

# Abstract Substitutions as Monoid Actions

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## Abstract

This entry provides a small, reusable, theory that specifies the abstract concept of substitution as monoid action. Both the substitution type and the object type are kept abstract. The theory provides multiple useful definitions and lemmas. Two example usages are provided for first order terms: one for terms from the AFP/First\_Order\_Terms session and one for terms from the Isabelle/HOL-ex session.

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```

theory Monoid-Action
imports Main
begin

```

## 1 General Results on Groups

```

lemma (in monoid) right-inverse-idem:
fixes inv
assumes right-inverse:  $\bigwedge a. a * \text{inv } a = \mathbf{1}$ 
shows  $\bigwedge a. \text{inv } (\text{inv } a) = a$ 
⟨proof⟩

lemma (in monoid) left-inverse-if-right-inverse:
fixes inv
assumes
right-inverse:  $\bigwedge a. a * \text{inv } a = \mathbf{1}$ 
shows  $\text{inv } a * a = \mathbf{1}$ 
⟨proof⟩

lemma (in monoid) group-wrt-right-inverse:
fixes inv
assumes right-inverse:  $\bigwedge a. a * \text{inv } a = \mathbf{1}$ 
shows group (*)  $\mathbf{1} \text{ inv}$ 
⟨proof⟩

```

## 2 Monoid

```

definition (in monoid) is-left-invertible where
is-left-invertible a  $\longleftrightarrow$  ( $\exists a\text{-inv}. a\text{-inv} * a = \mathbf{1}$ )

definition (in monoid) is-right-invertible where
is-right-invertible a  $\longleftrightarrow$  ( $\exists a\text{-inv}. a * a\text{-inv} = \mathbf{1}$ )

definition (in monoid) left-inverse where
is-left-invertible a  $\implies$  left-inverse a = (SOME a-inv. a-inv * a =  $\mathbf{1}$ )

definition (in monoid) right-inverse where
is-right-invertible a  $\implies$  right-inverse a = (SOME a-inv. a * a-inv =  $\mathbf{1}$ )

lemma (in monoid) comp-left-inverse [simp]:
is-left-invertible a  $\implies$  left-inverse a * a =  $\mathbf{1}$ 
⟨proof⟩

lemma (in monoid) comp-right-inverse [simp]:
is-right-invertible a  $\implies$  a * right-inverse a =  $\mathbf{1}$ 
⟨proof⟩

lemma (in monoid) neutral-is-left-invertible [simp]:

```

```
is-left-invertible 1
⟨proof⟩
```

```
lemma (in monoid) neutral-is-right-invertible [simp]:
is-right-invertible 1
⟨proof⟩
```

### 3 Semigroup Action

We define both left and right semigroup actions. Left semigroup actions seem to be prevalent in algebra, but right semigroup actions directly uses the usual notation of term/atom/literal/clause substitution.

```
locale left-semigroup-action = semigroup +
  fixes action :: 'a ⇒ 'b ⇒ 'b (infix ‘70’)
  assumes action-compatibility[simp]:  $\bigwedge a b x. (a * b) \cdot x = a \cdot (b \cdot x)$ 
```

```
locale right-semigroup-action = semigroup +
  fixes action :: 'b ⇒ 'a ⇒ 'b (infix ‘70’)
  assumes action-compatibility[simp]:  $\bigwedge x a b. x \cdot (a * b) = (x \cdot a) \cdot b$ 
```

We then instantiate the right action in the context of the left action in order to get access to any lemma proven in the context of the other locale. We do analogously in the context of the right locale.

```
sublocale left-semigroup-action ⊆ right: right-semigroup-action where
  f =  $\lambda x y. f y x$  and action =  $\lambda x y. action y x$ 
⟨proof⟩
```

```
sublocale right-semigroup-action ⊆ left: left-semigroup-action where
  f =  $\lambda x y. f y x$  and action =  $\lambda x y. action y x$ 
⟨proof⟩
```

```
lemma (in right-semigroup-action) lifting-semigroup-action-to-set:
right-semigroup-action (*) ( $\lambda X a. (\lambda x. action x a) ` X$ )
⟨proof⟩
```

```
lemma (in right-semigroup-action) lifting-semigroup-action-to-list:
right-semigroup-action (*) ( $\lambda xs a. map (\lambda x. action x a) xs$ )
⟨proof⟩
```

### 4 Monoid Action

```
locale left-monoid-action = monoid +
  fixes action :: 'a ⇒ 'b ⇒ 'b (infix ‘70’)
  assumes
    monoid-action-compatibility:  $\bigwedge a b x. (a * b) \cdot x = a \cdot (b \cdot x)$  and
    action-neutral[simp]:  $\bigwedge x. 1 \cdot x = x$ 
```

```

locale right-monoid-action = monoid +
  fixes action :: 'b ⇒ 'a ⇒ 'b (infix ← 70)
  assumes
    monoid-action-compatibility:  $\bigwedge x a b. x \cdot (a * b) = (x \cdot a) \cdot b$  and
    action-neutral[simp]:  $\bigwedge x. x \cdot \mathbf{1} = x$ 

sublocale left-monoid-action ⊆ left-semigroup-action
  ⟨proof⟩

sublocale right-monoid-action ⊆ right-semigroup-action
  ⟨proof⟩

sublocale left-monoid-action ⊆ right: right-monoid-action where
  f =  $\lambda x y. f y x$  and action =  $\lambda x y. action y x$ 
  ⟨proof⟩

sublocale right-monoid-action ⊆ left: left-monoid-action where
  f =  $\lambda x y. f y x$  and action =  $\lambda x y. action y x$ 
  ⟨proof⟩

lemma (in right-monoid-action) lifting-monoid-action-to-set:
  right-monoid-action (*)  $\mathbf{1} (\lambda X a. (\lambda x. action x a) ` X)$ 
  ⟨proof⟩

lemma (in right-monoid-action) lifting-monoid-action-to-list:
  right-monoid-action (*)  $\mathbf{1} (\lambda xs a. map (\lambda x. action x a) xs)$ 
  ⟨proof⟩

```

## 5 Group Action

```

locale left-group-action = group +
  fixes action :: 'a ⇒ 'b ⇒ 'b (infix ← 70)
  assumes
    group-action-compatibility:  $\bigwedge a b x. (a * b) \cdot x = a \cdot (b \cdot x)$  and
    group-action-neutral:  $\bigwedge x. \mathbf{1} \cdot x = x$ 

locale right-group-action = group +
  fixes action :: 'b ⇒ 'a ⇒ 'b (infixl ← 70)
  assumes
    group-action-compatibility:  $\bigwedge x a b. x \cdot (a * b) = (x \cdot a) \cdot b$  and
    group-action-neutral:  $\bigwedge x. x \cdot \mathbf{1} = x$ 

sublocale left-group-action ⊆ left-monoid-action
  ⟨proof⟩

sublocale right-group-action ⊆ right-monoid-action
  ⟨proof⟩

sublocale left-group-action ⊆ right: right-group-action where

```

```

 $f = \lambda x y. f y x$  and  $action = \lambda x y. action y x$ 
⟨proof⟩

sublocale right-group-action ⊆ left: left-group-action where
 $f = \lambda x y. f y x$  and  $action = \lambda x y. action y x$ 
⟨proof⟩

```

```

end
theory Substitution
  imports Monoid-Action
begin

```

```

abbreviation set-prod where
  set-prod ≡  $\lambda(t, t'). \{t, t'\}$ 

```

## 6 Assumption-free Substitution

```

locale substitution-ops =
  fixes
    subst :: ' $x \Rightarrow s \Rightarrow x$ ' (infixl  $\leftrightarrow$  67) and
    id-subst :: ' $s$ ' and
    comp-subst :: ' $s \Rightarrow s \Rightarrow s$ ' (infixl  $\circ\circ$  67) and
    is-ground :: ' $x \Rightarrow \text{bool}$ '
begin

definition subst-set :: ' $x \text{ set} \Rightarrow s \Rightarrow x \text{ set}$ ' where
  subst-set  $X \sigma = (\lambda x. subst x \sigma) ` X$ 

definition subst-list :: ' $x \text{ list} \Rightarrow s \Rightarrow x \text{ list}$ ' where
  subst-list  $xs \sigma = map (\lambda x. subst x \sigma) xs$ 

definition is-ground-set :: ' $x \text{ set} \Rightarrow \text{bool}$ ' where
  is-ground-set  $X \longleftrightarrow (\forall x \in X. is-ground x)$ 

definition is-ground-subst :: ' $s \Rightarrow \text{bool}$ ' where
  is-ground-subst  $\gamma \longleftrightarrow (\forall x. is-ground (x \cdot \gamma))$ 

definition generalizes :: ' $x \Rightarrow s \Rightarrow \text{bool}$ ' where
  generalizes  $x y \longleftrightarrow (\exists \sigma. x \cdot \sigma = y)$ 

definition specializes :: ' $x \Rightarrow s \Rightarrow \text{bool}$ ' where
  specializes  $x y \equiv generalizes y x$ 

definition strictly-generalizes :: ' $x \Rightarrow s \Rightarrow \text{bool}$ ' where
  strictly-generalizes  $x y \longleftrightarrow generalizes x y \wedge \neg generalizes y x$ 

definition strictly-specializes :: ' $x \Rightarrow s \Rightarrow \text{bool}$ ' where
  strictly-specializes  $x y \equiv strictly-generalizes y x$ 

```

```

definition instances :: ' $x \Rightarrow 'x \text{ set}$  where
  instances  $x = \{y. \text{generalizes } x \ y\}$ 

definition instances-set :: ' $x \text{ set} \Rightarrow 'x \text{ set}$  where
  instances-set  $X = (\bigcup_{x \in X}. \text{instances } x)$ 

definition ground-instances :: ' $x \Rightarrow 'x \text{ set}$  where
  ground-instances  $x = \{x_G \in \text{instances } x. \text{is-ground } x_G\}$ 

definition ground-instances-set :: ' $x \text{ set} \Rightarrow 'x \text{ set}$  where
  ground-instances-set  $X = \{x_G \in \text{instances-set } X. \text{is-ground } x_G\}$ 

lemma ground-instances-set-eq-Union-ground-instances:
  ground-instances-set  $X = (\bigcup_{x \in X}. \text{ground-instances } x)$ 
   $\langle \text{proof} \rangle$ 

lemma ground-instances-eq-Collect-subst-grounding:
  ground-instances  $x = \{x \cdot \gamma \mid \gamma. \text{is-ground } (x \cdot \gamma)\}$ 
   $\langle \text{proof} \rangle$ 

lemma mem-ground-instancesE[elim]:
  fixes  $x \ x_G :: 'x$ 
  assumes  $x_G \in \text{ground-instances } x$ 
  obtains  $\gamma :: 's \text{ where } x_G = x \cdot \gamma \text{ and is-ground } (x \cdot \gamma)$ 
   $\langle \text{proof} \rangle$ 

lemma mem-ground-instances-setE[elim]:
  fixes  $x_G :: 'x \text{ and } X :: 'x \text{ set}$ 
  assumes  $x_G \in \text{ground-instances-set } X$ 
  obtains  $x :: 'x \text{ and } \gamma :: 's \text{ where } x \in X \text{ and } x_G = x \cdot \gamma \text{ and is-ground } (x \cdot \gamma)$ 
   $\langle \text{proof} \rangle$ 

definition is-unifier :: ' $s \Rightarrow 'x \text{ set} \Rightarrow \text{bool}$  where
  is-unifier  $v \ X \longleftrightarrow \text{card } (\text{subst-set } X \ v) \leq 1$ 

definition is-unifier-set :: ' $s \Rightarrow 'x \text{ set set} \Rightarrow \text{bool}$  where
  is-unifier-set  $v \ XX \longleftrightarrow (\forall X \in XX. \text{is-unifier } v \ X)$ 

definition is-mgu :: ' $s \Rightarrow 'x \text{ set set} \Rightarrow \text{bool}$  where
  is-mgu  $\mu \ XX \longleftrightarrow \text{is-unifier-set } \mu \ XX \wedge (\forall v. \text{is-unifier-set } v \ XX \longrightarrow (\exists \sigma. \mu \odot \sigma = v))$ 

definition is-imgu :: ' $s \Rightarrow 'x \text{ set set} \Rightarrow \text{bool}$  where
  is-imgu  $\mu \ XX \longleftrightarrow \text{is-unifier-set } \mu \ XX \wedge (\forall \tau. \text{is-unifier-set } \tau \ XX \longrightarrow \mu \odot \tau = \tau)$ 

lemma is-unifier-iff-if-finite:
  assumes finite  $X$ 
  shows is-unifier  $\sigma \ X \longleftrightarrow (\forall x \in X. \forall y \in X. x \cdot \sigma = y \cdot \sigma)$ 

```

$\langle proof \rangle$

**lemma** *is-unifier-singleton*[simp]: *is-unifier*  $v \{x\}$   
 $\langle proof \rangle$

**lemma** *is-unifier-set-empty*[simp]:  
*is-unifier-set*  $\sigma \{\}$   
 $\langle proof \rangle$

**lemma** *is-unifier-set-insert*:  
*is-unifier-set*  $\sigma (\text{insert } X XX) \longleftrightarrow \text{is-unifier } \sigma X \wedge \text{is-unifier-set } \sigma XX$   
 $\langle proof \rangle$

**lemma** *is-unifier-set-insert-singleton*[simp]:  
*is-unifier-set*  $\sigma (\text{insert } \{x\} XX) \longleftrightarrow \text{is-unifier-set } \sigma XX$   
 $\langle proof \rangle$

**lemma** *is-mgu-insert-singleton*[simp]: *is-mgu*  $\mu (\text{insert } \{x\} XX) \longleftrightarrow \text{is-mgu } \mu XX$   
 $\langle proof \rangle$

**lemma** *is-imgu-insert-singleton*[simp]: *is-imgu*  $\mu (\text{insert } \{x\} XX) \longleftrightarrow \text{is-imgu } \mu XX$   
 $\langle proof \rangle$

**lemma** *subst-set-empty*[simp]: *subst-set*  $\{\} \sigma = \{\}$   
 $\langle proof \rangle$

**lemma** *subst-set-insert*[simp]: *subst-set*  $(\text{insert } x X) \sigma = \text{insert } (x \cdot \sigma)$  (*subst-set*  $X \sigma$ )  
 $\langle proof \rangle$

**lemma** *subst-set-union*[simp]: *subst-set*  $(X1 \cup X2) \sigma = \text{subst-set } X1 \sigma \cup \text{subst-set } X2 \sigma$   
 $\langle proof \rangle$

**lemma** *subst-list-Nil*[simp]: *subst-list*  $[] \sigma = []$   
 $\langle proof \rangle$

**lemma** *subst-list-insert*[simp]: *subst-list*  $(x \# xs) \sigma = (x \cdot \sigma) \# (\text{subst-list } xs \sigma)$   
 $\langle proof \rangle$

**lemma** *subst-list-append*[simp]: *subst-list*  $(xs1 @ xs2) \sigma = \text{subst-list } xs1 \sigma @ \text{subst-list } xs2 \sigma$   
 $\langle proof \rangle$

**lemma** *is-unifier-set-union*:  
*is-unifier-set*  $v (XX_1 \cup XX_2) \longleftrightarrow \text{is-unifier-set } v XX_1 \wedge \text{is-unifier-set } v XX_2$   
 $\langle proof \rangle$

```

lemma is-unifier-subset: is-unifier v A  $\Rightarrow$  finite A  $\Rightarrow$  B  $\subseteq$  A  $\Rightarrow$  is-unifier v
B
    ⟨proof⟩

lemma is-ground-set-subset: is-ground-set A  $\Rightarrow$  B  $\subseteq$  A  $\Rightarrow$  is-ground-set B
    ⟨proof⟩

lemma is-ground-set-ground-instances[simp]: is-ground-set (ground-instances x)
    ⟨proof⟩

lemma is-ground-set-ground-instances-set[simp]: is-ground-set (ground-instances-set
x)
    ⟨proof⟩

end

```

## 7 Basic Substitution

```

locale substitution-monoid = monoid comp-subst id-subst
  for
    comp-subst :: 's  $\Rightarrow$  's  $\Rightarrow$  's and
    id-subst :: 's
begin

  abbreviation is-renaming where
    is-renaming  $\equiv$  is-right-invertible

  lemmas is-renaming-def = is-right-invertible-def

  abbreviation renaming-inverse where
    renaming-inverse  $\equiv$  right-inverse

  lemmas renaming-inverse-def = right-inverse-def

  lemmas is-renaming-id-subst = neutral-is-right-invertible

  definition is-idem :: 's  $\Rightarrow$  bool where
    is-idem a  $\longleftrightarrow$  comp-subst a a = a

  lemma is-idem-id-subst [simp]: is-idem id-subst
    ⟨proof⟩

end

locale substitution =
  substitution-monoid comp-subst id-subst +
  comp-subst: right-monoid-action comp-subst id-subst subst +
  substitution-ops subst id-subst comp-subst is-ground

```

```

for
  comp-subst :: 's ⇒ 's ⇒ 's (infixl ⟨ $\odot$ ⟩ 70) and
  id-subst :: 's and
  subst :: 'x ⇒ 's ⇒ 'x (infixl ⟨ $\leftrightarrow$ ⟩ 69) and

  — Predicate identifying the fixed elements w.r.t. the monoid action
  is-ground :: 'x ⇒ bool +
assumes
  all-subst-ident-if-ground: is-ground x  $\implies$  ( $\forall \sigma. x \cdot \sigma = x$ )
begin

sublocale comp-subst-set: right-monoid-action comp-subst id-subst subst-set
  ⟨proof⟩

sublocale comp-subst-list: right-monoid-action comp-subst id-subst subst-list
  ⟨proof⟩

```

## 7.1 Substitution Composition

```

lemmas subst-comp-subst = comp-subst.action-compatibility
lemmas subst-set-comp-subst = comp-subst-set.action-compatibility
lemmas subst-list-comp-subst = comp-subst-list.action-compatibility

```

## 7.2 Substitution Identity

```

lemmas subst-id-subst = comp-subst.action-neutral
lemmas subst-set-id-subst = comp-subst-set.action-neutral
lemmas subst-list-id-subst = comp-subst-list.action-neutral

```

```

lemma is-unifier-id-subst-empty[simp]: is-unifier id-subst {}
  ⟨proof⟩

```

```

lemma is-unifier-set-id-subst-empty[simp]: is-unifier-set id-subst {}
  ⟨proof⟩

```

```

lemma is-mgu-id-subst-empty[simp]: is-mgu id-subst {}
  ⟨proof⟩

```

```

lemma is-imgu-id-subst-empty[simp]: is-imgu id-subst {}
  ⟨proof⟩

```

```

lemma is-unifier-id-subst: is-unifier id-subst  $X \longleftrightarrow \text{card } X \leq 1$ 
  ⟨proof⟩

```

```

lemma is-unifier-set-id-subst: is-unifier-set id-subst  $XX \longleftrightarrow (\forall X \in XX. \text{card } X \leq 1)$ 
  ⟨proof⟩

```

```

lemma is-mgu-id-subst: is-mgu id-subst  $XX \longleftrightarrow (\forall X \in XX. \text{card } X \leq 1)$ 
  ⟨proof⟩

```

```
lemma is-imgu-id-subst: is-imgu id-subst  $XX \longleftrightarrow (\forall X \in XX. \text{card } X \leq 1)$ 
   $\langle \text{proof} \rangle$ 
```

### 7.3 Generalization

```
sublocale generalizes: preorder generalizes strictly-generalizes
   $\langle \text{proof} \rangle$ 
```

```
lemma generalizes-antisym-if:
  assumes  $\bigwedge \sigma_1 \sigma_2 x. x \cdot (\sigma_1 \odot \sigma_2) = x \implies x \cdot \sigma_1 = x$ 
  shows  $\bigwedge x y. \text{generalizes } x y \implies \text{generalizes } y x \implies x = y$ 
   $\langle \text{proof} \rangle$ 
```

```
lemma order-generalizes-if:
  assumes  $\bigwedge \sigma_1 \sigma_2 x. x \cdot (\sigma_1 \odot \sigma_2) = x \implies x \cdot \sigma_1 = x$ 
  shows class.order generalizes strictly-generalizes
   $\langle \text{proof} \rangle$ 
```

### 7.4 Substituting on Ground Expressions

```
lemma subst-ident-if-ground[simp]: is-ground  $x \implies x \cdot \sigma = x$ 
   $\langle \text{proof} \rangle$ 
```

```
lemma subst-set-ident-if-ground[simp]: is-ground-set  $X \implies \text{subst-set } X \sigma = X$ 
   $\langle \text{proof} \rangle$ 
```

### 7.5 Instances of Ground Expressions

```
lemma instances-ident-if-ground[simp]: is-ground  $x \implies \text{instances } x = \{x\}$ 
   $\langle \text{proof} \rangle$ 
```

```
lemma instances-set-ident-if-ground[simp]: is-ground-set  $X \implies \text{instances-set } X = X$ 
   $\langle \text{proof} \rangle$ 
```

```
lemma ground-instances-ident-if-ground[simp]: is-ground  $x \implies \text{ground-instances } x = \{x\}$ 
   $\langle \text{proof} \rangle$ 
```

```
lemma ground-instances-set-ident-if-ground[simp]: is-ground-set  $X \implies \text{ground-instances-set } X = X$ 
   $\langle \text{proof} \rangle$ 
```

### 7.6 Unifier of Ground Expressions

```
lemma ground-eq-ground-if-unifiable:
  assumes is-unifier  $v \{t_1, t_2\}$  and is-ground  $t_1$  and is-ground  $t_2$ 
  shows  $t_1 = t_2$ 
   $\langle \text{proof} \rangle$ 
```

**lemma** *ball-eq-constant-if-unifier*:  
**assumes** *finite X and*  $x \in X$  **and** *is-unifier v X and* *is-ground-set X*  
**shows**  $\forall y \in X. y = x$   
*{proof}*

**lemma** *is-mgu-unifies*:  
**assumes** *is-mgu  $\mu$  XX*  $\forall X \in XX. finite X$   
**shows**  $\forall X \in XX. \forall t \in X. \forall t' \in X. t \cdot \mu = t' \cdot \mu$   
*{proof}*

**corollary** *is-mgu-unifies-pair*:  
**assumes** *is-mgu  $\mu \{\{t, t'\}\}$*   
**shows**  $t \cdot \mu = t' \cdot \mu$   
*{proof}*

**lemmas** *subst-mgu-eq-subst-mgu = is-mgu-unifies-pair*

**lemma** *is-imgu-unifies*:  
**assumes** *is-imgu  $\mu$  XX*  $\forall X \in XX. finite X$   
**shows**  $\forall X \in XX. \forall t \in X. \forall t' \in X. t \cdot \mu = t' \cdot \mu$   
*{proof}*

**corollary** *is-imgu-unifies-pair*:  
**assumes** *is-imgu  $\mu \{\{t, t'\}\}$*   
**shows**  $t \cdot \mu = t' \cdot \mu$   
*{proof}*

**lemmas** *subst-imgu-eq-subst-imgu = is-imgu-unifies-pair*

## 7.7 Ground Substitutions

**lemma** *is-ground-subst-comp-left*: *is-ground-subst  $\sigma \implies is-ground-subst (\sigma \odot \tau)$*   
*{proof}*

**lemma** *is-ground-subst-comp-right*: *is-ground-subst  $\tau \implies is-ground-subst (\sigma \odot \tau)$*   
*{proof}*

**lemma** *is-ground-subst-is-ground*:  
**assumes** *is-ground-subst  $\gamma$*   
**shows** *is-ground ( $t \cdot \gamma$ )*  
*{proof}*

## 7.8 IMGU is Idempotent and an MGU

**lemma** *is-imgu-iff-is-idem-and-is-mgu*: *is-imgu  $\mu$  XX  $\longleftrightarrow is-idem \mu \wedge is-mgu \mu$*   
*XX*  
*{proof}*

## 7.9 IMGU can be used before unification

```
lemma subst-imgu-subst-unifier:  
  assumes unif: is-unifier v X and imgu: is-imgu μ {X} and x ∈ X  
  shows x · μ · v = x · v  
(proof)
```

## 7.10 Groundings Idempotence

```
lemma image-ground-instances-ground-instances:  
  ground-instances ` ground-instances x = (λx. {x}) ` ground-instances x  
(proof)
```

```
lemma grounding-of-set-grounding-of-set-idem[simp]:  
  ground-instances-set (ground-instances-set X) = ground-instances-set X  
(proof)
```

## 7.11 Instances of Substitution

```
lemma instances-subst:  
  instances (x · σ) ⊆ instances x  
(proof)
```

```
lemma instances-set-subst-set:  
  instances-set (subst-set X σ) ⊆ instances-set X  
(proof)
```

```
lemma ground-instances-subst:  
  ground-instances (x · σ) ⊆ ground-instances x  
(proof)
```

```
lemma ground-instances-set-subst-set:  
  ground-instances-set (subst-set X σ) ⊆ ground-instances-set X  
(proof)
```

## 7.12 Instances of Renamed Expressions

```
lemma instances-subst-ident-if-renaming[simp]:  
  is-renaming ρ ==> instances (x · ρ) = instances x  
(proof)
```

```
lemma instances-set-subst-set-ident-if-renaming[simp]:  
  is-renaming ρ ==> instances-set (subst-set X ρ) = instances-set X  
(proof)
```

```
lemma ground-instances-subst-ident-if-renaming[simp]:  
  is-renaming ρ ==> ground-instances (x · ρ) = ground-instances x  
(proof)
```

```
lemma ground-instances-set-subst-set-ident-if-renaming[simp]:
```

```

is-renaming  $\varrho \implies \text{ground-instances-set} (\text{subst-set } X \varrho) = \text{ground-instances-set } X$ 
⟨proof⟩

end

end
theory Substitution-First-Order-Term
imports
  Substitution
  First-Order-Terms.Unification
begin

abbreviation is-ground-trm where
  is-ground-trm t ≡ vars-term t = {}

lemma is-ground-iff: is-ground-trm (t · γ)  $\longleftrightarrow (\forall x \in \text{vars-term } t. \text{is-ground-trm}(\gamma x))$ 
⟨proof⟩

lemma is-ground-trm-iff-ident-forall-subst: is-ground-trm t  $\longleftrightarrow (\forall \sigma. t \cdot \sigma = t)$ 
⟨proof⟩

global-interpretation term-subst: substitution where
  subst = subst-apply-term and id-subst = Var and comp-subst = subst-compose
and
  is-ground = is-ground-trm
⟨proof⟩

lemma term-subst-is-unifier-iff-unifiers-Times:
assumes finite X
shows term-subst.is-unifier μ X  $\longleftrightarrow \mu \in \text{unifiers}(X \times X)$ 
⟨proof⟩

lemma term-subst-is-unifier-set-iff-unifiers-Union-Times:
assumes  $\forall X \in XX. \text{finite } X$ 
shows term-subst.is-unifier-set μ XX  $\longleftrightarrow \mu \in \text{unifiers}(\bigcup X \in XX. X \times X)$ 
⟨proof⟩

lemma term-subst-is-mgu-iff-is-mgu-Union-Times:
assumes fin:  $\forall X \in XX. \text{finite } X$ 
shows term-subst.is-mgu μ XX  $\longleftrightarrow \text{is-mgu } \mu (\bigcup X \in XX. X \times X)$ 
⟨proof⟩

lemma term-subst-is-imgu-iff-is-imgu-Union-Times:
assumes  $\forall X \in XX. \text{finite } X$ 
shows term-subst.is-imgu μ XX  $\longleftrightarrow \text{is-imgu } \mu (\bigcup X \in XX. X \times X)$ 
⟨proof⟩

```

```

lemma range-vars-subset-if-is-imgu:
  assumes term-subst.is-imgu  $\mu$   $XX \forall X \in XX. \text{finite } X \text{ finite } XX$ 
  shows range-vars  $\mu \subseteq (\bigcup t \in \bigcup XX. \text{vars-term } t)$ 
  ⟨proof⟩

lemma term-subst-is-renaming-iff:
  term-subst.is-renaming  $\varrho \longleftrightarrow \text{inj } \varrho \wedge (\forall x. \text{is-Var } (\varrho x))$ 
  ⟨proof⟩

lemma term-subst-is-renaming-iff-ex-inj-fun-on-vars:
  term-subst.is-renaming  $\varrho \longleftrightarrow (\exists f. \text{inj } f \wedge \varrho = \text{Var } \circ f)$ 
  ⟨proof⟩

lemma ground-imgu-equals:
  assumes is-ground-trm  $t_1$  and is-ground-trm  $t_2$  and term-subst.is-imgu  $\mu \{\{t_1, t_2\}\}$ 
  shows  $t_1 = t_2$ 
  ⟨proof⟩

lemma is-unifier-the-mgu:
  assumes  $t \cdot \text{the-mgu } t t' = t' \cdot \text{the-mgu } t t'$ 
  shows term-subst.is-unifier (the-mgu  $t t' \{\{t, t'\}$ )
  ⟨proof⟩

lemma obtains-imgu-from-unifier-and-the-mgu:
  fixes  $v :: ('f, 'v) \text{ subst}$ 
  assumes  $t \cdot v = t' \cdot v P t t' (\text{Unification.the-mgu } t t')$ 
  obtains  $\mu :: ('f, 'v) \text{ subst}$ 
  where  $v = \mu \circ_s v \text{ term-subst.is-imgu } \mu \{\{t, t'\}\} P t t' \mu$ 
  ⟨proof⟩

lemma obtains-imgu:
  fixes  $v :: ('f, 'v) \text{ subst}$ 
  assumes  $t \cdot v = t' \cdot v$ 
  obtains  $\mu :: ('f, 'v) \text{ subst}$ 
  where  $v = \mu \circ_s v \text{ term-subst.is-imgu } \mu \{\{t, t'\}\}$ 
  ⟨proof⟩

lemma is-renaming-if-term-subst-is-renaming:
  assumes term-subst.is-renaming  $\varrho$ 
  shows Term.is-renaming  $\varrho$ 
  ⟨proof⟩

lemma is-mgu-iff-term-subst-is-imgu-image-set-prod:
  fixes  $\mu :: ('f, 'v) \text{ subst}$  and  $X :: (('f, 'v) \text{ term} \times ('f, 'v) \text{ term}) \text{ set}$ 
  shows Unifiers.is-imgu  $\mu X \longleftrightarrow \text{term-subst.is-imgu } \mu (\text{set-prod } ' X)$ 
  ⟨proof⟩

```

```

lemma the-mgu-term-subst-is-imgu:
  fixes v :: ('f, 'v) subst
  assumes s · v = t · v
  shows term-subst.is-imgu (Unification.the-mgu s t) {{s, t}}
  ⟨proof⟩

end

theory Substitution-HOL-ex-Unification
imports
  Substitution
  HOL-ex.Unification
begin

no-notation Comb (infix `·` 60)

quotient-type 'a subst = ('a × 'a trm) list / (⊣)
⟨proof⟩

lift-definition subst-comp :: 'a subst ⇒ 'a subst ⇒ 'a subst (infixl `○` 67)
  is Unification.comp
  ⟨proof⟩

definition subst-id :: 'a subst where
  subst-id = abs-subst []

global-interpretation subst-comp: monoid subst-comp subst-id
⟨proof⟩

lift-definition subst-apply :: 'a trm ⇒ 'a subst ⇒ 'a trm
  is Unification.subst
  ⟨proof⟩

abbreviation is-ground-trm where
  is-ground-trm t ≡ vars-of t = {}

global-interpretation term-subst: substitution where
  subst = subst-apply and id-subst = subst-id and comp-subst = subst-comp and
  is-ground = is-ground-trm
⟨proof⟩

end

theory Functional-Substitution
imports
  Substitution
  HOL-Library.FSet
begin

locale functional-substitution = substitution where
  subst = subst and is-ground = λexpr. vars expr = {}

```

```

for
  subst :: 'expr  $\Rightarrow$  ('var  $\Rightarrow$  'base)  $\Rightarrow$  'expr (infixl · 69) and
  vars :: 'expr  $\Rightarrow$  'var set +
assumes
  subst-eq:  $\bigwedge \text{expr } \sigma \tau. (\bigwedge x. x \in \text{vars expr} \implies \sigma x = \tau x) \implies \text{expr} \cdot \sigma = \text{expr} \cdot \tau$ 
begin

abbreviation is-ground where is-ground expr  $\equiv$  vars expr = {}

definition vars-set :: 'expr set  $\Rightarrow$  'var set where
  vars-set exprs  $\equiv$   $\bigcup \text{expr} \in \text{exprs. vars expr}$ 

lemma subst-redundant-upd [simp]:
assumes var  $\notin$  vars expr
shows expr  $\cdot$   $\sigma(\text{var} := \text{update}) = \text{expr} \cdot \sigma$ 
<proof>

lemma subst-redundant-if [simp]:
assumes vars expr  $\subseteq$  vars'
shows expr  $\cdot$   $(\lambda \text{var. if } \text{var} \in \text{vars}' \text{ then } \sigma \text{ var else } \sigma' \text{ var}) = \text{expr} \cdot \sigma$ 
<proof>

lemma subst-redundant-if' [simp]:
assumes vars expr  $\cap$  vars' = {}
shows expr  $\cdot$   $(\lambda \text{var. if } \text{var} \in \text{vars}' \text{ then } \sigma' \text{ var else } \sigma \text{ var}) = \text{expr} \cdot \sigma$ 
<proof>

lemma subst-cannot-unground:
assumes  $\neg \text{is-ground} (\text{expr} \cdot \sigma)$ 
shows  $\neg \text{is-ground} \text{ expr}$ 
<proof>

definition subst-domain :: ('var  $\Rightarrow$  'base)  $\Rightarrow$  'var set where
  subst-domain σ = {x. σ x  $\neq$  id-subst x}

abbreviation subst-range :: ('var  $\Rightarrow$  'base)  $\Rightarrow$  'base set where
  subst-range σ  $\equiv$  σ ` subst-domain σ

lemma subst-inv:
assumes σ ⊕ σ-inv = id-subst
shows expr  $\cdot$  σ  $\cdot$  σ-inv = expr
<proof>

definition rename where
  is-renaming ρ  $\implies$  rename ρ x  $\equiv$  SOME x'. ρ x = id-subst x'

end

```

```

locale all-subst-ident-iff-ground =
  functional-substitution +
  assumes
    all-subst-ident-iff-ground:  $\bigwedge \text{expr. is-ground expr} \longleftrightarrow (\forall \sigma. \text{subst expr } \sigma = \text{expr})$ 
  and
    exists-non-ident-subst:
       $\bigwedge \text{expr } S. \text{finite } S \implies \neg \text{is-ground expr} \implies \exists \sigma. \text{subst expr } \sigma \neq \text{expr} \wedge \text{subst expr } \sigma \notin S$ 

locale finite-variables = functional-substitution where vars = vars
  for vars :: 'expr  $\Rightarrow$  'var set +
  assumes finite-vars [intro]:  $\bigwedge \text{expr. finite (vars expr)}$ 
begin

abbreviation finite-vars :: 'expr  $\Rightarrow$  'var fset where
  finite-vars expr  $\equiv$  Abs-fset (vars expr)

lemma fset-finite-vars [simp]: fset (finite-vars expr) = vars expr
   $\langle \text{proof} \rangle$ 

end

locale renaming-variables = functional-substitution +
  assumes
    is-renaming-iff:  $\bigwedge \varrho. \text{is-renaming } \varrho \longleftrightarrow \text{inj } \varrho \wedge (\forall x. \exists x'. \varrho x = \text{id-subst } x')$ 
  and
    rename-variables:  $\bigwedge \text{expr } \varrho. \text{is-renaming } \varrho \implies \text{vars (expr} \cdot \varrho) = \text{rename } \varrho` (vars expr)$ 
  begin

lemma renaming-range-id-subst:
  assumes is-renaming  $\varrho$ 
  shows  $\varrho x \in \text{range id-subst}$ 
   $\langle \text{proof} \rangle$ 

lemma obtain-renamed-variable:
  assumes is-renaming  $\varrho$ 
  obtains  $x'$  where  $\varrho x = \text{id-subst } x'$ 
   $\langle \text{proof} \rangle$ 

lemma id-subst-rename [simp]:
  assumes is-renaming  $\varrho$ 
  shows id-subst (rename  $\varrho x$ ) =  $\varrho x$ 
   $\langle \text{proof} \rangle$ 

lemma rename-variables-id-subst:
  assumes is-renaming  $\varrho$ 
  shows id-subst ` vars (expr  $\cdot$   $\varrho$ ) =  $\varrho` (vars expr)$ 
   $\langle \text{proof} \rangle$ 

```

```

lemma surj-inv-renaming:
  assumes is-renaming  $\varrho$ 
  shows surj ( $\lambda x. \text{inv } \varrho (\text{id-subst } x)$ )
   $\langle \text{proof} \rangle$ 

lemma renaming-range:
  assumes is-renaming  $\varrho$   $x \in \text{vars}$  ( $\text{expr} \cdot \varrho$ )
  shows id-subst  $x \in \text{range } \varrho$ 
   $\langle \text{proof} \rangle$ 

lemma renaming-inv-into:
  assumes is-renaming  $\varrho$   $x \in \text{vars}$  ( $\text{expr} \cdot \varrho$ )
  shows  $\varrho(\text{inv } \varrho (\text{id-subst } x)) = \text{id-subst } x$ 
   $\langle \text{proof} \rangle$ 

lemma inv-renaming:
  assumes is-renaming  $\varrho$ 
  shows  $\text{inv } \varrho (\varrho x) = x$ 
   $\langle \text{proof} \rangle$ 

lemma renaming-inv-in-vars:
  assumes is-renaming  $\varrho$   $x \in \text{vars}$  ( $\text{expr} \cdot \varrho$ )
  shows  $\text{inv } \varrho (\text{id-subst } x) \in \text{vars}$   $\text{expr}$ 
   $\langle \text{proof} \rangle$ 

end

locale grounding = functional-substitution where vars = vars and id-subst =
id-subst
  for vars :: 'expr  $\Rightarrow$  'var set and id-subst :: 'var  $\Rightarrow$  'base +
  fixes to-ground :: 'expr  $\Rightarrow$  'exprG and from-ground :: 'exprG  $\Rightarrow$  'expr
  assumes
    range-from-ground-iff-is-ground: {expr. is-ground expr} = range from-ground
  and
    from-ground-inverse [simp]:  $\bigwedge \text{expr}_G. \text{to-ground}(\text{from-ground } \text{expr}_G) = \text{expr}_G$ 
begin

  definition ground-instances' :: 'expr  $\Rightarrow$  'exprG set where
    ground-instances' expr = { to-ground (expr  $\cdot$   $\gamma$ ) |  $\gamma$ . is-ground (expr  $\cdot$   $\gamma$ ) }

  lemma ground-instances'-eq-ground-instances:
    ground-instances' expr = (to-ground ` ground-instances expr)
     $\langle \text{proof} \rangle$ 

  lemma to-ground-from-ground-id [simp]: to-ground  $\circ$  from-ground = id
   $\langle \text{proof} \rangle$ 

  lemma surj-to-ground: surj to-ground

```

```

⟨proof⟩

lemma inj-from-ground: inj-on from-ground domainG
⟨proof⟩

lemma inj-on-to-ground: inj-on to-ground (from-ground ‘ domainG)
⟨proof⟩

lemma bij-betw-to-ground: bij-betw to-ground (from-ground ‘ domainG) domainG
⟨proof⟩

lemma bij-betw-from-ground: bij-betw from-ground domainG (from-ground ‘ do-
mainG)
⟨proof⟩

lemma ground-is-ground [simp, intro]: is-ground (from-ground exprG)
⟨proof⟩

lemma is-ground-iff-range-from-ground: is-ground expr  $\longleftrightarrow$  expr  $\in$  range from-ground
⟨proof⟩

lemma to-ground-inverse [simp]:
assumes is-ground expr
shows from-ground (to-ground expr) = expr
⟨proof⟩

corollary obtain-grounding:
assumes is-ground expr
obtains exprG where from-ground exprG = expr
⟨proof⟩

lemma from-ground-eq [simp]:
from-ground expr = from-ground expr'  $\longleftrightarrow$  expr = expr'
⟨proof⟩

lemma to-ground-eq [simp]:
assumes is-ground expr is-ground expr'
shows to-ground expr = to-ground expr'  $\longleftrightarrow$  expr = expr'
⟨proof⟩

end

locale base-functional-substitution = functional-substitution
where id-subst = id-subst and vars = vars
for id-subst :: 'var  $\Rightarrow$  'expr and vars :: 'expr  $\Rightarrow$  'var set +
assumes
vars-subst-vars:  $\bigwedge$ expr  $\varrho$ . vars (expr  $\cdot$   $\varrho$ ) =  $\bigcup$  (vars ‘  $\varrho$  ‘ vars expr) and
base-ground-exists:  $\exists$ expr. is-ground expr and
vars-id-subst:  $\bigwedge$ x. vars (id-subst x) = {x} and

```

```

comp-subst-iff:  $\bigwedge \sigma \sigma' x. (\sigma \odot \sigma') x = \sigma x \cdot \sigma'$ 

locale based-functional-substitution =
  base: base-functional-substitution where subst = base-subst and vars = base-vars
+
  functional-substitution where vars = vars
for
  base-subst :: 'base  $\Rightarrow$  ('var  $\Rightarrow$  'base)  $\Rightarrow$  'base and
  base-vars and
  vars :: 'expr  $\Rightarrow$  'var set +
assumes
  ground-subst-iff-base-ground-subst [simp]:  $\bigwedge \gamma. \text{is-ground-subst } \gamma \longleftrightarrow \text{base.is-ground-subst } \gamma$  and
  vars-subst:  $\bigwedge \text{expr } \varrho. \text{vars } (\text{expr} \cdot \varrho) = \bigcup (\text{base-vars } ' \varrho ' \text{ vars expr})$ 
begin

lemma is-grounding-iff-vars-grounded:
  is-ground (expr  $\cdot$   $\gamma$ )  $\longleftrightarrow$  ( $\forall \text{ var } \in \text{vars expr}. \text{base.is-ground } (\gamma \text{ var})$ )
  ⟨proof⟩

lemma obtain-ground-subst:
  obtains  $\gamma$ 
  where is-ground-subst  $\gamma$ 
  ⟨proof⟩

lemma exists-ground-subst [intro]:  $\exists \gamma. \text{is-ground-subst } \gamma$ 
  ⟨proof⟩

lemma ground-subst-extension:
  assumes is-ground (expr  $\cdot$   $\gamma$ )
  obtains  $\gamma'$ 
  where expr  $\cdot$   $\gamma$  = expr  $\cdot$   $\gamma'$  and is-ground-subst  $\gamma'$ 
  ⟨proof⟩

lemma ground-subst-extension':
  assumes is-ground (expr  $\cdot$   $\gamma$ )
  obtains  $\gamma'$ 
  where expr  $\cdot$   $\gamma$  = expr  $\cdot$   $\gamma'$  and base.is-ground-subst  $\gamma'$ 
  ⟨proof⟩

lemma ground-subst-update [simp]:
  assumes base.is-ground update is-ground (expr  $\cdot$   $\gamma$ )
  shows is-ground (expr  $\cdot$   $\gamma$ (var := update))
  ⟨proof⟩

lemma ground-exists:  $\exists \text{expr}. \text{is-ground } \text{expr}$ 
  ⟨proof⟩

lemma variable-grounding:

```

```

assumes is-ground (expr · γ) var ∈ vars expr
shows base.is-ground (γ var)
⟨proof⟩

definition range-vars :: ('var ⇒ 'base) ⇒ 'var set where
  range-vars σ = ⋃(base-vars ` subst-range σ)

lemma vars-subst-subset: vars (expr · σ) ⊆ (vars expr – subst-domain σ) ∪
range-vars σ
⟨proof⟩

end

locale variables-in-base-imgu = based-functional-substitution +
assumes variables-in-base-imgu:
  ⋀expr μ unifications.
    base.is-imgu μ unifications ==>
    finite unifications ==>
    ∀ unification ∈ unifications. finite unification ==>
    vars (expr · μ) ⊆ vars expr ∪ (⋃(base-vars ` ⋃ unifications))

context base-functional-substitution
begin

sublocale based-functional-substitution
  where base-subst = subst and base-vars = vars
  ⟨proof⟩

declare ground-subst-iff-base-ground-subst [simp del]

end

hide-fact base-functional-substitution.base-ground-exists
hide-fact base-functional-substitution.vars-subst-vars

end
theory Functional-Substitution-Example
  imports Functional-Substitution Substitution-First-Order-Term
begin

A selection of substitution properties for terms.

locale term-subst-properties =
  base-functional-substitution +
  finite-variables

global-interpretation term-subst: term-subst-properties where
  subst = subst-apply-term and id-subst = Var and comp-subst = subst-compose
and
  vars = vars-term :: ('f, 'v) term ⇒ 'v set

```

```

rewrites  $\bigwedge t. \text{term-subst.is-ground } t \longleftrightarrow \text{ground } t$ 

"rewrites" enables us to use our own equivalent definitions.

⟨proof⟩

Examples of generated lemmas and definitions

thm
 $\text{term-subst.subst-redundant-upd}$ 
 $\text{term-subst.subst-redundant-if}$ 

 $\text{term-subst.subst-domain-def}$ 
 $\text{term-subst.range-vars-def}$ 

 $\text{term-subst.vars-subst-subset}$ 

end

theory Natural-Magma
  imports Main
begin

locale natural-magma =
  fixes
    to-set :: ' $b \Rightarrow 'a \text{ set}$  and
    plus :: ' $b \Rightarrow 'b \Rightarrow 'b$  and
    wrap :: ' $a \Rightarrow 'b$  and
      add
  defines  $\bigwedge a b. \text{add } a b \equiv \text{plus } (\text{wrap } a) b$ 
  assumes
    to-set-plus [simp]:  $\bigwedge b b'. \text{to-set } (\text{plus } b b') = (\text{to-set } b) \cup (\text{to-set } b')$  and
    to-set-wrap [simp]:  $\bigwedge a. \text{to-set } (\text{wrap } a) = \{a\}$ 
begin

lemma to-set-add [simp]:  $\text{to-set } (\text{add } a b) = \text{insert } a (\text{to-set } b)$ 
  ⟨proof⟩

end

locale natural-magma-with-empty = natural-magma +
  fixes empty
  assumes to-set-empty [simp]:  $\text{to-set } \text{empty} = \{\}$ 

end

theory Natural-Functor
  imports Main
begin

locale natural-functor =
  fixes

```

```

map :: ('a ⇒ 'a) ⇒ 'b ⇒ 'b and
to-set :: 'b ⇒ 'a set
assumes
  map-comp [simp]: ∀b f g. map f (map g b) = map (λx. f (g x)) b and
  map-ident [simp]: ∀b. map (λx. x) b = b and
  map-cong0 [cong]: ∀b f g. (∀a. a ∈ to-set b ⇒ f a = g a) ⇒ map f b = map
g b and
  to-set-map [simp]: ∀b f. to-set (map f b) = f ` to-set b and
  exists-functor [intro]: ∀a. ∃b. a ∈ to-set b
begin

lemma map-id [simp]: map id b = b
  ⟨proof⟩

lemma map-cong [cong]:
  assumes b = b' ∀a. a ∈ to-set b' ⇒ f a = g a
  shows map f b = map g b'
  ⟨proof⟩

end

locale finite-natural-functor = natural-functor +
  assumes finite-to-set [intro]: ∀b. finite (to-set b)

locale natural-functor-conversion =
  natural-functor +
  functor': natural-functor where map = map' and to-set = to-set'
  for map' :: ('b ⇒ 'b) ⇒ 'd ⇒ 'd and to-set' :: 'd ⇒ 'b set +
  fixes
    map-to :: ('a ⇒ 'b) ⇒ 'c ⇒ 'd and
    map-from :: ('b ⇒ 'a) ⇒ 'd ⇒ 'c
  assumes
    to-set-map-from [simp]: ∀f d. to-set (map-from f d) = f ` to-set' d and
    to-set-map-to [simp]: ∀f c. to-set' (map-to f c) = f ` to-set c and
    conversion-map-comp [simp]: ∀c f g. map-from f (map-to g c) = map (λx. f (g
x)) c and
    conversion-map-comp' [simp]: ∀d f g. map-to f (map-from g d) = map' (λx. f
(g x)) d

end
theory Natural-Magma-Functor
  imports Natural-Magma Natural-Functor
begin

locale natural-magma-functor = natural-magma + natural-functor +
  assumes
    map-wrap: ∀f a. map f (wrap a) = wrap (f a) and
    map-plus: ∀f b b'. map f (plus b b') = plus (map f b) (map f b')

```

```

begin

lemma map-add:  $\bigwedge f a b. \text{map } f (\text{add } a b) = \text{add } (f a) (\text{map } f b)$ 
   $\langle \text{proof} \rangle$ 

end

end
theory Functional-Substitution-Lifting
imports Functional-Substitution Natural-Magma-Functor
begin

locale functional-substitution-lifting =
  sub: functional-substitution where subst = sub-subst and vars = sub-vars +
  natural-functor where map = map and to-set = to-set
  for
    sub-vars :: 'sub  $\Rightarrow$  'var set and
    sub-subst :: 'sub  $\Rightarrow$  ('var  $\Rightarrow$  'base)  $\Rightarrow$  'sub and
    map :: ('sub  $\Rightarrow$  'sub)  $\Rightarrow$  'expr  $\Rightarrow$  'expr and
    to-set :: 'expr  $\Rightarrow$  'sub set
begin

definition vars :: 'expr  $\Rightarrow$  'var set where
  vars expr  $\equiv$   $\bigcup$  (sub-vars ` to-set expr)

notation sub-subst (infixl  $\cdot_s$  70)

definition subst :: 'expr  $\Rightarrow$  ('var  $\Rightarrow$  'base)  $\Rightarrow$  'expr (infixl  $\cdot$  70) where
  expr  $\cdot$   $\sigma$   $\equiv$  map ( $\lambda$  sub. sub  $\cdot_s$   $\sigma$ ) expr

lemma map-id-cong [simp]:
  assumes  $\bigwedge$  sub. sub  $\in$  to-set expr  $\implies$  f sub = sub
  shows map f expr = expr
   $\langle \text{proof} \rangle$ 

lemma to-set-map-not-ident:
  assumes sub  $\in$  to-set expr f sub  $\notin$  to-set expr
  shows map f expr  $\neq$  expr
   $\langle \text{proof} \rangle$ 

lemma subst-in-to-set-subst [intro]:
  assumes sub  $\in$  to-set expr
  shows sub  $\cdot_s$   $\sigma$   $\in$  to-set (expr  $\cdot$   $\sigma$ )
   $\langle \text{proof} \rangle$ 

sublocale functional-substitution where subst = (.) and vars = vars
   $\langle \text{proof} \rangle$ 

lemma ground-subst-iff-sub-ground-subst [simp]: is-ground-subst  $\gamma$   $\longleftrightarrow$  sub.is-ground-subst

```

$\gamma$   
 $\langle proof \rangle$

**lemma** *to-set-is-ground* [intro]:

**assumes** *sub*  $\in$  *to-set expr* *is-ground expr*  
**shows** *sub.is-ground sub*  
 $\langle proof \rangle$

**lemma** *to-set-is-ground-subst*:

**assumes** *sub*  $\in$  *to-set expr* *is-ground (expr · γ)*  
**shows** *sub.is-ground (sub ·<sub>s</sub> γ)*  
 $\langle proof \rangle$

**lemma** *subst-empty*:

**assumes** *to-set expr'*  $=$   $\{\}$   
**shows** *expr · σ = expr' ↔ expr = expr'*  
 $\langle proof \rangle$

**lemma** *empty-is-ground*:

**assumes** *to-set expr*  $=$   $\{\}$   
**shows** *is-ground expr*  
 $\langle proof \rangle$

**lemma** *to-set-image*: *to-set (expr · σ) = (λa. a ·<sub>s</sub> σ) ` to-set expr*

$\langle proof \rangle$

**lemma** *to-set-subset-vars-subset*:

**assumes** *to-set expr*  $\subseteq$  *to-set expr'*  
**shows** *vars expr*  $\subseteq$  *vars expr'*  
 $\langle proof \rangle$

**lemma** *to-set-subset-is-ground*:

**assumes** *to-set expr'*  $\subseteq$  *to-set expr* *is-ground expr*  
**shows** *is-ground expr'*  
 $\langle proof \rangle$

**end**

**locale** *based-functional-substitution-lifting* =

*functional-substitution-lifting* +

*base: base-functional-substitution* **where** *subst = base-subst* **and** *vars = base-vars*

+

*sub: based-functional-substitution* **where** *subst = sub-subst* **and** *vars = sub-vars*

**begin**

**sublocale** *based-functional-substitution* **where** *subst = subst* **and** *vars = vars*  
 $\langle proof \rangle$

**end**

```

locale finite-variables-lifting =
  sub: finite-variables where vars = sub-vars :: 'sub ⇒ 'var set and subst =
  sub-subst +
    finite-natural-functor where to-set = to-set :: 'expr ⇒ 'sub set +
      functional-substitution-lifting
begin

abbreviation to-fset :: 'expr ⇒ 'sub fset where
  to-fset expr ≡ Abs-fset (to-set expr)

sublocale finite-variables where vars = vars and subst = subst
  ⟨proof⟩

lemma fset-to-fset [simp]: fset (to-fset expr) = to-set expr
  ⟨proof⟩

lemma to-fset-map: to-fset (map f expr) = f |` to-fset expr
  ⟨proof⟩

lemma to-fset-is-ground-subst:
  assumes sub |∈| to-fset expr is-ground (subst expr γ)
  shows sub.is-ground (sub ·s γ)
  ⟨proof⟩

end

locale renaming-variables-lifting =
  functional-substitution-lifting +
  sub: renaming-variables where vars = sub-vars and subst = sub-subst
begin

sublocale renaming-variables where subst = subst and vars = vars
  ⟨proof⟩

end

locale based-renaming-variables-lifting =
  renaming-variables-lifting +
  based-functional-substitution-lifting +
  base: renaming-variables where vars = base-vars and subst = base-subst

locale variables-in-base-imgu-lifting =
  based-functional-substitution-lifting +
  sub: variables-in-base-imgu where vars = sub-vars and subst = sub-subst
begin

sublocale variables-in-base-imgu where subst = subst and vars = vars
  ⟨proof⟩

```

```

end

locale grounding-lifting =
  functional-substitution-lifting where sub-vars = sub-vars and sub-subst = sub-subst
and
  map = map +
  sub: grounding where vars = sub-vars and subst = sub-subst and to-ground =
  sub-to-ground and
    from-ground = sub-from-ground +
    natural-functor-conversion where map = map and map-to = to-ground-map and
    map-from = from-ground-map and map' = ground-map and to-set' = to-set-ground
  for
    sub-to-ground :: 'sub => 'subG and
    sub-from-ground :: 'subG => 'sub and
    sub-vars :: 'sub => 'var set and
    sub-subst :: 'sub => ('var => 'base) => 'sub and
    map :: ('sub => 'sub) => 'expr => 'expr and
    to-ground-map :: ('sub => 'subG) => 'expr => 'exprG and
    from-ground-map :: ('subG => 'sub) => 'exprG => 'expr and
    ground-map :: ('subG => 'subG) => 'exprG => 'exprG and
    to-set-ground :: 'exprG => 'subG set
begin

definition to-ground :: 'expr => 'exprG where
  to-ground expr ≡ to-ground-map sub-to-ground expr

definition from-ground :: 'exprG => 'expr where
  from-ground exprG ≡ from-ground-map sub-from-ground exprG

sublocale grounding
  where vars = vars and subst = subst and to-ground = to-ground and from-ground =
  = from-ground
  ⟨proof⟩

lemma to-set-from-ground: to-set (from-ground expr) = sub-from-ground ` (to-set-ground
expr)
  ⟨proof⟩

lemma sub-in-ground-is-ground:
  assumes sub ∈ to-set (from-ground expr)
  shows sub.is-ground sub
  ⟨proof⟩

lemma ground-sub-in-ground:
  sub ∈ to-set-ground expr ↔ sub-from-ground sub ∈ to-set (from-ground expr)
  ⟨proof⟩

lemma ground-sub:

```

```


$$(\forall sub \in \text{to-set } (\text{from-ground } expr_G). P sub) \longleftrightarrow$$


$$(\forall sub_G \in \text{to-set-ground } expr_G. P (\text{sub-from-ground } sub_G))$$


$$\langle proof \rangle$$


end

locale all-subst-ident-iff-ground-lifting =
  finite-variables-lifting where map = map +
    sub: all-subst-ident-iff-ground where subst = sub-subst and vars = sub-vars
  for map :: ('sub  $\Rightarrow$  'sub)  $\Rightarrow$  'expr  $\Rightarrow$  'expr
  begin

    sublocale all-subst-ident-iff-ground where subst = subst and vars = vars
     $\langle proof \rangle$ 

  end

locale natural-magma-functional-substitution-lifting =
  functional-substitution-lifting + natural-magma
  begin

    lemma vars-add [simp]:
      vars (add sub expr) = sub-vars sub  $\cup$  vars expr
       $\langle proof \rangle$ 

    lemma vars-plus [simp]:
      vars (plus expr expr') = vars expr  $\cup$  vars expr'
       $\langle proof \rangle$ 

    lemma is-ground-add [simp]:
      is-ground (add sub expr)  $\longleftrightarrow$  sub.is-ground sub  $\wedge$  is-ground expr
       $\langle proof \rangle$ 

  end

locale natural-magma-functor-functional-substitution-lifting =
  natural-magma-functional-substitution-lifting + natural-magma-functor
  begin

    lemma add-subst [simp]:
      (add sub expr)  $\cdot$   $\sigma$  = add (sub  $\cdot_s$   $\sigma$ ) (expr  $\cdot$   $\sigma$ )
       $\langle proof \rangle$ 

    lemma plus-subst [simp]: (plus expr expr')  $\cdot$   $\sigma$  = plus (expr  $\cdot$   $\sigma$ ) (expr'  $\cdot$   $\sigma$ )
       $\langle proof \rangle$ 

  end

locale natural-magma-grounding-lifting =

```

```

grounding-lifting +
natural-magma-functional-substitution-lifting +
ground: natural-magma where
  to-set = to-set-ground and plus = plus-ground and wrap = wrap-ground and
  add = add-ground
for plus-ground wrap-ground add-ground +
assumes
  to-ground-plus [simp]:
     $\wedge \text{expr expr'. to-ground (plus expr expr')} = \text{plus-ground} (\text{to-ground expr}) (\text{to-ground expr'})$  and
    wrap-from-ground:  $\wedge \text{sub. from-ground (wrap-ground sub)} = \text{wrap} (\text{sub-from-ground sub})$  and
    wrap-to-ground:  $\wedge \text{sub. to-ground (wrap sub)} = \text{wrap-ground} (\text{sub-to-ground sub})$ 
begin

lemma from-ground-plus [simp]:
  from-ground (plus-ground expr expr') = plus (from-ground expr) (from-ground expr')
  ⟨proof⟩

lemma from-ground-add [simp]:
  from-ground (add-ground sub expr) = add (sub-from-ground sub) (from-ground expr)
  ⟨proof⟩

lemma to-ground-add [simp]:
  to-ground (add sub expr) = add-ground (sub-to-ground sub) (to-ground expr)
  ⟨proof⟩

lemma ground-add:
  assumes from-ground expr = add sub expr'
  shows expr = add-ground (sub-to-ground sub) (to-ground expr')
  ⟨proof⟩

end

locale natural-magma-with-empty-grounding-lifting =
natural-magma-grounding-lifting +
natural-magma-with-empty +
ground: natural-magma-with-empty where
  to-set = to-set-ground and plus = plus-ground and wrap = wrap-ground and
  add = add-ground and
  empty = empty-ground
for empty-ground +
assumes to-ground-empty [simp]: to-ground empty = empty-ground
begin

lemmas empty-magma-is-ground [simp] = empty-is-ground[OF to-set-empty]

```

```

lemmas magma-subst-empty [simp] =
  subst-ident-if-ground[OF empty-magma-is-ground]
  subst-empty[OF to-set-empty]

lemma from-ground-empty [simp]: from-ground empty-ground = empty
  <proof>

lemma to-ground-empty' [simp]: to-ground expr = empty-ground  $\longleftrightarrow$  expr = empty
  <proof>

lemma from-ground-empty' [simp]: from-ground expr = empty  $\longleftrightarrow$  expr = empty-ground
  <proof>

end

end
theory Functional-Substitution-Lifting-Example
imports
  Functional-Substitution-Lifting
  Functional-Substitution-Example
begin

Lifting of properties from term to equations (modelled as pairs)
type-synonym ('f,'v) equation = ('f, 'v) term  $\times$  ('f, 'v) term

All property locales have corresponding lifting locales

locale lifting-term-subst-properties =
  based-functional-substitution-lifting where
  id-subst = Var and comp-subst = subst-compose and base-subst = subst-apply-term
  and
  base-vars = vars-term :: ('f, 'v) term  $\Rightarrow$  'v set +
  finite-variables-lifting where id-subst = Var and comp-subst = subst-compose

global-interpretation equation-subst:
  lifting-term-subst-properties where
  sub-vars = vars-term and sub-subst = subst-apply-term and map =  $\lambda f. map\text{-prod}$ 
  f f and
  to-set = set-prod
  <proof>

Lifted lemmas and definitions

thm
  equation-subst.subst-redundant-upd
  equation-subst.subst-redundant-if
  equation-subst.vars-subst-subset

  equation-subst.vars-def
  equation-subst.subst-def

```

We can lift multiple levels

```
global-interpretation equation-set-subst:  
    lifting-term-subst-properties where  
        sub-vars = equation-subst.vars and sub-subst = equation-subst.subst and map =  
        fimage and  
        to-set = fset  
        ⟨proof⟩
```

Lifted lemmas and definitions

```
thm  
    equation-set-subst.subst-redundant-upd  
    equation-set-subst.subst-redundant-if  
    equation-set-subst.vars-subst-subset  
  
    equation-set-subst.vars-def  
    equation-set-subst.subst-def  
  
end
```