

# Mechanization of the Algebra for Wireless Networks (AWN)

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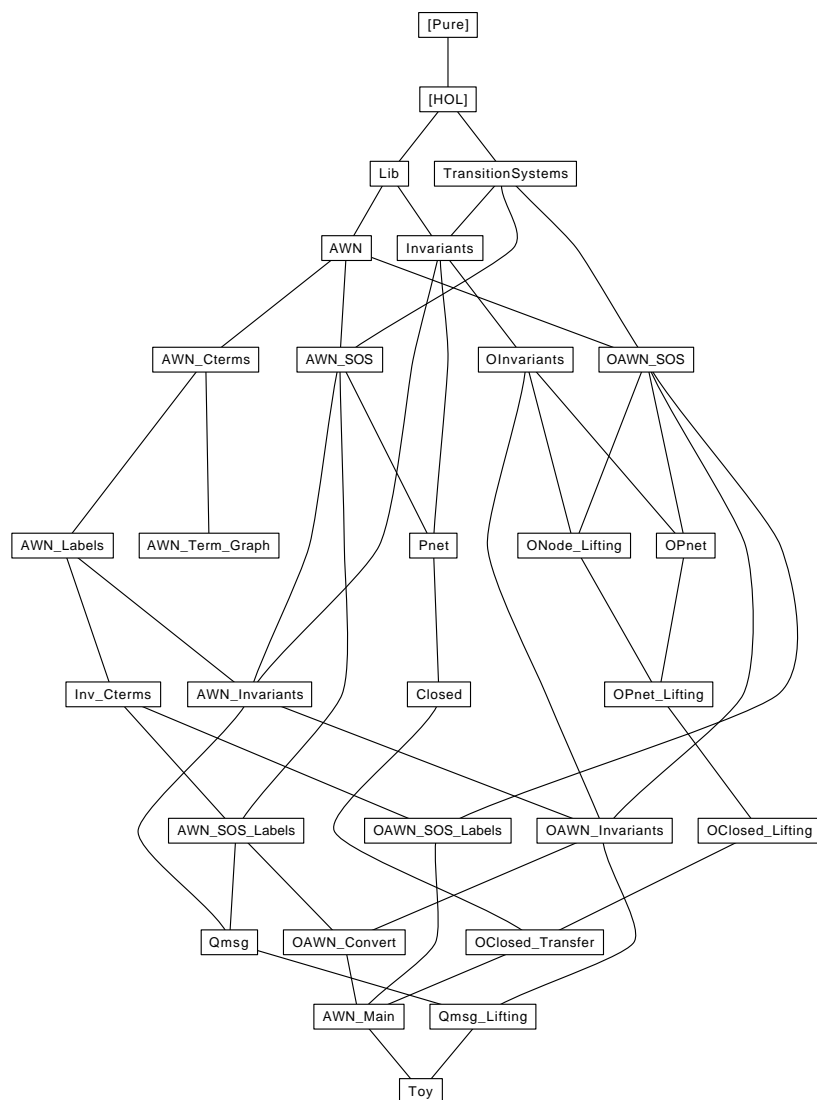
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## Abstract

AWN is a process algebra developed for modelling and analysing protocols for Mobile Ad hoc Networks (MANETs) and Wireless Mesh Networks (WMNs) [2, §4]. AWN models comprise five distinct layers: sequential processes, local parallel compositions, nodes, partial networks, and complete networks.

This development mechanises the original operational semantics of AWN and introduces a variant ‘open’ operational semantics that enables the compositional statement and proof of invariants across distinct network nodes. It supports labels (for weakening invariants) and (abstract) data state manipulations. A framework for compositional invariant proofs is developed, including a tactic (`inv_cterms`) for inductive invariant proofs of sequential processes, lifting rules for the open versions of the higher layers, and a rule for transferring lifted properties back to the standard semantics. A notion of ‘control terms’ reduces proof obligations to the subset of subterms that act directly (in contrast to operators for combining terms and joining processes).

Further documentation is available in [1].



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# 1 Generic functions and lemmas

```
theory Lib
imports Main
begin

definition
  TT :: "'a ⇒ bool"
where
  "TT = (λ_. True)"

lemma TT_True [intro, simp]: "TT a"
  unfolding TT_def by simp

lemma in_set_tl: "x ∈ set (tl xs) ⇒ x ∈ set xs"
  by (metis Nil_tl insert_iff list.collapse set_simps(2))

lemma nat_le_eq_or_lt [elim]:
  fixes x :: nat
  assumes "x ≤ y"
  and eq: "x = y ⇒ P x y"
  and lt: "x < y ⇒ P x y"
  shows "P x y"
  using assms unfolding nat_less_le by auto

lemma disjoint_commute:
  "(A ∩ B = {}) ⇒ (B ∩ A = {})"
  by auto

definition
  default :: "('i ⇒ 's) ⇒ ('i ⇒ 's option) ⇒ ('i ⇒ 's)"
where
  "default df f = (λi. case f i of None ⇒ df i | Some s ⇒ s)"

end
```

# 2 Transition systems (automata)

```
theory TransitionSystems
imports Main
begin

type_synonym ('s, 'a) transition = "'s × 'a × 's"

record ('s, 'a) automaton =
  init :: "'s set"
  trans :: "('s, 'a) transition set"

end
```

# 3 Reachability and Invariance

```
theory Invariants
imports Lib TransitionSystems
begin
```

## 3.1 Reachability

A state is ‘reachable’ under  $I$  if either it is the initial state, or it is the destination of a transition whose action satisfies  $I$  from a reachable state. The ‘standard’ definition of reachability is recovered by setting  $I$  to  $TT$ .

```
inductive_set reachable
  for A :: "('s, 'a) automaton"
```

```

and I :: "'a ⇒ bool"
where
  reachable_init: "s ∈ init A ⇒ s ∈ reachable A I"
  / reachable_step: "[[ s ∈ reachable A I; (s, a, s') ∈ trans A; I a ]] ⇒ s' ∈ reachable A I"

inductive_cases reachable_icas: "s ∈ reachable A I"

lemma reachable_pair_induct [consumes, case_names init step]:
  assumes "(ξ, p) ∈ reachable A I"
    and "∧ξ p. (ξ, p) ∈ init A ⇒ P ξ p"
    and "(∧ξ p ξ' p' a. [[ (ξ, p) ∈ reachable A I; P ξ p;
      ((ξ, p), a, (ξ', p')) ∈ trans A; I a ]] ⇒ P ξ' p'"
  shows "P ξ p"
using assms(1) proof (induction "(ξ, p)" arbitrary: ξ p)
  fix ξ p
  assume "(ξ, p) ∈ init A"
  with assms(2) show "P ξ p" .
next
  fix s a ξ' p'
  assume "s ∈ reachable A I"
    and tr: "(s, a, (ξ', p')) ∈ trans A"
    and "I a"
    and IH: "∧ξ p. s = (ξ, p) ⇒ P ξ p"
  from this(1) obtain ξ p where "s = (ξ, p)"
    and "(ξ, p) ∈ reachable A I"
  by (metis prod.collapse)
  note this(2)
  moreover from IH and ⟨s = (ξ, p)⟩ have "P ξ p" .
  moreover from tr and ⟨s = (ξ, p)⟩ have "((ξ, p), a, (ξ', p')) ∈ trans A" by simp
  ultimately show "P ξ' p'"
    using ⟨I a⟩ by (rule assms(3))
qed

lemma reachable_weakenE [elim]:
  assumes "s ∈ reachable A P"
    and PQ: "∧a. P a ⇒ Q a"
  shows "s ∈ reachable A Q"
using assms(1)
proof (induction)
  fix s assume "s ∈ init A"
  thus "s ∈ reachable A Q" ..
next
  fix s a s'
  assume "s ∈ reachable A P"
    and "s ∈ reachable A Q"
    and "(s, a, s') ∈ trans A"
    and "P a"
  from ⟨P a⟩ have "Q a" by (rule PQ)
  with ⟨s ∈ reachable A Q⟩ and ⟨(s, a, s') ∈ trans A⟩ show "s' ∈ reachable A Q" ..
qed

lemma reachable_weaken_TT [elim]:
  assumes "s ∈ reachable A I"
  shows "s ∈ reachable A TT"
using assms by rule simp

lemma init_empty_reachable_empty:
  assumes "init A = {}"
  shows "reachable A I = {}"
proof (rule ccontr)
  assume "reachable A I ≠ {}"
  then obtain s where "s ∈ reachable A I" by auto
  thus False
proof (induction rule: reachable.induct)

```

```

fix s
assume "s ∈ init A"
with ⟨init A = {⟩ show False by simp
qed
qed

```

## 3.2 Invariance

definition invariant

```

:: "('s, 'a) automaton ⇒ ('a ⇒ bool) ⇒ ('s ⇒ bool) ⇒ bool"
("_ ⊨ (1'(_ →')/ _)" [100, 0, 9] 8)

```

where

```

"(A ⊨ (I →) P) = (∀s∈reachable A I. P s)"

```

abbreviation

```

any_invariant
:: "('s, 'a) automaton ⇒ ('s ⇒ bool) ⇒ bool"
("_ ⊨ _" [100, 9] 8)

```

where

```

"(A ⊨ P) ≡ (A ⊨ (TT →) P)"

```

lemma invariantI [intro]:

```

assumes init: "∧s. s ∈ init A ⇒ P s"
and step: "∧s a s'. [ s ∈ reachable A I; P s; (s, a, s') ∈ trans A; I a ] ⇒ P s'"
shows "A ⊨ (I →) P"

```

unfolding invariant\_def

proof

```

fix s
assume "s ∈ reachable A I"
thus "P s"
proof induction
fix s assume "s ∈ init A"
thus "P s" by (rule init)
next
fix s a s'
assume "s ∈ reachable A I"
and "P s"
and "(s, a, s') ∈ trans A"
and "I a"
thus "P s'" by (rule step)
qed
qed

```

lemma invariant\_pairI [intro]:

```

assumes init: "∧ξ p. (ξ, p) ∈ init A ⇒ P (ξ, p)"
and step: "∧ξ p ξ' p' a.
[ (ξ, p) ∈ reachable A I; P (ξ, p); ((ξ, p), a, (ξ', p')) ∈ trans A; I a ]
⇒ P (ξ', p'"
shows "A ⊨ (I →) P"
using assms by auto

```

lemma invariant\_arbitraryI:

```

assumes "∧s. s ∈ reachable A I ⇒ P s"
shows "A ⊨ (I →) P"
using assms unfolding invariant_def by simp

```

lemma invariantD [dest]:

```

assumes "A ⊨ (I →) P"
and "s ∈ reachable A I"
shows "P s"
using assms unfolding invariant_def by blast

```

lemma invariant\_initE [elim]:

```

assumes invP: "A ⊨ (I →) P"

```

```

    and init: "s ∈ init A"
    shows "P s"
  proof -
    from init have "s ∈ reachable A I" ..
    with invP show ?thesis ..
  qed

lemma invariant_weakenE [elim]:
  fixes T σ P Q
  assumes invP: "A ⊨ (PI →) P"
    and PQ: "∧s. P s ⇒ Q s"
    and QIPI: "∧a. QI a ⇒ PI a"
  shows "A ⊨ (QI →) Q"
  proof
    fix s
    assume "s ∈ init A"
    with invP have "P s" ..
    thus "Q s" by (rule PQ)
  next
    fix s a s'
    assume "s ∈ reachable A QI"
      and "(s, a, s') ∈ trans A"
      and "QI a"
    from ⟨QI a⟩ have "PI a" by (rule QIPI)
    from ⟨s ∈ reachable A QI⟩ and QIPI have "s ∈ reachable A PI" ..
    hence "s' ∈ reachable A PI" using ⟨(s, a, s') ∈ trans A⟩ and ⟨PI a⟩ ..
    with invP have "P s'" ..
    thus "Q s'" by (rule PQ)
  qed

definition
  step_invariant
  :: "('s, 'a) automaton ⇒ ('a ⇒ bool) ⇒ (( 's, 'a) transition ⇒ bool) ⇒ bool"
  ("_ ⊨A (1'(_ →'))/_" [100, 0, 0] 8)
where
  "(A ⊨A (I →) P) = (∀a. I a → (∀s∈reachable A I. (∀s'.(s, a, s') ∈ trans A → P (s, a, s'))))"

lemma invariant_restrict_inD [dest]:
  assumes "A ⊨ (TT →) P"
  shows "A ⊨ (QI →) P"
  using assms by auto

abbreviation
  any_step_invariant
  :: "('s, 'a) automaton ⇒ (( 's, 'a) transition ⇒ bool) ⇒ bool"
  ("_ ⊨A _" [100, 9] 8)
where
  "(A ⊨A P) ≡ (A ⊨A (TT →) P)"

lemma step_invariant_true:
  "p ⊨A (λ(s, a, s'). True)"
  unfolding step_invariant_def by simp

lemma step_invariantI [intro]:
  assumes *: "∧s a s'. [ s∈reachable A I; (s, a, s')∈trans A; I a ] ⇒ P (s, a, s')"
  shows "A ⊨A (I →) P"
  unfolding step_invariant_def
  using assms by auto

lemma step_invariantD [dest]:
  assumes "A ⊨A (I →) P"
    and "s∈reachable A I"
    and "(s, a, s') ∈ trans A"
    and "I a"

```

```

shows "P (s, a, s')"
using assms unfolding step_invariant_def by blast

lemma step_invariantE [elim]:
  fixes T σ P I s a s'
  assumes "A  $\models_A$  (I  $\rightarrow$ ) P"
    and "s ∈ reachable A I"
    and "(s, a, s') ∈ trans A"
    and "I a"
    and "P (s, a, s')  $\implies$  Q"
  shows "Q"
using assms by auto

lemma step_invariant_pairI [intro]:
  assumes *: " $\bigwedge \xi p \xi' p' a.$ 
     $\llbracket (\xi, p) \in \text{reachable } A \ I; ((\xi, p), a, (\xi', p')) \in \text{trans } A; I \ a \rrbracket$ 
     $\implies P ((\xi, p), a, (\xi', p'))$ "
  shows "A  $\models_A$  (I  $\rightarrow$ ) P"
using assms by auto

lemma step_invariant_arbitraryI:
  assumes " $\bigwedge \xi p a \xi' p'. \llbracket (\xi, p) \in \text{reachable } A \ I; ((\xi, p), a, (\xi', p')) \in \text{trans } A; I \ a \rrbracket$ 
     $\implies P ((\xi, p), a, (\xi', p'))$ "
  shows "A  $\models_A$  (I  $\rightarrow$ ) P"
using assms by auto

lemma step_invariant_weakenE [elim!]:
  fixes T σ P Q
  assumes invP: "A  $\models_A$  (PI  $\rightarrow$ ) P"
    and PQ: " $\bigwedge t. P \ t \implies Q \ t$ "
    and QIPI: " $\bigwedge a. QI \ a \implies PI \ a$ "
  shows "A  $\models_A$  (QI  $\rightarrow$ ) Q"
proof
  fix s a s'
  assume "s ∈ reachable A QI"
    and "(s, a, s') ∈ trans A"
    and "QI a"
  from ⟨QI a⟩ have "PI a" by (rule QIPI)
  from ⟨s ∈ reachable A QI⟩ have "s ∈ reachable A PI" using QIPI ..
  with invP have "P (s, a, s')" using ⟨(s, a, s') ∈ trans A⟩ ⟨PI a⟩ ..
  thus "Q (s, a, s')" by (rule PQ)
qed

lemma step_invariant_weaken_with_invariantE [elim]:
  assumes pinv: "A  $\models$  (I  $\rightarrow$ ) P"
    and qinv: "A  $\models_A$  (I  $\rightarrow$ ) Q"
    and wr: " $\bigwedge s \ a \ s'. \llbracket P \ s; P \ s'; Q \ (s, a, s'); I \ a \rrbracket \implies R \ (s, a, s')$ "
  shows "A  $\models_A$  (I  $\rightarrow$ ) R"
proof
  fix s a s'
  assume sr: "s ∈ reachable A I"
    and tr: "(s, a, s') ∈ trans A"
    and "I a"
  hence "s' ∈ reachable A I" ..
  with pinv have "P s'" ..
  from pinv and sr have "P s" ..
  from qinv sr tr ⟨I a⟩ have "Q (s, a, s'" ..
  with ⟨P s⟩ and ⟨P s'⟩ show "R (s, a, s'" using ⟨I a⟩ by (rule wr)
qed

lemma step_to_invariantI:
  assumes sinv: "A  $\models_A$  (I  $\rightarrow$ ) Q"
    and init: " $\bigwedge s. s \in \text{init } A \implies P \ s$ "
    and step: " $\bigwedge s \ s' \ a.$ "

```



```

      [[ s ∈ reachable A I;
        P s;
        Q (s, a, s');
        I a ]] ⇒ P s'"
  shows "A ⊨ (I →) P"
proof
  fix s assume "s ∈ init A" thus "P s" by (rule init)
next
  fix s s' a
  assume "s ∈ reachable A I"
    and "P s"
    and "(s, a, s') ∈ trans A"
    and "I a"
  show "P s'"
proof -
  from sinv and ⟨s∈reachable A I⟩ and ⟨(s, a, s')∈trans A⟩ and ⟨I a⟩ have "Q (s, a, s')" ..
with ⟨s∈reachable A I⟩ and ⟨P s⟩ show "P s'" using ⟨I a⟩ by (rule step)
qed
qed
end

```

## 4 Open reachability and invariance

```

theory OInvariants
imports Invariants
begin

```

### 4.1 Open reachability

By convention, the states of an open automaton are pairs. The first component is considered to be the global state and the second is the local state.

A state is ‘open reachable’ under  $S$  and  $U$  if it is the initial state, or it is the destination of a transition—where the global components satisfy  $S$ —from an open reachable state, or it is the destination of an interleaved environment step where the global components satisfy  $U$ .

```

inductive_set oreachable
  :: "('g × 'l, 'a) automaton
     ⇒ ('g ⇒ 'g ⇒ 'a ⇒ bool)
     ⇒ ('g ⇒ 'g ⇒ bool)
     ⇒ ('g × 'l) set"
for A :: "('g × 'l, 'a) automaton"
and S :: "'g ⇒ 'g ⇒ 'a ⇒ bool"
and U :: "'g ⇒ 'g ⇒ bool"
where
  oreachable_init: "s ∈ init A ⇒ s ∈ oreachable A S U"
| oreachable_local: "[[ s ∈ oreachable A S U; (s, a, s') ∈ trans A; S (fst s) (fst s') a ]
                    ⇒ s' ∈ oreachable A S U"
| oreachable_other: "[[ s ∈ oreachable A S U; U (fst s) σ' ]
                    ⇒ (σ', snd s) ∈ oreachable A S U"

```

```

lemma oreachable_local' [elim]:
  assumes "(σ, p) ∈ oreachable A S U"
    and "((σ, p), a, (σ', p')) ∈ trans A"
    and "S σ σ' a"
  shows "(σ', p') ∈ oreachable A S U"
using assms by (metis fst_conv oreachable.oreachable_local)

```

```

lemma oreachable_other' [elim]:
  assumes "(σ, p) ∈ oreachable A S U"
    and "U σ σ'"
  shows "(σ', p) ∈ oreachable A S U"
proof -

```

```

from ⟨U σ σ'⟩ have "U (fst (σ, p)) σ'" by simp
with ⟨(σ, p) ∈ oreachable A S U⟩ have "(σ', snd (σ, p)) ∈ oreachable A S U"
  by (rule oreachable_other)
thus "(σ', p) ∈ oreachable A S U" by simp
qed

```

lemma oreachable\_pair\_induct [consumes, case\_names init other local]:

```

assumes "(σ, p) ∈ oreachable A S U"
  and "∧σ p. (σ, p) ∈ init A ⇒ P σ p"
  and "(∧σ p σ'. [ (σ, p) ∈ oreachable A S U; P σ p; U σ σ' ] ⇒ P σ' p)"
  and "(∧σ p σ' p' a. [ (σ, p) ∈ oreachable A S U; P σ p;
    ((σ, p), a, (σ', p')) ∈ trans A; S σ σ' a ] ⇒ P σ' p)"

```

```

shows "P σ p"
using assms (1) proof (induction "(σ, p)" arbitrary: σ p)
  fix σ p
  assume "(σ, p) ∈ init A"
  with assms(2) show "P σ p" .
next
  fix s σ'
  assume "s ∈ oreachable A S U"
  and "U (fst s) σ'"
  and IH: "∧σ p. s = (σ, p) ⇒ P σ p"
  from this(1) obtain σ p where "s = (σ, p)"
  and "(σ, p) ∈ oreachable A S U"
  by (metis surjective_pairing)
  note this(2)
  moreover from IH and ⟨s = (σ, p)⟩ have "P σ p" .
  moreover from ⟨U (fst s) σ'⟩ and ⟨s = (σ, p)⟩ have "U σ σ'" by simp
  ultimately have "P σ' p" by (rule assms(3))
  with ⟨s = (σ, p)⟩ show "P σ' (snd s)" by simp
next
  fix s a σ' p'
  assume "s ∈ oreachable A S U"
  and tr: "(s, a, (σ', p')) ∈ trans A"
  and "S (fst s) (fst (σ', p')) a"
  and IH: "∧σ p. s = (σ, p) ⇒ P σ p"
  from this(1) obtain σ p where "s = (σ, p)"
  and "(σ, p) ∈ oreachable A S U"
  by (metis surjective_pairing)
  note this(2)
  moreover from IH ⟨s = (σ, p)⟩ have "P σ p" .
  moreover from tr and ⟨s = (σ, p)⟩ have "((σ, p), a, (σ', p')) ∈ trans A" by simp
  moreover from ⟨S (fst s) (fst (σ', p')) a⟩ and ⟨s = (σ, p)⟩ have "S σ σ' a" by simp
  ultimately show "P σ' p'" by (rule assms(4))
qed

```

lemma oreachable\_weakenE [elim]:

```

assumes "s ∈ oreachable A PS PU"
  and PSQS: "∧s s' a. PS s s' a ⇒ QS s s' a"
  and PUQU: "∧s s'. PU s s' ⇒ QU s s'"
shows "s ∈ oreachable A QS QU"
using assms(1)
proof (induction)
  fix s assume "s ∈ init A"
  thus "s ∈ oreachable A QS QU" ..
next
  fix s a s'
  assume "s ∈ oreachable A QS QU"
  and "(s, a, s') ∈ trans A"
  and "PS (fst s) (fst s') a"
  from ⟨PS (fst s) (fst s') a⟩ have "QS (fst s) (fst s') a" by (rule PSQS)
  with ⟨s ∈ oreachable A QS QU⟩ and ⟨(s, a, s') ∈ trans A⟩ show "s' ∈ oreachable A QS QU" ..
next
  fix s g'

```

```

assume "s ∈ oreachable A QS QU"
and "PU (fst s) g'"
from ⟨PU (fst s) g'⟩ have "QU (fst s) g'" by (rule PUQU)
with ⟨s ∈ oreachable A QS QU⟩ show "(g', snd s) ∈ oreachable A QS QU" ..
qed

```

**definition**

```
act :: "('a ⇒ bool) ⇒ 's ⇒ 's ⇒ 'a ⇒ bool"
```

where

```
"act I ≡ (λ_ _. I)"
```

```
lemma act_simp [iff]: "act I s s' a = I a"
unfolding act_def ..
```

```
lemma reachable_in_oreachable [elim]:
```

```

fixes s
assumes "s ∈ reachable A I"
shows "s ∈ oreachable A (act I) U"
unfolding act_def using assms proof induction
fix s
assume "s ∈ init A"
thus "s ∈ oreachable A (λ_ _. I) U" ..

```

next

```

fix s a s'
assume "s ∈ oreachable A (λ_ _. I) U"
and "(s, a, s') ∈ trans A"
and "I a"
thus "s' ∈ oreachable A (λ_ _. I) U"
by (rule oreachable_local)

```

qed

## 4.2 Open Invariance

**definition** oinvariant

```

:: "('g × 'l, 'a) automaton
⇒ ('g ⇒ 'g ⇒ 'a ⇒ bool) ⇒ ('g ⇒ 'g ⇒ bool)
⇒ (('g × 'l) ⇒ bool) ⇒ bool"
("_ ⊨ (1'((1_)/ (1_ →'))/ _)" [100, 0, 0, 9] 8)

```

where

```
"(A ⊨ (S, U →) P) = (∀s∈oreachable A S U. P s)"
```

```
lemma oinvariantI [intro]:
```

```

fixes T TI S U P
assumes init: "∧s. s ∈ init A ⇒ P s"
and other: "∧g g' l.
[[ (g, l) ∈ oreachable A S U; P (g, l); U g g' ]] ⇒ P (g', l)"
and local: "∧s a s'.
[[ s ∈ oreachable A S U; P s; (s, a, s') ∈ trans A; S (fst s) (fst s') a ]] ⇒ P s'"

```

```
shows "A ⊨ (S, U →) P"
```

```
unfolding oinvariant_def
```

**proof**

```

fix s
assume "s ∈ oreachable A S U"
thus "P s"
proof induction
fix s assume "s ∈ init A"
thus "P s" by (rule init)
next
fix s a s'
assume "s ∈ oreachable A S U"
and "P s"
and "(s, a, s') ∈ trans A"
and "S (fst s) (fst s') a"
thus "P s'" by (rule local)

```

```

next
  fix s g'
  assume "s ∈ oreachable A S U"
    and "P s"
    and "U (fst s) g'"
  thus "P (g', snd s)"
    by - (rule other [where g="fst s"], simp_all)
qed
qed

lemma oinvariant_oreachableI:
  assumes "∧σ s. (σ, s) ∈ oreachable A S U ⇒ P (σ, s)"
  shows "A ⊨ (S, U →) P"
  using assms unfolding oinvariant_def by auto

lemma oinvariant_pairI [intro]:
  assumes init: "∧σ p. (σ, p) ∈ init A ⇒ P (σ, p)"
    and local: "∧σ p σ' p' a.
      [ (σ, p) ∈ oreachable A S U; P (σ, p); ((σ, p), a, (σ', p')) ∈ trans A;
        S σ σ' a ] ⇒ P (σ', p)"
    and other: "∧σ σ' p.
      [ (σ, p) ∈ oreachable A S U; P (σ, p); U σ σ' ] ⇒ P (σ', p)"
  shows "A ⊨ (S, U →) P"
  by (rule oinvariantI)
  (clarsimp | erule init | erule(3) local | erule(2) other)+

lemma oinvariantD [dest]:
  assumes "A ⊨ (S, U →) P"
    and "s ∈ oreachable A S U"
  shows "P s"
  using assms unfolding oinvariant_def
  by clarsimp

lemma oinvariant_initD [dest, elim]:
  assumes invP: "A ⊨ (S, U →) P"
    and init: "s ∈ init A"
  shows "P s"
  proof -
    from init have "s ∈ oreachable A S U" ..
    with invP show ?thesis ..
  qed

lemma oinvariant_weakenE [elim!]:
  assumes invP: "A ⊨ (PS, PU →) P"
    and PQ: "∧s. P s ⇒ Q s"
    and QSPS: "∧s s' a. QS s s' a ⇒ PS s s' a"
    and QUPU: "∧s s'. QU s s' ⇒ PU s s'"
  shows "A ⊨ (QS, QU →) Q"
  proof
    fix s
    assume "s ∈ init A"
    with invP have "P s" ..
    thus "Q s" by (rule PQ)
  next
    fix σ p σ' p' a
    assume "(σ, p) ∈ oreachable A QS QU"
      and "((σ, p), a, (σ', p')) ∈ trans A"
      and "QS σ σ' a"
    from this(3) have "PS σ σ' a" by (rule QSPS)
    from ((σ, p) ∈ oreachable A QS QU) and QSPS QUPU have "(σ, p) ∈ oreachable A PS PU" ..
    hence "(σ', p') ∈ oreachable A PS PU" using ((σ, p), a, (σ', p')) ∈ trans A and (PS σ σ' a) ..
    with invP have "P (σ', p')" ..
    thus "Q (σ', p')" by (rule PQ)
  next

```

```

fix  $\sigma \sigma' p$ 
assume " $(\sigma, p) \in \text{oreachable } A \text{ QS } QU$ "
  and " $Q (\sigma, p)$ "
  and " $QU \sigma \sigma'$ "
from  $\langle QU \sigma \sigma' \rangle$  have " $PU \sigma \sigma'$ " by (rule QUPU)
from  $\langle (\sigma, p) \in \text{oreachable } A \text{ QS } QU \rangle$  and QSPS QUPU have " $(\sigma, p) \in \text{oreachable } A \text{ PS } PU$ " ..
hence " $(\sigma', p) \in \text{oreachable } A \text{ PS } PU$ " using  $\langle PU \sigma \sigma' \rangle$  ..
with invP have " $P (\sigma', p)$ " ..
thus " $Q (\sigma', p)$ " by (rule PQ)
qed

```

lemma oinvariant\_weakenD [dest]:

```

assumes " $A \models (S', U' \rightarrow) P$ "
  and " $(\sigma, p) \in \text{oreachable } A \text{ S } U$ "
  and weakenS: " $\bigwedge s s' a. S s s' a \implies S' s s' a$ "
  and weakenU: " $\bigwedge s s'. U s s' \implies U' s s'$ "
shows " $P (\sigma, p)$ "
proof -
  from  $\langle (\sigma, p) \in \text{oreachable } A \text{ S } U \rangle$  have " $(\sigma, p) \in \text{oreachable } A \text{ S}' U''$ "
  by (rule oreachable_weakenE)
  (erule weakenS, erule weakenU)
  with  $\langle A \models (S', U' \rightarrow) P \rangle$  show " $P (\sigma, p)$ " ..
qed

```

lemma close\_open\_invariant:

```

assumes oinv: " $A \models (\text{act } I, U \rightarrow) P$ "
shows " $A \models (I \rightarrow) P$ "
proof
  fix  $s$ 
  assume " $s \in \text{init } A$ "
  with oinv show " $P s$ " ..
next
  fix  $\xi p \xi' p' a$ 
  assume sr: " $(\xi, p) \in \text{reachable } A \text{ I}$ "
  and step: " $((\xi, p), a, (\xi', p')) \in \text{trans } A$ "
  and " $I a$ "
  hence " $(\xi', p') \in \text{reachable } A \text{ I}$ " ..
  hence " $(\xi', p') \in \text{oreachable } A (\text{act } I) U$ " ..
  with oinv show " $P (\xi', p')$ " ..
qed

```

definition local\_steps :: " $((('i \Rightarrow 's1) \times 'l1) \times 'a \times ('i \Rightarrow 's2) \times 'l2) \text{ set} \Rightarrow 'i \text{ set} \Rightarrow \text{bool}$ "

where "local\_steps T J  $\equiv$

$$(\forall \sigma \zeta s a \sigma' s'. ((\sigma, s), a, (\sigma', s')) \in T \wedge (\forall j \in J. \zeta j = \sigma j) \rightarrow (\exists \zeta'. (\forall j \in J. \zeta' j = \sigma' j) \wedge ((\zeta, s), a, (\zeta', s')) \in T))"$$

lemma local\_stepsI [intro!]:

```

assumes " $\bigwedge \sigma \zeta s a \sigma' \zeta' s'. [ ((\sigma, s), a, (\sigma', s')) \in T; \forall j \in J. \zeta j = \sigma j ] \implies (\exists \zeta'. (\forall j \in J. \zeta' j = \sigma' j) \wedge ((\zeta, s), a, (\zeta', s')) \in T)$ "
shows "local_steps T J"
unfolding local_steps_def using assms by clarsimp

```

lemma local\_stepsE [elim, dest]:

```

assumes "local_steps T J"
  and " $((\sigma, s), a, (\sigma', s')) \in T$ "
  and " $\forall j \in J. \zeta j = \sigma j$ "
shows " $\exists \zeta'. (\forall j \in J. \zeta' j = \sigma' j) \wedge ((\zeta, s), a, (\zeta', s')) \in T$ "
using assms unfolding local_steps_def by blast

```

definition other\_steps :: " $((('i \Rightarrow 's) \Rightarrow ('i \Rightarrow 's) \Rightarrow \text{bool}) \Rightarrow 'i \text{ set} \Rightarrow \text{bool}$ "

where "other\_steps U J  $\equiv \forall \sigma \sigma'. U \sigma \sigma' \rightarrow (\forall j \in J. \sigma' j = \sigma j)$ "

lemma other\_stepsI [intro!]:

```

assumes " $\bigwedge \sigma \sigma' j. [ U \sigma \sigma'; j \in J ] \implies \sigma' j = \sigma j$ "

```

```

shows "other_steps U J"
using assms unfolding other_steps_def by simp

lemma other_stepsE [elim]:
  assumes "other_steps U J"
    and "U  $\sigma$   $\sigma'$ "
  shows " $\forall j \in J. \sigma' j = \sigma j$ "
using assms unfolding other_steps_def by simp

definition subreachable
where "subreachable A U J  $\equiv \forall I. \forall s \in \text{oreachable } A (\lambda s s'. I) U.
      (\exists \sigma. (\forall j \in J. \sigma j = (\text{fst } s) j) \wedge (\sigma, \text{snd } s) \in \text{reachable } A I)"$ "

lemma subreachableI [intro]:
  assumes "local_steps (trans A) J"
    and "other_steps U J"
  shows "subreachable A U J"
unfolding subreachable_def
proof (rule, rule)
  fix I s
  assume "s  $\in$  oreachable A ( $\lambda s s'. I$ ) U"
  thus " $(\exists \sigma. (\forall j \in J. \sigma j = (\text{fst } s) j) \wedge (\sigma, \text{snd } s) \in \text{reachable } A I)$ "
  proof induction
    fix s
    assume "s  $\in$  init A"
    hence " $(\text{fst } s, \text{snd } s) \in \text{reachable } A I$ "
      by simp (rule reachable_init)
    moreover have " $\forall j \in J. (\text{fst } s) j = (\text{fst } s) j$ "
      by simp
    ultimately show " $\exists \sigma. (\forall j \in J. \sigma j = (\text{fst } s) j) \wedge (\sigma, \text{snd } s) \in \text{reachable } A I$ "
      by auto
  next
    fix s a s'
    assume " $\exists \sigma. (\forall j \in J. \sigma j = (\text{fst } s) j) \wedge (\sigma, \text{snd } s) \in \text{reachable } A I$ "
      and " $(s, a, s') \in \text{trans } A$ "
      and "I a"
    then obtain  $\zeta$  where " $\forall j \in J. \zeta j = (\text{fst } s) j$ "
      and " $(\zeta, \text{snd } s) \in \text{reachable } A I$ " by auto

    from  $\langle (s, a, s') \in \text{trans } A \rangle$  have " $((\text{fst } s, \text{snd } s), a, (\text{fst } s', \text{snd } s')) \in \text{trans } A$ "
      by simp
    with  $\langle \text{local\_steps } (\text{trans } A) J \rangle$  obtain  $\zeta'$  where " $\forall j \in J. \zeta' j = (\text{fst } s') j$ "
      and " $((\zeta, \text{snd } s), a, (\zeta', \text{snd } s')) \in \text{trans } A$ "
      using " $\forall j \in J. \zeta j = (\text{fst } s) j$ " by - (drule(2) local_stepsE, clarsimp)
    from  $\langle (\zeta, \text{snd } s) \in \text{reachable } A I \rangle$ 
      and  $\langle ((\zeta, \text{snd } s), a, (\zeta', \text{snd } s')) \in \text{trans } A \rangle$ 
      and  $\langle I a \rangle$ 
      have " $(\zeta', \text{snd } s') \in \text{reachable } A I$ " ..

    with " $\forall j \in J. \zeta' j = (\text{fst } s') j$ "
      show " $\exists \sigma. (\forall j \in J. \sigma j = (\text{fst } s') j) \wedge (\sigma, \text{snd } s') \in \text{reachable } A I$ " by auto
  next
    fix s  $\sigma'$ 
    assume " $\exists \sigma. (\forall j \in J. \sigma j = (\text{fst } s) j) \wedge (\sigma, \text{snd } s) \in \text{reachable } A I$ "
      and "U (fst s)  $\sigma'$ "
    then obtain  $\sigma$  where " $\forall j \in J. \sigma j = (\text{fst } s) j$ "
      and " $(\sigma, \text{snd } s) \in \text{reachable } A I$ " by auto
    from  $\langle \text{other\_steps } U J \rangle$  and  $\langle U (\text{fst } s) \sigma' \rangle$  have " $\forall j \in J. \sigma' j = (\text{fst } s) j$ "
      by - (erule(1) other_stepsE)
    with " $\forall j \in J. \sigma j = (\text{fst } s) j$ " have " $\forall j \in J. \sigma j = \sigma' j$ "
      by clarsimp
    with  $\langle (\sigma, \text{snd } s) \in \text{reachable } A I \rangle$ 
      show " $\exists \sigma. (\forall j \in J. \sigma j = \text{fst } (\sigma', \text{snd } s) j) \wedge (\sigma, \text{snd } (\sigma', \text{snd } s)) \in \text{reachable } A I$ "
      by auto
  end
end

```

qed  
qed

```
lemma subreachableE [elim]:  
  assumes "subreachable A U J"  
    and "s ∈ oreachable A (λs s'. I) U"  
  shows "∃σ. (∀j∈J. σ j = (fst s) j) ∧ (σ, snd s) ∈ reachable A I"  
  using assms unfolding subreachable_def by simp
```

```
lemma subreachableE_pair [elim]:  
  assumes "subreachable A U J"  
    and "(σ, s) ∈ oreachable A (λs s'. I) U"  
  shows "∃ζ. (∀j∈J. ζ j = σ j) ∧ (ζ, s) ∈ reachable A I"  
  using assms unfolding subreachable_def by (metis fst_conv snd_conv)
```

```
lemma subreachable_otherE [elim]:  
  assumes "subreachable A U J"  
    and "(σ, l) ∈ oreachable A (λs s'. I) U"  
    and "U σ σ'"  
  shows "∃ζ'. (∀j∈J. ζ' j = σ' j) ∧ (ζ', l) ∈ reachable A I"  
proof -  
  from ⟨(σ, l) ∈ oreachable A (λs s'. I) U⟩ and ⟨U σ σ'⟩  
  have "(σ', l) ∈ oreachable A (λs s'. I) U"  
  by - (rule oreachable_other')  
  with ⟨subreachable A U J⟩ show ?thesis  
  by auto  
qed
```

```
lemma open_closed_invariant:  
  fixes J  
  assumes "A ⊨ (I →) P"  
    and "subreachable A U J"  
    and localp: "∧σ σ' s. [ [∀j∈J. σ' j = σ j; P (σ', s) ] ] ⇒ P (σ, s)"  
  shows "A ⊨ (act I, U →) P"  
proof (rule, simp_all only: act_def)  
  fix s  
  assume "s ∈ init A"  
  with ⟨A ⊨ (I →) P⟩ show "P s" ..  
next  
  fix s a s'  
  assume "s ∈ oreachable A (λ_ _ . I) U"  
    and "P s"  
    and "(s, a, s') ∈ trans A"  
    and "I a"  
  hence "s' ∈ oreachable A (λ_ _ . I) U"  
  by (metis oreachable_local)  
  with ⟨subreachable A U J⟩ obtain σ'  
  where "∀j∈J. σ' j = (fst s') j"  
    and "(σ', snd s') ∈ reachable A I"  
  by (metis subreachableE)  
  from ⟨A ⊨ (I →) P⟩ and ⟨(σ', snd s') ∈ reachable A I⟩ have "P (σ', snd s')" ..  
  with ⟨∀j∈J. σ' j = (fst s') j⟩ show "P s'"  
  by (metis localp prod.collapse)  
next  
  fix g g' l  
  assume or: "(g, l) ∈ oreachable A (λs s'. I) U"  
    and "U g g'"  
    and "P (g, l)"  
  from ⟨subreachable A U J⟩ and or and ⟨U g g'⟩  
  obtain gg' where "∀j∈J. gg' j = g' j"  
    and "(gg', l) ∈ reachable A I"  
  by (auto dest!: subreachable_otherE)  
  from ⟨A ⊨ (I →) P⟩ and ⟨(gg', l) ∈ reachable A I⟩  
  have "P (gg', l)" ..
```

with  $\langle \forall j \in J. gg' j = g' j \rangle$  show "P (g', l)"  
 by (rule localp)  
 qed

lemma oinvariant\_anyact:  
 assumes "A  $\models$  (act TT, U  $\rightarrow$ ) P"  
 shows "A  $\models$  (S, U  $\rightarrow$ ) P"  
 using assms by rule auto

definition

ostep\_invariant  
 :: "('g  $\times$  'l, 'a) automaton  
 $\Rightarrow$  ('g  $\Rightarrow$  'g  $\Rightarrow$  'a  $\Rightarrow$  bool)  $\Rightarrow$  ('g  $\Rightarrow$  'g  $\Rightarrow$  bool)  
 $\Rightarrow$  (('g  $\times$  'l, 'a) transition  $\Rightarrow$  bool)  $\Rightarrow$  bool"  
 ("\_  $\models_A$  (1'((1\_)/ (1\_)  $\rightarrow$ '))/ \_)" [100, 0, 0, 9] 8)

where

"(A  $\models_A$  (S, U  $\rightarrow$ ) P) =  
 ( $\forall s \in \text{oreachable } A \ S \ U. (\forall a \ s'. (s, a, s') \in \text{trans } A \wedge S \ (\text{fst } s) \ (\text{fst } s') \ a \ \longrightarrow P \ (s, a, s'))$ )"

lemma ostep\_invariant\_def':  
 "(A  $\models_A$  (S, U  $\rightarrow$ ) P) = ( $\forall s \in \text{oreachable } A \ S \ U.$   
 $(\forall a \ s'. (s, a, s') \in \text{trans } A \wedge S \ (\text{fst } s) \ (\text{fst } s') \ a \ \longrightarrow P \ (s, a, s'))$ )"  
 unfolding ostep\_invariant\_def by auto

lemma ostep\_invariantI [intro]:  
 assumes \*: " $\bigwedge \sigma \ s \ a \ \sigma' \ s'. \llbracket (\sigma, s) \in \text{oreachable } A \ S \ U; ((\sigma, s), a, (\sigma', s')) \in \text{trans } A; S \ \sigma \ \sigma' \ a \rrbracket$   
 $\implies P \ ((\sigma, s), a, (\sigma', s'))$ "  
 shows "A  $\models_A$  (S, U  $\rightarrow$ ) P"  
 unfolding ostep\_invariant\_def  
 using assms by auto

lemma ostep\_invariantD [dest]:  
 assumes "A  $\models_A$  (S, U  $\rightarrow$ ) P"  
 and " $(\sigma, s) \in \text{oreachable } A \ S \ U$ "  
 and " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "  
 and "S  $\sigma \ \sigma' \ a$ "  
 shows "P  $((\sigma, s), a, (\sigma', s'))$ "  
 using assms unfolding ostep\_invariant\_def' by clarsimp

lemma ostep\_invariantE [elim]:  
 assumes "A  $\models_A$  (S, U  $\rightarrow$ ) P"  
 and " $(\sigma, s) \in \text{oreachable } A \ S \ U$ "  
 and " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "  
 and "S  $\sigma \ \sigma' \ a$ "  
 and "P  $((\sigma, s), a, (\sigma', s')) \implies Q$ "  
 shows "Q"  
 using assms by auto

lemma ostep\_invariant\_weakenE [elim!]:  
 assumes invP: "A  $\models_A$  (PS, PU  $\rightarrow$ ) P"  
 and PQ: " $\bigwedge t. P \ t \implies Q \ t$ "  
 and QSPS: " $\bigwedge \sigma \ \sigma' \ a. QS \ \sigma \ \sigma' \ a \implies PS \ \sigma \ \sigma' \ a$ "  
 and QUPU: " $\bigwedge \sigma \ \sigma'. QU \ \sigma \ \sigma' \implies PU \ \sigma \ \sigma'$ "  
 shows "A  $\models_A$  (QS, QU  $\rightarrow$ ) Q"

proof

fix  $\sigma \ s \ \sigma' \ s' \ a$   
 assume " $(\sigma, s) \in \text{oreachable } A \ QS \ QU$ "  
 and " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "  
 and "QS  $\sigma \ \sigma' \ a$ "  
 from  $\langle QS \ \sigma \ \sigma' \ a \rangle$  have "PS  $\sigma \ \sigma' \ a$ " by (rule QSPS)  
 from  $\langle (\sigma, s) \in \text{oreachable } A \ QS \ QU \rangle$  have " $(\sigma, s) \in \text{oreachable } A \ PS \ PU$ " using QSPS QUPU ..  
 with invP have "P  $((\sigma, s), a, (\sigma', s'))$ " using  $\langle ((\sigma, s), a, (\sigma', s')) \in \text{trans } A \rangle$   $\langle PS \ \sigma \ \sigma' \ a \rangle$  ..  
 thus "Q  $((\sigma, s), a, (\sigma', s'))$ " by (rule PQ)

qed



```

lemma ostep_invariant_weaken_with_invariantE [elim]:
  assumes pinv: "A  $\models$  (S, U  $\rightarrow$ ) P"
    and qinv: "A  $\models_A$  (S, U  $\rightarrow$ ) Q"
    and wr: " $\bigwedge \sigma s a \sigma' s'. \llbracket P(\sigma, s); P(\sigma', s'); Q((\sigma, s), a, (\sigma', s')); S \sigma \sigma' a \rrbracket$ 
             $\implies R((\sigma, s), a, (\sigma', s'))$ "
  shows "A  $\models_A$  (S, U  $\rightarrow$ ) R"
proof
  fix  $\sigma s a \sigma' s'$ 
  assume sr: " $(\sigma, s) \in \text{oreachable } A \ S \ U$ "
    and tr: " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "
    and "S  $\sigma \sigma' a$ "
  hence " $(\sigma', s') \in \text{oreachable } A \ S \ U$ " ..
  with pinv have "P ( $\sigma', s'$ )" ..
  from pinv and sr have "P ( $\sigma, s$ )" ..
  from qinv sr tr  $\langle S \sigma \sigma' a \rangle$  have "Q ( $(\sigma, s), a, (\sigma', s')$ )" ..
  with  $\langle P(\sigma, s) \rangle$  and  $\langle P(\sigma', s') \rangle$  show "R ( $(\sigma, s), a, (\sigma', s')$ )" using  $\langle S \sigma \sigma' a \rangle$  by (rule wr)
qed

```

```

lemma ostep_to_invariantI:
  assumes sinv: "A  $\models_A$  (S, U  $\rightarrow$ ) Q"
    and init: " $\bigwedge \sigma s. (\sigma, s) \in \text{init } A \implies P(\sigma, s)$ "
    and local: " $\bigwedge \sigma s \sigma' s' a. \llbracket (\sigma, s) \in \text{oreachable } A \ S \ U;$ 
                P ( $\sigma, s$ );
                Q ( $(\sigma, s), a, (\sigma', s')$ );
                S  $\sigma \sigma' a \rrbracket \implies P(\sigma', s')$ "
    and other: " $\bigwedge \sigma \sigma' s. \llbracket (\sigma, s) \in \text{oreachable } A \ S \ U; U \sigma \sigma'; P(\sigma, s) \rrbracket \implies P(\sigma', s)$ "
  shows "A  $\models$  (S, U  $\rightarrow$ ) P"

```

```

proof
  fix  $\sigma s$  assume " $(\sigma, s) \in \text{init } A$ " thus "P ( $\sigma, s$ )" by (rule init)
next
  fix  $\sigma s \sigma' s' a$ 
  assume " $(\sigma, s) \in \text{oreachable } A \ S \ U$ "
    and "P ( $\sigma, s$ )"
    and " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "
    and "S  $\sigma \sigma' a$ "
  show "P ( $\sigma', s')$ "
proof -
  from sinv and  $\langle (\sigma, s) \in \text{oreachable } A \ S \ U \rangle$  and  $\langle ((\sigma, s), a, (\sigma', s')) \in \text{trans } A \rangle$  and  $\langle S \sigma \sigma' a \rangle$ 
    have "Q ( $(\sigma, s), a, (\sigma', s')$ )" ..
  with  $\langle (\sigma, s) \in \text{oreachable } A \ S \ U \rangle$  and  $\langle P(\sigma, s) \rangle$  show "P ( $\sigma', s')$ "
    using  $\langle S \sigma \sigma' a \rangle$  by (rule local)
qed

```

```

next
  fix  $\sigma \sigma' l$ 
  assume " $(\sigma, l) \in \text{oreachable } A \ S \ U$ "
    and "U  $\sigma \sigma'$ "
    and "P ( $\sigma, l$ )"
  thus "P ( $\sigma', l$ )" by (rule other)
qed

```

```

lemma open_closed_step_invariant:
  assumes "A  $\models_A$  (I  $\rightarrow$ ) P"
    and "local_steps (trans A) J"
    and "other_steps U J"
    and localp: " $\bigwedge \sigma \zeta a \sigma' \zeta' s s'. \llbracket \forall j \in J. \sigma j = \zeta j; \forall j \in J. \sigma' j = \zeta' j; P((\sigma, s), a, (\sigma', s')) \rrbracket$ 
                 $\implies P((\zeta, s), a, (\zeta', s'))$ "
  shows "A  $\models_A$  (act I, U  $\rightarrow$ ) P"

```

```

proof
  fix  $\sigma s a \sigma' s'$ 
  assume or: " $(\sigma, s) \in \text{oreachable } A \ (\text{act } I) \ U$ "
    and tr: " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "

```

```

    and "act I  $\sigma$   $\sigma'$  a"
  from ⟨act I  $\sigma$   $\sigma'$  a⟩ have "I a" ..
  from ⟨local_steps (trans A) J⟩ and ⟨other_steps U J⟩ have "subreachable A U J" ..
  then obtain  $\zeta$  where " $\forall j \in J. \zeta j = \sigma j$ "
    and " $(\zeta, s) \in \text{reachable A I}$ "
  using or unfolding act_def
  by (auto dest!: subreachableE_pair)

  from ⟨local_steps (trans A) J⟩ and tr and ⟨ $\forall j \in J. \zeta j = \sigma j$ ⟩
  obtain  $\zeta'$  where " $\forall j \in J. \zeta' j = \sigma' j$ "
    and " $((\zeta, s), a, (\zeta', s')) \in \text{trans A}$ "
  by auto

  from ⟨A  $\models_A$  (I  $\rightarrow$ ) P⟩ and ⟨ $(\zeta, s) \in \text{reachable A I}$ ⟩
    and ⟨ $((\zeta, s), a, (\zeta', s')) \in \text{trans A}$ ⟩
    and ⟨I a⟩
  have "P  $((\zeta, s), a, (\zeta', s'))$ " ..
  with ⟨ $\forall j \in J. \zeta j = \sigma j$ ⟩ and ⟨ $\forall j \in J. \zeta' j = \sigma' j$ ⟩ show "P  $((\sigma, s), a, (\sigma', s'))$ "
  by (rule localp)
qed

```

```

lemma oinvariant_step_anyact:
  assumes "p  $\models_A$  (act TT, U  $\rightarrow$ ) P"
  shows "p  $\models_A$  (S, U  $\rightarrow$ ) P"
  using assms by rule auto

```

### 4.3 Standard assumption predicates

otherwith

```

definition otherwith :: "('s  $\Rightarrow$  's  $\Rightarrow$  bool)
   $\Rightarrow$  'i set
   $\Rightarrow$  (('i  $\Rightarrow$  's)  $\Rightarrow$  'a  $\Rightarrow$  bool)
   $\Rightarrow$  ('i  $\Rightarrow$  's)  $\Rightarrow$  ('i  $\Rightarrow$  's)  $\Rightarrow$  'a  $\Rightarrow$  bool"
where "otherwith Q I P  $\sigma$   $\sigma'$  a  $\equiv$  ( $\forall i. i \notin I \rightarrow Q (\sigma i) (\sigma' i)$ )  $\wedge$  P  $\sigma$  a"

```

```

lemma otherwithI [intro]:
  assumes other: " $\bigwedge j. j \notin I \implies Q (\sigma j) (\sigma' j)$ "
  and sync: "P  $\sigma$  a"
  shows "otherwith Q I P  $\sigma$   $\sigma'$  a"
  unfolding otherwith_def using assms by simp

```

```

lemma otherwithE [elim]:
  assumes "otherwith Q I P  $\sigma$   $\sigma'$  a"
  and "[ P  $\sigma$  a;  $\forall j. j \notin I \rightarrow Q (\sigma j) (\sigma' j)$  ]  $\implies$  R  $\sigma$   $\sigma'$  a"
  shows "R  $\sigma$   $\sigma'$  a"
  using assms unfolding otherwith_def by simp

```

```

lemma otherwith_actionD [dest]:
  assumes "otherwith Q I P  $\sigma$   $\sigma'$  a"
  shows "P  $\sigma$  a"
  using assms by auto

```

```

lemma otherwith_syncD [dest]:
  assumes "otherwith Q I P  $\sigma$   $\sigma'$  a"
  shows " $\forall j. j \notin I \rightarrow Q (\sigma j) (\sigma' j)$ "
  using assms by auto

```

```

lemma otherwithEI [elim]:
  assumes "otherwith P I P0  $\sigma$   $\sigma'$  a"
  and " $\bigwedge \sigma a. P0 \sigma a \implies Q0 \sigma a$ "
  shows "otherwith P I Q0  $\sigma$   $\sigma'$  a"
  using assms(1) unfolding otherwith_def
  by (clarsimp elim!: assms(2))

```

```

lemma all_but:
  assumes " $\bigwedge \xi. S \xi \xi$ "
    and " $\sigma' i = \sigma i$ "
    and " $\forall j. j \neq i \longrightarrow S (\sigma j) (\sigma' j)$ "
  shows " $\forall j. S (\sigma j) (\sigma' j)$ "
  using assms by metis

lemma all_but_eq [dest]:
  assumes " $\sigma' i = \sigma i$ "
    and " $\forall j. j \neq i \longrightarrow \sigma j = \sigma' j$ "
  shows " $\sigma = \sigma'$ "
  using assms by - (rule ext, metis)

other

definition other :: " $('s \Rightarrow 's \Rightarrow \text{bool}) \Rightarrow 'i \text{ set} \Rightarrow ('i \Rightarrow 's) \Rightarrow ('i \Rightarrow 's) \Rightarrow \text{bool}$ "
where "other P I  $\sigma \sigma' \equiv \forall i. \text{if } i \in I \text{ then } \sigma' i = \sigma i \text{ else } P (\sigma i) (\sigma' i)$ "

lemma otherI [intro]:
  assumes local: " $\bigwedge i. i \in I \Longrightarrow \sigma' i = \sigma i$ "
    and other: " $\bigwedge j. j \notin I \Longrightarrow P (\sigma j) (\sigma' j)$ "
  shows "other P I  $\sigma \sigma'$ "
  using assms unfolding other_def by clarsimp

lemma otherE [elim]:
  assumes "other P I  $\sigma \sigma'$ "
    and "[ $\forall i \in I. \sigma' i = \sigma i; \forall j. j \notin I \longrightarrow P (\sigma j) (\sigma' j)$ ]  $\Longrightarrow R \sigma \sigma'$ "
  shows " $R \sigma \sigma'$ "
  using assms unfolding other_def by simp

lemma other_localD [dest]:
  "other P {i}  $\sigma \sigma' \Longrightarrow \sigma' i = \sigma i$ "
  by auto

lemma other_otherD [dest]:
  "other P {i}  $\sigma \sigma' \Longrightarrow \forall j. j \neq i \longrightarrow P (\sigma j) (\sigma' j)$ "
  by auto

lemma other_bothE [elim]:
  assumes "other P {i}  $\sigma \sigma'$ "
  obtains " $\sigma' i = \sigma i$ " and " $\forall j. j \neq i \longrightarrow P (\sigma j) (\sigma' j)$ "
  using assms by auto

lemma weaken_local [elim]:
  assumes "other P I  $\sigma \sigma'$ "
    and PQ: " $\bigwedge \xi \xi'. P \xi \xi' \Longrightarrow Q \xi \xi'$ "
  shows "other Q I  $\sigma \sigma'$ "
  using assms unfolding other_def by auto

definition global :: " $((\text{nat} \Rightarrow 's) \Rightarrow \text{bool}) \Rightarrow (\text{nat} \Rightarrow 's) \times 'local \Rightarrow \text{bool}$ "
where "global P  $\equiv (\lambda(\sigma, _). P \sigma)$ "

lemma globalsimp [simp]: "global P s = P (fst s)"
  unfolding global_def by (simp split: prod.split)

definition globala :: " $((\text{nat} \Rightarrow 's, 'action) \text{ transition} \Rightarrow \text{bool}) \Rightarrow ((\text{nat} \Rightarrow 's) \times 'local, 'action) \text{ transition} \Rightarrow \text{bool}$ "
where "globala P  $\equiv (\lambda((\sigma, _), a, (\sigma', _)). P (\sigma, a, \sigma'))$ "

lemma globalasimp [simp]: "globala P s = P (fst (fst s), fst (snd s), fst (snd (snd s)))"
  unfolding globala_def by (simp split: prod.split)

end

```

## 5 Terms of the Algebra for Wireless Networks

```
theory AWN
imports Lib
begin
```

### 5.1 Sequential Processes

```
type_synonym ip = nat
type_synonym data = nat
```

Most of AWN is independent of the type of messages, but the closed layer turns newpkt actions into the arrival of newpkt messages. We use a type class to maintain some abstraction (and independence from the definition of particular protocols).

```
class msg =
  fixes newpkt :: "data × ip ⇒ 'a"
  and eq_newpkt :: "'a ⇒ bool"
  assumes eq_newpkt_eq [simp]: "eq_newpkt (newpkt (d, i))"
```

Sequential process terms abstract over the types of data states ('s), messages ('m), process names ('p), and labels ('l).

```
datatype (dead 's, dead 'm, dead 'p, 'l) seqp =
  GUARD "'l" "'s ⇒ 's set" "('s, 'm, 'p, 'l) seqp"
| ASSIGN "'l" "'s ⇒ 's" "('s, 'm, 'p, 'l) seqp"
| CHOICE "('s, 'm, 'p, 'l) seqp" "('s, 'm, 'p, 'l) seqp"
| UCAST "'l" "'s ⇒ ip" "'s ⇒ 'm" "('s, 'm, 'p, 'l) seqp" "('s, 'm, 'p, 'l) seqp"
| BCAST "'l" "'s ⇒ 'm" "('s, 'm, 'p, 'l) seqp"
| GCAST "'l" "'s ⇒ ip set" "'s ⇒ 'm" "('s, 'm, 'p, 'l) seqp"
| SEND "'l" "'s ⇒ 'm" "('s, 'm, 'p, 'l) seqp"
| DELIVER "'l" "'s ⇒ data" "('s, 'm, 'p, 'l) seqp"
| RECEIVE "'l" "'m ⇒ 's ⇒ 's" "('s, 'm, 'p, 'l) seqp"
| CALL 'p
for map: labelmap
```

syntax

```
"_guard"    :: "[ 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨unbreakable⟩⟨_⟩⟩//_" [0, 60] 60)
"_lguard"   :: "[ 'a, 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨{ }⟨unbreakable⟩⟨_⟩⟩//_" [0, 0, 60] 60)
"_ifguard"  :: "[ pptrn, bool, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨unbreakable⟩⟨_ . _⟩⟩//_" [0, 0, 60] 60)

"_bassign"  :: "[ pptrn, 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨unbreakable⟩⟨[_ . _]⟩⟩//_" [0, 0, 60] 60)
"_lbassign" :: "[ 'a, pptrn, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("⟨{ }⟨unbreakable⟩⟨[_ . _]⟩⟩//_" [0, 0, 0, 60] 60)

"_assign"   :: "[ 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨unbreakable⟩⟨[_]⟩⟩//_" [0, 60] 60)
"_lassign"  :: "[ 'a, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("⟨{ }⟨unbreakable⟩⟨[_]⟩⟩//_" [0, 0, 60] 60)

"_unicast"  :: "[ 'a, 'a, ('s, 'm, 'p, unit) seqp, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨3unicast'((1(3_)/ (3_))') .//(_)/ (2> _))⟩" [0, 0, 60, 60] 60)
"_lunicast" :: "[ 'a, 'a, 'a, ('s, 'm, 'p, 'a) seqp, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("⟨⟨3{ }unicast'((1(3_)/ (3_))') .//(_)/ (2> _))⟩" [0, 0, 0, 60, 60] 60)

"_bcast"    :: "[ 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨3broadcast'((1(_))') .//_⟩⟩" [0, 60] 60)
"_lbcast"   :: "[ 'a, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("⟨⟨3{ }broadcast'((1(_))') .//_⟩⟩" [0, 0, 60] 60)

"_gcast"    :: "[ 'a, 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("⟨⟨3groupcast'((1(_)/ (_))') .//_⟩⟩" [0, 0, 60] 60)
```

```

"_lgcast"   :: "[ 'a, 'a, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("(3{ }groupcast'((1(_),/ (_)))' )//_" [0, 0, 0, 60] 60)

"_send"     :: "[ 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("(3send'((_)') )//_" [0, 60] 60)

"_lsend"    :: "[ 'a, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("(3{ }send'((_)') )//_" [0, 0, 60] 60)

"_deliver"  :: "[ 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("(3deliver'((_)') )//_" [0, 60] 60)

"_ldeliver" :: "[ 'a, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("(3{ }deliver'((_)') )//_" [0, 0, 60] 60)

"_receive"  :: "[ 'a, ('s, 'm, 'p, unit) seqp ] ⇒ ('s, 'm, 'p, unit) seqp"
              ("(3receive'((_)') )//_" [0, 60] 60)

"_lreceive" :: "[ 'a, 'a, ('s, 'm, 'p, 'a) seqp ] ⇒ ('s, 'm, 'p, 'a) seqp"
              ("(3{ }receive'((_)') )//_" [0, 0, 60] 60)

```

### translations

```

"_guard f p"      ⇒ "CONST GUARD () f p"
"_lguard l f p"   ⇒ "CONST GUARD l f p"
"_ifguard ξ e p"  ⇐ "CONST GUARD () (λξ. if e then {ξ} else {}) p"

"_assign f p"     ⇒ "CONST ASSIGN () f p"
"_lassign l f p"  ⇒ "CONST ASSIGN l f p"

"_bassign ξ e p"  ⇒ "CONST ASSIGN () (λξ. e) p"
"_lbassign l ξ e p" ⇒ "CONST ASSIGN l (λξ. e) p"

"_unicast fip fmsg p q" ⇒ "CONST UCAST () fip fmsg p q"
"_lunicast l fip fmsg p q" ⇒ "CONST UCAST l fip fmsg p q"

"_bcast fmsg p"   ⇒ "CONST BCAST () fmsg p"
"_lbcast l fmsg p" ⇒ "CONST BCAST l fmsg p"

"_gcast fipset fmsg p" ⇒ "CONST GCAST () fipset fmsg p"
"_lgcast l fipset fmsg p" ⇒ "CONST GCAST l fipset fmsg p"

"_send fmsg p"    ⇒ "CONST SEND () fmsg p"
"_lsend l fmsg p" ⇒ "CONST SEND l fmsg p"

"_deliver fdata p" ⇒ "CONST DELIVER () fdata p"
"_ldeliver l fdata p" ⇒ "CONST DELIVER l fdata p"

"_receive fmsg p" ⇒ "CONST RECEIVE () fmsg p"
"_lreceive l fmsg p" ⇒ "CONST RECEIVE l fmsg p"

```

notation "CHOICE" ("((\_)//⊕//(\_))" [56, 55] 55)  
and "CALL" ("(3call'((3\_))' )" [0] 60)

definition not\_call :: "('s, 'm, 'p, 'l) seqp ⇒ bool"  
where "not\_call p ≡ ∀pn. p ≠ call(pn)"

lemma not\_call\_simps [simp]:

```

"∧ l fg p.      not_call ({l}⟨fg⟩ p)"
"∧ l fa p.      not_call ({l}[fa] p)"
"∧ p1 p2.       not_call (p1 ⊕ p2)"
"∧ l fip fmsg p q. not_call ({l}unicast(fip, fmsg).p ▷ q)"
"∧ l fmsg p.    not_call ({l}broadcast(fmsg).p)"
"∧ l fips fmsg p. not_call ({l}groupcast(fips, fmsg).p)"
"∧ l fmsg p.    not_call ({l}send(fmsg).p)"
"∧ l fdata p.  not_call ({l}deliver(fdata).p)"
"∧ l fmsg p.    not_call ({l}receive(fmsg).p)"
"∧ l pn.       ¬(not_call (call(pn)))"

```

unfolding not\_call\_def by auto

definition not\_choice :: "('s, 'm, 'p, 'l) seqp  $\Rightarrow$  bool"  
 where "not\_choice p  $\equiv \forall p1 p2. p \neq p1 \oplus p2$ "

lemma not\_choice\_simps [simp]:

" $\bigwedge l fg p. \text{not\_choice } (\{l\}\langle fg \rangle p)$ "  
 " $\bigwedge l fa p. \text{not\_choice } (\{l\}\llbracket fa \rrbracket p)$ "  
 " $\bigwedge p1 p2. \neg(\text{not\_choice } (p1 \oplus p2))$ "  
 " $\bigwedge l fip fmsg p q. \text{not\_choice } (\{l\}\text{unicast}(fip, fmsg).p \triangleright q)$ "  
 " $\bigwedge l fmsg p. \text{not\_choice } (\{l\}\text{broadcast}(fmsg).p)$ "  
 " $\bigwedge l fips fmsg p. \text{not\_choice } (\{l\}\text{groupcast}(fips, fmsg).p)$ "  
 " $\bigwedge l fmsg p. \text{not\_choice } (\{l\}\text{send}(fmsg).p)$ "  
 " $\bigwedge l fdata p. \text{not\_choice } (\{l\}\text{deliver}(fdata).p)$ "  
 " $\bigwedge l fmsg p. \text{not\_choice } (\{l\}\text{receive}(fmsg).p)$ "  
 " $\bigwedge l pn. \text{not\_choice } (\text{call}(pn))$ "

unfolding not\_choice\_def by auto

lemma seqp\_congs:

" $\bigwedge l fg p. \{l\}\langle fg \rangle p = \{l\}\langle fg \rangle p$ "  
 " $\bigwedge l fa p. \{l\}\llbracket fa \rrbracket p = \{l\}\llbracket fa \rrbracket p$ "  
 " $\bigwedge p1 p2. p1 \oplus p2 = p1 \oplus p2$ "  
 " $\bigwedge l fip fmsg p q. \{l\}\text{unicast}(fip, fmsg).p \triangleright q = \{l\}\text{unicast}(fip, fmsg).p \triangleright q$ "  
 " $\bigwedge l fmsg p. \{l\}\text{broadcast}(fmsg).p = \{l\}\text{broadcast}(fmsg).p$ "  
 " $\bigwedge l fips fmsg p. \{l\}\text{groupcast}(fips, fmsg).p = \{l\}\text{groupcast}(fips, fmsg).p$ "  
 " $\bigwedge l fmsg p. \{l\}\text{send}(fmsg).p = \{l\}\text{send}(fmsg).p$ "  
 " $\bigwedge l fdata p. \{l\}\text{deliver}(fdata).p = \{l\}\text{deliver}(fdata).p$ "  
 " $\bigwedge l fmsg p. \{l\}\text{receive}(fmsg).p = \{l\}\text{receive}(fmsg).p$ "  
 " $\bigwedge l pn. \text{call}(pn) = \text{call}(pn)$ "

by auto

Remove data expressions from process terms.

fun seqp\_skeleton :: "('s, 'm, 'p, 'l) seqp  $\Rightarrow$  (unit, unit, 'p, 'l) seqp"

where

"seqp\_skeleton ( $\{l\}\langle \_ \rangle p$ ) =  $\{l\}\langle \lambda_. \{()\} \rangle$  (seqp\_skeleton p)"  
 | "seqp\_skeleton ( $\{l\}\llbracket \_ \rrbracket p$ ) =  $\{l\}\llbracket \lambda_. () \rrbracket$  (seqp\_skeleton p)"  
 | "seqp\_skeleton ( $p \oplus q$ ) = (seqp\_skeleton p)  $\oplus$  (seqp\_skeleton q)"  
 | "seqp\_skeleton ( $\{l\}\text{unicast}(\_, \_). p \triangleright q$ ) =  $\{l\}\text{unicast}(\lambda_. 0, \lambda_. ())$ . (seqp\_skeleton p)  $\triangleright$  (seqp\_skeleton q)"  
 | "seqp\_skeleton ( $\{l\}\text{broadcast}(\_). p$ ) =  $\{l\}\text{broadcast}(\lambda_. ())$ . (seqp\_skeleton p)"  
 | "seqp\_skeleton ( $\{l\}\text{groupcast}(\_, \_). p$ ) =  $\{l\}\text{groupcast}(\lambda_. \{ \}, \lambda_. ())$ . (seqp\_skeleton p)"  
 | "seqp\_skeleton ( $\{l\}\text{send}(\_). p$ ) =  $\{l\}\text{send}(\lambda_. ())$ . (seqp\_skeleton p)"  
 | "seqp\_skeleton ( $\{l\}\text{deliver}(\_). p$ ) =  $\{l\}\text{deliver}(\lambda_. 0)$ . (seqp\_skeleton p)"  
 | "seqp\_skeleton ( $\{l\}\text{receive}(\_). p$ ) =  $\{l\}\text{receive}(\lambda_. \_)$ . (seqp\_skeleton p)"  
 | "seqp\_skeleton (call(pn)) = call(pn)"

Calculate the subterms of a term.

fun subterms :: "('s, 'm, 'p, 'l) seqp  $\Rightarrow$  ('s, 'm, 'p, 'l) seqp set"

where

"subterms ( $\{l\}\langle fg \rangle p$ ) =  $\{\{l\}\langle fg \rangle p\} \cup \text{subterms } p$ "  
 | "subterms ( $\{l\}\llbracket fa \rrbracket p$ ) =  $\{\{l\}\llbracket fa \rrbracket p\} \cup \text{subterms } p$ "  
 | "subterms ( $p1 \oplus p2$ ) =  $\{p1 \oplus p2\} \cup \text{subterms } p1 \cup \text{subterms } p2$ "  
 | "subterms ( $\{l\}\text{unicast}(fip, fmsg). p \triangleright q$ ) =  
    $\{\{l\}\text{unicast}(fip, fmsg). p \triangleright q\} \cup \text{subterms } p \cup \text{subterms } q$ "  
 | "subterms ( $\{l\}\text{broadcast}(fmsg). p$ ) =  $\{\{l\}\text{broadcast}(fmsg). p\} \cup \text{subterms } p$ "  
 | "subterms ( $\{l\}\text{groupcast}(fips, fmsg). p$ ) =  $\{\{l\}\text{groupcast}(fips, fmsg). p\} \cup \text{subterms } p$ "  
 | "subterms ( $\{l\}\text{send}(fmsg). p$ ) =  $\{\{l\}\text{send}(fmsg). p\} \cup \text{subterms } p$ "  
 | "subterms ( $\{l\}\text{deliver}(fdata). p$ ) =  $\{\{l\}\text{deliver}(fdata). p\} \cup \text{subterms } p$ "  
 | "subterms ( $\{l\}\text{receive}(fmsg). p$ ) =  $\{\{l\}\text{receive}(fmsg). p\} \cup \text{subterms } p$ "  
 | "subterms (call(pn)) = {call(pn)}"

lemma subterms\_refl [simp]: "p  $\in$  subterms p"

by (cases p) simp\_all

```

lemma subterms_trans [elim]:
  assumes "q ∈ subterms p"
    and "r ∈ subterms q"
  shows "r ∈ subterms p"
  using assms by (induction p) auto

```

```

lemma root_in_subterms [simp]:
  "∧Γ pn. ∃pn'. Γ pn ∈ subterms (Γ pn')"
  by (rule_tac x=pn in exI) simp

```

```

lemma deriv_in_subterms [elim, dest]:
  "∧l f p q. {l}(f) q ∈ subterms p ⇒ q ∈ subterms p"
  "∧l fa p q. {l}[fa] q ∈ subterms p ⇒ q ∈ subterms p"
  "∧p1 p2 p. p1 ⊕ p2 ∈ subterms p ⇒ p1 ∈ subterms p"
  "∧p1 p2 p. p1 ⊕ p2 ∈ subterms p ⇒ p2 ∈ subterms p"
  "∧l fip fmsg p q r. {l}unicast(fip, fmsg). q ▷ r ∈ subterms p ⇒ q ∈ subterms p"
  "∧l fip fmsg p q r. {l}unicast(fip, fmsg). q ▷ r ∈ subterms p ⇒ r ∈ subterms p"
  "∧l fmsg p q. {l}broadcast(fmsg). q ∈ subterms p ⇒ q ∈ subterms p"
  "∧l fips fmsg p q. {l}groupcast(fips, fmsg). q ∈ subterms p ⇒ q ∈ subterms p"
  "∧l fmsg p q. {l}send(fmsg). q ∈ subterms p ⇒ q ∈ subterms p"
  "∧l fdata p q. {l}deliver(fdata). q ∈ subterms p ⇒ q ∈ subterms p"
  "∧l fmsg p q. {l}receive(fmsg). q ∈ subterms p ⇒ q ∈ subterms p"
  by auto

```

## 5.2 Actions

There are two sorts of  $\tau$  actions in AWN: one at the level of individual processes (within nodes), and one at the network level (outside nodes). We define a class so that we can ignore this distinction whenever it is not critical.

```

class tau =
  fixes tau :: "'a" ("τ")

```

### 5.2.1 Sequential Actions (and related predicates)

```

datatype 'm seq_action =
  broadcast 'm
  | groupcast "ip set" 'm
  | unicast ip 'm
  | notunicast ip          ("¬unicast _" [1000] 60)
  | send 'm
  | deliver data
  | receive 'm
  | seq_tau                ("τs")

```

```

instantiation "seq_action" :: (type) tau
begin
definition step_seq_tau [simp]: "τ ≡ τs"
instance ..
end

```

```

definition recvmsg :: "('m ⇒ bool) ⇒ 'm seq_action ⇒ bool"
where "recvmsg P a ≡ case a of receive m ⇒ P m
      | _ ⇒ True"

```

```

lemma recvmsg_simps[simp]:
  "∧m. recvmsg P (broadcast m) = True"
  "∧ips m. recvmsg P (groupcast ips m) = True"
  "∧ip m. recvmsg P (unicast ip m) = True"
  "∧ip. recvmsg P (notunicast ip) = True"
  "∧m. recvmsg P (send m) = True"
  "∧d. recvmsg P (deliver d) = True"
  "∧m. recvmsg P (receive m) = P m"
  "recvmsg P τs = True"
unfolding recvmsg_def by simp_all

```

lemma recvmsgTT [simp]: "recvmsg TT a"  
 by (cases a) simp\_all

lemma recvmsgE [elim]:  
 assumes "recvmsg (R  $\sigma$ ) a"  
 and " $\bigwedge m. R \sigma m \implies R \sigma' m$ "  
 shows "recvmsg (R  $\sigma'$ ) a"  
 using assms(1) by (cases a) (auto elim!: assms(2))

definition anycast :: "('m  $\Rightarrow$  bool)  $\Rightarrow$  'm seq\_action  $\Rightarrow$  bool"  
 where "anycast P a  $\equiv$  case a of broadcast m  $\Rightarrow$  P m  
 | groupcast \_ m  $\Rightarrow$  P m  
 | unicast \_ m  $\Rightarrow$  P m  
 | \_  $\Rightarrow$  True"

lemma anycast\_simps [simp]:  
 " $\bigwedge m. anycast P$  (broadcast m) = P m"  
 " $\bigwedge ips m. anycast P$  (groupcast ips m) = P m"  
 " $\bigwedge ip m. anycast P$  (unicast ip m) = P m"  
 " $\bigwedge ip. anycast P$  (notunicast ip) = True"  
 " $\bigwedge m. anycast P$  (send m) = True"  
 " $\bigwedge d. anycast P$  (deliver d) = True"  
 " $\bigwedge m. anycast P$  (receive m) = True"  
 "anycast P  $\tau_s$  = True"  
 unfolding anycast\_def by simp\_all

definition orecvmsg :: "((ip  $\Rightarrow$  's)  $\Rightarrow$  'm  $\Rightarrow$  bool)  $\Rightarrow$  (ip  $\Rightarrow$  's)  $\Rightarrow$  'm seq\_action  $\Rightarrow$  bool"  
 where "orecvmsg P  $\sigma$  a  $\equiv$  (case a of receive m  $\Rightarrow$  P  $\sigma$  m  
 | \_  $\Rightarrow$  True)"

lemma orecvmsg\_simps [simp]:  
 " $\bigwedge m. orecvmsg P \sigma$  (broadcast m) = True"  
 " $\bigwedge ips m. orecvmsg P \sigma$  (groupcast ips m) = True"  
 " $\bigwedge ip m. orecvmsg P \sigma$  (unicast ip m) = True"  
 " $\bigwedge ip. orecvmsg P \sigma$  (notunicast ip) = True"  
 " $\bigwedge m. orecvmsg P \sigma$  (send m) = True"  
 " $\bigwedge d. orecvmsg P \sigma$  (deliver d) = True"  
 " $\bigwedge m. orecvmsg P \sigma$  (receive m) = P  $\sigma$  m"  
 "orecvmsg P  $\sigma$   $\tau_s$  = True"  
 unfolding orecvmsg\_def by simp\_all

lemma orecvmsgEI [elim]:  
 "[[ orecvmsg P  $\sigma$  a;  $\bigwedge \sigma a. P \sigma a \implies Q \sigma a$  ]  $\implies$  orecvmsg Q  $\sigma$  a"  
 by (cases a) simp\_all

lemma orecvmsg\_stateless\_recvmsg [elim]:  
 "orecvmsg ( $\lambda_. P$ )  $\sigma$  a  $\implies$  recvmsg P a"  
 by (cases a) simp\_all

lemma orecvmsg\_recv\_weaken [elim]:  
 "[[ orecvmsg P  $\sigma$  a;  $\bigwedge \sigma a. P \sigma a \implies Q a$  ]  $\implies$  recvmsg Q a"  
 by (cases a) simp\_all

lemma orecvmsg\_recvmsg [elim]:  
 "orecvmsg P  $\sigma$  a  $\implies$  recvmsg (P  $\sigma$ ) a"  
 by (cases a) simp\_all

definition sendmsg :: "('m  $\Rightarrow$  bool)  $\Rightarrow$  'm seq\_action  $\Rightarrow$  bool"  
 where "sendmsg P a  $\equiv$  case a of send m  $\Rightarrow$  P m | \_  $\Rightarrow$  True"

lemma sendmsg\_simps [simp]:  
 " $\bigwedge m. sendmsg P$  (broadcast m) = True"  
 " $\bigwedge ips m. sendmsg P$  (groupcast ips m) = True"  
 " $\bigwedge ip m. sendmsg P$  (unicast ip m) = True"



```

" $\wedge$ ip.    sendmsg P (notunicast ip)    = True"
" $\wedge$ m.    sendmsg P (send m)           = P m"
" $\wedge$ d.    sendmsg P (deliver d)        = True"
" $\wedge$ m.    sendmsg P (receive m)        = True"
"        sendmsg P  $\tau_s$               = True"
unfolding sendmsg_def by simp_all

```

```
type_synonym ('s, 'm, 'p, 'l) seqp_env = "'p  $\Rightarrow$  ('s, 'm, 'p, 'l) seqp"
```

## 5.2.2 Node Actions (and related predicates)

```

datatype 'm node_action =
  node_cast "ip set" 'm          ("_*cast'(_)"          [200, 200] 200)
| node_deliver ip data          ("_:_deliver'(_)"          [200, 200] 200)
| node_arrive "ip set" "ip set" 'm ("_¬_:arrive'(_)"          [200, 200, 200] 200)
| node_connect ip ip            ("connect'(_, _)"          [200, 200] 200)
| node_disconnect ip ip        ("disconnect'(_, _)"      [200, 200] 200)
| node_newpkt ip data ip       ("_:_newpkt'(_, _)"      [200, 200, 200] 200)
| node_tau                      (" $\tau_n$ ")

```

```
instantiation "node_action" :: (type) tau
```

```
begin
```

```
definition step_node_tau [simp]: " $\tau \equiv \tau_n$ "
```

```
instance ..
```

```
end
```

```
definition arrivemsg :: "ip  $\Rightarrow$  ('m  $\Rightarrow$  bool)  $\Rightarrow$  'm node_action  $\Rightarrow$  bool"
```

```
where "arrivemsg i P a  $\equiv$  case a of node_arrive ii ni m  $\Rightarrow$  ((ii = {i}  $\longrightarrow$  P m)
| _  $\Rightarrow$  True"
```

```
lemma arrivemsg_simps[simp]:
```

```

" $\wedge$ R m.    arrivemsg i P (R:*cast(m))      = True"
" $\wedge$ d m.    arrivemsg i P (d:deliver(m))      = True"
" $\wedge$ i ii ni m. arrivemsg i P (ii¬ni:arrive(m)) = (ii = {i}  $\longrightarrow$  P m)"
" $\wedge$ i1 i2.  arrivemsg i P (connect(i1, i2))   = True"
" $\wedge$ i1 i2.  arrivemsg i P (disconnect(i1, i2)) = True"
" $\wedge$ i i' d di. arrivemsg i P (i':newpkt(d, di)) = True"
"          arrivemsg i P  $\tau_n$               = True"

```

```
unfolding arrivemsg_def by simp_all
```

```
lemma arrivemsgTT [simp]: "arrivemsg i TT = TT"
```

```
by (rule ext) (clarsimp simp: arrivemsg_def split: node_action.split)
```

```
definition oarrivemsg :: "((ip  $\Rightarrow$  's)  $\Rightarrow$  'm  $\Rightarrow$  bool)  $\Rightarrow$  (ip  $\Rightarrow$  's)  $\Rightarrow$  'm node_action  $\Rightarrow$  bool"
```

```
where "oarrivemsg P  $\sigma$  a  $\equiv$  case a of node_arrive ii ni m  $\Rightarrow$  P  $\sigma$  m | _  $\Rightarrow$  True"
```

```
lemma oarrivemsg_simps[simp]:
```

```

" $\wedge$ R m.    oarrivemsg P  $\sigma$  (R:*cast(m))    = True"
" $\wedge$ d m.    oarrivemsg P  $\sigma$  (d:deliver(m))  = True"
" $\wedge$ i ii ni m. oarrivemsg P  $\sigma$  (ii¬ni:arrive(m)) = P  $\sigma$  m"
" $\wedge$ i1 i2.  oarrivemsg P  $\sigma$  (connect(i1, i2)) = True"
" $\wedge$ i1 i2.  oarrivemsg P  $\sigma$  (disconnect(i1, i2)) = True"
" $\wedge$ i i' d di. oarrivemsg P  $\sigma$  (i':newpkt(d, di)) = True"
"          oarrivemsg P  $\sigma$   $\tau_n$           = True"

```

```
unfolding oarrivemsg_def by simp_all
```

```
lemma oarrivemsg_True [simp, intro]: "oarrivemsg ( $\lambda$  _ . True)  $\sigma$  a"
```

```
by (cases a) auto
```

```
definition castmsg :: "('m  $\Rightarrow$  bool)  $\Rightarrow$  'm node_action  $\Rightarrow$  bool"
```

```
where "castmsg P a  $\equiv$  case a of _:*cast(m)  $\Rightarrow$  P m
| _  $\Rightarrow$  True"
```

```
lemma castmsg_simps[simp]:
```

```

"∧R m.      castmsg P (R:*cast(m))      = P m"
"∧d m.      castmsg P (d:deliver(m))    = True"
"∧i ii ni m. castmsg P (ii~ni:arrive(m)) = True"
"∧i1 i2.    castmsg P (connect(i1, i2)) = True"
"∧i1 i2.    castmsg P (disconnect(i1, i2)) = True"
"∧i i' d di. castmsg P (i':newpkt(d, di)) = True"
"          castmsg P τn                = True"
unfolding castmsg_def by simp_all

```

### 5.3 Networks

```

datatype net_tree =
  Node ip "ip set"          ("⟨_; _⟩")
  | Subnet net_tree net_tree (infixl "||" 90)

declare net_tree.induct [[induct del]]
lemmas net_tree.induct [induct type: net_tree] = net_tree.induct [rename_abs i R p1 p2]

datatype 's net_state =
  NodeS ip 's "ip set"
  | SubnetS "'s net_state" "'s net_state"

fun net_ips :: "'s net_state ⇒ ip set"
where
  "net_ips (NodeS i s R) = {i}"
  | "net_ips (SubnetS n1 n2) = net_ips n1 ∪ net_ips n2"

fun net_tree_ips :: "net_tree ⇒ ip set"
where
  "net_tree_ips (p1 || p2) = net_tree_ips p1 ∪ net_tree_ips p2"
  | "net_tree_ips (⟨i; R⟩) = {i}"

lemma net_tree_ips_commute:
  "net_tree_ips (p1 || p2) = net_tree_ips (p2 || p1)"
  by simp (rule Un_commute)

fun wf_net_tree :: "net_tree ⇒ bool"
where
  "wf_net_tree (p1 || p2) = (net_tree_ips p1 ∩ net_tree_ips p2 = {})
    ∧ wf_net_tree p1 ∧ wf_net_tree p2"
  | "wf_net_tree (⟨i; R⟩) = True"

lemma wf_net_tree_children [elim]:
  assumes "wf_net_tree (p1 || p2)"
  obtains "wf_net_tree p1"
    and "wf_net_tree p2"
  using assms by simp

fun netmap :: "'s net_state ⇒ ip ⇒ 's option"
where
  "netmap (NodeS i p Ri) = [i ↦ p]"
  | "netmap (SubnetS s t) = netmap s ++ netmap t"

lemma not_in_netmap [simp]:
  assumes "i ∉ net_ips ns"
  shows "netmap ns i = None"
  using assms by (induction ns) simp_all

lemma netmap_none_not_in_net_ips:
  assumes "netmap ns i = None"
  shows "i ∉ net_ips ns"
  using assms by (induction ns) auto

lemma net_ips_is_dom_netmap: "net_ips s = dom(netmap s)"

```

```

proof (induction s)
  fix i Ri and p :: 's
  show "net_ips (NodeS i p Ri) = dom (netmap (NodeS i p Ri))"
    by auto
next
  fix s1 s2 :: "'s net_state"
  assume "net_ips s1 = dom (netmap s1)"
    and "net_ips s2 = dom (netmap s2)"
  thus "net_ips (SubnetS s1 s2) = dom (netmap (SubnetS s1 s2))"
    by auto
qed

lemma in_netmap [simp]:
  assumes "i ∈ net_ips ns"
  shows "netmap ns i ≠ None"
  using assms by (auto simp add: net_ips_is_dom_netmap)

lemma netmap_subnets_same:
  assumes "netmap s1 i = x"
    and "netmap s2 i = x"
  shows "netmap (SubnetS s1 s2) i = x"
  using assms by simp (metis map_add_dom_app_simps(1) map_add_dom_app_simps(3))

lemma netmap_subnets_samef:
  assumes "netmap s1 = f"
    and "netmap s2 = f"
  shows "netmap (SubnetS s1 s2) = f"
  using assms by simp (metis map_add_le_mapI map_le_antisym map_le_map_add map_le_refl)

lemma netmap_add_disjoint [elim]:
  assumes "∀i∈net_ips s1 ∪ net_ips s2. the ((netmap s1 ++ netmap s2) i) = σ i"
    and "net_ips s1 ∩ net_ips s2 = {}"
  shows "∀i∈net_ips s1. the (netmap s1 i) = σ i"
proof
  fix i
  assume "i ∈ net_ips s1"
  hence "i ∈ dom(netmap s1)" by (simp add: net_ips_is_dom_netmap)
  moreover with assms(2) have "i ∉ dom(netmap s2)" by (auto simp add: net_ips_is_dom_netmap)
  ultimately have "the (netmap s1 i) = the ((netmap s1 ++ netmap s2) i)"
    by (simp add: map_add_dom_app_simps)
  with assms(1) and ⟨i∈net_ips s1⟩ show "the (netmap s1 i) = σ i" by simp
qed

lemma netmap_add_disjoint2 [elim]:
  assumes "∀i∈net_ips s1 ∪ net_ips s2. the ((netmap s1 ++ netmap s2) i) = σ i"
  shows "∀i∈net_ips s2. the (netmap s2 i) = σ i"
  using assms by (simp add: net_ips_is_dom_netmap)
    (metis Un_iff map_add_dom_app_simps(1))

lemma net_ips_netmap_subnet [elim]:
  assumes "net_ips s1 ∩ net_ips s2 = {}"
    and "∀i∈net_ips (SubnetS s1 s2). the (netmap (SubnetS s1 s2) i) = σ i"
  shows "∀i∈net_ips s1. the (netmap s1 i) = σ i"
    and "∀i∈net_ips s2. the (netmap s2 i) = σ i"
proof -
  from assms(2) have "∀i∈net_ips s1 ∪ net_ips s2. the ((netmap s1 ++ netmap s2) i) = σ i" by auto
  with assms(1) show "∀i∈net_ips s1. the (netmap s1 i) = σ i"
    by - (erule(1) netmap_add_disjoint)
next
  from assms(2) have "∀i∈net_ips s1 ∪ net_ips s2. the ((netmap s1 ++ netmap s2) i) = σ i" by auto
  thus "∀i∈net_ips s2. the (netmap s2 i) = σ i"
    by - (erule netmap_add_disjoint2)
qed

```

```

fun inoclosed :: "'s ⇒ 'm::msg node_action ⇒ bool"
where
  "inoclosed _ (node_arrive ii ni m) = eq_newpkt m"
  | "inoclosed _ (node_newpkt i d di) = False"
  | "inoclosed _ _ = True"

lemma inclosed_simps [simp]:
  "∧σ ii ni. inoclosed σ (ii~ni:arrive(m)) = eq_newpkt m"
  "∧σ d di. inoclosed σ (i:newpkt(d, di)) = False"
  "∧σ R m. inoclosed σ (R:*cast(m)) = True"
  "∧σ i d. inoclosed σ (i:deliver(d)) = True"
  "∧σ i i'. inoclosed σ (connect(i, i')) = True"
  "∧σ i i'. inoclosed σ (disconnect(i, i')) = True"
  "∧σ. inoclosed σ (τ) = True"
by auto

definition
  netmask :: "ip set ⇒ ((ip ⇒ 's) × 'l) ⇒ ((ip ⇒ 's option) × 'l)"
where
  "netmask I s ≡ (λi. if i∈I then Some (fst s i) else None, snd s)"

lemma netmask_def' [simp]:
  "netmask I (σ, ζ) = (λi. if i∈I then Some (σ i) else None, ζ)"
unfolding netmask_def by auto

fun netgmap :: "('s ⇒ 'g × 'l) ⇒ 's net_state ⇒ (nat ⇒ 'g option) × 'l net_state"
where
  "netgmap sr (NodeS i s R) = ([i ↦ fst (sr s)], NodeS i (snd (sr s)) R)"
  | "netgmap sr (SubnetS s1 s2) = (let (σ1, ss) = netgmap sr s1 in
    let (σ2, tt) = netgmap sr s2 in
    (σ1 ++ σ2, SubnetS ss tt))"

lemma dom_fst_netgmap [simp, intro]: "dom (fst (netgmap sr n)) = net_ips n"
proof (induction n)
  fix i s R
  show "dom (fst (netgmap sr (NodeS i s R))) = net_ips (NodeS i s R)"
  by simp
next
  fix n1 n2
  assume a1: "dom (fst (netgmap sr n1)) = net_ips n1"
  and a2: "dom (fst (netgmap sr n2)) = net_ips n2"
  obtain σ1 ζ1 σ2 ζ2 where nm1: "netgmap sr n1 = (σ1, ζ1)"
  and nm2: "netgmap sr n2 = (σ2, ζ2)"
  by (metis surj_pair)
  hence "netgmap sr (SubnetS n1 n2) = (σ1 ++ σ2, SubnetS ζ1 ζ2)" by simp
  hence "dom (fst (netgmap sr (SubnetS n1 n2))) = dom (σ1 ++ σ2)" by simp
  also from a1 a2 nm1 nm2 have "dom (σ1 ++ σ2) = net_ips (SubnetS n1 n2)" by auto
  finally show "dom (fst (netgmap sr (SubnetS n1 n2))) = net_ips (SubnetS n1 n2)" .
qed

lemma netgmap_pair_dom [elim]:
  obtains σ ζ where "netgmap sr n = (σ, ζ)"
  and "dom σ = net_ips n"
  by (metis dom_fst_netgmap surjective_pairing)

lemma net_ips_netgmap [simp]:
  "net_ips (snd (netgmap sr s)) = net_ips s"
proof (induction s)
  fix s1 s2
  assume "net_ips (snd (netgmap sr s1)) = net_ips s1"
  and "net_ips (snd (netgmap sr s2)) = net_ips s2"
  thus "net_ips (snd (netgmap sr (SubnetS s1 s2))) = net_ips (SubnetS s1 s2)"
  by (cases "netgmap sr s1", cases "netgmap sr s2") auto
qed simp

```

```

lemma some_the_fst_netgmap:
  assumes "i ∈ net_ips s"
  shows "Some (the (fst (netgmap sr s) i)) = fst (netgmap sr s) i"
  using assms by (metis domIff dom_fst_netgmap option.collapse)

lemma fst_netgmap_none [simp]:
  assumes "i ∉ net_ips s"
  shows "fst (netgmap sr s) i = None"
  using assms by (metis domIff dom_fst_netgmap)

lemma fst_netgmap_subnet [simp]:
  "fst (case netgmap sr s1 of (σ1, ss) ⇒
    case netgmap sr s2 of (σ2, tt) ⇒
      (σ1 ++ σ2, SubnetS ss tt)) = (fst (netgmap sr s1) ++ fst (netgmap sr s2))"
  by (metis (mono_tags) fst_conv netgmap_pair_dom split_conv)

lemma snd_netgmap_subnet [simp]:
  "snd (case netgmap sr s1 of (σ1, ss) ⇒
    case netgmap sr s2 of (σ2, tt) ⇒
      (σ1 ++ σ2, SubnetS ss tt)) = (SubnetS (snd (netgmap sr s1)) (snd (netgmap sr s2)))"
  by (metis (lifting, no_types) Pair_inject split_beta' surjective_pairing)

lemma fst_netgmap_not_none [simp]:
  assumes "i ∈ net_ips s"
  shows "fst (netgmap sr s) i ≠ None"
  using assms by (induction s) auto

lemma netgmap_netgmap_not_rhs [simp]:
  assumes "i ∉ net_ips s2"
  shows "(fst (netgmap sr s1) ++ fst (netgmap sr s2)) i = (fst (netgmap sr s1)) i"
  proof -
    from assms(1) have "i ∉ dom (fst (netgmap sr s2))" by simp
    thus ?thesis by (simp add: map_add_dom_app_simps)
  qed

lemma netgmap_netgmap_rhs [simp]:
  assumes "i ∈ net_ips s2"
  shows "(fst (netgmap sr s1) ++ fst (netgmap sr s2)) i = (fst (netgmap sr s2)) i"
  using assms by (simp add: map_add_dom_app_simps)

lemma netgmap_netmask_subnets [elim]:
  assumes "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
    and "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
  shows "fst (netgmap sr (SubnetS s1 s2))
    = fst (netmask (net_tree_ips (n1 || n2)) (σ, snd (netgmap sr (SubnetS s1 s2))))"
  proof (rule ext)
    fix i
    have "i ∈ net_tree_ips n1 ∨ i ∈ net_tree_ips n2 ∨ (i ∉ net_tree_ips n1 ∪ net_tree_ips n2)"
      by auto
    thus "fst (netgmap sr (SubnetS s1 s2)) i
      = fst (netmask (net_tree_ips (n1 || n2)) (σ, snd (netgmap sr (SubnetS s1 s2)))) i"
  proof (elim disjE)
    assume "i ∈ net_tree_ips n1"
    with ⟨netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))⟩
      ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
    show ?thesis
      by (cases "netgmap sr s1", cases "netgmap sr s2", clarsimp)
        (metis (lifting, mono_tags) map_add_Some_iff)
  next
    assume "i ∈ net_tree_ips n2"
    with ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
    show ?thesis
  end
end

```

```

    by simp (metis (lifting, mono_tags) fst_conv map_add_find_right)
next
assume "i ∈ net_tree_ips n1 ∪ net_tree_ips n2"
with ⟨netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))⟩
  ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
show ?thesis
  by simp (metis (lifting, mono_tags) fst_conv)
qed
qed

lemma netgmap_netmask_subnets' [elim]:
assumes "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
  and "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
  and "s = SubnetS s1 s2"
shows "netgmap sr s = netmask (net_tree_ips (n1 || n2)) (σ, snd (netgmap sr s))"
by (simp only: assms(3))
  (rule prod_eqI [OF netgmap_netmask_subnets [OF assms(1-2)]], simp)

lemma netgmap_subnet_split1:
assumes "netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) (σ, ζ)"
  and "net_tree_ips n1 ∩ net_tree_ips n2 = {}"
  and "net_ips s1 = net_tree_ips n1"
  and "net_ips s2 = net_tree_ips n2"
shows "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
proof (rule prod_eqI)
show "fst (netgmap sr s1) = fst (netmask (net_tree_ips n1) (σ, snd (netgmap sr s1)))"
proof (rule ext, simp, intro conjI impI)
fix i
assume "i ∈ net_tree_ips n1"
with ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩ have "i ∉ net_tree_ips n2"
  by auto
from assms(1) [simplified prod_eq_iff]
  have "(fst (netgmap sr s1) ++ fst (netgmap sr s2)) i =
    (if i ∈ net_tree_ips n1 ∨ i ∈ net_tree_ips n2 then Some (σ i) else None)"
  by simp
also from ⟨i ∉ net_tree_ips n2⟩ and ⟨net_ips s2 = net_tree_ips n2⟩
  have "(fst (netgmap sr s1) ++ fst (netgmap sr s2)) i = fst (netgmap sr s1) i"
  by (metis dom_fst_netgmap map_add_dom_app_simps(3))
finally show "fst (netgmap sr s1) i = Some (σ i)"
  using ⟨i ∈ net_tree_ips n1⟩ by simp
next
fix i
assume "i ∉ net_tree_ips n1"
with ⟨net_ips s1 = net_tree_ips n1⟩ have "i ∉ net_ips s1" by simp
thus "fst (netgmap sr s1) i = None" by simp
qed
qed simp

lemma netgmap_subnet_split2:
assumes "netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) (σ, ζ)"
  and "net_ips s1 = net_tree_ips n1"
  and "net_ips s2 = net_tree_ips n2"
shows "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
proof (rule prod_eqI)
show "fst (netgmap sr s2) = fst (netmask (net_tree_ips n2) (σ, snd (netgmap sr s2)))"
proof (rule ext, simp, intro conjI impI)
fix i
assume "i ∈ net_tree_ips n2"
from assms(1) [simplified prod_eq_iff]
  have "(fst (netgmap sr s1) ++ fst (netgmap sr s2)) i =
    (if i ∈ net_tree_ips n1 ∨ i ∈ net_tree_ips n2 then Some (σ i) else None)"
  by simp
also from ⟨i ∈ net_tree_ips n2⟩ and ⟨net_ips s2 = net_tree_ips n2⟩
  have "(fst (netgmap sr s1) ++ fst (netgmap sr s2)) i = fst (netgmap sr s2) i"

```

```

    by (metis dom_fst_netgmap map_add_dom_app_simps(1))
  finally show "fst (netgmap sr s2) i = Some ( $\sigma$  i)"
    using (i ∈ net_tree_ips n2) by simp
next
  fix i
  assume "i ∉ net_tree_ips n2"
  with (net_ips s2 = net_tree_ips n2) have "i ∉ net_ips s2" by simp
  thus "fst (netgmap sr s2) i = None" by simp
qed
qed simp

lemma netmap_fst_netgmap_rel:
  shows "( $\lambda$ i. map_option (fst ∘ sr) (netmap s i)) = fst (netgmap sr s)"
  proof (induction s)
    fix ii s R
    show "( $\lambda$ i. map_option (fst ∘ sr) (netmap (NodeS ii s R) i)) = fst (netgmap sr (NodeS ii s R))"
      by auto
  next
    fix s1 s2
    assume a1: "( $\lambda$ i. map_option (fst ∘ sr) (netmap s1 i)) = fst (netgmap sr s1)"
      and a2: "( $\lambda$ i. map_option (fst ∘ sr) (netmap s2 i)) = fst (netgmap sr s2)"
    show "( $\lambda$ i. map_option (fst ∘ sr) (netmap (SubnetS s1 s2) i)) = fst (netgmap sr (SubnetS s1 s2))"
      proof (rule ext)
        fix i
        from a1 a2 have "map_option (fst ∘ sr) ((netmap s1 ++ netmap s2) i)
          = (fst (netgmap sr s1) ++ fst (netgmap sr s2)) i"
          by (metis fst_conv map_add_dom_app_simps(1) map_add_dom_app_simps(3)
            net_ips_is_dom_netmap netgmap_pair_dom)
        thus "map_option (fst ∘ sr) (netmap (SubnetS s1 s2) i) = fst (netgmap sr (SubnetS s1 s2)) i"
          by simp
      qed
    qed
  qed

lemma netmap_is_fst_netgmap:
  assumes "netmap s' = netmap s"
  shows "fst (netgmap sr s') = fst (netgmap sr s)"
  using assms by (metis netmap_fst_netgmap_rel)

lemma netmap_is_fst_netgmap':
  assumes "netmap s' i = netmap s i"
  shows "fst (netgmap sr s') i = fst (netgmap sr s) i"
  using assms by (metis netmap_fst_netgmap_rel)

lemma fst_netgmap_pair_fst [simp]:
  "fst (netgmap ( $\lambda$ (p, q). (fst p, snd p, q)) s) = fst (netgmap fst s)"
  by (induction s) auto

Introduce streamlined alternatives to netgmap to simplify certain property statements and thus make them easier
to understand and to present.

fun netlift :: "('s ⇒ 'g × 'l) ⇒ 's net_state ⇒ (nat ⇒ 'g option)"
  where
    "netlift sr (NodeS i s R) = [i ↦ fst (sr s)]"
  | "netlift sr (SubnetS s t) = (netlift sr s) ++ (netlift sr t)"

lemma fst_netgmap_netlift:
  "fst (netgmap sr s) = netlift sr s"
  by (induction s) simp_all

fun netliftl :: "('s ⇒ 'g × 'l) ⇒ 's net_state ⇒ 'l net_state"
  where
    "netliftl sr (NodeS i s R) = NodeS i (snd (sr s)) R"
  | "netliftl sr (SubnetS s t) = SubnetS (netliftl sr s) (netliftl sr t)"

lemma snd_netgmap_netliftl:

```

```
"snd (netgmap sr s) = netliftl sr s"
by (induction s) simp_all
```

```
lemma netgmap_netlift_netliftl: "netgmap sr s = (netlift sr s, netliftl sr s)"
by rule (simp_all add: fst_netgmap_netlift snd_netgmap_netliftl)
```

end

## 6 Semantics of the Algebra of Wireless Networks

```
theory Awn_SOS
imports TransitionSystems Awn
begin
```

### 6.1 Table 1: Structural operational semantics for sequential process expressions

```
inductive_set
```

```
seqp_sos
```

```
:: "('s, 'm, 'p, 'l) seqp_env  $\Rightarrow$  ('s  $\times$  ('s, 'm, 'p, 'l) seqp, 'm seq_action) transition set"
```

```
for  $\Gamma$  :: "('s, 'm, 'p, 'l) seqp_env"
```

```
where
```

```
  broadcastT: "(( $\xi$ , {l}broadcast(smsg).p), broadcast (smsg  $\xi$ ), ( $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "
| groupcastT: "(( $\xi$ , {l}groupcast(sips, smsg).p), groupcast (sips  $\xi$ ) (smsg  $\xi$ ), ( $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "
| unicastT: "(( $\xi$ , {l}unicast(sip, smsg).p  $\triangleright$  q), unicast (sip  $\xi$ ) (smsg  $\xi$ ), ( $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "
| notunicastT: "(( $\xi$ , {l}unicast(sip, smsg).p  $\triangleright$  q),  $\neg$ unicast (sip  $\xi$ ), ( $\xi$ , q))  $\in$  seqp_sos  $\Gamma$ "
| sendT: "(( $\xi$ , {l}send(smsg).p), send (smsg  $\xi$ ), ( $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "
| deliverT: "(( $\xi$ , {l}deliver(sdata).p), deliver (sdata  $\xi$ ), ( $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "
| receiveT: "(( $\xi$ , {l}receive(umsg).p), receive msg, (umsg msg  $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "
| assignT: "(( $\xi$ , {l}[u] p),  $\tau$ , (u  $\xi$ , p))  $\in$  seqp_sos  $\Gamma$ "

| callT: "[[ ( $\xi$ ,  $\Gamma$  pn), a, ( $\xi'$ , p')  $\in$  seqp_sos  $\Gamma$  ]  $\Longrightarrow$ 
  (( $\xi$ , call(pn)), a, ( $\xi'$ , p'))  $\in$  seqp_sos  $\Gamma$ ]"

| choiceT1: "(( $\xi$ , p), a, ( $\xi'$ , p'))  $\in$  seqp_sos  $\Gamma$   $\Longrightarrow$  (( $\xi$ , p  $\oplus$  q), a, ( $\xi'$ , p'))  $\in$  seqp_sos  $\Gamma$ "
| choiceT2: "(( $\xi$ , q), a, ( $\xi'$ , q'))  $\in$  seqp_sos  $\Gamma$   $\Longrightarrow$  (( $\xi$ , p  $\oplus$  q), a, ( $\xi'$ , q'))  $\in$  seqp_sos  $\Gamma$ "

| guardT: " $\xi' \in g \xi \Longrightarrow$  (( $\xi$ , {l}<g> p),  $\tau$ , ( $\xi'$ , p))  $\in$  seqp_sos  $\Gamma$ "
```

```
inductive_cases
```

```
seqp_callTE [elim]: "(( $\xi$ , call(pn)), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ "
and seqp_choiceTE [elim]: "(( $\xi$ , p1  $\oplus$  p2), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ "
```

```
lemma seqp_broadcastTE [elim]:
```

```
"[ (( $\xi$ , {l}broadcast(smsg). p), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ ;
  [a = broadcast (smsg  $\xi$ );  $\xi' = \xi$ ; q = p]  $\Longrightarrow$  P ]  $\Longrightarrow$  P"
by (ind_cases "(( $\xi$ , {l}broadcast(smsg). p), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ ") simp
```

```
lemma seqp_groupcastTE [elim]:
```

```
"[ (( $\xi$ , {l}groupcast(sips, smsg). p), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ ;
  [a = groupcast (sips  $\xi$ ) (smsg  $\xi$ );  $\xi' = \xi$ ; q = p]  $\Longrightarrow$  P ]  $\Longrightarrow$  P"
by (ind_cases "(( $\xi$ , {l}groupcast(sips, smsg). p), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ ") simp
```

```
lemma seqp_unicastTE [elim]:
```

```
"[ (( $\xi$ , {l}unicast(sip, smsg). p  $\triangleright$  q), a, ( $\xi'$ , r))  $\in$  seqp_sos  $\Gamma$ ;
  [a = unicast (sip  $\xi$ ) (smsg  $\xi$ );  $\xi' = \xi$ ; r = p]  $\Longrightarrow$  P;
  [a =  $\neg$ unicast (sip  $\xi$ );  $\xi' = \xi$ ; r = q]  $\Longrightarrow$  P ]  $\Longrightarrow$  P"
by (ind_cases "(( $\xi$ , {l}unicast(sip, smsg). p  $\triangleright$  q), a, ( $\xi'$ , r))  $\in$  seqp_sos  $\Gamma$ ") simp_all
```

```
lemma seqp_sendTE [elim]:
```

```
"[ (( $\xi$ , {l}send(smsg). p), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ ;
  [a = send (smsg  $\xi$ );  $\xi' = \xi$ ; q = p]  $\Longrightarrow$  P ]  $\Longrightarrow$  P"
by (ind_cases "(( $\xi$ , {l}send(smsg). p), a, ( $\xi'$ , q))  $\in$  seqp_sos  $\Gamma$ ") simp
```

```
lemma seqp_deliverTE [elim]:
```



$\llbracket ((\xi, \{l\}deliver(s_{data}). p), a, (\xi', q)) \in seqp\_sos \Gamma; \llbracket a = deliver(s_{data} \xi); \xi' = \xi; q = p \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((\xi, \{l\}deliver(s_{data}). p), a, (\xi', q)) \in seqp\_sos \Gamma$ ") simp

lemma seqp\_receiveTE [elim]:

$\llbracket ((\xi, \{l\}receive(u_{msg}). p), a, (\xi', q)) \in seqp\_sos \Gamma; \bigwedge msg. \llbracket a = receive \ msg; \xi' = u_{msg} \ msg \ \xi; q = p \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((\xi, \{l\}receive(u_{msg}). p), a, (\xi', q)) \in seqp\_sos \Gamma$ ") simp

lemma seqp\_assignTE [elim]:

$\llbracket ((\xi, \{l\}[u] p), a, (\xi', q)) \in seqp\_sos \Gamma; \llbracket a = \tau; \xi' = u \ \xi; q = p \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((\xi, \{l\}[u] p), a, (\xi', q)) \in seqp\_sos \Gamma$ ") simp

lemma seqp\_guardTE [elim]:

$\llbracket ((\xi, \{l\}\langle g \rangle p), a, (\xi', q)) \in seqp\_sos \Gamma; \llbracket a = \tau; \xi' \in g \ \xi; q = p \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((\xi, \{l\}\langle g \rangle p), a, (\xi', q)) \in seqp\_sos \Gamma$ ") simp

lemmas seqpTEs =

seqp\_broadcastTE  
 seqp\_groupcastTE  
 seqp\_unicastTE  
 seqp\_sendTE  
 seqp\_deliverTE  
 seqp\_receiveTE  
 seqp\_assignTE  
 seqp\_callTE  
 seqp\_choiceTE  
 seqp\_guardTE

declare seqp\_sos.intros [intro]

## 6.2 Table 2: Structural operational semantics for parallel process expressions

inductive\_set

parp\_sos :: "('s1, 'm seq\_action) transition set  
 $\implies$  ('s2, 'm seq\_action) transition set  
 $\implies$  ('s1  $\times$  's2, 'm seq\_action) transition set"

for S :: "('s1, 'm seq\_action) transition set"

and T :: "('s2, 'm seq\_action) transition set"

where

parleft: " $\llbracket (s, a, s') \in S; \bigwedge m. a \neq receive \ m \rrbracket \implies ((s, t), a, (s', t)) \in parp\_sos \ S \ T$ "  
 / parright: " $\llbracket (t, a, t') \in T; \bigwedge m. a \neq send \ m \rrbracket \implies ((s, t), a, (s, t')) \in parp\_sos \ S \ T$ "  
 / parboth: " $\llbracket (s, receive \ m, s') \in S; (t, send \ m, t') \in T \rrbracket \implies ((s, t), \tau, (s', t')) \in parp\_sos \ S \ T$ "

lemma par\_broadcastTE [elim]:

$\llbracket ((s, t), broadcast \ m, (s', t')) \in parp\_sos \ S \ T; \llbracket (s, broadcast \ m, s') \in S; t' = t \rrbracket \implies P; \llbracket (t, broadcast \ m, t') \in T; s' = s \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((s, t), broadcast \ m, (s', t')) \in parp\_sos \ S \ T$ ") simp\_all

lemma par\_groupcastTE [elim]:

$\llbracket ((s, t), groupcast \ ips \ m, (s', t')) \in parp\_sos \ S \ T; \llbracket (s, groupcast \ ips \ m, s') \in S; t' = t \rrbracket \implies P; \llbracket (t, groupcast \ ips \ m, t') \in T; s' = s \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((s, t), groupcast \ ips \ m, (s', t')) \in parp\_sos \ S \ T$ ") simp\_all

lemma par\_unicastTE [elim]:

$\llbracket ((s, t), unicast \ i \ m, (s', t')) \in parp\_sos \ S \ T; \llbracket (s, unicast \ i \ m, s') \in S; t' = t \rrbracket \implies P; \llbracket (t, unicast \ i \ m, t') \in T; s' = s \rrbracket \implies P \rrbracket \implies P$   
 by (ind\_cases " $((s, t), unicast \ i \ m, (s', t')) \in parp\_sos \ S \ T$ ") simp\_all

lemma par\_notunicastTE [elim]:

```

"[(s, t), notunicast i, (s', t')] ∈ parp_sos S T;
  [(s, notunicast i, s') ∈ S; t' = t] ⇒ P;
  [(t, notunicast i, t') ∈ T; s' = s] ⇒ P] ⇒ P"
by (ind_cases "(s, t), notunicast i, (s', t') ∈ parp_sos S T") simp_all

lemma par_sendTE [elim]:
"[(s, t), send m, (s', t')] ∈ parp_sos S T;
  [(s, send m, s') ∈ S; t' = t] ⇒ P] ⇒ P"
by (ind_cases "(s, t), send m, (s', t') ∈ parp_sos S T") auto

lemma par_deliverTE [elim]:
"[(s, t), deliver d, (s', t')] ∈ parp_sos S T;
  [(s, deliver d, s') ∈ S; t' = t] ⇒ P;
  [(t, deliver d, t') ∈ T; s' = s] ⇒ P] ⇒ P"
by (ind_cases "(s, t), deliver d, (s', t') ∈ parp_sos S T") simp_all

lemma par_receiveTE [elim]:
"[(s, t), receive m, (s', t')] ∈ parp_sos S T;
  [(t, receive m, t') ∈ T; s' = s] ⇒ P] ⇒ P"
by (ind_cases "(s, t), receive m, (s', t') ∈ parp_sos S T") auto

inductive_cases par_tauTE: "(s, t), τ, (s', t') ∈ parp_sos S T"

lemmas parpTEs =
  par_broadcastTE
  par_groupcastTE
  par_unicastTE
  par_notunicastTE
  par_sendTE
  par_deliverTE
  par_receiveTE

lemma parp_sos_cases [elim]:
  assumes "(s, t), a, (s', t') ∈ parp_sos S T"
    and "[ (s, a, s') ∈ S; ∧m. a ≠ receive m; t' = t ] ⇒ P"
    and "[ (t, a, t') ∈ T; ∧m. a ≠ send m; s' = s ] ⇒ P"
    and "∧m. [ (s, receive m, s') ∈ S; (t, send m, t') ∈ T ] ⇒ P"
  shows "P"
  using assms by cases auto

definition
  par_comp :: "('s1, 'm seq_action) automaton
    ⇒ ('s2, 'm seq_action) automaton
    ⇒ ('s1 × 's2, 'm seq_action) automaton"
  ("(_ << _)" [102, 103] 102)
where
  "s << t ≡ (| init = init s × init t, trans = parp_sos (trans s) (trans t) |)"

lemma trans_par_comp [simp]:
  "trans (s << t) = parp_sos (trans s) (trans t)"
  unfolding par_comp_def by simp

lemma init_par_comp [simp]:
  "init (s << t) = init s × init t"
  unfolding par_comp_def by simp



### 6.3 Table 3: Structural operational semantics for node expressions



inductive_set
  node_sos :: "('s, 'm seq_action) transition set ⇒ ('s net_state, 'm node_action) transition set"
  for S :: "('s, 'm seq_action) transition set"
where
  node_bcast:
    "(s, broadcast m, s') ∈ S ⇒ (NodeS i s R, R:*cast(m), NodeS i s' R) ∈ node_sos S"

```

```

/ node_gcast:
"(s, groupcast D m, s') ∈ S ⇒ (NodeS i s R, (R∩D):*cast(m), NodeS i s' R) ∈ node_sos S"
/ node_ucast:
"[(s, unicast d m, s') ∈ S; d∈R] ⇒ (NodeS i s R, {d}:*cast(m), NodeS i s' R) ∈ node_sos S"
/ node_notucast:
"[(s, ¬unicast d, s') ∈ S; d∉R] ⇒ (NodeS i s R, τ, NodeS i s' R) ∈ node_sos S"
/ node_deliver:
"(s, deliver d, s') ∈ S ⇒ (NodeS i s R, i:deliver(d), NodeS i s' R) ∈ node_sos S"
/ node_receive:
"(s, receive m, s') ∈ S ⇒ (NodeS i s R, {i}¬{i}:arrive(m), NodeS i s' R) ∈ node_sos S"
/ node_tau:
"(s, τ, s') ∈ S ⇒ (NodeS i s R, τ, NodeS i s' R) ∈ node_sos S"
/ node_arrive:
"(NodeS i s R, {i}¬{i}:arrive(m), NodeS i s R) ∈ node_sos S"
/ node_connect1:
"(NodeS i s R, connect(i, i'), NodeS i s (R ∪ {i'})) ∈ node_sos S"
/ node_connect2:
"(NodeS i s R, connect(i', i), NodeS i s (R ∪ {i'})) ∈ node_sos S"
/ node_disconnect1:
"(NodeS i s R, disconnect(i, i'), NodeS i s (R - {i'})) ∈ node_sos S"
/ node_disconnect2:
"(NodeS i s R, disconnect(i', i), NodeS i s (R - {i'})) ∈ node_sos S"
/ node_connect_other:
"[(i ≠ i'; i ≠ i'') ⇒ (NodeS i s R, connect(i', i''), NodeS i s R) ∈ node_sos S"
/ node_disconnect_other:
"[(i ≠ i'; i ≠ i'') ⇒ (NodeS i s R, disconnect(i', i''), NodeS i s R) ∈ node_sos S"

inductive_cases node_arriveTE: "(NodeS i s R, ii¬ni:arrive(m), NodeS i s' R) ∈ node_sos S"
and node_arriveTE': "(NodeS i s R, H¬K:arrive(m), s') ∈ node_sos S"
and node_castTE: "(NodeS i s R, RM:*cast(m), NodeS i s' R') ∈ node_sos S"
and node_castTE': "(NodeS i s R, RM:*cast(m), s') ∈ node_sos S"
and node_deliverTE: "(NodeS i s R, i:deliver(d), NodeS i s' R) ∈ node_sos S"
and node_deliverTE': "(s, i:deliver(d), s') ∈ node_sos S"
and node_deliverTE'': "(NodeS ii s R, i:deliver(d), s') ∈ node_sos S"
and node_tauTE: "(NodeS i s R, τ, NodeS i s' R) ∈ node_sos S"
and node_tauTE': "(NodeS i s R, τ, s') ∈ node_sos S"
and node_connectTE: "(NodeS ii s R, connect(i, i'), NodeS ii s' R') ∈ node_sos S"
and node_connectTE': "(NodeS ii s R, connect(i, i'), s') ∈ node_sos S"
and node_disconnectTE: "(NodeS ii s R, disconnect(i, i'), NodeS ii s' R') ∈ node_sos S"
and node_disconnectTE': "(NodeS ii s R, disconnect(i, i'), s') ∈ node_sos S"

lemma node_sos_never_newpkt [simp]:
assumes "(s, a, s') ∈ node_sos S"
shows "a ≠ i:newpkt(d, di)"
using assms by cases auto

lemma arrives_or_not:
assumes "(NodeS i s R, ii¬ni:arrive(m), NodeS i' s' R') ∈ node_sos S"
shows "(ii = {i} ∧ ni = {}) ∨ (ii = {} ∧ ni = {i})"
using assms by rule simp_all

definition
node_comp :: "ip ⇒ ('s, 'm seq_action) automaton ⇒ ip set
⇒ ('s net_state, 'm node_action) automaton"
("⟨_ : _⟩" [0, 0, 0] 104)
where
"⟨i : np : R_i⟩ ≡ (| init = {NodeS i s R_i | s. s ∈ init np}, trans = node_sos (trans np) |)"

lemma trans_node_comp:
"trans (⟨i : np : R_i⟩) = node_sos (trans np)"
unfolding node_comp_def by simp

lemma init_node_comp:
"init (⟨i : np : R_i⟩) = {NodeS i s R_i | s. s ∈ init np}"

```

unfolding node\_comp\_def by simp

lemmas node\_comps = trans\_node\_comp init\_node\_comp

lemma trans\_par\_node\_comp [simp]:

"trans ((i : s << t : R)) = node\_sos (parp\_sos (trans s) (trans t))"  
unfolding node\_comp\_def by simp

lemma snd\_par\_node\_comp [simp]:

"init ((i : s << t : R)) = {NodeS i st R | st. st ∈ init s × init t}"  
unfolding node\_comp\_def by simp

lemma node\_sos\_dest\_is\_net\_state:

assumes "(s, a, s') ∈ node\_sos S"  
shows "∃ i' P' R'. s' = NodeS i' P' R'"  
using assms by induct auto

lemma node\_sos\_dest:

assumes "(NodeS i p R, a, s') ∈ node\_sos S"  
shows "∃ P' R'. s' = NodeS i P' R'"  
using assms assms [THEN node\_sos\_dest\_is\_net\_state]  
by - (erule node\_sos.cases, auto)

lemma node\_sos\_states [elim]:

assumes "(ns, a, ns') ∈ node\_sos S"  
obtains i s R s' R' where "ns = NodeS i s R"  
and "ns' = NodeS i s' R'"

proof -

assume [intro!]: " $\bigwedge i s R s' R'. ns = NodeS i s R \implies ns' = NodeS i s' R' \implies thesis$ "  
from assms(1) obtain i s R where "ns = NodeS i s R"  
by (cases ns) auto  
moreover with assms(1) obtain s' R' where "ns' = NodeS i s' R'"  
by (metis node\_sos\_dest)  
ultimately show thesis ..  
qed

lemma node\_sos\_cases [elim]:

"(NodeS i p R, a, NodeS i p' R') ∈ node\_sos S  $\implies$   
 $(\bigwedge m . \quad \llbracket a = R.*cast(m); \quad R' = R; (p, broadcast\ m, p') \in S \rrbracket \implies P) \implies$   
 $(\bigwedge m D. \quad \llbracket a = (R \cap D).*cast(m); \quad R' = R; (p, groupcast\ D\ m, p') \in S \rrbracket \implies P) \implies$   
 $(\bigwedge d m. \quad \llbracket a = \{d\}.*cast(m); \quad R' = R; (p, unicast\ d\ m, p') \in S; d \in R \rrbracket \implies P) \implies$   
 $(\bigwedge d. \quad \llbracket a = \tau; \quad R' = R; (p, \neg unicast\ d, p') \in S; d \notin R \rrbracket \implies P) \implies$   
 $(\bigwedge d. \quad \llbracket a = i:deliver(d); \quad R' = R; (p, deliver\ d, p') \in S \rrbracket \implies P) \implies$   
 $(\bigwedge m. \quad \llbracket a = \{i\}\neg\{j\}:arrive(m); \quad R' = R; (p, receive\ m, p') \in S \rrbracket \implies P) \implies$   
 $(\quad \llbracket a = \tau; \quad R' = R; (p, \tau, p') \in S \rrbracket \implies P) \implies$   
 $(\bigwedge m. \quad \llbracket a = \{j\}\neg\{i\}:arrive(m); \quad R' = R; p = p' \rrbracket \implies P) \implies$   
 $(\bigwedge i i'. \quad \llbracket a = connect(i, i'); \quad R' = R \cup \{i'\}; p = p' \rrbracket \implies P) \implies$   
 $(\bigwedge i i'. \quad \llbracket a = connect(i', i); \quad R' = R \cup \{i'\}; p = p' \rrbracket \implies P) \implies$   
 $(\bigwedge i i'. \quad \llbracket a = disconnect(i, i'); \quad R' = R - \{i'\}; p = p' \rrbracket \implies P) \implies$   
 $(\bigwedge i i'. \quad \llbracket a = disconnect(i', i); \quad R' = R - \{i'\}; p = p' \rrbracket \implies P) \implies$   
 $(\bigwedge i i' i''. \quad \llbracket a = connect(i', i''); \quad R' = R; p = p'; i \neq i'; i \neq i'' \rrbracket \implies P) \implies$   
 $(\bigwedge i i' i''. \quad \llbracket a = disconnect(i', i''); \quad R' = R; p = p'; i \neq i'; i \neq i'' \rrbracket \implies P) \implies$   
P"  
by (erule node\_sos.cases) simp\_all

## 6.4 Table 4: Structural operational semantics for partial network expressions

inductive\_set

pnet\_sos :: "('s net\_state, 'm node\_action) transition set  
 $\implies$  ('s net\_state, 'm node\_action) transition set  
 $\implies$  ('s net\_state, 'm node\_action) transition set"

for S :: "('s net\_state, 'm node\_action) transition set"

and T :: "('s net\_state, 'm node\_action) transition set"

where

```

pnet_cast1: "[ (s, R:*cast(m), s') ∈ S; (t, H¬K:arrive(m), t') ∈ T; H ⊆ R; K ∩ R = {} ]
⇒ (SubnetS s t, R:*cast(m), SubnetS s' t') ∈ pnet_sos S T"

/ pnet_cast2: "[ (s, H¬K:arrive(m), s') ∈ S; (t, R:*cast(m), t') ∈ T; H ⊆ R; K ∩ R = {} ]
⇒ (SubnetS s t, R:*cast(m), SubnetS s' t') ∈ pnet_sos S T"

/ pnet_arrive: "[ (s, H¬K:arrive(m), s') ∈ S; (t, H'¬K':arrive(m), t') ∈ T ]
⇒ (SubnetS s t, (H ∪ H')¬(K ∪ K'):arrive(m), SubnetS s' t') ∈ pnet_sos S T"

/ pnet_deliver1: "(s, i:deliver(d), s') ∈ S
⇒ (SubnetS s t, i:deliver(d), SubnetS s' t) ∈ pnet_sos S T"
/ pnet_deliver2: "[ (t, i:deliver(d), t') ∈ T ]
⇒ (SubnetS s t, i:deliver(d), SubnetS s' t') ∈ pnet_sos S T"

/ pnet_tau1: "(s, τ, s') ∈ S ⇒ (SubnetS s t, τ, SubnetS s' t) ∈ pnet_sos S T"
/ pnet_tau2: "(t, τ, t') ∈ T ⇒ (SubnetS s t, τ, SubnetS s' t') ∈ pnet_sos S T"

/ pnet_connect: "[ (s, connect(i, i'), s') ∈ S; (t, connect(i, i'), t') ∈ T ]
⇒ (SubnetS s t, connect(i, i'), SubnetS s' t') ∈ pnet_sos S T"

/ pnet_disconnect: "[ (s, disconnect(i, i'), s') ∈ S; (t, disconnect(i, i'), t') ∈ T ]
⇒ (SubnetS s t, disconnect(i, i'), SubnetS s' t') ∈ pnet_sos S T"

inductive_cases partial_castTE [elim]: "(s, R:*cast(m), s') ∈ pnet_sos S T"
and partial_arriveTE [elim]: "(s, H¬K:arrive(m), s') ∈ pnet_sos S T"
and partial_deliverTE [elim]: "(s, i:deliver(d), s') ∈ pnet_sos S T"
and partial_tauTE [elim]: "(s, τ, s') ∈ pnet_sos S T"
and partial_connectTE [elim]: "(s, connect(i, i'), s') ∈ pnet_sos S T"
and partial_disconnectTE [elim]: "(s, disconnect(i, i'), s') ∈ pnet_sos S T"

lemma pnet_sos_never_newpkt:
assumes "(st, a, st') ∈ pnet_sos S T"
and "∧i d di a s s'. (s, a, s') ∈ S ⇒ a ≠ i:newpkt(d, di)"
and "∧i d di a t t'. (t, a, t') ∈ T ⇒ a ≠ i:newpkt(d, di)"
shows "a ≠ i:newpkt(d, di)"
using assms(1) by cases (auto dest!: assms(2-3))

fun pnet :: "('s, 'm seq_action) automaton
⇒ net_tree ⇒ ('s net_state, 'm node_action) automaton"

where
"pnet np (<i; R_i) = <i : np i : R_i>"
/ "pnet np (p1 || p2) = (| init = {SubnetS s1 s2 | s1 s2. s1 ∈ init (pnet np p1)
∧ s2 ∈ init (pnet np p2)},
trans = pnet_sos (trans (pnet np p1)) (trans (pnet np p2)) |)"

lemma pnet_node_init [elim, simp]:
assumes "s ∈ init (pnet np <i; R>)"
shows "s ∈ {NodeS i s R | s. s ∈ init (np i)}"
using assms by (simp add: node_comp_def)

lemma pnet_node_init' [elim]:
assumes "s ∈ init (pnet np <i; R>)"
obtains ns where "s = NodeS i ns R"
and "ns ∈ init (np i)"
using assms by (auto simp add: node_comp_def)

lemma pnet_node_trans [elim, simp]:
assumes "(s, a, s') ∈ trans (pnet np <i; R>)"
shows "(s, a, s') ∈ node_sos (trans (np i))"
using assms by (simp add: trans_node_comp)

lemma pnet_never_newpkt':
assumes "(s, a, s') ∈ trans (pnet np n)"
shows "∀i d di. a ≠ i:newpkt(d, di)"

```

```

using assms proof (induction n arbitrary: s a s')
  fix n1 n2 s a s'
  assume IH1: " $\bigwedge s a s'. (s, a, s') \in \text{trans (pnet np n1)} \implies \forall i d di. a \neq i:\text{newpkt}(d, di)$ "
    and IH2: " $\bigwedge s a s'. (s, a, s') \in \text{trans (pnet np n2)} \implies \forall i d di. a \neq i:\text{newpkt}(d, di)$ "
    and "(s, a, s')  $\in \text{trans (pnet np (n1 || n2))}$ "
  show " $\forall i d di. a \neq i:\text{newpkt}(d, di)$ "
  proof (intro allI)
    fix i d di
    from  $\langle (s, a, s') \in \text{trans (pnet np (n1 || n2))} \rangle$ 
      have "(s, a, s')  $\in \text{pnet\_sos (trans (pnet np n1)) (trans (pnet np n2))}$ "
        by simp
    thus "a  $\neq i:\text{newpkt}(d, di)$ "
      by (rule pnet_sos_never_newpkt) (auto dest!: IH1 IH2)
  qed
qed (simp add: node_comps)

```

```

lemma pnet_never_newpkt:
  assumes "(s, a, s')  $\in \text{trans (pnet np n)}$ "
  shows "a  $\neq i:\text{newpkt}(d, di)$ "
  proof -
    from assms have " $\forall i d di. a \neq i:\text{newpkt}(d, di)$ "
      by (rule pnet_never_newpkt')
    thus ?thesis by clarsimp
  qed

```

## 6.5 Table 5: Structural operational semantics for complete network expressions

inductive\_set

```

cnet_sos :: "('s, ('m::msg) node_action) transition set
           $\Rightarrow$  ('s, 'm node_action) transition set"

```

```

for S :: "('s, 'm node_action) transition set"

```

where

```

  cnet_connect: "(s, connect(i, i'), s')  $\in S \implies (s, \text{connect}(i, i'), s') \in \text{cnet\_sos } S"$ 
| cnet_disconnect: "(s, disconnect(i, i'), s')  $\in S \implies (s, \text{disconnect}(i, i'), s') \in \text{cnet\_sos } S"$ 
| cnet_cast: "(s, R:*cast(m), s')  $\in S \implies (s, \tau, s') \in \text{cnet\_sos } S"$ 
| cnet_tau: "(s,  $\tau$ , s')  $\in S \implies (s, \tau, s') \in \text{cnet\_sos } S"$ 
| cnet_deliver: "(s, i:deliver(d), s')  $\in S \implies (s, i:\text{deliver}(d), s') \in \text{cnet\_sos } S"$ 
| cnet_newpkt: "(s, {i}¬K:arrive(newpkt(d, di)), s')  $\in S \implies (s, i:\text{newpkt}(d, di), s') \in \text{cnet\_sos } S"$ 

```

```

inductive_cases connect_completeTE: "(s, connect(i, i'), s')  $\in \text{cnet\_sos } S"$ 
and disconnect_completeTE: "(s, disconnect(i, i'), s')  $\in \text{cnet\_sos } S"$ 
and tau_completeTE: "(s,  $\tau$ , s')  $\in \text{cnet\_sos } S"$ 
and deliver_completeTE: "(s, i:deliver(d), s')  $\in \text{cnet\_sos } S"$ 
and newpkt_completeTE: "(s, i:newpkt(d, di), s')  $\in \text{cnet\_sos } S"$ 

```

```

lemmas completeTEs = connect_completeTE
                    disconnect_completeTE
                    tau_completeTE
                    deliver_completeTE
                    newpkt_completeTE

```

lemma complete\_no\_cast [simp]:

```

"(s, R:*cast(m), s')  $\notin \text{cnet\_sos } T"$ 

```

proof

```

  assume "(s, R:*cast(m), s')  $\in \text{cnet\_sos } T"$ 

```

```

  hence "R:*cast(m)  $\neq R:*cast(m)"$ 

```

```

  by (rule cnet_sos.cases) auto

```

```

  thus False by simp

```

qed

lemma complete\_no\_arrive [simp]:

```

"(s, ii¬ni:arrive(m), s')  $\notin \text{cnet\_sos } T"$ 

```

proof

```

  assume "(s, ii¬ni:arrive(m), s')  $\in \text{cnet\_sos } T"$ 

```

```

hence "ii¬ni:arrive(m) ≠ ii¬ni:arrive(m)"
  by (rule cnet_sos.cases) auto
thus False by simp
qed

```

abbreviation

```

closed :: "('s net_state, ('m::msg) node_action) automaton ⇒ ('s net_state, 'm node_action) automaton"
where
  "closed ≡ (λA. A (| trans := cnet_sos (trans A) |))"

```

end

## 7 Control terms and well-definedness of sequential processes

```

theory Awn_Cterms
imports Awn
begin

```

### 7.1 Microsteps

We distinguish microsteps from ‘external’ transitions (observable or not). Here, they are a kind of ‘hypothetical computation’, since, unlike  $\tau$ -transitions, they do not make choices but rather ‘compute’ which choices are possible.

inductive

```

microstep :: "('s, 'm, 'p, 'l) seqp_env
  ⇒ ('s, 'm, 'p, 'l) seqp
  ⇒ ('s, 'm, 'p, 'l) seqp
  ⇒ bool"

```

```

for  $\Gamma$  :: "('s, 'm, 'p, 'l) seqp_env"

```

where

```

  microstep_choiceI1 [intro, simp]: "microstep  $\Gamma$  (p1  $\oplus$  p2) p1"
  | microstep_choiceI2 [intro, simp]: "microstep  $\Gamma$  (p1  $\oplus$  p2) p2"
  | microstep_callI [intro, simp]: "microstep  $\Gamma$  (call(pn)) ( $\Gamma$  pn)"

```

abbreviation microstep\_rtcl

```

where "microstep_rtcl  $\Gamma$  p q ≡ (microstep  $\Gamma$ )** p q"

```

abbreviation microstep\_tcl

```

where "microstep_tcl  $\Gamma$  p q ≡ (microstep  $\Gamma$ )++ p q"

```

syntax

```

"_microstep"
  :: "[('s, 'm, 'p, 'l) seqp, ('s, 'm, 'p, 'l) seqp_env, ('s, 'm, 'p, 'l) seqp] ⇒ bool"
  ("(_  $\rightsquigarrow$ _ (_)" [61, 0, 61] 50)
"_microstep_rtcl"
  :: "[('s, 'm, 'p, 'l) seqp, ('s, 'm, 'p, 'l) seqp_env, ('s, 'm, 'p, 'l) seqp] ⇒ bool"
  ("(_  $\rightsquigarrow$ * (_)" [61, 0, 61] 50)
"_microstep_tcl"
  :: "[('s, 'm, 'p, 'l) seqp, ('s, 'm, 'p, 'l) seqp_env, ('s, 'm, 'p, 'l) seqp] ⇒ bool"
  ("(_  $\rightsquigarrow$ + (_)" [61, 0, 61] 50)

```

translations

```

"p1  $\rightsquigarrow$  $\Gamma$  p2" ≡ "CONST microstep  $\Gamma$  p1 p2"
"p1  $\rightsquigarrow$  $\Gamma$ * p2" ≡ "CONST microstep_rtcl  $\Gamma$  p1 p2"
"p1  $\rightsquigarrow$  $\Gamma$ + p2" ≡ "CONST microstep_tcl  $\Gamma$  p1 p2"

```

lemma microstep\_choiceD [dest]:

```

"(p1  $\oplus$  p2)  $\rightsquigarrow$  $\Gamma$  p  $\implies$  p = p1  $\vee$  p = p2"
by (ind_cases "(p1  $\oplus$  p2)  $\rightsquigarrow$  $\Gamma$  p") auto

```

lemma microstep\_choiceE [elim]:

```

"[ (p1  $\oplus$  p2)  $\rightsquigarrow$  $\Gamma$  p;
  (p1  $\oplus$  p2)  $\rightsquigarrow$  $\Gamma$  p1  $\implies$  P;
  (p1  $\oplus$  p2)  $\rightsquigarrow$  $\Gamma$  p2  $\implies$  P ]  $\implies$  P"

```

```

by (blast)

lemma microstep_callD [dest]:
  "(call(pn))  $\rightsquigarrow_{\Gamma}$  p  $\implies$  p =  $\Gamma$  pn"
  by (ind_cases "(call(pn))  $\rightsquigarrow_{\Gamma}$  p")

lemma microstep_callE [elim]:
  "[[ (call(pn))  $\rightsquigarrow_{\Gamma}$  p; p =  $\Gamma$ (pn)  $\implies$  P ]  $\implies$  P"
  by auto

lemma no_microstep_guard: " $\neg$  (({1}⟨g⟩ p)  $\rightsquigarrow_{\Gamma}$  q)"
  by (rule notI) (ind_cases "{1}⟨g⟩ p  $\rightsquigarrow_{\Gamma}$  q")

lemma no_microstep_assign: " $\neg$  ({1}[[f]] p)  $\rightsquigarrow_{\Gamma}$  q"
  by (rule notI) (ind_cases "{1}[[f]] p  $\rightsquigarrow_{\Gamma}$  q")

lemma no_microstep_unicast: " $\neg$  ({1}unicast(sip, smsg).p ▷ q)  $\rightsquigarrow_{\Gamma}$  r)"
  by (rule notI) (ind_cases "{1}unicast(sip, smsg).p ▷ q  $\rightsquigarrow_{\Gamma}$  r")

lemma no_microstep_broadcast: " $\neg$  ({1}broadcast(smsg).p)  $\rightsquigarrow_{\Gamma}$  q)"
  by (rule notI) (ind_cases "{1}broadcast(smsg).p  $\rightsquigarrow_{\Gamma}$  q")

lemma no_microstep_groupcast: " $\neg$  ({1}groupcast(sips, smsg).p)  $\rightsquigarrow_{\Gamma}$  q)"
  by (rule notI) (ind_cases "{1}groupcast(sips, smsg).p  $\rightsquigarrow_{\Gamma}$  q")

lemma no_microstep_send: " $\neg$  ({1}send(smsg).p)  $\rightsquigarrow_{\Gamma}$  q)"
  by (rule notI) (ind_cases "{1}send(smsg).p  $\rightsquigarrow_{\Gamma}$  q")

lemma no_microstep_deliver: " $\neg$  ({1}deliver(sdata).p)  $\rightsquigarrow_{\Gamma}$  q)"
  by (rule notI) (ind_cases "{1}deliver(sdata).p  $\rightsquigarrow_{\Gamma}$  q")

lemma no_microstep_receive: " $\neg$  ({1}receive(umsg).p)  $\rightsquigarrow_{\Gamma}$  q)"
  by (rule notI) (ind_cases "{1}receive(umsg).p  $\rightsquigarrow_{\Gamma}$  q")

lemma microstep_call_or_choice [dest]:
  assumes "p  $\rightsquigarrow_{\Gamma}$  q"
  shows "( $\exists$ pn. p = call(pn))  $\vee$  ( $\exists$ p1 p2. p = p1  $\oplus$  p2)"
  using assms by clarsimp (metis microstep.simps)

lemmas no_microstep [intro,simp] =
  no_microstep_guard
  no_microstep_assign
  no_microstep_unicast
  no_microstep_broadcast
  no_microstep_groupcast
  no_microstep_send
  no_microstep_deliver
  no_microstep_receive

```

## 7.2 Wellformed process specifications

A process specification  $\Gamma$  is wellformed if its *microstep*  $\Gamma$  relation is free of loops and infinite chains.

For example, these specifications are not wellformed:  $\Gamma_1$   $p1 = \text{call}(p1)$

$\Gamma_2$   $p1 = \text{send}(msg) . \text{call}(p1) \oplus \text{call}(p1)$

$\Gamma_3$   $p1 = \text{send}(msg) . \text{call}(p2)$   $\Gamma_3$   $p2 = \text{call}(p3)$   $\Gamma_3$   $p3 = \text{call}(p4)$   $\Gamma_3$   $p4 = \text{call}(p5) \dots$

**definition**

```
wellformed :: "('s, 'm, 'p, 'l) seqp_env  $\implies$  bool"
```

where

```
"wellformed  $\Gamma$  = wf {(q, p). p  $\rightsquigarrow_{\Gamma}$  q}"
```

```
lemma wellformed_defP: "wellformed  $\Gamma$  = wfP ( $\lambda$ q p. p  $\rightsquigarrow_{\Gamma}$  q)"
```

```
unfolding wellformed_def wfP_def by simp
```



The induction rule for *wellformed*  $\Gamma$  is stronger than  $\llbracket \bigwedge x1\ x2\ x3. ?P\ x3 \implies ?P\ (\{x1\}\langle x2 \rangle\ x3); \bigwedge x1\ x2\ x3. ?P\ x3 \implies ?P\ (\{x1\}\llbracket x2 \rrbracket\ x3); \bigwedge x1\ x2. \llbracket ?P\ x1; ?P\ x2 \rrbracket \implies ?P\ (x1 \oplus x2); \bigwedge x1\ x2\ x3\ x4\ x5. \llbracket ?P\ x4; ?P\ x5 \rrbracket \implies ?P\ (\{x1\}\text{unicast}(x2, x3) . x4 \triangleright x5); \bigwedge x1\ x2\ x3. ?P\ x3 \implies ?P\ (\{x1\}\text{broadcast}(x2) . x3); \bigwedge x1\ x2\ x3\ x4. ?P\ x4 \implies ?P\ (\{x1\}\text{groupcast}(x2, x3) . x4); \bigwedge x1\ x2\ x3. ?P\ x3 \implies ?P\ (\{x1\}\text{send}(x2) . x3); \bigwedge x1\ x2\ x3. ?P\ x3 \implies ?P\ (\{x1\}\text{deliver}(x2) . x3); \bigwedge x1\ x2\ x3. ?P\ x3 \implies ?P\ (\{x1\}\text{receive}(x2) . x3); \bigwedge x. ?P\ (\text{call}(x)) \rrbracket \implies ?P\ ?\text{seqp}$  because the case for  $\text{call}(pn)$  can be shown with the assumption on  $\Gamma\ pn$ .

**lemma** *wellformed\_induct*

[consumes 1, case\_names ASSIGN CHOICE CALL GUARD UCAST BCAST GCAST SEND DELIVER RECEIVE,  
induct set: wellformed]:

assumes "wellformed  $\Gamma$ "

and ASSIGN: " $\bigwedge l\ f\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\llbracket f \rrbracket\ p)$ "  
and GUARD: " $\bigwedge l\ f\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\langle f \rangle\ p)$ "  
and UCAST: " $\bigwedge l\ fip\ fmsg\ p\ q. \text{wellformed } \Gamma \implies P\ (\{l\}\text{unicast}(fip, fmsg). p \triangleright q)$ "  
and BCAST: " $\bigwedge l\ fmsg\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\text{broadcast}(fmsg). p)$ "  
and GCAST: " $\bigwedge l\ fips\ fmsg\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\text{groupcast}(fips, fmsg). p)$ "  
and SEND: " $\bigwedge l\ fmsg\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\text{send}(fmsg). p)$ "  
and DELIVER: " $\bigwedge l\ fdata\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\text{deliver}(fdata). p)$ "  
and RECEIVE: " $\bigwedge l\ fmsg\ p. \text{wellformed } \Gamma \implies P\ (\{l\}\text{receive}(fmsg). p)$ "  
and CHOICE: " $\bigwedge p1\ p2. \llbracket \text{wellformed } \Gamma; P\ p1; P\ p2 \rrbracket \implies P\ (p1 \oplus p2)$ "  
and CALL: " $\bigwedge pn. \llbracket \text{wellformed } \Gamma; P\ (\Gamma\ pn) \rrbracket \implies P\ (\text{call}(pn))$ "

shows "P a"

using *assms(1) unfolding wellformed\_defP*

**proof** (rule *wfP\_induct\_rule, case\_tac x, simp\_all*)

fix  $p1\ p2$

assume " $\bigwedge q. (p1 \oplus p2) \rightsquigarrow_{\Gamma} q \implies P\ q$ "

then obtain "P  $p1$ " and "P  $p2$ " by (auto *intro!: microstep.intros*)

thus "P  $(p1 \oplus p2)$ " by (rule CHOICE [OF  $\langle \text{wellformed } \Gamma \rangle$ ])

next

fix  $pn$

assume " $\bigwedge q. (\text{call}(pn)) \rightsquigarrow_{\Gamma} q \implies P\ q$ "

hence "P  $(\Gamma\ pn)$ " by (auto *intro!: microstep.intros*)

thus "P  $(\text{call}(pn))$ " by (rule CALL [OF  $\langle \text{wellformed } \Gamma \rangle$ ])

qed (auto *intro: assms*)

### 7.3 Start terms (sterms)

Formulate sets of local subterms from which an action is directly possible. Since the process specification  $\Gamma$  is not considered, only choice terms  $p1 \oplus p2$  are traversed, and not  $\text{call}(p)$  terms.

**fun** *stermsl* :: "(*s*, *m*, *p*, *l*) *seqp*  $\Rightarrow$  (*s*, *m*, *p*, *l*) *seqp set*"

where

"*stermsl*  $(p1 \oplus p2) = \text{stermsl } p1 \cup \text{stermsl } p2$ "  
| "*stermsl*  $p = \{p\}$ "

**lemma** *stermsl\_nobigger*: " $q \in \text{stermsl } p \implies \text{size } q \leq \text{size } p$ "

by (induct  $p$ ) auto

**lemma** *stermsl\_no\_choice*[*simp*]: " $p1 \oplus p2 \notin \text{stermsl } p$ "

by (induct  $p$ ) *simp\_all*

**lemma** *stermsl\_choice\_disj*[*simp*]:

" $p \in \text{stermsl } (p1 \oplus p2) = (p \in \text{stermsl } p1 \vee p \in \text{stermsl } p2)$ "  
by *simp*

**lemma** *stermsl\_in\_branch*[*elim*]:

" $\llbracket p \in \text{stermsl } (p1 \oplus p2); p \in \text{stermsl } p1 \implies P; p \in \text{stermsl } p2 \implies P \rrbracket \implies P$ "  
by auto

**lemma** *stermsl\_commute*:

"*stermsl*  $(p1 \oplus p2) = \text{stermsl } (p2 \oplus p1)$ "  
by *simp* (rule *Un\_commute*)

**lemma** *stermsl\_not\_empty*:

```

"stermsl p ≠ {}"
by (induct p) auto

lemma stermsl_idem [simp]:
"(⋃q∈stermsl p. stermsl q) = stermsl p"
by (induct p) simp_all

lemma stermsl_in_wfpf:
assumes AA: "A ⊆ {(q, p). p ↘Γ q} ‘‘ A"
and *: "p ∈ A"
shows "∃r∈stermsl p. r ∈ A"
using *
proof (induction p)
fix p1 p2
assume IH1: "p1 ∈ A ⇒ ∃r∈stermsl p1. r ∈ A"
and IH2: "p2 ∈ A ⇒ ∃r∈stermsl p2. r ∈ A"
and *: "p1 ⊕ p2 ∈ A"
from * and AA have "p1 ⊕ p2 ∈ {(q, p). p ↘Γ q} ‘‘ A" by auto
hence "p1 ∈ A ∨ p2 ∈ A" by auto
hence "(∃r∈stermsl p1. r ∈ A) ∨ (∃r∈stermsl p2. r ∈ A)"
proof
assume "p1 ∈ A" hence "∃r∈stermsl p1. r ∈ A" by (rule IH1) thus ?thesis ..
next
assume "p2 ∈ A" hence "∃r∈stermsl p2. r ∈ A" by (rule IH2) thus ?thesis ..
qed
hence "∃r∈stermsl p1 ∪ stermsl p2. r ∈ A" by blast
thus "∃r∈stermsl (p1 ⊕ p2). r ∈ A" by simp
next case UCAST from UCAST.prem show ?case by auto
qed auto

lemma nocall_stermsl_max:
assumes "r ∈ stermsl p"
and "not_call r"
shows "¬ (r ↘Γ q)"
using assms
by (induction p) auto

theorem wf_no_direct_calls[intro]:
fixes Γ :: "('s, 'm, 'p, 'l) seqp_env"
assumes no_calls: "∧pn'. call(pn') ∉ stermsl(Γ(pn))"
shows "wellformed Γ"
unfolding wellformed_def wfP_def
proof (rule wfI_pf)
fix A
assume ARA: "A ⊆ {(q, p). p ↘Γ q} ‘‘ A"
hence hasnext: "∧p. p ∈ A ⇒ ∃q. p ↘Γ q ∧ q ∈ A" by auto
show "A = {}"
proof (rule Set.equalsOI)
fix p assume "p ∈ A" thus "False"
proof (induction p)
fix l f p'
assume *: "{l}⟨f⟩ p' ∈ A"
from hasnext [OF *] have "∃q. ({l}⟨f⟩ p') ↘Γ q" by simp
thus "False" by simp
next
fix p1 p2
assume *: "p1 ⊕ p2 ∈ A"
and IH1: "p1 ∈ A ⇒ False"
and IH2: "p2 ∈ A ⇒ False"
have "∃q. (p1 ⊕ p2) ↘Γ q ∧ q ∈ A" by (rule hasnext [OF *])
hence "p1 ∈ A ∨ p2 ∈ A" by auto
thus "False" by (auto dest: IH1 IH2)
next
fix pn

```

```

assume "call(pn) ∈ A"
hence "∃q. (call(pn)) ↘Γ q ∧ q ∈ A" by (rule hasnext)
hence "Γ(pn) ∈ A" by auto

with ARA [THEN stermsl_in_wfpf] obtain q where "q ∈ stermsl (Γ pn)" and "q ∈ A" by metis
hence "not_call q" using no_calls [of pn]
  unfolding not_call_def by auto

from hasnext [OF ⟨q ∈ A⟩] obtain q' where "q ↘Γ q'" by auto
moreover from ⟨q ∈ stermsl (Γ pn)⟩ ⟨not_call q⟩ have "¬ (q ↘Γ q)"
  by (rule nocall_stermsl_max)
ultimately show "False" by simp
qed (auto dest: hasnext)
qed
qed

```

## 7.4 Start terms

The start terms are those terms, relative to a wellformed process specification  $\Gamma$ , from which transitions can occur directly.

```

function (domintros, sequential) sterms
:: "('s, 'm, 'p, 'l) seqp_env ⇒ ('s, 'm, 'p, 'l) seqp ⇒ ('s, 'm, 'p, 'l) seqp set"
where
  sterms_choice: "sterms Γ (p1 ⊕ p2) = sterms Γ p1 ∪ sterms Γ p2"
| sterms_call: "sterms Γ (call(pn)) = sterms Γ (Γ pn)"
| sterms_other: "sterms Γ p = {p}"
by pat_completeness auto

```

```

lemma sterms_dom_basic[simp]:
assumes "not_call p"
  and "not_choice p"
shows "sterms_dom (Γ, p)"
proof (rule accpI)
  fix y
  assume "sterms_rel y (Γ, p)"
  with assms show "sterms_dom y"
  by (cases p) (auto simp: sterms_rel.simps)
qed

```

```

lemma sterms_termination:
assumes "wellformed Γ"
shows "sterms_dom (Γ, p)"
proof -
  have sterms_rel':
    "sterms_rel = (λgq gp. (gq, gp) ∈ {((Γ, q), (Γ', p)). Γ = Γ' ∧ p ↘Γ q})"
  by (rule ext)+ (auto simp: sterms_rel.simps elim: microstep.cases)

  from assms have "∀x. x ∈ Wellfounded.acc {(q, p). p ↘Γ q}"
  unfolding wellformed_def by (simp add: wf_acc_iff)
  hence "p ∈ Wellfounded.acc {(q, p). p ↘Γ q}" ..

  hence "(Γ, p) ∈ Wellfounded.acc {((Γ, q), (Γ', p)). Γ = Γ' ∧ p ↘Γ q}"
  by (rule acc_induct) (auto intro: accI)

  thus "sterms_dom (Γ, p)" unfolding sterms_rel' accp_acc_eq .
qed

```

```

declare sterms.psimps [simp]

```

```

lemmas sterms_psimps[simp] = sterms.psimps [OF sterms_termination]
and sterms_pinduct = sterms.pinduct [OF sterms_termination]

```

```

lemma sterms_reflD [dest]:

```

```

assumes "q ∈ sterms Γ p"
  and "not_choice p" "not_call p"
  shows "q = p"
using assms by (cases p) auto

lemma sterms_choice_disj [simp]:
  assumes "wellformed Γ"
  shows "p ∈ sterms Γ (p1 ⊕ p2) = (p ∈ sterms Γ p1 ∨ p ∈ sterms Γ p2)"
using assms by (simp)

lemma sterms_no_choice [simp]:
  assumes "wellformed Γ"
  shows "p1 ⊕ p2 ∉ sterms Γ p"
using assms by induction auto

lemma sterms_not_choice [simp]:
  assumes "wellformed Γ"
  and "q ∈ sterms Γ p"
  shows "not_choice q"
using assms unfolding not_choice_def
by (auto dest: sterms_no_choice)

lemma sterms_no_call [simp]:
  assumes "wellformed Γ"
  shows "call(pn) ∉ sterms Γ p"
using assms by induction auto

lemma sterms_not_call [simp]:
  assumes "wellformed Γ"
  and "q ∈ sterms Γ p"
  shows "not_call q"
using assms unfolding not_call_def
by (auto dest: sterms_no_call)

lemma sterms_in_branch:
  assumes "wellformed Γ"
  and "p ∈ sterms Γ (p1 ⊕ p2)"
  and "p ∈ sterms Γ p1 ⇒ P"
  and "p ∈ sterms Γ p2 ⇒ P"
  shows "P"
using assms by auto

lemma sterms_commute:
  assumes "wellformed Γ"
  shows "sterms Γ (p1 ⊕ p2) = sterms Γ (p2 ⊕ p1)"
using assms by simp (rule Un_commute)

lemma sterms_not_empty:
  assumes "wellformed Γ"
  shows "sterms Γ p ≠ {}"
using assms
by (induct p rule: sterms_pinduct [OF ⟨wellformed Γ⟩]) simp_all

lemma sterms_sterms [simp]:
  assumes "wellformed Γ"
  shows "(⋃ x ∈ sterms Γ p. sterms Γ x) = sterms Γ p"
using assms by induction simp_all

lemma sterms_stermsl:
  assumes "ps ∈ sterms Γ p"
  and "wellformed Γ"
  shows "ps ∈ stermsl p ∨ (∃ pn. ps ∈ stermsl (Γ pn))"
using assms by (induction p rule: sterms_pinduct [OF ⟨wellformed Γ⟩]) auto

```

```

lemma stermsl_sterms [elim]:
  assumes "q ∈ stermsl p"
    and "not_call q"
    and "wellformed Γ"
  shows "q ∈ sterms Γ p"
  using assms by (induct p) auto

lemma sterms_stermsl_heads:
  assumes "ps ∈ sterms Γ (Γ pn)"
    and "wellformed Γ"
  shows "∃pn. ps ∈ stermsl (Γ pn)"
proof -
  from assms have "ps ∈ stermsl (Γ pn) ∨ (∃pn'. ps ∈ stermsl (Γ pn'))"
  by (rule sterms_stermsl)
  thus ?thesis by auto
qed

lemma sterms_subterms [dest]:
  assumes "wellformed Γ"
    and "∃pn. p ∈ subterms (Γ pn)"
    and "q ∈ sterms Γ p"
  shows "∃pn. q ∈ subterms (Γ pn)"
  using assms by (induct p) auto

lemma no_microsteps_sterms_refl:
  assumes "wellformed Γ"
  shows "(¬(∃q. p ⇨Γ q)) = (sterms Γ p = {p})"
proof (cases p)
  fix p1 p2
  assume "p = p1 ⊕ p2"
  from ⟨wellformed Γ⟩ have "p1 ⊕ p2 ∉ sterms Γ (p1 ⊕ p2)" by simp
  hence "sterms Γ (p1 ⊕ p2) ≠ {p1 ⊕ p2}" by auto
  moreover have "∃q. (p1 ⊕ p2) ⇨Γ q" by auto
  ultimately show ?thesis
  using ⟨p = p1 ⊕ p2⟩ by simp
next
  fix pn
  assume "p = call(pn)"
  from ⟨wellformed Γ⟩ have "call(pn) ∉ sterms Γ (call(pn))" by simp
  hence "sterms Γ (call(pn)) ≠ {call(pn)}" by auto
  moreover have "∃q. (call(pn)) ⇨Γ q" by auto
  ultimately show ?thesis
  using ⟨p = call(pn)⟩ by simp
qed simp_all

lemma sterms_maximal [elim]:
  assumes "wellformed Γ"
    and "q ∈ sterms Γ p"
  shows "sterms Γ q = {q}"
  using assms by (cases q) auto

lemma microstep_rtranscl_equal:
  assumes "not_call p"
    and "not_choice p"
    and "p ⇨Γ* q"
  shows "q = p"
  using assms(3) proof (rule converse_rtranclpE)
  fix p'
  assume "p ⇨Γ p'"
  with assms(1-2) show "q = p"
  by (cases p) simp_all
qed simp

lemma microstep_rtranscl_singleton [simp]:

```

```

assumes "not_call p"
  and "not_choice p"
  shows "{q. p  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}} = {p}"
proof (rule set_eqI)
  fix p'
  show "(p'  $\in$  {q. p  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}) = (p'  $\in$  {p})"
  proof
    assume "p'  $\in$  {q. p  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    hence "(microstep  $\Gamma$ )** p p'" and "sterms  $\Gamma$  p' = {p}'" by auto
    from this(1) have "p' = p"
    proof (rule converse_rtranclpE)
      fix q assume "p  $\rightsquigarrow_{\Gamma}$  q"
      with <not_call p> and <not_choice p> have False
      by (cases p) auto
      thus "p' = p" ..
    qed simp
    thus "p'  $\in$  {p}" by simp
  next
    assume "p'  $\in$  {p}"
    hence "p' = p" ..
    with <not_call p> and <not_choice p> show "p'  $\in$  {q. p  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    by (cases p) simp_all
  qed
qed

```

theorem sterms\_maximal\_microstep:

```

assumes "wellformed  $\Gamma$ "
  shows "sterms  $\Gamma$  p = {q. p  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$   $\neg$ ( $\exists$ q'. q  $\rightsquigarrow_{\Gamma}$  q')}"
proof
  from <wellformed  $\Gamma$ > have "sterms  $\Gamma$  p  $\subseteq$  {q. p  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
  proof induction
    fix p1 p2
    assume IH1: "sterms  $\Gamma$  p1  $\subseteq$  {q. p1  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
      and IH2: "sterms  $\Gamma$  p2  $\subseteq$  {q. p2  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    have "sterms  $\Gamma$  p1  $\subseteq$  {q. (p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    proof
      fix p'
      assume "p'  $\in$  sterms  $\Gamma$  p1"
      with IH1 have "p1  $\rightsquigarrow_{\Gamma}^*$  p'" by auto
      moreover have "(p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}$  p1" ..
      ultimately have "(p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  p'"
      by - (rule converse_rtranclp_into_rtranclp)
      moreover from <wellformed  $\Gamma$ > and <p'  $\in$  sterms  $\Gamma$  p1> have "sterms  $\Gamma$  p' = {p}'" ..
      ultimately show "p'  $\in$  {q. (p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
      by simp
    qed
    moreover have "sterms  $\Gamma$  p2  $\subseteq$  {q. (p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    proof
      fix p'
      assume "p'  $\in$  sterms  $\Gamma$  p2"
      with IH2 have "p2  $\rightsquigarrow_{\Gamma}^*$  p'" and "sterms  $\Gamma$  p' = {p}'" by auto
      moreover have "(p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}$  p2" ..
      ultimately have "(p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  p'"
      by - (rule converse_rtranclp_into_rtranclp)
      with <sterms  $\Gamma$  p' = {p}'> show "p'  $\in$  {q. (p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
      by simp
    qed
    ultimately show "sterms  $\Gamma$  (p1  $\oplus$  p2)  $\subseteq$  {q. (p1  $\oplus$  p2)  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    using <wellformed  $\Gamma$ > by simp
  next
    fix pn
    assume IH: "sterms  $\Gamma$  ( $\Gamma$  pn)  $\subseteq$  {q.  $\Gamma$  pn  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    show "sterms  $\Gamma$  (call(pn))  $\subseteq$  {q. (call(pn))  $\rightsquigarrow_{\Gamma}^*$  q  $\wedge$  sterms  $\Gamma$  q = {q}}"
    proof

```

```

fix p'
assume "p' ∈ sterms Γ (call(pn))"
with ⟨wellformed Γ⟩ have "p' ∈ sterms Γ (Γ pn)" by simp
with IH have "Γ pn  $\rightsquigarrow_{\Gamma}^*$  p'" and "sterms Γ p' = {p'}" by auto
note this(1)
moreover have "(call(pn))  $\rightsquigarrow_{\Gamma}$  Γ pn" by simp
ultimately have "(call(pn))  $\rightsquigarrow_{\Gamma}^*$  p'"
  by - (rule converse_rtranclp_into_rtranclp)
with ⟨sterms Γ p' = {p'}⟩ show "p' ∈ {q. (call(pn))  $\rightsquigarrow_{\Gamma}^*$  q} ∧ sterms Γ q = {q}"
  by simp
qed
qed simp_all
with ⟨wellformed Γ⟩ show "sterms Γ p ⊆ {q. p  $\rightsquigarrow_{\Gamma}^*$  q ∧ ¬(∃q'. q  $\rightsquigarrow_{\Gamma}$  q')}"
  by (simp only: no_microsteps_sterms_refl)
next
from ⟨wellformed Γ⟩ have "{q. p  $\rightsquigarrow_{\Gamma}^*$  q ∧ sterms Γ q = {q}} ⊆ sterms Γ p"
proof (induction)
  fix p1 p2
  assume IH1: "{q. p1  $\rightsquigarrow_{\Gamma}^*$  q ∧ sterms Γ q = {q}} ⊆ sterms Γ p1"
    and IH2: "{q. p2  $\rightsquigarrow_{\Gamma}^*$  q ∧ sterms Γ q = {q}} ⊆ sterms Γ p2"
  show "{q. (p1 ⊕ p2)  $\rightsquigarrow_{\Gamma}^*$  q ∧ sterms Γ q = {q}} ⊆ sterms Γ (p1 ⊕ p2)"
  proof (rule, drule CollectD, erule conjE)
    fix q'
    assume "(p1 ⊕ p2)  $\rightsquigarrow_{\Gamma}^*$  q'"
      and "sterms Γ q' = {q'}"
    with ⟨wellformed Γ⟩ have "(p1 ⊕ p2)  $\rightsquigarrow_{\Gamma}^+$  q'"
      by (auto dest!: rtranclpD sterms_no_choice)
    hence "p1  $\rightsquigarrow_{\Gamma}^*$  q' ∨ p2  $\rightsquigarrow_{\Gamma}^*$  q'"
      by (auto dest: tranclpD)
    thus "q' ∈ sterms Γ (p1 ⊕ p2)"
  proof
    assume "p1  $\rightsquigarrow_{\Gamma}^*$  q'"
    with IH1 and ⟨sterms Γ q' = {q'}⟩ have "q' ∈ sterms Γ p1" by auto
    with ⟨wellformed Γ⟩ show ?thesis by auto
  next
    assume "p2  $\rightsquigarrow_{\Gamma}^*$  q'"
    with IH2 and ⟨sterms Γ q' = {q'}⟩ have "q' ∈ sterms Γ p2" by auto
    with ⟨wellformed Γ⟩ show ?thesis by auto
  qed
qed
qed
next
fix pn
assume IH: "{q. Γ pn  $\rightsquigarrow_{\Gamma}^*$  q ∧ sterms Γ q = {q}} ⊆ sterms Γ (Γ pn)"
show "{q. (call(pn))  $\rightsquigarrow_{\Gamma}^*$  q ∧ sterms Γ q = {q}} ⊆ sterms Γ (call(pn))"
proof (rule, drule CollectD, erule conjE)
  fix q'
  assume "(call(pn))  $\rightsquigarrow_{\Gamma}^*$  q'"
    and "sterms Γ q' = {q'}"
  with ⟨wellformed Γ⟩ have "(call(pn))  $\rightsquigarrow_{\Gamma}^+$  q'"
    by (auto dest!: rtranclpD sterms_no_call)
  moreover have "(call(pn))  $\rightsquigarrow_{\Gamma}$  Γ pn" ..
  ultimately have "Γ pn  $\rightsquigarrow_{\Gamma}^*$  q'"
    by (auto dest!: tranclpD)
  with ⟨sterms Γ q' = {q'}⟩ and IH have "q' ∈ sterms Γ (Γ pn)" by auto
  with ⟨wellformed Γ⟩ show "q' ∈ sterms Γ (call(pn))" by simp
qed
qed simp_all
with ⟨wellformed Γ⟩ show "{q. p  $\rightsquigarrow_{\Gamma}^*$  q ∧ ¬(∃q'. q  $\rightsquigarrow_{\Gamma}$  q')} ⊆ sterms Γ p"
  by (simp only: no_microsteps_sterms_refl)
qed

```

## 7.5 Derivative terms

The derivatives of a term are those *sterms* potentially reachable by taking a transition, relative to a wellformed process specification  $\Gamma$ . These terms overapproximate the reachable terms, since the truth of guards is not considered.

```
function (domintros) dterms
  :: "('s, 'm, 'p, 'l) seqp_env  $\Rightarrow$  ('s, 'm, 'p, 'l) seqp  $\Rightarrow$  ('s, 'm, 'p, 'l) seqp set"
  where
    "dterms  $\Gamma$  ( $\{1\}\langle g \rangle$  p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  ( $\{1\}\llbracket u \rrbracket$  p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  (p1  $\oplus$  p2) = dterms  $\Gamma$  p1  $\cup$  dterms  $\Gamma$  p2"
  | "dterms  $\Gamma$  ( $\{1\}$ unicast( $s_{ip}$ ,  $s_{msg}$ ).p  $\triangleright$  q) = sterms  $\Gamma$  p  $\cup$  sterms  $\Gamma$  q"
  | "dterms  $\Gamma$  ( $\{1\}$ broadcast( $s_{msg}$ ). p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  ( $\{1\}$ groupcast( $s_{ips}$ ,  $s_{msg}$ ). p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  ( $\{1\}$ send( $s_{msg}$ ).p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  ( $\{1\}$ deliver( $s_{data}$ ).p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  ( $\{1\}$ receive( $u_{msg}$ ).p) = sterms  $\Gamma$  p"
  | "dterms  $\Gamma$  (call(pn)) = dterms  $\Gamma$  ( $\Gamma$  pn)"
  by pat_completeness auto
```

```
lemma dterms_dom_basic [simp]:
  assumes "not_call p"
    and "not_choice p"
  shows "dterms_dom ( $\Gamma$ , p)"
proof (rule accpI)
  fix y
  assume "dterms_rel y ( $\Gamma$ , p)"
  with assms show "dterms_dom y"
  by (cases p) (auto simp: dterms_rel.simps)
qed
```

```
lemma dterms_termination:
  assumes "wellformed  $\Gamma$ "
  shows "dterms_dom ( $\Gamma$ , p)"
proof -
  have dterms_rel': "dterms_rel = ( $\lambda$ gq gp. (gq, gp)  $\in$   $\{((\Gamma, q), (\Gamma', p)). \Gamma = \Gamma' \wedge p \rightsquigarrow_{\Gamma} q\}$ )"
  by (rule ext)+ (auto simp: dterms_rel.simps elim: microstep.cases)
  from  $\langle$ wellformed( $\Gamma$ ) $\rangle$  have " $\forall x. x \in$  Wellfounded.acc  $\{(q, p). p \rightsquigarrow_{\Gamma} q\}$ "
  unfolding wellformed_def by (simp add: wf_acc_iff)
  hence "p  $\in$  Wellfounded.acc  $\{(q, p). p \rightsquigarrow_{\Gamma} q\}$ " ..
  hence " $(\Gamma, p) \in$  Wellfounded.acc  $\{((\Gamma, q), \Gamma', p). \Gamma = \Gamma' \wedge p \rightsquigarrow_{\Gamma} q\}$ "
  by (rule acc_induct) (auto intro: accI)
  thus "dterms_dom ( $\Gamma$ , p)"
  unfolding dterms_rel' by (subst accp_acc_eq)
qed
```

```
lemmas dterms_psimps [simp] = dterms.psimps [OF dterms_termination]
  and dterms_pinduct = dterms.pinduct [OF dterms_termination]
```

```
lemma sterms_after_dterms [simp]:
  assumes "wellformed  $\Gamma$ "
  shows " $(\bigcup x \in$  dterms  $\Gamma$  p. sterms  $\Gamma$  x) = dterms  $\Gamma$  p"
  using assms by (induction p) simp_all
```

```
lemma sterms_before_dterms [simp]:
  assumes "wellformed  $\Gamma$ "
  shows " $(\bigcup x \in$  sterms  $\Gamma$  p. dterms  $\Gamma$  x) = dterms  $\Gamma$  p"
  using assms by (induction p) simp_all
```

```
lemma dterms_choice_disj [simp]:
  assumes "wellformed  $\Gamma$ "
  shows "p  $\in$  dterms  $\Gamma$  (p1  $\oplus$  p2) = (p  $\in$  dterms  $\Gamma$  p1  $\vee$  p  $\in$  dterms  $\Gamma$  p2)"
  using assms by (simp)
```



```

lemma dterms_in_branch:
  assumes "wellformed  $\Gamma$ "
    and " $p \in \text{dterms } \Gamma (p1 \oplus p2)$ "
    and " $p \in \text{dterms } \Gamma p1 \implies P$ "
    and " $p \in \text{dterms } \Gamma p2 \implies P$ "
  shows "P"
  using assms by auto

lemma dterms_no_choice:
  assumes "wellformed  $\Gamma$ "
  shows " $p1 \oplus p2 \notin \text{dterms } \Gamma p$ "
  using assms by induction simp_all

lemma dterms_not_choice [simp]:
  assumes "wellformed  $\Gamma$ "
    and " $q \in \text{dterms } \Gamma p$ "
  shows "not_choice q"
  using assms unfolding not_choice_def
  by (auto dest: dterms_no_choice)

lemma dterms_no_call:
  assumes "wellformed  $\Gamma$ "
  shows " $\text{call}(pn) \notin \text{dterms } \Gamma p$ "
  using assms by induction simp_all

lemma dterms_not_call [simp]:
  assumes "wellformed  $\Gamma$ "
    and " $q \in \text{dterms } \Gamma p$ "
  shows "not_call q"
  using assms unfolding not_call_def
  by (auto dest: dterms_no_call)

lemma dterms_subterms:
  assumes wf: "wellformed  $\Gamma$ "
    and " $\exists pn. p \in \text{subterms } (\Gamma pn)$ "
    and " $q \in \text{dterms } \Gamma p$ "
  shows " $\exists pn. q \in \text{subterms } (\Gamma pn)$ "
  using assms
  proof (induct p)
    fix p1 p2
    assume IH1: " $\exists pn. p1 \in \text{subterms } (\Gamma pn) \implies q \in \text{dterms } \Gamma p1 \implies \exists pn. q \in \text{subterms } (\Gamma pn)$ "
      and IH2: " $\exists pn. p2 \in \text{subterms } (\Gamma pn) \implies q \in \text{dterms } \Gamma p2 \implies \exists pn. q \in \text{subterms } (\Gamma pn)$ "
      and *: " $\exists pn. p1 \oplus p2 \in \text{subterms } (\Gamma pn)$ "
      and " $q \in \text{dterms } \Gamma (p1 \oplus p2)$ "
    from * obtain pn where " $p1 \oplus p2 \in \text{subterms } (\Gamma pn)$ "
      by auto
    hence " $p1 \in \text{subterms } (\Gamma pn)$ " and " $p2 \in \text{subterms } (\Gamma pn)$ "
      by auto
    from  $\langle q \in \text{dterms } \Gamma (p1 \oplus p2) \rangle$  wf have " $q \in \text{dterms } \Gamma p1 \vee q \in \text{dterms } \Gamma p2$ "
      by auto
    thus " $\exists pn. q \in \text{subterms } (\Gamma pn)$ "
      proof
        assume " $q \in \text{dterms } \Gamma p1$ "
        with  $\langle p1 \in \text{subterms } (\Gamma pn) \rangle$  show ?thesis
          by (auto intro: IH1)
        next
        assume " $q \in \text{dterms } \Gamma p2$ "
        with  $\langle p2 \in \text{subterms } (\Gamma pn) \rangle$  show ?thesis
          by (auto intro: IH2)
      qed
  qed auto

```

Note that the converse of  $\llbracket \text{wellformed } ?\Gamma; \exists pn. ?p \in \text{subterms } (?\Gamma pn); ?q \in \text{dterms } ?\Gamma ?p \rrbracket \implies \exists pn. ?q \in \text{subterms } (?\Gamma pn)$  is not true because *dterms* are an over-approximation; i.e., we cannot show, in general, that

guards return a non-empty set of post-states.

## 7.6 Control terms

The control terms of a process specification  $\Gamma$  are those subterms from which transitions are directly possible. We can omit  $call(pn)$  terms, since the root terms of all processes are considered, and also  $p1 \oplus p2$  terms since they effectively combine the transitions of the subterms  $p1$  and  $p2$ .

It will be shown that only the control terms, rather than all subterms, need be considered in invariant proofs.

**inductive\_set**

```
cterm :: "('s, 'm, 'p, 'l) seqp_env  $\Rightarrow$  ('s, 'm, 'p, 'l) seqp set"
for  $\Gamma$  :: "('s, 'm, 'p, 'l) seqp_env"
```

where

```
ctermSI[intro]: "p  $\in$  sterm  $\Gamma$  ( $\Gamma$  pn)  $\implies$  p  $\in$  cterm  $\Gamma$ "
| ctermDI[intro]: "[[ pp  $\in$  cterm  $\Gamma$ ; p  $\in$  dterm  $\Gamma$  pp ]  $\implies$  p  $\in$  cterm  $\Gamma$ "
```

**lemma cterms\_not\_choice [simp]:**

```
assumes "wellformed  $\Gamma$ "
and "p  $\in$  cterm  $\Gamma$ "
shows "not_choice p"
using assms
proof (cases p)
case CHOICE from (p  $\in$  cterm  $\Gamma$ ) show ?thesis
using (wellformed  $\Gamma$ ) by cases simp_all
qed simp_all
```

**lemma cterms\_no\_choice [simp]:**

```
assumes "wellformed  $\Gamma$ "
shows "p1  $\oplus$  p2  $\notin$  cterm  $\Gamma$ "
using assms by (auto dest: cterms_not_choice)
```

**lemma cterms\_not\_call [simp]:**

```
assumes "wellformed  $\Gamma$ "
and "p  $\in$  cterm  $\Gamma$ "
shows "not_call p"
using assms
proof (cases p)
case CALL from (p  $\in$  cterm  $\Gamma$ ) show ?thesis
using (wellformed  $\Gamma$ ) by cases simp_all
qed simp_all
```

**lemma cterms\_no\_call [simp]:**

```
assumes "wellformed  $\Gamma$ "
shows "call(pn)  $\notin$  cterm  $\Gamma$ "
using assms by (auto dest: cterms_not_call)
```

**lemma sterm\_cterm [elim]:**

```
assumes "p  $\in$  cterm  $\Gamma$ "
and "q  $\in$  sterm  $\Gamma$  p"
and "wellformed  $\Gamma$ "
shows "q  $\in$  cterm  $\Gamma$ "
using assms by - (cases p, auto)
```

**lemma dterm\_cterm [elim]:**

```
assumes "p  $\in$  cterm  $\Gamma$ "
and "q  $\in$  dterm  $\Gamma$  p"
and "wellformed  $\Gamma$ "
shows "q  $\in$  cterm  $\Gamma$ "
using assms by (cases p) auto
```

**lemma derivs\_in\_cterm [simp]:**

```
" $\bigwedge$  l f p. {l}<f> p  $\in$  cterm  $\Gamma$   $\implies$  sterm  $\Gamma$  p  $\subseteq$  cterm  $\Gamma$ "
" $\bigwedge$  l f p. {l}[f] p  $\in$  cterm  $\Gamma$   $\implies$  sterm  $\Gamma$  p  $\subseteq$  cterm  $\Gamma$ "
" $\bigwedge$  l fip fmsg q p. {l}unicast(fip, fmsg). p  $\triangleright$  q  $\in$  cterm  $\Gamma$ "
```

$$\implies \text{sterms } \Gamma \ p \subseteq \text{cterm s } \Gamma \wedge \text{sterms } \Gamma \ q \subseteq \text{cterm s } \Gamma$$

```

"∧1 fmsg p.      {1}broadcast(fmsg).p ∈ cterms Γ      ⇒ sterms Γ p ⊆ cterms Γ"
"∧1 fips fmsg p. {1}groupcast(fips, fmsg).p ∈ cterms Γ ⇒ sterms Γ p ⊆ cterms Γ"
"∧1 fmsg p.      {1}send(fmsg).p ∈ cterms Γ           ⇒ sterms Γ p ⊆ cterms Γ"
"∧1 fdata p.     {1}deliver(fdata).p ∈ cterms Γ       ⇒ sterms Γ p ⊆ cterms Γ"
"∧1 fmsg p.      {1}receive(fmsg).p ∈ cterms Γ        ⇒ sterms Γ p ⊆ cterms Γ"
by (auto simp: dterms.psimps)

```

## 7.7 Local control terms

We introduce a ‘local’ version of `cterm s` that does not step through calls and, thus, that is defined independently of a process specification  $\Gamma$ . This allows an alternative, terminating characterisation of `cterm s` as a set of subterms. Including `call(pn)`s in the set makes for a simpler relation with `stermsl`, even if they must be filtered out for the desired characterisation.

**function**

```
cterm sl :: "('s, 'm, 'p, 'l) seqp ⇒ ('s, 'm, 'p, 'l) seqp set"
```

**where**

```

"cterm sl ({1}(g) p)          = insert ({1}(g) p) (cterm sl p)"
| "cterm sl ({1}[u] p)       = insert ({1}[u] p) (cterm sl p)"
| "cterm sl ({1}unicast(sip, smsg). p ▷ q) = insert ({1}unicast(sip, smsg). p ▷ q)
                                     (cterm sl p ∪ cterm sl q)"
| "cterm sl ({1}broadcast(smsg). p)      = insert ({1}broadcast(smsg). p) (cterm sl p)"
| "cterm sl ({1}groupcast(sips, smsg). p) = insert ({1}groupcast(sips, smsg). p) (cterm sl p)"
| "cterm sl ({1}send(smsg). p)           = insert ({1}send(smsg). p) (cterm sl p)"
| "cterm sl ({1}deliver(sdata). p)       = insert ({1}deliver(sdata). p) (cterm sl p)"
| "cterm sl ({1}receive(umsg). p)       = insert ({1}receive(umsg). p) (cterm sl p)"
| "cterm sl (p1 ⊕ p2)                   = cterm sl p1 ∪ cterm sl p2"
| "cterm sl (call(pn))                  = {call(pn)}"

```

**by** `pat_completeness auto`

**termination by** `(relation "measure(size)") (auto dest: stermsl_nobigger)`

**lemmas** `cterm sl_induct =`

```

cterm sl.induct [case_names GUARD ASSIGN UCAST BCAST GCAST
                  SEND DELIVER RECEIVE CHOICE CALL]

```

**lemma** `cterm sl_refl [intro]: "not_choice p ⇒ p ∈ cterm sl p"`

**by** `(cases p) auto`

**lemma** `cterm sl_subterms:`

```
"cterm sl p = {q. q ∈ subterms p ∧ not_choice q}" (is "?lhs = ?rhs")
```

**proof**

```
show "?lhs ⊆ ?rhs" by (induct p, auto) next
```

```
show "?rhs ⊆ ?lhs" by (induct p, auto)
```

**qed**

**lemma** `cterm sl_trans [elim]:`

```
assumes "q ∈ cterm sl p"
```

```
and "r ∈ cterm sl q"
```

```
shows "r ∈ cterm sl p"
```

**using** `assms`

**proof** `(induction p rule: cterm sl_induct)`

```
case (CHOICE p1 p2)
```

```
have "(q ∈ cterm sl p1) ∨ (q ∈ cterm sl p2)"
```

```
using CHOICE.prem s(1) by simp
```

```
hence "r ∈ cterm sl p1 ∨ r ∈ cterm sl p2"
```

```
proof (rule disj_forward)
```

```
assume "q ∈ cterm sl p1"
```

```
thus "r ∈ cterm sl p1" using (r ∈ cterm sl q) by (rule CHOICE.IH)
```

```
next
```

```
assume "q ∈ cterm sl p2"
```

```
thus "r ∈ cterm sl p2" using (r ∈ cterm sl q) by (rule CHOICE.IH)
```

```
qed
```

```
thus "r ∈ cterm sl (p1 ⊕ p2)" by simp
```

```

qed auto

lemma ctermsl_ex_trans [elim]:
  assumes "∃q ∈ ctermsl p. r ∈ ctermsl q"
  shows "r ∈ ctermsl p"
  using assms by auto

lemma call_ctermsl_empty [elim]:
  "[[ p ∈ ctermsl p'; not_call p ]] ⇒ not_call p'"
  unfolding not_call_def by (cases p) auto

lemma stermsl_ctermsl_choice1 [simp]:
  assumes "q ∈ stermsl p1"
  shows "q ∈ ctermsl (p1 ⊕ p2)"
  using assms by (induction p1) auto

lemma stermsl_ctermsl_choice2 [simp]:
  assumes "q ∈ stermsl p2"
  shows "q ∈ ctermsl (p1 ⊕ p2)"
  using assms by (induction p2) auto

lemma stermsl_ctermsl [elim]:
  assumes "q ∈ stermsl p"
  shows "q ∈ ctermsl p"
  using assms
  proof (cases p)
    case (CHOICE p1 p2)
    hence "q ∈ stermsl (p1 ⊕ p2)" using assms by simp
    hence "q ∈ stermsl p1 ∨ q ∈ stermsl p2" by simp
    hence "q ∈ ctermsl (p1 ⊕ p2)" by (rule) (simp_all del: ctermsl.simps)
    thus "q ∈ ctermsl p" using CHOICE by simp
  qed simp_all

lemma stermsl_after_ctermsl [simp]:
  "(⋃x∈ctermsl p. stermsl x) = ctermsl p"
  by (induct p) auto

lemma stermsl_before_ctermsl [simp]:
  "(⋃x∈stermsl p. ctermsl x) = ctermsl p"
  by (induct p) simp_all

lemma ctermsl_no_choice: "p1 ⊕ p2 ∉ ctermsl p"
  by (induct p) simp_all

lemma ctermsl_ex_stermsl: "q ∈ ctermsl p ⇒ ∃ps∈stermsl p. q ∈ ctermsl ps"
  by (induct p) auto

lemma dterms_ctermsl [intro]:
  assumes "q ∈ dterms Γ p"
  and "wellformed Γ"
  shows "q ∈ ctermsl p ∨ (∃pn. q ∈ ctermsl (Γ pn))"
  using assms(1-2)
  proof (induction p rule: dterms_pinduct [OF ⟨wellformed Γ⟩])
    fix Γ l fg p
    assume "q ∈ dterms Γ ({l}⟨fg⟩ p)"
    and "wellformed Γ"
    hence "q ∈ sterms Γ p" by simp
    hence "q ∈ stermsl p ∨ (∃pn. q ∈ stermsl (Γ pn))"
    using ⟨wellformed Γ⟩ by (rule sterms_stermsl)
    thus "q ∈ ctermsl ({l}⟨fg⟩ p) ∨ (∃pn. q ∈ ctermsl (Γ pn))"
  proof
    assume "q ∈ stermsl p"
    hence "q ∈ ctermsl p" by (rule stermsl_ctermsl)
    hence "q ∈ ctermsl ({l}⟨fg⟩ p)" by simp
  end
end

```

```

    thus ?thesis ..
next
  assume "∃pn. q ∈ stermsl (Γ pn)"
  then obtain pn where "q ∈ stermsl (Γ pn)" by auto
  hence "q ∈ ctermsl (Γ pn)" by (rule stermsl_ctermsl)
  hence "∃pn. q ∈ ctermsl (Γ pn)" ..
  thus ?thesis ..
qed
next
  fix Γ p1 p2
  assume "q ∈ dterms Γ (p1 ⊕ p2)"
    and IH1: "⊥ q ∈ dterms Γ p1; wellformed Γ ] ⇒ q ∈ ctermsl p1 ∨ (∃pn. q ∈ ctermsl (Γ pn))"
    and IH2: "⊥ q ∈ dterms Γ p2; wellformed Γ ] ⇒ q ∈ ctermsl p2 ∨ (∃pn. q ∈ ctermsl (Γ pn))"
    and "wellformed Γ"
  thus "q ∈ ctermsl (p1 ⊕ p2) ∨ (∃pn. q ∈ ctermsl (Γ pn))"
    by auto
next
  fix Γ pn
  assume "q ∈ dterms Γ (call(pn))"
    and "wellformed Γ"
    and "⊥ q ∈ dterms Γ (Γ pn); wellformed Γ ] ⇒ q ∈ ctermsl (Γ pn) ∨ (∃pn. q ∈ ctermsl (Γ pn))"
  thus "q ∈ ctermsl (call(pn)) ∨ (∃pn. q ∈ ctermsl (Γ pn))"
    by auto
qed (simp_all, (metis sterms_stermsl stermsl_ctermsl)+)

lemma ctermsl_cterms [elim]:
  assumes "q ∈ ctermsl p"
    and "not_call q"
    and "sterms Γ p ⊆ cterms Γ"
    and "wellformed Γ"
  shows "q ∈ cterms Γ"
  using assms by (induct p rule: ctermsl.induct) auto

```

## 7.8 Local deriviative terms

We define local *dterms* for use in the theorem that relates *cterms* and sets of *ctermsl*.

```

function dtermsl
  :: "('s, 'm, 'p, 'l) seqp ⇒ ('s, 'm, 'p, 'l) seqp set"
  where
    "dtermsl ({}<fg> p) = stermsl p"
  | "dtermsl ({}[[fa]] p) = stermsl p"
  | "dtermsl (p1 ⊕ p2) = dtermsl p1 ∪ dtermsl p2"
  | "dtermsl ({}unicast(fip, fmsg).p ▷ q) = stermsl p ∪ stermsl q"
  | "dtermsl ({}broadcast(fmsg). p) = stermsl p"
  | "dtermsl ({}groupcast(fips, fmsg). p) = stermsl p"
  | "dtermsl ({}send(fmsg).p) = stermsl p"
  | "dtermsl ({}deliver(fdata).p) = stermsl p"
  | "dtermsl ({}receive(fmsg).p) = stermsl p"
  | "dtermsl (call(pn)) = {}"
  by pat_completeness auto
  termination by (relation "measure(size)") (auto dest: stermsl_nobigger)

```

```

lemma stermsl_after_dtermsl [simp]:
  shows "(⋃x∈dtermsl p. stermsl x) = dtermsl p"
  by (induct p) simp_all

```

```

lemma stermsl_before_dtermsl [simp]:
  "(⋃x∈stermsl p. dtermsl x) = dtermsl p"
  by (induct p) simp_all

```

```

lemma dtermsl_no_choice [simp]: "p1 ⊕ p2 ∉ dtermsl p"
  by (induct p) simp_all

```

```

lemma dtermsl_choice_disj [simp]:

```

```


$p \in dtermsl (p1 \oplus p2) = (p \in dtermsl p1 \vee p \in dtermsl p2)$ "



by simp



lemma dtermsl_in_branch [elim]:



" $\llbracket p \in dtermsl (p1 \oplus p2); p \in dtermsl p1 \implies P; p \in dtermsl p2 \implies P \rrbracket \implies P$ "



by auto



lemma ctermsl_dtermsl [elim]:



assumes " $q \in dtermsl p$ "



shows " $q \in ctermsl p$ "



using assms by (induct p) (simp_all, (metis stermsl_ctermsl)+)



lemma dtermsl_dterms [elim]:



assumes " $q \in dtermsl p$ "



and "not_call q"



and "wellformed  $\Gamma$ "



shows " $q \in dterms \Gamma p$ "



using assms



using assms by (induct p) (simp_all, (metis stermsl_sterms)+)



lemma ctermsl_stermsl_or_dtermsl:



assumes " $q \in ctermsl p$ "



shows " $q \in stermsl p \vee (\exists p' \in dtermsl p. q \in ctermsl p')$ "



using assms by (induct p) (auto dest: ctermsl_ex_stermsl)



lemma dtermsl_add_stermsl_beforeD:



assumes " $q \in dtermsl p$ "



shows " $\exists ps \in stermsl p. q \in dtermsl ps$ "



proof -



from assms have " $q \in (\bigcup x \in stermsl p. dtermsl x)$ " by auto



thus ?thesis



by (rule UN_E) auto



qed



lemma call_dtermsl_empty [elim]:



" $q \in dtermsl p \implies not\_call p$ "



by (cases p) simp_all


```

## 7.9 More properties of control terms

We now show an alternative definition of *cterms* based on sets of local control terms. While the original definition has convenient induction and simplification rules, useful for proving properties like *cterms\_includes\_sterms\_of\_seq\_reachable*, this definition makes it easier to systematically generate the set of control terms of a process specification.

```

theorem cterms_def':
  assumes wfg: "wellformed  $\Gamma$ "
  shows " $cterms \Gamma = \{ p \mid p \text{ pn. } p \in ctermsl (\Gamma \text{ pn}) \wedge not\_call p \}$ "
    (is " $\_ = ?ctermsl\_set$ ")
  proof (rule iffI [THEN set_eqI])
    fix p
    assume " $p \in cterms \Gamma$ "
    thus " $p \in ?ctermsl\_set$ "
    proof (induction p)
      fix p pn
      assume " $p \in sterms \Gamma (\Gamma \text{ pn})$ "
      then obtain pn' where " $p \in stermsl (\Gamma \text{ pn}')$ " using wfg
        by (blast dest: sterms_stermsl_heads)
      hence " $p \in ctermsl (\Gamma \text{ pn}')$ " ..
      moreover from  $\langle p \in sterms \Gamma (\Gamma \text{ pn}) \rangle$  wfg have "not_call p" by simp
      ultimately show " $p \in ?ctermsl\_set$ " by auto
    next
      fix pp p
      assume " $pp \in cterms \Gamma$ "
      and IH: " $pp \in ?ctermsl\_set$ "

```

```

    and *: "p ∈ dterms Γ pp"
  from * have "p ∈ ctermsl pp ∨ (∃pn. p ∈ ctermsl (Γ pn))"
    using wfg by (rule dterms_termsl)
  hence "∃pn. p ∈ ctermsl (Γ pn)"
  proof
    assume "p ∈ ctermsl pp"
    from ⟨pp ∈ cterms Γ⟩ and IH obtain pn' where "pp ∈ ctermsl (Γ pn')"
      by auto
    with ⟨p ∈ ctermsl pp⟩ have "p ∈ ctermsl (Γ pn')" by auto
    thus "∃pn. p ∈ ctermsl (Γ pn)" ..
  qed -
  moreover from ⟨p ∈ dterms Γ pp⟩ wfg have "not_call p" by simp
  ultimately show "p ∈ ?ctermssl_set" by auto
qed
next
fix p
assume "p ∈ ?ctermssl_set"
then obtain pn where *: "p ∈ ctermsl (Γ pn)" and "not_call p" by auto
from * have "p ∈ stermsl (Γ pn) ∨ (∃p'∈dtermsl (Γ pn). p ∈ ctermsl p')"
  by (rule ctermsl_stermsl_or_dtermsl)
thus "p ∈ cterms Γ"
proof
  assume "p ∈ stermsl (Γ pn)"
  hence "p ∈ sterms Γ (Γ pn)" using ⟨not_call p⟩ wfg ..
  thus "p ∈ cterms Γ" ..
next
assume "∃p'∈dtermsl (Γ pn). p ∈ ctermsl p'"
then obtain p' where p'1: "p' ∈ dtermsl (Γ pn)"
  and p'2: "p ∈ ctermsl p'" ..
from p'2 and ⟨not_call p⟩ have "not_call p'" ..
from p'1 obtain ps where ps1: "ps ∈ stermsl (Γ pn)"
  and ps2: "p' ∈ dtermsl ps"
  by (blast dest: dtermsl_add_stermsl_beforeD)
from ps2 have "not_call ps" ..
with ps1 have "ps ∈ cterms Γ" using wfg by auto
with ⟨p' ∈ dtermsl ps⟩ and ⟨not_call p'⟩ have "p' ∈ cterms Γ" using wfg by auto
hence "sterms Γ p' ⊆ cterms Γ" using wfg by auto
with ⟨p ∈ ctermsl p'⟩ ⟨not_call p⟩ show "p ∈ cterms Γ" using wfg ..
qed
qed

lemma ctermsE [elim]:
  assumes "wellformed Γ"
    and "p ∈ cterms Γ"
  obtains pn where "p ∈ ctermsl (Γ pn)"
    and "not_call p"
  using assms(2) unfolding cterms_def' [OF assms(1)] by auto

corollary cterms_subterms:
  assumes "wellformed Γ"
  shows "cterm Γ = {p | p pn. p ∈ subterms (Γ pn) ∧ not_call p ∧ not_choice p}"
  by (subst cterms_def' [OF assms(1)], subst ctermsl_subterms) auto

lemma subterms_in_cterms [elim]:
  assumes "wellformed Γ"
    and "p ∈ subterms (Γ pn)"
    and "not_call p"
    and "not_choice p"
  shows "p ∈ cterms Γ"
  using assms unfolding cterms_subterms [OF ⟨wellformed Γ⟩] by auto

lemma subterms_stermsl_termsl:
  assumes "q ∈ subterms p"
    and "r ∈ stermsl q"

```

```

  shows "r ∈ ctermsl p"
using assms
proof (induct p)
  fix p1 p2
  assume IH1: "q ∈ subterms p1 ⇒ r ∈ stermsl q ⇒ r ∈ ctermsl p1"
    and IH2: "q ∈ subterms p2 ⇒ r ∈ stermsl q ⇒ r ∈ ctermsl p2"
    and *: "q ∈ subterms (p1 ⊕ p2)"
    and "r ∈ stermsl q"
  from * have "q ∈ {p1 ⊕ p2} ∪ subterms p1 ∪ subterms p2" by simp
  thus "r ∈ ctermsl (p1 ⊕ p2)"
  proof (elim UnE)
    assume "q ∈ {p1 ⊕ p2}" with ⟨r ∈ stermsl q⟩ show ?thesis
    by simp (metis stermsl_ctermsl)
  next
    assume "q ∈ subterms p1" hence "r ∈ ctermsl p1" using ⟨r ∈ stermsl q⟩ by (rule IH1)
    thus ?thesis by simp
  next
    assume "q ∈ subterms p2" hence "r ∈ ctermsl p2" using ⟨r ∈ stermsl q⟩ by (rule IH2)
    thus ?thesis by simp
  qed
qed auto

```

lemma subterms\_sterms\_cterms:

```

assumes wf: "wellformed Γ"
  and "p ∈ subterms (Γ pn)"
  shows "sterms Γ p ⊆ cterms Γ"
using assms(2)
proof (induct p)
  fix p
  assume "call(p) ∈ subterms (Γ pn)"
  from wf have "sterms Γ (call(p)) = sterms Γ (Γ p)" by simp
  thus "sterms Γ (call(p)) ⊆ cterms Γ" by auto
next
  fix p1 p2
  assume IH1: "p1 ∈ subterms (Γ pn) ⇒ sterms Γ p1 ⊆ cterms Γ"
    and IH2: "p2 ∈ subterms (Γ pn) ⇒ sterms Γ p2 ⊆ cterms Γ"
    and *: "p1 ⊕ p2 ∈ subterms (Γ pn)"
  from * have "p1 ∈ subterms (Γ pn)" by auto
  hence "sterms Γ p1 ⊆ cterms Γ" by (rule IH1)
  moreover from * have "p2 ∈ subterms (Γ pn)" by auto
  hence "sterms Γ p2 ⊆ cterms Γ" by (rule IH2)
  ultimately show "sterms Γ (p1 ⊕ p2) ⊆ cterms Γ" using wf by simp
qed (auto elim!: subterms_in_cterms [OF ⟨wellformed Γ⟩])

```

lemma subterms\_sterms\_in\_cterms:

```

assumes "wellformed Γ"
  and "p ∈ subterms (Γ pn)"
  and "q ∈ sterms Γ p"
  shows "q ∈ cterms Γ"
using assms
by (auto dest!: subterms_sterms_cterms [OF ⟨wellformed Γ⟩])

```

end

## 8 Labelling sequential processes

```

theory AWN_Labels
imports AWN AWN_Cterms
begin

```



## 8.1 Labels

Labels serve two main purposes. They allow the substitution of *stems* in *invariant* proofs. They also allow the strengthening (control state dependent) of invariants.

```

function (domintros) labels
  :: "(s, m, p, l) seqp_env  $\Rightarrow$  (s, m, p, l) seqp  $\Rightarrow$  l set"
  where
    "labels  $\Gamma$  (fg) p" = {l}"
  | "labels  $\Gamma$  (fa) p" = {l}"
  | "labels  $\Gamma$  (p1  $\oplus$  p2)" = labels  $\Gamma$  p1  $\cup$  labels  $\Gamma$  p2"
  | "labels  $\Gamma$  (unicast(fip, fmsg).p  $\triangleright$  q)" = {l}"
  | "labels  $\Gamma$  (broadcast(fmsg). p)" = {l}"
  | "labels  $\Gamma$  (groupcast(fips, fmsg). p)" = {l}"
  | "labels  $\Gamma$  (send(fmsg).p)" = {l}"
  | "labels  $\Gamma$  (deliver(fdata).p)" = {l}"
  | "labels  $\Gamma$  (receive(fmsg).p)" = {l}"
  | "labels  $\Gamma$  (call(pn))" = labels  $\Gamma$  ( $\Gamma$  pn)"
  by pat_completeness auto

lemma labels_dom_basic [simp]:
  assumes "not_call p"
    and "not_choice p"
  shows "labels_dom ( $\Gamma$ , p)"
  proof (rule accpI)
    fix y
    assume "labels_rel y ( $\Gamma$ , p)"
    with assms show "labels_dom y"
    by (cases p) (auto simp: labels_rel.simps)
  qed

lemma labels_termination:
  fixes  $\Gamma$  p
  assumes "wellformed( $\Gamma$ )"
  shows "labels_dom ( $\Gamma$ , p)"
  proof -
    have labels_rel': "labels_rel = ( $\lambda$ gq gp. (gq, gp)  $\in$  {(( $\Gamma$ , q), ( $\Gamma'$ , p)).  $\Gamma$  =  $\Gamma'$   $\wedge$  p  $\rightsquigarrow_{\Gamma}$  q})"
    by (rule ext)+ (auto simp: labels_rel.simps intro: microstep.intros elim: microstep.cases)
    from (wellformed( $\Gamma$ )) have " $\forall$ x. x  $\in$  Wellfounded.acc {(q, p). p  $\rightsquigarrow_{\Gamma}$  q}"
    unfolding wellformed_def by (simp add: wf_acc_iff)
    hence "p  $\in$  Wellfounded.acc {(q, p). p  $\rightsquigarrow_{\Gamma}$  q}" ..
    hence "( $\Gamma$ , p)  $\in$  Wellfounded.acc {(( $\Gamma$ , q), ( $\Gamma'$ , p)).  $\Gamma$  =  $\Gamma'$   $\wedge$  p  $\rightsquigarrow_{\Gamma}$  q}"
    by (rule acc_induct) (auto intro: accI)
    thus "labels_dom ( $\Gamma$ , p)"
    unfolding labels_rel' by (subst accp_acc_eq)
  qed

declare labels.psimps[simp]

lemmas labels_pinduct = labels.pinduct [OF labels_termination]
  and labels.psimps[simp] = labels.psimps [OF labels_termination]

lemma labels_not_empty:
  fixes  $\Gamma$  p
  assumes "wellformed  $\Gamma$ "
  shows "labels  $\Gamma$  p  $\neq$  {}"
  by (induct p rule: labels_pinduct [OF (wellformed  $\Gamma$ )] simp_all)

lemma has_label [dest]:
  fixes  $\Gamma$  p
  assumes "wellformed  $\Gamma$ "
  shows " $\exists$ l. l  $\in$  labels  $\Gamma$  p"
  using labels_not_empty [OF assms] by auto

lemma singleton_labels [simp]:

```

```

" $\bigwedge \Gamma \ 1 \ 1' \ f \ p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\langle f \rangle p)$        $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ f \ p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\llbracket f \rrbracket p)$        $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fip } \text{fmsg } p \ q.$   $1 \in \text{labels } \Gamma \ (\{1'\}\text{unicast}(\text{fip}, \text{fmsg}).p \triangleright q)$   $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fmsg } p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\text{broadcast}(\text{fmsg}).p)$        $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fips } \text{fmsg } p.$   $1 \in \text{labels } \Gamma \ (\{1'\}\text{groupcast}(\text{fips}, \text{fmsg}).p)$   $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fmsg } p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\text{send}(\text{fmsg}).p)$        $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fdata } p.$       $1 \in \text{labels } \Gamma \ (\{1'\}\text{deliver}(\text{fdata}).p)$       $= (1 = 1')$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fmsg } p.$       $1 \in \text{labels } \Gamma \ (\{1'\}\text{receive}(\text{fmsg}).p)$       $= (1 = 1')$ "
by auto

```

lemma in\_labels\_singletons [dest!]:

```

" $\bigwedge \Gamma \ 1 \ 1' \ f \ p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\langle f \rangle p)$        $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ f \ p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\llbracket f \rrbracket p)$        $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fip } \text{fmsg } p \ q.$   $1 \in \text{labels } \Gamma \ (\{1'\}\text{unicast}(\text{fip}, \text{fmsg}).p \triangleright q)$   $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fmsg } p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\text{broadcast}(\text{fmsg}).p)$        $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fips } \text{fmsg } p.$   $1 \in \text{labels } \Gamma \ (\{1'\}\text{groupcast}(\text{fips}, \text{fmsg}).p)$   $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fmsg } p.$        $1 \in \text{labels } \Gamma \ (\{1'\}\text{send}(\text{fmsg}).p)$        $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fdata } p.$       $1 \in \text{labels } \Gamma \ (\{1'\}\text{deliver}(\text{fdata}).p)$       $\implies 1 = 1'$ "
" $\bigwedge \Gamma \ 1 \ 1' \ \text{fmsg } p.$       $1 \in \text{labels } \Gamma \ (\{1'\}\text{receive}(\text{fmsg}).p)$       $\implies 1 = 1'$ "
by auto

```

definition

```
simple_labels :: "( 's, 'm, 'p, 'l) seqp_env  $\Rightarrow$  bool"
```

where

```
"simple_labels  $\Gamma \equiv \forall pn. \forall p \in \text{subterms } (\Gamma \ pn). (\exists !l. \text{labels } \Gamma \ p = \{l\})"$ 
```

lemma simple\_labelsI [intro]:

```

assumes " $\bigwedge pn \ p. p \in \text{subterms } (\Gamma \ pn) \implies \exists !l. \text{labels } \Gamma \ p = \{l\}"$ 
shows "simple_labels  $\Gamma$ "
using assms unfolding simple_labels_def by auto

```

The *simple\_labels*  $\Gamma$  property is necessary to transfer results shown over the *cterms* of a process specification  $\Gamma$  to the reachable actions of that process.

Consider the process  $\{l_1\}\text{send}(m1) . p1 \oplus \{l_2\}\text{send}(m2) . p2$ . The iteration over *cterms*  $\Gamma$  will cover the two transitions  $(l_1, \text{send } m1, p1)$  and  $(l_2, \text{send } m2, p2)$ , but reachability requires the four transitions  $(l_1, \text{send } m1, p1)$ ,  $(l_1, \text{send } m2, p2)$ ,  $(l_2, \text{send } m1, p1)$ , and  $(l_2, \text{send } m2, p2)$ .

In a simply labelled process, the former is sufficient to show the latter, since  $l_1 = l_2$ .

This requirement seems really only to be restrictive for processes where a *call*(*pn*) occurs as a direct subterm of a choice operator. Consider, for instance,  $\{l_1\}\llbracket e \rrbracket p \oplus \text{call}(pn)$ . Here  $l_1$  must equal the label of  $\Gamma \ pn$ , which can then not be distinguished from any other subterm that calls *pn* in any other process.

This limitation stems from the fact that the "call points" of a process are effectively treated as the root of the called process. This is by design; we try to treat call sites as "syntactic pastings" of process terms, giving rise, conceptually, to an infinite tree structure. But this prejudices the alternative view that process calls are used as "join points" of "process threads", in complement to the "fork points" of the  $p1 \oplus p2$  operator.

lemma simple\_labels\_in\_sterms:

```

fixes  $\Gamma \ 1 \ p$ 
assumes "simple_labels  $\Gamma$ "
      and "wellformed  $\Gamma$ "
      and " $\exists pn. p \in \text{subterms } (\Gamma \ pn)"$ 
      and " $1 \in \text{labels } \Gamma \ p$ "
shows " $\forall p' \in \text{sterms } \Gamma \ p. 1 \in \text{labels } \Gamma \ p'$ "
using assms
proof (induct p rule: labels_pinduct [OF (wellformed  $\Gamma$ )])
  fix  $\Gamma \ p1 \ p2$ 
  assume s1: "simple_labels  $\Gamma$ "
      and wf: "wellformed  $\Gamma$ "
      and IH1: " $\llbracket \text{simple_labels } \Gamma; \text{wellformed } \Gamma; \exists pn. p1 \in \text{subterms } (\Gamma \ pn); 1 \in \text{labels } \Gamma \ p1 \rrbracket \implies \forall p' \in \text{sterms } \Gamma \ p1. 1 \in \text{labels } \Gamma \ p'$ "
      and IH2: " $\llbracket \text{simple_labels } \Gamma; \text{wellformed } \Gamma; \exists pn. p2 \in \text{subterms } (\Gamma \ pn); 1 \in \text{labels } \Gamma \ p2 \rrbracket \implies \forall p' \in \text{sterms } \Gamma \ p2. 1 \in \text{labels } \Gamma \ p'$ "

```

```

    and ein: "∃pn. p1 ⊕ p2 ∈ subterms (Γ pn)"
    and l12: "l ∈ labels Γ (p1 ⊕ p2)"
  from sl ein l12 have "labels Γ (p1 ⊕ p2) = {l}"
    unfolding simple_labels_def by (metis empty_iff insert_iff)
  with wf have "labels Γ p1 ∪ labels Γ p2 = {l}" by simp
  moreover have "labels Γ p1 ≠ {}" and "labels Γ p2 ≠ {}"
    using wf by (metis labels_not_empty)+
  ultimately have "l ∈ labels Γ p1" and "l ∈ labels Γ p2"
    by (metis Un_iff empty_iff insert_iff set_eqI)+
  moreover from ein have "∃pn. p1 ∈ subterms (Γ pn)"
    and "∃pn. p2 ∈ subterms (Γ pn)"

  by auto
  ultimately show "∀p' ∈ sterms Γ (p1 ⊕ p2). l ∈ labels Γ p'"
    using wf IH1 [OF sl wf] IH2 [OF sl wf] by auto
qed auto

```

lemma labels\_in\_sterms:

```

  fixes Γ l p
  assumes "wellformed Γ"
    and "l ∈ labels Γ p"
  shows "∃p' ∈ sterms Γ p. l ∈ labels Γ p'"
  using assms
  by (induct p rule: labels_pinduct [OF <wellformed Γ>]) (auto intro: Un_iff)

```

lemma labels\_sterms\_labels:

```

  fixes Γ p p' l
  assumes "wellformed Γ"
    and "p' ∈ sterms Γ p"
    and "l ∈ labels Γ p'"
  shows "l ∈ labels Γ p"
  using assms
  by (induct p rule: labels_pinduct [OF <wellformed Γ>]) auto

```

primrec labelfrom :: "int ⇒ int ⇒ ('s, 'm, 'p, 'a) seqp ⇒ int × ('s, 'm, 'p, int) seqp"
where

```

"labelfrom n nn ({}_<f> p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}<f> p'))"
| "labelfrom n nn ({}_[[f]] p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}[[f]] p'))"
| "labelfrom n nn (p ⊕ q) =
  (let (nn', p') = labelfrom n nn p in
  let (nn'', q') = labelfrom n nn' q in
  (nn'', p' ⊕ q'))"
| "labelfrom n nn ({}_unicast(fip, fmsg). p ▷ q) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  let (nn'', q') = labelfrom nn' (nn' + 1) q in
  (nn'', {n}unicast(fip, fmsg). p' ▷ q'))"
| "labelfrom n nn ({}_broadcast(fmsg). p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}broadcast(fmsg). p'))"
| "labelfrom n nn ({}_groupcast(fipset, fmsg). p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}groupcast(fipset, fmsg). p'))"
| "labelfrom n nn ({}_send(fmsg). p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}send(fmsg). p'))"
| "labelfrom n nn ({}_deliver(fdata). p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}deliver(fdata). p'))"
| "labelfrom n nn ({}_receive(fmsg). p) =
  (let (nn', p') = labelfrom nn (nn + 1) p in
  (nn', {n}receive(fmsg). p'))"

```

```

| "labelfrom n nn (call(fargs)) = (nn - 1, call(fargs))"

datatype 'pn label =
  LABEL 'pn int ("_-" [1000, 1000] 999)

instantiation "label" :: (ord) ord
begin

fun less_eq_label :: "'a label ⇒ 'a label ⇒ bool"
where "(l1-:n1) ≤ (l2-:n2) = (l1 = l2 ∧ n1 ≤ n2)"

definition less_label: "(l1::'a label) < l2 ⟷ l1 ≤ l2 ∧ ¬ (l1 ≤ l2)"

instance ..
end

abbreviation labelled :: "'p ⇒ ('s, 'm, 'p, 'a) seqp ⇒ ('s, 'm, 'p, 'p label) seqp"
where "labelled pn p ≡ labelmap (λl. LABEL pn l) (snd (labelfrom 0 1 p))"

end

```

## 9 A custom tactic for showing invariants via control terms

```

theory Inv_Cterms
imports AWN_Labels
begin

```

This tactic tries to solve a goal by reducing it to a problem over (local) cterms (using one of the cterms.intros intro rules); expanding those to consider all process names (using one of the ctermssl\_cases destruction rules); simplifying each (using the cterms.env simplification rules); splitting them up into separate subgoals; replacing the derivative term with a variable; ‘executing’ a transition of each term; and then simplifying.

The tactic can stop after applying introduction rule (“inv\_cterms (intro\_only)”), or after having generated the verification condition subgoals and before having simplified them (“inv\_cterms (vcs\_only)”). It takes arguments to add or remove simplification rules (“simp add: lemmanames”), to add forward rules on assumptions (to introduce previously proved invariants; “inv add: lemmanames”), or to add elimination rules that solve any remaining subgoals (“solve: lemmanames”).

To configure the tactic for a set of transition rules:

1. add elimination rules: declare seqpTEs [ctermssl\_cases]
2. add rules to replace derivative terms: declare elimders [ctermssl\_cases]

To configure the tactic for a process environment ( $\Gamma$ ):

1. add simp rules: declare  $\Gamma$ .simps [ctermssl\_cases]
2. add case rules: declare aadv\_proc\_cases [ctermssl\_cases]
3. add invariant intros declare seq\_invariant\_ctermssl [OF aadv\_wf aadv\_control\_within aadv\_simple\_labels, cterms\_intros] seq\_step\_invariant\_ctermssl [OF aadv\_wf aadv\_control\_within aadv\_simple\_labels, cterms\_intros]

```

lemma has_ctermssl: "p ∈ ctermssl Γ ⟹ p ∈ ctermssl Γ" .

```

```

named_theorems cterms_elimders "rules for truncating sequential process terms"
named_theorems cterms_seqte "elimination rules for sequential process terms"
named_theorems cterms_env "simplification rules for sequential process environments"
named_theorems ctermssl_cases "destruction rules for case splitting ctermssl"
named_theorems cterms_intros "introduction rules from cterms"
named_theorems cterms_invs "invariants to try to apply at each vc"
named_theorems cterms_final "elimination rules to try on each vc after simplification"

```

ML  $\langle$

```

fun simp_only thms ctxt =
  asm_full_simp_tac
    (ctxt |> Raw_Simplifier.clear_simpset |> fold Simplifier.add_simp thms)

(* shallow_simp is useful for mopping up assumptions before really trying to simplify.
   Perhaps surprisingly, this saves minutes in some of the proofs that use a lot of
   invariants of the form (l = P-:n --> P). *)
fun shallow_simp ctxt =
  let val ctxt' = Config.put simp_depth_limit 2 ctxt in
    TRY o safe_asm_full_simp_tac ctxt'
  end

fun create_vcs ctxt i =
  let val main_simp_thms = rev (Named_Theorems.get ctxt @{named_theorems cterms_env})
      val ctermssl_cases = rev (Named_Theorems.get ctxt @{named_theorems ctermssl_cases})
  in
    dresolve_tac ctxt @{thms has_ctermssl} i
    THEN_ELSE (dmatch_tac ctxt ctermssl_cases i
      THEN
      TRY (REPEAT_ALL_NEW (ematch_tac ctxt [@[thm disjE]]) i)
      THEN
      PARALLEL_ALLGOALS
        (fn i => simp_only main_simp_thms ctxt i
          THEN TRY (REPEAT_ALL_NEW (ematch_tac ctxt [@[thm disjE]]) i)), all_tac)
  end

fun try_invs ctxt =
  let val inv_thms = rev (Named_Theorems.get ctxt @{named_theorems cterms_invs})
      fun fapp thm =
        TRY o (EVERY' (forward_tac ctxt [thm] :: replicate (Thm.nprems_of thm - 1) (assume_tac ctxt)))
  in
    EVERY' (map fapp inv_thms)
  end

fun try_final ctxt =
  let val final_thms = rev (Named_Theorems.get ctxt @{named_theorems cterms_final})
      fun eapp thm = EVERY' (eresolve_tac ctxt [thm] :: replicate (Thm.nprems_of thm - 1) (assume_tac ctxt))
  in
    TRY o (FIRST' (map eapp final_thms))
  end

fun each ctxt =
  (EVERY' ((ematch_tac ctxt (rev (Named_Theorems.get ctxt @{named_theorems cterms_elimders})) ::
    replicate 2 (assume_tac ctxt)))
  THEN' simp_only @{thms labels_psimps} ctxt
  THEN' (ematch_tac ctxt (rev (Named_Theorems.get ctxt @{named_theorems cterms_seqte}))
    THEN_ALL_NEW
      (fn j => simp_only [@[thm mem_Collect_eq]] ctxt j
        THEN REPEAT (eresolve_tac ctxt @{thms exE} j)
        THEN REPEAT (eresolve_tac ctxt @{thms conjE} j))))
  ORELSE' (SOLVED' (clarsimp_tac ctxt))

fun simp_all ctxt =
  let val ctxt' =
      ctxt |> fold Splitter.add_split [@[thm if_split_asm]]
  in
    PARALLEL_ALLGOALS (shallow_simp ctxt)
    THEN
    TRY (CHANGED_PROP (PARALLEL_ALLGOALS (asm_full_simp_tac ctxt' THEN' try_final ctxt)))
  end

fun intro_and_invs ctxt i =
  let val cterms_intros = rev (Named_Theorems.get ctxt @{named_theorems cterms_intros})
  in
    match_tac ctxt cterms_intros i
  end

```

```

    THEN PARALLEL_ALLGOALS (try_invs ctxt)
  end

fun process_vcs ctxt _ =
  ALLGOALS (create_vcs ctxt ORELSE' (SOLVED' (clarsimp_tac ctxt)))
  THEN PARALLEL_ALLGOALS (TRY o each ctxt)
)

method_setup inv_ctypems = (
  let
    val intro_onlyN = "intro_only"
    val vcs_onlyN = "vcs_only"
    val invN = "inv"
    val solveN = "solve"

    val inv_ctypems_options =
      (Args.parens (Args.$$$ intro_onlyN) >> K intro_and_invs ||
       Args.parens (Args.$$$ vcs_onlyN) >> K (fn ctxt => intro_and_invs ctxt
                                               THEN' process_vcs ctxt) ||
       Scan.succeed (fn ctxt => intro_and_invs ctxt
                     THEN' process_vcs ctxt
                     THEN' K (simp_all ctxt)))
  in
    (Scan.lift inv_ctypems_options --| Method.sections
     ((Args.$$$ invN -- Args.add -- Args.colon >>
      K (Method.modifier (Named_Theorems.add @{named_theorems cterms_invs}) here))
     :: (Args.$$$ solveN -- Args.colon >>
      K (Method.modifier (Named_Theorems.add @{named_theorems cterms_final}) here))
     :: Simplifier.simp_modifiers)
     >> (fn tac => SIMPLE_METHOD' o tac))
  end
) "solve invariants by considering all (interesting) control terms"

declare
  insert_iff [ctypems_env]
  Un_insert_right [ctypems_env]
  sup_bot_right [ctypems_env]
  Product_Type.prod_cases [ctypems_env]
  cterms1.simps [ctypems_env]

```

end

## 10 Configure the inv-ctypems tactic for sequential processes

```

theory Awn_SOS_Labels
imports Awn_SOS Inv_Cterms
begin

```

```

lemma elimder_guard:
  assumes "p = {l}<fg> qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}<fg> p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemma elimder_assign:
  assumes "p = {l}[[fa]] qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}[[fa]] p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemma elimder_ucast:
  assumes "p = {l}unicast(fip, fmsg).q1 ▷ q2"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' pp' where "p = {l}unicast(fip, fmsg).p' ▷ pp'"
    and "case a of unicast _ _ ⇒ l' ∈ labels Γ q1
          | _ ⇒ l' ∈ labels Γ q2"
  using assms by simp (erule seqpTEs, auto)

```

```

lemma elimder_bcast:
  assumes "p = {l}broadcast(fmsg).qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}broadcast(fmsg).p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemma elimder_gcast:
  assumes "p = {l}groupcast(fips, fmsg).qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}groupcast(fips, fmsg).p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemma elimder_send:
  assumes "p = {l}send(fmsg).qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}send(fmsg).p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemma elimder_deliver:
  assumes "p = {l}deliver(fdata).qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}deliver(fdata).p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemma elimder_receive:
  assumes "p = {l}receive(fmsg).qq"
    and "l' ∈ labels Γ q"
    and "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  obtains p' where "p = {l}receive(fmsg).p'"
    and "l' ∈ labels Γ qq"
  using assms by auto

```

```

lemmas elimders =
  elimder_guard
  elimder_assign
  elimder_ucast
  elimder_bcast
  elimder_gcast
  elimder_send
  elimder_deliver
  elimder_receive

```

```

declare
  seqpTEs [cterms_seqte]
  elimders [cterms_elimders]

```

```

end

```

## 11 Lemmas for partial networks

```
theory Pnet
imports AWN_SOS Invariants
begin
```

These lemmas mostly concern the preservation of node structure by `pnet_sos` transitions.

```
lemma pnet_maintains_dom:
  assumes "(s, a, s') ∈ trans (pnet np p)"
  shows "net_ips s = net_ips s'"
  using assms proof (induction p arbitrary: s a s')
    fix i R σ s a s'
    assume "(s, a, s') ∈ trans (pnet np ⟨i; R⟩)"
    hence "(s, a, s') ∈ node_sos (trans (np i))" ..
    thus "net_ips s = net_ips s'"
      by (rule node_sos.cases) simp_all
  next
    fix p1 p2 s a s'
    assume "∧s a s'. (s, a, s') ∈ trans (pnet np p1) ⇒ net_ips s = net_ips s'"
      and "∧s a s'. (s, a, s') ∈ trans (pnet np p2) ⇒ net_ips s = net_ips s'"
      and "(s, a, s') ∈ trans (pnet np (p1 || p2))"
    thus "net_ips s = net_ips s'"
      by simp (erule pnet_sos.cases, simp_all)
  qed
```

```
lemma pnet_net_ips_net_tree_ips [elim]:
  assumes "s ∈ reachable (pnet np p) I"
  shows "net_ips s = net_tree_ips p"
  using assms proof induction
    fix s
    assume "s ∈ init (pnet np p)"
    thus "net_ips s = net_tree_ips p"
    proof (induction p arbitrary: s)
      fix i R s
      assume "s ∈ init (pnet np ⟨i; R⟩)"
      then obtain ns where "s = NodeS i ns R" ..
      thus "net_ips s = net_tree_ips ⟨i; R⟩"
        by simp
    next
      fix p1 p2 s
      assume IH1: "∧s. s ∈ init (pnet np p1) ⇒ net_ips s = net_tree_ips p1"
        and IH2: "∧s. s ∈ init (pnet np p2) ⇒ net_ips s = net_tree_ips p2"
        and "s ∈ init (pnet np (p1 || p2))"
      from this(3) obtain s1 s2 where "s1 ∈ init (pnet np p1)"
        and "s2 ∈ init (pnet np p2)"
        and "s = SubnetS s1 s2" by auto
      from this(1-2) have "net_ips s1 = net_tree_ips p1"
        and "net_ips s2 = net_tree_ips p2"
        using IH1 IH2 by auto
      with ⟨s = SubnetS s1 s2⟩ show "net_ips s = net_tree_ips (p1 || p2)" by auto
    qed
  next
    fix s a s'
    assume "(s, a, s') ∈ trans (pnet np p)"
      and "net_ips s = net_tree_ips p"
    from this(1) have "net_ips s = net_ips s'"
      by (rule pnet_maintains_dom)
    with ⟨net_ips s = net_tree_ips p⟩ show "net_ips s' = net_tree_ips p"
      by simp
  qed
```

```
lemma pnet_init_net_ips_net_tree_ips:
  assumes "s ∈ init (pnet np p)"
  shows "net_ips s = net_tree_ips p"
```



```

using assms(1) by (rule reachable_init [THEN pnet_net_ips_net_tree_ips])

lemma pnet_init_in_net_ips_in_net_tree_ips [elim]:
  assumes "s ∈ init (pnet np p)"
    and "i ∈ net_ips s"
  shows "i ∈ net_tree_ips p"
using assms by (clarsimp dest!: pnet_init_net_ips_net_tree_ips)

lemma pnet_init_in_net_tree_ips_in_net_ips [elim]:
  assumes "s ∈ init (pnet np p)"
    and "i ∈ net_tree_ips p"
  shows "i ∈ net_ips s"
using assms by (clarsimp dest!: pnet_init_net_ips_net_tree_ips)

lemma pnet_init_not_in_net_tree_ips_not_in_net_ips [elim]:
  assumes "s ∈ init (pnet np p)"
    and "i ∉ net_tree_ips p"
  shows "i ∉ net_ips s"
proof
  assume "i ∈ net_ips s"
  with assms(1) have "i ∈ net_tree_ips p" ..
  with assms(2) show False ..
qed

lemma net_node_reachable_is_node:
  assumes "st ∈ reachable (pnet np ⟨ii; Ri⟩) I"
  shows "∃ ns R. st = NodeS ii ns R"
using assms proof induct
  fix s
  assume "s ∈ init (pnet np ⟨ii; Ri⟩)"
  thus "∃ ns R. s = NodeS ii ns R"
    by (rule pnet_node_init') simp
next
  fix s a s'
  assume "s ∈ reachable (pnet np ⟨ii; Ri⟩) I"
    and "∃ ns R. s = NodeS ii ns R"
    and "(s, a, s') ∈ trans (pnet np ⟨ii; Ri⟩)"
    and "I a"
  thus "∃ ns R. s' = NodeS ii ns R"
    by (auto simp add: trans_node_comp dest!: node_sos_dest)
qed

lemma partial_net_preserves_subnets:
  assumes "(SubnetS s t, a, st') ∈ pnet_sos (trans (pnet np p1)) (trans (pnet np p2))"
  shows "∃ s' t'. st' = SubnetS s' t'"
using assms by cases simp_all

lemma net_par_reachable_is_subnet:
  assumes "st ∈ reachable (pnet np (p1 || p2)) I"
  shows "∃ s t. st = SubnetS s t"
using assms by induct (auto dest!: partial_net_preserves_subnets)

lemma reachable_par_subnet_induct [consumes, case_names init step]:
  assumes "SubnetS s t ∈ reachable (pnet np (p1 || p2)) I"
    and init: "∧ s t. SubnetS s t ∈ init (pnet np (p1 || p2)) ⇒ P s t"
    and step: "∧ s t s' t' a. [
      SubnetS s t ∈ reachable (pnet np (p1 || p2)) I;
      P s t; (SubnetS s t, a, SubnetS s' t') ∈ (trans (pnet np (p1 || p2))); I a ]
      ⇒ P s' t'"
  shows "P s t"
using assms(1) proof (induction "SubnetS s t" arbitrary: s t)
  fix s t
  assume "SubnetS s t ∈ init (pnet np (p1 || p2))"
  with init show "P s t" .

```

```

next
  fix st a s' t'
  assume "st ∈ reachable (pnet np (p1 || p2)) I"
    and tr: "(st, a, SubnetS s' t') ∈ trans (pnet np (p1 || p2))"
    and "I a"
    and IH: "∧s t. st = SubnetS s t ⇒ P s t"
  from this(1) obtain s t where "st = SubnetS s t"
    and str: "SubnetS s t ∈ reachable (pnet np (p1 || p2)) I"
  by (metis net_par_reachable_is_subnet)
note this(2)
moreover from IH and ⟨st = SubnetS s t⟩ have "P s t" .
moreover from ⟨st = SubnetS s t⟩ and tr
  have "(SubnetS s t, a, SubnetS s' t') ∈ trans (pnet np (p1 || p2))" by simp
ultimately show "P s' t'"
  using ⟨I a⟩ by (rule step)
qed

```

lemma subnet\_reachable:

```

assumes "SubnetS s1 s2 ∈ reachable (pnet np (p1 || p2)) TT"
shows "s1 ∈ reachable (pnet np p1) TT"
      "s2 ∈ reachable (pnet np p2) TT"

```

proof -

```

from assms have "s1 ∈ reachable (pnet np p1) TT"
  ∧ s2 ∈ reachable (pnet np p2) TT"

```

proof (induction rule: reachable\_par\_subnet\_induct)

fix s1 s2

```
assume "SubnetS s1 s2 ∈ init (pnet np (p1 || p2))"
```

```
thus "s1 ∈ reachable (pnet np p1) TT"
```

```
  ∧ s2 ∈ reachable (pnet np p2) TT"
```

```
  by (auto dest: reachable_init)
```

next

```
case (step s1 s2 s1' s2' a)
```

```
hence "SubnetS s1 s2 ∈ reachable (pnet np (p1 || p2)) TT"
```

```
  and sr1: "s1 ∈ reachable (pnet np p1) TT"
```

```
  and sr2: "s2 ∈ reachable (pnet np p2) TT"
```

```
  and "(SubnetS s1 s2, a, SubnetS s1' s2') ∈ trans (pnet np (p1 || p2))" by auto
```

```
from this(4)
```

```
  have "(SubnetS s1 s2, a, SubnetS s1' s2') ∈ pnet_sos (trans (pnet np p1)) (trans (pnet np p2))"
```

```
  by simp
```

```
thus "s1' ∈ reachable (pnet np p1) TT"
```

```
  ∧ s2' ∈ reachable (pnet np p2) TT"
```

```
  by cases (insert sr1 sr2, auto elim: reachable_step)
```

qed

```
thus "s1 ∈ reachable (pnet np p1) TT"
```

```
  "s2 ∈ reachable (pnet np p2) TT" by auto
```

qed

lemma delivered\_to\_node [elim]:

```
assumes "s ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
```

```
  and "(s, i:deliver(d), s') ∈ trans (pnet np ⟨ii; Ri⟩)"
```

```
shows "i = ii"
```

proof -

```
from assms(1) obtain P R where "s = NodeS ii P R"
```

```
  by (metis net_node_reachable_is_node)
```

```
with assms(2) show "i = ii"
```

```
  by (clarsimp simp add: trans_node_comp elim!: node_deliverTE')
```

qed

lemma delivered\_to\_net\_ips:

```
assumes "s ∈ reachable (pnet np p) TT"
```

```
  and "(s, i:deliver(d), s') ∈ trans (pnet np p)"
```

```
shows "i ∈ net_ips s"
```

```
using assms proof (induction p arbitrary: s s')
```

```
  fix ii Ri s s'
```

```

assume sr: "s ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
  and "(s, i:deliver(d), s') ∈ trans (pnet np ⟨ii; Ri⟩)"
from this(2) have tr: "(s, i:deliver(d), s') ∈ node_sos (trans (np ii))" by simp
from sr obtain P R where [simp]: "s = NodeS ii P R"
  by (metis net_node_reachable_is_node)
moreover from tr obtain P' R' where [simp]: "s' = NodeS ii P' R'"
  by simp (metis node_sos_dest)
ultimately have "i = ii" using tr by auto
thus "i ∈ net_ips s" by simp
next
fix p1 p2 s s'
assume IH1: "∧s s'. [ s ∈ reachable (pnet np p1) TT;
  (s, i:deliver(d), s') ∈ trans (pnet np p1) ] ⇒ i ∈ net_ips s"
  and IH2: "∧s s'. [ s ∈ reachable (pnet np p2) TT;
  (s, i:deliver(d), s') ∈ trans (pnet np p2) ] ⇒ i ∈ net_ips s"
  and sr: "s ∈ reachable (pnet np (p1 || p2)) TT"
  and tr: "(s, i:deliver(d), s') ∈ trans (pnet np (p1 || p2))"
from tr have "(s, i:deliver(d), s') ∈ pnet_sos (trans (pnet np p1)) (trans (pnet np p2))"
  by simp
thus "i ∈ net_ips s"
proof (rule partial_deliverTE)
  fix s1 s1' s2
  assume "s = SubnetS s1 s2"
    and "s' = SubnetS s1' s2"
    and tr: "(s1, i:deliver(d), s1') ∈ trans (pnet np p1)"
  from sr have "s1 ∈ reachable (pnet np p1) TT"
    by (auto simp only: ⟨s = SubnetS s1 s2⟩ elim: subnet_reachable)
  hence "i ∈ net_ips s1" using tr by (rule IH1)
  thus "i ∈ net_ips s" by (simp add: ⟨s = SubnetS s1 s2⟩)
next
fix s2 s2' s1
assume "s = SubnetS s1 s2"
  and "s' = SubnetS s1 s2'"
  and tr: "(s2, i:deliver(d), s2') ∈ trans (pnet np p2)"
  from sr have "s2 ∈ reachable (pnet np p2) TT"
    by (auto simp only: ⟨s = SubnetS s1 s2⟩ elim: subnet_reachable)
  hence "i ∈ net_ips s2" using tr by (rule IH2)
  thus "i ∈ net_ips s" by (simp add: ⟨s = SubnetS s1 s2⟩)
qed
qed

lemma wf_net_tree_net_ips_disjoint [elim]:
  assumes "wf_net_tree (p1 || p2)"
    and "s1 ∈ reachable (pnet np p1) S"
    and "s2 ∈ reachable (pnet np p2) S"
  shows "net_ips s1 ∩ net_ips s2 = {}"
proof -
  from ⟨wf_net_tree (p1 || p2)⟩ have "net_tree_ips p1 ∩ net_tree_ips p2 = {}" by auto
  moreover from assms(2) have "net_ips s1 = net_tree_ips p1" ..
  moreover from assms(3) have "net_ips s2 = net_tree_ips p2" ..
  ultimately show ?thesis by simp
qed

lemma init_mapstate_Some_aadv_init [elim]:
  assumes "s ∈ init (pnet np p)"
    and "netmap s i = Some v"
  shows "v ∈ init (np i)"
using assms proof (induction p arbitrary: s)
  fix ii R s
  assume "s ∈ init (pnet np ⟨ii; R⟩)"
    and "netmap s i = Some v"
  from this(1) obtain ns where s: "s = NodeS ii ns R"
    and ns: "ns ∈ init (np ii)" ..
  from s and ⟨netmap s i = Some v⟩ have "i = ii"

```

```

    by simp (metis domI domIff)
  with s ns show "v ∈ init (np i)"
    using ⟨netmap s i = Some v⟩ by simp
next
  fix p1 p2 s
  assume IH1: "∧s. s ∈ init (pnet np p1) ⇒ netmap s i = Some v ⇒ v ∈ init (np i)"
    and IH2: "∧s. s ∈ init (pnet np p2) ⇒ netmap s i = Some v ⇒ v ∈ init (np i)"
    and "s ∈ init (pnet np (p1 || p2))"
    and "netmap s i = Some v"
  from this(3) obtain s1 s2 where "s = SubnetS s1 s2"
    and "s1 ∈ init (pnet np p1)"
    and "s2 ∈ init (pnet np p2)" by auto
  from this(1) and ⟨netmap s i = Some v⟩
    have "netmap s1 i = Some v ∨ netmap s2 i = Some v" by auto
  thus "v ∈ init (np i)"
  proof
    assume "netmap s1 i = Some v"
    with ⟨s1 ∈ init (pnet np p1)⟩ show ?thesis by (rule IH1)
  next
    assume "netmap s2 i = Some v"
    with ⟨s2 ∈ init (pnet np p2)⟩ show ?thesis by (rule IH2)
  qed
qed

```

lemma reachable\_connect\_netmap [elim]:

```

  assumes "s ∈ reachable (pnet np n) TT"
    and "(s, connect(i, i'), s') ∈ trans (pnet np n)"
  shows "netmap s' = netmap s"
  using assms proof (induction n arbitrary: s s')
    fix ii Ri s s'
    assume sr: "s ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
      and "(s, connect(i, i'), s') ∈ trans (pnet np ⟨ii; Ri⟩)"
    from this(2) have tr: "(s, connect(i, i'), s') ∈ node_sos (trans (np ii))" ..
    from sr obtain p R where "s = NodeS ii p R"
      by (metis net_node_reachable_is_node)
    with tr show "netmap s' = netmap s"
      by (auto elim!: node_sos.cases)
  next
    fix p1 p2 s s'
    assume IH1: "∧s s'. [ s ∈ reachable (pnet np p1) TT;
      (s, connect(i, i'), s') ∈ trans (pnet np p1) ] ⇒ netmap s' = netmap s"
      and IH2: "∧s s'. [ s ∈ reachable (pnet np p2) TT;
      (s, connect(i, i'), s') ∈ trans (pnet np p2) ] ⇒ netmap s' = netmap s"
      and sr: "s ∈ reachable (pnet np (p1 || p2)) TT"
      and tr: "(s, connect(i, i'), s') ∈ trans (pnet np (p1 || p2))"
    from tr have "(s, connect(i, i'), s') ∈ pnet_sos (trans (pnet np p1)) (trans (pnet np p2))"
      by simp
    thus "netmap s' = netmap s"
  proof cases
    fix s1 s1' s2 s2'
    assume "s = SubnetS s1 s2"
      and "s' = SubnetS s1' s2'"
      and tr1: "(s1, connect(i, i'), s1') ∈ trans (pnet np p1)"
      and tr2: "(s2, connect(i, i'), s2') ∈ trans (pnet np p2)"
    from this(1) and sr
      have "SubnetS s1 s2 ∈ reachable (pnet np (p1 || p2)) TT" by simp
    hence sr1: "s1 ∈ reachable (pnet np p1) TT"
      and sr2: "s2 ∈ reachable (pnet np p2) TT"
      by (auto intro: subnet_reachable)
    from sr1 tr1 have "netmap s1' = netmap s1" by (rule IH1)
    moreover from sr2 tr2 have "netmap s2' = netmap s2" by (rule IH2)
    ultimately show "netmap s' = netmap s"
      using ⟨s = SubnetS s1 s2⟩ and ⟨s' = SubnetS s1' s2'⟩ by simp
  qed simp_all

```

qed

```
lemma reachable_disconnect_netmap [elim]:
  assumes "s ∈ reachable (pnet np n) TT"
    and "(s, disconnect(i, i'), s') ∈ trans (pnet np n)"
  shows "netmap s' = netmap s"
using assms proof (induction n arbitrary: s s')
  fix ii Ri s s'
  assume sr: "s ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
    and "(s, disconnect(i, i'), s') ∈ trans (pnet np ⟨ii; Ri⟩)"
  from this(2) have tr: "(s, disconnect(i, i'), s') ∈ node_sos (trans (np ii))" ..
  from sr obtain p R where "s = NodeS ii p R"
    by (metis net_node_reachable_is_node)
  with tr show "netmap s' = netmap s"
    by (auto elim!: node_sos.cases)
next
  fix p1 p2 s s'
  assume IH1: "∧s s'. [ s ∈ reachable (pnet np p1) TT;
    (s, disconnect(i, i'), s') ∈ trans (pnet np p1) ] ⇒ netmap s' = netmap s"
    and IH2: "∧s s'. [ s ∈ reachable (pnet np p2) TT;
    (s, disconnect(i, i'), s') ∈ trans (pnet np p2) ] ⇒ netmap s' = netmap s"
    and sr: "s ∈ reachable (pnet np (p1 || p2)) TT"
    and tr: "(s, disconnect(i, i'), s') ∈ trans (pnet np (p1 || p2))"
  from tr have "(s, disconnect(i, i'), s') ∈ pnet_sos (trans (pnet np p1)) (trans (pnet np p2))"
    by simp
  thus "netmap s' = netmap s"
proof cases
  fix s1 s1' s2 s2'
  assume "s = SubnetS s1 s2"
    and "s' = SubnetS s1' s2'"
    and tr1: "(s1, disconnect(i, i'), s1') ∈ trans (pnet np p1)"
    and tr2: "(s2, disconnect(i, i'), s2') ∈ trans (pnet np p2)"
  from this(1) and sr
    have "SubnetS s1 s2 ∈ reachable (pnet np (p1 || p2)) TT" by simp
  hence sr1: "s1 ∈ reachable (pnet np p1) TT"
    and sr2: "s2 ∈ reachable (pnet np p2) TT"
    by (auto intro: subnet_reachable)
  from sr1 tr1 have "netmap s1' = netmap s1" by (rule IH1)
  moreover from sr2 tr2 have "netmap s2' = netmap s2" by (rule IH2)
  ultimately show "netmap s' = netmap s"
    using ⟨s = SubnetS s1 s2⟩ and ⟨s' = SubnetS s1' s2'⟩ by simp
qed simp_all
qed
```

```
fun net_ip_action :: "(ip ⇒ ('s, 'm seq_action) automaton)
  ⇒ 'm node_action ⇒ ip ⇒ net_tree ⇒ 's net_state ⇒ 's net_state ⇒ bool"
```

where

```
"net_ip_action np a i (p1 || p2) (SubnetS s1 s2) (SubnetS s1' s2') =
  ((i ∈ net_ips s1 → ((s1, a, s1') ∈ trans (pnet np p1)
    ∧ s2' = s2 ∧ net_ip_action np a i p1 s1 s1'))
  ∧ (i ∈ net_ips s2 → ((s2, a, s2') ∈ trans (pnet np p2)
    ∧ s1' = s1 ∧ net_ip_action np a i p2 s2 s2')))"
| "net_ip_action np a i p s s' = True"
```

```
lemma pnet_tau_single_node [elim]:
```

```
  assumes "wf_net_tree p"
    and "s ∈ reachable (pnet np p) TT"
    and "(s, τ, s') ∈ trans (pnet np p)"
  shows "∃i∈net_ips s. ((∀j. j≠i → netmap s' j = netmap s j)
    ∧ net_ip_action np τ i p s s')"
```

using assms proof (induction p arbitrary: s s')

```
  fix ii Ri s s'
  assume "s ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
    and "(s, τ, s') ∈ trans (pnet np ⟨ii; Ri⟩)"
```

```

from this obtain p R p' R' where "s = NodeS ii p R" and "s' = NodeS ii p' R'"
  by (metis (hide_lams, no_types) TT_True net_node_reachable_is_node
        reachable_step)

hence "net_ips s = {ii}"
  and "net_ips s' = {ii}" by simp_all
hence " $\exists i \in \text{dom}(\text{netmap } s). \forall j. j \neq i \rightarrow \text{netmap } s' j = \text{netmap } s j$ "
  by (simp add: net_ips_is_dom_netmap)
thus " $\exists i \in \text{net\_ips } s. (\forall j. j \neq i \rightarrow \text{netmap } s' j = \text{netmap } s j)
      \wedge \text{net\_ip\_action } np \tau i (\langle ii; R_i \rangle) s s'$ "
  by (simp add: net_ips_is_dom_netmap)

next
fix p1 p2 s s'
assume IH1: " $\wedge s s'. \llbracket \text{wf\_net\_tree } p1;
      s \in \text{reachable } (\text{pnet } np p1) TT;
      (s, \tau, s') \in \text{trans } (\text{pnet } np p1) \rrbracket
      \Rightarrow \exists i \in \text{net\_ips } s. (\forall j. j \neq i \rightarrow \text{netmap } s' j = \text{netmap } s j)
      \wedge \text{net\_ip\_action } np \tau i p1 s s'$ "

  and IH2: " $\wedge s s'. \llbracket \text{wf\_net\_tree } p2;
      s \in \text{reachable } (\text{pnet } np p2) TT;
      (s, \tau, s') \in \text{trans } (\text{pnet } np p2) \rrbracket
      \Rightarrow \exists i \in \text{net\_ips } s. (\forall j. j \neq i \rightarrow \text{netmap } s' j = \text{netmap } s j)
      \wedge \text{net\_ip\_action } np \tau i p2 s s'$ "

  and sr: "s  $\in$  reachable (pnet np (p1 || p2)) TT"
  and "wf_net_tree (p1 || p2)"
  and tr: "(s,  $\tau$ , s')  $\in$  trans (pnet np (p1 || p2))"
from (wf_net_tree (p1 || p2)) have "net_tree_ips p1  $\cap$  net_tree_ips p2 = {}"
  and "wf_net_tree p1"
  and "wf_net_tree p2" by auto
from tr have "(s,  $\tau$ , s')  $\in$  pnet_sos (trans (pnet np p1)) (trans (pnet np p2))" by simp
thus " $\exists i \in \text{net\_ips } s. (\forall j. j \neq i \rightarrow \text{netmap } s' j = \text{netmap } s j)
      \wedge \text{net\_ip\_action } np \tau i (p1 || p2) s s'$ "

proof cases
fix s1 s1' s2
assume subs: "s = SubnetS s1 s2"
  and subs': "s' = SubnetS s1' s2'"
  and tr1: "(s1,  $\tau$ , s1')  $\in$  trans (pnet np p1)"
from sr have sr1: "s1  $\in$  reachable (pnet np p1) TT"
  and "s2  $\in$  reachable (pnet np p2) TT"
  by (simp_all only: subs) (erule subnet_reachable)+
with (net_tree_ips p1  $\cap$  net_tree_ips p2 = {}) have "dom(netmap s1)  $\cap$  dom(netmap s2) = {}"
  by (metis net_ips_is_dom_netmap pnet_net_ips_net_tree_ips)
from (wf_net_tree p1) sr1 tr1 obtain i where "i  $\in$  dom(netmap s1)"
  and *: " $\forall j. j \neq i \rightarrow \text{netmap } s1' j = \text{netmap } s1 j$ "
  and "net_ip_action np  $\tau$  i p1 s1 s1'"

  by (auto simp add: net_ips_is_dom_netmap dest!: IH1)
from this(1) and (dom(netmap s1)  $\cap$  dom(netmap s2) = {}) have "i  $\notin$  dom(netmap s2)"
  by auto
with subs subs' tr1 (net_ip_action np  $\tau$  i p1 s1 s1') have "net_ip_action np  $\tau$  i (p1 || p2) s s'"
  by (simp add: net_ips_is_dom_netmap)
moreover have " $\forall j. j \neq i \rightarrow (\text{netmap } s1' ++ \text{netmap } s2) j = (\text{netmap } s1 ++ \text{netmap } s2) j$ "
proof (intro allI impI)
  fix j
  assume "j  $\neq$  i"
  with * have "netmap s1' j = netmap s1 j" by simp
  thus "(netmap s1' ++ netmap s2) j = (netmap s1 ++ netmap s2) j"
    by (metis (hide_lams, mono_tags) map_add_dom_app_simps(1) map_add_dom_app_simps(3))
qed
ultimately show ?thesis using (i  $\in$  dom(netmap s1)) subs subs'
  by (auto simp add: net_ips_is_dom_netmap)

next
fix s2 s2' s1
assume subs: "s = SubnetS s1 s2"
  and subs': "s' = SubnetS s1 s2'"
  and tr2: "(s2,  $\tau$ , s2')  $\in$  trans (pnet np p2)"

```

```

from sr have "s1 ∈ reachable (pnet np p1) TT"
  and sr2: "s2 ∈ reachable (pnet np p2) TT"
  by (simp_all only: subs) (erule subnet_reachable)+
with ⟨net_tree_ips p1 ∩ net_tree_ips p2 = {}⟩ have "dom(netmap s1) ∩ dom(netmap s2) = {}"
  by (metis net_ips_is_dom_netmap pnet_net_ips_net_tree_ips)
from ⟨wf_net_tree p2⟩ sr2 tr2 obtain i where "i ∈ dom(netmap s2)"
  and *: "∀j. j ≠ i → netmap s2' j = netmap s2 j"
  and "net_ip_action np τ i p2 s2 s2'"
  by (auto simp add: net_ips_is_dom_netmap dest!: IH2)
from this(1) and ⟨dom(netmap s1) ∩ dom(netmap s2) = {}⟩ have "i ∉ dom(netmap s1)"
  by auto
with subs subs' tr2 ⟨net_ip_action np τ i p2 s2 s2'⟩ have "net_ip_action np τ i (p1 ∥ p2) s s'"
  by (simp add: net_ips_is_dom_netmap)
moreover have "∀j. j ≠ i → (netmap s1 ++ netmap s2') j = (netmap s1 ++ netmap s2) j"
proof (intro allI impI)
  fix j
  assume "j ≠ i"
  with * have "netmap s2' j = netmap s2 j" by simp
  thus "(netmap s1 ++ netmap s2') j = (netmap s1 ++ netmap s2) j"
  by (metis (hide_lams, mono_tags) domD map_add_Some_iff map_add_dom_app_simps(3))
qed
ultimately show ?thesis using ⟨i ∈ dom(netmap s2)⟩ subs subs'
  by (clarsimp simp add: net_ips_is_dom_netmap)
  (metis domI dom_map_add map_add_find_right)
qed simp_all
qed

```

lemma pnet\_deliver\_single\_node [elim]:

```

assumes "wf_net_tree p"
  and "s ∈ reachable (pnet np p) TT"
  and "(s, i:deliver(d), s') ∈ trans (pnet np p)"
shows "(∀j. j ≠ i → netmap s' j = netmap s j) ∧ net_ip_action np (i:deliver(d)) i p s s'"
  (is "?P p s s'")
using assms proof (induction p arbitrary: s s')
  fix ii Ri s s'
  assume sr: "s ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
  and tr: "(s, i:deliver(d), s') ∈ trans (pnet np ⟨ii; Ri⟩)"
  from this obtain p R p' R' where "s = NodeS ii p R" and "s' = NodeS ii p' R'"
  by (metis (hide_lams, no_types) TT_True net_node_reachable_is_node
    reachable_step)
  hence "net_ips s = {ii}"
  and "net_ips s' = {ii}" by simp_all
  hence "∀j. j ≠ ii → netmap s' j = netmap s j"
  by simp
  moreover from sr tr have "i = ii" by (rule delivered_to_node)
  ultimately show "(∀j. j ≠ i → netmap s' j = netmap s j)
    ∧ net_ip_action np (i:deliver(d)) i (⟨ii; Ri⟩) s s'"
  by simp

```

next

```

fix p1 p2 s s'
assume IH1: "∧s s'. [ wf_net_tree p1;
  s ∈ reachable (pnet np p1) TT;
  (s, i:deliver(d), s') ∈ trans (pnet np p1) ]
  ⇒ (∀j. j ≠ i → netmap s' j = netmap s j)
  ∧ net_ip_action np (i:deliver(d)) i p1 s s'"
and IH2: "∧s s'. [ wf_net_tree p2;
  s ∈ reachable (pnet np p2) TT;
  (s, i:deliver(d), s') ∈ trans (pnet np p2) ]
  ⇒ (∀j. j ≠ i → netmap s' j = netmap s j)
  ∧ net_ip_action np (i:deliver(d)) i p2 s s'"
and sr: "s ∈ reachable (pnet np (p1 ∥ p2)) TT"
and "wf_net_tree (p1 ∥ p2)"
and tr: "(s, i:deliver(d), s') ∈ trans (pnet np (p1 ∥ p2))"
from ⟨wf_net_tree (p1 ∥ p2)⟩ have "net_tree_ips p1 ∩ net_tree_ips p2 = {}"

```

```

        and "wf_net_tree p1"
        and "wf_net_tree p2" by auto
from tr have "(s, i:deliver(d), s') ∈ pnet_sos (trans (pnet np p1)) (trans (pnet np p2))" by simp
thus "(∀j. j ≠ i → netmap s' j = netmap s j)
  ∧ net_ip_action np (i:deliver(d)) i (p1 || p2) s s'"
proof cases
  fix s1 s1' s2
  assume subs: "s = SubnetS s1 s2"
    and subs': "s' = SubnetS s1' s2"
    and tr1: "(s1, i:deliver(d), s1') ∈ trans (pnet np p1)"
  from sr have sr1: "s1 ∈ reachable (pnet np p1) TT"
    and "s2 ∈ reachable (pnet np p2) TT"
  by (simp_all only: subs) (erule subnet_reachable)+
  with ⟨net_tree_ips p1 ∩ net_tree_ips p2 = {}⟩ have "dom(netmap s1) ∩ dom(netmap s2) = {}"
  by (metis net_ips_is_dom_netmap pnet_net_ips_net_tree_ips)
  moreover from sr1 tr1 have "i ∈ net_ips s1" by (rule delivered_to_net_ips)
  ultimately have "i ∉ dom(netmap s2)" by (auto simp add: net_ips_is_dom_netmap)

  from ⟨wf_net_tree p1⟩ sr1 tr1 have *: "∀j. j ≠ i → netmap s1' j = netmap s1 j"
    and "net_ip_action np (i:deliver(d)) i p1 s1 s1'"
    by (auto dest!: IH1)
  from subs subs' tr1 this(2) ⟨i ∉ dom(netmap s2)⟩
  have "net_ip_action np (i:deliver(d)) i (p1 || p2) s s'"
    by (simp add: net_ips_is_dom_netmap)
  moreover have "∀j. j ≠ i → (netmap s1' ++ netmap s2) j = (netmap s1 ++ netmap s2) j"
  proof (intro allI impI)
    fix j
    assume "j ≠ i"
    with * have "netmap s1' j = netmap s1 j" by simp
    thus "(netmap s1' ++ netmap s2) j = (netmap s1 ++ netmap s2) j"
      by (metis (hide_lams, mono_tags) map_add_dom_app_simps(1) map_add_dom_app_simps(3))
  qed
  ultimately show ?thesis using ⟨i ∈ net_ips s1⟩ subs subs' by auto
next
  fix s2 s2' s1
  assume subs: "s = SubnetS s1 s2"
    and subs': "s' = SubnetS s1 s2'"
    and tr2: "(s2, i:deliver(d), s2') ∈ trans (pnet np p2)"
  from sr have "s1 ∈ reachable (pnet np p1) TT"
    and sr2: "s2 ∈ reachable (pnet np p2) TT"
  by (simp_all only: subs) (erule subnet_reachable)+
  with ⟨net_tree_ips p1 ∩ net_tree_ips p2 = {}⟩ have "dom(netmap s1) ∩ dom(netmap s2) = {}"
  by (metis net_ips_is_dom_netmap pnet_net_ips_net_tree_ips)
  moreover from sr2 tr2 have "i ∈ net_ips s2" by (rule delivered_to_net_ips)
  ultimately have "i ∉ dom(netmap s1)" by (auto simp add: net_ips_is_dom_netmap)

  from ⟨wf_net_tree p2⟩ sr2 tr2 have *: "∀j. j ≠ i → netmap s2' j = netmap s2 j"
    and "net_ip_action np (i:deliver(d)) i p2 s2 s2'"
    by (auto dest!: IH2)
  from subs subs' tr2 this(2) ⟨i ∉ dom(netmap s1)⟩
  have "net_ip_action np (i:deliver(d)) i (p1 || p2) s s'"
    by (simp add: net_ips_is_dom_netmap)
  moreover have "∀j. j ≠ i → (netmap s1 ++ netmap s2') j = (netmap s1 ++ netmap s2) j"
  proof (intro allI impI)
    fix j
    assume "j ≠ i"
    with * have "netmap s2' j = netmap s2 j" by simp
    thus "(netmap s1 ++ netmap s2') j = (netmap s1 ++ netmap s2) j"
      by (metis (hide_lams, mono_tags) domD map_add_Some_iff map_add_dom_app_simps(3))
  qed
  ultimately show ?thesis using ⟨i ∈ net_ips s2⟩ subs subs' by auto
qed simp_all
qed

```



end

## 12 Lemmas for closed networks

```
theory Closed
imports Pnet
begin
```

```
lemma complete_net_preserves_subnets:
  assumes "(SubnetS s t, a, st') ∈ cnet_sos (pnet_sos (trans (pnet np p1)) (trans (pnet np p2)))"
  shows "∃ s' t'. st' = SubnetS s' t'"
  using assms by cases (auto dest: partial_net_preserves_subnets)

lemma complete_net_reachable_is_subnet:
  assumes "st ∈ reachable (closed (pnet np (p1 || p2))) I"
  shows "∃ s t. st = SubnetS s t"
  using assms by induction (auto dest!: complete_net_preserves_subnets)

lemma closed_reachable_par_subnet_induct [consumes, case_names init step]:
  assumes "SubnetS s t ∈ reachable (closed (pnet np (p1 || p2))) I"
  and init: "∧ s t. SubnetS s t ∈ init (closed (pnet np (p1 || p2))) ⇒ P s t"
  and step: "∧ s t s' t' a. [
    SubnetS s t ∈ reachable (closed (pnet np (p1 || p2))) I;
    P s t; (SubnetS s t, a, SubnetS s' t') ∈ trans (closed (pnet np (p1 || p2))); I a ]
    ⇒ P s' t'"
  shows "P s t"
  using assms(1) proof (induction "SubnetS s t" arbitrary: s t)
  fix s t
  assume "SubnetS s t ∈ init (closed (pnet np (p1 || p2)))"
  with init show "P s t" .
  next
  fix st a s' t'
  assume "st ∈ reachable (closed (pnet np (p1 || p2))) I"
  and tr: "(st, a, SubnetS s' t') ∈ trans (closed (pnet np (p1 || p2)))"
  and "I a"
  and IH: "∧ s t. st = SubnetS s t ⇒ P s t"
  from this(1) obtain s t where "st = SubnetS s t"
  and "SubnetS s t ∈ reachable (closed (pnet np (p1 || p2))) I"
  by (metis complete_net_reachable_is_subnet)
  note this(2)
  moreover from IH and ⟨st = SubnetS s t⟩ have "P s t" .
  moreover from tr and ⟨st = SubnetS s t⟩
  have "(SubnetS s t, a, SubnetS s' t') ∈ trans (closed (pnet np (p1 || p2)))" by simp
  ultimately show "P s' t'"
  using ⟨I a⟩ by (rule assms(3))
  qed
qed

lemma reachable_closed_reachable_pnet [elim]:
  assumes "s ∈ reachable (closed (pnet np n)) TT"
  shows "s ∈ reachable (pnet np n) TT"
  using assms proof (induction rule: reachable.induct)
  fix s s' a
  assume sr: "s ∈ reachable (pnet np n) TT"
  and "(s, a, s') ∈ trans (closed (pnet np n))"
  from this(2) have "(s, a, s') ∈ cnet_sos (trans (pnet np n))" by simp
  thus "s' ∈ reachable (pnet np n) TT"
  by cases (insert sr, auto elim!: reachable_step)
  qed (auto elim: reachable_init)

lemma closed_node_net_state [elim]:
  assumes "st ∈ reachable (closed (pnet np ⟨ii; Ri⟩)) TT"
  obtains ξ p q R where "st = NodeS ii ((ξ, p), q) R"
  using assms by (metis net_node_reachable_is_node reachable_closed_reachable_pnet surj_pair)
```

```

lemma closed_subnet_net_state [elim]:
  assumes "st ∈ reachable (closed (pnet np (p1 || p2))) TT"
  obtains s t where "st = SubnetS s t"
  using assms by (metis reachable_closed_reachable_pnet net_par_reachable_is_subnet)

lemma closed_imp_pnet_trans [elim, dest]:
  assumes "(s, a, s') ∈ trans (closed (pnet np n))"
  shows "∃ a'. (s, a', s') ∈ trans (pnet np n)"
  using assms by (auto elim!: cnet_sos.cases)

lemma reachable_not_in_net_tree_ips [elim]:
  assumes "s ∈ reachable (closed (pnet np n)) TT"
  and "i ∉ net_tree_ips n"
  shows "netmap s i = None"
  using assms proof induction
  fix s
  assume "s ∈ init (closed (pnet np n))"
  and "i ∉ net_tree_ips n"
  thus "netmap s i = None"
  proof (induction n arbitrary: s)
  fix ii R s
  assume "s ∈ init (closed (pnet np ⟨ii; R⟩))"
  and "i ∉ net_tree_ips ⟨ii; R⟩"
  from this(2) have "i ≠ ii" by simp
  moreover from ⟨s ∈ init (closed (pnet np ⟨ii; R⟩))⟩ obtain p where "s = NodeS ii p R"
  by simp (metis pnet_simps(1) pnet_node_init')
  ultimately show "netmap s i = None" by simp
  next
  fix p1 p2 s
  assume IH1: "∧s. s ∈ init (closed (pnet np p1)) ⇒ i ∉ net_tree_ips p1
  ⇒ netmap s i = None"
  and IH2: "∧s. s ∈ init (closed (pnet np p2)) ⇒ i ∉ net_tree_ips p2
  ⇒ netmap s i = None"
  and "s ∈ init (closed (pnet np (p1 || p2)))"
  and "i ∉ net_tree_ips (p1 || p2)"
  from this(3) obtain s1 s2 where "s = SubnetS s1 s2"
  and "s1 ∈ init (closed (pnet np p1))"
  and "s2 ∈ init (closed (pnet np p2))" by simp metis
  moreover from ⟨i ∉ net_tree_ips (p1 || p2)⟩ have "i ∉ net_tree_ips p1"
  and "i ∉ net_tree_ips p2" by auto
  ultimately have "netmap s1 i = None"
  and "netmap s2 i = None"
  using IH1 IH2 by auto
  with ⟨s = SubnetS s1 s2⟩ show "netmap s i = None" by simp
  qed
  next
  fix s a s'
  assume sr: "s ∈ reachable (closed (pnet np n)) TT"
  and tr: "(s, a, s') ∈ trans (closed (pnet np n))"
  and IH: "i ∉ net_tree_ips n ⇒ netmap s i = None"
  and "i ∉ net_tree_ips n"
  from this(3-4) have "i ∉ net_ips s" by auto
  with tr have "i ∉ net_ips s'"
  by simp (erule cnet_sos.cases, (metis net_ips_is_dom_netmap pnet_maintains_dom)+)
  thus "netmap s' i = None" by simp
  qed

lemma closed_pnet_aadv_init [elim]:
  assumes "s ∈ init (closed (pnet np n))"
  and "i ∈ net_tree_ips n"
  shows "the (netmap s i) ∈ init (np i)"
  using assms proof (induction n arbitrary: s)
  fix ii R s
  assume "s ∈ init (closed (pnet np ⟨ii; R⟩))"

```

```

    and "i ∈ net_tree_ips ⟨ii; R⟩"
  hence "s ∈ init (pnet np ⟨i; R⟩)" by simp
  then obtain p where "s = NodeS i p R"
    and "p ∈ init (np i)" ..
  with ⟨s = NodeS i p R⟩ have "netmap s = [i ↦ p]" by simp
  with ⟨p ∈ init (np i)⟩ show "the (netmap s i) ∈ init (np i)" by simp
next
fix p1 p2 s
assume IH1: "∧s. s ∈ init (closed (pnet np p1)) ⇒
  i ∈ net_tree_ips p1 ⇒ the (netmap s i) ∈ init (np i)"
  and IH2: "∧s. s ∈ init (closed (pnet np p2)) ⇒
  i ∈ net_tree_ips p2 ⇒ the (netmap s i) ∈ init (np i)"
  and "s ∈ init (closed (pnet np (p1 || p2)))"
  and "i ∈ net_tree_ips (p1 || p2)"
from this(3) obtain s1 s2 where "s = SubnetS s1 s2"
  and "s1 ∈ init (closed (pnet np p1))"
  and "s2 ∈ init (closed (pnet np p2))"

  by auto
from this(2) have "net_tree_ips p1 = net_ips s1"
  by (clarsimp dest!: pnet_init_net_ips_net_tree_ips)
from ⟨s2 ∈ init (closed (pnet np p2))⟩ have "net_tree_ips p2 = net_ips s2"
  by (clarsimp dest!: pnet_init_net_ips_net_tree_ips)
show "the (netmap s i) ∈ init (np i)"
proof (cases "i ∈ net_tree_ips p2")
  assume "i ∈ net_tree_ips p2"
  with ⟨s2 ∈ init (closed (pnet np p2))⟩ have "the (netmap s2 i) ∈ init (np i)"
    by (rule IH2)
  moreover from ⟨i ∈ net_tree_ips p2⟩ and ⟨net_tree_ips p2 = net_ips s2⟩
    have "i ∈ net_ips s2" by simp
  ultimately show ?thesis
    using ⟨s = SubnetS s1 s2⟩ by (auto simp add: net_ips_is_dom_netmap)
next
assume "i ∉ net_tree_ips p2"
with ⟨i ∈ net_tree_ips (p1 || p2)⟩ have "i ∈ net_tree_ips p1" by simp
with ⟨s1 ∈ init (closed (pnet np p1))⟩ have "the (netmap s1 i) ∈ init (np i)"
  by (rule IH1)
moreover from ⟨i ∈ net_tree_ips p1⟩ and ⟨net_tree_ips p1 = net_ips s1⟩
  have "i ∈ net_ips s1" by simp
moreover from ⟨i ∉ net_tree_ips p2⟩ and ⟨net_tree_ips p2 = net_ips s2⟩
  have "i ∉ net_ips s2" by simp
ultimately show ?thesis
  using ⟨s = SubnetS s1 s2⟩
  by (simp add: map_add_dom_app_simps net_ips_is_dom_netmap)
qed
qed

```

end

## 13 Open semantics of the Algebra of Wireless Networks

```

theory OAWN_SOS
imports TransitionSystems AWN
begin

```

These are variants of the SOS rules that work against a mixed global/local context, where the global context is represented by a function  $\sigma$  mapping ip addresses to states.

### 13.1 Open structural operational semantics for sequential process expressions

```

inductive_set
  oseqp_sos
  :: "('s, 'm, 'p, 'l) seqp_env ⇒ ip
  ⇒ ((ip ⇒ 's) × ('s, 'm, 'p, 'l) seqp, 'm seq_action) transition set"

```

```

for  $\Gamma :: "(s, m, p, l) \text{ seqp\_env}"$ 
and  $i :: ip$ 
where
  obroadcastT: " $\sigma' i = \sigma i \implies$ 
    ( $(\sigma, \{l\}\text{broadcast}(s_{msg}).p)$ ),          broadcast ( $s_{msg} (\sigma i)$ ),          ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "
  / ogroupcastT: " $\sigma' i = \sigma i \implies$ 
    ( $(\sigma, \{l\}\text{groupcast}(s_{ips}, s_{msg}).p)$ ), groupcast ( $s_{ips} (\sigma i)$ ) ( $s_{msg} (\sigma i)$ ), ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "
  / ounicastT: " $\sigma' i = \sigma i \implies$ 
    ( $(\sigma, \{l\}\text{unicast}(s_{ip}, s_{msg}).p \triangleright q)$ ), unicast ( $s_{ip} (\sigma i)$ ) ( $s_{msg} (\sigma i)$ ),          ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "
  / onotunicastT: " $\sigma' i = \sigma i \implies$ 
    ( $(\sigma, \{l\}\text{unicast}(s_{ip}, s_{msg}).p \triangleright q)$ ),  $\neg$ unicast ( $s_{ip} (\sigma i)$ ),          ( $\sigma', q$ )  $\in$  oseqp
 $\Gamma i$ "
  / osendT: " $\sigma' i = \sigma i \implies$ 
    ( $(\sigma, \{l\}\text{send}(s_{msg}).p)$ ),          send ( $s_{msg} (\sigma i)$ ),          ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "
  / odeliverT: " $\sigma' i = \sigma i \implies$ 
    ( $(\sigma, \{l\}\text{deliver}(s_{data}).p)$ ),          deliver ( $s_{data} (\sigma i)$ ),          ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "
  / oreceiveT: " $\sigma' i = u_{msg} \text{ msg} (\sigma i) \implies$ 
    ( $(\sigma, \{l\}\text{receive}(u_{msg}).p)$ ),          receive msg,          ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "
  / oassignT: " $\sigma' i = u (\sigma i) \implies$ 
    ( $(\sigma, \{l\}[u] p)$ ),           $\tau$ ,          ( $\sigma', p$ )  $\in$  oseqp
 $\Gamma i$ "

  / ocallT: " $((\sigma, \Gamma pn), a, (\sigma', p')) \in \text{oseqp\_sos } \Gamma i \implies$ 
    ( $(\sigma, \text{call}(pn)), a, (\sigma', p')$ )  $\in$  oseqp\_sos  $\Gamma i$ "

  / ochoiceT1: " $((\sigma, p), a, (\sigma', p')) \in \text{oseqp\_sos } \Gamma i \implies$ 
    ( $(\sigma, p \oplus q), a, (\sigma', p')$ )  $\in$  oseqp\_sos  $\Gamma i$ "
  / ochoiceT2: " $((\sigma, q), a, (\sigma', q')) \in \text{oseqp\_sos } \Gamma i \implies$ 
    ( $(\sigma, p \oplus q), a, (\sigma', q')$ )  $\in$  oseqp\_sos  $\Gamma i$ "

  / oguardT: " $\sigma' i \in g (\sigma i) \implies ((\sigma, \{l\}\langle g \rangle p), \tau, (\sigma', p)) \in \text{oseqp\_sos } \Gamma i$ "

inductive_cases
  oseq_callTE [elim]: " $((\sigma, \text{call}(pn)), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
  and oseq_choiceTE [elim]: " $((\sigma, p1 \oplus p2), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "

lemma oseq_broadcastTE [elim]:
  " $[((\sigma, \{l\}\text{broadcast}(s_{msg}). p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i;$ 
   $[a = \text{broadcast } (s_{msg} (\sigma i)); \sigma' i = \sigma i; q = p] \implies P] \implies P$ "
  by (ind_cases " $((\sigma, \{l\}\text{broadcast}(s_{msg}). p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ ") simp

lemma oseq_groupcastTE [elim]:
  " $[((\sigma, \{l\}\text{groupcast}(s_{ips}, s_{msg}). p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i;$ 
   $[a = \text{groupcast } (s_{ips} (\sigma i)) (s_{msg} (\sigma i)); \sigma' i = \sigma i; q = p] \implies P] \implies P$ "
  by (ind_cases " $((\sigma, \{l\}\text{groupcast}(s_{ips}, s_{msg}). p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ ") simp

lemma oseq_unicastTE [elim]:
  " $[((\sigma, \{l\}\text{unicast}(s_{ip}, s_{msg}). p \triangleright q), a, (\sigma', r)) \in \text{oseqp\_sos } \Gamma i;$ 
   $[a = \text{unicast } (s_{ip} (\sigma i)) (s_{msg} (\sigma i)); \sigma' i = \sigma i; r = p] \implies P;$ 
   $[a = \neg \text{unicast } (s_{ip} (\sigma i)); \sigma' i = \sigma i; r = q] \implies P] \implies P$ "
  by (ind_cases " $((\sigma, \{l\}\text{unicast}(s_{ip}, s_{msg}). p \triangleright q), a, (\sigma', r)) \in \text{oseqp\_sos } \Gamma i$ ") simp_all

lemma oseq_sendTE [elim]:
  " $[((\sigma, \{l\}\text{send}(s_{msg}). p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i;$ 
   $[a = \text{send } (s_{msg} (\sigma i)); \sigma' i = \sigma i; q = p] \implies P] \implies P$ "
  by (ind_cases " $((\sigma, \{l\}\text{send}(s_{msg}). p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ ") simp

lemma oseq_deliverTE [elim]:

```

```

"[( $(\sigma, \{l\}deliver(s_{data}). p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ;
 $\llbracket a = deliver (s_{data} (\sigma i)); \sigma' i = \sigma i; q = p \rrbracket \implies P \rrbracket \implies P$ "
by (ind_cases " $(\sigma, \{l\}deliver(s_{data}). p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ") simp

```

lemma oseq\_receiveTE [elim]:

```

"[( $(\sigma, \{l\}receive(u_{msg}). p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ;
 $\bigwedge msg. \llbracket a = receive msg; \sigma' i = u_{msg} msg (\sigma i); q = p \rrbracket \implies P \rrbracket \implies P$ "
by (ind_cases " $(\sigma, \{l\}receive(u_{msg}). p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ") simp

```

lemma oseq\_assignTE [elim]:

```

"[( $(\sigma, \{l\}\llbracket u \rrbracket p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ;  $\llbracket a = \tau; \sigma' i = u (\sigma i); q = p \rrbracket \implies P \rrbracket \implies P$ "
by (ind_cases " $(\sigma, \{l\}\llbracket u \rrbracket p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ") simp

```

lemma oseq\_guardTE [elim]:

```

"[( $(\sigma, \{l\}\langle g \rangle p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ;  $\llbracket a = \tau; \sigma' i \in g (\sigma i); q = p \rrbracket \implies P \rrbracket \implies P$ "
by (ind_cases " $(\sigma, \{l\}\langle g \rangle p), a, (\sigma', q)) \in oseqp\_sos \Gamma i$ ") simp

```

lemmas oseqpTEs =

```

oseq_broadcastTE
oseq_groupcastTE
oseq_unicastTE
oseq_sendTE
oseq_deliverTE
oseq_receiveTE
oseq_assignTE
oseq_callTE
oseq_choiceTE
oseq_guardTE

```

declare oseqp\_sos.intros [intro]

## 13.2 Open structural operational semantics for parallel process expressions

inductive\_set

```

oparp_sos :: "ip
 $\Rightarrow ((ip \Rightarrow 's) \times 's1, 'm seq\_action)$  transition set
 $\Rightarrow ('s2, 'm seq\_action)$  transition set
 $\Rightarrow ((ip \Rightarrow 's) \times ('s1 \times 's2), 'm seq\_action)$  transition set"

```

for i :: ip

and S :: " $(ip \Rightarrow 's) \times 's1, 'm seq\_action)$  transition set"

and T :: " $('s2, 'm seq\_action)$  transition set"

where

```

oparleft: " $\llbracket ((\sigma, s), a, (\sigma', s')) \in S; \bigwedge m. a \neq receive\ m \rrbracket \implies$ 
 $((\sigma, (s, t)), a, (\sigma', (s', t))) \in oparp\_sos\ i\ S\ T$ "
/ oparright: " $\llbracket (t, a, t') \in T; \bigwedge m. a \neq send\ m; \sigma' i = \sigma i \rrbracket \implies$ 
 $((\sigma, (s, t)), a, (\sigma', (s, t'))) \in oparp\_sos\ i\ S\ T$ "
/ oparboth: " $\llbracket ((\sigma, s), receive\ m, (\sigma', s')) \in S; (t, send\ m, t') \in T \rrbracket \implies$ 
 $((\sigma, (s, t)), \tau, (\sigma', (s', t'))) \in oparp\_sos\ i\ S\ T$ "

```

lemma opar\_broadcastTE [elim]:

```

"[( $(\sigma, (s, t)), broadcast\ m, (\sigma', (s', t')) \in oparp\_sos\ i\ S\ T$ ;
 $\llbracket ((\sigma, s), broadcast\ m, (\sigma', s')) \in S; t' = t \rrbracket \implies P$ ;
 $\llbracket (t, broadcast\ m, t') \in T; s' = s; \sigma' i = \sigma i \rrbracket \implies P \rrbracket \implies P$ "
by (ind_cases " $(\sigma, (s, t)), broadcast\ m, (\sigma', (s', t')) \in oparp\_sos\ i\ S\ T$ ") simp_all

```

lemma opar\_groupcastTE [elim]:

```

"[( $(\sigma, (s, t)), groupcast\ ips\ m, (\sigma', (s', t')) \in oparp\_sos\ i\ S\ T$ ;
 $\llbracket ((\sigma, s), groupcast\ ips\ m, (\sigma', s')) \in S; t' = t \rrbracket \implies P$ ;
 $\llbracket (t, groupcast\ ips\ m, t') \in T; s' = s; \sigma' i = \sigma i \rrbracket \implies P \rrbracket \implies P$ "
by (ind_cases " $(\sigma, (s, t)), groupcast\ ips\ m, (\sigma', (s', t')) \in oparp\_sos\ i\ S\ T$ ") simp_all

```

lemma opar\_unicastTE [elim]:

```

"[( $(\sigma, (s, t)), unicast\ i\ m, (\sigma', (s', t')) \in oparp\_sos\ i\ S\ T$ ;
 $\llbracket ((\sigma, s), unicast\ i\ m, (\sigma', s')) \in S; t' = t \rrbracket \implies P$ ;

```

$\llbracket (t, \text{unicast } i \ m, t') \in T; s' = s; \sigma' i = \sigma i \rrbracket \implies P \implies P''$   
 by (ind\_cases " $((\sigma, (s, t)), \text{unicast } i \ m, (\sigma', (s', t')))) \in \text{oparp\_sos } i \ S \ T''$ ") simp\_all

lemma opar\_notunicastTE [elim]:

$\llbracket ((\sigma, (s, t)), \text{notunicast } i, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T;$   
 $\llbracket ((\sigma, s), \text{notunicast } i, (\sigma', s')) \in S; t' = t \rrbracket \implies P;$   
 $\llbracket (t, \text{notunicast } i, t') \in T; s' = s; \sigma' i = \sigma i \rrbracket \implies P \rrbracket \implies P''$   
 by (ind\_cases " $((\sigma, (s, t)), \text{notunicast } i, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T''$ ") simp\_all

lemma opar\_sendTE [elim]:

$\llbracket ((\sigma, (s, t)), \text{send } m, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T;$   
 $\llbracket ((\sigma, s), \text{send } m, (\sigma', s')) \in S; t' = t \rrbracket \implies P \rrbracket \implies P''$   
 by (ind\_cases " $((\sigma, (s, t)), \text{send } m, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T''$ ") auto

lemma opar\_deliverTE [elim]:

$\llbracket ((\sigma, (s, t)), \text{deliver } d, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T;$   
 $\llbracket ((\sigma, s), \text{deliver } d, (\sigma', s')) \in S; t' = t \rrbracket \implies P;$   
 $\llbracket (t, \text{deliver } d, t') \in T; s' = s; \sigma' i = \sigma i \rrbracket \implies P \rrbracket \implies P''$   
 by (ind\_cases " $((\sigma, (s, t)), \text{deliver } d, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T''$ ") simp\_all

lemma opar\_receiveTE [elim]:

$\llbracket ((\sigma, (s, t)), \text{receive } m, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T;$   
 $\llbracket (t, \text{receive } m, t') \in T; s' = s; \sigma' i = \sigma i \rrbracket \implies P \rrbracket \implies P''$   
 by (ind\_cases " $((\sigma, (s, t)), \text{receive } m, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T''$ ") auto

inductive\_cases opar\_tauTE: " $((\sigma, (s, t)), \tau, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T''$ "

lemmas oparpTEs =

opar\_broadcastTE  
 opar\_groupcastTE  
 opar\_unicastTE  
 opar\_notunicastTE  
 opar\_sendTE  
 opar\_deliverTE  
 opar\_receiveTE

lemma oparp\_sos\_cases [elim]:

assumes " $((\sigma, (s, t)), a, (\sigma', (s', t'))) \in \text{oparp\_sos } i \ S \ T''$ "  
 and " $\llbracket ((\sigma, s), a, (\sigma', s')) \in S; \bigwedge m. a \neq \text{receive } m; t' = t \rrbracket \implies P''$ "  
 and " $\llbracket (t, a, t') \in T; \bigwedge m. a \neq \text{send } m; s' = s; \sigma' i = \sigma i \rrbracket \implies P''$ "  
 and " $\bigwedge m. \llbracket a = \tau; ((\sigma, s), \text{receive } m, (\sigma', s')) \in S; (t, \text{send } m, t') \in T \rrbracket \implies P''$ "  
 shows "P"  
 using assms by cases auto

definition extg ::  $'a \times 'b \times 'c \Rightarrow 'a \times 'b \times 'c$

where "extg  $\equiv \lambda((\sigma, l1), l2). (\sigma, (l1, l2))$ "

lemma extgsimp [simp]:

"extg  $((\sigma, l1), l2) = (\sigma, (l1, l2))$ "  
 unfolding extg\_def by simp

lemma extg\_range\_prod: "extg  $' (i1 \times i2) = \{(\sigma, (s1, s2)) \mid \sigma \ s1 \ s2. (\sigma, s1) \in i1 \ \wedge \ s2 \in i2\}$ "

unfolding image\_def extg\_def  
 by (rule Collect\_cong) (auto split: prod.split)

definition

opar\_comp ::  $((ip \Rightarrow 's) \times 's1, 'm \ \text{seq\_action}) \ \text{automaton}$   
 $\Rightarrow ip$   
 $\Rightarrow ('s2, 'm \ \text{seq\_action}) \ \text{automaton}$   
 $\Rightarrow ((ip \Rightarrow 's) \times 's1 \times 's2, 'm \ \text{seq\_action}) \ \text{automaton}$   
 (" $\_ \llbracket \_ \rrbracket$ " [102, 0, 103] 102)

where

$'s \llbracket i \ t \equiv (\llbracket \text{init} = \text{extg } ' (\text{init } s \times \text{init } t), \text{trans} = \text{oparp\_sos } i \ (\text{trans } s) \ (\text{trans } t) \rrbracket)$ "

```

lemma opar_comp_def':
  "s <<_i t = (| init = {(σ, (sl, tl) | σ sl tl. (σ, sl) ∈ init s ∧ tl ∈ init t},
              trans = oparp_sos i (trans s) (trans t) |)"
  unfolding opar_comp_def extg_def image_def by (auto split: prod.split)

```

```

lemma trans_opar_comp [simp]:
  "trans (s <<_i t) = oparp_sos i (trans s) (trans t)"
  unfolding opar_comp_def by simp

```

```

lemma init_opar_comp [simp]:
  "init (s <<_i t) = extg ' (init s × init t)"
  unfolding opar_comp_def by simp

```

### 13.3 Open structural operational semantics for node expressions

inductive\_set

```

onode_sos :: "(ip ⇒ 's) × 'l, 'm seq_action) transition set
           ⇒ ((ip ⇒ 's) × 'l net_state, 'm node_action) transition set"
for S :: "(ip ⇒ 's) × 'l, 'm seq_action) transition set"

```

where

```

  onode_bcast:
    "((σ, s), broadcast m, (σ', s')) ∈ S ⇒ ((σ, NodeS i s R), R:*cast(m), (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_gcast:
    "((σ, s), groupcast D m, (σ', s')) ∈ S ⇒ ((σ, NodeS i s R), (R∩D):*cast(m), (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_ucast:
    "[[ (σ, s), unicast d m, (σ', s') ∈ S; d∈R ]] ⇒ ((σ, NodeS i s R), {d}:*cast(m), (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_notucast: "[[ (σ, s), ¬unicast d, (σ', s') ∈ S; d∉R; ∀j. j≠i → σ' j = σ j ]]
                    ⇒ ((σ, NodeS i s R), τ, (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_deliver: "[[ (σ, s), deliver d, (σ', s') ∈ S; ∀j. j≠i → σ' j = σ j ]]
                    ⇒ ((σ, NodeS i s R), i:deliver(d), (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_tau: "[[ (σ, s), τ, (σ', s') ∈ S; ∀j. j≠i → σ' j = σ j ]]
                ⇒ ((σ, NodeS i s R), τ, (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_receive:
    "((σ, s), receive m, (σ', s')) ∈ S ⇒ ((σ, NodeS i s R), {i}¬{i}:arrive(m), (σ', NodeS i s' R)) ∈ onode_sos S"

```

```

  | onode_arrive:
    "σ' i = σ i ⇒ ((σ, NodeS i s R), {i}¬{i}:arrive(m), (σ', NodeS i s R)) ∈ onode_sos S"

```

```

  | onode_connect1:
    "σ' i = σ i ⇒ ((σ, NodeS i s R), connect(i, i'), (σ', NodeS i s (R ∪ {i'}))) ∈ onode_sos S"

```

```

  | onode_connect2:
    "σ' i = σ i ⇒ ((σ, NodeS i s R), connect(i', i), (σ', NodeS i s (R ∪ {i'}))) ∈ onode_sos S"

```

```

  | onode_disconnect1:
    "σ' i = σ i ⇒ ((σ, NodeS i s R), disconnect(i, i'), (σ', NodeS i s (R - {i'}))) ∈ onode_sos S"

```

```

  | onode_disconnect2:
    "σ' i = σ i ⇒ ((σ, NodeS i s R), disconnect(i', i), (σ', NodeS i s (R - {i'}))) ∈ onode_sos S"

```

```

  | onode_connect_other:
    "[[ i ≠ i'; i ≠ i''; σ' i = σ i ]] ⇒ ((σ, NodeS i s R), connect(i', i''), (σ', NodeS i s R)) ∈ onode_sos S"

```

| onode\_disconnect\_other:  
 "[[ i ≠ i'; i ≠ i''; σ' i = σ i ]] ⇒ ((σ, NodeS i s R), disconnect(i', i''), (σ', NodeS i s R)) ∈ onode\_sos S"

inductive\_cases

onode\_arriveTE [elim]: "((σ, NodeS i s R), ii~ni:arrive(m), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 and onode\_castTE [elim]: "((σ, NodeS i s R), RR:\*cast(m), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 and onode\_deliverTE [elim]: "((σ, NodeS i s R), ii:deliver(d), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 and onode\_connectTE [elim]: "((σ, NodeS i s R), connect(ii, ii'), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 and onode\_disconnectTE [elim]: "((σ, NodeS i s R), disconnect(ii, ii'), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 and onode\_newpktTE [elim]: "((σ, NodeS i s R), ii:newpkt(d, di), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 and onode\_tauTE [elim]: "((σ, NodeS i s R), τ, (σ', NodeS i' s' R')) ∈ onode\_sos S"

lemma oarrives\_or\_not:

assumes "((σ, NodeS i s R), ii~ni:arrive(m), (σ', NodeS i' s' R')) ∈ onode\_sos S"  
 shows "(ii = {i} ∧ ni = {}) ∨ (ii = {} ∧ ni = {i})"  
 using assms by rule simp\_all

definition

onode\_comp :: "ip  
 ⇒ ((ip ⇒ 's) × 'l, 'm seq\_action) automaton  
 ⇒ ip set  
 ⇒ ((ip ⇒ 's) × 'l net\_state, 'm node\_action) automaton"  
 ("⟨⟨\_ : (\_ : \_)⟩\_⟩" [0, 0, 0] 104)

where

"⟨i : onp : R\_i⟩\_o ≡ (| init = {(σ, NodeS i s R\_i) | σ s. (σ, s) ∈ init onp},  
 trans = onode\_sos (trans onp) |)"

lemma trans\_onode\_comp:

"trans (⟨i : S : R⟩\_o) = onode\_sos (trans S)"  
 unfolding onode\_comp\_def by simp

lemma init\_onode\_comp:

"init (⟨i : S : R⟩\_o) = {(σ, NodeS i s R) | σ s. (σ, s) ∈ init S}"  
 unfolding onode\_comp\_def by simp

lemmas onode\_comps = trans\_onode\_comp init\_onode\_comp

lemma fst\_par\_onode\_comp [simp]:

"trans (⟨i : s ⟨⟨I t : R⟩\_o⟩ = onode\_sos (oparp\_sos I (trans s) (trans t))"  
 unfolding onode\_comp\_def by simp

lemma init\_par\_onode\_comp [simp]:

"init (⟨i : s ⟨⟨I t : R⟩\_o⟩ = {(σ, NodeS i (s1, s2) R) | σ s1 s2. ((σ, s1), s2) ∈ init s × init t}"  
 unfolding onode\_comp\_def by (simp add: extg\_range\_prod)

lemma onode\_sos\_dest\_is\_net\_state:

assumes "((σ, p), a, s') ∈ onode\_sos S"  
 shows "∃σ' i' ζ' R'. s' = (σ', NodeS i' ζ' R)"  
 using assms proof -  
 assume "((σ, p), a, s') ∈ onode\_sos S"  
 then obtain σ' i' ζ' R' where "s' = (σ', NodeS i' ζ' R)"  
 by (cases s') (auto elim!: onode\_sos.cases)  
 thus ?thesis by simp

qed



```

lemma onode_sos_dest_is_net_state':
  assumes "((σ, NodeS i p R), a, s') ∈ onode_sos S"
  shows "∃σ' ζ' R'. s' = (σ', NodeS i ζ' R')"
  using assms proof -
    assume "((σ, NodeS i p R), a, s') ∈ onode_sos S"
    then obtain σ' ζ' R' where "s' = (σ', NodeS i ζ' R'"
      by (cases s') (auto elim!: onode_sos.cases)
    thus ?thesis by simp
  qed

lemma onode_sos_dest_is_net_state'':
  assumes "((σ, NodeS i p R), a, (σ', s')) ∈ onode_sos S"
  shows "∃ζ' R'. s' = NodeS i ζ' R'"
  proof -
    define ns' where "ns' = (σ', s'"
    with assms have "((σ, NodeS i p R), a, ns') ∈ onode_sos S" by simp
    then obtain σ'' ζ' R' where "ns' = (σ'', NodeS i ζ' R'"
      by (metis onode_sos_dest_is_net_state')
    hence "s' = NodeS i ζ' R'" by (simp add: ns'_def)
    thus ?thesis by simp
  qed

lemma onode_sos_src_is_net_state:
  assumes "((σ, p), a, s') ∈ onode_sos S"
  shows "∃i ζ R. p = NodeS i ζ R"
  using assms proof -
    assume "((σ, p), a, s') ∈ onode_sos S"
    then obtain i ζ R where "p = NodeS i ζ R"
      by (cases s') (auto elim!: onode_sos.cases)
    thus ?thesis by simp
  qed

lemma onode_sos_net_states:
  assumes "((σ, s), a, (σ', s')) ∈ onode_sos S"
  shows "∃i ζ R ζ' R'. s = NodeS i ζ R ∧ s' = NodeS i ζ' R'"
  proof -
    from assms obtain i ζ R where "s = NodeS i ζ R"
      by (metis onode_sos_src_is_net_state)
    moreover with assms obtain ζ' R' where "s' = NodeS i ζ' R'"
      by (auto dest!: onode_sos_dest_is_net_state')
    ultimately show ?thesis by simp
  qed

lemma node_sos_cases [elim]:
  "((σ, NodeS i p R), a, (σ', NodeS i p' R')) ∈ onode_sos S ⇒
  (∧m .      [ a = R:*cast(m);      R' = R; ((σ, p), broadcast m, (σ', p')) ∈ S ] ⇒ P) ⇒
  (∧m D.     [ a = (R ∩ D):*cast(m);  R' = R; ((σ, p), groupcast D m, (σ', p')) ∈ S ] ⇒ P) ⇒
  (∧d m.     [ a = {d}:*cast(m);      R' = R; ((σ, p), unicast d m, (σ', p')) ∈ S; d ∈ R ] ⇒ P)
  ⇒
  (∧d.       [ a = τ;                  R' = R; ((σ, p), ¬unicast d, (σ', p')) ∈ S; d ∉ R ] ⇒ P)
  ⇒
  (∧d.       [ a = i:deliver(d);       R' = R; ((σ, p), deliver d, (σ', p')) ∈ S ] ⇒ P) ⇒
  (∧m.       [ a = {i}¬{j}:arrive(m);  R' = R; ((σ, p), receive m, (σ', p')) ∈ S ] ⇒ P) ⇒
  (         [ a = τ;                  R' = R; ((σ, p), τ, (σ', p')) ∈ S ] ⇒ P) ⇒
  (∧m.       [ a = {j}¬{i}:arrive(m);  R' = R; p = p'; σ' i = σ i ] ⇒ P) ⇒
  (∧i i'.    [ a = connect(i, i');      R' = R ∪ {i'}; p = p'; σ' i = σ i ] ⇒ P) ⇒
  (∧i i'.    [ a = connect(i', i);     R' = R ∪ {i'}; p = p'; σ' i = σ i ] ⇒ P) ⇒
  (∧i i'.    [ a = disconnect(i, i');   R' = R - {i'}; p = p'; σ' i = σ i ] ⇒ P) ⇒
  (∧i i'.    [ a = disconnect(i', i);  R' = R - {i'}; p = p'; σ' i = σ i ] ⇒ P) ⇒
  (∧i i' i''. [ a = connect(i', i'');   R' = R; p = p'; i ≠ i'; i ≠ i''; σ' i = σ i ] ⇒ P) ⇒
  (∧i i' i''. [ a = disconnect(i', i''); R' = R; p = p'; i ≠ i'; i ≠ i''; σ' i = σ i ] ⇒ P)
  P"
  by (erule onode_sos.cases) (simp | metis)+

```

## 13.4 Open structural operational semantics for partial network expressions

inductive\_set

```
opnet_sos :: "(ip ⇒ 's) × 'l net_state, 'm node_action) transition set
           ⇒ ((ip ⇒ 's) × 'l net_state, 'm node_action) transition set
           ⇒ ((ip ⇒ 's) × 'l net_state, 'm node_action) transition set"
for S :: "(ip ⇒ 's) × 'l net_state, 'm node_action) transition set"
and T :: "(ip ⇒ 's) × 'l net_state, 'm node_action) transition set"
```

where

opnet\_cast1:

```
"[( (σ, s), R:*cast(m), (σ', s')) ∈ S; ((σ, t), H¬K:arrive(m), (σ', t')) ∈ T; H ⊆ R; K ∩ R = {} ]
 ⇒ ((σ, SubnetS s t), R:*cast(m), (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

/ opnet\_cast2:

```
"[( (σ, s), H¬K:arrive(m), (σ', s')) ∈ S; ((σ, t), R:*cast(m), (σ', t')) ∈ T; H ⊆ R; K ∩ R = {} ]
 ⇒ ((σ, SubnetS s t), R:*cast(m), (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

/ opnet\_arrive:

```
"[( (σ, s), H¬K:arrive(m), (σ', s')) ∈ S; ((σ, t), H'¬K':arrive(m), (σ', t')) ∈ T ]
 ⇒ ((σ, SubnetS s t), (H ∪ H')¬(K ∪ K'):arrive(m), (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

/ opnet\_deliver1:

```
"((σ, s), i:deliver(d), (σ', s')) ∈ S
 ⇒ ((σ, SubnetS s t), i:deliver(d), (σ', SubnetS s' t)) ∈ opnet_sos S T"
```

/ opnet\_deliver2:

```
"[( (σ, t), i:deliver(d), (σ', t')) ∈ T ]
 ⇒ ((σ, SubnetS s t), i:deliver(d), (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

/ opnet\_tau1:

```
"((σ, s), τ, (σ', s')) ∈ S ⇒ ((σ, SubnetS s t), τ, (σ', SubnetS s' t)) ∈ opnet_sos S T"
```

/ opnet\_tau2:

```
"((σ, t), τ, (σ', t')) ∈ T ⇒ ((σ, SubnetS s t), τ, (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

/ opnet\_connect:

```
"[( (σ, s), connect(i, i'), (σ', s')) ∈ S; ((σ, t), connect(i, i'), (σ', t')) ∈ T ]
 ⇒ ((σ, SubnetS s t), connect(i, i'), (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

/ opnet\_disconnect:

```
"[( (σ, s), disconnect(i, i'), (σ', s')) ∈ S; ((σ, t), disconnect(i, i'), (σ', t')) ∈ T ]
 ⇒ ((σ, SubnetS s t), disconnect(i, i'), (σ', SubnetS s' t')) ∈ opnet_sos S T"
```

```
inductive_cases opartial_castTE [elim]:      "((σ, s), R:*cast(m), (σ', s')) ∈ opnet_sos S T"
and opartial_arriveTE [elim]:              "((σ, s), H¬K:arrive(m), (σ', s')) ∈ opnet_sos S T"
and opartial_deliverTE [elim]:            "((σ, s), i:deliver(d), (σ', s')) ∈ opnet_sos S T"
and opartial_tauTE [elim]:                "((σ, s), τ, (σ', s')) ∈ opnet_sos S T"
and opartial_connectTE [elim]:            "((σ, s), connect(i, i'), (σ', s')) ∈ opnet_sos S T"
and opartial_disconnectTE [elim]:         "((σ, s), disconnect(i, i'), (σ', s')) ∈ opnet_sos S T"
and opartial_newpktTE [elim]:             "((σ, s), i:newpkt(d, di), (σ', s')) ∈ opnet_sos S T"
```

```
fun opnet :: "(ip ⇒ ((ip ⇒ 's) × 'l, 'm seq_action) automaton)
             ⇒ net_tree ⇒ ((ip ⇒ 's) × 'l net_state, 'm node_action) automaton"
```

where

```
"opnet onp ⟨i; Ri⟩ = ⟨i : onp i : Ri⟩o"
```

```
/ "opnet onp (p1 || p2) = (| init = {(σ, SubnetS s1 s2) | σ s1 s2.
```

```
(σ, s1) ∈ init (opnet onp p1)
```

```
∧ (σ, s2) ∈ init (opnet onp p2)
```

```
∧ net_ips s1 ∩ net_ips s2 = {},
```

```
trans = opnet_sos (trans (opnet onp p1)) (trans (opnet onp p2)) |)"
```

lemma opnet\_node\_init [elim, simp]:

```
assumes "(σ, s) ∈ init (opnet onp ⟨i; R⟩)"
```

```

shows " $(\sigma, s) \in \{ (\sigma, \text{NodeS } i \text{ ns } R) \mid \sigma \text{ ns. } (\sigma, \text{ns}) \in \text{init } (\text{onp } i)\}$ "
using assms by (simp add: onode_comp_def)

```

```

lemma opnet_node_init' [elim]:
  assumes " $(\sigma, s) \in \text{init } (\text{opnet } \text{onp } \langle i; R \rangle)$ "
  obtains ns where " $s = \text{NodeS } i \text{ ns } R$ "
    and " $(\sigma, \text{ns}) \in \text{init } (\text{onp } i)$ "
  using assms by (auto simp add: onode_comp_def)

```

```

lemma opnet_node_trans [elim, simp]:
  assumes " $(s, a, s') \in \text{trans } (\text{opnet } \text{onp } \langle i; R \rangle)$ "
  shows " $(s, a, s') \in \text{onode\_sos } (\text{trans } (\text{onp } i))$ "
  using assms by (simp add: trans_onode_comp)

```

### 13.5 Open structural operational semantics for complete network expressions

inductive\_set

```

ocnet_sos :: " $((ip \Rightarrow 's) \times 'l \text{ net\_state}, 'm::\text{msg node\_action}) \text{ transition set}$ 
            $\Rightarrow ((ip \Rightarrow 's) \times 'l \text{ net\_state}, 'm \text{ node\_action}) \text{ transition set}$ "

```

```

for S :: " $((ip \Rightarrow 's) \times 'l \text{ net\_state}, 'm \text{ node\_action}) \text{ transition set}$ "

```

where

```

ocnet_connect:

```

```

" $\llbracket ((\sigma, s), \text{connect}(i, i'), (\sigma', s')) \in S; \forall j. j \notin \text{net\_ips } s \longrightarrow (\sigma' j = \sigma j) \rrbracket$ 
 $\implies ((\sigma, s), \text{connect}(i, i'), (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

/ ocnet_disconnect:

```

```

" $\llbracket ((\sigma, s), \text{disconnect}(i, i'), (\sigma', s')) \in S; \forall j. j \notin \text{net\_ips } s \longrightarrow (\sigma' j = \sigma j) \rrbracket$ 
 $\implies ((\sigma, s), \text{disconnect}(i, i'), (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

/ ocnet_cast:

```

```

" $\llbracket ((\sigma, s), R:*\text{cast}(m), (\sigma', s')) \in S; \forall j. j \notin \text{net\_ips } s \longrightarrow (\sigma' j = \sigma j) \rrbracket$ 
 $\implies ((\sigma, s), \tau, (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

/ ocnet_tau:

```

```

" $\llbracket ((\sigma, s), \tau, (\sigma', s')) \in S; \forall j. j \notin \text{net\_ips } s \longrightarrow (\sigma' j = \sigma j) \rrbracket$ 
 $\implies ((\sigma, s), \tau, (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

/ ocnet_deliver:

```

```

" $\llbracket ((\sigma, s), i:\text{deliver}(d), (\sigma', s')) \in S; \forall j. j \notin \text{net\_ips } s \longrightarrow (\sigma' j = \sigma j) \rrbracket$ 
 $\implies ((\sigma, s), i:\text{deliver}(d), (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

/ ocnet_newpkt:

```

```

" $\llbracket ((\sigma, s), \{i\}\neg K:\text{arrive}(\text{newpkt}(d, di)), (\sigma', s')) \in S; \forall j. j \notin \text{net\_ips } s \longrightarrow (\sigma' j = \sigma j) \rrbracket$ 
 $\implies ((\sigma, s), i:\text{newpkt}(d, di), (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

inductive_cases oconnect_completeTE: " $((\sigma, s), \text{connect}(i, i'), (\sigma', s')) \in \text{ocnet\_sos } S$ "
  and odisconnect_completeTE: " $((\sigma, s), \text{disconnect}(i, i'), (\sigma', s')) \in \text{ocnet\_sos } S$ "
  and otau_completeTE: " $((\sigma, s), \tau, (\sigma', s')) \in \text{ocnet\_sos } S$ "
  and odeliver_completeTE: " $((\sigma, s), i:\text{deliver}(d), (\sigma', s')) \in \text{ocnet\_sos } S$ "
  and onewpkt_completeTE: " $((\sigma, s), i:\text{newpkt}(d, di), (\sigma', s')) \in \text{ocnet\_sos } S$ "

```

```

lemmas ocompleteTEs = oconnect_completeTE
  odisconnect_completeTE
  otau_completeTE
  odeliver_completeTE
  onewpkt_completeTE

```

lemma ocomplete\_no\_cast [simp]:

```

" $((\sigma, s), R:*\text{cast}(m), (\sigma', s')) \notin \text{ocnet\_sos } T$ "

```

proof

```

  assume " $((\sigma, s), R:*\text{cast}(m), (\sigma', s')) \in \text{ocnet\_sos } T$ "

```

```

  hence " $R:*\text{cast}(m) \neq R:*\text{cast}(m)$ "

```

```

  by (rule ocnet_sos.cases) auto

```

```

  thus False by simp

```

qed

```
lemma ocomplete_no_arrive [simp]:  
  "(( $\sigma$ , s), ii¬ni:arrive(m), ( $\sigma'$ , s'))  $\notin$  ocnet_sos T"  
proof  
  assume "(( $\sigma$ , s), ii¬ni:arrive(m), ( $\sigma'$ , s'))  $\in$  ocnet_sos T"  
  hence "ii¬ni:arrive(m)  $\neq$  ii¬ni:arrive(m)"  
  by (rule ocnet_sos.cases) auto  
  thus False by simp  
qed
```

```
lemma ocomplete_no_change [elim]:  
  assumes "(( $\sigma$ , s), a, ( $\sigma'$ , s'))  $\in$  ocnet_sos T"  
  and "j  $\notin$  net_ips s"  
  shows " $\sigma'$  j =  $\sigma$  j"  
using assms by cases simp_all
```

```
lemma ocomplete_transE [elim]:  
  assumes "(( $\sigma$ ,  $\zeta$ ), a, ( $\sigma'$ ,  $\zeta'$ ))  $\in$  ocnet_sos (trans (opnet onp n))"  
  obtains a' where "(( $\sigma$ ,  $\zeta$ ), a', ( $\sigma'$ ,  $\zeta'$ ))  $\in$  trans (opnet onp n)"  
  using assms by (cases a) (auto elim!: ocompleteTEs [simplified])
```

abbreviation

```
oclosed :: "((ip  $\Rightarrow$  's)  $\times$  'l net_state, ('m::msg) node_action) automaton  
           $\Rightarrow$  ((ip  $\Rightarrow$  's)  $\times$  'l net_state, 'm node_action) automaton"
```

where

```
"oclosed  $\equiv$  ( $\lambda$ A. A ( $\mid$  trans := ocnet_sos (trans A)  $\mid$ ))"
```

end

## 14 Configure the inv-cterms tactic for open sequential processes

```
theory OAWN_SOS_Labels  
imports OAWN_SOS Inv_Cterms  
begin
```

```
lemma oelinder_guard:  
  assumes "p = {l}<fg> qq"  
  and "l'  $\in$  labels  $\Gamma$  q"  
  and "(( $\sigma$ , p), a, ( $\sigma'$ , q))  $\in$  oseqp_sos  $\Gamma$  i"  
  obtains p' where "p = {l}<fg> p'"  
  and "l'  $\in$  labels  $\Gamma$  qq"  
using assms by auto
```

```
lemma oelinder_assign:  
  assumes "p = {l}[[fa]] qq"  
  and "l'  $\in$  labels  $\Gamma$  q"  
  and "(( $\sigma$ , p), a, ( $\sigma'$ , q))  $\in$  oseqp_sos  $\Gamma$  i"  
  obtains p' where "p = {l}[[fa]] p'"  
  and "l'  $\in$  labels  $\Gamma$  qq"  
using assms by auto
```

```
lemma oelinder_ucast:  
  assumes "p = {l}unicast(fip, fmsg).q1  $\triangleright$  q2"  
  and "l'  $\in$  labels  $\Gamma$  q"  
  and "(( $\sigma$ , p), a, ( $\sigma'$ , q))  $\in$  oseqp_sos  $\Gamma$  i"  
  obtains p' pp' where "p = {l}unicast(fip, fmsg).p'  $\triangleright$  pp'"  
  and "case a of unicast _ _  $\Rightarrow$  l'  $\in$  labels  $\Gamma$  q1  
  | _  $\Rightarrow$  l'  $\in$  labels  $\Gamma$  q2"  
using assms by simp (erule oseqpTEs, auto)
```

```
lemma oelinder_bcast:  
  assumes "p = {l}broadcast(fmsg).qq"  
  and "l'  $\in$  labels  $\Gamma$  q"
```

```

    and " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
obtains  $p'$  where " $p = \{l\}\text{broadcast}(fmsg). p'$ "
    and " $l' \in \text{labels } \Gamma qq$ "
using assms by auto

lemma oelimder_gcast:
  assumes " $p = \{l\}\text{groupcast}(fips, fmsg).qq$ "
    and " $l' \in \text{labels } \Gamma q$ "
    and " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
  obtains  $p'$  where " $p = \{l\}\text{groupcast}(fips, fmsg). p'$ "
    and " $l' \in \text{labels } \Gamma qq$ "
  using assms by auto

lemma oelimder_send:
  assumes " $p = \{l\}\text{send}(fmsg).qq$ "
    and " $l' \in \text{labels } \Gamma q$ "
    and " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
  obtains  $p'$  where " $p = \{l\}\text{send}(fmsg). p'$ "
    and " $l' \in \text{labels } \Gamma qq$ "
  using assms by auto

lemma oelimder_deliver:
  assumes " $p = \{l\}\text{deliver}(fdata).qq$ "
    and " $l' \in \text{labels } \Gamma q$ "
    and " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
  obtains  $p'$  where " $p = \{l\}\text{deliver}(fdata). p'$ "
    and " $l' \in \text{labels } \Gamma qq$ "
  using assms by auto

lemma oelimder_receive:
  assumes " $p = \{l\}\text{receive}(fmsg).qq$ "
    and " $l' \in \text{labels } \Gamma q$ "
    and " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
  obtains  $p'$  where " $p = \{l\}\text{receive}(fmsg). p'$ "
    and " $l' \in \text{labels } \Gamma qq$ "
  using assms by auto

lemmas oelimders =
  oelimder_guard
  oelimder_assign
  oelimder_ucast
  oelimder_bcast
  oelimder_gcast
  oelimder_send
  oelimder_deliver
  oelimder_receive

declare
  oseqpTEs [cterms_seqte]
  oelimders [cterms_elimders]

end

```

## 15 Lemmas for open partial networks

```

theory OPnet
imports OAWN_SOS OInvariants
begin

```

These lemmas mostly concern the preservation of node structure by *opnet\_sos* transitions.

```

lemma opnet_maintains_dom:
  assumes " $((\sigma, ns), a, (\sigma', ns')) \in \text{trans } (\text{opnet } np p)$ "
  shows " $\text{net\_ips } ns = \text{net\_ips } ns'$ "
  using assms proof (induction  $p$  arbitrary:  $\sigma ns a \sigma' ns'$ )

```

```

fix i R σ ns a σ' ns'
assume "((σ, ns), a, (σ', ns')) ∈ trans (opnet np ⟨i; R⟩)"
hence "((σ, ns), a, (σ', ns')) ∈ onode_sos (trans (np i))" ..
thus "net_ips ns = net_ips ns'"
  by (simp add: net_ips_is_dom_netmap)
    (erule onode_sos.cases, simp_all)
next
fix p1 p2 σ ns a σ' ns'
assume "∧σ ns a σ' ns'. ((σ, ns), a, (σ', ns')) ∈ trans (opnet np p1) ⇒ net_ips ns = net_ips ns'"
  and "∧σ ns a σ' ns'. ((σ, ns), a, (σ', ns')) ∈ trans (opnet np p2) ⇒ net_ips ns = net_ips ns'"
  and "((σ, ns), a, (σ', ns')) ∈ trans (opnet np (p1 || p2))"
thus "net_ips ns = net_ips ns'"
  by simp (erule opnet_sos.cases, simp_all)
qed

lemma opnet_net_ips_net_tree_ips:
  assumes "(σ, ns) ∈ oreachable (opnet np p) S U"
  shows "net_ips ns = net_tree_ips p"
  using assms proof (induction rule: oreachable_pair_induct)
    fix σ s
    assume "(σ, s) ∈ init (opnet np p)"
    thus "net_ips s = net_tree_ips p"
  proof (induction p arbitrary: σ s)
    fix p1 p2 σ s
    assume IH1: "(∧σ s. (σ, s) ∈ init (opnet np p1) ⇒ net_ips s = net_tree_ips p1)"
      and IH2: "(∧σ s. (σ, s) ∈ init (opnet np p2) ⇒ net_ips s = net_tree_ips p2)"
      and "(σ, s) ∈ init (opnet np (p1 || p2))"
    thus "net_ips s = net_tree_ips (p1 || p2)"
      by (clarsimp simp add: net_ips_is_dom_netmap)
        (metis Un_commute)
  qed (clarsimp simp add: onode_comps)
next
fix σ s σ' s' a
assume "(σ, s) ∈ oreachable (opnet np p) S U"
  and "net_ips s = net_tree_ips p"
  and "((σ, s), a, (σ', s')) ∈ trans (opnet np p)"
  and "S σ σ' a"
thus "net_ips s' = net_tree_ips p"
  by (simp add: net_ips_is_dom_netmap)
    (metis net_ips_is_dom_netmap opnet_maintains_dom)
qed simp

lemma opnet_net_ips_net_tree_ips_init:
  assumes "(σ, ns) ∈ init (opnet np p)"
  shows "net_ips ns = net_tree_ips p"
  using assms(1) by (rule oreachable_init [THEN opnet_net_ips_net_tree_ips])

lemma opartial_net_preserves_subnets:
  assumes "((σ, SubnetS s t), a, (σ', st')) ∈ opnet_sos (trans (opnet np p1)) (trans (opnet np p2))"
  shows "∃s' t'. st' = SubnetS s' t'"
  using assms by cases simp_all

lemma net_par_oreachable_is_subnet:
  assumes "(σ, st) ∈ oreachable (opnet np (p1 || p2)) S U"
  shows "∃s t. st = SubnetS s t"
  proof -
    define p where "p = (σ, st)"
    with assms have "p ∈ oreachable (opnet np (p1 || p2)) S U" by simp
    hence "∃σ s t. p = (σ, SubnetS s t)"
      by induct (auto dest!: opartial_net_preserves_subnets)
    with p_def show ?thesis by simp
  qed
end

```

## 16 Lifting rules for (open) nodes

theory *ONode\_Lifting*

imports *AWN OAWN\_SOS OInvariants*

begin

lemma *node\_net\_state'*:

assumes " $s \in \text{oreachable } (\langle i : T : R_i \rangle_o) S U$ "

shows " $\exists \sigma \zeta R. s = (\sigma, \text{NodeS } i \zeta R)$ "

using *assms proof induction*

fix  $s$

assume " $s \in \text{init } (\langle i : T : R_i \rangle_o)$ "

then obtain  $\sigma \zeta$  where " $s = (\sigma, \text{NodeS } i \zeta R_i)$ "

by (*auto simp: onode\_comps*)

thus " $\exists \sigma \zeta R. s = (\sigma, \text{NodeS } i \zeta R)$ " by *auto*

next

fix  $s$  a  $\sigma'$

assume *rt*: " $s \in \text{oreachable } (\langle i : T : R_i \rangle_o) S U$ "

and *ih*: " $\exists \sigma \zeta R. s = (\sigma, \text{NodeS } i \zeta R)$ "

and " $U (\text{fst } s) \sigma'$ "

then obtain  $\sigma \zeta R$

where " $(\sigma, \text{NodeS } i \zeta R) \in \text{oreachable } (\langle i : T : R_i \rangle_o) S U$ "

and " $U \sigma \sigma'$ " and " $\text{snd } s = \text{NodeS } i \zeta R$ " by *auto*

from *this(1-2)*

have " $(\sigma', \text{NodeS } i \zeta R) \in \text{oreachable } (\langle i : T : R_i \rangle_o) S U$ "

by - (*erule(1) oreachable\_other'*)

with  $\langle \text{snd } s = \text{NodeS } i \zeta R \rangle$  show " $\exists \sigma \zeta R. (\sigma', \text{snd } s) = (\sigma, \text{NodeS } i \zeta R)$ " by *simp*

next

fix  $s$  a  $s'$

assume *rt*: " $s \in \text{oreachable } (\langle i : T : R_i \rangle_o) S U$ "

and *ih*: " $\exists \sigma \zeta R. s = (\sigma, \text{NodeS } i \zeta R)$ "

and *tr*: " $(s, a, s') \in \text{trans } (\langle i : T : R_i \rangle_o)$ "

and " $S (\text{fst } s) (\text{fst } s') a$ "

from *ih* obtain  $\sigma \zeta R$  where " $s = (\sigma, \text{NodeS } i \zeta R)$ " by *auto*

with *tr* have " $((\sigma, \text{NodeS } i \zeta R), a, s') \in \text{onode\_sos } (\text{trans } T)$ "

by (*simp add: onode\_comps*)

then obtain  $\sigma' \zeta' R'$  where " $s' = (\sigma', \text{NodeS } i \zeta' R')$ "

using *onode\_sos\_dest\_is\_net\_state'* by *metis*

with *tr*  $\langle s = (\sigma, \text{NodeS } i \zeta R) \rangle$  show " $\exists \sigma \zeta R. s' = (\sigma, \text{NodeS } i \zeta R)$ "

by *simp*

qed

lemma *node\_net\_state*:

assumes " $(\sigma, s) \in \text{oreachable } (\langle i : T : R_i \rangle_o) S U$ "

shows " $\exists \zeta R. s = \text{NodeS } i \zeta R$ "

using *assms*

by (*metis Pair\_inject node\_net\_state'*)

lemma *node\_net\_state\_trans [elim]*:

assumes *sor*: " $(\sigma, s) \in \text{oreachable } (\langle i : \zeta_i : R_i \rangle_o) S U$ "

and *str*: " $((\sigma, s), a, (\sigma', s')) \in \text{trans } (\langle i : \zeta_i : R_i \rangle_o)$ "

obtains  $\zeta R \zeta' R'$

where " $s = \text{NodeS } i \zeta R$ "

and " $s' = \text{NodeS } i \zeta' R'$ "

proof -

assume \*: " $\bigwedge \zeta R \zeta' R'. s = \text{NodeS } i \zeta R \implies s' = \text{NodeS } i \zeta' R' \implies \text{thesis}$ "

from *sor* obtain  $\zeta R$  where " $s = \text{NodeS } i \zeta R$ "

by (*metis node\_net\_state*)

moreover with *str* obtain  $\zeta' R'$  where " $s' = \text{NodeS } i \zeta' R'$ "

by (*simp only: onode\_comps*)

(*metis onode\_sos\_dest\_is\_net\_state''*)

ultimately show *thesis* by (*rule \**)

qed

```

lemma nodemap_induct' [consumes, case_names init other local]:
  assumes "( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    and init: " $\bigwedge \sigma \zeta$ . ( $\sigma$ , NodeS ii  $\zeta$  R $_i$ )  $\in$  init ( $\langle$ ii : T : R $\rangle_o$ )  $\implies$  P ( $\sigma$ , NodeS ii  $\zeta$  R)"
    and other: " $\bigwedge \sigma \zeta R \sigma' a$ .
      [ ( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U;
        U  $\sigma \sigma'$ ; P ( $\sigma$ , NodeS ii  $\zeta$  R) ]  $\implies$  P ( $\sigma'$ , NodeS ii  $\zeta$  R)"
    and local: " $\bigwedge \sigma \zeta R \sigma' \zeta' R' a$ .
      [ ( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U;
        (( $\sigma$ , NodeS ii  $\zeta$  R), a, ( $\sigma'$ , NodeS ii  $\zeta'$  R'))  $\in$  trans ( $\langle$ ii : T : R $\rangle_o$ );
        S  $\sigma \sigma' a$ ; P ( $\sigma$ , NodeS ii  $\zeta$  R) ]  $\implies$  P ( $\sigma'$ , NodeS ii  $\zeta'$  R)"
  shows "P ( $\sigma$ , NodeS ii  $\zeta$  R)"
using assms(1) proof induction
  fix s
  assume "s  $\in$  init ( $\langle$ ii : T : R $\rangle_o$ )"
  hence "s  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    by (rule oreachable_init)
  with (s  $\in$  init ( $\langle$ ii : T : R $\rangle_o$ )) obtain  $\sigma \zeta$  where "s = ( $\sigma$ , NodeS ii  $\zeta$  R)"
    by (simp add: onode_comps) metis
  with (s  $\in$  init ( $\langle$ ii : T : R $\rangle_o$ )) and init show "P s" by simp
next
  fix s a  $\sigma'$ 
  assume sr: "s  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    and "U (fst s)  $\sigma'$ "
    and "P s"
  from sr obtain  $\sigma \zeta R$  where "s = ( $\sigma$ , NodeS ii  $\zeta$  R)"
    using node_net_state' by metis
  with sr (U (fst s)  $\sigma'$ ) (P s) show "P ( $\sigma'$ , snd s)"
    by simp (metis other)
next
  fix s a s'
  assume sr: "s  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    and tr: "(s, a, s')  $\in$  trans ( $\langle$ ii : T : R $\rangle_o$ )"
    and "S (fst s) (fst s') a"
    and "P s"
  from this(1-3) have "s'  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    by - (erule(2) oreachable_local)
  then obtain  $\sigma' \zeta' R'$  where [simp]: "s' = ( $\sigma'$ , NodeS ii  $\zeta'$  R)"
    using node_net_state' by metis
  from sr and (P s) obtain  $\sigma \zeta R$ 
    where [simp]: "s = ( $\sigma$ , NodeS ii  $\zeta$  R)"
    and A1: "( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    and A4: "P ( $\sigma$ , NodeS ii  $\zeta$  R)"
    using node_net_state' by metis
  with tr and (S (fst s) (fst s') a)
  have A2: "(( $\sigma$ , NodeS ii  $\zeta$  R), a, ( $\sigma'$ , NodeS ii  $\zeta'$  R'))  $\in$  trans ( $\langle$ ii : T : R $\rangle_o$ )"
    and A3: "S  $\sigma \sigma' a$ " by simp_all
  from A1 A2 A3 A4 have "P ( $\sigma'$ , NodeS ii  $\zeta'$  R)" by (rule local)
  thus "P s'" by simp
qed

```

```

lemma nodemap_induct [consumes, case_names init step]:
  assumes "( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U"
    and init: " $\bigwedge \sigma \zeta$ . ( $\sigma$ , NodeS ii  $\zeta$  R $_i$ )  $\in$  init ( $\langle$ ii : T : R $\rangle_o$ )  $\implies$  P  $\sigma \zeta R_i$ "
    and other: " $\bigwedge \sigma \zeta R \sigma' a$ .
      [ ( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U;
        U  $\sigma \sigma'$ ; P  $\sigma \zeta R$  ]  $\implies$  P  $\sigma' \zeta R$ "
    and local: " $\bigwedge \sigma \zeta R \sigma' \zeta' R' a$ .
      [ ( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  oreachable ( $\langle$ ii : T : R $\rangle_o$ ) S U;
        (( $\sigma$ , NodeS ii  $\zeta$  R), a, ( $\sigma'$ , NodeS ii  $\zeta'$  R'))  $\in$  trans ( $\langle$ ii : T : R $\rangle_o$ );
        S  $\sigma \sigma' a$ ; P  $\sigma \zeta R$  ]  $\implies$  P  $\sigma' \zeta' R'$ "
  shows "P  $\sigma \zeta R$ "
using assms(1) proof (induction "( $\sigma$ , NodeS ii  $\zeta$  R)" arbitrary:  $\sigma \zeta R$ )
  fix  $\sigma \zeta R$ 
  assume a1: "( $\sigma$ , NodeS ii  $\zeta$  R)  $\in$  init ( $\langle$ ii : T : R $\rangle_o$ )"

```



```

hence "R = Ri" by (simp add: init_onode_comp)
with a1 have "(σ, NodeS ii ζ Ri) ∈ init ((ii : T : Ri)o)" by simp
with init and (R = Ri) show "P σ ζ R" by simp
next
fix st a σ' ζ' R'
assume "st ∈ oreachable ((ii : T : Ri)o) S U"
  and tr: "(st, a, (σ', NodeS ii ζ' R')) ∈ trans ((ii : T : Ri)o)"
  and "S (fst st) (fst (σ', NodeS ii ζ' R')) a"
  and IH: "∧σ ζ R. st = (σ, NodeS ii ζ R) ⇒ P σ ζ R"
from this(1) obtain σ ζ R where "st = (σ, NodeS ii ζ R)"
  and "(σ, NodeS ii ζ R) ∈ oreachable ((ii : T : Ri)o) S U"
  by (metis node_net_state')
note this(2)
moreover from tr and (st = (σ, NodeS ii ζ R))
  have "((σ, NodeS ii ζ R), a, (σ', NodeS ii ζ' R')) ∈ trans ((ii : T : Ri)o)" by simp
moreover from (S (fst st) (fst (σ', NodeS ii ζ' R')) a) and (st = (σ, NodeS ii ζ R))
  have "S σ σ' a" by simp
moreover from IH and (st = (σ, NodeS ii ζ R)) have "P σ ζ R" .
ultimately show "P σ' ζ' R'" by (rule local)
next
fix st σ' ζ R
assume "st ∈ oreachable ((ii : T : Ri)o) S U"
  and "U (fst st) σ'"
  and "snd st = NodeS ii ζ R"
  and IH: "∧σ ζ R. st = (σ, NodeS ii ζ R) ⇒ P σ ζ R"
from this(1,3) obtain σ where "st = (σ, NodeS ii ζ R)"
  and "(σ, NodeS ii ζ R) ∈ oreachable ((ii : T : Ri)o) S U"
  by (metis surjective_pairing)
note this(2)
moreover from (U (fst st) σ') and (st = (σ, NodeS ii ζ R)) have "U σ σ'" by simp
moreover from IH and (st = (σ, NodeS ii ζ R)) have "P σ ζ R" .
ultimately show "P σ' ζ R" by (rule other)
qed

lemma node_addressD [dest, simp]:
  assumes "(σ, NodeS i ζ R) ∈ oreachable ((ii : T : Ri)o) S U"
  shows "i = ii"
  using assms by (clarsimp dest!: node_net_state')

lemma node_proc_reachable [dest]:
  assumes "(σ, NodeS i ζ R) ∈ oreachable ((ii : T : Ri)o)
    (otherwith S {ii} (oarrivmsg I)) (other U {ii})"
  and sgivesu: "∧ξ ξ'. S ξ ξ' ⇒ U ξ ξ'"
  shows "(σ, ζ) ∈ oreachable T (otherwith S {ii} (orecvmsg I)) (other U {ii})"
proof -
  from assms(1) have "(σ, NodeS ii ζ R) ∈ oreachable ((ii : T : Ri)o)
    (otherwith S {ii} (oarrivmsg I)) (other U {ii})"
  by - (frule node_addressD, simp)
  thus ?thesis
proof (induction rule: nodemap_induct)
  fix σ ζ
  assume "(σ, NodeS ii ζ Ri) ∈ init ((ii : T : Ri)o)"
  hence "(σ, ζ) ∈ init T" by (auto simp: onode_comps)
  thus "(σ, ζ) ∈ oreachable T (otherwith S {ii} (orecvmsg I)) (other U {ii})"
  by (rule oreachable_init)
next
  fix σ ζ R σ' ζ' R' a
  assume "other U {ii} σ σ'"
  and "(σ, ζ) ∈ oreachable T (otherwith S {ii} (orecvmsg I)) (other U {ii})"
  thus "(σ', ζ) ∈ oreachable T (otherwith S {ii} (orecvmsg I)) (other U {ii})"
  by - (rule oreachable_other')
next
  fix σ ζ R σ' ζ' R' a
  assume rs: "(σ, NodeS ii ζ R) ∈ oreachable ((ii : T : Ri)o)"

```

```

                                (otherwith S {ii} (oarrivemsg I)) (other U {ii})"
and tr: "(( $\sigma$ , NodeS ii  $\zeta$  R), a, ( $\sigma'$ , NodeS ii  $\zeta'$  R'))  $\in$  trans ( $\langle ii : T : R_i \rangle_o$ )"
and ow: "otherwith S {ii} (oarrivemsg I)  $\sigma$   $\sigma'$  a"
and ih: "( $\sigma$ ,  $\zeta$ )  $\in$  oreachable T (otherwith S {ii} (orecvmsg I)) (other U {ii})"

from ow have *: " $\sigma'$  ii =  $\sigma$  ii  $\implies$  other U {ii}  $\sigma$   $\sigma'$ "
  by (clarsimp elim!: otherwithE) (rule otherI, simp_all, metis sgivesu)
from tr have "(( $\sigma$ , NodeS ii  $\zeta$  R), a, ( $\sigma'$ , NodeS ii  $\zeta'$  R'))  $\in$  onode_sos (trans T)"
  by (simp add: onode_comps)
thus "( $\sigma'$ ,  $\zeta'$ )  $\in$  oreachable T (otherwith S {ii} (orecvmsg I)) (other U {ii})"
proof cases
  case onode_bcast
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case onode_gcast
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case onode_ucast
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case onode_notucast
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case onode_deliver
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case onode_tau
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case onode_receive
  with ih and ow show ?thesis
  by (auto elim!: oreachable_local' otherwithE)
next
  case (onode_arrive m)
  hence " $\zeta' = \zeta$ " and " $\sigma'$  ii =  $\sigma$  ii" by auto
  from this(2) have "other U {ii}  $\sigma$   $\sigma'$ " by (rule *)
  with ih and ( $\zeta' = \zeta$ ) show ?thesis by auto
next
  case onode_connect1
  hence " $\zeta' = \zeta$ " and " $\sigma'$  ii =  $\sigma$  ii" by auto
  from this(2) have "other U {ii}  $\sigma$   $\sigma'$ " by (rule *)
  with ih and ( $\zeta' = \zeta$ ) show ?thesis by auto
next
  case onode_connect2
  hence " $\zeta' = \zeta$ " and " $\sigma'$  ii =  $\sigma$  ii" by auto
  from this(2) have "other U {ii}  $\sigma$   $\sigma'$ " by (rule *)
  with ih and ( $\zeta' = \zeta$ ) show ?thesis by auto
next
  case onode_connect_other
  hence " $\zeta' = \zeta$ " and " $\sigma'$  ii =  $\sigma$  ii" by auto
  from this(2) have "other U {ii}  $\sigma$   $\sigma'$ " by (rule *)
  with ih and ( $\zeta' = \zeta$ ) show ?thesis by auto
next
  case onode_disconnect1
  hence " $\zeta' = \zeta$ " and " $\sigma'$  ii =  $\sigma$  ii" by auto
  from this(2) have "other U {ii}  $\sigma$   $\sigma'$ " by (rule *)
  with ih and ( $\zeta' = \zeta$ ) show ?thesis by auto
next

```

```

case onode_disconnect2
hence " $\zeta' = \zeta$ " and " $\sigma' \text{ ii} = \sigma \text{ ii}$ " by auto
from this(2) have "other U {ii}  $\sigma \sigma'$ " by (rule *)
with ih and " $\zeta' = \zeta$ " show ?thesis by auto
next
case onode_disconnect_other
hence " $\zeta' = \zeta$ " and " $\sigma' \text{ ii} = \sigma \text{ ii}$ " by auto
from this(2) have "other U {ii}  $\sigma \sigma'$ " by (rule *)
with ih and " $\zeta' = \zeta$ " show ?thesis by auto
qed
qed
qed

```

lemma node\_proc\_reachable\_statelessassm [dest]:

```

assumes " $(\sigma, \text{NodeS } i \ \zeta \ R) \in \text{oreachable } (\langle \text{ii} : T : R_i \rangle_o)$ "
      (otherwith  $(\lambda \_ \_ . \text{True}) \{ii\} (\text{oarrivemsg } I)$ )
      (other  $(\lambda \_ \_ . \text{True}) \{ii\}$ )"
shows " $(\sigma, \zeta) \in \text{oreachable } T$ "
      (otherwith  $(\lambda \_ \_ . \text{True}) \{ii\} (\text{orecvmsg } I)$ ) (other  $(\lambda \_ \_ . \text{True}) \{ii\}$ )"
using assms
by (rule node_proc_reachable) simp_all

```

lemma node\_lift:

```

assumes " $T \models (\text{otherwith } S \{ii\} (\text{orecvmsg } I), \text{ other } U \{ii\} \rightarrow) \text{ global } P$ "
      and " $\bigwedge \xi \xi'. S \xi \xi' \implies U \xi \xi'$ "
shows " $\langle \text{ii} : T : R_i \rangle_o \models (\text{otherwith } S \{ii\} (\text{oarrivemsg } I), \text{ other } U \{ii\} \rightarrow) \text{ global } P$ "
proof (rule oinvariant_oreachableI)
fix  $\sigma \ \zeta$ 
assume " $(\sigma, \zeta) \in \text{oreachable } (\langle \text{ii} : T : R_i \rangle_o) (\text{otherwith } S \{ii\} (\text{oarrivemsg } I)) (\text{other } U \{ii\})$ "
moreover then obtain  $i \ s \ R$  where " $\zeta = \text{NodeS } i \ s \ R$ "
  by (metis node_net_state)
ultimately have " $(\sigma, \text{NodeS } i \ s \ R) \in \text{oreachable } (\langle \text{ii} : T : R_i \rangle_o)$ "
      (otherwith  $S \{ii\} (\text{oarrivemsg } I)$ ) (other  $U \{ii\}$ )"
  by simp
hence " $(\sigma, s) \in \text{oreachable } T (\text{otherwith } S \{ii\} (\text{orecvmsg } I)) (\text{other } U \{ii\})$ "
  by - (erule node_proc_reachable, erule assms(2))
with assms(1) show "global P  $(\sigma, \zeta)$ "
  by (metis fst_conv globalsimp oinvariantD)
qed

```

lemma node\_lift\_step [intro]:

```

assumes pinv: " $T \models_A (\text{otherwith } S \{i\} (\text{orecvmsg } I), \text{ other } U \{i\} \rightarrow) \text{ globala } (\lambda(\sigma, \_, \sigma'). Q \sigma \sigma')$ "
      and other: " $\bigwedge \sigma \sigma'. \text{ other } U \{i\} \sigma \sigma' \implies Q \sigma \sigma'$ "
      and sgivesu: " $\bigwedge \xi \xi'. S \xi \xi' \implies U \xi \xi'$ "
shows " $\langle i : T : R_i \rangle_o \models_A (\text{otherwith } S \{i\} (\text{oarrivemsg } I), \text{ other } U \{i\} \rightarrow)$ "
      globala  $(\lambda(\sigma, \_, \sigma'). Q \sigma \sigma')$ "
(is " $\_ \models_A (?S, ?U \rightarrow) \_$ ")
proof (rule ostep_invariantI, simp)
fix  $\sigma \ s \ a \ \sigma' \ s'$ 
assume rs: " $(\sigma, s) \in \text{oreachable } (\langle i : T : R_i \rangle_o) ?S ?U$ "
      and tr: " $((\sigma, s), a, (\sigma', s')) \in \text{trans } (\langle i : T : R_i \rangle_o)$ "
      and ow: " $?S \sigma \sigma' \ a$ "
from ow have *: " $\sigma' \ i = \sigma \ i \implies \text{ other } U \{i\} \sigma \sigma'$ "
  by (clarsimp elim!: otherwithE) (rule otherI, simp_all, metis sgivesu)
from rs tr obtain  $\zeta \ R$ 
  where [simp]: " $s = \text{NodeS } i \ \zeta \ R$ "
      and " $(\sigma, \text{NodeS } i \ \zeta \ R) \in \text{oreachable } (\langle i : T : R_i \rangle_o) ?S ?U$ "
  by (metis node_net_state)
from this(2) have or: " $(\sigma, \zeta) \in \text{oreachable } T (\text{otherwith } S \{i\} (\text{orecvmsg } I)) ?U$ "
  by (rule node_proc_reachable [OF _ assms(3)])
from tr have " $((\sigma, \text{NodeS } i \ \zeta \ R), a, (\sigma', s')) \in \text{onode\_sos } (\text{trans } T)$ "
  by (simp add: onode_comps)
thus " $Q \sigma \sigma'$ "
proof cases

```

```

fix m ζ'
assume "a = R:*cast(m)"
  and tr': "((σ, ζ), broadcast m, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' (broadcast m)"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix D m ζ'
assume "a = (R ∩ D):*cast(m)"
  and tr': "((σ, ζ), groupcast D m, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' (groupcast D m)"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix d m ζ'
assume "a = {d}:*cast(m)"
  and tr': "((σ, ζ), unicast d m, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' (unicast d m)"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix d ζ'
assume "a = τ"
  and tr': "((σ, ζ), ¬unicast d, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' (¬unicast d)"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix d ζ'
assume "a = i:deliver(d)"
  and tr': "((σ, ζ), deliver d, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' (deliver d)"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix ζ'
assume "a = τ"
  and tr': "((σ, ζ), τ, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' τ"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix m ζ'
assume "a = {i}¬{i}:arrive(m)"
  and tr': "((σ, ζ), receive m, (σ', ζ')) ∈ trans T"
from this(1) and ⟨?S σ σ' a⟩ have "otherwith S {i} (orecvmsg I) σ σ' (receive m)"
  by (auto elim!: otherwithE)
with or tr' show ?thesis by (rule ostep_invariantD [OF pinv, simplified])
next
fix m
assume "a = {i}¬{i}:arrive(m)"
  and "σ' i = σ i"
from this(2) have "other U {i} σ σ'" by (rule *)
thus ?thesis by (rule other)
next
fix i'
assume "a = connect(i, i'"
  and "σ' i = σ i"
from this(2) have "other U {i} σ σ'" by (rule *)
thus ?thesis by (rule other)
next
fix i'
assume "a = connect(i', i)"
  and "σ' i = σ i"

```

```

    from this(2) have "other U {i}  $\sigma$   $\sigma'$ " by (rule *)
  thus ?thesis by (rule other)
next
  fix i' i''
  assume "a = connect(i', i'')"
    and " $\sigma'$  i =  $\sigma$  i"
  from this(2) have "other U {i}  $\sigma$   $\sigma'$ " by (rule *)
  thus ?thesis by (rule other)
next
  fix i'
  assume "a = disconnect(i, i')"
    and " $\sigma'$  i =  $\sigma$  i"
  from this(2) have "other U {i}  $\sigma$   $\sigma'$ " by (rule *)
  thus ?thesis by (rule other)
next
  fix i'
  assume "a = disconnect(i', i)"
    and " $\sigma'$  i =  $\sigma$  i"
  from this(2) have "other U {i}  $\sigma$   $\sigma'$ " by (rule *)
  thus ?thesis by (rule other)
next
  fix i' i''
  assume "a = disconnect(i', i'')"
    and " $\sigma'$  i =  $\sigma$  i"
  from this(2) have "other U {i}  $\sigma$   $\sigma'$ " by (rule *)
  thus ?thesis by (rule other)
qed
qed

```

lemma node\_lift\_step\_statelessassm [intro]:

```

  assumes "T  $\models_A$  ( $\lambda \sigma \_ . \text{orecvmsg } I \sigma, \text{other } (\lambda \_ \_ . \text{True}) \{i\} \rightarrow$ )
    globala ( $\lambda(\sigma, \_, \sigma'). Q (\sigma i) (\sigma' i)$ )"
  and " $\bigwedge \xi. Q \xi \xi$ "
  shows " $\langle i : T : R_i \rangle_o \models_A$  ( $\lambda \sigma \_ . \text{oarrivmsg } I \sigma, \text{other } (\lambda \_ \_ . \text{True}) \{i\} \rightarrow$ )
    globala ( $\lambda(\sigma, \_, \sigma'). Q (\sigma i) (\sigma' i)$ )"

```

proof -

```

  from assms(1)
  have "T  $\models_A$  ( $\text{otherwith } (\lambda \_ \_ . \text{True}) \{i\} (\text{orecvmsg } I), \text{other } (\lambda \_ \_ . \text{True}) \{i\} \rightarrow$ )
    globala ( $\lambda(\sigma, \_, \sigma'). Q (\sigma i) (\sigma' i)$ )"
  by rule auto
  with assms(2) have " $\langle i : T : R_i \rangle_o \models_A$  ( $\text{otherwith } (\lambda \_ \_ . \text{True}) \{i\} (\text{oarrivmsg } I),$ 
     $\text{other } (\lambda \_ \_ . \text{True}) \{i\} \rightarrow$ )
    globala ( $\lambda(\sigma, \_, \sigma'). Q (\sigma i) (\sigma' i)$ )"
  by - (rule node_lift_step, auto)
  thus ?thesis by rule auto
qed

```

lemma node\_lift\_anycast [intro]:

```

  assumes pinv: "T  $\models_A$  ( $\text{otherwith } S \{i\} (\text{orecvmsg } I), \text{other } U \{i\} \rightarrow$ )
    globala ( $\lambda(\sigma, a, \sigma'). \text{anycast } (Q \sigma \sigma') a$ )"
  and " $\bigwedge \xi \xi'. S \xi \xi' \implies U \xi \xi'$ "
  shows " $\langle i : T : R_i \rangle_o \models_A$  ( $\text{otherwith } S \{i\} (\text{oarrivmsg } I), \text{other } U \{i\} \rightarrow$ )
    globala ( $\lambda(\sigma, a, \sigma'). \text{castmsg } (Q \sigma \sigma') a$ )"
  (is " $\_ \models_A$  (?S, ?U  $\rightarrow$ )  $\_$ ")

```

proof (rule ostep\_invariantI, simp)

```

  fix  $\sigma s a \sigma' s'$ 
  assume rs: " $(\sigma, s) \in \text{oreachable } (\langle i : T : R_i \rangle_o) ?S ?U$ "
    and tr: " $((\sigma, s), a, (\sigma', s')) \in \text{trans } (\langle i : T : R_i \rangle_o)$ "
    and " $?S \sigma \sigma' a$ "
  from this(1-2) obtain  $\zeta R$ 
  where [simp]: " $s = \text{NodeS } i \zeta R$ "
    and " $(\sigma, \text{NodeS } i \zeta R) \in \text{oreachable } (\langle i : T : R_i \rangle_o) ?S ?U$ "
  by (metis node_net_state)
  from this(2) have " $(\sigma, \zeta) \in \text{oreachable } T (\text{otherwith } S \{i\} (\text{orecvmsg } I)) ?U$ "

```

```

    by (rule node_proc_reachable [OF _ assms(2)])
  moreover from tr have "((σ, NodeS i ζ R), a, (σ', s')) ∈ onode_sos (trans T)"
    by (simp add: onode_comps)
  ultimately show "castmsg (Q σ σ') a" using ⟨?S σ σ' a⟩
    by - (erule onode_sos.cases, auto elim!: otherwithE dest!: ostep_invariantD [OF pinv])
qed

```

lemma node\_lift\_anycast\_statelessassm [intro]:

```

assumes pinv: "T ⊢A (λσ _. orecvmsg I σ, other (λ_ _. True) {i} →)
              globala (λ(σ, a, σ'). anycast (Q σ σ') a)"
shows "⟨i : T : Ri⟩o ⊢A (λσ _. oarrivemsg I σ, other (λ_ _. True) {i} →)
              globala (λ(σ, a, σ'). castmsg (Q σ σ') a)"
  (is "_ ⊢A (?S, _ →) _")
proof -
  from assms(1)
  have "T ⊢A (otherwith (λ_ _. True) {i} (orecvmsg I), other (λ_ _. True) {i} →)
          globala (λ(σ, a, σ'). anycast (Q σ σ') a)"
    by rule auto
  hence "⟨i : T : Ri⟩o ⊢A (otherwith (λ_ _. True) {i} (oarrivemsg I), other (λ_ _. True) {i} →)
          globala (λ(σ, a, σ'). castmsg (Q σ σ') a)"
    by (rule node_lift_anycast) simp_all
  thus ?thesis
    by rule auto
qed

```

lemma node\_local\_deliver:

```

"⟨i : ζi : Ri⟩o ⊢A (S, U →) globala (λ(_, a, _). ∀j. j≠i → (∀d. a ≠ j:deliver(d)))"
proof (rule ostep_invariantI, simp)
  fix σ s a σ' s'
  assume 1: "(σ, s) ∈ oreachable (⟨i : ζi : Ri⟩o) S U"
    and 2: "((σ, s), a, (σ', s')) ∈ trans (⟨i : ζi : Ri⟩o)"
    and "S σ σ' a"
  moreover from 1 2 obtain ζ R ζ' R' where "s = NodeS i ζ R" and "s' = NodeS i ζ' R'" ..
  ultimately show "∀j. j≠i → (∀d. a ≠ j:deliver(d))"
    by (cases a) (auto simp add: onode_comps)
qed

```

lemma node\_tau\_deliver\_unchanged:

```

"⟨i : ζi : Ri⟩o ⊢A (S, U →) globala (λ(σ, a, σ'). a = τ ∨ (∃i d. a = i:deliver(d))
              → (∀j. j≠i → σ' j = σ j))"
proof (rule ostep_invariantI, clarsimp simp only: globalasimp snd_conv fst_conv)
  fix σ s a σ' s' j
  assume 1: "(σ, s) ∈ oreachable (⟨i : ζi : Ri⟩o) S U"
    and 2: "((σ, s), a, (σ', s')) ∈ trans (⟨i : ζi : Ri⟩o)"
    and "S σ σ' a"
    and "a = τ ∨ (∃i d. a = i:deliver(d))"
    and "j ≠ i"
  moreover from 1 2 obtain ζ R ζ' R' where "s = NodeS i ζ R" and "s' = NodeS i ζ' R'" ..
  ultimately show "σ' j = σ j"
    by (cases a) (auto simp del: step_node_tau simp add: onode_comps)
qed

```

end

## 17 Lifting rules for (open) partial networks

theory OPnet\_Lifting

imports ONode\_Lifting OAWN\_SOS OPnet

begin

lemma oreachable\_par\_subnet\_induct [consumes, case\_names init other local]:

```

assumes "(σ, SubnetS s t) ∈ oreachable (opnet onp (p1 || p2)) S U"
  and init: "∧σ s t. (σ, SubnetS s t) ∈ init (opnet onp (p1 || p2)) ⇒ P σ s t"
  and other: "∧σ s t σ'. [(σ, SubnetS s t) ∈ oreachable (opnet onp (p1 || p2)) S U;

```

```

      U σ σ'; P σ s t ] ⇒ P σ' s t"
and local: "∧σ s t σ' s' t' a. [ (σ, SubnetS s t) ∈ oreachable (opnet onp (p1 || p2)) S U;
      ((σ, SubnetS s t), a, (σ', SubnetS s' t')) ∈ trans (opnet onp (p1 || p2));
      S σ σ' a; P σ s t ] ⇒ P σ' s' t'"
shows "P σ s t"
using assms(1) proof (induction "(σ, SubnetS s t)" arbitrary: s t σ)
  fix s t σ
  assume "(σ, SubnetS s t) ∈ init (opnet onp (p1 || p2))"
  with init show "P σ s t" .
next
  fix st a s' t' σ'
  assume "st ∈ oreachable (opnet onp (p1 || p2)) S U"
    and tr: "(st, a, (σ', SubnetS s' t')) ∈ trans (opnet onp (p1 || p2))"
    and "S (fst st) (fst (σ', SubnetS s' t')) a"
    and IH: "∧s t σ. st = (σ, SubnetS s t) ⇒ P σ s t"
  from this(1) obtain s t σ where "st = (σ, SubnetS s t)"
    and "(σ, SubnetS s t) ∈ oreachable (opnet onp (p1 || p2)) S U"
  by (metis net_par_oreachable_is_subnet prod.collapse)
  note this(2)
  moreover from tr and ⟨st = (σ, SubnetS s t)⟩
  have "((σ, SubnetS s t), a, (σ', SubnetS s' t')) ∈ trans (opnet onp (p1 || p2))" by simp
  moreover from ⟨S (fst st) (fst (σ', SubnetS s' t')) a⟩ and ⟨st = (σ, SubnetS s t)⟩
  have "S σ σ' a" by simp
  moreover from IH and ⟨st = (σ, SubnetS s t)⟩ have "P σ s t" .
  ultimately show "P σ' s' t'" by (rule local)
next
  fix st σ' s t
  assume "st ∈ oreachable (opnet onp (p1 || p2)) S U"
    and "U (fst st) σ'"
    and "snd st = SubnetS s t"
    and IH: "∧s t σ. st = (σ, SubnetS s t) ⇒ P σ s t"
  from this(1,3) obtain σ where "st = (σ, SubnetS s t)"
    and "(σ, SubnetS s t) ∈ oreachable (opnet onp (p1 || p2)) S U"
  by (metis prod.collapse)
  note this(2)
  moreover from ⟨U (fst st) σ'⟩ and ⟨st = (σ, SubnetS s t)⟩ have "U σ σ'" by simp
  moreover from IH and ⟨st = (σ, SubnetS s t)⟩ have "P σ s t" .
  ultimately show "P σ' s t" by (rule other)
qed

lemma other_net_tree_ips_par_left:
  assumes "other U (net_tree_ips (p1 || p2)) σ σ'"
    and "∧ξ. U ξ ξ"
  shows "other U (net_tree_ips p1) σ σ'"
proof -
  from assms(1) obtain ineq: "∀i∈net_tree_ips (p1 || p2). σ' i = σ i"
    and outU: "∀j. j∉net_tree_ips (p1 || p2) → U (σ j) (σ' j)" ..
  show ?thesis
  proof (rule otherI)
    fix i
    assume "i∈net_tree_ips p1"
    hence "i∈net_tree_ips (p1 || p2)" by simp
    with ineq show "σ' i = σ i" ..
  next
    fix j
    assume "j∉net_tree_ips p1"
    show "U (σ j) (σ' j)"
    proof (cases "j∈net_tree_ips p2")
      assume "j∈net_tree_ips p2"
      hence "j∈net_tree_ips (p1 || p2)" by simp
      with ineq have "σ' j = σ j" ..
      thus "U (σ j) (σ' j)"
      by simp (rule ⟨∧ξ. U ξ ξ⟩)
    next
  end
end

```

```

    assume "j∉net_tree_ips p₂"
    with ⟨j∉net_tree_ips p₁⟩ have "j∉net_tree_ips (p₁ || p₂)" by simp
    with outU show "U (σ j) (σ' j)" by simp
  qed
qed
qed

```

lemma other\_net\_tree\_ips\_par\_right:

```

assumes "other U (net_tree_ips (p₁ || p₂)) σ σ'"
  and "∧ξ. U ξ ξ"
shows "other U (net_tree_ips p₂) σ σ'"

```

proof -

```

from assms(1) have "other U (net_tree_ips (p₂ || p₁)) σ σ'"
  by (subst net_tree_ips_commute)
thus ?thesis using ⟨∧ξ. U ξ ξ⟩
  by (rule other_net_tree_ips_par_left)

```

qed

lemma ostep\_arrive\_invariantD [elim]:

```

assumes "p ⊨A (λσ -. oarrivemsg I σ, U →) P"
  and "(σ, s) ∈ oreachable p (otherwith S IPS (oarrivemsg I)) U"
  and "((σ, s), a, (σ', s')) ∈ trans p"
  and "oarrivemsg I σ a"
shows "P ((σ, s), a, (σ', s'))"

```

proof -

```

from assms(2) have "(σ, s) ∈ oreachable p (λσ -. a. oarrivemsg I σ a) U"
  by (rule oreachable_weakenE) auto
thus "P ((σ, s), a, (σ', s'))"
  using assms(3-4) by (rule ostep_invariantD [OF assms(1)])

```

qed

lemma opnet\_sync\_action\_subnet\_oreachable:

```

assumes "(σ, SubnetS s t) ∈ oreachable (opnet onp (p₁ || p₂))
  (λσ -. oarrivemsg I σ) (other U (net_tree_ips (p₁ || p₂)))"
  (is "_ ∈ oreachable _ (?S (p₁ || p₂)) (?U (p₁ || p₂))")

```

and "∧ξ. U ξ ξ"

```

and act1: "opnet onp p₁ ⊨A (λσ -. oarrivemsg I σ, other U (net_tree_ips p₁) →)
  globala (λ(σ, a, σ'). castmsg (I σ) a
    ∧ (a = τ ∨ (∃i d. a = i:deliver(d)) →
      ((∀i∈net_tree_ips p₁. U (σ i) (σ' i))
        ∧ (∀i. i∉net_tree_ips p₁ → σ' i = σ i))))"

```

```

and act2: "opnet onp p₂ ⊨A (λσ -. oarrivemsg I σ, other U (net_tree_ips p₂) →)
  globala (λ(σ, a, σ'). castmsg (I σ) a
    ∧ (a = τ ∨ (∃i d. a = i:deliver(d)) →
      ((∀i∈net_tree_ips p₂. U (σ i) (σ' i))
        ∧ (∀i. i∉net_tree_ips p₂ → σ' i = σ i))))"

```

```

shows "(σ, s) ∈ oreachable (opnet onp p₁) (λσ -. oarrivemsg I σ) (other U (net_tree_ips p₁))
  ∧ (σ, t) ∈ oreachable (opnet onp p₂) (λσ -. oarrivemsg I σ) (other U (net_tree_ips p₂))
  ∧ net_tree_ips p₁ ∩ net_tree_ips p₂ = {}"

```

using assms(1)

proof (induction rule: oreachable\_par\_subnet\_induct)

case (init σ s t)

hence sinit: "(σ, s) ∈ init (opnet onp p₁)"

and tinit: "(σ, t) ∈ init (opnet onp p₂)"

and "net\_ips s ∩ net\_ips t = {}" by auto

moreover from sinit have "net\_ips s = net\_tree\_ips p₁"

by (rule opnet\_net\_ips\_net\_tree\_ips\_init)

moreover from tinit have "net\_ips t = net\_tree\_ips p₂"

by (rule opnet\_net\_ips\_net\_tree\_ips\_init)

ultimately show ?case by (auto elim: oreachable\_init)



```

next
case (other  $\sigma$  s t  $\sigma'$ )
hence "other U (net_tree_ips (p1 || p2))  $\sigma$   $\sigma'$ "
  and IHs: " $(\sigma, s) \in \text{oreachable (opnet onp p}_1) (?S p_1) (?U p_1)$ "
  and Iht: " $(\sigma, t) \in \text{oreachable (opnet onp p}_2) (?S p_2) (?U p_2)$ "
  and "net_tree_ips p1  $\cap$  net_tree_ips p2 = {}" by auto

have " $(\sigma', s) \in \text{oreachable (opnet onp p}_1) (?S p_1) (?U p_1)$ "
proof -
  from (?U (p1 || p2)  $\sigma$   $\sigma'$ ) and  $\langle \bigwedge \xi. U \xi \xi \rangle$  have "?U p1  $\sigma$   $\sigma'$ "
  by (rule other_net_tree_ips_par_left)
  with IHs show ?thesis by - (erule(1) oreachable_other')
qed

moreover have " $(\sigma', t) \in \text{oreachable (opnet onp p}_2) (?S p_2) (?U p_2)$ "
proof -
  from (?U (p1 || p2)  $\sigma$   $\sigma'$ ) and  $\langle \bigwedge \xi. U \xi \xi \rangle$  have "?U p2  $\sigma$   $\sigma'$ "
  by (rule other_net_tree_ips_par_right)
  with Iht show ?thesis by - (erule(1) oreachable_other')
qed

ultimately show ?case using  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$  by simp
next
case (local  $\sigma$  s t  $\sigma'$  s' t' a)
hence stor: " $(\sigma, \text{SubnetS } s \ t) \in \text{oreachable (opnet onp (p}_1 \ || \ p_2)) (?S (p_1 \ || \ p_2)) (?U (p_1 \ || \ p_2))$ "
  and tr: " $((\sigma, \text{SubnetS } s \ t), a, (\sigma', \text{SubnetS } s' \ t')) \in \text{trans (opnet onp (p}_1 \ || \ p_2))$ "
  and "oarrivmsg I  $\sigma$  a"
  and sor: " $(\sigma, s) \in \text{oreachable (opnet onp p}_1) (?S p_1) (?U p_1)$ "
  and tor: " $(\sigma, t) \in \text{oreachable (opnet onp p}_2) (?S p_2) (?U p_2)$ "
  and "net_tree_ips p1  $\cap$  net_tree_ips p2 = {}" by auto
from tr have " $((\sigma, \text{SubnetS } s \ t), a, (\sigma', \text{SubnetS } s' \ t'))$ 
   $\in \text{opnet\_sos (trans (opnet onp p}_1)) (\text{trans (opnet onp p}_2))$ " by simp
hence " $(\sigma', s') \in \text{oreachable (opnet onp p}_1) (?S p_1) (?U p_1)$ 
   $\wedge (\sigma', t') \in \text{oreachable (opnet onp p}_2) (?S p_2) (?U p_2)$ "
proof (cases)
  fix H K m H' K'
  assume "a = (H  $\cup$  H')  $\neg$  (K  $\cup$  K'):arrive(m)"
  and str: " $((\sigma, s), H \neg K:\text{arrive}(m), (\sigma', s')) \in \text{trans (opnet onp p}_1)$ "
  and ttr: " $((\sigma, t), H' \neg K':\text{arrive}(m), (\sigma', t')) \in \text{trans (opnet onp p}_2)$ "
  from this(1) and  $\langle \text{oarrivmsg } I \ \sigma \ a \rangle$  have "I  $\sigma$  m" by simp

  with sor str
  have " $(\sigma', s') \in \text{oreachable (opnet onp p}_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
  moreover from  $\langle I \ \sigma \ m \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable (opnet onp p}_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
  ultimately show ?thesis ..
next
fix R m H K
assume str: " $((\sigma, s), R:\text{*cast}(m), (\sigma', s')) \in \text{trans (opnet onp p}_1)$ "
  and ttr: " $((\sigma, t), H \neg K:\text{arrive}(m), (\sigma', t')) \in \text{trans (opnet onp p}_2)$ "
from sor str have "I  $\sigma$  m"
  by - (drule(1) ostep_invariantD [OF act1], simp_all)
with sor str
  have " $(\sigma', s') \in \text{oreachable (opnet onp p}_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
  moreover from  $\langle I \ \sigma \ m \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable (opnet onp p}_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
  ultimately show ?thesis ..
next
fix R m H K
assume str: " $((\sigma, s), H \neg K:\text{arrive}(m), (\sigma', s')) \in \text{trans (opnet onp p}_1)$ "

```

```

    and ttr: " $((\sigma, t), R:\text{*cast}(m), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "
  from tor ttr have "I  $\sigma$  m"
  by - (drule(1) ostep_invariantD [OF act2], simp_all)
  with sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
  moreover from  $\langle I \sigma m \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
  ultimately show ?thesis ..
next
fix i i'
assume str: " $((\sigma, s), \text{connect}(i, i'), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), \text{connect}(i, i'), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "
  with sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
  moreover from tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
  ultimately show ?thesis ..
next
fix i i'
assume str: " $((\sigma, s), \text{disconnect}(i, i'), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), \text{disconnect}(i, i'), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "
  with sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
  moreover from tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
  ultimately show ?thesis ..
next
fix i d
assume " $t' = t$ "
  and str: " $((\sigma, s), i:\text{deliver}(d), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "

  from sor str have " $\forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j$ "
  by - (drule(1) ostep_invariantD [OF act1], simp_all)
  moreover with  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$ 
  have " $\forall j. j \in \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j$ " by auto
  moreover from sor str have " $\forall j \in \text{net\_tree\_ips } p_1. U(\sigma j) (\sigma' j)$ "
  by - (drule(1) ostep_invariantD [OF act1], simp_all)
  ultimately have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  using tor  $\langle t' = t \rangle$  by (clarsimp elim!: oreachable_other')
  (metis otherI  $\langle \bigwedge \xi. U \xi \xi \rangle$ +)

  moreover from sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
  ultimately show ?thesis by (rule conjI [rotated])
next
fix i d
assume " $s' = s$ "
  and ttr: " $((\sigma, t), i:\text{deliver}(d), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "

  from tor ttr have " $\forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j$ "
  by - (drule(1) ostep_invariantD [OF act2], simp_all)
  moreover with  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$ 
  have " $\forall j. j \in \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j$ " by auto
  moreover from tor ttr have " $\forall j \in \text{net\_tree\_ips } p_2. U(\sigma j) (\sigma' j)$ "
  by - (drule(1) ostep_invariantD [OF act2], simp_all)
  ultimately have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  using sor  $\langle s' = s \rangle$  by (clarsimp elim!: oreachable_other')

```

```

      (metis otherI (⋀ξ. U ξ ξ))+

moreover from tor ttr
  have "(σ', t') ∈ oreachable (opnet onp p₂) (?S p₂) (?U p₂)"
    by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
next
assume "t' = t"
  and str: "((σ, s), τ, (σ', s')) ∈ trans (opnet onp p₁)"

from sor str have "∀j. j ∉ net_tree_ips p₁ → σ' j = σ j"
  by - (drule(1) ostep_invariantD [OF act1], simp_all)
moreover with (net_tree_ips p₁ ∩ net_tree_ips p₂ = {})
  have "∀j. j ∈ net_tree_ips p₂ → σ' j = σ j" by auto
moreover from sor str have "∀j ∈ net_tree_ips p₁. U (σ j) (σ' j)"
  by - (drule(1) ostep_invariantD [OF act1], simp_all)
ultimately have "(σ', t') ∈ oreachable (opnet onp p₂) (?S p₂) (?U p₂)"
  using tor (t' = t) by (clarsimp elim!: oreachable_other')
      (metis otherI (⋀ξ. U ξ ξ))+

moreover from sor str
  have "(σ', s') ∈ oreachable (opnet onp p₁) (?S p₁) (?U p₁)"
    by - (erule(1) oreachable_local, auto)
ultimately show ?thesis by (rule conjI [rotated])
next
assume "s' = s"
  and ttr: "((σ, t), τ, (σ', t')) ∈ trans (opnet onp p₂)"

from tor ttr have "∀j. j ∉ net_tree_ips p₂ → σ' j = σ j"
  by - (drule(1) ostep_invariantD [OF act2], simp_all)
moreover with (net_tree_ips p₁ ∩ net_tree_ips p₂ = {})
  have "∀j. j ∈ net_tree_ips p₁ → σ' j = σ j" by auto
moreover from tor ttr have "∀j ∈ net_tree_ips p₂. U (σ j) (σ' j)"
  by - (drule(1) ostep_invariantD [OF act2], simp_all)
ultimately have "(σ', s') ∈ oreachable (opnet onp p₁) (?S p₁) (?U p₁)"
  using sor (s' = s) by (clarsimp elim!: oreachable_other')
      (metis otherI (⋀ξ. U ξ ξ))+

moreover from tor ttr
  have "(σ', t') ∈ oreachable (opnet onp p₂) (?S p₂) (?U p₂)"
    by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
qed
with (net_tree_ips p₁ ∩ net_tree_ips p₂ = {}) show ?case by simp
qed

```

'Splitting' reachability is trivial when there are no assumptions on interleavings, but this is useless for showing non-trivial properties, since the interleaving steps can do anything at all. This lemma is too weak.

**lemma subnet\_oreachable\_true\_true:**

```

assumes "(σ, SubnetS s₁ s₂) ∈ oreachable (opnet onp (p₁ || p₂)) (λ_ _ . True) (λ_ _ . True)"
shows "(σ, s₁) ∈ oreachable (opnet onp p₁) (λ_ _ . True) (λ_ _ . True)"
      "(σ, s₂) ∈ oreachable (opnet onp p₂) (λ_ _ . True) (λ_ _ . True)"
      (is "_ ∈ ?oreachable p₂")
using assms proof -
from assms have "(σ, s₁) ∈ ?oreachable p₁ ∧ (σ, s₂) ∈ ?oreachable p₂"
proof (induction rule: oreachable_par_subnet_induct)
  fix σ s₁ s₂
  assume "(σ, SubnetS s₁ s₂) ∈ init (opnet onp (p₁ || p₂))"
  thus "(σ, s₁) ∈ ?oreachable p₁ ∧ (σ, s₂) ∈ ?oreachable p₂"
    by (auto dest: oreachable_init)
next
case (local σ s₁ s₂ σ' s₁' s₂' a)
hence "(σ, SubnetS s₁ s₂) ∈ ?oreachable (p₁ || p₂)"
  and sr1: "(σ, s₁) ∈ ?oreachable p₁"

```

```

and sr2: "(σ, s₂) ∈ ?oreachable p₂"
and "((σ, SubnetS s₁ s₂), a, (σ', SubnetS s₁' s₂')) ∈ trans (opnet onp (p₁ || p₂))" by auto
from this(4)
  have "((σ, SubnetS s₁ s₂), a, (σ', SubnetS s₁' s₂'))
    ∈ opnet_sos (trans (opnet onp p₁)) (trans (opnet onp p₂))" by simp
thus "(σ', s₁') ∈ ?oreachable p₁ ∧ (σ', s₂') ∈ ?oreachable p₂"
proof cases
  fix R m H K
  assume "a = R:*cast(m)"
    and tr1: "((σ, s₁), R:*cast(m), (σ', s₁')) ∈ trans (opnet onp p₁)"
    and tr2: "((σ, s₂), H¬K:arrive(m), (σ', s₂')) ∈ trans (opnet onp p₂)"
  from sr1 and tr1 and TrueI have "(σ', s₁') ∈ ?oreachable p₁"
  by (rule ooreachable_local')
  moreover from sr2 and tr2 and TrueI have "(σ', s₂') ∈ ?oreachable p₂"
  by (rule ooreachable_local')
  ultimately show ?thesis ..
next
  assume "a = τ"
    and "s₂' = s₂"
    and tr1: "((σ, s₁), τ, (σ', s₁')) ∈ trans (opnet onp p₁)"
  from sr2 and this(2) have "(σ', s₂') ∈ ?oreachable p₂" by auto
  moreover have "(λ_ _ . True) σ σ'" by (rule TrueI)
  ultimately have "(σ', s₂') ∈ ?oreachable p₂"
  by (rule ooreachable_other')
  moreover from sr1 and tr1 and TrueI have "(σ', s₁') ∈ ?oreachable p₁"
  by (rule ooreachable_local')
qed (insert sr1 sr2, simp_all, (metis (no_types) ooreachable_local'
  ooreachable_other'+))

qed auto
thus "(σ, s₁) ∈ ?oreachable p₁"
  "(σ, s₂) ∈ ?oreachable p₂" by auto
qed

```

It may also be tempting to try splitting from the assumption  $(\sigma, \text{SubnetS } s_1 s_2) \in \text{oreachable } (\text{opnet onp } (p_1 \parallel p_2)) (\lambda\_ \_ \_ . \text{True}) (\lambda\_ \_ \_ . \text{False})$ , where the environment step would be trivially true (since the assumption is false), but the lemma cannot be shown when only one side acts, since it must guarantee the assumption for the other side.

**lemma lift\_opnet\_sync\_action:**

```

assumes "\ξ. U ξ ξ"
and act1: "\i R. ⟨i : onp i : R⟩_o ⊨_A (λσ _ . oarrivemsg I σ, other U {i} →)
  globala (λ(σ, a, _). castmsg (I σ) a)"
and act2: "\i R. ⟨i : onp i : R⟩_o ⊨_A (λσ _ . oarrivemsg I σ, other U {i} →)
  globala (λ(σ, a, σ'). (a ≠ τ ∧ (∀d. a ≠ i:deliver(d)) → S (σ i) (σ' i)))"
and act3: "\i R. ⟨i : onp i : R⟩_o ⊨_A (λσ _ . oarrivemsg I σ, other U {i} →)
  globala (λ(σ, a, σ'). (a = τ ∨ (∃d. a = i:deliver(d)) → U (σ i) (σ' i)))"
shows "opnet onp p ⊨_A (λσ _ . oarrivemsg I σ, other U (net_tree_ips p) →)
  globala (λ(σ, a, σ'). castmsg (I σ) a
    ∧ (a ≠ τ ∧ (∀i d. a ≠ i:deliver(d)) →
      (∀i∈net_tree_ips p. S (σ i) (σ' i)))
    ∧ (a = τ ∨ (∃i d. a = i:deliver(d)) →
      ((∀i∈net_tree_ips p. U (σ i) (σ' i))
        ∧ (∀i. i∉net_tree_ips p → σ' i = σ i))))"

```

(is "opnet onp p ⊨\_A (?I, ?U p →) ?inv (net\_tree\_ips p)")

**proof (induction p)**

fix i R

show "opnet onp ⟨i; R⟩ ⊨\_A (?I, ?U ⟨i; R⟩ →) ?inv (net\_tree\_ips ⟨i; R⟩)"

**proof (rule ostep\_invariantI, simp only: opnet.simps net\_tree\_ips.simps)**

fix σ s a σ' s'

assume sor: "(σ, s) ∈ ooreachable ((i : onp i : R)\_o) (λσ \_ . oarrivemsg I σ) (other U {i})"

and str: "((σ, s), a, (σ', s')) ∈ trans ((i : onp i : R)\_o)"

and oam: "oarrivemsg I σ a"

hence "castmsg (I σ) a"

by - (drule(2) ostep\_invariantD [OF act1], simp)

moreover from sor str oam have "a ≠ τ ∧ (∀i d. a ≠ i:deliver(d)) → S (σ i) (σ' i)"

```

    by - (drule(2) ostep_invariantD [OF act2], simp)
  moreover have "a =  $\tau \vee (\exists i d. a = i:\text{deliver}(d)) \longrightarrow U (\sigma i) (\sigma' i)"
  proof -
    from sor str oam have "a =  $\tau \vee (\exists d. a = i:\text{deliver}(d)) \longrightarrow U (\sigma i) (\sigma' i)"
      by - (drule(2) ostep_invariantD [OF act3], simp)
    moreover from sor str oam have " $\forall j. j \neq i \longrightarrow (\forall d. a \neq j:\text{deliver}(d))"$ "
      by - (drule(2) ostep_invariantD [OF node_local_deliver], simp)
    ultimately show ?thesis
      by clarsimp metis
  qed
  moreover from sor str oam have " $\forall j. j \neq i \longrightarrow (\forall d. a \neq j:\text{deliver}(d))"$ "
    by - (drule(2) ostep_invariantD [OF node_local_deliver], simp)
  moreover from sor str oam have "a =  $\tau \vee (\exists i d. a = i:\text{deliver}(d)) \longrightarrow (\forall j. j \neq i \longrightarrow \sigma' j = \sigma j)"
    by - (drule(2) ostep_invariantD [OF node_tau_deliver_unchanged], simp)
  ultimately show "?inv {i} (( $\sigma, s$ ), a, ( $\sigma', s'$ ))" by simp
  qed
next
  fix p1 p2
  assume inv1: "opnet onp p1  $\models_A$  (?I, ?U p1  $\rightarrow$ ) ?inv (net_tree_ips p1)"
    and inv2: "opnet onp p2  $\models_A$  (?I, ?U p2  $\rightarrow$ ) ?inv (net_tree_ips p2)"
  show "opnet onp (p1 || p2)  $\models_A$  (?I, ?U (p1 || p2)  $\rightarrow$ ) ?inv (net_tree_ips (p1 || p2))"
  proof (rule ostep_invariantI)
    fix  $\sigma$  st a  $\sigma'$  st'
    assume "( $\sigma, st$ )  $\in$  oreachable (opnet onp (p1 || p2)) ?I (?U (p1 || p2))"
      and "(( $\sigma, st$ ), a, ( $\sigma', st'$ ))  $\in$  trans (opnet onp (p1 || p2))"
      and "oarrivemsg I  $\sigma$  a"
    from this(1) obtain s t
      where "st = SubnetS s t"
      and *: "( $\sigma, \text{SubnetS } s t$ )  $\in$  oreachable (opnet onp (p1 || p2)) ?I (?U (p1 || p2))"
      by - (frule net_par_oreachable_is_subnet, metis)

    from this(2) and inv1 and inv2
      obtain sor: "( $\sigma, s$ )  $\in$  oreachable (opnet onp p1) ?I (?U p1)"
        and tor: "( $\sigma, t$ )  $\in$  oreachable (opnet onp p2) ?I (?U p2)"
        and "net_tree_ips p1  $\cap$  net_tree_ips p2 = {}"
        by - (drule opnet_sync_action_subnet_oreachable [OF _  $\langle \bigwedge \xi. U \xi \xi \rangle$ ], auto)

    from * and  $\langle ((\sigma, st), a, (\sigma', st')) \in \text{trans (opnet onp (p}_1 \parallel p_2)) \rangle$  and  $\langle st = \text{SubnetS } s t \rangle$ 
      obtain s' t' where "st' = SubnetS s' t'"
        and "(( $\sigma, \text{SubnetS } s t$ ), a, ( $\sigma', \text{SubnetS } s' t'$ ))
           $\in$  opnet_sos (trans (opnet onp p1)) (trans (opnet onp p2))"
        by clarsimp (frule opartial_net_preserves_subnets, metis)

    from this(2)
      have "castmsg (I  $\sigma$ ) a
         $\wedge$  (a  $\neq \tau \wedge (\forall i d. a \neq i:\text{deliver}(d)) \longrightarrow (\forall i \in \text{net\_tree\_ips (p}_1 \parallel p_2). S (\sigma i) (\sigma' i))$ )
         $\wedge$  (a =  $\tau \vee (\exists i d. a = i:\text{deliver}(d)) \longrightarrow (\forall i \in \text{net\_tree\_ips (p}_1 \parallel p_2). U (\sigma i) (\sigma' i))$ )
           $\wedge$  ( $\forall i. i \notin \text{net\_tree\_ips (p}_1 \parallel p_2) \longrightarrow \sigma' i = \sigma i$ )"

  proof cases
    fix R m H K
    assume "a = R:*cast(m)"
      and str: "(( $\sigma, s$ ), R:*cast(m), ( $\sigma', s'$ ))  $\in$  trans (opnet onp p1)"
      and ttr: "(( $\sigma, t$ ), H $\neg$ K:arrive(m), ( $\sigma', t'$ ))  $\in$  trans (opnet onp p2)"
    from sor and str have "I  $\sigma$  m  $\wedge (\forall i \in \text{net\_tree\_ips } p_1. S (\sigma i) (\sigma' i))"$ 
      by (auto dest: ostep_invariantD [OF inv1])
    moreover with tor and ttr have " $\forall i \in \text{net\_tree\_ips } p_2. S (\sigma i) (\sigma' i)"$ 
      by (auto dest: ostep_invariantD [OF inv2])
    ultimately show ?thesis
      using  $\langle a = R:*cast(m) \rangle$  by auto
  next
    fix R m H K
    assume "a = R:*cast(m)"
      and str: "(( $\sigma, s$ ), H $\neg$ K:arrive(m), ( $\sigma', s'$ ))  $\in$  trans (opnet onp p1)"
      and ttr: "(( $\sigma, t$ ), R:*cast(m), ( $\sigma', t'$ ))  $\in$  trans (opnet onp p2)"$$$ 
```

```

from tor and ttr have "I  $\sigma$  m  $\wedge$  ( $\forall i \in \text{net\_tree\_ips } p_2. S (\sigma i) (\sigma' i)$ )"
  by (auto dest: ostep_invariantD [OF inv2])
moreover with sor and str have " $\forall i \in \text{net\_tree\_ips } p_1. S (\sigma i) (\sigma' i)$ "
  by (auto dest: ostep_invariantD [OF inv1])
ultimately show ?thesis
  using  $\langle a = R:*cast(m) \rangle$  by auto
next
fix H K m H' K'
assume "a = (H  $\cup$  H')  $\neg$  (K  $\cup$  K'):arrive(m)"
  and str: " $((\sigma, s), H \neg K:arrive(m), (\sigma', s')) \in \text{trans (opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), H' \neg K':arrive(m), (\sigma', t')) \in \text{trans (opnet onp } p_2)$ "
from this(1) and  $\langle \text{oarrivemsg } I \sigma a \rangle$  have "I  $\sigma$  m" by simp
with sor and str have " $\forall i \in \text{net\_tree\_ips } p_1. S (\sigma i) (\sigma' i)$ "
  by (auto dest: ostep_invariantD [OF inv1])
moreover from tor and ttr and  $\langle I \sigma m \rangle$  have " $\forall i \in \text{net\_tree\_ips } p_2. S (\sigma i) (\sigma' i)$ "
  by (auto dest: ostep_invariantD [OF inv2])
ultimately show ?thesis
  using  $\langle a = (H \cup H') \neg (K \cup K'):arrive(m) \rangle$  by auto
next
fix i d
assume "a = i:deliver(d)"
  and str: " $((\sigma, s), i:deliver(d), (\sigma', s')) \in \text{trans (opnet onp } p_1)$ "
with sor have " $((\forall i \in \text{net\_tree\_ips } p_1. U (\sigma i) (\sigma' i))$ 
   $\wedge (\forall i. i \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' i = \sigma i))$ "
  by (auto dest!: ostep_invariantD [OF inv1])
with  $\langle a = i:deliver(d) \rangle$  and  $\langle \bigwedge \xi. U \xi \xi \rangle$  show ?thesis
  by auto
next
fix i d
assume "a = i:deliver(d)"
  and ttr: " $((\sigma, t), i:deliver(d), (\sigma', t')) \in \text{trans (opnet onp } p_2)$ "
with tor have " $((\forall i \in \text{net\_tree\_ips } p_2. U (\sigma i) (\sigma' i))$ 
   $\wedge (\forall i. i \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' i = \sigma i))$ "
  by (auto dest!: ostep_invariantD [OF inv2])
with  $\langle a = i:deliver(d) \rangle$  and  $\langle \bigwedge \xi. U \xi \xi \rangle$  show ?thesis
  by auto
next
assume "a =  $\tau$ "
  and str: " $((\sigma, s), \tau, (\sigma', s')) \in \text{trans (opnet onp } p_1)$ "
with sor have " $((\forall i \in \text{net\_tree\_ips } p_1. U (\sigma i) (\sigma' i))$ 
   $\wedge (\forall i. i \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' i = \sigma i))$ "
  by (auto dest!: ostep_invariantD [OF inv1])
with  $\langle a = \tau \rangle$  and  $\langle \bigwedge \xi. U \xi \xi \rangle$  show ?thesis
  by auto
next
assume "a =  $\tau$ "
  and ttr: " $((\sigma, t), \tau, (\sigma', t')) \in \text{trans (opnet onp } p_2)$ "
with tor have " $((\forall i \in \text{net\_tree\_ips } p_2. U (\sigma i) (\sigma' i))$ 
   $\wedge (\forall i. i \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' i = \sigma i))$ "
  by (auto dest!: ostep_invariantD [OF inv2])
with  $\langle a = \tau \rangle$  and  $\langle \bigwedge \xi. U \xi \xi \rangle$  show ?thesis
  by auto
next
fix i i'
assume "a = connect(i, i')"
  and str: " $((\sigma, s), \text{connect}(i, i'), (\sigma', s')) \in \text{trans (opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), \text{connect}(i, i'), (\sigma', t')) \in \text{trans (opnet onp } p_2)$ "
from sor and str have " $\forall i \in \text{net\_tree\_ips } p_1. S (\sigma i) (\sigma' i)$ "
  by (auto dest: ostep_invariantD [OF inv1])
moreover from tor and ttr have " $\forall i \in \text{net\_tree\_ips } p_2. S (\sigma i) (\sigma' i)$ "
  by (auto dest: ostep_invariantD [OF inv2])
ultimately show ?thesis
  using  $\langle a = \text{connect}(i, i') \rangle$  by auto
next

```

```

fix i i'
assume "a = disconnect(i, i')"
  and str: "((σ, s), disconnect(i, i'), (σ', s')) ∈ trans (opnet onp p₁)"
  and ttr: "((σ, t), disconnect(i, i'), (σ', t')) ∈ trans (opnet onp p₂)"
from sor and str have "∀i∈net_tree_ips p₁. S (σ i) (σ' i)"
  by (auto dest: ostep_invariantD [OF inv1])
moreover from tor and ttr have "∀i∈net_tree_ips p₂. S (σ i) (σ' i)"
  by (auto dest: ostep_invariantD [OF inv2])
ultimately show ?thesis
  using ⟨a = disconnect(i, i')⟩ by auto
qed
thus "?inv (net_tree_ips (p₁ || p₂)) ((σ, st), a, (σ', st'))" by simp
qed
qed

```

theorem subnet\_reachable:

```

assumes "(σ, SubnetS s t) ∈ reachable (opnet onp (p₁ || p₂))
  (otherwith S (net_tree_ips (p₁ || p₂)) (oarrivemsg I))
  (other U (net_tree_ips (p₁ || p₂)))"
  (is "_ ∈ reachable _ (?S (p₁ || p₂)) (?U (p₁ || p₂))")

and "∧ξ. S ξ ξ"
and "∧ξ. U ξ ξ"

and node1: "∧i R. ⟨i : onp i : R⟩ₒ ⊨ₐ (λσ _. oarrivemsg I σ, other U {i} →)
  globala (λ(σ, a, _). castmsg (I σ) a)"
and node2: "∧i R. ⟨i : onp i : R⟩ₒ ⊨ₐ (λσ _. oarrivemsg I σ, other U {i} →)
  globala (λ(σ, a, σ'). (a ≠ τ ∧ (∀d. a ≠ i:deliver(d)) → S (σ i) (σ' i)))"
and node3: "∧i R. ⟨i : onp i : R⟩ₒ ⊨ₐ (λσ _. oarrivemsg I σ, other U {i} →)
  globala (λ(σ, a, σ'). (a = τ ∨ (∃d. a = i:deliver(d)) → U (σ i) (σ' i)))"

```

```

shows "(σ, s) ∈ reachable (opnet onp p₁)
  (otherwith S (net_tree_ips p₁) (oarrivemsg I))
  (other U (net_tree_ips p₁))
  ∧ (σ, t) ∈ reachable (opnet onp p₂)
  (otherwith S (net_tree_ips p₂) (oarrivemsg I))
  (other U (net_tree_ips p₂))
  ∧ net_tree_ips p₁ ∩ net_tree_ips p₂ = {}"

```

using assms(1) proof (induction rule: reachable\_par\_subnet\_induct)

```

case (init σ s t)
hence sinit: "(σ, s) ∈ init (opnet onp p₁)"
  and tinit: "(σ, t) ∈ init (opnet onp p₂)"
  and "net_ips s ∩ net_ips t = {}" by auto
moreover from sinit have "net_ips s = net_tree_ips p₁"
  by (rule opnet_net_ips_net_tree_ips_init)
moreover from tinit have "net_ips t = net_tree_ips p₂"
  by (rule opnet_net_ips_net_tree_ips_init)
ultimately show ?case by (auto elim: reachable_init)

```

next

```

case (other σ s t σ')
hence "other U (net_tree_ips (p₁ || p₂)) σ σ'"
  and IHs: "(σ, s) ∈ reachable (opnet onp p₁) (?S p₁) (?U p₁)"
  and IHt: "(σ, t) ∈ reachable (opnet onp p₂) (?S p₂) (?U p₂)"
  and "net_tree_ips p₁ ∩ net_tree_ips p₂ = {}" by auto

```

have "(σ', s) ∈ reachable (opnet onp p₁) (?S p₁) (?U p₁)"

proof -

```

from ⟨?U (p₁ || p₂) σ σ'⟩ and ⟨∧ξ. U ξ ξ⟩ have "?U p₁ σ σ'"
  by (rule other_net_tree_ips_par_left)
with IHs show ?thesis by (erule(1) reachable_other')

```

qed

moreover have "(σ', t) ∈ reachable (opnet onp p₂) (?S p₂) (?U p₂)"

proof -

```

from ⟨?U (p1 || p2) σ σ'⟩ and ⟨∧ξ. U ξ ξ⟩ have "?U p2 σ σ'"
  by (rule other_net_tree_ips_par_right)
with IHt show ?thesis by - (erule(1) oreachable_other')
qed

ultimately show ?case using ⟨net_tree_ips p1 ∩ net_tree_ips p2 = {}⟩ by simp
next
case (local σ s t σ' s' t' a)
hence stor: "(σ, SubnetS s t) ∈ oreachable (opnet onp (p1 || p2)) (?S (p1 || p2)) (?U (p1 || p2))"
  and tr: "((σ, SubnetS s t), a, (σ', SubnetS s' t')) ∈ trans (opnet onp (p1 || p2))"
  and "?S (p1 || p2) σ σ' a"
  and sor: "(σ, s) ∈ oreachable (opnet onp p1) (?S p1) (?U p1)"
  and tor: "(σ, t) ∈ oreachable (opnet onp p2) (?S p2) (?U p2)"
  and "net_tree_ips p1 ∩ net_tree_ips p2 = {}" by auto

have act: "∧p. opnet onp p ⊨A (λσ -. oarrivemsg I σ, other U (net_tree_ips p) →)
  globala (λ(σ, a, σ'). castmsg (I σ) a
    ∧ (a ≠ τ ∧ (∀i d. a ≠ i:deliver(d)) →
      (∀i∈net_tree_ips p. S (σ i) (σ' i)))
    ∧ (a = τ ∨ (∃i d. a = i:deliver(d)) →
      ((∀i∈net_tree_ips p. U (σ i) (σ' i))
        ∧ (∀i. i∉net_tree_ips p → σ' i = σ i))))"
  by (rule lift_opnet_sync_action [OF assms(3-6)])

from ⟨?S (p1 || p2) σ σ' a⟩ have "∀j. j ∉ net_tree_ips (p1 || p2) → S (σ j) (σ' j)"
  and "oarrivemsg I σ a"
  by (auto elim!: otherwithE)
from tr have "((σ, SubnetS s t), a, (σ', SubnetS s' t'))
  ∈ opnet_sos (trans (opnet onp p1)) (trans (opnet onp p2))" by simp
hence "(σ', s') ∈ oreachable (opnet onp p1) (?S p1) (?U p1)
  ∧ (σ', t') ∈ oreachable (opnet onp p2) (?S p2) (?U p2)"
proof (cases)
fix H K m H' K'
assume "a = (H ∪ H') - (K ∪ K'):arrive(m)"
  and str: "((σ, s), H-K:arrive(m), (σ', s')) ∈ trans (opnet onp p1)"
  and ttr: "((σ, t), H'-K':arrive(m), (σ', t')) ∈ trans (opnet onp p2)"
from this(1) and ⟨?S (p1 || p2) σ σ' a⟩ have "I σ m" by auto

with sor str have "∀i∈net_tree_ips p1. S (σ i) (σ' i)"
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
moreover from ⟨I σ m⟩ tor ttr have "∀i∈net_tree_ips p2. S (σ i) (σ' i)"
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
ultimately have "∀i. S (σ i) (σ' i)"
  using ⟨∀j. j ∉ net_tree_ips (p1 || p2) → S (σ j) (σ' j)⟩ by auto

with ⟨I σ m⟩ sor str
  have "(σ', s') ∈ oreachable (opnet onp p1) (?S p1) (?U p1)"
  by - (erule(1) oreachable_local, auto)
moreover from ⟨∀i. S (σ i) (σ' i)⟩ ⟨I σ m⟩ tor ttr
  have "(σ', t') ∈ oreachable (opnet onp p2) (?S p2) (?U p2)"
  by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
next
fix R m H K
assume str: "((σ, s), R:*cast(m), (σ', s')) ∈ trans (opnet onp p1)"
  and ttr: "((σ, t), H-K:arrive(m), (σ', t')) ∈ trans (opnet onp p2)"
from sor str have "I σ m"
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
with sor str tor ttr have "∀i. S (σ i) (σ' i)"
  using ⟨∀j. j ∉ net_tree_ips (p1 || p2) → S (σ j) (σ' j)⟩
  by (fastforce dest!: ostep_arrive_invariantD [OF act] ostep_arrive_invariantD [OF act])
with ⟨I σ m⟩ sor str
  have "(σ', s') ∈ oreachable (opnet onp p1) (?S p1) (?U p1)"
  by - (erule(1) oreachable_local, auto)

```



```

moreover from  $\langle \forall i. S(\sigma i)(\sigma' i) \rangle \langle I \sigma m \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
next
fix R m H K
assume str: " $((\sigma, s), H \neg K : \text{arrive}(m), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), R : \text{*cast}(m), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "
from tor ttr have " $I \sigma m$ "
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
with sor str tor ttr have " $\forall i. S(\sigma i)(\sigma' i)$ "
  using  $\langle \forall j. j \notin \text{net\_tree\_ips}(p_1 \parallel p_2) \longrightarrow S(\sigma j)(\sigma' j) \rangle$ 
  by (fastforce dest!: ostep_arrive_invariantD [OF act] ostep_arrive_invariantD [OF act])
with  $\langle I \sigma m \rangle$  sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
moreover from  $\langle \forall i. S(\sigma i)(\sigma' i) \rangle \langle I \sigma m \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
next
fix i i'
assume str: " $((\sigma, s), \text{connect}(i, i'), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), \text{connect}(i, i'), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "
with sor tor have " $\forall i. S(\sigma i)(\sigma' i)$ "
  using  $\langle \forall j. j \notin \text{net\_tree\_ips}(p_1 \parallel p_2) \longrightarrow S(\sigma j)(\sigma' j) \rangle$ 
  by (fastforce dest!: ostep_arrive_invariantD [OF act] ostep_arrive_invariantD [OF act])
with sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
moreover from  $\langle \forall i. S(\sigma i)(\sigma' i) \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
next
fix i i'
assume str: " $((\sigma, s), \text{disconnect}(i, i'), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "
  and ttr: " $((\sigma, t), \text{disconnect}(i, i'), (\sigma', t')) \in \text{trans}(\text{opnet onp } p_2)$ "
with sor tor have " $\forall i. S(\sigma i)(\sigma' i)$ "
  using  $\langle \forall j. j \notin \text{net\_tree\_ips}(p_1 \parallel p_2) \longrightarrow S(\sigma j)(\sigma' j) \rangle$ 
  by (fastforce dest!: ostep_arrive_invariantD [OF act] ostep_arrive_invariantD [OF act])
with sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
moreover from  $\langle \forall i. S(\sigma i)(\sigma' i) \rangle$  tor ttr
  have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
  by - (erule(1) oreachable_local, auto)
ultimately show ?thesis ..
next
fix i d
assume " $t' = t$ "
  and str: " $((\sigma, s), i : \text{deliver}(d), (\sigma', s')) \in \text{trans}(\text{opnet onp } p_1)$ "
from sor str have " $\forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j$ "
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
hence " $\forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow S(\sigma j)(\sigma' j)$ "
  by (auto intro:  $\langle \bigwedge \xi. S \xi \xi \rangle$ )
with sor str
  have " $(\sigma', s') \in \text{oreachable}(\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
  by - (erule(1) oreachable_local, auto)
moreover have " $(\sigma', t') \in \text{oreachable}(\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
proof -
  from  $\langle \forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j \rangle$  and  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$ 
  have " $\forall j. j \in \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j$ " by auto

```

```

    moreover from sor str have " $\forall j \in \text{net\_tree\_ips } p_1. U (\sigma j) (\sigma' j)$ "
      by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
    ultimately show ?thesis
      using tor  $\langle t' = t \rangle \langle \forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j \rangle$ 
      by (clarsimp elim!: oreachable_other')
        (metis otherI  $\langle \bigwedge \xi. U \xi \xi \rangle$ )+
qed
ultimately show ?thesis ..
next
fix i d
assume "s' = s"
  and ttr: " $((\sigma, t), i:\text{deliver}(d), (\sigma', t')) \in \text{trans } (\text{opnet onp } p_2)$ "
from tor ttr have " $\forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j$ "
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
hence " $\forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow S (\sigma j) (\sigma' j)$ "
  by (auto intro:  $\langle \bigwedge \xi. S \xi \xi \rangle$ )
with tor ttr
  have " $(\sigma', t') \in \text{oreachable } (\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
    by - (erule(1) oreachable_local, auto)

moreover have " $(\sigma', s') \in \text{oreachable } (\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
proof -
  from  $\langle \forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j \rangle$  and  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$ 
    have " $\forall j. j \in \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j$ " by auto
  moreover from tor ttr have " $\forall j \in \text{net\_tree\_ips } p_2. U (\sigma j) (\sigma' j)$ "
    by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
  ultimately show ?thesis
    using sor  $\langle s' = s \rangle \langle \forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j \rangle$ 
    by (clarsimp elim!: oreachable_other')
      (metis otherI  $\langle \bigwedge \xi. U \xi \xi \rangle$ )+
qed
ultimately show ?thesis by - (rule conjI)
next
assume "s' = s"
  and ttr: " $((\sigma, t), \tau, (\sigma', t')) \in \text{trans } (\text{opnet onp } p_2)$ "
from tor ttr have " $\forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j$ "
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
hence " $\forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow S (\sigma j) (\sigma' j)$ "
  by (auto intro:  $\langle \bigwedge \xi. S \xi \xi \rangle$ )
with tor ttr
  have " $(\sigma', t') \in \text{oreachable } (\text{opnet onp } p_2) (?S p_2) (?U p_2)$ "
    by - (erule(1) oreachable_local, auto)

moreover have " $(\sigma', s') \in \text{oreachable } (\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "
proof -
  from  $\langle \forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j \rangle$  and  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$ 
    have " $\forall j. j \in \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j$ " by auto
  moreover from tor ttr have " $\forall j \in \text{net\_tree\_ips } p_2. U (\sigma j) (\sigma' j)$ "
    by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
  ultimately show ?thesis
    using sor  $\langle s' = s \rangle \langle \forall j. j \notin \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j \rangle$ 
    by (clarsimp elim!: oreachable_other')
      (metis otherI  $\langle \bigwedge \xi. U \xi \xi \rangle$ )+
qed
ultimately show ?thesis by - (rule conjI)
next
assume "t' = t"
  and str: " $((\sigma, s), \tau, (\sigma', s')) \in \text{trans } (\text{opnet onp } p_1)$ "
from sor str have " $\forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j$ "
  by - (drule(1) ostep_arrive_invariantD [OF act], simp_all)
hence " $\forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow S (\sigma j) (\sigma' j)$ "
  by (auto intro:  $\langle \bigwedge \xi. S \xi \xi \rangle$ )
with sor str
  have " $(\sigma', s') \in \text{oreachable } (\text{opnet onp } p_1) (?S p_1) (?U p_1)$ "

```

by - (erule(1) oreachable\_local, auto)

moreover have " $(\sigma', t') \in \text{oreachable } (\text{opnet } \text{onp } p_2) (?S p_2) (?U p_2)$ "

proof -

from  $\langle \forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j \rangle$  and  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$

have " $\forall j. j \in \text{net\_tree\_ips } p_2 \longrightarrow \sigma' j = \sigma j$ " by auto

moreover from sor str have " $\forall j \in \text{net\_tree\_ips } p_1. U (\sigma j) (\sigma' j)$ "

by - (drule(1) ostep\_arrive\_invariantD [OF act], simp\_all)

ultimately show ?thesis

using tor  $\langle t' = t \rangle \langle \forall j. j \notin \text{net\_tree\_ips } p_1 \longrightarrow \sigma' j = \sigma j \rangle$

by (clarsimp elim!: oreachable\_other')

(metis otherI  $\langle \bigwedge \xi. U \xi \xi \rangle$ )

qed

ultimately show ?thesis ..

qed

with  $\langle \text{net\_tree\_ips } p_1 \cap \text{net\_tree\_ips } p_2 = \{\} \rangle$  show ?case by simp

qed

lemmas subnet\_oreachable1 [dest] = subnet\_oreachable [THEN conjunct1, rotated 1]

lemmas subnet\_oreachable2 [dest] = subnet\_oreachable [THEN conjunct2, THEN conjunct1, rotated 1]

lemmas subnet\_oreachable\_disjoint [dest] = subnet\_oreachable

[THEN conjunct2, THEN conjunct2, rotated 1]

corollary pnet\_lift:

assumes " $\bigwedge ii R_i. \langle ii : \text{onp } ii : R_i \rangle_o$   
 $\models (\text{otherwith } S \{ii\} (\text{oarrivemsg } I), \text{ other } U \{ii\} \rightarrow \text{global } (P ii))$ "

and " $\bigwedge \xi. S \xi \xi$ "

and " $\bigwedge \xi. U \xi \xi$ "

and node1: " $\bigwedge i R. \langle i : \text{onp } i : R \rangle_o \models_A (\lambda \sigma \_ . \text{oarrivemsg } I \sigma, \text{ other } U \{i\} \rightarrow$   
 $\text{globala } (\lambda (\sigma, a, \_). \text{castmsg } (I \sigma) a))$ "

and node2: " $\bigwedge i R. \langle i : \text{onp } i : R \rangle_o \models_A (\lambda \sigma \_ . \text{oarrivemsg } I \sigma, \text{ other } U \{i\} \rightarrow$   
 $\text{globala } (\lambda (\sigma, a, \sigma'). (a \neq \tau \wedge (\forall d. a \neq i:\text{deliver}(d)) \rightarrow S (\sigma i) (\sigma' i)))$ "

and node3: " $\bigwedge i R. \langle i : \text{onp } i : R \rangle_o \models_A (\lambda \sigma \_ . \text{oarrivemsg } I \sigma, \text{ other } U \{i\} \rightarrow$   
 $\text{globala } (\lambda (\sigma, a, \sigma'). (a = \tau \vee (\exists d. a = i:\text{deliver}(d)) \rightarrow U (\sigma i) (\sigma' i)))$ "

shows " $\text{opnet } \text{onp } p \models (\text{otherwith } S (\text{net\_tree\_ips } p) (\text{oarrivemsg } I),$   
 $\text{ other } U (\text{net\_tree\_ips } p) \rightarrow \text{global } (\lambda \sigma. \forall i \in \text{net\_tree\_ips } p. P i \sigma)$ "

(is " $\_ \models (?owS p, ?U p \rightarrow \_)$ "

proof (induction p)

fix ii R<sub>i</sub>

from assms(1) show " $\text{opnet } \text{onp } \langle ii; R_i \rangle \models (?owS \langle ii; R_i \rangle, ?U \langle ii; R_i \rangle \rightarrow$   
 $\text{global } (\lambda \sigma. \forall i \in \text{net\_tree\_ips } \langle ii; R_i \rangle. P i \sigma)$ " by auto

next

fix p<sub>1</sub> p<sub>2</sub>

assume ih1: " $\text{opnet } \text{onp } p_1 \models (?owS p_1, ?U p_1 \rightarrow \text{global } (\lambda \sigma. \forall i \in \text{net\_tree\_ips } p_1. P i \sigma))$ "

and ih2: " $\text{opnet } \text{onp } p_2 \models (?owS p_2, ?U p_2 \rightarrow \text{global } (\lambda \sigma. \forall i \in \text{net\_tree\_ips } p_2. P i \sigma))$ "

show " $\text{opnet } \text{onp } (p_1 \parallel p_2) \models (?owS (p_1 \parallel p_2), ?U (p_1 \parallel p_2) \rightarrow$

$\text{global } (\lambda \sigma. \forall i \in \text{net\_tree\_ips } (p_1 \parallel p_2). P i \sigma)$ "

unfolding oinvariant\_def

proof

fix pq

assume " $pq \in \text{oreachable } (\text{opnet } \text{onp } (p_1 \parallel p_2)) (?owS (p_1 \parallel p_2)) (?U (p_1 \parallel p_2))$ "

moreover then obtain  $\sigma s t$  where " $pq = (\sigma, \text{SubnetS } s t)$ "

by (metis net\_par\_oreachable\_is\_subnet surjective\_pairing)

ultimately have " $(\sigma, \text{SubnetS } s t) \in \text{oreachable } (\text{opnet } \text{onp } (p_1 \parallel p_2))$

$(?owS (p_1 \parallel p_2)) (?U (p_1 \parallel p_2))$ " by simp

then obtain sor: " $(\sigma, s) \in \text{oreachable } (\text{opnet } \text{onp } p_1) (?owS p_1) (?U p_1)$ "

and tor: " $(\sigma, t) \in \text{oreachable } (\text{opnet } \text{onp } p_2) (?owS p_2) (?U p_2)$ "

by - (drule subnet\_oreachable [OF \_ \_ \_ node1 node2 node3], auto intro: assms(2-3))

from sor have " $\forall i \in \text{net\_tree\_ips } p_1. P i \sigma$ "

by (auto dest: oinvariantD [OF ih1])

moreover from tor have " $\forall i \in \text{net\_tree\_ips } p_2. P i \sigma$ "

```

    by (auto dest: oinvariantD [OF ih2])
    ultimately have "∀i∈net_tree_ips (p1 || p2). P i σ" by auto
    with ⟨pq = (σ, SubnetS s t)⟩ show "global (λσ. ∀i∈net_tree_ips (p1 || p2). P i σ) pq" by simp
qed
qed

```

end

## 18 Lifting rules for (open) closed networks

```

theory OClosed_Lifting
imports OPnet_Lifting
begin

```

```

lemma trans_fst_oclosed_fst1 [dest]:

```

```

  "(s, connect(i, i'), s') ∈ ocnet_sos (trans p) ⇒ (s, connect(i, i'), s') ∈ trans p"
  by (metis prod.exhaust oconnect_completeTE)

```

```

lemma trans_fst_oclosed_fst2 [dest]:

```

```

  "(s, disconnect(i, i'), s') ∈ ocnet_sos (trans p) ⇒ (s, disconnect(i, i'), s') ∈ trans p"
  by (metis prod.exhaust odisconnect_completeTE)

```

```

lemma trans_fst_oclosed_fst3 [dest]:

```

```

  "(s, i:deliver(d), s') ∈ ocnet_sos (trans p) ⇒ (s, i:deliver(d), s') ∈ trans p"
  by (metis prod.exhaust odeliver_completeTE)

```

```

lemma oclosed_oreachable_inclosed:

```

```

  assumes "(σ, ζ) ∈ oreachable (oclosed (opnet np p)) (λ_ _ . True) U"
  shows "(σ, ζ) ∈ oreachable (opnet np p) (otherwith ((=)) (net_tree_ips p) inoclosed) U"
  (is "_ ∈ oreachable _ ?owS _")

```

```

  using assms proof (induction rule: oreachable_pair_induct)

```

```

    fix σ ζ

```

```

    assume "(σ, ζ) ∈ init (oclosed (opnet np p))"

```

```

    hence "(σ, ζ) ∈ init (opnet np p)" by simp

```

```

    thus "(σ, ζ) ∈ oreachable (opnet np p) ?owS U" ..

```

```

  next

```

```

    fix σ ζ σ'

```

```

    assume "(σ, ζ) ∈ oreachable (opnet np p) ?owS U"

```

```

    and "U σ σ'"

```

```

    thus "(σ', ζ) ∈ oreachable (opnet np p) ?owS U"

```

```

    by - (rule oreachable_other')

```

```

  next

```

```

    fix σ ζ σ' ζ' a

```

```

    assume zor: "(σ, ζ) ∈ oreachable (opnet np p) ?owS U"

```

```

    and ztr: "((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet np p))"

```

```

    from this(1) have [simp]: "net_ips ζ = net_tree_ips p"

```

```

    by (rule opnet_net_ips_net_tree_ips)

```

```

    from ztr have "((σ, ζ), a, (σ', ζ')) ∈ ocnet_sos (trans (opnet np p))" by simp

```

```

    thus "(σ', ζ') ∈ oreachable (opnet np p) ?owS U"

```

```

  proof cases

```

```

    fix i K d di

```

```

    assume "a = i:newpkt(d, di)"

```

```

    and tr: "((σ, ζ), {i}→K:arrive(msg_class.newpkt (d, di)), (σ', ζ')) ∈ trans (opnet np p)"

```

```

    and "∀j. j ∉ net_ips ζ → σ' j = σ j"

```

```

    from this(3) have "∀j. j ∉ net_tree_ips p → σ' j = σ j"

```

```

    using (net_ips ζ = net_tree_ips p) by auto

```

```

    hence "otherwith ((=)) (net_tree_ips p) inoclosed σ σ' ({i}→K:arrive(msg_class.newpkt (d, di)))"

```

```

    by auto

```

```

    with zor tr show ?thesis

```

```

    by - (rule oreachable_local')

```

```

  next

```

```

    assume "a = τ"

```

```

    and tr: "((σ, ζ), τ, (σ', ζ')) ∈ trans (opnet np p)"

```

```

    and "∀j. j ∉ net_ips ζ → σ' j = σ j"

```

```

from this(3) have "∀j. j ∉ net_tree_ips p → σ' j = σ j"
  using ⟨net_ips ζ = net_tree_ips p⟩ by auto
hence "otherwith ((=)) (net_tree_ips p) inclosed σ σ' τ"
  by auto
with zor tr show ?thesis by - (rule oreachable_local')
qed (insert ⟨net_ips ζ = net_tree_ips p⟩,
      auto elim!: oreachable_local' [OF zor])
qed

```

```

lemma oclosed_oreachable_oreachable [elim]:
  assumes "(σ, ζ) ∈ oreachable (oclosed (opnet onp p)) (λ_ _ . True) U"
  shows "(σ, ζ) ∈ oreachable (opnet onp p) (λ_ _ . True) U"
  using assms by (rule oclosed_oreachable_inclosed [THEN oreachable_weakenE]) simp

```

```

lemma inclosed_closed [intro]:
  assumes cinv: "opnet np p ⊨ (otherwith ((=)) (net_tree_ips p) inclosed, U →) P"
  shows "oclosed (opnet np p) ⊨ (λ_ _ . True, U →) P"
  using assms unfolding oinvariant_def
  by (clarsimp dest!: oclosed_oreachable_inclosed)

```

end

## 19 Generic invariants on sequential AWN processes

```

theory AWN_Invariants
imports Invariants AWN_SOS AWN_Labels
begin

```

### 19.1 Invariants via labelled control terms

Used to state that the initial control-state of an automaton appears within a process specification  $\Gamma$ , meaning that its transitions, and those of its subterms, are subsumed by those of  $\Gamma$ .

**definition**

```

control_within :: "('s, 'm, 'p, 'l) seqp_env ⇒ ('z × ('s, 'm, 'p, 'l) seqp) set ⇒ bool"
where
  "control_within Γ σ ≡ ∀(ξ, p) ∈ σ. ∃pn. p ∈ subterms (Γ pn)"

```

```

lemma control_withinI [intro]:
  assumes "∧p. p ∈ Range σ ⇒ ∃pn. p ∈ subterms (Γ pn)"
  shows "control_within Γ σ"
  using assms unfolding control_within_def by auto

```

```

lemma control_withinD [dest]:
  assumes "control_within Γ σ"
  and "(ξ, p) ∈ σ"
  shows "∃pn. p ∈ subterms (Γ pn)"
  using assms unfolding control_within_def by blast

```

```

lemma control_within_topI [intro]:
  assumes "∧p. p ∈ Range σ ⇒ ∃pn. p = Γ pn"
  shows "control_within Γ σ"
  using assms unfolding control_within_def
  byclarsimp (metis Range.RangeI subterms_refl)

```

```

lemma seqp_sos_subterms:
  assumes "wellformed Γ"
  and "∃pn. p ∈ subterms (Γ pn)"
  and "((ξ, p), a, (ξ', p')) ∈ seqp_sos Γ"
  shows "∃pn. p' ∈ subterms (Γ pn)"
  using assms
  proof (induct p)
    fix p1 p2
    assume IH1: "∃pn. p1 ∈ subterms (Γ pn) ⇒

```

```

      ((ξ, p1), a, (ξ', p')) ∈ seqp_sos Γ ⇒
      ∃pn. p' ∈ subterms (Γ pn)"
and IH2: "∃pn. p2 ∈ subterms (Γ pn) ⇒
      ((ξ, p2), a, (ξ', p')) ∈ seqp_sos Γ ⇒
      ∃pn. p' ∈ subterms (Γ pn)"
and "∃pn. p1 ⊕ p2 ∈ subterms (Γ pn)"
and "((ξ, p1 ⊕ p2), a, (ξ', p')) ∈ seqp_sos Γ"
from (∃pn. p1 ⊕ p2 ∈ subterms (Γ pn)) obtain pn
      where "p1 ∈ subterms (Γ pn)"
      and "p2 ∈ subterms (Γ pn)" by auto
from ((ξ, p1 ⊕ p2), a, (ξ', p')) ∈ seqp_sos Γ
  have "((ξ, p1), a, (ξ', p')) ∈ seqp_sos Γ
    ∨ ((ξ, p2), a, (ξ', p')) ∈ seqp_sos Γ" by auto
thus "∃pn. p' ∈ subterms (Γ pn)"
proof
  assume "((ξ, p1), a, (ξ', p')) ∈ seqp_sos Γ"
  with ⟨p1 ∈ subterms (Γ pn)⟩ show ?thesis by (auto intro: IH1)
next
  assume "((ξ, p2), a, (ξ', p')) ∈ seqp_sos Γ"
  with ⟨p2 ∈ subterms (Γ pn)⟩ show ?thesis by (auto intro: IH2)
qed
qed auto

```

lemma reachable\_subterms:

```

assumes "wellformed Γ"
  and "control_within Γ (init A)"
  and "trans A = seqp_sos Γ"
  and "(ξ, p) ∈ reachable A I"
shows "∃pn. p ∈ subterms (Γ pn)"
using assms(4)
proof (induct rule: reachable_pair_induct)
  fix ξ p
  assume "(ξ, p) ∈ init A"
  with ⟨control_within Γ (init A)⟩ show "∃pn. p ∈ subterms (Γ pn)" ..
next
  fix ξ p a ξ' p'
  assume "(ξ, p) ∈ reachable A I"
    and "∃pn. p ∈ subterms (Γ pn)"
    and *: "((ξ, p), a, (ξ', p')) ∈ trans A"
    and "I a"
  moreover from * and assms(3) have "((ξ, p), a, (ξ', p')) ∈ seqp_sos Γ" by simp
  ultimately show "∃pn. p' ∈ subterms (Γ pn)"
  using ⟨wellformed Γ⟩
    by (auto elim: seqp_sos_subterms)
qed

```

definition

```

onl :: "('s, 'm, 'p, 'l) seqp_env
      ⇒ ('z × 'l ⇒ bool)
      ⇒ 'z × ('s, 'm, 'p, 'l) seqp
      ⇒ bool"

```

where

```

"onl Γ P ≡ (λ(ξ, p). ∀l∈labels Γ p. P (ξ, l))"

```

lemma onlI [intro]:

```

assumes "∧l. l∈labels Γ p ⇒ P (ξ, l)"
shows "onl Γ P (ξ, p)"
using assms unfolding onl_def by simp

```

lemmas onlI' [intro] = onlI [simplified atomize\_ball]

lemma onlD [dest]:

```

assumes "onl Γ P (ξ, p)"
shows "∀l∈labels Γ p. P (ξ, l)"

```

```

using assms unfolding onl_def by simp

lemma onl_invariantI [intro]:
  assumes init: " $\bigwedge \xi p l. \llbracket (\xi, p) \in \text{init } A; l \in \text{labels } \Gamma p \rrbracket \implies P (\xi, l)$ "
    and step: " $\bigwedge \xi p a \xi' p' l'. \llbracket (\xi, p) \in \text{reachable } A I; \forall l \in \text{labels } \Gamma p. P (\xi, l); ((\xi, p), a, (\xi', p')) \in \text{trans } A; l' \in \text{labels } \Gamma p'; I a \rrbracket \implies P (\xi', l')$ "
  shows "A  $\models (I \rightarrow) \text{onl } \Gamma P$ "
proof (rule invariant_pairI)
  fix  $\xi p$ 
  assume " $(\xi, p) \in \text{init } A$ "
  hence " $\forall l \in \text{labels } \Gamma p. P (\xi, l)$ " using init by simp
  thus "onl  $\Gamma P (\xi, p)$ " ..
next
  fix  $\xi p a \xi' p'$ 
  assume rp: " $(\xi, p) \in \text{reachable } A I$ "
    and "onl  $\Gamma P (\xi, p)$ "
    and tr: " $((\xi, p), a, (\xi', p')) \in \text{trans } A$ "
    and "I a"
  from (onl  $\Gamma P (\xi, p)$ ) have " $\forall l \in \text{labels } \Gamma p. P (\xi, l)$ " ..
  with rp tr (I a) have " $\forall l' \in \text{labels } \Gamma p'. P (\xi', l')$ " by (auto elim: step)
  thus "onl  $\Gamma P (\xi', p')$ " ..
qed

lemma onl_invariantD [dest]:
  assumes "A  $\models (I \rightarrow) \text{onl } \Gamma P$ "
    and " $(\xi, p) \in \text{reachable } A I$ "
    and " $l \in \text{labels } \Gamma p$ "
  shows "P ( $\xi, l$ )"
using assms unfolding onl_def by auto

lemma onl_invariant_initD [dest]:
  assumes invP: "A  $\models (I \rightarrow) \text{onl } \Gamma P$ "
    and init: " $(\xi, p) \in \text{init } A$ "
    and pnl: " $l \in \text{labels } \Gamma p$ "
  shows "P ( $\xi, l$ )"
proof -
  from init have " $(\xi, p) \in \text{reachable } A I$ " ..
  with invP show ?thesis using pnl ..
qed

lemma onl_invariant_sterms:
  assumes wf: "wellformed  $\Gamma$ "
    and il: "A  $\models (I \rightarrow) \text{onl } \Gamma P$ "
    and rp: " $(\xi, p) \in \text{reachable } A I$ "
    and "p'  $\in \text{sterms } \Gamma p$ "
    and "l  $\in \text{labels } \Gamma p'$ "
  shows "P ( $\xi, l$ )"
proof -
  from wf (p'  $\in \text{sterms } \Gamma p$ ) (l  $\in \text{labels } \Gamma p'$ ) have "l  $\in \text{labels } \Gamma p$ "
    by (rule labels_sterms_labels)
  with il rp show "P ( $\xi, l$ )" ..
qed

lemma onl_invariant_sterms_weaken:
  assumes wf: "wellformed  $\Gamma$ "
    and il: "A  $\models (I \rightarrow) \text{onl } \Gamma P$ "
    and rp: " $(\xi, p) \in \text{reachable } A I$ "
    and "p'  $\in \text{sterms } \Gamma p$ "
    and "l  $\in \text{labels } \Gamma p'$ "
    and weaken: " $\bigwedge a. I' a \implies I a$ "

```

```

shows "P (ξ, l)"
proof -
  from ⟨(ξ, p) ∈ reachable A I'⟩ have "(ξ, p) ∈ reachable A I"
    by (rule reachable_weakenE)
      (erule weaken)
  with assms(1-2) show ?thesis using assms(4-5) by (rule onl_invariant_sterms)
qed

lemma onl_invariant_sterms_TT:
  assumes wf: "wellformed Γ"
    and il: "A ⊨ onl Γ P"
    and rp: "(ξ, p) ∈ reachable A I"
    and "p' ∈ sterms Γ p"
    and "l ∈ labels Γ p'"
  shows "P (ξ, l)"
  using assms by (rule onl_invariant_sterms_weaken) simp

lemma trans_from_sterms:
  assumes "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
    and "wellformed Γ"
  shows "∃ p' ∈ sterms Γ p. ((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ"
  using assms by (induction p rule: sterms_pinduct [OF ⟨wellformed Γ⟩]) auto

lemma trans_from_sterms':
  assumes "((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ"
    and "wellformed Γ"
    and "p' ∈ sterms Γ p"
  shows "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  using assms by (induction p rule: sterms_pinduct [OF ⟨wellformed Γ⟩]) auto

lemma trans_to_dterms:
  assumes "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
    and "wellformed Γ"
  shows "∀ r ∈ sterms Γ q. r ∈ dterms Γ p"
  using assms by (induction q) auto

theorem cterms_includes_sterms_of_seq_reachable:
  assumes "wellformed Γ"
    and "control_within Γ (init A)"
    and "trans A = seqp_sos Γ"
  shows "⋃ (sterms Γ ' snd ' reachable A I) ⊆ cterms Γ"
proof
  fix qs
  assume "qs ∈ ⋃ (sterms Γ ' snd ' reachable A I)"
  then obtain ξ and q where *: "(ξ, q) ∈ reachable A I"
    and **: "qs ∈ sterms Γ q" by auto
  from * have "∧ x. x ∈ sterms Γ q ⇒ x ∈ cterms Γ"
  proof (induction rule: reachable_pair_induct)
    fix ξ p q
    assume "(ξ, p) ∈ init A"
      and "q ∈ sterms Γ p"
    from ⟨control_within Γ (init A)⟩ and ⟨(ξ, p) ∈ init A⟩
      obtain pn where "p ∈ subterms (Γ pn)" by auto
    with ⟨wellformed Γ⟩ show "q ∈ cterms Γ" using ⟨q ∈ sterms Γ p⟩
      by (rule subterms_sterms_in_cterms)
  next
    fix p ξ a ξ' q x
    assume "(ξ, p) ∈ reachable A I"
      and IH: "∧ x. x ∈ sterms Γ p ⇒ x ∈ cterms Γ"
      and "((ξ, p), a, (ξ', q)) ∈ trans A"
      and "x ∈ sterms Γ q"
    from this(3) and ⟨trans A = seqp_sos Γ⟩ have "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ" by simp
    from this and ⟨wellformed Γ⟩ obtain ps
      where ps: "ps ∈ sterms Γ p"

```



```

    and step: "((ξ, ps), a, (ξ', q)) ∈ seqp_sos Γ"
  by (rule trans_from_sterms [THEN bexE])
from ps have "ps ∈ cterms Γ" by (rule IH)
moreover from step ⟨wellformed Γ⟩ ⟨x ∈ sterms Γ q⟩ have "x ∈ dterms Γ ps"
  by (rule trans_to_dterms [rule_format])
ultimately show "x ∈ cterms Γ" by (rule ctermsDI)
qed
thus "qs ∈ cterms Γ" using **.
qed

```

corollary seq\_reachable\_in\_cterms:

```

assumes "wellformed Γ"
  and "control_within Γ (init A)"
  and "trans A = seqp_sos Γ"
  and "(ξ, p) ∈ reachable A I"
  and "p' ∈ sterms Γ p"
shows "p' ∈ cterms Γ"
using assms(1-3)
proof (rule cterms_includes_sterms_of_seq_reachable [THEN set_mp])
  from assms(4-5) show "p' ∈ ⋃ (sterms Γ ' snd ' reachable A I)"
    by (auto elim!: rev_bexI)
qed

```

lemma seq\_invariant\_ctermI:

```

assumes wf: "wellformed Γ"
  and cw: "control_within Γ (init A)"
  and sl: "simple_labels Γ"
  and sp: "trans A = seqp_sos Γ"
  and init: "⋀ξ p l. [
    (ξ, p) ∈ init A;
    l ∈ labels Γ p
  ] ⇒ P (ξ, l)"
  and step: "⋀p l ξ a q l' ξ' pp. [
    p ∈ cterms Γ;
    l ∈ labels Γ p;
    P (ξ, l);
    ((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ;
    ((ξ, p), a, (ξ', q)) ∈ trans A;
    l' ∈ labels Γ q;
    (ξ, pp) ∈ reachable A I;
    p ∈ sterms Γ pp;
    (ξ', q) ∈ reachable A I;
    I a
  ] ⇒ P (ξ', l'"
shows "A ⊨ (I →) onl Γ P"

```

proof

```

  fix ξ p l
  assume "(ξ, p) ∈ init A"
  and *: "l ∈ labels Γ p"
  with init show "P (ξ, l)" by auto

```

next

```

  fix ξ p a ξ' q l'
  assume sr: "(ξ, p) ∈ reachable A I"
  and pl: "⋀l ∈ labels Γ p. P (ξ, l)"
  and tr: "((ξ, p), a, (ξ', q)) ∈ trans A"
  and A6: "l' ∈ labels Γ q"
  and "I a"
  from this(3) and ⟨trans A = seqp_sos Γ⟩ have tr': "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ" by simp
  show "P (ξ', l'"

```

proof -

```

  from sr and tr and ⟨I a⟩ have A7: "(ξ', q) ∈ reachable A I" ..
  from tr' obtain p' where "p' ∈ sterms Γ p"
    and "((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ"
  by (blast dest: trans_from_sterms [OF _ wf])

```

```

from wf cw sp sr this(1) have A1: "p' ∈ cterms Γ"
  by (rule seq_reachable_in_cterms)
from labels_not_empty [OF wf] obtain l1 where A2: "l1 ∈ labels Γ p'"
  by blast
with ⟨p' ∈ sterms Γ p⟩ have "l1 ∈ labels Γ p"
  by (rule labels_sterms_labels [OF wf])
with p1 have A3: "P (ξ, l1)" by simp
from ⟨((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ⟩ and sp
  have A5: "((ξ, p'), a, (ξ', q)) ∈ trans A" by simp
with sp have A4: "((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ" by simp
from sr ⟨p' ∈ sterms Γ p⟩
  obtain pp where A7: "(ξ, pp) ∈ reachable A I"
    and A8: "p' ∈ sterms Γ pp"

  by auto
from sr tr ⟨I a⟩ have A9: "(ξ', q) ∈ reachable A I" ..
from A1 A2 A3 A4 A5 A6 A7 A8 A9 ⟨I a⟩ show ?thesis by (rule step)
qed
qed

```

lemma seq\_invariant\_ctermsI:

```

assumes wf: "wellformed Γ"
  and "control_within Γ (init A)"
  and "simple_labels Γ"
  and "trans A = seqp_sos Γ"
  and init: "∧ξ p l. [
    (ξ, p) ∈ init A;
    l ∈ labels Γ p
  ] ⇒ P (ξ, l)"
  and step: "∧p l ξ a q l' ξ' pp pn. [
    wellformed Γ;
    p ∈ ctermsl (Γ pn);
    not_call p;
    l ∈ labels Γ p;
    P (ξ, l);
    ((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ;
    ((ξ, p), a, (ξ', q)) ∈ trans A;
    l' ∈ labels Γ q;
    (ξ, pp) ∈ reachable A I;
    p ∈ sterms Γ pp;
    (ξ', q) ∈ reachable A I;
    I a
  ] ⇒ P (ξ', l'"
shows "A ⊨ (I →) onl Γ P"

```

using assms(1-4) proof (rule seq\_invariant\_ctermI)

```

fix ξ p l
assume "(ξ, p) ∈ init A"
  and "l ∈ labels Γ p"
thus "P (ξ, l)" by (rule init)

```

next

```

fix p l ξ a q l' ξ' pp
assume "p ∈ cterms Γ"
  and otherassms: "l ∈ labels Γ p"
  "P (ξ, l)"
  "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  "((ξ, p), a, (ξ', q)) ∈ trans A"
  "l' ∈ labels Γ q"
  "(ξ, pp) ∈ reachable A I"
  "p ∈ sterms Γ pp"
  "(ξ', q) ∈ reachable A I"
  "I a"
from this(1) obtain pn where "p ∈ ctermsl (Γ pn)"
  and "not_call p"
  unfolding cterms_def' [OF wf] by auto
with wf show "P (ξ', l'"

```

using otherassms by (rule step)  
qed

## 19.2 Step invariants via labelled control terms

definition

```
onll :: "('s, 'm, 'p, 'l) seqp_env
      => (('z × 'l, 'a) transition => bool)
      => ('z × ('s, 'm, 'p, 'l) seqp, 'a) transition => bool"
```

where

```
"onll Γ P ≡ (λ((ξ, p), a, (ξ', p')). ∀l∈labels Γ p. ∀l'∈labels Γ p'. P ((ξ, l), a, (ξ', l')))"
```

lemma onllI [intro]:

```
assumes "∧l l'. [ l∈labels Γ p; l'∈labels Γ p' ] => P ((ξ, l), a, (ξ', l'))"
shows "onll Γ P ((ξ, p), a, (ξ', p'))"
using assms unfolding onll_def by simp
```

lemma onllII [intro]:

```
assumes "∀l∈labels Γ p. ∀l'∈labels Γ p'. P ((ξ, l), a, (ξ', l'))"
shows "onll Γ P ((ξ, p), a, (ξ', p'))"
using assms by auto
```

lemma onllD [dest]:

```
assumes "onll Γ P ((ξ, p), a, (ξ', p'))"
shows "∀l∈labels Γ p. ∀l'∈labels Γ p'. P ((ξ, l), a, (ξ', l'))"
using assms unfolding onll_def by simp
```

lemma onl\_weaken [elim!]: "∧Γ P Q s. [ onl Γ P s; ∧s. P s => Q s ] => onl Γ Q s"  
by (clarsimp dest!: onlD intro!: onlI)

lemma onll\_weaken [elim!]: "∧Γ P Q s. [ onll Γ P s; ∧s. P s => Q s ] => onll Γ Q s"  
by (clarsimp dest!: onllD intro!: onllI)

lemma onll\_weaken' [elim!]: "∧Γ P Q s. [ onll Γ P ((ξ, p), a, (ξ', p'));  
∧l l'. P ((ξ, l), a, (ξ', l')) => Q ((ξ, l), a, (ξ', l')) ]  
=> onll Γ Q ((ξ, p), a, (ξ', p'))"  
by (clarsimp dest!: onllD intro!: onllI)

lemma onll\_step\_invariantI [intro]:

```
assumes *: "∧ξ p l a ξ' p' l'. [ (ξ, p)∈reachable A I;  
((ξ, p), a, (ξ', p')) ∈ trans A;  
I a;  
l ∈labels Γ p;  
l'∈labels Γ p' ]  
=> P ((ξ, l), a, (ξ', l'))"
```

shows "A ⊨<sub>A</sub> (I →) onll Γ P"

proof

fix ξ p ξ' p' a

assume "(ξ, p) ∈ reachable A I"

and "((ξ, p), a, (ξ', p')) ∈ trans A"

and "I a"

hence "∀l∈labels Γ p. ∀l'∈labels Γ p'. P ((ξ, l), a, (ξ', l'))" by (auto elim!: \*)

thus "onll Γ P ((ξ, p), a, (ξ', p'))" ..

qed

lemma onll\_step\_invariantE [elim]:

assumes "A ⊨<sub>A</sub> (I →) onll Γ P"

and "(ξ, p) ∈ reachable A I"

and "((ξ, p), a, (ξ', p')) ∈ trans A"

and "I a"

and lp: "l ∈labels Γ p"

and lp': "l'∈labels Γ p'"

shows "P ((ξ, l), a, (ξ', l'))"

proof -

```

    from assms(1-4) have "onll  $\Gamma$  P (( $\xi$ , p), a, ( $\xi'$ , p'))" ..
    with lp lp' show "P (( $\xi$ , l), a, ( $\xi'$ , l'))" by auto
qed

lemma onll_step_invariantD [dest]:
  assumes "A  $\models_A$  (I  $\rightarrow$ ) onll  $\Gamma$  P"
    and "( $\xi$ , p)  $\in$  reachable A I"
    and "(( $\xi$ , p), a, ( $\xi'$ , p'))  $\in$  trans A"
    and "I a"
  shows " $\forall l \in \text{labels } \Gamma p. \forall l' \in \text{labels } \Gamma p'. P ((\xi, l), a, (\xi', l'))$ "
  using assms by auto

lemma onll_step_to_invariantI [intro]:
  assumes sinv: "A  $\models_A$  (I  $\rightarrow$ ) onll  $\Gamma$  Q"
    and wf: "wellformed  $\Gamma$ "
    and init: " $\bigwedge \xi l p. [(\xi, p) \in \text{init } A; l \in \text{labels } \Gamma p] \implies P (\xi, l)$ "
    and step: " $\bigwedge \xi p l \xi' l' a. [(\xi, p) \in \text{reachable } A I; l \in \text{labels } \Gamma p; P (\xi, l); Q ((\xi, l), a, (\xi', l')); I a] \implies P (\xi', l')$ "
  shows "A  $\models$  (I  $\rightarrow$ ) onl  $\Gamma$  P"
proof
  fix  $\xi p l$ 
  assume "( $\xi$ , p)  $\in$  init A" and "l  $\in$  labels  $\Gamma$  p"
  thus "P ( $\xi$ , l)" by (rule init)
next
  fix  $\xi p a \xi' p' l'$ 
  assume sr: "( $\xi$ , p)  $\in$  reachable A I"
    and lp: " $\forall l \in \text{labels } \Gamma p. P (\xi, l)$ "
    and tr: "(( $\xi$ , p), a, ( $\xi'$ , p'))  $\in$  trans A"
    and "I a"
    and lp': "l'  $\in$  labels  $\Gamma$  p'"
  show "P ( $\xi'$ , l')"
proof -
  from lp obtain l where "l  $\in$  labels  $\Gamma$  p" and "P ( $\xi$ , l)"
  using labels_not_empty [OF wf] by auto
  from sinv sr tr (I a) this(1) lp' have "Q (( $\xi$ , l), a, ( $\xi'$ , l'))" ..
  with sr (l  $\in$  labels  $\Gamma$  p) (P ( $\xi$ , l)) show "P ( $\xi'$ , l')" using (I a) by (rule step)
qed
qed

lemma onll_step_invariant_sterms:
  assumes wf: "wellformed  $\Gamma$ "
    and si: "A  $\models_A$  (I  $\rightarrow$ ) onll  $\Gamma$  P"
    and sr: "( $\xi$ , p)  $\in$  reachable A I"
    and sos: "(( $\xi$ , p), a, ( $\xi'$ , q))  $\in$  trans A"
    and "I a"
    and "l'  $\in$  labels  $\Gamma$  q"
    and "p'  $\in$  sterms  $\Gamma$  p"
    and "l  $\in$  labels  $\Gamma$  p'"
  shows "P (( $\xi$ , l), a, ( $\xi'$ , l'))"
proof -
  from wf (p'  $\in$  sterms  $\Gamma$  p) (l  $\in$  labels  $\Gamma$  p') have "l  $\in$  labels  $\Gamma$  p"
  by (rule labels_sterms_labels)
  with si sr sos (I a) show "P (( $\xi$ , l), a, ( $\xi'$ , l'))" using (l'  $\in$  labels  $\Gamma$  q) ..
qed

lemma seq_step_invariant_sterms:
  assumes inv: "A  $\models_A$  (I  $\rightarrow$ ) onll  $\Gamma$  P"
    and wf: "wellformed  $\Gamma$ "
    and sp: "trans A = seqp_sos  $\Gamma$ "
    and "l'  $\in$  labels  $\Gamma$  q"

```

```

and sr: "(ξ, p) ∈ reachable A I"
and tr: "((ξ, p'), a, (ξ', q)) ∈ trans A"
and "I a"
and "p' ∈ sterms Γ p"
shows "∀ l ∈ labels Γ p'. P ((ξ, l), a, (ξ', l'))"
proof
from tr and sp have "((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ" by simp
hence "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
using wf ⟨p' ∈ sterms Γ p⟩ by (rule trans_from_sterms')
with sp have trp: "((ξ, p), a, (ξ', q)) ∈ trans A" by simp
fix l assume "l ∈ labels Γ p'"
with wf inv sr trp ⟨I a⟩ ⟨l' ∈ labels Γ q⟩ ⟨p' ∈ sterms Γ p⟩
show "P ((ξ, l), a, (ξ', l'))" by (rule onll_step_invariant_sterms)
qed

```

lemma seq\_step\_invariant\_sterms\_weaken:

```

assumes "A ⊨A (I →) onll Γ P"
and "wellformed Γ"
and "trans A = seqp_sos Γ"
and "l' ∈ labels Γ q"
and "(ξ, p) ∈ reachable A I'"
and "((ξ, p'), a, (ξ', q)) ∈ trans A"
and "I' a"
and "p' ∈ sterms Γ p"
and weaken: "∧ a. I' a ⇒ I a"
shows "∀ l ∈ labels Γ p'. P ((ξ, l), a, (ξ', l'))"

```

proof -

```

from ⟨I' a⟩ have "I a" by (rule weaken)
from ⟨(ξ, p) ∈ reachable A I'⟩ have Ir: "(ξ, p) ∈ reachable A I"
by (rule reachable_weakenE) (erule weaken)
with assms(1-4) show ?thesis
using ⟨((ξ, p'), a, (ξ', q)) ∈ trans A⟩ ⟨I a⟩ and ⟨p' ∈ sterms Γ p⟩
by (rule seq_step_invariant_sterms)
qed

```

lemma seq\_step\_invariant\_sterms\_TT:

```

assumes "A ⊨A onll Γ P"
and "wellformed Γ"
and "trans A = seqp_sos Γ"
and "l' ∈ labels Γ q"
and "(ξ, p) ∈ reachable A I"
and "((ξ, p'), a, (ξ', q)) ∈ trans A"
and "I a"
and "p' ∈ sterms Γ p"
shows "∀ l ∈ labels Γ p'. P ((ξ, l), a, (ξ', l'))"
using assms by (rule seq_step_invariant_sterms_weaken) simp

```

lemma onll\_step\_invariant\_any\_sterms:

```

assumes "wellformed Γ"
and "A ⊨A (I →) onll Γ P"
and "(ξ, p) ∈ reachable A I"
and "((ξ, p), a, (ξ', q)) ∈ trans A"
and "I a"
and "l' ∈ labels Γ q"
shows "∀ p' ∈ sterms Γ p. ∀ l ∈ labels Γ p'. P ((ξ, l), a, (ξ', l'))"
by (intro ballI) (rule onll_step_invariant_sterms [OF assms])

```

lemma seq\_step\_invariant\_ctermI [intro]:

```

assumes wf: "wellformed Γ"
and cw: "control_within Γ (init A)"
and sl: "simple_labels Γ"
and sp: "trans A = seqp_sos Γ"
and step: "∧ p pp l ξ a q l' ξ'. [
p ∈ cterms Γ;

```

```

    l ∈ labels Γ p;
    ((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ;
    ((ξ, p), a, (ξ', q)) ∈ trans A;
    l' ∈ labels Γ q;
    (ξ, pp) ∈ reachable A I;
    p ∈ sterms Γ pp;
    (ξ', q) ∈ reachable A I;
    I a
  ] ⇒ P ((ξ, l), a, (ξ', l'))"
shows "A ⊨A (I →) onll Γ P"
proof
  fix ξ p l a ξ' q l'
  assume sr: "(ξ, p) ∈ reachable A I"
  and tr: "((ξ, p), a, (ξ', q)) ∈ trans A"
  and "I a"
  and pl: "l ∈ labels Γ p"
  and A5: "l' ∈ labels Γ q"
  from this(2) and sp have tr': "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ" by simp
  then obtain p' where "p' ∈ sterms Γ p"
    and A3: "((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ"
  by (blast dest: trans_from_sterms [OF _ wf])
  from wf cw sp sr this(1) have A1: "p' ∈ cterms Γ"
  by (rule seq_reachable_in_cterms)
  from ((ξ, p'), a, (ξ', q)) ∈ seqp_sos Γ and sp
  have A4: "((ξ, p'), a, (ξ', q)) ∈ trans A" by simp
  from sr (p' ∈ sterms Γ p) obtain pp where A6: "(ξ, pp) ∈ reachable A I"
    and A7: "p' ∈ sterms Γ pp"

  by auto
  from sr tr (I a) have A8: "(ξ', q) ∈ reachable A I" ..
  from wf cw sp sr have "∃pn. p ∈ subterms (Γ pn)"
  by (rule reachable_subterms)
  with sl wf have "∀p' ∈ sterms Γ p. l ∈ labels Γ p'"
  using pl by (rule simple_labels_in_sterms)
  with (p' ∈ sterms Γ p) have "l ∈ labels Γ p'" by simp
  with A1 show "P ((ξ, l), a, (ξ', l'))" using A3 A4 A5 A6 A7 A8 (I a)
  by (rule step)
qed

```

```

lemma seq_step_invariant_ctermsI [intro]:
  assumes wf: "wellformed Γ"
  and cw: "control_within Γ (init A)"
  and sl: "simple_labels Γ"
  and sp: "trans A = seqp_sos Γ"
  and step: "∧p l ξ a q l' ξ' pp pn. [
    wellformed Γ;
    p ∈ ctermssl (Γ pn);
    not_call p;
    l ∈ labels Γ p;
    ((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ;
    ((ξ, p), a, (ξ', q)) ∈ trans A;
    l' ∈ labels Γ q;
    (ξ, pp) ∈ reachable A I;
    p ∈ sterms Γ pp;
    (ξ', q) ∈ reachable A I;
    I a
  ] ⇒ P ((ξ, l), a, (ξ', l'))"
shows "A ⊨A (I →) onll Γ P"
using assms(1-4) proof (rule seq_step_invariant_ctermI)
  fix p pp l ξ a q l' ξ'
  assume "p ∈ cterms Γ"
  and otherassms: "l ∈ labels Γ p"
  "((ξ, p), a, (ξ', q)) ∈ seqp_sos Γ"
  "((ξ, p), a, (ξ', q)) ∈ trans A"
  "l' ∈ labels Γ q"

```

```

      "(\xi, pp) \in reachable A I"
      "p \in sterms \Gamma pp"
      "(\xi', q) \in reachable A I"
      "I a"
    from this(1) obtain pn where "p \in cterms1(\Gamma pn)"
      and "not_call p"
      unfolding cterms_def' [OF wf] by auto
    with wf show "P ((\xi, l), a, (\xi', l'))"
      using otherassms by (rule step)
  qed

end

```

## 20 Generic open invariants on sequential AWN processes

```

theory OAWN_Invariants
imports Invariants OInvariants
        Awn_Cterms Awn_Labels Awn_Invariants
        OAWN_SOS

begin

```

### 20.1 Open invariants via labelled control terms

```

lemma oseqp_sos_subterms:
  assumes "wellformed \Gamma"
    and "\exists pn. p \in subterms (\Gamma pn)"
    and "((\sigma, p), a, (\sigma', p')) \in oseqp_sos \Gamma i"
  shows "\exists pn. p' \in subterms (\Gamma pn)"
  using assms
  proof (induct p)
    fix p1 p2
    assume IH1: "\exists pn. p1 \in subterms (\Gamma pn) \implies
      ((\sigma, p1), a, (\sigma', p')) \in oseqp_sos \Gamma i \implies
      \exists pn. p' \in subterms (\Gamma pn)"
      and IH2: "\exists pn. p2 \in subterms (\Gamma pn) \implies
      ((\sigma, p2), a, (\sigma', p')) \in oseqp_sos \Gamma i \implies
      \exists pn. p' \in subterms (\Gamma pn)"
      and "\exists pn. p1 \oplus p2 \in subterms (\Gamma pn)"
      and "((\sigma, p1 \oplus p2), a, (\sigma', p')) \in oseqp_sos \Gamma i"
    from (\exists pn. p1 \oplus p2 \in subterms (\Gamma pn)) obtain pn
      where "p1 \in subterms (\Gamma pn)"
      and "p2 \in subterms (\Gamma pn)" by auto
    from ((\sigma, p1 \oplus p2), a, (\sigma', p')) \in oseqp_sos \Gamma i
      have "((\sigma, p1), a, (\sigma', p')) \in oseqp_sos \Gamma i
        \vee ((\sigma, p2), a, (\sigma', p')) \in oseqp_sos \Gamma i" by auto
    thus "\exists pn. p' \in subterms (\Gamma pn)"
  proof
    assume "((\sigma, p1), a, (\sigma', p')) \in oseqp_sos \Gamma i"
    with (p1 \in subterms (\Gamma pn)) show ?thesis by (auto intro: IH1)
  next
    assume "((\sigma, p2), a, (\sigma', p')) \in oseqp_sos \Gamma i"
    with (p2 \in subterms (\Gamma pn)) show ?thesis by (auto intro: IH2)
  qed
  qed auto

```

```

lemma oreachable_subterms:
  assumes "wellformed \Gamma"
    and "control_within \Gamma (init A)"
    and "trans A = oseqp_sos \Gamma i"
    and "(\sigma, p) \in oreachable A S U"
  shows "\exists pn. p \in subterms (\Gamma pn)"
  using assms(4)
  proof (induct rule: oreachable_pair_induct)
    fix \sigma p

```

```

assume "(σ, p) ∈ init A"
with ⟨control_within Γ (init A)⟩ show "∃pn. p ∈ subterms (Γ pn)" ..
next
fix σ p a σ' p'
assume "(σ, p) ∈ oreachable A S U"
and "∃pn. p ∈ subterms (Γ pn)"
and 3: "((σ, p), a, (σ', p')) ∈ trans A"
and "S σ σ' a"
moreover from 3 and ⟨trans A = oseqp_sos Γ i⟩
have "((σ, p), a, (σ', p')) ∈ oseqp_sos Γ i" by simp
ultimately show "∃pn. p' ∈ subterms (Γ pn)"
using ⟨wellformed Γ⟩
by (auto elim: oseqp_sos_subterms)
qed

```

lemma onl\_oinvariantI [intro]:

```

assumes init: "∧σ p l. [ (σ, p) ∈ init A; l ∈ labels Γ p ] ⇒ P (σ, l)"
and other: "∧σ σ' p l. [ (σ, p) ∈ oreachable A S U;
                        ∀l∈labels Γ p. P (σ, l);
                        U σ σ' ] ⇒ ∀l∈labels Γ p. P (σ', l)"
and step: "∧σ p a σ' p' l'.
           [ (σ, p) ∈ oreachable A S U;
             ∀l∈labels Γ p. P (σ, l);
             ((σ, p), a, (σ', p')) ∈ trans A;
             l' ∈ labels Γ p';
             S σ σ' a ] ⇒ P (σ', l'"
shows "A ⊨ (S, U →) onl Γ P"

```

proof

```

fix σ p
assume "(σ, p) ∈ init A"
hence "∀l∈labels Γ p. P (σ, l)" using init by simp
thus "onl Γ P (σ, p)" ..

```

next

```

fix σ p a σ' p'
assume rp: "(σ, p) ∈ oreachable A S U"
and "onl Γ P (σ, p)"
and tr: "((σ, p), a, (σ', p')) ∈ trans A"
and "S σ σ' a"
from ⟨onl Γ P (σ, p)⟩ have "∀l∈labels Γ p. P (σ, l)" ..
with rp tr ⟨S σ σ' a⟩ have "∀l'∈labels Γ p'. P (σ', l'" by (auto elim: step)
thus "onl Γ P (σ', p)" ..

```

next

```

fix σ σ' p
assume "(σ, p) ∈ oreachable A S U"
and "onl Γ P (σ, p)"
and "U σ σ'"
from ⟨onl Γ P (σ, p)⟩ have "∀l∈labels Γ p. P (σ, l)" by auto
with ⟨(σ, p) ∈ oreachable A S U⟩ have "∀l∈labels Γ p. P (σ', l)"
using ⟨U σ σ'⟩ by (rule other)
thus "onl Γ P (σ', p)" by auto
qed

```

lemma global\_oinvariantI [intro]:

```

assumes init: "∧σ p. (σ, p) ∈ init A ⇒ P σ"
and other: "∧σ σ' p l. [ (σ, p) ∈ oreachable A S U; P σ; U σ σ' ] ⇒ P σ'"
and step: "∧σ p a σ' p'.
           [ (σ, p) ∈ oreachable A S U;
             P σ;
             ((σ, p), a, (σ', p')) ∈ trans A;
             S σ σ' a ] ⇒ P σ'"
shows "A ⊨ (S, U →) (λ(σ, _). P σ)"

```

proof

```

fix σ p
assume "(σ, p) ∈ init A"

```



```

thus "(λ(σ, _). P σ) (σ, p)"
  by simp (erule init)
next
fix σ p a σ' p'
assume rp: "(σ, p) ∈ oreachable A S U"
  and "(λ(σ, _). P σ) (σ, p)"
  and tr: "((σ, p), a, (σ', p')) ∈ trans A"
  and "S σ σ' a"
from ⟨(λ(σ, _). P σ) (σ, p)⟩ have "P σ" by simp
with rp have "P σ'"
  using tr ⟨S σ σ' a⟩ by (rule step)
thus "(λ(σ, _). P σ) (σ', p)" by simp
next
fix σ σ' p
assume "(σ, p) ∈ oreachable A S U"
  and "(λ(σ, _). P σ) (σ, p)"
  and "U σ σ'"
hence "P σ'" by simp (erule other)
thus "(λ(σ, _). P σ) (σ', p)" by simp
qed

lemma onl_oinvariantD [dest]:
  assumes "A ⊨ (S, U →) onl Γ P"
    and "(σ, p) ∈ oreachable A S U"
    and "l ∈ labels Γ p"
  shows "P (σ, l)"
  using assms unfolding onl_def by auto

lemma onl_oinvariant_weakenD [dest]:
  assumes "A ⊨ (S', U' →) onl Γ P"
    and "(σ, p) ∈ oreachable A S U"
    and "l ∈ labels Γ p"
    and weakenS: "∧ s s' a. S s s' a ⇒ S' s s' a"
    and weakenU: "∧ s s'. U s s' ⇒ U' s s'"
  shows "P (σ, l)"
proof -
  from ⟨(σ, p) ∈ oreachable A S U⟩ have "(σ, p) ∈ oreachable A S' U'"
    by (rule oreachable_weakenE)
    (erule weakenS, erule weakenU)
  with ⟨A ⊨ (S', U' →) onl Γ P⟩ show "P (σ, l)"
    using ⟨l ∈ labels Γ p⟩ ..
qed

lemma onl_oinvariant_initD [dest]:
  assumes invP: "A ⊨ (S, U →) onl Γ P"
    and init: "(σ, p) ∈ init A"
    and pnl: "l ∈ labels Γ p"
  shows "P (σ, l)"
proof -
  from init have "(σ, p) ∈ oreachable A S U" ..
  with invP show ?thesis using pnl ..
qed

lemma onl_oinvariant_sterms:
  assumes wf: "wellformed Γ"
    and il: "A ⊨ (S, U →) onl Γ P"
    and rp: "(σ, p) ∈ oreachable A S U"
    and "p' ∈ sterms Γ p"
    and "l ∈ labels Γ p'"
  shows "P (σ, l)"
proof -
  from wf ⟨p' ∈ sterms Γ p⟩ ⟨l ∈ labels Γ p'⟩ have "l ∈ labels Γ p"
    by (rule labels_sterms_labels)
  with il rp show "P (σ, l)" ..

```

qed

lemma onl\_oinvariant\_sterms\_weaken:

assumes wf: "wellformed  $\Gamma$ "  
and il: " $A \models (S', U' \rightarrow) \text{ onl } \Gamma P$ "  
and rp: " $(\sigma, p) \in \text{oreachable } A S U$ "  
and "p'  $\in$  sterms  $\Gamma p$ "  
and "l  $\in$  labels  $\Gamma p$ "  
and weakenS: " $\bigwedge \sigma \sigma' a. S \sigma \sigma' a \implies S' \sigma \sigma' a$ "  
and weakenU: " $\bigwedge \sigma \sigma'. U \sigma \sigma' \implies U' \sigma \sigma'$ "  
shows " $P(\sigma, l)$ "

proof -

from  $(\sigma, p) \in \text{oreachable } A S U$  have " $(\sigma, p) \in \text{oreachable } A S' U'$ "  
by (rule oreachable\_weakenE)  
(erule weakenS, erule weakenU)  
with assms(1-2) show ?thesis using assms(4-5)  
by (rule onl\_oinvariant\_sterms)

qed

lemma otrans\_from\_sterms:

assumes " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "  
and "wellformed  $\Gamma$ "  
shows " $\exists p' \in \text{sterms } \Gamma p. ((\sigma, p'), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "  
using assms by (induction p rule: sterms\_pinduct [OF <wellformed  $\Gamma$ >]) auto

lemma otrans\_from\_sterms':

assumes " $((\sigma, p'), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "  
and "wellformed  $\Gamma$ "  
and " $p' \in \text{sterms } \Gamma p$ "  
shows " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "  
using assms by (induction p rule: sterms\_pinduct [OF <wellformed  $\Gamma$ >]) auto

lemma otrans\_to\_dterms:

assumes " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "  
and "wellformed  $\Gamma$ "  
shows " $\forall r \in \text{sterms } \Gamma q. r \in \text{dterms } \Gamma p$ "  
using assms by (induction q) auto

theorem cterms\_includes\_sterms\_of\_oseq\_reachable:

assumes "wellformed  $\Gamma$ "  
and "control\_within  $\Gamma$  (init A)"  
and "trans A = oseqp\_sos  $\Gamma i$ "  
shows " $\bigcup (\text{sterms } \Gamma \text{ ' snd ' oreachable } A S U) \subseteq \text{cterms } \Gamma$ "

proof

fix qs  
assume " $qs \in \bigcup (\text{sterms } \Gamma \text{ ' snd ' oreachable } A S U)$ "  
then obtain  $\xi$  and q where \*: " $(\xi, q) \in \text{oreachable } A S U$ "  
and \*\*: " $qs \in \text{sterms } \Gamma q$ " by auto  
from \* have " $\bigwedge x. x \in \text{sterms } \Gamma q \implies x \in \text{cterms } \Gamma$ "  
proof (induction rule: oreachable\_pair\_induct)  
fix  $\sigma p q$   
assume " $(\sigma, p) \in \text{init } A$ "  
and " $q \in \text{sterms } \Gamma p$ "  
from <control\_within  $\Gamma$  (init A)> and  $(\sigma, p) \in \text{init } A$   
obtain pn where " $p \in \text{subterms } (\Gamma pn)$ " by auto  
with <wellformed  $\Gamma$ > show " $q \in \text{cterms } \Gamma$ " using <q  $\in$  sterms  $\Gamma p$ >  
by (rule subterms\_sterms\_in\_cterms)

next

fix p  $\sigma a \sigma' q x$   
assume " $(\sigma, p) \in \text{oreachable } A S U$ "  
and IH: " $\bigwedge x. x \in \text{sterms } \Gamma p \implies x \in \text{cterms } \Gamma$ "  
and " $((\sigma, p), a, (\sigma', q)) \in \text{trans } A$ "  
and " $x \in \text{sterms } \Gamma q$ "  
from this(3) and <trans A = oseqp\_sos  $\Gamma i$ >

```

  have step: " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i$ " by simp
from step <wellformed  $\Gamma$ > obtain ps
  where ps: " $ps \in \text{sterms } \Gamma \ p$ "
  and step': " $((\sigma, ps), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i$ "
  by (rule otrans_from_sterms [THEN bexE])
from ps have " $ps \in \text{cterms } \Gamma$ " by (rule IH)
moreover from step' <wellformed  $\Gamma$ > < $x \in \text{sterms } \Gamma \ q$ > have " $x \in \text{dterms } \Gamma \ ps$ "
  by (rule otrans_to_dterms [rule_format])
ultimately show " $x \in \text{cterms } \Gamma$ " by (rule ctermsDI)
qed
thus " $qs \in \text{cterms } \Gamma$ " using **.
qed

```

corollary oseq\_reachable\_in\_cterms:

```

assumes "wellformed  $\Gamma$ "
  and "control_within  $\Gamma$  (init A)"
  and "trans A = oseqp_sos  $\Gamma \ i$ "
  and " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "
  and " $p' \in \text{sterms } \Gamma \ p$ "
shows " $p' \in \text{cterms } \Gamma$ "
using assms(1-3)
proof (rule cterms_includes_sterms_of_oseq_reachable [THEN set_mp])
  from assms(4-5) show " $p' \in \bigcup (\text{sterms } \Gamma \ ' \ \text{snd } ' \ \text{oreachable } A \ S \ U)$ "
    by (auto elim!: rev_bexI)
qed

```

lemma oseq\_invariant\_ctermI:

```

assumes wf: "wellformed  $\Gamma$ "
  and cw: "control_within  $\Gamma$  (init A)"
  and sl: "simple_labels  $\Gamma$ "
  and sp: "trans A = oseqp_sos  $\Gamma \ i$ "
  and init: " $\bigwedge \sigma \ p \ l. \llbracket$ 
     $(\sigma, p) \in \text{init } A;$ 
     $l \in \text{labels } \Gamma \ p$ 
   $\rrbracket \implies P (\sigma, l)$ "
  and other: " $\bigwedge \sigma \ \sigma' \ p \ l. \llbracket$ 
     $(\sigma, p) \in \text{oreachable } A \ S \ U;$ 
     $l \in \text{labels } \Gamma \ p;$ 
     $P (\sigma, l);$ 
     $U \ \sigma \ \sigma' \rrbracket \implies P (\sigma', l)$ "
  and local: " $\bigwedge p \ l \ \sigma \ a \ q \ l' \ \sigma' \ pp. \llbracket$ 
     $p \in \text{cterms } \Gamma;$ 
     $l \in \text{labels } \Gamma \ p;$ 
     $P (\sigma, l);$ 
     $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i;$ 
     $((\sigma, p), a, (\sigma', q)) \in \text{trans } A;$ 
     $l' \in \text{labels } \Gamma \ q;$ 
     $(\sigma, pp) \in \text{oreachable } A \ S \ U;$ 
     $p \in \text{sterms } \Gamma \ pp;$ 
     $(\sigma', q) \in \text{oreachable } A \ S \ U;$ 
     $S \ \sigma \ \sigma' \ a$ 
   $\rrbracket \implies P (\sigma', l')$ "
shows " $A \models (S, U \rightarrow) \text{ onl } \Gamma \ P$ "

```

proof

```

  fix  $\sigma \ p \ l$ 
  assume " $(\sigma, p) \in \text{init } A$ "
  and *: " $l \in \text{labels } \Gamma \ p$ "
  with init show " $P (\sigma, l)$ " by auto

```

next

```

  fix  $\sigma \ p \ a \ \sigma' \ q \ l'$ 
  assume sr: " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "
  and pl: " $\forall l \in \text{labels } \Gamma \ p. P (\sigma, l)$ "
  and tr: " $((\sigma, p), a, (\sigma', q)) \in \text{trans } A$ "
  and A6: " $l' \in \text{labels } \Gamma \ q$ "

```

```

    and "S  $\sigma \sigma' a$ "
    thus "P ( $\sigma', l'$ )"
proof -
  from sr and tr and  $\langle S \sigma \sigma' a \rangle$  have A7: " $(\sigma', q) \in \text{oreachable } A S U$ "
    by - (rule oreachable_local')
  from tr and sp have tr': " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ " by simp
  then obtain p' where "p'  $\in$  sterms  $\Gamma p$ "
    and A4: " $((\sigma, p'), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i$ "
    by (blast dest: otrans_from_sterms [OF wf])
  from wf cw sp sr this(1) have A1: "p'  $\in$  cterms  $\Gamma$ "
    by (rule oseq_reachable_in_cterms)
  from labels_not_empty [OF wf] obtain l1 where A2: "l1  $\in$  labels  $\Gamma p'$ "
    by blast
  with  $\langle p' \in \text{sterms } \Gamma p \rangle$  have "l1  $\in$  labels  $\Gamma p$ "
    by (rule labels_sterms_labels [OF wf])
  with p1 have A3: "P ( $\sigma, l1$ )" by simp
  from sr  $\langle p' \in \text{sterms } \Gamma p \rangle$ 
    obtain pp where A7: " $(\sigma, pp) \in \text{oreachable } A S U$ "
      and A8: "p'  $\in$  sterms  $\Gamma pp$ "
    by auto
  from sr tr  $\langle S \sigma \sigma' a \rangle$  have A9: " $(\sigma', q) \in \text{oreachable } A S U$ "
    by - (rule oreachable_local')
  from sp and  $\langle ((\sigma, p'), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma i \rangle$ 
    have A5: " $((\sigma, p'), a, (\sigma', q)) \in \text{trans } A$ " by simp
  from A1 A2 A3 A4 A5 A6 A7 A8 A9  $\langle S \sigma \sigma' a \rangle$  show ?thesis by (rule local)
qed
next
fix  $\sigma \sigma' p l$ 
assume sr: " $(\sigma, p) \in \text{oreachable } A S U$ "
  and "  $\forall l \in \text{labels } \Gamma p. P (\sigma, l)$ "
  and "U  $\sigma \sigma'$ "
show "  $\forall l \in \text{labels } \Gamma p. P (\sigma', l)$ "
proof
  fix l
  assume "l  $\in$  labels  $\Gamma p$ "
  with  $\langle \forall l \in \text{labels } \Gamma p. P (\sigma, l) \rangle$  have "P ( $\sigma, l$ )" ..
  with sr and  $\langle l \in \text{labels } \Gamma p \rangle$ 
    show "P ( $\sigma', l$ )" using  $\langle U \sigma \sigma' \rangle$  by (rule other)
qed
qed
lemma oseq_invariant_ctermsI:
  assumes wf: "wellformed  $\Gamma$ "
    and cw: "control_within  $\Gamma$  (init A)"
    and sl: "simple_labels  $\Gamma$ "
    and sp: "trans A = oseqp_sos  $\Gamma i$ "
    and init: " $\bigwedge \sigma p l. \llbracket$ 
      ( $\sigma, p$ )  $\in$  init A;
      l  $\in$  labels  $\Gamma p$ 
     $\rrbracket \implies P (\sigma, l)$ "
    and other: " $\bigwedge \sigma \sigma' p l. \llbracket$ 
      wellformed  $\Gamma$ ;
      ( $\sigma, p$ )  $\in$  oreachable A S U;
      l  $\in$  labels  $\Gamma p$ ;
      P ( $\sigma, l$ );
      U  $\sigma \sigma'$ 
     $\rrbracket \implies P (\sigma', l)$ "
    and local: " $\bigwedge p l \sigma a q l' \sigma' pp pn. \llbracket$ 
      wellformed  $\Gamma$ ;
      p  $\in$  ctermsl ( $\Gamma pn$ );
      not_call p;
      l  $\in$  labels  $\Gamma p$ ;
      P ( $\sigma, l$ );
      (( $\sigma, p$ ), a, ( $\sigma', q$ ))  $\in$  oseqp_sos  $\Gamma i$ ;
      (( $\sigma, p$ ), a, ( $\sigma', q$ ))  $\in$  trans A;
     $\rrbracket$ "

```

```

      l' ∈ labels Γ q;
      (σ, pp) ∈ oreachable A S U;
      p ∈ sterms Γ pp;
      (σ', q) ∈ oreachable A S U;
      S σ σ' a
    ]] ⇒ P (σ', l')"
  shows "A ⊨ (S, U →) onl Γ P"
proof (rule oseq_invariant_ctermI [OF wf cw sl sp])
  fix σ p l
  assume "(σ, p) ∈ init A"
    and "l ∈ labels Γ p"
  thus "P (σ, l)" by (rule init)
next
  fix σ σ' p l
  assume "(σ, p) ∈ oreachable A S U"
    and "l ∈ labels Γ p"
    and "P (σ, l)"
    and "U σ σ'"
  with wf show "P (σ', l)" by (rule other)
next
  fix p l σ a q l' σ' pp
  assume "p ∈ cterms Γ"
    and otherassms: "l ∈ labels Γ p"
      "P (σ, l)"
      "((σ, p), a, (σ', q)) ∈ oseqp_sos Γ i"
      "((σ, p), a, (σ', q)) ∈ trans A"
      "l' ∈ labels Γ q"
      "(σ, pp) ∈ oreachable A S U"
      "p ∈ sterms Γ pp"
      "(σ', q) ∈ oreachable A S U"
      "S σ σ' a"
  from this(1) obtain pn where "p ∈ ctermssl(Γ pn)"
    and "not_call p"
  unfolding cterms_def' [OF wf] by auto
  with wf show "P (σ', l)"
  using otherassms by (rule local)
qed

```

## 20.2 Open step invariants via labelled control terms

lemma onll\_ostep\_invariantI [intro]:

```

  assumes *: "∧σ p l a σ' p' l'. [| (σ, p) ∈ oreachable A S U;
    ((σ, p), a, (σ', p')) ∈ trans A;
    S σ σ' a;
    l ∈ labels Γ p;
    l' ∈ labels Γ p' |]
    ⇒ P ((σ, l), a, (σ', l'))"

```

shows "A ⊨<sub>A</sub> (S, U →) onll Γ P"

proof

```

  fix σ p σ' p' a
  assume "(σ, p) ∈ oreachable A S U"
    and "((σ, p), a, (σ', p')) ∈ trans A"
    and "S σ σ' a"
  hence "∀l ∈ labels Γ p. ∀l' ∈ labels Γ p'. P ((σ, l), a, (σ', l'))" by (auto elim!: *)
  thus "onll Γ P ((σ, p), a, (σ', p'))" ..
qed

```

lemma onll\_ostep\_invariantE [elim]:

```

  assumes "A ⊨A (S, U →) onll Γ P"
    and "(σ, p) ∈ oreachable A S U"
    and "((σ, p), a, (σ', p')) ∈ trans A"
    and "S σ σ' a"
    and lp: "l ∈ labels Γ p"
    and lp': "l' ∈ labels Γ p'"

```

```

  shows "P ((σ, l), a, (σ', l'))"
proof -
  from assms(1-4) have "onll Γ P ((σ, p), a, (σ', p'))" ..
  with lp lp' show "P ((σ, l), a, (σ', l'))" by auto
qed

lemma onll_ostep_invariantD [dest]:
  assumes "A ⊨A (S, U →) onll Γ P"
    and "(σ, p) ∈ oreachable A S U"
    and "((σ, p), a, (σ', p')) ∈ trans A"
    and "S σ σ' a"
  shows "∀l∈labels Γ p. ∀l'∈labels Γ p'. P ((σ, l), a, (σ', l'))"
  using assms by auto

lemma onll_ostep_invariant_weakenD [dest]:
  assumes "A ⊨A (S', U' →) onll Γ P"
    and "(σ, p) ∈ oreachable A S U"
    and "((σ, p), a, (σ', p')) ∈ trans A"
    and "S' σ σ' a"
    and weakenS: "∧s s' a. S s s' a ⇒ S' s s' a"
    and weakenU: "∧s s'. U s s' ⇒ U' s s'"
  shows "∀l∈labels Γ p. ∀l'∈labels Γ p'. P ((σ, l), a, (σ', l'))"
proof -
  from ⟨(σ, p) ∈ oreachable A S U⟩ have "(σ, p) ∈ oreachable A S' U'"
  by (rule oreachable_weakenE)
  (erule weakenS, erule weakenU)
  with ⟨A ⊨A (S', U' →) onll Γ P⟩ show ?thesis
  using ⟨((σ, p), a, (σ', p')) ∈ trans A⟩ and ⟨S' σ σ' a⟩ ..
qed

lemma onll_ostep_to_invariantI [intro]:
  assumes sinv: "A ⊨A (S, U →) onll Γ Q"
    and wf: "wellformed Γ"
    and init: "∧σ l p. [ (σ, p) ∈ init A; l∈labels Γ p ] ⇒ P (σ, l)"
    and other: "∧σ σ' p l.
    [ (σ, p) ∈ oreachable A S U;
      l∈labels Γ p;
      P (σ, l);
      U σ σ' ] ⇒ P (σ', l)"
    and local: "∧σ p l σ' l' a.
    [ (σ, p) ∈ oreachable A S U;
      l∈labels Γ p;
      P (σ, l);
      Q ((σ, l), a, (σ', l'));
      S σ σ' a ] ⇒ P (σ', l'"
  shows "A ⊨ (S, U →) onll Γ P"
proof
  fix σ p l
  assume "(σ, p) ∈ init A" and "l∈labels Γ p"
  thus "P (σ, l)" by (rule init)
next
  fix σ p a σ' p' l'
  assume sr: "(σ, p) ∈ oreachable A S U"
    and lp: "∀l∈labels Γ p. P (σ, l)"
    and tr: "((σ, p), a, (σ', p')) ∈ trans A"
    and "S σ σ' a"
    and lp': "l' ∈ labels Γ p'"
  show "P (σ', l'"
proof -
  from lp obtain l where "l∈labels Γ p" and "P (σ, l)"
  using labels_not_empty [OF wf] by auto
  from sinv sr tr ⟨S σ σ' a⟩ this(1) lp' have "Q ((σ, l), a, (σ', l'))" ..
  with sr ⟨l∈labels Γ p⟩ ⟨P (σ, l)⟩ show "P (σ', l'" using ⟨S σ σ' a⟩ by (rule local)
qed

```

```

next
  fix  $\sigma \sigma' p l$ 
  assume " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "
    and " $\forall l \in \text{labels } \Gamma \ p. \ P(\sigma, l)$ "
    and " $U \ \sigma \ \sigma'$ "
  show " $\forall l \in \text{labels } \Gamma \ p. \ P(\sigma', l)$ "
proof
  fix  $l$ 
  assume " $l \in \text{labels } \Gamma \ p$ "
  with  $\langle \forall l \in \text{labels } \Gamma \ p. \ P(\sigma, l) \rangle$  have " $P(\sigma, l)$ " ..
  with  $\langle (\sigma, p) \in \text{oreachable } A \ S \ U \rangle$  and  $\langle l \in \text{labels } \Gamma \ p \rangle$ 
  show " $P(\sigma', l)$ " using  $\langle U \ \sigma \ \sigma' \rangle$  by (rule other)
qed
qed

```

lemma onll\_ostep\_invariant\_sterms:

```

assumes wf: "wellformed  $\Gamma$ "
  and si: " $A \models_A (S, U \rightarrow) \text{onll } \Gamma \ P$ "
  and sr: " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "
  and sos: " $((\sigma, p), a, (\sigma', q)) \in \text{trans } A$ "
  and " $S \ \sigma \ \sigma' \ a$ "
  and " $l' \in \text{labels } \Gamma \ q$ "
  and " $p' \in \text{sterms } \Gamma \ p$ "
  and " $l \in \text{labels } \Gamma \ p'$ "
shows " $P((\sigma, l), a, (\sigma', l'))$ "
proof -
  from wf  $\langle p' \in \text{sterms } \Gamma \ p \rangle$   $\langle l \in \text{labels } \Gamma \ p' \rangle$  have " $l \in \text{labels } \Gamma \ p$ "
  by (rule labels_sterms_labels)
  with si sr sos  $\langle S \ \sigma \ \sigma' \ a \rangle$  show " $P((\sigma, l), a, (\sigma', l'))$ " using  $\langle l' \in \text{labels } \Gamma \ q \rangle$  ..
qed

```

lemma oseq\_step\_invariant\_sterms:

```

assumes inv: " $A \models_A (S, U \rightarrow) \text{onll } \Gamma \ P$ "
  and wf: "wellformed  $\Gamma$ "
  and sp: " $\text{trans } A = \text{oseqp\_sos } \Gamma \ i$ "
  and " $l' \in \text{labels } \Gamma \ q$ "
  and sr: " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "
  and tr: " $((\sigma, p'), a, (\sigma', q)) \in \text{trans } A$ "
  and " $S \ \sigma \ \sigma' \ a$ "
  and " $p' \in \text{sterms } \Gamma \ p$ "
shows " $\forall l \in \text{labels } \Gamma \ p'. \ P((\sigma, l), a, (\sigma', l'))$ "
proof
  from assms(3, 6) have " $((\sigma, p'), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i$ " by simp
  hence " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i$ "
  using wf  $\langle p' \in \text{sterms } \Gamma \ p \rangle$  by (rule otrans_from_sterms')
  with assms(3) have trp: " $((\sigma, p), a, (\sigma', q)) \in \text{trans } A$ " by simp
  fix  $l$  assume " $l \in \text{labels } \Gamma \ p'$ "
  with wf inv sr trp  $\langle S \ \sigma \ \sigma' \ a \rangle$   $\langle l' \in \text{labels } \Gamma \ q \rangle$   $\langle p' \in \text{sterms } \Gamma \ p \rangle$ 
  show " $P((\sigma, l), a, (\sigma', l'))$ "
  by - (erule(7) onll_ostep_invariant_sterms)
qed

```

lemma oseq\_step\_invariant\_sterms\_weaken:

```

assumes inv: " $A \models_A (S, U \rightarrow) \text{onll } \Gamma \ P$ "
  and wf: "wellformed  $\Gamma$ "
  and sp: " $\text{trans } A = \text{oseqp\_sos } \Gamma \ i$ "
  and " $l' \in \text{labels } \Gamma \ q$ "
  and sr: " $(\sigma, p) \in \text{oreachable } A \ S' \ U''$ "
  and tr: " $((\sigma, p'), a, (\sigma', q)) \in \text{trans } A$ "
  and " $S' \ \sigma \ \sigma' \ a$ "
  and " $p' \in \text{sterms } \Gamma \ p$ "
  and weakenS: " $\bigwedge \sigma \ \sigma' \ a. \ S' \ \sigma \ \sigma' \ a \implies S \ \sigma \ \sigma' \ a$ "
  and weakenU: " $\bigwedge \sigma \ \sigma'. \ U' \ \sigma \ \sigma' \implies U \ \sigma \ \sigma'$ "
shows " $\forall l \in \text{labels } \Gamma \ p'. \ P((\sigma, l), a, (\sigma', l'))$ "

```

proof -  
 from  $\langle S' \sigma \sigma' a \rangle$  have "S  $\sigma \sigma' a$ " by (rule weakenS)  
 from  $\langle (\sigma, p) \in \text{oreachable } A \ S \ U' \rangle$   
 have Ir: " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "  
 by (rule oreachable\_weakenE)  
 (erule weakenS, erule weakenU)  
 with assms(1-4) show ?thesis  
 using tr  $\langle S \sigma \sigma' a \rangle \langle p' \in \text{sterms } \Gamma \ p \rangle$   
 by (rule oseq\_step\_invariant\_sterms)  
 qed

lemma onll\_ostep\_invariant\_any\_sterms:  
 assumes wf: "wellformed  $\Gamma$ "  
 and si: " $A \models_A (S, U \rightarrow) \text{onll } \Gamma \ P$ "  
 and sr: " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "  
 and sos: " $((\sigma, p), a, (\sigma', q)) \in \text{trans } A$ "  
 and "S  $\sigma \sigma' a$ "  
 and " $l' \in \text{labels } \Gamma \ q$ "  
 shows " $\forall p' \in \text{sterms } \Gamma \ p. \forall l \in \text{labels } \Gamma \ p'. P ((\sigma, l), a, (\sigma', l'))$ "  
 by (intro ballI) (rule onll\_ostep\_invariant\_sterms [OF assms])

lemma oseq\_step\_invariant\_ctermI [intro]:  
 assumes wf: "wellformed  $\Gamma$ "  
 and cw: "control\_within  $\Gamma$  (init A)"  
 and sl: "simple\_labels  $\Gamma$ "  
 and sp: "trans A = oseqp\_sos  $\Gamma \ i$ "  
 and local: " $\bigwedge p \ l \ \sigma \ a \ q \ l' \ \sigma' \ pp. \llbracket$   
    $p \in \text{cterms } \Gamma;$   
    $l \in \text{labels } \Gamma \ p;$   
    $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i;$   
    $((\sigma, p), a, (\sigma', q)) \in \text{trans } A;$   
    $l' \in \text{labels } \Gamma \ q;$   
    $(\sigma, pp) \in \text{oreachable } A \ S \ U;$   
    $p \in \text{sterms } \Gamma \ pp;$   
    $(\sigma', q) \in \text{oreachable } A \ S \ U;$   
   S  $\sigma \sigma' a$   
    $\rrbracket \implies P ((\sigma, l), a, (\sigma', l'))$ "  
 shows " $A \models_A (S, U \rightarrow) \text{onll } \Gamma \ P$ "

proof  
 fix  $\sigma \ p \ l \ a \ \sigma' \ q \ l'$   
 assume sr: " $(\sigma, p) \in \text{oreachable } A \ S \ U$ "  
 and tr: " $((\sigma, p), a, (\sigma', q)) \in \text{trans } A$ "  
 and "S  $\sigma \sigma' a$ "  
 and pl: " $l \in \text{labels } \Gamma \ p$ "  
 and A5: " $l' \in \text{labels } \Gamma \ q$ "  
 from this(2) and sp have " $((\sigma, p), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i$ " by simp  
 then obtain  $p'$  where " $p' \in \text{sterms } \Gamma \ p$ "  
   and A3: " $((\sigma, p'), a, (\sigma', q)) \in \text{oseqp\_sos } \Gamma \ i$ "  
 by (blast dest: otrans\_from\_sterms [OF \_ wf])  
 from this(2) and sp have A4: " $((\sigma, p'), a, (\sigma', q)) \in \text{trans } A$ " by simp  
 from wf cw sp sr  $\langle p' \in \text{sterms } \Gamma \ p \rangle$  have A1: " $p' \in \text{cterms } \Gamma$ "  
 by (rule oseq\_reachable\_in\_cterms)  
 from sr  $\langle p' \in \text{sterms } \Gamma \ p \rangle$   
 obtain  $pp$  where A6: " $(\sigma, pp) \in \text{oreachable } A \ S \ U$ "  
   and A7: " $p' \in \text{sterms } \Gamma \ pp$ "  
 by auto  
 from sr tr  $\langle S \sigma \sigma' a \rangle$  have A8: " $(\sigma', q) \in \text{oreachable } A \ S \ U$ "  
 by - (erule(2) oreachable\_local')  
 from wf cw sp sr have " $\exists pn. p \in \text{subterms } (\Gamma \ pn)$ "  
 by (rule oreachable\_subterms)  
 with sl wf have " $\forall p' \in \text{sterms } \Gamma \ p. l \in \text{labels } \Gamma \ p'$ "  
 using pl by (rule simple\_labels\_in\_sterms)  
 with  $\langle p' \in \text{sterms } \Gamma \ p \rangle$  have " $l \in \text{labels } \Gamma \ p'$ " by simp  
 with A1 show " $P ((\sigma, l), a, (\sigma', l'))$ " using A3 A4 A5 A6 A7 A8  $\langle S \sigma \sigma' a \rangle$



```

    by (rule local)
qed

lemma oseq_step_invariant_ctermsI [intro]:
  assumes wf: "wellformed  $\Gamma$ "
    and "control_within  $\Gamma$  (init A)"
    and "simple_labels  $\Gamma$ "
    and "trans A = oseqp_sos  $\Gamma$  i"
    and local: " $\bigwedge p\ l\ \sigma\ a\ q\ l'\ \sigma'\ pp\ pn.$  [
      wellformed  $\Gamma$ ;
       $p \in cterms1\ (\Gamma\ pn)$ ;
      not_call p;
       $l \in labels\ \Gamma\ p$ ;
       $((\sigma, p), a, (\sigma', q)) \in oseqp_sos\ \Gamma\ i$ ;
       $((\sigma, p), a, (\sigma', q)) \in trans\ A$ ;
       $l' \in labels\ \Gamma\ q$ ;
       $(\sigma, pp) \in oreachable\ A\ S\ U$ ;
       $p \in sterms\ \Gamma\ pp$ ;
       $(\sigma', q) \in oreachable\ A\ S\ U$ ;
       $S\ \sigma\ \sigma'\ a$ 
    ]  $\implies P\ ((\sigma, l), a, (\sigma', l'))$ "
  shows " $A \models_A (S, U \rightarrow) onll\ \Gamma\ P$ "
using assms(1-4) proof (rule oseq_step_invariant_ctermI)
  fix p l  $\sigma$  a q l'  $\sigma'$  pp
  assume "p  $\in cterms\ \Gamma$ "
    and otherassms: "l  $\in labels\ \Gamma\ p$ "
      " $((\sigma, p), a, (\sigma', q)) \in oseqp_sos\ \Gamma\ i$ "
      " $((\sigma, p), a, (\sigma', q)) \in trans\ A$ "
      "l'  $\in labels\ \Gamma\ q$ "
      " $(\sigma, pp) \in oreachable\ A\ S\ U$ "
      "p  $\in sterms\ \Gamma\ pp$ "
      " $(\sigma', q) \in oreachable\ A\ S\ U$ "
      "S  $\sigma\ \sigma'\ a$ "
  from this(1) obtain pn where "p  $\in cterms1(\Gamma\ pn)$ "
    and "not_call p"
  unfolding cterms_def' [OF wf] by auto
  with wf show "P  $((\sigma, l), a, (\sigma', l'))$ "
  using otherassms by (rule local)
qed

lemma open_seqp_action [elim]:
  assumes "wellformed  $\Gamma$ "
    and " $((\sigma\ i, p), a, (\sigma'\ i, p')) \in seqp_sos\ \Gamma$ "
  shows " $((\sigma, p), a, (\sigma', p')) \in oseqp_sos\ \Gamma\ i$ "
proof -
  from assms obtain ps where " $ps \in sterms\ \Gamma\ p$ "
    and " $((\sigma\ i, ps), a, (\sigma'\ i, p')) \in seqp_sos\ \Gamma$ "
  by - (drule trans_from_sterms, auto)
  thus ?thesis
proof (induction p)
  fix p1 p2
  assume "[ ps  $\in sterms\ \Gamma\ p1$ ;  $((\sigma\ i, ps), a, \sigma'\ i, p') \in seqp_sos\ \Gamma$  ]"
     $\implies ((\sigma, p1), a, (\sigma', p')) \in oseqp_sos\ \Gamma\ i$ 
    and "[ ps  $\in sterms\ \Gamma\ p2$ ;  $((\sigma\ i, ps), a, \sigma'\ i, p') \in seqp_sos\ \Gamma$  ]"
     $\implies ((\sigma, p2), a, (\sigma', p')) \in oseqp_sos\ \Gamma\ i$ 
    and "ps  $\in sterms\ \Gamma\ (p1 \oplus p2)$ "
    and " $((\sigma\ i, ps), a, (\sigma'\ i, p')) \in seqp_sos\ \Gamma$ "
  with assms(1) show " $((\sigma, p1 \oplus p2), a, (\sigma', p')) \in oseqp_sos\ \Gamma\ i$ "
  by simp (metis oseqp_sos.ochoiceT1 oseqp_sos.ochoiceT2)
next
  fix l fip fmsg p1 p2
  assume IH1: "[ ps  $\in sterms\ \Gamma\ p1$ ;  $((\sigma\ i, ps), a, \sigma'\ i, p') \in seqp_sos\ \Gamma$  ]"
     $\implies ((\sigma, p1), a, (\sigma', p')) \in oseqp_sos\ \Gamma\ i$ 
    and IH2: "[ ps  $\in sterms\ \Gamma\ p2$ ;  $((\sigma\ i, ps), a, \sigma'\ i, p') \in seqp_sos\ \Gamma$  ]"

```

```

      ⇒ ((σ, p2), a, (σ', p')) ∈ oseqp_sos Γ i"
    and "ps ∈ sterms Γ ({1}unicast(fip, fmsg). p1 ▷ p2)"
    and "((σ i, ps), a, (σ' i, p')) ∈ seqp_sos Γ"
  from this(3-4) have "((σ i, {1}unicast(fip, fmsg). p1 ▷ p2), a, (σ' i, p')) ∈ seqp_sos Γ"
  by simp
  thus "((σ, {1}unicast(fip, fmsg). p1 ▷ p2), a, (σ', p')) ∈ oseqp_sos Γ i"
  proof (rule seqp_unicastTE)
    assume "a = unicast (fip (σ i)) (fmsg (σ i))"
      and "σ' i = σ i"
      and "p' = p1"
    thus ?thesis by auto
  next
    assume "a = ¬unicast (fip (σ i))"
      and "σ' i = σ i"
      and "p' = p2"
    thus ?thesis by auto
  qed
next
fix p
assume "ps ∈ sterms Γ (call(p))"
  and "((σ i, ps), a, (σ' i, p')) ∈ seqp_sos Γ"
with assms(1) have "((σ, ps), a, (σ', p')) ∈ oseqp_sos Γ i"
  by (cases ps) auto
with assms(1) (ps ∈ sterms Γ (call(p))) have "((σ, Γ p), a, (σ', p')) ∈ oseqp_sos Γ i"
  by - (rule otrans_from_sterms', simp_all)
thus "((σ, call(p)), a, (σ', p')) ∈ oseqp_sos Γ i" by auto
qed auto
qed

```

end

## 21 Transfer standard invariants into open invariants

```

theory OAWN_Convert
imports AWN_SOS_Labels AWN_Invariants
        OAWN_SOS OAWN_Invariants
begin

```

```

definition initiali :: "'i ⇒ (('i ⇒ 'g) × 'l) set ⇒ ('g × 'l) set ⇒ bool"
where "initiali i OI CI ≡ ({(σ i, p) | σ p. (σ, p) ∈ OI} = CI)"

```

```

lemma initialiI [intro]:
  assumes OICI: "∧σ p. (σ, p) ∈ OI ⇒ (σ i, p) ∈ CI"
  and CIOI: "∧ξ p. (ξ, p) ∈ CI ⇒ ∃σ. ξ = σ i ∧ (σ, p) ∈ OI"
  shows "initiali i OI CI"
  unfolding initiali_def
  by (intro set_eqI iffI) (auto elim!: OICI CIOI)

```

```

lemma open_from_initialiD [dest]:
  assumes "initiali i OI CI"
  and "(σ, p) ∈ OI"
  shows "∃ξ. σ i = ξ ∧ (ξ, p) ∈ CI"
  using assms unfolding initiali_def by auto

```

```

lemma closed_from_initialiD [dest]:
  assumes "initiali i OI CI"
  and "(ξ, p) ∈ CI"
  shows "∃σ. σ i = ξ ∧ (σ, p) ∈ OI"
  using assms unfolding initiali_def by auto

```

```

definition
  seq1 :: "'i ⇒ (('s × 'l) ⇒ bool) ⇒ (('i ⇒ 's) × 'l) ⇒ bool"
where
  "seq1 i P ≡ (λ(σ, p). P (σ i, p))"

```

```

lemma seqI [intro]:
  "P (fst s i, snd s)  $\implies$  seqI i P s"
  by (clarsimp simp: seqI_def)

lemma same_seqI [elim]:
  assumes " $\forall j \in \{i\}. \sigma' j = \sigma j$ "
    and "seqI i P ( $\sigma'$ , s)"
  shows "seqI i P ( $\sigma$ , s)"
  using assms unfolding seqI_def by (clarsimp)

lemma seqI_simp:
  "seqI i P ( $\sigma$ , p) = P ( $\sigma$  i, p)"
  unfolding seqI_def by simp

lemma other_steps_resp_local [intro!, simp]: "other_steps (other A I) I"
  by (clarsimp elim!: otherE)

lemma seqI_onI_swap:
  "seqI i (onI  $\Gamma$  P) = onI  $\Gamma$  (seqI i P)"
  unfolding seqI_def onI_def by simp

lemma oseqp_sos_resp_local_steps [intro!, simp]:
  fixes  $\Gamma$  :: "'p  $\implies$  ('s, 'm, 'p, 'l) seqp"
  shows "local_steps (oseqp_sos  $\Gamma$  i) {i}"
  proof
    fix  $\sigma \sigma' \zeta \zeta' ::$  "nat  $\implies$  's" and s a s'
    assume tr: " $((\sigma, s), a, \sigma', s') \in$  oseqp_sos  $\Gamma$  i"
      and " $\forall j \in \{i\}. \zeta j = \sigma j$ "
    thus " $\exists \zeta'. (\forall j \in \{i\}. \zeta' j = \sigma' j) \wedge ((\zeta, s), a, (\zeta', s')) \in$  oseqp_sos  $\Gamma$  i"
  proof induction
    fix  $\sigma \sigma' l ms p$ 
    assume " $\sigma' i = \sigma i$ "
      and " $\forall j \in \{i\}. \zeta j = \sigma j$ "
    hence " $((\zeta, \{l\}broadcast(ms).p), broadcast (ms (\sigma i)), (\sigma', p)) \in$  oseqp_sos  $\Gamma$  i"
      by (metis obroadcastT singleton_iff)
    with " $\forall j \in \{i\}. \zeta j = \sigma j$ " show " $\exists \zeta'. (\forall j \in \{i\}. \zeta' j = \sigma' j) \wedge$ 
       $((\zeta, \{l\}broadcast(ms).p), broadcast (ms (\sigma i)), (\zeta', p)) \in$  oseqp_sos  $\Gamma$  i"
      by blast
  next
    fix  $\sigma \sigma' ::$  "nat  $\implies$  's" and fmsg :: "'m  $\implies$  's  $\implies$  's" and msg l p
    assume *: " $\sigma' i = fmsg msg (\sigma i)$ "
      and **: " $\forall j \in \{i\}. \zeta j = \sigma j$ "
    hence " $\forall j \in \{i\}. (\zeta(i := fmsg msg (\zeta i))) j = \sigma' j$ " by clarsimp
    moreover from * **
      have " $((\zeta, \{l\}receive(fmsg).p), receive msg, (\zeta(i := fmsg msg (\zeta i)), p)) \in$  oseqp_sos  $\Gamma$  i"
      by (metis fun_upd_same oreceiveT)
    ultimately show " $\exists \zeta'. (\forall j \in \{i\}. \zeta' j = \sigma' j) \wedge$ 
       $((\zeta, \{l\}receive(fmsg).p), receive msg, (\zeta', p)) \in$  oseqp_sos  $\Gamma$  i"
      by blast
  next
    fix  $\sigma' \sigma l p$  and fas :: "'s  $\implies$  's"
    assume *: " $\sigma' i = fas (\sigma i)$ "
      and **: " $\forall j \in \{i\}. \zeta j = \sigma j$ "
    hence " $\forall j \in \{i\}. (\zeta(i := fas (\zeta i))) j = \sigma' j$ " by clarsimp
    moreover from * ** have " $((\zeta, \{l\}[fas] p), \tau, (\zeta(i := fas (\zeta i)), p)) \in$  oseqp_sos  $\Gamma$  i"
      by (metis fun_upd_same oassignT)
    ultimately show " $\exists \zeta'. (\forall j \in \{i\}. \zeta' j = \sigma' j) \wedge ((\zeta, \{l\}[fas] p), \tau, (\zeta', p)) \in$  oseqp_sos  $\Gamma$  i"
      by blast
  next
    fix g :: "'s  $\implies$  's set" and  $\sigma \sigma' l p$ 
    assume *: " $\sigma' i \in g (\sigma i)$ "
      and **: " $\forall j \in \{i\}. \zeta j = \sigma j$ "
    hence " $\forall j \in \{i\}. (SOME \zeta'. \zeta' i = \sigma' i) j = \sigma' j$ " by simp (metis (lifting, full_types) some_eq_ex)
  
```

```

moreover with * ** have "((ζ, {l}(g) p), τ, (SOME ζ'. ζ' i = σ' i, p)) ∈ oseqp_sos Γ i"
  by simp (metis oguardT step_seq_tau)
ultimately show "∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, {l}(g) p), τ, (ζ', p)) ∈ oseqp_sos Γ i"
  by blast
next
fix σ pn a σ' p'
assume "((σ, Γ pn), a, (σ', p')) ∈ oseqp_sos Γ i"
  and IH: "∀j∈{i}. ζ j = σ j ⇒ ∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, Γ pn), a, (ζ', p')) ∈ oseqp_
Γ i"
  and "∀j∈{i}. ζ j = σ j"
then obtain ζ' where "∀j∈{i}. ζ' j = σ' j"
  and "((ζ, Γ pn), a, (ζ', p')) ∈ oseqp_sos Γ i"
  by blast
thus "∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, call(pn)), a, (ζ', p')) ∈ oseqp_sos Γ i"
  by blast
next
fix σ p a σ' p' q
assume "((σ, p), a, (σ', p')) ∈ oseqp_sos Γ i"
  and "∀j∈{i}. ζ j = σ j ⇒ ∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, p), a, (ζ', p')) ∈ oseqp_sos
Γ i"
  and "∀j∈{i}. ζ j = σ j"
then obtain ζ' where "∀j∈{i}. ζ' j = σ' j"
  and "((ζ, p), a, (ζ', p')) ∈ oseqp_sos Γ i"
  by blast
thus "∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, p ⊕ q), a, (ζ', p')) ∈ oseqp_sos Γ i"
  by blast
next
fix σ p a σ' q q'
assume "((σ, q), a, (σ', q')) ∈ oseqp_sos Γ i"
  and "∀j∈{i}. ζ j = σ j ⇒ ∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, q), a, (ζ', q')) ∈ oseqp_sos
Γ i"
  and "∀j∈{i}. ζ j = σ j"
then obtain ζ' where "∀j∈{i}. ζ' j = σ' j"
  and "((ζ, q), a, (ζ', q')) ∈ oseqp_sos Γ i"
  by blast
thus "∃ζ'. (∀j∈{i}. ζ' j = σ' j) ∧ ((ζ, p ⊕ q), a, (ζ', q')) ∈ oseqp_sos Γ i"
  by blast
qed (simp_all, (metis ogrouppcastT ounicastT onotunicastT osendT odeliverT)+)
qed

```

```

lemma oseqp_sos_subreachable [intro!, simp]:
  assumes "trans OA = oseqp_sos Γ i"
  shows "subreachable OA (other ANY {i}) {i}"
  by rule (clarsimp simp add: assms(1))+

```

```

lemma oseq_step_is_seq_step:
  fixes σ :: "'ip ⇒ 's"
  assumes "((σ, p), a :: 'm seq_action, (σ', p')) ∈ oseqp_sos Γ i"
  and "σ i = ξ"
  shows "∃ξ'. σ' i = ξ' ∧ ((ξ, p), a, (ξ', p')) ∈ seqp_sos Γ i"
using assms proof induction
  fix σ σ' l ms p
  assume "σ' i = σ i"
  and "σ i = ξ"
  hence "σ' i = ξ" by simp
  have "((ξ, {l}broadcast(ms).p), broadcast (ms ξ), (ξ, p)) ∈ seqp_sos Γ i"
  by auto
  with ⟨σ i = ξ⟩ and ⟨σ' i = ξ⟩ show "∃ξ'. σ' i = ξ'"
  ∧ ((ξ, {l}broadcast(ms).p), broadcast (ms (σ i)), (ξ', p)) ∈ seqp_sos Γ i"
  by clarsimp
next
fix fmsg :: "'m ⇒ 's ⇒ 's" and msg :: 'm and σ σ' l p
assume "σ' i = fmsg msg (σ i)"
and "σ i = ξ"

```

```

have "((ξ, {l}receive(fmsg).p), receive msg, (fmsg msg ξ, p)) ∈ seqp_sos Γ"
  by auto
with ⟨σ' i = fmsg msg (σ i)⟩ and ⟨σ i = ξ⟩
  show "∃ξ'. σ' i = ξ' ∧ ((ξ, {l}receive(fmsg).p), receive msg, (ξ', p)) ∈ seqp_sos Γ"
  by clarsimp
qed (simp_all, (metis assignT choiceT1 choiceT2 groupcastT guardT
  callT unicastT notunicastT sendT deliverT step_seq_tau)+)

```

lemma reachable\_oseq\_seqp\_sos:

```

assumes "(σ, p) ∈ reachable OA I"
  and "initiali i (init OA) (init A)"
  and spo: "trans OA = oseqp_sos Γ i"
  and sp: "trans A = seqp_sos Γ"
  shows "∃ξ. σ i = ξ ∧ (ξ, p) ∈ reachable A I"
using assms(1) proof (induction rule: reachable_pair_induct)
  fix σ p
  assume "(σ, p) ∈ init OA"
  with ⟨initiali i (init OA) (init A)⟩ obtain ξ where "σ i = ξ"
    and "(ξ, p) ∈ init A"
  by auto
  from ⟨(ξ, p) ∈ init A⟩ have "(ξ, p) ∈ reachable A I" ..
  with ⟨σ i = ξ⟩ show "∃ξ. σ i = ξ ∧ (ξ, p) ∈ reachable A I"
  by auto
next
  fix σ p σ' p' a
  assume "(σ, p) ∈ reachable OA I"
  and IH: "∃ξ. σ i = ξ ∧ (ξ, p) ∈ reachable A I"
  and otr: "((σ, p), a, (σ', p')) ∈ trans OA"
  and "I a"
  from IH obtain ξ where "σ i = ξ"
    and cr: "(ξ, p) ∈ reachable A I"
  by clarsimp
  from otr and spo have "((σ, p), a, (σ', p')) ∈ oseqp_sos Γ i" by simp
  with ⟨σ i = ξ⟩ obtain ξ' where "σ' i = ξ'"
    and "((ξ, p), a, (ξ', p')) ∈ seqp_sos Γ"
  by (auto dest!: oseq_step_is_seq_step)
  from this(2) and sp have ctr: "((ξ, p), a, (ξ', p')) ∈ trans A" by simp
  from ⟨(ξ, p) ∈ reachable A I⟩ and ctr and ⟨I a⟩
  have "(ξ', p') ∈ reachable A I" ..
  with ⟨σ' i = ξ'⟩ show "∃ξ. σ' i = ξ ∧ (ξ, p') ∈ reachable A I"
  by blast
qed

```

lemma reachable\_oseq\_seqp\_sos':

```

assumes "s ∈ reachable OA I"
  and "initiali i (init OA) (init A)"
  and "trans OA = oseqp_sos Γ i"
  and "trans A = seqp_sos Γ"
  shows "∃ξ. (fst s) i = ξ ∧ (ξ, snd s) ∈ reachable A I"
using assms
  by - (cases s, auto dest: reachable_oseq_seqp_sos)

```

Any invariant shown in the (simpler) closed semantics can be transferred to an invariant in the open semantics.

theorem open\_seq\_invariant [intro]:

```

assumes "A ⊨ (I →) P"
  and "initiali i (init OA) (init A)"
  and spo: "trans OA = oseqp_sos Γ i"
  and sp: "trans A = seqp_sos Γ"
  shows "OA ⊨ (act I, other ANY {i} →) (seq1 i P)"
proof -
  have "OA ⊨ (I →) (seq1 i P)"
  proof (rule invariant_arbitraryI)
    fix s
    assume "s ∈ reachable OA I"

```

```

with ⟨initiali i (init OA) (init A)⟩ obtain ξ where "(fst s) i = ξ"
                                and "(ξ, snd s) ∈ reachable A I"
  by (auto dest: reachable_oseq_seqp_sos' [OF _ _ spo sp])
with ⟨A ⊨ (I →) P⟩ have "P (ξ, snd s)" by auto
with ⟨(fst s) i = ξ⟩ show "seq1 i P s" by auto
qed
moreover from spo have "subreachable OA (other ANY {i}) {i}" ..
ultimately show ?thesis
proof (rule open_closed_invariant)
  fix σ σ' s
  assume "∀j∈{i}. σ' j = σ j"
    and "seq1 i P (σ', s)"
  thus "seq1 i P (σ, s)" ..
qed
qed

definition
  seqll :: "'i ⇒ (((s × 'l) × 'a × (s × 'l)) ⇒ bool)
           ⇒ (((i ⇒ 's) × 'l) × 'a × ((i ⇒ 's) × 'l)) ⇒ bool"
where
  "seqll i P ≡ (λ((σ, p), a, (σ', p')). P ((σ i, p), a, (σ' i, p')))"

lemma same_seqll [elim]:
  assumes "∀j∈{i}. σ1' j = σ1 j"
    and "∀j∈{i}. σ2' j = σ2 j"
    and "seqll i P ((σ1', s), a, (σ2', s'))"
  shows "seqll i P ((σ1, s), a, (σ2, s'))"
  using assms unfolding seqll_def by (clarsimp)

lemma seqllI [intro!]:
  assumes "P ((σ i, p), a, (σ' i, p'))"
  shows "seqll i P ((σ, p), a, (σ', p'))"
  using assms unfolding seqll_def by simp

lemma seqllD [dest]:
  assumes "seqll i P ((σ, p), a, (σ', p'))"
  shows "P ((σ i, p), a, (σ' i, p'))"
  using assms unfolding seqll_def by simp

lemma seqllsimp:
  "seqll i P ((σ, p), a, (σ', p')) = P ((σ i, p), a, (σ' i, p'))"
  unfolding seqll_def by simp

lemma seqll_onll_swap:
  "seqll i (onll Γ P) = onll Γ (seqll i P)"
  unfolding seqll_def onll_def by simp

theorem open_seq_step_invariant [intro]:
  assumes "A ⊨A (I →) P"
    and "initiali i (init OA) (init A)"
    and spo: "trans OA = oseqp_sos Γ i"
    and sp: "trans A = seqp_sos Γ"
  shows "OA ⊨A (act I, other ANY {i} →) (seqll i P)"
proof -
  have "OA ⊨A (I →) (seqll i P)"
  proof (rule step_invariant_arbitraryI)
    fix σ p a σ' p'
    assume or: "(σ, p) ∈ reachable OA I"
      and otr: "((σ, p), a, (σ', p')) ∈ trans OA"
      and "I a"
    from or ⟨initiali i (init OA) (init A)⟩ spo sp obtain ξ where "σ i = ξ"
      and cr: "(ξ, p) ∈ reachable A I"
    by - (drule(3) reachable_oseq_seqp_sos', auto)
    from otr and spo have "((σ, p), a, (σ', p')) ∈ oseqp_sos Γ i" by simp
  qed

```

```

with ⟨σ i = ξ⟩ obtain ξ' where "σ' i = ξ'"
  and ctr: "((ξ, p), a, (ξ', p')) ∈ seqp_sos Γ"
  by (auto dest!: oseq_step_is_seq_step)
with sp have "((ξ, p), a, (ξ', p')) ∈ trans A" by simp
with ⟨A ⊨A (I →) P⟩ cr have "P ((ξ, p), a, (ξ', p'))" using ⟨I a⟩ ..
with ⟨σ i = ξ⟩ and ⟨σ' i = ξ'⟩ have "P ((σ i, p), a, (σ' i, p'))" by simp
thus "seqll i P ((σ, p), a, (σ', p'))" ..
qed
moreover from spo have "local_steps (trans OA) {i}" by simp
moreover have "other_steps (other ANY {i}) {i}" ..
ultimately show ?thesis
proof (rule open_closed_step_invariant)
  fix σ ζ a σ' ζ' s s'
  assume "∀j∈{i}. σ j = ζ j"
    and "∀j∈{i}. σ' j = ζ' j"
    and "seqll i P ((σ, s), a, (σ', s'))"
  thus "seqll i P ((ζ, s), a, (ζ', s'))" ..
qed
qed
end

```

## 22 Model the standard queuing model

```

theory Qmsg
imports Awn_SOS_Labels Awn_Invariants
begin

Define the queue process

fun ΓQMSG :: "('m list, 'm, unit, unit label) seqp_env"
where
  "ΓQMSG () = labelled () (receive(λmsg msgs. msgs @ [msg]). call())
    ⊕ ⟨msgs. msgs ≠ []⟩
      (send(λmsgs. hd msgs).
        ([msgs. tl msgs] call())
        ⊕ receive(λmsg msgs. tl msgs @ [msg]). call())
    ⊕ receive(λmsg msgs. msgs @ [msg]). call())"

definition σQMSG :: "((m::msg) list × ('m list, 'm, unit, unit label) seqp) set"
where "σQMSG ≡ {[[], ΓQMSG ()]}"

abbreviation qmsg
  :: "((m::msg) list × ('m list, 'm, unit, unit label) seqp, 'm seq_action) automaton"
where
  "qmsg ≡ (| init = σQMSG, trans = seqp_sos ΓQMSG |)"

declare ΓQMSG.simps [simp del, code del]
lemmas ΓQMSG.simps [simp, code] = ΓQMSG.simps [simplified]

lemma σQMSG_not_empty [simp, intro]: "σQMSG ≠ {}"
  unfolding σQMSG_def by simp

lemma σQMSG_exists [simp]: "∃qmsg q. (qmsg, q) ∈ σQMSG"
  unfolding σQMSG_def by simp

lemma qmsg_wf [simp]: "wellformed ΓQMSG"
  by (rule wf_no_direct_calls) auto

lemmas qmsg_labels_not_empty [simp] = labels_not_empty [OF qmsg_wf]

lemma qmsg_control_within [simp]: "control_within ΓQMSG (init qmsg)"
  unfolding σQMSG_def by (rule control_withinI) (auto simp del: ΓQMSG.simps)

lemma qmsg_simple_labels [simp]: "simple_labels ΓQMSG"

```

```

unfolding simple_labels_def by auto

lemma qmsg_trans: "trans qmsg = seqp_sos  $\Gamma_{QMSG}$ "
  by simp

lemma  $\sigma_{QMSG\_labels}$  [simp]: " $(\xi, q) \in \sigma_{QMSG} \implies labels \Gamma_{QMSG} q = \{()\}$ "
  unfolding  $\sigma_{QMSG\_def}$  by simp

lemma qmsg_proc_cases [dest]:
  fixes p pn
  shows " $p \in ctermsl (\Gamma_{QMSG} pn) \implies p \in ctermsl (\Gamma_{QMSG} ())$ "
  by simp

declare
   $\Gamma_{QMSG\_simps}$  [ctermenv]
  qmsg_proc_cases [ctermcases]
  seq_invariant_ctermI [OF qmsg_wf qmsg_control_within qmsg_simple_labels qmsg_trans, cterms_intros]
  seq_step_invariant_ctermI [OF qmsg_wf qmsg_control_within qmsg_simple_labels qmsg_trans, cterms_intros]

end

```

## 23 Lifting rules for parallel compositions with QMSG

```

theory Qmsg_Lifting
imports Qmsg OAWN_SOS Inv_Cterms OAWN_Invariants
begin

lemma oseq_no_change_on_send:
  fixes  $\sigma$  s a  $\sigma'$  s'
  assumes " $((\sigma, s), a, (\sigma', s')) \in oseqp\_sos \Gamma i$ "
  shows "case a of
    broadcast m       $\implies \sigma' i = \sigma i$ 
  | groupcast ips m  $\implies \sigma' i = \sigma i$ 
  | unicast ips m    $\implies \sigma' i = \sigma i$ 
  |  $\neg$ unicast ips    $\implies \sigma' i = \sigma i$ 
  | send m           $\implies \sigma' i = \sigma i$ 
  | deliver m        $\implies \sigma' i = \sigma i$ 
  | _  $\implies True$ "
  using assms by induction simp_all

lemma qmsg_no_change_on_send_or_receive:
  fixes  $\sigma$  s a  $\sigma'$  s'
  assumes " $((\sigma, s), a, (\sigma', s')) \in oparp\_sos i (oseqp\_sos \Gamma i) (seqp\_sos \Gamma_{QMSG})$ "
    and "a  $\neq \tau$ "
  shows " $\sigma' i = \sigma i$ "
  proof -
    from assms(1) obtain p q p' q'
      where " $((\sigma, (p, q)), a, (\sigma', (p', q')) \in oparp\_sos i (oseqp\_sos \Gamma i) (seqp\_sos \Gamma_{QMSG})$ "
      by (cases s, cases s', simp)
    thus ?thesis
  proof
    assume " $((\sigma, p), a, (\sigma', p')) \in oseqp\_sos \Gamma i$ "
      and " $\bigwedge m. a \neq receive\ m$ "
    with (a  $\neq \tau$ ) show " $\sigma' i = \sigma i$ "
      by - (drule oseq_no_change_on_send, cases a, auto)
  next
    assume " $(q, a, q') \in seqp\_sos \Gamma_{QMSG}$ "
      and " $\sigma' i = \sigma i$ "
    thus " $\sigma' i = \sigma i$ " by simp
  next
    assume "a =  $\tau$ " with (a  $\neq \tau$ ) show ?thesis by auto
  qed
qed

```



```

lemma qmsg_msgs_not_empty:
  "qmsg  $\models$  onl  $\Gamma_{QMSG}$  ( $\lambda$ (msgs, l). l = ()-:1  $\longrightarrow$  msgs  $\neq$  [])"
  by inv_cterms

lemma qmsg_send_from_queue:
  "qmsg  $\models_A$  ( $\lambda$ ((msgs, q), a, _). sendmsg ( $\lambda$ m. m $\in$ set msgs) a)"
  proof -
    have "qmsg  $\models_A$  onll  $\Gamma_{QMSG}$  ( $\lambda$ ((msgs, _), a, _). sendmsg ( $\lambda$ m. m $\in$ set msgs) a)"
      by (inv_cterms inv add: onl_invariant_sterms [OF qmsg_wf qmsg_msgs_not_empty])
    thus ?thesis
      by (rule step_invariant_weakenE) (auto dest!: onllD)
  qed

lemma qmsg_queue_contents:
  "qmsg  $\models_A$  ( $\lambda$ ((msgs, q), a, (msgs', q')). case a of
    receive m  $\Rightarrow$  set msgs'  $\subseteq$  set (msgs @ [m])
    | _  $\Rightarrow$  set msgs'  $\subseteq$  set msgs)"
  proof -
    have "qmsg  $\models_A$  onll  $\Gamma_{QMSG}$  ( $\lambda$ ((msgs, q), a, (msgs', q')).
      case a of
        receive m  $\Rightarrow$  set msgs'  $\subseteq$  set (msgs @ [m])
        | _  $\Rightarrow$  set msgs'  $\subseteq$  set msgs)"
      by (inv_cterms) (clarsimp simp add: in_set_tl)+
    thus ?thesis
      by (rule step_invariant_weakenE) (auto dest!: onllD)
  qed

lemma qmsg_send_receive_or_tau:
  "qmsg  $\models_A$  ( $\lambda$ (_, a, _).  $\exists$ m. a = send m  $\vee$  a = receive m  $\vee$  a =  $\tau$ )"
  proof -
    have "qmsg  $\models_A$  onll  $\Gamma_{QMSG}$  ( $\lambda$ (_, a, _).  $\exists$ m. a = send m  $\vee$  a = receive m  $\vee$  a =  $\tau$ )"
      by inv_cterms
    thus ?thesis
      by rule (auto dest!: onllD)
  qed

lemma par_qmsg_oreachable:
  assumes "( $\sigma$ ,  $\zeta$ )  $\in$  oreachable (A  $\langle\langle$ i qmsg) (otherwith S {i} (orecvmsg R)) (other U {i})"
    (is "_  $\in$  oreachable _ ?owS _")
  and pinv: "A  $\models_A$  (otherwith S {i} (orecvmsg R), other U {i}  $\rightarrow$ )
    globala ( $\lambda$ ( $\sigma$ , _,  $\sigma'$ ). U ( $\sigma$  i) ( $\sigma'$  i))"
  and ustutter: " $\bigwedge$  $\xi$ . U  $\xi$   $\xi$ "
  and sgivesu: " $\bigwedge$  $\xi$   $\xi'$ . S  $\xi$   $\xi'$   $\implies$  U  $\xi$   $\xi'$ "
  and upreservesq: " $\bigwedge$  $\sigma$   $\sigma'$  m.  $\llbracket \forall j. U (\sigma j) (\sigma' j); R \sigma m \rrbracket \implies R \sigma' m$ "
  shows "( $\sigma$ , fst  $\zeta$ )  $\in$  oreachable A ?owS (other U {i})
     $\wedge$  snd  $\zeta$   $\in$  reachable qmsg (recvmsg (R  $\sigma$ ))
     $\wedge$  ( $\forall m \in$ set (fst (snd  $\zeta$ )). R  $\sigma$  m)"
  using assms(1) proof (induction rule: oreachable_pair_induct)
    fix  $\sigma$  pq
    assume "( $\sigma$ , pq)  $\in$  init (A  $\langle\langle$ i qmsg)"
    then obtain p ms q where "pq = (p, (ms, q))"
      and "( $\sigma$ , p)  $\in$  init A"
      and "(ms, q)  $\in$  init qmsg"
    by (clarsimp simp del:  $\Gamma_{QMSG\_simps}$ )
    from this(2) have "( $\sigma$ , p)  $\in$  oreachable A ?owS (other U {i})" ..
    moreover from  $\langle$ (ms, q)  $\in$  init qmsg have "(ms, q)  $\in$  reachable qmsg (recvmsg (R  $\sigma$ ))" ..
    moreover from  $\langle$ (ms, q)  $\in$  init qmsg have "ms = []"
      unfolding  $\sigma_{QMSG\_def}$  by simp
    ultimately show "( $\sigma$ , fst pq)  $\in$  oreachable A ?owS (other U {i})
       $\wedge$  snd pq  $\in$  reachable qmsg (recvmsg (R  $\sigma$ ))
       $\wedge$  ( $\forall m \in$ set (fst (snd pq)). R  $\sigma$  m)"
      using  $\langle$ pq = (p, (ms, q)) $\rangle$  by simp
  next
    note  $\Gamma_{QMSG\_simps}$  [simp del]
  end

```

```

case (other  $\sigma$  pq  $\sigma'$ )
hence " $(\sigma, \text{fst } pq) \in \text{oreachable } A \text{ ?owS } (\text{other } U \{i\})$ "
  and " $\text{other } U \{i\} \sigma \sigma'$ "
  and qr: " $\text{snd } pq \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \sigma))$ "
  and " $\forall m \in \text{set } (\text{fst } (\text{snd } pq)). R \sigma m$ "
  by simp_all
from (other  $U \{i\} \sigma \sigma'$ ) and ustutter have " $\forall j. U (\sigma j) (\sigma' j)$ "
  by (clarsimp elim!: otherE) metis
from (other  $U \{i\} \sigma \sigma'$ )
  and  $(\sigma, \text{fst } pq) \in \text{oreachable } A \text{ ?owS } (\text{other } U \{i\})$ 
  have " $(\sigma', \text{fst } pq) \in \text{oreachable } A \text{ ?owS } (\text{other } U \{i\})$ "
  by - (rule oreachable_other')
moreover have " $\forall m \in \text{set } (\text{fst } (\text{snd } pq)). R \sigma' m$ "
proof
  fix m assume " $m \in \text{set } (\text{fst } (\text{snd } pq))$ "
  with  $(\forall m \in \text{set } (\text{fst } (\text{snd } pq)). R \sigma m)$  have " $R \sigma m$ " ..
  with  $(\forall j. U (\sigma j) (\sigma' j))$  show " $R \sigma' m$ " by (rule upreservesq)
qed
moreover from qr have " $\text{snd } pq \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \sigma'))$ "
proof
  fix a
  assume " $\text{recvmsg } (R \sigma) a$ "
  thus " $\text{recvmsg } (R \sigma') a$ "
  proof (rule recvmsgE [where R=R])
    fix m assume " $R \sigma m$ "
    with  $(\forall j. U (\sigma j) (\sigma' j))$  show " $R \sigma' m$ " by (rule upreservesq)
  qed
qed
ultimately show ?case using qr by simp
next
case (local  $\sigma$  pq  $\sigma'$  pq' a)
obtain p ms q p' ms' q' where " $\text{pq} = (p, (\text{ms}, q))$ "
  and " $\text{pq}' = (p', (\text{ms}', q'))$ "
  by (cases pq, cases pq') metis
with local.hyps local.IH
  have pptr: " $((\sigma, (p, (\text{ms}, q))), a, (\sigma', (p', (\text{ms}', q'))))$ "
     $\in \text{oparp\_sos } i (\text{trans } A) (\text{seqp\_sos } \Gamma_{QMSG})$ "
  and por: " $(\sigma, p) \in \text{oreachable } A \text{ ?owS } (\text{other } U \{i\})$ "
  and qr: " $(\text{ms}, q) \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \sigma))$ "
  and " $\forall m \in \text{set } \text{ms}. R \sigma m$ "
  and " $\text{?owS } \sigma \sigma' a$ "
  by (simp_all del:  $\Gamma_{QMSG}$ _simps)
from (owS  $\sigma \sigma' a$ ) have " $\forall j. j \neq i \rightarrow S (\sigma j) (\sigma' j)$ "
  by (clarsimp dest!: otherwith_syncD)
with sgivesu have " $\forall j. j \neq i \rightarrow U (\sigma j) (\sigma' j)$ " by simp
from (owS  $\sigma \sigma' a$ ) have "orecvmsg  $R \sigma a$ " by (rule otherwithE)
hence " $\text{recvmsg } (R \sigma) a$ " ..
from pptr have " $(\sigma', p') \in \text{oreachable } A \text{ ?owS } (\text{other } U \{i\})$ "
   $\wedge (\text{ms}', q') \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \sigma'))$ 
   $\wedge (\forall m \in \text{set } \text{ms}'. R \sigma' m)$ "
proof
  assume " $((\sigma, p), a, (\sigma', p')) \in \text{trans } A$ "
  and " $\bigwedge m. a \neq \text{receive } m$ "
  and " $(\text{ms}', q') = (\text{ms}, q)$ "
  from this(1) have ptr: " $((\sigma, p), a, (\sigma', p')) \in \text{trans } A$ " by simp
  with pinv por and (owS  $\sigma \sigma' a$ ) have " $U (\sigma i) (\sigma' i)$ "
  by (auto dest!: ostep_invariantD)
  with  $(\forall j. j \neq i \rightarrow U (\sigma j) (\sigma' j))$  have " $\forall j. U (\sigma j) (\sigma' j)$ " by auto
  hence recvmsg': " $\bigwedge a. \text{recvmsg } (R \sigma) a \implies \text{recvmsg } (R \sigma') a$ "
  by (auto elim!: recvmsgE [where R=R] upreservesq)

```

from por ptr (?owS  $\sigma$   $\sigma'$  a) have " $(\sigma', p') \in \text{oreachable } A \text{ ?owS (other } U \{i\})$ "  
 by - (rule oreachable\_local')

moreover have " $(ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma'))$ "

proof -

from qr and  $\langle (ms', q') = (ms, q) \rangle$

have " $(ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma))$ " by simp

thus ?thesis by (rule reachable\_weakenE) (erule recvmsg')

qed

moreover have " $\forall m \in \text{set } ms'. R \sigma' m$ "

proof

fix m

assume " $m \in \text{set } ms'$ "

with  $\langle (ms', q') = (ms, q) \rangle$  have " $m \in \text{set } ms$ " by simp

with  $\langle \forall m \in \text{set } ms. R \sigma m \rangle$  have " $R \sigma m$ " ..

with  $\langle \forall j. U (\sigma j) (\sigma' j) \rangle$  show " $R \sigma' m$ "

by (rule upreservesq)

qed

ultimately show

" $(\sigma', p') \in \text{oreachable } A \text{ ?owS (other } U \{i\})$ "

$\wedge (ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma'))$

$\wedge (\forall m \in \text{set } ms'. R \sigma' m)$ " by simp\_all

next

assume qtr: " $((ms, q), a, (ms', q')) \in \text{seqp\_sos } \Gamma_{QMSG}$ "

and " $\bigwedge m. a \neq \text{send } m$ "

and " $p' = p$ "

and " $\sigma' i = \sigma i$ "

from this(4) and  $\langle \bigwedge \xi. U \xi \xi \rangle$  have " $U (\sigma i) (\sigma' i)$ " by simp

with  $\langle \forall j. j \neq i \rightarrow U (\sigma j) (\sigma' j) \rangle$  have " $\forall j. U (\sigma j) (\sigma' j)$ " by auto

hence recvmsg': " $\bigwedge a. \text{recvmsg (R } \sigma) a \implies \text{recvmsg (R } \sigma') a$ "

by (auto elim!: recvmsgE [where R=R] upreservesq)

from qtr have tqtr: " $((ms, q), a, (ms', q')) \in \text{trans } \text{qmsg}$ " by simp

from  $\langle \forall j. U (\sigma j) (\sigma' j) \rangle$  and  $\langle \sigma' i = \sigma i \rangle$  have " $\text{other } U \{i\} \sigma \sigma'$ " by auto

with por and  $\langle p' = p \rangle$

have " $(\sigma', p') \in \text{oreachable } A \text{ ?owS (other } U \{i\})$ "

by (auto dest: oreachable\_other)

moreover have " $(ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma'))$ "

proof (rule reachable\_weakenE [where P="recvmsg (R  $\sigma$ )"])

from qr tqtr (recvmsg (R  $\sigma$ ) a) show " $(ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma))$ " ..

qed (rule recvmsg')

moreover have " $\forall m \in \text{set } ms'. R \sigma' m$ "

proof

fix m

assume " $m \in \text{set } ms'$ "

moreover have " $\text{case } a \text{ of receive } m \Rightarrow \text{set } ms' \subseteq \text{set } (ms @ [m]) \mid \_ \Rightarrow \text{set } ms' \subseteq \text{set } ms$ "

proof -

from qr have " $(ms, q) \in \text{reachable } \text{qmsg } TT$ " ..

thus ?thesis using tqtr

by (auto dest!: step\_invariantD [OF qmsg\_queue\_contents])

qed

ultimately have " $R \sigma m$ " using  $\langle \forall m \in \text{set } ms. R \sigma m \rangle$  and  $\langle \text{recvmsg } R \sigma a \rangle$

by (cases a) auto

with  $\langle \forall j. U (\sigma j) (\sigma' j) \rangle$  show " $R \sigma' m$ "

by (rule upreservesq)

qed

ultimately show " $(\sigma', p') \in \text{oreachable } A \text{ ?owS (other } U \{i\})$   
 $\wedge (ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma'))$   
 $\wedge (\forall m \in \text{set } ms'. R \sigma' m)$ " by simp

next

fix m  
 assume "a =  $\tau$ "  
 and " $((\sigma, p), \text{receive } m, (\sigma', p')) \in \text{trans } A$ "  
 and " $((ms, q), \text{send } m, (ms', q')) \in \text{seqp\_sos } \Gamma_{QMSG}$ "  
 from this(2-3)  
 have ptr: " $((\sigma, p), \text{receive } m, (\sigma', p')) \in \text{trans } A$ "  
 and qtr: " $((ms, q), \text{send } m, (ms', q')) \in \text{trans } \text{qmsg}$ " by simp\_all

from qr have " $(ms, q) \in \text{reachable } \text{qmsg } TT$ " ..

with qtr have " $m \in \text{set } ms$ "

by (auto dest!: step\_invariantD [OF qmsg\_send\_from\_queue])

with  $\langle \forall m \in \text{set } ms. R \sigma m \rangle$  have " $R \sigma m$ " ..

hence " $\text{orecvmsg } R \sigma (\text{receive } m)$ " by simp

with  $\langle \forall j. j \neq i \rightarrow S (\sigma j) (\sigma' j) \rangle$  have " $\text{?owS } \sigma \sigma' (\text{receive } m)$ "

by (auto intro!: otherwithI)

with pinv por ptr have " $U (\sigma i) (\sigma' i)$ "

by (auto dest!: ostep\_invariantD)

with  $\langle \forall j. j \neq i \rightarrow U (\sigma j) (\sigma' j) \rangle$  have " $\forall j. U (\sigma j) (\sigma' j)$ " by auto

hence  $\text{recvmsg}'$ : " $\bigwedge a. \text{recvmsg (R } \sigma) a \implies \text{recvmsg (R } \sigma') a$ "

by (auto elim!: recvmsgE [where R=R] upreservesq)

from por ptr have " $(\sigma', p') \in \text{oreachable } A \text{ ?owS (other } U \{i\})$ "

using  $\langle \text{?owS } \sigma \sigma' (\text{receive } m) \rangle$  by - (erule(1) oreachable\_local, simp)

moreover have " $(ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma'))$ "

proof (rule reachable\_weakenE [where P="recvmsg (R  $\sigma$ )"])

have " $\text{recvmsg (R } \sigma) (\text{send } m)$ " by simp

with qr qtr show " $(ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma))$ " ..

qed (rule recvmsg')

moreover have " $\forall m \in \text{set } ms'. R \sigma' m$ "

proof

fix m

assume " $m \in \text{set } ms'$ "

moreover have " $\text{set } ms' \subseteq \text{set } ms$ "

proof -

from qr have " $(ms, q) \in \text{reachable } \text{qmsg } TT$ " ..

thus ?thesis using qtr

by (auto dest!: step\_invariantD [OF qmsg\_queue\_contents])

qed

ultimately have " $R \sigma m$ " using  $\langle \forall m \in \text{set } ms. R \sigma m \rangle$  by auto

with  $\langle \forall j. U (\sigma j) (\sigma' j) \rangle$  show " $R \sigma' m$ "

by (rule upreservesq)

qed

ultimately show " $(\sigma', p') \in \text{oreachable } A \text{ ?owS (other } U \{i\})$ "

$\wedge (ms', q') \in \text{reachable } \text{qmsg (recvmsg (R } \sigma'))$

$\wedge (\forall m \in \text{set } ms'. R \sigma' m)$ " by simp

qed

with  $\langle pq = (p, (ms, q)) \rangle$  and  $\langle pq' = (p', (ms', q')) \rangle$  show ?case

by (simp\_all del:  $\Gamma_{QMSG\_simps}$ )

qed

lemma par\_qmsg\_oreachable\_statelessassm:

assumes " $(\sigma, \zeta) \in \text{oreachable } (A \langle \langle_i \text{qmsg} \rangle \rangle$ "

$(\lambda \sigma \_ . \text{orecvmsg } (\lambda \_ . R) \sigma) (\text{other } (\lambda \_ \_ . \text{True}) \{i\})$ "

and ustutter: " $\bigwedge \xi. U \xi \xi$ "

shows " $(\sigma, \text{fst } \zeta) \in \text{oreachable } A (\lambda \sigma \_ . \text{orecvmsg } (\lambda \_ . R) \sigma) (\text{other } (\lambda \_ \_ . \text{True}) \{i\})$ "

$\wedge \text{snd } \zeta \in \text{reachable\_qmsg } (\text{recvmsg } R)$   
 $\wedge (\forall m \in \text{set } (\text{fst } (\text{snd } \zeta))). R m$ "

proof -

from *assms*(1)

have " $(\sigma, \zeta) \in \text{oreachable } (A \langle\langle_i \text{qmsg}\rangle\rangle$   
 $(\text{otherwith } (\lambda\_ \_ . \text{True}) \{i\} (\text{orecvmsg } (\lambda\_ . R))))$   
 $(\text{other } (\lambda\_ \_ . \text{True}) \{i\})$ " by auto

moreover

have " $A \models_A (\text{otherwith } (\lambda\_ \_ . \text{True}) \{i\} (\text{orecvmsg } (\lambda\_ . R))),$   
 $\text{other } (\lambda\_ \_ . \text{True}) \{i\} \rightarrow \text{globala } (\lambda(\sigma, \_, \sigma') . \text{True})$ "

by auto

ultimately

obtain " $(\sigma, \text{fst } \zeta) \in \text{oreachable } A$   
 $(\text{otherwith } (\lambda\_ \_ . \text{True}) \{i\} (\text{orecvmsg } (\lambda\_ . R))) (\text{other } (\lambda\_ \_ . \text{True}) \{i\})$ "  
and \*: " $\text{snd } \zeta \in \text{reachable\_qmsg } (\text{recvmsg } R)$ "  
and \*\*: " $(\forall m \in \text{set } (\text{fst } (\text{snd } \zeta))). R m$ "  
by (auto dest!: *par\_qmsg\_oreachable*)

from *this*(1)

have " $(\sigma, \text{fst } \zeta) \in \text{oreachable } A (\lambda\sigma \_ . \text{orecvmsg } (\lambda\_ . R) \sigma) (\text{other } (\lambda\_ \_ . \text{True}) \{i\})$ "  
by rule auto

thus ?thesis using \* \*\* by simp

qed

lemma *lift\_into\_qmsg*:

assumes " $A \models (\text{otherwith } S \{i\} (\text{orecvmsg } R), \text{other } U \{i\} \rightarrow) \text{global } P$ "  
and " $\bigwedge \xi . U \xi \xi$ "  
and " $\bigwedge \xi \xi' . S \xi \xi' \implies U \xi \xi'$ "  
and " $\bigwedge \sigma \sigma' m . \llbracket \forall j . U (\sigma j) (\sigma' j); R \sigma m \rrbracket \implies R \sigma' m$ "  
and " $A \models_A (\text{otherwith } S \{i\} (\text{orecvmsg } R), \text{other } U \{i\} \rightarrow)$   
 $\text{globala } (\lambda(\sigma, \_, \sigma') . U (\sigma i) (\sigma' i))$ "

shows " $A \langle\langle_i \text{qmsg}\rangle\rangle \models (\text{otherwith } S \{i\} (\text{orecvmsg } R), \text{other } U \{i\} \rightarrow) \text{global } P$ "

proof (rule *oinvariant\_oreachableI*)

fix  $\sigma \zeta$

assume " $(\sigma, \zeta) \in \text{oreachable } (A \langle\langle_i \text{qmsg}\rangle\rangle (\text{otherwith } S \{i\} (\text{orecvmsg } R)) (\text{other } U \{i\}))$ "

then obtain *s* where " $(\sigma, s) \in \text{oreachable } A (\text{otherwith } S \{i\} (\text{orecvmsg } R)) (\text{other } U \{i\})$ "

by (auto dest!: *par\_qmsg\_oreachable [OF \_ assms(5,2-4)]*)

with *assms*(1) show " $\text{global } P (\sigma, \zeta)$ "

by (auto dest: *oinvariant\_weakenD [OF assms(1)]*)

qed

lemma *lift\_step\_into\_qmsg*:

assumes *inv*: " $A \models_A (\text{otherwith } S \{i\} (\text{orecvmsg } R), \text{other } U \{i\} \rightarrow) \text{globala } P$ "  
and *ustutter*: " $\bigwedge \xi . U \xi \xi$ "  
and *sgivesu*: " $\bigwedge \xi \xi' . S \xi \xi' \implies U \xi \xi'$ "  
and *upreservesq*: " $\bigwedge \sigma \sigma' m . \llbracket \forall j . U (\sigma j) (\sigma' j); R \sigma m \rrbracket \implies R \sigma' m$ "  
and *self\_sync*: " $A \models_A (\text{otherwith } S \{i\} (\text{orecvmsg } R), \text{other } U \{i\} \rightarrow)$   
 $\text{globala } (\lambda(\sigma, \_, \sigma') . U (\sigma i) (\sigma' i))$ "

and *recv\_stutter*: " $\bigwedge \sigma \sigma' m . \llbracket \forall j . U (\sigma j) (\sigma' j); \sigma' i = \sigma i \rrbracket \implies P (\sigma, \text{receive } m, \sigma')$ "

and *receive\_right*: " $\bigwedge \sigma \sigma' m . P (\sigma, \text{receive } m, \sigma') \implies P (\sigma, \tau, \sigma')$ "

shows " $A \langle\langle_i \text{qmsg}\rangle\rangle \models_A (\text{otherwith } S \{i\} (\text{orecvmsg } R), \text{other } U \{i\} \rightarrow) \text{globala } P$ "

(is " $\_ \models_A (?owS, ?U \rightarrow) \_$ ")

proof (rule *ostep\_invariantI*)

fix  $\sigma \zeta a \sigma' \zeta'$

assume *or*: " $(\sigma, \zeta) \in \text{oreachable } (A \langle\langle_i \text{qmsg}\rangle\rangle ?owS ?U$ "

and *otr*: " $((\sigma, \zeta), a, (\sigma', \zeta')) \in \text{trans } (A \langle\langle_i \text{qmsg}\rangle\rangle)$ "

and " $?owS \sigma \sigma' a$ "

from *this*(2) have " $((\sigma, \zeta), a, (\sigma', \zeta')) \in \text{oparp\_sos } i (\text{trans } A) (\text{seqp\_sos } \Gamma_{QMSG})$ "

by simp

then obtain *s msgs q s' msgs' q'*

where " $\zeta = (s, (\text{msgs}, q))$ " " $\zeta' = (s', (\text{msgs}', q'))$ "

and " $((\sigma, (s, (\text{msgs}, q))), a, (\sigma', (s', (\text{msgs}', q'))))$

$\in \text{oparp\_sos } i (\text{trans } A) (\text{seqp\_sos } \Gamma_{QMSG})$ "

by (*metis prod\_cases3*)

```

from this(1-2) and or
  obtain " $(\sigma, s) \in \text{oreachable } A \text{ ?owS ?U}$ "
    " $(\text{msgs}, q) \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \ \sigma))$ "
    " $(\forall m \in \text{set } \text{msgs}. R \ \sigma \ m)$ "
  by (auto dest: par_qmsg_oreachable [OF _ self_sync ustutter sgivesu]
      elim!: upreservesq)
from otr  $\langle \zeta = (s, (\text{msgs}, q)) \rangle \langle \zeta' = (s', (\text{msgs}', q')) \rangle$ 
  have " $((\sigma, (s, (\text{msgs}, q))), a, (\sigma', (s', (\text{msgs}', q'))))$ "
     $\in \text{oparp\_sos } i \ (\text{trans } A) \ (\text{seqp\_sos } \Gamma_{QMSG})$ "
  by simp
hence "globala P  $((\sigma, s), a, (\sigma', s'))$ "
proof
  assume " $((\sigma, s), a, (\sigma', s')) \in \text{trans } A$ "
  with  $\langle (\sigma, s) \in \text{oreachable } A \text{ ?owS ?U} \rangle$ 
  show "globala P  $((\sigma, s), a, (\sigma', s'))$ "
    using  $\langle \text{?owS } \sigma \ \sigma' \ a \rangle$  by (rule ostep_invariantD [OF inv])
next
  assume " $((\text{msgs}, q), a, (\text{msgs}', q')) \in \text{seqp\_sos } \Gamma_{QMSG}$ "
  and " $\bigwedge m. a \neq \text{send } m$ "
  and " $\sigma' \ i = \sigma \ i$ "
  from this(3) and ustutter have " $U \ (\sigma \ i) \ (\sigma' \ i)$ " by simp
  with  $\langle \text{?owS } \sigma \ \sigma' \ a \rangle$  and sgivesu have " $\forall j. U \ (\sigma \ j) \ (\sigma' \ j)$ "
  by (clarsimp dest!: otherwith_syncD) metis
  moreover have " $(\exists m. a = \text{receive } m) \vee (a = \tau)$ "
  proof -
    from  $\langle (\text{msgs}, q) \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \ \sigma)) \rangle$ 
    have " $(\text{msgs}, q) \in \text{reachable } \text{qmsg } TT$ " ..
    moreover from  $\langle ((\text{msgs}, q), a, (\text{msgs}', q')) \in \text{seqp\_sos } \Gamma_{QMSG} \rangle$ 
    have " $((\text{msgs}, q), a, (\text{msgs}', q')) \in \text{trans } \text{qmsg}$ " by simp
    ultimately show ?thesis
      using  $\langle \bigwedge m. a \neq \text{send } m \rangle$ 
      by (auto dest!: step_invariantD [OF qmsg_send_receive_or_tau])
  qed
  ultimately show "globala P  $((\sigma, s), a, (\sigma', s'))$ "
    using  $\langle \sigma' \ i = \sigma \ i \rangle$ 
    by simp (metis receive_right recv_stutter step_seq_tau)
next
  fix m
  assume "a =  $\tau$ "
  and " $((\sigma, s), \text{receive } m, (\sigma', s')) \in \text{trans } A$ "
  and " $((\text{msgs}, q), \text{send } m, (\text{msgs}', q')) \in \text{seqp\_sos } \Gamma_{QMSG}$ "

  from  $\langle (\text{msgs}, q) \in \text{reachable } \text{qmsg } (\text{recvmsg } (R \ \sigma)) \rangle$ 
  have " $(\text{msgs}, q) \in \text{reachable } \text{qmsg } TT$ " ..
  moreover from  $\langle ((\text{msgs}, q), \text{send } m, (\text{msgs}', q')) \in \text{seqp\_sos } \Gamma_{QMSG} \rangle$ 
  have " $((\text{msgs}, q), \text{send } m, (\text{msgs}', q')) \in \text{trans } \text{qmsg}$ " by simp
  ultimately have " $m \in \text{set } \text{msgs}$ "
  by (auto dest!: step_invariantD [OF qmsg_send_from_queue])

  with  $\langle \forall m \in \text{set } \text{msgs}. R \ \sigma \ m \rangle$  have " $R \ \sigma \ m$ " ..
  with  $\langle \text{?owS } \sigma \ \sigma' \ a \rangle$  have " $\text{?owS } \sigma \ \sigma' \ (\text{receive } m)$ "
  by (auto dest!: otherwith_syncD)

  with  $\langle ((\sigma, s), \text{receive } m, (\sigma', s')) \in \text{trans } A \rangle$ 
  have "globala P  $((\sigma, s), \text{receive } m, (\sigma', s'))$ "
  using  $\langle (\sigma, s) \in \text{oreachable } A \text{ ?owS ?U} \rangle$ 
  by - (rule ostep_invariantD [OF inv])
  hence "P  $(\sigma, \text{receive } m, \sigma')$ " by simp
  hence "P  $(\sigma, \tau, \sigma')$ " by (rule receive_right)
  with  $\langle a = \tau \rangle$  show "globala P  $((\sigma, s), a, (\sigma', s'))$ " by simp
qed
with  $\langle \zeta = (s, (\text{msgs}, q)) \rangle$  and  $\langle \zeta' = (s', (\text{msgs}', q')) \rangle$  show "globala P  $((\sigma, \zeta), a, (\sigma', \zeta'))$ "
  by simp
qed

```

```

lemma lift_step_into_qmsg_statelessassm:
  assumes "A  $\models_A (\lambda\sigma \_ . \text{orecvmsg } (\lambda\_ . R) \sigma, \text{other } (\lambda\_ \_ . \text{True}) \{i\} \rightarrow) \text{globala } P$ "
    and " $\bigwedge\sigma \sigma' m. \sigma' i = \sigma i \implies P (\sigma, \text{receive } m, \sigma')$ "
    and " $\bigwedge\sigma \sigma' m. P (\sigma, \text{receive } m, \sigma') \implies P (\sigma, \tau, \sigma')$ "
  shows "A  $\langle\langle_i \text{qmsg} \models_A (\lambda\sigma \_ . \text{orecvmsg } (\lambda\_ . R) \sigma, \text{other } (\lambda\_ \_ . \text{True}) \{i\} \rightarrow) \text{globala } P$ "
proof -
  from assms(1) have *: "A  $\models_A (\text{otherwith } (\lambda\_ \_ . \text{True}) \{i\} (\text{orecvmsg } (\lambda\_ . R)),$ 
    other  $(\lambda\_ \_ . \text{True}) \{i\} \rightarrow) \text{globala } P$ "
  by rule auto
  hence "A  $\langle\langle_i \text{qmsg} \models_A$ 
    otherwith  $(\lambda\_ \_ . \text{True}) \{i\} (\text{orecvmsg } (\lambda\_ . R)), \text{other } (\lambda\_ \_ . \text{True}) \{i\} \rightarrow) \text{globala } P$ "
  by (rule lift_step_into_qmsg)
  (auto elim!: assms(2-3) simp del: step_seq_tau)
  thus ?thesis by rule auto
qed

```

end

## 24 Transfer open results onto closed models

```

theory OClosed_Transfer
imports Closed OClosed_Lifting
begin

```

```

locale openproc =
  fixes np :: "'ip  $\Rightarrow$  ('s, ('m::msg) seq_action) automaton"
    and onp :: "'ip  $\Rightarrow$  ((ip  $\Rightarrow$  'g)  $\times$  'l, 'm seq_action) automaton"
    and sr :: "'s  $\Rightarrow$  ('g  $\times$  'l)"
  assumes init: "{ $(\sigma, \zeta) \mid \sigma \zeta s. s \in \text{init } (np \ i)$ 
     $\wedge (\sigma \ i, \zeta) = \text{sr } s$ 
     $\wedge (\forall j. j \neq i \longrightarrow \sigma \ j \in (\text{fst } o \ \text{sr}) \ ' \ \text{init } (np \ j)) \}$   $\subseteq \text{init } (onp \ i)$ "
    and init_notempty: " $\forall j. \text{init } (np \ j) \neq \{\}$ "
    and trans: " $\bigwedge s \ a \ s' \ \sigma \ \sigma'. \llbracket \sigma \ i = \text{fst } (\text{sr } s);$ 
     $\sigma' \ i = \text{fst } (\text{sr } s');$ 
     $(s, a, s') \in \text{trans } (np \ i) \rrbracket$ 
     $\implies ((\sigma, \text{snd } (\text{sr } s)), a, (\sigma', \text{snd } (\text{sr } s')))) \in \text{trans } (onp \ i)$ "
begin

```

```

lemma init_pnet_p_NodeS:
  assumes "NodeS i s R  $\in \text{init } (\text{pnet } np \ p)$ "
  shows "p =  $\langle i; R \rangle$ "
  using assms by (cases p) (auto simp add: node_comps)

```

```

lemma init_pnet_p_SubnetS:
  assumes "SubnetS s1 s2  $\in \text{init } (\text{pnet } np \ p)$ "
  obtains p1 p2 where "p = (p1  $\parallel$  p2)"
    and "s1  $\in \text{init } (\text{pnet } np \ p1)$ "
    and "s2  $\in \text{init } (\text{pnet } np \ p2)$ "
  using assms by (cases p) (auto simp add: node_comps)

```

```

lemma init_pnet_fst_sr_netgmap:
  assumes "s  $\in \text{init } (\text{pnet } np \ p)$ "
    and "i  $\in \text{net\_ips } s$ "
    and "wf_net_tree p"
  shows "the (fst (netgmap sr s) i)  $\in (\text{fst } o \ \text{sr}) \ ' \ \text{init } (np \ i)$ "
  using assms proof (induction s arbitrary: p)
  fix ii s Ri p
  assume "NodeS ii s Ri  $\in \text{init } (\text{pnet } np \ p)$ "
    and "i  $\in \text{net\_ips } (\text{NodeS } ii \ s \ R_i)$ "
    and "wf_net_tree p"
  note this(1)
  moreover then have "p =  $\langle ii; R_i \rangle$ "
  by (rule init_pnet_p_NodeS)

```

```

ultimately have "s ∈ init (np ii)"
  by (clarsimp simp: node_comps)
with ⟨i ∈ net_ips (NodeS ii s R_i)⟩
  show "the (fst (netgmap sr (NodeS ii s R_i)) i) ∈ (fst ∘ sr) ‘ init (np i)'"
  by clarsimp
next
fix s1 s2 p
assume IH1: "∧p. s1 ∈ init (pnet np p)
  ⇒ i ∈ net_ips s1
  ⇒ wf_net_tree p
  ⇒ the (fst (netgmap sr s1) i) ∈ (fst ∘ sr) ‘ init (np i)'"
and IH2: "∧p. s2 ∈ init (pnet np p)
  ⇒ i ∈ net_ips s2
  ⇒ wf_net_tree p
  ⇒ the (fst (netgmap sr s2) i) ∈ (fst ∘ sr) ‘ init (np i)'"
and "SubnetS s1 s2 ∈ init (pnet np p)"
and "i ∈ net_ips (SubnetS s1 s2)"
and "wf_net_tree p"
from this(3) obtain p1 p2 where "p = (p1 || p2)"
  and "s1 ∈ init (pnet np p1)"
  and "s2 ∈ init (pnet np p2)"
  by (rule init_pnet_p_SubnetS)
from this(1) and ⟨wf_net_tree p⟩ have "wf_net_tree p1"
  and "wf_net_tree p2"
  and "net_tree_ips p1 ∩ net_tree_ips p2 = {}"

  by auto
from ⟨i ∈ net_ips (SubnetS s1 s2)⟩ have "i ∈ net_ips s1 ∨ i ∈ net_ips s2"
  by simp
thus "the (fst (netgmap sr (SubnetS s1 s2)) i) ∈ (fst ∘ sr) ‘ init (np i)'"
proof
  assume "i ∈ net_ips s1"
  hence "i ∉ net_ips s2"
  proof -
    from ⟨s1 ∈ init (pnet np p1)⟩ and ⟨i ∈ net_ips s1⟩ have "i ∈ net_tree_ips p1" ..
    with ⟨net_tree_ips p1 ∩ net_tree_ips p2 = {}⟩ have "i ∉ net_tree_ips p2" by auto
    with ⟨s2 ∈ init (pnet np p2)⟩ show ?thesis ..
  qed
  moreover from ⟨s1 ∈ init (pnet np p1)⟩ ⟨i ∈ net_ips s1⟩ and ⟨wf_net_tree p1⟩
    have "the (fst (netgmap sr s1) i) ∈ (fst ∘ sr) ‘ init (np i)'"
    by (rule IH1)
  ultimately show ?thesis by simp
next
  assume "i ∈ net_ips s2"
  moreover with ⟨s2 ∈ init (pnet np p2)⟩ have "the (fst (netgmap sr s2) i) ∈ (fst ∘ sr) ‘ init (np
i)'"
    using ⟨wf_net_tree p2⟩ by (rule IH2)
  moreover from ⟨s2 ∈ init (pnet np p2)⟩ and ⟨i ∈ net_ips s2⟩ have "i ∈ net_tree_ips p2" ..
  ultimately show ?thesis by simp
qed
qed

lemma init_lifted:
  assumes "wf_net_tree p"
  shows "{(σ, snd (netgmap sr s)) | σ s. s ∈ init (pnet np s)}
    ∧ (∀i. if i ∈ net_tree_ips p then σ i = the (fst (netgmap sr s) i)
      else σ i ∈ (fst ∘ sr) ‘ init (np i)) } ⊆ init (opnet onp p)"
  using assms proof (induction p)
    fix i R
    assume "wf_net_tree ⟨i; R⟩"
    show "{(σ, snd (netgmap sr s)) | σ s. s ∈ init (pnet np ⟨i; R⟩)}
      ∧ (∀j. if j ∈ net_tree_ips ⟨i; R⟩ then σ j = the (fst (netgmap sr s) j)
        else σ j ∈ (fst ∘ sr) ‘ init (np j)) } ⊆ init (opnet onp ⟨i; R⟩)"
    by (clarsimp simp add: node_comps onode_comps)
    (rule set_mp [OF init], auto)
  qed

```



```

next
  fix p1 p2
  assume IH1: "wf_net_tree p1
    ⇒ {(σ, snd (netgmap sr s)) | σ s. s ∈ init (pnet np p1)
      ∧ (∀i. if i ∈ net_tree_ips p1 then σ i = the (fst (netgmap sr s) i)
        else σ i ∈ (fst ∘ sr) ‘ init (np i))} ⊆ init (opnet onp p1)"
    (is "_ ⇒ ?S1 ⊆ _)"
  and IH2: "wf_net_tree p2
    ⇒ {(σ, snd (netgmap sr s)) | σ s. s ∈ init (pnet np p2)
      ∧ (∀i. if i ∈ net_tree_ips p2 then σ i = the (fst (netgmap sr s) i)
        else σ i ∈ (fst ∘ sr) ‘ init (np i))} ⊆ init (opnet onp p2)"
    (is "_ ⇒ ?S2 ⊆ _)"
  and "wf_net_tree (p1 || p2)"
from this(3) have "wf_net_tree p1"
  and "wf_net_tree p2"
  and "net_tree_ips p1 ∩ net_tree_ips p2 = {}" by auto
show "{(σ, snd (netgmap sr s)) | σ s. s ∈ init (pnet np (p1 || p2))
  ∧ (∀i. if i ∈ net_tree_ips (p1 || p2) then σ i = the (fst (netgmap sr s) i)
    else σ i ∈ (fst ∘ sr) ‘ init (np i))} ⊆ init (opnet onp (p1 || p2))"
proof (rule, clarsimp simp only: split_paired_all pnet.simps automaton.simps)
  fix σ s1 s2
  assume σ_desc: "∀i. if i ∈ net_tree_ips (p1 || p2)
    then σ i = the (fst (netgmap sr (SubnetS s1 s2)) i)
    else σ i ∈ (fst ∘ sr) ‘ init (np i)"
  and "s1 ∈ init (pnet np p1)"
  and "s2 ∈ init (pnet np p2)"
from this(2-3) have "net_ips s1 = net_tree_ips p1"
  and "net_ips s2 = net_tree_ips p2" by auto
have "(σ, snd (netgmap sr s1)) ∈ ?S1"
proof -
  { fix i
    assume "i ∈ net_tree_ips p1"
    with (net_tree_ips p1 ∩ net_tree_ips p2 = {}) have "i ∉ net_tree_ips p2" by auto
    with (s2 ∈ init (pnet np p2)) have "i ∉ net_ips s2" ..
    hence "the ((fst (netgmap sr s1) ++ fst (netgmap sr s2)) i) = the (fst (netgmap sr s1) i)"
      by simp
  }
  moreover
  { fix i
    assume "i ∉ net_tree_ips p1"
    have "σ i ∈ (fst ∘ sr) ‘ init (np i)"
    proof (cases "i ∈ net_tree_ips p2")
      assume "i ∉ net_tree_ips p2"
      with (i ∉ net_tree_ips p1) and σ_desc show ?thesis
        by simp
    next
      assume "i ∈ net_tree_ips p2"
      with (s2 ∈ init (pnet np p2)) have "i ∈ net_ips s2" ..
      with (s2 ∈ init (pnet np p2)) have "the (fst (netgmap sr s2) i) ∈ (fst ∘ sr) ‘ init (np i)"
        using (wf_net_tree p2) by (rule init_pnet_fst_sr_netgmap)
      with (i ∈ net_tree_ips p2) and (i ∈ net_ips s2) show ?thesis
        using σ_desc by simp
    qed
  }
  ultimately show ?thesis
    using (s1 ∈ init (pnet np p1)) and σ_desc by auto
qed
hence "(σ, snd (netgmap sr s1)) ∈ init (opnet onp p1)"
  by (rule set_mp [OF IH1 [OF (wf_net_tree p1)]])

have "(σ, snd (netgmap sr s2)) ∈ ?S2"
proof -
  { fix i
    assume "i ∈ net_tree_ips p2"

```

```

with ⟨s2 ∈ init (pnet np p2)⟩ have "i ∈ net_ips s2" ..
hence "the ((fst (netgmap sr s1) ++ fst (netgmap sr s2)) i) = the (fst (netgmap sr s2) i)"
  by simp
}
moreover
{ fix i
  assume "i ∉ net_tree_ips p2"
  have "σ i ∈ (fst o sr) ' init (np i)"
  proof (cases "i ∈ net_tree_ips p1")
    assume "i ∉ net_tree_ips p1"
    with ⟨i ∉ net_tree_ips p2⟩ and σ_desc show ?thesis
      by simp
  next
    assume "i ∈ net_tree_ips p1"
    with ⟨s1 ∈ init (pnet np p1)⟩ have "i ∈ net_ips s1" ..
    with ⟨s1 ∈ init (pnet np p1)⟩ have "the (fst (netgmap sr s1) i) ∈ (fst o sr) ' init (np i)"
      using ⟨wf_net_tree p1⟩ by (rule init_pnet_fst_sr_netgmap)
    moreover from ⟨s2 ∈ init (pnet np p2)⟩ and ⟨i ∉ net_tree_ips p2⟩ have "i ∉ net_ips s2" ..
    ultimately show ?thesis
      using ⟨i ∈ net_tree_ips p1⟩ ⟨i ∈ net_ips s1⟩ and ⟨i ∉ net_tree_ips p2⟩ σ_desc by simp
  qed
}
ultimately show ?thesis
  using ⟨s2 ∈ init (pnet np p2)⟩ and σ_desc by auto
qed
hence "(σ, snd (netgmap sr s2)) ∈ init (opnet onp p2)"
  by (rule set_mp [OF IH2 [OF ⟨wf_net_tree p2⟩]])

with ⟨(σ, snd (netgmap sr s1)) ∈ init (opnet onp p1)⟩
show "(σ, snd (netgmap sr (SubnetS s1 s2))) ∈ init (opnet onp (p1 || p2))"
  using ⟨net_tree_ips p1 ∩ net_tree_ips p2 = {}⟩
    ⟨net_ips s1 = net_tree_ips p1⟩
    ⟨net_ips s2 = net_tree_ips p2⟩ by simp
qed
qed

```

lemma init\_pnet\_opnet [elim]:

```

assumes "wf_net_tree p"
  and "s ∈ init (pnet np p)"
shows "netgmap sr s ∈ netmask (net_tree_ips p) ' init (opnet onp p)"
proof -
  from ⟨wf_net_tree p⟩
  have "{ (σ, snd (netgmap sr s)) | σ s. s ∈ init (pnet np p)
    ∧ (∀ i. if i ∈ net_tree_ips p then σ i = the (fst (netgmap sr s) i)
      else σ i ∈ (fst o sr) ' init (np i)) } ⊆ init (opnet onp p)"
    (is "?S ⊆ _")
  by (rule init_lifted)
hence "netmask (net_tree_ips p) ' ?S ⊆ netmask (net_tree_ips p) ' init (opnet onp p)"
  by (rule image_mono)
moreover have "netgmap sr s ∈ netmask (net_tree_ips p) ' ?S"
proof -
  { fix i
    from init_notempty have "∃ s. s ∈ (fst o sr) ' init (np i)" by auto
    hence "(SOME x. x ∈ (fst o sr) ' init (np i)) ∈ (fst o sr) ' init (np i)" ..
  }
  with ⟨s ∈ init (pnet np p)⟩ and init_notempty
  have "(λ i. if i ∈ net_tree_ips p
    then the (fst (netgmap sr s) i)
    else SOME x. x ∈ (fst o sr) ' init (np i), snd (netgmap sr s)) ∈ ?S"
    (is "?s ∈ ?S") by auto
moreover have "netgmap sr s = netmask (net_tree_ips p) ?s"
proof (intro prod_eqI ext)
  fix i
  show "fst (netgmap sr s) i = fst (netmask (net_tree_ips p) ?s) i"

```

```

proof (cases "i ∈ net_tree_ips p")
  assume "i ∈ net_tree_ips p"
  with ⟨s ∈ init (pnet np p)⟩ have "i ∈ net_ips s" ..
  hence "Some (the (fst (netgmap sr s) i)) = fst (netgmap sr s) i"
    by (rule some_the_fst_netgmap)
  with ⟨i ∈ net_tree_ips p⟩ show ?thesis
    by simp
next
  assume "i ∉ net_tree_ips p"
  moreover with ⟨s ∈ init (pnet np p)⟩ have "i ∉ net_ips s" ..
  ultimately show ?thesis
    by simp
qed
qed simp
ultimately show ?thesis
  by (rule rev_image_eqI)
qed
ultimately show ?thesis
  by (rule set_rev_mp [rotated])
qed

lemma transfer_connect:
  assumes "(s, connect(i, i'), s') ∈ trans (pnet np n)"
  and "s ∈ reachable (pnet np n) TT"
  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "wf_net_tree n"
  obtains σ' ζ' where "((σ, ζ), connect(i, i'), (σ', ζ')) ∈ trans (opnet onp n)"
  and "∀j. j ∉ net_ips ζ → σ' j = σ j"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ')"
proof atomize_elim
  from assms have "((σ, snd (netgmap sr s)), connect(i, i'), (σ, snd (netgmap sr s'))) ∈ trans (opnet
onp n)
  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ, snd (netgmap sr s'))"
proof (induction n arbitrary: s s' ζ)
  fix ii Ri ns ns' ζ
  assume "(ns, connect(i, i'), ns') ∈ trans (pnet np ⟨ii; Ri⟩)"
  and "netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)"
  from this(1) have "(ns, connect(i, i'), ns') ∈ node_sos (trans (np ii))"
  by (simp add: node_comps)
  moreover then obtain ni s s' R R' where "ns = NodeS ni s R"
  and "ns' = NodeS ni s' R'" ..
  ultimately have "(NodeS ni s R, connect(i, i'), NodeS ni s' R') ∈ node_sos (trans (np ii))"
  by simp
  moreover then have "s' = s" by auto
  ultimately have "((σ, NodeS ni (snd (sr s)) R), connect(i, i'), (σ, NodeS ni (snd (sr s)) R'))
  ∈ onode_sos (trans (onp ii))"
  by - (rule node_connectTE', auto intro!: onode_sos.intros [simplified])
  with ⟨ns = NodeS ni s R⟩ ⟨ns' = NodeS ni s' R'⟩ ⟨s' = s⟩
  and ⟨netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)⟩
  show "((σ, snd (netgmap sr ns)), connect(i, i'), (σ, snd (netgmap sr ns'))) ∈ trans (opnet onp
⟨ii; Ri⟩)
  ∧ netgmap sr ns' = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, snd (netgmap sr ns'))"
  by (simp add: onode_comps)
next
  fix n1 n2 s s' ζ
  assume IH1: "∧s s' ζ. (s, connect(i, i'), s') ∈ trans (pnet np n1)
  ⇒ s ∈ reachable (pnet np n1) TT
  ⇒ netgmap sr s = netmask (net_tree_ips n1) (σ, ζ)
  ⇒ wf_net_tree n1
  ⇒ ((σ, snd (netgmap sr s)), connect(i, i'), (σ, snd (netgmap sr s'))) ∈ trans (opne
onp n1)
  ∧ netgmap sr s' = netmask (net_tree_ips n1) (σ, snd (netgmap sr s'))"
  and IH2: "∧s s' ζ. (s, connect(i, i'), s') ∈ trans (pnet np n2)
  ⇒ s ∈ reachable (pnet np n2) TT

```

```

    ⇒ netgmap sr s = netmask (net_tree_ips n2) (σ, ζ)
    ⇒ wf_net_tree n2
    ⇒ ((σ, snd (netgmap sr s)), connect(i, i'), (σ, snd (netgmap sr s'))) ∈ trans (opne
onp n2)

    ∧ netgmap sr s' = netmask (net_tree_ips n2) (σ, snd (netgmap sr s'))"
  and tr: "(s, connect(i, i'), s') ∈ trans (pnet np (n1 || n2))"
  and sr: "s ∈ reachable (pnet np (n1 || n2)) TT"
  and nm: "netgmap sr s = netmask (net_tree_ips (n1 || n2)) (σ, ζ)"
  and "wf_net_tree (n1 || n2)"
from this(3) have "(s, connect(i, i'), s') ∈ pnet_sos (trans (pnet np n1))
                  (trans (pnet np n2))"

  by simp
then obtain s1 s1' s2 s2' where "s = SubnetS s1 s2"
                              and "s' = SubnetS s1' s2'"
                              and "(s1, connect(i, i'), s1') ∈ trans (pnet np n1)"
                              and "(s2, connect(i, i'), s2') ∈ trans (pnet np n2)"

  by (rule partial_connectTE) auto
from this(1) and nm have "netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) (σ, ζ)"
  by simp

from ⟨wf_net_tree (n1 || n2)⟩ have "wf_net_tree n1" and "wf_net_tree n2"
  and "net_tree_ips n1 ∩ net_tree_ips n2 = {}" by auto

from sr ⟨s = SubnetS s1 s2⟩ have "s1 ∈ reachable (pnet np n1) TT" by (metis subnet_reachable(1))
  hence "net_ips s1 = net_tree_ips n1" by (rule pnet_net_ips_net_tree_ips)

from sr ⟨s = SubnetS s1 s2⟩ have "s2 ∈ reachable (pnet np n2) TT" by (metis subnet_reachable(2))
  hence "net_ips s2 = net_tree_ips n2" by (rule pnet_net_ips_net_tree_ips)

from nm ⟨s = SubnetS s1 s2⟩
  have "netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) (σ, ζ)" by simp
  hence "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
    using ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩ ⟨net_ips s1 = net_tree_ips n1⟩
    and ⟨net_ips s2 = net_tree_ips n2⟩ by (rule netgmap_subnet_split1)
with ⟨(s1, connect(i, i'), s1') ∈ trans (pnet np n1)⟩
  and ⟨s1 ∈ reachable (pnet np n1) TT⟩
  have "((σ, snd (netgmap sr s1)), connect(i, i'), (σ, snd (netgmap sr s1')))) ∈ trans (opnet onp
n1)"

  and "netgmap sr s1' = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1'))"
    using ⟨wf_net_tree n1⟩ unfolding atomize_conj by (rule IH1)

from ⟨netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) (σ, ζ)⟩
  ⟨net_ips s1 = net_tree_ips n1⟩ and ⟨net_ips s2 = net_tree_ips n2⟩
  have "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
    by (rule netgmap_subnet_split2)
with ⟨(s2, connect(i, i'), s2') ∈ trans (pnet np n2)⟩
  and ⟨s2 ∈ reachable (pnet np n2) TT⟩
  have "((σ, snd (netgmap sr s2)), connect(i, i'), (σ, snd (netgmap sr s2')))) ∈ trans (opnet onp
n2)"

  and "netgmap sr s2' = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2'))"
    using ⟨wf_net_tree n2⟩ unfolding atomize_conj by (rule IH2)

have "((σ, snd (netgmap sr s)), connect(i, i'), (σ, snd (netgmap sr s'))))
    ∈ trans (opnet onp (n1 || n2))"
proof -
  from ⟨((σ, snd (netgmap sr s1)), connect(i, i'), (σ, snd (netgmap sr s1')))) ∈ trans (opnet onp
n1)⟩
  and ⟨((σ, snd (netgmap sr s2)), connect(i, i'), (σ, snd (netgmap sr s2')))) ∈ trans (opnet onp
n2)⟩
  have "((σ, SubnetS (snd (netgmap sr s1)) (snd (netgmap sr s2))), connect(i, i'),
    (σ, SubnetS (snd (netgmap sr s1')) (snd (netgmap sr s2'))))
    ∈ opnet_sos (trans (opnet onp n1)) (trans (opnet onp n2))"
    by (rule opnet_connect)
  with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ show ?thesis by simp

```

qed

moreover from  $\langle \text{netgmap sr } s1' = \text{netmask (net\_tree\_ips } n1) (\sigma, \text{snd (netgmap sr } s1')) \rangle$   
 $\langle \text{netgmap sr } s2' = \text{netmask (net\_tree\_ips } n2) (\sigma, \text{snd (netgmap sr } s2')) \rangle$   
 $\langle s' = \text{SubnetS } s1' s2' \rangle$   
have "netgmap sr s' = netmask (net\_tree\_ips (n1 || n2)) ( $\sigma$ , snd (netgmap sr s'))" ..

ultimately show " $((\sigma, \text{snd (netgmap sr } s)), \text{connect}(i, i'), (\sigma, \text{snd (netgmap sr } s')))$   
 $\in \text{trans (opnet onp (n1 || n2))}$   
 $\wedge \text{netgmap sr } s' = \text{netmask (net\_tree\_ips (n1 || n2)) (\sigma, \text{snd (netgmap sr } s'))$ " ..

qed

moreover from  $\langle \text{netgmap sr } s = \text{netmask (net\_tree\_ips } n) (\sigma, \zeta) \rangle$  have " $\zeta = \text{snd (netgmap sr } s)$ " by simp  
ultimately show " $\exists \sigma' \zeta'. ((\sigma, \zeta), \text{connect}(i, i'), (\sigma', \zeta')) \in \text{trans (opnet onp } n)$   
 $\wedge (\forall j. j \notin \text{net\_ips } \zeta \rightarrow \sigma' j = \sigma j)$   
 $\wedge \text{netgmap sr } s' = \text{netmask (net\_tree\_ips } n) (\sigma', \zeta')$ " by auto

qed

lemma transfer\_disconnect:

assumes " $(s, \text{disconnect}(i, i'), s') \in \text{trans (pnet np } n)$ "  
and " $s \in \text{reachable (pnet np } n) TT$ "  
and " $\text{netgmap sr } s = \text{netmask (net\_tree\_ips } n) (\sigma, \zeta)$ "  
and " $\text{wf\_net\_tree } n$ "

obtains  $\sigma' \zeta'$  where " $((\sigma, \zeta), \text{disconnect}(i, i'), (\sigma', \zeta')) \in \text{trans (opnet onp } n)$ "  
and " $\forall j. j \notin \text{net\_ips } \zeta \rightarrow \sigma' j = \sigma j$ "  
and " $\text{netgmap sr } s' = \text{netmask (net\_tree\_ips } n) (\sigma', \zeta')$ "

proof atomize\_elim

from assms have " $((\sigma, \text{snd (netgmap sr } s)), \text{disconnect}(i, i'), (\sigma, \text{snd (netgmap sr } s')) \in \text{trans (opnet onp } n)$

$\wedge \text{netgmap sr } s' = \text{netmask (net\_tree\_ips } n) (\sigma, \text{snd (netgmap sr } s'))$ "

proof (induction n arbitrary: s s'  $\zeta$ )

fix ii  $R_i$  ns ns'  $\zeta$

assume " $(ns, \text{disconnect}(i, i'), ns') \in \text{trans (pnet np } \langle ii; R_i \rangle)$ "  
and " $\text{netgmap sr } ns = \text{netmask (net\_tree\_ips } \langle ii; R_i \rangle) (\sigma, \zeta)$ "

from this(1) have " $(ns, \text{disconnect}(i, i'), ns') \in \text{node\_sos (trans (np } ii))$ "  
by (simp add: node\_comps)

moreover then obtain  $ni s' R R'$  where " $ns = \text{NodeS } ni s R$ "

and " $ns' = \text{NodeS } ni s' R'$ " ..

ultimately have " $(\text{NodeS } ni s R, \text{disconnect}(i, i'), \text{NodeS } ni s' R') \in \text{node\_sos (trans (np } ii))$ "

by simp

moreover then have " $s' = s$ " by auto

ultimately have " $((\sigma, \text{NodeS } ni (\text{snd (sr } s)) R), \text{disconnect}(i, i'), (\sigma, \text{NodeS } ni (\text{snd (sr } s)) R')) \in \text{onode\_sos (trans (onp } ii))$ "

by - (rule node\_disconnectTE', auto intro!: onode\_sos.intros [simplified])

with  $\langle ns = \text{NodeS } ni s R \rangle \langle ns' = \text{NodeS } ni s' R' \rangle \langle s' = s \rangle$

and  $\langle \text{netgmap sr } ns = \text{netmask (net\_tree\_ips } \langle ii; R_i \rangle) (\sigma, \zeta) \rangle$

show " $((\sigma, \text{snd (netgmap sr } ns)), \text{disconnect}(i, i'), (\sigma, \text{snd (netgmap sr } ns')) \in \text{trans (opnet onp } \langle ii; R_i \rangle)$

$\wedge \text{netgmap sr } ns' = \text{netmask (net\_tree\_ips } \langle ii; R_i \rangle) (\sigma, \text{snd (netgmap sr } ns'))$ "

by (simp add: onode\_comps)

next

fix n1 n2 s s'  $\zeta$

assume IH1: " $\bigwedge s s' \zeta. (s, \text{disconnect}(i, i'), s') \in \text{trans (pnet np } n1)$

$\Rightarrow s \in \text{reachable (pnet np } n1) TT$

$\Rightarrow \text{netgmap sr } s = \text{netmask (net\_tree\_ips } n1) (\sigma, \zeta)$

$\Rightarrow \text{wf\_net\_tree } n1$

$\Rightarrow ((\sigma, \text{snd (netgmap sr } s)), \text{disconnect}(i, i'), (\sigma, \text{snd (netgmap sr } s')) \in \text{trans$

(opnet onp n1)

$\wedge \text{netgmap sr } s' = \text{netmask (net\_tree\_ips } n1) (\sigma, \text{snd (netgmap sr } s'))$ "

and IH2: " $\bigwedge s s' \zeta. (s, \text{disconnect}(i, i'), s') \in \text{trans (pnet np } n2)$

$\Rightarrow s \in \text{reachable (pnet np } n2) TT$

$\Rightarrow \text{netgmap sr } s = \text{netmask (net\_tree\_ips } n2) (\sigma, \zeta)$

$\Rightarrow \text{wf\_net\_tree } n2$

$\Rightarrow ((\sigma, \text{snd (netgmap sr } s)), \text{disconnect}(i, i'), (\sigma, \text{snd (netgmap sr } s')) \in \text{trans$

(opnet onp n2)

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       $\wedge$  netgmap sr s' = netmask (net_tree_ips n2) ( $\sigma$ , snd (netgmap sr s'))"
    and tr: "(s, disconnect(i, i'), s')  $\in$  trans (pnet np (n1 || n2))"
    and sr: "s  $\in$  reachable (pnet np (n1 || n2)) TT"
    and nm: "netgmap sr s = netmask (net_tree_ips (n1 || n2)) ( $\sigma$ ,  $\zeta$ )"
    and "wf_net_tree (n1 || n2)"
  from this(3) have "(s, disconnect(i, i'), s')  $\in$  pnet_sos (trans (pnet np n1))
                    (trans (pnet np n2))"

  by simp
  then obtain s1 s1' s2 s2' where "s = SubnetS s1 s2"
    and "s' = SubnetS s1' s2'"
    and "(s1, disconnect(i, i'), s1')  $\in$  trans (pnet np n1)"
    and "(s2, disconnect(i, i'), s2')  $\in$  trans (pnet np n2)"

  by (rule partial_disconnectTE) auto
  from this(1) and nm have "netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) ( $\sigma$ ,  $\zeta$ )"
  by simp

  from  $\langle$ wf_net_tree (n1 || n2) $\rangle$  have "wf_net_tree n1" and "wf_net_tree n2"
    and "net_tree_ips n1  $\cap$  net_tree_ips n2 = {}" by auto

  from sr  $\langle$ s = SubnetS s1 s2 $\rangle$  have "s1  $\in$  reachable (pnet np n1) TT" by (metis subnet_reachable(1))
  hence "net_ips s1 = net_tree_ips n1" by (rule pnet_net_ips_net_tree_ips)

  from sr  $\langle$ s = SubnetS s1 s2 $\rangle$  have "s2  $\in$  reachable (pnet np n2) TT" by (metis subnet_reachable(2))
  hence "net_ips s2 = net_tree_ips n2" by (rule pnet_net_ips_net_tree_ips)

  from nm  $\langle$ s = SubnetS s1 s2 $\rangle$ 
    have "netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) ( $\sigma$ ,  $\zeta$ )" by simp
  hence "netgmap sr s1 = netmask (net_tree_ips n1) ( $\sigma$ , snd (netgmap sr s1))"
    using  $\langle$ net_tree_ips n1  $\cap$  net_tree_ips n2 = {} $\rangle$   $\langle$ net_ips s1 = net_tree_ips n1 $\rangle$ 
    and  $\langle$ net_ips s2 = net_tree_ips n2 $\rangle$  by (rule netgmap_subnet_split1)
  with  $\langle$ (s1, disconnect(i, i'), s1')  $\in$  trans (pnet np n1) $\rangle$ 
    and  $\langle$ s1  $\in$  reachable (pnet np n1) TT $\rangle$ 
    have " $\langle$ ( $\sigma$ , snd (netgmap sr s1)), disconnect(i, i'), ( $\sigma$ , snd (netgmap sr s1')) $\rangle$   $\in$  trans (opnet
  onp n1)"
    and "netgmap sr s1' = netmask (net_tree_ips n1) ( $\sigma$ , snd (netgmap sr s1'))"
    using  $\langle$ wf_net_tree n1 $\rangle$  unfolding atomize_conj by (rule IH1)

  from  $\langle$ netgmap sr (SubnetS s1 s2) = netmask (net_tree_ips (n1 || n2)) ( $\sigma$ ,  $\zeta$ ) $\rangle$ 
     $\langle$ net_ips s1 = net_tree_ips n1 $\rangle$  and  $\langle$ net_ips s2 = net_tree_ips n2 $\rangle$ 
    have "netgmap sr s2 = netmask (net_tree_ips n2) ( $\sigma$ , snd (netgmap sr s2))"
    by (rule netgmap_subnet_split2)
  with  $\langle$ (s2, disconnect(i, i'), s2')  $\in$  trans (pnet np n2) $\rangle$ 
    and  $\langle$ s2  $\in$  reachable (pnet np n2) TT $\rangle$ 
    have " $\langle$ ( $\sigma$ , snd (netgmap sr s2)), disconnect(i, i'), ( $\sigma$ , snd (netgmap sr s2')) $\rangle$   $\in$  trans (opnet
  onp n2)"
    and "netgmap sr s2' = netmask (net_tree_ips n2) ( $\sigma$ , snd (netgmap sr s2'))"
    using  $\langle$ wf_net_tree n2 $\rangle$  unfolding atomize_conj by (rule IH2)

  have " $\langle$ ( $\sigma$ , snd (netgmap sr s)), disconnect(i, i'), ( $\sigma$ , snd (netgmap sr s')) $\rangle$ 
         $\in$  trans (opnet onp (n1 || n2))"

  proof -
    from  $\langle$ ( $\sigma$ , snd (netgmap sr s1)), disconnect(i, i'), ( $\sigma$ , snd (netgmap sr s1')) $\rangle$   $\in$  trans (opnet
  onp n1) $\rangle$ 
    and  $\langle$ ( $\sigma$ , snd (netgmap sr s2)), disconnect(i, i'), ( $\sigma$ , snd (netgmap sr s2')) $\rangle$   $\in$  trans (opnet
  onp n2) $\rangle$ 
    have " $\langle$ ( $\sigma$ , SubnetS (snd (netgmap sr s1)) (snd (netgmap sr s2))), disconnect(i, i'),
        ( $\sigma$ , SubnetS (snd (netgmap sr s1')) (snd (netgmap sr s2'))) $\rangle$ 
         $\in$  opnet_sos (trans (opnet onp n1)) (trans (opnet onp n2))"
    by (rule opnet_disconnect)
    with  $\langle$ s = SubnetS s1 s2 $\rangle$   $\langle$ s' = SubnetS s1' s2' $\rangle$  show ?thesis by simp
  qed

  moreover from  $\langle$ netgmap sr s1' = netmask (net_tree_ips n1) ( $\sigma$ , snd (netgmap sr s1')) $\rangle$ 
     $\langle$ netgmap sr s2' = netmask (net_tree_ips n2) ( $\sigma$ , snd (netgmap sr s2')) $\rangle$ 

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      (s' = SubnetS s1' s2')
    have "netgmap sr s' = netmask (net_tree_ips (n1 || n2)) (σ, snd (netgmap sr s'))" ..

    ultimately show "((σ, snd (netgmap sr s)), disconnect(i, i'), (σ, snd (netgmap sr s'))))
      ∈ trans (opnet onp (n1 || n2))
      ∧ netgmap sr s' = netmask (net_tree_ips (n1 || n2)) (σ, snd (netgmap sr s'))" ..
  qed
  moreover from (netgmap sr s = netmask (net_tree_ips n) (σ, ζ)) have "ζ = snd (netgmap sr s)" by simp
  ultimately show "∃σ' ζ'. ((σ, ζ), disconnect(i, i'), (σ', ζ')) ∈ trans (opnet onp n)
    ∧ (∀j. j ∉ net_ips ζ → σ' j = σ j)
    ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'" by auto
  qed

lemma transfer_tau:
  assumes "(s, τ, s') ∈ trans (pnet np n)"
    and "s ∈ reachable (pnet np n) TT"
    and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
    and "wf_net_tree n"
  obtains σ' ζ' where "((σ, ζ), τ, (σ', ζ')) ∈ trans (opnet onp n)"
    and "∀j. j ∉ net_ips ζ → σ' j = σ j"
    and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"

proof atomize_elim
  from assms(4,2,1) obtain i where "i ∈ net_ips s"
    and "∀j. j ≠ i → netmap s' j = netmap s j"
    and "net_ip_action np τ i n s s'"
  by (metis pnet_tau_single_node)
  from this(2) have "∀j. j ≠ i → fst (netgmap sr s') j = fst (netgmap sr s) j"
  by (clarsimp intro!: netmap_is_fst_netgmap')
  from (s, τ, s') ∈ trans (pnet np n) have "net_ips s' = net_ips s"
  by (rule pnet_maintains_dom [THEN sym])
  define σ' where "σ' j = (if j = i then the (fst (netgmap sr s') i) else σ j)" for j
  from (∀j. j ≠ i → fst (netgmap sr s') j = fst (netgmap sr s) j)
    and (netgmap sr s = netmask (net_tree_ips n) (σ, ζ))
  have "∀j. j ≠ i → σ' j = σ j"
  unfolding σ'_def by clarsimp

  from assms(2) have "net_ips s = net_tree_ips n"
  by (rule pnet_net_ips_net_tree_ips)

  from (netgmap sr s = netmask (net_tree_ips n) (σ, ζ))
  have "ζ = snd (netgmap sr s)" by simp

  from (∀j. j ≠ i → fst (netgmap sr s') j = fst (netgmap sr s) j) (i ∈ net_ips s)
    (net_ips s = net_tree_ips n) (net_ips s' = net_ips s)
    (netgmap sr s = netmask (net_tree_ips n) (σ, ζ))
  have "fst (netgmap sr s') = fst (netmask (net_tree_ips n) (σ', snd (netgmap sr s')))"
  unfolding σ'_def [abs_def] by - (rule ext, clarsimp)

  hence "netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"
  by (rule prod_eqI, simp)

  with assms(1, 3)
  have "((σ, snd (netgmap sr s)), τ, (σ', snd (netgmap sr s')))) ∈ trans (opnet onp n)"
  using assms(2,4) (i ∈ net_ips s) and (net_ip_action np τ i n s s')
  proof (induction n arbitrary: s s' ζ)
    fix ii R_i ns ns' ζ
    assume "(ns, τ, ns') ∈ trans (pnet np (ii; R_i))"
      and nsr: "ns ∈ reachable (pnet np (ii; R_i)) TT"
      and "netgmap sr ns = netmask (net_tree_ips (ii; R_i)) (σ, ζ)"
      and "netgmap sr ns' = netmask (net_tree_ips (ii; R_i)) (σ', snd (netgmap sr ns'))"
      and "i ∈ net_ips ns"
    from this(1) have "(ns, τ, ns') ∈ node_sos (trans (np ii))"
    by (simp add: node_comps)
    moreover with nsr obtain s s' R R' where "ns = NodeS ii s R"

```

and "ns' = NodeS ii s' R'"

by (metis net\_node\_reachable\_is\_node node\_tauTE')

moreover from  $\langle i \in \text{net\_ips } ns \rangle$  and  $\langle ns = \text{NodeS } ii \ s \ R \rangle$  have "ii = i" by simp  
ultimately have ntr: " $(\text{NodeS } i \ s \ R, \tau, \text{NodeS } i \ s' \ R') \in \text{node\_sos } (\text{trans } (np \ i))$ "

by simp

hence "R' = R" by (metis net\_state.inject(1) node\_tauTE')

from ntr obtain a where " $(s, a, s') \in \text{trans } (np \ i)$ "

and " $(\exists d. a = \neg \text{unicast } d \wedge d \notin R) \vee (a = \tau)$ "

by (rule node\_tauTE') auto

from  $\langle \text{netgmap } sr \ ns = \text{netmask } (\text{net\_tree\_ips } \langle ii; R_i \rangle) (\sigma, \zeta) \rangle$   $\langle ns = \text{NodeS } ii \ s \ R \rangle$  and  $\langle ii = i \rangle$   
have " $\sigma \ i = \text{fst } (sr \ s)$ " by simp (metis map\_upd\_Some\_unfold)

moreover from  $\langle \text{netgmap } sr \ ns' = \text{netmask } (\text{net\_tree\_ips } \langle ii; R_i \rangle) (\sigma', \text{snd } (\text{netgmap } sr \ ns')) \rangle$   
 $\langle ns' = \text{NodeS } ii \ s' \ R' \rangle$  and  $\langle ii = i \rangle$

have " $\sigma' \ i = \text{fst } (sr \ s')$ "

unfolding  $\sigma'_{\text{def}}$  [abs\_def] by clarsimp (hypsubst\_thin,  
metis (full\_types, lifting) fun\_upd\_same option.sel)

ultimately have " $((\sigma, \text{snd } (sr \ s)), a, (\sigma', \text{snd } (sr \ s')))) \in \text{trans } (\text{onp } i)$ "

using  $\langle (s, a, s') \in \text{trans } (np \ i) \rangle$  by (rule trans)

from  $\langle (\exists d. a = \neg \text{unicast } d \wedge d \notin R) \vee (a = \tau) \rangle \langle \forall j. j \neq i \longrightarrow \sigma' \ j = \sigma \ j \rangle \langle R'=R \rangle$

and  $\langle ((\sigma, \text{snd } (sr \ s)), a, (\sigma', \text{snd } (sr \ s')))) \in \text{trans } (\text{onp } i) \rangle$

have " $((\sigma, \text{NodeS } i \ (\text{snd } (sr \ s)) \ R), \tau, (\sigma', \text{NodeS } i \ (\text{snd } (sr \ s')) \ R')) \in \text{onode\_sos } (\text{trans } (\text{onp } i))$ "

by (metis onode\_sos.onode\_notucast onode\_sos.onode\_tau)

with  $\langle ns = \text{NodeS } ii \ s \ R \rangle$   $\langle ns' = \text{NodeS } ii \ s' \ R' \rangle$   $\langle ii = i \rangle$

show " $((\sigma, \text{snd } (\text{netgmap } sr \ ns)), \tau, (\sigma', \text{snd } (\text{netgmap } sr \ ns')))) \in \text{trans } (\text{opnet } \text{onp } \langle ii; R_i \rangle)$ "

by (simp add: onode\_comps)

next

fix n1 n2 s s'  $\zeta$

assume IH1: " $\bigwedge s \ s' \ \zeta. (s, \tau, s') \in \text{trans } (\text{pnet } np \ n1)$ "

$\implies \text{netgmap } sr \ s = \text{netmask } (\text{net\_tree\_ips } n1) (\sigma, \zeta)$

$\implies \text{netgmap } sr \ s' = \text{netmask } (\text{net\_tree\_ips } n1) (\sigma', \text{snd } (\text{netgmap } sr \ s'))$

$\implies s \in \text{reachable } (\text{pnet } np \ n1) \ \text{TT}$

$\implies \text{wf\_net\_tree } n1$

$\implies i \in \text{net\_ips } s$

$\implies \text{net\_ip\_action } np \ \tau \ i \ n1 \ s \ s'$

$\implies ((\sigma, \text{snd } (\text{netgmap } sr \ s)), \tau, (\sigma', \text{snd } (\text{netgmap } sr \ s')))) \in \text{trans } (\text{opnet } \text{onp } n1)$ "

and IH2: " $\bigwedge s \ s' \ \zeta. (s, \tau, s') \in \text{trans } (\text{pnet } np \ n2)$ "

$\implies \text{netgmap } sr \ s = \text{netmask } (\text{net\_tree\_ips } n2) (\sigma, \zeta)$

$\implies \text{netgmap } sr \ s' = \text{netmask } (\text{net\_tree\_ips } n2) (\sigma', \text{snd } (\text{netgmap } sr \ s'))$

$\implies s \in \text{reachable } (\text{pnet } np \ n2) \ \text{TT}$

$\implies \text{wf\_net\_tree } n2$

$\implies i \in \text{net\_ips } s$

$\implies \text{net\_ip\_action } np \ \tau \ i \ n2 \ s \ s'$

$\implies ((\sigma, \text{snd } (\text{netgmap } sr \ s)), \tau, (\sigma', \text{snd } (\text{netgmap } sr \ s')))) \in \text{trans } (\text{opnet } \text{onp } n2)$ "

and tr: " $(s, \tau, s') \in \text{trans } (\text{pnet } np \ (n1 \parallel n2))$ "

and sr: " $s \in \text{reachable } (\text{pnet } np \ (n1 \parallel n2)) \ \text{TT}$ "

and nm: " $\text{netgmap } sr \ s = \text{netmask } (\text{net\_tree\_ips } (n1 \parallel n2)) (\sigma, \zeta)$ "

and nm': " $\text{netgmap } sr \ s' = \text{netmask } (\text{net\_tree\_ips } (n1 \parallel n2)) (\sigma', \text{snd } (\text{netgmap } sr \ s'))$ "

and " $\text{wf\_net\_tree } (n1 \parallel n2)$ "

and " $i \in \text{net\_ips } s$ "

and " $\text{net\_ip\_action } np \ \tau \ i \ (n1 \parallel n2) \ s \ s'$ "

from tr have " $(s, \tau, s') \in \text{pnet\_sos } (\text{trans } (\text{pnet } np \ n1)) (\text{trans } (\text{pnet } np \ n2))$ " by simp

then obtain s1 s1' s2 s2' where " $s = \text{SubnetS } s1 \ s2$ "

and " $s' = \text{SubnetS } s1' \ s2'$ "

by (rule partial\_tauTE) auto

from this(1) and nm have " $\text{netgmap } sr \ (\text{SubnetS } s1 \ s2) = \text{netmask } (\text{net\_tree\_ips } (n1 \parallel n2)) (\sigma, \zeta)$ "

by simp

from  $\langle s' = \text{SubnetS } s1' \ s2' \rangle$  and nm'

have " $\text{netgmap } sr \ (\text{SubnetS } s1' \ s2') = \text{netmask } (\text{net\_tree\_ips } (n1 \parallel n2)) (\sigma', \text{snd } (\text{netgmap } sr \ s'))$ "



```

by simp

from ⟨wf_net_tree (n1 || n2)⟩ have "wf_net_tree n1"
  and "wf_net_tree n2"
  and "net_tree_ips n1 ∩ net_tree_ips n2 = {}" by auto

from sr [simplified ⟨s = SubnetS s1 s2⟩] have "s1 ∈ reachable (pnet np n1) TT"
  by (rule subnet_reachable(1))
hence "net_ips s1 = net_tree_ips n1" by (rule pnet_net_ips_net_tree_ips)

from sr [simplified ⟨s = SubnetS s1 s2⟩] have "s2 ∈ reachable (pnet np n2) TT"
  by (rule subnet_reachable(2))
hence "net_ips s2 = net_tree_ips n2" by (rule pnet_net_ips_net_tree_ips)

from nm [simplified ⟨s = SubnetS s1 s2⟩]
  ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
  ⟨net_ips s1 = net_tree_ips n1⟩
  ⟨net_ips s2 = net_tree_ips n2⟩
  have "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
  by (rule netgmap_subnet_split1)

from nm [simplified ⟨s = SubnetS s1 s2⟩]
  ⟨net_ips s1 = net_tree_ips n1⟩
  ⟨net_ips s2 = net_tree_ips n2⟩
  have "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
  by (rule netgmap_subnet_split2)

from ⟨i ∈ net_ips s⟩ and ⟨s = SubnetS s1 s2⟩ have "i ∈ net_ips s1 ∨ i ∈ net_ips s2" by auto
thus "((σ, snd (netgmap sr s)), τ, (σ', snd (netgmap sr s'))) ∈ trans (opnet onp (n1 || n2))"
proof
  assume "i ∈ net_ips s1"
  with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨net_ip_action np τ i (n1 || n2) s s'⟩
  have "(s1, τ, s1') ∈ trans (pnet np n1)"
    and "net_ip_action np τ i n1 s1 s1'"
    and "s2' = s2" by simp_all

  from ⟨net_ips s1 = net_tree_ips n1⟩ and ⟨(s1, τ, s1') ∈ trans (pnet np n1)⟩
  have "net_ips s1' = net_tree_ips n1" by (metis pnet_maintains_dom)

  from nm' [simplified ⟨s' = SubnetS s1' s2'⟩ ⟨s2' = s2⟩]
    ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
    ⟨net_ips s1' = net_tree_ips n1⟩
    ⟨net_ips s2 = net_tree_ips n2⟩
  have "netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))"
  by (rule netgmap_subnet_split1)

  from ⟨(s1, τ, s1') ∈ trans (pnet np n1)⟩
    ⟨netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))⟩
    ⟨netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))⟩
    ⟨s1 ∈ reachable (pnet np n1) TT⟩
    ⟨wf_net_tree n1⟩
    ⟨i ∈ net_ips s1⟩
    ⟨net_ip_action np τ i n1 s1 s1'⟩
  have "((σ, snd (netgmap sr s1)), τ, (σ', snd (netgmap sr s1')))) ∈ trans (opnet onp n1)"
  by (rule IH1)

  with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨s2' = s2⟩ show ?thesis
  by (simp del: step_node_tau) (erule opnet_tau1)
next
  assume "i ∈ net_ips s2"
  with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨net_ip_action np τ i (n1 || n2) s s'⟩
  have "(s2, τ, s2') ∈ trans (pnet np n2)"
    and "net_ip_action np τ i n2 s2 s2'"
    and "s1' = s1" by simp_all

```

```

from ⟨net_ips s2 = net_tree_ips n2⟩ and ⟨(s2, τ, s2') ∈ trans (pnet np n2)⟩
  have "net_ips s2' = net_tree_ips n2" by (metis pnet_maintains_dom)

from nm' [simplified ⟨s' = SubnetS s1' s2'⟩ ⟨s1' = s1⟩]
  ⟨net_ips s1 = net_tree_ips n1⟩
  ⟨net_ips s2' = net_tree_ips n2⟩
  have "netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))"
    by (rule netgmap_subnet_split2)

from ⟨(s2, τ, s2') ∈ trans (pnet np n2)⟩
  ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
  ⟨netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))⟩
  ⟨s2 ∈ reachable (pnet np n2) TT⟩
  ⟨wf_net_tree n2⟩
  ⟨i ∈ net_ips s2⟩
  ⟨net_ip_action np τ i n2 s2 s2'⟩
  have "((σ, snd (netgmap sr s2)), τ, (σ', snd (netgmap sr s2')))) ∈ trans (opnet onp n2)"
    by (rule IH2)

```

```

with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨s1' = s1⟩ show ?thesis
  by (simp del: step_node_tau) (erule opnet_tau2)

```

qed

qed

```

with ⟨ζ = snd (netgmap sr s)⟩ have "((σ, ζ), τ, (σ', snd (netgmap sr s')))) ∈ trans (opnet onp n)"
  by simp

```

```

moreover from ⟨∀j. j ≠ i → σ' j = σ j⟩ ⟨i ∈ net_ips s⟩ ⟨ζ = snd (netgmap sr s)⟩

```

```

  have "∀j. j ∉ net_ips ζ → σ' j = σ j" by (metis net_ips_netgmap)

```

```

ultimately have "((σ, ζ), τ, (σ', snd (netgmap sr s')))) ∈ trans (opnet onp n)

```

```

  ∧ (∀j. j ∉ net_ips ζ → σ' j = σ j)

```

```

  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"

```

```

  using ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))⟩ by simp

```

```

thus "∃σ' ζ'. ((σ, ζ), τ, (σ', ζ')) ∈ trans (opnet onp n)

```

```

  ∧ (∀j. j ∉ net_ips ζ → σ' j = σ j)

```

```

  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'" by auto

```

qed

lemma transfer\_deliver:

```

assumes "(s, i:deliver(d), s') ∈ trans (pnet np n)"

```

```

  and "s ∈ reachable (pnet np n) TT"

```

```

  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"

```

```

  and "wf_net_tree n"

```

```

obtains σ' ζ' where "((σ, ζ), i:deliver(d), (σ', ζ')) ∈ trans (opnet onp n)"

```

```

  and "∀j. j ∉ net_ips ζ → σ' j = σ j"

```

```

  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"

```

proof atomize\_elim

```

from assms(4,2,1) obtain "i ∈ net_ips s"

```

```

  and "∀j. j ≠ i → netmap s' j = netmap s j"

```

```

  and "net_ip_action np (i:deliver(d)) i n s s'"

```

```

  by (metis delivered_to_net_ips pnet_deliver_single_node)

```

```

from this(2) have "∀j. j ≠ i → fst (netgmap sr s') j = fst (netgmap sr s) j"

```

```

  by (clarsimp intro!: netmap_is_fst_netgmap')

```

```

from ⟨(s, i:deliver(d), s') ∈ trans (pnet np n)⟩ have "net_ips s' = net_ips s"

```

```

  by (rule pnet_maintains_dom [THEN sym])

```

```

define σ' where "σ' j = (if j = i then the (fst (netgmap sr s')) i) else σ j" for j

```

```

from ⟨∀j. j ≠ i → fst (netgmap sr s') j = fst (netgmap sr s) j⟩

```

```

  and ⟨netgmap sr s = netmask (net_tree_ips n) (σ, ζ)⟩

```

```

  have "∀j. j ≠ i → σ' j = σ j"

```

```

  unfolding σ'_def by clarsimp

```

```

from assms(2) have "net_ips s = net_tree_ips n"

```

```

  by (rule pnet_net_ips_net_tree_ips)

```

```

from ⟨netgmap sr s = netmask (net_tree_ips n) (σ, ζ)⟩

```

```

have "ζ = snd (netgmap sr s)" by simp

from (∀ j. j ≠ i → fst (netgmap sr s') j = fst (netgmap sr s) j) (i ∈ net_ips s)
  (net_ips s = net_tree_ips n) (net_ips s' = net_ips s)
  (netgmap sr s = netmask (net_tree_ips n) (σ, ζ))
have "fst (netgmap sr s') = fst (netmask (net_tree_ips n) (σ', snd (netgmap sr s')))"
  unfolding σ'_def [abs_def] by - (rule ext, clarsimp)

hence "netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"
  by (rule prod_eqI, simp)

with assms(1, 3)
  have "((σ, snd (netgmap sr s)), i:deliver(d), (σ', snd (netgmap sr s'))) ∈ trans (opnet onp n)"
    using assms(2,4) (i ∈ net_ips s) and (net_ip_action np (i:deliver(d)) i n s s')
proof (induction n arbitrary: s s' ζ)
  fix ii R_i ns ns' ζ
  assume "(ns, i:deliver(d), ns') ∈ trans (pnet np ⟨ii; R_i⟩)"
    and nsr: "ns ∈ reachable (pnet np ⟨ii; R_i⟩) TT"
    and "netgmap sr ns = netmask (net_tree_ips ⟨ii; R_i⟩) (σ, ζ)"
    and "netgmap sr ns' = netmask (net_tree_ips ⟨ii; R_i⟩) (σ', snd (netgmap sr ns'))"
    and "i ∈ net_ips ns"
  from this(1) have "(ns, i:deliver(d), ns') ∈ node_sos (trans (np ii))"
    by (simp add: node_comps)
  moreover with nsr obtain s s' R R' where "ns = NodeS ii s R"
    and "ns' = NodeS ii s' R'"
    by (metis net_node_reachable_is_node node_sos_dest)
  moreover from (i ∈ net_ips ns) and (ns = NodeS ii s R) have "ii = i" by simp
  ultimately have ntr: "(NodeS i s R, i:deliver(d), NodeS i s' R') ∈ node_sos (trans (np i))"
    by simp
  hence "R' = R" by (metis net_state.inject(1) node_deliverTE')

  from ntr have "(s, deliver d, s') ∈ trans (np i)"
    by (rule node_deliverTE') simp

  from (netgmap sr ns = netmask (net_tree_ips ⟨ii; R_i⟩) (σ, ζ)) (ns = NodeS ii s R) and (ii = i)
    have "σ i = fst (sr s)" by simp (metis map_upd_Some_unfold)

  moreover from (netgmap sr ns' = netmask (net_tree_ips ⟨ii; R_i⟩) (σ', snd (netgmap sr ns')))
    (ns' = NodeS ii s' R') and (ii = i)
    have "σ' i = fst (sr s')"
      unfolding σ'_def [abs_def] by clarsimp (hypsubst_thin,
        metis (lifting, full_types) fun_upd_same option.sel)
  ultimately have "((σ, snd (sr s)), deliver d, (σ', snd (sr s'))) ∈ trans (onp i)"
    using ((s, deliver d, s') ∈ trans (np i)) by (rule trans)

  with (∀ j. j ≠ i → σ' j = σ j) (R'=R)
    have "((σ, NodeS i (snd (sr s)) R), i:deliver(d), (σ', NodeS i (snd (sr s')) R'))
      ∈ onode_sos (trans (onp i))"
      by (metis onode_sos.onode_deliver)

  with (ns = NodeS ii s R) (ns' = NodeS ii s' R') (ii = i)
    show "((σ, snd (netgmap sr ns)), i:deliver(d), (σ', snd (netgmap sr ns'))) ∈ trans (opnet onp ⟨ii;
R_i⟩)"
    by (simp add: onode_comps)
next
fix n1 n2 s s' ζ
assume IH1: "∧ s s' ζ. (s, i:deliver(d), s') ∈ trans (pnet np n1)
  ⇒ netgmap sr s = netmask (net_tree_ips n1) (σ, ζ)
  ⇒ netgmap sr s' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s'))
  ⇒ s ∈ reachable (pnet np n1) TT
  ⇒ wf_net_tree n1
  ⇒ i ∈ net_ips s
  ⇒ net_ip_action np (i:deliver(d)) i n1 s s'
  ⇒ ((σ, snd (netgmap sr s)), i:deliver(d), (σ', snd (netgmap sr s'))) ∈ trans (opnet

```

```

onp n1)"
  and IH2: " $\bigwedge s s' \zeta. (s, i:\text{deliver}(d), s') \in \text{trans (pnet np n2)}$ 
     $\Rightarrow \text{netgmap sr } s = \text{netmask (net\_tree\_ips n2)} (\sigma, \zeta)$ 
     $\Rightarrow \text{netgmap sr } s' = \text{netmask (net\_tree\_ips n2)} (\sigma', \text{snd (netgmap sr } s'))$ 
     $\Rightarrow s \in \text{reachable (pnet np n2)} TT$ 
     $\Rightarrow \text{wf\_net\_tree n2}$ 
     $\Rightarrow i \in \text{net\_ips } s$ 
     $\Rightarrow \text{net\_ip\_action np (i:\text{deliver}(d)) i n2 s'}$ 
     $\Rightarrow ((\sigma, \text{snd (netgmap sr } s)), i:\text{deliver}(d), (\sigma', \text{snd (netgmap sr } s')))) \in \text{trans (opnet$ 
onp n2))"
  and tr: "(s, i:\text{deliver}(d), s') \in \text{trans (pnet np (n1 || n2))}"
  and sr: "s \in \text{reachable (pnet np (n1 || n2)) TT}"
  and nm: "netgmap sr s = netmask (net\_tree\_ips (n1 || n2)) (\sigma, \zeta)"
  and nm': "netgmap sr s' = netmask (net\_tree\_ips (n1 || n2)) (\sigma', \text{snd (netgmap sr } s')))"
  and "wf\_net\_tree (n1 || n2)"
  and "i \in \text{net\_ips } s"
  and "net\_ip\_action np (i:\text{deliver}(d)) i (n1 || n2) s s'"
from tr have "(s, i:\text{deliver}(d), s') \in \text{pnet\_sos (trans (pnet np n1)) (trans (pnet np n2))}" by simp
then obtain s1 s1' s2 s2' where "s = SubnetS s1 s2"
  and "s' = SubnetS s1' s2'"
  by (rule partial\_deliverTE) auto
from this(1) and nm have "netgmap sr (SubnetS s1 s2) = netmask (net\_tree\_ips (n1 || n2)) (\sigma, \zeta)"
  by simp
from (s' = SubnetS s1' s2') and nm'
  have "netgmap sr (SubnetS s1' s2') = netmask (net\_tree\_ips (n1 || n2)) (\sigma', \text{snd (netgmap sr } s')))"
  by simp

from (wf\_net\_tree (n1 || n2)) have "wf\_net\_tree n1"
  and "wf\_net\_tree n2"
  and "net\_tree\_ips n1 \cap net\_tree\_ips n2 = {}" by auto

from sr [simplified (s = SubnetS s1 s2)] have "s1 \in \text{reachable (pnet np n1)} TT"
  by (rule subnet\_reachable(1))
hence "net\_ips s1 = net\_tree\_ips n1" by (rule pnet\_net\_ips\_net\_tree\_ips)

from sr [simplified (s = SubnetS s1 s2)] have "s2 \in \text{reachable (pnet np n2)} TT"
  by (rule subnet\_reachable(2))
hence "net\_ips s2 = net\_tree\_ips n2" by (rule pnet\_net\_ips\_net\_tree\_ips)

from nm [simplified (s = SubnetS s1 s2)]
  (net\_tree\_ips n1 \cap net\_tree\_ips n2 = {})
  (net\_ips s1 = net\_tree\_ips n1)
  (net\_ips s2 = net\_tree\_ips n2)
  have "netgmap sr s1 = netmask (net\_tree\_ips n1) (\sigma, \text{snd (netgmap sr } s1))"
  by (rule netgmap\_subnet\_split1)

from nm [simplified (s = SubnetS s1 s2)]
  (net\_ips s1 = net\_tree\_ips n1)
  (net\_ips s2 = net\_tree\_ips n2)
  have "netgmap sr s2 = netmask (net\_tree\_ips n2) (\sigma, \text{snd (netgmap sr } s2))"
  by (rule netgmap\_subnet\_split2)

from (i \in \text{net\_ips } s) and (s = SubnetS s1 s2) have "i \in \text{net\_ips } s1 \vee i \in \text{net\_ips } s2" by auto
thus "((\sigma, \text{snd (netgmap sr } s)), i:\text{deliver}(d), (\sigma', \text{snd (netgmap sr } s')))) \in \text{trans (opnet onp (n1 || n2))}"
proof
  assume "i \in \text{net\_ips } s1"
  with (s = SubnetS s1 s2) (s' = SubnetS s1' s2') (net\_ip\_action np (i:\text{deliver}(d)) i (n1 || n2) s s')
    have "(s1, i:\text{deliver}(d), s1') \in \text{trans (pnet np n1)}"
      and "net\_ip\_action np (i:\text{deliver}(d)) i n1 s1 s1'"
      and "s2' = s2" by simp_all

  from (net\_ips s1 = net\_tree\_ips n1) and ((s1, i:\text{deliver}(d), s1') \in \text{trans (pnet np n1)})
    have "net\_ips s1' = net\_tree\_ips n1" by (metis pnet\_maintains\_dom)

```

```

from nm' [simplified ⟨s' = SubnetS s1' s2'⟩ ⟨s2' = s2⟩]
  ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
  ⟨net_ips s1' = net_tree_ips n1⟩
  ⟨net_ips s2 = net_tree_ips n2⟩
have "netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))"
  by (rule netgmap_subnet_split1)

from ⟨(s1, i:deliver(d), s1') ∈ trans (pnet np n1)⟩
  ⟨netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))⟩
  ⟨netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))⟩
  ⟨s1 ∈ reachable (pnet np n1) TT⟩
  ⟨wf_net_tree n1⟩
  ⟨i ∈ net_ips s1⟩
  ⟨net_ip_action np (i:deliver(d)) i n1 s1 s1'⟩
have "((σ, snd (netgmap sr s1)), i:deliver(d), (σ', snd (netgmap sr s1')))) ∈ trans (opnet onp
n1)"
  by (rule IH1)

with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨s2' = s2⟩ show ?thesis
  by simp (erule opnet_deliver1)
next
assume "i ∈ net_ips s2"
with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨net_ip_action np (i:deliver(d)) i (n1 || n2) s s'⟩
  have "(s2, i:deliver(d), s2') ∈ trans (pnet np n2)"
  and "net_ip_action np (i:deliver(d)) i n2 s2 s2'"
  and "s1' = s1" by simp_all

from ⟨net_ips s2 = net_tree_ips n2⟩ and ⟨(s2, i:deliver(d), s2') ∈ trans (pnet np n2)⟩
  have "net_ips s2' = net_tree_ips n2" by (metis pnet_maintains_dom)

from nm' [simplified ⟨s' = SubnetS s1' s2'⟩ ⟨s1' = s1⟩]
  ⟨net_ips s1 = net_tree_ips n1⟩
  ⟨net_ips s2' = net_tree_ips n2⟩
have "netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))"
  by (rule netgmap_subnet_split2)

from ⟨(s2, i:deliver(d), s2') ∈ trans (pnet np n2)⟩
  ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
  ⟨netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))⟩
  ⟨s2 ∈ reachable (pnet np n2) TT⟩
  ⟨wf_net_tree n2⟩
  ⟨i ∈ net_ips s2⟩
  ⟨net_ip_action np (i:deliver(d)) i n2 s2 s2'⟩
have "((σ, snd (netgmap sr s2)), i:deliver(d), (σ', snd (netgmap sr s2')))) ∈ trans (opnet onp
n2)"
  by (rule IH2)

with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨s1' = s1⟩ show ?thesis
  by simp (erule opnet_deliver2)
qed
qed
with ⟨ζ = snd (netgmap sr s)⟩ have "((σ, ζ), i:deliver(d), (σ', snd (netgmap sr s')))) ∈ trans (opnet
onp n)"
  by simp
moreover from ⟨∀j. j ≠ i → σ' j = σ j⟩ ⟨i ∈ net_ips s⟩ ⟨ζ = snd (netgmap sr s)⟩
  have "∀j. j ∉ net_ips ζ → σ' j = σ j" by (metis net_ips_netgmap)
ultimately have "((σ, ζ), i:deliver(d), (σ', snd (netgmap sr s')))) ∈ trans (opnet onp n)
  ∧ (∀j. j ∉ net_ips ζ → σ' j = σ j)
  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"
  using ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))⟩ by simp
thus "∃σ' ζ'. ((σ, ζ), i:deliver(d), (σ', ζ')) ∈ trans (opnet onp n)
  ∧ (∀j. j ∉ net_ips ζ → σ' j = σ j)
  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'" by auto

```

qed

lemma transfer\_arrive':

```
assumes "(s, H-K:arrive(m), s') ∈ trans (pnet np n)"
  and "s ∈ reachable (pnet np n) TT"
  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"
  and "wf_net_tree n"
shows "((σ, ζ), H-K:arrive(m), (σ', ζ')) ∈ trans (opnet onp n)"
proof -
```

```
from assms(2) have "net_ips s = net_tree_ips n" ..
with assms(1) have "net_ips s' = net_tree_ips n"
  by (metis pnet_maintains_dom)
```

```
from ⟨netgmap sr s = netmask (net_tree_ips n) (σ, ζ)⟩
  have "ζ = snd (netgmap sr s)" by simp
```

```
from ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', ζ')⟩
  have "ζ' = snd (netgmap sr s'"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"
  by simp_all
```

```
from assms(1-3) ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))⟩ assms(5)
  have "((σ, snd (netgmap sr s)), H-K:arrive(m), (σ', snd (netgmap sr s'))) ∈ trans (opnet onp n)"
proof (induction n arbitrary: s s' ζ H K)
```

```
  fix ii Ri ns ns' ζ H K
```

```
  assume "(ns, H-K:arrive(m), ns') ∈ trans (pnet np ⟨ii; Ri⟩)"
  and nsr: "ns ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
  and "netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)"
  and "netgmap sr ns' = netmask (net_tree_ips ⟨ii; Ri⟩) (σ', snd (netgmap sr ns'))"
```

```
from this(1) have "(ns, H-K:arrive(m), ns') ∈ node_sos (trans (np ii))"
  by (simp add: node_comps)
```

```
moreover with nsr obtain s s' R where "ns = NodeS ii s R"
  and "ns' = NodeS ii s' R"
```

```
  by (metis net_node_reachable_is_node_node_arriveTE')
```

```
ultimately have "(NodeS ii s R, H-K:arrive(m), NodeS ii s' R) ∈ node_sos (trans (np ii))"
  by simp
```

```
from this(1) have "((σ, NodeS ii (snd (sr s)) R), H-K:arrive(m), (σ', NodeS ii (snd (sr s')) R))
  ∈ onode_sos (trans (onp ii))"
```

```
proof (rule node_arriveTE)
```

```
  assume "(s, receive m, s') ∈ trans (np ii)"
  and "H = {ii}"
  and "K = {}"
```

```
from ⟨netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)⟩ and (ns = NodeS ii s R)
  have "σ ii = fst (sr s)"
```

```
  by simp (metis map_upd_Some_unfold)
```

```
moreover from ⟨netgmap sr ns' = netmask (net_tree_ips ⟨ii; Ri⟩) (σ', snd (netgmap sr ns'))⟩
  and (ns' = NodeS ii s' R)
```

```
  have "σ' ii = fst (sr s'" by simp (metis map_upd_Some_unfold)
```

```
ultimately have "((σ, snd (sr s)), receive m, (σ', snd (sr s')) ∈ trans (onp ii)"
```

```
  using ⟨(s, receive m, s') ∈ trans (np ii)⟩ by (rule trans)
```

```
hence "((σ, NodeS ii (snd (sr s)) R), {ii}-{}:arrive(m), (σ', NodeS ii (snd (sr s')) R))
  ∈ onode_sos (trans (onp ii))"
```

```
  by (rule onode_receive)
```

```
with ⟨H={ii}⟩ and ⟨K={}⟩
```

```
  show "((σ, NodeS ii (snd (sr s)) R), H-K:arrive(m), (σ', NodeS ii (snd (sr s')) R))
  ∈ onode_sos (trans (onp ii))"
```

```
  by simp
```

```
next
```

```
  assume "H = {}"
```

```
  and "s = s'"
```

```
  and "K = {ii}"
```

```
from ⟨s = s'⟩ ⟨netgmap sr ns' = netmask (net_tree_ips ⟨ii; Ri⟩) (σ', snd (netgmap sr ns'))⟩
```

```

      ⟨netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)⟩
      ⟨ns = NodeS ii s R⟩ and ⟨ns' = NodeS ii s' R⟩
    have "σ' ii = σ ii" by simp (metis option.sel)
  hence "((σ, NodeS ii (snd (sr s)) R), {}¬{ii}:arrive(m), (σ', NodeS ii (snd (sr s)) R))
      ∈ onode_sos (trans (onp ii)))"
    by (rule onode_arrive)
  with ⟨H={⟩ ⟨K={ii}⟩ and ⟨s = s'⟩
    show "((σ, NodeS ii (snd (sr s)) R), H¬K:arrive(m), (σ', NodeS ii (snd (sr s')) R))
      ∈ onode_sos (trans (onp ii)))"
    by simp
qed
with ⟨ns = NodeS ii s R⟩ ⟨ns' = NodeS ii s' R⟩
  show "((σ, snd (netgmap sr ns)), H¬K:arrive(m), (σ', snd (netgmap sr ns')))
      ∈ trans (opnet onp ⟨ii; Ri⟩)"
    by (simp add: onode_comps)
next
fix n1 n2 s s' ζ H K
assume IH1: "∧s s' ζ H K. (s, H¬K:arrive(m), s') ∈ trans (pnet np n1)
      ⇒ s ∈ reachable (pnet np n1) TT
      ⇒ netgmap sr s = netmask (net_tree_ips n1) (σ, ζ)
      ⇒ netgmap sr s' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s'))
      ⇒ wf_net_tree n1
      ⇒ ((σ, snd (netgmap sr s)), H¬K:arrive(m), σ', snd (netgmap sr s'))
          ∈ trans (opnet onp n1)"
and IH2: "∧s s' ζ H K. (s, H¬K:arrive(m), s') ∈ trans (pnet np n2)
      ⇒ s ∈ reachable (pnet np n2) TT
      ⇒ netgmap sr s = netmask (net_tree_ips n2) (σ, ζ)
      ⇒ netgmap sr s' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s'))
      ⇒ wf_net_tree n2
      ⇒ ((σ, snd (netgmap sr s)), H¬K:arrive(m), σ', snd (netgmap sr s'))
          ∈ trans (opnet onp n2)"
and "(s, H¬K:arrive(m), s') ∈ trans (pnet np (n1 || n2))"
and sr: "s ∈ reachable (pnet np (n1 || n2)) TT"
and nm: "netgmap sr s = netmask (net_tree_ips (n1 || n2)) (σ, ζ)"
and nm': "netgmap sr s' = netmask (net_tree_ips (n1 || n2)) (σ', snd (netgmap sr s'))"
and "wf_net_tree (n1 || n2)"
from this(3) have "(s, H¬K:arrive(m), s') ∈ pnet_sos (trans (pnet np n1))
      (trans (pnet np n2))"
  by simp
thus "((σ, snd (netgmap sr s)), H¬K:arrive(m), (σ', snd (netgmap sr s')))
      ∈ trans (opnet onp (n1 || n2))"
proof (rule partial_arriveTE)
  fix s1 s1' s2 s2' H1 H2 K1 K2
  assume "s = SubnetS s1 s2"
    and "s' = SubnetS s1' s2'"
    and tr1: "(s1, H1¬K1:arrive(m), s1') ∈ trans (pnet np n1)"
    and tr2: "(s2, H2¬K2:arrive(m), s2') ∈ trans (pnet np n2)"
    and "H = H1 ∪ H2"
    and "K = K1 ∪ K2"

  from ⟨wf_net_tree (n1 || n2)⟩ have "wf_net_tree n1"
    and "wf_net_tree n2"
    and "net_tree_ips n1 ∩ net_tree_ips n2 = {}" by auto

  from sr [simplified ⟨s = SubnetS s1 s2⟩] have "s1 ∈ reachable (pnet np n1) TT"
    by (rule subnet_reachable(1))
  hence "net_ips s1 = net_tree_ips n1" by (rule pnet_net_ips_net_tree_ips)
  with tr1 have "net_ips s1' = net_tree_ips n1" by (metis pnet_maintains_dom)

  from sr [simplified ⟨s = SubnetS s1 s2⟩] have "s2 ∈ reachable (pnet np n2) TT"
    by (rule subnet_reachable(2))
  hence "net_ips s2 = net_tree_ips n2" by (rule pnet_net_ips_net_tree_ips)
  with tr2 have "net_ips s2' = net_tree_ips n2" by (metis pnet_maintains_dom)

```

```

from ⟨(s1, H1→K1:arrive(m), s1') ∈ trans (pnet np n1)⟩
  ⟨s1 ∈ reachable (pnet np n1) TT⟩
  have "((σ, snd (netgmap sr s1)), H1→K1:arrive(m), (σ', snd (netgmap sr s1'))))
      ∈ trans (opnet onp n1)"

proof (rule IH1 [OF _ _ _ _ ⟨wf_net_tree n1⟩])
  from nm [simplified ⟨s = SubnetS s1 s2⟩]
    ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
    ⟨net_ips s1 = net_tree_ips n1⟩
    ⟨net_ips s2 = net_tree_ips n2⟩
  show "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
  by (rule netgmap_subnet_split1)
next
  from nm' [simplified ⟨s' = SubnetS s1' s2'⟩]
    ⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
    ⟨net_ips s1' = net_tree_ips n1⟩
    ⟨net_ips s2' = net_tree_ips n2⟩
  show "netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))"
  by (rule netgmap_subnet_split1)
qed

moreover from ⟨(s2, H2→K2:arrive(m), s2') ∈ trans (pnet np n2)⟩
  ⟨s2 ∈ reachable (pnet np n2) TT⟩
  have "((σ, snd (netgmap sr s2)), H2→K2:arrive(m), (σ', snd (netgmap sr s2'))))
      ∈ trans (opnet onp n2)"

proof (rule IH2 [OF _ _ _ _ ⟨wf_net_tree n2⟩])
  from nm [simplified ⟨s = SubnetS s1 s2⟩]
    ⟨net_ips s1 = net_tree_ips n1⟩
    ⟨net_ips s2 = net_tree_ips n2⟩
  show "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
  by (rule netgmap_subnet_split2)
next
  from nm' [simplified ⟨s' = SubnetS s1' s2'⟩]
    ⟨net_ips s1' = net_tree_ips n1⟩
    ⟨net_ips s2' = net_tree_ips n2⟩
  show "netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))"
  by (rule netgmap_subnet_split2)
qed
ultimately show "((σ, snd (netgmap sr s)), H→K:arrive(m), (σ', snd (netgmap sr s'))))
  ∈ trans (opnet onp (n1 || n2))"
  using ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ ⟨H = H1 ∪ H2⟩ ⟨K = K1 ∪ K2⟩
  by simp (rule opnet_sos.opnet_arrive)
qed
qed
with ⟨ζ = snd (netgmap sr s)⟩ and ⟨ζ' = snd (netgmap sr s')⟩
  show "((σ, ζ), H→K:arrive(m), (σ', ζ')) ∈ trans (opnet onp n)"
  by simp
qed

```

lemma transfer\_arrive:

```

assumes "(s, H→K:arrive(m), s') ∈ trans (pnet np n)"
  and "s ∈ reachable (pnet np n) TT"
  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "wf_net_tree n"

```

```

obtains σ' ζ' where "((σ, ζ), H→K:arrive(m), (σ', ζ')) ∈ trans (opnet onp n)"
  and "∀j. j ∉ net_ips ζ → σ' j = σ j"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ')"

```

proof atomize\_elim

```

define σ' where "σ' i = (if i ∈ net_tree_ips n then the (fst (netgmap sr s')) i else σ i)" for i

```

```

from assms(2) have "net_ips s = net_tree_ips n"
  by (rule pnet_net_ips_net_tree_ips)

```

```

with assms(1) have "net_ips s' = net_tree_ips n"
  by (metis pnet_maintains_dom)

```



```

have "netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"
proof (rule prod_eqI)
  from ⟨net_ips s' = net_tree_ips n⟩
  show "fst (netgmap sr s') = fst (netmask (net_tree_ips n) (σ', snd (netgmap sr s')))"
  unfolding σ'_def [abs_def] by - (rule ext, clarsimp)
qed simp

moreover with assms(1-3)
have "((σ, ζ), H¬K:arrive(m), (σ', snd (netgmap sr s')))) ∈ trans (opnet onp n)"
  using ⟨wf_net_tree n⟩ by (rule transfer_arrive')

moreover have "∀j. j∉net_ips ζ → σ' j = σ j"
proof -
  have "∀j. j∉net_tree_ips n → σ' j = σ j" unfolding σ'_def by simp
  with assms(3) and ⟨net_ips s = net_tree_ips n⟩
  show ?thesis
  by clarsimp (metis (mono_tags) net_ips_netgmap snd_conv)
qed

ultimately show "∃σ' ζ'. ((σ, ζ), H¬K:arrive(m), (σ', ζ')) ∈ trans (opnet onp n)
  ∧ (∀j. j∉net_ips ζ → σ' j = σ j)
  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'" by auto
qed

lemma transfer_cast:
assumes "(s, mR:*cast(m), s') ∈ trans (pnet np n)"
  and "s ∈ reachable (pnet np n) TT"
  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "wf_net_tree n"
obtains σ' ζ' where "((σ, ζ), mR:*cast(m), (σ', ζ')) ∈ trans (opnet onp n)"
  and "∀j. j∉net_ips ζ → σ' j = σ j"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"
proof atomize_elim
  define σ' where "σ' i = (if i∈net_tree_ips n then the (fst (netgmap sr s')) i) else σ i)" for i

  from assms(2) have "net_ips s = net_tree_ips n" ..
  with assms(1) have "net_ips s' = net_tree_ips n"
  by (metis pnet_maintains_dom)
  have "netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))"
  proof (rule prod_eqI)
    from ⟨net_ips s' = net_tree_ips n⟩
    show "fst (netgmap sr s') = fst (netmask (net_tree_ips n) (σ', snd (netgmap sr s')))"
    unfolding σ'_def [abs_def] by - (rule ext, clarsimp simp add: some_the_fst_netgmap)
  qed simp

  from ⟨net_ips s' = net_tree_ips n⟩ and ⟨net_ips s = net_tree_ips n⟩
  have "∀j. j∉net_ips (snd (netgmap sr s)) → σ' j = σ j"
  unfolding σ'_def by simp

  from ⟨netgmap sr s = netmask (net_tree_ips n) (σ, ζ)⟩
  have "ζ = snd (netgmap sr s)" by simp

  from assms(1-3) ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', snd (netgmap sr s'))⟩ assms(4)
  have "((σ, snd (netgmap sr s)), mR:*cast(m), (σ', snd (netgmap sr s')))) ∈ trans (opnet onp n)"
  proof (induction n arbitrary: s s' ζ mR)
    fix ii Ri ns ns' ζ mR
    assume "(ns, mR:*cast(m), ns') ∈ trans (pnet np ⟨ii; Ri⟩)"
      and nsr: "ns ∈ reachable (pnet np ⟨ii; Ri⟩) TT"
      and "netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)"
      and "netgmap sr ns' = netmask (net_tree_ips ⟨ii; Ri⟩) (σ', snd (netgmap sr ns'))"
    from this(1) have "(ns, mR:*cast(m), ns') ∈ node_sos (trans (np ii))"
    by (simp add: node_comps)
    moreover with nsr obtain s s' R where "ns = NodeS ii s R"
      and "ns' = NodeS ii s' R"

```

```

    by (metis net_node_reachable_is_node node_castTE')
ultimately have "(NodeS ii s R, mR:*cast(m), NodeS ii s' R) ∈ node_sos (trans (np ii))"
  by simp

from ⟨netgmap sr ns = netmask (net_tree_ips ⟨ii; Ri⟩) (σ, ζ)⟩ and ⟨ns = NodeS ii s R⟩
  have "σ ii = fst (sr s)"
    by simp (metis map_upd_Some_unfold)
from ⟨netgmap sr ns' = netmask (net_tree_ips ⟨ii; Ri⟩) (σ', snd (netgmap sr ns'))⟩
  and ⟨ns' = NodeS ii s' R⟩
  have "σ' ii = fst (sr s')" by simp (metis map_upd_Some_unfold)

from ⟨(NodeS ii s R, mR:*cast(m), NodeS ii s' R) ∈ node_sos (trans (np ii))⟩
  have "((σ, NodeS ii (snd (sr s)) R), mR:*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"

proof (rule node_castTE)
  assume "(s, broadcast m, s') ∈ trans (np ii)"
    and "R = mR"
  from ⟨σ ii = fst (sr s)⟩ ⟨σ' ii = fst (sr s')⟩ and this(1)
    have "((σ, snd (sr s)), broadcast m, (σ', snd (sr s'))) ∈ trans (onp ii)"
      by (rule trans)
  hence "((σ, NodeS ii (snd (sr s)) R), R:*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"
    by (rule onode_bcast)
  with ⟨R=mR⟩ show "((σ, NodeS ii (snd (sr s)) R), mR:*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"
    by simp
next
fix D
assume "(s, groupcast D m, s') ∈ trans (np ii)"
  and "mR = R ∩ D"
from ⟨σ ii = fst (sr s)⟩ ⟨σ' ii = fst (sr s')⟩ and this(1)
  have "((σ, snd (sr s)), groupcast D m, (σ', snd (sr s'))) ∈ trans (onp ii)"
    by (rule trans)
  hence "((σ, NodeS ii (snd (sr s)) R), (R ∩ D):*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"
    by (rule onode_gcast)
  with ⟨mR = R ∩ D⟩ show "((σ, NodeS ii (snd (sr s)) R), mR:*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"
    by simp
next
fix d
assume "(s, unicast d m, s') ∈ trans (np ii)"
  and "d ∈ R"
  and "mR = {d}"
from ⟨σ ii = fst (sr s)⟩ ⟨σ' ii = fst (sr s')⟩ and this(1)
  have "((σ, snd (sr s)), unicast d m, (σ', snd (sr s'))) ∈ trans (onp ii)"
    by (rule trans)
  hence "((σ, NodeS ii (snd (sr s)) R), {d}:*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"
    using ⟨d∈R⟩ by (rule onode_ucast)
  with ⟨mR={d}⟩ show "((σ, NodeS ii (snd (sr s)) R), mR:*cast(m), (σ', NodeS ii (snd (sr s')) R))
    ∈ onode_sos (trans (onp ii)))"
    by simp
qed
with ⟨ns = NodeS ii s R⟩ ⟨ns' = NodeS ii s' R⟩
  show "((σ, snd (netgmap sr ns)), mR:*cast(m), (σ', snd (netgmap sr ns')))
    ∈ trans (opnet onp ⟨ii; Ri⟩)"
    by (simp add: onode_comps)
next
fix n1 n2 s s' ζ mR
assume IH1: "∧s s' ζ mR. (s, mR:*cast(m), s') ∈ trans (pnet np n1)
  ⇒ s ∈ reachable (pnet np n1) TT
  ⇒ netgmap sr s = netmask (net_tree_ips n1) (σ, ζ)"

```

```

    ⇒ netgmap sr s' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s'))
    ⇒ wf_net_tree n1
    ⇒ ((σ, snd (netgmap sr s)), mR:*cast(m), σ', snd (netgmap sr s'))
      ∈ trans (opnet onp n1)"
and IH2: "∧s s' ζ mR. (s, mR:*cast(m), s') ∈ trans (pnet np n2)
    ⇒ s ∈ reachable (pnet np n2) TT
    ⇒ netgmap sr s = netmask (net_tree_ips n2) (σ, ζ)
    ⇒ netgmap sr s' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s'))
    ⇒ wf_net_tree n2
    ⇒ ((σ, snd (netgmap sr s)), mR:*cast(m), σ', snd (netgmap sr s'))
      ∈ trans (opnet onp n2)"

and "(s, mR:*cast(m), s') ∈ trans (pnet np (n1 || n2))"
and sr: "s ∈ reachable (pnet np (n1 || n2)) TT"
and nm: "netgmap sr s = netmask (net_tree_ips (n1 || n2)) (σ, ζ)"
and nm': "netgmap sr s' = netmask (net_tree_ips (n1 || n2)) (σ', snd (netgmap sr s'))"
and "wf_net_tree (n1 || n2)"
from this(3) have "(s, mR:*cast(m), s') ∈ pnet_sos (trans (pnet np n1)) (trans (pnet np n2))"
by simp
then obtain s1 s1' s2 s2' H K
where "s = SubnetS s1 s2"
and "s' = SubnetS s1' s2'"
and "H ⊆ mR"
and "K ∩ mR = {}"
and trtr: "((s1, mR:*cast(m), s1') ∈ trans (pnet np n1)
    ∧ (s2, H¬K:arrive(m), s2') ∈ trans (pnet np n2))
  ∨ ((s1, H¬K:arrive(m), s1') ∈ trans (pnet np n1)
    ∧ (s2, mR:*cast(m), s2') ∈ trans (pnet np n2))"
by (rule partial_castTE) metis+

from ⟨wf_net_tree (n1 || n2)⟩ have "wf_net_tree n1"
and "wf_net_tree n2"
and "net_tree_ips n1 ∩ net_tree_ips n2 = {}" by auto

from sr [simplified ⟨s = SubnetS s1 s2⟩] have "s1 ∈ reachable (pnet np n1) TT"
by (rule subnet_reachable(1))
hence "net_ips s1 = net_tree_ips n1" by (rule pnet_net_ips_net_tree_ips)
with trtr have "net_ips s1' = net_tree_ips n1" by (metis pnet_maintains_dom)

from sr [simplified ⟨s = SubnetS s1 s2⟩] have "s2 ∈ reachable (pnet np n2) TT"
by (rule subnet_reachable(2))
hence "net_ips s2 = net_tree_ips n2" by (rule pnet_net_ips_net_tree_ips)
with trtr have "net_ips s2' = net_tree_ips n2" by (metis pnet_maintains_dom)

from nm [simplified ⟨s = SubnetS s1 s2⟩]
⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
⟨net_ips s1 = net_tree_ips n1⟩
⟨net_ips s2 = net_tree_ips n2⟩
have "netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))"
by (rule netgmap_subnet_split1)

from nm' [simplified ⟨s' = SubnetS s1' s2'⟩]
⟨net_tree_ips n1 ∩ net_tree_ips n2 = {}⟩
⟨net_ips s1' = net_tree_ips n1⟩
⟨net_ips s2' = net_tree_ips n2⟩
have "netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))"
by (rule netgmap_subnet_split1)

from nm [simplified ⟨s = SubnetS s1 s2⟩]
⟨net_ips s1 = net_tree_ips n1⟩
⟨net_ips s2 = net_tree_ips n2⟩
have "netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))"
by (rule netgmap_subnet_split2)

from nm' [simplified ⟨s' = SubnetS s1' s2'⟩]

```

```

    ⟨net_ips s1' = net_tree_ips n1⟩
    ⟨net_ips s2' = net_tree_ips n2⟩
  have "netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))"
    by (rule netgmap_subnet_split2)

from trtr show "((σ, snd (netgmap sr s)), mR:*cast(m), (σ', snd (netgmap sr s')))
  ∈ trans (opnet onp (n1 || n2))"

proof (elim disjE conjE)
  assume "(s1, mR:*cast(m), s1') ∈ trans (pnet np n1)"
  and "(s2, H¬K:arrive(m), s2') ∈ trans (pnet np n2)"
  from ⟨(s1, mR:*cast(m), s1') ∈ trans (pnet np n1)⟩
    ⟨s1 ∈ reachable (pnet np n1) TT⟩
    ⟨netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))⟩
    ⟨netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))⟩
    ⟨wf_net_tree n1⟩
  have "((σ, snd (netgmap sr s1)), mR:*cast(m), (σ', snd (netgmap sr s1'))) ∈ trans (opnet onp
n1)"
    by (rule IH1)

  moreover from ⟨(s2, H¬K:arrive(m), s2') ∈ trans (pnet np n2)⟩
    ⟨s2 ∈ reachable (pnet np n2) TT⟩
    ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
    ⟨netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))⟩
    ⟨wf_net_tree n2⟩
  have "((σ, snd (netgmap sr s2)), H¬K:arrive(m), (σ', snd (netgmap sr s2'))) ∈ trans (opnet onp
n2)"
    by (rule transfer_arrive')

  ultimately have "((σ, SubnetS (snd (netgmap sr s1)) (snd (netgmap sr s2))), mR:*cast(m),
    (σ', SubnetS (snd (netgmap sr s1')) (snd (netgmap sr s2'))))
    ∈ opnet_sos (trans (opnet onp n1)) (trans (opnet onp n2))"
    using ⟨H ⊆ mR⟩ and ⟨K ∩ mR = {}⟩ by (rule opnet_sos.intros(1))
  with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ show ?thesis by simp
next
  assume "(s1, H¬K:arrive(m), s1') ∈ trans (pnet np n1)"
  and "(s2, mR:*cast(m), s2') ∈ trans (pnet np n2)"
  from ⟨(s1, H¬K:arrive(m), s1') ∈ trans (pnet np n1)⟩
    ⟨s1 ∈ reachable (pnet np n1) TT⟩
    ⟨netgmap sr s1 = netmask (net_tree_ips n1) (σ, snd (netgmap sr s1))⟩
    ⟨netgmap sr s1' = netmask (net_tree_ips n1) (σ', snd (netgmap sr s1'))⟩
    ⟨wf_net_tree n1⟩
  have "((σ, snd (netgmap sr s1)), H¬K:arrive(m), (σ', snd (netgmap sr s1'))) ∈ trans (opnet onp
n1)"
    by (rule transfer_arrive')

  moreover from ⟨(s2, mR:*cast(m), s2') ∈ trans (pnet np n2)⟩
    ⟨s2 ∈ reachable (pnet np n2) TT⟩
    ⟨netgmap sr s2 = netmask (net_tree_ips n2) (σ, snd (netgmap sr s2))⟩
    ⟨netgmap sr s2' = netmask (net_tree_ips n2) (σ', snd (netgmap sr s2'))⟩
    ⟨wf_net_tree n2⟩
  have "((σ, snd (netgmap sr s2)), mR:*cast(m), (σ', snd (netgmap sr s2'))) ∈ trans (opnet onp
n2)"
    by (rule IH2)

  ultimately have "((σ, SubnetS (snd (netgmap sr s1)) (snd (netgmap sr s2))), mR:*cast(m),
    (σ', SubnetS (snd (netgmap sr s1')) (snd (netgmap sr s2'))))
    ∈ opnet_sos (trans (opnet onp n1)) (trans (opnet onp n2))"
    using ⟨H ⊆ mR⟩ and ⟨K ∩ mR = {}⟩ by (rule opnet_sos.intros(2))
  with ⟨s = SubnetS s1 s2⟩ ⟨s' = SubnetS s1' s2'⟩ show ?thesis by simp
qed
qed
with ⟨ζ = snd (netgmap sr s)⟩ have "((σ, ζ), mR:*cast(m), (σ', snd (netgmap sr s'))) ∈ trans (opnet
onp n)"
  by simp

```

```

moreover from ⟨ $\forall j. j \notin \text{net\_ips } (\text{snd } (\text{netgmap } sr s)) \longrightarrow \sigma' j = \sigma j \rangle \langle \zeta = \text{snd } (\text{netgmap } sr s) \rangle$ 
  have " $\forall j. j \notin \text{net\_ips } \zeta \longrightarrow \sigma' j = \sigma j$ " by simp
moreover note ⟨ $\text{netgmap } sr s' = \text{netmask } (\text{net\_tree\_ips } n) (\sigma', \text{snd } (\text{netgmap } sr s')) \rangle$ 
ultimately show " $\exists \sigma' \zeta'. ((\sigma, \zeta), \text{mR}:\text{*cast}(m), (\sigma', \zeta')) \in \text{trans } (\text{opnet } \text{opn } n)$ 
   $\wedge (\forall j. j \notin \text{net\_ips } \zeta \longrightarrow \sigma' j = \sigma j)$ 
   $\wedge \text{netgmap } sr s' = \text{netmask } (\text{net\_tree\_ips } n) (\sigma', \zeta')$ "

  by auto
qed

```

lemma transfer\_pnet\_action:

```

assumes "s ∈ reachable (pnet np n) TT"
  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "wf_net_tree n"
  and "(s, a, s') ∈ trans (pnet np n)"
obtains σ' ζ' where "((σ, ζ), a, (σ', ζ')) ∈ trans (opnet opn n)"
  and " $\forall j. j \notin \text{net\_ips } \zeta \longrightarrow \sigma' j = \sigma j$ "
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ)'"

```

proof atomize\_elim

```

show " $\exists \sigma' \zeta'. ((\sigma, \zeta), a, (\sigma', \zeta')) \in \text{trans } (\text{opnet } \text{opn } n)$ 
   $\wedge (\forall j. j \notin \text{net\_ips } \zeta \longrightarrow \sigma' j = \sigma j)$ 
   $\wedge \text{netgmap } sr s' = \text{netmask } (\text{net\_tree\_ips } n) (\sigma', \zeta')$ "

```

proof (cases a)

```

case node_cast
with assms(4) show ?thesis
  by (auto elim!: transfer_cast [OF _ assms(1-3)])

```

next

```

case node_deliver
with assms(4) show ?thesis
  by (auto elim!: transfer_deliver [OF _ assms(1-3)])

```

next

```

case node_arrive
with assms(4) show ?thesis
  by (auto elim!: transfer_arrive [OF _ assms(1-3)])

```

next

```

case node_connect
with assms(4) show ?thesis
  by (auto elim!: transfer_connect [OF _ assms(1-3)])

```

next

```

case node_disconnect
with assms(4) show ?thesis
  by (auto elim!: transfer_disconnect [OF _ assms(1-3)])

```

next

```

case node_newpkt
with assms(4) have False by (metis pnet_never_newpkt)
thus ?thesis ..

```

next

```

case node_tau
with assms(4) show ?thesis
  by (auto elim!: transfer_tau [OF _ assms(1-3), simplified])

```

qed

qed

lemma transfer\_action\_pnet\_closed:

```

assumes "(s, a, s') ∈ trans (closed (pnet np n))"
obtains a' where "(s, a', s') ∈ trans (pnet np n)"
  and " $\wedge \sigma \zeta \sigma' \zeta'. \llbracket ((\sigma, \zeta), a', (\sigma', \zeta')) \in \text{trans } (\text{opnet } \text{opn } n);$ 
   $(\forall j. j \notin \text{net\_ips } \zeta \longrightarrow \sigma' j = \sigma j) \rrbracket$ 
   $\implies ((\sigma, \zeta), a, (\sigma', \zeta')) \in \text{trans } (\text{oclosed } (\text{opnet } \text{opn } n))$ "

```

proof (atomize\_elim)

```

from assms have "(s, a, s') ∈ cnet_sos (trans (pnet np n))" by simp

```

```

thus " $\exists a'. (s, a', s') \in \text{trans } (\text{pnet } \text{np } n)$ "

```

```

   $\wedge (\forall \sigma \zeta \sigma' \zeta'. ((\sigma, \zeta), a', (\sigma', \zeta')) \in \text{trans } (\text{opnet } \text{opn } n)$ 
     $\longrightarrow (\forall j. j \notin \text{net\_ips } \zeta \longrightarrow \sigma' j = \sigma j)$ 
     $\longrightarrow ((\sigma, \zeta), a, (\sigma', \zeta')) \in \text{trans } (\text{oclosed } (\text{opnet } \text{opn } n)))$ "

```

```

proof cases
  case (cnet_cast R m) thus ?thesis
  by (auto intro!: exI [where x="R:*cast(m)"] dest!: ocnet_cast)
qed (auto intro!: ocnet_sos.intros [simplified])
qed

lemma transfer_action:
  assumes "s ∈ reachable (closed (pnet np n)) TT"
  and "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "wf_net_tree n"
  and "(s, a, s') ∈ trans (closed (pnet np n))"
  obtains σ' ζ' where "((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet onp n))"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ')"
proof atomize_elim
  from assms(1) have "s ∈ reachable (pnet np n) TT" ..
  from assms(4)
  show "∃σ' ζ'. ((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet onp n))
  ∧ netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"
  by (cases a)
  ((elim transfer_action_pnet_closed
  transfer_pnet_action [OF (s ∈ reachable (pnet np n) TT) assms(2-3)]?)?,
  (auto intro!: exI)[1])+)
qed

lemma pnet_reachable_transfer':
  assumes "wf_net_tree n"
  and "s ∈ reachable (closed (pnet np n)) TT"
  shows "netgmap sr s ∈ netmask (net_tree_ips n) ' oreachable (oclosed (opnet onp n)) (λ_ _ . True)
  U"
  (is " _ ∈ ?f ' ?oreachable n")
using assms(2) proof induction
  fix s
  assume "s ∈ init (closed (pnet np n))"
  hence "s ∈ init (pnet np n)" by simp
  with (wf_net_tree n) have "netgmap sr s ∈ netmask (net_tree_ips n) ' init (opnet onp n)"
  by (rule init_pnet_opnet)
  hence "netgmap sr s ∈ netmask (net_tree_ips n) ' init (oclosed (opnet onp n))"
  by simp
  moreover have "netmask (net_tree_ips n) ' init (oclosed (opnet onp n))
  ⊆ netmask (net_tree_ips n) ' ?oreachable n"
  by (intro image_mono subsetI) (rule oreachable_init)
  ultimately show "netgmap sr s ∈ netmask (net_tree_ips n) ' ?oreachable n"
  by (rule set_rev_mp)
next
  fix s a s'
  assume "s ∈ reachable (closed (pnet np n)) TT"
  and "netgmap sr s ∈ netmask (net_tree_ips n) ' ?oreachable n"
  and "(s, a, s') ∈ trans (closed (pnet np n))"
  from this(2) obtain σ ζ where "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
  and "(σ, ζ) ∈ ?oreachable n"
  by clarsimp
  from (s ∈ reachable (closed (pnet np n)) TT) this(1) (wf_net_tree n)
  and ((s, a, s') ∈ trans (closed (pnet np n)))
  obtain σ' ζ' where "((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet onp n))"
  and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"
  by (rule transfer_action)
  from ((σ, ζ) ∈ ?oreachable n) and this(1) have "(σ', ζ') ∈ ?oreachable n"
  by (rule oreachable_local) simp
  with (netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'))
  show "netgmap sr s' ∈ netmask (net_tree_ips n) ' ?oreachable n" by (rule image_eqI)
qed

definition
  someinit :: "nat ⇒ 'g"

```

where

```
"someinit i ≡ SOME x. x ∈ (fst o sr) ‘ init (np i)"
```

definition

```
initmissing :: "(nat ⇒ 'g option) × 'a ⇒ (nat ⇒ 'g) × 'a"
```

where

```
"initmissing σ = (λi. case (fst σ) i of None ⇒ someinit i | Some s ⇒ s, snd σ)"
```

lemma initmissing\_def':

```
"initmissing = apfst (default someinit)"
```

```
by (auto simp add: initmissing_def default_def)
```

lemma netmask\_initmissing\_netgmap:

```
"netmask (net_ips s) (initmissing (netgmap sr s)) = netgmap sr s"
```

```
proof (intro prod_eqI ext)
```

```
fix i
```

```
show "fst (netmask (net_ips s) (initmissing (netgmap sr s))) i = fst (netgmap sr s) i"
```

```
unfolding initmissing_def by (clarsimp split: option.split)
```

```
qed (simp add: initmissing_def)
```

lemma snd\_initmissing [simp]:

```
"snd (initmissing x) = snd x"
```

```
unfolding initmissing_def by simp
```

lemma initmissing\_snd\_netgmap [simp]:

```
assumes "initmissing (netgmap sr s) = (σ, ζ)"
```

```
shows "snd (netgmap sr s) = ζ"
```

```
using assms unfolding initmissing_def by simp
```

lemma in\_net\_ips\_fst\_init\_missing [simp]:

```
assumes "i ∈ net_ips s"
```

```
shows "fst (initmissing (netgmap sr s)) i = the (fst (netgmap sr s) i)"
```

```
using assms unfolding initmissing_def by (clarsimp split: option.split)
```

lemma not\_in\_net\_ips\_fst\_init\_missing [simp]:

```
assumes "i ∉ net_ips s"
```

```
shows "fst (initmissing (netgmap sr s)) i = someinit i"
```

```
using assms unfolding initmissing_def by (clarsimp split: option.split)
```

lemma initmissing\_oreachable\_netmask [elim]:

```
assumes "initmissing (netgmap sr s) ∈ oreachable (oclosed (opnet onp n)) (λ_ _ . True) U"  
(is "_ ∈ ?oreachable n")
```

```
and "net_ips s = net_tree_ips n"
```

```
shows "netgmap sr s ∈ netmask (net_tree_ips n) ‘ ?oreachable n"
```

```
proof -
```

```
obtain σ ζ where "initmissing (netgmap sr s) = (σ, ζ)" by (metis surj_pair)
```

```
with assms(1) have "(σ, ζ) ∈ ?oreachable n" by simp
```

```
have "netgmap sr s = netmask (net_ips s) (σ, ζ)"
```

```
proof (intro prod_eqI ext)
```

```
fix i
```

```
show "fst (netgmap sr s) i = fst (netmask (net_ips s) (σ, ζ)) i"
```

```
proof (cases "i ∈ net_ips s")
```

```
assume "i ∈ net_ips s"
```

```
hence "fst (initmissing (netgmap sr s)) i = the (fst (netgmap sr s) i)"
```

```
by (rule in_net_ips_fst_init_missing)
```

```
moreover from ⟨i ∈ net_ips s⟩ have "Some (the (fst (netgmap sr s) i)) = fst (netgmap sr s) i"
```

```
by (rule some_the_fst_netgmap)
```

```
ultimately show ?thesis
```

```
using ⟨initmissing (netgmap sr s) = (σ, ζ)⟩ by simp
```

```
qed simp
```

```
next
```

```
from ⟨initmissing (netgmap sr s) = (σ, ζ)⟩
```

```

    show "snd (netgmap sr s) = snd (netmask (net_ips s) (σ, ζ))"
      by simp
qed
with assms(2) have "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)" by simp
moreover from ⟨(σ, ζ) ∈ ?oreachable n⟩
  have "netmask (net_ips s) (σ, ζ) ∈ netmask (net_ips s) ‘ ?oreachable n"
    by (rule imageI)
ultimately show ?thesis
  by (simp only: assms(2))
qed

lemma pnet_reachable_transfer:
  assumes "wf_net_tree n"
    and "s ∈ reachable (closed (pnet np n)) TT"
  shows "initmissing (netgmap sr s) ∈ oreachable (oclosed (opnet onp n)) (λ_ _ . True) U"
    (is " _ ∈ ?oreachable n")
using assms(2) proof induction
  fix s
  assume "s ∈ init (closed (pnet np n))"
  hence "s ∈ init (pnet np n)" by simp

  from ⟨wf_net_tree n⟩ have "initmissing (netgmap sr s) ∈ init (opnet onp n)"
  proof (rule init_lifted [THEN set_mp], intro CollectI exI conjI allI)
    show "initmissing (netgmap sr s) = (fst (initmissing (netgmap sr s)), snd (netgmap sr s))"
      by (metis snd_initmissing surjective_pairing)
  next
  from ⟨s ∈ init (pnet np n)⟩ show "s ∈ init (pnet np n)" ..
  next
  fix i
  show "if i ∈ net_tree_ips n
    then (fst (initmissing (netgmap sr s))) i = the (fst (netgmap sr s) i)
    else (fst (initmissing (netgmap sr s))) i ∈ (fst ∘ sr) ‘ init (np i)"
  proof (cases "i ∈ net_tree_ips n", simp_all only: if_True if_False)
    assume "i ∈ net_tree_ips n"
    with ⟨s ∈ init (pnet np n)⟩ have "i ∈ net_ips s" ..
    thus "fst (initmissing (netgmap sr s)) i = the (fst (netgmap sr s) i)" by simp
  next
  assume "i ∉ net_tree_ips n"
  with ⟨s ∈ init (pnet np n)⟩ have "i ∉ net_ips s" ..
  hence "fst (initmissing (netgmap sr s)) i = someinit i" by simp
  moreover have "someinit i ∈ (fst ∘ sr) ‘ init (np i)"
  unfolding someinit_def proof (rule someI_ex)
    from init_notempty show "∃x. x ∈ (fst ∘ sr) ‘ init (np i)" by auto
  qed
  ultimately show "fst (initmissing (netgmap sr s)) i ∈ (fst ∘ sr) ‘ init (np i)"
    by simp
  qed
  qed
  hence "initmissing (netgmap sr s) ∈ init (oclosed (opnet onp n))" by simp
  thus "initmissing (netgmap sr s) ∈ ?oreachable n" ..
next
fix s a s'
assume "s ∈ reachable (closed (pnet np n)) TT"
  and "(s, a, s') ∈ trans (closed (pnet np n))"
  and "initmissing (netgmap sr s) ∈ ?oreachable n"
from this(1) have "s ∈ reachable (pnet np n) TT" ..
hence "net_ips s = net_tree_ips n" by (rule pnet_net_ips_net_tree_ips)
with ⟨initmissing (netgmap sr s) ∈ ?oreachable n⟩
  have "netgmap sr s ∈ netmask (net_tree_ips n) ‘ ?oreachable n"
    by (rule initmissing_oreachable_netmask)

obtain σ ζ where "(σ, ζ) = initmissing (netgmap sr s)" by (metis surj_pair)
with ⟨initmissing (netgmap sr s) ∈ ?oreachable n⟩
  have "(σ, ζ) ∈ ?oreachable n" by simp

```



```

from ⟨(σ, ζ) = initmissing (netgmap sr s)⟩ and ⟨net_ips s = net_tree_ips n⟩ [symmetric]
  have "netgmap sr s = netmask (net_tree_ips n) (σ, ζ)"
    by (clarsimp simp add: netmask_initmissing_netgmap)

with ⟨s ∈ reachable (closed (pnet np n)) TT⟩
  obtain σ' ζ' where "((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet onp n))"
    and "netgmap sr s' = netmask (net_tree_ips n) (σ', ζ'"
  using ⟨wf_net_tree n⟩ and ⟨(s, a, s') ∈ trans (closed (pnet np n))⟩
  by (rule transfer_action)

from ⟨(σ, ζ) ∈ ?oreachable n⟩ have "net_ips ζ = net_tree_ips n"
  by (rule opnet_net_ips_net_tree_ips [OF oclosed_oreachable_oreachable])
with ⟨((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet onp n))⟩
  have "∀j. j ∉ net_tree_ips n ⟶ σ' j = σ j"
    by (clarsimp elim!: ocomplete_no_change)
have "initmissing (netgmap sr s') = (σ', ζ'"
proof (intro prod_eqI ext)
  fix i
  from ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', ζ')⟩
    ⟨∀j. j ∉ net_tree_ips n ⟶ σ' j = σ j⟩
    ⟨(σ, ζ) = initmissing (netgmap sr s)⟩
    ⟨net_ips s = net_tree_ips n⟩
  show "fst (initmissing (netgmap sr s')) i = fst (σ', ζ') i"
    unfolding initmissing_def by simp
next
  from ⟨netgmap sr s' = netmask (net_tree_ips n) (σ', ζ')⟩
    show "snd (initmissing (netgmap sr s')) = snd (σ', ζ'" by simp
qed
moreover from ⟨(σ, ζ) ∈ ?oreachable n⟩ ⟨((σ, ζ), a, (σ', ζ')) ∈ trans (oclosed (opnet onp n))⟩
  have "(σ', ζ') ∈ ?oreachable n"
    by (rule oreachable_local) (rule TrueI)

ultimately show "initmissing (netgmap sr s') ∈ ?oreachable n"
  by simp
qed

definition
  netglobal :: "(nat ⇒ 'g) ⇒ bool" ⇒ 's net_state ⇒ bool"
where
  "netglobal P ≡ (λs. P (fst (initmissing (netgmap sr s))))"

lemma netglobalsimp [simp]:
  "netglobal P s = P (fst (initmissing (netgmap sr s)))"
  unfolding netglobal_def by simp

lemma netglobalE [elim]:
  assumes "netglobal P s"
    and "∧σ. [ P σ; fst (initmissing (netgmap sr s)) = σ ] ⟹ Q σ"
  shows "netglobal Q s"
  using assms by simp

lemma netglobal_weakenE [elim]:
  assumes "p || netglobal P"
    and "∧σ. P σ ⟹ Q σ"
  shows "p || netglobal Q"
  using assms(1) proof (rule invariant_weakenE)
    fix s
    assume "netglobal P s"
    thus "netglobal Q s"
      by (rule netglobalE) (erule assms(2))
  qed

lemma close_opnet:
  assumes "wf_net_tree n"

```

```

    and "oclosed (opnet onp n)  $\models$  ( $\lambda\_ \_ \_ . \text{True}, U \rightarrow$ ) global P"
  shows "closed (pnet np n)  $\models$  netglobal P"
unfolding invariant_def proof
  fix s
  assume "s  $\in$  reachable (closed (pnet np n)) TT"
  with assms(1)
    have "initmissing (netgmap sr s)  $\in$  oreachable (oclosed (opnet onp n)) ( $\lambda\_ \_ \_ . \text{True}$ ) U"
    by (rule pnet_reachable_transfer)
  with assms(2) have "global P (initmissing (netgmap sr s))" ..
  thus "netglobal P s" by simp
qed

```

end

```

locale openproc_parq =
  op?: openproc np onp sr for np :: "ip  $\Rightarrow$  ('s, ('m::msg) seq_action) automaton" and onp sr
  + fixes qp :: "('t, 'm seq_action) automaton"
  assumes init_qp_notempty: "init qp  $\neq$  {}"

```

```

sublocale openproc_parq  $\subseteq$  openproc "λi. np i  $\langle\langle$  qp"
                                     "λi. onp i  $\langle\langle$  i qp"
                                     "λ(p, q). (fst (sr p), (snd (sr p), q))"

```

proof unfold\_locales

```

  fix i
  show "{ (σ, ζ) | σ ζ s. s  $\in$  init (np i  $\langle\langle$  qp)
         $\wedge$  (σ i, ζ) = ((λ(p, q). (fst (sr p), (snd (sr p), q))) s)
         $\wedge$  ( $\forall$ j. j $\neq$ i  $\rightarrow$  σ j  $\in$  (fst o (λ(p, q). (fst (sr p), (snd (sr p), q))))
        ' init (np j  $\langle\langle$  qp)) }  $\subseteq$  init (onp i  $\langle\langle$  i qp)"

```

(is "?S  $\subseteq$  \_")

proof

```

  fix s
  assume "s  $\in$  ?S"
  then obtain σ p lq
    where "s = (σ, (snd (sr p), lq))"
    and "lq  $\in$  init qp"
    and "p  $\in$  init (np i)"
    and "σ i = fst (sr p)"
    and " $\forall$ j. j  $\neq$  i  $\rightarrow$  σ j  $\in$  (fst o (λ(p, q). (fst (sr p), snd (sr p), q)))
        ' (init (np j)  $\times$  init qp)"

```

by auto

from this(5) have " $\forall$ j. j  $\neq$  i  $\rightarrow$  σ j  $\in$  (fst o sr) ' init (np j)"

by auto

with ⟨p  $\in$  init (np i)⟩ and ⟨σ i = fst (sr p)⟩ have "(σ, snd (sr p))  $\in$  init (onp i)"

by - (rule init [THEN set\_mp], auto)

with ⟨lq  $\in$  init qp⟩ have "((σ, snd (sr p)), lq)  $\in$  init (onp i)  $\times$  init qp"

by simp

hence "(σ, (snd (sr p), lq))  $\in$  extg ' (init (onp i)  $\times$  init qp)"

by (rule rev\_image\_eqI) simp

with ⟨s = (σ, (snd (sr p), lq))⟩ show "s  $\in$  init (onp i  $\langle\langle$  i qp)"

by simp

qed

next

fix i s a s' σ σ'

assume "σ i = fst ((λ(p, q). (fst (sr p), (snd (sr p), q))) s)"

and "σ' i = fst ((λ(p, q). (fst (sr p), (snd (sr p), q))) s')"

and "(s, a, s')  $\in$  trans (np i  $\langle\langle$  qp)"

then obtain p q p' q' where "s = (p, q)"

and "s' = (p', q')"

and "σ i = fst (sr p)"

and "σ' i = fst (sr p')"

by (clarsimp split: prod.split\_asm)

from this(1-2) and ⟨(s, a, s')  $\in$  trans (np i  $\langle\langle$  qp)⟩

have "((p, q), a, (p', q'))  $\in$  parp\_sos (trans (np i)) (trans qp)" by simp

hence "((σ, (snd (sr p), q)), a, (σ', (snd (sr p'), q')))  $\in$  trans (onp i  $\langle\langle$  i qp)"

proof cases

```
assume "q' = q"
  and "(p, a, p') ∈ trans (np i)"
  and "∧m. a ≠ receive m"
from ⟨σ i = fst (sr p)⟩ and ⟨σ' i = fst (sr p')⟩ this(2)
  have "((σ, snd (sr p)), a, (σ', snd (sr p'))) ∈ trans (onp i)" by (rule trans)
with ⟨q' = q⟩ and ∧m. a ≠ receive m
  show "((σ, snd (sr p), q), a, (σ', (snd (sr p'), q'))) ∈ trans (onp i ⟨⟨i qp⟩⟩)"
    by (auto elim!: oparleft)
```

next

```
assume "p' = p"
  and "(q, a, q') ∈ trans qp"
  and "∧m. a ≠ send m"
with ⟨σ i = fst (sr p)⟩ and ⟨σ' i = fst (sr p')⟩
  show "((σ, snd (sr p), q), a, (σ', (snd (sr p'), q'))) ∈ trans (onp i ⟨⟨i qp⟩⟩)"
    by (auto elim!: oparright)
```

next

```
fix m
assume "a = τ"
  and "(p, receive m, p') ∈ trans (np i)"
  and "(q, send m, q') ∈ trans qp"
from ⟨σ i = fst (sr p)⟩ and ⟨σ' i = fst (sr p')⟩ this(2)
  have "((σ, snd (sr p)), receive m, (σ', snd (sr p'))) ∈ trans (onp i)"
    by (rule trans)
with ⟨(q, send m, q') ∈ trans qp⟩ and ⟨a = τ⟩
  show "((σ, snd (sr p), q), a, (σ', (snd (sr p'), q'))) ∈ trans (onp i ⟨⟨i qp⟩⟩)"
    by (simp del: step_seq_tau) (rule oparboth)
```

qed

```
with ⟨s = (p, q)⟩ ⟨s' = (p', q')⟩
  show "((σ, snd ((λ(p, q). (fst (sr p), (snd (sr p), q))) s)), a,
    (σ', snd ((λ(p, q). (fst (sr p), (snd (sr p), q))) s'))) ∈ trans (onp i ⟨⟨i qp⟩⟩)"
    by simp
```

next

```
show "∀j. init (np j ⟨⟨ qp ⟩⟩ ≠ {}"
  by (clarsimp simp add: init_notempty init_qp_notempty)
```

qed

end

## 25 Import all AWN-related theories

```
theory AWN_Main
imports AWN_SOS AWN_SOS_Labels OAWN_SOS_Labels AWN_Invariants
        OAWN_Convert OClosed_Transfer
```

begin

end

## 26 Simple toy example

```
theory Toy
imports Main AWN_Main Qmsg_Lifting
begin
```

### 26.1 Messages used in the protocol

```
datatype msg =
  Pkt data ip
  | Newpkt data ip
```

```
instantiation msg :: msg
```

begin

```
definition newpkt_def [simp]: "newpkt ≡ λ(d,did). Newpkt d did"
```

```
definition eq_newpkt_def: "eq_newpkt m  $\equiv$  case m of Newpkt d did  $\Rightarrow$  True | _  $\Rightarrow$  False"
```

```
instance by intro_classes (simp add: eq_newpkt_def)
end
```

```
definition pkt :: "nat  $\times$  nat  $\Rightarrow$  msg"
where "pkt  $\equiv$   $\lambda$ (no, sid). Pkt no sid"
```

```
lemma pkt_simp [simp]:
  "pkt(no, sid) = Pkt no sid"
  unfolding pkt_def by simp
```

```
lemma not_eq_newpkt_pkt [simp]: " $\neg$ eq_newpkt (Pkt no sid)"
  unfolding eq_newpkt_def by simp
```

## 26.2 Protocol model

```
record state =
  id    :: "nat"
  no    :: "nat"
  nhid  :: "nat"
```

```
msg    :: "msg"
num    :: "nat"
sid    :: "nat"
```

```
abbreviation toy_init :: "ip  $\Rightarrow$  state"
```

```
where "toy_init i  $\equiv$  (|
  id = i,
  no = 0,
  nhid = i,

  msg = (SOME x. True),
  num = (SOME x. True),
  sid = (SOME x. True)
|)"
```

```
lemma some_neq_not_eq [simp]: " $\neg$ ((SOME x :: nat. x  $\neq$  i) = i)"
  by (subst some_eq_ex) (metis zero_neq_numeral)
```

```
definition clear_locals :: "state  $\Rightarrow$  state"
```

```
where "clear_locals  $\xi$  =  $\xi$  (|
  msg := (SOME x. True),
  num := (SOME x. True),
  sid := (SOME x. True)
|)"
```

```
lemma clear_locals_but_not_globals [simp]:
```

```
"id (clear_locals  $\xi$ ) = id  $\xi$ "
"no (clear_locals  $\xi$ ) = no  $\xi$ "
"nhid (clear_locals  $\xi$ ) = nhid  $\xi$ "
unfolding clear_locals_def by auto
```

```
definition is_newpkt
```

```
where "is_newpkt  $\xi$   $\equiv$  case msg  $\xi$  of
  Newpkt data did  $\Rightarrow$  {  $\xi$ (num := data) }
  | _  $\Rightarrow$  {}"
```

```
definition is_pkt
```

```
where "is_pkt  $\xi$   $\equiv$  case msg  $\xi$  of
  Pkt num' sid'  $\Rightarrow$  {  $\xi$ ( num := num', sid := sid' ) }
  | _  $\Rightarrow$  {}"
```

```
lemmas is_msg_defs =
```

```

is_pkt_def is_newpkt_def

lemma is_msg_inv_id [simp]:
  "ξ' ∈ is_pkt ξ    ⇒ id ξ' = id ξ"
  "ξ' ∈ is_newpkt ξ ⇒ id ξ' = id ξ"
  unfolding is_msg_defs
  by (cases "msg ξ", clarsimp+)+

lemma is_msg_inv_sid [simp]:
  "ξ' ∈ is_newpkt ξ ⇒ sid ξ' = sid ξ"
  unfolding is_msg_defs
  by (cases "msg ξ", clarsimp+)+

lemma is_msg_inv_no [simp]:
  "ξ' ∈ is_pkt ξ    ⇒ no ξ' = no ξ"
  "ξ' ∈ is_newpkt ξ ⇒ no ξ' = no ξ"
  unfolding is_msg_defs
  by (cases "msg ξ", clarsimp+)+

lemma is_msg_inv_nhid [simp]:
  "ξ' ∈ is_pkt ξ    ⇒ nhid ξ' = nhid ξ"
  "ξ' ∈ is_newpkt ξ ⇒ nhid ξ' = nhid ξ"
  unfolding is_msg_defs
  by (cases "msg ξ", clarsimp+)+

lemma is_msg_inv_msg [simp]:
  "ξ' ∈ is_pkt ξ    ⇒ msg ξ' = msg ξ"
  "ξ' ∈ is_newpkt ξ ⇒ msg ξ' = msg ξ"
  unfolding is_msg_defs
  by (cases "msg ξ", clarsimp+)+

datatype pseqp =
  PToy

fun nat_of_seqp :: "pseqp ⇒ nat"
where
  "nat_of_seqp PToy = 1"

instantiation "pseqp" :: ord
begin
definition less_eq_seqp [iff]: "l1 ≤ l2 = (nat_of_seqp l1 ≤ nat_of_seqp l2)"
definition less_seqp [iff]:    "l1 < l2 = (nat_of_seqp l1 < nat_of_seqp l2)"
instance ..
end

abbreviation Toy
where
  "Toy ≡ λ_. [[clear_locals]] call(PToy)"

fun ΓTOY :: "(state, msg, pseqp, pseqp label) seqp_env"
where
  "ΓTOY PToy = labelled PToy (
    receive(λmsg' ξ. ξ (| msg := msg' |)).
    [[ξ. ξ (|nhid := id ξ|)]
    ( <is_newpkt>
      (
        [[ξ. ξ (|no := max (no ξ) (num ξ)|)]
        broadcast(λξ. pkt(no ξ, id ξ)). Toy()
      )
    )
    ⊕ <is_pkt>
    (
      <ξ. num ξ > no ξ>
      [[ξ. ξ (|no := num ξ|)]
      [[ξ. ξ (|nhid := sid ξ|)]
    )
  )

```

```

        broadcast( $\lambda\xi$ . pkt(no  $\xi$ , id  $\xi$ )). Toy()
     $\oplus$  ( $\xi$ . num  $\xi \leq$  no  $\xi$ )
        Toy()
    )
))"

```

```

declare  $\Gamma_{TOY}$ .simps [simp del, code del]
lemmas  $\Gamma_{TOY}$ .simps [simp, code] =  $\Gamma_{TOY}$ .simps [simplified]

```

```

fun  $\Gamma_{TOY}$ _skeleton
where " $\Gamma_{TOY}$ _skeleton PToy = seqp_skeleton ( $\Gamma_{TOY}$  PToy)"

```

```

lemma  $\Gamma_{TOY}$ _skeleton_wf [simp]:
"wellformed  $\Gamma_{TOY}$ _skeleton"
proof (rule, intro allI)
  fix pn pn'
  show "call(pn')  $\notin$  sterms1 ( $\Gamma_{TOY}$ _skeleton pn)"
  by (cases pn) simp_all
qed

```

```

declare  $\Gamma_{TOY}$ _skeleton.simps [simp del, code del]
lemmas  $\Gamma_{TOY}$ _skeleton.simps [simp, code] =  $\Gamma_{TOY}$ _skeleton.simps [simplified  $\Gamma_{TOY}$ .simps seqp_skeleton.simps]

```

```

lemma toy_proc_cases [dest]:
  fixes p pn
  assumes "p  $\in$  cterms1 ( $\Gamma_{TOY}$  pn)"
  shows "p  $\in$  cterms1 ( $\Gamma_{TOY}$  PToy)"
  using assms
  by (cases pn) simp_all

```

```

definition  $\sigma_{TOY}$  :: "ip  $\Rightarrow$  (state  $\times$  (state, msg, pseq, pseq label) seqp) set"
where " $\sigma_{TOY}$  i  $\equiv$  {(toy_init i,  $\Gamma_{TOY}$  PToy)}"

```

```

abbreviation ptoy
  :: "ip  $\Rightarrow$  (state  $\times$  (state, msg, pseq, pseq label) seqp, msg seq_action) automaton"
where
  "ptoy i  $\equiv$  ( $\mid$  init =  $\sigma_{TOY}$  i, trans = seqp_sos  $\Gamma_{TOY}$   $\mid$ )"

```

```

lemma toy_trans: "trans (ptoy i) = seqp_sos  $\Gamma_{TOY}$ "
  by simp

```

```

lemma toy_control_within [simp]: "control_within  $\Gamma_{TOY}$  (init (ptoy i))"
  unfolding  $\sigma_{TOY}$ _def by (rule control_withinI) (auto simp del:  $\Gamma_{TOY}$ .simps)

```

```

lemma toy_wf [simp]:
"wellformed  $\Gamma_{TOY}$ "
proof (rule, intro allI)
  fix pn pn'
  show "call(pn')  $\notin$  sterms1 ( $\Gamma_{TOY}$  pn)"
  by (cases pn) simp_all
qed

```

```

lemmas toy_labels_not_empty [simp] = labels_not_empty [OF toy_wf]

```

```

lemma toy_ex_label [intro]: " $\exists$  l. l $\in$ labels  $\Gamma_{TOY}$  p"
  by (metis toy_labels_not_empty all_not_in_conv)

```

```

lemma toy_ex_labelE [elim]:
  assumes " $\forall$  l $\in$ labels  $\Gamma_{TOY}$  p. P l p"
  and " $\exists$  p l. P l p  $\implies$  Q"
  shows "Q"
  using assms by (metis toy_ex_label)

```

```

lemma toy_simple_labels [simp]: "simple_labels  $\Gamma_{TOY}$ "

```

```

proof
  fix pn p
  assume "p ∈ subterms(ΓTOY pn)"
  thus "∃!l. labels ΓTOY p = {l}"
  by (cases pn) (simp_all cong: seqp_congs | elim disjE)+
qed

```

```

lemma σTOY_labels [simp]: "(ξ, p) ∈ σTOY i ⇒ labels ΓTOY p = {PToy-:0}"
  unfolding σTOY_def by simp

```

By default, we no longer let the simplifier descend into process terms.

```

declare seqp_congs [cong]

```

```

declare

```

```

  ΓTOY_simps [ctermenv]
  toy_proc_cases [ctermenv_cases]
  seq_invariant_ctermenvI [OF toy_wf toy_control_within toy_simple_labels toy_trans, ctermenv_intros]
  seq_step_invariant_ctermenvI [OF toy_wf toy_control_within toy_simple_labels toy_trans, ctermenv_intros]

```

### 26.3 Define an open version of the protocol

```

definition σOTOY :: "(ip ⇒ state) × ((state, msg, pseqp, pseqp label) seqp) set"
where "σOTOY ≡ {(toy_init, ΓTOY PToy)}"

```

```

abbreviation optoy

```

```

  :: "ip ⇒ ((ip ⇒ state) × (state, msg, pseqp, pseqp label) seqp, msg seq_action) automaton"

```

```

where

```

```

  "optoy i ≡ (| init = σOTOY, trans = oseqp_sos ΓTOY i |)"

```

```

lemma initiali_toy [intro!, simp]: "initiali i (init (optoy i)) (init (ptoy i))"
  unfolding σTOY_def σOTOY_def by rule simp_all

```

```

lemma oaadv_control_within [simp]: "control_within ΓTOY (init (optoy i))"
  unfolding σOTOY_def by (rule control_withinI) (auto simp del: ΓTOY_simps)

```

```

lemma σOTOY_labels [simp]: "(σ, p) ∈ σOTOY ⇒ labels ΓTOY p = {PToy-:0}"
  unfolding σOTOY_def by simp

```

```

lemma otoy_trans: "trans (optoy i) = oseqp_sos ΓTOY i"
  by simp

```

```

declare

```

```

  oseq_invariant_ctermenvI [OF toy_wf oaadv_control_within toy_simple_labels otoy_trans, ctermenv_intros]
  oseq_step_invariant_ctermenvI [OF toy_wf oaadv_control_within toy_simple_labels otoy_trans, ctermenv_intros]

```

### 26.4 Predicates

```

definition msg_sender :: "msg ⇒ ip"

```

```

where "msg_sender m ≡ case m of Pkt _ ipc ⇒ ipc"

```

```

lemma msg_sender_simps [simp]:

```

```

  "∧d did sid. msg_sender (Pkt d sid) = sid"

```

```

  unfolding msg_sender_def by simp_all

```

```

abbreviation not_Pkt :: "msg ⇒ bool"

```

```

where "not_Pkt m ≡ case m of Pkt _ _ ⇒ False | _ ⇒ True"

```

```

definition nos_inc :: "state ⇒ state ⇒ bool"

```

```

where "nos_inc ξ ξ' ≡ (no ξ ≤ no ξ')"

```

```

definition msg_ok :: "(ip ⇒ state) ⇒ msg ⇒ bool"

```

```

where "msg_ok σ m ≡ case m of Pkt num' sid' ⇒ num' ≤ no (σ sid') | _ ⇒ True"

```

```

lemma msg_okI [intro]:
  assumes " $\bigwedge \text{num}' \text{ sid}'. m = \text{Pkt num}' \text{ sid}' \implies \text{num}' \leq \text{no} (\sigma \text{ sid}')$ "
  shows "msg_ok  $\sigma$  m"
  using assms unfolding msg_ok_def
  by (auto split: msg.split)

```

```

lemma msg_ok_Pkt [simp]:
  "msg_ok  $\sigma$  (Pkt data src) = (data  $\leq$  no ( $\sigma$  src))"
  unfolding msg_ok_def by simp

```

```

lemma msg_ok_pkt [simp]:
  "msg_ok  $\sigma$  (pkt(data, src)) = (data  $\leq$  no ( $\sigma$  src))"
  unfolding msg_ok_def by simp

```

```

lemma msg_ok_Newpkt [simp]:
  "msg_ok  $\sigma$  (Newpkt data dst)"
  unfolding msg_ok_def by simp

```

```

lemma msg_ok_newpkt [simp]:
  "msg_ok  $\sigma$  (newpkt(data, dst))"
  unfolding msg_ok_def by simp

```

## 26.5 Sequential Invariants

```

lemma seq_no_leq_num:
  "ptoy i  $\models$  onl  $\Gamma_{TOY} (\lambda(\xi, l). l \in \{\text{PToy-:7..PToy-:8}\} \longrightarrow \text{no } \xi \leq \text{num } \xi)$ "
  by inv_cterms

```

```

lemma seq_nos_incs:
  "ptoy i  $\models_A$  onll  $\Gamma_{TOY} (\lambda((\xi, _), _, (\xi', _)). \text{nos\_inc } \xi \xi')$ "
  unfolding nos_inc_def
  by (inv_cterms inv add: onl_invariant_sterms [OF toy_wf seq_no_leq_num])

```

```

lemma seq_nos_incs':
  "ptoy i  $\models_A$  ( $\lambda((\xi, _), _, (\xi', _)). \text{nos\_inc } \xi \xi')$ "
  by (rule step_invariant_weakenE [OF seq_nos_incs]) (auto dest!: onllD)

```

```

lemma sender_ip_valid:
  "ptoy i  $\models_A$  onll  $\Gamma_{TOY} (\lambda((\xi, _), a, _). \text{anycast } (\lambda m. \text{msg\_sender } m = \text{id } \xi) a)$ "
  by inv_cterms

```

```

lemma id_constant:
  "ptoy i  $\models$  onl  $\Gamma_{TOY} (\lambda(\xi, _). \text{id } \xi = i)$ "
  by inv_cterms (simp add:  $\sigma_{TOY\_def}$ )

```

```

lemma nhid_eq_id:
  "ptoy i  $\models$  onl  $\Gamma_{TOY} (\lambda(\xi, l). l \in \{\text{PToy-:2..PToy-:8}\} \longrightarrow \text{nhid } \xi = \text{id } \xi)$ "
  by inv_cterms

```

```

lemma seq_msg_ok:
  "ptoy i  $\models_A$  onll  $\Gamma_{TOY} (\lambda((\xi, _), a, _).
    \text{anycast } (\lambda m. \text{case } m \text{ of Pkt num}' \text{ sid}' \Rightarrow \text{num}' = \text{no } \xi \wedge \text{sid}' = i \mid \_ \Rightarrow \text{True}) a)$ "
  by (inv_cterms inv add: onl_invariant_sterms [OF toy_wf id_constant])

```

```

lemma nhid_eq_i:
  "ptoy i  $\models$  onl  $\Gamma_{TOY} (\lambda(\xi, l). l \in \{\text{PToy-:2..PToy-:8}\} \longrightarrow \text{nhid } \xi = i)$ "
  proof (rule invariant_arbitraryI, clarify intro!: onll impI)
    fix  $\xi$  p l n
    assume " $(\xi, p) \in \text{reachable } (\text{ptoy } i) \text{ TT}$ "
      and " $l \in \text{labels } \Gamma_{TOY} p$ "
      and " $l \in \{\text{PToy-:2..PToy-:8}\}$ "
    from this(1-3) have "nhid  $\xi = \text{id } \xi$ "
    by - (drule invariantD [OF nhid_eq_id], auto)
  qed

```



```

    moreover with  $\langle (\xi, p) \in \text{reachable } (\text{ptoy } i) \text{ TT} \rangle$  and  $\langle l \in \text{labels } \Gamma_{\text{TOY}} p \rangle$  have "id  $\xi = i$ "
    by (auto dest: invariantD [OF id_constant])
    ultimately show "nhid  $\xi = i$ "
    by simp
qed

```

## 26.6 Global Invariants

```

lemma nos_incD [dest]:
  assumes "nos_inc  $\xi \xi'$ "
  shows "no  $\xi \leq \text{no } \xi'$ "
  using assms unfolding nos_inc_def .

```

```

lemma nos_inc_simp [simp]:
  "nos_inc  $\xi \xi' = (\text{no } \xi \leq \text{no } \xi')$ "
  unfolding nos_inc_def ..

```

```

lemmas oseq_nos_incs =
  open_seq_step_invariant [OF seq_nos_incs initiali_toy otoy_trans toy_trans,
    simplified seqll_onll_swap]

```

```

lemmas oseq_no_leq_num =
  open_seq_invariant [OF seq_no_leq_num initiali_toy otoy_trans toy_trans,
    simplified seql_onl_swap]

```

```

lemma all_nos_inc:
  shows "optoy  $i \models_A (\text{otherwith nos\_inc } \{i\} S,$ 
    other nos\_inc  $\{i\} \rightarrow)$ 
    onll  $\Gamma_{\text{TOY}} (\lambda((\sigma, \_), a, (\sigma', \_)). (\forall j. \text{nos\_inc } (\sigma j) (\sigma' j)))"$ 
  proof -
    have *: " $\bigwedge \sigma \sigma' a. [\text{otherwith nos\_inc } \{i\} S \sigma \sigma' a; \text{no } (\sigma i) \leq \text{no } (\sigma' i)]$ 
       $\implies \forall j. \text{no } (\sigma j) \leq \text{no } (\sigma' j)"$ 
    by (auto dest!: otherwith_syncD)
    show ?thesis
    by (inv_cterms
      inv add: oseq_step_invariant_sterms [OF oseq_nos_incs [THEN oinvariant_step_anyact]
        toy_wf otoy_trans]
      simp add: seqllsimp) (auto elim!: *)
  qed

```

```

lemma oreceived_msg_inv:
  assumes other: " $\bigwedge \sigma \sigma' m. [P \sigma m; \text{other } Q \{i\} \sigma \sigma'] \implies P \sigma' m"$ 
    and local: " $\bigwedge \sigma m. P \sigma m \implies P (\sigma(i := \sigma i(\text{msg} := m))) m"$ 
  shows "optoy  $i \models (\text{otherwith } Q \{i\} (\text{orecvmsg } P), \text{other } Q \{i\} \rightarrow)$ 
    onl  $\Gamma_{\text{TOY}} (\lambda(\sigma, l). l \in \{\text{PToy-:1}\} \longrightarrow P \sigma (\text{msg } (\sigma i)))"$ 
  proof (inv_cterms, intro impI)
    fix  $\sigma \sigma' l$ 
    assume "l = PToy-:1  $\longrightarrow P \sigma (\text{msg } (\sigma i))"$ 
    and "l = PToy-:1"
    and "other  $Q \{i\} \sigma \sigma'$ "
    from this(1-2) have " $P \sigma (\text{msg } (\sigma i))"$  ..
    hence " $P \sigma' (\text{msg } (\sigma i))"$  using  $\langle \text{other } Q \{i\} \sigma \sigma' \rangle$ 
    by (rule other)
    moreover from  $\langle \text{other } Q \{i\} \sigma \sigma' \rangle$  have " $\sigma' i = \sigma i"$  ..
    ultimately show " $P \sigma' (\text{msg } (\sigma' i))"$  by simp
  next
    fix  $\sigma \sigma' \text{msg}$ 
    assume "otherwith  $Q \{i\} (\text{orecvmsg } P) \sigma \sigma' (\text{receive msg})"$ 
    and " $\sigma' i = \sigma i(\text{msg} := \text{msg})"$ 
    from this(1) have " $P \sigma \text{msg}$ "
    and " $\forall j. j \neq i \longrightarrow Q (\sigma j) (\sigma' j)"$  by auto
    from this(1) have " $P (\sigma(i := \sigma i(\text{msg} := \text{msg}))) \text{msg}$ " by (rule local)
    thus " $P \sigma' \text{msg}$ "
  proof (rule other)

```

```

    from ⟨σ' i = σ i(msg := msg)⟩ and ⟨∀j. j≠i → Q (σ j) (σ' j)⟩
      show "other Q {i} (σ(i := σ i(msg := msg))) σ'"
        by - (rule otherI, auto)
  qed
qed

lemma msg_ok_other_nos_inc [elim]:
  assumes "msg_ok σ m"
    and "other nos_inc {i} σ σ'"
  shows "msg_ok σ' m"
proof (cases m)
  fix num sid
  assume "m = Pkt num sid"
  with ⟨msg_ok σ m⟩ have "num ≤ no (σ sid)" by simp
  also from ⟨other nos_inc {i} σ σ'⟩ have "no (σ sid) ≤ no (σ' sid)"
    by (rule otherE) (metis eq_iff nos_incD)
  finally have "num ≤ no (σ' sid)" .
  with ⟨m = Pkt num sid⟩ show ?thesis
    by simp
qed simp

lemma msg_ok_no_leq_no [simp, elim]:
  assumes "msg_ok σ m"
    and "∀j. no (σ j) ≤ no (σ' j)"
  shows "msg_ok σ' m"
using assms(1) proof (cases m)
  fix num sid
  assume "m = Pkt num sid"
  with ⟨msg_ok σ m⟩ have "num ≤ no (σ sid)" by simp
  also from ⟨∀j. no (σ j) ≤ no (σ' j)⟩ have "no (σ sid) ≤ no (σ' sid)"
    by simp
  finally have "num ≤ no (σ' sid)" .
  with ⟨m = Pkt num sid⟩ show ?thesis
    by simp
qed (simp add: assms(1))

lemma oreceived_msg_ok:
  "optoy i ⊨ (otherwith nos_inc {i} (orecvmsg msg_ok),
    other nos_inc {i} →)
    onl ΓTOY (λ(σ, l). l ∈ {PToy-:1..} → msg_ok σ (msg (σ i)))"
(is "_ ⊨ (?S, ?U →) _")
proof (inv_cterms inv add: oseq_step_invariant_sterms [OF all_nos_inc toy_wf otoy_trans],
  intro impI, elim impE)
  fix σ σ'
  assume "msg_ok σ (msg (σ i))"
    and other: "other nos_inc {i} σ σ'"
  moreover from other have "msg (σ' i) = msg (σ i)"
    by (clarsimp elim!: otherE)
  ultimately show "msg_ok σ' (msg (σ' i))"
    by (auto)
next
  fix p l σ a q l' σ' pp p' m
  assume a1: "(σ', p') ∈ oreachable (optoy i) ?S ?U"
    and a2: "PToy-:1 ∈ labels ΓTOY p'"
    and a3: "σ' i = σ i(msg := m)"
  have inv: "optoy i ⊨ (?S, ?U →) onl ΓTOY (λ(σ, l). l ∈ {PToy-:1} → msg_ok σ (msg (σ i)))"
  proof (rule oreceived_msg_inv)
    fix σ σ' m
    assume "msg_ok σ m"
      and "other nos_inc {i} σ σ'"
    thus "msg_ok σ' m" ..
  next
  fix σ m
  assume "msg_ok σ m"

```

```

    thus "msg_ok ( $\sigma(i := \sigma i(\text{msg} := m))$ ) m"
      by (cases m) auto
qed
from a1 a2 a3 show "msg_ok  $\sigma'$  m"
  by (clarsimp dest!: oinvariantD [OF inv] onID)
qed simp

lemma is_pkt_handler_num_leq_no:
  shows "optoy i  $\models$  (otherwith nos_inc {i} (orecvmsg msg_ok),
    other nos_inc {i}  $\rightarrow$ )
    on1  $\Gamma_{TOY} (\lambda(\sigma, l). l \in \{PToy-:6..PToy-:10\} \rightarrow \text{num}(\sigma i) \leq \text{no}(\sigma(\text{sid}(\sigma i))))$ "

proof -
  { fix  $\sigma \sigma'$ 
    assume " $\forall j. \text{no}(\sigma j) \leq \text{no}(\sigma' j)$ "
      and " $\text{num}(\sigma i) \leq \text{no}(\sigma(\text{sid}(\sigma i)))$ "
    have " $\text{num}(\sigma i) \leq \text{no}(\sigma'(\text{sid}(\sigma i)))$ "
    proof -
      note  $\langle \text{num}(\sigma i) \leq \text{no}(\sigma(\text{sid}(\sigma i))) \rangle$ 
      also from  $\langle \forall j. \text{no}(\sigma j) \leq \text{no}(\sigma' j) \rangle$  have " $\text{no}(\sigma(\text{sid}(\sigma i))) \leq \text{no}(\sigma'(\text{sid}(\sigma i)))$ "
      by auto
      finally show ?thesis .
    qed
  } note solve_step = this
show ?thesis
proof (inv_terms inv add: oseq_step_invariant_sterms [OF all_nos_inc toy_wf otoy_trans]
  on1_oinvariant_sterms [OF toy_wf oreceived_msg_ok]
  solve: solve_step, intro impI, elim impE)
  fix  $\sigma \sigma'$ 
  assume *: " $\text{num}(\sigma i) \leq \text{no}(\sigma(\text{sid}(\sigma i)))$ "
    and "other nos_inc {i}  $\sigma \sigma'$ "
  from this(2) obtain " $\forall i \in \{i\}. \sigma' i = \sigma i$ "
    and " $\forall j. j \notin \{i\} \rightarrow \text{nos\_inc}(\sigma j)(\sigma' j)$ " ..
  show " $\text{num}(\sigma' i) \leq \text{no}(\sigma'(\text{sid}(\sigma' i)))$ "
  proof (cases "sid  $\sigma i = i$ ")
    assume "sid  $\sigma i = i$ "
    with *  $\langle \forall i \in \{i\}. \sigma' i = \sigma i \rangle$ 
    show ?thesis by simp
  next
    assume "sid  $\sigma i \neq i$ "
    with  $\langle \forall j. j \notin \{i\} \rightarrow \text{nos\_inc}(\sigma j)(\sigma' j) \rangle$ 
    have " $\text{no}(\sigma(\text{sid}(\sigma i))) \leq \text{no}(\sigma'(\text{sid}(\sigma i)))$ " by simp
    with *  $\langle \forall i \in \{i\}. \sigma' i = \sigma i \rangle$ 
    show ?thesis by simp
  qed
next
  fix p l  $\sigma a q l' \sigma' pp p'$ 
  assume "msg_ok  $\sigma(\text{msg}(\sigma i))$ "
    and " $\forall j. \text{no}(\sigma j) \leq \text{no}(\sigma' j)$ "
    and " $\sigma' i \in \text{is\_pkt}(\sigma i)$ "
  show " $\text{num}(\sigma' i) \leq \text{no}(\sigma'(\text{sid}(\sigma' i)))$ "
  proof (cases "msg  $\sigma i$ ")
    fix num' sid'
    assume "msg  $\sigma i = \text{Pkt num}' \text{sid}'$ "
    with  $\langle \sigma' i \in \text{is\_pkt}(\sigma i) \rangle$  obtain " $\text{num}(\sigma' i) = \text{num}'$ "
      and " $\text{sid}(\sigma' i) = \text{sid}'$ "
    unfolding is_pkt_def by auto
    with  $\langle \text{msg}(\sigma i) = \text{Pkt num}' \text{sid}' \rangle$  and  $\langle \text{msg\_ok} \sigma(\text{msg}(\sigma i)) \rangle$ 
    have " $\text{num}(\sigma' i) \leq \text{no}(\sigma(\text{sid}(\sigma' i)))$ "
    by simp
    also from  $\langle \forall j. \text{no}(\sigma j) \leq \text{no}(\sigma' j) \rangle$  have " $\text{no}(\sigma(\text{sid}(\sigma' i))) \leq \text{no}(\sigma'(\text{sid}(\sigma' i)))$ " ..
    finally show ?thesis .
  next
    fix num' sid'
    assume "msg  $\sigma i = \text{Newpkt num}' \text{sid}'$ "

```

```

with (σ' i ∈ is_pkt (σ i)) have False
  unfolding is_pkt_def by simp
  thus ?thesis ..
qed
qed
qed

```

```

lemmas oseq_id_constant =
  open_seq_invariant [OF id_constant initiali_toy otoy_trans toy_trans,
    simplified seql_onl_swap]

```

```

lemmas oseq_nhid_eq_i =
  open_seq_invariant [OF nhid_eq_i initiali_toy otoy_trans toy_trans,
    simplified seql_onl_swap]

```

```

lemmas oseq_nhid_eq_id =
  open_seq_invariant [OF nhid_eq_id initiali_toy otoy_trans toy_trans,
    simplified seql_onl_swap]

```

```

lemma oseq_bigger_than_next:
  shows "optoy i ⊨ (otherwith nos_inc {i} (orecvmsg msg_ok),
    other nos_inc {i} →) global (λσ. no (σ i) ≤ no (σ (nhid (σ i))))"
  (is "_ ⊨ (?S, ?U →) ?P")
  proof -
    have nhidinv: "optoy i ⊨ (?S, ?U →)
      onl ΓTOY (λ(σ, l). l ∈ {PToy:-2..PToy:-8}
        → nhid (σ i) = id (σ i))"
      by (rule oinvariant_weakenE [OF oseq_nhid_eq_id]) (auto simp: seqlsimp)
    have idinv: "optoy i ⊨ (?S, ?U →) onl ΓTOY (λ(σ, l). id (σ i) = i)"
      by (rule oinvariant_weakenE [OF oseq_id_constant]) (auto simp: seqlsimp)
    { fix σ σ' a
      assume "no (σ i) ≤ no (σ (nhid (σ i)))"
        and "∀j. nos_inc (σ j) (σ' j)"
      note this(1)
      also from (∀j. nos_inc (σ j) (σ' j)) have "no (σ (nhid (σ i))) ≤ no (σ' (nhid (σ i)))"
        by auto
      finally have "no (σ i) ≤ no (σ' (nhid (σ i)))" ..
    } note * = this
    have "optoy i ⊨ (otherwith nos_inc {i} (orecvmsg msg_ok),
      other nos_inc {i} →)
      onl ΓTOY (λ(σ, l). no (σ i) ≤ no (σ (nhid (σ i))))"
    proof (inv_terms
      inv add: onl_oinvariant_sterms [OF toy_wf oseq_no_leq_num [THEN oinvariant_anyact]]
        oseq_step_invariant_sterms [OF all_nos_inc toy_wf otoy_trans]
        onl_oinvariant_sterms [OF toy_wf is_pkt_handler_num_leq_no]
        onl_oinvariant_sterms [OF toy_wf nhidinv]
        onl_oinvariant_sterms [OF toy_wf idinv]
      simp add: seqlsimp seqllsimp
      simp del: nos_inc_simp
      solve: *)
      fix σ p l
      assume "(σ, p) ∈ σOTOY"
      thus "no (σ i) ≤ no (σ (nhid (σ i)))"
        by (simp add: σOTOY_def)
    next
      fix σ σ' p l
      assume or: "(σ, p) ∈ oreachable (optoy i) ?S ?U"
        and "l ∈ labels ΓTOY p"
        and "no (σ i) ≤ no (σ (nhid (σ i)))"
        and "other nos_inc {i} σ σ'"
      show "no (σ' i) ≤ no (σ' (nhid (σ' i)))"
      proof (cases "nhid (σ' i) = i")
        assume "nhid (σ' i) = i"
        with (no (σ i) ≤ no (σ (nhid (σ i)))) show ?thesis

```

```

    by simp
next
  assume "nhid ( $\sigma'$  i)  $\neq$  i"
  moreover from  $\langle$ other nos_inc {i}  $\sigma$   $\sigma'$  $\rangle$  [THEN other_localD] have " $\sigma'$  i =  $\sigma$  i"
  by simp
  ultimately have "no ( $\sigma$  (nhid ( $\sigma$  i)))  $\leq$  no ( $\sigma'$  (nhid ( $\sigma'$  i)))"
  using  $\langle$ other nos_inc {i}  $\sigma$   $\sigma'$  $\rangle$  and  $\langle$  $\sigma'$  i =  $\sigma$  i $\rangle$  by (auto)
  with  $\langle$ no ( $\sigma$  i)  $\leq$  no ( $\sigma$  (nhid ( $\sigma$  i))) $\rangle$  and  $\langle$  $\sigma'$  i =  $\sigma$  i $\rangle$  show ?thesis
  by simp
qed
next
  fix p l  $\sigma$  a q l'  $\sigma'$  pp p'
  assume "no ( $\sigma$  i)  $\leq$  num ( $\sigma$  i)"
    and "num ( $\sigma$  i)  $\leq$  no ( $\sigma$  (sid ( $\sigma$  i)))"
    and " $\forall j$ . nos_inc ( $\sigma$  j) ( $\sigma'$  j)"
  from this(1-2) have "no ( $\sigma$  i)  $\leq$  no ( $\sigma$  (sid ( $\sigma$  i)))"
  by (rule le_trans)
  also from  $\langle$  $\forall j$ . nos_inc ( $\sigma$  j) ( $\sigma'$  j) $\rangle$ 
  have "no ( $\sigma$  (sid ( $\sigma$  i)))  $\leq$  no ( $\sigma'$  (sid ( $\sigma$  i)))"
  by auto
  finally show "no ( $\sigma$  i)  $\leq$  no ( $\sigma'$  (sid ( $\sigma$  i)))" ..
qed
thus ?thesis
  by (rule oinvariant_weakenE)
  (auto simp: onl_def)
qed

lemma anycast_weakenE [elim]:
  assumes "anycast P a"
    and " $\bigwedge m$ . P m  $\implies$  Q m"
  shows "anycast Q a"
  using assms unfolding anycast_def
  by (auto split: seq_action.split)

lemma oseq_msg_ok:
  "optoy i  $\models_A$  (act TT, other U {i}  $\rightarrow$ ) globala ( $\lambda(\sigma, a, \_)$ . anycast (msg_ok  $\sigma$ ) a)"
  by (rule ostep_invariant_weakenE [OF open_seq_step_invariant
    [OF seq_msg_ok initiali_toy otoy_trans toy_trans, simplified seq_onl_swap]])
  (auto simp: seqllsimp dest!: onlD elim!: anycast_weakenE intro!: msg_okI)



## 26.7 Lifting



lemma opar_bigger_than_next:
  shows "optoy i  $\langle\langle$ i qmsg  $\models$  (otherwith nos_inc {i} (orecvmsg msg_ok),
    other nos_inc {i}  $\rightarrow$ ) global ( $\lambda\sigma$ . no ( $\sigma$  i)  $\leq$  no ( $\sigma$  (nhid ( $\sigma$  i)))) $\rangle\rangle$ "
  proof (rule lift_into_qmsg [OF oseq_bigger_than_next])
    fix  $\sigma$   $\sigma'$  m
    assume " $\forall j$ . nos_inc ( $\sigma$  j) ( $\sigma'$  j)"
      and "msg_ok  $\sigma$  m"
    from this(2) show "msg_ok  $\sigma'$  m"
  proof (cases m, simp only: msg_ok_Pkt)
    fix num' sid'
    assume "num'  $\leq$  no ( $\sigma$  sid'"
    also from  $\langle$  $\forall j$ . nos_inc ( $\sigma$  j) ( $\sigma'$  j) $\rangle$  have "no ( $\sigma$  sid')  $\leq$  no ( $\sigma'$  sid'"
    by simp
    finally show "num'  $\leq$  no ( $\sigma'$  sid'" .
  qed simp
next
  show "optoy i  $\models_A$  (otherwith nos_inc {i} (orecvmsg msg_ok), other nos_inc {i}  $\rightarrow$ )
    globala ( $\lambda(\sigma, \_, \sigma')$ . nos_inc ( $\sigma$  i) ( $\sigma'$  i))"
  by (rule ostep_invariant_weakenE [OF open_seq_step_invariant
    [OF seq_nos_incs initiali_toy otoy_trans toy_trans]])
  (auto simp: seqllsimp dest!: onlD)
qed simp

```

lemma onode\_bigger\_than\_next:

```
"⟨i : optoy i ⟨⟨i qmsg : R_i⟩_o
  ⟦ (otherwith nos_inc {i} (oarrivemsg msg_ok), other nos_inc {i} →)
    global (λσ. no (σ i) ≤ no (σ (nhid (σ i)))) ⟧"
by (rule node_lift [OF opar_bigger_than_next])"
```

lemma node\_local\_nos\_inc:

```
"⟨i : optoy i ⟨⟨i qmsg : R_i⟩_o ⟦_A (λσ _ . oarrivemsg (λ_ _ . True) σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, _, σ'). nos_inc (σ i) (σ' i))"
proof (rule node_lift_step_statelessassm)
```

```
  have "optoy i ⟦_A (λσ _ . orecvmsg (λ_ _ . True) σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, _, σ'). nos_inc (σ i) (σ' i))"
  by (rule ostep_invariant_weakenE [OF oseq_nos_incs])
  (auto simp: seqllsimp dest!: onlld)
  thus "optoy i ⟨⟨i qmsg ⟦_A (λσ _ . orecvmsg (λ_ _ . True) σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, _, σ'). nos_inc (σ i) (σ' i))"
  by (rule lift_step_into_qmsg_statelessassm) auto
qed simp
```

lemma opnet\_bigger\_than\_next:

```
"opnet (λi. optoy i ⟨⟨i qmsg⟩ n
  ⟦ (otherwith nos_inc (net_tree_ips n) (oarrivemsg msg_ok),
    other nos_inc (net_tree_ips n) →)
    global (λσ. ∀i∈net_tree_ips n. no (σ i) ≤ no (σ (nhid (σ i))))"
proof (rule pnet_lift [OF onode_bigger_than_next])
```

```
  fix i R_i
  have "⟨i : optoy i ⟨⟨i qmsg : R_i⟩_o ⟦_A (λσ _ . oarrivemsg msg_ok σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, a, _). castmsg (msg_ok σ) a)"
proof (rule node_lift_anycast_statelessassm)
```

```
  have "optoy i ⟦_A (λσ _ . orecvmsg (λ_ _ . True) σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, a, _). anycast (msg_ok σ) a)"
  by (rule ostep_invariant_weakenE [OF oseq_msg_ok]) auto
  hence "optoy i ⟨⟨i qmsg ⟦_A (λσ _ . orecvmsg (λ_ _ . True) σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, a, _). anycast (msg_ok σ) a)"
  by (rule lift_step_into_qmsg_statelessassm) auto
  thus "optoy i ⟨⟨i qmsg ⟦_A (λσ _ . orecvmsg msg_ok σ, other (λ_ _ . True) {i} →)
    globala (λ(σ, a, _). anycast (msg_ok σ) a)"
  by (rule ostep_invariant_weakenE) auto
qed
  thus "⟨i : optoy i ⟨⟨i qmsg : R_i⟩_o ⟦_A (λσ _ . oarrivemsg msg_ok σ, other nos_inc {i} →)
    globala (λ(σ, a, _). castmsg (msg_ok σ) a)"
  by (rule ostep_invariant_weakenE) auto
```

next

```
  fix i R_i
  show "⟨i : optoy i ⟨⟨i qmsg : R_i⟩_o ⟦_A (λσ _ . oarrivemsg msg_ok σ,
    other nos_inc {i} →)
    globala (λ(σ, a, σ'). a ≠ τ ∧ (∀d. a = i:deliver(d)) → nos_inc (σ i) (σ' i))"
  by (rule ostep_invariant_weakenE [OF node_local_nos_inc]) auto
```

next

```
  fix i R
  show "⟨i : optoy i ⟨⟨i qmsg : R⟩_o ⟦_A (λσ _ . oarrivemsg msg_ok σ,
    other nos_inc {i} →)
    globala (λ(σ, a, σ'). a = τ ∨ (∃d. a = i:deliver(d)) → nos_inc (σ i) (σ' i))"
  by (rule ostep_invariant_weakenE [OF node_local_nos_inc]) auto
qed simp_all
```

lemma ocnet\_bigger\_than\_next:

```
"oclosed (opnet (λi. optoy i ⟨⟨i qmsg⟩ n
  ⟦ (λ_ _ _ . True, other nos_inc (net_tree_ips n) →)
    global (λσ. ∀i∈net_tree_ips n. no (σ i) ≤ no (σ (nhid (σ i))))"
proof (rule inclosed_closed)
```

```
  show "opnet (λi. optoy i ⟨⟨i qmsg⟩ n
  ⟦ (otherwith (=) (net_tree_ips n) inclosed, other nos_inc (net_tree_ips n) →)
```

```

      global ( $\lambda\sigma. \forall i \in \text{net\_tree\_ips } n. \text{no } (\sigma \ i) \leq \text{no } (\sigma \ (\text{nhid } (\sigma \ i))))$ )"
proof (rule oinvariant_weakenE [OF opnet_bigger_than_next])
  fix s s' :: "nat  $\Rightarrow$  state" and a :: "msg node_action"
  assume "otherwith (=) (net_tree_ips n) inoclosed s s' a"
  thus "otherwith nos_inc (net_tree_ips n) (oarrivemsg msg_ok) s s' a"
  proof (rule otherwithE, intro otherwithI)
    assume "inoclosed s a"
      and " $\forall j. j \notin \text{net\_tree\_ips } n \longrightarrow s \ j = s' \ j$ "
      and "otherwith ((=)) (net_tree_ips n) inoclosed s s' a"
    thus "oarrivemsg msg_ok s a"
      by (cases a) auto
  qed auto
qed simp
qed

```

## 26.8 Transfer

definition

```

initmissing :: "(nat  $\Rightarrow$  state option)  $\times$  'a  $\Rightarrow$  (nat  $\Rightarrow$  state)  $\times$  'a"

```

where

```

"initmissing  $\sigma = (\lambda i. \text{case } (\text{fst } \sigma) \ i \ \text{of } \text{None} \Rightarrow \text{toy\_init } i \ | \ \text{Some } s \Rightarrow s, \text{snd } \sigma)$ "

```

lemma not\_in\_net\_ips\_fst\_init\_missing [simp]:

```

assumes "i  $\notin$  net_ips  $\sigma$ "

```

```

shows "fst (initmissing (netgmap fst  $\sigma$ )) i = toy_init i"

```

```

using assms unfolding initmissing_def by simp

```

lemma fst\_initmissing\_netgmap\_pair\_fst [simp]:

```

"fst (initmissing (netgmap ( $\lambda(p, q). (\text{fst } (\text{Fun.id } p), \text{snd } (\text{Fun.id } p), q)) \ s)))$ 
  = fst (initmissing (netgmap fst s))"

```

```

unfolding initmissing_def by auto

```

interpretation toy\_openproc: openproc ptoy optoy Fun.id

```

rewrites "toy_openproc.initmissing = initmissing"

```

proof -

```

show "openproc ptoy optoy Fun.id"

```

```

proof unfold_locales

```

```

  fix i :: ip

```

```

  have "{( $\sigma, \zeta$ ). ( $\sigma \ i, \zeta$ )  $\in$   $\sigma_{TOY} \ i \wedge (\forall j. j \neq i \longrightarrow \sigma \ j \in \text{fst } ' \ \sigma_{TOY} \ j)}$ }  $\subseteq \sigma_{OTOY}$ "

```

```

    unfolding  $\sigma_{TOY\_def} \ \sigma_{OTOY\_def}$ 

```

```

    proof (rule equalityD1)

```

```

      show " $\bigwedge f \ p. \{(\sigma, \zeta). (\sigma \ i, \zeta) \in \{(f \ i, p)\} \wedge (\forall j. j \neq i$ 
         $\longrightarrow \sigma \ j \in \text{fst } ' \ \{(f \ j, p)\})\} = \{(f, p)\}$ "

```

```

        by (rule set_eqI) auto

```

```

    qed

```

```

  thus "{( $\sigma, \zeta$ ) |  $\sigma \ \zeta \ s. s \in \text{init } (\text{ptoy } i)$ 

```

```

     $\wedge (\sigma \ i, \zeta) = \text{Fun.id } s$ 

```

```

     $\wedge (\forall j. j \neq i \longrightarrow \sigma \ j \in (\text{fst } o \ \text{Fun.id}) \ ' \ \text{init } (\text{ptoy } j)) \} \subseteq \text{init } (\text{optoy } i)"$ 

```

```

    by simp

```

next

```

  show " $\forall j. \text{init } (\text{ptoy } j) \neq \{\}$ "

```

```

    unfolding  $\sigma_{TOY\_def}$  by simp

```

next

```

  fix i s a s'  $\sigma \ \sigma'$ 

```

```

  assume " $\sigma \ i = \text{fst } (\text{Fun.id } s)$ "

```

```

    and " $\sigma' \ i = \text{fst } (\text{Fun.id } s')$ "

```

```

    and " $(s, a, s') \in \text{trans } (\text{ptoy } i)"$ 

```

```

  then obtain q q' where " $s = (\sigma \ i, q)$ "

```

```

    and " $s' = (\sigma' \ i, q')$ "

```

```

    and " $((\sigma \ i, q), a, (\sigma' \ i, q')) \in \text{trans } (\text{ptoy } i)"$ 

```

```

    by (cases s, cases s') auto

```

```

  from this(3) have " $((\sigma, q), a, (\sigma', q')) \in \text{trans } (\text{optoy } i)"$ 

```

```

    by simp (rule open_seqp_action [OF toy_wf])

```

```

with ⟨s = (σ i, q)⟩ and ⟨s' = (σ' i, q')⟩
  show "((σ, snd (Fun.id s)), a, (σ', snd (Fun.id s'))) ∈ trans (optoy i)"
    by simp
qed
then interpret op0: openproc ptoy optoy Fun.id .
have [simp]: "∧i. (SOME x. x ∈ (fst o Fun.id) ' init (ptoy i)) = toy_init i"
  unfolding σ_TOY_def by simp
hence "∧i. openproc.initmissing ptoy Fun.id i = initmissing i"
  unfolding op0.initmissing_def op0.someinit_def initmissing_def
  by (auto split: option.split)
thus "openproc.initmissing ptoy Fun.id = initmissing" ..
qed

lemma fst_initmissing_netgmap_default_toy_init_netlift:
  "fst (initmissing (netgmap sr s)) = default toy_init (netlift sr s)"
  unfolding initmissing_def default_def
  by (simp add: fst_netgmap_netlift del: One_nat_def)

definition
  netglobal :: "(nat ⇒ state) ⇒ bool ⇒ ((state × 'b) × 'c) net_state ⇒ bool"
where
  "netglobal P ≡ (λs. P (default toy_init (netlift fst s)))"

interpretation toy_openproc_par_qmsg: openproc_parq ptoy optoy Fun.id qmsg
rewrites "toy_openproc_par_qmsg.netglobal = netglobal"
  and "toy_openproc_par_qmsg.initmissing = initmissing"
proof -
  show "openproc_parq ptoy optoy Fun.id qmsg"
    by (unfold_locales) simp
  then interpret opq: openproc_parq ptoy optoy Fun.id qmsg .

  have im: "∧σ. openproc.initmissing (λi. ptoy i ⟨⟨ qmsg⟩ (λ(p, q). (fst (Fun.id p), snd (Fun.id p), q))⟩ σ
                                                    = initmissing σ"
    unfolding opq.initmissing_def opq.someinit_def initmissing_def
    unfolding σ_TOY_def σ_QMSG_def by (clarsimp cong: option.case_cong)
  thus "openproc.initmissing (λi. ptoy i ⟨⟨ qmsg⟩ (λ(p, q). (fst (Fun.id p), snd (Fun.id p), q))⟩ = initmissing
    by (rule ext)

  have "∧P σ. openproc.netglobal (λi. ptoy i ⟨⟨ qmsg⟩ (λ(p, q). (fst (Fun.id p), snd (Fun.id p), q))⟩
P σ
                                                    = netglobal P σ"
    unfolding opq.netglobal_def netglobal_def opq.initmissing_def initmissing_def opq.someinit_def
    unfolding σ_TOY_def σ_QMSG_def
    by (clarsimp cong: option.case_cong
        simp del: One_nat_def
        simp add: fst_initmissing_netgmap_default_toy_init_netlift
        [symmetric, unfolded initmissing_def])
  thus "openproc.netglobal (λi. ptoy i ⟨⟨ qmsg⟩ (λ(p, q). (fst (Fun.id p), snd (Fun.id p), q))⟩ = netglobal
    by auto
qed

```

## 26.9 Final result

```

lemma bigger_than_next:
  assumes "wf_net_tree n"
  shows "closed (pnet (λi. ptoy i ⟨⟨ qmsg⟩ n) ⟩ = netglobal (λσ. ∀i. no (σ i) ≤ no (σ (nhid (σ i))))"
    (is "_ = netglobal (λσ. ∀i. ?inv σ i)")
proof -
  from ⟨wf_net_tree n⟩
  have proto: "closed (pnet (λi. ptoy i ⟨⟨ qmsg⟩ n)
    = netglobal (λσ. ∀i∈net_tree_ips n. no (σ i) ≤ no (σ (nhid (σ i))))"
    by (rule toy_openproc_par_qmsg.close_opnet [OF _ ocnnet_bigger_than_next])
  show ?thesis

```



```

unfolding invariant_def opnet_sos.opnet_tau1
proof (rule, simp only: toy_openproc_par_qmsg.netglobalsimp
                    fst_initmissing_netgmap_pair_fst, rule allI)
  fix  $\sigma$  i
  assume sr: " $\sigma \in \text{reachable} (\text{closed} (\text{pnet} (\lambda i. \text{ptoy } i \langle\langle \text{qmsg} \rangle \rangle n)) \text{TT})$ "
  hence " $\forall i \in \text{net\_tree\_ips } n. ?\text{inv} (\text{fst} (\text{initmissing} (\text{netgmap } \text{fst } \sigma))) i$ "
    by - (drule invariantD [OF proto],
          simp only: toy_openproc_par_qmsg.netglobalsimp
                    fst_initmissing_netgmap_pair_fst)
  thus "?inv (fst (initmissing (netgmap fst  $\sigma$ ))) i"
  proof (cases "i ∈ net_tree_ips n")
    assume "i ∉ net_tree_ips n"
    from sr have " $\sigma \in \text{reachable} (\text{pnet} (\lambda i. \text{ptoy } i \langle\langle \text{qmsg} \rangle \rangle n) \text{TT})$ " ..
    hence "net_ips  $\sigma = \text{net\_tree\_ips } n$ " ..
    with (i ∉ net_tree_ips n) have "i ∉ net_ips  $\sigma$ " by simp
    hence "(fst (initmissing (netgmap fst  $\sigma$ ))) i = toy_init i"
      by simp
    thus ?thesis by simp
  qed metis
qed
qed
end

```

## 27 Acknowledgements

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## References

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