentiable-on-nonpos[simp]:
 $(cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x))) differentiable-on \{..0\}$ 
by (rewrite differentiable-on-cong[where  $f=\lambda t. 0$ ]; simp add: cdfTx-nonpos-0)

lemma cdfTx-has-real-derivative-0-at-neg:
 $(cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x)) has-real-derivative 0) (at t) \text{ if } t < 0$  for
 $t::real$ 
apply (rewrite has-real-derivative-iff-has-vector-derivative)
apply (rule has-vector-derivative-transform-within-open[of  $\lambda -. 0 - - \{.. < 0\}$ ])
using cdfTx-nonpos-0 that by simp-all

lemma cdfTx-integrable-Icc:
 $set-integrable lborel \{a..b\} (cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x)))$  for  $a b :: real$ 
proof -
  have set-integrable lborel \{a..b\} ( $\lambda -. 1::real$ )
  unfolding set-integrable-def
  using emeasure-compact-finite by (simp, intro integrable-real-indicator; force)
  thus ?thesis
    apply (rewrite distrTx-RD.cdf-ccdf, rewrite distrTx-RD.prob-space)
    using ccdfTx-integrable-Icc by (rewrite set-integral-diff; simp)
  qed

corollary cdfTx-integrable-on-Icc:
 $cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x)) integrable-on \{a..b\}$  for  $a b :: real$ 
using cdfTx-integrable-Icc set-borel-integral-eq-integral by force

lemma cdfTx-PX:
 $cdf(distr(\mathfrak{M} \downarrow alive x) borel(T x)) t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x < X \xi \wedge X \xi \leq x+t) /$ 
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ 
for  $t::real$ 
apply (rewrite cdfTx-cond-prob)
unfolding cond-prob-def futr-life-def PXx-pos by (smt (verit) Collect-cong)

lemma cdfT0-eq-cdfX:  $cdf(distr(\mathfrak{M} \downarrow alive 0) borel(T 0)) = cdf(distr \mathfrak{M} borel X)$ 
proof
  interpret alive0-PS: prob-space  $\mathfrak{M} \downarrow alive 0$ 
  apply (rule MM-PS.cond-prob-space-correct, simp)
  using PXx-pos alive-def psi-pos' by presburger
  interpret distrT0-RD: real-distribution distr( $\mathfrak{M} \downarrow alive 0$ ) borel(T 0) by simp

```

```

show  $\wedge x. \text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } 0) \text{ borel } (T 0)) x = \text{cdf}(\text{distr } \mathfrak{M} \text{ borel } X) x$ 
  using  $\text{ccdfT0-eq-ccdfX distrX-RD.ccdf-cdf distrT0-RD.ccdf-cdf}$ 
  by (smt (verit, best)  $\text{distrT0-RD.prob-space distrX-RD.prob-space psi-pos}'$ )
qed

```

```

lemma continuous-cdfX-cdfTx:
  continuous(at(x+t) within {x..}) (cdf(distr  $\mathfrak{M}$  borel X))  $\longleftrightarrow$ 
    continuous(at t within {0..}) (cdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x)))
  if  $t \geq 0$  for  $t:\text{real}$ 
proof -
  have continuous(at(x+t) within {x..}) (cdf(distr  $\mathfrak{M}$  borel X))  $\longleftrightarrow$ 
    continuous(at(x+t) within {x..}) (ccdf(distr  $\mathfrak{M}$  borel X))
  by (rule  $\text{distrX-RD.continuous-cdf-ccdf}$ )
  also have ...  $\longleftrightarrow$  continuous(at t within {0..}) (ccdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x)))
  using continuous-ccdfX-ccdfTx that by simp
  also have ...  $\longleftrightarrow$  continuous(at t within {0..}) (cdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x)))
  using  $\text{distrTx-RD.continuous-cdf-ccdf}$  by simp
  finally show ?thesis .
qed

```

```

lemma isCont-cdfX-cdfTx:
  isCont(cdf(distr  $\mathfrak{M}$  borel X)) (x+t)  $\longleftrightarrow$ 
  isCont(cdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x))) t
  if  $t > 0$  for  $t:\text{real}$ 
  apply (rewrite  $\text{distrX-RD.isCont-cdf-ccdf}$ )
  apply (rewrite  $\text{isCont-ccdfX-ccdfTx}$ , simp-all add: that)
  by (rule  $\text{distrTx-RD.isCont-cdf-ccdf}[\text{THEN sym}]$ )

```

```

lemma has-real-derivative-cdfX-cdfTx:
  ((cdf(distr  $\mathfrak{M}$  borel X)) has-real-derivative D) (at (x+t))  $\longleftrightarrow$ 
  ((cdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x))) has-real-derivative (D /  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ ) (at t))
  if  $t > 0$  for  $t D :: \text{real}$ 
proof -
  have ((cdf(distr  $\mathfrak{M}$  borel X)) has-real-derivative D) (at (x+t))  $\longleftrightarrow$ 
    (ccdf(distr  $\mathfrak{M}$  borel X) has-real-derivative -D) (at (x+t))
  using  $\text{distrX-RD.has-real-derivative-cdf-ccdf}$  by force
  also have ...  $\longleftrightarrow$ 
    ((ccdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x))) has-real-derivative (-D /  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ ) (at t))
  using has-real-derivative-ccdfX-ccdfTx that by simp
  also have ...  $\longleftrightarrow$ 
    ((cdf(distr( $\mathfrak{M} \downarrow \text{alive } x$ ) borel (T x))) has-real-derivative (D /  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ ) (at t))
  by (simp add:  $\text{distrTx-RD.has-real-derivative-cdf-ccdf}$ )
  finally show ?thesis .
qed

```

```

lemma differentiable-cdfX-cdfTx:
  (cdf (distr M borel X)) differentiable at (x+t)  $\longleftrightarrow$ 
  (cdf (distr (M \ alive x) borel (T x))) differentiable at t
  if t > 0 for t::real
  apply (rewrite differentiable-eq-field-differentiable-real)+
  unfolding field-differentiable-def using has-real-derivative-cdfX-cdfTx that
  by (meson differentiable-ccdfX-ccdfTx distrTx-RD.finite-borel-measure-axioms
    distrX-RD.finite-borel-measure-axioms finite-borel-measure.differentiable-cdf-ccdf
    real-differentiable-def x-lt-psi)

```

5.1.9 Properties of $\$q\{t \& x\}$

```

lemma q-nonpos-0:  $\$q\{t \& x\} = 0$  if t ≤ 0 for t::real
  unfolding die-def using that cdfTx-nonpos-0 by simp

```

```

corollary q-0-0:  $\$q\{0 \& x\} = 0$ 
  using q-nonpos-0 by simp

```

```

lemma q-nonneg[simp]:  $\$q\{t \& x\} \geq 0$  for t::real
  unfolding die-def using distrTx-RD.cdf-nonneg by simp

```

```

lemma q-le-1[simp]:  $\$q\{t \& x\} \leq 1$  for t::real
  unfolding die-def using distrTx-RD.cdf-bounded-prob by force

```

```

lemma q-1-equiv:  $\$q\{t \& x\} = 1 \longleftrightarrow x + t \geq \$\psi$  for t::real
  unfolding die-def using cdfTx-1-equiv by simp

```

```

lemma q-PTx:  $\$q\{t \& x\} = \mathcal{P}(\xi \text{ in } M. T x \xi \leq t \mid T x \xi > 0)$  for t::real
  unfolding die-def using cdfTx-cond-prob by simp

```

```

lemma q-PX:  $\$q\{t \& x\} = \mathcal{P}(\xi \text{ in } M. x < X \xi \wedge X \xi \leq x + t) / \mathcal{P}(\xi \text{ in } M. X \xi > x)$ 
  unfolding die-def using cdfTx-PX by simp

```

```

lemma q-defer-0-q[simp]:  $\$q\{0 | t \& x\} = \$q\{t \& x\}$  for t::real
  unfolding die-defer-def using q-0-0 by simp

```

```

lemma q-defer-0-0:  $\$q\{f | 0 \& x\} = 0$  for f::real
  unfolding die-defer-def by simp

```

```

lemma q-defer-nonneg[simp]:  $\$q\{f | t \& x\} \geq 0$  for f t :: real
  unfolding die-defer-def by simp

```

```

lemma q-defer-q:  $\$q\{f | t \& x\} = \$q\{f + t \& x\} - \$q\{f \& x\}$  if t ≥ 0 for f t :: real
  unfolding die-defer-def die-def using distrTx-RD.cdf-nondecreasing that by
  simp

```

```

corollary q-defer-le-1[simp]:  $\$q\{f | t \& x\} \leq 1$  if t ≥ 0 for f t :: real

```

by (*smt (verit, ccfv-SIG) q-defer-q q-le-1 q-nonneg that*)

lemma *q-defer-PTx*: $\$q\{-f|t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t \mid T x \xi > 0)$
if $t \geq 0$ **for** $f t :: \text{real}$

proof –

have $\$q\{-f|t\&x\} = \$q\{-f+t \& x\} - \$q\{-f\&x\}$ **using** *q-defer-q* **that by** *simp*
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq f + t \mid T x \xi > 0) - \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq f \mid T x \xi > 0)$
using *q-PTx* **by** *simp*
also have $\dots = \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f + t) - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f)$
using *MM-PS.cond-prob-space-cond-prob alive-T* **by** *simp*
also have $\dots = \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). f < T x \xi \wedge T x \xi \leq f + t)$

proof –

have $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f + t\} - \{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f\} =$
 $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). f < T x \xi \wedge T x \xi \leq f + t\}$
using *that* **by** *force*
hence *alivex-PS.prob*
 $(\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f + t\} - \{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f\}) =$
 $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). f < T x \xi \wedge T x \xi \leq f + t)$
by *simp*
moreover have $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f\} \subseteq \{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \leq f + t\}$
using *that* **by** *force*
ultimately show *?thesis* **by** (*rewrite alivex-PS.finite-measure-Diff[THEN sym]; simp*)
qed
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t \mid T x \xi > 0)$
using *MM-PS.cond-prob-space-cond-prob alive-T* **by** *simp*
finally show *?thesis*.

qed

lemma *q-defer-PX*: $\$q\{-f|t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x + f < X \xi \wedge X \xi \leq x + f + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
if $f \geq 0 t \geq 0$ **for** $f t :: \text{real}$

proof –

have $\$q\{-f|t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t \wedge T x \xi > 0) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0)$
apply (*rewrite q-defer-PTx; (simp add: that)?*)
unfolding *cond-prob-def* **by** *simp*
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0)$

proof –

have $\bigwedge \xi. \xi \in \text{space } \mathfrak{M} \implies f < T x \xi \wedge T x \xi \leq f + t \wedge T x \xi > 0 \longleftrightarrow f < T x \xi \wedge T x \xi \leq f + t$
using *that* **by** *auto*

```

hence  $\{\xi \in space \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t \wedge T x \xi > 0\} =$ 
       $\{\xi \in space \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t\}$  by blast
thus ?thesis by simp
qed
also have ... =  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. x + f < X \xi \wedge X \xi \leq x + f + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi$ 
> x)
  unfolding futr-life-def by (smt (verit) Collect-cong)
  finally show ?thesis .
qed

```

lemma q-defer-old-0: $\$q\{f|t\&x\} = 0$ if $x+f \geq \$\psi t \geq 0$ for $f t :: real$

proof -

have $\$q\{f|t\&x\} = \$q\{f+t \& x\} - \$q\{f\&x\}$ using q-defer-q that by simp

moreover have $\$q\{f+t \& x\} = 1$ using q-1-equiv that le-ereal-le by auto

moreover have $\$q\{f\&x\} = 1$ using q-1-equiv that by simp

ultimately show ?thesis by simp

qed

end

5.1.10 Properties of Cumulative Distributive Function for X

```

lemma cdfX-continuous-unborn[simp]: continuous-on {..0} (cdf (distr  $\mathfrak{M}$  borel  $X$ ))
  using cdfTx-continuous-on-nonpos by (metis cdfT0-eq-cdfX psi-pos')

lemma cdfX-differentiable-unborn[simp]: (cdf (distr  $\mathfrak{M}$  borel  $X$ )) differentiable-on
{..0}
  using cdfTx-differentiable-on-nonpos by (metis cdfT0-eq-cdfX psi-pos')

lemma cdfX-has-real-derivative-0-unborn:
  (cdf (distr  $\mathfrak{M}$  borel  $X$ ) has-real-derivative 0) (at x) if  $x < 0$  for  $x::real$ 
  using cdfTx-has-real-derivative-0-at-neg by (metis cdfT0-eq-cdfX psi-pos' that)

lemma cdfX-integrable-Icc:
  set-integrable lborel {a..b} (cdf (distr  $\mathfrak{M}$  borel  $X$ )) for  $a b :: real$ 
  using cdfTx-integrable-Icc by (metis cdfT0-eq-cdfX psi-pos')

corollary cdfX-integrable-on-Icc:
  cdf (distr  $\mathfrak{M}$  borel  $X$ ) integrable-on {a..b} for  $a b :: real$ 
  using cdfX-integrable-Icc set-borel-integral-eq-integral by force

lemma cdfX-q: cdf (distr  $\mathfrak{M}$  borel  $X$ )  $x = \$q\{x\&0\}$  if  $x \geq 0$  for  $x::real$ 
  by (metis cdfT0-eq-cdfX die-def psi-pos')

```

5.1.11 Relations between $\$p\{t\&x\}$ and $\$q\{t\&x\}$

```

context
  fixes  $x::real$ 
  assumes x-lt-psi[simp]:  $x < \$\psi$ 
begin

```

interpretation *alive*_{PS}: prob-space $\mathfrak{M} \downharpoonright \text{alive } x$
by (rule MM-PS.cond-prob-space-correct, simp-all add: alive-def)

interpretation *distrTx-RD*: real-distribution *distr* ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$) by simp

lemma *p-q-1*: $\$p\{-t \& x\} + \$q\{-t \& x\} = 1$ **for** $t :: \text{real}$
unfolding *survive-def die-def* **using** *distrTx-RD.add-cdf-ccdf*
by (smt (verit) *distrTx-RD.prob-space x-lt-psi*)

lemma *q-defer-p*: $\$q\{-f \mid t \& x\} = \$p\{-f \& x\} - \$p\{-f+t \& x\}$ **if** $t \geq 0$ **for** $f t :: \text{real}$
using *q-defer-q p-q-1* **that** *x-lt-psi* **by** (smt (verit))

lemma *q-defer-p-q-defer*: $\$p\{-f \& x\} * \$q\{-f' \mid t \& x+f\} = \$q\{-f+f' \mid t \& x\}$
if $x+f < \$\psi f \geq 0 f' \geq 0 t \geq 0$ **for** $f f' t :: \text{real}$

proof –

have $\$p\{-f \& x\} * \$q\{-f' \mid t \& x+f\} =$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+f) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) *$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. x+f+f' < X \xi \wedge X \xi \leq x+f+f'+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+f)$
apply (rewrite *p-PX*, (simp-all add: that)[2])
by (rewrite *survival-model.q-defer-PX*, simp-all add: that *survival-model-axioms*)
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x+f+f' < X \xi \wedge X \xi \leq x+f+f'+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
using *survival-model.PXx-pos*[of $\mathfrak{M} X x+f$] nonzero-mult-div-cancel-left that
by (smt (verit, ccfv-SIG) *survival-model-axioms times-divide-eq-left times-divide-eq-right*)
also have $\dots = \$q\{-f+f' \mid t \& x\}$
by (rewrite *q-defer-PX*; simp add: that group-cancel.add1)
finally show ?thesis .

qed

lemma *q-defer-pq*: $\$q\{-f \mid t \& x\} = \$p\{-f \& x\} * \$q\{-t \& x+f\}$
if $x+f < \$\psi t \geq 0 f \geq 0$ **for** $f t :: \text{real}$
using *q-defer-p-q-defer*[where $f'=0$] **that**
by (simp add: *survival-model.q-defer-0-q survival-model-axioms*)

5.1.12 Properties of Life Expectation

lemma *e-nonneg*: $\$e^{\circ}-x \geq 0$
unfolding *life-expect-def*
by (rule Bochner-Integration.integral-nonneg, simp add: less-eq-real-def)

lemma *e-P*: $\$e^{\circ}-x =$
MM-PS.expectation ($\lambda \xi. \text{indicator} (\text{alive } x) \xi * T x \xi$) / $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0)$
unfolding *life-expect-def*
by (rewrite *MM-PS.integral-cond-prob-space-nn*, auto simp add: alive-T)

proposition *nn-integral-T-p*:
 $(\int^+ \xi. \text{ennreal} (T x \xi) \partial(\mathfrak{M} \downharpoonright \text{alive } x)) = (\int^+ t \in \{0..\}. \text{ennreal} (\$p\{-t \& x\}) \partial \text{borel})$

```

apply (rewrite alivex-PS.expectation-nonneg-tail, simp-all add: less-imp-le)
apply (rule nn-integral-cong)
unfolding survive-def using distrTx-RD.prob-space distrTx-RD.ccdf-cdf by pres-
burger

lemma nn-integral-T-pos: ( $\int^+ \xi. \text{ennreal } (T x \xi) \partial(\mathfrak{M} \downharpoonright \text{alive } x)$ ) > 0
proof -
let ?f =  $\lambda t. - \text{ccdf } (\text{distr } (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) t$ 
have  $\bigwedge t u. t \leq u \implies ?f t \leq ?f u$  using distrTx-RD.ccdf-nonincreasing by simp
moreover have continuous (at-right 0) ?f
  using distrTx-RD.ccdf-is-right-cont by (intro continuous-intros)
ultimately have  $\forall e > 0. \exists d > 0. ?f (0 + d) - ?f 0 < e$ 
  using continuous-at-right-real-increasing by simp
hence  $\exists d > 0. ?f (0 + d) - ?f 0 < 1/2$  by (smt (verit, del-insts) field-sum-of-halves)
from this obtain d where d-pos:  $d > 0$  and  $\$p\{-d \& x\} \geq 1/2$ 
  using p-0-1 unfolding survive-def by auto
hence  $\bigwedge t. t \in \{0..d\} \implies \$p\{-t \& x\} \geq 1/2$ 
  unfolding survive-def using distrTx-RD.ccdf-nonincreasing by force
hence  $(\int^+ t \in \{0..d\}. \text{ennreal } (\$p\{-t \& x\}) \partial borel) \geq (\int^+ t \in \{0..d\}. \text{ennreal } (1/2) \partial borel)$ 
  apply (intro nn-set-integral-mono, simp-all)
  unfolding survive-def using Tx-alivex-measurable apply force
  by (rule AE-I2) (smt (verit) ennreal-half ennreal-leI half-bounded-equal)
  moreover have  $(\int^+ t \in \{0..\}. \text{ennreal } (\$p\{-t \& x\}) \partial borel) \geq (\int^+ t \in \{0..d\}. \text{ennreal } (\$p\{-t \& x\}) \partial borel)$ 
    by (rule nn-set-integral-set-mono) simp
  moreover have  $(\int^+ t \in \{0..d\}. \text{ennreal } (1/2) \partial borel) > 0$ 
    apply (rewrite nn-integral-cmult-indicator, simp-all)
    using d-pos emeasure-lborel-Icc ennreal-zero-less-mult-iff by fastforce
  ultimately show ?thesis using nn-integral-T-p by simp
qed

lemma e-pos-Tx:  $\$e^{\circ} \cdot x > 0$  if integrable ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) ( $T x$ )
  unfolding life-expect-def
  apply (rewrite integral-eq-nn-integral, simp-all)
  apply (smt (verit, ccfv-SIG) AE-I2 alivex-Tx-pos)
  using nn-integral-T-pos that
  by (smt (verit) AE-I2 alivex-Tx-pos enn2real-ennreal ennreal-less-zero-iff
    nn-integral-cong nn-integral-eq-integral)

proposition e-LBINT-p:  $\$e^{\circ} \cdot x = (\text{LBINT } t : \{0..\}. \$p\{-t \& x\})$ 
  — Note that  $0 = 0$  holds when the integral diverges.
  unfolding life-expect-def apply (rewrite integral-eq-nn-integral, simp-all add: less-imp-le)
  unfolding set-lebesgue-integral-def apply (rewrite integral-eq-nn-integral, simp-all)
  apply (measurable, simp add: survive-def)
  by (rewrite nn-integral-T-p) (simp add: indicator-mult-ennreal mult.commute)

corollary e-integral-p:  $\$e^{\circ} \cdot x = \text{integral } \{0..\} (\lambda t. \$p\{-t \& x\})$ 

```

— Note that $\theta = 0$ holds when the integral diverges.

proof —

```
have $e^{\circ}-x = (LBINT t:{0..}. \$p-\{t&x\}) using e-LBINT-p by simp
also have ... = integral {0..} (\lambda t. \$p-\{t&x\})
  apply (rule set-borel-integral-eq-integral-nonneg, simp-all)
  unfolding survive-def by simp
finally show ?thesis .
```

qed

lemma e-pos: $\$e^{\circ}-x > 0$ if set-integrable lborel {0..} ($\lambda t. \$p-\{t&x\}$)

proof —

```
have ( $\int^+ t \in \{0..\}. ennreal (\$p-\{t&x\}) \partial lborel$ ) = ennreal ( $\int t \in \{0..\}. \$p-\{t&x\} \partial borel$ )
  by (intro set-nn-integral-eq-set-integral; simp add: that)
also have ... <  $\infty$  using that by simp
finally have ( $\int^+ \xi. ennreal (T x \xi) \partial (\mathfrak{M} \downarrow alive x)$ ) <  $\infty$  using nn-integral-T-p
  by simp
hence integrable ( $\mathfrak{M} \downarrow alive x$ ) ( $T x$ )
  by (smt (verit) alivex-Tx-pos integrableI-bounded nn-integral-cong real-norm-def
       survival-model.Tx-alivex-measurable survival-model-axioms)
thus ?thesis by (rule e-pos-Tx)
qed
```

corollary e-pos': $\$e^{\circ}-x > 0$ if $(\lambda t. \$p-\{t&x\})$ integrable-on {0..}

apply (rule e-pos)

using that apply (rewrite integrable-on-iff-set-integrable-nonneg; simp)

unfolding survive-def by simp

lemma e-LBINT-p-Icc: $\$e^{\circ}-x = (LBINT t:\{0..n\}. \$p-\{t&x\})$ if $x+n \geq \$\psi$ for $n::real$

proof —

have [simp]: $\{0..n\} \cap \{n<..\} = \{\}$ using ivl-disj-int-one(7) by blast

have [simp]: $\{0..n\} \cup \{n<..\} = \{0..\}$

by (smt (verit) ereal-less-le ivl-disj-un-one(7) leD that x-lt-ps)

have [simp]: $\bigwedge t. n < t \implies 0 \leq t$ using that x-lt-ps by (smt (verit) ereal-less-le leD)

have [simp]: $\bigwedge t. n < t \implies \$\psi \leq ereal (x+t)$ using that by (simp add: le-ereal-le)

have gt-n-0: has-bochner-integral lborel ($\lambda t. indicator-real \{n<..\} t * \$p-\{t&x\}$) 0

apply (rewrite has-bochner-integral-cong[where N=lborel and g= $\lambda t. 0$ and y=0], simp-all)

using p-0-equiv that x-lt-ps

apply (smt (verit, ccfv-SIG) greaterThan-iff indicator-simps le-ereal-le linorder-not-le)

by (rule has-bochner-integral-zero)

hence gt-n: set-integrable lborel {n<..} ($\lambda t. \$p-\{t&x\}$)

unfolding set-integrable-def using integrable.simps by auto

moreover have le-n: set-integrable lborel {0..n} ($\lambda t. \$p-\{t&x\}$)

unfolding survive-def by (intro ccdfTx-integrable-Icc) simp

ultimately have set-integrable lborel ({0..n} \cup {n<..}) ($\lambda t. \$p-\{t&x\}$)

using set-integrable-Un by force

```

hence set-integrable lborel {0..} ( $\lambda t. \$p\{-t\&x\}$ ) by force
thus ?thesis
  apply (rewrite e-LBINT-p, simp)
  apply (rewrite set-integral-Un[of {0..n} {n<..}, simplified], simp-all add: gt-n
le-n)
  unfolding set-lebesgue-integral-def using gt-n-0 has-bochner-integral-integral-eq
by fastforce
qed

lemma e-integral-p-Icc: $e`o-x = integral {0..n} ( $\lambda t. \$p\{-t\&x\}$ ) if x+n ≥ $ψ for
n::real
  using that apply (rewrite e-LBINT-p-Icc, simp-all)
  using ccdfTx-integrable-Icc unfolding survive-def
  by (rewrite set-borel-integral-eq-integral; simp)

lemma temp-e-le-n: $e`o-{x:n} ≤ n if n ≥ 0 for n::real
proof -
  have nni-n: ( $\int^{+}_{-} ennreal n \partial(\mathfrak{M} \downharpoonright alive x)$ ) = ennreal n
    by (rewrite nn-integral-const, rewrite alivex-PS.emeasure-space-1) simp
  hence hbi-n: has-bochner-integral ( $\mathfrak{M} \downharpoonright alive x$ ) ( $\lambda_{-} n$ ) n
    by (intro has-bochner-integral-nn-integral; simp add: that)
  hence integrable ( $\mathfrak{M} \downharpoonright alive x$ ) ( $\lambda_{-} n$ ) by simp
  moreover have integrable ( $\mathfrak{M} \downharpoonright alive x$ ) ( $\lambda\xi. min (T x \xi) n$ )
  proof -
    have ( $\int^{+}_{-} \xi. ennreal (norm (min (T x \xi) n)) \partial(\mathfrak{M} \downharpoonright alive x)$ ) ≤  $\int^{+}_{-} ennreal n \partial(\mathfrak{M} \downharpoonright alive x)$ 
      apply (rule nn-integral-mono, rule ennreal-leI)
      apply (rewrite real-norm-def, rewrite abs-of-nonneg; simp add: that)
      by (smt (verit) alivex-Tx-pos)
    also have ... < ∞ using nni-n by simp
    finally have ( $\int^{+}_{-} \xi. ennreal (norm (min (T x \xi) n)) \partial(\mathfrak{M} \downharpoonright alive x)$ ) < ∞ .
    thus ?thesis by (intro integrableI-bounded; simp)
  qed
  ultimately have $e`o-{x:n} ≤ integralL ( $\mathfrak{M} \downharpoonright alive x$ ) ( $\lambda_{-} n$ )
    unfolding temp-life-expect-def by (intro integral-mono; simp)
  also have ... = n using hbi-n has-bochner-integral-iff by blast
  finally show ?thesis .
qed

lemma temp-e-P: $e`o-{x:n} =
  MM-PS.expectation ( $\lambda\xi. indicator (alive x) \xi * min (T x \xi) n$ ) /  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x$ 
 $\xi > 0)$ 
  if n ≥ 0 for n::real
  unfolding temp-life-expect-def
  by (rewrite MM-PS.integral-cond-prob-space-nn; simp add: alive-T that)

lemma temp-e-LBINT-p: $e`o-{x:n} = (LBINT t:{0..n}. $p{-t&x}) if n ≥ 0 for
n::real
proof -

```

```

let ?minTxn =  $\lambda \xi. \min(T x \xi) n$ 
let ?F = cdf (distr ( $\mathfrak{M} \downarrow \text{alive } x$ ) borel ( $T x$ ))
let ?Fn = cdf (distr ( $\mathfrak{M} \downarrow \text{alive } x$ ) borel ?minTxn)
interpret distrTxn-RD: real-distribution distr ( $\mathfrak{M} \downarrow \text{alive } x$ ) borel ?minTxn by
(simp add: that)
  have [simp]:  $\forall \xi. \xi \in \text{space}(\mathfrak{M} \downarrow \text{alive } x) \implies 0 \leq T x \xi$  by (smt (verit)
alivex-Tx-pos)
  have  $(\int^+ \xi. \text{ennreal}(\min(T x \xi) n) \partial(\mathfrak{M} \downarrow \text{alive } x)) = (\int^+ t \in \{0..\}. \text{ennreal}(1 - ?F t) \partial borel)$ 
    by (rewrite alivex-PS.expectation-nonneg-tail; simp add: that)
  also have ... =  $(\int^+ t \in \{0..\}. (\text{ennreal}(1 - ?F t) * \text{indicator}\{\dots < n\} t) \partial borel)$ 
    apply (rule nn-integral-cong)
  by (rewrite alivex-PS.cdf-distr-min; simp add: indicator-mult-ennreal mult.commute)
  also have ... =  $(\int^+ t \in \{0..<n\}. \text{ennreal}(1 - ?F t) \partial borel)$ 
    apply (rule nn-integral-cong) using nn-integral-set-ennreal
    by (smt (verit, best) Int-def atLeastLessThan-def ennreal-mult-right-cong
      indicator-simps mem-Collect-eq mult.commute mult-1)
  also have ... =  $(\int^+ t \in \{0..n\}. \text{ennreal}(1 - ?F t) \partial borel)$ 
proof -
  have sym-diff  $\{0..<n\} \{0..n\} = \{n\}$  using that by force
  thus ?thesis by (intro nn-integral-null-delta; force)
qed
also have ... = ennreal (LBINT t:{0..n}. $p-{t&x})
proof -
  have set-integrable lborel  $\{0..n\} (\lambda t. \$p-\{t&x\})$ 
    unfolding survive-def by (intro ccdfTx-integrable-Icc) simp
  thus ?thesis
    unfolding set-lebesgue-integral-def unfolding set-integrable-def
    apply (rewrite nn-integral-eq-integral[THEN sym]; simp)
    apply (rule nn-integral-cong, simp)
    unfolding survive-def using distrTx-RD.ccdf-cdf distrTx-RD.prob-space
nn-integral-set-ennreal
    by (simp add: indicator-mult-ennreal mult.commute)
  qed
  finally have  $(\int^+ \xi. \text{ennreal}(\min(T x \xi) n) \partial(\mathfrak{M} \downarrow \text{alive } x)) =$ 
    ennreal (LBINT t:{0..n}. $p-{t&x}) .
  thus ?thesis
    unfolding temp-life-expect-def by (rewrite integral-eq-nn-integral; simp add:
that)
  qed

lemma temp-e-integral-p:  $\$e^{\circ}-\{x:n\} = \text{integral}\{0..n\} (\lambda t. \$p-\{t&x\})$  if  $n \geq 0$ 
for  $n::\text{real}$ 
using that apply (rewrite temp-e-LBINT-p, simp-all)
using ccdfTx-integrable-Icc unfolding survive-def
by (rewrite set-borel-integral-eq-integral; simp)

lemma e-eq-temp:  $\$e^{\circ}-x = \$e^{\circ}-\{x:n\}$  if  $n \geq 0$   $x+n \geq \$\psi$  for  $n::\text{real}$ 
using that e-LBINT-p-Icc temp-e-LBINT-p by simp

```

```

lemma curt-e-P: $e-x =
  MM-PS.expectation (λξ. indicator (alive x) ξ * ⌊T x ξ⌋) / P(ξ in M. T x ξ > 0)
unfolding curt-life-expect-def
apply (rewrite MM-PS.integral-cond-prob-space-nn; simp add: alive-T)
  by (metis (no-types, lifting) Bochner-Integration.integral-cong indicator-simps
of-int-0 of-int-1)

lemma curt-e-sum-P: $e-x = (∑ k. P(ξ in M. T x ξ ≥ k + 1 | T x ξ > 0))
  if summable (λk. P(ξ in M. T x ξ ≥ k + 1 | T x ξ > 0))

proof -
  let ?F-flrTx = cdf (distr (M ⊢ alive x) borel (λξ. ⌊T x ξ⌋))
  have [simp]: ∀ξ. ξ ∈ space (M ⊢ alive x) ⇒ 0 ≤ T x ξ by (smt (verit)
alivex-Tx-pos)
  have integralN (M ⊢ alive x) (λξ. ennreal ⌊T x ξ⌋) =
    (∫+t∈{0..}. ennreal (1 - ?F-flrTx t) ∂lborel)
    by (rewrite alivex-PS.expectation-nonneg-tail; simp)
  also have ... = (∫+t∈{0::real..}. ennreal P(ξ in M. T x ξ ≥ ⌊t⌋ + 1 | T x ξ
> 0) ∂lborel)
  proof -
    { fix t::real assume t ≥ 0
      hence 1 - ?F-flrTx t = P(ξ in M. T x ξ ≥ real-of-int ⌊t⌋ + 1 | T x ξ > 0)
      proof -
        have 1 - ?F-flrTx t = 1 - P(ξ in (M ⊢ alive x). T x ξ < real-of-int ⌊t⌋
+ 1)
          by (rewrite alivex-PS.cdf-distr-floor-P; simp)
        also have ... = 1 - P(ξ in M. T x ξ < real-of-int ⌊t⌋ + 1 | T x ξ > 0)
          using alive-T by (rewrite MM-PS.cond-prob-space-cond-prob; simp)
        also have ... = P(ξ in M. T x ξ ≥ real-of-int ⌊t⌋ + 1 | T x ξ > 0)
          by (rewrite not-le[THEN sym], rewrite MM-PS.cond-prob-neg; simp)
        finally show ?thesis .
    }
    qed
  thus ?thesis
    apply -
    by (rule nn-set-integral-cong2, rule AE-I2) simp
  qed
  also have ... = (∑ k. ∫+t∈{k..<k+1}. ennreal P(ξ in M. T x ξ ≥ ⌊t⌋ + 1 |
T x ξ > 0) ∂lborel)
  apply (rewrite nn-integral-disjoint-family[THEN sym]; simp)
    apply (rewrite add.commute, rule Ico-nat-disjoint)
    by (rewrite Ico-nat-union[THEN sym], simp add: add.commute)
  also have ... = (∑ k. ∫+t∈{k..<k+1::nat}. ennreal P(ξ in M. T x ξ ≥ k + 1
| T x ξ > 0) ∂lborel)
  proof -
    { fix k::nat and t::real
      assume real k ≤ t and t < k + 1 + real k
      hence real-of-int ⌊t⌋ = real k
        by (metis add.commute floor-eq2 of-int-of-nat-eq)
    }

```

```

hence  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq \text{real-of-int } \lfloor t \rfloor + 1 \mid T x \xi > 0) =$ 
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq 1 + \text{real } k \mid T x \xi > 0)$ 
  by (simp add: add.commute) }
thus ?thesis
apply -
apply (rule suminf-cong, rule nn-set-integral-cong2, rule AE-I2)
by (rule impI) simp
qed
also have ... = ( $\sum k. \text{ennreal } \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$ 
  by (rewrite nn-integral-cmult-indicator; simp add: add.commute)
also have ... = ennreal ( $\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$ 
  by (rewrite suminf-ennreal2; simp add: that)
finally have integralN ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) ( $\lambda \xi. \text{ennreal } \lfloor T x \xi \rfloor$ ) =
  ennreal ( $\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$ .
hence integralL ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) ( $\lambda \xi. \lfloor T x \xi \rfloor$ ) = ( $\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1$ 
   $\mid T x \xi > 0))$ 
  apply (rewrite integral-eq-nn-integral; simp)
  apply (rewrite enn2real-ennreal; simp add: add.commute)
  apply (rule suminf-nonneg; simp?)
  by (rewrite add.commute, simp add: that)
thus ?thesis unfolding curt-life-expect-def by (simp add: add.commute)
qed

lemma curt-e-sum-P-finite: $\e-x = ( $\sum k < n. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi >$ 
   $0))$ 
  if  $x+n+1 > \psi$  for  $n::nat$ 
proof -
  from that have psi-fin:  $\psi < \infty$  by force
  let ?P =  $\lambda k::nat. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$ 
  let ?P-fin =  $\lambda k::nat. \text{if } k \in \{.. < n\} \text{ then } ?P k \text{ else } 0$ 
  have  $\bigwedge k. ?P k = ?P\text{-fin } k$ 
  proof -
    fix k
    show ?P k = ?P-fin k
    proof (cases 'k ∈ {.. < n}')
      case True
      thus ?thesis by simp
    next
      case False
      hence  $\neg k < n$  by simp
      hence  $x + k + 1 > \text{real-of-ereal } \psi$ 
        using that psi-nonneg real-of-ereal-ord-simps(4) by fastforce
      hence  $\{\xi \in \text{space } \mathfrak{M}. T x \xi \geq k + 1 \wedge T x \xi > 0\} \subseteq \{\xi \in \text{space } \mathfrak{M}. X \xi >$ 
         $\text{real-of-ereal } \psi\}$ 
        unfolding futr-life-def using that less-ereal-le of-nat-1 of-nat-add by force
      hence  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \wedge T x \xi > 0) \leq \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > \text{real-of-ereal }$ 
         $\psi)$ 
        by (intro MM-PSFINITE-MEASURE-MONO, simp-all)
      also have ... = 0 using MM-PSCDF-DISTR-P_X-RV_CCDF_X-PSI-0_PSI-FIN by

```

presburger
finally have $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \wedge T x \xi > 0) = 0$ **using measure-le-0-iff**
by blast
hence $?P k = 0$ **unfolding cond-prob-def by** (*simp add: add.commute*)
thus $?thesis$ **by simp**
qed
qed
moreover have $?P\text{-fin sums } (\sum k < n. ?P k)$ **using sums-If-finite-set by force**
ultimately have $\star: ?P \text{ sums } (\sum k < n. ?P k)$ **using sums-cong by simp**
moreover hence summable $?P$ **using sums-summable by blast**
ultimately have $?P \text{ sums } \$e\text{-}x$ **using curt-e-sum-P by force**
hence $\$e\text{-}x = (\sum k < n. ?P k)$ **by** (*rewrite sums-unique2[of ?P]; simp add: \star*)
thus $?thesis$ **by** (*simp add: add.commute*)
qed
lemma *curt-e-sum-p*: $\$e\text{-}x = (\sum k. \$p\{-k+1\&x\})$
if *summable* $(\lambda k. \$p\{-k+1\&x\}) \wedge k::nat. isCont (\lambda t. \$p\{-t\&x\}) (k+1)$
proof –
have $\bigwedge k::nat. \$p\{-k+1\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$
apply (*rewrite p-PTx-ge-ccdf-isCont, simp-all*)
using *that(2)* *isCont-ccdfX-ccdfTx unfolding survive-def by simp*
thus $?thesis$ **using** *that p-PTx-ge-ccdf-isCont curt-e-sum-P by presburger*
qed
lemma *curt-e-rec*: $\$e\text{-}x = \$p\text{-}x * (1 + \$e\text{-}(x+1))$
if *summable* $(\lambda k. \$p\{-k+1\&x\}) \wedge k::nat. isCont (\lambda t. \$p\{-t\&x\}) (real k + 1) x+1 < \ψ
proof –
have *px-neq-0[simp]*: $\$p\text{-}x \neq 0$ **using** *p-0-equiv that by auto*
have $(\lambda k. \$p\{-k+1\&x\}) \text{ sums } \$e\text{-}x$
using *that apply (rewrite curt-e-sum-p, simp-all add: add.commute)*
by (*rule summable-sums, simp add: that*)
hence $(\lambda k. \$p\text{-}x * \$p\{-k\&x+1\}) \text{ sums } \$e\text{-}x$
apply (*rewrite sums-cong[where g=λk. \$p\{-k+1\&x\}]; simp?*)
using *p-mult by (smt (verit) of-nat-0-le-iff that(3) x-lt-psi)*
hence $(\lambda k. \$p\{-k\&x+1\}) \text{ sums } (\$e\text{-}x / \$p\text{-}x)$
using *sums-mult-D that by (smt (verit, best) linorder-not-le p-0-equiv sums-cong x-lt-psi)*
hence *p-e-p*: $(\lambda k. \$p\{-Suc k \& x+1\}) \text{ sums } (\$e\text{-}x / \$p\text{-}x - \$p\{-0\&x+1\})$
using *sums-split-initial-segment[where n=1] by force*
moreover have $(\lambda k. \$p\{-Suc k \& x+1\}) \text{ sums } \$e\text{-}(x+1)$
proof –
have [*simp*]: *summable* $(\lambda k::nat. \$p\{-real k + 1 \& x + 1\})$
apply (*intro sums-summable[where l=\$e\text{-}x / \$p\text{-}x - \$p\{-0\&x+1\}]*)
using *p-e-p by (simp add: add.commute)*
have [*simp*]: $\bigwedge k::nat. isCont (\lambda t. \$p\{-t\&x+1\}) (real k + 1)$
proof –
fix $k::nat$
have *isCont* $(\lambda t. \$p\text{-}x * \$p\{-t-1\&x+1\}) (real k + 2)$

```

proof -
  let ?S={real k + 1 <..< real k + 3}
  have open ?S by simp
  moreover have real k + 2 ∈ ?S by simp
  moreover have ∏t. t ∈ ?S ⇒ $p-x * $p-{t-1&x+1} = $p-{t&x}
    using p-mult
    by (smt (verit, del-insts) greaterThanLessThan-iff of-nat-0-le-iff that(3)
      x-lt-ps)
  ultimately show ?thesis
    apply (rewrite isCont-cong[where g=λt. $p-{t&x}])
    apply (rewrite eventually-nhds, blast)
    using that by (smt (verit) of-nat-1 of-nat-add)
  qed
  hence isCont (λt. $p-x * $p-{t-1&x+1} / $p-x) (real k + 2)
    by (intro isCont-divide[where g=λt. $p-x], auto)
  hence isCont ((λt. $p-{t-1&x+1}) o (λt. t+1)) (real k + 1)
    by simp (rule continuous-at-compose, simp-all add: add.commute)
  thus isCont (λt. $p-{t&x+1}) (real k + 1) unfolding comp-def by simp
  qed
  show ?thesis
    apply (rewrite survival-model.curt-e-sum-p; simp add: survival-model-axioms
      that)
    using summable-sums by (rewrite add.commute) force
  qed
  ultimately have $e-x / $p-x - $p-{0&x+1} = $e-(x+1) by (rule sums-unique2)
  thus ?thesis
    using p-0-1 that
    by (smt (verit) px-neq-0 divide-mult-cancel mult.commute mult-cancel-left2
      p-mult that(3))
  qed

lemma curt-e-sum-p-finite: $e-x = (∑ k< n. $p-{k+1&x})
  if ∏k::nat. k < n ⇒ isCont (λt. $p-{t&x}) (real k + 1) x+n+1 > $ψ for
  n::nat
proof -
  have ∏k::nat. k < n ⇒ $p-{k+1&x} = P(ξ in M. T x ξ ≥ k + 1 | T x ξ >
  0)
    apply (rewrite p-PTx-ge-ccdf-isCont, simp-all)
    using that isCont-ccdfX-ccdfTx unfolding survive-def by (smt (verit) of-nat-0-le-iff
      x-lt-ps)
  thus ?thesis using that p-PTx-ge-ccdf-isCont curt-e-sum-P-finite by simp
  qed

lemma temp-curt-e-P: $e-{x:n} =
  MM-PS.expectation (λξ. indicator (alive x) ξ * ⌊ min (T x ξ) n ⌋) / P(ξ in M. T
  x ξ > 0)
  if n ≥ 0 for n::real
  unfolding temp-curt-life-expect-def
  apply (rewrite MM-PS.integral-cond-prob-space-nn; simp add: alive-T that)

```

```

apply (rule disjI2, rule Bochner-Integration.integral-cong; simp)
using indicator-simps of-int-0 of-int-1 by (smt (verit))

lemma temp-curt-e-sum-P: $e-{x:n} = ( $\sum k < n. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$ ) for n::nat
proof -
  let ?F-flrminTx = cdf (distr ( $\mathfrak{M} \downarrow \text{alive } x$ ) borel ( $\lambda \xi. \lfloor \min(T x \xi) n \rfloor$ ))
  have [simp]:  $\forall \xi. \xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x) \implies 0 \leq T x \xi$  by (smt (verit) alivex-Tx-pos)
  have integralN ( $\mathfrak{M} \downarrow \text{alive } x$ ) ( $\lambda \xi. \text{ennreal } \lfloor \min(T x \xi) n \rfloor$ ) =
    ( $\int^+ t \in \{0..\}. \text{ennreal } (1 - ?F\text{-flrminTx } t) \partial borel$ )
    by (rewrite alivex-PS.expectation-nonneg-tail; simp)
  also have ... = ( $\int^+ t \in \{0::real..\}. \text{ennreal } ((1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) * \text{of-bool } (\lfloor t \rfloor + 1 \leq n)) \partial borel$ )
  proof -
    have  $\forall t. t \geq 0 \implies \text{ennreal } (1 - ?F\text{-flrminTx } t) =$ 
       $\text{ennreal } ((1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) * \text{of-bool } (\lfloor t \rfloor + 1 \leq n))$ 
    proof -
      fix t::real assume t ≥ 0
      thus ennreal (1 - ?F-flrminTx t) =
        ennreal ((1 - P(ξ in (M ⊢ alive x). T x ξ < ⌊t⌋ + 1)) * of-bool (⌊t⌋ + 1 ≤ n))
      proof (cases ⌊t⌋ + 1 ≤ n)
        case True
        thus ?thesis
          apply (rewrite alivex-PS.cdf-distr-floor-P; simp)
          using min-less-iff-disj
          by (smt (verit, ccfv-SIG) Collect-cong floor-eq floor-less-cancel floor-of-nat of-int-floor-le)
      qed
      qed
      thus ?thesis
        apply (rewrite alivex-PS.cdf-distr-floor-P; simp)
        using min-less-iff-disj
        by (smt (verit, ccfv-SIG) Collect-cong MM-PS.space-cond-prob-space alive-T alive-event
          alivex-PS.prob-space mem-Collect-eq of-int-of-nat-eq of-nat-less-of-int-iff)
      qed
      qed
      thus ?thesis
        by (intro nn-set-integral-cong2, intro AE-I2) auto
      qed
      also have ... = ( $\int^+ t \in \{0..n\}. \text{ennreal } (1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) \partial borel$ )
      proof -
        { fix t::real
          have (⌊t⌋ + 1 ≤ n) = (t < n) by linarith
          hence  $\bigwedge r::real.$ 
        }
      qed
    qed
  qed
next
  case False
  thus ?thesis
    apply (rewrite alivex-PS.cdf-distr-floor-P; simp)
    using min-less-iff-disj
    by (smt (verit, ccfv-SIG) Collect-cong MM-PS.space-cond-prob-space alive-T alive-event
      alivex-PS.prob-space mem-Collect-eq of-int-of-nat-eq of-nat-less-of-int-iff)
  qed
qed

```

```

ennreal (r * of_bool (|t| + 1 ≤ n)) * indicator {0..} t = ennreal r * indicator
{0..} t
  unfolding atLeastLessThan-def by (simp add: indicator-def) }
  thus ?thesis by simp
qed
also have ... = (ʃ+t∈{0..}. ennreal P(ξ in (ℳ ⊥ alive x). T x ξ ≥ |t| +
1) ∂lborel)
  by (rewrite alivex-PS.prob-neg[THEN sym]; simp add: not-less)
also have ... = (∑ k<n. ∫+t∈{k..}. ennreal P(ξ in (ℳ ⊥ alive x). T x ξ
≥ |t| + 1) ∂lborel)
  apply (rewrite Ico-nat-union-finite[of n, THEN sym])
  apply (rewrite nn-integral-disjoint-family-on-finite; simp add: add.commute)
  apply (rule disjoint-family-on-mono[of - {0..}]; simp)
  by (rewrite add.commute, rule Ico-nat-disjoint)
also have ... = (∑ k<n. ennreal P(ξ in ℳ. T x ξ ≥ k + 1 | T x ξ > 0))
proof -
  { fix k::nat
    assume k < n
    hence (ʃ+t∈{k..<(1 + real k)}. ennreal P(ξ in (ℳ ⊥ alive x). T x ξ ≥
real-of-int |t| + 1) ∂lborel) =
      P(ξ in ℳ. T x ξ ≥ 1 + real k | T x ξ > 0) (is ?LHS = ?RHS)
    proof -
      have ?LHS = (ʃ+t∈{k..<(1 + real k)}. ennreal P(ξ in (ℳ ⊥ alive x). T x
ξ ≥ k + 1) ∂lborel)
      proof -
        { fix t::real
          assume k ≤ t t < 1 + real k
          hence real-of-int |t| = real k by (smt (verit) floor-eq2 of-int-of-nat-eq)
          hence P(ξ in (ℳ ⊥ alive x). T x ξ ≥ real-of-int |t| + 1) =
            P(ξ in (ℳ ⊥ alive x). T x ξ ≥ 1 + real k)
            by (simp add: add.commute) }
        thus ?thesis by (intro nn-set-integral-cong2, intro AE-I2) auto
      qed
      also have ... = ennreal P(ξ in (ℳ ⊥ alive x). T x ξ ≥ k + 1)
        by (rewrite nn-integral-cmult-indicator; simp)
      also have ... = ?RHS
        by (rewrite MM-PS.cond-prob-space-cond-prob, simp-all add: alive-T)
      finally show ?thesis .
    qed
    thus ?thesis by (intro sum.cong) auto
  qed
  also have ... = ennreal (∑ k<n. P(ξ in ℳ. T x ξ ≥ k + 1 | T x ξ > 0)) by
  simp
  finally have integralN (ℳ ⊥ alive x) (λξ. ennreal [min (T x ξ) n]) =
    ennreal (∑ k<n. P(ξ in ℳ. T x ξ ≥ k + 1 | T x ξ > 0)) .
  hence integralL (ℳ ⊥ alive x) (λξ. [min (T x ξ) n]) =
    (∑ k<n. P(ξ in ℳ. T x ξ ≥ k + 1 | T x ξ > 0))
  apply (intro nn-integral-eq-integrable[THEN iffD1, THEN conjunct2]; simp)
  using MM-PS.cond-prob-nonneg by (meson sum-nonneg)

```

thus *?thesis unfolding temp-curt-life-expect-def by simp*
qed

corollary *curt-e-eq-temp: \$e-x = \$e-{x:n} if $x+n+1 > \psi$ for $n:\text{nat}$*
using curt-e-sum-P-finite temp-curt-e-sum-P that by simp

lemma *temp-curt-e-sum-p: \$e-{x:n} = (\sum k < n. \\$p-\{k+1&x\})*
if $\bigwedge k:\text{nat}. k < n \implies \text{isCont } (\lambda t. \$p-\{t&x\})$ (real $k + 1$) for $n:\text{nat}$
proof –
have $\bigwedge k:\text{nat}. k < n \implies \$p-\{k+1&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$
apply (rewrite p-PTx-ge-ccdf-isCont, simp-all)
using that isCont-ccdfX-ccdfTx unfolding survive-def by (smt (verit) of-nat-0-le-iff x-lt-psi)
thus *?thesis*
apply (rewrite temp-curt-e-sum-P)
by (rule sum.cong) simp-all
qed

lemma *temp-curt-e-rec: \$e-{x:n} = \$p-x * (1 + \$e-{x+1:n-1})*
if $\bigwedge k:\text{nat}. k < n \implies \text{isCont } (\lambda t. \$p-\{t&x\})$ (real $k + 1$) $x+1 < \psi$ $n \neq 0$ for $n:\text{nat}$
proof –
have [simp]: $\$p-x \neq 0$ using p-0-equiv that by auto
have [simp]: $\bigwedge k. k < n-1 \implies \text{isCont } (\lambda t. \$p-\{t&x+1\})$ (real $k + 1$)
proof –
fix k:nat assume k-n: $k < n-1$
*have isCont (λt. \$p-x * \$p-{t-1&x+1}) (real k + 2)*
proof –
let ?S={real k + 1 <..< real k + 3}
have open ?S by simp
moreover have real k + 2 ∈ ?S by simp
*moreover have $\bigwedge t. t \in ?S \implies \$p-x * \$p-\{t-1&x+1\} = \$p-\{t&x\}$*
using p-mult
by (smt (verit, del-insts) greaterThanLessThan-iff of-nat-0-le-iff that(2))
x-lt-psi)
ultimately show *?thesis*
apply (rewrite isCont-cong[where g=λt. \$p-\{t&x\}])
apply (rewrite eventually-nhds, blast)
using that k-n by (smt (verit) less-diff-conv of-nat-1 of-nat-add)
qed
*hence isCont (λt. \$p-x * \$p-{t-1&x+1} / \$p-x) (real k + 2)*
by (intro isCont-divide[where g=λt. \$p-x], auto)
hence isCont ((λt. \$p-\{t-1&x+1\}) ∘ (λt. t+1)) (real k + 1)
by simp (rule continuous-at-compose, simp-all add: add.commute)
thus isCont (λt. \$p-\{t&x+1\}) (real k + 1) unfolding comp-def by simp
qed
*have $\$p-x * (1 + \$e-{x+1:n-1}) = \$p-x * (1 + (\sum k < (n-1). \$p-\{k+1&x+1\}))$*
by (rewrite survival-model.temp-curt-e-sum-p; simp add: survival-model-axioms)

that)

also have ... = \$p-x + ($\sum k < (n-1)$. \$p-x * \$p-{k+1&x+1})
apply (rewrite distrib-left, simp)
by (rewrite vector-space-over-itself.scale-sum-right, simp)
also have ... = \$p-x + ($\sum k < (n-1)$. \$p-{k+2&x})
apply (rewrite sum.cong, simp-all add: add.commute)
using p-mult by (smt (verit) of-nat-0-le-iff that(2) x-lt-psi)
also have ... = ($\sum k < \text{Suc}(n-1)$. \$p-{k+1&x})
apply (rewrite lessThan-atLeast0)+
by (rewrite sum.atLeast0-lessThan-Suc-shift) simp
also have ... = \$e-{x:n} **using** that **by** (rewrite temp-curt-e-sum-p; simp)
finally show ?thesis **by** simp
qed

end

lemma p-set-integrable-shift:
set-integrable lborel {0..} (λt . \$p-{t&0}) \longleftrightarrow \text{set-integrable lborel } \{0..\} (\lambda t. \\$p-\{t&x\})

if $x < \psi$ **for** $x::real$

proof –

have set-integrable lborel {0..} (λt . \$p-{t&0}) \longleftrightarrow \text{set-integrable lborel } \{x..\} (\lambda t. \\$p-\{t&x\})

by (rule set-integrable-Ici-equiv)
(metis (no-types, lifting) ccdfX-integrable-Icc ccdfX-p set-integrable-cong)

also have ... \longleftrightarrow set-integrable lborel {0..} (λt . \$p-{x+t&0})

using set-integrable-Ici-shift[of x x] **by** force

also have ... \longleftrightarrow set-integrable lborel {0..} (λt . \$p-{x+t&0} / \\$p-\{x&0\})

using that p-0-equiv **by** (rewrite set-integrable-mult-divide-iff; simp)

also have ... \longleftrightarrow set-integrable lborel {0..} (λt . \$p-\{t&x\})

by (rule set-integrable-cong; simp) (simp add: ccdfTx-ccdfX ccdfX-p survive-def that)

finally show ?thesis .
qed

lemma e-p-e: \$e^{\circ}-x = \$e^{\circ}-{x:n} + \$p-\{n&x\} * \$e^{\circ}-(x+n)
if set-integrable lborel {0..} (λt . \$p-\{t&x\}) n \geq 0 x+n < \psi **for** $x n :: real$

proof –

have [simp]: ereal $x < \psi$ **using** that **by** (simp add: ereal-less-le)
hence \$e^{\circ}-x = (LBINT t:{0..}. \$p-\{t&x\}) **by** (simp add: e-LBINT-p)
also have ... = (LBINT t:{0..n}. \$p-\{t&x\}) + (LBINT t:{n..}. \$p-\{t&x\})

proof –

have AE t in lborel. $\neg(t \in \{0..n\} \wedge t \in \{n..\})$ **using** AE-lborel-singleton **by** force
moreover have {0..} = {0..n} \cup {n..} **using** that **by** auto
moreover have set-integrable lborel {0..n} (λt . \$p-\{t&x\})
using that
by (metis Icc-subset-Ici-iff atLeastAtMost-borel order.refl set-integrable-subset sets-lborel)
moreover have set-integrable lborel {n..} (λt . \$p-\{t&x\})

```

using that by (metis atLeast-borel atLeast-subset-iff set-integrable-subset
sets-lborel)
ultimately show ?thesis
  using set-integral-Un-AE
  by (metis (no-types, lifting) AE-cong atLeastAtMost-borel atLeast-borel sets-lborel)
qed
also have ... = (LBINT t:{0..n}. $p-{t&x}) + $p-{n&x} * (LBINT t:{0..}.
$p-{t & x+n})
proof -
  have (LBINT t:{n..}. $p-{t&x}) = (LBINT t:{0..}. $p-{n+t & x})
    using lborel-set-integral-Ici-shift[of n - n, simplified] by force
  also have ... = (LBINT t:{0..}. $p-{n&x} * $p-{t & x+n})
    apply (rule set-lebesgue-integral-cong; simp)
    using that p-mult by force
  finally show ?thesis by simp
qed
also have ... = $e^o-{x:n} + $p-{n&x} * $e^o-(x+n)
  apply (rewrite temp-e-LBINT-p, (simp-all add: that)[2])
  by (rewrite e-LBINT-p; simp add: that)
finally show ?thesis .
qed

proposition x-ex-mono: x + $e^o-x ≤ y + $e^o-y if x ≤ y y < $ψ for x y :: real
proof -
  have x-lt-psi[simp]: ereal x < $ψ using that ereal-less-le by simp
  show ?thesis
  proof (cases `set-integrable lborel {0..} (λt. $p-{t&x}))>
    case True
      hence $e^o-x = $e^o-{x:y-x} + $p-{y-x&x}*$e^o-y by (rewrite e-p-e[of x
y-x]; simp add: that)
      also have ... ≤ y - x + $e^o-y
    proof -
      have $e^o-{x:y-x} ≤ y - x using that by (intro temp-e-le-n; simp)
      moreover have $p-{y-x&x}*$e^o-y ≤ $e^o-y
        using p-le-1 x-lt-psi that
        by (smt (verit, ccfv-threshold) e-nonneg mult-less-cancel-right1)
      ultimately show ?thesis by simp
    qed
    finally show ?thesis by simp
  next
    case False
    hence $e^o-x = 0
      using e-LBINT-p not-integrable-integral-eq
      unfolding set-integrable-def set-lebesgue-integral-def
      by simp
    moreover have $e^o-y = 0
    proof -
      have ¬ set-integrable lborel {0..} (λt. $p-{t&y})
        using that False

```

```

apply (rewrite p-set-integrable-shift[THEN sym], simp)
  by (rewrite p-set-integrable-shift[of x]; simp)
thus ?thesis
  using e-LBINT-p not-integrable-integral-eq that
  unfolding set-integrable-def set-lebesgue-integral-def
  by simp
qed
ultimately show ?thesis using that by simp
qed
qed

proposition x-ex-const-equiv:  $x + \$e^{\circ}x = y + \$e^{\circ}y \longleftrightarrow \$q\{y-x&x\} = 0$ 
  if set-integrable lborel {0..} ( $\lambda t. \$p\{t&0\}$ )  $x \leq y$   $y < \$\psi$  for  $x y :: real$ 
proof -
  have ey: set-integrable lborel {0..} ( $\lambda t. \$p\{t&y\}$ ) using that p-set-integrable-shift
  by blast
  have x-lt-ps[i]simp]: ereal  $x < \$\psi$  using that ereal-less-le by simp
  hence ex: set-integrable lborel {0..} ( $\lambda t. \$p\{t&x\}$ ) using that p-set-integrable-shift
  by blast
  show ?thesis
  proof
    assume const:  $x + \$e^{\circ}x = y + \$e^{\circ}y$ 
    hence 0 =  $y - x - \$e^{\circ}x + \$e^{\circ}y$  by simp
    also have ... =  $y - x - \$e^{\circ}\{x:y-x\} - \$p\{y-x&x\} * \$e^{\circ}y + \$e^{\circ}y$ 
      using e-p-e[of x y-x] ex that by simp
    also have ... =  $(y - x - \$e^{\circ}\{x:y-x\}) + (1 - \$p\{y-x&x\}) * \$e^{\circ}y$ 
      by (simp add: left-diff-distrib)
    finally have 0 =  $(y - x - \$e^{\circ}\{x:y-x\}) + (1 - \$p\{y-x&x\}) * \$e^{\circ}y$  .
    moreover have  $y - x - \$e^{\circ}\{x:y-x\} \geq 0$  using temp-e-le-n that by simp
    ultimately have  $(1 - \$p\{y-x&x\}) * \$e^{\circ}y = 0$ 
      by (smt (verit, ccfv-threshold) e-nonneg mult-nonneg-nonneg p-le-1 that
        x-lt-ps[i])
    moreover have  $\$e^{\circ}y > 0$  using that e-pos ey by simp
    ultimately have  $1 - \$p\{y-x&x\} = 0$  by simp
    thus  $\$q\{y-x&x\} = 0$  by (smt (verit) p-q-1 x-lt-ps[i])
  next
    interpret alivex-PS: prob-space  $\mathfrak{M} \mid alive x$ 
      by (rule MM-PS.cond-prob-space-correct, simp-all add: alive-def)
    interpret distrTx-RD: real-distribution distr ( $\mathfrak{M} \mid alive x$ ) borel ( $T x$ ) by simp
    assume  $\$q\{y-x&x\} = 0$ 
    hence p1:  $\$p\{y-x&x\} = 1$  using p-q-1 by (metis add.right-neutral x-lt-ps[i])
    hence  $\bigwedge t. t \in \{0..y-x\} \implies \$p\{t&x\} = 1$ 
      unfolding survive-def using distrTx-RD.ccdf-nonincreasing
      by simp (smt (verit) distrTx-RD.ccdf-bounded-prob)
    hence  $\$e^{\circ}\{x:y-x\} = y - x$ 
      using that apply (rewrite temp-e-LBINT-p; simp)
      by (rewrite set-lebesgue-integral-cong[where g= $\lambda t. 1$ ]; simp)
    moreover have  $\$e^{\circ}x = \$e^{\circ}\{x:y-x\} + \$p\{y-x&x\} * \$e^{\circ}y$ 
      by (rewrite e-p-e[of x y-x]; simp add: that ex)

```

```

ultimately show $e^{\circ-x} = y + $e^{\circ-y} using p1 by simp
qed
qed
end

```

5.2 Piecewise Differentiable Survival Function

```

locale smooth-survival-function = survival-model +
assumes ccdfX-piecewise-differentiable[simp]:
  (ccdf (distr M borel X)) piecewise-differentiable-on UNIV
begin

```

```

interpretation distrX-RD: real-distribution distr M borel X
using MM-PS.real-distribution-distr by simp

```

5.2.1 Properties of Survival Function for X

```

lemma ccdfX-continuous[simp]: continuous-on UNIV (ccdf (distr M borel X))
  using ccdfX-piecewise-differentiable piecewise-differentiable-on-imp-continuous-on
by fastforce

```

```

corollary ccdfX-borel-measurable[measurable]: ccdf (distr M borel X) ∈ borel-measurable
borel
  by (rule borel-measurable-continuous-onI) simp

```

```

lemma ccdfX-nondifferentiable-finite-set[simp]:
  finite {x. ¬ ccdf (distr M borel X) differentiable at x}
proof -
  obtain S where
    fin: finite S and ∀x. x ∉ S ==> (ccdf (distr M borel X)) differentiable at x
    using ccdfX-piecewise-differentiable unfolding piecewise-differentiable-on-def
  by blast
  hence {x. ¬ ccdf (distr M borel X) differentiable at x} ⊆ S by blast
  thus ?thesis using finite-subset fin by blast
qed

```

```

lemma ccdfX-differentiable-open-set: open {x. ccdf (distr M borel X) differentiable
at x}
  using ccdfX-nondifferentiable-finite-set finite-imp-closed
  by (metis (mono-tags, lifting) Collect-cong open-Collect-neg)

```

```

lemma ccdfX-differentiable-borel-set[measurable, simp]:
  {x. ccdf (distr M borel X) differentiable at x} ∈ sets borel
  using ccdfX-differentiable-open-set by simp

```

```

lemma ccdfX-differentiable-AE:
  AE x in lborel. (ccdf (distr M borel X)) differentiable at x
  apply (rule AE-I'[of {x. ¬ ccdf (distr M borel X) differentiable at x}], simp-all)
  using ccdfX-nondifferentiable-finite-set by (simp add: finite-imp-null-set-lborel)

```

```

lemma deriv-ccdfX-measurable[measurable]: deriv (ccdf (distr M borel X)) ∈ borel-measurable
borel
proof –
  have set-borel-measurable borel UNIV (deriv (ccdf (distr M borel X)))
  by (rule piecewise-differentiable-on-deriv-measurable-real; simp)
  thus ?thesis unfolding set-borel-measurable-def by simp
qed

```

5.2.2 Properties of Cumulative Distributive Function for X

```

lemma cdfX-piecewise-differentiable[simp]:
  (cdf (distr M borel X)) piecewise-differentiable-on UNIV
  by (rewrite distrX-RD.cdf-ccdf) (rule piecewise-differentiable-diff; simp)

lemma cdfX-continuous[simp]: continuous-on UNIV (cdf (distr M borel X))
  using cdfX-piecewise-differentiable piecewise-differentiable-on-imp-continuous-on
  by fastforce

corollary cdfX-borel-measurable[measurable]: cdf (distr M borel X) ∈ borel-measurable
borel
  by (rule borel-measurable-continuous-onI) simp

lemma cdfX-nondifferentiable-finite-set[simp]:
  finite {x. ¬ cdf (distr M borel X) differentiable at x}
  using distrX-RD.differentiable-cdf-ccdf ccdFX-nondifferentiable-finite-set by simp

lemma cdfX-differentiable-open-set: open {x. cdf (distr M borel X) differentiable
at x}
  using distrX-RD.differentiable-cdf-ccdf ccdFX-differentiable-open-set by simp

lemma cdfX-differentiable-borel-set[measurable, simp]:
  {x. cdf (distr M borel X) differentiable at x} ∈ sets borel
  using cdfX-differentiable-open-set by simp

lemma cdfX-differentiable-AE:
  AE x in lborel. (cdf (distr M borel X)) differentiable at x
  using ccdFX-differentiable-AE distrX-RD.differentiable-cdf-ccdf AE-iffI by simp

lemma deriv-cdfX-measurable[measurable]: deriv (cdf (distr M borel X)) ∈ borel-measurable
borel
proof –
  have set-borel-measurable borel UNIV (deriv (cdf (distr M borel X)))
  by (rule piecewise-differentiable-on-deriv-measurable-real; simp)
  thus ?thesis unfolding set-borel-measurable-def by simp
qed

```

5.2.3 Introduction of Probability Density Functions of X and $T(x)$

definition $\text{pdf}X :: \text{real} \Rightarrow \text{real}$
where $\text{pdf}X x \equiv \text{if } \text{cdf}(\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } x$
 $\text{then deriv}(\text{cdf}(\text{distr } \mathfrak{M} \text{ borel } X)) x \text{ else } 0$
— This function is defined to be always nonnegative for future application.

definition $\text{pdf}T :: \text{real} \Rightarrow \text{real} \Rightarrow \text{real}$
where $\text{pdf}T x t \equiv \text{if } \text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) \text{ differentiable at } t$
 $\text{then deriv}(\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))) t \text{ else } 0$
— This function is defined to be always nonnegative for future application.

lemma $\text{pdf}X\text{-measurable}[measurable]: \text{pdf}X \in \text{borel-measurable borel}$
proof —
let $?cdfX = \text{cdf}(\text{distr } \mathfrak{M} \text{ borel } X)$
have $\text{countable } \{x. \neg ?cdfX \text{ differentiable at } x\}$
using $\text{cdf}X\text{-nondifferentiable-finite-set uncountable-infinite by force}$
thus $?thesis$
unfolding $\text{pdf}X\text{-def}$
apply —
by (*rule measurable-discrete-difference*
[**where** $X = \{x. \neg ?cdfX \text{ differentiable at } x\}$ **and** $f = \text{deriv } ?cdfX$; *simp*)
qed

lemma $\text{distributed-pdf}X: \text{distributed } \mathfrak{M} \text{ lborel } X \text{ pdf}X$
proof —
let $?cdfX = \text{cdf}(\text{distr } \mathfrak{M} \text{ borel } X)$
obtain S **where** $\text{fin}: \text{finite } S$ **and** $\text{diff}: \bigwedge t. t \notin S \implies ?cdfX \text{ differentiable at } t$
using $\text{cdf}X\text{-piecewise-differentiable unfolding piecewise-differentiable-on-def}$
by *blast*
{ **fix** $t::\text{real}$ **assume** $t\text{-notin}: t \notin S$
have $?cdfX \text{ differentiable at } t$ **using** $\text{diff } t\text{-notin by simp}$
hence $(?cdfX \text{ has-real-derivative pdf}X t) (\text{at } t)$
unfolding $\text{pdf}X\text{-def using DERIV-deriv-iff-real-differentiable by auto }$
thus $?thesis$
by (*intro MM-PS.distributed-deriv-cdf[where S=S]; simp add: fin*)
qed

lemma $\text{pdf}T0-X: \text{pdf}T 0 = \text{pdf}X$
unfolding $\text{pdf}T\text{-def pdf}X\text{-def using cdf}T0\text{-eq-cdf}X \text{ psi-pos' by fastforce}$

5.2.4 Properties of Survival Function for $T(x)$

context
fixes $x::\text{real}$
assumes $x\text{-lt-psi}[simp]: x < \ψ
begin

interpretation $\text{alive}x\text{-PS}: \text{prob-space } \mathfrak{M} \downarrow \text{alive } x$

by (rule MM-PS.cond-prob-space-correct; simp add: alive-def)

interpretation distrTx-RD: real-distribution distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$) by simp

lemma ccdfTx-continuous-on-nonneg[simp]:

continuous-on {0..} (ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)))

apply (rewrite continuous-on-eq-continuous-within, auto)

apply (rewrite continuous-ccdfX-ccdfTx[THEN sym], auto)

by (metis UNIV-I ccdfX-continuous continuous-at-imp-continuous-at-within continuous-on-eq-continuous-within)

lemma ccdfTx-continuous[simp]: continuous-on UNIV (ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)))

proof –

have [simp]: {.. $0::\text{real}$ } \cup {0..} = UNIV by auto

have continuous-on {.. 0 } (ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)))

by (rule ccdfTx-continuous-on-nonpos) simp

moreover have continuous-on {0..} (ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$))) by simp

ultimately have continuous-on ({.. 0 } \cup {0..}) (ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)))

by (intro continuous-on-closed-Un) simp-all

thus ?thesis by simp

qed

corollary ccdfTx-borel-measurable[measurable]:

ccdf (distr ($\mathfrak{M} \downharpoonright \text{alive } x$) borel ($T x$)) \in borel-measurable borel

by (rule borel-measurable-continuous-onI) simp

lemma ccdfTx-nondifferentiable-finite-set[simp]:

finite { $t. \neg \text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) \text{ differentiable at } t$ }

proof –

let ?P = $\lambda t. \text{ccdf} (\text{distr} (\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) \text{ differentiable at } t$

have { $t. t < 0 \wedge \neg ?P t$ } = {}

proof (rule equals0I)

fix t **assume** asm: $t \in \{t. t < 0 \wedge \neg ?P t\}$

hence ?P t **using** ccdfTx-has-real-derivative-0-at-neg real-differentiable-def x-lt-psi **by** blast

with asm **show** False **by** simp

qed

hence { $t. \neg ?P t$ } \subseteq insert 0 { $t. t > 0 \wedge \neg ?P t$ } **by** force

moreover have finite { $t. t > 0 \wedge \neg ?P t$ }

proof –

have { $t. \neg \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } (x+t)$ } =

plus (-x) ‘ { $s. \neg \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } s$ }’

by force

hence finite { $t. \neg \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } (x+t)$ }

using ccdfX-nondifferentiable-finite-set **by** simp

```

thus ?thesis
  using finite-subset differentiable-ccdfX-ccdfTx
  by (metis (no-types, lifting) mem-Collect-eq subsetI x-lt-psi)
qed
ultimately show ?thesis using finite-subset by auto
qed

lemma ccdfTx-differentiable-open-set:
open {t. ccdf (distr (M ` alive x) borel (T x)) differentiable at t}
using ccdfTx-nondifferentiable-finite-set finite-imp-closed
by (metis (mono-tags, lifting) Collect-cong open-Collect-neg)

lemma ccdfTx-differentiable-borel-set[measurable, simp]:
{t. ccdf (distr (M ` alive x) borel (T x)) differentiable at t} ∈ sets borel
using ccdfTx-differentiable-open-set by simp

lemma ccdfTx-differentiable-AE:
AE t in lborel. (ccdf (distr (M ` alive x) borel (T x))) differentiable at t
apply (rule AE-I'[of {t. ¬ ccdf (distr (M ` alive x) borel (T x)) differentiable at t}]; simp?)
using ccdfTx-nondifferentiable-finite-set by (simp add: finite-imp-null-set-lborel)

lemma ccdfTx-piecewise-differentiable[simp]:
(ccdf (distr (M ` alive x) borel (T x))) piecewise-differentiable-on UNIV
proof -
have ∀ t ∈ UNIV - {t. ¬ ccdf (distr (M ` alive x) borel (T x)) differentiable at t}.
  ccdf (distr (M ` alive x) borel (T x)) differentiable at t
  by auto
thus ?thesis
  unfolding piecewise-differentiable-on-def
  using ccdfTx-continuous ccdfTx-nondifferentiable-finite-set by blast
qed

lemma deriv-ccdfTx-measurable[measurable]:
deriv (ccdf (distr (M ` alive x) borel (T x))) ∈ borel-measurable borel
proof -
have set-borel-measurable borel UNIV (deriv (ccdf (distr (M ` alive x) borel (T x))))
  by (rule piecewise-differentiable-on-deriv-measurable-real; simp)
thus ?thesis unfolding set-borel-measurable-def by simp
qed

```

5.2.5 Properties of Cumulative Distributive Function for $T(x)$

```

lemma cdfTx-continuous[simp]:
continuous-on UNIV (cdf (distr (M ` alive x) borel (T x)))
using distrTx-RD.cdf-ccdf ccdfTx-continuous by (simp add: continuous-on-eq-continuous-within)

corollary cdfTx-borel-measurable[measurable]:

```

*cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$)) ∈ borel-measurable borel
by (rule borel-measurable-continuous-onI) simp*

lemma *cdfTx-nondifferentiable-finite-set*[simp]:
finite {t. ¬ cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$)) differentiable at t}
using *distrTx-RD.differentiable-cdf-ccdf* **by** simp

lemma *cdfTx-differentiable-open-set*:
open {t. cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$)) differentiable at t}
using *distrTx-RD.differentiable-cdf-ccdf ccdfTx-differentiable-open-set* **by** simp

lemma *cdfTx-differentiable-borel-set*[measurable, simp]:
{t. cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$)) differentiable at t} ∈ sets borel
using *cdfTx-differentiable-open-set* **by** simp

lemma *cdfTx-differentiable-AE*:
AE t in l borel. (cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$))) differentiable at t
by (rewrite *distrTx-RD.differentiable-cdf-ccdf*, simp add: *ccdfTx-differentiable-AE*)

lemma *cdfTx-piecewise-differentiable*[simp]:
(cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$))) piecewise-differentiable-on UNIV
using *piecewise-differentiable-diff piecewise-differentiable-const ccdfTx-piecewise-differentiable*
by (rewrite *distrTx-RD.cdf-ccdf*) blast

lemma *deriv-cdfTx-measurable*[measurable]:
deriv (cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$))) ∈ borel-measurable borel
proof –
have *set-borel-measurable borel UNIV (deriv (cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$))))*
by (rule *piecewise-differentiable-on-deriv-measurable-real*; simp)
thus ?thesis **unfolding** *set-borel-measurable-def* **by** simp
qed

5.2.6 Properties of Probability Density Function of $T(x)$

lemma *pdfTx-nonneg*: *pdfT x t ≥ 0* **for** *t::real*
proof –
fix *t*
show *pdfT x t ≥ 0*
proof (cases *<cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$)) differentiable at t>*)
case *True*
have *mono-on UNIV (cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$)))*
unfolding *mono-on-def* **using** *distrTx-RD.cdf-nondecreasing* **by** blast
moreover have *(cdf (distr ($\mathfrak{M} \downarrow \text{alive } x$) borel ($T x$))) has-real-derivative pdfT x t (at t)*
unfolding *pdfT-def* **using** *True DERIV-deriv-iff-real-differentiable* **by** presburger
ultimately show ?thesis **by** (rule *mono-on-imp-deriv-nonneg*) simp
next

```

case False
  thus ?thesis unfolding pdfT-def by simp
qed
qed

lemma pdfTx-neg-0: pdfT x t = 0 if t < 0 for t::real
proof -
  let ?e = -t/2
  have (cdf (distr (?M \ alive x) borel (T x)) has-real-derivative 0) (at t)
    apply (rewrite DERIV-cong-ev[of t t - λ-. 0 0 0], simp-all add: that)
    apply (rewrite eventually-nhds)
    apply (rule exI[of - ball t ?e], auto simp add: that)
    by (rule cdfTx-nonpos-0; simp add: dist-real-def)
  thus ?thesis unfolding pdfT-def by (meson DERIV-imp-deriv)
qed

lemma pdfTx-0-0: pdfT x 0 = 0
proof (cases ‹cdf (distr (?M \ alive x) borel (T x)) differentiable at 0›)
  case True
  let ?cdfx = cdf (distr (?M \ alive x) borel (T x))
  have (?cdfx has-real-derivative 0) (at 0)
  proof -
    from True obtain c where cdfx-deriv: (?cdfx has-real-derivative c) (at 0)
      unfolding differentiable-def using has-real-derivative by blast
    hence (?cdfx has-real-derivative c) (at 0 within {..0})
      by (rule has-field-derivative-at-within)
    moreover have (?cdfx has-real-derivative 0) (at 0 within {..0})
    proof -
      have ∀ F t in at 0 within {..0}. ?cdfx t = 0
      proof -
        { fix t assume t ∈ {..0::real} t ≠ 0 dist t 0 < 1
          hence ?cdfx t = 0 using cdfTx-nonpos-0 x-lt-ps by blast }
        hence ∃ d>0::real. ∀ t∈{..0}. t≠0 ∧ dist t 0 < d → ?cdfx t = 0 by (smt
        (verit))
        thus ?thesis by (rewrite eventually-at) simp
      qed
      moreover have ?cdfx 0 = 0
      proof -
        have continuous (at 0 within {..0}) ?cdfx
        using True differentiable-imp-continuous-within differentiable-subset by
        blast
        thus ?thesis by (simp add: cdfTx-nonpos-0)
      qed
      ultimately show ?thesis
        by (rewrite has-field-derivative-cong-eventually[of - λ-. 0::real 0 {..0} 0];
      simp)
    qed
    ultimately have c = 0
      using has-real-derivative-iff-has-vector-derivative

```

```

apply (intro vector-derivative-unique-within[of 0 {..0} ?cdfx c 0]; blast?)
by (rewrite at-within-eq-bot-iff)
  (metis closure-lessThan islimpt-in-closure limpt-of-closure
    trivial-limit-at-left-real trivial-limit-within)
thus (?cdfx has-real-derivative 0) (at 0) using cdfx-deriv by simp
qed
thus ?thesis unfolding pdfT-def by (meson DERIV-imp-deriv)
next
  case False
  thus ?thesis unfolding pdfT-def by simp
qed

lemma pdfTx-nonpos-0: pdfT x t = 0 if t ≤ 0 for t::real
using pdfTx-neg-0 pdfTx-0-0 that by force

lemma pdfTx-beyond-0: pdfT x t = 0 if x+t ≥ $ψ for t::real
proof (cases ‹t ≤ 0›)
  case True
  thus ?thesis using pdfTx-nonpos-0 by simp
next
  let ?cdfTx = cdf (distr (ℳ ⊢ alive x) borel (T x))
  case False
  hence t-pos: t > 0 by simp
  thus ?thesis
  proof -
    have (?cdfTx has-real-derivative 0) (at-right t)
    apply (rewrite has-field-derivative-cong-eventually[where g=λ_. 1], simp-all)
      apply (rewrite eventually-at-right-field)
    using that cdfTx-1-equiv
    by (intro exI[of ‹t+1›], auto simp add: le-ereal-less less-eq-ereal-def)
    thus ?thesis unfolding pdfT-def
    by (meson has-real-derivative-iff-has-vector-derivative has-vector-derivative-at-within
      DERIV-deriv-iff-real-differentiable trivial-limit-at-right-real
      vector-derivative-unique-within)
  qed
qed

lemma pdfTx-pdfX: pdfT x t = pdfX (x+t) / ℙ(ξ in ℳ. X ξ > x) if t > 0 for
t::real
proof (cases ‹cdf (distr ℳ borel X) differentiable at (x+t)›)
  case True
  let ?cdfX = cdf (distr ℳ borel X)
  let ?cdfTx = cdf (distr (ℳ ⊢ alive x) borel (T x))
  have [simp]: ?cdfTx differentiable at t using differentiable-cdfX-cdfTx that True
  by simp
  have pdfT x t = deriv ?cdfTx t unfolding pdfT-def using that differentiable-cdfX-cdfTx
  by simp
  hence (?cdfTx has-field-derivative (pdfT x t)) (at t)
  using True DERIV-deriv-iff-real-differentiable by simp

```

moreover have $\bigwedge u. \text{dist } u t < t \implies$
 $?cdfX(x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - 1) = ?cdfTx u$
proof –
fix $u::\text{real}$
assume $\text{dist } u t < t$
hence [simp]: $u > 0$ **using** *dist-real-def* **by** *fastforce*
have $?cdfX(x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - 1) =$
 $(1 - \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+u)) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - 1)$
using *MM-PS.ccdf-distr-P X-RV distrX-RD.cdf-ccdf distrX-RD.prob-space* **by** *presburger*
also have $\dots = 1 - \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
by (simp add: diff-divide-distrib)
also have $\dots = ?cdfTx u$
apply (rewrite *ccdfTx-PX[THEN sym]*, simp-all add: less-eq-real-def)
using *distrTx-RD.cdf-ccdf distrTx-RD.prob-space* **by** *presburger*
finally show $?cdfX(x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - 1) = ?cdfTx u$.
qed
ultimately have $((\lambda u. ?cdfX(x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) - 1))$
has-field-derivative (*pdfT x t*) (at *t*)
apply –
by (rule *has-field-derivative-transform-within[where d=t]*; simp add: that)
hence $((\lambda u. ?cdfX(x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ *has-field-derivative* (*pdfT x t*) (at *t*)
unfolding *has-field-derivative-def*
using *has-derivative-add-const[where c=1 / P(x in M. X xi > x) - 1]* **by** *force*
hence $((\lambda u. ?cdfX(x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$
has-field-derivative
 $\text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ (at *t*)
using *DERIV-cmult-right[where c=P(x in M. X xi > x)]* **by** *force*
hence $((\lambda u. ?cdfX(x+u))$ *has-field-derivative* $\text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ (at *t*)
unfolding *has-field-derivative-def* **using** *has-derivative-transform PXx-pos x-lt-ps*
by (smt (verit) Collect-cong UNIV-I nonzero-mult-div-cancel-right times-divide-eq-left)
hence $(?cdfX \text{ has-field-derivative } \text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ (at *(x+t)*)
using *DERIV-at-within-shift* **by** *force*
moreover have (?cdfX *has-field-derivative deriv* ?cdfX (*x+t*)) (at *(x+t)*)
using *True DERIV-deriv-iff-real-differentiable* **by** *blast*
ultimately have $\text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) = \text{deriv } ?cdfX(x+t)$
by (simp add: *DERIV-imp-deriv*)
thus $\text{pdfT } x \ t = \text{pdfX } (x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
unfolding *pdfX-def* **using** *True*
by simp (metis *PXx-pos nonzero-mult-div-cancel-right x-lt-ps zero-less-measure-iff*)
next
case *False*

hence [simp]: $\neg \text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))$ differentiable at t
using differentiable-cdfX-cdfTx that **by** simp
thus $\text{pdf}T x t = \text{pdf}X(x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ **unfolding** pdfT-def pdfX-def
using False **by** simp
qed

lemma pdfTx-measurable[measurable]: $\text{pdf}T x \in \text{borel-measurable borel}$
proof –
let ?cdfTx = $\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))$
have countable { $x. \neg ?\text{cdf}Tx$ differentiable at x }
using cdfX-nondifferentiable-finite-set uncountable-infinite **by** force
thus ?thesis
unfolding pdfT-def
apply –
by (rule measurable-discrete-difference
[where $X = \{x. \neg ?\text{cdf}Tx$ differentiable at $x\}$ and $f = \text{deriv } ?\text{cdf}Tx$]; simp)
qed

lemma distributed-pdfTx: distributed $(\mathfrak{M} \downarrow \text{alive } x) \text{ lborel } (T x) (\text{pdf}T x)$
proof –
let ?cdfTx = $\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))$
obtain S where fin: finite S and diff: $\bigwedge t. t \notin S \implies ?\text{cdf}Tx$ differentiable at t
using cdfTx-piecewise-differentiable **unfolding** piecewise-differentiable-on-def
by blast
{ **fix** $t::\text{real}$ **assume** t-notin: $t \notin S$
have ?cdfTx differentiable at t **using** diff t-notin **by** simp
hence (?cdfTx has-real-derivative pdfT x t) (at t)
unfolding pdfT-def **using** DERIV-deriv-iff-real-differentiable **by** force }
thus ?thesis
by (intro alivex-PS.distributed-deriv-cdf[where S=S]; simp add: fin)
qed

lemma nn-integral-pdfTx-1: $(\int^+ s. \text{pdf}T x s) \partial\text{lborel} = 1$
proof –
have $(\int^+ s. \text{pdf}T x s) \partial\text{lborel} = \text{emeasure}(\text{density lborel } (\text{pdf}T x)) \text{ UNIV}$
by (rewrite emeasure-density) simp-all
also have ... = $\text{emeasure}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ lborel } (T x)) \text{ UNIV}$
using distributed-pdfTx **unfolding** distributed-def **by** simp
also have ... = 1
by (metis distrTx-RD.emeasure-space-1 distrTx-RD.space-eq-univ distr-cong sets-lborel)
finally show ?thesis .
qed

corollary has-bochner-integral-pdfTx-1: has-bochner-integral lborel (pdfT x) 1
by (rule has-bochner-integral-nn-integral; simp add: pdfTx-nonneg nn-integral-pdfTx-1)
corollary LBINT-pdfTx-1: $(\text{LBINT } s. \text{pdf}T x s) = 1$
using has-bochner-integral-pdfTx-1 **by** (simp add: has-bochner-integral-integral-eq)

```

corollary pdfTx-has-integral-1: (pdfT x has-integral 1) UNIV
  by (rule nn-integral-has-integral; simp add: pdfTx-nonneg nn-integral-pdfTx-1)

lemma set-nn-integral-pdfTx-1: ( $\int^+ s \in \{0..\}. pdfT x s \partial borel$ ) = 1
  apply (rewrite nn-integral-pdfTx-1[THEN sym])
  apply (rule nn-integral-cong)
  using pdfTx-nonpos-0
  by (metis atLeast-iff ennreal-0 indicator-simps(1) linorder-le-cases
        mult.commute mult-1 mult-zero-left)

corollary has-bochner-integral-pdfTx-1-nonpos:
  has-bochner-integral lborel ( $\lambda s. pdfT x s * indicator \{0..\} s$ ) 1
  apply (rule has-bochner-integral-nn-integral, simp-all)
  using pdfTx-nonneg apply simp
  using set-nn-integral-pdfTx-1 by (simp add: nn-integral-set-ennreal)

corollary set-LBINT-pdfTx-1: (LBINT s:{0..}. pdfT x s) = 1
  unfolding set-lebesgue-integral-def
  using has-bochner-integral-pdfTx-1-nonpos has-bochner-integral-integral-eq
  apply (simp, rewrite mult.commute)
  by (smt (verit) Bochner-Integration.integral-cong has-bochner-integral-integral-eq)

corollary pdfTx-has-integral-1-nonpos: (pdfT x has-integral 1) {0..}
proof -
  have set-integrable lebesgue {0..} (pdfT x)
  unfolding set-integrable-def
  apply (rewrite integrable-completion, simp-all)
  using has-bochner-integral-pdfTx-1-nonpos by (rewrite mult.commute, rule integrable.intros)
  moreover have (LINT s:{0..}|lebesgue. pdfT x s) = 1
  using set-LBINT-pdfTx-1 unfolding set-lebesgue-integral-def
  by (rewrite integral-completion; simp)
  ultimately show ?thesis using has-integral-set-lebesgue by force
qed

lemma set-nn-integral-pdfTx-PTx: ( $\int^+ s \in A. pdfT x s \partial borel$ ) =  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \in A \mid T x \xi > 0)$ 
  if A ∈ sets lborel for A :: real set
proof -
  have [simp]: A ∈ sets borel using that by simp
  have ( $\int^+ s \in A. pdfT x s \partial borel$ ) = emeasure (density lborel (pdfT x)) A
  using that by (rewrite emeasure-density; force)
  also have ... = emeasure (distr ( $\mathfrak{M} \downarrow alive x$ ) lborel (T x)) A
  using distributed-pdfTx unfolding distributed-def by simp
  also have ... = ennreal  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow alive x). T x \xi \in A)$ 
  apply (rewrite emeasure-distr, simp-all)
  apply (rewrite finite-measure.emeasure-eq-measure, simp)
  by (smt (verit) Collect-cong vimage-eq Int-def)

```

```

also have ... = ennreal P(ξ in ℳ. T x ξ ∈ A | T x ξ > 0)
  unfolding alive-def
  apply (rewrite MM-PS.cond-prob-space-cond-prob[THEN sym], simp-all add:
pred-def futr-life-def)
  using borel-measurable-diff X-RV that by measurable
  finally show ?thesis .
qed

lemma pdfTx-set-integrable: set-integrable lborel A (pdfT x) if A ∈ sets lborel
  unfolding set-integrable-def
  using that pdfTx-nonneg apply (intro integrableI-nonneg, simp-all)
  apply (rewrite mult.commute)
  using set-nn-integral-pdfTx-PTx that
  by (metis (no-types, lifting) ennreal-indicator ennreal-less-top ennreal-mult' nn-integral-cong)

lemma set-integral-pdfTx-PTx: (LBINT s:A. pdfT x s) = P(ξ in ℳ. T x ξ ∈ A | T x ξ > 0)
  if A ∈ sets lborel for A :: real set
  unfolding set-lebesgue-integral-def
  apply (rewrite integral-eq-nn-integral)
  using that apply (simp-all add: pdfTx-nonneg)
  apply (rewrite indicator-mult-ennreal[THEN sym], rewrite mult.commute)
  by (rewrite set-nn-integral-pdfTx-PTx; simp)

lemma pdfTx-has-integral-PTx: (pdfT x has-integral P(ξ in ℳ. T x ξ ∈ A | T x ξ > 0)) A
  if A ∈ sets lborel for A :: real set
proof -
  have ((λs. pdfT x s * indicat-real A s) has-integral P(ξ in ℳ. T x ξ ∈ A | T x ξ > 0)) UNIV
    using that pdfTx-nonneg apply (intro nn-integral-has-integral, simp-all)
    using set-nn-integral-pdfTx-PTx that by (simp add: nn-integral-set-ennreal)
  thus ?thesis
    by (smt (verit) has-integral-cong has-integral-restrict-UNIV
        indicator-eq-0-iff indicator-simps(1) mult-cancel-left2 mult-eq-0-iff)
qed

corollary pdfTx-has-integral-PTx-Icc:
  (pdfT x has-integral P(ξ in ℳ. a ≤ T x ξ ∧ T x ξ ≤ b | T x ξ > 0)) {a..b} for
  a b :: real
  using pdfTx-has-integral-PTx[of {a..b}] by simp

corollary pdfTx-integrable-on-Icc: pdfT x integrable-on {a..b} for a b :: real
  using pdfTx-has-integral-PTx-Icc by auto

end

```

5.2.7 Properties of Probability Density Function of X

```

lemma pdfX-nonneg: pdfX  $x \geq 0$  for  $x::real$ 
  using pdfTx-nonneg pdfT0-X psi-pos' by (smt (verit))

lemma pdfX-nonpos-0: pdfX  $x = 0$  if  $x \leq 0$  for  $x::real$ 
  using pdfTx-nonpos-0 pdfT0-X psi-pos' that by (smt (verit))

lemma pdfX-beyond-0: pdfX  $x = 0$  if  $x \geq \$\psi$  for  $x::real$ 
  using pdfTx-beyond-0 pdfT0-X psi-pos' that by (smt (verit))

lemma nn-integral-pdfX-1: integralN lborel pdfX = 1
  using nn-integral-pdfTx-1 pdfT0-X psi-pos' by metis

corollary has-bochner-integral-pdfX-1: has-bochner-integral lborel pdfX 1
  by (rule has-bochner-integral-nn-integral; simp add: pdfX-nonneg nn-integral-pdfX-1)

corollary LBINT-pdfX-1: (LBINT s. pdfX s) = 1
  using has-bochner-integral-pdfX-1 by (simp add: has-bochner-integral-integral-eq)

corollary pdfX-has-integral-1: (pdfX has-integral 1) UNIV
  by (rule nn-integral-has-integral; simp add: pdfX-nonneg nn-integral-pdfX-1)

lemma set-nn-integral-pdfX-PX: set-nn-integral lborel A pdfX =  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)$ 
  if  $A \in \text{sets lborel}$  for  $A :: \text{real set}$ 
  using PT0-eq-PX-lborel that
  by (rewrite pdfT0-X[THEN sym], rewrite set-nn-integral-pdfTx-PTx; simp)

lemma pdfX-set-integrable: set-integrable lborel A pdfX if  $A \in \text{sets lborel}$  for  $A :: \text{real set}$ 
  using pdfTx-set-integrable pdfT0-X psi-pos' that by (smt (verit))

lemma set-integral-pdfX-PX: (LBINT s:A. pdfX s) =  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)$ 
  if  $A \in \text{sets lborel}$  for  $A :: \text{real set}$ 
  using PT0-eq-PX-lborel that by (rewrite pdfT0-X[THEN sym], rewrite set-integral-pdfTx-PTx; simp)

lemma pdfX-has-integral-PX: (pdfX has-integral  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)$ ) A
  if  $A \in \text{sets lborel}$  for  $A :: \text{real set}$ 
  using that apply (rewrite pdfT0-X[THEN sym], rewrite PT0-eq-PX-lborel[THEN sym], simp)
  by (rule pdfTx-has-integral-PTx; simp)

corollary pdfX-has-integral-PX-Icc: (pdfX has-integral  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. a \leq X \xi \wedge X \xi \leq b)) \{a..b\}$ 
  for  $a b :: \text{real}$ 
  using pdfX-has-integral-PX[of {a..b}] by simp

corollary pdfX-integrable-on-Icc: pdfX integrable-on  $\{a..b\}$  for  $a b :: \text{real}$ 

```

using $\text{pdf}X\text{-has-integral-}PX\text{-Icc}$ by auto

5.2.8 Relations between Life Expectation and Probability Density Function

```

context
fixes x::real
assumes x-lt-psi[simp]:  $x < \psi$ 
begin

interpretation alivePS: prob-space  $\mathfrak{M} \downharpoonright \text{alive } x$ 
by (rule MM-PS.cond-prob-space-correct; simp add: alive-def)

interpretation distrTx-RD: real-distribution distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) borel ( $T x$ ) by
simp

proposition nn-integral-T-pdfT:
 $(\int^+ \xi. \text{ennreal } (g (T x \xi)) \partial(\mathfrak{M} \downharpoonright \text{alive } x)) = (\int^+ s \in \{0..\}. \text{ennreal } (\text{pdf}T x s * g s) \partial \text{lborel})$ 
if  $g \in \text{borel-measurable lborel}$  for  $g :: \text{real} \Rightarrow \text{real}$ 
proof -
have  $(\int^+ \xi. \text{ennreal } (g (T x \xi)) \partial(\mathfrak{M} \downharpoonright \text{alive } x)) = \int^+ s. \text{ennreal } (\text{pdf}T x s) *$ 
 $\text{ennreal } (g s) \partial \text{lborel}$ 
proof -
have distributed ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) lborel ( $T x$ )  $(\lambda s. \text{ennreal } (\text{pdf}T x s))$ 
by (intro distributed-pdfTx) simp
moreover have  $(\lambda s. \text{ennreal } (g s)) \in \text{borel-measurable borel}$  using that by
measurable
ultimately show ?thesis by (rewrite distributed-nn-integral; simp)
qed
also have ... =  $\int^+ s. \text{ennreal } (\text{pdf}T x s * g s) \partial \text{lborel}$  using ennreal-mult'
pdfTx-nonneg by force
also have ... =  $(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdf}T x s * g s) \partial \text{lborel})$ 
apply (rule nn-integral-cong, simp)
by (metis atLeast-iff ennreal-0 indicator-simps linorder-not-le mult-1 mult-commute-abs
mult-zero-left pdfTx-neg-0 x-lt-ps)
finally show ?thesis .
qed

lemma expectation-LBINT-pdfT-nonneg:
alivex-PS.expectation  $(\lambda \xi. g (T x \xi)) = (\text{LBINT } s \in \{0..\}. \text{pdf}T x s * g s)$ 
if  $\bigwedge s. s \geq 0 \implies g s \geq 0$   $g \in \text{borel-measurable lborel}$  for  $g :: \text{real} \Rightarrow \text{real}$ 
— Note that  $0 = 0$  holds when the integral diverges.
using that apply (rewrite integral-eq-nn-integral, simp)
apply (rule AE-I2, metis alivex-Tx-pos less-imp-le)
unfolding set-lebesgue-integral-def apply (rewrite integral-eq-nn-integral, simp-all)
apply (rule AE-I2,
metis indicator-pos-le pdfTx-nonpos-0 x-lt-ps zero-le-mult-iff zero-le-square
pdfTx-nonneg)

```

by (rewrite nn-integral-T-pdfT) (simp-all add: indicator-mult-ennreal mult.commute)

corollary expectation-integral-pdfT-nonneg:

alivex-PS.expectation ($\lambda\xi. g(T x \xi)) = \text{integral } \{0..\} (\lambda s. \text{pdfT } x s * g s)$

if $\bigwedge s. s \geq 0 \implies g s \geq 0$ $g \in \text{borel-measurable lborel}$ **for** $g :: \text{real} \Rightarrow \text{real}$

— Note that $0 = 0$ holds when the integral diverges.

proof —

have alivex-PS.expectation ($\lambda\xi. g(T x \xi)) = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * g s)$

using expectation-LBINT-pdfT-nonneg **that** **by** simp

also have ... = $\text{integral } \{0..\} (\lambda s. \text{pdfT } x s * g s)$

using that pdfTx-nonneg **by** (intro set-borel-integral-eq-integral-nonneg; simp)

finally show ?thesis .

qed

proposition expectation-LBINT-pdfT:

alivex-PS.expectation ($\lambda\xi. g(T x \xi)) = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * g s)$

if set-integrable lborel $\{0..\}$ ($\lambda s. \text{pdfT } x s * g s)$ $g \in \text{borel-measurable lborel}$

for $g :: \text{real} \Rightarrow \text{real}$

proof —

have pg-fin: $(\int^+ \xi. \text{ennreal } (g(T x \xi)) \partial(\mathfrak{M} \downarrow \text{alive } x)) \neq \infty$

using that apply (rewrite nn-integral-T-pdfT, simp)

using that unfolding set-integrable-def apply (rewrite in asm real-integrable-def, simp)

by (simp add: indicator-mult-ennreal mult.commute)

moreover have mg-fin: $(\int^+ \xi. \text{ennreal } (-g(T x \xi)) \partial(\mathfrak{M} \downarrow \text{alive } x)) \neq \infty$

using that apply (rewrite nn-integral-T-pdfT[of $\lambda s. -g s$, simp])

using that unfolding set-integrable-def apply (rewrite in asm real-integrable-def, simp)

by (simp add: indicator-mult-ennreal mult.commute)

ultimately have [simp]: integrable ($\mathfrak{M} \downarrow \text{alive } x$) ($\lambda\xi. g(T x \xi))$

using that by (rewrite real-integrable-def; simp)

have $(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (g s)) \partial \text{lborel}) =$

$(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * g s) \partial \text{lborel})$

using SPMF.ennreal-max-0 ennreal-mult' pdfTx-nonneg x-lt-ps ψ **by** presburger

also have ... $< \infty$

using pg-fin nn-integral-T-pdfT that top.not-eq-extremum **by** auto

finally have $(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (g s)) \partial \text{lborel}) < \infty$.

hence [simp]: set-integrable lborel $\{0..\}$ ($\lambda s. \text{pdfT } x s * \max 0 (g s))$

unfolding set-integrable-def **using** that apply (intro integrableI-nonneg, simp-all)

using pdfTx-nonneg apply (intro AE-I2, force)

by (metis (no-types, lifting) indicator-mult-ennreal mult.commute nn-integral-cong)

have $(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (-g s)) \partial \text{lborel}) =$

$(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * -g s) \partial \text{lborel})$

using SPMF.ennreal-max-0 ennreal-mult' pdfTx-nonneg x-lt-ps ψ **by** presburger

also have ... $< \infty$

using mg-fin nn-integral-T-pdfT[of $\lambda s. -g s$] that top.not-eq-extremum **by**

force

finally have $(\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (-g s)) \partial \text{lborel}) < \infty$.

hence [simp]: set-integrable lborel $\{0..\}$ ($\lambda s. \text{pdfT } x s * \max 0 (-g s))$

```

unfolding set-integrable-def using that apply (intro integrableI-nonneg, simp-all)
using pdfTx-nonneg apply (intro AE-I2, force)
by (metis (no-types, lifting) indicator-mult-ennreal mult.commute nn-integral-cong)
have alivex-PS.expectation ( $\lambda\xi. g(T x \xi)) =$ 
    alivex-PS.expectation ( $\lambda\xi. \max_0(g(T x \xi))) - \text{alivex-PS.expectation}(\lambda\xi. \max_0(-g(T x \xi)))$ 
by (rewrite Bochner-Integration.integral-cong
    [where N=Μ ↳ alive x and  $g=\lambda\xi. \max_0(g(T x \xi)) - \max_0(-g(T x \xi))$ ]; simp)
moreover have alivex-PS.expectation ( $\lambda\xi. \max_0(g(T x \xi))) =$ 
    (LBINT s:{0..}. pdfT x s * max_0(g s))
using that by (rewrite expectation-LBINT-pdfT-nonneg[where  $g=\lambda s. \max_0(g s)$ ]) simp-all
moreover have alivex-PS.expectation ( $\lambda\xi. \max_0(-g(T x \xi))) =$ 
    (LBINT s:{0..}. pdfT x s * max_0(-g s))
using that by (rewrite expectation-LBINT-pdfT-nonneg[where  $g=\lambda s. \max_0(-g s)$ ]) simp-all
ultimately have alivex-PS.expectation ( $\lambda\xi. g(T x \xi)) =$ 
    (LBINT s:{0..}. pdfT x s * max_0(g s)) - (LBINT s:{0..}. pdfT x s * max_0(-g s)) by simp
also have ... = (LBINT s:{0..}. (pdfT x s * max_0(g s) - pdfT x s * max_0(-g s)))
by (rewrite set-integral-diff; simp)
also have ... = (LBINT s:{0..}. pdfT x s * (max_0(g s) - max_0(-g s)))
by (simp add: right-diff-distrib)
also have ... = (LBINT s:{0..}. pdfT x s * g s)
using minus-max-eq-min
by (metis (no-types, opaque-lifting) diff-zero max-def min-def minus-diff-eq)
finally show ?thesis .
qed

```

```

corollary expectation-integral-pdfT:
    alivex-PS.expectation ( $\lambda\xi. g(T x \xi)) = \text{integral}\{0..\}(\lambda s. pdfT x s * g s)$ 
    if  $(\lambda s. pdfT x s * g s)$  absolutely-integrable-on  $\{0..\}$   $g \in$  borel-measurable lborel
    for  $g :: \text{real} \Rightarrow \text{real}$ 
proof -
    have [simp]: set-integrable lborel  $\{0..\}(\lambda s. pdfT x s * g s)$ 
    using that by (rewrite absolutely-integrable-on-iff-set-integrable; simp)
    show ?thesis
    apply (rewrite set-borel-integral-eq-integral(2)[THEN sym], simp)
    using that by (rewrite expectation-LBINT-pdfT; simp)
qed

```

```

corollary e-LBINT-pdfT: $e`o-x = (LBINT s:{0..}. pdfT x s * s)
— Note that  $0 = 0$  holds when the life expectation diverges.
unfolding life-expect-def using expectation-LBINT-pdfT-nonneg by force

```

```

corollary e-integral-pdfT: $e`o-x = integral  $\{0..\}(\lambda s. pdfT x s * s)$ 
— Note that  $0 = 0$  holds when the life expectation diverges.

```

```
unfolded life-expect-def using expectation-integral-pdfT-nonneg by force
```

```
end
```

```
corollary e-LBINT-pdfX: $e'`0 = (LBINT x:{0..}. pdfX x * x)
```

— Note that $0 = 0$ holds when the life expectation diverges.

```
using e-LBINT-pdfT pdfT0-X psi-pos' by presburger
```

```
corollary e-integral-pdfX: $e'`0 = integral {0..} (\lambda x. pdfX x * x)
```

— Note that $0 = 0$ holds when the life expectation diverges.

```
using e-integral-pdfT pdfT0-X psi-pos' by presburger
```

5.2.9 Introduction of Force of Mortality

```
definition force-mortal :: real ⇒ real (⟨$μ'--> [101] 200)
```

```
where $μ-x ≡ MM-PS.hazard-rate X x
```

```
lemma mu-pdfX: $μ-x = pdfX x / ccdf (distr M borel X) x
```

```
if (cdf (distr M borel X)) differentiable at x for x::real
```

```
unfolded force-mortal-def pdfX-def
```

```
by (rewrite MM-PS.hazard-rate-deriv-cdf, simp-all add: that)
```

```
lemma mu-unborn-0: $μ-x = 0 if x < 0 for x::real
```

```
apply (rewrite mu-pdfX)
```

```
using cdfX-has-real-derivative-0-unborn real-differentiable-def that apply blast
```

```
using pdfX-nonpos-0 that by auto
```

```
lemma mu-beyond-0: $μ-x = 0 if x ≥ $ψ for x::real
```

— Note that division by 0 is defined as 0 in Isabelle/HOL.

```
unfolded force-mortal-def using MM-PS.hazard-rate-0-ccdf-0 ccdfX-0-equiv that  
by simp
```

```
lemma mu-nonneg-differentiable: $μ-x ≥ 0
```

```
if (cdf (distr M borel X)) differentiable at x for x::real
```

```
apply (rewrite mu-pdfX, simp add: that)
```

```
using pdfX-nonneg distrX-RD.ccdf-nonneg by simp
```

```
lemma mu-nonneg-AE: AE x in lborel. $μ-x ≥ 0
```

```
using cdfX-differentiable-AE mu-nonneg-differentiable by auto
```

```
lemma mu-measurable[measurable]: (λx. $μ-x) ∈ borel-measurable borel
```

```
proof –
```

```
obtain S where
```

```
finite S and ∃x. x ∉ S ⇒ (cdf (distr M borel X)) differentiable at x
```

```
using cdfX-piecewise-differentiable unfolding piecewise-differentiable-on-def
```

```
by blast
```

```
thus ?thesis
```

```
apply (rewrite measurable-discrete-difference)
```

```
[where f=λx. pdfX x / ccdf (distr M borel X) x and X=S], simp-all)
```

```

by (simp-all add: MM-PS.ccdf-distr-measurable borel-measurable-divide countable-finite mu-pdfX)
qed

lemma mu-deriv-ccdf:  $\mu_x = - \text{deriv}(\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)) x / \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x$ 
  if (ccdf(distr M borel X)) differentiable at  $x$   $x < \psi$  for  $x::real$ 
  unfolding force-mortal-def
  by (rewrite MM-PS.hazard-rate-deriv-ccdf; simp add: that)

lemma mu-deriv-ln:  $\mu_x = - \text{deriv}(\lambda x. \ln(\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x)) x$ 
  if (ccdf(distr M borel X)) differentiable at  $x$   $x < \psi$  for  $x::real$ 
  unfolding force-mortal-def
  apply (rewrite MM-PS.hazard-rate-deriv-ln-ccdf; simp-all add: that)
  using ccdfX-0-equiv that by force

lemma p-exp-integral-mu:  $p\{t & x\} = \exp(-\text{integral}\{x..x+t\} (\lambda y. \mu_y))$ 
  if  $x \geq 0$   $t \geq 0$   $x+t < \psi$  for  $x t :: real$ 
proof -
  have [simp]:  $x < \psi$  using that by (simp add: ereal-less-le)
  have  $p\{t & x\} = (\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)(x+t)) / (\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x)$ 
  apply (rewrite p-PX; simp-all add: that)
  by (rewrite MM-PS.ccdf-distr-P; simp)+
  also have ... =  $\exp(-\text{integral}\{x..x+t\} (\text{MM-PS.hazard-rate } X))$ 
  apply (rule MM-PS.ccdf-exp-cumulative-hazard; simp-all add: that)
  using ccdfX-piecewise-differentiable piecewise-differentiable-on-subset apply
  blast
  using ccdfX-continuous continuous-on-subset apply blast
  using ccdfX-0-equiv that ereal-less-le linorder-not-le by force
  also have ... =  $\exp(-\text{integral}\{x..x+t\} (\lambda y. \mu_y))$  unfolding force-mortal-def
  by simp
  finally show ?thesis .
qed

corollary ccdfX-exp-integral-mu:  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x = \exp(-\text{integral}\{0..x\} (\lambda y. \mu_y))$ 
  if  $0 \leq x \wedge x < \psi$  for  $x::real$ 
  by (rewrite p-exp-integral-mu[where t=x and x=0, simplified, THEN sym]; simp add: that cdfX-p)

```

5.2.10 Properties of Force of Mortality

```

context
  fixes  $x::real$ 
  assumes  $x < \psi$  [simp]
begin

```

```

interpretation alivePS: prob-space  $\mathfrak{M} \mid alive x$ 
  by (rule MM-PS.cond-prob-space-correct; simp add: alive-def)

```

interpretation *distrTx-RD*: *real-distribution distr* ($\mathfrak{M} \downharpoonright \text{alive } x$) *borel* ($T x$) **by** *simp*

lemma *hazard-rate-Tx-mu*: *alivex-PS.hazard-rate* ($T x$) $t = \$\mu-(x+t)$

if $t \geq 0$ $x+t < \$\psi$ **for** $t::\text{real}$

proof –

have [simp]: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) > 0$ **by** *force*

moreover have [simp]: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x + t) > 0$ **using** *that* **by** *force*

moreover have [simp]: $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \xi > x + t\} = \{\xi \in \text{space } \mathfrak{M}. X \xi > x + t\}$

unfolding *alive-def* **using** *that* **by** (rewrite *MM-PS.space-cond-prob-space*, *simp-all*, *force*)

hence [simp]: $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \xi > x + t\} \in \text{MM-PS.events}$ **by** *simp*
ultimately have $\star[\text{simp}]$: $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). X \xi > x + t) > 0$

unfolding *alive-def*

apply (rewrite *MM-PS.cond-prob-space-cond-prob*[*THEN sym*], (*simp-all add: pred-def*)[3])

unfolding *cond-prob-def* **by** (*smt (verit) Collect-cong divide-pos-pos*)

have *alivex-PS.hazard-rate* ($T x$) $t =$

$\text{Lim}(\text{at-right } 0) (\lambda dt. \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). t < T x \xi \wedge T x \xi \leq t + dt \mid T x \xi > t) / dt)$

unfolding *alivex-PS.hazard-rate-def* **by** *simp*

also have ... = $\text{Lim}(\text{at-right } 0)$

$(\lambda dt. \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) / dt)$

apply (rule *Lim-cong*, rule *eventually-at-rightI*[of 0 1], *simp-all*)

unfolding *futr-life-def cond-prob-def* **using** *Collect-cong* **by** (*smt (verit)*)

also have ... =

$\text{Lim}(\text{at-right } 0) (\lambda dt. \mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) / dt)$

proof –

have $\bigwedge dt. \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) =$

$\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). ?rngX \xi \mid X \xi > x + t) =$

$\mathcal{P}(\xi \text{ in } ((\mathfrak{M} \downarrow \text{alive } x) \downharpoonright \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\}). ?rngX \xi)$

using \star **by** (rewrite *alivex-PS.cond-prob-space-cond-prob*) *simp-all*

also have ... = $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \{\eta \in \text{space } \mathfrak{M}. X \eta > x + t\}). ?rngX \xi)$

proof –

have ($\mathfrak{M} \downarrow \text{alive } x$) $\downharpoonright \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\} =$

$\mathfrak{M} \downarrow \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\}$

apply (rewrite *MM-PS.cond-cond-prob-space*, *simp-all*)

unfolding *alive-def* **using** *that* **by** *force*

also have ... = $\mathfrak{M} \downarrow \{\eta \in \text{space } \mathfrak{M}. X \eta > x + t\}$ **by** *simp*

finally have ($\mathfrak{M} \downarrow \text{alive } x$) $\downharpoonright \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\} =$

```

 $\mathfrak{M} \downharpoonright \{\eta \in \text{space } \mathfrak{M}. X \eta > x + t\} .$ 
thus ?thesis by simp
qed
also have ... =  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t)$ 
by (rewrite MM-PS.cond-prob-space-cond-prob, simp-all add: pred-def
sets.sets-Collect-conj)
finally show  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downharpoonright \text{alive } x). ?rngX \xi \mid X \xi > x + t) =$ 
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) .$ 
qed
thus ?thesis
apply -
by (rule Lim-cong, rule eventually-at-rightI[of 0 1]; simp)
qed
also have ... = $μ-(x+t) unfolding force-mortal-def MM-PS.hazard-rate-def
by simp
finally show ?thesis .
qed

lemma pdfTx-p-mu: pdfT x t = $p-{t&x} * $μ-(x+t)
if (cdf (distr (mathfrak{M} \downharpoonright alive x) borel (T x))) differentiable at t t > 0 for t::real
proof (cases <x+t < $ψ>)
case True
hence [simp]: t ≥ 0 and (ccdf (distr (mathfrak{M} \downharpoonright alive x) borel (T x))) t ≠ 0
using that p-0-equiv unfolding survive-def by auto
have deriv (cdf (distr (mathfrak{M} \downharpoonright alive x) borel (T x))) t =
ccdf (distr (mathfrak{M} \downharpoonright alive x) borel (T x)) t * alivex-PS.hazard-rate (T x) t
by (rule alivex-PS.deriv-cdf-hazard-rate; simp add: that)
thus ?thesis unfolding survive-def pdfT-def using hazard-rate-Tx-mu True that
by simp
next
case False
hence x+t ≥ $ψ by simp
thus ?thesis using pdfTx-beyond-0 mu-beyond-0 by simp
qed

lemma deriv-t-p-mu: deriv (λs. $p-{s&x}) t = - $p-{t&x} * $μ-(x+t)
if (λs. $p-{s&x}) differentiable at t t > 0 for t::real
proof -
let ?cdfTx = cdf (distr (mathfrak{M} \downharpoonright alive x) borel (T x))
have diff: ?cdfTx differentiable at t
using that distrTx-RD.differentiable-cdf-ccdf unfolding survive-def by blast
hence deriv ?cdfTx t = $p-{t&x} * $μ-(x+t) using that pdfTx-p-mu pdfT-def
by simp
moreover have deriv ?cdfTx t = - deriv (λs. $p-{s&x}) t
using that diff distrTx-RD.deriv-cdf-ccdf unfolding survive-def by presburger
ultimately show ?thesis by simp
qed

lemma pdfTx-p-mu-AE: AE s in lborel. s > 0 —→ pdfT x s = $p-{s&x} * $μ-(x+s)

```

```

proof -
let ?cdfX = cdf (distr M borel X)
let ?cdfTx = cdf (distr (M \ alive x) borel (T x))
from cdfX-differentiable-AE obtain N
  where diff:  $\bigwedge r. r \in space lborel - N \implies ?cdfX$  differentiable at r
    and null:  $N \in null\text{-sets lborel}$ 
  using AE-E3 by blast
let ?N' = {s. x+s ∈ N}
have ?N' ∈ null\text{-sets lborel}
  using null\text{-sets-translation}[of N - x, simplified] null by (simp add: add.commute)
  moreover have {s ∈ space lborel. s > 0 → pdfT x s = $p-{s&x} * $μ-(x+s)} ⊆ ?N'
  proof (rewrite Compl-subset-Compl-iff[THEN sym], safe)
    fix s::real
    assume s ∈ space lborel and x+s ∉ N and s > 0
    thus pdfT x s = $p-{s&x} * $μ-(x+s)
      apply (intro pdfTx-p-mu, simp-all)
      by (rewrite differentiable-cdfX-cdfTx[THEN sym]; simp add: diff)
  qed
  ultimately show ?thesis using AE-I'[of ?N'] by simp
qed

lemma LBINT-p-mu-q-defer: (LBINT s:{f<..f+t}. $p-{s&x} * $μ-(x+s)) = $q-{f|t&x}
  if t ≥ 0 f ≥ 0 for t f :: real
proof -
  have (LBINT s:{f<..f+t}. $p-{s&x} * $μ-(x+s)) = (LBINT s:{f<..f+t}. pdfT x s)
    apply (rule set-lebesgue-integral-cong-AE; simp)
    apply (simp add: survive-def)
    using pdfTx-p-mu-AE apply (rule AE-mp)
    using that by (intro always-eventually; simp add: ereal-less-le)
  also have ... = enn2real ( $\int^+ s \in \{f < .. f+t\}. ennreal (pdfT x s)$ ) ∂lborel
  proof -
    have ( $\int^+ s \in \{f < .. f+t\}. ennreal (pdfT x s)$ ) ∂lborel < ⊤
    proof -
      have ( $\int^+ s \in \{f < .. f+t\}. ennreal (pdfT x s)$ ) ∂lborel ≤  $\int^+ s. ennreal (pdfT x s)$ 
        by (smt (verit) indicator-simps le-zero-eq linorder-le-cases
          mult.right-neutral mult-zero-right nn-integral-mono)
      also have ... < ⊤ using nn-integral-pdfTx-1 by simp
      finally show ?thesis .
    qed
    thus ?thesis
      unfolding set-lebesgue-integral-def
      by (rewrite enn2real-nn-integral-eq-integral[where g=λs. pdfT x s * indicator
        {f<..f+t} s])
        (simp-all add: pdfTx-nonneg mult.commute ennreal-indicator ennreal-mult')
  qed
  also have ... = measure (density lborel (pdfT x)) {f<..f+t}

```

```

unfolding measure-def by (rewrite emeasure-density; simp)
also have ... = measure (distr ( $\mathfrak{M} \downharpoonright \text{alive } x$ ) lborel ( $T x$ ))  $\{f <.. f+t\}$ 
  using distributed-pdfTx unfolding distributed-def by simp
also have ... =
   $cdf (distr (\mathfrak{M} \downharpoonright \text{alive } x) lborel (T x)) (f+t) - cdf (distr (\mathfrak{M} \downharpoonright \text{alive } x) lborel (T x)) f$ 
  using that finite-borel-measure.cdf-diff-eq2
  by (smt (verit) distrTx-RD.finite-borel-measure-axioms distr-cong sets-lborel)
also have ... =  $\$q\{-f|t\&x\}$ 
  using q-defer-q-die-def that by (metis distr-cong sets-lborel x-lt-psi)
  finally show ?thesis .
qed

corollary LBINT-p-mu-q: ( $LBINT s;\{0 <.. t\}. \$p\{-s\&x\} * \$\mu\-(x+s)$ ) =  $\$q\{-t\&x\}$ 
if  $t \geq 0$  for  $t::real$ 
  using LBINT-p-mu-q-defer that by force

lemma set-integrable-p-mu: set-integrable lborel  $\{f <.. f+t\}$   $(\lambda s. \$p\{-s\&x\} * \$\mu\-(x+s))$ 
  if  $t \geq 0 f \geq 0$  for  $t f :: real$ 
proof -
  have  $(\lambda s. \$p\{-s\&x\}) \in \text{borel-measurable borel}$  unfolding survive-def by simp
  moreover have AE s in lborel.  $f < s \wedge s \leq f + t \longrightarrow \$p\{-s\&x\} * \$\mu\-(x+s) = pdfT x s$ 
  using pdfTx-p-mu-AE apply (rule AE-mp)
  using that by (intro always-eventually; simp add: ereal-less-le)
  moreover have set-integrable lborel  $\{f <.. f+t\}$  (pdfT x) using pdfTx-set-integrable
  by simp
  ultimately show ?thesis by (rewrite set-integrable-cong-AE[where g=pdfT x];
  simp)
qed

lemma p-mu-has-integral-q-defer-Ioc:
   $((\lambda s. \$p\{-s\&x\} * \$\mu\-(x+s)) \text{ has-integral } \$q\{-f|t\&x\}) \{f <.. f+t\}$ 
  if  $t \geq 0 f \geq 0$  for  $t f :: real$ 
  apply (rewrite LBINT-p-mu-q-defer[THEN sym], simp-all add: that)
  apply (rewrite set-borel-integral-eq-integral, simp add: set-integrable-p-mu that)
  by (rewrite has-integral-integral[THEN sym];
  simp add: set-borel-integral-eq-integral set-integrable-p-mu that)

lemma p-mu-has-integral-q-defer-Icc:
   $((\lambda s. \$p\{-s\&x\} * \$\mu\-(x+s)) \text{ has-integral } \$q\{-f|t\&x\}) \{f..f+t\}$  if  $t \geq 0 f \geq 0$  for
   $t f :: real$ 
proof -
  have negligible {f} by simp
  hence [simp]: negligible  $(\{f..f+t\} - \{f <.. f+t\})$ 
  by (smt (verit) Diff-iff atLeastAtMost-iff greaterThanAtMost-iff insertI1
    negligible-sing negligible-subset subsetI)
  have [simp]: negligible  $(\{f <.. f+t\} - \{f..f+t\})$  by (simp add: subset-eq)
  show ?thesis

```

```

apply (rewrite has-integral-spike-set-eq[where T={f<..f+t}])
  apply (rule negligible-subset[of {f..f+t} - {f<..f+t}], simp, blast)
  apply (rule negligible-subset[of {f<..f+t} - {f..f+t}], simp, blast)
  using p-mu-has-integral-q-defer-Icc that by simp
qed

```

corollary *p-mu-has-integral-q-Icc*:
 $((\lambda s. \$p\{-s\&x\} * \$\mu\{-(x+s)\}) \{0..t\}$ if $t \geq 0$ for $t::real$
 using *p-mu-has-integral-q-defer-Icc*[where $f=0$] that by *simp*

corollary *p-mu-integrable-on-Icc*:
 $(\lambda s. \$p\{-s\&x\} * \$\mu\{-(x+s)\})$ integrable-on $\{0..t\}$ if $t \geq 0$ for $t::real$
 using *p-mu-has-integral-q-Icc* that by *auto*

lemma *e-ennreal-p-mu*: $(\int^+ \xi. ennreal (T x \xi) \partial(\mathfrak{M} \downarrow alive x)) =$
 $(\int^+ s \in \{0..\}. ennreal (\$p\{-s\&x\} * \$\mu\{-(x+s)\} * s) \partial borel)$
proof –
 have [simp]: $sym-diff \{0..\} \{0<..\} = \{0::real\}$ by force
 have $(\int^+ \xi. ennreal (T x \xi) \partial(\mathfrak{M} \downarrow alive x)) = (\int^+ s \in \{0..\}. ennreal (pdfT x s * s) \partial borel)$
 by (rewrite nn-integral-T-pdfT[where $g=\lambda s. s$; *simp*])
 also have ... = $(\int^+ s \in \{0<..\}. ennreal (pdfT x s * s) \partial borel)$
 by (rewrite nn-integral-null-delta; force)
 also have ... = $(\int^+ s \in \{0<..\}. ennreal (\$p\{-s\&x\} * \$\mu\{-(x+s)\} * s) \partial borel)$
 apply (rule nn-integral-cong-AE)
 using pdfTx-p-mu-AE apply (rule AE-mp, intro AE-I2) by force
 also have ... = $(\int^+ s \in \{0..\}. ennreal (\$p\{-s\&x\} * \$\mu\{-(x+s)\} * s) \partial borel)$
 by (rewrite nn-integral-null-delta[THEN sym]; force)
 finally show ?thesis .
qed

lemma *e-LBINT-p-mu*: $\$e^{\circ} \circ x = (LBINT s; \{0..\}. \$p\{-s\&x\} * \$\mu\{-(x+s)\} * s)$
 — Note that $0 = 0$ holds when the life expectation diverges.
proof –
 let ?f = $\lambda s. \$p\{-s\&x\} * \$\mu\{-(x+s)\} * s$
 have [simp]: $(\lambda s. ?f s * indicat-real \{0..\} s) \in borel-measurable borel$
 by measurable (simp-all add: survive-def)
 hence [simp]: $(\lambda s. indicat-real \{0..\} s * ?f s) \in borel-measurable borel$
 by (meson measurable-cong mult.commute)
 have [simp]: AE s in borel. ?f s * indicat-real \{0..\} s ≥ 0
proof –
 have AE s in borel. pdfT x s * s * indicat-real \{0..\} s ≥ 0
 using pdfTx-nonneg pdfTx-nonpos-0 x-lt-psi
 by (intro AE-I2) (smt (verit, del-insts) indicator-pos-le mult-eq-0-iff mult-nonneg-nonneg)
 thus ?thesis
 apply (rule AE-mp)
 using pdfTx-p-mu-AE apply (rule AE-mp)
 by (rule AE-I2) (metis atLeast-iff indicator-simps(2) mult-eq-0-iff order-le-less)
qed

```

hence [simp]: AE s in lborel. indicat-real {0..} s * ?f s ≥ 0
  by (metis (no-types, lifting) AE-cong mult.commute)
show ?thesis
proof (cases `integrable (M ⊥ alive x) (T x)`)
  case True
  hence ennreal ($e `o` x) = ∫⁺ξ. ennreal (T x ξ) ∂(M ⊥ alive x)
    unfolding life-expect-def apply (rewrite nn-integral-eq-integral, simp-all)
    by (smt (verit) AE-I2 alivex-Tx-pos)
  also have ... = ∫⁺s. ennreal (?f s * indicat-real {0..} s) ∂lborel
    by (simp add: nn-integral-set-ennreal e-ennreal-p-mu)
  also have ... = ennreal (LBINT s:{0..}. ?f s)
  proof -
    have integrable lborel (λs. ?f s * indicat-real {0..} s)
    proof -
      have (∫⁺s. ennreal (?f s * indicat-real {0..} s) ∂lborel) < ∞
        by (metis calculation ennreal-less-top infinity-ennreal-def)
      thus ?thesis by (intro integrableI-nonneg; simp)
    qed
    thus ?thesis
      unfolding set-lebesgue-integral-def
      by (rewrite nn-integral-eq-integral, simp-all) (meson mult.commute)
  qed
  finally have ennreal ($e `o` x) = ennreal (LBINT s:{0..}. ?f s) .
  moreover have (LBINT s:{0..}. ?f s) ≥ 0
    unfolding set-lebesgue-integral-def by (rule integral-nonneg-AE) simp
  ultimately show ?thesis using e-nonneg by simp
next
  case False
  hence $e `o` x = 0 unfolding life-expect-def using not-integrable-integral-eq by
  force
  also have ... = (LBINT s:{0..}. ?f s)
  proof -
    have ∞ = ∫⁺ξ. ennreal (T x ξ) ∂(M ⊥ alive x)
      using nn-integral-nonneg-infinite False
    by (smt (verit) AE-cong Tx-alivex-measurable alivex-PS.AE-prob-1 alivex-PS.prob-space
        alivex-Tx-pos nn-integral-cong)
    hence 0 = enn2real (∫⁺s:{0..}. ennreal (?f s) ∂lborel) using e-ennreal-p-mu
  by simp
  also have ... = (LBINT s:{0..}. ?f s)
    unfolding set-lebesgue-integral-def apply (rewrite integral-eq-nn-integral,
    simp-all)
    by (simp add: indicator-mult-ennreal mult.commute)
    finally show ?thesis by simp
  qed
  finally show ?thesis .
qed
qed

```

lemma e-integral-p-mu: \$e `o` x = integral {0..} (λs. \$p-{s&x} * \$μ-(x+s) * s)

— Note that $\theta = 0$ holds when the life expectation diverges.

proof —

have $(LBINT s:\{0..\}. \$p\{-s\&x\} * \$\mu\(-(x+s) * s) = integral \{0..\} (\lambda s. \$p\{-s\&x\} * \$\mu\(-(x+s) * s)$

proof —

have $AE s \text{ in } lborel. s \geq 0 \rightarrow \$p\{-s\&x\} * \$\mu\(-(x+s) * s \geq 0$

proof —

have $AE s \text{ in } lborel. \$\mu\(-(x+s) \geq 0 \text{ by (rule AE-translation, rule mu-nonneg-AE)}$

with $p\text{-nonneg}$ show $?thesis$ by force

qed

moreover have $(\lambda s. \$p\{-s\&x\} * \$\mu\(-(x+s) * s) \in borel\text{-measurable borel}$

unfolding $survive\text{-def}$ by simp

ultimately show $?thesis$ by (intro set-borel-integral-eq-integral-nonneg-AE; simp)

qed

thus $?thesis$ using $e\text{-LBINT-}p\text{-mu}$ by simp

qed

end

lemma $p\text{-has-real-derivative-}x\text{-ccdf}X$:

$((\lambda y. \$p\{-t\&y\}) \text{ has-real-derivative}$

$((deriv svl (x+t) * svl x - svl (x+t) * deriv svl x) / (svl x)^2) \text{ (at } x\text{)}$

if $svl \equiv ccdf$ ($distr \mathfrak{M} borel X$) svl differentiable at x svl differentiable at $(x+t)$

$t \geq 0 x < \$\psi$ for $x t :: real$

proof —

have $open \{y. svl \text{ differentiable at } y\}$ using $ccdfX\text{-differentiable-open-set}$ that by simp

with that obtain $e0$ where $e0\text{-pos}: e0 > 0$ and $ball\text{-}e0: ball x e0 \subseteq \{y. svl \text{ differentiable at } y\}$

using $open\text{-contains-ball}$ by blast

define e where $e \equiv if \$\psi < \infty \text{ then } min e0 \text{ (real-of-ereal } \$\psi - x) \text{ else } e0$

have $e > 0 \wedge ball x e \subseteq \{y. svl y \neq 0 \wedge svl \text{ differentiable at } y\}$

proof (cases $\langle \$\psi < \infty \rangle$)

case $True$

hence $e > 0$

proof —

have $real\text{-of-ereal } \$\psi - x > 0$ using $not\text{-inftyI} True$ that by force

with $e0\text{-pos}$ show $?thesis$ unfolding $e\text{-def}$ using $True$ by simp

qed

moreover have $ball x e \subseteq \{y. svl y \neq 0\}$

proof —

have $e \leq real\text{-of-ereal } \$\psi - x$ unfolding $e\text{-def}$ using $True$ by simp

hence $ball x e \subseteq \{.. < real\text{-of-ereal } \$\psi\}$ unfolding $ball\text{-def}$ $dist\text{-real}\text{-def}$ by force

thus $?thesis$ using that $ccdfX\text{-}0\text{-equiv}$

using $True$ $ereal\text{-less-real}\text{-iff}$ $\psi\text{-nonneg}$ by auto

qed

moreover have $ball x e \subseteq \{y. svl \text{ differentiable at } y\}$

```

proof -
  have  $e \leq e0$  unfolding  $e\text{-def}$  using  $\text{True}$  by  $\text{simp}$ 
  hence  $\text{ball } x \ e \subseteq \text{ball } x \ e0$  by  $\text{force}$ 
    with  $\text{ball-}e0$  show  $?thesis$  by  $\text{simp}$ 
  qed
  ultimately show  $?thesis$  by  $\text{blast}$ 
next
  case  $\text{False}$ 
  hence  $\text{ball } x \ e \subseteq \{y. \text{svl } y \neq 0\}$  using  $\text{ccdfX-0-equiv}$  that by  $\text{simp}$ 
    with  $\text{False}$  show  $?thesis$  unfolding  $e\text{-def}$  using  $e0\text{-pos ball-}e0$  by  $\text{force}$ 
  qed
  hence  $e\text{-pos: } e > 0$  and  $\text{ball-}e: \text{ball } x \ e \subseteq \{y. \text{svl } y \neq 0 \wedge \text{svl differentiable at } y\}$ 
  by  $\text{auto}$ 
  hence  $\text{ball-svl: } \bigwedge y. y \in \text{ball } x \ e \implies \text{svl } y \neq 0 \wedge \text{svl field-differentiable at } y$ 
    using  $\text{differentiable-eq-field-differentiable-real}$  by  $\text{blast}$ 
  hence  $\bigwedge y. y \in \text{ball } x \ e \implies \$p\{-t\&y\} = \text{svl } (y+t) / \text{svl } y$ 
    unfolding  $\text{survive-def}$  using that  $\text{ccdfX-0-equiv}$  by (rewrite  $\text{ccdfTx-ccdfX}$ ,
 $\text{simp-all}$ ) force
  moreover from  $\text{ball-svl}$  have  $((\lambda y. \text{svl } (y+t) / \text{svl } y) \text{ has-real-derivative}$ 
 $((\text{deriv svl } (x+t) * \text{svl } x - \text{svl } (x+t) * \text{deriv svl } x) / (\text{svl } x)^2)$  (at  $x$ )
  apply (rewrite  $\text{power2-eq-square}$ )
  apply (rule  $\text{DERIV-divide}$ )
  using  $\text{DERIV-deriv-iff-real-differentiable DERIV-shift}$  that apply  $\text{blast}$ 
  using that  $\text{DERIV-deriv-iff-real-differentiable}$  apply  $\text{simp}$ 
  by ( $\text{simp add: } e\text{-pos}$ )
  ultimately show  $?thesis$ 
    using  $e\text{-pos}$ 
    apply (rewrite  $\text{has-field-derivative-cong-eventually}$ [where  $g=\lambda y. \text{svl } (y+t) / \text{svl } y$ ],  $\text{simp-all}$ )
      by (smt (verit)  $\text{dist-commute eventually-at}$ )
  qed

lemma  $p\text{-has-real-derivative-}x\text{-}p\text{-mu}:$ 
 $((\lambda y. \$p\{-t\&y\}) \text{ has-real-derivative } \$p\{-t\&x\} * (\$p\{-x\} - \$p\{-x+t\}))$  (at  $x$ )
  if  $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } x$   $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } (x+t)$ 
   $t \geq 0 \ x < \$\psi$  for  $x \ t :: \text{real}$ 
proof (cases  $x+t < \$\psi$ )
  case  $\text{True}$ 
  let  $?svl = \text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X)$ 
  have [simp]:  $?svl \neq 0$  using that  $\text{ccdfX-0-equiv}$  by (smt (verit)  $\text{le-ereal-le}$ 
 $\text{linorder-not-le}$ )
  have [simp]:  $?svl \text{ field-differentiable at } (x+t)$ 
    using that  $\text{differentiable-eq-field-differentiable-real}$  by  $\text{simp}$ 
  have  $((\lambda y. \$p\{-t\&y\}) \text{ has-real-derivative}$ 
 $((\text{deriv } ?svl \ (x+t) * ?svl \ x - ?svl \ (x+t) * \text{deriv } ?svl \ x) / (?svl \ x)^2)$  (at  $x$ )
    using  $p\text{-has-real-derivative-}x\text{-}ccdfX$  that by  $\text{simp}$ 
  moreover have  $(\text{deriv } ?svl \ (x+t) * ?svl \ x - ?svl \ (x+t) * \text{deriv } ?svl \ x) / (?svl \ x)^2 =$ 

```

$\$p\{-t \& x\} * (\$μ - x - \$μ - (x + t))$ (**is** $?LHS = ?RHS$)
proof –
have $deriv ?svl (x + t) = deriv (\lambda y. ?svl (y + t)) x$
using that by (metis DERIV-deriv-iff-real-differentiable DERIV-imp-deriv DERIV-shift)
hence $?LHS = (deriv (\lambda y. ?svl (y + t)) x * ?svl x - ?svl (x + t) * deriv ?svl x) / (?svl x)^2$
by simp
also have $\dots = (deriv (\lambda y. ?svl (y + t)) x - ?svl (x + t) * deriv ?svl x) / ?svl x$
by (simp add: add-divide-eq-if-simps(4) power2-eq-square)
also have $\dots = (- ?svl (x + t) * \$μ - (x + t) + ?svl (x + t) * \$μ - x) / ?svl x$
proof –
have $deriv (\lambda y. ?svl (y + t)) x = - ?svl (x + t) * \$μ - (x + t)$
apply (rewrite add.commute, rewrite deriv-shift[THEN sym], rewrite add.commute, simp)
using add.commute that **by** (metis MM-PS.deriv-ccdf-hazard-rate X-RV force-mortal-def)
moreover have $- ?svl (x + t) * deriv ?svl x / ?svl x = ?svl (x + t) * \$μ - x$
using that by (simp add: MM-PS.deriv-ccdf-hazard-rate force-mortal-def)
ultimately show ?thesis by simp
qed
also have $\dots = ?svl (x + t) * (\$μ - x - \$μ - (x + t)) / ?svl x$ **by** (simp add: mult-diff-mult)
also have $\dots = ?RHS$ unfolding survive-def using ccdfTx-ccdfX that by simp
ultimately show ?thesis by simp
qed
ultimately show ?thesis by simp
next
case False
hence ptx-0: $\$p\{-t \& x\} = 0$ **using** p-0-equiv that **by** simp
moreover have $((\lambda y. \$p\{-t \& y\}) has-real-derivative 0)$ (at x)
proof –
have $\lambda y. x < y \implies y < \$ψ \implies \$p\{-t \& y\} = 0$
using False p-0-equiv that **by** (smt (verit, ccfv-SIG) ereal-less-le linorder-not-le)
hence $\forall F x \text{ in at-right } x. \$p\{-t \& x\} = 0$
apply (rewrite eventually-at-right-field)
using that by (meson ereal-dense2 ereal-le-less less-eq-ereal-def less-ereal.simps)
hence $((\lambda y. \$p\{-t \& y\}) has-real-derivative 0)$ (at-right x)
using ptx-0 **by** (rewrite has-field-derivative-cong-eventually[where g=λ-. 0]; simp)
thus ?thesis
using vector-derivative-unique-within p-has-real-derivative-x-ccdfX that
by (metis has-field-derivative-at-within has-real-derivative-iff-has-vector-derivative trivial-limit-at-right-real)
qed
ultimately show ?thesis by simp
qed

corollary deriv-x-p-mu: deriv ($\lambda y. \$p - \{t \& y\}$) $x = \$p - \{t \& x\} * (\$mu - x - \$mu - (x + t))$
if ccdf (distr \mathfrak{M} borel X) differentiable at x ccdf (distr \mathfrak{M} borel X) differentiable at $(x + t)$
 $t \geq 0$ $x < \$\psi$ **for** $x t :: real$
using that p-has-real-derivative-x-p-mu DERIV-imp-deriv **by** blast

lemma e-has-derivative-mu-e: $((\lambda x. \$e^{\circ} - x) has-real-derivative (\$mu - x * \$e^{\circ} - x - 1))$ (at x)
if $\bigwedge x. x \in \{a < .. < b\} \implies$ set-integrable lborel $\{x..\}$ (ccdf (distr \mathfrak{M} borel X))
 $ccdf (distr \mathfrak{M} borel X)$ differentiable at x $x \in \{a < .. < b\}$ $b \leq \$\psi$
for $a b x :: real$

proof –

let ?svl = ccdf (distr \mathfrak{M} borel X)
have $x < \$\psi$ **using** that ereal-le-less **by** simp
hence svlx-neq0[simp]: $?svl x \neq 0$ **by** (simp add: ccdfX-0-equiv linorder-not-le)
define $d :: real$ **where** $d \equiv \min(b - x) (x - a)$
have d-pos: $d > 0$ **unfolding** d-def **using** that ereal-less-real-iff **by** force
have e-svl: $\bigwedge y. y < \$\psi \implies \$e^{\circ} - y = (LBINT t:\{0..\}. ?svl (y+t)) / ?svl y$
apply (rewrite e-LBINT-p, simp)
apply (rewrite set-integral-divide-zero[THEN sym])
apply (rule set-lebesgue-integral-cong, simp-all)
unfolding survive-def **using** ccdfTx-ccdfX **by** force
have $((\lambda y. LBINT t:\{0..\}. ?svl (y+t)) has-real-derivative (- ?svl x))$ (at x)

proof –

{ fix y **assume** dist $y x < d$
hence $y - ab: y \in \{a < .. < b\}$ **unfolding** d-def dist-real-def **by** force
hence set-integrable lborel $\{y..\}$?svl **using** that **by** simp
hence set-integrable lborel (einterval $y \infty$) ?svl
by (rewrite set-integrable-discrete-difference[**where** $X = \{y\}$]; simp) force
moreover have $\bigwedge u. ((\lambda u. u - y) has-real-derivative (1 - 0))$ (at u)
by (rule derivative-intros) auto
moreover have $\bigwedge u. y < u \implies isCont (\lambda t. ?svl (y+t)) (u - y)$
apply (rewrite add.commute, rewrite isCont-shift, simp)
using ccdfX-continuous continuous-on-eq-continuous-within **by** blast
moreover have $((ereal \circ (\lambda u. u - y) \circ real-of-ereal) \longrightarrow 0)$ (at-right (ereal y))
by (smt (verit, ccfv-SIG) LIM-zero Lim-cong-within ereal-tendsto-simps1(2)
ereal-tendsto-simps2(1) tendsto-ident-at zero-ereal-def)
moreover have $((ereal \circ (\lambda u. u - y) \circ real-of-ereal) \longrightarrow \infty)$ (at-left ∞)

proof –

have $\bigwedge r. r + y + 1 \leq u \implies r < u - y$ **by** auto
hence $\bigwedge r. \forall_F u \text{ in at-top. } r < u - y$ **by** (rule eventually-at-top-linorderI)
thus ?thesis **by** (rewrite ereal-tendsto-simps, rewrite tendsto-PInfty) simp
qed

ultimately have $(LBINT t=0..\infty. ?svl (y+t)) = (LBINT u=y..\infty. ?svl (y+(u-y)) * 1)$
using distrX-RD.ccdf-nonneg **by** (intro interval-integral-substitution-nonneg(2);
simp)

moreover have $(LBINT t:\{0..\}. ?svl (y+t)) = (LBINT t=0..\infty. ?svl (y+t))$
unfolding *interval-lebesgue-integral-def einterval-def* **apply** *simp*
by (*rule set-integral-discrete-difference[where X={0}]; force*)
moreover have $(LBINT u=y..\infty. ?svl (y+(u-y)*1)) = (LBINT u:\{y..\}.$
 $?svl u)$
unfolding *interval-lebesgue-integral-def einterval-def* **apply** *simp*
by (*rule set-integral-discrete-difference[where X={y}]; force*)
ultimately have $(LBINT t:\{0..\}. ?svl (y+t)) = (LBINT u:\{y..\}. ?svl u)$ **by**
simp }
hence $\forall_F y \text{ in nhds } x. (LBINT t:\{0..\}. ?svl (y+t)) = (LBINT u:\{y..\}. ?svl u)$
using *d-pos* **by** (*rewrite eventually-nhds-metric*) *auto*
moreover have $((\lambda y. LBINT u:\{y..\}. ?svl u) \text{ has-real-derivative } (- ?svl x))$
 $(\text{at } x)$
proof –
have $((\lambda y. \text{integral } \{y..b\} ?svl) \text{ has-real-derivative } (- ?svl x))$ $(\text{at } x \text{ within } \{a..b\})$
using that **apply** (*intro integral-has-real-derivative'*; *simp*)
using *ccdfX-continuous continuous-on-subset* **by** *blast*
hence $((\lambda y. \text{integral } \{y..b\} ?svl) \text{ has-real-derivative } (- ?svl x))$ $(\text{at } x)$
using that **apply** (*rewrite at-within-open[where S={a<..<b}, THEN sym]*,
simp-all)
by (*rule DERIV-subset[where s={a..b}]*) *auto*
moreover have $\forall_F y \text{ in nhds } x. (LBINT u:\{y..b\}. ?svl u) = \text{integral } \{y..b\}$
 $?svl$
apply (*rewrite eventually-nhds-metric*)
using *d-pos* **by** (*metis ccdfX-integrable-Icc set-borel-integral-eq-integral(2)*)
ultimately have $((\lambda y. LBINT u:\{y..b\}. ?svl u) \text{ has-real-derivative } (- ?svl x))$ $(\text{at } x)$
by (*rewrite DERIV-cong-ev*; *simp*)
hence $((\lambda y. (LBINT u:\{y..b\}. ?svl u) + (LBINT u:\{b<..\}. ?svl u)) \text{ has-real-derivative } (- ?svl x))$ $(\text{at } x)$
by (*rewrite to - ?svl x + 0 add-0-right[THEN sym]*, *rule DERIV-add*; *simp*)
moreover have $\forall_F y \text{ in nhds } x.$
 $(LBINT u:\{y..\}. ?svl u) = (LBINT u:\{y..b\}. ?svl u) + (LBINT u:\{b<..\}.$
 $?svl u)$
proof –
{ fix y **assume** *dist y x < d*
hence $y-ab: y \in \{a<..<b\}$ **unfolding** *d-def dist-real-def* **by** *force*
hence *set-integrable lborel {y..}* $?svl$ **using that** **by** *simp*
hence *set-integrable lborel {y..b}* $?svl$ *set-integrable lborel {b<..}* $?svl$
apply (*rule set-integrable-subset*, *simp-all*)+
using *y-ab* **by** *force*
moreover have $\{y..b\} \cap \{b<..\} = \{\} \{y..\} = \{y..b\} \cup \{b<..\}$ **using** *y-ab*
by *force+*
ultimately have
 $(LBINT u:\{y..\}. ?svl u) = (LBINT u:\{y..b\}. ?svl u) + (LBINT u:\{b<..\}.$
 $?svl u)$
using *set-integral-Un* **by** *simp* }
thus *?thesis* **using** *d-pos* **by** (*rewrite eventually-nhds-metric*) *blast*

```

qed
ultimately show ?thesis by (rewrite has-field-derivative-cong-ev; simp)
qed
ultimately show ?thesis by (rewrite DERIV-cong-ev; simp)
qed
moreover have (?svl has-real-derivative (deriv ?svl x)) (at x)
  using that DERIV-deriv-iff-real-differentiable by blast
ultimately have ((λy. (LBINT t:{0..}. ?svl (y+t)) / ?svl y) has-real-derivative
  ((- ?svl x) * ?svl x - (LBINT t:{0..}. ?svl (x+t)) * deriv ?svl x) / (?svl x *
?svl x)) (at x)
  by (rule DERIV-divide) simp
moreover have eventually (λy. (LBINT t:{0..}. ?svl (y+t)) / ?svl y = $e‘o-y)
(at x)
proof -
  { fix y assume dist y x < d y ≠ x
    hence y < $ψ
      unfolding dist-real-def d-def using that ereal-le-less by fastforce
    hence $e‘o-y = (LBINT t:{0..}. ?svl (y+t)) / ?svl y by (rule e-svl) }
thus ?thesis
  apply (rewrite eventually-at-filter, rewrite eventually-nhds-metric)
  using d-pos that by metis
qed
ultimately have ((λy. $e‘o-y) has-real-derivative
  ((- ?svl x) * ?svl x - (LBINT t:{0..}. ?svl (x+t)) * deriv ?svl x) / (?svl x *
?svl x)) (at x)
  using e-svl by (rewrite has-field-derivative-cong-eventually[THEN sym]; simp)
moreover have
  ((- ?svl x) * ?svl x - (LBINT t:{0..}. ?svl (x+t)) * deriv ?svl x) / (?svl x *
?svl x) =
  $μ-x * $e‘o-x - 1 (is ?LHS = ?RHS)
proof -
  have ?LHS = -1 + (LBINT t:{0..}. ?svl (x+t)) / ?svl x * (- deriv ?svl x /
?svl x)
    by simp (smt (verit) svlx-neq0 add-divide-distrib divide-eq-minus-1-iff
      mult-minus-left real-divide-square-eq)
  also have ... = -1 + $e‘o-x * $μ-x using e-svl mu-deriv-ccdf that by force
  also have ... = ?RHS by simp
  finally show ?thesis .
qed
ultimately show ?thesis by simp
qed

corollary e-has-derivative-mu-e': ((λx. $e‘o-x) has-real-derivative ($μ-x * $e‘o-x
- 1)) (at x)
if ∫x. x ∈ {a < .. < b} ⇒ ccdf (distr M borel X) integrable-on {x..}
  ccdf (distr M borel X) differentiable at x x ∈ {a < .. < b} b ≤ $ψ
for a b x :: real
using that apply (intro e-has-derivative-mu-e[where a=a], simp-all)
using distrX-RD.ccdf-nonneg by (rewrite integrable-on-iff-set-integrable-nonneg;

```

simp)

5.2.11 Properties of Curtate Life Expectation

context

fixes $x::real$

assumes $x-lt-\psi[simp]: x < \psi$

begin

lemma $isCont-p-nat: isCont (\lambda t. \$p-\{t&x\}) (k + (1::real))$ **for** $k::nat$

proof –

fix $k::nat$

have $continuous-on \{0 < ..\} (\lambda t. \$p-\{t&x\})$

unfolding $survive-def$

using $ccdfTx-continuous-on-nonneg continuous-on-subset Ioi-le-Ico x-lt-\psi$ **by** $blast$

hence $\forall t \in \{0 < ..\}. isCont (\lambda t. \$p-\{t&x\}) t$

using $continuous-on-eq-continuous-at open-greaterThan$ **by** $blast$

thus $isCont (\lambda t. \$p-\{t&x\}) (real k+1)$ **by** $force$

qed

lemma $curt-e-sum-p-smooth: \$e-x = (\sum k. \$p-\{k+1&x\})$ **if** $summable (\lambda k. \$p-\{k+1&x\})$

using $curt-e-sum-p$ $isCont-p-nat$ **that by** $simp$

lemma $curt-e-rec-smooth: \$e-x = \$p-x * (1 + \$e-(x+1))$ **if** $summable (\lambda k. \$p-\{k+1&x\})$

$x+1 < \psi$

using $curt-e-rec$ $isCont-p-nat$ **that by** $simp$

lemma $curt-e-sum-p-finite-smooth: \$e-x = (\sum k < n. \$p-\{k+1&x\})$ **if** $x+n+1 > \psi$ **for** $n::nat$

using $curt-e-sum-p-finite$ $isCont-p-nat$ **that by** $simp$

lemma $temp-curt-e-sum-p-smooth: \$e-\{x:n\} = (\sum k < n. \$p-\{k+1&x\})$ **for** $n::nat$

using $temp-curt-e-sum-p$ $isCont-p-nat$ **by** $simp$

lemma $temp-curt-e-rec-smooth: \$e-\{x:n\} = \$p-x * (1 + \$e-\{x+1:n-1\})$

if $x+1 < \psi$ $n \neq 0$ **for** $n::nat$

using $temp-curt-e-rec$ $isCont-p-nat$ **that by** $simp$

end

end

5.3 Limited Survival Function

locale $limited-survival-function = survival-model +$

assumes $\psi-limited[simp]: \psi < \infty$

begin

definition $ult-age :: nat (\langle \$\omega \rangle)$

where $\omega \equiv \text{LEAST } x::\text{nat}. \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x = 0$
— the conventional notation for ultimate age

```

lemma ccdfX-ceil-psi-0:  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) [\text{real-of-ereal } \psi] = 0$ 
proof -
  have  $\text{real-of-ereal } \psi \leq [\text{real-of-ereal } \psi]$  by simp
  thus ?thesis using ccdfX-0-equiv psi-limited ccdfX-psi-0 le-ereal-le by presburger
qed

lemma ccdfX-omega-0:  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) \omega = 0$ 
  unfolding ult-age-def
proof (rule LeastI-ex)
  have  $[\text{real-of-ereal } \psi] \geq 0$  using psi-nonneg real-of-ereal-pos by fastforce
  thus  $\exists x::\text{nat}. \text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)(\text{real } x) = 0$ 
    using ccdfX-ceil-psi-0 by (metis of-int-of-nat-eq zero-le-imp-eq-int)
qed

corollary psi-le-omega:  $\psi \leq \omega$ 
  using ccdfX-0-equiv ccdfX-omega-0 by blast

corollary omega-pos:  $\omega > 0$ 
  using psi-le-omega order.strict-iff-not by fastforce

lemma omega-ceil-psi:  $\omega = [\text{real-of-ereal } \psi]$ 
proof (rule antisym)
  let ?cpsi =  $[\text{real-of-ereal } \psi]$ 
  have  $\star: ?\text{cpsi} \geq 0$  using psi-nonneg real-of-ereal-pos by fastforce
  moreover have  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) ?\text{cpsi} = 0$  by (rule ccdfX-ceil-psi-0)
  ultimately have  $\omega \leq \text{nat } ?\text{cpsi}$ 
    unfolding ult-age-def using wellorder-Least-lemma of-nat-nat by (smt (verit))
    thus  $\text{int } \omega \leq ?\text{cpsi}$  using le-nat-iff  $\star$  by blast
next
  show  $[\text{real-of-ereal } \psi] \leq \text{int } \omega$ 
    using psi-le-omega by (simp add: ceiling-le-iff real-of-ereal-ord-simps(2))
qed

lemma ccdfX-0-equiv-nat:  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) x = 0 \longleftrightarrow x \geq \omega$  for  $x::\text{nat}$ 
proof
  assume  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)(\text{real } x) = 0$ 
  thus  $x \geq \omega$  unfolding ult-age-def using wellorder-Least-lemma by fastforce
next
  assume  $x \geq \omega$ 
  hence  $\text{ereal } (\text{real } x) \geq \psi$  using psi-le-omega le-ereal-le of-nat-mono by blast
  thus  $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)(\text{real } x) = 0$  using ccdfX-0-equiv by simp
qed

lemma psi-le-iff-omega-le:  $\psi \leq x \longleftrightarrow \omega \leq x$  for  $x::\text{nat}$ 
  using ccdfX-0-equiv ccdfX-0-equiv-nat by presburger

```

```

context
  fixes  $x::nat$ 
  assumes  $x\text{-lt-}\omega[\text{simp}]: x < \$\omega$ 
begin

lemma  $x\text{-lt-}\psi[\text{simp}]: x < \$\psi$ 
  using  $x\text{-lt-}\omega \psi\text{-le-iff-}\omega\text{-le}$  by (meson linorder-not-less)

lemma  $p\text{-0-1-nat}: \$p\{-0\&x\} = 1$ 
  using  $p\text{-0-1}$  by simp

lemma  $p\text{-0-equiv-nat}: \$p\{-t\&x\} = 0 \longleftrightarrow x+t \geq \$\omega$  for  $t::nat$ 
  using  $\psi\text{-le-iff-}\omega\text{-le } p\text{-0-equiv}$  by (metis of-nat-add  $x\text{-lt-}\psi$ )

lemma  $q\text{-0-0-nat}: \$q\{-0\&x\} = 0$ 
  using  $p\text{-q-1 } p\text{-0-1-nat}$  by (smt (verit)  $x\text{-lt-}\psi$ )

lemma  $q\text{-1-equiv-nat}: \$q\{-t\&x\} = 1 \longleftrightarrow x+t \geq \$\omega$  for  $t::nat$ 
  using  $p\text{-q-1 } p\text{-0-equiv-nat}$  by (smt (verit)  $x\text{-lt-}\psi$ )

lemma  $q\text{-defer-old-0-nat}: \$q\{-f|t\&x\} = 0$  if  $\$w \leq x+f$  for  $f t :: nat$ 
  using  $q\text{-defer-old-0 } \psi\text{-le-iff-}\omega\text{-le that}$  by (metis of-nat-0-le-iff of-nat-add  $x\text{-lt-}\psi$ )

lemma  $curt-e\text{-sum-}P\text{-finite-nat}: \$e\text{-}x = (\sum k < n. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$ 
  if  $x+n \geq \$\omega$  for  $n::nat$ 
  apply (rule  $curt-e\text{-sum-}P\text{-finite}$ , simp)
  using that  $\psi\text{-le-iff-}\omega\text{-le}$  by (smt (verit) le-ereal-less of-nat-add)

lemma  $curt-e\text{-sum-}p\text{-finite-nat}: \$e\text{-}x = (\sum k < n. \$p\{-k+1\&x\})$ 
  if  $\bigwedge k::nat. k < n \implies isCont(\lambda t. \$p\{-t\&x\})$  (real  $k + 1$ )  $x+n \geq \$\omega$  for  $n::nat$ 
  apply (rule  $curt-e\text{-sum-}p\text{-finite}$ , simp-all add: that)
  using that  $\psi\text{-le-iff-}\omega\text{-le}$  by (smt (verit) le-ereal-less of-nat-add)

end

lemma  $q\text{-}\omega\text{-1}: \$q\{(\$w - 1)\} = 1$ 
  using  $q\text{-1-equiv-nat}$ 
  by (metis diff-less dual-order.refl le-diff-conv of-nat-1 omega-pos zero-less-one)

end

end
theory Life-Table
  imports Survival-Model
begin

```

6 Life Table

Define a life table axiomatically.

```
locale life-table =
  fixes l :: real ⇒ real (|$l'--> [101] 200)
  assumes l-0-pos: 0 < l 0
  and l-neg-nil: ∀x. x ≤ 0 ⇒ l x = l 0
  and l-PInfty-0: (l —> 0) at-top
  and l-antimono: antimono l
  and l-right-continuous: ∀x. continuous (at-right x) l
begin
```

6.1 Basic Properties of Life Table

```
lemma l-0-neq-0[simp]: $l-0 ≠ 0
  using l-0-pos by simp
```

```
lemma l-nonneg[simp]: $l-x ≥ 0 for x::real
  using l-antimono l-PInfty-0 by (rule antimono-at-top-le)
```

```
lemma l-bounded[simp]: $l-x ≤ $l-0 for x::real
  using l-neg-nil l-antimono by (smt (verit) antimonoD)
```

```
lemma l-measurable[measurable, simp]: l ∈ borel-measurable borel
  by (rule borel-measurable-antimono, rule l-antimono)
```

```
lemma l-left-continuous-nonpos: continuous (at-left x) l if x ≤ 0 for x::real
proof -
  have $l-x = $l-0 using l-neg-nil that by blast
  moreover have continuous (at-left x) (λ-. $l-0) by simp
  moreover have eventually (λy. $l-y = $l-0) (at-left x)
    apply (rule eventually-at-leftI[of x-1], simp-all)
    using that l-neg-nil by (smt (verit))
  ultimately show ?thesis by (rewrite continuous-at-within-cong[where g=λ-.
$l-0]; simp)
qed
```

```
lemma l-integrable-Icc: set-integrable lborel {a..b} l for a b :: real
  unfolding set-integrable-def
  apply (rule integrableI-bounded-set[where A={a..b} and B=$l-0], simp-all)
  using emeasure-compact-finite by auto
```

```
corollary l-integrable-on-Icc: l integrable-on {a..b} for a b :: real
  using l-integrable-Icc by (rewrite integrable-on-iff-set-integrable-nonneg[THEN
sym]; simp)
```

```
lemma l-integrable-Icc-shift: set-integrable lborel {a..b} (λt. $l-(x+t)) for a b x :: real
  using set-integrable-Icc-shift l-integrable-Icc by (metis (full-types) add-diff-cancel-right')
```

```

corollary l-integrable-on-Icc-shift: ( $\lambda t. \$l\cdot(x+t)$ ) integrable-on {a..b} for x a b :: real
  using l-integrable-Icc-shift by (rewrite integrable-on-iff-set-integrable-nonneg[THEN sym]; simp)

lemma l-normal-antimono: antimono ( $\lambda x. \$l\cdot x / \$l\cdot 0$ )
  using divide-le-cancel l-0-pos l-antimono unfolding antimono-def by fastforce

lemma compl-l-normal-right-continuous: continuous (at-right x) ( $\lambda x. 1 - \$l\cdot x / \$l\cdot 0$ ) for x::real
  using l-0-pos l-right-continuous by (intro continuous-intros; simp)

lemma compl-l-normal-NInfty-0: (( $\lambda x. 1 - \$l\cdot x / \$l\cdot 0$ ) —> 0) at-bot
  apply (rewrite tendsto-cong[where g= $\lambda \cdot. 0$ ], simp-all)
  by (smt (verit) div-self eventually-at-bot-linorder l-0-pos l-neg-nil)

lemma compl-l-normal-PInfty-1: (( $\lambda x. 1 - \$l\cdot x / \$l\cdot 0$ ) —> 1) at-top
  using l-0-pos l-PInfty-0 by (intro tendsto-eq-intros) simp-all+

lemma compl-l-real-distribution: real-distribution (interval-measure ( $\lambda x. 1 - \$l\cdot x / \$l\cdot 0$ ))
  using l-normal-antimono compl-l-normal-right-continuous
    compl-l-normal-NInfty-0 compl-l-normal-PInfty-1
  by (intro real-distribution-interval-measure; simp add: antimono-def)

definition total :: real  $\Rightarrow$  real ( $\langle \$T' \rightarrow [101] 200 \rangle$  where  $\$T\cdot x \equiv LBINT y:\{x..\}.$ 
 $\$l\cdot y$ 
  — the number of lives older than the ones aged x
  — The parameter x must be nonnegative.

lemma T-nonneg[simp]:  $\$T\cdot x \geq 0$  for x::real
  unfolding total-def by simp

definition total-finite  $\equiv$  set-integrable lborel {0..} l

lemma total-finite-iff-set-integrable-Ici:
  total-finite  $\longleftrightarrow$  set-integrable lborel {x..} l for x::real
  unfolding total-finite-def using set-integrable-Ici-equiv l-integrable-Icc by blast

lemma total-finite-iff-integrable-on-Ici: total-finite  $\longleftrightarrow$  l integrable-on {x..} for x::real
  using total-finite-iff-set-integrable-Ici integrable-on-iff-set-integrable-nonneg l-nonneg
  by (metis atLeast-borel l-measurable measurable-lborel2 sets-lborel)

lemma total-finite-iff-summable: total-finite  $\longleftrightarrow$  summable ( $\lambda k. \$l\cdot(x+k)$ ) for x::real
  apply (rewrite total-finite-iff-set-integrable-Ici)
  apply (rule set-integrable-iff-summable[of x, simplified], simp-all)
  using l-antimono unfolding antimono-def monotone-on-def by simp

```

```

lemma T-tendsto-0: (( $\lambda x. \$T\cdot x$ ) —> 0) at-top if total-finite
proof -
have  $\bigwedge x. x \geq 0 \implies \$T\cdot x = \$T\cdot 0 - (\text{LBINT } y:\{0..x\}. \$l\cdot y)$ 
proof -
fix  $x::\text{real}$  assume  $asm: x \geq 0$ 
let  $?A = \{x..\}$  and  $?B = \{0..x\}$ 
have  $\{0..\} = ?A \cup ?B$  using  $asm$  by auto
thus  $\$T\cdot x = \$T\cdot 0 - (\text{LBINT } y:\{0..x\}. \$l\cdot y)$ 
  unfolding total-def apply (rewrite eq-diff-eq)
  using that total-finite-iff-set-integrable-Ici l-integrable-Icc
  apply (rewrite set-integral-Un-AE[THEN sym], simp-all)
  using AE-lborel-singleton add-0 asm le-add-same-cancel2 le-numeral-extra(3)
by force
qed
hence  $\forall_F x \text{ in at-top}. \$T\cdot x = \$T\cdot 0 - (\text{LBINT } y:\{0..x\}. \$l\cdot y)$ 
  by (rule eventually-at-top-linorderI[of 0])
moreover have  $((\lambda x. \text{LBINT } y:\{0..x\}. \$l\cdot y) \longrightarrow \$T\cdot 0)$  at-top
  using that unfolding total-def total-finite-def
  by (intro tendsto-set-lebesgue-integral-at-top; simp)
ultimately show ?thesis
apply (rewrite tendsto-cong, simp-all)
using LIM-zero-iff' by force
qed

definition lives :: real  $\Rightarrow$  real  $\Rightarrow$  real ( $\langle\$L'\cdot\{\cdot\&\cdot\}\rangle [0,0] 200$ )
where  $\$L\cdot\{n\&x\} \equiv \text{LBINT } y:\{x..x+n\}. \$l\cdot y$ 
— the number of lives between ages  $x$  and  $x+n$ 
— The parameter  $x$  must be nonnegative.
— The parameter  $n$  is usually nonnegative, but theoretically it can be negative.

abbreviation lives-1 :: real  $\Rightarrow$  real ( $\langle\$L'\cdot\{\cdot\}\rangle [101] 200$ )
where  $\$L\cdot x \equiv \$L\cdot\{1\&x\}$ 

lemma l-has-integral-L: ( $l$  has-integral  $\$L\cdot\{n\&x\}$ )  $\{x..x+n\}$  for  $x n :: \text{real}$ 
unfolding lives-def by (rule has-integral-set-integral-real) (rule l-integrable-Icc)

lemma L-neg-0[simp]:  $\$L\cdot\{n\&x\} = 0$  if  $n < 0$  for  $x n :: \text{real}$ 
unfolding lives-def using that by (rewrite to {} atLeastAtMost-empty; simp)

lemma L-nonneg[simp]:  $\$L\cdot\{n\&x\} \geq 0$  for  $x n :: \text{real}$ 
unfolding lives-def by simp

lemma L-T:  $\$L\cdot\{n\&x\} = \$T\cdot x - \$T\cdot(x+n)$  if total-finite  $n \geq 0$  for  $x n :: \text{real}$ 
proof -
have  $\{x..x+n\} \cup \{x+n..\} = \{x..\}$  using that by force
moreover have
   $(\text{LBINT } y:\{x..x+n\} \cup \{x+n..\}. \$l\cdot y) = (\text{LBINT } y:\{x..x+n\}. \$l\cdot y) + (\text{LBINT } y:\{x+n..\}. \$l\cdot y)$ 

```

proof –

have $\text{AE } y \text{ in } \text{lborel}. \neg (y \in \{x..x+n\} \wedge y \in \{x+n..\})$ by (rule $\text{AE-}I'[\text{where } N=\{x+n\}]; \text{force}$)
moreover have $\text{set-integrable } \text{lborel } \{x..x+n\} l$ by (rule $\text{l-integrable-}Icc$)
moreover have $\text{set-integrable } \text{lborel } \{x+n..\} l$
using that $\text{total-finite-iff-set-integrable-}Ici$ by simp
ultimately show ?thesis by (intro $\text{set-integral-}Un-\text{AE}$; simp)
qed
ultimately show ?thesis unfolding $\text{total-def lives-def}$ by simp
qed

lemma $L\text{-sums-}T: (\lambda k. \$L-(x+k)) \text{ sums } \$T-x$ **if** total-finite **for** $x::\text{real}$
proof –

have $(\lambda k::\text{nat}. \$T-(x+k)) \longrightarrow 0$
using $T\text{-tendsto-}0$
apply (rule $\text{filterlim-compose}[\text{where } f=\lambda k::\text{nat}. x+k \text{ and } g=\text{total}], \text{simp add: that})$
using $\text{filterlim-real-sequentially filterlim-tendsto-add-at-top}$ by blast
hence $(\lambda k. \$T-(x+k) - \$T-(x + \text{Suc } k)) \text{ sums } \$T-x$
by (simp) (rule $\text{telescope-sums}'[\text{of } \lambda k. \$T-(x+k) 0, \text{ simplified}]$)
thus ?thesis using that $L\text{-}T$ by (rewrite $\text{sums-cong}, \text{simp-all}$) (smt (verit))
qed

definition $\text{death} :: \text{real} \Rightarrow \text{real} (\langle \$d' \rangle [0,0] 200)$
where $\$d\{-t\&x\} \equiv \max 0 (\$l-x - \$l-(x+t))$
— the number of deaths between ages x and $x+t$
— The parameter t is usually nonnegative, but theoretically it can be negative.

abbreviation $\text{death1} :: \text{real} \Rightarrow \text{real} (\langle \$d' \rangle [101] 200)$
where $\$d-x \equiv \$d\{-1\&x\}$

lemma $\text{death-def-nonneg}: \$d\{-t\&x\} = \$l-x - \$l-(x+t)$ **if** $t \geq 0$ **for** $t x :: \text{real}$
using that l-antimono unfolding $\text{death-def antimono-def}$ by simp

lemma $d\text{-nonpos-}0: \$d\{-t\&x\} = 0$ **if** $t \leq 0$ **for** $t x :: \text{real}$
using that l-antimono unfolding $\text{death-def antimono-def}$ by simp

corollary $d\text{-}0\text{-}0: \$d\{-0\&x\} = 0$ **for** $x::\text{real}$
using $d\text{-nonpos-}0$ by simp

lemma $d\text{-nonneg}[\text{simp}]: \$d\{-t\&x\} \geq 0$ **for** $t x :: \text{real}$
unfolding death-def by simp

lemma $dx\text{-}l: \$d-x = \$l-x - \$l-(x+1)$ **for** $x::\text{real}$
using death-def-nonneg by simp

lemma $sum\text{-}dx\text{-}l: (\sum k < n. \$d-(x+k)) = \$l-x - \$l-(x+n)$ **for** $x::\text{real}$ **and** $n::\text{nat}$
proof (induction n)
case 0

```

thus ?case by simp
next
  case (Suc n)
  thus ?case
    using dx-l
    by (metis Set-Interval.comm-monoid-add-class.sum.lessThan-Suc
        add-diff-cancel-left' diff-diff-eq2 of-nat-Suc)
qed

corollary d-sums-l: ( $\lambda k. \$d\{x+k\}$ ) sums $l-x for x::real
  unfolding sums-def
  apply (rewrite sum-dx-l)
  apply (rule tendsto-diff[where b=0, simplified], simp)
  using l-PInfty-0 filterlim-compose filterlim-real-sequentially filterlim-tendsto-add-at-top
        tendsto-const by blast

lemma add-d:  $\$d\{t \& x\} + \$d\{t' \& x+t\} = \$d\{t+t' \& x\}$  if  $t \geq 0$   $t' \geq 0$  for t
 $t' :: real$ 
  using death-def-nonneg that by (smt (verit))

definition die-central :: real  $\Rightarrow$  real  $\Rightarrow$  real ( $\langle \$m' \{-\&\- \} \rangle [0,0] 200$ )
  where  $\$m\{n \& x\} \equiv \$d\{n \& x\} / \$L\{n \& x\}$ 
    — central death rate

abbreviation die-central-1 :: real  $\Rightarrow$  real ( $\langle \$m' \rightarrow [101] 200$ )
  where  $\$m\{x\} \equiv \$m\{1 \& x\}$ 

```

6.2 Construction of Survival Model from Life Table

```

definition life-table-measure :: real measure ( $\langle \mathfrak{M} \rangle$ )
  where  $\mathfrak{M} \equiv interval\text{-measure } (\lambda x. 1 - \$l\{x\} / \$l\{0\})$ 

lemma prob-space-actuary-MM: prob-space-actuary  $\mathfrak{M}$ 
  unfolding life-table-measure-def using compl-l-real-distribution real-distribution-def
  by (intro prob-space-actuary.intro) force

definition survival-model-X :: real  $\Rightarrow$  real ( $\langle X \rangle$ ) where  $X \equiv \lambda x. x$ 

lemma survival-model-MM-X: survival-model  $\mathfrak{M} X$ 
proof -
  let ?F =  $\lambda x. 1 - \$l\{x\} / \$l\{0\}$ 
  show survival-model  $\mathfrak{M} X$ 
    unfolding life-table-measure-def survival-model-X-def
    proof (rule survival-model.intro)
      show prob-space-actuary (interval-measure ?F)
        using prob-space-actuary-MM unfolding life-table-measure-def by simp
        show survival-model-axioms (interval-measure ?F) ( $\lambda x. x$ )
      proof -
        have [simp]:  $\{\xi :: real. \xi \leq 0\} = \{\dots 0\}$  by blast
      qed
    qed
  qed
qed

```

```

have measure (interval-measure ( $\lambda x. 1 - \$l-x / \$l-0$ ))  $\{\dots\} = 0$ 
using l-normal-antimono compl-l-normal-right-continuous compl-l-normal-NInfty-0
  by (rewrite measure-interval-measure-Iic, simp-all add: antimono-def)
hence emeasure (interval-measure ( $\lambda x. 1 - \$l-x / \$l-0$ ))  $\{\dots\} = ennreal 0$ 
  apply (rewrite finite-measure.emeasure-eq-measure, simp-all)
  using compl-l-real-distribution prob-space-def real-distribution-def by blast
thus ?thesis
  apply (intro survival-model-axioms.intro, simp)
  apply (rewrite AE-iff-null, simp)
  by (rewrite not-less) auto
qed
qed
qed

end

sublocale life-table  $\subseteq$  survival-model  $\mathfrak{M} X$ 
by (rule survival-model-MM-X)

context life-table
begin

interpretation distrX-RD: real-distribution distr  $\mathfrak{M}$  borel  $X$ 
using MM-PS.real-distribution-distr by simp

```

6.2.1 Relations between Life Table and Survival Function for X

```

lemma ccdfX-l-normal: ccdf (distr  $\mathfrak{M}$  borel  $X$ ) = ( $\lambda x. \$l-x / \$l-0$ )
proof (rule ext)
let ?F =  $\lambda x. 1 - \$l-x / \$l-0$ 
interpret F-FBM: finite-borel-measure interval-measure ?F
  using compl-l-real-distribution real-distribution.finite-borel-measure-M by blast
show  $\lambda x. ccdf (distr \mathfrak{M} borel X) x = \$l-x / \$l-0$ 
  unfolding ccdf-def life-table-measure-def survival-model-X-def
  apply (rewrite measure-distr, simp-all)
  using l-normal-antimono compl-l-normal-right-continuous
    compl-l-normal-NInfty-0 compl-l-normal-PInfty-1
  by (rewrite F-FBM.measure-interval-measure-Ioi; simp add: antimono-def)
qed

corollary deriv-ccdfX-l: deriv (ccdf (distr  $\mathfrak{M}$  borel  $X$ ))  $x = deriv l x / \$l-0$ 
  if  $l$  differentiable at  $x$  for  $x::real$ 
  using differentiable-eq-field-differentiable-real that
  by (rewrite ccdfX-l-normal, rewrite deriv-cdivide-right; simp)

notation death-pt ( $\langle \$\psi \rangle$ )

lemma l-0-equiv:  $\$l-x = 0 \longleftrightarrow x \geq \$\psi$  for  $x::real$ 
using ccdfX-l-normal ccdfX-0-equiv by simp

```

```

lemma d-old-0: $d-{t&x} = 0 if x ≥ $ψ t ≥ 0 for x t :: real
  unfolding death-def using l-0-equiv that by (smt (verit) le-ereal-le)

lemma d-l-equiv: $d-{t&x} = $l-x ↔ x+t ≥ $ψ if t ≥ 0 for x t :: real
  using death-def-nonneg l-0-equiv that by simp

lemma continuous-ccdfX-l: continuous F (ccdf (distr ℳ borel X)) ↔ continuous
F l
  for F :: real filter
proof -
  have continuous F (ccdf (distr ℳ borel X)) ↔ continuous F (λx. $l-x / $l-0)
    using ccdfX-l-normal by simp
  also have ... ↔ continuous F l using continuous-cdivide-iff l-0-neq-0 by blast
  finally show ?thesis .
qed

lemma has-real-derivative-ccdfX-l:
(ccdf (distr ℳ borel X) has-real-derivative D) (at x) ↔
(l has-real-derivative $l-0 * D) (at x)
for D x :: real
proof -
  have (ccdf (distr ℳ borel X) has-real-derivative D) (at x) ↔
    ((λx. $l-x / $l-0) has-real-derivative D) (at x)
    by (rule has-field-derivative-cong-eventually; simp add: ccdfX-l-normal)
  also have ... ↔ ((λx. $l-0 * ($l-x / $l-0)) has-real-derivative $l-0 * D) (at x)
    by (rule DERIV-cmult-iff, simp)
  also have ... ↔ (l has-real-derivative $l-0 * D) (at x) by simp
  finally show ?thesis .
qed

corollary differentiable-ccdfX-l:
ccdf (distr ℳ borel X) differentiable (at x) ↔ l differentiable (at x)
for D x :: real
using has-real-derivative-ccdfX-l
by (metis l-0-neq-0 mult.commute nonzero-divide-eq-eq real-differentiable-def)

lemma PX-l-normal: ℙ(ξ in ℳ. X ξ > x) = $l-x / $l-0 for x::real
  using MM-PS.ccdf-distr-P ccdfX-l-normal X-RV by (metis (mono-tags, lifting)
Collect-cong)

lemma set-integrable-ccdfX-l:
set-integrable lborel A (ccdf (distr ℳ borel X)) ↔ set-integrable lborel A l
  if A ∈ sets lborel for A :: real set
proof -
  have set-integrable lborel A (ccdf (distr ℳ borel X)) ↔
    set-integrable lborel A (λx. $l-x / $l-0)
    by (rule set-integrable-cong; simp add: ccdfX-l-normal)
  also have ... ↔ set-integrable lborel A l by simp

```

```

finally show ?thesis .
qed

corollary integrable-ccdfX-l: integrable lborel (ccdf (distr M borel X))  $\longleftrightarrow$  integrable lborel l
  using set-integrable-ccdfX-l[where A=UNIV] by (simp add: set-integrable-def)

lemma integrable-on-ccdfX-l:
  ccdf (distr M borel X) integrable-on A  $\longleftrightarrow$  l integrable-on A for A :: real set
proof -
  have ccdf (distr M borel X) integrable-on A  $\longleftrightarrow$  ( $\lambda x. \$l\cdot x / \$l\cdot 0$ ) integrable-on A
    by (rule integrable-cong) (simp add: ccdfX-l-normal)
  also have ...  $\longleftrightarrow$  l integrable-on A
    using integrable-on-cdivide-iff[of \$l·0 l] by simp
  finally show ?thesis .
qed

```

6.2.2 Relations between Life Table and Cumulative Distributive Function for X

```

lemma cdfX-l-normal: cdf (distr M borel X) = ( $\lambda x. 1 - \$l\cdot x / \$l\cdot 0$ ) for x::real
  using ccdfX-l-normal distrX-RD.cdf-ccdf distrX-RD.prob-space by presburger

```

```

lemma deriv-cdfX-l: deriv (cdf (distr M borel X)) x = - deriv l x / \$l·0
  if l differentiable at x for x::real
  using distrX-RD.cdf-ccdf differentiable-eq-field-differentiable-real that differentiable-ccdfX-l
  deriv-diff deriv-ccdfX-l that by simp

```

```

lemma continuous-cdfX-l: continuous F (cdf (distr M borel X))  $\longleftrightarrow$  continuous F l
  for F :: real filter
  using distrX-RD.continuous-cdf-ccdf continuous-ccdfX-l by simp

```

```

lemma has-real-derivative-cdfX-l:
  (cdf (distr M borel X) has-real-derivative D) (at x)  $\longleftrightarrow$ 
  (l has-real-derivative - (\$l·0 * D)) (at x)
  for D x :: real
  using distrX-RD.has-real-derivative-cdf-ccdf has-real-derivative-ccdfX-l by simp

```

```

lemma differentiable-cdfX-l:
  cdf (distr M borel X) differentiable (at x)  $\longleftrightarrow$  l differentiable (at x) for D x :: real
  using differentiable-eq-field-differentiable-real distrX-RD.differentiable-cdf-ccdf
  differentiable-ccdfX-l by simp

```

```

lemma PX-compl-l-normal: P(ξ in M. X ξ ≤ x) = 1 - \$l·x / \$l·0 for x::real
  using PX-l-normal by (metis MM-PS.prob-compl X-compl-gt-le X-gt-event)

```

6.2.3 Relations between Life Table and Survival Function for $T(x)$

context

fixes $x:\text{real}$
assumes $x\text{-lt-psi}[\text{simp}]: x < \ψ

begin

notation $futr-life (\langle T \rangle)$

interpretation alivex-PS : prob-space $\mathfrak{M} \downharpoonright \text{alive } x$
by (rule $\text{MM-PS}.cond\text{-prob\text{-}space\text{-}correct}$, $\text{simp\text{-}all add: alive\text{-}def}$)

interpretation distrTx-RD : real-distribution $\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)$ **by** simp

lemma $lx\text{-neq-0}[\text{simp}]: \$l\text{-}x \neq 0$
using $l\text{-0-equiv } x\text{-lt-psi linorder-not-less by blast}$

corollary $lx\text{-pos}[\text{simp}]: \$l\text{-}x > 0$
using $lx\text{-neq-0 } l\text{-nonneg by (smt (verit))}$

lemma ccdfTx-l-normal : $\text{ccdf}(\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) t = \$l\text{-}(x+t) / \$l\text{-}x$
if $t \geq 0$ **for** $t:\text{real}$
using $\text{ccdfTx-PX } PX\text{-l-normal } l\text{-0-neq-0}$ **that by** simp

lemma deriv-ccdfTx-l :

$\text{deriv}(\text{ccdf}(\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x))) t = \text{deriv}(\lambda t. \$l\text{-}(x+t) / \$l\text{-}x) t$
if $t > 0$ l differentiable at $(x+t)$ **for** $t:\text{real}$

proof –

have $\forall_F s \text{ in nhds } t. \text{ccdf}(\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) s = \$l\text{-}(x+s) / \$l\text{-}x$
apply (rewrite eventually-nhds-metric)
using that $\text{ccdfTx-l-normal dist-real-def by (intro exI[of - t]) auto}$
thus $?thesis$ **by** (rule deriv-cong-ev) simp

qed

lemma $\text{continuous-at-within-ccdfTx-l}$:

$\text{continuous}(\text{at } t \text{ within } \{0..\}) (\text{ccdf}(\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x))) \longleftrightarrow$
 $\text{continuous}(\text{at } (x+t) \text{ within } \{...\}) l$
if $t \geq 0$ **for** $t:\text{real}$
using $\text{continuous-ccdfX-ccdfTx}$ **that continuous-ccdfX-l by force**

lemma isCont-ccdfTx-l :

$\text{isCont}(\text{ccdf}(\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x))) t \longleftrightarrow \text{isCont } l (x+t) \text{ if } t > 0 \text{ for } t:\text{real}$
using that $\text{continuous-ccdfX-l isCont-ccdfX-ccdfTx by force}$

lemma $\text{has-real-derivative-ccdfTx-l}$:

$(\text{ccdf}(\text{distr}(\mathfrak{M} \downharpoonright \text{alive } x) \text{ borel } (T x)) \text{ has-real-derivative } D) (\text{at } t) \longleftrightarrow$
 $(l \text{ has-real-derivative } \$l\text{-}x * D) (\text{at } (x+t))$

```

if  $t > 0$  for  $t D :: \text{real}$ 
proof -
  have ( $\text{ccdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) \text{ has-real-derivative } D$ ) (at  $t$ )  $\longleftrightarrow$ 
    ( $\text{ccdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) \text{ has-real-derivative }$ 
      $(\$l-x / \$l-0 * D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)))$  (at  $t$ )
  using  $PX\text{-}l\text{-normal}$  by force
  also have ... = ( $\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X) \text{ has-real-derivative } (\$l-x / \$l-0 * D)$ )
(at  $(x+t)$ )
  using  $\text{has-real-derivative-ccdf}X\text{-ccdf}Tx$  that by simp
  also have ... = ( $l \text{ has-real-derivative } (\$l-x * D)$ ) (at  $(x+t)$ )
  using  $\text{has-real-derivative-ccdf}X\text{-}l$  by simp
  finally show ?thesis .
qed

```

```

lemma  $\text{differentiable-ccdf}Tx\text{-}l$ :
   $\text{ccdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) \text{ differentiable at } t \longleftrightarrow l \text{ differentiable (at } (x+t))$ 
  if  $t > 0$  for  $t::\text{real}$ 
  using  $\text{differentiable-ccdf}X\text{-ccdf}Tx$   $\text{differentiable-ccdf}X\text{-}l$  that by force

lemma  $P Tx\text{-}l\text{-normal}$ :  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > t \mid T x \xi > 0) = \$l-(x+t) / \$l-x$  if  $t \geq 0$  for  $t::\text{real}$ 
  using  $\text{ccdf}Tx\text{-}l\text{-normal}$  that by (simp add:  $\text{ccdf}Tx\text{-cond-prob}$ )

```

6.2.4 Relations between Life Table and Cumulative Distributive Function for $T(x)$

```

lemma  $\text{cdf}Tx\text{-compl-}l\text{-normal}$ :  $\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t = 1 - \$l-(x+t) / \$l-x$ 
  if  $t \geq 0$  for  $t::\text{real}$ 
  using  $\text{distr}Tx\text{-}RD.\text{cdf-ccdf}$   $\text{ccdf}Tx\text{-}l\text{-normal}$  that  $\text{distr}Tx\text{-}RD.\text{prob-space}$  by auto

lemma  $\text{deriv-cdf}Tx\text{-}l$ :
   $\text{deriv}(\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))) t = - \text{deriv}(\lambda t. \$l-(x+t) / \$l-x) t$ 
  if  $t > 0$   $l$   $\text{differentiable at } (x+t)$  for  $t::\text{real}$ 
  using  $\text{deriv-ccdf}Tx\text{-}l$   $\text{differentiable-cdf}X\text{-}cdf}Tx$   $\text{differentiable-cdf}X\text{-}l$   $\text{distr}Tx\text{-}RD.\text{deriv-cdf-ccdf}$ 
  that by fastforce

lemma  $\text{continuous-at-within-cdf}Tx\text{-}l$ :
   $\text{continuous}(\text{at } t \text{ within } \{0..\}) (\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))) \longleftrightarrow$ 
   $\text{continuous}(\text{at } (x+t) \text{ within } \{..x\}) l$ 
  if  $t \geq 0$  for  $t::\text{real}$ 
  using  $\text{that continuous-cdf}X\text{-}l$   $\text{continuous-cdf}X\text{-}cdf}Tx$  by force

lemma  $\text{isCont-cdf}Tx\text{-}l$ :
   $\text{isCont}(\text{cdf}(\text{distr}(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))) t \longleftrightarrow \text{isCont } l (x+t)$  if  $t > 0$  for  $t::\text{real}$ 
  using  $\text{that continuous-cdf}X\text{-}l$   $\text{isCont-cdf}X\text{-}cdf}Tx$  by force

```

```

lemma has-real-derivative-cdfTx-l:
  (cdf (distr ( $\mathfrak{M} \downarrow \text{alive } x$ ) borel ( $T x$ )) has-real-derivative  $D$ ) (at  $t$ )  $\longleftrightarrow$ 
    ( $l$  has-real-derivative  $- \$l\cdot x * D$ ) (at  $(x+t)$ )
  if  $t > 0$  for  $t D :: \text{real}$ 
  using has-real-derivative-ccdfTx-l that distrTx-RD.has-real-derivative-cdf-ccdf by
  auto

lemma differentiable-cdfTx-l:
  cdf (distr ( $\mathfrak{M} \downarrow \text{alive } x$ ) borel ( $T x$ )) differentiable at  $t \longleftrightarrow l$  differentiable (at
   $(x+t)$ )
  if  $t > 0$  for  $t::\text{real}$ 
  using differentiable-cdfX-l that differentiable-cdfX-cdfTx by auto

lemma PTx-compl-l-normal:  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq t \mid T x \xi > 0) = 1 - \$l\cdot(x+t) / \$l\cdot x$ 
  if  $t \geq 0$  for  $t::\text{real}$ 
  using cdfTx-compl-l-normal that by (simp add: cdfTx-cond-prob)

```

6.2.5 Life Table and Actuarial Notations

```

notation survive ( $\langle \$p' \{-\&\- \} \rangle [0,0] 200$ )
notation survive-1 ( $\langle \$p' \rightarrow [101] 200$ )
notation die ( $\langle \$q' \{-\&\- \} \rangle [0,0] 200$ )
notation die-1 ( $\langle \$q' \rightarrow [101] 200$ )
notation die-defer ( $\langle \$q' \{-|\&\- \} \rangle [0,0,0] 200$ )
notation die-defer-1 ( $\langle \$q' \{-|\&\- \} \rangle [0,0] 200$ )
notation life-expect ( $\langle \$e^{\circ} \rightarrow [101] 200$ )
notation temp-life-expect ( $\langle \$e^{\circ} \{-:-\} \rangle [0,0] 200$ )
notation curt-life-expect ( $\langle \$e' \rightarrow [101] 200$ )
notation temp-curt-life-expect ( $\langle \$e' \{-:-\} \rangle [0,0] 200$ )

```

```

lemma p-l:  $\$p\{-t\&x\} = \$l\cdot(x+t) / \$l\cdot x$  if  $t \geq 0$  for  $t::\text{real}$ 
  unfolding survive-def using ccdfTx-l-normal that by simp

```

```

corollary p-1-l:  $\$p\cdot x = \$l\cdot(x+1) / \$l\cdot x$ 
  using p-l by simp

```

```

lemma isCont-p-l: isCont ( $\lambda s. \$p\{-s\&x\}$ )  $t \longleftrightarrow$  isCont  $l$  ( $x+t$ ) if  $t > 0$  for  $t::\text{real}$ 
proof -
  have  $\forall F s \text{ in nhds } t. \$p\{-s\&x\} = \$l\cdot(x+s) / \$l\cdot x$ 
  apply (rewrite eventually-nhds-metric)
  apply (rule exI[of - t], auto simp add: that)
  by (rewrite p-l; simp add: dist-real-def)
  hence isCont ( $\lambda s. \$p\{-s\&x\}$ )  $t \longleftrightarrow$  isCont ( $\lambda s. \$l\cdot(x+s) / \$l\cdot x$ )  $t$  by (rule
  isCont-cong)
  also have ...  $\longleftrightarrow$  isCont ( $\lambda s. \$l\cdot(x+s)$ )  $t$  using continuous-cdivide-iff lx-neq-0
  by metis
  also have ...  $\longleftrightarrow$  isCont  $l$  ( $x+t$ ) using isCont-shift by (force simp add:
  add.commute)

```

```

finally show ?thesis .
qed

lemma total-finite-iff-p-set-integrable-Ici:
  total-finite  $\longleftrightarrow$  set-integrable lborel {0..} ( $\lambda t. \$p\{t \& x\}$ )
  apply (rewrite set-integrable-cong-AE[where g= $\lambda t. \$l(x+t) / \$l-x$ ; simp])
  using survive-def apply simp
  using p-l apply (intro AE-I2, simp)
  by (metis l-integrable-Icc-shift set-integrable-Ici-equiv set-integrable-Ici-shift
       total-finite-iff-set-integrable-Ici)

lemma p-PTx-ge-l-isCont:  $\$p\{t \& x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq t \mid T x \xi > 0)$ 
  if isCont l (x+t) t > 0 for t::real
  using p-PTx-ge-ccdf-isCont that continuous-ccdfX-l by force

lemma q-defer-l:  $\$q\{f \mid t \& x\} = (\$l(x+f) - \$l(x+f+t)) / \$l-x$  if  $f \geq 0$   $t \geq 0$  for
 $f t :: real$ 
  apply (rewrite q-defer-p, simp-all add: that)
  using that by (rewrite p-l, simp)+ (smt (verit) diff-divide-distrib)

corollary q-defer-d-l:  $\$q\{f \mid t \& x\} = \$d\{t \& x+f\} / \$l-x$  if  $f \geq 0$   $t \geq 0$  for  $f t :: real$ 
  using q-defer-l that death-def-nonneg by simp

corollary q-defer-1-d-l:  $\$q\{f \mid t \& x\} = \$d(x+f) / \$l-x$  if  $f \geq 0$  for  $f :: real$ 
  using q-defer-d-l that by simp

lemma q-d-l:  $\$q\{t \& x\} = \$d\{t \& x\} / \$l-x$  for  $t :: real$ 
proof (cases t ≥ 0)
  case True
  thus ?thesis using q-defer-d-l[of 0] by simp
next
  case False
  thus ?thesis using q-nonpos-0 d-nonpos-0 by simp
qed

corollary q-1-d-l:  $\$q\{x\} = \$d\{x\} / \$l-x$ 
  using q-d-l by simp

lemma LBINT-p-l:  $(LBINT t:A. \$p\{t \& x\}) = (LBINT t:A. \$l(x+t)) / \$l-x$ 
  if  $A \subseteq \{0..\}$   $A \in sets$  lborel for  $A :: real$  set
  — Note that  $0 = 0$  holds when the integral diverges.
proof -
  have [simp]:  $\bigwedge t. t \in A \implies \$p\{t \& x\} = \$l(x+t) / \$l-x$  using p-l that by blast
  hence  $(LBINT t:A. \$p\{t \& x\}) = (LBINT t:A. \$l(x+t)) / \$l-x$ 
    using that by (rewrite set-lebesgue-integral-cong[where g= $\lambda t. \$l(x+t) / \$l-x$ ; simp])
  also have ... =  $(LBINT t:A. \$l(x+t)) / \$l-x$  by (rewrite set-integral-divide-zero)
  simp

```

finally show ?thesis .

qed

corollary e-LBINT-l: \$e`o-x = (LBINT t:{0..}. \$l-(x+t)) / \$l-x

— Note that $0 = 0$ holds when the integral diverges.

by (simp add: e-LBINT-p LBINT-p-l)

corollary e-LBINT-l-Icc: \$e`o-x = (LBINT t:{0..n}. \$l-(x+t)) / \$l-x **if** $x+n \geq \psi$ **for** $n::real$

using e-LBINT-p-Icc **by** (rewrite LBINT-p-l[THEN sym]; simp add: that)

lemma temp-e-LBINT-l: \$e`o-{x:n} = (LBINT t:{0..n}. \$l-(x+t)) / \$l-x **if** $n \geq 0$ **for** $n::real$

using temp-e-LBINT-p **by** (rewrite LBINT-p-l[THEN sym]; simp add: that)

lemma integral-p-l: integral A ($\lambda t. \$p\{-t\&x\}$) = (integral A ($\lambda t. \$l-(x+t)$)) / \$l-x
if $A \subseteq \{0..\}$ $A \in sets$ lborel **for** $A :: real$ set

— Note that $0 = 0$ holds when the integral diverges.

using that apply (rewrite set-borel-integral-eq-integral-nonneg[THEN sym], simp-all)

apply (simp add: survive-def)

apply (rewrite set-borel-integral-eq-integral-nonneg[THEN sym], simp-all)

by (rule LBINT-p-l; simp)

corollary e-integral-l: \$e`o-x = integral {0..} ($\lambda t. \$l-(x+t)$) / \$l-x

— Note that $0 = 0$ holds when the integral diverges.

by (simp add: e-integral-p integral-p-l)

corollary e-integral-l-Icc:

\$e`o-x = integral {0..n} ($\lambda t. \$l-(x+t)$) / \$l-x **if** $x+n \geq \psi$ **for** $n::real$

using e-integral-p-Icc **by** (rewrite integral-p-l[THEN sym]; simp add: that)

lemma e-pos-total-finite: \$e`o-x > 0 **if** total-finite

using e-pos total-finite-iff-p-set-integrable-Ici that **by** simp

lemma temp-e-integral-l:

\$e`o-{x:n} = integral {0..n} ($\lambda t. \$l-(x+t)$) / \$l-x **if** $n \geq 0$ **for** $n::real$

using temp-e-integral-p **by** (rewrite integral-p-l[THEN sym]; simp add: that)

lemma curt-e-sum-l: \$e-x = ($\sum k. \$l-(x+k+1)$) / \$l-x **if** total-finite $\wedge k::nat$. is-Cont l (x+k+1)

proof —

have summable ($\lambda k. \$l-(x+(k+1::nat))$)

using that total-finite-iff-summable **by** (rewrite summable-iff-shift[of $\lambda k. \$l-(x+k+1)$]) simp

moreover hence summable ($\lambda k. \$p\{-k+1\&x\}$) **by** (rewrite p-l, simp-all add: add.commute)

moreover have $\bigwedge k::nat$. isCont ($\lambda t. \$p\{-t\&x\}$) (k+1)

using isCont-p-l that **by** (simp add: add.assoc)

ultimately show ?thesis

```

apply (rewrite curt-e-sum-p, simp-all)
apply (rewrite p-l, simp)
by (rewrite suminf-divide) (simp add: add.commute, simp add: add.assoc)
qed

lemma curt-e-sum-l-finite: $e-x = ( $\sum k < n. \$l-(x+k+1)$ ) / $l-x
if  $\bigwedge k : \text{nat}. k < n \implies \text{isCont } l(x+k+1) x+n+1 > \$\psi$  for  $n : \text{nat}$ 
apply (rewrite curt-e-sum-p-finite[of x n], simp-all add: that)
using isCont-p-l that apply (simp add: add.assoc)
apply (rewrite sum-divide-distrib, rule sum.cong, simp)
using p-l by (smt (verit) of-nat-0-le-iff)

lemma temp-curt-e-sum-p: $e-\{x:n\} = ( $\sum k < n. \$l-(x+k+1)$ ) / $l-x
if  $\bigwedge k : \text{nat}. k < n \implies \text{isCont } l(x+k+1)$  for  $n : \text{nat}$ 
apply (rewrite temp-curt-e-sum-p[of x n], simp-all add: that)
using isCont-p-l that apply (simp add: add.assoc)
apply (rewrite sum-divide-distrib, rule sum.cong, simp)
using p-l by (smt (verit) of-nat-0-le-iff)

lemma e-T-l: $e^{\circ}-x = $T-x / $l-x
unfolding total-def
apply (rewrite e-LBINT-l, simp-all)
by (metis add-cancel-left-left diff-add-cancel lborel-set-integral-Ici-shift)

lemma temp-e-L-l: $e^{\circ}-\{x:n\} = $L-\{n&x\} / $l-x if  $n \geq 0$  for  $n : \text{real}$ 
unfolding lives-def using that
apply (rewrite temp-e-LBINT-l, simp-all)
using diff-self add-diff-cancel-left' lborel-set-integral-Icc-shift by metis

lemma m-q-e: $m-\{n&x\} = $q-\{n&x\} / $e^{\circ}-\{x:n\} if  $n \geq 0$  for  $n : \text{real}$ 
proof -
have $m-\{n&x\} = ($d-\{n&x\} / $l-x) / ($L-\{n&x\} / $l-x) unfolding die-central-def
by simp
thus ?thesis using q-d-l temp-e-L-l that by simp
qed

end

lemma l-p: $l-x / $l-0 = $p-\{x&0\} for x::real
using ccdfX-l-normal ccdfX-p by force

lemma e-p-e-total-finite: $e^{\circ}-x = $e^{\circ}-\{x:n\} + $p-\{n&x\} * $e^{\circ}-(x+n)
if total-finite  $n \geq 0$   $x+n < \$\psi$  for  $x n :: \text{real}$ 
using e-p-e that total-finite-iff-p-set-integrable-Ici by (smt (verit) ereal-less-le)

proposition x-ex-const-equiv-total-finite:  $x + \$e^{\circ}-x = y + \$e^{\circ}-y \longleftrightarrow \$q-\{y-x&x\} = 0$ 
if total-finite  $x \leq y$   $y < \$\psi$  for  $x y :: \text{real}$ 
using x-ex-const-equiv that total-finite-iff-p-set-integrable-Ici p-set-integrable-shift

```

```

by blast

corollary x-ex-const-iff-l-const: x + $e'^o-x = y + $e'^o-y  $\longleftrightarrow$  $l-x = $l-y
  if total-finite x ≤ y y < $ψ for x y :: real
  using x-ex-const-equiv-total-finite that
  by (smt (verit, ccfv-threshold) divide-cancel-right ereal-less-le
    l-0-equiv life-table.death-def-nonneg life-table.q-d-l life-table-axioms q-1-equiv)

end

```

6.3 Piecewise Differentiable Life Table

```

locale smooth-life-table = life-table +
  assumes l-piecewise-differentiable[simp]: l piecewise-differentiable-on UNIV
begin

lemma smooth-survival-function-MM-X: smooth-survival-function ℳ X
  proof (rule smooth-survival-function.intro)
    show survival-model ℳ X by (rule survival-model-axioms)
    show smooth-survival-function-axioms ℳ X
    proof
      show ccdf (distr ℳ borel X) piecewise-differentiable-on UNIV
        apply (rewrite ccdfX-l-normal)
        apply (rewrite divide-inverse, rewrite mult.commute)
        using l-piecewise-differentiable piecewise-differentiable-scaleR[of l] by simp
    qed
  qed

end

sublocale smooth-life-table ⊆ smooth-survival-function ℳ X
  by (rule smooth-survival-function-MM-X)

context smooth-life-table
begin

notation force-mortal (⟨$μ'--> [101] 200)

lemma l-continuous[simp]: continuous-on UNIV l
  using l-piecewise-differentiable piecewise-differentiable-on-imp-continuous-on by
  fastforce

lemma l-nondifferentiable-finite-set[simp]: finite {x. ¬ l differentiable at x}
  using differentiable-ccdfX-l ccdfX-nondifferentiable-finite-set by simp

lemma l-differentiable-borel-set[measurable, simp]: {x. l differentiable at x} ∈ sets borel
  using differentiable-ccdfX-l ccdfX-differentiable-borel-set by simp

```

lemma *l-differentiable-AE*: *AE x in lborel. l differentiable at x*
using *differentiable-ccdfX-l ccdfX-differentiable-AE* **by** *simp*

lemma *deriv-l-measurable[measurable]*: *deriv l ∈ borel-measurable borel*
proof –
 let $?S = \{x. \neg l \text{ differentiable at } x\}$
 have $\bigwedge x. x \notin ?S \implies \$l\text{-}0 * \text{deriv}(\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)) x = \text{deriv } l x$
using *deriv-ccdfX-l* **by** *simp*
 thus $?thesis$
 apply –
 by (rule measurable-discrete-difference
 [where $X=?S$ and $f=\lambda x. \$l\text{-}0 * \text{deriv}(\text{ccdf}(\text{distr } \mathfrak{M} \text{ borel } X)) x$])
 (simp-all add: countable-finite)
qed

lemma *pdfX-l-normal*:
 $\text{pdf}X x = (\text{if } l \text{ differentiable at } x \text{ then } -\text{deriv } l x / \$l\text{-}0 \text{ else } 0)$ **for** $x:\text{real}$
unfolding *pdfX-def*
using *differentiable-eq-field-differentiable-real differentiable-cdfX-l deriv-cdfX-l* **by** *simp*

lemma *mu-deriv-l*: $\$μ\text{-}x = -\text{deriv } l x / \$l\text{-}x$ **if** *l differentiable at x* **for** $x:\text{real}$
using *mu-pdfX* that *ccdfX-l-normal* that *pdfX-l-normal* **by** (simp add: differentiable-cdfX-l)

lemma *mu-nonneg-differentiable-l*: $\$μ\text{-}x \geq 0$ **if** *l differentiable at x* **for** $x:\text{real}$
using *differentiable-cdfX-l mu-nonneg-differentiable* that **by** *simp*

lemma *mu-deriv-ln-l*:
 $\$μ\text{-}x = -\text{deriv}(\lambda x. \ln(\$l\text{-}x)) x$ **if** *l differentiable at x* $x < \$ψ$ **for** $x:\text{real}$
proof –
 have $\forall_F x \text{ in nhds } x. \ln(\$l\text{-}x / \$l\text{-}0) = \ln(\$l\text{-}x) - \ln(\$l\text{-}0)$
proof (cases $\langle \$ψ < ∞ \rangle$)
 case *True*
 thus $?thesis$
 apply (rewrite eventually-nhds-metric)
 apply (intro exI[of - real-of-ereal \$ψ - x], auto)
 using that *True* not-inftyI apply fastforce
 apply (rewrite ln-div, simp-all)
 using lx-pos dist-real-def not-inftyI that(2) **by** fastforce
 next
 case *False*
 hence $\bigwedge x. \$l\text{-}x > 0$ **using** *l-0-equiv* **by** force
 thus $?thesis$
 by (simp add: ln-divide-pos)
qed
 hence $\text{deriv}(\lambda x. \ln(\$l\text{-}x / \$l\text{-}0)) x = \text{deriv}(\lambda x. \ln(\$l\text{-}x)) x$
 apply (rewrite deriv-cong-ev[of - λx. ln(\$l-x) - ln(\$l-0)], simp-all)
 apply (rewrite deriv-diff, simp-all)

unfolding field-differentiable-def **using** that
by (metis DERIV-ln-divide-chain lx-pos real-differentiable-def)
thus ?thesis **using** ccdfX-l-normal mu-deriv-ln that differentiable-ccdfX-l **by** force
qed

lemma deriv-l-shift: deriv l (x+t) = deriv ($\lambda t. \$l(x+t)$) t
if l differentiable at (x+t) **for** x t :: real
using deriv-shift differentiable-eq-field-differentiable-real that **by** simp

context

fixes x::real
assumes x-lt-psi[simp]: $x < \$\psi$
begin

lemma p-mu-l: $\$p\{t \& x\} * \$\mu-(x+t) = - \text{deriv } l(x+t) / \$l-x$
if l differentiable at (x+t) t > 0 x+t < \$ψ **for** t::real
using p-l mu-deriv-l that **by** simp

lemma p-mu-l-AE: AE s in lborel. $0 < s \wedge x+s < \$\psi \longrightarrow \$p\{s \& x\} * \$\mu-(x+s) = - \text{deriv } l(x+s) / \$l-x$

proof –

have AE s in lborel. l differentiable at (x+s)
apply (rule AE-borel-affine[of 1 $\lambda u. l$ differentiable at u x, simplified])
unfolding pred-def **using** l-differentiable-AE **by** simp-all
moreover have AE s in lborel.
 l differentiable at (x+s) $\longrightarrow 0 < s \wedge x+s < \$\psi \longrightarrow \$p\{s \& x\} * \$\mu-(x+s) = - \text{deriv } l(x+s) / \$l-x$
using p-mu-l **by** (intro AE-I2) simp
ultimately show ?thesis **by** (rule AE-mp)
qed

lemma LBINT-l-mu-q: $(\text{LBINT } s:\{f <.. f+t\}. \$l-(x+s) * \$\mu-(x+s)) / \$l-x = \$q\{f | t \& x\}$
if t ≥ 0 f ≥ 0 **for** t f :: real

proof –

have $\bigwedge s. s \in \{f <.. f+t\} \implies \$p\{s \& x\} = \$l-(x+s) / \$l-x$ **using** p-l that **by** simp
hence $\$q\{f | t \& x\} = (\text{LBINT } s:\{f <.. f+t\}. \$l-(x+s) / \$l-x * \$\mu-(x+s))$
using LBINT-p-mu-q-defer
by (smt (verit) greaterThanAtMost-borel set-lebesgue-integral-cong sets-lborel
 that x-lt-psi)
also have ... = $(\text{LBINT } s:\{f <.. f+t\}. \$l-(x+s) * \$\mu-(x+s)) / \$l-x$
using set-integral-divide-zero **by** simp
finally show ?thesis **by** simp
qed

lemma set-integrable-l-mu: set-integrable lborel {f <.. f+t} ($\lambda s. \$l-(x+s) * \$\mu-(x+s)$)
if t ≥ 0 f ≥ 0 **for** t f :: real

proof –

have set-integrable lborel {f <.. f+t} ($\lambda s. \$l-(x+s) * \$\mu-(x+s) / \$l-x$)
using p-l set-integrable-p-mu that

by (rewrite set-integrable-cong[**where** $f' = \lambda s. \$p\{-s\&x\} * \$\mu\-(x+s)$]) simp-all+
thus ?thesis **by** simp
qed

lemma l-mu-has-integral-q-defer:

(($\lambda s. \$l\-(x+s) * \$\mu\-(x+s) / \$l\-x$) has-integral $\$q\-\{f|t\&x\}$) { $f..f+t$ }
if $t \geq 0$ $f \geq 0$ **for** $t f :: real$
using p-l that p-mu-has-integral-q-defer-Icc
by (rewrite has-integral-cong[of - - $\lambda s. \$p\{-s\&x\} * \$\mu\-(x+s)$]; simp)

corollary l-mu-has-integral-q:

(($\lambda s. \$l\-(x+s) * \$\mu\-(x+s) / \$l\-x$) has-integral $\$q\-\{t\&x\}$) { $0..t$ } **if** $t \geq 0$ **for** $t::real$
using l-mu-has-integral-q-defer[**where** $f=0$] that **by** simp

lemma l-mu-has-integral-d:

(($\lambda s. \$l\-(x+s) * \$\mu\-(x+s)$) has-integral $\$d\-\{t \& x+f\}$) { $f..f+t$ }
if $t \geq 0$ $f \geq 0$ **for** $t f :: real$

proof –

have (($\lambda s. \$l\-x * (\$p\{-s\&x\} * \$\mu\-(x+s))$) has-integral $\$l\-x * \$q\-\{f|t\&x\}$) { $f..f+t$ }
apply (rule has-integral-mult-right)
by (rule p-mu-has-integral-q-defer-Icc; simp add: that)
thus ?thesis
using that **apply** (rewrite in asm q-defer-d-l, simp-all)
apply (rewrite has-integral-cong[**where** $g = \lambda s. \$l\-x * (\$p\{-s\&x\} * \$\mu\-(x+s))$])
by (rewrite p-l; simp)

qed

corollary l-mu-has-integral-d-1:

(($\lambda s. \$l\-(x+s) * \$\mu\-(x+s)$) has-integral $\$d\-(x+f)$) { $f..f+1$ } **if** $t \geq 0$ $f \geq 0$ **for** t
 $f :: real$
using l-mu-has-integral-d[**where** $t=1$] that **by** simp

lemma e-LBINT-l: $\$e^{\circ}\-x = (LBINT s:\{0..\}. \$l\-(x+s) * \$\mu\-(x+s) * s) / \$l\-x$

— Note that $0 = 0$ holds when the life expectation diverges.

proof –

have $\bigwedge s. s \in \{0..\} \implies \$p\{-s\&x\} = \$l\-(x+s) / \$l\-x$ **using** p-l **by** simp
hence $\$e^{\circ}\-x = (LBINT s:\{0..\}. \$l\-(x+s) / \$l\-x * \$\mu\-(x+s) * s)$
using e-LBINT-p-mu
by (smt (verit) atLeast-borel set-lebesgue-integral-cong sets-lborel x-lt-psi)
also have ... = ($L\int s:\{0..\}. \$l\-(x+s) * \$\mu\-(x+s) * s) / \$l\-x$
using set-integral-divide-zero **by** simp
finally show ?thesis .

qed

lemma e-integral-l: $\$e^{\circ}\-x = \text{integral } \{0..\} (\lambda s. \$l\-(x+s) * \$\mu\-(x+s) * s) / \$l\-x$

— Note that $0 = 0$ holds when the life expectation diverges.

proof –

have AE s in lborel. $\$mu\-(x+s) \geq 0$ **by** (rule AE-translation, rule mu-nonneg-AE)
hence ($L\int s:\{0..\}. \$l\-(x+s) * \$\mu\-(x+s) * s$) = $\text{integral } \{0..\} (\lambda s. \$l\-(x+s)$

```

* $μ-(x+s) * s)
  by (intro set-borel-integral-eq-integral-nonneg-AE; force)
  thus ?thesis using e-LBINT-l by simp
qed

lemma m-LBINT-p-mu: $m-{n&x} = (LBINT t:{0<..n}. $p-{t&x} * $μ-(x+t))
/ (LBINT t:{0..n}. $p-{t&x})
  if n ≥ 0 for n::real
  using that
  apply (rewrite m-q-e, simp-all)
  apply (rewrite LBINT-p-mu-q[simplified], simp-all)
  by (rewrite temp-e-LBINT-p; simp)

lemma m-integral-p-mu:
  $m-{n&x} = integral {0..n} (λt. $p-{t&x} * $μ-(x+t)) / integral {0..n} (λt.
  $p-{t&x})
  if n ≥ 0 for n::real
  using that
  apply (rewrite m-q-e, simp-all)
  apply (rewrite integral-unique[OF p-mu-has-integral-q-Icc])
  apply simp-all[2]
  by (rewrite temp-e-integral-p; simp)

end

lemma deriv-x-p-mu-l: deriv (λy. $p-{t&y}) x = $p-{t&x} * ($μ-x - $μ-(x+t))
  if l differentiable at x l differentiable at (x+t) t ≥ 0 x < $ψ for x t :: real
  using deriv-x-p-mu that differentiable-ccdfX-l by blast

lemma e-has-derivative-mu-e-l: ((λx. $e^{'o}-x) has-real-derivative ($μ-x * $e^{'o}-x -
1)) (at x)
  if total-finite l differentiable at x x ∈ {a < .. < b} b ≤ $ψ for a b x :: real
  using total-finite-iff-set-integrable-Ici that
  e-has-derivative-mu-e differentiable-ccdfX-l set-integrable-ccdfX-l
  by force

corollary e-has-derivative-mu-e-l': ((λx. $e^{'o}-x) has-real-derivative ($μ-x * $e^{'o}-x
- 1)) (at x)
  if total-finite l differentiable at x x ∈ {a < .. < b} b ≤ $ψ for a b x :: real
  using that by (intro e-has-derivative-mu-e-l[where a=a]; simp)

context
  fixes x::real
  assumes x-lt-psi[simp]: x < $ψ
begin

lemma curt-e-sum-l-smooth: $e-x = (∑ k. $l-(x+k+1)) / $l-x if total-finite
proof -
  have [simp]: summable (λk. $l-(x+k+1))

```

```

using total-finite-iff-summable[of x+1] that
by (metis (no-types, lifting) add.commute add.left-commute summable-def sums-cong)
hence summable ( $\lambda k. \$p\{k+1 \& x\}$ ) by (rewrite p-l; simp add: add.assoc)
hence  $\$e\cdot x = (\sum k. \$p\{k+1 \& x\})$  using curt-e-sum-p-smooth by simp
also have ... =  $(\sum k. \$l\cdot(x+k+1)) / \$l\cdot x$  by (rewrite p-l; simp add: add.assoc)
also have ... =  $(\sum k. \$l\cdot(x+k+1)) / \$l\cdot x$  by (rewrite suminf-divide; simp)
finally show ?thesis .
qed

lemma curt-e-sum-l-finite-smooth:  $\$e\cdot x = (\sum k < n. \$l\cdot(x+k+1)) / \$l\cdot x$  if  $x+n+1 > \$\psi$  for  $n::nat$ 
apply (rewrite curt-e-sum-p-finite-smooth[of x n], simp-all add: that)
apply (rewrite p-l, simp-all)
by (smt (verit) sum.cong sum-divide-distrib)

lemma temp-curt-e-sum-l-smooth:  $\$e\{-x:n\} = (\sum k < n. \$l\cdot(x+k+1)) / \$l\cdot x$  for  $n::nat$ 
apply (rewrite temp-curt-e-sum-p-smooth[of x n], simp)
apply (rewrite p-l, simp-all)
by (smt (verit) sum.cong sum-divide-distrib)

end

end

```

6.4 Interpolations

```

context life-table
begin

definition linear-interpolation ≡
   $\forall (x::nat)(t::real). 0 \leq t \wedge t \leq 1 \longrightarrow \$l\cdot(x+t) = (1-t)*\$l\cdot x + t*\$l\cdot(x+1)$ 

lemma linear-l:  $\$l\cdot(x+t) = (1-t)*\$l\cdot x + t*\$l\cdot(x+1)$ 
  if linear-interpolation  $0 \leq t \wedge t \leq 1$  for  $x::nat$  and  $t::real$ 
  using that unfolding linear-interpolation-def by (metis of-nat-1 of-nat-add)

lemma linear-l-d:  $\$l\cdot(x+t) = \$l\cdot x - t*\$d\cdot x$ 
  if linear-interpolation  $0 \leq t \wedge t \leq 1$  for  $x::nat$  and  $t::real$ 
  using death-def-nonneg that unfolding linear-interpolation-def
  by (smt (verit) distrib-left left-diff-distrib')

lemma linear-p-q:  $\$p\{-t \& x\} = 1 - t*\$q\cdot x$ 
  if linear-interpolation  $0 \leq t \wedge t \leq 1$   $x < \$\psi$  for  $x::nat$  and  $t::real$ 
  using that
  apply (rewrite p-l, simp-all)
  apply (rewrite q-d-l, simp-all)
  using divide-self[of  $\$l\cdot(\text{real } x)$ ] linear-l-d
  by (smt (verit, ccfv-SIG) add-divide-distrib lx-neq-0)

```

```

lemma linear-q: $q-{t&x} = t*$q-x
  if linear-interpolation 0 ≤ t t ≤ 1 x < $ψ for x::nat and t::real
  using that linear-p-q p-q-1 by (smt (verit))

lemma linear-L-l-d: $L-x = $l-x - $d-x / 2 if linear-interpolation for x::nat
proof -
  have $L-(real x) = (LBINT t:{0..1}. $l-(real x + t))
  unfolding lives-def using lborel-set-integral-Icc-shift[of real x real x + 1 l real
x]
  by (simp add: add.commute)
  also have ... = (LBINT t:{0..1}. $l-(real x) - t*$d-(real x))
  using linear-l-d that by (intro set-lebesgue-integral-cong; simp)
  also have ... = $l-(real x) - $d-(real x) / 2
  proof -
    have (LBINT t:{0::real..1}. t) = 1/2
    unfolding set-lebesgue-integral-def using integral-power[of 0 1 1] by (simp
add: mult.commute)
    hence (LBINT t:{0..1}. t*$d-(real x)) = $d-(real x) / 2 by auto
    moreover have set-integrable lborel {0..1} (λt. t*$d-(real x))
    apply (rule set-integrable-mult-left[where f=id and a=$d-(real x), simplified])
    unfolding set-integrable-def using integrable-power[of 0 1 1] by (simp add:
mult.commute)
    moreover have (LBINT t:{0::real..1}. $l-(real x)) = $l-(real x)
    unfolding set-lebesgue-integral-def by simp
    ultimately show ?thesis using set-integrable-def by (rewrite set-integral-diff;
force)
  qed
  finally show ?thesis .
qed

lemma linear-L-l-d': $L-x = $l-(x+1) + $d-x / 2 if linear-interpolation for x::nat
proof -
  have $L-(real x) = $l-(real x) - $d-(real x) + $d-(real x) / 2 using that lin-
ear-L-l-d by simp
  also have ... = $l-(real (x+1)) + $d-(real x) / 2 using dx-l by (smt (verit)
of-nat-1 of-nat-add)
  finally show ?thesis .
qed

lemma linear-l-continuous: continuous-on UNIV l if linear-interpolation
  unfolding continuous-on-def
proof
  fix u::real
  show l -u→ $l-u
  proof (cases 'u ≤ 0')
    case True
    hence (l —→ $l-u) (at-left u) using l-left-continuous-nonpos continuous-within
    by auto

```

```

thus ?thesis
  apply (rule filterlim-split-at-real)
  using l-right-continuous continuous-within by auto
next
  case False
  hence u-pos:  $u > 0$  by simp
  thus ?thesis
  proof (cases ‹u = real-of-int ⌊u⌋›)
    case True
    from this u-pos obtain x::nat where ux:  $u = Suc x$ 
      by (metis gr0-implies-Suc of-int-0-less-iff of-int-of-nat-eq pos-int-cases)
    have  $((\lambda t. (1-t)*\$l(real x) + t*\$l(real x + 1)) \longrightarrow \$l(real x + 1))$  (at-left
1)
      apply (rewrite in (- \longrightarrow □) - add-0[THEN sym], rule tendsto-add)
      apply (simp add: LIM-zero-iff' tendsto-mult-left-zero)
      by (rewrite in (- \longrightarrow □) - mult-1[THEN sym], rule tendsto-mult-right)
    simp
    hence  $((\lambda t. \$l(real x + t)) \longrightarrow \$l(real x + 1))$  (at-left 1)
      apply (rewrite tendsto-cong; simp)
      apply (rule eventually-at-leftI[of 0]; simp)
      using that by (rewrite linear-l; simp add: add.commute)
    moreover have  $((\lambda t. \$l(real x + t)) \longrightarrow \$l(real x + 1))$  (at-right 1)
      using l-right-continuous apply (rule continuous-within-tendsto-compose,
simp-all)
      apply (rule eventually-at-right-less)
      by (rule tendsto-intros, simp-all)
    ultimately show ?thesis
      apply (rewrite ux)+
      apply (rewrite filterlim-shift-iff[where d=x, THEN sym])
      by (rule filterlim-split-at-real; simp add: comp-def add.commute)
next
  case False
  let ?x = nat ⌊u⌋
  let ?t =  $u - real ?x$ 
  let ?e = min ?t ( $1 - ?t$ )
  from False have ?t > 0 using u-pos by linarith
  moreover have ?t < 1 by linarith
  ultimately have e-pos: ?e > 0 by simp
  hence
     $\forall F v \text{ in nhds } u. \$l-v = (1 - (v - real ?x))*\$l(real ?x) + (v - real ?x)*\$l(real ?x + 1)$ 
    proof -
      { fix v::real assume vu-e: dist v u < ?e
        hence v - real ?x ≥ 0 using dist-real-def by force
        moreover have v - real ?x ≤ 1 using dist-real-def vu-e by force
        ultimately have \$l-v =  $(1 - (v - real ?x))*\$l(real ?x) + (v - real ?x)*\$l(real ?x + 1)$ 
          using linear-l that by (smt (verit, ccfv-threshold) linear-interpolation-def)
      }

```

```

thus ?thesis using eventually-nhds-metric e-pos by blast
qed
moreover have
  isCont (?v. (1 - (v - real ?x))*$l-(real ?x) + (v - real ?x)*$l-(real ?x +
1)) u
  by (rule continuous-intros)+
ultimately have isCont l u using isCont-cong by force
thus ?thesis by (simp add: isContD)
qed
qed
qed

lemma linear-l-sums-T-l: (?k. $l-(x + Suc k)) sums ($T-x - $l-x / 2)
  if linear-interpolation total-finite for x::nat
proof -
  have ?k::nat. $l-(real (x + Suc k)) = $L-(real (x+k)) - $d-(real (x+k)) / 2
    using linear-L-l-d' that by (smt (verit) Suc-eq-plus1 add-Suc-right)
  moreover have (?k::nat. $L-(real (x+k))) sums $T-x using L-sums-T that by
  simp
  moreover have (?k::nat. $d-(real (x+k)) / 2) sums ($l-(real x) / 2)
    using sums-divide d-sums-l by auto
  ultimately show ?thesis
    apply (rewrite sums-cong, simp)
    by (rule sums-diff; simp)
qed

corollary linear-T-suminf-l: $T-x = (?k. $l-(x+k+1)) + $l-x / 2
  if linear-interpolation total-finite for x::nat
  using linear-l-sums-T-l that sums-unique
  by (metis (no-types, lifting) diff-diff-eq2 nat-arith.add1 right-minus-eq semiring-norm(174) suminf-cong)

lemma linear-mx-q: $m-x = $q-x / (1 - $q-x / 2) if linear-interpolation x < $ψ
for x::nat
proof -
  have [simp]: $l-(real x) ≠ 0 using that by simp
  have $m-(real x) = $d-(real x) / ($l-(real x) - $d-(real x) / 2)
    unfolding die-central-def using linear-L-l-d that by simp
  also have ... = ($d-(real x) / $l-(real x)) / ((($l-(real x) - $d-(real x) / 2) /
  $l-(real x)))
    by simp
  also have ... = ($d-(real x) / $l-(real x)) / (1 - ($d-(real x) / $l-(real x)) / 2)
    by (rewrite diff-divide-distrib) simp
  also have ... = $q-(real x) / (1 - $q-(real x) / 2) using that q-d-l by simp
  finally show ?thesis .
qed

lemma linear-e-curt-e: $e^o-x = $e-x + 1/2
  if linear-interpolation total-finite x < $ψ for x::nat

```

```

proof -
  have $e'◦-(real x) = (( $\sum k::nat. \$l-(real (x+k+1))$ ) +  $\$l-(real x) / 2$ ) /  $\$l-(real x)$ 
  using e-T-l linear-T-suminf-l that by simp
  also have ... = ( $\sum k::nat. \$l-(real (x+k+1))$ ) /  $\$l-(real x) + (\$l-(real x) / 2)$ 
  /  $\$l-(real x)$ 
  using add-divide-distrib by blast
  also have ... = $e-(real x) + 1/2
  using that apply (rewrite curt-e-sum-l, simp-all)
  using linear-l-continuous by (rule continuous-on-interior, simp-all add: that
add.commute)
  finally show ?thesis .
qed

end

context smooth-life-table
begin

lemma linear-l-has-derivative-at-frac:
  (( $\lambda s. \$l-(x+s)$ ) has-real-derivative – $d-x) (at t)
  if linear-interpolation  $0 < t < 1$  for x::nat and t::real
proof –
  let ?x = real x
  have (( $\lambda s. \$l-?x - s*\$d-?x$ ) has-real-derivative (0 – $d-?x)) (at t)
  apply (rule derivative-intros, simp)
  apply (rule DERIV-cmult-right[of id 1, simplified])
  by (metis DERIV-ident eq-id-iff)
  moreover have  $\forall F s$  in nhds t.  $\$l-(?x + s) = \$l-?x - s*\$d-?x$ 
proof –
  let ?r = min t (1-t)
  have ?r > 0 using that by simp
  moreover
  { fix s assume dist s t < ?r
    hence  $\$l-(?x + s) = \$l-?x - s*\$d-?x$  using linear-l-d that dist-real-def by
    force }
  ultimately show ?thesis using eventually-nhds-metric that by blast
  qed
  ultimately show ?thesis by (rewrite DERIV-cong-ev; simp)
qed

lemma linear-l-has-derivative-at-frac':
  ( $l$  has-real-derivative – $d-x) (at y)
  if linear-interpolation  $x < y$   $y < x+1$  for x::nat and y::real
  apply (rewrite DERIV-at-within-shift[where x=y – real x and z=real x and
S=UNIV, simplified])
  using linear-l-has-derivative-at-frac that by simp

lemma linear-l-differentiable-on-frac:

```

```

l differentiable-on {x<..<x+1} if linear-interpolation for x::nat
proof -
{ fix y::real assume y ∈ {real x <..< real (x+1)}
  hence l differentiable at y
  using linear-l-has-derivative-at-frac' that real-differentiable-def by auto }
  thus ?thesis unfolding differentiable-on-def by (metis differentiable-at-withinI)
qed

lemma linear-l-has-right-derivative-at-nat:
(l has-real-derivative - $d-x) (at-right x) if linear-interpolation for x::nat
proof -
let ?x = real x
have [simp]: plus ?x ‘ {0<..} = {?x<..}
  unfolding image-def greaterThan-def apply simp
  by (metis Groups.ab-semigroup-add-class.add.commute
       add-minus-cancel neg-less-iff-less real-0-less-add-iff)
have ((λs. $l-(?x + s)) has-real-derivative - $d-?x) (at-right 0)
  apply (rewrite has-field-derivative-cong-eventually[where g=λs. $l-?x - s*$d-?x])
  using linear-l-d that apply (intro eventually-at-rightI[of 0 1], simp-all)
  apply (rule has-field-derivative-at-within)
  apply (rewrite diff-0[of $d-?x, THEN sym])
  apply (rule DERIV-diff, simp)
  apply (rule DERIV-cmult-right[of id 1, simplified])
  by (metis DERIV-ident eq-id-iff)
thus ?thesis
  by (rewrite DERIV-at-within-shift[where z=?x and x=0 and S={0<..},
simplified]) simp
qed

lemma linear-l-has-left-derivative-at-nat:
(l has-real-derivative - $d-(real x - 1)) (at-left x) if linear-interpolation for
x::nat
proof (cases x)
  case 0
  hence (l has-real-derivative 0) (at-left (real x))
    apply (rewrite has-field-derivative-cong-eventually[where g=λ-. $l-0]; simp)
    apply (rule eventually-at-leftI[of -1]; simp)
    using l-neg-nil less-eq-real-def by blast
  moreover have $d-(real x - 1) = 0 using 0 dx-l l-neg-nil less-eq-real-def by
fastforce
  ultimately show ?thesis by auto
next
  case (Suc y)
  let ?x = real x and ?y = real y
  have [simp]: ?y + 1 = ?x using Suc by simp
  have [simp]: plus ?y ‘ {..<1} = {..<?x}
    using Suc unfolding image-def lessThan-def apply simp
    by (metis (no-types, opaque-lifting) Groups.ab-semigroup-add-class.add.commute
         add-less-cancel-right diff-add-cancel)

```

```

have $l-?x = $l-real y - $d-real y using Suc by (simp add: dx-l)
moreover have ∀ F s in at-left 1. $l-(?y + s) = $l-?y - s*$d-?y
  apply (rule eventually-at-leftI[of 0], simp-all)
  using Suc linear-l-d that by simp
moreover have ((λs. $l-?y - s*$d-?y) has-real-derivative = $d-?y) (at-left 1)
  apply (rule has-field-derivative-at-within)
  apply (rewrite diff-0[of $d-?y, THEN sym])
  apply (rule DERIV-diff, simp)
  apply (rule DERIV-cmult-right[of id 1, simplified])
  by (metis DERIV-ident eq-id-iff)
ultimately have ((λs. $l-(?y + s)) has-real-derivative = $d-?y) (at-left 1)
  by (rewrite has-field-derivative-cong-eventually[where g=λs. $l-?y - s*$d-?y];
      simp)
thus ?thesis
  apply (rewrite DERIV-at-within-shift[where S={..<1} and z=?y and x=1,
  simplified])
  using Suc by simp
qed

lemma linear-l-has-derivative-at-nat-iff-d:
  (l has-real-derivative = $d-x) (at x) ↔ $d-x = $d-(real x - 1)
  if linear-interpolation for x::nat
proof -
  let ?x = real x
  have (l has-real-derivative = $d-?x) (at ?x) ↔
    (l has-real-derivative = $d-?x) (at-right ?x) ∧
    (l has-real-derivative = $d-?x) (at-left ?x)
  using has-real-derivative-at-split by auto
also have ... ↔ $d-?x = $d-(?x - 1) (is ?LHS = ?RHS)
proof
  assume ?LHS
  hence (l has-real-derivative = $d-?x) (at-left ?x) by simp
  moreover have (l has-real-derivative = $d-(?x - 1)) (at-left ?x)
    using that linear-l-has-left-derivative-at-nat by simp
  ultimately show ?RHS
  using has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
    trivial-limit-at-left-real by (smt (verit, ccfv-SIG))
next
  assume ?RHS
  thus ?LHS
    using that linear-l-has-right-derivative-at-nat linear-l-has-left-derivative-at-nat
  by metis
  qed
  finally show ?thesis .
qed

lemma linear-l-differentiable-at-nat-iff-d:
  l differentiable at x ↔ $d-x = $d-(real x - 1)
  if linear-interpolation for x::nat

```

```

proof
  let ?x = real x
  assume l differentiable at x
  from this obtain D where DERIV-l: (l has-real-derivative D) (at ?x)
    using real-differentiable-def by blast
  hence (l has-real-derivative D) (at-right ?x)
    using has-field-derivative-at-within by blast
  moreover have at ?x within {real x<..} ≠ ⊥ by simp
  moreover have (l has-real-derivative – $d-?x) (at-right ?x)
    using linear-l-has-right-derivative-at-nat that by simp
  ultimately have D = – $d-?x
  using that has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
  by blast
  thus $d-?x = $d-(?x – 1)
  using linear-l-has-derivative-at-nat-iff-d that DERIV-l by blast
next
  assume $d-(real x) = $d-(real x – 1)
  thus l differentiable at (real x)
    using linear-l-has-derivative-at-nat-iff-d that real-differentiable-def by blast
qed

lemma linear-l-limited: $ψ < ∞ if linear-interpolation
proof –
  let ?ND = {y. ¬ l differentiable at y}
  obtain xn::nat where xn-def: Max ?ND < real xn using reals-Archimedean2
  by blast
  hence xn-dif: ∀x::nat. x ≥ xn ⇒ l differentiable at (real x)
  proof –
    fix x::nat assume x ≥ xn
    with xn-def have real x > Max ?ND by simp
    hence real x ∉ ?ND using notI Max.coboundedI l-nondifferentiable-finite-set
    leD by blast
    thus l differentiable at (real x) by simp
  qed
  hence d-const: ∀x::nat. x ≥ xn ⇒ $d-(real x) = $d-(real xn)
  proof –
    fix x::nat assume x ≥ xn
    moreover have
      ∀y::nat. xn ≤ y ⇒ $d-(real y) = $d-(real xn) ⇒ $d-(real (Suc y)) = $d-(real
      xn)
      using linear-l-differentiable-at-nat-iff-d that
      by (smt (verit, best) of-nat-Suc of-nat-le-iff xn-dif)
    ultimately show $d-(real x) = $d-(real xn)
      using nat-induct-at-least[where P=λx::nat. $d-(real x) = $d-(real xn)] by
      simp
  qed
  have ¬ $d-(real xn) > 0
  proof (rule notI)
    assume $d-(real xn) > 0

```

```

from this obtain N::nat where N-def:  $N * \$d-(real xn) > \$l-(real xn)$ 
  using reals-Archimedean3 by blast
hence  $\$l-(real (xn+N)) < 0$ 
proof -
  have  $\$l-(real (xn+N)) = \$l-(real xn) - (\sum k < N. \$d-(real xn + real k))$  using
  sum-dx-l by simp
  also have ... =  $\$l-(real xn) - (\sum k < N. \$d-(real xn))$ 
    using d-const by (metis le-add1 of-nat-add)
  also have ... =  $\$l-(real xn) - N * \$d-(real xn)$  by simp
  also have ... < 0 using N-def by simp
  finally show ?thesis .
qed
thus False by (smt (verit, ccfv-SIG) l-nonneg)
qed
hence dxn0:  $\$d-(real xn) = 0$  by (smt (verit) d-nonneg)
hence  $\bigwedge x::nat. x \geq xn \implies \$l-(real x) = \$l-(real xn)$ 
proof -
  fix x::nat
  assume xn ≤ x
  moreover have
     $\bigwedge y::nat. xn \leq y \implies \$l-(real y) = \$l-(real xn) \implies \$l-(real (Suc y)) = \$l-(real xn)$ 
    by (smt (verit, ccfv-threshold) dxn0 d-const dx-l of-nat-Suc)
  ultimately show  $\$l-(real x) = \$l-(real xn)$ 
    using nat-induct-at-least[where P=λx::nat.  $\$l-(real x) = \$l-(real xn)$ ] by
  simp
qed
hence  $(\lambda x::nat. \$l-(real x)) \longrightarrow \$l-(real xn)$ 
  using eventually-sequentially by (intro tends-to-eventually) blast
moreover have  $(\lambda x::nat. \$l-(real x)) \longrightarrow 0$ 
  using l-PInfty-0 by (simp add: filterlim-compose filterlim-real-sequentially)
ultimately have  $\$l-(real xn) = 0$  by (simp add: LIMSEQ-unique)
thus ?thesis by force
qed

lemma linear-mu-q:  $\$μ-(x+t) = \$q-x / (1 - t * \$q-x)$ 
  if linear-interpolation l differentiable at (x+t)  $0 < t < 1$   $x+t < \$ψ$ 
  for x::nat and t::real
proof -
  have [simp]: ereal x < $ψ using that by (simp add: ereal-less-le)
  have [simp]:  $\$l-(real x) \neq 0$  by simp
  have [simp]: l field-differentiable at (real x + t)
    using differentiable-eq-field-differentiable-real that by simp
  define d where d ≡ min t (1-t)
  have d-pos: d > 0 unfolding d-def using that by simp
  have  $\$p\{t \& real x\} \neq 0$  using that by (simp add: ereal-less-le p-0-equiv)
  moreover have  $(\lambda s. \$p\{s \& real x\})$  differentiable at t
  proof -
    have  $(\lambda s. \$l-(real x + s) / \$l-(real x))$  field-differentiable at t

```

```

using that apply (intro derivative-intros)
  apply (rewrite add.commute, rewrite field-differentiable-shift[THEN sym])
  by (rewrite add.commute) simp-all
thus ?thesis
  apply (rewrite differentiable-eq-field-differentiable-real)
  apply (rule field-differentiable-transform-within[where d=d], simp-all add:
d-pos)
    apply (rewrite p-l, simp-all) unfolding d-def using dist-real-def by auto
qed
ultimately have $μ-(real x + t) = - deriv (λs. $p-{s & real x}) t / $p-{t &
real x}
  using that deriv-t-p-mu by simp
also have ... = $q-(real x) / (1 - t*$q-(real x))
proof -
  have ∀s. dist s t < d ==> $p-{s & real x} = 1 - s*$q-(real x)
  proof -
    fix s assume dist s t < d
    hence 0 ≤ s s ≤ 1 unfolding d-def using that dist-real-def by auto
    thus $p-{s & real x} = 1 - s*$q-(real x) by (intro linear-p-q; simp add: that)
  qed
  hence deriv (λs. $p-{s & real x}) t = deriv (λs. 1 - s*$q-(real x)) t
    using d-pos apply (intro deriv-cong-ev, simp-all)
    by (rewrite eventually-nhds-metric) auto
  also have ... = - $q-(real x)
    apply (rewrite deriv-diff, simp-all)
    by (rule derivative-intros) auto
  finally have deriv (λs. $p-{s & real x}) t = - $q-(real x) .
  thus ?thesis using linear-p-q that by simp
qed
finally show ?thesis .
qed

```

definition exponential-interpolation ≡
 $\forall(x::nat)(t::real). x+1 < \psi \longrightarrow 0 \leq t \wedge t < 1 \longrightarrow \mu-(x+t) = \mu-x$
— Without $x+1 < \psi$, the smooth life table could not be limited.

lemma exponential-mu: $\mu-(x+t) = \mu-x$
if exponential-interpolation $x+1 < \psi \ 0 \leq t \ t < 1$ for $x::nat$ and $t::real$
using that unfolding exponential-interpolation-def by simp

corollary exponential-mu': $\mu-y = \mu-x$
if exponential-interpolation $x \leq y \ y < x+1 \ x+1 < \psi$ for $x::nat$ and $y::real$
proof -
let ?t = y - real x
have 0 ≤ ?t and ?t < 1 using that by simp-all
moreover have $\mu-y = \mu-(real x + ?t)$ by simp
ultimately show ?thesis using exponential-mu that by presburger
qed

lemma *exponential-integral-mu*: $\text{integral} \{x..<x+t\} (\lambda y. \$\mu-y) = \$\mu-x * t$
if exponential-interpolation $x+1 < \$\psi$ $0 \leq t \leq 1$ **for** $x::nat$ **and** $t::real$
proof –
have $\text{integral} \{real x ..< real x + t\} (\lambda y. \$\mu-y) =$
 $\text{integral} \{real x ..< real x + t\} (\lambda y. \$\mu-(real x))$
using *exponential-mu'* **that by** (*intro integral-cong; simp*)
also have $\dots = \text{integral} \{real x .. real x + t\} (\lambda y. \$\mu-(real x))$
apply (*rule integral-subset-negligible, force*)
using **that by** (*rewrite Icc-minus-Ico; simp*)
also have $\dots = \$\mu-x * t$ **using** **that by** (*rewrite integral-const-real*) *simp*
finally show $\text{integral} \{real x ..< real x + t\} (\lambda y. \$\mu-y) = \$\mu-(real x) * t$.
qed

lemma *exponential-p-mu*: $\$p-x = \exp(-\$\mu-x)$ **if exponential-interpolation** $x+1 < \$\psi$ **for** $x::nat$
proof –
have $\$p-x = \exp(-\text{integral} \{real x ..< real x + 1\} (\lambda y. \$\mu-y))$
using **that apply** (*rewrite p-exp-integral-mu; simp add: add.commute*)
apply (*rule integral-subset-negligible[THEN sym], force*)
by (*rewrite Icc-minus-Ico; simp*)
also have $\dots = \exp(-\$\mu-(real x))$ **using** **that by** (*rewrite exponential-integral-mu; simp*)
finally show ?thesis.
qed

corollary *exponential-mu-p*: $\$μ-x = -\ln(\$p-x)$ **if exponential-interpolation** $x+1 < \$ψ$ **for** $x::nat$
using *exponential-p-mu* **that by** *simp*

corollary *exponential-mu-xt-p*: $\$μ-(x+t) = -\ln(\$p-x)$
if exponential-interpolation $x+1 < \$ψ$ $0 \leq t \leq 1$ **for** $x::nat$ **and** $t::real$
using **that exponential-mu exponential-mu-p by presburger**

corollary *exponential-q-mu*: $\$q-x = 1 - \exp(-\$μ-x)$
if exponential-interpolation $x+1 < \$ψ$ **for** $x::nat$
using *exponential-p-mu* **that p-q-1**
by (*smt (verit, ccfv-SIG) ereal-less-le not-add-less1 of-nat-less-imp-less*)

lemma *exponential-p*: $\$p-\{t&x\} = (\$p-x).^\wedge t$
if exponential-interpolation $x+1 < \$ψ$ $0 \leq t \leq 1$ **for** $x::nat$ **and** $t::real$
proof –
have [*simp*]: $real x + t < \$ψ$ **using** **that ereal-less-le by auto**
have $\$p-\{t & real x\} = \exp(-\text{integral} \{real x ..< real x + t\} (\lambda y. \$\mu-y))$
using **that apply** (*rewrite p-exp-integral-mu, simp-all*)
apply (*rule integral-subset-negligible[THEN sym], force*)
by (*rewrite Icc-minus-Ico; simp*)
also have $\dots = \exp(-\$μ-(real x) * t)$
using **that by** (*rewrite exponential-integral-mu; simp*)
also have $\dots = (\$p-(real x)).^\wedge t$

using exponential-p-mu that

by (smt (verit) exp-not-eq-zero exponential-mu-p mult.commute powr-def)

finally show ?thesis .

qed

lemma exponential-q: $\$q\{t \& x\} = 1 - (1 - \$q\cdot x)\cdot \hat{t}$

if exponential-interpolation $x+1 < \$\psi$ $0 \leq t \leq 1$ for $x::nat$ and $t::real$

proof –

have $\$q\{t \& real x\} = 1 - \$p\{t \& real x\}$

using p-q-1 that by (smt (verit) ereal-less-le le-add1 of-nat-mono)

also have $\dots = 1 - (\$p\cdot(real x))\cdot \hat{t}$ using that by (rewrite exponential-p; simp)

also have $\dots = 1 - (1 - \$q\cdot(real x))\cdot \hat{t}$

using p-q-1 that by (smt (verit) ereal-less-le not-add-less1 of-nat-less-imp-less)

finally show ?thesis .

qed

lemma exponential-l-p: $\$l\cdot(x+t) = \$l\cdot x * (\$p\cdot x)\cdot \hat{t}$

if exponential-interpolation $x+1 < \$\psi$ $0 \leq t \leq 1$ for $x::nat$ and $t::real$

proof –

have ereal (real x) < $\$ψ$ using that ereal-less-le by auto

hence $\$l\cdot(real x + t) = \$l\cdot x * \$p\{t \& real x\}$ using that by (rewrite p-l; simp)

also have $\dots = \$l\cdot(real x) * (\$p\cdot(real x))\cdot \hat{t}$ using that by (rewrite exponential-p; simp)

finally show ?thesis .

qed

lemma exponential-l-has-derivative-at-fraction:

((λs. $\$l\cdot(x+s)$) has-real-derivative ($-\ $l\cdot x * \$μ\cdot x * (\$p\cdot x)\cdot \hat{t}$)) (at t)

if exponential-interpolation $x+1 < \$\psi$ $0 < t \leq 1$ for $x::nat$ and $t::real$

proof –

let ?x = real x

have ((λs. $(\$p\cdot ?x)\cdot \hat{s}$) has-real-derivative ($-\ $μ\cdot ?x * (\$p\cdot ?x)\cdot \hat{t}$)) (at t)

using that exponential-p-mu has-real-derivative-powr2

by (metis Groups.ab-semigroup-mult-class.mult.commute exp-gt-zero ln-exp)

hence ((λs. $\$l\cdot ?x * (\$p\cdot ?x)\cdot \hat{s}$) has-real-derivative ($-\ $l\cdot ?x * \$μ\cdot ?x * (\$p\cdot ?x)\cdot \hat{t}$))

(at t)

by (rewrite minus-mult-commute, rewrite mult.assoc) (rule DERIV-cmult)

moreover have $\forall F s$ in nhds t. $\$l\cdot(?x + s) = \$l\cdot ?x * (\$p\cdot ?x)\cdot \hat{s}$

proof –

let ?r = min t (1-t)

have ?r > 0 using that by simp

moreover have $\bigwedge s. dist s t < ?r \implies \$l\cdot(?x + s) = \$l\cdot ?x * (\$p\cdot ?x)\cdot \hat{s}$

using dist-real-def that exponential-l-p by force

ultimately show ?thesis using eventually-nhds-metric by blast

qed

ultimately show ?thesis by (rewrite DERIV-cong-ev[where g=λs. $\$l\cdot ?x * (\$p\cdot ?x)\cdot \hat{s}$]; simp)

qed

```

lemma exponential-l-has-derivative-at-frac':
  (l has-real-derivative ( $-\$l \cdot x * \$\mu \cdot x * (\$p \cdot x) \cdot \widehat{(y-x)}$ ) (at y)
   if exponential-interpolation  $x+1 < \$\psi$   $x < y$   $y < x+1$  for  $x::nat$  and  $y::real$ 
   apply (rewrite DERIV-at-within-shift[where  $x=y - real\ x$  and  $z=real\ x$  and
    $S=UNIV$ , simplified])
   using exponential-l-has-derivative-at-frac that by simp

lemma exponential-l-differentiable-on-frac:
  l differentiable-on { $x <.. < x+1$ } if exponential-interpolation  $x+1 < \$\psi$  for  $x::nat$ 
proof -
  { fix  $y::real$  assume  $y \in \{real\ x <.. < real\ (x+1)\}$ 
    hence l differentiable at y
    using exponential-l-has-derivative-at-frac' that real-differentiable-def by auto
  }
  thus ?thesis unfolding differentiable-on-def by (metis differentiable-at-withinI)
qed

lemma exponential-l-has-right-derivative-at-nat:
  (l has-real-derivative ( $-\$l \cdot x * \$\mu \cdot x$ ) (at-right x)
   if exponential-interpolation  $x+1 < \$\psi$  for  $x::nat$ 
proof -
  let ?x = real x
  have [simp]: plus ?x ‘ {0<..} = {?x<..}
  unfolding image-def greaterThan-def apply simp
  by (metis Groups.ab-semigroup-add-class.add.commute
        add-minus-cancel neg-less-iff-less real-0-less-add-iff)
  have [simp]:  $\$p \cdot x > 0$  using exponential-p-mu that by auto
  hence [simp]:  $\$p \cdot x \neq 0$  by force
  have (( $\lambda s. \$l \cdot (?x + s)$ ) has-real-derivative ( $-\$l \cdot ?x * \$\mu \cdot ?x$ ) (at-right 0)
  apply (rewrite has-field-derivative-cong-eventually[where g= $\lambda s. \$l \cdot ?x * (\$p \cdot ?x) \cdot \widehat{s}$ ])
  using exponential-l-p that apply (intro eventually-at-rightI[of 0 1]; simp)
  using powr-zero-eq-one apply simp
  apply (rewrite minus-mult-commute, rule DERIV-cmult)
  apply (rule has-field-derivative-at-within)
  using that apply (rewrite exponential-mu-p; simp)
  using has-real-derivative-powr2[of  $\$p \cdot x 0$ ] powr-zero-eq-one by force
  thus ?thesis
  by (rewrite DERIV-at-within-shift[where z=?x and x=0 and S={0<..},
        simplified]) simp
qed

lemma exponential-l-has-left-derivative-at-nat:
  (l has-real-derivative ( $-\$l \cdot x * \$\mu \cdot (real\ x - 1)$ ) (at-left x)
   if exponential-interpolation  $x < \$\psi$  for  $x::nat$ 
proof (cases x)
  case 0
  hence (l has-real-derivative 0) (at-left (real x))
  apply (rewrite has-field-derivative-cong-eventually[where g= $\lambda -. \$l \cdot 0$ ; simp)
  apply (rule eventually-at-leftI[of -1]; simp)

```

```

using l-neg-nil less-eq-real-def by blast
moreover have - $l-x * $μ-(real x - 1) = 0 using mu-unborn-0 0 by simp
ultimately show ?thesis by auto
next
let ?x = real x
case (Suc y)
let ?y = real y
have [simp]: ?y + 1 = ?x using Suc by simp
hence [simp]: ereal ?y < $ψ using that by (smt (verit) ereal-less-le)
have [simp]: $p-?y > 0 using Suc exponential-p-mu that by auto
have [simp]: plus ?y ` {..<?x} = {..<?x}
    using Suc unfolding image-def lessThan-def apply simp
by (metis (no-types, opaque-lifting) Groups.ab-semigroup-add-class.add.commute
    add-less-cancel-right diff-add-cancel)
have ((λt. ($p-?y).^t) has-real-derivative ($p-?y * -$μ-?y)) (at-left 1)
    apply (rule DERIV-subset[where s=UNIV]; simp)
    using that apply (rewrite exponential-mu-p; simp add: add.commute[of 1 ?y])
    by (rule has-real-derivative-powr2[of $p-?y 1, simplified])
    hence ((λt. $l-?y * ($p-?y).^t) has-real-derivative ($l-?y * $p-?y * -$μ-?y))
(at-left 1)
    by (metis DERIV-cmult mult-ac(1))
moreover have $l-?y * $p-?y = $l-?x using Suc p-1-l by simp
ultimately have ((λt. $l-?y * ($p-?y).^t) has-real-derivative (- $l-?x * $μ-?y))
(at-left 1)
    by simp
moreover have ∀ F t in at-left 1. $l-?y + t = $l-?y * ($p-?y).^t
    apply (rule eventually-at-leftI[of 0]; simp)
    by (rewrite exponential-l-p; simp add: that add.commute[of 1 ?y])
    ultimately have ((λt. $l-?y + t) has-real-derivative (- $l-?x * $μ-?y)) (at-left 1)
    by (rewrite has-field-derivative-cong-eventually[where g=λt. $l-?y * ($p-?y).^t];
        simp add: p-1-l)
thus ?thesis
    apply (rewrite DERIV-at-within-shift[where S={..<1} and z=?y and x=1,
simplified])
    using Suc by simp
qed

```

```

lemma exponential-l-has-derivative-at-nat-iff-mu:
(l has-real-derivative (- $l-x * $μ-x)) (at x) ↔ $μ-x = $μ-(real x - 1)
if exponential-interpolation x+1 < $ψ for x::nat
proof -
let ?x = real x
have [simp]: ?x < $ψ using that by (simp add: ereal-less-le)
hence [simp]: $l-?x ≠ 0 by simp
have (l has-real-derivative (- $l-?x * $μ-?x)) (at ?x) ↔
(l has-real-derivative (- $l-?x * $μ-?x)) (at-right ?x) ∧
(l has-real-derivative (- $l-?x * $μ-?x)) (at-left ?x)
using has-real-derivative-at-split by auto

```

```

also have ...  $\longleftrightarrow - \$l\cdot?x * \$\mu\cdot?x = - \$l\cdot?x * \$\mu\cdot(?x - 1)$  (is ?LHS  $\longleftrightarrow$  ?RHS)
proof
assume ?LHS
hence (l has-real-derivative ( $- \$l\cdot?x * \$\mu\cdot?x$ ) (at-left ?x) by simp
moreover have (l has-real-derivative ( $- \$l\cdot?x * \$\mu\cdot(?x - 1)$ ) (at-left ?x)
    using that exponential-l-has-left-derivative-at-nat by force
ultimately show ?RHS
using has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
    trivial-limit-at-left-real by blast
next
assume ?RHS
thus ?LHS
using that exponential-l-has-right-derivative-at-nat exponential-l-has-left-derivative-at-nat
    by force
qed
also have ...  $\longleftrightarrow \$\mu\cdot?x = \$\mu\cdot(?x - 1)$  by simp
finally show ?thesis .
qed

lemma exponential-l-differentiable-at-nat-iff-mu:
l differentiable at x  $\longleftrightarrow \$\mu\cdot x = \$\mu\cdot(\text{real } x - 1)$ 
if exponential-interpolation  $x+1 < \$\psi$  for  $x::nat$ 
proof
let ?x = real x
assume l differentiable at ?x
from this obtain D where DERIV-l: (l has-real-derivative D) (at ?x)
    using real-differentiable-def by blast
hence (l has-real-derivative D) (at-right ?x)
    using has-field-derivative-at-within by blast
moreover have at ?x within {real x < ..}  $\neq \perp$  by simp
moreover have (l has-real-derivative ( $- \$l\cdot\text{real } x * \$\mu\cdot\text{real } x$ ) (at-right ?x)
    using exponential-l-has-right-derivative-at-nat that by simp
ultimately have D =  $- \$l\cdot?x * \$\mu\cdot?x$ 
    using that has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
    by blast
thus  $\$μ\cdot?x = \$μ\cdot(?x - 1)$ 
    using exponential-l-has-derivative-at-nat-iff-mu that DERIV-l by blast
next
assume  $\$μ\cdot(\text{real } x) = \$μ\cdot(\text{real } x - 1)$ 
thus l differentiable at (real x)
    using exponential-l-has-derivative-at-nat-iff-mu that real-differentiable-def by
blast
qed

lemma exponential-L-d-mu:  $\$L\cdot x = \$d\cdot x / \$μ\cdot x$ 
if exponential-interpolation  $\$μ\cdot x \neq 0$   $x+1 < \$ψ$  for  $x::nat$ 
proof -
have [simp]: ereal (real x)  $< \$ψ$  using that ereal-less-le by simp

```

```

have [simp]: $l-(real x) ≠ 0 by simp
have p-pos[simp]: $p-(real x) > 0 using that by (simp add: exponential-p-mu)
have [simp]: $p-(real x) ≠ 1 using that exponential-p-mu by simp
have $L-(real x) = (LBINT t:{0..1}. $l-(real x + t))
  unfolding lives-def by (rewrite lborel-set-integral-Icc-shift[where t=x]) simp
also have ... = integral {0..1} (λt. $l-(real x + t))
  by (rule set-borel-integral-eq-integral-nonneg; simp)
also have ... = integral {0..1} (λt. $l-(real x) * ($p-(real x)).^t)
  apply (rule integral-cong)
  using that by (rewrite exponential-l-p; simp)
also have ... = $l-(real x) * integral {0..1} (λt. ($p-(real x)).^t)
  using integral-mult-right by blast
also have ... = $l-(real x) * ($q-(real x) / - ln ($p-(real x)))
proof -
  have integral {0..1} (λt. ($p-(real x)).^t) = (($p-(real x)).^1 - 1) / ln ($p-(real x))
    apply (rule integral-unique)
    by (intro has-integral-powr2-from-0; simp)
  also have ... = $q-(real x) / - ln ($p-(real x))
    using p-pos apply (rewrite powr-one, linarith)
    using that by (rewrite in -- □ p-q-1[of x 1, THEN sym]; simp)
  finally show ?thesis by simp
qed
also have ... = $d-(real x) / $μ-(real x) using that exponential-mu-p by (rewrite q-d-l; simp)
finally show ?thesis .
qed

lemma exponential-mx-mu: $m-x = $μ-x if exponential-interpolation x+1 < $ψ
for x::nat
proof (cases '$μ-(real x) = 0')
  have lx-neq0: $l-(real x) ≠ 0 using ereal-less-le that by simp
  case True
  hence $q-(real x) = 0 using exponential-q-mu that by simp
  hence $d-(real x) = 0
    using q-1-d-l that by (metis lx-neq0 d-old-0 divide-eq-0-iff linorder-not-less zero-le-one)
  hence $m-(real x) = 0 unfolding die-central-def by simp
  also have ... = $μ-(real x) using True by simp
  finally show ?thesis .
next
  case False
  thus ?thesis unfolding die-central-def using exponential-L-d-mu that
    by (smt (verit) divide-eq-0-iff divide-mult-cancel exp-eq-one-iff exponential-q-mu linorder-not-less mu-beyond-0 nonzero-mult-div-cancel-left q-1-d-l)
qed

lemma exponential-d-mu-sums-T: (λk. $d-(x+k) / $μ-(x+k)) sums $T-x
  if exponential-interpolation total-finite ∧ k::nat. $μ-(x+k) ≠ 0 for x::nat

```

```

proof -
  have  $\neg \psi < \infty$ 
  proof
    assume  $\psi < \infty$ 
    from this obtain  $y::nat$  where  $\psi < ereal (real y)$  using less-PInf-Ex-of-nat
    by fastforce
    hence  $xy: \psi < ereal (real (x+y))$  by (simp add: less-ereal-le)
    hence  $\mu-(real (x+y)) \neq 0$  using that by simp
    moreover have  $\mu-(real (x+y)) = 0$  using xy mu-beyond-0 by simp
    ultimately show False ..
  qed
  hence  $\psi = \infty$  by simp
  moreover hence  $\bigwedge k::nat. d-(real (x+k)) / \mu-(real (x+k)) = L-(real (x+k))$ 
    using that by (rewrite exponential-L-d-mu; simp)
  ultimately show ?thesis
    apply (rewrite sums-cong; simp)
    by (rule L-sums-T, simp add: that)
  qed

lemma exponential-e-d-l-mu:  $(\lambda k. d-(x+k) / (\$l-x * \mu-(x+k)))$  sums  $e^{\circ x}$ 
  if exponential-interpolation total-finite  $\bigwedge k::nat. \mu-(x+k) \neq 0$  for  $x::nat$ 
proof -
  let ?x = real x
  have  $\neg ereal ?x \geq \psi$  using that mu-beyond-0 by (metis add-cancel-right-right)
  hence [simp]:  $ereal ?x < \psi$  by simp
  have  $(\lambda k. d-(real (x+k)) / \mu-(real (x+k)) / \$l-?x)$  sums  $(T-?x / \$l-?x)$ 
    using sums-divide exponential-d-mu-sums-T that by force
  thus ?thesis by (rewrite e-T-l; simp add: mult.commute)
  qed

end

```

6.5 Limited Life Table

```

locale limited-life-table = life-table +
  assumes l-limited:  $\exists x::real. \$l-x = 0$ 
begin

lemma limited-survival-function-MM-X: limited-survival-function  $\mathfrak{M} X$ 
proof (rule limited-survival-function.intro)
  show survival-model  $\mathfrak{M} X$  by (rule survival-model-MM-X)
  show limited-survival-function-axioms  $\mathfrak{M} X$ 
    unfolding limited-survival-function-axioms-def using l-limited death-pt-def by
  fastforce
qed

end

sublocale limited-life-table  $\subseteq$  limited-survival-function  $\mathfrak{M} X$ 

```

```

by (rule limited-survival-function-MM-X)

context limited-life-table
begin

notation ult-age (<$ω>)

lemma l-omega-0: $l-$ω = 0
  using ccdfX-l-normal ccdfX-omega-0 by simp

lemma l-0-equiv-nat: $l-x = 0 ↔ x ≥ $ω for x::nat
  using ccdfX-l-normal ccdfX-0-equiv-nat by simp

lemma d-l-equiv-nat: $d-{t&x} = $l-x ↔ x+t ≥ $ω if t ≥ 0 for x t :: nat
  by (metis d-l-equiv of-nat-0-le-iff of-nat-add psi-le-iff-omega-le)

corollary d-1-omega-l: $d-($ω-1) = $l-($ω-1)
  using d-l-equiv-nat[of 1 $ω-1] omega-pos by simp

lemma limited-life-table-imp-total-finite: total-finite
proof -
  have {0..} = {0 .. real $ω} ∪ {real $ω <..} by force
  moreover have set-integrable lborel {0 .. real $ω} l by (rule l-integrable-Icc)
  moreover have set-integrable lborel {real $ω <..} l
    apply (rewrite set-integrable-cong[where f'=λ_. 0], simp-all)
    using l-0-equiv-nat apply (meson l-0-equiv le-ereal-le order-le-less)
    unfolding set-integrable-def by simp
  ultimately have set-integrable lborel {0..} l
    using set-integrable-Un
    by (smt (verit, del-insts) Ici-subset-Loi-iff add-mono-thms-linordered-field(1)
      atLeast-borel l-0-pos set-integrable-subset sets-lborel total-finite-iff-set-integrable-Ici)
    thus ?thesis unfolding total-finite-def by simp
qed

context
  fixes x::nat
  assumes x-lt-omega[simp]: x < $ω
begin

lemma curt-e-sum-l-finite-nat: $e-x = (∑ k< n. $l-(x+k+1)) / $l-x
  if ∀k::nat. k < n ==> isCont l (x+k+1) x+n ≥ $ω for n::nat
  apply (rewrite curt-e-sum-l-finite[of x n], simp)
  using that le-ereal-less psi-le-omega apply (metis of-nat-1 of-nat-add, force)
  by (simp add: add.commute)

end

end

```

```

end
theory Examples
imports Life-Table
begin

```

7 Examples

The following lemma is a verification of the solution to the multiple choice question No. 3 of Exam LTAM Spring 2022 by Society of Actuaries.

```

context smooth-survival-function
begin

```

```

lemma SoA-LTAM-2022-Spring-MCQ-No3:
assumes A:x::real. 0 ≤ x ⟹ x ≤ 100 ⟹ ccdf (distr M borel X) x = (1 −
0.01*x).^0.5
shows |1000*$μ-25 − 6.7| < 0.05
proof −
let ?svl = ccdf (distr M borel X)
have [simp]: ereal 25 < $ψ
apply (rewrite not-le[THEN sym])
using assms by (rewrite cdfX-0-equiv[THEN sym]) simp
have ∗: ((λx. ln (1 − x/100)) has-real-derivative (−1/75)) (at 25)
proof −
have ((λx. (1 − x/100)) has-real-derivative (0 − 1/100)) (at 25)
apply (intro derivative-intros)
by (rule DERIV-cdivide) simp
hence ((λx. ln (1 − x/100)) has-real-derivative (1 / (1 − 25/100) * (−1/100)))
(at 25)
by (intro derivative-intros) auto
thus ?thesis by simp
qed
have exp o (λx::real. 0.5 * ln (1 − 0.01*x)) field-differentiable at 25
apply (rule field-differentiable-compose, simp-all)
apply (rule derivative-intros, simp-all)
using ∗ field-differentiable-def apply blast
using field-differentiable-within-exp by blast
hence ?svl differentiable at 25
apply (rewrite differentiable-eq-field-differentiable-real)
by (rule field-differentiable-transform-within[where d=1])
(simp-all add: powr-def dist-real-def assms)
hence $μ-25 = − deriv (λx. ln (?svl x)) 25 by (rule mu-deriv-ln; simp)
also have ... = 0.005 / (1 − 0.01*25)
proof −
have ∀ F x in nhds 25. ln (?svl x) * 2 = ln (1 − x/100)
proof −
{ fix x::real assume dist x 25 < 1
hence asm: 0 < x x < 100 using dist-real-def by auto
hence ln (?svl x) * 2 = ln ((1 − 0.01*x).^0.5) * 2 using assms by (smt

```

```

(verit))
  also have ... = (0.5 * ln (1 - 0.01*x)) * 2 using asm by (rewrite
ln-powr) auto
    finally have ln (?svl x) * 2 = ln (1 - x/100) by simp }
    thus ?thesis by (rewrite eventually-nhds-metric) (smt (verit, del-insts))
qed
hence ((λx. ln (?svl x)) has-real-derivative (-0.005 / (1 - 0.01*25))) (at 25)
  using * apply (rewrite DERIV-cong-ev[where g=λx. 0.5 * ln (1 - 0.01*x)],
simp-all)
    by (rule derivative-eq-intros) auto
    thus ?thesis using DERIV-imp-deriv by force
qed
finally show ?thesis by simp
qed

end

```

The following lemma is a verification of the solution to the problem No. 2. (1)-1 of Life Insurance Mathematics 2016 by the Institute of Actuaries of Japan, slightly modified; see the remark below.

```

context smooth-life-table
begin

```

```

lemma IAJ-Life-Insurance-Math-2016-2-1-1:
  fixes a b :: real
  assumes -1 < a a < 0 0 < b -b/a ≤ $ψ and
    total-finite and
    ∀x. 0 < x ⇒ x < -b/a ⇒ l differentiable at x and
    ∀x. 0 ≤ x ⇒ x < -b/a ⇒ $e^x = a*x + b
  shows ∀x. 0 ≤ x ⇒ x < -b/a ⇒ $l-x = $l-0 * (b / (a*x + b)).~((a+1)/a)

proof -
  fix x assume asm-x: 0 ≤ x x < -b/a
  hence x-lt-psi[simp]: ereal x < $ψ using assms ereal-le-less by presburger
  from asm-x have axb-pos[simp]: a*x + b > 0
    using assms by (smt (verit, ccfv-threshold) mult.commute neg-less-divide-eq)
  have mu: ∀t. t ∈ {0 <.. < -b/a} ⇒ $μ-t = (a + 1) / (a*t + b)
  proof -
    fix t assume asm-t: t ∈ {0 <.. < -b/a}
    with assms have ((λu. $e^u) has-real-derivative ($μ-t * $e^t - 1)) (at t)
      by (intro e-has-derivative-mu-e-l'[where a=0]; simp)
    moreover have ((λu. $e^u) has-real-derivative a) (at t)
    proof -
      let ?d = min t (-b/a - t)
      have ?d > 0 using assms asm-t by simp
      moreover have ∀u. dist u t < ?d ⇒ $e^u = a*u + b using assms
        dist-real-def by auto
      ultimately have ∀F u in nhds t. $e^u = a*u + b by (rewrite eventually-nhds-metric) blast
      thus ?thesis
    qed
  qed

```

```

using assms apply (rewrite DERIV-cong-ev[where g=λt. a*t + b], simp-all)
  by (intro derivative-eq-intros) auto
qed
ultimately have $μ-t * $e^o-t - 1 = a using DERIV-unique by blast
moreover have $e^o-t = a*t + b using assms asm-t by simp
ultimately show $μ-t = (a + 1) / (a*t + b)
  using assms by (smt (verit) mult-eq-0-iff nonzero-mult-div-cancel-right)
qed
hence $p-{x&0} = (b / (a*x + b)).~((a+1)/a)
proof -
  have integ-mu: integral {0..x} (λt. $μ-t) = (a + 1) / a * ln ((a*x + b) / b)
  proof -
    have integral {0..x} (λt. $μ-t) = integral {0<..x} (λt. $μ-t)
      apply (rule integral-spike-set)
      apply (rule negligible-subset[where s={0}]; force)
      by (rule negligible-subset[where s={}]; simp)
    also have ... = integral {0<..x} (λt. ((a + 1) / a) * (a / (a*t + b)))
      using asm-x assms by (intro integral-cong) (rewrite mu; simp)
    also have ... = (a + 1) / a * integral {0<..x} (λt. a / (a*t + b))
      by (metis integral-mult-right)
    also have ... = (a + 1) / a * ln ((a*x + b) / b)
    proof -
      have integral {0<..x} (λt. a / (a*t + b)) = integral {0..x} (λt. a / (a*t +
b))
        apply (rule integral-spike-set)
        apply (rule negligible-subset[where s={}]; simp)
        by (rule negligible-subset[where s={0}]; force)
      also have ... = ln (a*x + b) - ln (a*0 + b)
        apply (rule integral-unique)
        using assms asm-x apply (intro inverse-fun-has-integral-ln, simp-all)
        using axb-pos assms apply (smt (verit) mult-less-cancel-left)
        apply (intro continuous-intros)
        by (intro derivative-eq-intros) auto
      also have ... = ln ((a*x + b) / b) using assms by (simp add: ln-divide-pos)
      finally have integral {0<..x} (λt. a / (a*t + b)) = ln ((a*x + b) / b) .
      thus ?thesis by simp
    qed
    finally show ?thesis .
  qed
  thus ?thesis
    apply (rewrite p-exp-integral-mu, simp-all add: asm-x)
    unfolding powr-def using assms
    by simp (smt (verit) axb-pos ln-div
      nonzero-minus-divide-divide nonzero-minus-divide-right right-diff-distrib')
  qed
  thus $l-x = $l-0 * (b / (a*x + b)).~((a+1)/a)
    using assms asm-x apply (rewrite in asm p-l, simp-all)
    by (metis divide-mult-cancel l-0-neq-0 mult.commute)
qed

```

REMARK. The original problem lacks the following hypotheses: (i) $0 < b$,
(ii) $-b/a \leq \$\psi$, (iii) $\forall x. 0 < x < -b/a \rightarrow l$ differentiable at x , (iv) $\forall x. 0 \leq x < -b/a \rightarrow l$ integrable-on $\{x..\}$. Moreover, the hypothesis $\forall x. 0 \leq x < -b/a$ is originally $\forall x. 0 \leq x \leq -b/a$.

end

end